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Secure Estimation Using Partially Homomorphic Encryption for Unmanned Aerial Systems in the Presence of Eavesdroppers

Xinhao Yan, Guanzhong Zhou, Yue Huang, Wei Meng, Anh-Tu Nguyen, Hailong Huang

Abstract—Unmanned aerial systems (UASs) are attracting increasing attention thanks to the great mobility and flexibility of unmanned aerial vehicles (UAVs). This paper considers a typical UAS, which consists of a UAV, a sensing device that provides some sensed data to the UAV, and an end-user that operates the UAV. However, the information exchanged between these parties is vulnerable to eavesdropping attacks, emphasizing the need to develop privacy-preserving approaches. The cryptographic methods are undoubtedly effective, but their high computational overhead may adversely impact the normal operations of UASs. Additionally, the dynamic of a UAV has a high dimension, which is disadvantageous for both estimation and encryption. Therefore, this paper proposes a secure distributed estimation protocol with partially homomorphic encryption by encrypting the transmitted measurements and estimates. Attribute to distributed structure and partial homomorphism, the computation amount for secure estimation is greatly reduced. At the same time, the raw data that needs to be encrypted is transferred into the space of plaintexts by a uniform quantizer and a mapping strategy. Finally, the effectiveness of the proposed method is verified by computer simulation.

Index Terms—Unmanned aerial vehicles (UAVs), state estimation, homomorphic encryption, uniform quantization

I. INTRODUCTION

A. Background

Due to the flexibility of unmanned aerial vehicles (UAVs), unmanned aerial systems (UASs) that consist of UAVs and other equipment have found a wide range of applications, including task offloading [1], covert surveillance [2], multiple access [3], and relay communication [4]. As one of the most important issues, the trajectories of UAVs need to be tracked and some state estimation methods are applied to decrease the influence of system noises [5], [6]. Unfortunately, the UASs are facing many safety problems [7], and particularly, there exist many malicious eavesdroppers that want to intercept the information of UASs [8]. For example, they require the precise position of a UAV to launch some vicious attacks, such as spoofing attacks [9] and even physical attacks [10]. A structure of UAS in the presence of eavesdroppers is shown in Fig. 1. In this case, the eavesdroppers have access to the user-UAV channel and sensor-UAV channel, and they can speculate the accurate states of UAVs based on the wiretapped measurements, estimates, and other information. Because of the great harm brought by eavesdropping, it is of great significance to protect the privacy of the transmitted data in UASs.

B. Related Work

Homomorphic encryption is a classical cryptographic approach, which allows certain operations on the encrypted data [20]. It has been widely used for various scenarios, including cloud-based systems [21], Internet of Things (IoT) [22], and image processing [23]. In general, there are two homomorphisms: additive homomorphism and multiplicative homomorphism, which respectively means addition and multiplication are operable on ciphertexts. Then, the homomorphic scheme can be divided into two main kinds: fully
homomorphic encryption [24] and partially homomorphic encryption [25]. The fully homomorphic encryption [24] can simultaneously support both additive and multiplicative homomorphisms, while the partially homomorphic encryption can only support one of them. Paillier [25] is a kind of typically additive homomorphic encryption, RSA [26] and ElGamal [27] are typically multiplicative homomorphic encryption approaches. Note that the fully homomorphic encryption requires a long computation time, which has been proved by the practical experiments in [28]. It shows that the estimation with partially homomorphic encryption only requires several milliseconds, while that with fully homomorphic encryption takes several hours and that with garbled circuits takes several minutes. Hence, partially homomorphic encryption is more practical for real-time estimation. Besides, there also exists a special kind called hybrid homomorphic encryption [29], where two different homomorphic encryption approaches are combined to achieve both multiplicative and additive homomorphisms.

Moreover, it should be pointed out that the data going to be processed by the homomorphic encryption must lie in the space of plaintext. This is because the encryption and decryption functions are based on the modular arithmetic that only deals with integers. This condition is an important preliminary, but it is always ignored. For example, the performance degradation is neglected for theoretical analysis [30]. To solve this problem, the float number is expressed by a positive exponent and a mantissa in [31], and the representations are all integers that can be encrypted. Such encoding method has already been employed to state estimation field [32]. Instead of the detailed representation of the whole float number, quantization is another method to transfer the data into the space of plaintext. It only preserves the integer part and directly deletes the mantissa. A probabilistic uniform quantization is adopted in [29], where the output is chosen randomly in an interval.

Recently, the differentially private estimator was proposed in [33], while the legitimate estimation performance was degraded due to the injection of random noises. The homomorphic encryption-based estimator was studied in [29], but the hybrid homomorphism consumed a great computation amount. Further, the additive homomorphic encryption was applied to protect the estimator in [30], but the mapping process was ignored. Notice that most traditional estimators are centralized [5,6]. In this case, the great computation burden will be large due to the augmentation of all the local components. Moreover, although the confidentiality problem of UAS is always a significant issue, homomorphic encryption is rarely studied for UAS. Except for the normal operations on control and estimation, the extra computation on UAS should also be as little as possible, because real-time performance is important for UAS. Meanwhile, some parameters do not require privacy preservation, because the eavesdroppers cannot speculate the state from them. Therefore, partially homomorphic encryption is considered in this paper, which requires less computation amount than fully and hybrid homomorphic encryption methods.

C. Contributions

According to the above analysis, the main contributions of this paper are summarized as follows.

1) The globally high-dimensional system of a UAV is separated into a translational subsystem and a rotational subsystem, and the two subsystems are respectively estimated by their own low-dimensional estimators. Such a distributed estimation structure significantly decreases the computational complexity of both estimation and encryption.

2) A secure estimation protocol is proposed for the UAS against eavesdroppers, where the privacy of the translational state estimates is protected by a partially homomorphic encryption approach with a uniform quantizer. The partial encryption and time-invariant quantization further reduce the computation amount.

The rest of this paper is organized as follows. Section II discusses the complete dynamics of the UAV. Then, the distributed estimation models are presented by dividing the global system into a translational subsystem and a rotational subsystem. In Section III, the Kalman-like distributed estimator is designed in the minimum variance sense. Next, Paillier homomorphic encryption and uniform quantization methods are introduced to preserve the privacy of UAS, and the complete protocol is proposed. Afterwards, the simulation results are shown in Section IV, including the tracking trajectories, estimation performance, and time cost. Finally, the conclusion is given in Section V. The notations frequently used throughout the paper are summarized in Table I.

### II. Problem Formulation

#### A. Quadrotor UAV Dynamics

The UAS considered in this paper contains a quadcopter UAV with 6 degrees of freedom (6-DoF), where the attitude is depicted by Euler angles [10], [34]–[36]. The schematic model of a 6-DoF quadrotor UAV is shown in Fig. 2. Here, we use $\mathcal{I}$ to denote the inertial frame and $\mathcal{B}$ to denote the body frame, and the time index $t$ is neglected in this subsection for brevity.

---

**Table I: Notations**

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>$E$</td>
<td>define</td>
</tr>
<tr>
<td>$I$</td>
<td>identity matrix with appropriate dimension</td>
</tr>
<tr>
<td>$\mathbb{Z}$</td>
<td>set of integers</td>
</tr>
<tr>
<td>$\mathbb{Z}^{n \times m}$</td>
<td>set of $n \times m$ integer matrix</td>
</tr>
<tr>
<td>$\mathbb{Z}_N$</td>
<td>set of integers modulo $N$</td>
</tr>
<tr>
<td>$\mathbb{Z}_N^*$</td>
<td>set of invertible integers modulo $N$</td>
</tr>
<tr>
<td>$\mathbb{R}$</td>
<td>set of real number</td>
</tr>
<tr>
<td>$\mathbb{R}^m$</td>
<td>set of $m$-dimensional real vectors</td>
</tr>
<tr>
<td>$\mathbb{R}^{n \times m}$</td>
<td>set of $n \times m$ real matrices</td>
</tr>
<tr>
<td>$\mathbb{R}^n$</td>
<td>mathematical expectation</td>
</tr>
<tr>
<td>$\mathbb{x}$</td>
<td>derivative of $x$</td>
</tr>
<tr>
<td>$A^T$</td>
<td>transpose of matrix $A$</td>
</tr>
<tr>
<td>$[x]_i$</td>
<td>$i$-th component of vector $x$</td>
</tr>
<tr>
<td>$[A]_{ij}$</td>
<td>$i$-th row and $j$-th column of matrix $A$</td>
</tr>
<tr>
<td>$\text{col} { \cdot }$</td>
<td>column vector</td>
</tr>
<tr>
<td>$\text{Tr} { \cdot }$</td>
<td>trace of a matrix</td>
</tr>
<tr>
<td>$X &gt; (\leq) 0$</td>
<td>positive-definite (negative-definite)</td>
</tr>
<tr>
<td>$X \geq (\leq) 0$</td>
<td>non-negative definite (non-positive definite)</td>
</tr>
<tr>
<td>$\text{sgn}(\cdot)$</td>
<td>sign function</td>
</tr>
<tr>
<td>$Q(\cdot)$</td>
<td>quantization function</td>
</tr>
<tr>
<td>$x_q$</td>
<td>quantized output of $x$</td>
</tr>
<tr>
<td>$M(\cdot)$</td>
<td>mapping function</td>
</tr>
<tr>
<td>$x_m$</td>
<td>mapped output of $x$</td>
</tr>
<tr>
<td>$\text{Enc}(\cdot)$</td>
<td>encryption function</td>
</tr>
<tr>
<td>$\text{Dec}(\cdot)$</td>
<td>decryption function</td>
</tr>
<tr>
<td>$pk$</td>
<td>public key</td>
</tr>
<tr>
<td>$sk$</td>
<td>private key</td>
</tr>
<tr>
<td>$[x]$</td>
<td>encrypted value of $x$</td>
</tr>
</tbody>
</table>
Note that a 6-DoF quadrotor dynamic model is nonlinear and the global state consists of 12 dimensions. It is a complex and high-dimensional system that will increase the computation burden of estimation. As a result, the UAV dynamic system is separated into translational and rotational subsystems in this paper. First, the translational dynamics of UAV can be described by the Newton equation [34]:

\[ \dot{v}_T = \frac{1}{m} R_{BZ} T_B - g_Z, \]

where \( p_T(t) \triangleq \text{col}\{p_x, p_y, p_z\} \) and \( v_T(t) \triangleq \text{col}\{v_x, v_y, v_z\} \) respectively denotes the 3-dimensional (3D) position vector and velocity vector of UAV in the inertial frame. \( T_B(t) = \text{col}\{0, 0, F\} \) is the thrust force, \( g_Z = \text{col}\{0, 0, g\} \) is the gravitational acceleration in inertial frame, and \( m \) is the mass of UAV. The matrix \( R_{BZ} \) denoting the transformation from the body frame to the inertial frame can be described by [10]:

\[
R_{BZ} = \begin{bmatrix}
    c\theta c\psi & s\theta c\psi & -s\phi c\theta c\psi + s\phi s\theta s\psi & -s\phi s\theta c\psi + s\phi c\theta s\psi \\
    c\phi s\psi & c\phi c\psi & s\phi s\psi & c\phi s\theta c\psi + s\phi c\theta s\psi \\
    s\phi c\psi & c\phi s\psi & s\phi s\theta c\psi + s\phi c\theta s\psi & -s\phi s\theta s\psi + c\phi c\theta s\psi \\
    -s\phi & c\phi & 0 & 0
\end{bmatrix},
\]

where \( s\phi, c\phi, l\phi \) respectively stand for the sine, cosine and tangent of angle \( \phi \). Combining the transformation matrix (2), the scalar form of translational subsystem (1) can be expressed by

\[
\begin{align*}
\dot{v}_x &= \frac{1}{m} \left( c\psi s\phi + s\psi s\theta \right) F \\
\dot{v}_y &= \frac{1}{m} \left( c\phi c\psi - s\phi s\theta \right) F \\
\dot{v}_z &= \frac{1}{m} \left( c\phi s\psi \right) F g
\end{align*}
\]

Second, the rotational dynamics of UAV can be modeled by the Euler equation [34]:

\[ J_B \ddot{\omega}_B = M_B - \omega_B \times J_B \omega_B, \]

where \( M_B = \text{col}\{M_x, M_y, M_z\} \) is the moment, \( \omega_B = \text{col}\{\dot{\phi}, \dot{\theta}, \dot{\psi}\} \) is the angular speed and \( J_B = \begin{bmatrix}
    J_{xx} & 0 & 0 \\
    0 & J_{yy} & 0 \\
    0 & 0 & J_{zz}
\end{bmatrix} \) is the moment of inertia matrix. Here, \( \{\phi, \theta, \psi\} \) stand for the Euler angles in the inertial frame, where \( \phi \) is the roll angle, \( \theta \) is the pitch angle, and \( \psi \) is the yaw angle. \( \{p, q, r\} \) are the corresponding angular velocity in the body frame. Then, the transformation between \( \{\phi, \theta, \psi\} \) and \( \{p, q, r\} \) is given by [10]:

\[
\begin{bmatrix}
    \dot{\phi} \\
    \dot{\theta} \\
    \dot{\psi}
\end{bmatrix} = \begin{bmatrix}
    1 & s\phi q & c\phi q \\
    0 & c\phi & -s\phi & 0 \\
    0 & s\phi & c\phi & 0
\end{bmatrix} \begin{bmatrix}
    p \\
    q \\
    r
\end{bmatrix}.
\]

Based on the small-angle approximation [37], one has \( \phi = 0, \theta = 0, \psi = 0 \). Thus, the above relationship can be simplified as

\[
\begin{bmatrix}
    \dot{\phi} \\
    \dot{\theta} \\
    \dot{\psi}
\end{bmatrix} = \begin{bmatrix}
    1 & 0 & 0 \\
    0 & 1 & 0 \\
    0 & 0 & 1
\end{bmatrix} \begin{bmatrix}
    p \\
    q \\
    r
\end{bmatrix}.
\]

Then, the scalar form of dynamic model in (4) can be expressed by resorting to the approximation (6):

\[
\begin{align*}
\dot{p} &= \frac{1}{J_{xx}} J_{yy} - J_{zz} q r \\
\dot{q} &= \frac{1}{J_{yy}} J_{xx} - J_{zz} p r \\
\dot{r} &= \frac{1}{J_{zz}} J_{xx} + J_{yy} p q
\end{align*}
\]

Moreover, the force \( T_B \) and moments \( M_B \) in above dynamics are generated with 4 rotors [35, 36]:

\[
\begin{align*}
M_x &= k_F (m_1 + m_2 + m_3 + m_4) \\
M_y &= k_F (m_3 - m_4) \\
M_z &= k_F (m_1 - m_2) \\
M_2 &= k_M (m_1 + m_2 - m_3 - m_4)
\end{align*}
\]

where \( k_F \) is the thrust coefficient, \( k_M \) is the drag coefficient, and \( m_j (j = 1, 2, 3, 4) \) are the respective pulse width modulation (PWM) inputs for 4 rotors. \( l \) is the distance from the axis of rotation of the rotors to the center of the UAV. Besides, the PWM inputs can be derived as

\[
\begin{align*}
m_1 &= \frac{F}{4k_F + \frac{M_y}{2k_F} \frac{M_z}{4k_M}} \\
m_2 &= \frac{F}{4k_F - \frac{M_y}{2k_F} \frac{M_z}{4k_M}} \\
m_3 &= \frac{F}{4k_F + \frac{M_y}{2k_F} \frac{M_z}{4k_M}} \\
m_4 &= \frac{F}{4k_F - \frac{M_y}{2k_F} \frac{M_z}{4k_M}}
\end{align*}
\]

B. Problem Formulation

Except for a UAV, the discussed UAS in this paper also consists of a user and some sensors. Here, the user and sensor will send the control and measurement signals to the UAV, respectively. Then, the UAV estimates the real-time states and feeds certain information back to the user. Particularly, the controller in user will send 4-channel signals \( \{F, M_x, M_y, M_z\} \) at each time slot, which is determined by the demand of the user.

In the beginning, we should further modify the aforementioned UAS models. In fact, practical models of UASs are generally in discrete-time domain due to their running on certain computation units, thus the above continuous-time systems should be discretized. According to the above analysis, the global state \( x(t) = \text{col}\{x_1(t), x_2(t)\} \) can be divided into the translational state \( x_1(t) = \text{col}\{p_x(t), v_x(t), p_y(t), v_y(t), p_z(t), v_z(t)\} \) and the rotational state \( x_2(t) = \text{col}\{\phi(t), \theta(t), \psi(t), q(t), r(t)\} \). In this case, the discrete-time model of dynamics (3) and (7) can be derived by using the first-order Runge-Kutta method:

\[
\begin{align*}
x_1(t+1) &= A_1 x_1(t) + f_1(x_2(t), w_1(t)) + w_1(t) \\
x_2(t+1) &= f_2(x_2(t)) + B_2(t) w_2(t) + w_2(t)
\end{align*}
\]

where \( w_1(t) \in \mathbb{R}^6 \) and \( w_2(t) \in \mathbb{R}^6 \) are the system noises coming from the modeling error and other internal disturbances. Here, they are assumed to be mutually independent.
white Gaussian noises (WGNs), which is one of the most widely used methods. Then, their covariances $Q_{w_1}$ and $Q_{w_2}$ are defined as follows:

$$E\{w_i(t_1)w_j^T(t_2)\} = \delta(t_1, t_2)\delta(i, j)Q_{w_i} \quad (i, j = 1, 2),$$

where $\delta(i, j)$ is the indicator function such that $\delta(i, j) = 1$ if $i = j$; otherwise, $\delta(i, j) = 0$. Besides, the concrete values of covariance can be designed by resorting to the experimental experience. The 4-channel control input $u(t) \triangleq \text{col}\{u_1(t), u_2(t)\}$ is composed of the force $u_1(t) = F(t)$ and the moment $u_2(t) = \text{col}\{M_x(t), M_y(t), M_z(t)\}$. Moreover, other state-space matrices and nonlinearities of model (10) are given by

$$A_1 = \begin{bmatrix} 1 & T_0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & T_s & 0 \\ 0 & 0 & 0 & 1 & T_s \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}, \quad B_2(t) = \begin{bmatrix} 0 & 0 & 0 \\ \frac{1}{J_{x_0}} & 0 & 0 \\ 0 & \frac{1}{J_{y_0}} & 0 \\ 0 & 0 & \frac{1}{J_{z_0}} \end{bmatrix},$$

$$f_{12}(x_2(t), u_1(t)) = \begin{bmatrix} 0 & \frac{F(t)T_m}{m}(c_0s_{0}\psi + s_0\psi) \\ \frac{F(t)T_m}{m}(c_0\phi s_{0}\psi - s_0\phi \psi) \\ \frac{F(t)T_m}{m}(c_0\psi - s_0\phi) - g T_s \\ \phi(t) + p(t)T_s \\ \theta(t) + q(t)T_s \\ \psi(t) + r(t)T_s \\ \frac{J_{z_0} - J_{y_0}}{J_{x_0}}q(t)p(t)T_s \end{bmatrix},$$

$$f_2(x_2(t)) = \begin{bmatrix} \phi(t) + p(t)T_s \\ \theta(t) + q(t)T_s \\ \psi(t) + r(t)T_s \\ \frac{J_{z_0} - J_{y_0}}{J_{x_0}}q(t)p(t)T_s \end{bmatrix}$$

where $T_s$ is the sampling period for discretization.

In order to estimate the UAV states, some sensors are always deployed to observe the real dynamics of the UAV. There may be many measurement outputs from several sensors, but we finally consider an augmented measurement in this paper. Such modeling can also be treated as a kind of centralized fusion method, where all the local measurements are gathered by one party. Then, the measurement equation of the sensor can be expressed by

$$\begin{cases} y_1(t) = C_1x_1(t) + v_1(t) \\ y_2(t) = C_2x_2(t) + v_2(t) \end{cases}$$

where $y_1(t)(i = 1, 2)$ represent the measured outputs and $C_i(i = 1, 2)$ represent the measurement matrices. The disturbances $v_i(t)(i = 1, 2)$ are also WGNs that satisfy

$$E\{v_i(t_1)v_j^T(t_2)\} = \delta(t_1, t_2)\delta(i, j)Q_{v_i} \quad (i, j = 1, 2),$$

where $Q_{v_i}(i = 1, 2)$ are the corresponding covariances. Since $w_i(t)$ and $v_j(t)$ are mutually independent, one has

$$E\{w_i(t_1)v_j^T(t_2)\} = 0 \quad (\forall i, j, t_1, t_2).$$

In general, there exist many sensors that are suitable for observing the UAV, such as global position system (GPS), range sensors, and inertial measurement units (IMUs). In the UAS of this paper, the measurements provided by GPS and accelerometer are chosen as an example, which contains the 3D position and 3D angle of the UAV. It means that the measurement matrices are expressed in the following form:

$$C_1 = \begin{bmatrix} C_x & 0 & 0 \\ 0 & C_y & 0 \\ 0 & 0 & C_z \end{bmatrix}, \quad C_2 = \begin{bmatrix} C_p & 0 & 0 \\ 0 & C_q & 0 \\ 0 & 0 & C_r \end{bmatrix},$$

where $C_x = C_y = C_z = [1 \ 0], C_p = C_q = C_r = [1 \ 0]$. The above measurement system represents a general and simple structure, and the performance can be further improved by introducing more high-precision sensors.

In traditional methods, the translational subsystem and the rotational subsystem are estimated together, which is a centralized estimation structure. In this paper, we consider a distributed estimation scheme, where the two subsystems are estimated individually. One advantage of such decoupling is that the computation amount can be reduced. Based on all the measurements $\{y_1(t), \ldots, y_2(t)\}$ at discrete time slot $t$, the recursive estimators are designed based on the Kalman-like filter [38], [39]:

$$\begin{cases} \hat{x}_1(t) = A_1\hat{x}_1(t - 1) + f_{12}(\hat{x}_2(t - 1), u_1(t - 1)) \\ \hat{x}_2(t) = (I - K_1(t)C_1)\hat{x}_1(t) + K_1(t)y_1(t) \end{cases},$$

$$\hat{x}_2(t) = (I - K_2(t)C_2)\hat{x}_2(t) + K_2(t)y_2(t),$$

where $\hat{x}_1(t)$ and $\hat{x}_2(t)$ are the distributed state estimates. $K_1(t)$ and $K_2(t)$ are the estimator gains to be designed.

Notice that there exist many eavesdroppers that can silently overhear the communication channels and then estimate the state system with the wiretapped data. In some cases, they can employ the wiretapped measurements $y(t)$ to estimate the state as what the legitimate system does. On the other hand, they can also directly wiretap the estimates $\hat{x}(t)$. Accordingly, the privacy of all the above-mentioned signals should be preserved. Meanwhile, the ciphertexts may also be cracked and utilized. Here, the detailed definition of the eavesdropper is given as follows.

**Definition 2.1 (Eavesdropping Model):** The eavesdropper has access to both the user-UAV channel and the sensor-UAV channel, and it understands all the meanings of the wiretapped data. When it acquires plaintexts, a state estimator will be employed to calculate the final estimate with the raw data. On the other hand, when the eavesdropper acquires ciphertexts, it will launch chosen plaintext attack (CPA), which means that it will encrypt likely plaintexts with the public key and then test if they are equal to the ciphertexts.

For countering the potential eavesdroppers in UAV’s estimation systems, certain information is required to be encrypted. Since the fundamental operations of the proposed estimators (17) are related to addition and multiplication, the homomorphic encryption method is feasible and suitable for privacy preservation. Then, the main problem to be tackled in this paper is how to apply the homomorphic encryption method to protect the privacy of the state estimates in the considered UAS. Besides, the calculation amount should be reduced as much as possible, because the complexity of cryptographic methods is high on a certain level.

**III. SECURE ESTIMATION PROTOCOL**

**A. Overview**

In this section, the homomorphic encryption-based secure estimation protocol is presented, where the design of the
about Taylor expansion, and the Jacobian matrix
rotational state, \( \Delta \) its estimate (17), the estimation error
According to the dynamics of the system state in (10) and
f
classical method in the minimum variance sense. However,
proposed method, where Paillier homomorphic encryption is
complete protocol is given to show the realization of the
protocol is composed of two main parts: estimation and
encryption. First, based on the decoupled subsystems, we design
the distributed estimators in the minimum variance sense to
accurately track the states of UAV. Second, we present the
definitions and properties of Paillier homomorphic
encryption. Meanwhile, the detailed quantization method is
discussed, which is a preliminary of encryption. Finally, the
complete protocol is given to show the realization of the
proposed method, where Paillier homomorphic encryption is
embedded in the distributed estimator.

B. Estimator Design
To design the estimator gain, Kalman filter [40] is a
classical method in the minimum variance sense. However,
there exist nonlinear items \( f_{12}(x_2(t)) \) and \( f_{2}(x_2(t)) \)
that cannot be directly processed. Thus, by resorting to the main
idea of extended Kalman filter (EKF) [41], the nonlinear
functions are respectively linearized at the values of certain
estimates.

Since the rotational subsystem has its own dynamics and
is not affected by the translational subsystem, we first discuss
the rotational state estimator \( \hat{x}_2(t) \). As what is done in EKF,
we expand the function \( f_2(x_2(t)) \) in the Taylor series about
\( \hat{x}_2(t) \):

\[
f_2(x_2(t)) = f_2(\hat{x}_2(t)) + A_{2,1}(t)\Delta f_2(\hat{x}_2(t)), \quad (18)
\]
where \( \Delta f_2(\hat{x}_2(t)) \) is the estimation error for the
rotational state, \( \Delta f_2(\hat{x}_2(t)) \) represents the high-order terms
about Taylor expansion, and the Jacobian matrix \( A_{2,1}(t) \) is
calculated by

\[
A_{2,1}(t) =
\begin{pmatrix}
1 & T_s & 0 \\
0 & 1 & 0 \\
0 & 0 & 1 \\
0 & J_{x_2,x_2}q(T_s) - J_{x_2,x_2}r(T_s) T_s & 0 \\
0 & J_{x_2,x_2}q(T_s) & 0 \\
0 & 0 & 1 \\
\end{pmatrix}
\]

where \( \Delta f_2(\hat{x}_2(t)) \) is the high-order terms and the Jacobian
matrix \( A_{12,1}(t) = \frac{\partial f_{12}(\hat{x}_2(t))}{\partial x_2(t)} \bigg|_{x_2(t) = \hat{x}_2(t)} \)
is

\[
A_{12,1}(t) =
\begin{pmatrix}
0 & 0 & 0 \\
-\phi_x \phi_y & \phi_x \phi_y & 0 \\
0 & 0 & 0 \\
-\phi_x \phi_y & -\phi_y \phi_x & 0 \\
0 & 0 & 0 \\
-\phi_x \phi_y & 0 & -\phi_y \phi_x \\
0 & 0 & 0 \\
\end{pmatrix}
\]

Then, the estimation error \( \hat{x}_1(t) \) defined by
translational subsystem can be expressed as

\[
\hat{x}_1(t) = \hat{x}_1(t) + \hat{x}_2(t) \]

Define the covariances \( P_{12}(t) \) and \( P_{21}(t) \) for the
translational and rotational subsystems, respectively,

\[
P_{12}(t) = \frac{\partial f_{12}(\hat{x}_2(t))}{\partial x_2(t)} \bigg|_{x_2(t) = \hat{x}_2(t)}
\]

By minimizing the trace of the estimation error covariance
\( P_{21}(t) \) for the estimator gain \( K_1(t) \) is derived as

\[
P_{21}(t) = \frac{\partial f_{21}(\hat{x}_2(t))}{\partial x_2(t)} \bigg|_{x_2(t) = \hat{x}_2(t)}
\]

C. Paillier Homomorphic Encryption
The homomorphic cryptosystem is a kind of system that
allows computation on encrypted data [20]. This means
that the operations on ciphertext can be reflected on the
original plaintext. Since UAV requires fast response, we
consider the partially homomorphic encryption that consumes
less computational resources when compared with fully and
hybrid homomorphic encryption methods. Hence, a classical
kind of partially homomorphic encryption called Paillier
homomorphic encryption [25] is applied to preserve privacy for
UAS. Meanwhile, we assume that the UAS has enough
computation ability to implement such an encryption approach.
The complete procedures related to Paillier homomorphic
encryption are summarized as follows.

1) Key Generation: Given two primes \( p \) and \( q \). Let \( N = pq \)
and the least common multiple (LCM) is \( \lambda = \text{lcm}(p-1, q- \)
1. $g$ satisfies $\text{gcd}(L(q^x \mod N^2), N) = 1$, where $L(x) = (x - 1)/N$ and $\text{gcd}$ stand for the greatest common divisor (GCD). Then, the secret keys are generated as follows:

$$\text{pk} = (N, g), \text{sk} = \lambda,$$

where pk denotes the public key and sk denotes the private key.

2) Quantization and Mapping: In fact, the most common type of data in UAS is float or double, but not all types of numbers can be treated as the input of the Paillier cryptosystem. The data to be encrypted should lie in the plaintext space $M = \mathbb{Z}_N$. Hence, these original data should be quantized and mapped into positive integers. The general form of quantization is denoted as

$$x_q = Q(x),$$

and the mapping procedure is given by [31]:

$$x_m = M(x_q) = \begin{cases} x_q, & \text{if } x_q \geq 0 \\ x_q + N, & \text{if } x_q < 0 \end{cases}$$

In the rest of this paper, $Q(\cdot)$ denotes the quantization function and $M(\cdot)$ denotes the mapping function. The variable with subscript “q” represents the quantized data and that with “m” represents the mapped data.

3) Encryption: In this step, the raw message $m$ will be encrypted into the ciphertext $c$ with public key pk. The notation $[[\cdot]]$ represents the encrypted message and $c \triangleq [[m]]$. Then, the detailed encryption process is expressed by

$$c = \text{Enc}(m, \text{pk}) = g^{m \cdot N} \mod N^2,$$

where $\text{Enc}(\cdot)$ is the encryption function and $r \in \mathbb{Z}_N$ is a non-zero random integer.

4) Decryption: When legitimate users receive the encrypted signals, they can successfully decrypt them. This means that the ciphertext $c$ will be decrypted into the original message $m$ based on the private key sk, and the concrete decryption process is given by

$$m = \text{Dec}(c, \text{sk}) = \frac{L(c^\lambda \mod N^2)}{L(q^x \mod N^2)} \mod N,$$

where $\text{Dec}(\cdot)$ is the decryption function.

5) Retrieval: After decryption, the obtained result becomes

$$x_m = x_{m1} + x_{m2} \text{ or } x_m = x_{m1} \times x_{m2},$$

which is the value after mapping. Thus, the quantized result $x_q = x_{q1} + x_{q2}$ or $x_q = x_{q1} \times x_{q2}$ should be retrieved by the following criterion [29]:

$$x_q = R(x_m) = \begin{cases} x_m \mod N, & \text{if } x_m \mod N \leq \frac{N}{2} \\ x_m \mod N - N, & \text{if } \frac{N}{2} \leq x_m \mod N < N \end{cases}.$$

Based on the above procedures, the Paillier cryptosystem can provide two fundamental homomorphisms for scalars: addition and constant multiplication. The two operations are respectively denoted by “⊕” and “⊗”, and the detailed expressions are given as follows.

1) Addition: The addition of two encrypted values satisfies

$$\text{Dec}(\text{Enc}(m_1, \text{pk}) \oplus \text{Enc}(m_2, \text{pk}), \text{sk}) = m_1 + m_2.$$  

2) Constant Multiplication: The multiplication of a constant $c$ in plaintext and an encrypted value satisfies

$$\text{Dec}(c \otimes \text{Enc}(m_1, \text{pk}), \text{sk}) = cm_1.$$  

### D. Quantized Estimator

As mentioned before, the quantization strategy is required in the cryptosystem. The following uniform quantizer $Q : \mathbb{R} \rightarrow \mathbb{Z}$ is considered [42]:

$$Q(x) = \bar{x} + \text{sgn}(x - \bar{x}) \cdot \frac{|x - \bar{x}| + 1}{2}.$$  

Here, $x$ denotes the input while $\bar{x}$ denotes the mid-value of the corresponding interval. The parameter $\delta = \frac{1}{2}$ is the sensitivity, i.e., the maximum error, where $l$ is the length of each quantization interval. $\text{sgn}(\cdot)$ is the sign function that means $\text{sgn}(x) = \begin{cases} 1, & \text{if } x > 0 \\ 0, & \text{if } x = 0 \end{cases}$. Furthermore, the quantization function can be rewritten by

$$Q(x) = x + \Delta,$$

where $\Delta$ is the quantization error satisfying $\Delta \leq \delta$. Nevertheless, it is a function related to a scalar, and the quantization $Q : \mathbb{R}^{m \times n} \rightarrow \mathbb{Z}^{m \times n}$ on matrix $A \in \mathbb{R}^{m \times n}$ is denoted by

$$Q(A) \triangleq \begin{bmatrix} Q(A_{1,1}) & \cdots & Q(A_{1,n}) \\ \vdots & \ddots & \vdots \\ Q(A_{m,1}) & \cdots & Q(A_{m,n}) \end{bmatrix}.$$  

Generally, attitude sensors are deployed in the body of the UAV such that they can observe the rotational state. Such communication is difficult to be overheard by eavesdroppers because it is always propriated and wired. On the other hand, the base stations are deployed in a fixed position on the ground. Therefore, only the privacy of translational estimator $\hat{x}_1(t)$ should be protected. This is also a reason why we divide the global system into two subsystems since the encryption cost will be halved.

Now, the measurements and estimates in the translational subsystem should be quantized for encryption. Although the system parameters do not require privacy preservation, they should also be quantized and mapped into the plaintext space due to the criterion about constant multiplication (33).

It should be pointed out that the accuracy of sensed and estimated data cannot be directly improved by modifying the sensitivity of the quantizer. If the sensitivity decreases, the outputs of the quantizer may not be integers, which is disadvantageous for encryption. To solve this problem, the accuracy is adjusted by introducing a scaling factor $L_q$ in this paper, and the final quantized output can be described as follows:

$$\begin{align*}
\hat{x}_1(t) & \triangleq Q(\hat{x}_1(t)) = (\hat{x}_1(t) + \Delta_x(t))/L_q \\
\hat{y}_1(t) & \triangleq Q(y_1(t)) = (y_1(t) + \Delta_y(t))/L_q \\
K_{C,1}(t) & \triangleq Q(K_{1,1}(t)) = (K_{1,1}(t) + \Delta_C(t))/L_q \\
\hat{K}_{1,1}(t) & \triangleq Q(K_{1,1}(t)) = (K_{1,1}(t) + \Delta_K(t))/L_q
\end{align*}$$

where $K_{1,1}(t) = I - K_{1}(t)C_1$. Then, the quantized estimator is proposed in the following form:

$$\hat{x}_{1,q}(t) = L_q^2 K_{C,q}(t) \hat{x}_{1,q}(t) + L_q K_{1,q}(t) \hat{y}_{1,q}(t).$$

Similar to the analysis in [29], [38], the estimation error under quantization is bounded, because the quantization error is bounded, i.e., $\Delta_x(t), \Delta_y(t), \Delta_C(t), \Delta_K(t) \leq \delta$.

**Remark 3.1:** The probabilistic quantization has been utilized in [29], which randomly selects the output from two boundaries of certain intervals. In this paper, the output is fixed when the input lies in a certain interval in the proposed
quantizer (34). The differences are summarized as follows. 1) Due to the extra randomness, the error of the probabilistic quantizer is larger. Particularly, its error upper bound is twice that of the fixed quantizer. 2) The computation amount of the probabilistic quantizer is also larger because it is time-varying and requires random values at each time. Note that such computation reduction can be negligible in this paper because most computation is used for encryption, but it may work for other applications.

E. Complete Protocol

Based on the above analysis, the complete protocol is summarized in this subsection. The notations “$\oplus$” and “$\otimes$” are omitted for brevity. Firstly, we shall initialize the whole UAS. The initial values of estimation systems, such as system transition parameters, state estimates, and covariances, should be synchronized for all parties, including the user, UAV, and sensor. Meanwhile, another important procedure is generating secret keys for the user as (26), including a private key and a public key. Then, the user distributes public key pk and corresponding system information to the UAV and the sensor.

After initialization, all the parties will work in real-time. Except for the original control input $u(t)$, the user will send encrypted prediction $[[\hat{x}_{1,m}^{t}(t)]]$ to the UAV, which has been processed with quantization, mapping, and encryption:

$$[[\hat{x}_{1,m}^{t}(t)]] = \text{Enc}(M(Q(\hat{x}^{t}_1(t))), pk).$$  (38)

On the other hand, the sensor sends the encrypted measurements $[[y_{1,m}(t)]]$ to the UAV, that is

$$[[y_{1,m}(t)]] = \text{Enc}(M(Q(y(t))), pk).$$  (39)

After gathering these data in UAV, the real-time encrypted estimate $[[\hat{x}_{1,m}(t)]]$ will be computed by

$$[[\hat{x}_{1,m}(t)]] = L_{q}^{2}K_{C,q}(t)[[[\hat{x}_{1,m}(t)]]] + L_{q}^{2}K_{Q,q}(t)=[[y_{1,m}(t)]]$$  (40)

Then, the above encrypted estimate will be directly sent to the user. Besides, the UAV will control the motors with the control signal $u(t)$ as (9). Eventually, after receiving the encrypted signal, the user will apply private key sk for decryption. More concretely, by resorting to the decryption process (30) and retrieval function (31), the final estimate for the user is given as follows, which is equivalent to the quantized result:

$$\hat{x}_{1,q}(t) = R(\text{Dec}([[\hat{x}_{1,m}(t)]]), sk).$$  (41)

The procedures of the proposed protocol mentioned above is summarized in Protocol 1, and a flowchart is given in Fig. 3 to show it more clearly. Before analyzing the security of this protocol, we first give the following definitions.

**Definition 3.1 (Negligible Function [13]):** A function $\epsilon(x)$ is said to be negligible if for every positive polynomial $p(x)$, there exists $\epsilon(x) < 1/p(x)$, $\exists X \in \mathbb{N}$, $\forall x > X$.

**Definition 3.2 (Semantic Security [13]):** Given arbitrary two messages $m_1$ and $m_2$ and randomly encrypt $m_i(i = 1, 2)$ into $[[m_i]]$. A cryptosystem is semantically secure when the result of any probabilistic polynomial-time (PPT) algorithm $i' = D([[m_i]])$ is negligible, i.e., $\Pr(i = i') < \epsilon(t)$.

Then, the security of the proposed protocol is analyzed as follows. Notice that the strongest eavesdropper can simultaneously wiretap both the user-UAV channel and the sensor-UAV channel as stated in Definition 2.1. Therefore, the parameters acquired by the eavesdropper consist of the encrypted measurements $[[y_{1,m}(t)]]$, encrypted predictions $[[\hat{x}_{1,m}(t)]]$, encrypted estimates $[[\hat{x}_{1,m}(t)]]$, control input $u(t)$, and rotational estimate $\hat{x}_2(t)$. Since the Paillier cryptosystem satisfies the semantic security, the eavesdropper that launches CPA cannot directly speculate the decrypted values. Moreover, based on the control input $u(t)$ and rotational estimate $\hat{x}_2(t)$, the translational estimate can also not be calculated with an estimator due to the lack of process and measurement noises.

IV. SIMULATION EXAMPLES

The performances of the proposed methods are studied via computer simulations. The end-user controls the UAV to deliver some important objects at a certain secret position. At the start, the initial states of UAS are all set as 0, which means the UAV is stationary on the ground and its initial position is the origin of the axes. Next, the interval length of the uniform quantizer is set as $l = 1$ which is able to...
TABLE II

<table>
<thead>
<tr>
<th>time slot</th>
<th>10</th>
<th>50</th>
</tr>
</thead>
<tbody>
<tr>
<td>real position</td>
<td>(2.37,0.06,59.64)</td>
<td>(16.05,3.32,177.49)</td>
</tr>
<tr>
<td>quantized measurement</td>
<td>(220,-13,5978)</td>
<td>(1594,346,17739)</td>
</tr>
<tr>
<td>encrypted measurement</td>
<td>(8.8689144e+17+17,1.87115528e+18,4.36727027e+17)</td>
<td>(2.64780208e+18+18,3.52588274e+18,8.18185766e+17)</td>
</tr>
<tr>
<td>encrypted estimate</td>
<td>(4.91360600e+18,4.87347306e+18,2.08898016e+18)</td>
<td>(4.82501600e+18,3.76336150e+18,2.28017677e+17)</td>
</tr>
<tr>
<td>decrypted estimate</td>
<td>(2.40,0.10,59.66)</td>
<td>(16.08,3.34,177.55)</td>
</tr>
</tbody>
</table>

Fig. 4. The real and estimated 3D trajectories of the UAV.

Fig. 5. The comparison of errors between the measurements and estimates.

transfer the float into an integer. The scaling factor is set as $L_q = 0.01$, which reserves two decimal fractions.

We commend the UAV for a period of 100 seconds, and the real and estimated 3D trajectories of the UAV are plotted in Fig. 4 to show the real-time tracking performance. It can be intuitively seen from this figure that the proposed distributed estimator can track the target well. Then, the comparison of errors between measurements and estimates is shown in Fig. 5. We can see that the estimation error is less than the measurement error, which demonstrates the advantage of the proposed estimator on track. To more concretely show the procedures, measurements, estimates, and their encrypted results at certain time slots are given in Table. II. Obviously, the encrypted results are strings of numbers with no meaning, and the eavesdroppers cannot infer valid information from them without the private key.

Moreover, the trajectories of 3D position, attitude, and their estimates are presented in Fig. 6. It is directly seen from this figure that all the sub-states are tracked well. Notice that the estimation error in one test is not convincing enough. In order to assess the estimation performance more accurately, we consider the mean square error (MSE) for quantitative analysis. Here, the theoretical MSE is approximated by applying the Monte Carlo method with 100 interdependent runs. The MSEs of every component of the estimates are given in Fig. 7. All the errors in this figure are bounded, which demonstrates the boundedness of certain components.

Furthermore, the overall MSEs of translational and rotational estimates are displayed in Fig. 8 to show the global stability of the proposed estimator. By introducing the scaling factor, we can achieve almost the same accuracy as the optimal estimator without quantization. Moreover, we also compare the proposed performance with the differentially private mechanism. To achieve differential privacy (DP), the noises with covariance $Q_a = \text{diag}(1.00,1.00,1.00)$ are inserted into the translational measurements, while the value of covariance is determined by the privacy parameters. Obviously, the MSE with DP is larger than that with homomorphic encryption, which demonstrates the proposed estimator can preserve more data utility. Also, the comparison of time costs is shown in Fig. 9. As plotted in this figure, the proposed distributed protocol costs less running time when compared with centralized protocol, because the encryption and ciphertexts-based operations are simplified. Specifically, the time cost of the centralized method is about twice as much as that of the proposed distributed method.

V. CONCLUSION & DISCUSSION

This paper studies the confidentiality problem for UAS and the homomorphic encryption approach is considered to protect the privacy of the UAV’s location against eavesdroppers. First, the globally high-dimensional system is separated into the translational subsystem and the rotational subsystem, and each subsystem is estimated by its own estimator in the
minimum variance sense. Then, a secure estimation protocol with Paillier homomorphic encryption is proposed to preserve the privacy of translational states. Under such protocol, the eavesdropper cannot infer the location of the UAV due to the encryption of measurements and estimates, while the user can obtain the desired data based on homomorphisms and the private key. Meanwhile, the uniform quantization method with fixed sensitivity is applied to change the type of transmitted data from the float into integers. Compared with centralized estimation and full homomorphism, the computation amount of the proposed protocol is greatly reduced thanks to the design of the distributed structure and partial encryption.

It is important to acknowledge that there are inherent trade-offs among various system performance metrics, including the demand for privacy, the volume of computation, and the accuracy of estimation. These factors often counterbalance each other, necessitating careful consideration and strategic decision-making. The homomorphic encryption adopted in this paper provides a higher privacy guarantee, while this comes at the cost of a significantly higher computational load. Although the distributed structure is proposed to mitigate such adverse impact, the calculation is still large data perturbation approaches, which only requires simple randomness. Therefore, the selection of an appropriate method should be guided by practical requirements and the specific context of the situation.

REFERENCES


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