

# **An Assignment-Based Dispatching Heuristic For Agglomeration-Scale AMoD: Sensitivity Analysis Of The Demand Prediction Horizon On The Benefits For Supplier And Customers**

Mélanie Cortina, Nicolas Chiabaut, Ludovic Leclercq

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# TRB Annual Meeting

## An Assignment-Based Dispatching Heuristic For Agglomeration-Scale AMoD: Sensitivity Analysis Of The Demand Prediction Horizon On The Benefits For Supplier And Customers --Manuscript Draft--



## AN ASSIGNMENT-BASED DISPATCHING HEURISTIC FOR AGGLOMERATION-SCALE AMOD: SENSITIVITY ANALYSIS OF THE DEMAND PREDICTION HORIZON ON THE BENEFITS FOR SUPPLIER AND CUSTOMERS Mélanie Cortina, Corresponding Author Univ. Lyon, Univ. Gustave Eiffel, ENTPE, LICIT-ECO7 UMR T 9401, F-69518, LYON, France Tel: +33 4 72 04 77 11 Email: melanie.cortina@entpe.fr Nicolas Chiabaut Citec Ingénieur Conseils SAS F-69342, LYON, France Email: nicolas.chiabaut@citec.ch Ludovic Leclercq Univ. Lyon, Univ. Gustave Eiffel, ENTPE, LICIT-ECO7 UMR T 9401, F-69518, LYON, France Tel: +33 4 72 04 77 16 Email: ludovic.leclercq@univ-eiffel.fr 25 Word Count: 7052 words + 0 table(s)  $\times$  250 = 7052 words Submission Date: July 30, 2022

### ABSTRACT

- As shared mobility solutions are developing, the next generation of e-ride-hailing systems are an-
- nounced to be based on self-driving vehicles. Autonomous Mobility on Demand (AMoD) could
- enhance the level of service experienced by travelers thanks to lower repositioning costs and cen-
- tralized control of compliant vehicles. The dispatching strategy determines system performance
- from AMoD and travelers' points of view. Rule-based heuristics do not exploit the full potential of the fleet because they are myopic to future requests. The rolling horizon approach allows to
- anticipate and make better dispatch decisions. In this study, we question the benefits for AMoD
- and travelers to include a short/mid-term demand prediction in an assignment-based heuristic. We
- define four variants of the assignment-based approach, with order matching only or joint order
- matching + fleet management, with impatient or productivist AVs. In a theoretical agglomeration,
- we test the sensitivity of horizon length on AMoD's profit, order response rate and, travelers' wait-
- ing times under the four strategies. We carefully study the outcomes of the profit-oriented heuristic
- in terms of equity.
- 
- *Keywords*: rolling horizon, dispatching, heuristic, autonomous mobility on demand, morning com-
- mute, equity

#### INTRODUCTION

 Recently, shared mobility solutions are booming and take several forms: carsharing, bikesharing, scooter sharing, ridesharing, ridesourcing, and e-ride-hailing (*[1](#page-27-0)*). The latter service type uses smartphone apps to connect drivers with passengers (e.g., Uber). It involves three stakeholders: the transportation network company (TNC), drivers, and travelers.

 The next generation of e-ride-hailing systems may remove drivers from the equation. Au- tonomous Mobility-on-Demand (AMoD, (*[2](#page-27-1)*)) systems, based on autonomous vehicles (AVs), at- tract investments from car manufacturers, mobility companies and, digital corporations (*[3](#page-27-2)*). Com- pared to traditional e-ride-hailing, AMoD operational cost is awaited to be reduced thanks to sav- ings on wages (*[4](#page-27-3)*). Repositioning is facilitated, and surge-pricing, traditionally operated by TNC to balance supply and demand, is no more relevant with compliant vehicles. Static fares and a sufficient level of service encourage adoption by users for recurrent trips. It is an opportunity to in-crease the vehicle utilization rate and reduce car ownership in urban areas. The centralized control

<span id="page-4-0"></span>of vehicles allows to adopt fleet-wide strategies and make more optimal dispatch decisions.



FIGURE 1: AMoD dispatching.

 Dispatching encompasses fleet management for rebalancing AVs from oversupplied to un- dersupplied areas and order matching for assigning customers to AVs, as shown in figure [1.](#page-4-0) The dispatcher module is the cornerstone of a system composed of AMoD and demand. It is under the management of the private AMoD operator. In practice, operators' first interest may not be to deliver the highest quality of service equally over its coverage area but to maximize its profit. Trav- elers' objective is to minimize their travel costs. The dispatcher determines AMoD and travelers' outcomes by receiving AVs' positions, service plannings, capacities, travelers' requests, timing constraints, and connecting their respective objectives. Understanding the impact of the dispatch- ing strategy chosen on these metrics is worthwhile. Real-time dispatching is the only option when the environment is highly dynamic and

 AMoD has no knowledge about future requests. Rule-based heuristics are used in this case. The most largely used heuristics are presented in (*[5](#page-27-4)*). The *nearest-idle-vehicle* assigns the nearest cur- rently idle vehicle to the request that has just been emitted in the oversupply regime and assigns the oldest request to the vehicle that has just turned available in the undersupply regime. The *nearest- vehicle* matches the nearest vehicle for which future availability time is known, to the oldest open request. The *nearest-idle-vehicle/nearest-open-request* matches the nearest open request with the vehicle that has just turned available in the undersupply regime approaches (*[6](#page-27-5)*). The first heuristic

 is nearly random in the undersupply regime, the second sticks to FIFO rule for requests, the third relaxes FIFO rule but remains myopic. They all proceed to local optimization by considering a single request or vehicle. They try to maximize the number of customers served by chasing each new request without considering the profitability associated with a match.

 To improve dispatching decisions, we need an estimation of future requests via enabling reservation, predicting coming requests with historical data, or both. (*[7](#page-27-6)*) has highlighted the gap between exact offline optimization methods, with full knowledge of demand on the studied period, and rule-based heuristics. In practice for AMoD, the horizon of prediction cannot be so long. Given the uncertainties inherent to on-demand systems, a short/mid-term horizon is more reasonable. The rolling horizon, introduced by (*[8](#page-27-7)*), is the state-of-the-art approach to deal with sequentially growing knowledge about future requests. Exact optimization methods have employed it.

 The most popular formulation enabling exact optimization of dispatch decisions is the Dial- a-Ride Problem (DARP) (*[9](#page-27-8)*). This formulation builds a set of routes to maximize a certain objective function. They are at least two decision variables: binary variables indicating if a vehicle travels between two pick-up/drop-off nodes and the times at which each vehicle starts to serve a node (*[10](#page-27-9)*). The sophisticated algorithms required to solve DARPs lack scalability. They hardly adapt to large-scale and real-time applications in highly dynamic scenarios. Indeed, the more uncertain the prediction on the horizon is, the more often routes should be re-optimized. Notably, (*[11](#page-27-10)*) review refers to (*[12](#page-27-11)*) (300 customers) and (*[13](#page-27-12)*) (900 customers).

 We propose a new assignment-based heuristic using the rolling horizon principle in this study. It is far less myopic than the rule-based heuristics presented above while remaining much simpler and scalable compared to routes optimization addressed by the DARP. Figure [2](#page-5-0) summa- rizes our positioning. Inspired by (*[14](#page-27-13)*), our approach stands out due to the anticipation of future requests. In the latter reference, the dispatcher remains reactive. It reacts to the receipt of newly emitted requests. The availability of vehicles is the only element shortly anticipated. Our approach cyclically recalculates the best assignment of current and future open requests with vehicles that are or will be soon available. At each call, it solves a maximization problem where the objective function reflects its total expected profit on the coming horizon.

<span id="page-5-0"></span>

#### FIGURE 2: Positioning.

29 Primarily, we question the gains emerging from the knowledge of future requests in a the-30 oretical scenario where consequent repositioning efforts are required from AVs. Gains are in-31 vestigated from the AMoD side (profit) and travelers side (availability of service, pick-up time), Cortina, Chiabaut, and Leclercq 5

 considering that the dispatching process favors only the former. On travelers' side, the experi- enced level of service is weighted against traveler's dependency on AMoD, which is approached thanks to the modeling of the whole multimodal network in our simulation platform. Two profit estimation strategies are investigated: one where early arrival of AV at a pick-up point is penalized, and another where the profit earned per elapsed time is preferred. We also explore two types of matching: one with permanent matches on the whole horizon (order matching only), another with permanent matches on the first part of the horizon, and temporary matches for repositioning vehi- cles on pick-up points awaited for the second part of the horizon (joint order matching and fleet management). A sensitivity analysis of the horizon length on these gains under the different dis- patching strategies allows to define the best configuration for each stakeholder. Finally, a pricing scheme to conciliate AMoD and travelers' interests is inspected. Briefly, our contributions are: (i) the implementation of a simple and scalable assignment-

based heuristic configured with four strategies to render AMoD's profit-oriented dispatcher; (ii)

the development of a platform to simulate customers and AMoD interactions, where alternative

itineraries are modeled to assess for travelers willingness to wait for being picked-up; (iii) in a

Manhattan-like network, the sensitivity analysis of horizon length, fleet size and planning length

on travelers-side service efficiency and AMoD-side profits.

#### AN ASSIGNMENT-BASED HEURISTIC ANTICIPATING FUTURE REQUESTS

#### <span id="page-7-1"></span>Hypotheses on AMoD

First, AMoD manages every single AV composing the fleet. The dispatching decisions are made

 in a centralized manner. Considering a centralized dispatcher allows better decisions than a de- centralized system where each AV would be responsible for its routing, relocating, and pick-up actions.

 As announced, we assume AMoD has access to an estimation of future requests. Demand forecast can be achieved by allowing reservation and analyzing historical data (*[15–](#page-27-14)[17](#page-27-15)*). The fore- caster module is out of the scope of this study. We suppose that it provides a perfect knowledge of demand within the coming horizon. As our main focus is the benefits that can emerge from the introduction of a horizon, this hypothesis allows us to draw the upper bound of what each proposed heuristic strategy can achieve.

 Another assumption is that AMoD has no long-term economic strategy. It has a unique fixed fare scheme over the greater urban area. This fare scheme is made to be competitive while ensuring lucrative rides for AVs. It is expected to be more attractive economically for travelers than today's (e-)ride-hailing services. One of the most popular fare scheme in literature (*[18–](#page-28-0)[21](#page-28-1)*) is used here. Values for base, distance-based, and time-based fares are chosen to overcome the production cost (*[22](#page-28-2)*) of any ride an AV performs. Money is effectively earned for a ride.

#### <span id="page-7-0"></span>Algorithm



FIGURE 3: Rolling horizon principle for *permanent* and *temporary* modes.

 Let *t* be the current time. At *t*, AMoD is aware of all requests already emitted and requests that 21 will be emitted between *t* and  $t + H$ . The set of known open requests is noted  $R_H$ . A request is said open when it has not been permanently matched and has not been canceled by traveler yet. Each 23 open request  $j \in R$ *H* is characterized by an earliest pick-up time  $e_j$  (which is equal to the request 24 time in our case) and a latest pick-up time  $l_j$ . A maximum waiting time  $w_j$  is defined for each traveler  $(l_j = e_j + w_j)$ . As we focus on ride-hailing, the capacity of all AVs is set to one. We call *A* 26 the set of available AVs. An AV is said to be available when it has no more than  $K \geq 0$  requests in its service list. The service list corresponds to an activity planning for the AV. It comprises service missions only. The rolling horizon principle is represented on figure [3.](#page-7-0)

 We implement two modes in the algorithm: a permanent matching mode (*permanent*) and, a short-term permanent matching / long-term temporary repositioning mode (*temporary*). In *per-* Cortina, Chiabaut, and Leclercq 7

- 1 *manent* mode, all matches identified by the optimal assignment resolution method are made per-
- 2 manent. In *temporary* mode, only matches implying a request *j* satisfying  $e_i \le t + \alpha H$  ( $\alpha \in [0,1]$ )
- 3 are made permanent. Other matches  $(i, j)$  identified are so that  $e_j > t + \alpha H$ . They lead to the ini-
- 4 tiation of repositioning movements by concerned AVs. A repositioning movement is not a service
- 5 mission. The match is not registered in AV's service list. At the next optimal assignment resolu-
- 6 tion method call, relocating AVs are thus considered available. Note that repositioning movements 7 are only accessible to idle AVs (having no mission in their service list). The *temporary* mode in-
- 8 creases AMoD flexibility by putting back into play a part of the decisions taken. In this mode, the
- 9 dispatcher does achieve order matching jointly with fleet management. In both modes, the horizon
- 10 is rolling, i.e., the optimal assignment algorithm is called each  $\alpha H$ . Since demand prediction is
- 11 perfect, it is unnecessary to call it more often. The dispatch algorithm is detailed below.

## Algorithm 1: Dispatch algorithm

1 Initialize *A* with all AVs,  $R_H$  to an empty list,  $t = 0$ 

- 2 Set horizon to  $[t, t + H]$ , gather requests *j* emitted within the horizon and append them to *RH*, gather AVs *i* with at most *K* requests in their service list to form *A*
- 3 Calculate utilities  $u_{ij}$  for each  $(i, j)$  pair
- 4 Form a bipartite graph with  $u_{ij}$  as cost on arcs and compute the matching maximizing the sum of utilities
	- <sup>5</sup> *permanent*: Make permanent all assignments decided in step 4. *temporary*: Make permanent all assignments decided in step 4 for which  $e_j \in [t, t + \alpha H]$ , and initiate temporary repositioning movements for the other matches decided. Both: Remove permanently matched *j* from *R<sup>H</sup>*
	- 6 Set  $t = t + \alpha H$  and go to step 2.

#### 13 Impatient and productivist utilities

- A request  $j \in R_H$  is characterized by a brut revenue  $p_j$ , which is the price traveler pays to the AV.
- 15 The AMoD fare scheme is defined as:

$$
p_j = f_0^a + f_d^a d_j + f_t^a t_j \tag{1}
$$

16

12

where  $f_0^a$  $\int_0^a$  is the base fare,  $f_d^a$ 17 where  $f_0^a$  is the base fare,  $f_d^a$  is the distance-based fare  $f_t^a$  is the time-based fare,  $d_j$  is the travel 18 distance of the path traveler wants to ride onboard AV,  $t_i$  the travel time on this path.

The pair  $(i, j)$  designates a potential match between *i* and *j*. It is characterized by  $d_{ij}$ , the 20 distance on the shortest path between *i* and *j*'s pick-up point,  $t_{ij}$ , the travel time on this path, and  $u_{ij}$  the utility of match computed in step 3 of the algorithm. We define two strategies AMoD can adopt to maximize its profit. The *impatient* strategy (equation [2\)](#page-8-0) takes into account the repositioning cost of *i* toward *j*'s pick-up point, the expected revenue for serving *j*, the eventual time lost by *i* while 24 waiting for traveler departure  $e_j$ . The waiting time AV experience by arriving early at a pick-up point is penalized. The *productivist* strategy (equation [3\)](#page-9-0) takes into account the repositioning cost and the expected revenue into a profit per time unit. Earning more profit in less time is favored.

<span id="page-8-0"></span>
$$
u_{ij}^{impatient} = -C_d^a d_{ij} + (p_j - C_d^a d_j) \xi_{ij} - f_t^a max(0, e_j - (t_i + t_{ij}))
$$
\n(2)

<span id="page-9-0"></span>
$$
u_{ij}^{productivist} = \frac{-C_d^a d_{ij} + (p_j - C_d^a d_j) \xi_{ij}}{t_{ij} + max(0, e_j - (t_i + t_{ij})) + t_j}
$$
\n
$$
\tag{3}
$$

1

where  $C_d^a$ 2 where  $C_d^a$  designates the distance-based operation cost of an AV,  $t_i$  is the time at which *i* will drop-

3 off the last traveler currently appearing in its service list  $(t_i = t$  if there is no mission in *i*'s service 4 list),  $\xi_{ij}$  equals 1 if  $t_i + t_{ij} \leq l_j$ , 0 otherwise (no profit can be earned if AV arrives too late). Note 5 that AMoD is assumed to know  $l_j$  and  $e_j$ .

#### <span id="page-9-3"></span>6 Optimal assignment

7 In step 4 of the algorithm, the following optimization problem is solved:

$$
\max_{x_{ij}} \sum_{i \in A} \sum_{j \in R_H} u_{ij} x_{ij} \tag{4a}
$$

subject to 
$$
x_{ij} \in \{0, 1\}, \forall i \in A, \forall j \in R_H
$$
 (4b)

$$
\sum_{i \in A} x_{ij} \le 1, \forall j \in R_H
$$
\n<sup>(4c)</sup>

<span id="page-9-2"></span><span id="page-9-1"></span>
$$
\sum_{j \in R_H} x_{ij} \le 1, \forall i \in A
$$
\n
$$
(4d)
$$

8

9 where  $x_{ij}$  are the binary decision variables, equal to 1 if  $(i, j)$  match is decided, 0 otherwise. To prevent unfavorable interactions between two assignments for one AV decided during the same dispatch round, constraint [4c](#page-9-1) imposes that an AV can be assigned to at most one request per optimal assignment resolution call. Indeed, when several matches are allowed for one AV *i*, utility  $u_{ij}$ ,  $j \in$ *R<sub>H</sub>*, is directly dependent on the assignment of other requests to the same AV. Separately, matches  $(i,1)$  and  $(i, 2)$  can have high utilities, but it does not mean that utility of a route including customers 1 and 2 has a good utility. Here is the main difference between routes building and assignment approaches. The former computes routes utilities, the latter computes assignment utilities. Due to constraint [4c,](#page-9-1) there is a risk for under-exploited AVs. Some AVs may finish their mission before the next call of the optimal assignment problem and remain idle for a few minutes. To prevent 19 this phenomenon, we properly set  $\alpha$  to and enable AVs to have more than one service mission in their planning (at most *K*, where missions are appended in different calls of the optimal assignment 21 resolution). The value of  $\alpha$  should not be too small to keep the scalability of the method and not too big to limit the number of lost resources between two calls. It should be consistent with the average riding time of a service mission. Finally, constraint [4d](#page-9-2) ensures that a request is matched with at most one AV.

#### <span id="page-9-4"></span>25 A benchmark heuristic

- 26 To evaluate the benefits of including a demand prediction, we want to compare our assignment-
- 27 based anticipative dispatch with a rule-based myopic heuristic. We implement a *nearest-vehicle/nearest-*
- 28 *open-request* dispatcher. Contrary to our algorithm, which regularly calls the optimal assignment
- 29 resolution, this is an event-based approach. Dispatcher is triggered whenever a new request is
- 30 received in the oversupply regime and whenever an AV turns available (it has just dropped off a
- 31 traveler so that its service list has strictly less than *K* requests) in the undersupply regime. The
- 32 dispatcher is completely myopic to future requests. It does not completely stick to FIFO rule con-
- 33 trary to *nearest-idle-vehicle* and *nearest-vehicle* strategies and, shortly anticipates AVs availability
- contrary to *nearest-idle-vehicle/nearest-open-request* heuristic. The selection is done according to
- the following equations:

$$
i^* = \underset{i \in A}{\operatorname{argmin}} (t_i + t_{ij}) \tag{5}
$$

where 
$$
j
$$
 has just been received.

$$
j^* = \underset{j \in R_H}{\operatorname{argmin}} (t_i + t_{ij}) \tag{6}
$$

where *i* has just turned available.

#### SIMULATING AMOD-TRAVELERS INTERACTIONS

- AMoD and travelers interact within an environment representing a theoretical agglomeration. We
- developed a simple simulation platform that focuses on the elementary bricks required to study the
- impact of the horizon on supplier and customers' benefits.

#### Realistic distances and travel times

Since distances and travel times determine indicators we want to assess, consistent parameters

should be provided to the optimal assignment problem. The simulator takes a road network as an

input. In our theoretical agglomeration, we chose a Manhattan road network with three mesh sizes

- to account for topology differences over the greater urban area. The smaller mesh size stands in the
- 13 center of the city  $(U_1)$ , the medium one in the suburbs  $(U_2)$ , the biggest one in the extended suburbs
- and the close rural area surrounding the city (*U*3, figures [4a](#page-11-0) and [4\)](#page-11-1). Two ring roads are added on
- 15 the boundaries between  $U_1$ - $U_2$  ( $U_{12}$ ) and  $U_2$ - $U_3$  ( $U_{23}$ ). A fixed mean speed ( $v_{U_k}$ ) for cars and AVs
- is defined in each urban zone. Vehicles travel arcs of the network node by node at mean speed.
- Traffic dynamics are not modeled here. In addition, since the number of AVs circulating remains
- limited compared to background traffic, we assume they have no impact on the known mean speed. Adding a macroscopic traffic model such as MFD (*[23](#page-28-3)*) to account for traffic dynamics and AVs'
- 
- contribution to congestion will be studied in future research.

### A consistent demand addressed to AMoD

- Applying an itinerary choice on a global demand pattern allows to better render the distribution
- of pick-up and drop-off points over the agglomeration compared to a random generation. We thus
- model the entire multimodal network.

#### *Multimodal network*

 To represent the different transportation modes including walk (*w*), car (*c*), AV (*a*), train (*r*), sub-27 way (*s*) and bus (*b*), a digraph  $G = (V, A)$ , as the one shown in figure [5,](#page-13-0) is used. *V* and *A* are, respectively, the sets of vertices and arcs. Each mode has an associated layer in this graph. Thus,  $G_w = (V_w, A_w)$  is the walking layer,  $G_c = (V_c, A_c)$  is the personal car layer,  $G_a = (V_a, A_a)$  is the AV 30 layer,  $G_r = (V_r, A_r)$  is the train layer,  $G_s = (V_s, A_s)$  is the subway layer and  $G_b = (V_b, A_b)$  is the bus layer. The walk layer is based on walkable streets. The car and AV layers are based on the road network: vertices and arcs correspond to road intersections and links. The train, subway, and bus layers are based on the public transportation network: vertices and arcs correspond to transit stations and itineraries between two stations. A set of transfer arcs connect the walk layer to the 35 other layers. Transfer arcs types are: *starter* (walk  $\rightarrow$  car), *park* (car  $\rightarrow$  walk), *pick-up* (walk  $\rightarrow$ 

36 Av), *drop-off* (AV  $\rightarrow$  walk), *board* (walk  $\rightarrow$  train, walk  $\rightarrow$  subway, walk  $\rightarrow$  bus), *alight* (train  $\rightarrow$ 

<span id="page-11-1"></span><span id="page-11-0"></span>

#0-0	$#0-1$	$#0-2$	$#0-3$	$U_3$ $#0-4$	train line 15 bus line bhns line subway line transit station ٠ $10$
$#1 - 0$	$#1 - 1$	$#1 - 2$	$\,U_2\,$ $#1-3$	$#1 - 4$	road network 5.
$#2 - 0$	$#2-1$	$U_1$ $#2-2$	$#2-3$	$#2 - 4$	$\forall$ (km) $\mathbf{0}$ nna na n
$#3 - 0$	$#3-1$	$#3-2$	$#3-3$	$#3-4$	$-5$ $-10$
$#4-0$	$#4-1$	$#4-2$	$#4-3$	#4-4	$^{\rm -15}$ $-15$ $-10$ 10 15 $-5$ 5 $\circ$ $x$ (km)
(b) Multimodal network instance. (a) Urban (in colors) and service zoning of the agglomer-					

(a) Urban (in colors) and service zoning of the agglomeration.

<span id="page-11-2"></span>FIGURE 4: A theoretical urban area.

- 1 walk, subway  $\rightarrow$  walk, bus  $\rightarrow$  walk). Another set of transfer arcs of type *alighboard* are internal
- 2 to the urban public transportation network and connect the subway and bus layers.

#### 3 *Travel times*

4 Intra-layer arcs belonging to  $A_c$  and  $A_a$  are parameterized with a mean speed. A fixed walking 5 speed  $v_w$  applies on all walking arcs. Travel times on all arcs in  $A_c \cup A_a \cup A_w$  are deduced from 6 speed and distance. Each transit line is characterized by a cruising speed  $(v_r, v_s, v_b)$ , a headway 7  $(h_r, h_s, h_b)$  and a time lost per station  $(\tau_r, \tau_s, \tau_b)$ . The travel time between two stations of the 8 same transit line is the sum of time lost per stop and travel time at cruising speed. The travel time 9 between two stations of different transit lines located at the same coordinates is approximated to 10 be half the headway, such as the travel time between a walk node and a transit station with the 11 same coordinates. The travel times on *starter*, *park*, *drop-off*, *alight* arcs are considered null and, 12 travel times on *pick-up* arcs depend on the behavior of the AMoD but a reference pick-up time  $(\tau_a)$ 

13 is used for travelers itinerary choice.

#### 14 *Travel costs*

15 Itinerary choice is considered deterministic and based on generalized travel costs. Travel cost 16 includes a monetary cost term, a travel time cost term, and a transfer penalty term. The monetary

- cost associated with a car itinerary includes a distance-based cost  $(C_d^c)$ 17 cost associated with a car itinerary includes a distance-based cost  $(C_d^c)$ , carried by the *car* arcs, and
- 18 a parking cost carried by the *park* arcs. On each *park* arcs, a certain parking capacity  $(K^c)$  and a
- 19 price  $(p^c)$  are defined.
- The train pricing scheme comprises a base fare  $(f_0^r)$  $\binom{r}{0}$  and a distance-based fare  $(f_d^r)$ 20 The train pricing scheme comprises a base fare  $(f_0^r)$  and a distance-based fare  $(f_d^r)$ . The urban public transport network pricing scheme only includes a base fare  $(f_0^{s,b})$ 21 urban public transport network pricing scheme only includes a base fare  $(f_0^{s,b})$  that should be paid
- 22 once at the network's entrance. The AMoD pricing scheme has been presented in section [3.1.](#page-7-1)

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 To homogenize time and money, we use several values of time depending on the activity 2 ( $\beta_{ridge}$ ,  $\beta_{drive}$ ,  $\beta_{wait}$ ). In such an intermodal network, the transfer cost cannot be ignored ([24](#page-28-4)). The disutility associated with transferring from one mode to another has largely been studied, particularly in transit systems (*[25,](#page-28-5) [26](#page-28-6)*). This is taken into account through a monetized transfer 5 penalty *γ<sub>trans fer*. An adapted version of the Dijkstra algorithm is used to account for transfer</sub> penalties while exploring the graph.

#### *Relevant willingness to wait*

 Resulting from the itinerary choices, a set of requests is addressed to AMoD where each request 9 has an origin, a destination ( $\in V_w$ ), and timing constraints. A critical parameter is the maximum 10 waiting time  $w_j$ . Here,  $w_j$  is computed following equation [7.](#page-12-0)

<span id="page-12-0"></span>
$$
w_j = \frac{c_j^{\overline{a}} - c_j^a}{\beta_{wait}} \tag{7}
$$

11 where  $c_j^a$  is the travel cost of *j* on the currently followed itinerary starting from the arrival node of

12 *pick-up* edge,  $c_j^{\overline{a}}$  is the travel cost of *j* from the origin node of *pick-up* edge to his destination on the

best itinerary that does not rely on AMoD. This equation corresponds to a simple behavior: when

traveler's waiting cost reaches the extra cost of not using AV during his trip, he cancels his request

and diverts to this alternative itinerary. We assume that after having canceled once, a traveler will

not try to use AV later in his journey: he follows the alternative path until his destination.

#### An unfavorable scenario to challenge heuristics

Finally, as we are looking to evaluate the benefits of introducing a horizon to an assignment-based

heuristic, we stick to an unfavorable scenario for AMoD. By unfavorable, we mean: (i) operating

within a large coverage zone, (ii) with a limited number of vehicles, (iii) during morning commute

in a monocentric city. Due to (i) and (iii), more repositioning efforts are required by vehicles,

especially for distant requests from the center. Due to (ii), the dispatcher has effectively choices to

make between requests. Some will be prioritize to the detriment of others, and we investigate the

impact of dispatching strategy on this prioritization.

 An ideal hybrid transit system, as the one proposed by (*[27,](#page-28-7) [28](#page-28-8)*), is chosen for this mono-centric city. It comprises two train lines, two subway lines, twelve bus lines (figure [4b\)](#page-11-2).

 The morning commute demand pattern is a many-to-one: all travelers' destinations are in the center. Regarding network topology and travel costs, most of them prefer to request AV for a door-to-door ride. A marginal part of them requests AV for the first-mile, to join a transit station.

To analyse AMoD's level of service, the area is split into 25 service zones called from #0-0 to

#4-4 (figure [4a\)](#page-11-0). Figure [6](#page-14-0) shows the demand-supply imbalance. In total, 4000 requests are emitted

- during 3 hours.
- All parameters can be found in appendix [8.](#page-29-0)

## AMoD-travelers interactions

AMoD and travelers interact within this environment following figure [7.](#page-15-0) A traveler *j* who chooses

an itinerary including an AV ride sends his request to AMoD. It contains timing constraints com-

37 puted with  $w_j$ , obtained from other transportation alternatives. The dispatcher receives requests *H* 

in advance and, in real-time, AVs states, positions, and service plannings. It resolves the optimal

assignment problem and decides to match *j* with *i*. For *permanent* strategy, *j* is immediately added

<span id="page-13-0"></span>

FIGURE 5: The multimodal network.

<span id="page-14-0"></span>

FIGURE 6: Number of pick-ups minus number of drop-offs requested per zone. A many-toone demand pattern is considered in the theoretical instance to provoke tough imbalance between supply and demand.

1 to *i*'s service list. For *temporary*, if  $e_i$  belongs to the second part of the horizon, *i* only receives

the order to reposition toward *j*'s pick-up point till the next assignment round. For a permanent

3 match, *j* is picked-up by *i* as soon as *i* arrives, except if  $e_i$  is not yet reached. In this case, *i* waits

for *j*. Then, *i* picks *j* up, carries him to his destination, and drops him off. For non-AV legs of

the itinerary, customer travels node by node on the other graph layers. Travelers who have cho-

sen car mode book parking in advance to prevent overcoming capacity. Behaviors are detailed in

appendices [9](#page-30-0) and [10.](#page-31-0)

### Benefits for AMoD and travelers

Seven specific indicators are tracked.



<span id="page-15-0"></span>

FIGURE 7: AMoD-travelers interactions.

## 1 RESULTS

2 We present the results obtained on the instance described above.

## 3 Sensitivity of *H* on supplier and customer

 To better highlight differences between the proposed strategies, we let more freedom to the dis- patcher by focusing on patient requesters. In this section, all travelers have a large willingness 6 to wait  $W = 20$ min. The maximum planning length *K* is set to 2 by default. As discussed in sec-7 tion [3.4,](#page-9-3) a proper value of  $\alpha$  should be chosen. We preliminary compare  $\alpha H = 5,10$ min,  $H \in [0,60]$ 8 under the four strategies. We choose to set  $\alpha H = 10$ min if  $H \le 10$ , and *H* otherwise because the variation range of simulation time is smaller with a similar number of idle AVs o, average. Our four assignment strategies are compared on their sensitivity to the horizon parameter *H*. Results are represented on figure [8.](#page-16-0) *Introduction of a horizon enables AMoD to serve more requests whatever the strategy*

- 
- 13 ORR starts by increasing with *H*. The fleet is underused at  $H = 0$ , with at least a hundred AVs
- 14 being idle at any time. The lack of anticipation prevents AVs from being dispatched to distant
- 15 requests. The repositioning time required to join pick-up points of corner zones (#0-0, #0-4, #4-0,
- 16 #4-4) from #2-2 (where most AVs finish their service mission), is higher than *W*. Consequently, 17 inequalities in ORR among zones are significant, as shown in figure [9a](#page-17-0) and confirmed by G, above
- 18 0.2 for all strategies.

<span id="page-16-0"></span>

FIGURE 8: Sensitivity of *H* on metrics for the four dispatching strategies.

<span id="page-17-0"></span>

<span id="page-17-2"></span><span id="page-17-1"></span>FIGURE 9: ORR per zone.

- 1 *The productivist strategies are better to maximize ORR, to the detriment of G*
- 2 A horizon of 15-25min maximizes the number of matches. The maximal value of ORR reached
- 3 depends on the strategy. The *productivist* are better than *impatient*. In parallel, G decreases but not
- 4 to the same extent for the four strategies. The *impatient* utilities erase more ORR disparities than
- 5 *productivist* ones. Figures [9b](#page-17-1) and [9c](#page-17-2) represent the ORR per zone at  $H = 20$  for the two extreme
- 6 strategies in terms of G. With the same knowledge of future requests, the improvement of ORR in
- 7 *U*<sup>3</sup> corner zones is better for *permanent-impatient*.

## 8 *Two distinct behaviors emerge from the two utilities*

 For *permanent-impatient*, distant requests for which AVs can arrive just in time are more attrac- tive than closer requests for which AVs would have to wait for the customer. For *temporary- productivist*, distant requests prior notices are not convincing, and AMoD prefers to dispatch AVs to closer customers. Similarly, for *permanent-productivist*, AVs are matched in priority with closer 13 requests even if they must wait for the customer. AV waiting time is included in  $u_{ij}^{productivist}$  defi- nition, contributing to the time required to earn a certain profit. But it does not have as much im- portance as in the *impatient* utility. The average number of waiting AVs for *H* = 20 in *permanent- productivist* has a higher score (50 AVs) than in the other three (around 20 AVs). For *temporary- productivist*, we do not find a consequent number since only AVs waiting for a customer involved in a permanent match are counted.

## 19 *Extending knowledge even more has no effect for temporary strategies*

- 20 The decreasing phase with *H* higher than 15-25min does not affect the four strategies similarly. The
- 21 *temporary* ones are less impacted, with a slight decrease in the number of matches and early stabi-
- 22 lization of metrics. Providing more information to the dispatcher does not change fleet's behavior.
- 23 The dispatcher has exploited the maximum amount of information possible. The additional notices
- 24 included within the horizon are candidates for repositioning in the *temporary* strategies. Yet, only
- 25 idle AVs are likely to be temporarily matched with a candidate for repositioning. Since fleet is not 26 oversized, most AVs are occupied with a service mission at each instant. No supplementary idle
- 27 AV is available for being matched with the additional prior notices. The value of *H* from which
- 28 stabilization happens depends mainly on fleet size (oversupply allows to keep more AVs available

 for repositioning missions) and *K* (longer planning may free some AVs by loading AVs' schedules more).

 For *permanent* strategies, we do not observe the same stabilization when *H* increases over *H* = 25. From *H* = 25, *permanent-impatient* looses 200 matches and *permanent-productivist* looses 280 matches.

## *With permanent-productivist, extending H too much favours AMoD and puts customers at a disad-vantage*

 We find that for *H* ≥ 25, TDTS and TDTE decrease at the same pace. AVs serve fewer travelers, provide less repositioning effort and serve fewer kilometers. In parallel, G reaches its minimum at  $H = 25$  and then increases. AMoD is more focused on central zones and serve shorter and closer requests. Extending the horizon makes more of these requests visible to the AMoD, and AVs that 13 were sent in  $U_3$  are kept in  $U_1$  instead. As more and more AVs desert  $U_3$  to wait for new customers in *U*<sup>1</sup> and *U*2, the number of matches decreases. The average number of waiting AVs doubles 15 between  $H = 20$  and  $H = 60$ . 16 Meanwhile, P increases until  $H = 40$ , when the maximum profit score is reached (11112) EUR). By focusing more on shorter and closer requests while ignoring more distant and longer

requests, AMoD's profit increases until the time lost due to AVs arriving ahead of time becomes

disadvantageous.

#### *Anticipating too much is negative for ORR but dummy for G and P with permanent-impatient*

 Both TDTS and TDTE grow at the same pace. AVs serve fewer customers but make more reposi- tioning efforts for longer rides. The time to achieve a mission, from repositioning to drop-off, is greater, and AVs are monopolized longer for a given mission. Moreover, the number of requests candidates for a match increases with a wider horizon. Match being permanent, more AVs are nec-25 essary to answer a batch of requests. The average number of idle AVs decreases from 40 at  $H = 25$ 26 to 20 at  $H = 60$ . By over-anticipating, *permanent-impatient* AMoD behaves too impatiently to be able to maintain global ORR. It assigns AVs with distant requests that have an earliest pick-up time in far future but requiring enough repositioning effort to allow the just-in-time arrival of AV at the pick-up point. These far requests, being associated with longer rides, are profitable enough to be prioritized by the optimal assignment method. Increasing *K* jointly with *H* may absorb, to some extent, the reduction of global ORR by making AVs currently monopolized by long rides, available to receive more service mission orders. In figures [10,](#page-19-0) we effectively have more matches 33 when  $K = 3$  but, within the same simulated duration, it does not mean higher TDTS. At the end 34 of the studied period, we have more AVs that have not finished their service missions. At  $H = 60$ , TDTS of both scenarios join. This confirms that within a 0-60min range for the horizon, 2 is a proper value for the planning length.

 A wider horizon does not negatively impact P nor G. P variation range is quite narrow (9856-10215 EUR) since TDTS and TDTE increase jointly: repositioning efforts and service dis-39 tance benefits are balanced. G maintains from  $H = 30$ .

<span id="page-19-0"></span>

**FIGURE 10:** Comparison of sensibilities to *H* for  $K = 2$  and  $K = 3$  for the *permanent-impatient* strategy.

- 1 *The temporary-impatient strategy is the worse for profit due to unproductive repositioning*
- 2 For *temporary-impatient*, we notice that TDTE keeps increasing while TDTS follows ORR from
- $3 \text{ } H = 15$ . AVs are making more repositioning efforts without any reward in terms of TDTS. More-
- 4 over, counting the average number of idle AVs shows that more AVs are running (serving or repo-
- 5 sitioning) during the simulation. Consequently, the profit curve is the worst among all strategies,
- 6 reaching a minimum of 8395 EUR. A temporary match with a distant request requires more than
- $7 \alpha H$  (10min) to be achieved. Then, it is put back into play and is exposed to interruption. Around
- 8 70-80% of repositioning missions are interrupted. The repositioning AV is either permanently
- 9 matched with another request than its previous target or assigned another repositioning target. The
- 10 *temporary* strategy is at the origin of unproductive repositioning and, sub-optimality in dispatching
- 11 decisions here. When demand prediction is not exact, it may improve system's flexibility.
- 12 *Repositioning in temporary-productivist is more flourishing and leads to a satisfying and stable P*
- 13 *for supplier but higher spatial disparities in ORR*
- 14 For *temporary-productivist*, TDTS, TDTE and P are similar to *permanent-productivist* until  $H =$
- 15 15. From then, the three indicators and G stabilize, letting this strategy below *permanent-productivist*
- 16 in terms of P. The maximal profit reached is the second best (10622 EUR). We note that using the
- 17 *productivist* utility rather than the *impatient* one, allows to reach a higher ratio of repositioning
- 18 missions achieved (70-80%). Indeed, the repositioning missions are shorter so have less occasions
- 19 to be interrupted.
- 20 *The permanent strategies are better for improving TWT and mean realized pick-up time as H*
- 21 *increases*
- 22 The *permanent* strategies have similar monotonically decreasing TWT profiles. Regarding the total
- 23 realized pick-up time, it reduces as *H* increases. The mean realized pick-up time starts around
- 24 12min for  $H = 0$  and reaches 2min at  $H = 60$  for both strategies. If the number of matches
- 25 decreases, the quality of service for matched travelers largely increases.
- 26 The *temporary* strategies also have similar TWT profiles: it starts by decreasing, then sta-27 bilizes just as other indicators. The mean realized pick-up time variation range is narrow, reaching

1 9min for  $H = 15$  and remaining around 12min for shorter and wider horizon.

 The advantage of *productivist* utility for the supplier is evident here: waiting passively in the center is a more profitable strategy for AMoD compared to actively looking for more distant rides. This behavior is the origin of deeper disparities in terms of ORR (G).

# Extension to traveler-specific willingness to wait

- In this section, we wonder how strategies adapt to more impatient travelers. We also go further in
- <span id="page-20-0"></span>the analysis of travelers-side indicators.



FIGURE 11: Mean maximum waiting time per zone.

 When the maximum waiting time is not constant but based on available alternatives on the transportation network, it becomes traveler-specific. Figure [11](#page-20-0) represents the mean maximum waiting time per zone. It shows smaller values in zones better supplied by transit. Travelers who 12 emit a request from a corner zone of  $U_3$  are more willing to wait (16min) than travelers who emit a request from the central zone (7min). The more traveler is willing to wait, the more he depends on AVs. Here, we investigate the inequalities in terms of quality of service, represented by the realized pick-up time. We wonder if inequalities in quality of service are related to AV-dependency and to what extent. As *permanent* strategies have performed better than *temporary* strategies for both utilities (in terms of P, ORR, TWT), we select them for this analysis.

<span id="page-21-1"></span>

<span id="page-21-0"></span>FIGURE 12: Sensitivity of *H* with constant (*W*) and traveler-specific (*wj*) maximum waiting time.

 Figure [12b](#page-21-0) shows that mean realized pick-up time is always lower in the traveler-specific 2 maximum waiting time scenario  $(w_i)$ . It is consistent since travelers are more demanding  $(w_i \leq W)$ . The difference in mean realized pick-up time with *W* scenario reduces as *H* increases. Figure [12a](#page-21-1) shows that with a short horizon, AMoD does not succeed in answering as many requests as in the *W*. When anticipation is sufficient, all indicators of *w<sup>j</sup>* scenario (except TWT) join those of *W* scenario. It demonstrates the resilience of the *permanent* anticipative heuristics proposed under tougher timing constraints.

8 Figure [13](#page-22-0) presents the concentration curves (*[29](#page-28-9)*) for quality of service. To obtain it, we 9 have defined five classes of AV-dependency from the less dependent to the most:

10 1.  $w_j \le 7$ 

11 2.  $8 \leq w_i \leq 10$ 

12 3.  $11 \leq w_i \leq 13$ 

$$
13 \qquad \qquad 4. \ 14 \le w_j \le 16
$$

14 5.  $17 \leq w_i$ 

 On the x-axis, it plots the cumulative proportion of travelers (all on [13a,](#page-22-1) only those that have been matched on [13b\)](#page-22-2), beginning with the less AV-dependent, ending with the most AV-dependent. On the y-axis of figure [13a](#page-22-1) (resp. figure [13b\)](#page-22-2), it plots the cumulative share of travelers being matched (resp. of total waiting time). If matches (resp. waiting times) are equally undergone across classes, the concentration curve would coincide with the diagonal. The degree of inequality 20 can be represented by the concentration index  $C_i$ , which equals 0 for perfect equality. It is worth -1 (resp. 1) for a perfect inequality to the disadvantage of the less AV-dependent, i.e., all matches are attributed to the less AV-dependent (resp. the most AV-dependent travelers undergo all the waiting 23 time imposed by AMoD). Here, computing  $c_i$  and looking at the curves indicates that:

- 24  $C_i$ (*matches*, *impatient*) =  $-0.015$
- 25  $C_i(waiting times, impairment) = 0.093$
- 26  $C_i$ (*matches*, *productivist*) =  $-0.088$
- 27  $C_i(waiting times, productivist) = 0.180$
- 28 *productivist* favors the less AV-dependent travelers in terms of matches and pick-up time 29 more than *impatient*
- 30 with *productivist*, the 50% of travelers the less AV-dependent attract 60% of the matches

<span id="page-22-1"></span><span id="page-22-0"></span>

<span id="page-22-2"></span>FIGURE 13: Concentration curves for quality of service related to AV-dependency.

1 • with *productivist*, the 20% of matched travelers the most AV-dependent undergo 1/3 of 2 the total realized pick-up time while the 20% of matched travelers the less AV-dependent 3 undergo 12% of this quantity

### 4 Comparison with heuristic

 We have shown that a horizon helps improve AMoD-side and travelers-side benefits. Yet, we have stuck to a slightly undersized fleet. In this section, we compare our anticipative assignment-based approach with the rule-based heuristic presented in section [3.5](#page-9-4) for undersized and oversized fleets. We still focus on *permanent* strategies and, set *H* to 25min (where ORR is maximal for both *permanent* strategies). Figure [14](#page-23-0) present the results for a traveler-specific maximum waiting time. The (rule-based) *heuristic* strategy is less sensible to fleet size than *permanent* in terms of ORR and P. There are more and more idle AVs that do not succeed in answering the remaining unmatched requests. In our unfavorable instance, the lack of anticipation prevents idle vehicles, mostly located in the central zone, from satisfying timing constraints of corner zones' requests. The horizon allows the other two strategies to take advantage of supplementary resources: ORR keeps increasing till the maximum fleet size. On average, for 600 AVs, there are 380 idle AVs with the *heuristic* strategy, 80 with the *permanent-impatient* and, 75 with the *permanent-productivist*. The ORR (resp. P) gap is around 0.45 (resp. 4000 EUR). Note also that as fleet size increases, the gap between *impatient* and *productivist* strategies narrows. Regarding G, *heuristic* remains around 0.2-0.3 which is better than *permanent-productivist* for a fleet containing less than 350 AVs, but worse for greater fleets. The tendency to focus on less distant and shorter rides highlighted for *permanent-productivist* is close to what *heuristic* emulates. For assignment-based heuristics, when more AVs are available, ORRs of corner zones increase, and G is almost null, as shown in figure [15.](#page-23-1)

24 Regarding mean realized pick-up time, *heuristic* and *permanent-impatient* have similar

<span id="page-23-0"></span>

<span id="page-23-1"></span>FIGURE 14: Sensitivity of fleet size for *permanent* and *heuristic* strategies with traveler-specific  $(w_j)$  maximum waiting time ( $H = 25$ )



FIGURE 15: ORR per zone for *permanent-productivist* strategy with 600 AVs (*H*=25).

<span id="page-24-0"></span>

FIGURE 16: PSC1 where prices are positive and subsidies are negative (EUR).

- decreasing profiles. In the case of *heuristic*, it decreases from 8min for 150 AVs to 6min for 600
- Avs, but this is at the price of numerous idle AVs. The *permanent-productivist* strategy is better than the other two, with a variation range extending from 3min to 5.5min.

## Inclusion of a regulator

As highlighted above, *permanent-productivist* is the strategy maximizing the profit of AMoD, but

it does to the detriment of equity among travelers under a limited fleet size. In this section, we add

 a regulator stakeholder in the system. It aims at restoring equity under this strategy by setting up a pick-up pricing scheme. We test two intuitive pricing schemes:

- PSC1: Subsidies are attributed to AVs for picking up a traveler in an AV-dependent zone, and prices are imposed on AVs for picking up a traveler in a non-AV-dependent zone. Subsidies and prices are proportional to: (i) the AV-dependency of the zone, (ii) the distance of this zone from the center of the city. Figure [16](#page-24-0) shows prices per zone.
- PSC2: Subsidies are the same as in PSC1, but there are no prices.

 Figure [17](#page-25-0) shows that compared to no pick-up price scenario, PSC2 reaches equality in terms of matches while PSC1 brings equity. In the latter scenario, the 40% travelers more AV- dependent get 50% of the matches. Due to the definition of schemes providing more subsidies than collecting prices, P increases by 2715 EUR with PSC1 and by 3757 EUR with PSC2. As expected, ORR slightly decreases.

## CONCLUSION

In this study, we have investigated the effect of demand prediction horizon on AMoD-side and

- travelers-side indicators for an assignment-based dispatching heuristic. We have tested four strate-
- gies and analyzed the sensitivity of each to horizon length in an unfavorable theoretical environ-
- ment. Introducing a short/mid-term horizon (5-25min) allows AMoD to increase its profit by
- serving more requests, improves the mean realized pick-up time for travelers and reduces spatial
- disparities in terms of ORR. For *temporary* mode, extending more the horizon does not change

<span id="page-25-0"></span>

(a) Concentration curves for the three pricing schemes



**FIGURE 17:** Effect of pricing schemes for *permanent-productivist* ( $w_j$ ,  $H = 25$ , 350 AVs)

- travelers' outcomes but degrades AMoD's profit in case of unproductive repositioning, as it hap-
- pens in *temporary-impatient* strategy. For *permanent* mode, anticipating long-term requests de-
- grades ORR. Two distinct behaviors emerge depending on the strategy. With an *impatient* utility,
- AVs are rewarded for making more repositioning efforts and maintain a good profit level for higher
- horizon lengths. With a *productivist* utility, AVs succeed in earning more profit by waiting for fu-ture customers in the center. Comparison with a non-anticipative heuristic has shown the ability
- of our dispatching to exploit supplementary vehicles. Even when fleet is oversized, a severe im-
- balance between supply and demand prevents the rule-based heuristic from performing similarly.
- The *permanent-productivist* maximizes AMoD's profit but worsen spatial inequalities. Especially,
- the most AV-dependent travelers are less matched. Those who are matched undergo a bigger part
- of the total waiting time compared to the least AV-dependent. One pricing scheme for restoring
- equality and another for advantaging the most AV-dependent have been proposed.
- Our study can be extended in several ways. A real-world instance would be a more nuanced
- scenario to test our dispatching strategies. Removing the monocentric hypothesis, the many-to-one
- filter would lead to a lighter imbalance between supply and demand and may highlight new behav-
- iors. Considering the entire demand on a real agglomeration would enable us to perform a simu-
- lation time analysis and check the scalability of the approach. Finally, we could go further on de-
- signing pricing schemes that benefit AMoD, travelers and regulator. For example, prices/subsidies
- could be passed on AVs and travelers jointly to change AMoD usage from a majority of door-to-
- door rides to more first/last-mile rides.

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# <span id="page-29-0"></span>1 THEORETICAL INSTANCE PARAMETERS

## 2 Road network



## 3 Transit



## 4 AMoD



# 5 Travelers



## <span id="page-30-0"></span>1 TRAVELER BEHAVIOR





## <span id="page-31-0"></span>1 AV BEHAVIOR



FIGURE 19: Flow chart of AV.