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## Polyhedral Reductions for Petri nets

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## Polyhedral equivalence

$$(N_1, m_1) \equiv_E (N_2, m_2)$$

Polyhedral equivalence defines a *linear dependence relation* ( $E$ ) between the reachable markings of two nets; where  $E$  is a Presburger formula, or a system of linear Diophantine equations. This equivalence, denoted  $(N_1, m_1) \equiv_E (N_2, m_2)$ , preserves enough information in  $E$  so that we can rebuild the state space of  $N_1$  knowing only the one of  $N_2$ .

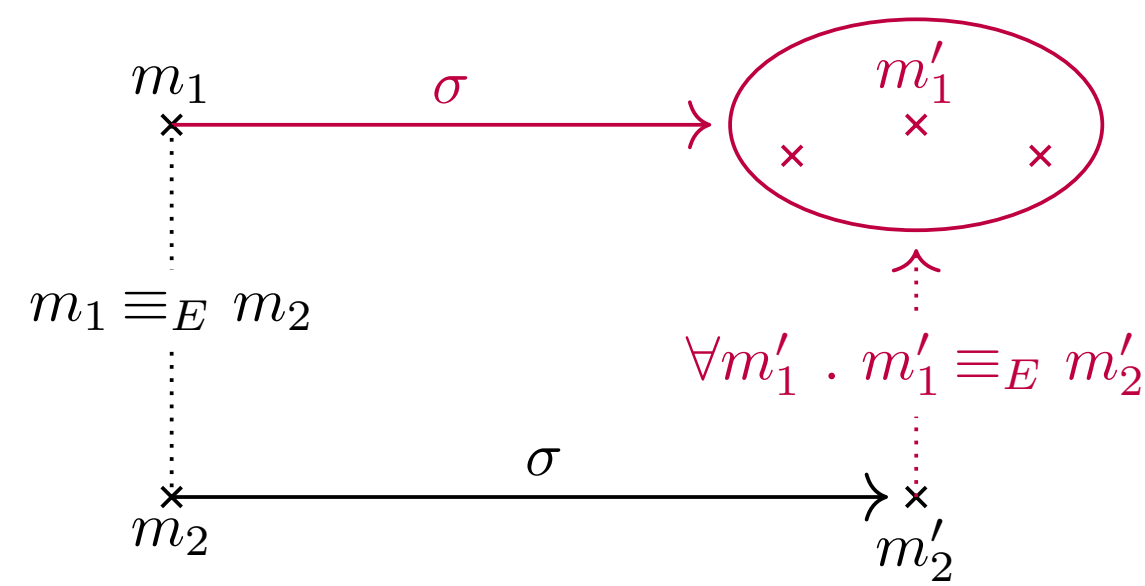
We say that two markings  $m_1, m_2$  are compatible up-to  $E$ , denoted  $m_1 \equiv_E m_2$ , when formula  $E \wedge m_1 \wedge m_2$  is satisfiable, with  $m$  the conjunction  $\bigwedge_{p \in P} p = m(p)$  (e.g.,  $(p_0 = 5) \wedge (p_1 = 0) \wedge \dots$ ).

Nets  $(N_1, m_1)$  and  $(N_2, m_2)$  are  $E$ -equivalent when:

- (A1)  $E \wedge m$  is satisfiable for all markings  $m$  in  $R(N_1, m_1)$  or  $R(N_2, m_2)$ ;
- (A2) initial markings are *compatible*:  $m_1 \equiv_E m_2$ ;
- (A3) assume  $m'_1, m'_2$  are markings of  $N_1, N_2$ , such that  $m'_1 \equiv_E m'_2$ , then  $m'_1$  is reachable iff  $m'_2$  is reachable.

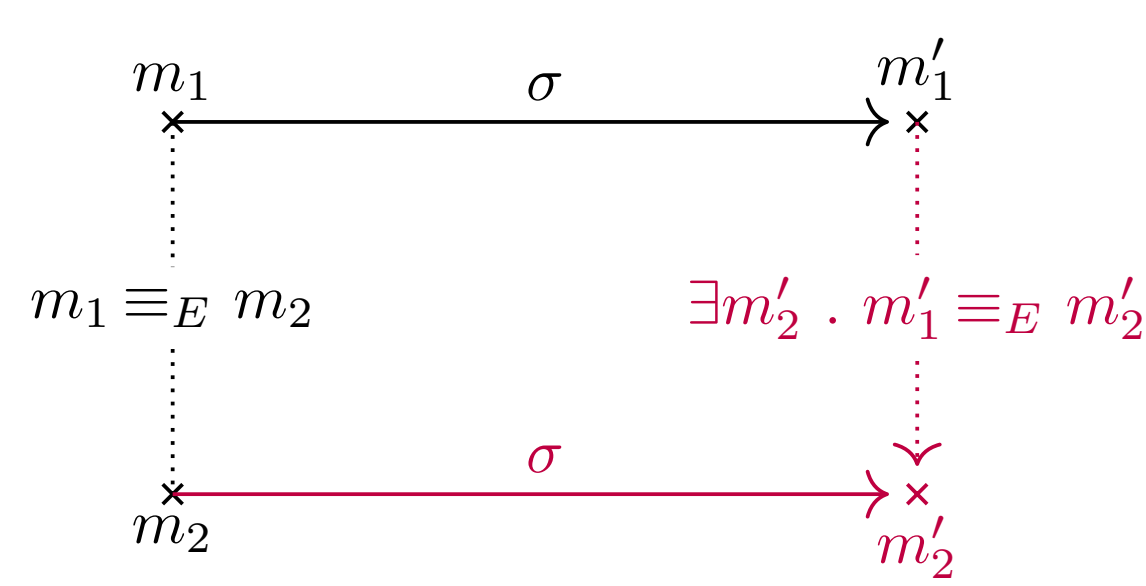
### Lemma (Reachability Checking).

Given  $m'_2$  reachable in  $N_2$  then  $m'_1 \equiv_E m'_2$  implies  $m'_1 \in R(N_1, m_1)$ .



### Lemma (Invariance Checking).

For all  $m'_1$  in  $R(N_1, m_1)$  there is  $m'_2$  in  $R(N_2, m_2)$  such that  $m'_1 \equiv_E m'_2$ .



## Counting using SDD



```
$ sift AirplaneLD-PT-0050.pnml
4471223 marking(s), 19756224 transitions(s)
67.569s

$ tedd -q AirplaneLD-PT-0050.pnml
4471223 state(s)
2.691s

$ tedd -q AirplaneLD-PT-4000.pnml
2.1773922960e12 state(s)
33.574s

$ tedd -rg,redundant,compact AirplaneLD-PT-4000.pnml
2.1773922960e12 state(s)
0.000s
```

## Reachability using SMT



To check if property  $F_1$  is reachable in  $N_1$ , it is enough to check its  $E$ -transform,  $F_2$ , on  $N_2$ .

$$F_2(\mathbf{y}) \triangleq \exists \mathbf{x} . \tilde{E}(\mathbf{x}, \mathbf{y}) \wedge F_1(\mathbf{x})$$

We use this property in a SMT-based model-checker with structural reductions. It supports generalized reachability formulas of the form:  $AG((p_1 + p_3 \geq 5) \vee (p_0 < p_4))$

```
$ python3 -m smpt -n SmallOpSys-MT0016.pnml
--methods BMC K-INDUCTION PDR-REACH STATE-EQUATION
--show-time --show-techniques ...
FORMULA RC4 FALSE TECHNIQUES ... BMC TIME 0.04s

$ python3 -m smpt -n SmallOpSys-MT2048.pnml ...
timeout

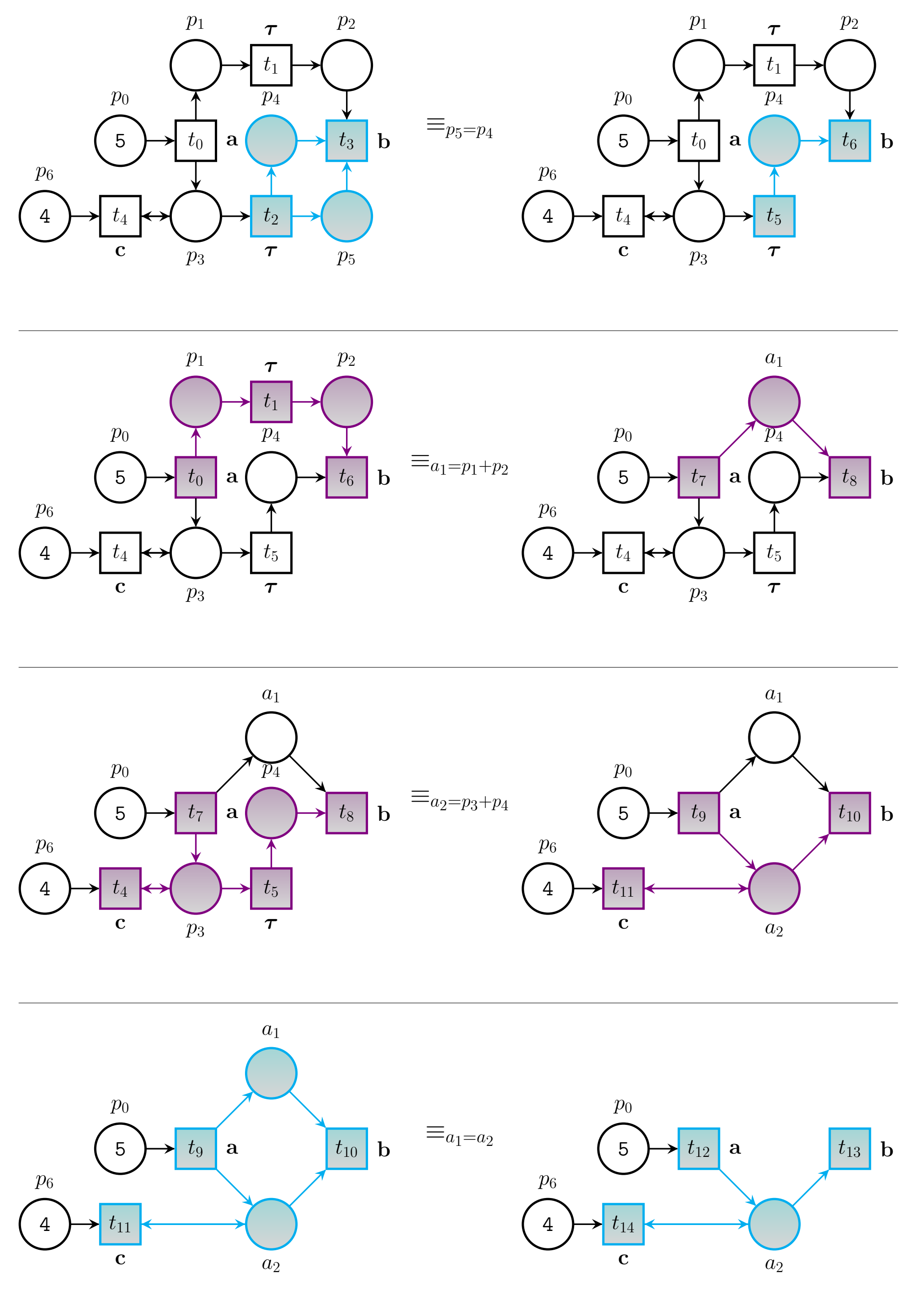
$ same with option --auto-reduce or --mcc
# Reduction time (Full): 0.11s
FORMULA RC4 FALSE TIME 0.05s

$ or with options --project --show-projection
# Reduction time (TFG): 0.09s
# Reduction ratio (TFG) 44%
# Projection of RC4: [ ] - (1 ≤ p0 + p2 + a2) ∧ ...
```

## Polyhedral reductions



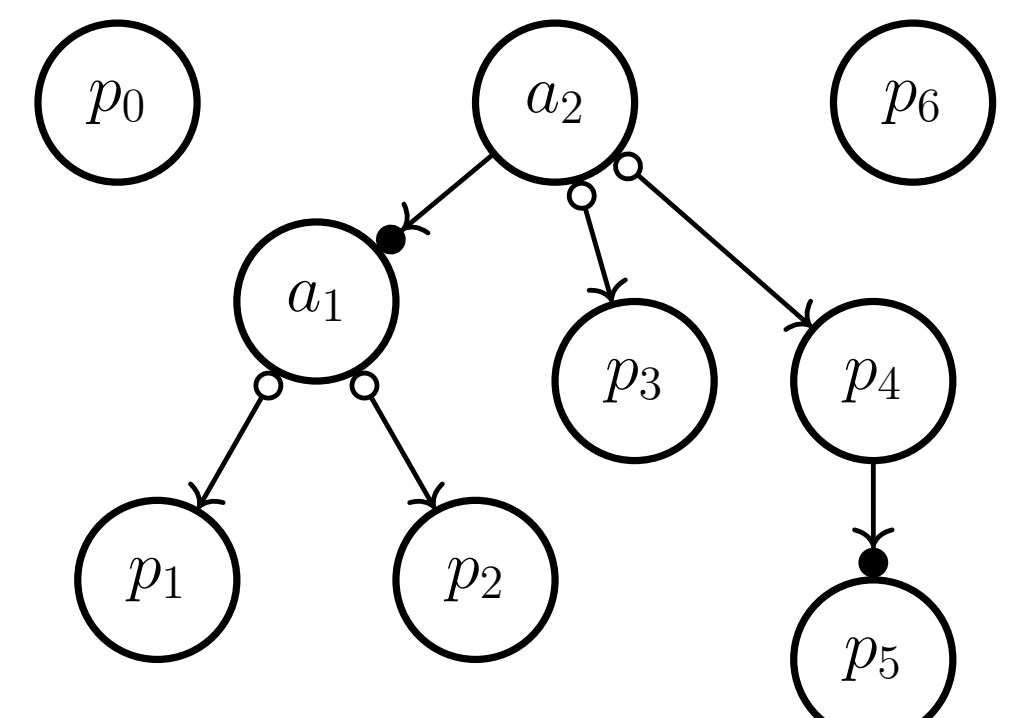
We can compute instances of polyhedral equivalences using structural reductions.



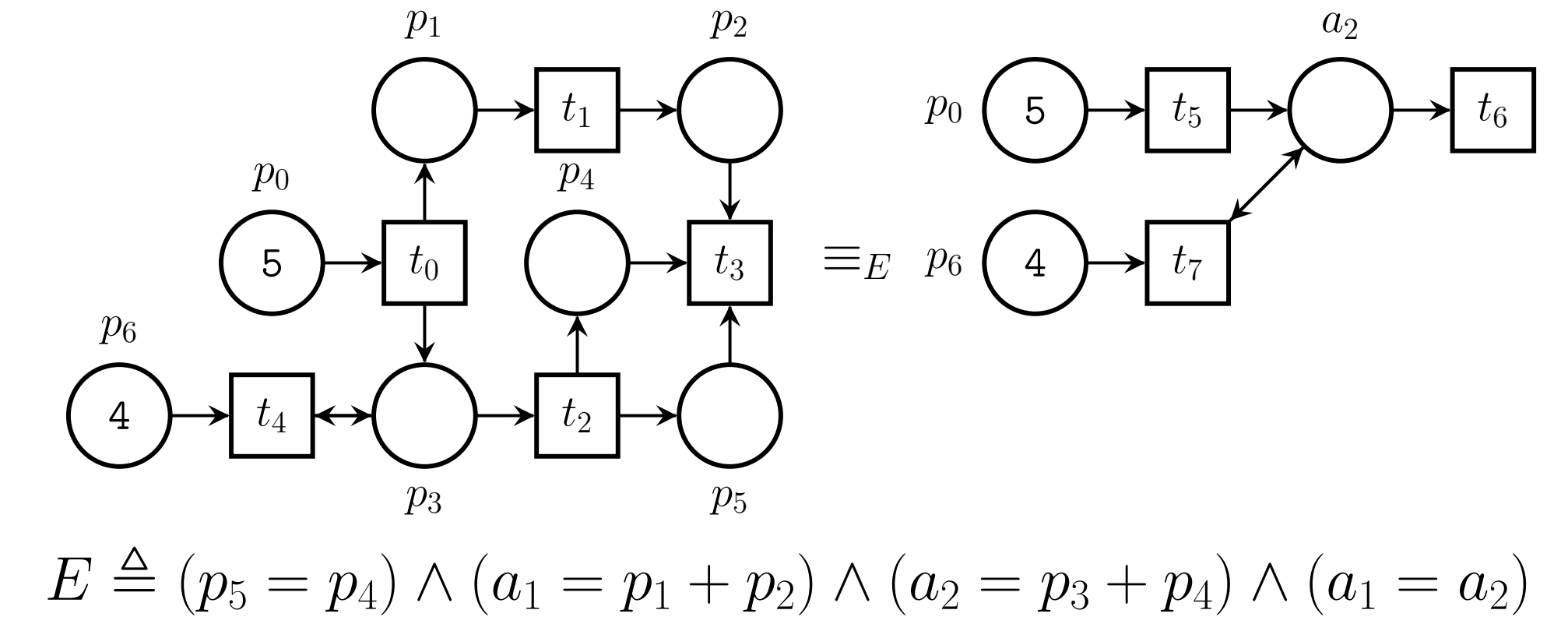
## Token Flow Graph

- Octant
- Kong
- Polycount

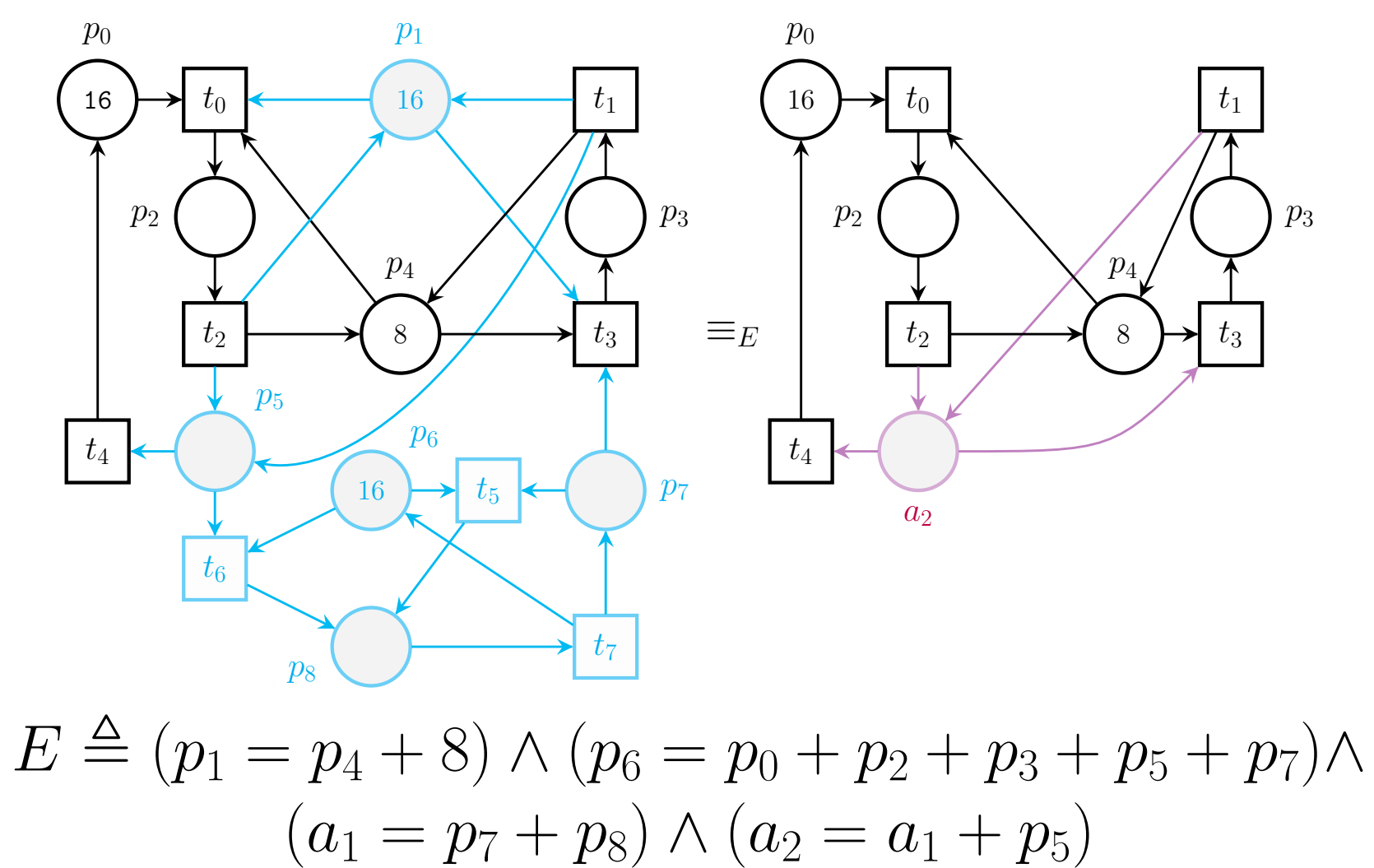
- #  $R \vdash p_5 = p_4$
- #  $A \vdash a_1 = p_2 + p_1$
- #  $A \vdash a_2 = p_4 + p_3$
- #  $R \vdash a_1 = a_2$



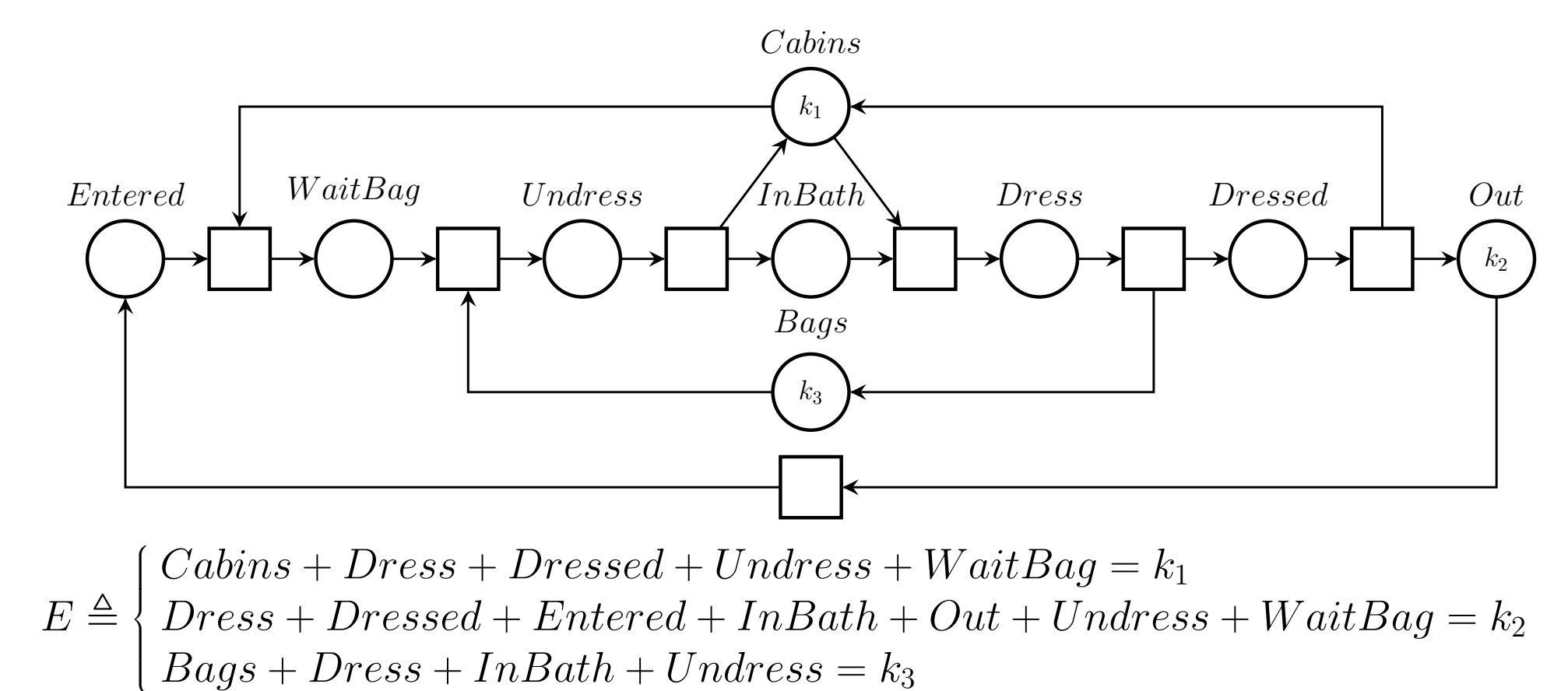
## Some examples



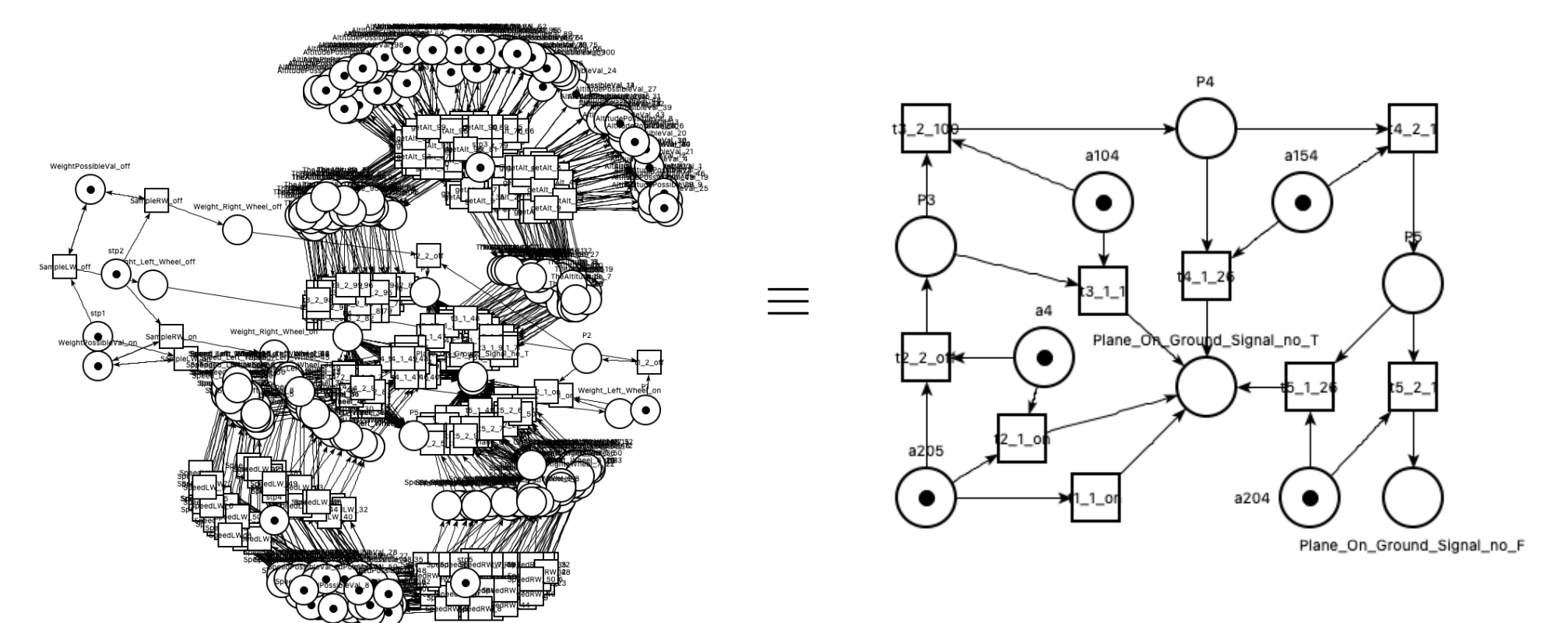
### SmallOperatingSystem-PT-MT0016DC0008



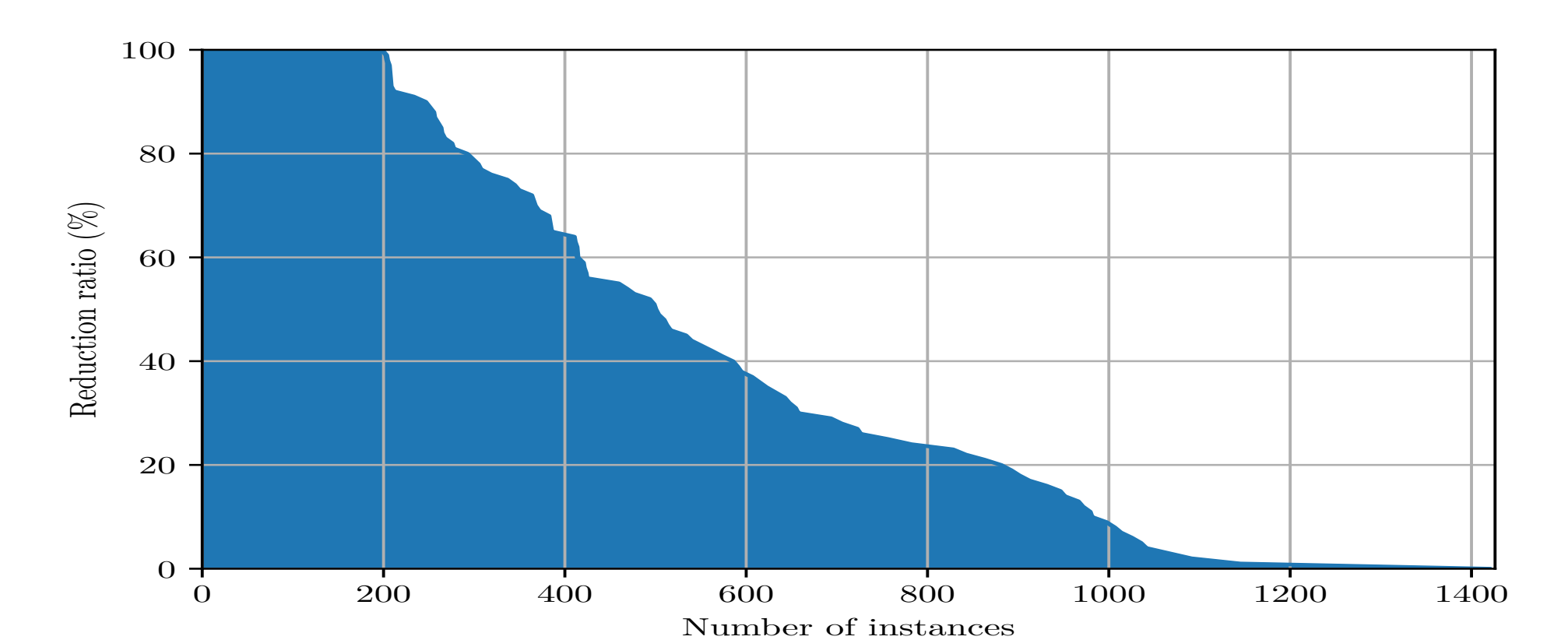
### SwimmingPool (fully reduced)



### AirplaneLD-PT-0050

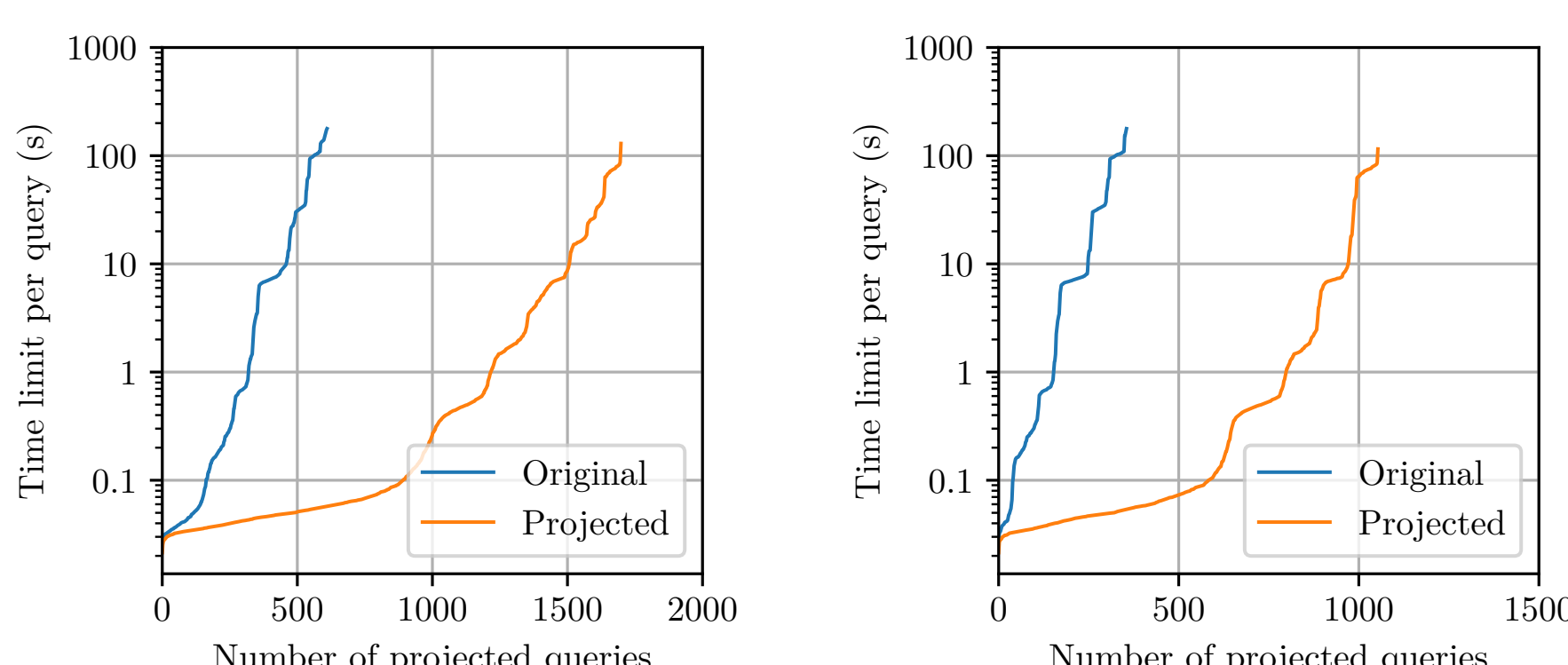


## Prevalence of reductions in the MCC



## Automatic equivalence checking

The problem of checking whether two nets are equivalent is undecidable. But we have implemented a sound, complete procedure in the case of equivalence between *parametric nets*:  $(N_1, C_1) \approx_E (N_2, C_2)$ .



$$r \in [0.01, 0.5[ \quad r \in [0.5, 1]$$

$k$ -induction with and without reductions

