

Securing a High-Level Language Virtual Machine Through its ISA: Pharo as a Case Study

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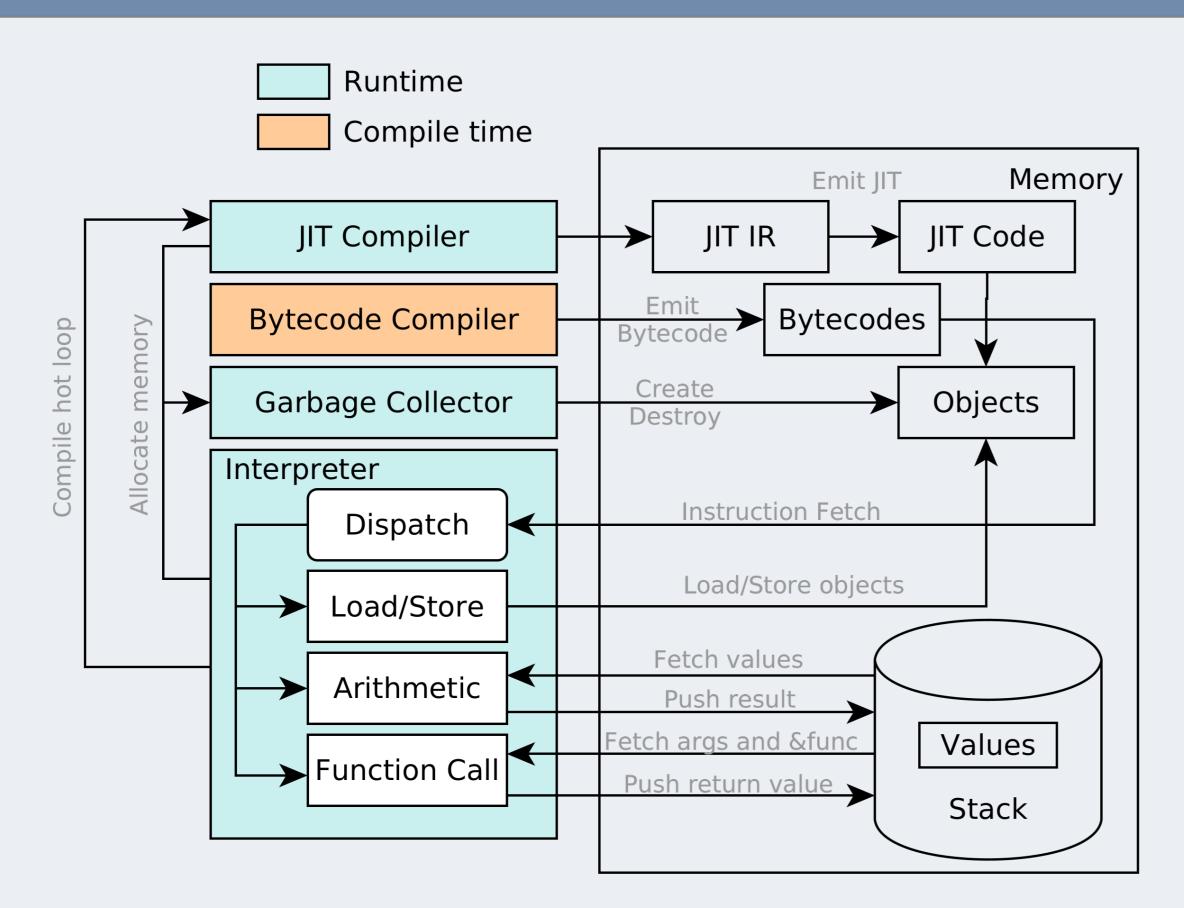
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Language Virtual Machines JIT Compilation



- Bytecode Compiler: transforms source code in a high-level intermediate representation, bytecodes
- Interpreter: process the bytecodes and monitor instruction hotness
- JIT compiler: compiles hot methods/blocks/constants into native code
- Garbage Collector: allocates and deallocates memory

Memory Protection and Process Isolation

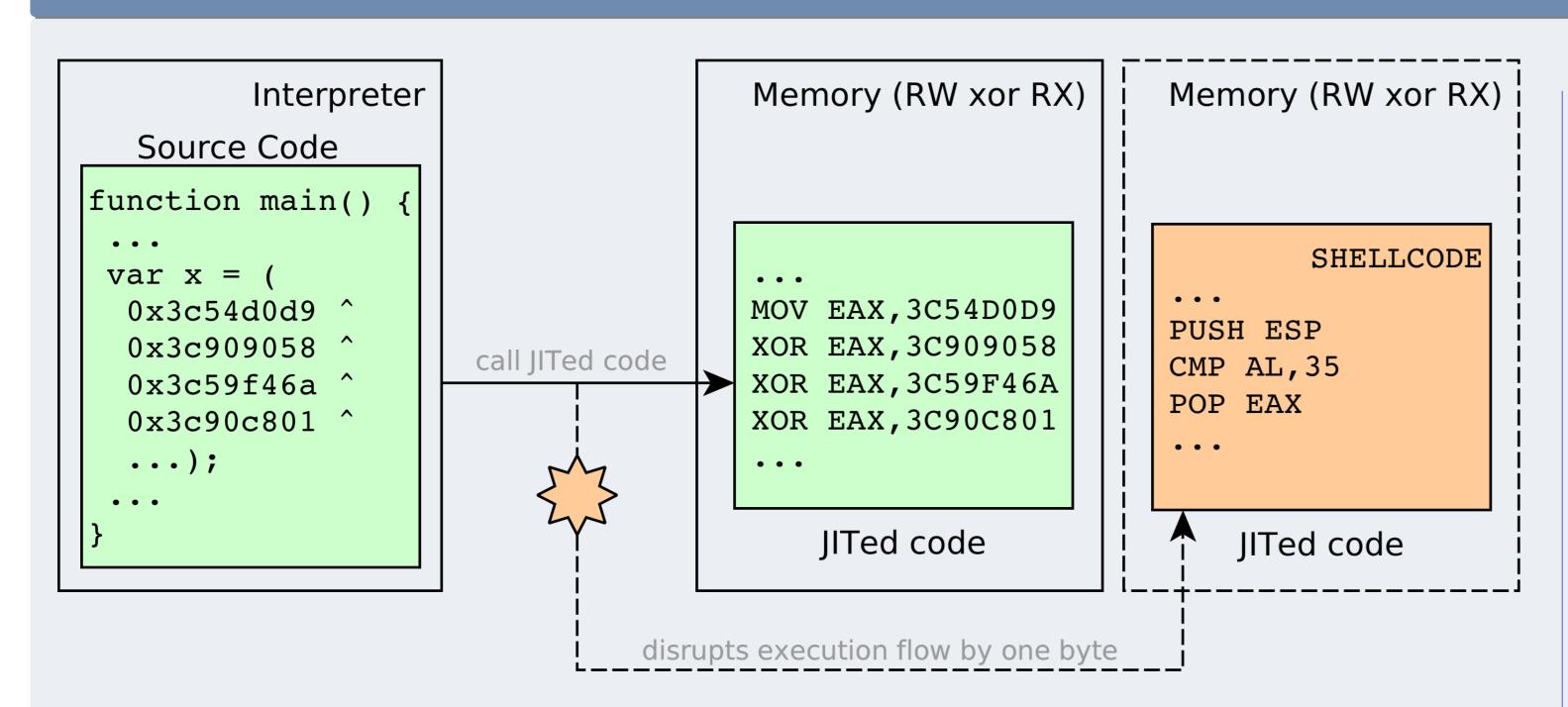
Memory protections are present at different level and granularity:

- Memory Page Permissions: Read (R), Write (W) or Execute (X) granted to a given memory page (4kB)
- Memory Protection Keys: Userspace hardware mechanism to control page table permissions with key tags.
- **Domain Keys:** Per *process isolation* that can be split through several memory pages.
- Trusted Execution Environments: Hardware extension enabling isolated execution environments called *enclaves*.

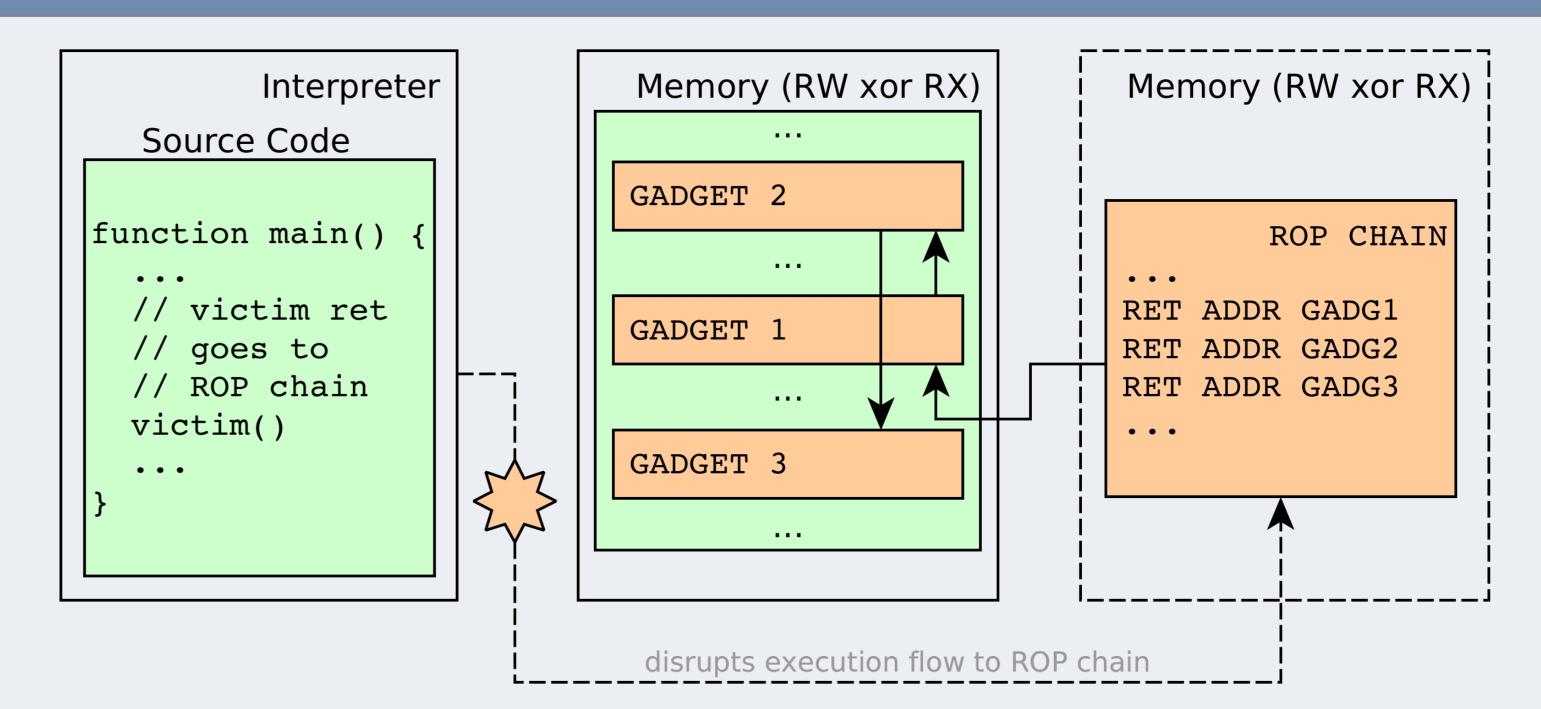
Nanufacturers propose their TEEs such as ARM Trustzone or Intel Software Guard Extension (SGX) as well as RISC-V Keystone [1].

Protections are enforced on JIT engines using $W \oplus X$, Data Execution Prevention (DEP), Address Space Layout Randomization (ASLR) or Control-Flow Integrity (CFI). More specific protections have also been developed by integrating MPK [2] or SGX [3].

Attack on Virtual Machines and their JIT Engines



JIT Spraying [4] forces the JIT compilation of a **XOR chain** to produce a set of constants in JITed memory then disrupts the execution flow by one byte to reveal the hidden shellcode.



JIT-ROP [5] creates the payload from available **gadgets** in memory, either already present or JIT compiled directly by the attacker. The execution flow is disrupted to a return chain to launch the shellcode.

Pharo as a Case Study Virtual Machine for Security

Why choose Pharo?

- Migration in Process or Completed to other ISAs (ARM v8).
- Simpler than heavily engineered Java or JavaScript Virtual Machines.
- **Complex enough** to be in commercial use and require the security aspect.
- High-level Language Test Harness of hundreds of *ISA-agnostic* tests.

A VM implementation for a new ISA runs through three main phases:

- Unit-test a simulation of the VM in the development environment itself.
- **Emulate** the whole system on an architecture emulator (QEMU/Unicorn).
- **Execute** the system on the **real hardware**.

Objectives

- Port the Pharo VM to the RISC-V ISA.
- Replay well-known attacks in a controlled environment.
- Design enclave behaviour on critical components using Keystone.
- Design instructions to handle isolation.
- Test the implementation in hardware (Beagle-V, CVA6).

Open Questions

- Are the attacks replayable on RISC-V?
- What defines a pertinent instruction for JIT engines?

References

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At a glance

- **Problem:** JIT engines are powerful but vulnerable pieces of software.
- **Approach:** involving the ISA in a high-level component provides more security guarantees.
- Implementation: Pharo VM on RISC-V adding instruction extensions and enclave support.

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