Proof assistants for teaching: a survey

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Motivations

Starting point

phD thesis about the use of proof assistants for teaching

• As far as we know : no survey ?

| | Theorem provers | | | | | | |
|------------------------------|-------------------------|------------------|--|--|--|--|--|
| Computer Science World | Interactive (ITP) | Automatic(ATP) | | | | | |
| World | Use of ITP for teaching | ATP for teaching | | | | | |

| | Theorem prove | ers | |
|------------------------------|-------------------------------------|----------------------------------|--|
| Computer Science World | Interactive (ITP) | Automatic(ATP) | |
| | Use of ITP for teaching | ATP for teaching | |
| Math | interactive proof Software for t | automatic proof eaching proof | |
| Educ | Intelligent Tutorir | oring Systems | |
| World | Digital Media in | Educ | |





Contents



- 2 Different categories of proof assistants for teaching
- Input language
- 4 Feedback
- **5** Underlying theory
- 6 Teaching environment
- Conclusion
- 8 Bibliography











Different categories of proof assistants for teaching

- Dedicated proof assistants : teaching community (Didactics, ITS)
 - For logic and undergraduate maths : CPT, CalcCheck, Tutch, XBarnacle/CLAM, WinKE, ProveEasy, Jape, ProofLab, Omega Tutor, PhoX, TPE (Epgy), EasyProve, Lurch, Diproche, Edukera
 - For Geometry: AgentGeom, Baghera, Chypre, Cabri Euclide, Geometrix, Geometry Tutor, Geometry Explanation Tutor, Angle, Mentonizeh, QED-Tutrix, Turing, Advanced Geometry Tutor, Geometrix

• General purpose proof assistants : CS community

Mizar, Coq, Isabelle/Isar, Lean, HOL-Light, ...

• Adaptations of general purpose proof assistants for general use

- On top of Coq : CoqIDE, jsCoq, ProofGeneral
- On top of Isabelle : XIsabelle, Isabelle/jEdit, ProofGeneral
- On top of Lean : Lean Web Editor, ProofWidgets, Paperproof
- Adaptations of general purpose proof assistants for teaching
 - On top of Coq : PCoq, Papuq, CoqWeb, Waterproof, ProofWeb, TryLogic, Geoview, GeoProof
 - On top of Lean : Deaduction, Lean Verbose
 - Other : ETPS on top of TPS

(The Automath language) is difficult to write and to read. It's true, but - in the sense of programming languages - we can compare it to machine language, and we hope that we will manage to build simpler languages, that can be written and read more easily, and that can be translated by a computer in our "machine language". De Bruijn 1969

• Bare proof terms

4. Chaque suite constante est convergente.

| 4.11 | n | | aa | := | $w(p,c,delta,n_o,n)$ | bool | |
|------|----------------|---|----|-----|------------------------------------|-----------------|-----------------------|
| 4.12 | n | | b | := | then 2([u,TRUE(a)] abbrev 2,[u,TR | UE(a)]then) TR | WE(aa) |
| 4.13 | n _o | | h | := | then 2([s,nat] TRUE(aa(s)), [s,na | t] b(s)) | |
| | | | | | | TRUE(z(p,c,de | lta,n _o)) |
| 4.14 | 8.55 | d | | := | then 13*(nat,[x,nat]z(p,c,delta,1 |),1,h(1)) | |
| | | | | | | TRUE(y(p,c,de | lta)) |
| 4.15 | delta | e | | := | then 2([x,abbrev 1]TRUE(y(p,c,del | ta)), | |
| | | | | | [x,abbrev 1]d(x)) | TRUE(q(p,c,de | lta)) |
| 4.16 | c | f | | i ÷ | then 2([s,real]TRUE(q(p,c,s)),[s,: | real]e(s)) TRU | E(lim(p,c)) |
| 4.17 | с | g | | := | then 13*(real,[s,real]lim(p,s),c, | f) TRUE(conv(p) |) |

Proof terms in Automath language (De Bruijn 1969).

• Bare proof terms

```
Lemma and iff of imp : \forall P Q :Prop, (P \rightarrow Q) \rightarrow ((P \land Q) \leftrightarrow P) :=
  \lambda P Q : Prop \mapsto
     \lambda h PimpQ : P \rightarrow Q \mapsto
      Iff.intro
         \lambda h PandQ : P \wedge Q \mapsto
             (h PandQ.left : P)
          \lambda h P : P \mapsto
            have h Q : Q := h PimpQ h P
           (And.intro (h P : P) (h Q : Q) : (P \land Q) )
```

Lean Proof term mode.

- Bare proof terms
- Core tactic-language

```
1 Lemma and iff of imp : \forall P Q : P, (P → Q) → ((P ∧ Q) ↔ P).
 2 Proof.
 3
     intros.
 4
     constructor.
 5
6
     intro H PandQ.
 7
8
     destruct H PandQ as [HP HQ].
     exact HP.
 9
10
     intro HP.
11
   constructor.
12 exact HP.
13
   apply H.
14 exact HP.
15 Qed.
```

Coq Procedural tactic language.

- Bare proof terms
- Core tactic-language
- Declarative style

```
proof
51
      assume h: "P \rightarrow Q"
52
      have h1: "(P \land Q) \longrightarrow P"
53
      proof
54
        assume "P ∧ Q"
55
        thus "P" ..
56
      ged
57
      have h2: "P \rightarrow (P \wedge Q)"
58
      proof
59
        assume hP : P
60
        with h have 0 ...
61
        with hP show "P \land Q" ...
62
      qed
63
      from h1 and h2 show "((P \land Q) \leftrightarrow P)" by blast
64
    aed
65
```

- Mizar Matuszewski and Rudnicki 2005
- Isabelle/Isar Wenzel 2007

Isabelle/Isar declarative language.

- Bare proof terms
- Core tactic-language
- Declarative style
- Controlled Natural Languages (CNL)

Example Diproche Text 1

Es sei x eine ganze Zahl. Zeige: Wenn x gerade ist, dann ist 2-3*x gerade.

Beweis: Es sei x gerade. Dann gibt es eine ganze Zahl k mit x=2*k. Sei k eine ganze Zahl mit x=2*k. Dann ist 2-3*x=2-3*(2*k)=2*(1-3*k). Also ist 2-3*x gerade. qed []

^{*a*}Let *x* be an integer. Prove: If *x* is even, then 2-3x is even.

Proof: Let *x* be even. Then there is an integer *k* such that x = 2k. Let *k* be an integer with x = 2k. Then we have $2 - 3x = 2 - 3 \cdot (2k) = 2(1 - 3k)$. Hence 2 - 3x is even. qed.

- SAD and ForTheL, Naproche [Lyaletski and Verchinine 2010 Kühlwein et al. 2009]
- CNL as tactic languages
 - $Plat\Omega / \Omega mega$

[Wagner, Autexier, and Benzmüller 2007]

- Diproche [Carl, Lorenzen, and Schmi
- Lean Verbose [Massot 2021]
- Waterproof [Portegies et al. 2022

Diproche

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```
Example "The squeeze theorem."
   Given: (u v w : \mathbb{N} \to \mathbb{R}) (1 : \mathbb{R})
   Assume: (hu : u converges to 1) (hw : w converges to
             (h: \forall n, un < vn) (h': \forall n, vn < wn)
   Conclusion: v converges to 1
 Proof:
   Fix \in > 0
   Since u converges to 1 and \varepsilon > 0 we get N such that
      hN: \forall n > N, |un - 1| < \varepsilon
   Since w converges to 1 and \varepsilon > 0 we get N' such that
      hN' : \forall n \geq N', |w n - 1| \leq \varepsilon
   Let's prove that max N N' works : \forall n \geq max N N', |v
   Fix n > max N N'
   Since n > max N N' we get (hn : n > N) and (hn' : n
   Since \forall n \geq N, |u n - 1| \leq \varepsilon and n \geq N we get
     (hNl : -\varepsilon < u n - 1) and (hNd : u n - 1 < \varepsilon)
   Since \forall n \geq N', |w n - 1| \leq \varepsilon and n \geq N' we get
      (hN'1 : -\varepsilon < w n - 1) and (hN'd : w n - 1 < \varepsilon)
    Lean Verbose
```

- SAD and ForTheL, Naproche [Lyaletski and Verchinine 2010 Kühlwein et al. 2009]
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```
Lemma exercise 1 : 2 is the infimum of [2, 5).
Proof.
We need to show that
                                                               SAD and ForTheL, Naproche
  (2 is a lower bound for [2, 5) \land
                                                                    [Lyaletski and Verchinine 2010
     (for all m : \mathbb{R}, m is a lower bound for [2, 5) \Rightarrow m \leq 2)).
                                                                   Kühlwein et al. 2009]
We show both statements.
 - We need to show that (2 is a lower bound for [2, 5)).
                                                                   CNL as tactic languages
  We need to show that (for all x : \mathbb{R}, x : [2, 5) \Rightarrow 2 \leq x).
  Take x : \mathbb{R}. Assume that (x : [2, 5)).

    PlatΩ / Ωmega

  We conclude that (2 \leq x).
 - We need to show that
                                                                     [Wagner, Autexier, and Benzmüller 2007]
    (for all m : \mathbb{R}, m is a lower bound for [2, 5) \Rightarrow m \leq 2).

    Diproche [Carl, Lorenzen, and Schmi

  Take m : R. Assume that (m is a lower bound for [2, 5)).

    Lean Verbose [Massot 2021]

  It holds that (2 : [2, 5)).
  We conclude that (m \leq 2).
                                                                       • Waterproof [Portegies et al. 2022
```

Qed.

Waterproof

- Bare proof terms
- Core tactic-language
- Declarative style
- Controlled Natural Languages (CNL)
- Point-and-click graphical input

| - | Command + State + Search T | heorems <1> | | | |
|--|------------------------------|-------------|-------|--|--|
| File Display Edit | Selections Editing-Tools Coq | Help | | | |
| Theorem second_degree_decomposition: $\forall a_b \ b_x \ discr! A_discr=a_i < a > b = (discr \ge 0) \Rightarrow$ $x^i + a \cdot x + b = (x + (a + \sqrt{(discr)}) / 2) \cdot (x + (a - \sqrt{(discr)}) / 2).$ $I: ntros a b x discr H^I H0.$ | | | | | |
| Auto 🖂 Do it | < 1 of 1 > | Discard | Abort | | |
| $ \begin{array}{l} R \mid x^{*} + a + x + b = (x + a \mid 2 + \sqrt{[discr] \mid 2}) \cdot (x + (a - \sqrt{[discr]}) \mid 2) \\ H^{i} \circ : (discr \geq 0) \\ H^{i} : (AP \ discr = a^{i} - 4 + b) \\ discr : R \end{array} $ | | | | | |

Figure 3: Notations for real number algebraic calculus.

• CtCoq / PCoq [Bertot 1999; Amerkad et a

CtCoq

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- Bare proof terms
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| A : Prop B : Prop (1/1) A /\ B | Previous step Hints Goal is a conjunction. Both parts should be proved Prove by contradiction. Previous step | Next | CtCoq / PCoq [Bertot 1999; Amerkad et a |
|---|---|------------------------------|---|
| A : Prop B : Prop (1/1) A \/ B | Hirds Geal is a disjunction. Prove the high disjunct Prove that high disjunct Prove that negation of the night disjunct implies the right one Prove that negation of the night disjunct implies the left one Prove by contradiction. Previous step | Next Next Next Next | Predicate Prover [Abrial and Cansell 2003] TPE / Epgy [Sommer and Nuckols 2004] Papuq [Chrzaszcz and Sakowicz 2007] |
| T : Type P : T -> Prop (1/1) exists x:T, P x | Hirts Show an element of type T satisfying P x Apply ex_intro Prove by contradiction. | Next Next Show Next | |

Papuq

- Bare proof terms
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CoqWeb

- Bare proof terms
- Core tactic-language
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- Point-and-click graphical input

| | ProofLab | • • • | |
|---|---|----------------------|--|
| ProofLab Edit Window Help | | 5 | |
| The ProofLab U.S. / East: Set Theory / Set Theoretic Operat / Quantifiers and identity: Practice Problem 4 | ons / Set Theoretic Operations | tranmini 2 | CtCog / PCog [Bertot 1990: Americad et a |
| Derivation Info * Edit * Options * | | Problem Score: 0/1 🗶 | CLOQ / I COQ [Deltor 1999, Allerkau et a |
| Rule Preview Rules Definition Elimination Logic Rules | 1. $(\forall x \in s)x \in b$ 2. $s \subseteq b$ | Prem Defi(Subset | Predicate Prover [Abrial and Cansell 2003] |
| Definition: depends on selectio | . kER 3. s ∈ p(b) | Goal | TPE / Epgy [Sommer and Nuckols 2004] |
| | + ER | • | Papuq [Chrzaszcz and Sakowicz 2007] |
| Forward: | | • | Coqweb [Blanc et al. 2007] |
| 2. Press Apply to derive its definitens. Backward: Special Rules | | • | ProofLab [Sieg 2007] |
| DefE: cannot be applied backward. | E E | | • • • |
| DefE cannot be applied strategically. Goal Creating: Clear A | | | |
| DefE cannot be applied in order to create goals. | | | |
| | | | |

ProofLab

- Bare proof terms
- Core tactic-language
- Declarative style
- Controlled Natural Languages (CNL)
- Point-and-click graphical input

| E | Term edi | tor | | | | | | | |
|----|---------------|----------|-----|---------------|---|--------|--------|---|--|
| | Logic | and sets | Var | iables | | | | • | CtCoq / PCoq [Bertot 1999; Amerkad et a |
| | А | Е | т | \perp | ^ | V | 7 | • | Predicate Prover [Abrial and Cansell 2003] |
| | \Rightarrow | ⇔ | = | ≠ | E | ∉ | С | • | TPE / Epgy [Sommer and Nuckols 2004] |
| | ⊄ | | ⊈ | U | U | | \cap | • | Papuq [Chrzaszcz and Sakowicz 2007] |
| | | Ø | | \mathcal{D} | | | 0 | • | Coqweb [Blanc et al. 2007] |
| | | | ~ | / | | \/ | 0 | • | ProofLab [Sieg 2007] |
| ŀ | ⊐∃A ∀B | (□⇔□) | | | | | | • | EasyProve [Materzok 2015] |
| | | | | | | Cancel | OK | | |
| Ea | syprov | e | | | | | | | |

- Bare proof terms
- Core tactic-language
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Edukera



Edukera

- Bare proof terms
- Core tactic-language
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Feedback

• Basic feedback : proof check, proof navigation, proof state

```
def converges_to (u:sequence) (ℓ : ℝ) : Prop :=
      \forall \ \varepsilon: \mathbb{R}, \ \varepsilon \ > 0 \rightarrow \exists \ n0 \ : \ \mathbb{N}, \ \forall \ n: \mathbb{N}, \ n > n0 \rightarrow |u \ n \ - \ell| < \varepsilon
theorem uniqueness_of_limits (u : \mathbb{N} \to \mathbb{R}) (\ell \ \ell' : \mathbb{R})
 (h : converges_to u \ell) (h': converges_to u \ell') : \ell = \ell' :=
 have h0 : \forall \varepsilon > 0, |\ell - \ell'| < \varepsilon :=
  \lambda (\varepsilon: \mathbb{R}) (h\varepsilon : \varepsilon >0 ) \mapsto
     have h\varepsilon 2 : \varepsilon/2 > 0 := by linarith
     Exists elim (h (\varepsilon/2) h\varepsilon2)
     (
                                               Proof state
        sorry
                                               \mathbf{n}: \mathbb{N} \to \mathbb{R}
                                               ℓ ℓ': ℝ
                                               h: converges_to u \ell
                                               h': converges_to u \ell'
                                               \varepsilon : \mathbb{R}
                                               he: \epsilon > 0
                                               he2; \varepsilon / 2 > 0
                                               \vdash \forall (a : N), (\forall n > a, |u n - \ell| < \varepsilon / 2) \rightarrow |\ell - \ell'| < \varepsilon
```

Feedback

- Basic feedback : proof check, proof navigation, proof state
- Proof planning : hints, analysis of granularity

- Hints, proof planning, proof sketches
 [F. Wiedijk 2003; Huang 1999; Edward William Ayers 2021]
- Proof planning and adaptation of proof step granularity [Schiller, Dietrich, and Benzmüller 2008]

Feedback

- Basic feedback : proof check, proof navigation, proof state
- Proof planning : hints, analysis of granularity
- Error diagnosis

Models simulating student misconceptions :

Brown and Burton 1978; Farrell, Anderson, and Reiser 1984; Zinn 2006

- Carl : anti-ATP (Diproche) : introduction of false rules [Carl 2020]
 - logical fallacies (ex deduce $\exists y, \forall x, \phi \text{ from } \forall x, \exists y, \phi$)
 - false analogies (ex : assume $(a + b)^2 = a^2 + b^2$ by false analogy with $(a + b) \cdot 2 = a \cdot 2 + b \cdot 2$)

Feedback

- Basic feedback : proof check, proof navigation, proof state
- Proof planning : hints, analysis of granularity
- Error diagnosis
- Output in NL + Formatting (LaTeX...) + selectable level of detail

| | | Theorem (continuous_of_dense). Let X be a topological space and let Y be a | Current proof state: |
|---|---------------------------------|--|---|
| | | regular topological space. Let A be a dense subset of $X.$ Let $f:X\to Y$ be a | X is a topological space |
| | | function. Assume that for all elements x of X , f is continuous at x within A . Then | Y is a regular topological spa |
| | | f is continuous. | \boldsymbol{A} is a dense subset of \boldsymbol{X} |
| ٩ | Proverb [Huang 1999] | <i>Proof.</i> \bigcirc ⊕Let x be an element of X . \bigcirc One can see it suffices to prove that for all closed neighborhoods V' of $f(x)$, there exists a neighborhood U of x such that $f[U] \subseteq V'$. \bigcirc \bigcirc Let V' be a closed neighborhood of $f(x)$. \bigcirc \bigcirc We obtain an open | $f: X \rightarrow Y$ For all elements x' of X , f is within A |
| ٩ | Theorema[Buchberger et al. 2016 | neighborhood V of x such that $f[V \cap A] \subseteq V'$. We will show that V is suitable by proving that V is a neighborhood of x and $f[V] \subseteq V'$. By assumption, V is a neighborhood of x . \Box Let z be an element of V . \Box | x is an element of XV' is a closed neighborhood |
| ٥ | Ganesalingam and Gowers | \odot Claim: $f(z_1)$ tends to $f(z)$ as z_1 tends to z within $V \cap A$. | V is an open neighborhood of $f[V \cap A] \subseteq V'$ |
| | experience | • By definition it suffices (1) to prove that $z \in \overline{V \cap A}$ and (2) to prove that for all neighborhoods V'' of $f(z)$, there exists a neighborhood U of z such that | z is an element of V |
| | [Ganesalingam and Gowers 2013] | $ f[U \cap (V \cap A)] \subseteq V''.$ Sixpand all Collapse all nat $z \in \overline{V \cap A}.$ | Goal: $f(z_1)$ tends to $f(z)$ as $z \cap A$ |
| ۲ | [Massot 2023] | | |

• WTT [Nederpelt 2002]

Feedback

- Basic feedback : proof check, proof navigation, proof state
- Proof planning : hints, analysis of granularity
- Error diagnosis

9/27

- Output in NL + Formatting (LaTeX...) + selectable level of detail
- Visual feedback : display of proof trees, underlying objects



Feedback

- Basic feedback : proof check, proof navigation, proof state
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- Proof Visualisation
 - Paperproof
 - [Karunus and Kovsharov 2024]
 - Deaduction [Kerjean et al. 2022]
- Underlying objects

(geometry diagrams, venn diagrams, ...)

ProofWidgets

[Nawrocki, Edward W. Ayers, and Ebner 2023]

Feedback

- Basic feedback : proof check, proof navigation, proof state
- Proof planning : hints, analysis of granularity
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Underlying theory

Fundations

Sentential Logic, FOL, HOL, Dependent Type theory... : consequences on teaching

• Naive Type theory in Papuq [Kozubek and Urzyczyn 2008; Chrzaszcz and Sakowicz 2007]

Covered topics

. . .

logics, sets, relations, functions, sequences, analysis, polynomials,

• Comprehensiveness and adaptation of the libraries

Handling and formalising partial functions and undefined terms

- Proof Obligations in TPE/Epgy [Sommer and Nuckols 2004]
- Partial setoids in Matita [Coen and Zoli 2007]

Teaching environment related features

Type of installation

- Standalone, locally installed
- Plugin (VsCode), locally installed
- Web app, (running on client browser or on server)

Teaching software environment

- Exercise Database (Edukera, ProofLab, Deaduction,)
- Grading feature (edukera)
- Student progress tracking, Moodle access (edukera, trylogic)
- Teacher mode (axiom and rule set customization, hint level parametrization) (waterproof, coqweb, edukera)
- Data collection (proofBuddy Karsten et al. 2023)
- Mixed Document feature (jsCoq, Waterproof)

Use of proof assistants for teaching :

Opportunity for the student to receive quick, frequent, individualized feedback

Use of proof assistants for teaching :

Opportunity for the student to receive quick, frequent, individualized feedback

Logic, CS : Widespread

Undergraduate maths : remained in the community of researchers who developped tools

Use of proof assistants for teaching :

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Why?

| Use of proof assistants for teaching : | | | | |
|---|----------------------------------|--|--|--|
| Opportunity for the student to receive quick, frequent, individualized feedback | | | | |
| Logic, CS : Widespread Undergraduate maths : remained in the community of | | | | |
| Why ? | researchers who developped tools | | | |
| Still Theoretical issues | | | | |
| lack of suitable libraries adapted to the classroom | | | | |
| • formalization : partial functions | | | | |

| Use of proof assistants for teaching : | | | | | |
|--|--|--|--|--|--|
| Opportunity for the student to receive quick, frequent, individualized feedback | | | | | |
| Logic, CS : Widespread Undergraduate maths : remained in the community of Why ? researchers who developped tools | | | | | |
| Still Theoretical issues | | | | | |
| lack of suitable libraries adapted to the classroom formalization : partial functions | | | | | |
| What could help to spread ? | | | | | |
| Develop shared reference resources (like Software Foundations) Favor the development of easy to use tools, with no installation Involve didacticians to evaluate tools in the classroom Collaborate with teachers | | | | | |

- raising awareness of the role of logic in proof
- training in the use of proof assistants

Questions ?

Thank you for your attention.

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