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# New Algorithms for Multivalued Component Trees Nicolas Passat<sup>1</sup>, Romain Perrin<sup>2</sup>, Jimmy Francky Randrianasoa<sup>2,3</sup>, Camille Kurtz<sup>4</sup>, Benoît Naegel<sup>2</sup>

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#### Motivation

The component tree (CT) [3] can model grey-level images for various image processing / analysis purposes (filtering, segmentation, registration, retrieval...). Its generalized version, the multivalued component tree (MCT) [1] can model images with hierarchically organized values. We provide new tools to handle MCTs:

- a new algorithm for the construction of MCTs;
- two strategies for building hierarchical orders on values, required to further build MCTs.

### Multivalued Component Tree

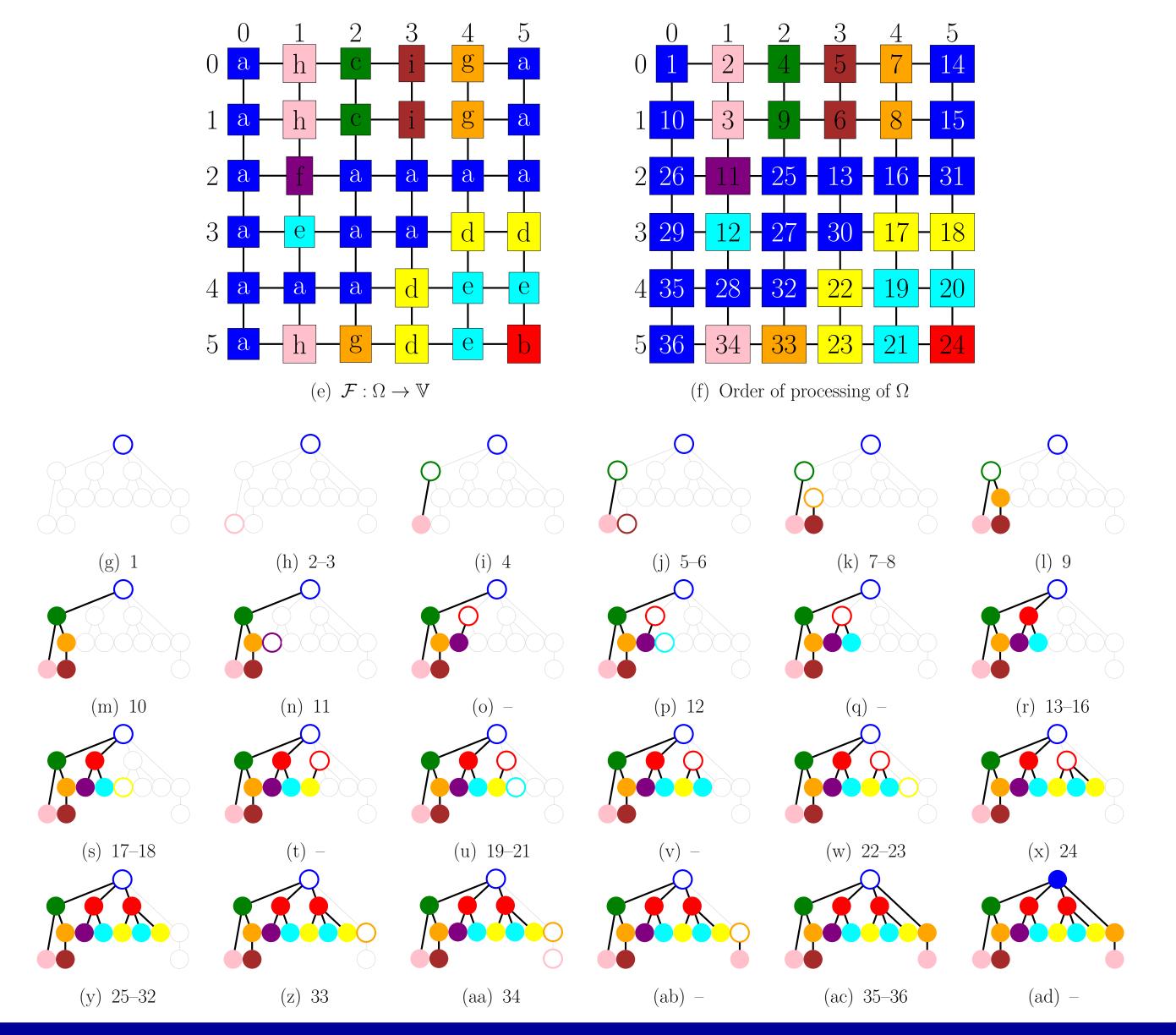
Let  $\mathfrak{G} = (\Omega, \gamma)$  be a non-directed graph. Let  $\mathcal{C}[X]$  be the set of the connected components of  $X \subseteq \Omega$ . Let  $\mathbb{V}$  be a finite set and  $\leq$  a hierarchical order on  $\mathbb{V}$ , i.e. an order (1) which admits a minimum (resp. a maximum) and (2) such that for any  $v \in \mathbb{V}$ , the subset of the elements lower (resp. greater) than v is totally ordered by  $\leq$ .

Let us consider an image  $\mathcal{F}: \Omega \to \mathbb{V}$ . The threshold set of  $\mathcal{F}$  at value  $v \in \mathbb{V}$  is defined by  $\Lambda_{v}(\mathcal{F}) = \mathbb{V}$  $\{\mathbf{x} \in \Omega \mid v \leq \mathcal{F}(\mathbf{x})\}$ . We define the set of nodes of the MCT as

$$\Theta = \left[ -\frac{1}{C} \left[ \Lambda_{\rm v}(\mathcal{F}) \right] \right] \tag{1}$$

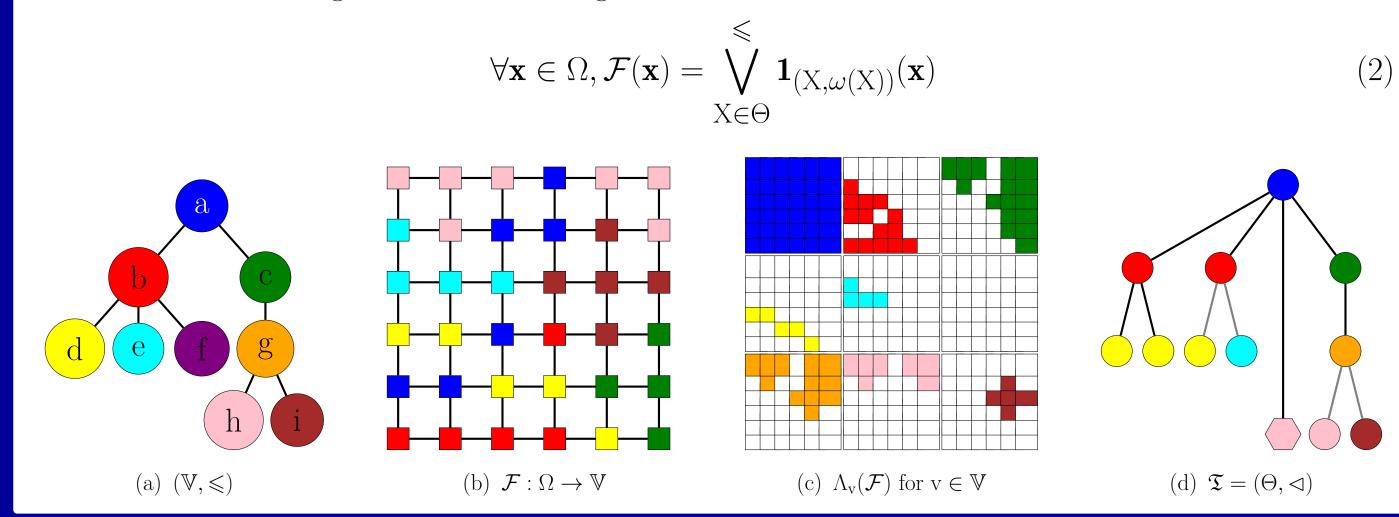
## Running the Construction Algorithm

Construction of the MCT of an image  $\mathcal{F}: \Omega \to \mathbb{V}$ . At a current stage: a plain coloured node is fully built; a contour-colored node is under construction; a non-colored node has not been considered yet; a black edge is built; a light gray edge is not built.



$$\bigcup_{v \in \mathbb{V}} \mathbb{V}^{[1,1]}(\mathbf{v})$$

with  $\mathbb{I}(X) = \{ v \in \mathbb{V} \mid X \in \mathcal{C}[\Lambda_v(\mathcal{F})] \}, \ \omega(X) = \bigvee^{\leq} \mathbb{I}(X) \text{ and } \tau(X) = |\mathbb{I}(X)| \text{ for any node } X \in \Theta.$ The inclusion relation  $\subseteq$  is a hierarchical order on  $\Theta$ . Let  $\triangleleft$  be the reflexive-transitive reduction of  $\subseteq$ . The Hasse diagram  $\mathfrak{T} = (\Theta, \triangleleft)$  of  $(\Theta, \subseteq)$  is the multivalued component tree of the image  $\mathcal{F}$ . For any node  $X \in \Theta$ , we set  $\rho(X) = X \setminus \bigcup_{Y \triangleleft X} Y = \{ \mathbf{x} \in \Omega \mid \mathcal{F}(\mathbf{x}) = \omega(X) \}.$ The MCT  $\mathfrak{T}$  is an image model of the image  $\mathcal{F}$ :



#### Building the Multivalued Component Tree

This construction algorithm derives from the CT construction of [3].

- **nodes**: stores the nodes of the multivalued component tree.
- **points**: stores the processed points of the image.
- **status**: stores the status of each point of the image.
- **nb\_nodes** and **index**: store the number of nodes already fully built and the index of the node currently built at each value of  $\mathbb{V}$ .
- **progress**: indicates if there exists a node at value v, currently under construction or to be built,

#### Hierarchical order construction (1/2): (Pre)ordering the value set

Building the MCT requires a hierarchical order on  $\mathbb{V}$ .

First, we can build a hierarchical preorder  $\leq_{\mathbb{V}}$  on  $\mathbb{V}$ . This can be done by building the CT of an "image" composed by the value set. It is only required that  $\mathbb{V}$  be endowed with:

- an adjacency  $\frown_{\mathbb{V}}$ , allowing to map a graph structure on  $\mathbb{V}$ ;
- a function  $\delta_{\mathbb{V}} : \mathbb{V} \to \mathbb{N}$ , allowing to associate to each element of  $\mathbb{V}$  a value within the totally ordered set  $(\mathbb{N}, \leq)$ .

The CT of  $\delta_{\mathbb{V}}$  is a hierarchical preorder  $\leq_{\mathbb{V}}$  on  $\mathbb{V}$ .

which is an ancestor of the node at value u currently being defined.

Algorithm 1: Build the multivalued component tree

```
Input: (\Omega, \frown), (\mathbb{V}, \leqslant), \mathcal{F} : \Omega \to \mathbb{V}
Output: \mathfrak{T} = (\Theta, \triangleleft)
Build nodes, points, status, nb_nodes, index, progress
v_{\min} := \bigwedge^{\leqslant} \mathbb{V}
Choose \mathbf{x}_{\min} \in \Omega such that \mathcal{F}(\mathbf{x}_{\min}) = v_{\min}
points[v_{min}].add(x_{min})
progress[v_{min}] := true
Flood(v_{min})
```

#### Function Flood

Input:  $u \in V$ : current level

Output:  $w \in \mathbb{V}$ : value of the parent node of the root of the built (partial) MCT at value u

```
while !(points[u].empty()) do
```

```
x := points[u].remove()
```

```
if index[u] > nb_nodes[u] then
```

```
nb_nodes[u] := index[u]
```

```
X := create_node()
```

```
nodes[u].insert(X)
```

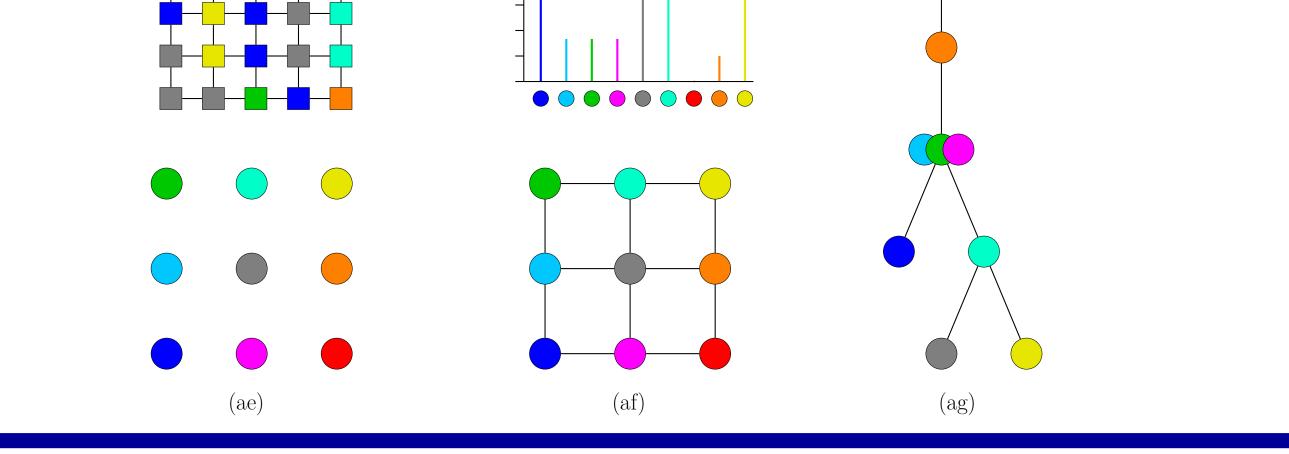
```
if \mathcal{F}(\mathbf{x}) \neq u then
```

```
\mathbf{w} := \mathcal{F}(\mathbf{x})
points[w].add(x)
progress[w] := true
```

```
while u < w do w := Flood(w)
```

#### else

```
status[x] := index[u]
nodes[u][index[u]].add_to_proper_part(x)
foreach \mathbf{y} \sim \mathbf{x} do
    \mathbf{w} := \mathcal{F}(\mathbf{y})
```

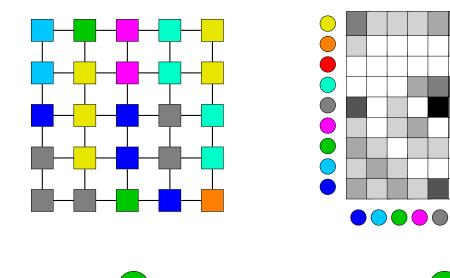


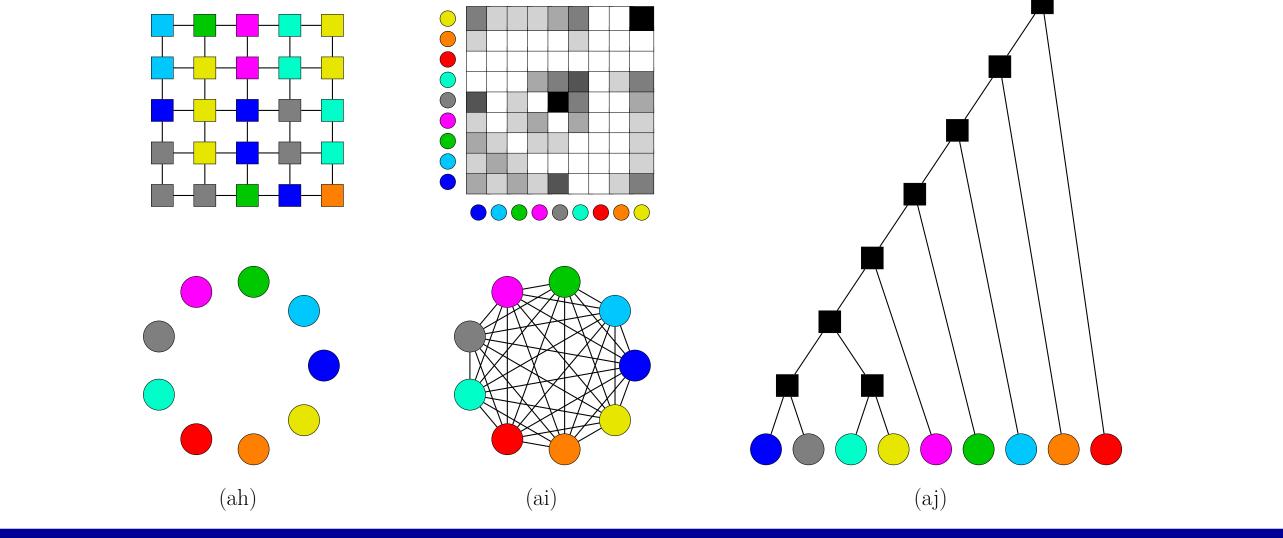
#### Hierarchical order construction (2/2): Ordering the enriched value set

Second, we can build hierarchical order  $\leq_{\mathbb{W}}$  on  $\mathbb{W} \supset \mathbb{V}$  so that the elements of  $\mathbb{V}$  are the maximal elements with respect to  $\leq_W$ .

This can be done by building the binary partition tree (BPT) [2] of an "image" composed by the value set  $\leq_{\mathbb{V}}$ . It is only required that  $\mathbb{V}$  be endowed with:

- an adjacency  $\frown_{\mathbb{V}}$ , allowing to map a graph structure on  $\mathbb{V}$ ;
- a priority function  $\delta_{\gamma_{\mathbb{W}}} : \gamma_{\mathbb{V}} \to \mathbb{N}$ , allowing to determine the couples of nodes to be merged in priority.





# if status[y] = -1 then if $u \leq w$ then $\widehat{w} := w$ else $\widehat{\mathbf{w}} := \bigwedge^{\leqslant} \{\mathbf{u}, \mathbf{w}\}$ $points[\widehat{w}].add(y)$ status[y] := 0 $progress[\widehat{w}] := true$ while $u < \widehat{w}$ do $\widehat{w} := Flood(\widehat{w})$ if $u = v_{\min}$ then $W := \varepsilon$ else $\mathbf{w} := \bigvee^{\leqslant} \{ \mathbf{w'} \in \mathbb{V} \mid \mathbf{w'} < \mathbf{u} \}$ while *progress[w]* = false do $w := \bigvee^{\leq} \{w' \in \mathbb{V} \mid w' < w\}$ create\_edge(nodes[u][index[u]],nodes[w][index[w]]) progress[u] = falseindex[u]++

// new edge in  $\triangleleft$ 

// new node in  $\Theta$ 

#### References

- [1] C. Kurtz, B. Naegel, and N. Passat. Connected filtering based on multivalued component-trees. IEEE Transactions on Image Processing, 23:5152-5164, 2014.
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- [3] P. Salembier, A. Oliveras, and L. Garrido. Anti-extensive connected operators for image and sequence processing. IEEE Transactions on Image Processing, 7:555–570, 1998.

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