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# Period collapse in Ehrhart quasi-polynomials of {1,3}-graphs \*

Cristina G. Fernandes † José C. de Pina † Jorge Luis Ramírez Alfonsín ‡ Sinai Robins †

September 19, 2022

#### Abstract

A graph whose nodes have degree 1 or 3 is called a  $\{1,3\}$ -graph. Liu and Osserman associated a polytope to each  $\{1,3\}$ -graph and studied the Ehrhart quasi-polynomials of these polytopes. They showed that the vertices of these polytopes have coordinates in the set  $\{0,\frac{1}{4},\frac{1}{2},1\}$ , which implies that the period of their Ehrhart quasi-polynomials is either 1, 2, or 4. We show that the period of the Ehrhart quasi-polynomial of these polytopes is 2 if the graph is a tree, the period is at most 2 if the graph is cubic, and the period is 4 otherwise.

In the process of proving this theorem, several interesting combinatorial and geometric properties of these polytopes were uncovered, arising from the structure of their associated graphs. The tools developed here may find other applications in the study of Ehrhart quasi-polynomials and enumeration problems for other polytopes that arise from graphs. Additionally, we have identified some interesting connections with triangulations of 3-manifolds.

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### 1 Introduction

A  $\{1,3\}$ -graph is a graph whose nodes have degree 1 or 3. If all nodes of a graph have degree three then it is *cubic*. We allow all of our graphs to have loops and parallel edges. Liu and Osserman [12] associated a polytope  $\mathcal{P}_G$  to each  $\{1,3\}$ -graph G and studied the Ehrhart quasi-polynomial arising from  $\mathcal{P}_G$ . They were mainly motivated by the relation of these quasi-polynomials to the study of dormant torally indigenous bundles on a general curve, objects arising in algebraic geometry [16]. This connection was further investigated in [20], and more properties of the polytope  $\mathcal{P}_G$  were presented in [6].

Specifically, Liu and Osserman [12, Theorem 3.9] observed that Mochizuki [16] implicitly proved that the value of the Ehrhart quasi-polynomial of  $\mathcal{P}_G$  on odd primes is the number of dormant torally indigenous bundles in a certain class of curves parametrized by the number of nodes and edges of G. Also, they had proved that the coordinates of all vertices of  $\mathcal{P}_G$  are in  $\{0, \frac{1}{4}, \frac{1}{2}, 1\}$ . This implies that the period of the Ehrhart quasi-polynomial of  $\mathcal{P}_G$  is either 1 or 2 or 4 [2, 5]. Using a result of Mochizuki [16], they concluded that the odd constituents of this Ehrhart quasi-polynomial are the same polynomial. Liu and Osserman [12, Question 4.3] then raised questions about the period of the Ehrhart quasi-polynomial of  $\mathcal{P}_G$ . We answer some of these questions in Theorem 4 if G is a tree, and in Theorem 5 if G is a general  $\{1,3\}$ -graph.

The polytope  $\mathcal{P}_G$  has nice geometric and combinatorial properties. For instance, Wakabayashi has proved [20, Proposition 5.3 and Corollary 5.4] that, for cubic graphs  $G_1$  and  $G_2$ , the polytopes  $\mathcal{P}_{G_1}$  and  $\mathcal{P}_{G_2}$  are isomorphic (that is, there is an  $\mathbb{R}$ -linear bijection  $f: \mathbb{R}^d \to \mathbb{R}^d$  such that  $\mathcal{P}_{G_2} = f(\mathcal{P}_{G_1})$ ) if and only if the graphs  $G_1$  and  $G_2$  are isomorphic. Let T be a  $\{1,3\}$ -tree. It turns out that  $\mathcal{P}_T$  has additional interesting geometric properties that can be combinatorially described. The latter might provide useful and attractive insights for appealing questions on  $\theta/1$  polytopes (that is, the convex hull of subsets of  $\{0,1\}^d$ , the vertices of the d-cube). All these will be discussed in Sections 2 and 7.

Our work is also closely connected with some invariants of 3-manifolds, investigated by Maria and Spreer [13]. They associated a linear system of inequalities to a triangulation T of a 3-manifold and studied admissible colourings of the edges of T with the aim to understand better Turaev-Viro type invariants. It turns out that admissible colourings correspond to integer points belonging to a polytope related to one of the two central polytopes investigated in this paper. This relationship as well as an application, with the same topological flavor as in [13], on a problem concerning non-intersecting closed curves in the plane will be explained towards the end of the paper.

## 1.1 Ehrhart quasi-polynomials and period collapse

Let  $\mathcal{L}$  be a sublattice of  $\mathbb{Z}^d$  and fix  $u \in \mathbb{Z}^d$ . We define the *discrete volume* of a polytope  $\mathcal{P}$  in  $\mathbb{R}^d$  with respect to the coset  $\mathcal{L} + u \subseteq \mathbb{Z}^d$  by

$$\operatorname{vol}_{\mathcal{L}+u}(\mathcal{P}) := |\mathcal{P} \cap (\mathcal{L}+u)|.$$

Here we use the standard notation for the translation of any set  $Y \subseteq \mathbb{R}^d$  by a fixed vector  $u \in \mathbb{R}^d$ , namely  $Y + u := \{y + u : y \in Y\} \subseteq \mathbb{R}^d$ . Ehrhart [5] introduced the function

$$L_{\mathcal{P}}(t) := \operatorname{vol}_{\mathbb{Z}^d}(t\mathcal{P}),$$

which is the number of lattice points in the dilated polytope  $t\mathcal{P} := \{tx : x \in \mathcal{P}\}$ , for a nonnegative integer dilation t. Ehrhart showed that if  $\mathcal{P}$  is an integral polytope, then this function is a polynomial in the integer parameter t. A quasi-polynomial f(t) is a function defined by a list  $p_0, p_1, \ldots, p_{\alpha-1}$  of polynomials such that

$$f(t) = \begin{cases} p_0(t) & \text{if } t \equiv 0 \pmod{\alpha}, \\ p_1(t) & \text{if } t \equiv 1 \pmod{\alpha}, \\ \vdots & \vdots \\ p_{\alpha-1}(t) & \text{if } t \equiv \alpha - 1 \pmod{\alpha}. \end{cases}$$

The minimal such  $\alpha$  is the *period* of f and  $p_0, p_1, \ldots, p_{\alpha-1}$  are the *constituents* of f. More generally, Ehrhart also showed that if  $\mathcal{P}$  is a rational polytope, then the function  $L_{\mathcal{P}}(t)$  is a quasi-polynomial, for integer values of t, whose period divides the least common multiple of the denominators in the coordinates of the vertices of  $\mathcal{P}$  [2, 5].

If a rational polytope  $\mathcal{P}$  has a quasi-polynomial  $L_{\mathcal{P}}(t)$  whose period is p, we also call p the period of  $\mathcal{P}$ . The denominator of  $\mathcal{P}$  is the minimum positive integer  $\alpha$  such that the vertices of the dilated polytope  $\alpha \mathcal{P}$  are integral. For a 'generic' rational polytope, one expects its period to be its denominator. Regarding the complexity of computing the periods, Woods [21] has shown that, for fixed dimension d, there is a polynomial-time algorithm which, given a rational polytope  $\mathcal{P} \subset \mathbb{R}^d$  and an integer n > 0, decides whether n is a multiple of the period of the quasi-polynomial  $L_{\mathcal{P}}(t)$ .

When we have a rational polytope whose period is smaller than its denominator, we refer to this situation as period collapse. Usually, researchers [3, 14, 15, 21, 22] prove that the phenomenon of period collapse exists for some particular rational polytope  $\mathcal{P}$  by exhibiting a decomposition of  $\mathcal{P}$  into rational simplices, and then applying a (different) unimodular affine linear transformation to each simplex. If, somewhat magically, the reassembling of all of the images of these simplices form an integral polytope, then the Ehrhart quasi-polynomial of  $\mathcal{P}$  is in fact a polynomial. It is an open problem to characterize those rational polytopes for which it is possible to carry out this reassembling process [10]. Turner and Wu [19] gave examples in  $\mathbb{R}^2$  of rational polygons with the same Ehrhart quasi-polynomial for which such process is not possible.

Here we develop a different technique for proving period collapse, building on an idea of Liu and Osserman. Namely, rather than decomposing the object into smaller polytopes, and using special types of unimodular transformations for each of them, we instead decompose the integer lattice into a certain sublattice  $\mathcal{L}$ , together with all of its cosets in the integer lattice. Then, for each fixed coset of  $\mathcal{L}$ , we count the number of points of the coset that belong to  $\mathcal{P}$ , and show that this number is in bijection with the number of integer points that belong to another, naturally-defined, integral polytope. This novel technique seems interesting in itself and might be useful in a wider context.

In this research, we have made extensive use of both of the following software packages: LattE [1], and polymake [7, 8].

#### 1.2 Liu and Osserman's polytopes

Let G be a  $\{1,3\}$ -graph. We say that a node of degree 1 is a *leaf node* and a node of degree 3 is an *internal node*. We denote by V(G), E(G), and I(G) the set of nodes, the set of edges, and the set of internal nodes of G, respectively. If the graph is clear from the context we write simply V, E, and I. A subgraph H of G is *internally Eulerian* if the degree in H of every internal node of G is equal to zero or two. In particular, the empty subgraph is internally Eulerian.

If X is a finite set, then we denote by  $\mathbb{1}_S: X \to \{0,1\}$  the characteristic vector of a set  $S \subseteq X$ . If  $S = \{e\}$  for some e in X, we write  $\mathbb{1}_e$ . If S = E(H) for some graph H, we write  $\mathbb{1}_H$ .

Liu and Osserman [12] associated to each  $\{1,3\}$ -graph G a polytope  $\mathcal{P}_G$  in  $\mathbb{R}^E$  as follows. For each internal node v of G, let a, b, and c be the three edges incident to v. If v is incident to a loop e, then e appears twice among the three previous edges. Denote by  $S^{\triangle}(v)$  the linear system of triangle inequalities defined on the variables  $w_a$ ,  $w_b$ , and  $w_c$  as follows:

$$w_a \leq w_b + w_c$$

$$w_b \leq w_a + w_c$$

$$w_c \leq w_a + w_b.$$

From  $S^{\triangle}(v)$ , one can derive that  $w_a$ ,  $w_b$ , and  $w_c$  are nonnegative. We denote by  $S_t^{\mathcal{P}}(v)$  the linear system of inequalities, at each internal node v of G, resulting from adding to  $S^{\triangle}(v)$  the perimeter inequality

$$w_a + w_b + w_c \le t. \tag{1}$$

Now, consider the union of all the linear systems  $S_1^{\mathcal{P}}(v)$ , taken over all internal nodes v of G. Add the constraint  $0 \leq w_e \leq 1/2$  for every edge e in G that alone is a component of G. The set  $\mathcal{P}_G$  consists of all real solutions for this linear system [12, Definition 2.3]. The constraints defining  $\mathcal{P}_G$  imply that  $0 \leq w_e \leq 1/2$  for every edge e in G, hence  $\mathcal{P}_G$  is a polytope. We denote the Ehrhart quasi-polynomial of  $\mathcal{P}_G$  by

$$L_G^{\mathcal{P}}(t) := \operatorname{vol}_{\mathbb{Z}^d}(t\mathcal{P}_G),$$
 (2)

where d = |E(G)|.

If G has connected components H and J, then  $\mathcal{P}_G$  would be the Cartesian product of  $\mathcal{P}_H$  and  $\mathcal{P}_J$ , and thus  $L_G^{\mathcal{P}}(t) = L_H^{\mathcal{P}}(t) \cdot L_J^{\mathcal{P}}(t)$  for all t. Therefore, henceforth we assume that G is a connected graph.

In order to derive properties of  $L_G^{\mathcal{P}}$  for a  $\{1,3\}$ -graph G, Liu and Osserman considered a polytope  $\mathcal{Q}_G$  in  $\mathbb{R}^E \times \mathbb{R}^I$  closely related to  $\mathcal{P}_G$ , where I is the set of internal nodes of G. To define  $\mathcal{Q}_G$ , for each internal node v of G, with a, b, and c being the three edges incident

to v, consider an auxiliary variable  $z_v$  and let  $S_t^{\mathcal{Q}}(v)$  be the previous linear system  $S^{\triangle}(v)$ , together with the additional parity constraint:

$$w_a + w_b + w_c = 2 z_v$$

$$z_v \le t.$$

$$(3)$$

We now consider the union of all the linear systems  $S_1^{\mathcal{Q}}(v)$ , taken over all internal nodes v of G. We add the constraint  $0 \leq w_e \leq 1$  for every edge e in G that alone is a component of G. The polytope  $\mathcal{Q}_G$  consists of the solutions for this linear system [12, Definition 3.1]. For each nonnegative integer t, the number of integer points in the polytope  $t\mathcal{Q}_G$  is denoted by  $L_G^{\mathcal{Q}}(t)$ , the Ehrhart quasi-polynomial in t, associated to the polytope  $\mathcal{Q}_G$ . That is,

$$L_G^{\mathcal{Q}}(t) := \operatorname{vol}_{\mathbb{Z}^d}(t\mathcal{Q}_G), \tag{4}$$

where d = |E| + |I|. Figure 1 shows an example.

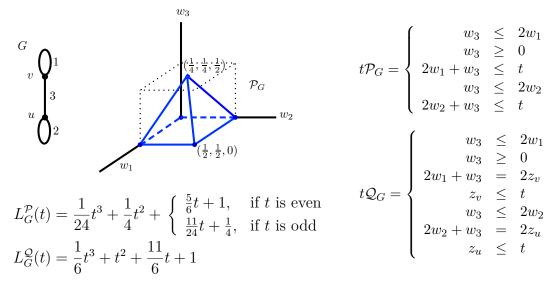


Figure 1: A cubic graph G, its polytope  $\mathcal{P}_G$ , the linear systems of  $t\mathcal{P}_G$  and  $t\mathcal{Q}_G$ , and the Ehrhart quasi-polynomials  $L_G^{\mathcal{P}}(t)$  and  $L_G^{\mathcal{Q}}(t)$  [1]. The polytope  $\mathcal{Q}_G$  lies in  $\mathbb{R}^3 \times \mathbb{R}^2$ . The points in  $t\mathcal{P}_G$  have coordinates  $(w_1, w_2, w_3)$  and the points in  $t\mathcal{Q}_G$  have coordinates  $((w_1, w_2, w_3), (z_v, z_u))$ .

**Example 1.** For the graph G in Figure 1,  $L_G^{\mathcal{P}}(t)$  has period 2 and  $L_G^{\mathcal{Q}}(t)$  has period 1, therefore the number of integer points in the polytope  $\mathcal{P}_G$  is  $L_G^{\mathcal{P}}(1) = \frac{1}{24} + \frac{1}{4} + \frac{11}{24} + \frac{1}{4} = 1$  and the number of integer points in the polytope  $\mathcal{Q}_G$  is  $L_G^{\mathcal{Q}}(1) = \frac{1}{6} + 1 + \frac{11}{6} + 1 = 4$ . Indeed, the unique integer point in  $\mathcal{P}_G$  is (0,0,0) and the integer points in  $\mathcal{Q}_G$  are ((0,0,0),(0,0)),((0,1,0),(0,1)),((1,0,0),(1,0)), and ((1,1,0),(1,1)).

We denote the constituent polynomials of  $L_G^{\mathcal{P}}(t)$  by  $p_0(t), p_1(t), p_2(t)$ , and  $p_3(t)$ , where

$$p_k(t) = L_G^{\mathcal{P}}(t)$$
, for  $t \equiv k \pmod{4}$ .

If  $L_G^{\mathcal{P}}(t)$  has period 1, then  $p_0 = p_1 = p_2 = p_3$ . If  $L_G^{\mathcal{P}}(t)$  has period 2, then  $p_0 = p_2$  and  $p_1 = p_3$ . Similarly, we denote by  $q_0(t)$  and  $q_1(t)$  the constituent polynomials of  $L_G^{\mathcal{Q}}(t)$ , where

$$q_k(t) = L_G^{\mathcal{Q}}(t)$$
, for  $t \equiv k \pmod{2}$ .

If  $L_G^{\mathcal{Q}}(t)$  has period 1, then  $q_0 = q_1$ .

**Example 2.** For the graph G in Figure 1, we have that

$$p_0(t) = p_2(t) = \frac{1}{24}t^3 + \frac{1}{4}t^2 + \frac{5}{6}t + 1$$
 and  $p_1(t) = p_3(t) = \frac{1}{24}t^3 + \frac{1}{4}t^2 + \frac{11}{24}t + \frac{1}{4}$ 

so that here  $L_G^{\mathcal{P}}(t)$  has period 2. The quasi-polynomial  $L_G^{\mathcal{Q}}(t)$  has period 1 because

$$q_0(t) = q_1(t) = \frac{1}{6}t^3 + t^2 + \frac{11}{6}t + 1$$
.

The vertices of  $Q_G$  are its integer points and the point  $((\frac{1}{2}, \frac{1}{2}, 1), (1, 1))$ . The least common multiple of denominators in the coordinates of the vertices of  $\mathcal{P}_G$  is 4 and of  $Q_G$  is 2, giving us examples of period collapse.

## 1.3 Hightlighting context of our main results

Liu and Osserman [12, Proposition 3.5] proved that, in general, the coordinates of all vertices of  $\mathcal{P}_G$  are in  $\{0, \frac{1}{4}, \frac{1}{2}, 1\}$ , while the coordinates of all vertices of  $\mathcal{Q}_G$  are in  $\{0, \frac{1}{2}, 1\}$ . Therefore, the period of  $L_G^{\mathcal{P}}(t)$  is either 1 or 2 or 4 and the period of  $L_G^{\mathcal{Q}}(t)$  is either 1 or 2. In particular, Liu and Osserman [12, Question 4.3] posed some questions related to the period of the Ehrhart quasi-polynomials  $L_G^{\mathcal{P}}(t)$  and  $L_G^{\mathcal{Q}}(t)$ , restated as follows.

Question 3 ([12, Question 4.3]).

- (a) Is it true that if G is cubic then the period of the Ehrhart quasi-polynomial  $L_G^{\mathcal{P}}(t)$  is 2?
- (b) For which  $\{1,3\}$ -graphs G is the period of  $L_G^{\mathcal{P}}(t)$  smaller than the least common multiple of the denominators of the vertices of  $\mathcal{P}_G$ ?
- (c) Is the period of the Ehrhart quasi-polynomial  $L_G^{\mathcal{Q}}(t)$  always half the period of  $L_G^{\mathcal{P}}(t)$  for every  $\{1,3\}$ -graph G?

This paper gives a partial answer to some of these questions. We say that T is trivial if it has no internal node. We prove the following.

**Theorem 4** (the period for  $\{1,3\}$ -trees). If T is a  $\{1,3\}$ -tree, then

- (a)  $Q_T$  is integral and hence  $L_T^{\mathcal{Q}}(t)$  has period 1.
- (b)  $\mathcal{P}_T$  is half-integral and  $L_T^{\mathcal{P}}(t)$  has period equal to 2.

Theorem 4 (a) implies that  $\{1,3\}$ -trees do not have the property addressed by Question 3(b), and that Theorem 4 (b) answers positively Question 3(c) for  $\{1,3\}$ -trees.

**Theorem 5** (the period for  $\{1,3\}$ -graphs). If G is a connected  $\{1,3\}$ -graph then the period of the Ehrhart quasi-polynomial  $L_G^{\mathcal{P}}(t)$  associated to the polytope  $\mathcal{P}_G$  is at most 2 if G is a tree or a cubic graph, and it is equal to 4 otherwise.

Theorem 5 partially answers Question 3(a) and is one of the main contributions of this paper. Liu and Osserman [12, Lemma 3.3] proved that the constituent polynomials  $p_1 = p_3$  by showing that, for every  $\{1,3\}$ -graph G,  $q_1(t) = N_G p_1(t) = N_G p_3(t)$ , where  $N_G$  is the number of internally Eulerian subgraphs of G.

**Example 6.** For the graph G in Figure 1, the internally Eulerian subgraphs are induced by the edge sets  $\emptyset$ ,  $\{1\}$ ,  $\{2\}$ , and  $\{1,2\}$ , so  $N_G = 4$ , and we have

$$q_1(t) = \frac{1}{6}t^3 + t^2 + \frac{11}{6}t + 1 = N_G p_1(t) = 4\left(\frac{1}{24}t^3 + \frac{1}{4}t^2 + \frac{11}{24}t + \frac{1}{4}\right).$$

Their proof [12, Lemma 3.3] is based on a partition, for every nonnegative odd integer t, of the  $L_G^{\mathcal{Q}}(t) = q_1(t)$  integer points of the polytope  $t\mathcal{Q}_G$  into  $N_G$  parts of size  $L_G^{\mathcal{P}}(t)$ . So, answering Question 3(a) boils down to deciding whether we have equality between the two constituent polynomials  $p_0$  and  $p_2$  for cubic graphs.

We answer Question 3(a) positively by also presenting, for every cubic graph G and nonnegative even integer t, a partition of the  $L_G^{\mathcal{Q}}(t) = q_0(t)$  integer points of the polytope  $t\mathcal{Q}_G$ . This partition has  $N_G$  parts, one part of size  $L_G^{\mathcal{P}}(t)$  and  $N_G-1$  parts of size  $L_G^{\mathcal{P}}(t)-(\frac{t}{2}+1)^{k-1}$ , where m and n are the number of edges and nodes of G, respectively, and k=m-n+1 is the cyclomatic number of G. From this it follows that

$$q_0(t) \ = \ N_G \, p_0(t) - (N_G - 1) \left(\frac{t}{2} + 1\right)^{k-1} = \ N_G \, p_2(t) - (N_G - 1) \left(\frac{t}{2} + 1\right)^{k-1},$$

implying that  $p_0 = p_2$  and that the period of  $L_G^{\mathcal{P}}(t)$  is at most 2. This also shows that the class of cubic graphs is a class of  $\{1,3\}$ -graphs as sought-after in Question 3(b).

**Example 7.** Inspecting the example in Figure 1 a bit further, we see that k = 3-2+1 = 2 and, because  $N_G = 4$ , we have that

$$q_0(t) = \frac{1}{6}t^3 + t^2 + \frac{11}{6}t + 1$$

$$= N_G p_0(t) - (N_G - 1) \left(\frac{t}{2} + 1\right)^{k-1}$$

$$= 4\left(\frac{1}{24}t^3 + \frac{1}{4}t^2 + \frac{5}{6}t + 1\right) - 3\left(\frac{t}{2} + 1\right).$$

We were not able to show that, for a cubic graph G, the period of  $L_G^{\mathcal{P}}(t)$  is exactly 2, that is, that  $p_0 \neq p_1$ . However, we point out that the periods of  $L_G^{\mathcal{P}}(t)$  and  $L_G^{\mathcal{Q}}(t)$  are different. Indeed, from the above,

$$q_0(1) = N_G p_0(1) - (N_G - 1) \left(\frac{3}{2}\right)^{k-1} = 2^k p_0(1) - (2^k - 1) \left(\frac{3}{2}\right)^{k-1}$$
 and  $q_1(1) = N_G p_1(1) = N_G = 2^k$ .

Therefore, if  $L_G^{\mathcal{Q}}(t)$  has period 1, then  $q_0 = q_1$ , hence  $q_0(1) = q_1(1)$ , which implies that  $p_0(1) = 1 + \frac{2^k - 1}{2^k} (\frac{3}{2})^{k-1} \neq 1 = p_1(1)$ , because  $k \geq 2$ , thus  $L_G^{\mathcal{P}}(t)$  has period 2. Similarly, if  $L_G^{\mathcal{P}}(t)$  has period 1, then  $L_G^{\mathcal{Q}}(t)$  has period 2. This tackles Question 3(c).

Finally, we derive that the period of  $L_G^{\mathcal{P}}(t)$  is 4 for any  $\{1,3\}$ -graph G that is not a tree or a cubic graph, applying the same strategy used for cubic graphs. This shows that cubic graphs are the only ones that satisfy Question 3(b). Together with Theorem 4, this result leads to Theorem 5.

Although our approach allows us to control the behavior for the difference of polynomials, we were not able, despite much effort, to find a method to compute the desired polynomials explicitly. This seems a challenging task even for  $\{1,3\}$ -trees as stated in [6, Problem 6.5].

#### 1.4 Paper organization

The paper is organized as follows. In Section 2, we show that, for every  $\{1,3\}$ -tree T, the coordinates of all vertices of  $\mathcal{P}_T$  are in  $\{0,\frac{1}{2}\}$  (Theorem 11). We also derive that  $\mathcal{Q}_T$  is a 0/1 polytope (Theorem 12) yielding Theorem 4. In Section 3 we carry out the decomposition of the integer points of the polytope  $\mathcal{Q}_G$  that will lead to Theorem 5. We do so by studying  $L_G^{\mathcal{Q}}$  via cosets of certain lattices and based on internally Eulerian subgraphs of G. The latter allows us to establish a connection with  $L_G^{\mathcal{P}}$  through a simple class of representative trees called caterpillars. To this end, we use the NNI machinery (an easy local move performed on the edges of G). Section 4 presents the proof of Theorem 5 using two main technical lemmas that are proved in Sections 5 and 6. Finally, in Section 7, we discuss further geometric properties of the polytopes  $\mathcal{P}_T$  and a related topological connection.

# 2 The structure and period for $\{1,3\}$ -trees

Let T be a  $\{1,3\}$ -tree. We say that T is trivial if it has no internal node. Hence, a trivial  $\{1,3\}$ -tree is formed by an edge connecting two distinct nodes. A leaf-edge is an edge that is incident to a leaf. A path in T is a leaf-path if its extreme edges are leaf-edges. An internally Eulerian subgraph in T is simply a (possibly empty) collection of disjoint leaf-paths.

If x is a point in  $\mathcal{P}_T$ , then  $0 \le x_e \le \frac{1}{2}$  for each edge e. Suppose that all coordinates of x are in  $\{0, \frac{1}{2}\}$ . If a, b, and c are edges incident to a node, then either zero or two of  $x_a$ ,  $x_b$ , and  $x_c$  are equal to  $\frac{1}{2}$ . This implies that the support

of x is the set of edges of a collection of disjoint leaf-paths. We shall prove that  $V_{\mathcal{P}_T} := \{\frac{1}{2}\mathbb{1}_H : H \text{ is a collection of disjoint leaf-paths in } T\}$  is the set of vertices of  $\mathcal{P}_T$ .

**Example 8.** For the  $\{1,3\}$ -tree T in Figure 2 there are four collections of disjoint leaf-paths: the empty collection, the collection with the leaf-path whose edge set is  $\{1,2\}$ , the collection with the leaf-path whose edge set is  $\{2,3\}$ , and the collection with the leaf-path whose edge set is  $\{1,3\}$ . Therefore  $V_{\mathcal{P}_T} = \{(0,0,0), (\frac{1}{2},\frac{1}{2},0), (0,\frac{1}{2},\frac{1}{2}), (\frac{1}{2},0,\frac{1}{2})\}$  is the vertex set of  $\mathcal{P}_T$ .

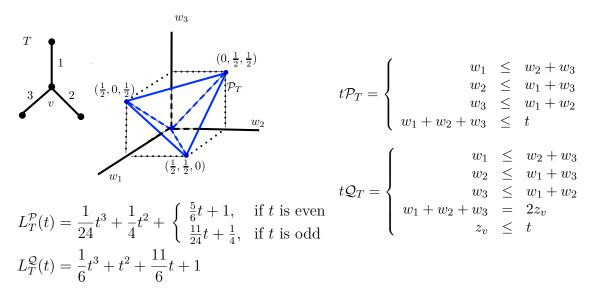


Figure 2: A  $\{1,3\}$ -tree T, its polytope  $\mathcal{P}_T$ , the linear systems of  $t\mathcal{P}_T$  and  $t\mathcal{Q}_T$ , and the Ehrhart quasi-polynomials  $L_T^{\mathcal{P}}(t)$  and  $L_T^{\mathcal{Q}}(t)$  [1]. The Ehrhart quasi-polynomials associated to T are equal to the Ehrhart quasi-polynomials associated to the cubic graph G of Figure 1. The points in  $t\mathcal{P}_G$  have coordinates  $(w_1, w_2, w_3)$  and the points in  $t\mathcal{Q}_G$  have coordinates  $((w_1, w_2, w_3), (z_v))$ .

In Lemma 9 we show that any point in  $\mathcal{P}_T$  with all coordinates in  $\{0, \frac{1}{2}\}$  is uniquely determined by the coordinates associated to leaf-edges. In the statement of the lemma we take into account the fact that any collection of disjoint leaf-paths of a  $\{1, 3\}$ -tree has an even number of leaves.

**Lemma 9.** If T is a  $\{1,3\}$ -tree,  $X_T$  is the set of leaf-edges of T, and X is a subset of  $X_T$  incident to an even number of leaves, then there exists a unique point x in  $\mathcal{P}_T$  such that

$$x_e = \begin{cases} \frac{1}{2} & \text{if } e \in X\\ 0 & \text{if } e \in X_T \setminus X. \end{cases}$$
 (5)

Moreover, all coordinates of x are in  $\{0, \frac{1}{2}\}$  and its support is the set of edges of a collection of disjoint leaf-paths.

*Proof.* The proof is by induction on the number of nodes of T. If T is trivial or has only one internal node, then all edges of T are leaf-edges and x is completely defined by Equation (5) and satisfies all conditions of the lemma. Thus, we may assume that T has at least two internal nodes.

Let r be an internal node of T adjacent to two leaf-edges, and let a, b, and c be the three edges incident to r. We may assume that b is incident to a leaf u and c is incident to a leaf v. Take  $T' := T - \{u, v\}$  and  $X' := X \setminus \{b, c\} \cup \{a\}$  if  $|X \cap \{b, c\}| = 1$  and  $X' := X \setminus \{b, c\}$  otherwise. By induction, there exists a unique point x' in  $\mathcal{P}_{T'}$  with coordinates in  $\{0, \frac{1}{2}\}$  such that for each leaf-edge e of T' we have that  $x'_e = \frac{1}{2}$  if and only if e belongs to X' and such that the support of x' is the set of edges of a collection of disjoint leaf-paths of T'. One can verify that x' can be uniquely extended to a point x in  $\mathcal{P}_T$  with coordinates in  $\{0, \frac{1}{2}\}$  satisfying the conditions of the lemma.

Next, we show that any point in  $\mathcal{P}_T$  with coordinates in  $\{0, \frac{1}{2}\}$  whose support is a collection of disjoint leaf-paths is a vertex of  $\mathcal{P}_T$ .

**Lemma 10.** If T is a  $\{1,3\}$ -tree and H is a collection of disjoint leaf-paths in T, then  $\frac{1}{2}\mathbb{1}_H$  is a vertex of  $\mathcal{P}_T$ .

Proof. It is clear that  $\frac{1}{2}\mathbb{1}_H$  is in  $\mathcal{P}_T$ . Let  $X_T$  be the set of leaf-edges of T and X be the edges of H that are leaf-edges of T. The hyperplane  $2\sum_{e\in X}w_e - 2\sum_{e\in X_T\setminus X}w_e = |X|$  is a supporting hyperplane of  $\mathcal{P}_T$ . Indeed, for each point x in  $\mathcal{P}_T$  and each edge e, we have that  $0 \le x_e \le \frac{1}{2}$ . Thus,  $2\sum_{e\in X}x_e - 2\sum_{e\in X_T\setminus X}x_e \le 2\sum_{e\in X}x_e \le |X|$  and equality holds if and only if  $x_e = 1/2$  for each e in X and  $x_e = 0$  for each e in  $X_T\setminus X$ . By Lemma 9,  $\frac{1}{2}\mathbb{1}_H$  is the unique point in  $\mathcal{P}_T$  for which equality holds and therefore it is a vertex of  $\mathcal{P}_T$ .  $\square$ 

**Theorem 11.** If T is a  $\{1,3\}$ -tree, then  $V_{\mathcal{P}_T}$  is the set of vertices of the polytope  $\mathcal{P}_T$ .

*Proof.* Let  $\mathcal{R}_T$  be the set of points that are a convex combination of points in  $V_{\mathcal{P}_T}$ . By Lemma 10 we have that  $\mathcal{R}_T \subseteq \mathcal{P}_T$ . In order to prove the lemma it suffices to show the converse inclusion.

Suppose that  $\mathcal{P}_T \not\subseteq \mathcal{R}_T$  and let T be the smallest  $\{1,3\}$ -tree with  $\mathcal{P}_T \not\subseteq \mathcal{R}_T$ ; that is, with the number of nodes as small as possible. It is clear that T is nontrivial. Let x be a vertex of  $\mathcal{P}_T$  not contained in  $\mathcal{R}_T$ .

Since x is a vertex of  $\mathcal{P}_T$ , there are |E| linearly independent inequalities among the ones that determine  $\mathcal{P}_T$  satisfied by x with equality. As |E| = 2|I| + 1 there must exist an internal node v such that three inequalities in  $S_1^{\mathcal{P}}(v)$  are satisfied by x with equality. Let a, b, and c be the three edges incident to v. We have two possibilities:

- either the three inequalities in  $S^{\triangle}(v)$  are satisfied by x with equality and therefore  $x_a = x_b = x_c = 0$ ;
- or the perimeter inequality and two inequalities in  $S^{\triangle}(v)$  are satisfied by x with equality and therefore two of  $x_a$ ,  $x_b$ , and  $x_c$  are  $\frac{1}{2}$  and the other is 0.

The tree T can be partitioned into  $\{1,3\}$ -trees  $T_a$ ,  $T_b$ , and  $T_c$  that share the node v and have a, b, and c as leaf-edges, respectively. Let  $x_{T_a}$ ,  $x_{T_b}$ , and  $x_{T_c}$  be the corresponding 'projections' of x onto the edges of  $T_a$ ,  $T_b$ , and  $T_c$ , respectively. One can verify that in both possibilities ( $x_a = 0$  or  $x_a = \frac{1}{2}$ )  $x_{T_a}$  is in  $\mathcal{P}_{T_a}$ . Since  $T_a$  is smaller than T, it follows that  $x_{T_a}$  is in  $\mathcal{R}_{T_a}$  and therefore  $x_{T_a}$  can be decomposed as a convex combination of points in  $V_{\mathcal{P}_{T_a}}$ . Similarly, we can decompose  $x_{T_b}$  and  $x_{T_c}$  as a convex combination of points in  $V_{\mathcal{P}_{T_c}}$ , respectively. These decompositions can be easily glued together to form a decomposition of x as a convex combination of points in  $V_{\mathcal{P}_T}$ , contradicting our assumption.

Using the same strategy, one can derive the following about  $\mathcal{Q}_T$ . Recalling that I is the set of internal nodes of G, let  $\mathbb{1}'_H$  denote the pair  $(\mathbb{1}_H, z_H)$  in  $\mathbb{Z}^E \times \mathbb{Z}^I$ , where  $(z_H)_v = 1$  if v is an internal node of a leaf-path in H, and  $(z_H)_v = 0$  otherwise. Let  $V_{\mathcal{Q}_T} := \{\mathbb{1}'_H : H \text{ is a collection of disjoint leaf-paths in } T\}$ . This discussion is summarized in the following theorem.

**Theorem 12.** If T is a  $\{1,3\}$ -tree, then  $V_{\mathcal{Q}_T}$  is the set of vertices of the polytope  $\mathcal{Q}_T$ .

**Example 13.** For the 
$$\{1,3\}$$
-tree  $T$  in Figure 2 we have that the vertex set of  $\mathcal{Q}_T$  is  $V_{\mathcal{Q}_T} = \{((0,0,0),(0)),((1,1,0),(1)),((0,1,1),(1)),((1,0,1),(1))\}.$ 

Now we can give a proof for Theorem 4, which characterizes the periods of the Ehrhart quasi-polynomials for  $\{1,3\}$ -trees.

Proof of Theorem 4 (the period for  $\{1,3\}$ -trees). By Theorem 12, we know that  $\mathcal{Q}_T$  is integral. It follows from Ehrhart's theorem [2, Theorem 3.8] that  $L_T^{\mathcal{Q}}(t)$  has period 1, which proves part (a) of Theorem 4. To prove part (b), first note that  $\mathcal{P}_T$  is half-integral by Theorem 11, so  $L_T^{\mathcal{P}}(t)$  has period at most 2. Let  $p_0$  and  $p_1$  be the constituents of  $L_T^{\mathcal{P}}(t)$ . It is clear that  $p_0(0) = 1$  and the constant coefficient of  $p_0$  is 1. Liu and Osserman denote by  $q_0$  and  $q_1$  the constituents of  $L_T^{\mathcal{Q}}(t)$ , as they address the more general case in which  $L_T^{\mathcal{Q}}(t)$  might have period 2. Here, as  $L_T^{\mathcal{Q}}(t)$  has period 1, we have that  $q_0 = q_1 = L_T^{\mathcal{Q}}(t)$ . Liu and Osserman [12, Proposition 3.5] proved that  $q_1 = N_T p_1$ , where  $N_T$  is the number of internally Eulerian subgraphs of T. Therefore, the constant coefficient of  $p_1$  is  $1/N_T$  and, because  $N_T \geq 2$  for each  $\{1,3\}$ -tree, it follows that  $p_0 \neq p_1$ . Hence, the period of  $L_T^{\mathcal{P}}(t)$  is exactly 2.

## 3 Cosets and combinatorial tools

The intent of this section is to derive a relation between  $L_G^{\mathcal{Q}}(t)$  and  $L_G^{\mathcal{P}}(t)$  for a nonnegative integer t. In the next section we shall use this relation to derive the period of  $L_G^{\mathcal{P}}$  from the known facts concerning the period of  $L_G^{\mathcal{Q}}$ . For this task, internally Eulerian subgraphs, a class of graphs that resemble caterpillars, and a local move performed on graphs called nearest neighbor interchange (NNI) play central roles. The general idea is partitioning the integer points of the polytope  $t\mathcal{Q}$  into parts of size having  $L_G^{\mathcal{P}}(t)$  as our unit of standard measurement.

#### 3.1 Eulerianicity

Let G be a  $\{1,3\}$ -graph and  $\mathcal{L}_G := (2\mathbb{Z}^E) \times \mathbb{Z}^I$ . For each integer  $t \geq 0$  and each internally Eulerian subgraph H of G, we define

$$\operatorname{part}_{t}(G, H) := t \mathcal{Q}_{G} \cap (\mathcal{L}_{G} + (\mathbb{1}_{H}, 0)) \text{ and } \operatorname{vol}_{t}(G, H) := |\operatorname{part}_{t}(G, H)|,$$
 (6)

where  $\mathcal{L}_G + (\mathbb{1}_H, 0)$  is a coset of the lattice  $\mathcal{L}_G$ , in the lattice  $\mathbb{Z}^E \times \mathbb{Z}^I$ .

In words, we are now counting integer points in cosets of a lattice, not just in a lattice. The term "part<sub>t</sub>" originates from the fact that these sets correspond to parts of a partition of the set of integer points of the polytope  $tQ_G$ . We have that

$$tQ_G \cap (\mathbb{Z}^E \times \mathbb{Z}^I) = \bigcup_H \operatorname{part}_t(G, H) \text{ and } L_G^{\mathcal{Q}}(t) = \sum_H \operatorname{vol}_t(G, H),$$
 (7)

where the union and the summation are over all internally Eulerian subgraphs H of G. Indeed, because of the parity constraint (3), for every integer point (w, z) in  $tQ_G$ , the set of edges e for which  $w_e$  is odd induces an internally Eulerian subgraph of G.

**Example 14.** Consider the polytope  $tQ_G$  in Figure 1. The graph G has four internally Eulerian subgraphs, namely, the subgraphs  $\emptyset$ ,  $H_{\{1\}}$ ,  $H_{\{2\}}$ , and  $H_{\{1,2\}}$  induced by the edge sets  $\emptyset$ ,  $\{1\}$ ,  $\{2\}$ , and  $\{1,2\}$ , respectively. One can verify that

$$\begin{array}{ll} \operatorname{part}_2(G,\emptyset) &= \{((0,0,0),(0,0)),((2,0,0),(2,0)),((0,2,0),(0,2)),((2,2,0),(2,2))\},\\ \operatorname{part}_2(G,H_{\{1\}}) &= \{((1,0,0),(1,0)),((1,2,0),(1,2))\},\\ \operatorname{part}_2(G,H_{\{2\}}) &= \{((0,1,0),(0,1)),((2,1,0),(2,1))\}, \ and\\ \operatorname{part}_2(G,H_{\{1,2\}}) &= \{((1,1,0),(1,1)),((1,1,2),(2,2))\}. \end{array}$$

Therefore, 
$$L_G^{\mathcal{Q}}(2) = 10 = \operatorname{vol}_2(G, \emptyset) + \operatorname{vol}_2(G, H_{\{1\}}) + \operatorname{vol}_2(G, H_{\{2\}}) + \operatorname{vol}_2(G, H_{\{1,2\}}) = 4 + 2 + 2 + 2.$$

The number of internally Eulerian subgraphs of G is denoted by  $N_G$ . If G has n nodes and m edges, then its cyclomatic number is m-n+1. Liu and Osserman [12, Remark 3.11] observed that if G has h leaves and cyclomatic number k, then  $N_G = 2^k$  if h = 0 and  $N_G = 2^{k+h-1}$  if h > 0.

The plan for this section is to measure the terms in the summation in (7) having  $L_G^{\mathcal{P}}(t)$  as standard. An important special term is  $\operatorname{vol}_t(G,\emptyset)$ . This term is associated to the set  $\operatorname{part}_t(G,\emptyset)$  formed by all integer points (w,z) in  $t\mathcal{Q}_G$  for which all coordinates of w are even. Equivalently,

$$\operatorname{part}_t(G,\emptyset) = \{(w,z) \in t\mathcal{Q}_G : w = 2x \text{ for an integer point } x \text{ in } t\mathcal{P}_G\}$$

since z is uniquely determined by w in order for (w,z) to be in  $t\mathcal{Q}_G$ . Therefore

$$\operatorname{vol}_{t}(G,\emptyset) = |\operatorname{part}_{t}(G,\emptyset)| = L_{G}^{\mathcal{P}}(t). \tag{8}$$

For instance, for the graph G appearing in Figure 1,  $\operatorname{vol}_2(G,\emptyset) = 4 = L_G^{\mathcal{P}}(2)$ .

#### 3.2 Caterpillars

Wakabayashi [20, Theorem A(ii)] proved that polytopes  $\mathcal{P}_G$  associated to any connected  $\{1,3\}$ -graph G with a given number n of nodes and a given number m of edges have the same Ehrhart quasi-polynomial  $L_G^{\mathcal{P}}(t)$ . This gives us the freedom to elect a convenient representative for every equivalence class of connected graphs with n nodes and m edges. It turns out that, in order to derive the period of  $L_G^{\mathcal{P}}(t)$ , a particularly convenient choice of a representative graph for G resembles a caterpillar. We define this convenient class of representative graphs in the sequel.

A caterpillar is a tree for which the removal of all leaves results in a path, called its central path, or in the empty graph (Figure 3(a)). When the result is the empty graph, we set the central path to be the empty path. The edges that are not in the central path of the caterpillar are its legs. The nodes that are not in the central path and are adjacent to an end node of the central path are its roots. When the central path is empty, both nodes are roots. The legs adjacent to the roots are the stems of the caterpillar.

For nonnegative integers h and k such that  $h+k \geq 2$ , let T be a  $\{1,3\}$ -tree which is a caterpillar with h+k leaves. A graph obtained from T by adding a loop to k of its leaves is called an (h,k)-caterpillar (Figure 3(b)). The number of nodes in T and in an (h,k)-caterpillar obtained from T is the same, that is, 2(h+k-1). The number of edges in an (h,k)-caterpillar is 2(h+k-1)-1+k=2h+3k-3.

We say that the leaf-edges of an (h, k)-caterpillar are consecutive if their nodes in the central path are consecutive. We say that the loops of an (h, k)-caterpillar are consecutive if their neighbors in the central path are consecutive. We denoted by  $G_{h,k}$  the (h, k)-caterpillar whose loops are all consecutive and leaf-edges are also all consecutive (Figure 3(c)). In particular, neither  $G_{1,0}$  nor  $G_{0,1}$  exist, and the smallest (h, k)-caterpillar is  $G_{2,0}$ , which corresponds to the trivial  $\{1,3\}$ -tree.

Due to the discussion above, henceforth all of our ambient graphs will be (h, k)caterpillars  $G_{h,k}$ . To make the proofs that follow easier to read, we use the notation

$$L_{h,k}^{\mathcal{P}}(t) := L_{G_{h,k}}^{\mathcal{P}}(t) \quad \text{and} \quad L_{h,k}^{\mathcal{Q}}(t) := L_{G_{h,k}}^{\mathcal{Q}}(t).$$

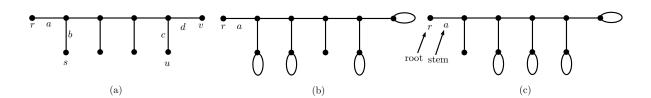


Figure 3: (a) Caterpillar with 6 legs, with roots r, s, u, v and corresponding stems a, b, c, d. (b) (2, 4)-caterpillar with root r and corresponding stem a. (c) (2, 4)-caterpillar  $G_{2,4}$  with root r and corresponding stem a.

#### 3.3 Nearest neighbor interchanges

This section is devoted to scrutinizing the terms in the summation in (7). In the following two paragraphs we outline how we deal with the left and right side of the summation in (7), respectively.

If G is a connected graph with n nodes and m edges then its cyclomatic number is k = m - n + 1. In addition, if G is a connected  $\{1,3\}$ -graph with m edges and h leaves then 2m = 3n - 2h. Due to these two linearly independent equalities we conclude that any connected  $\{1,3\}$ -graph with h leaves and cyclomatic number k has the same number n of nodes and the same number m of edges as the (h,k)-caterpillar  $G_{h,k}$ . Therefore, from Wakabayashi [20, Theorem A(ii)] it follows that if G is a connected  $\{1,3\}$ -graph with h leaves and cyclomatic number k then  $L_G^p = L_{h,k}^p$ . The first step to turn the summation in (7) more amenable to our purposes is to prove a similar identity to polytope  $\mathcal{Q}$ , namely  $L_G^{\mathcal{Q}} = L_{h,k}^{\mathcal{Q}}$ .

An internally Eulerian subgraph H of any (h, k)-caterpillar consists of a disjoint collection of leaf-paths and loops. The next step is to prove that the value of  $\operatorname{vol}_t(G_{h,k}, H)$  depends only on the number of leaf-paths and loops in H. Finally, we show that  $\operatorname{vol}_t(G_{h,k}, H)$  actually depends on the number of leaf-paths in H and on whether or not H has a loop. For all these tasks we shall need the machinery of the nearest neighbor interchange.

A nearest neighbor interchange (NNI) is a local invertible move performed in G on a trail of length three, marked by dark edges in Figure 4. This move interchanges the ends of the two extreme edges of the trail on the central edge.

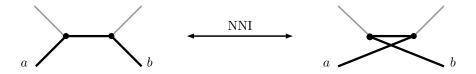


Figure 4: An NNI move on the trail marked by bold darker line segments. One of the incidences of the edges a and b were interchanged.

We refer to the central edge of the trail as the *pivot* of the NNI move. The result of the move is another  $\{1,3\}$ -graph G' on the same number of nodes, edges, and connected components. We call  $\{G,G'\}$  an *NNI pair*.

Consider an NNI pair  $\{G, G'\}$ . Let w and z be weight functions defined on the edges and internal nodes of G, respectively. A weighted NNI is a local invertible move performed on (G, w, z). It is induced by the NNI pair. The result of the move is (G', w', z') where w' and z' are weight functions on the edges and internal nodes of G' defined as follows. Suppose that e = uv is the pivot of the NNI move. Let a and b be the first and last edges in the trail, and c and d be the remaining edges adjacent to e, with their corresponding weights depicted in Figure 5. Note that a, b, c, and d are not necessarily pairwise distinct. The weight function w', defined as in [6], is such that  $w'_f = w_f$  for every  $f \neq e$  and

$$w'_e = w_e + \max\{w_a + w_c, w_b + w_d\} - \max\{w_b + w_c, w_a + w_d\}.$$

The weight function z is such that  $z'_x = z_x$  for every  $x \notin \{u, v\}$ ,

$$z'_u = (w'_b + w'_c + w'_e)/2$$
, and  $z'_v = (w'_a + w'_d + w'_e)/2$ .

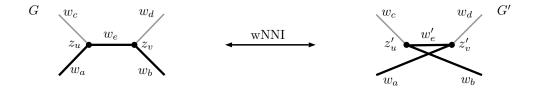


Figure 5: Weighted NNI move on the trail marked by bold darker line segments. Weights modified are the ones on the pivot edge and its end nodes.

Let us restrict our attention to integer points satisfying the parity constraint (3). If (w, z) is one of these points, then so is (w', z'). In fact, for any internally Eulerian subgraph H of G, a weighted NNI move maps points in  $\operatorname{part}_t(G, H)$  into  $\operatorname{part}_t(G', H')$ , where H' is uniquely determined by  $\{G, G'\}$  and H. Indeed, a weighted NNI acts on the parity (of the weight) of the NNI pivot based uniquely on the parity of the edges incident to it (Figure 6). Therefore, as a weighted NNI move is invertible, there is a bijection between  $\operatorname{part}_t(G, H)$  and  $\operatorname{part}_t(G', H')$  for every corresponding pair H and H'. We call  $\{(G, H), (G', H')\}$  a weighted NNI pair.

The following lemma that summarizes the previous discussion is applied extensively, explicitly and implicitly, in several proofs ahead.

**Lemma 15.** Let  $\{(G, H), (G', H')\}$  be a weighted NNI pair. The corresponding weighted NNI is a bijection between  $\operatorname{part}_t(G, H)$  and  $\operatorname{part}_t(G', H')$  and therefore

$$\operatorname{vol}_t(G, H) = \operatorname{vol}_t(G', H').$$

**Example 16.** The graph G in Figure 1 and G' in Figure 7 form an NNI pair  $\{G, G'\}$ . Consider the polytope  $tQ_{G'}$  in Figure 7. The graph G' has four internally Eulerian subgraphs, namely, the subgraphs  $\emptyset$ ,  $H'_{\{1,2\}}$ ,  $H'_{\{1,3\}}$  and  $H'_{\{2,3\}}$  induced by the edge sets  $\emptyset$ ,  $\{1,2\}$ ,  $\{2,3\}$ , and  $\{1,3\}$ , respectively. One can verify that

$$\begin{array}{ll} \mathrm{part}_2(G',\emptyset) &= \{((0,0,0),(0,0)),((2,2,0),(2,2)),((2,0,2),(2,2)),((0,2,2),(2,2))\},\\ \mathrm{part}_2(G',H'_{\{1,2\}}) &= \{((1,1,0),(1,1)),((1,1,2),(2,2))\},\\ \mathrm{part}_2(G',H'_{\{1,3\}}) &= \{((1,0,1),(1,1)),((1,2,1),(2,2))\}, \ and\\ \mathrm{part}_2(G',H'_{\{2,3\}}) &= \{((0,1,1),(1,1)),((2,1,1),(2,2))\}. \end{array}$$

The corresponding weighted NNI pairs are  $\{(G,\emptyset), (G',\emptyset)\}, \{(G,H_{\{1\}}), (G',H'_{\{1,3\}})\}, \{(G,H_{\{2\}}), (G',H'_{\{2,3\}})\}, \{(G,H_{\{1,2\}}), (G',H'_{\{1,2\}})\}.$  Therefore,  $\operatorname{vol}_t(G,H) = \operatorname{vol}_t(G',H')$  for each weighted NNI pair  $\{(G,H), (G',H')\},$  confirming the previous Lemma 15.

Having the tool of weighted NNIs in hand, we are ready to follow the steps delineated.

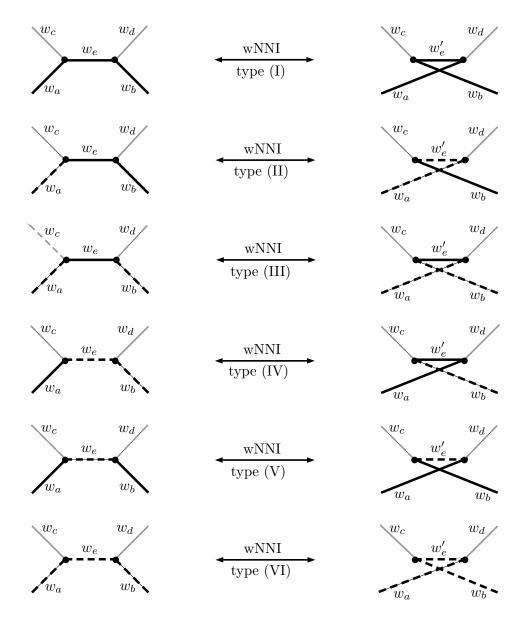


Figure 6: Parity of weights on edges modified by a weighted NNI on a trail marked by bold darker line segments. Dashed line segments indicate edges with odd weights, and continuous line segments indicate edges with even weights.

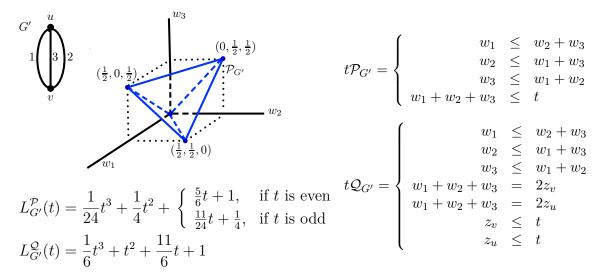


Figure 7: The cubic graph G' is obtained from applying an NNI to the graph G in Figure 1. The pivot of the NNI was the edge 3. The polytope  $\mathcal{P}_{G'}$ , the linear systems of  $t\mathcal{P}_{G'}$  and  $t\mathcal{Q}_{G'}$ , and the Ehrhart quasi-polynomials associated to G' are identical to the ones associated to the  $\{1,3\}$ -tree T in Figure 2. The points in  $t\mathcal{P}_{G'}$  have coordinates  $(w_1, w_2, w_3)$  and the points in  $t\mathcal{Q}_{G'}$  have coordinates  $((w_1, w_2, w_3), (z_v, z_u))$ .

**Lemma 17.** If G is a connected  $\{1,3\}$ -graph with h leaves and cyclomatic number k then  $L_G^{\mathcal{Q}}(t) = L_{h,k}^{\mathcal{Q}}(t)$  for each nonnegative integer t.

*Proof.* The graph G can be transformed into  $G_{h,k}$  through a series of NNI moves [6, Theorem 1]. We have that

$$L_G^{\mathcal{Q}}(t) = \sum_{H} \operatorname{vol}_t(G, H) \tag{9}$$

$$= \sum_{H} \operatorname{vol}_{t}(G_{h,k}, H) \tag{10}$$

$$=L_{h,k}^{\mathcal{Q}}(t),\tag{11}$$

where the summations in (9) and in (10) are over the internally Eulerian subgraphs H of G and  $G_{h,k}$ , respectively. Equalities (9) and (11) are due to the summation in (7), and Equality (10) follows from Lemma 15 by induction on the number of NNI moves.

Now, we prove that  $\operatorname{vol}_t(G_{h,k}, H)$  depends only on the number of leaf-paths and loops in H.

**Lemma 18.** If H and H' are internally Eulerian subgraphs of  $G_{h,k}$  with the same number of leaf-paths and loops, then

$$\operatorname{vol}_t(G_{h,k}, H) = \operatorname{vol}_t(G_{h,k}, H'),$$

for each nonnegative integer t.

*Proof.* Consider a series of weighted NNIs taking the loops in H to the ones in H', by swapping consecutive legs of  $G_{h,k}$  and an appropriated relabeling. Similarly, consider a series of weighted NNIs taking the leaf-edges of leaf-paths in H to the ones in H', by swapping consecutive legs of  $G_{h,k}$  and an appropriated relabeling. The composition of the corresponding bijections is a bijection between  $\operatorname{part}_t(G_{h,k}, H)$  and  $\operatorname{part}_t(G_{h,k}, H')$ . Indeed, the used NNIs are as in Corollary 10 from [6], and preserve the parity of the values on the corresponding leaf-edges and loops. So the lemma follows.

Any internally Eulerian subgraph of  $G_{h,k}$  consists of a disjoint collection of leaf-paths and loops. Inspired by Lemma 18, we shall start to use the following notation:  $H_{i,j}$  denotes any internally Eulerian subgraph with i leaf-paths and j loops. Recalling the definitions (6) in terms of the present notation, we have

$$\operatorname{part}_{t}(G_{h,k}, H_{i,j}) := t \mathcal{Q}_{G_{h,k}} \cap (\mathcal{L}_{G_{h,k}} + (\mathbb{1}_{H_{i,j}}, 0)), \text{ and }$$
  
 $\operatorname{vol}_{t}(G_{h,k}, H_{i,j}) := |\operatorname{part}_{t}(G_{h,k}, H_{i,j})|.$ 

Thus we may rewrite the summation in (7), namely that  $L_G^{\mathcal{Q}}(t) = \sum_H \operatorname{vol}_t(G, H)$ , to obtain

$$L_{h,k}^{\mathcal{Q}}(t) = \sum_{H} \operatorname{vol}_{t}(G_{h,k}, H) = \sum_{i=0}^{\lfloor h/2 \rfloor} \binom{h}{2i} \sum_{j=0}^{k} \binom{k}{j} \operatorname{vol}_{t}(G_{h,k}, H_{i,j}), \tag{12}$$

where the first summation is over all internally Eulerian subgraphs  $H_{i,j}$  of  $G_{h,k}$ .

Next we strengthen Lemma 18 and show that  $\operatorname{vol}_t(G_{h,k}, H)$  does not depend on the exact number of loops in H, but only on whether H has a loop or not. In other words, the value of  $\operatorname{vol}_t(G_{h,k}, H_{i,j})$  depends only on whether j = 0 or  $j \geq 1$ . This will allow us to simplify the summation in (12).

**Lemma 19.** Let h and k be nonnegative integers such that  $k \ge 1$  and  $h + k \ge 2$ . For every nonnegative integer t, i = 0, ..., |h/2|, and j = 1, ..., k, we have that

$$\operatorname{vol}_t(G_{h,k}, H_{i,j}) = \operatorname{vol}_t(G_{h,k}, H_{i,1}).$$

*Proof.* If h = i = 0 and k = j = 2, then the composition of the corresponding two weighted NNIs shown in Figure 8(a) and a relabeling is by Lemma 15 a bijection between  $\operatorname{part}_t(G_{0,2}, H_{0,2})$  and  $\operatorname{part}_t(G_{0,2}, H_{0,1})$ , where the Eulerian subgraphs  $H_{0,2}$  and  $H_{0,1}$  are induced by the edge sets  $\{a, c\}$  and  $\{c\}$ , respectively.

For the other cases, the proof is by induction on j. Suppose that  $j \geq 2$  and let H be an internally Eulerian subgraph of  $G_{h,k}$  with i leaf-paths and j loops. By applying weighted NNIs, we may assume that two of the loops in H are at distance 2. Call a one of these loops. The composition of the corresponding four weighted NNIs shown in Figure 8(b) and a relabeling is by Lemma 15 a bijection between  $\operatorname{part}_t(G_{h,k}, H)$  and  $\operatorname{part}_t(G_{h,k}, H-a)$ . Therefore,  $\operatorname{vol}_t(G_{h,k}, H_{i,j}) = \operatorname{vol}_t(G_{h,k}, H_{i,j-1}) = \operatorname{vol}_t(G_{h,k}, H_{i,1})$ , where the second equality follows from the induction hypothesis.

Finally, we focus special attention on the value of  $\operatorname{vol}_t(G_{h,k}, H_{i,0})$  for i > 0. This will allow us to simplify even further the summation in (12).

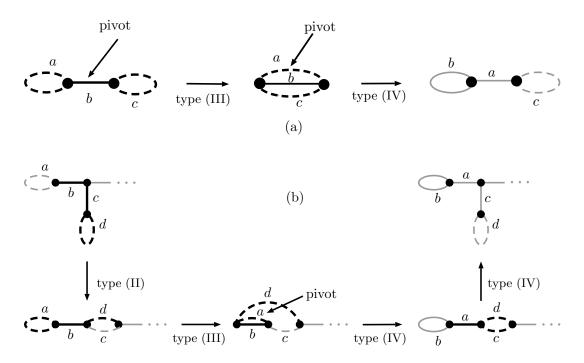


Figure 8: Weighted NNIs that bring two loops with odd weights to one loop with odd weight. The trail of each NNI is marked by bold darker line segments. Dashed line segments indicate the edges with odd weights.

**Lemma 20.** Let h and k be nonnegative integers such that  $k \ge 1$  and  $h + k \ge 2$ . For every nonnegative integer t,  $i = 1, ..., \lfloor h/2 \rfloor$ , we have that

$$\operatorname{vol}_t(G_{h,k}, H_{i,0}) = \operatorname{vol}_t(G_{h,k}, H_{i,1}).$$

*Proof.* Let H be an internally Eulerian subgraph of  $G_{h,k}$  with i leaf-paths and no loops. By applying weighted NNIs we may assume that there is a leaf-path of H at distance 1 from a loop of  $G_{h,k}$ . We denote the loop by a. The composition of the corresponding bijections of the two weighted NNIs illustrated in Figure 9 and a relabeling is a bijection between  $\operatorname{part}_t(G_{h,k}, H)$  and  $\operatorname{part}_t(G_{h,k}, H+a)$ . Therefore,  $\operatorname{vol}_t(G_{h,k}, H_{i,0}) = \operatorname{vol}_t(G_{h,k}, H_{i,1})$ .  $\square$ 

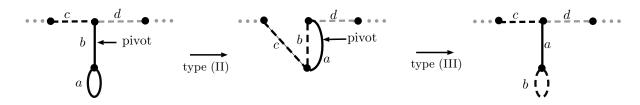


Figure 9: Weighted NNIs that, using a leaf-path with odd weights, change the weight of a loop from even to odd. The trail of each NNI is marked by bold darker line segments. Dashed line segments indicate the edges with odd weights.

#### 3.4 Key summation

Finally, we arrive at the targeted expression of this section by manipulating the summation in (12). If G is a  $\{1,3\}$ -graph with h leaves and cyclomatic number k then

$$L_{G}^{Q}(t) = L_{h,k}^{Q}(t)$$

$$= \sum_{i=0}^{\lfloor h/2 \rfloor} {h \choose 2i} \sum_{j=0}^{k} {k \choose j} \operatorname{vol}_{t}(G_{h,k}, H_{i,j})$$

$$= \sum_{i=0}^{\lfloor h/2 \rfloor} {h \choose 2i} \left( \operatorname{vol}_{t}(G_{h,k}, H_{i,0}) + \sum_{j=1}^{k} {k \choose j} \operatorname{vol}_{t}(G_{h,k}, H_{i,1}) \right)$$

$$= \sum_{i=0}^{\lfloor h/2 \rfloor} {h \choose 2i} \left( \operatorname{vol}_{t}(G_{h,k}, H_{i,0}) + (2^{k} - 1) \operatorname{vol}_{t}(G_{h,k}, H_{i,1}) \right)$$

$$= \sum_{i=0}^{\lfloor h/2 \rfloor} {h \choose 2i} \left( 2^{k} \operatorname{vol}_{t}(G_{h,k}, H_{i,0}) - (2^{k} - 1) (\operatorname{vol}_{t}(G_{h,k}, H_{i,0}) - \operatorname{vol}_{t}(G_{h,k}, H_{i,1})) \right)$$

$$= 2^{k} \sum_{i=0}^{\lfloor h/2 \rfloor} {h \choose 2i} \operatorname{vol}_{t}(G_{h,k}, H_{i,0})$$

$$- (2^{k} - 1) \sum_{i=0}^{\lfloor h/2 \rfloor} {h \choose 2i} \operatorname{vol}_{t}(G_{h,k}, H_{i,0}) - \operatorname{vol}_{t}(G_{h,k}, H_{i,1}),$$

$$= 2^{k} \sum_{i=0}^{\lfloor h/2 \rfloor} {h \choose 2i} \operatorname{vol}_{t}(G_{h,k}, H_{i,0})$$

$$- (2^{k} - 1) (\operatorname{vol}_{t}(G_{h,k}, H_{i,0})$$

$$- (2^{k} - 1) (\operatorname{vol}_{t}(G_{h,k}, H_{i,0}) - \operatorname{vol}_{t}(G_{h,k}, H_{0,1})),$$

$$(15)$$

where Equality (13) is due to Lemma 17, Equality (14) holds by Lemma 19, and Equality (16) follows from Lemma 20. We remind the reader that  $\operatorname{vol}_t(G_{h,k}, H_{0,0}) = \operatorname{vol}_t(G_{h,k}, \emptyset)$ .

## 4 The period for $\{1,3\}$ -graphs

In this section we prove the main theorem of this paper. We claim that we may henceforth consider only t an even integer, due to the following argument. Liu and Osserman [12, Lemma 3.3] proved that if t is a nonnegative odd integer, then  $\operatorname{vol}_t(G_{h,k}, H_{i,j}) = L_{h,k}^{\mathcal{P}}(t)$  for every  $i \geq 0$  and  $j \geq 0$ . Applying their result for t odd to the final relation of the previous section we see that the term (16) vanishes and from the summation in (15) we have that

$$L_G^{\mathcal{Q}}(t) = L_{h,k}^{\mathcal{Q}}(t) = 2^k \sum_{i=0}^{\lfloor h/2 \rfloor} \binom{h}{2i} \operatorname{vol}_t(G_{h,k}, H_{i,0})$$
$$= 2^k \sum_{i=0}^{\lfloor h/2 \rfloor} \binom{h}{2i} L_{h,k}^{\mathcal{P}}(t)$$

$$= N_{G_{h,k}} L_{h,k}^{\mathcal{P}}(t) = N_G L_G^{\mathcal{P}}(t), \tag{17}$$

where Equality (17) follows from the fact that the number of internal Eulerian subgraphs of  $G_{h,k}$  is  $N_{G_{h,k}} = 2^k$  if h = 0 and  $N_{G_{h,k}} = 2^{k+h-1}$  if h > 0. From the relation  $L_G^{\mathcal{Q}}(t) = N_G L_G^{\mathcal{P}}(t)$  for odd t, and the fact that  $L_G^{\mathcal{Q}}(t)$  has period 1 or 2, Liu and Osserman concluded that the odd constituent polynomials  $p_1$  and  $p_3$  of  $L_G^{\mathcal{P}}$  are equal.

#### 4.1 The lemmas of the leaf-paths and of the loops

For the purpose of determining the period of  $L_G^{\mathcal{P}}$ , we settle the values of the summation in (15) and the difference in the term (16) for nonnegative even integers t as stated in the two lemmas below. Lemma 21 is concerned with the summation in (15) and Lemma 22 deals with the difference in the term (16). We recall that for any (h, k)-caterpillar  $G_{h,k}$  we have  $h + k \geq 2$ , by definition.

**Lemma 21** (the leaf-paths). Let h and k be integers such that  $h \ge 0$  and  $k \ge 0$ . For every nonnegative even integer t, we have that

$$2^{k} \sum_{i=0}^{\lfloor h/2 \rfloor} \binom{h}{2i} \operatorname{vol}_{t}(G_{h,k}, H_{i,0}) = N_{G_{h,k}} \left( L_{h,k}^{\mathcal{P}}(t) + \sum_{j=1}^{\lfloor h/2 \rfloor} (-4)^{-j} \binom{h-j}{j} \frac{h}{h-j} L_{h-j,k}^{\mathcal{P}}(t) \right).$$

**Lemma 22** (the loops). Let h and k be integers such that  $h \ge 0$  and  $k \ge 1$ . For every nonnegative even integer t, we have that

$$\operatorname{vol}_{t}(G_{h,k},\emptyset) - \operatorname{vol}_{t}(G_{h,k}, H_{0,1}) = \begin{cases} (\frac{t}{2} + 1)^{k-1} & \text{if } t \equiv 0 \pmod{4} \text{ or } h = 0, \\ 0 & \text{if } t \equiv 2 \pmod{4}. \end{cases}$$

The proofs of these lemmas are technical. Therefore we shall defer the proof of Lemma 21 to Section 5 and of Lemma 22 to Section 6. In the remainder of this section we refer to these lemmas in order to deduce the period of the Ehrhart quasi-polynomial in t associated to the polytope  $\mathcal{P}_G$ .

## 4.2 Proof of the theorem of the period for $\{1,3\}$ -graphs

We have now all tools to deliver a theorem on the period collapse of the quasipolynomial  $L_G^{\mathcal{P}}(t)$ .

Proof of Theorem 5 (the period for  $\{1,3\}$ -graphs). We denote by  $q_0^{h,k}(t)$  and  $q_1^{h,k}(t)$  the constituents of  $L_{h,k}^{\mathcal{Q}}(t)$  for t even and t odd, respectively. We also denote by  $p_j^{h,k}(t)$  the constituent polynomials of  $L_{h,k}^{\mathcal{P}}(t)$  for t congruent to j (mod 4).

Say G has h leaves and cyclomatic number k. Note that  $h + k \geq 2$ . By Wakabayashi [20, Theorem A(ii)] and Lemma 17, we may assume that  $G = G_{h,k}$ . Liu and Osserman [12] proved that  $p_1^{h,k} = p_3^{h,k}$ . Therefore the period of  $L_{h,k}^{\mathcal{P}}$  is at most 2 if and only if  $p_0^{h,k} = p_2^{h,k}$ .

Case 1. k = 0.

Then  $G_{h,0}$  is a tree, and the period of  $L_{h,0}^{\mathcal{P}}(t)$  is 2 by Corollary 4.

Case 2. h = 0.

Then  $G_{0,k}$  is a cubic graph and (15) is equal to  $2^k \operatorname{vol}_t(G_{0,k}, H_{0,0}) = 2^k \operatorname{vol}_t(G_{0,k}, \emptyset)$  because the summation reduces to the term associated with i = 0. By Lemma 22, for every nonnegative even integer t, we have that

$$L_{0,k}^{\mathcal{Q}}(t) = q_0^{0,k}(t) = 2^k \operatorname{vol}_t(G_{0,k}, \emptyset) - (2^k - 1)(\operatorname{vol}_t(G_{0,k}, \emptyset) - \operatorname{vol}_t(G_{h,k}, H_{0,1}))$$
$$= N_{G_{0,k}} L_{0,k}^{\mathcal{P}}(t) - (N_{G_{0,k}} - 1) \left(\frac{t}{2} + 1\right)^{k-1}.$$

Therefore,  $q_0^{0,k}(t) = N_{G_{0,k}} p_0^{0,k}(t) - (N_{G_{0,k}} - 1)(\frac{t}{2} + 1)^{k-1}$  for  $t \equiv 0 \pmod{4}$ . Because both sides of the latter equation are polynomials in t, the equality holds for every real t. Similarly,  $q_0^{0,k}(t) = N_{G_{0,k}} p_2^{0,k}(t) - (N_{G_{0,k}} - 1)(\frac{t}{2} + 1)^{k-1}$  for every real t. Hence  $p_0^{0,k} = p_2^{0,k}$  and the period of  $L_{0,k}^{\mathcal{P}}$  is at most 2.

#### Case 3. $h \ge 1$ and $k \ge 1$ .

Now we tackle all other connected  $\{1,3\}$ -graphs. For  $t \equiv 0 \pmod{4}$ , applying Lemma 21 and Lemma 22 to the final relation of the previous section, we obtain that

$$L_{h,k}^{\mathcal{Q}}(t) = q_0^{h,k}(t) = N_{G_{h,k}} \left( L_{h,k}^{\mathcal{P}}(t) + \sum_{j=1}^{\lfloor h/2 \rfloor} (-4)^{-j} \binom{h-j}{j} \frac{h}{h-j} L_{h-j,k}^{\mathcal{P}}(t) \right) - (2^k - 1) \left( \frac{t}{2} + 1 \right)^{k-1}.$$
(18)

Similarly, for  $t \equiv 2 \pmod{4}$ , we can derive that

$$L_{h,k}^{\mathcal{Q}}(t) = q_0^{h,k}(t) = N_{G_{h,k}} \left( L_{h,k}^{\mathcal{P}}(t) + \sum_{j=1}^{\lfloor h/2 \rfloor} (-4)^{-j} \binom{h-j}{j} \frac{h}{h-j} L_{h-j,k}^{\mathcal{P}}(t) \right). \tag{19}$$

Using (18) with  $t \equiv 0 \pmod{4}$ , and substituting  $N_{G_{h,k}} = 2^{h+k-1}$ , we have

$$\frac{q_0^{h,k}(t)}{N_{G_{h,k}}} = p_0^{h,k}(t) + \sum_{j=1}^{\lfloor h/2 \rfloor} (-4)^{-j} \binom{h-j}{j} \frac{h}{h-j} p_0^{h-j,k}(t) - \frac{1}{2^{h-1}} (1 - \frac{1}{2^k}) \left(\frac{t}{2} + 1\right)^{k-1}. \tag{20}$$

Because both sides of Equality (20) are polynomials in t, it holds for every t. Similarly, from Equality (19), we deduce that, for every t,

$$\frac{q_0^{h,k}(t)}{N_{G_{h,k}}} = p_2^{h,k}(t) + \sum_{j=1}^{\lfloor h/2 \rfloor} (-4)^{-j} \binom{h-j}{j} \frac{h}{h-j} p_2^{h-j,k}(t). \tag{21}$$

By subtracting (20) from (21) and rearranging the terms, we deduce that

$$(p_0^{h,k} - p_2^{h,k})(t) + \sum_{j=1}^{\lfloor h/2 \rfloor} (-4)^{-j} \binom{h-j}{j} \frac{h}{h-j} (p_0^{h-j,k} - p_2^{h-j,k})(t)$$

$$= \frac{1}{2^{h-1}} \left( 1 - \frac{1}{2^k} \right) \left( \frac{t}{2} + 1 \right)^{k-1}. \tag{22}$$

We define

$$d(\alpha) := \frac{(-4)^{\alpha}}{\alpha} \frac{(p_0^{\alpha,k} - p_2^{\alpha,k})(t)}{(1 - \frac{1}{2^k})(\frac{t}{2} + 1)^{k-1}},$$
(23)

so that we may divide (22) by  $\frac{h}{(-4)^h}(1-\frac{1}{2^k})(\frac{t}{2}+1)^{k-1}$  and obtain

$$\sum_{j=0}^{h-1} \binom{h-j}{j} d(h-j) = \frac{(-4)^h}{h \, 2^{h-1}} = \frac{(-1)^h \, 2^{h+1}}{h}. \tag{24}$$

The  $j = \lfloor h/2 \rfloor + 1, \dots, h-1$  terms are all equal to zero.

From one of the so called simpler Chebyshev inverse relations [17, Table 2.3, item 5], we derive that

$$d(h) = \sum_{j=0}^{h-1} \left[ \binom{h+j-1}{j} - \binom{h+j-1}{j-1} \right] \frac{(-1)^j (-1)^{h-j} 2^{h-j+1}}{h-j}$$

$$= \sum_{j=0}^{h-1} \frac{h-j}{h+j} \binom{h+j}{j} \frac{(-1)^h 2^{h-j+1}}{h-j}$$

$$= (-1)^h \sum_{j=0}^{h-1} \frac{2^{h-j+1}}{h+j} \binom{h+j}{j}.$$

Thus, we have that d(h) > 0 for every positive even h and d(h) < 0 for every odd h. Finally, substituting in definition (23) with  $\alpha = h$ , we obtain that

$$p_0^{h,k}(t) - p_2^{h,k}(t) = \frac{h}{4^h} \left( 1 - \frac{1}{2^k} \right) \left( \frac{t}{2} + 1 \right)^{k-1} \sum_{j=0}^{h-1} \frac{2^{h-j+1}}{h+j} \binom{h+j}{j} > 0$$

and conclude that  $p_0^{h,k}(t) > p_2^{h,k}(t)$  for every  $t \geq 0$ . Consequently, the period of  $L_{h,k}^{\mathcal{P}}$  is 4.

## 5 Lemma of the leaf-paths

The purpose of this section is to prove Lemma 21, the lemma of the leaf-paths. To achieve this goal, we first establish the necessary tools in two supporting lemmas and a combinatorial identity. We defer the proofs of these lemmas to the end of this section.

Then, in Section 5.3, we give a tool that is purely set-theoretic, namely Lemma 26. This tool allows us to shift our lattice, in order to measure the difference between two discrete volumes. This is also the main tool that is employed in the proofs of the supporting lemmas of this section, and in the following section. Finally, we end this section by giving the proofs of the supporting lemmas.

Before continuing we would like to recall that  $H_{i,j}$  denotes any internally Eulerian subgraph with i leaf-paths and j loops and that for any (h, k)-caterpillar  $G_{h,k}$  we have  $h + k \ge 2$ , by definition.

#### 5.1 Supporting lemmas, and an identity

Here, we state, without proving, the lemmas and identity used in the proof of the lemma of the leaf-paths.

**Lemma 23** (one less leaf-path). Let h and k be integers such that  $h \ge 2$  and  $k \ge 1$ . For every nonnegative even integer t and for every integer i = 1, ..., |h/2|, we have that

$$\operatorname{vol}_{t}(G_{h,k}, H_{i,0}) = \operatorname{vol}_{t}(G_{h,k}, H_{i-1,0}) - \operatorname{vol}_{t}(G_{h-1,k}, H_{i-1,0}). \tag{25}$$

From this lemma, we can deduce the following identity.

**Lemma 24** (the size of coset  $H_{i,0}$ ). Let h and k be integers such that  $h \ge 0$  and  $k \ge 1$ . For every nonnegative even integer t and for every integer  $i = 0, ..., \lfloor h/2 \rfloor$ , we have that

$$\operatorname{vol}_{t}(G_{h,k}, H_{i,0}) = \sum_{j=0}^{i} (-1)^{j} {i \choose j} L_{h-j,k}^{\mathcal{P}}(t).$$

We shall derive the lemma of the leaf-paths from Lemma 24 and the following known combinatorial identity.

**Identity 25.** [17, Problem 18(a), Chap. 6] For every integer  $h \ge 1$  and  $j = 0, \ldots, \lfloor h/2 \rfloor$ ,

$$\sum_{i=j}^{\lfloor h/2\rfloor} \binom{h}{2i} \binom{i}{j} \ = \ \binom{h-j}{j} \frac{h}{h-j} \, 2^{h-1-2j}.$$

## 5.2 Proof of the lemma of the leaf-paths

We are now prepared to carry out the proof of Lemma 21, upon which the proof of Theorem 5 of the period of  $L_G^{\mathcal{P}}(t)$  rests.

Proof of Lemma 21 (the leaf-paths). The lemma clearly holds for h = 0. Thus we may assume  $h \ge 1$  and we have that

$$2^{k} \sum_{i=0}^{\lfloor h/2 \rfloor} \binom{h}{2i} \operatorname{vol}_{t}(G_{h,k}, H_{i,0}) = 2^{k} \sum_{i=0}^{\lfloor h/2 \rfloor} \binom{h}{2i} \sum_{j=0}^{i} (-1)^{j} \binom{i}{j} L_{h-j,k}^{\mathcal{P}}(t)$$

$$= 2^{k} \sum_{i=0}^{\lfloor h/2 \rfloor} \binom{h}{2i} \left( L_{h,k}^{\mathcal{P}}(t) + \sum_{j=1}^{i} (-1)^{j} \binom{i}{j} L_{h-j,k}^{\mathcal{P}}(t) \right)$$

$$= 2^{k+h-1} L_{h,k}^{\mathcal{P}}(t) + 2^{k} \sum_{i=1}^{\lfloor h/2 \rfloor} \binom{h}{2i} \sum_{j=1}^{i} (-1)^{j} \binom{i}{j} L_{h-j,k}^{\mathcal{P}}(t)$$

$$= N_{G_{h,k}} L_{h,k}^{\mathcal{P}}(t) + 2^{k} \sum_{j=1}^{\lfloor h/2 \rfloor} (-1)^{j} L_{h-j,k}^{\mathcal{P}}(t) \sum_{i=j}^{\lfloor h/2 \rfloor} \binom{h}{2i} \binom{i}{j}$$

$$(27)$$

$$= N_{G_{h,k}} L_{h,k}^{\mathcal{P}}(t) + 2^{k} \sum_{j=1}^{\lfloor h/2 \rfloor} (-1)^{j} L_{h-j,k}^{\mathcal{P}}(t) \binom{h-j}{j} \frac{h}{h-j} 2^{h-1-2j}$$

$$= N_{G_{h,k}} L_{h,k}^{\mathcal{P}}(t) + 2^{k+h-1} \sum_{j=1}^{\lfloor h/2 \rfloor} (-1)^{j} 2^{-2j} \binom{h-j}{j} \frac{h}{h-j} L_{h-j,k}^{\mathcal{P}}(t)$$

$$= N_{G_{h,k}} L_{h,k}^{\mathcal{P}}(t) + N_{G_{h,k}} \sum_{j=1}^{\lfloor h/2 \rfloor} (-4)^{-j} \binom{h-j}{j} \frac{h}{h-j} L_{h-j,k}^{\mathcal{P}}(t)$$

$$= N_{G_{h,k}} \left( L_{h,k}^{\mathcal{P}}(t) + \sum_{j=1}^{\lfloor h/2 \rfloor} (-4)^{-j} \binom{h-j}{j} \frac{h}{h-j} L_{h-j,k}^{\mathcal{P}}(t) \right),$$

$$= N_{G_{h,k}} \left( L_{h,k}^{\mathcal{P}}(t) + \sum_{j=1}^{\lfloor h/2 \rfloor} (-4)^{-j} \binom{h-j}{j} \frac{h}{h-j} L_{h-j,k}^{\mathcal{P}}(t) \right),$$

where (26) holds by Lemma 24, Equality (27) is the result of exchanging the two summations, and finally Equality (28) is due to Identity (25).  $\Box$ 

#### 5.3 Bijections by shifting

Here we derive a purely set-theoretic tool, Lemma 26, that gives a partial bijection by shifting. We will employ this tool repeatedly in the remaining proofs of this section, and in those of the next section.

We begin with some notation that is followed by the lemma. For an e in E and a set  $\mathcal{X} \subset \mathbb{R}^E$ , let

$$\operatorname{Out}_e(\mathcal{X}) := \{ w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is even and } w + \mathbb{1}_e \notin \mathcal{X} \} \text{ and } \operatorname{In}_e(\mathcal{X}) := \{ w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is odd and } w - \mathbb{1}_e \notin \mathcal{X} \}.$$

We think of the function  $w \mapsto w + \mathbb{1}_e$  as a shifting of the integer points with even e-coordinate to the integer points with odd e-coordinate.  $\operatorname{Out}_e(\mathcal{X})$  is the set of integer points of  $\mathcal{X}$  with even e-coordinate that are moved outside  $\mathcal{X}$  and  $\operatorname{In}_e(\mathcal{X})$  is the set of integer points of  $\mathcal{X}$  with odd e-coordinate that are images of points outside  $\mathcal{X}$  (Figure 10). Roughly, considering the integer points, the shifting enters  $\mathcal{X}$  through  $\operatorname{In}_e(\mathcal{X})$  and leaves  $\mathcal{X}$  through  $\operatorname{Out}_e(\mathcal{X})$ .

For any subset  $\mathcal{X} \subset \mathbb{R}^E$ , we define

$$\operatorname{vol}_{e \operatorname{even}}(\mathcal{X}) := |\{w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is even}\}|,$$

and

$$\operatorname{vol}_{e \operatorname{odd}}(\mathcal{X}) := |\{w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is odd}\}|.$$

**Lemma 26** (the bijection by shifting). If  $\mathcal{X} \subset \mathbb{R}^E$  is bounded and  $e \in E$ , then

$$\operatorname{vol}_{e \operatorname{even}}(\mathcal{X}) - \operatorname{vol}_{e \operatorname{odd}}(\mathcal{X}) = |\operatorname{Out}_e(\mathcal{X})| - |\operatorname{In}_e(\mathcal{X})|.$$

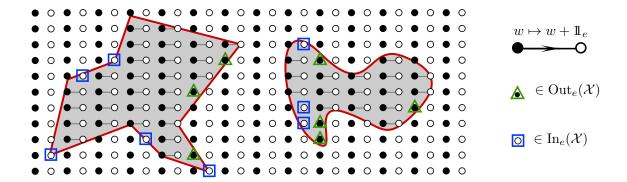


Figure 10: Filled circles represent points with an even e-coordinate, while empty circles represent points with an odd e-coordinate. Let  $\mathcal{X}$  be the set of points in the grey area. The circles marked by triangles represent the set  $\operatorname{Out}_e(\mathcal{X})$  and the circles marked by squares represent  $\operatorname{In}_e(\mathcal{X})$ . The difference between the number of filled and empty circles in  $\mathcal{X}$  is  $|\operatorname{Out}_e(\mathcal{X})| - |\operatorname{In}_e(\mathcal{X})| = 7 - 8 = -1$ .

*Proof.* We have that

$$\operatorname{Out}_e(\mathcal{X}) \subseteq \{w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is even}\}\ \text{and } \operatorname{In}_e(\mathcal{X}) \subseteq \{w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is odd}\}.$$

Therefore,

$$\operatorname{vol}_{e \operatorname{even}}(\mathcal{X}) = |\{w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is even}\} \setminus \operatorname{Out}_e(\mathcal{X})| + |\operatorname{Out}_e(\mathcal{X})|, \text{ and} \qquad (29)$$

$$\operatorname{vol}_{e \operatorname{odd}}(\mathcal{X}) = |\{w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is odd}\} \setminus \operatorname{In}_e(\mathcal{X})| + |\operatorname{In}_e(\mathcal{X})|. \qquad (30)$$

The function  $w \mapsto w + \mathbb{1}_e$  is a bijection from  $\{w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is even}\} \setminus \text{Out}_e(\mathcal{X})$  to  $\{w \in \mathcal{X} \cap \mathbb{Z}^E : w_e \text{ is odd}\} \setminus \text{In}_e(\mathcal{X})$ . As  $\mathcal{X}$  is bounded,  $\text{vol}_{e \text{ even}}(\mathcal{X})$  and  $\text{vol}_{e \text{ odd}}(\mathcal{X})$  are finite and the result follows from Equalities (29) and (30).

## 5.4 Proofs of the supporting lemmas

We begin this section by presenting the proof of the supporting Lemma 23 of one less leaf-path.

Proof of Lemma 23 (one less leaf-path). By Lemma 18 we may assume that one of the leaf-paths P in  $H_{i,0}$  consists of two stems a and b of  $G_{h,k}$  incident to a vertex r. Let u and v be the leaves adjacent to a and b, respectively (Figure 11(a)). Let  $G_{h-1,k} := G_{h,k} - \{u,v\}$  and let c be the edge incident to r other than a or b (Figure 11(b)). Hence,  $H_{i-1,0} := H_{i,0} - P$  is an internally Eulerian subgraph of  $G_{h,k}$  and  $G_{h-1,k}$ .

For conciseness, let  $W := \operatorname{part}_t(G_{h,k}, H_{i,0})$ ,  $Z := \operatorname{part}_t(G_{h,k}, H_{i-1,0})$ , and  $X := \operatorname{part}_t(G_{h-1,k}, H_{i-1,0})$ . Equality (25) is equivalent to |X| = |Z| - |W|. We partition Z and W into

$$Z_{<} := \{ (w, z) \in Z : z_r < t \},\$$

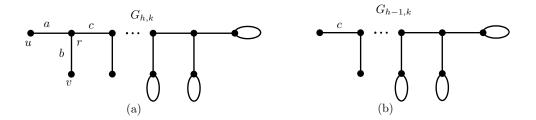


Figure 11: (a)  $G_{h,k}$  and the stems a and b. (b)  $G_{h-1,k}$ .

$$Z_{=} := \{(w, z) \in Z : z_{r} = t\},$$

$$W_{<} := \{(w, z) \in W : w_{c} < w_{a} + w_{b}\} \text{ and }$$

$$W_{=} := \{(w, z) \in W : w_{c} = w_{a} + w_{b}\}.$$

Let I be the set of internal nodes of  $G_{h,k}$ , and consider the function  $\phi_+: \mathbb{R}^E \times \mathbb{R}^I \to \mathbb{R}^E \times \mathbb{R}^I$ , given by the shifting operation

$$\phi_{+}(w,z) = (w,z) + (\mathbb{1}_{a,b}, \mathbb{1}_r).$$

One can verify that  $\phi_+$  is a bijection from  $Z_{<}$  to  $W_{<}$  and therefore  $|Z_{<}| = |W_{<}|$ . Consequently, in order to prove the lemma, it suffices to show that

$$|X| = |Z| - |W| = (|Z_{<}| + |Z_{=}|) - (|W_{<}| + |W_{=}|) = |Z_{=}| - |W_{=}|.$$

Conceptually, from the viewpoint of shiftings and Lemma 26,  $W_{=}$  is reminiscent of  $\operatorname{In}_{a,b}(Z)$  and  $Z_{=}$  is reminiscent of  $\operatorname{Out}_{a,b}(Z)$ , meaning that the shifting operation  $\phi_{+}$  enters Z through  $W_{=}$  and leaves Z through  $Z_{=}$ .

For j = 0, 1, ..., t/2, let

$$X_j := \{(w, z) \in X : w_c = 2j\}$$

and, for j = 0, 1, ..., t/2 and  $\ell = 0, 1, ..., j$ , let

$$Z_{j,\ell} := \{(w,z) \in Z_{=} : w_a = t - 2\ell, \ w_b = t - 2j + 2\ell, \ w_c = 2j\}.$$

The edge c is not in  $H_{i-1,0}$ , thus the sets  $X_0, \ldots, X_{t/2}$  form a partition of X. The sets  $Z_{0,0}, Z_{1,0}, \ldots, Z_{t/2,t/2}$  form a partition of  $Z_{=}$ . Moreover, for each j, we have that  $|Z_{j,\ell}| = |X_j|$ , for every  $\ell$ . Therefore,

$$|Z_{-}| = \sum_{j=0}^{t/2} \sum_{\ell=0}^{j} |Z_{j,\ell}| = \sum_{j=0}^{t/2} \sum_{\ell=0}^{j} |X_{j}| = \sum_{j=0}^{t/2} (j+1)|X_{j}|.$$

For  $j = 0, 1, \dots, t/2$  and  $\ell = 1, \dots, j$ , we define

$$W_{j,\ell} := \{(w,z) \in W_{=} : w_a = 2\ell - 1, \ w_b = 2j - 2\ell + 1, \ w_c = 2j\}.$$

The sets  $W_{0,1}, \ldots, W_{t/2,t/2}$  form a partition of  $W_{=}$  and  $|W_{j,\ell}| = |X_j|$  for each j and each  $\ell$ . Hence,

$$|W_{=}| = \sum_{j=0}^{t/2} \sum_{\ell=1}^{j} |W_{j,\ell}| = \sum_{j=0}^{t/2} \sum_{\ell=1}^{j} |X_{j}| = \sum_{j=0}^{t/2} j |X_{j}|.$$

Therefore,  $|Z_{=}| - |W_{=}| = \sum_{j=0}^{t/2} |X_{j}| = |X|$ .

Finishing this section we handle the proof of the supporting Lemma 24 of the size of coset  $H_{i,0}$ .

Proof of Lemma 24 (the size of coset  $H_{i,0}$ ). The proof is by induction on i, the number of leaf-paths in  $H_{i,0}$ . For i=0, the lemma follows from Equality (8). Hence, we may assume that  $i \geq 1$  and thus  $h \geq 2$ . We then have that

$$vol_t(G_{h,k}, H_{i,0}) = vol_t(G_{h,k}, H_{i-1,0}) - vol_t(G_{h-1,k}, H_{i-1,0})$$
(31)

$$= \sum_{j=0}^{i-1} (-1)^j \binom{i-1}{j} L_{h-j,k}^{\mathcal{P}}(t) - \sum_{j=0}^{i-1} (-1)^j \binom{i-1}{j} L_{h-1-j,k}^{\mathcal{P}}(t)$$
 (32)

$$= \sum_{j=0}^{i-1} (-1)^{j} {i-1 \choose j} L_{h-j,k}^{\mathcal{P}}(t) - \sum_{j=1}^{i} (-1)^{j-1} {i-1 \choose j-1} L_{h-j,k}^{\mathcal{P}}(t)$$

$$= \sum_{j=0}^{i-1} (-1)^{j} {i-1 \choose j} L_{h-j,k}^{\mathcal{P}}(t) + \sum_{j=1}^{i} (-1)^{j} {i-1 \choose j-1} L_{h-j,k}^{\mathcal{P}}(t)$$

$$= \sum_{j=0}^{i} (-1)^{j} {i \choose j} L_{h-j,k}^{\mathcal{P}}(t),$$
(33)

where Equality (31) holds by Lemma 23, Equality (32) follows by induction, and Equality (33) holds by changing j to j + 1 in the second sum.

## 6 Lemma of the loops

This section is dedicated to establishing Lemma 22 of the loops. As always,  $H_{i,j}$  denotes any internally Eulerian subgraph of  $G_{h,k}$  with i leaf-paths and j loops. To handle the boundary conditions on the indices, we recall that  $H_{0,0} = \emptyset$ . Moreover, for any (h, k)-caterpillar  $G_{h,k}$  we have  $h + k \ge 2$ , by definition.

## 6.1 Supporting lemmas

The proof of Lemma 22 of the loops relies on two supporting lemmas. For conciseness, we recall that if  $\mathcal{X}$  is a subset of  $\mathbb{R}^E$  we write  $\operatorname{vol}_{a\operatorname{even}}(\mathcal{X})$  to denote the number of integer points w in  $\mathcal{X}$  such that  $w_a$  is even and we write  $\operatorname{vol}_{a\operatorname{odd}}(\mathcal{X})$  to denote the number of integer points w in  $\mathcal{X}$  such that  $w_a$  is odd.

**Lemma 27** ( $\mathcal{Q}$  to  $\mathcal{P}$ ). Let h and k be integers such that  $h \geq 0$  and  $k \geq 1$  and let a be a leaf-edge in  $G_{h+1,k-1}$ . For every nonnegative even integer t, we have that

$$\operatorname{vol}_t(G_{h,k},\emptyset) - \operatorname{vol}_t(G_{h,k},H_{0,1}) = \operatorname{vol}_{a \operatorname{even}}(t\mathcal{P}_{G_{h+1,k-1}}) - \operatorname{vol}_{a \operatorname{odd}}(t\mathcal{P}_{G_{h+1,k-1}}).$$

**Lemma 28** (the evaluation). Let h and k be integers such that  $h \ge 1$  and  $k \ge 0$  and let a be a leaf-edge in  $G_{h,k}$ . For every nonnegative even integer t, we have that

$$\operatorname{vol}_{a \operatorname{even}}(t\mathcal{P}_{G_{h,k}}) - \operatorname{vol}_{a \operatorname{odd}}(t\mathcal{P}_{G_{h,k}}) = \begin{cases} (\frac{t}{2} + 1)^k & \text{if } t \equiv 0 \pmod{4} \text{ or } h = 1, \\ 0 & \text{if } t \equiv 2 \pmod{4}. \end{cases}$$

#### 6.2 Proof of the lemma of the loops

Lemma 22 is a straightforward consequence of the supporting lemmas.

Proof of Lemma 22 (the loops). Because h and k are integers such that  $h \ge 0$  and  $k \ge 1$ ,  $G_{h+1,k-1}$  has a leaf-edge a. We have that

$$\operatorname{vol}_{t}(G_{h,k}, \emptyset) - \operatorname{vol}_{t}(G_{h,k}, H_{0,1}) = \operatorname{vol}_{a \operatorname{even}}(t \mathcal{P}_{G_{h+1,k-1}}) - \operatorname{vol}_{a \operatorname{odd}}(t \mathcal{P}_{G_{h+1,k-1}}) 
= \begin{cases} (\frac{t}{2} + 1)^{k+1} & \text{if } t \equiv 0 \pmod{4} \text{ or } h = 0, \\ 0 & \text{if } t \equiv 2 \pmod{4}, \end{cases}$$

where the first equality is by Lemma 27 and, noting that  $h+1 \ge 1$  and  $k-1 \ge 0$ , the second equality is obtained by applying Lemma 28 to  $G_{h+1,k-1}$ .

## 6.3 Proofs of the supporting lemmas

We begin by proving Lemma 27, which moves our attention from the polytope  $\mathcal{Q}_{G_{h,k}}$  to the polytope  $\mathcal{P}_{G_{h,k}}$ .

Proof of Lemma 27 (Q to P). Let a be an edge of  $G_{h,k}$ , and let  $\ell$  be a loop of  $G_{h,k}$ , both incident with the same node r. By Lemma 18, we may assume that  $H_{0,1}$  is the subgraph induced by the loop  $\ell$  and that the configuration is as depicted in Figure 12(a). In the proof we write  $\ell$  instead of  $H_{0,1}$ . This lemma is concerned with the evaluation of

$$\operatorname{vol}_{t}(G_{h,k},\emptyset) - \operatorname{vol}_{t}(G_{h,k},\ell) = |\operatorname{part}_{t}(G_{h,k},\emptyset)| - |\operatorname{part}_{t}(G_{h,k},\ell)|.$$

For  $\alpha = 0, \ldots, t$ , let

$$Z_{\alpha}^{\ell \text{ even}} := \{(w, z) \in \text{part}_t(G_{h,k}, \emptyset) : z_r = \alpha\}, \text{ and}$$
  
 $Z_{\alpha}^{\ell \text{ odd}} := \{(w, z) \in \text{part}_t(G_{h,k}, \ell) : z_r = \alpha\}.$ 

The sets  $Z_0^{\ell \text{ even}}, \ldots, Z_t^{\ell \text{ even}}$  form a partition of  $\operatorname{part}_t(G_{h,k}, \emptyset)$  and the sets  $Z_0^{\ell \text{ odd}}, \ldots, Z_t^{\ell \text{ odd}}$  form a partition of  $\operatorname{part}_t(G_{h,k}, \ell)$ . Accordingly,

$$\operatorname{vol}_{t}(G_{h,k},\emptyset) = |Z_{0}^{\ell \operatorname{even}}| + \dots + |Z_{t}^{\ell \operatorname{even}}|, \text{ and}$$
(34)

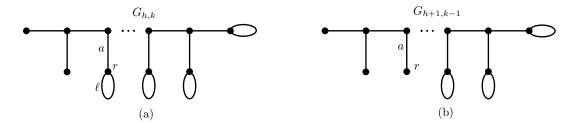


Figure 12: (a)  $G_{h,k}$  and the edge a, loop  $\ell$  and node r. (b)  $G_{h+1,k-1}$  obtained from  $G_{h,k}$ by deleting  $\ell$ .

$$\operatorname{vol}_{t}(G_{h,k},\ell) = |Z_{0}^{\ell \operatorname{odd}}| + \dots + |Z_{t}^{\ell \operatorname{odd}}|. \tag{35}$$

Let E be the set of edges and I be the set of internal nodes of  $G_{h,k}$ . Consider the function  $\phi_+: \mathbb{R}^E \times \mathbb{R}^I \to \mathbb{R}^E \times \mathbb{R}^I$ , given by the shifting operation

$$\phi_{+}(w,z) = (w,z) + (\mathbb{1}_{\ell}, \mathbb{1}_{r}).$$

One can verify that for every  $\alpha$  even,  $\alpha < t$ , the function  $\phi_+$  is a bijection between  $Z_{\alpha}^{\ell \, \text{even}}$ and  $Z_{\alpha+1}^{\ell \text{ odd}}$ . Therefore, for every  $\alpha$  even,  $\alpha < t$ , we have that  $|Z_{\alpha}^{\ell \text{ even}}| = |Z_{\alpha+1}^{\ell \text{ odd}}|$ . Similarly, consider the function  $\phi_{-} : \mathbb{R}^{E} \times \mathbb{R}^{I} \to \mathbb{R}^{E} \times \mathbb{R}^{I}$ , given by the reverse shifting

$$\phi_{-}(w,z) = (w,z) - (\mathbb{1}_{\ell}, \mathbb{1}_{r}).$$

One can also check that for every  $\alpha$  odd,  $\alpha \leq t$ , the function  $\phi_{-}$  is a bijection between  $Z_{\alpha}^{\ell \, \mathrm{even}}$  and  $Z_{\alpha-1}^{\ell \, \mathrm{odd}}$ . Consequently, for every  $\alpha \, \mathrm{odd}$ ,  $\alpha \leq t$ , we have that  $|Z_{\alpha}^{\ell \, \mathrm{even}}| = |Z_{\alpha-1}^{\ell \, \mathrm{odd}}|$ .

Keeping in mind the induced bijections given by the shiftings  $\phi_+$  and  $\phi_-$ , and using Equalities (34) and (35), we can see that for every nonnegative even integer t:

$$\operatorname{vol}_{t}(G_{h,k},\emptyset) - \operatorname{vol}_{t}(G_{h,k},\ell) = |Z_{0}^{\ell \operatorname{even}}| + \cdots + |Z_{t}^{\ell \operatorname{even}}| - |Z_{0}^{\ell \operatorname{odd}}| - \cdots - |Z_{t}^{\ell \operatorname{odd}}| 
= (|Z_{0}^{\ell \operatorname{even}}| - |Z_{1}^{\ell \operatorname{odd}}|) + (|Z_{1}^{\ell \operatorname{even}}| - |Z_{0}^{\ell \operatorname{odd}}|) + \cdots 
+ (|Z_{t-2}^{\ell \operatorname{even}}| - |Z_{t-1}^{\ell \operatorname{odd}}|) + (|Z_{t-1}^{\ell \operatorname{even}}| - |Z_{t-2}^{\ell \operatorname{odd}}|) 
+ (|Z_{t}^{\ell \operatorname{even}}| - |Z_{t}^{\ell \operatorname{odd}}|) 
= |Z_{t}^{\ell \operatorname{even}}| - |Z_{t}^{\ell \operatorname{odd}}|.$$

Our next and last step in this proof is showing that for every nonnegative even integer t:

$$|Z_t^{\ell \text{ even}}| - |Z_t^{\ell \text{ odd}}| = \text{vol}_{a \text{ even}}(t\mathcal{P}_{G_{h+1,k-1}}) - \text{vol}_{a \text{ odd}}(t\mathcal{P}_{G_{h+1,k-1}}).$$

$$(36)$$

For this we outline bijections between the sets  $Z_t^{\ell \text{ even}}$  and  $\{w' \in t\mathcal{P}_{G_{h+1,k-1}} : w'_a \text{ is even}\}$ and between the sets  $Z_t^{\ell \text{ odd}}$  and  $\{w' \in t\mathcal{P}_{G_{h+1,k-1}} : w'_a \text{ is odd}\}.$ 

The (h+1,k-1)-caterpillar  $G_{h+1,k-1}$  will be seen as resulting from the deletion of the loop  $\ell$  from  $G_{h,k}$ . Because we denote the edge set of  $G_{h,k}$  by E, then the edge set of  $G_{h+1,k-1}$  is  $E \setminus \{\ell\}$ . If (w,z) is a point in  $Z_t^{\ell \text{ even}} \cup Z_t^{\ell \text{ odd}}$ , then  $w_e$  is an even integer for each e in  $E \setminus \{\ell\}$ . Moreover, as each point  $(w,z) \in \mathcal{Q}_{G_{h,k}}$  satisfies  $w_a + 2w_\ell = 2z_r$ , if (w, z) is a point in  $Z_t^{\ell \text{ even}}$ , then  $w_a \equiv 0 \pmod{4}$  and if (w, z) is a point in  $Z_t^{\ell \text{ odd}}$ , then  $w_a \equiv 2 \pmod{4}$ .

Consider the function from  $\mathbb{R}^E \times \mathbb{R}^I$  to  $\mathbb{R}^{E\setminus\{\ell\}}$  given by  $(w,z)\mapsto w'$ , where  $w'_e=w_e/2$  for each e in  $E\setminus\{\ell\}$ . From the above observations concerning parity, one can verify that  $(w,z)\mapsto w'$  injectively maps  $Z^{\ell \text{ even}}_t$  into  $\{w'\in t\mathcal{P}_{G_{h+1,k-1}}: w'_a \text{ is even}\}$  and  $Z^{\ell \text{ odd}}_t$  into  $\{w'\in t\mathcal{P}_{G_{h+1,k-1}}: w'_a \text{ is odd}\}$ .

Conversely, consider the function from  $\mathbb{R}^{E\setminus\{\ell\}}$  to  $\mathbb{R}^E\times\mathbb{R}^I$  given by  $w'\mapsto (w,z)$  as follows. Firstly we set  $w_e=2w'_e$  for each e in  $E\setminus\{\ell\}$  and set  $w_\ell=t-w'_a$ . Secondly, for each internal node v of  $G_{h,k}$ , we set  $z_v$  to the sum of the components of w associated to the edges incident to v, adding twice the value associated to loops. Once more, one can check that  $w'\mapsto (w,z)$  provides the inverse injective function.

**Remark 29.** From the previous proof, one can derive that for every nonnegative odd integer t:

$$\operatorname{vol}_{t}(G_{h,k},\emptyset) - \operatorname{vol}_{t}(G_{h,k},\ell) = |Z_{0}^{\ell \operatorname{even}}| + \dots + |Z_{t}^{\ell \operatorname{even}}| - |Z_{0}^{\ell \operatorname{odd}}| - \dots - |Z_{t}^{\ell \operatorname{odd}}|$$

$$= (|Z_{0}^{\ell \operatorname{even}}| - |Z_{1}^{\ell \operatorname{odd}}|) + (|Z_{1}^{\ell \operatorname{even}}| - |Z_{0}^{\ell \operatorname{odd}}|) + \dots$$

$$+ (|Z_{t-1}^{\ell \operatorname{even}}| - |Z_{t}^{\ell \operatorname{odd}}|) + (|Z_{t}^{\ell \operatorname{even}}| - |Z_{t-1}^{\ell \operatorname{odd}}|)$$

$$= 0.$$

This equality also follows from the more generic lemma of Liu and Osserman [12, Lemma 3.3].  $\Box$ 

It remains to prove Lemma 28 of the evaluation. Consider an (h, k)-caterpillar  $G_{h,k}$  with  $h \geq 1$ . Let r be a leaf of  $G_{h,k}$ , let a be the edge incident to r, let q be the other end node of a. There are two other edges, not necessarily distinct, b and c incident to q. If w is a point in  $\operatorname{Out}_a(t\mathcal{P}_{G_{h,k}})$  then  $w_a + w_b + w_c = t$  or  $w_a = w_b + w_c$ . Therefore,  $w_a = \min\{t - w_b - w_c, w_b + w_c\}$ . Similarly, if w is a point in  $\operatorname{In}_a(t\mathcal{P}_{G_{h,k}})$ , then  $w_b = w_a + w_c$  or  $w_c = w_a + w_b$ . Thus,  $w_a = \max\{w_b - w_c, w_c - w_b\}$ . This digression is summarized in the following lemma.

**Lemma 30.** If a is a leaf-edge of  $G_{h,k}$  and b and c are two other edges incident to a, not necessarily distinct, then  $w_a = \min\{t - w_b - w_c, w_b + w_c\}$  for every  $w \in \operatorname{Out}_a(t\mathcal{P}_{G_{h,k}})$  and  $w_a = \max\{w_b - w_c, w_c - w_b\}$  for every  $w \in \operatorname{In}_a(t\mathcal{P}_{G_{h,k}})$ .

It is a consequence of Lemma 18 and Lemma 27 that the value of

$$\operatorname{vol}_{a \operatorname{even}}(t\mathcal{P}_{G_{h,k}}) - \operatorname{vol}_{a \operatorname{odd}}(t\mathcal{P}_{G_{h,k}})$$

does not depend on the chosen leaf-edge a; it depends only on the integers t, h and k. Thus, for conciseness, we may define

$$\Delta_{h,k}(t) := \operatorname{vol}_{a \operatorname{even}}(t\mathcal{P}_{G_{h,k}}) - \operatorname{vol}_{a \operatorname{odd}}(t\mathcal{P}_{G_{h,k}}),$$

where a is any leaf-edge of  $G_{h,k}$ . We now state and prove a lemma that supplies the inductive step used in the validation of Lemma 28.

**Lemma 31** (the inductive step). Let h and k be integers such that  $h \ge 1$ ,  $k \ge 0$ , and  $h + k \ge 3$ . For every nonnegative even integer t, we have that

$$\Delta_{h,k}(t) = \begin{cases} \Delta_{1,1}(t) \times \Delta_{1,k-1}(t) & \text{if } h = 1, \\ \Delta_{2,0}(t) \times \Delta_{h-1,k}(t) & \text{if } h > 1. \end{cases}$$

Proof. Let a be a leaf-edge of  $G_{h,k}$ . Let q be the other end vertex of a. Because  $h+k \geq 3$ , there are two other distinct edges b and c of  $G_{h,k}$  incident to q. We may assume that a is a stem of  $G_{h,k}$  and b is either a leaf-edge of  $G_{h,k}$  or is incident to a loop (Figure 13(a)). We shall conveniently delineate two edge-disjoint subgraphs of  $G_{h,k}$  having the vertex q as root. The first is the subgraph B that has q as its root and b as its stem. The second is the subgraph C having q as its root and c as a stem. If b is a leaf-edge of  $G_{h,k}$ , then  $B = G_{2,0}$  and  $C = G_{h-1,k}$ , otherwise  $B = G_{1,1}$  and  $C = G_{h,k-1}$  (Figure 13(b)). If we define

$$\Delta_B(t) := \operatorname{vol}_{b \operatorname{even}}(t\mathcal{P}_B) - \operatorname{vol}_{b \operatorname{odd}}(t\mathcal{P}_B) \text{ and } \Delta_C(t) := \operatorname{vol}_{c \operatorname{even}}(t\mathcal{P}_C) - \operatorname{vol}_{c \operatorname{odd}}(t\mathcal{P}_C),$$

then the assertion of this lemma translates into the equality  $\Delta_{h,k}(t) = \Delta_B(t) \times \Delta_C(t)$ .

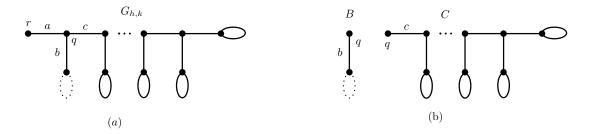


Figure 13: (a)  $G_{h,k}$  with stem a, edges b and c, and node q. The dotted loop may not exist. (b) The graph B with stem b and root q and the graph C with stem c and root q.

In order to prove the lemma it suffices to show that

$$|\operatorname{Out}_{a}(t\mathcal{P}_{G_{h,k}})| = \operatorname{vol}_{b \operatorname{even}}(t\mathcal{P}_{B}) \times \operatorname{vol}_{c \operatorname{even}}(t\mathcal{P}_{C}) + \operatorname{vol}_{b \operatorname{odd}}(t\mathcal{P}_{B}) \times \operatorname{vol}_{c \operatorname{odd}}(t\mathcal{P}_{C}), \text{ and } (37)$$

$$|\operatorname{In}_{a}(t\mathcal{P}_{G_{h,k}})| = \operatorname{vol}_{b \operatorname{even}}(t\mathcal{P}_{B}) \times \operatorname{vol}_{c \operatorname{odd}}(t\mathcal{P}_{C}) + \operatorname{vol}_{b \operatorname{odd}}(t\mathcal{P}_{B}) \times \operatorname{vol}_{c \operatorname{even}}(t\mathcal{P}_{C}). \tag{38}$$

Indeed, by setting  $\mathcal{X} = t\mathcal{P}_{G_{h,k}}$  and e = a in Lemma 26 of the bijection by shifting, we attain as a consequence that

$$\begin{split} \Delta_{h,k}(t) &= |\operatorname{Out}_a(t\mathcal{P}_{G_{h,k}})| - |\operatorname{In}_a(t\mathcal{P}_{G_{h,k}})| \\ &= \operatorname{vol}_{b \operatorname{even}}(t\mathcal{P}_B) \times \operatorname{vol}_{c \operatorname{even}}(t\mathcal{P}_C) + \operatorname{vol}_{b \operatorname{odd}}(t\mathcal{P}_B) \times \operatorname{vol}_{c \operatorname{odd}}(t\mathcal{P}_C) \\ &- \operatorname{vol}_{b \operatorname{even}}(t\mathcal{P}_B) \times \operatorname{vol}_{c \operatorname{odd}}(t\mathcal{P}_C) - \operatorname{vol}_{b \operatorname{odd}}(t\mathcal{P}_B) \times \operatorname{vol}_{c \operatorname{even}}(t\mathcal{P}_C) \\ &= (\operatorname{vol}_{b \operatorname{even}}(t\mathcal{P}_B) - \operatorname{vol}_{b \operatorname{odd}}(t\mathcal{P}_B)) \times (\operatorname{vol}_{c \operatorname{even}}(t\mathcal{P}_C) - \operatorname{vol}_{c \operatorname{odd}}(t\mathcal{P}_C)) \\ &= \Delta_B(t) \times \Delta_C(t), \end{split}$$

where the first equality is by Lemma 26, the second equality is due to (37) and (38) and the last equality follows by definition.

We start by showing Equality (37). For this we give maps between  $\operatorname{Out}_a(t\mathcal{P}_{G_{h,k}})$  and the collection of pairs  $(w^B, w^C)$  of integer points with  $w^B$  in  $t\mathcal{P}_B$  and  $w^C$  in  $t\mathcal{P}_C$  such that  $w^B_b + w^C_c$  is even, indicating that  $w^B_b$  and  $w^C_c$  are both even or both odd. By Lemma 30, for every w in  $\operatorname{Out}_a(t\mathcal{P}_{G_{h,k}})$  we have that  $w_a = \min\{t - w_b - w_c, w_b + w_c\}$ , implying that  $w_b + w_c$  is even because t and  $w_a$  are even. Thus, the restriction of each w in  $\operatorname{Out}_a(t\mathcal{P}_{G_{h,k}})$  to  $t\mathcal{P}_B$  and to  $t\mathcal{P}_C$  produces a unique pair  $(w^B, w^C)$  of integer points such that  $w^B_b + w^C_c$  is even. Conversely, given such a pair  $(w^B, w^C)$  of integers point with  $w^B_b + w^C_c$  even, we construct an integer point w by setting  $w_e = w^B_e$  for  $e \in E(B)$ , setting  $w_e = w^C_e$  for  $e \in E(C)$ , and setting  $w_a = \min\{t - w_b - w_c, w_b + w_c\}$ . Clearly w satisfies the system of inequalities that defines  $\mathcal{P}_{G_{h,k}}$  associated to each internal node v in  $G_{h,k}$  with  $v \neq q$ . Let us now examine the inequalities associated to the vertex q. Firstly, the value of  $w_a$  was set so that  $w_a + w_b + w_c \le t$  and  $w_a \le w_b + w_c$ . Secondly, because  $w_c = w^C_c \ge 0$  and  $w_b = w^B_b \ge 0$ , if  $w_a = w_b + w_c$  then  $w_b \le w_a + w_c$  and  $w_c \le w_a + w_b$ . Finally, because  $w_b = w^B_b \le t/2$  and  $w_c = w^C_c \le t/2$ , if  $w_a = t - w_b - w_c$ , then  $w_a + w_b = t - w_b - w_c + w_b = t - w_c \ge w_c$  and  $w_a + w_c = t - w_b - w_c + w_c = t - w_b \ge w_b$ . This concludes the proof of Equality (37).

We verify the validity of Equality (38) in the same fashion. We mount maps between  $\text{In}_a(t\mathcal{P}_{G_{h,k}})$  and the collection of pairs  $(w^B, w^C)$  of integer points with  $w^B$  in  $t\mathcal{P}_B$  and  $w^C$  in  $t\mathcal{P}_C$  such that  $w_b^B + w_c^C$  is odd, indicating that  $w_b^B$  and  $w_c^C$  have different parity. If w is in  $\text{In}_a(t\mathcal{P}_{G_{h,k}})$ , then  $w_a$  is odd and, by Lemma 30,  $w_a = \max\{w_b - w_c, w_c - w_b\}$ , implying that  $w_b + w_c$  is odd. The restriction of w to  $t\mathcal{P}_B$  and to  $t\mathcal{P}_C$  maps w to a pair  $(w^B, w^C)$  of integer points such that  $w_b^B + w_c^C$  is odd. This gives us one of the maps we want. The other map is the extention of such a pair  $(w^B, w^C)$  of integer points with  $w_b^B + w_c^C$  odd to the integer point w such that  $w_a = \max\{w_b^B - w_c^C, w_c^C - w_b^B\}$ . This gives us the converse map. It is clear that w satisfies the inequalities associated to each internal node v of  $G_{h,k}$  with  $v \neq q$ . By symmetry we may assume that  $w_a = w_b - w_c = w_b^B - w_c^C$ . Because  $0 \le w_b^B \le t/2$  and  $0 \le w_c^C \le t/2$ , then  $0 \le w_a \le t/2$ ,  $w_a \le w_b + w_c$ ,  $w_b = w_a + w_c$ ,  $w_c = w_b - w_a \le w_a + w_b$  and  $w_a + w_b + w_c = w_b^B - w_c^C + w_b^B + w_c^C = 2w_b^B \le t$ . This concludes the proof of the lemma.

We have now all the tools to prove Lemma 28 of the evaluation and close the proof of Theorem 5 of the period of  $L_G^{\mathcal{P}}(t)$ .

Proof of Lemma 28 (the evaluation). Using the abbreviation  $\Delta_{h,k}(t) = \text{vol}_{a \text{ even}}(t\mathcal{P}_{G_{h,k}}) - \text{vol}_{a \text{ odd}}(t\mathcal{P}_{G_{h,k}})$ , we have to prove that

$$\Delta_{h,k}(t) = \begin{cases} (\frac{t}{2} + 1)^k & \text{if } t \equiv 0 \pmod{4} \text{ or } h = 1, \\ 0 & \text{if } t \equiv 2 \pmod{4}. \end{cases}$$

The proof is by induction on h+k. The base case is h+k=2 and we have two possibilities, either (h,k)=(2,0) or (h,k)=(1,1). Let us first consider the case (h,k)=(2,0). The system that defines  $t\mathcal{P}_{G_{2,0}}$  is simply

$$0 \le w_a \le t/2,$$

where a is the unique edge of  $G_{2,0}$ . The integer points in  $t\mathcal{P}_{G_{2,0}}$  are  $w_a = 0, 1, \ldots, t/2$ . If  $t \equiv 0 \pmod{4}$  then t/2 is even and  $\Delta_{2,0}(t) = 1$ . If  $t \equiv 2 \pmod{4}$  then t/2 is odd and  $\Delta_{2,0}(t) = 0$ . Therefore,

$$\Delta_{2,0}(t) = \begin{cases} 1 & \text{if } t \equiv 0 \pmod{4}, \\ 0 & \text{if } t \equiv 2 \pmod{4}. \end{cases}$$
(39)

Now we consider the case (h, k) = (1, 1). The system that defines  $t\mathcal{P}_{G_{1,1}}$  is simply

$$w_a \leq t - 2w_\ell$$

$$w_a \leq 2w_\ell$$

$$w_a \geq 0,$$

where a is the leaf-edge and  $\ell$  is the loop of  $G_{1,1}$ . By Lemma 30, we have that  $w_a = 0$  for every w in  $\operatorname{In}_a(t\mathcal{P}_{G_{1,1}})$  and that  $w_a = \min\{t - 2w_\ell, 2w_\ell\}$  for every w in  $\operatorname{Out}_a(t\mathcal{P}_{G_{1,1}})$ . Thus,  $\operatorname{In}_a(t\mathcal{P}_{G_{1,1}}) = \emptyset$  and, for each  $w_\ell$  in  $\{0, \ldots, t/2\}$ , there is a unique point w in  $\operatorname{Out}_a(t\mathcal{P}_{G_{1,1}})$ . Therefore, by setting  $\mathcal{X} = t\mathcal{P}_{G_{1,1}}$  and e = a in Lemma 26 of the bijection by shifting, we attain as consequence that for every nonnegative even integer t

$$\Delta_{1,1}(t) = |\text{Out}_a(t\mathcal{P}_{G_{h,k}})| - |\text{In}_a(t\mathcal{P}_{G_{h,k}})| = \frac{t}{2} + 1.$$
(40)

Now, we may assume that  $h+k\geq 3$ . If h=1, then for every nonnegative even integer t

$$\Delta_{h,k}(t) = \Delta_{1,1}(t) \times \Delta_{1,k-1}(t) = \left(\frac{t}{2} + 1\right) \times \left(\frac{t}{2} + 1\right)^{k-1} = \left(\frac{t}{2} + 1\right)^k,$$

where the first equality follows from Lemma 31 and the second equality is due to (40) and the induction hypothesis.

If h > 1, then for every nonnegative even integer t

$$\Delta_{h,k}(t) = \Delta_{2,0}(t) \times \Delta_{h-1,k}(t) = \begin{cases} 1 \times \left(\frac{t}{2} + 1\right)^k = \left(\frac{t}{2} + 1\right)^k & \text{if } t \equiv 0 \pmod{4}, \\ 0 \times \Delta_{h-1,k}(t) = 0 & \text{if } t \equiv 2 \pmod{4}, \end{cases}$$

where the first equality follows from Lemma 31 and the second equality is due to (39) and the induction hypothesis. With this we finalize the proof of this lemma.

## 7 Further research directions

In this section, we present various notable results on the polytopes of Liu and Osserman, obtained in the course of our investigations, that seem interesting on their own. We also discuss a surprising relation between these polytopes and invariants used to distinguish 3-manifolds as well as an unexpected connection with arrangements of curves.

#### 7.1 The coset-based technique

It would be very interesting if the the coset-based method explored by Liu and Osserman [12], and futher developed in this paper, could be applied to derive some information about the Ehrhart quasi-polynomials for other families of polytopes. In the coset-based method, the integer points of a polytope are embedded into the collection of intersections of cosets of a lattice with another polytope, where we already have some knowledge about its Ehrhart quasi-polynomial. We note that there is a very high degree of symmetry in the polytopes  $\mathcal{P}_G$  and  $\mathcal{Q}_G$ . Families of polytopes with a certain degree of symmetry could be natural candidates for possibly applying the coset-based technique with some success.

A fascinating collection of polytopes are those integer polytopes that tile  $\mathbb{R}^d$  by translations with some lattice  $\mathcal{L}$ . Indeed such an approach seems promising for these tiling polytopes due to their symmetry and the recent results of Lev and Liu [11], as follows.

**Theorem 32** ([11, Proposition 1.3]). Let  $A, B \subset \mathbb{R}^d$  be two d-dimensional polytopes, and let  $\mathcal{L}$  be a d-dimensional lattice. Suppose A tiles  $\mathbb{R}^d$  by translations with the lattice  $\mathcal{L}$ . Then B also tiles  $\mathbb{R}^d$  by translations with the lattice  $\mathcal{L}$  if and only if A and B are equidecomposable with respect to vectors from  $\mathcal{L}$ .

If a rational polytope A tiles  $\mathbb{R}^d$  by translations with a lattice  $\mathcal{L} \subset \mathbb{Z}^d$ , then the cosets of  $\mathcal{L}$  in  $\mathbb{Z}^d$  are a natural candidate for the coset method we explored here. In particular, if B is a rational cube that has the same volume as A, then we might derive information about the Ehrhart quasi-polynomial of A by using the "trivial" Ehrhart quasi-polynomial of B.

When attempting to obtain some knowledge about the Ehrhart polynomial of a tiling polytope by using the latter theorem, one complication that may arise is that partially-closed polytopes might be required, due to the non-overlapping constraint inherent in such a geometric reconstruction. Nevertheless, perhaps such an approach for the family of tiling polytopes may be feasible.

In closing, we also mention the family of multi-tiling polytopes [11], which naturally extend the family of tiling polytopes, are necessarily symmetric and their facets are also symmetric [9].

## 7.2 Skeleton of the $\{1,3\}$ -trees polytopes

The next result was previously observed by Liu and Osserman [12, proof of Corollary 3.6].

**Lemma 33.** If T is a  $\{1,3\}$ -tree, then the polytope  $\mathcal{P}_T$  is full-dimensional.

Proof. If n is the number of degree 3 nodes in T, then we must prove that  $\mathcal{P}_T$  has dimension 2n+1. Let E be the set of edges of T. For an edge  $e \in E$ , let  $w^e = \frac{1}{3}(\mathbb{1}_T - \mathbb{1}_e)$ . It is easy to check that  $w^e \in \mathcal{P}_T$  for every e. Moreover, the set  $\{w^e : e \in E\}$  together with the origin forms a set of affinely independent vectors with 2n+2 vectors.

Here is a consequence of Theorem 11.

Corollary 34. For every  $\{1,3\}$ -tree T, there is a bijection between vertices of  $\mathcal{P}_T$  and subsets of the leaves of T with an even number of leaves. Thus, for every  $\{1,3\}$ -tree T with m edges,  $\mathcal{P}_T$  has  $2^{\frac{m+1}{2}}$  vertices.

*Proof.* Let S denote an arbitrary subset of the leaves of T such that |S| is even. As T is a tree, each such S corresponds to exactly one collection of |S|/2 disjoint leaf-paths in T whose ends are exactly the leaves in S. The converse is also true: to each collection H of disjoint leaf-paths in T, we can associate the set of ends of the paths in H, and this set contains only leaves, and clearly an even number of them.

Recall that a  $\{1,3\}$ -tree on m edges has  $\ell = \frac{m+3}{2}$  leaves. If  $\ell$  is even, then the number of subsets S of the leaves with |S| even is half of the total number of sets of leaves in T, that is, half of  $2^{\ell}$ . If  $\ell$  is odd, then the number of subsets S of the leaves with |S| even is also half of the total number of sets of leaves in T. Indeed, it is the sum of  $\binom{\ell}{2i}$  for  $i=0,1,\ldots,(\ell-1)/2$ , which is equal to the sum of  $\binom{\ell-1}{i}$  for  $i=0,\ldots,\ell-1$ , that is  $2^{\ell-1}=2^{\frac{m+1}{2}}$ .

Corollary 34 implies that there is a nontrivial bijection between vertices of the polytopes of different  $\{1,3\}$ -trees with the same number of edges.

The symmetric difference between sets A and B is the set  $A \triangle B = (A \cup B) \setminus (A \cap B)$ . For a vertex w of  $\mathcal{P}_T$ , let  $H_w$  denote the collection of disjoint leaf-paths such that  $w = \frac{1}{2} \mathbb{1}_{H_w}$ .

**Theorem 35.** Let w and w' be two distinct vertices of  $\mathcal{P}_T$ . Then w and w' are adjacent in the 1-skeleton of  $\mathcal{P}_T$  if and only if  $H_w \triangle H_{w'}$  is a leaf-path.

*Proof.* Note that  $H_w \triangle H_{w'}$  is a disjoint collection  $\{P_1, \ldots, P_k\}$  of leaf-paths, with  $k \ge 1$ . If k = 1, then let L denote the set of edges incident to leaves of T. Let  $L_w$  denote the edges of  $H_w$  in L and  $L_{w'}$  denote the edges of  $H_{w'}$  in L. Because k = 1, we have that

$$|L_w \triangle L_{w'}| = 2. (41)$$

The hyperplane  $h(x): \sum_{e \in L_w \cap L_{w'}} 2 x_e + \sum_{e \in L \setminus (L_w \cup L_{w'})} (1 - 2 x_e) = |L| - 2$  is a supporting hyperplane of  $\mathcal{P}_T$ , with w and w' being the only vertices of  $\mathcal{P}_T$  in this hyperplane. Indeed, for every  $x \in \mathcal{P}_T$ ,

$$\sum_{e \in L_w \cap L_{w'}} 2 x_e + \sum_{e \in L \setminus (L_w \cup L_{w'})} (1 - 2 x_e) \le |L_w \cap L_{w'}| + |L \setminus (L_w \cup L_{w'})|$$

$$= |L \setminus (L_w \triangle L_{w'})| = |L| - 2,$$
(42)

where the last equality is due to (41). Also, if x is a vertex of  $\mathcal{P}_T$ , then inequality (42) is tight if and only if the set  $L_x$  of leaves of  $H_x$  is such that  $L_w \cap L_{w'} \subseteq L_x \subseteq L_w \cup L_{w'}$ . There are exactly only two different such sets  $L_x$  with  $|L_x|$  even, namely,  $L_w$  and  $L_w'$ . Therefore, by Corollary 34, inequality (42) is tight only for vertices w and w', implying that w and w' are adjacent in the 1-skeleton of  $\mathcal{P}_t$ .

If k > 1, then we will show that the middle point m of the segment [w, w'] is a convex combination of other two vertices of  $\mathcal{P}_T$ , and therefore w and w' are not adjacent.

Let  $u = \mathbb{1}_{H_w \triangle P_1}$  and  $v = \mathbb{1}_{H_w \triangle (P_2 \cup \cdots \cup P_k)}$ . Since  $H_w \triangle P_1$  and  $H_w \triangle (P_2 \cup \cdots \cup P_k)$  are disjoint collections of leaf-paths, u and v are vertices by Theorem 11. Also, u and v are distinct from w and w', as k > 1. Now it is enough to note that  $m = \frac{w+w'}{2} = \frac{u+v}{2}$ .

Corollary 36. Let T be a  $\{1,3\}$ -tree with  $\ell$  leaves. Then the degree of each vertex of the polytope  $\mathcal{P}_T$  is  $\binom{\ell}{2}$ .

*Proof.* From Theorem 35, every vertex of  $\mathcal{P}_T$  has a neighbour for each leaf-path in T.  $\square$ 

**Question 37.** Is there a (combinatorial) characterization of the vertices and edges of the 1-skeleton of  $\mathcal{P}_G$  for an arbitrary  $\{1,3\}$ -graph G?

## 7.3 Symmetry of the $\{1,3\}$ -trees polytopes

We propose a family of involutive isometries of  $\mathcal{P}_T$  showing its high degree of symmetry.

**Theorem 38.** Let T be a  $\{1,3\}$ -tree with n degree 3 nodes and let H be a disjoint collection of leaf-paths in T. We define the function  $h_H: \mathbb{R}^{2n+1} \longrightarrow \mathbb{R}^{2n+1}$  by

$$h_H(w) = \begin{cases} \frac{1}{2} - w_e & if \ e \in E(H), \\ w_e & otherwise. \end{cases}$$

Then  $h_H$  is an isometry of  $\mathcal{P}_T$  to itself.

*Proof.* We notice that  $h_H$  is an involution, and that

$$h_H(w) = w \cdot B + \frac{1}{2} \mathbb{1}_H$$

where B is the  $((2n+1)\times(2n+1))$ -matrix where the entry  $(e,e)_{1\leq e\leq 2n+1}$  equals 1 (resp. -1) if  $e \notin E(H)$  (resp.  $e \in E(H)$ ) and zero elsewhere. It can be checked that  $\det(B) = \pm 1$  and also that  $B \cdot B^t = B^t \cdot B = I$  with  $B^t$  the transpose of B and I the identity. Therefore, B is a rotation matrix and thus an isometry. The  $\frac{1}{2}\mathbb{1}_H$  translation does not affect the isometry, so  $h_H$  is indeed an isometry.

Moreover, if v is a vertex of  $\mathcal{P}_T$  then  $h_H(v)$  is a vertex of  $h_H(\mathcal{P}_T) = \mathcal{P}_T$ . Indeed, if v is a vertex of  $\mathcal{P}_T$  then there is a collection of disjoint leaf-paths  $H_v$  such that  $v = \frac{1}{2} \mathbb{1}_{H_v}$ . We notice that

$$h_H(v) = \frac{1}{2} \mathbb{1}_{H \triangle H_v}$$

and, because  $H \triangle H_v$  is also a collection of disjoint leaf-paths,  $h_H(v)$  correspond to a vertex of  $\mathcal{P}_T$ . Hence  $h_H$  is an isometry of  $\mathcal{P}_T$  to itself.

**Remark 39.** Each function  $h_H$  is an involution, and it can be therefore thought of as a particular even permutation on the set of vertices of  $\mathcal{P}_T$ .

By Theorem 11, we clearly have that  $2\mathcal{P}_T$  is a 0/1 polytope. In view of the above combinatorial properties, it might be reasonable to consider  $2\mathcal{P}_T$  as a good candidate to study different questions in connection with 0/1 polytopes. For instance, a very basic (but difficult) problem is to count the minimal number of simplices needed to triangulate the d-dimensional cube. The following question is in the same spirit.

**Question 40.** Let T be a (1,3)-tree. What is the smallest number of simplices needed to triangulate  $2\mathcal{P}_T$ ?

#### 7.4 Graphs with the same degree sequence

Liu and Osserman [12, Remark 3.11] observed that, if G and H are two connected  $\{1, 3\}$ -graphs on n nodes and m edges, then G and H have the same number of internally Eulerian subgraphs. In other words, for connected  $\{1, 3\}$ -graphs, the number  $N_G$  depends only on the number of nodes and edges in G. Specifically, as we have mentioned just after Example 14, if k = m - n + 1 is the cyclomatic number of G and G is the number of leaves in G, then G if G and G is the number of leaves in G, then G if G and G if G and G if G if G and G if G i

We observe that, if G and H are two connected graphs on n nodes and the same degree sequence, then G and H have the same number of internally Eulerian subgraphs, that is,  $N_G = N_H$ . Indeed, a previous result [6, Theorem 1] states that G can be transformed into H by a series of NNIs. Specifically, an NNI move preserves the degree sequence of the graph. The discussion that precedes Lemma 15 establishes for connected  $\{1,3\}$ -graphs that an NNI move preserves (internally) Eulerian subgraphs; but NNI moves also preserve (internally) Eulerian subgraphs in connected graphs with the same degree sequence. Thus the NNIs transforming G into H naturally induce a bijection between (internally) Eulerian subgraphs of G and H. From the proof of [6, Theorem 1], we may strengthen Liu and Osserman's remark as follows.

**Remark 41.** Let G be a connected graph with n nodes and m edges. Let h be the number of leaves of G and k = m - n + 1 be the cyclomatic number of G. Then the number  $N_G$  of internally Eulerian subgraphs of G is  $2^k$  if h = 0 and  $2^{k+h-1}$  if h > 0.

This leads also to a purely combinatorial proof (instead of a linear algebraic approach) that the number of Eulerian subgraphs of a connected graph with n nodes and m edges is  $2^k$  where k = m - n + 1 is the cyclomatic number of the graph [4, Theorem 1.9.6].

## 7.5 Nonintersecting closed curves

A 3-regular hypergraph is a pair H = (V, E) where V is the set of nodes of H and each element of E is a hyperedge, and consists of exactly three elements of V. Let a, b, and c be the three nodes in a hyperedge  $e \in E$ . Let  $w_a$ ,  $w_b$ , and  $w_c$  be variables satisfying the following system of linear inequalities, which we refer to as  $S_t^H(e)$ :

$$w_a < w_b + w_c$$

$$w_b \leq w_a + w_c$$

$$w_c \leq w_a + w_b$$

$$w_a + w_b + w_c \leq t. \tag{43}$$

Let  $S_t^H$  be the union of all the linear systems  $S_t^H(e)$  taken over all hyperedges e of H and let  $\mathcal{P}_H$  be the set consisting of all real solutions to this linear system when t = 1. Because of the constraints (43),  $\mathcal{P}_H$  turns out to be a polytope.

Given a cubic graph G, we can naturally associate a 3-regular hypergraph  $H_G$  having as nodes the set of edges of G and each hyperedge is given by the edges incident to a node of G. In this case, we have that  $\mathcal{P}_{H_G} = \mathcal{P}_G$ .

For each hyperedge e of a hypergraph H, let us consider an auxiliary variable  $z_e$  and substitute constraint (43) in each system  $S_t^H(e)$  by the parity constraint:

$$w_a + w_b + w_c = 2 z_e$$
$$z_e \le t.$$

Let  $\bar{S}_t^H$  be the union of all these modified linear systems, taken over all hyperedges e of H. The polytope  $\mathcal{Q}_H$  consists of all real solutions to this linear system when t=1. We notice that  $\mathcal{Q}_{H_G} = \mathcal{Q}_G$  for every cubic graph G.

Let  $\mathcal{T}$  be a triangulation of a 3-manifold. Let  $H_{\mathcal{T}} = (V, E)$  be the hypergraph having as set V of nodes the edges of  $\mathcal{T}$  and the set E of hyperedges are the 3-sets corresponding to the triangles which are faces of the tetrahedra used in  $\mathcal{T}$ . We notice that a node of  $H_{\mathcal{T}}$  (that is, an edge of  $\mathcal{T}$ ) could belong to more than two hyperedges (that is, the corresponding edge is shared by two or more tetrahedra in  $\mathcal{T}$ ). Maria and Spreer [13] studied the notion of admissible colourings of the edges of  $\mathcal{T}$  with r-1 colours, that correspond to integer solutions of the linear system  $\bar{S}_t^{H_{\mathcal{T}}}$  for t=r-2, and thus correspond to integer points in the dilated polytope  $t\mathcal{Q}_{H_{\mathcal{T}}}$ . They interpret each admissible colouring as a surface embedded in the triangulated 3-manifold and use this to derive better algorithms to compute Turaev-Viro invariants of degree 4 for the 3-manifold. Note that Maria and Spreer [13, Section 2.3] defined a reduction of an admissible colouring that is an admissible colouring with two colours, that is, r=3, and they decomposed the invariants according to these reduced colourings. This plays very much the same role as the cosets we used in our results.

In the same spirit, we mimic the above construction of Maria and Spreer [13] for a triangulation T of the 2-sphere. We do so by taking T as a graph embedded in the plane and by considering its dual graph  $T^*$ . Note that  $T^*$  is a planar cubic graph. We shall see that the integer points in the dilated polytope  $t\mathcal{Q}_{T^*}$  have an intriguing geometric interpretation in terms of arrangements of pseudocircles.

A pseudocircle is a non-self-intersecting continuous closed curve in the plane. A T-arrangement of pseudocircles is a (possibly empty) set of nonintersecting pseudocircles C on the plane such that (i) C intersects T transversally in the interior of edges (not touching nodes) and (ii) when C enters into a facial triangle through an edge, it leaves the triangle through a different edge. The order of a T-arrangement is the maximum number of times a facial triangle of T is traversed by pseudocircles in the arrangement.

It turns out that each integer point in the rational polytope  $tQ_{T^*}$  induces a T-arrangement of pseudocircles of order at most t and vice-versa. Indeed, for each such integer point in  $tQ_{T^*}$ , we can construct systems of arcs in each facial triangle of T inducing such a T-arrangement of pseudocircles. To show this correspondence, we may proceed as follows.

Consider a facial triangle of T formed by edges  $\{a, b, c\}$  and let  $w'_a$ ,  $w'_b$ , and  $w'_c$  be the values of  $w_a$ ,  $w_b$ , and  $w_c$  in the solution  $\mathbf{w}$ , respectively. We may assume without loss of generality that  $w'_a \geq w'_b \geq w'_c$ . Draw  $w'_a$  points along a, and similarly for b and c. Recall that  $w'_a \leq w'_b + w'_c$ . Let  $x_{ab}$  be the common node of a and b. If  $w'_a = w'_b + w'_c$  then draw arcs joining the  $w'_b$  points in a closer to  $x_{ab}$  to the points in b, and draw arcs joining the remaining  $w'_c$  points in a to the points in c (Figure 14(a)). If  $w'_a < w'_b + w'_c$  then let  $x_{bc}$  be the common node of b and c. First draw arcs joining the  $(w'_b + w'_c - w'_a)/2$  points in c closer to c to the  $(w'_b + w'_c - w'_a)/2$  points in c closer to c to the points in c to the points in c closer to c and draw arcs joining the remaining points in c to the points farther from c and c (Figure 14(b)).

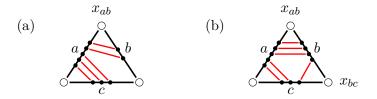


Figure 14: (a)  $w'_a = w'_b + w'_c$  and (b)  $w'_a < w'_b + w'_c$ .

**Example 42.** Consider a triangulation of the 2-sphere whose corresponding graph is a  $K_4$ . Note that the dual  $K_4^*$  of a  $K_4$  is isomorphic to  $K_4$  (Figure 15(a)). It can be checked that  $w_1 = w_4 = w_6 = 2, w_2 = w_3 = 3, w_5 = 1, z_{e_1} = 2, z_{e_2} = z_{e_3} = z_{e_4} = 3$  and  $w_1 = w_2 = w_3 = w_4 = w_5 = w_6 = 2, z_{e_1} = z_{e_2} = z_{e_3} = z_{e_4} = 3$  are two integer points in  $3Q_{K_4}$ . The corresponding induced nonintersecting pseudocircles are illustrated in Figure 15(b).

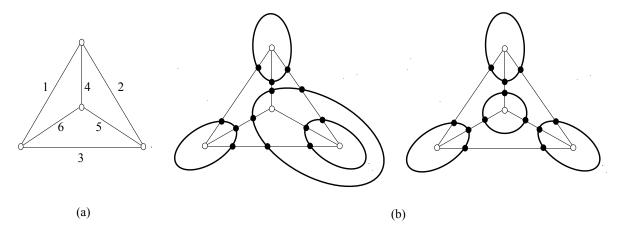


Figure 15: (a)  $K_4$  with labeled edges, representing a triangulation of the plane; (b) Two arrangements of curves of order 3.

As a consequence of the above discussion we conclude that the number  $L_{T^*}^{\mathcal{Q}}(t)$  of integer points in  $t\mathcal{Q}_{T^*}$  is exactly the number of T-arrangements of pseudocircles of order at most t.

**Lemma 43.** For any triangulation T of the 2-sphere, the number of T-arrangements of pseudocircles of order at most t is the number  $L_{T^*}^{\mathcal{Q}}(t)$  of integer points in the rational polytope  $t\mathcal{Q}_{T^*}$ .

We know by Ehrhart theory that the number of points in  $tQ_{T^*}$  grows as a quasi-polynomial in t, so that, by Lemma 43, the number of T-arrangements of pseudocircles of order at most t also grows as a quasi-polynomial in t. But we would like to know whether this quasi-polynomial collapses to a polynomial function of t.

**Question 44.** Let T be a triangulation of the plane and  $t \ge 0$  be an integer. Does the number of T-arrangements of pseudocircles of order t grow polynomially in t?

A positive answer to the previous question would imply that  $Q_{T^*}$  has period 1.

A well-known problem in dimension 1 asks for the number of ways to construct an admissible set of n parentheses for a word of length 2n [18, Problem 6.19(b)]. To put this problem into our context, this 1-dimensional counting problem is equivalent to the following question [18, Problem 6.19(o)]:

What is the number of ways of connecting 2n points lying on a horizontal line by n nonintersecting arcs, each arc connecting two of the points and lying above the points?

The answer to this problem is given by the Catalan numbers. Note that each configuration in Figure 16 naturally induces an arrangement of nonintersecting circles (by closing up arcs with their mirrors) which is closely connected to arrangements of pseudocircles.

Question 45. Would the above information shed light on the understanding of  $L_{Q_{T^*}}$ ?



Figure 16: Five ways to connect 6 points by 3 nonintersecting arcs.

## 8 Acknowledgements

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