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Low-Complexity Scheduled Expectation Propagation based on QRD and SIC

Adam Mekhiche, Antonio Maria Cipriano, Charly Poulliat

Abstract—This letter presents novel Expectation Propagation (EP) algorithms for Single User (SU) and Multiple Users (MU) Multiple-Input Multiple-Output (MIMO) detection that address the performance loss and slow convergence of classic Scalar EP (SEP) solutions with even lower computational complexity. The proposed algorithms employ QR Decomposition (QRD), optimized EP message scheduling, and Successive Interference Cancellation (SIC) techniques to achieve superior performance and convergence speed comparable to those of Vector Expectation Propagation (VEP). Our results are illustrated through error rate curves and analysis of performance-complexity trade-offs in SU and MU scenarios.

Index Terms—Expectation Propagation, QR Decomposition, Scheduling, Successive Interference Cancellation

I. INTRODUCTION

IMO with Orthogonal Frequency Division Multiplexing (OFDM) is a widely used technology at the core of many wireless telecommunication standards like Wireless Fidelity (WiFi), 4G - Long Term Evolution (LTE) and 5G - New Radio (NR), in part, due to the possible high spectral efficiency gain through Spatial Multiplexing (SM). One recent proposal is to boost performance by introducing tens of antennas, e.g. massive MIMO (mMIMO), which comes at the cost of high computational complexity for detection, i.e. an exponential increase with the number of antennas. Bit Interleaved Coded Modulation with Iterative Detection (BICM-ID) can achieve optimal performance through the iterative exchange of soft information, usually based on Log-Likelihood Ratios (LLR), between a Soft-Input Soft-Output (SISO) detector and a SISO decoder at the cost of an even higher complexity due to several instances of detection and decoding. To cope with the increase in complexity, new SISO detection algorithms have been proposed based on Message Passing (MP), such as Belief Propagation (BP) [1], Approximate Message Passing (AMP) [2], or Vector AMP (VAMP) [2], [3] algorithms. These MP Algorithms (MPA) aim at computing a low-complexity approximation of the symbol-wise Maximum A Posteriori (MAP) criterion, which is too costly in practice.

In this paper, we focus on MPA derived within the EP framework [4], which is another approach to deriving an efficient MPA. EP has been shown to be a good candidate for MIMO and mMIMO applications compared to other MPA. BP

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has a computational complexity slightly higher than MAP [1], while VAMP requires a costly Singular Value Decomposition (SVD) for a matrix inversion. Authors from [5] proposed to apply EP on a scalar Factor Graph (FG) to reduce the complexity, by removing the need to invert a full matrix, compared to EP used on a vector FG like in [6], [7]. Similarly to [8], authors from [9] proposed a QRD pre-processing with EP to reduce the number of messages exchanged during an EP iteration and improve performance. Even though better performance is achieved, there is still room for improvement to close the gap between the performance of a SEP [5], [9], [10] and a VEP [6], [7], [10].

In this paper, we propose an optimized scheduling of the exchange of EP messages to improve both convergence speed and performance for SU-MIMO and MU-MIMO. The article is organized as follows: In Section II, the system model is introduced, and then, the proposed scheduled SEP with QRD is explained in Section III. The related algorithm is detailed in Section III-A for the SU setting and in Section III-B for the MU setting. Its complexity is assessed in Section III-C and compared with other competitive MPA. Finally, simulation results are presented in Section IV.

II. SYSTEM MODEL

A spatially multiplexed $N_t \times N_r$ SU-MIMO-OFDM system using a Space-Time (ST) BICM-ID scheme [11] with N_t transmitting antennas and N_r receiving antennas is considered. At time $n \in [1, N_s]$, a stream of information bits $\mathbf{b}_n \in \mathbb{F}_2^{K_b}$, where \mathbb{F}_2 is the finite field of order 2, is encoded using an Error Correcting Code (ECC) of length N and rate $R = K_b/N$. The encoded bits $\mathbf{c}_n \in \mathbb{F}_2^N$ are interleaved with a random interleaver, and then Gray-mapped to symbols of a constellation $\mathcal{X} \subset \mathbb{C}$. A symbol x represents a tuple of m bits and there are $|\mathcal{X}| = M = 2^m$ symbols in the constellation. The mapping function is denoted $\varphi: \mathbb{F}_2^m \to \mathcal{X}$ while $\varphi^{-1}: \mathcal{X} \to \mathbb{F}_2^m$ is the demapping function that returns the m-tuple of coded bits. $\forall k \in [1, m], \varphi_k^{-1} : \mathcal{X} \to \mathbb{F}_2$ is the demapping function that returns the single k^{th} bit corresponding to a symbol x. The vector of symbols $\mathbf{x}_n = [x_{n,1}, \cdots, x_{n,N_t}]$ is then splitted across the N_t antennas. This system model employs a Space-Time (ST) [11] Bit-Interleaved Coded Modulation with Iterative Decoding (BICM-ID) scheme. The other model studied is single antenna MU-MIMO where, instead of a single information bits stream, there are $N_u = N_t$ user and one stream per user, which is equivalent to N_u SU-MIMO. At the time n, the signal \mathbf{x}_n is sent through N_t transmitting antennas on a non-correlated ergodic Rayleigh channel $\mathbf{H}_n \in \mathbb{C}^{N_r \times N_t}$

and the signal $\mathbf{y}_n \in \mathbb{C}^{N_r}$ is received using N_r antennas. The corresponding linear model is:

$$\mathbf{y}_n = \mathbf{H}_n \mathbf{x}_n + \mathbf{w}_n, \tag{1}$$

with $\mathbf{w}_n \in \mathbb{C}^{N_r}$ an Additive White Gaussian Noise (AWGN) sample with the properties $\mathbb{E}(\mathbf{w}_n) = 0$ and $\mathbb{E}(\mathbf{w}_n \mathbf{w}_n^H) = N_0 I_{N_r}$. The time index will be omitted for readability except when needed for comprehension. The Channel State Information (CSI) available on the receiver side is considered perfect. The turbo-iterated receiver aims at computing the probability $\mathbb{P}(\mathbf{x}|\mathbf{y},\mathbf{H},N_0)$ which can be factorized as:

$$\mathbb{P}(\mathbf{x}|\mathbf{y},\mathbf{H},N_0) \propto \prod_{j=1}^{N_r} \underbrace{\mathbb{P}(y_j|\mathbf{x},\mathbf{H},N_0)}_{f_i^{\text{EQU}}} \prod_{i=1}^{N_t} \underbrace{\mathbb{P}(x_i|\mathbf{c}_i)}_{f_i^{\text{DEM}}} \underbrace{\mathbb{P}(\mathbf{c}_i)}_{f_i^{\text{DEC}}}$$
(2)

where $f_j^{\rm EQU}$ are likelihood functions, $f_i^{\rm DEM}$ are the factors associated with the demodulation process, and $f_i^{\rm DEC}$ are priors associated with coded binary vector labels \mathbf{c}_i , which are provided by the ECC during the turbo-iterated process.

From the factorization (2), the corresponding Factor Graph (FG) can be represented as in Fig. 1, which is the typical scalar FG in [5], [9], [10]. This FG is the support on which Message Passing Algorithms (MPAs) are applied to approximate the Probability Density Function (PDF) $\mathbb{P}(\mathbf{x}|\mathbf{y},\mathbf{H},N_0)$. Authors from [5], [9] applied Expectation Propagation (EP) [4] on this FG. EP is an approximate Bayesian inference technique within the message-passing framework. It tracks a distribution p by approximating it with another distribution q from a specific set Q, from the exponential family. To project p on the set Q, i.e. choosing the statistically closest distribution $q \in Q$, EP uses the inclusive Kullback-Leibler divergence [12]:

$$D_{KL}(p||q) = \int p(x) \log\left(\frac{p(x)}{q(x)}\right) dx \tag{3}$$

so the projection can be written:

$$q = \operatorname{proj}_{\mathcal{Q}}[p] = \operatorname*{arg\,min}_{\tilde{q} \in \mathcal{Q}} D_{KL}(p||\tilde{q}). \tag{4}$$

EP exchanges messages, which are Gaussian distributions, along the edges of an FG, and the message from a factor node f_i to a variable node x_i is:

$$m_{f_j \to x_i}(x_i) = \frac{1}{m_{x_i \to f_j}(x_i)} \times$$

$$\operatorname{proj}_{\mathcal{Q}} \left[m_{x_i \to f_j}(x_i) \int_{\mathbf{x}} f_j(\mathbf{x}) \prod_{\substack{i' \in \mathcal{N}(f_j) \\ i' \neq i}} m_{x_{i'} \to f_j}(x_{i'}) \right]. \quad (5)$$

with $\mathcal{N}(f_j)$ the set of connected variable nodes to (f_j) . To improve EP on this FG, authors from [9] proposed a QR Decomposition as a pre-processing of the received signal, which helps reduce the number of edges in the graph and improves performance. This matrix decomposition was also studied with other MPAs like Belief Propagation [8], [13]. Applying QRD on \mathbf{H} , i.e. $\mathbf{H} = \mathbf{Q}\mathbf{R}$, gives the equivalent observation model

$$\tilde{\mathbf{y}} = \mathbf{Q}^H \mathbf{y} = \mathbf{R} \mathbf{x} + \tilde{\mathbf{w}} \tag{6}$$

with $\mathbf{Q} \in \mathbb{C}^{N_r \times N_r}$, a unitary matrix, $\mathbf{R} \in \mathbb{C}^{N_r \times N_t}$, an upper triangular matrix and $\tilde{\mathbf{w}} = \mathbf{Q}^H \mathbf{w}$. Statistical properties of the additive noise remain unchanged due to the unitary property of

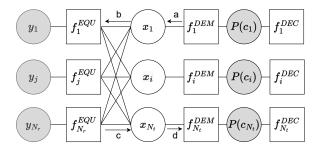


Fig. 1. Factor Graph representation of (2) with $a \triangleq \mathbb{C}\mathcal{N}(\overleftarrow{\mu}_i^d, \overleftarrow{\nu}_i^d), b \triangleq \mathbb{C}\mathcal{N}(\overleftarrow{\mu}_{i,j}^e, \overleftarrow{\nu}_{i,j}^e), c \triangleq \mathbb{C}\mathcal{N}(\overrightarrow{\mu}_{i,j}^e, \overrightarrow{\nu}_{i,j}^e), d \triangleq \mathbb{C}\mathcal{N}(\overrightarrow{\mu}_i^d, \overrightarrow{\nu}_i^d)$.

Q. The new FG after QR pre-processing is shown in Fig. 2. In the next section, we propose improved versions of EP applied to a scalar FG, which shows better performance and faster convergence than existing SEP algorithms.

III. PROPOSED SCHEDULED SCALAR EXPECTATION PROPAGATION WITH QRD

A. SSEP ORD in SU-MIMO

Authors from [5], [7], [9], [10] proposed a flooding scheduling of the EP messages on the graph. The proposed algorithm aims at computing $\mathbb{P}(\mathbf{x}|\mathbf{y},\mathbf{H})$ more efficiently than the classic Scalar EP (SEP) [5] or QRD SEP [9], thanks to better scheduling of the message exchange. The QRD pre-processing of the channel matrix generates a new sparser FG (Fig. 2), which can be more efficiently browsed in the same manner as a classic Block Decision Feedback Equalizer (BDFE) [14].

The algorithm starts the messages exchange at the bottom of the FG to compute an estimation of x_{N_t} , which is connected to all the equalizing nodes. It means, $\forall j \in [\![1,N_r]\!]$, it propagates the messages $m_{f_j^{\rm EQU} \to x_{N_t}}(x_{N_t})$ to gather all the available information about x_{N_t} at the equalizing side of the FG. Instead of simultaneously propagating messages toward the others x_i like in SEP [5] or QRD SEP [9], the following propagated message is $m_{x_{N_t} \to f_{N_t}^{\rm DEM}}(x_{N_t})$. Then, information coming from the demapper is sent to the variable node x_{N_t} with the message $m_{f_{N_t}^{\rm DEM} \to x_{N_t}}(x_{N_t})$.

At this point, the variable x_{N_t} has been estimated with all the available information. Hence, the following targeted variable node is x_{N_*-1} . It will be more accurately estimated than in a flooding scheduling since the equalizing nodes have a better estimation of x_{N_t} , as it comes from the demapper. Messages are exchanged similarly, propagating the estimation to the demapper and sending the demapping information to the equalizing nodes, up to node x_1 . The previous process can be eventually repeated multiple times (denoted as auto-iterations). After L auto-iterations, the messages $m_{x_i o f^{\mathrm{EOU}}}(x_i)$ are propagated once more time. Compared to [9], this step results in a slightly higher complexity, which is analyzed in Section III-C. The overall number of equalization steps and demapping steps is the same as in ORD SEP. This algorithm is presented in Alg. 1 and can be used for SUor MU-MIMO. More efficient scheduling is also possible for MU-MIMO and SU-MIMO with V-BLAST [15]. It will be presented in the next section.

Algorithm 1 Scheduled Scalar EP with QRD

```
Input: y, H, N_0, T = Turbo iteration, L = Auto iteration
Output: \hat{\mathbf{b}} - the estimated bits.
   1: \mathbf{H} = \mathbf{Q}\mathbf{R}, \quad \mathbf{y} \leftarrow \mathbf{Q}^H\mathbf{y}, \quad \mathbf{H} \leftarrow \mathbf{R}
   2: \forall i, k \in [\![1, N_t]\!], [\![1, m]\!], \lambda_{i,k}^a = 0
   3: for t = 0 : T do
                    \begin{array}{l} \forall i,j \in \llbracket 1,N_t \rrbracket, \llbracket 1,N_r \rrbracket, \\ m_{f_i^{\text{EQU}} \rightarrow x_i}(x_i) = m_{x_i \rightarrow f_i^{\text{DEM}}}(x_i) = \mathbb{C} \mathcal{N}(0,+\infty) \end{array}
   4:
   5:
                    for i = 1 : N_t do
   6:
                             Compute m_{f^{\text{DEM}} \to x_i}(x_i):
   7:
                           \begin{split} \tilde{q}(x_i) &\propto \exp\left(-\frac{|x_i - \mu_i^d|^2}{\nu_i^d} - \sum_{k=1}^m \varphi_k^{-1}(x_i)\lambda_{i,k}^a\right) \\ q(x_i) &\sim \mathbb{C}\mathcal{N}\left(\mu_i^d = \mathbb{E}\left[\tilde{q}(x_i)\right], \nu_i^d = \operatorname{Var}\left[\tilde{q}(x_i)\right]\right) \\ \leftarrow \mu_i^d &= \nu_i^d \left(\frac{\mu_i^d}{\nu_i^d} - \frac{\mu_i^d}{\nu_i^d}\right) \text{ and } \nu_i^d = \left(\frac{1}{\nu_i^d} - \frac{1}{\nu_i^d}\right)^{-1} \end{split}
   8:
   9:
 10.
 11:
                    for i = 1 : N_t and then for j = 1 : N_r do
 12:
                             Compute m_{x_i \to f_i^{\text{EQU}}}(x_i):
 13:
                           14:
 15:
 16:
                    for l = 0 : L and then i = N_t : -1 : 1 do
 17:
                            for j = 1 : N_r do
 18:
                                      Compute m_{f_{\cdot}^{\text{EQU}} \to x_i}(x_i):
 19:
                                     \overrightarrow{\nu_{i,j}^{e}} = \left(N_0 + \sum_{i' \neq i} |h_{j,i'}|^2 \overleftarrow{\nu_{i',j}^{e}}\right) / |h_{j,i}|^2
\overrightarrow{\mu_{i,j}^{e}} = \left(y_j - \sum_{i' \neq i} h_{j,i'} \overrightarrow{\mu_{i',j}^{e}}\right) / h_{j,i}
 20:
 21:
 22:
                            Compute m_{x_i \to f_i^{\text{DEM}}}(x_i):
23:
                            \overrightarrow{\nu_i^d} = \left(\sum_{j' \in \mathcal{N}(x_i)} \frac{1}{\overrightarrow{\nu_{i,j'}^e}}\right)^{-1}, \ \overrightarrow{\mu_i^d} = \overrightarrow{\nu_i^d} \left(\sum_{j' \in \mathcal{N}(x_i)} \frac{\overrightarrow{\mu_{i,j'}^e}}{\overrightarrow{\nu_{i,j'}^e}}\right)^{-1}
 24:
 25:
                                      Compute m_{f_i^{\text{DEM}} \to x_i}(x_i):
 26:
                                      Mean and variance as in Alg. 1 line. 7
 27:
 28:
                                      for j = 1 : N_r do
                                               Compute m_{x_i \to f_i^{\text{EQU}}}(x_i):
 29:
                                               Mean and variance as in Alg. 1 line. 13
 30.
                                      end for
31:
                            end if
 32:
 33:
                    end for
                   \begin{array}{l} \text{for } i=1:N_t \text{ and then for } k=1:m \text{ do} \\ \lambda_{i,k}^e = \log \left( \frac{\sum_{x_i \in \mathcal{X}: \varphi_k^{-1}(x_i) = 1} \bar{q}(x_i)}{\sum_{x_i \in \mathcal{X}: \varphi_k^{-1}(x_i) = 0} \bar{q}(x_i)} \right) - \lambda_{i,k}^a \end{array}
 34:
35:
 36:
                    Send \lambda^e to decoder and receive \lambda^a and \hat{\mathbf{b}}
38: end for
```

B. SSEP QRD SIC for MU-MIMO

In MU-MIMO, it is possible to benefit from the error correction capability of the ECC within a turbo-iteration through the use of Successive Interference Cancellation (SIC). Indeed, we propose another detector using a Scheduled Scalar Expectation Propagation with QR Decomposition and SIC (QRD SSEP SIC). Instead of propagating messages of each user up to the demapper, we propose to go through the decoder and use the extrinsic output of the decoder to refine the soft estimation of the associated symbol.

Fig. 2 shows the differences between QRD SSEP in SU-MIMO, where decoding is done concurrently since a single codeword is split across all antennas, and QRD SSEP SIC in MU-MIMO, where each user encodes its codeword which

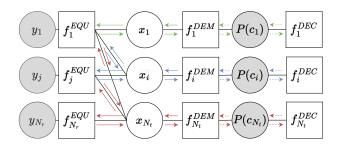


Fig. 2. Factor Graph representation of (2) after QRD. Red messages are propagated first, then the blue ones, and finally, the green ones. Dotted lines are done in parallel for SU-MIMO and serialized, in accordance with their color, for MU-MIMO.

Algorithm 2 Scheduled Scalar EP with QRD and SIC

```
Input: y, H, N_0, T = Turbo iteration, L = Auto iteration
Output: \hat{\mathbf{b}} - the estimated bits.
 1: \mathbf{H} = \mathbf{Q}\mathbf{R}, \quad \mathbf{y} \leftarrow \mathbf{Q}^H\mathbf{y}, \quad \mathbf{H} \leftarrow \mathbf{R}
2: \forall i, k \in [\![1, N_t]\!], [\![1, m]\!], \quad \lambda_{i,k}^a = 0
 3: for t = 0 : T do
            \begin{aligned} &\forall i,j \in [\![1,N_t]\!], [\![1,N_r]\!], \\ &m_{f_i^{\text{EQU}} \to x_i}(x_i) = m_{x_i \to f_i^{\text{DEM}}}(x_i) = \mathbb{C}\mathcal{N}(0,+\infty) \end{aligned}
  5:
             \mathbf{for}^{j} i = 1 : N_t \ \mathbf{do}
  6:
                  Compute m_{f^{\text{DEM}} \to x_i}(x_i) as in Alg. 1 line 7.
  7:
 8:
             for i = 1 : N_t and then for j = 1 : N_r do
 9:
                  Compute m_{x_i \to f_i^{\text{EQU}}}(x_i) as in Alg. 1 line 13.
10:
11:
            for l = 0 : L and then i = N_t : -1 : 1 do
12:
                  for j = 1 : N_r do
13:
14:
                        Compute m_{f_i^{\text{EQU}} \to x_i}(x_i) as in Alg. 1 line 19.
15:
16:
                  Compute m_{x_i \to f_i^{\text{DEM}}}(x_i) as in Alg. 1 line 23.
17:
                  if l = L then
                        Compute \lambda_i^e as in Alg. 1 line 35.
18:
                        Send \lambda^e to decoder and receive \lambda^a and \hat{\mathbf{b}}
19
20:
                  end if
                  Compute m_{f_i^{\text{DEM}} \to x_i}(x_i) as in Alg. 1 line 7.
21:
                  for j=1:N_r do
22:
23:
                        Compute m_{x_i \to f_i^{\text{EQU}}}(x_i) as in Alg. 1 line 13.
                  end for
24:
25:
             end for
26: end for
```

can be independently decoded. This scheduling enables the refinement of the detection process for the other users by obtaining a more accurate estimation of a user's symbol after a single decoding step. It eliminates the need to wait for all users to be decoded, i.e. the next turbo-iteration, to converge.

C. Complexity Analysis of SSEP QRD

In this section, a precise estimation of the computational complexity is performed to assess the gain of the proposed algorithm.

VEP [6] has an $\mathcal{O}(N^3)$ complexity as it requires a $N \times N$ matrix inversion, with $N = \min(N_t, N_r)$, to compute messages coming from the equalizer node. DEP [7] has the same complexity as VEP since both use the same FG and apply the same EP message rule on it. Their difference lies in saving messages $m_{x_i \to f_i^{\text{DEM}}}(x_i)$, $\forall i \in [1, N_t]$, from one turbo iteration to the next for DEP, and resetting these messages for VEP. Authors of VEP also proposed to keep the same EP messages in [16, Chapter 4, Section 3]. SEP [5] does

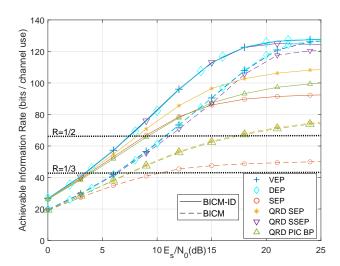


Fig. 3. Achievable Information Rate (AIR) of a 32×32 SU-MIMO with a 16QAM ST-BICM and ST-BICM-ID transmissions.

not require any matrix inversion, and its overall complexity is smaller than VEP even though there are more messages to be exchanged between equalizing nodes and variable nodes x_i , $N_t \times N_r$ instead of N_t . Its most computationally demanding messages are the messages from the equalizing nodes to variable node x_i . QRD SEP [9] must perform one QR decomposition per MIMO symbol, no matter the number of auto and turbo iterations, and QRD has a complexity $\mathcal{O}(N_t^3/3)$. Once this pre-processing is done, the messages are a bit less complex, fewer sums in (19), and there are $N_t(N_t+1)/2$ messages instead of $N_t \times N_r$. Finally, the proposed QRD SSEP has a slightly higher complexity than QRD SEP since there are $N_t(N_t+1)/2$ more messages from variable node x_i to equalizing node f_i^{EQU} . This additional complexity is minor compared to the overall complexity: it results in a 5% or less increase in complexity for QRD SSEP compared to QRD SEP.

Table I shows two relative computational complexities per MIMO scenario. The first percentage is the relative complexity of QRD SSEP over VEP to highlight the gain of the proposed algorithms compared to the optimal EP. In all the presented scenarios, QRD SSEP has much lower complexity than VEP, from -10% for small 4×4 MIMO without auto-iteration to -58% for massive 100×100 MIMO cases with high order constellations and two auto iterations which are the scenarios where VEP might be too complex to be implemented. The second percentage is the complexity of QRD SSEP over QRD SEP to show the additional marginal cost of the proposed algorithm compared to its closest competitor QRD SEP. The two algorithms have very similar complexities per auto iteration, but their performance is vastly different, as shown in the next section. Finally, QRD SSEP SIC has the same computational complexity as QRD SSEP since they compute as many messages for detection (complexity for decoding is omitted here). However, the serial activation of the decoders in the MU case can cause, in certain implementations, a higher decoding latency, which is typical for a SIC detector. Compared to (turbo)-MMSE-SIC [17], [18] or VEP-SIC, this detector does not require to do more operations than its non-SIC variant because of the scalar nature of its underlying FG.

TABLE I RELATIVE DETECTION COMPLEXITIES: QRD SSEP / VEP - QRD SSEP/QRD SEP

Auto iterations MIMO scenarios	0	1	2
4 x 4 - 16 QAM	93% - 103%	84% - 103%	81% - 103%
10 x 10 - 8 PSK	80% - 103%	62% - 104%	56% - 104%
32 x 32 - 16 QAM	73% - 101%	52% - 101%	46% - 102%
100 x 100 - 64 QAM	71% - 100%	50% - 100%	42% - 100%

IV. PERFORMANCE RESULTS

Fig. 3 shows the Achievable Information Rate (AIR), which is the maximum data rate in "bits/channel use" that could be achieved with ideal coding. The data rate upper bound is 128 bits/channel use for a 32×32 SU-MIMO with a 16 QAM constellation. The dashed lines represent the AIR of several detectors using BICM (without turbo-iteration), while the solid lines represent the BICM-ID (with turbo-iteration). SEP [5] (resp. QRD SEP [9]) cannot achieve the maximum data rate and is stuck at 50 (resp. 78) bits/channel use with BICM, whereas it can reach 93 (resp. 110) bits/channel use with BICM-ID. QRD PIC BP [13] can achieve the same performance as QRD SEP with BICM and saturate at 102 bits/channel use with BICM-ID. The proposed QRD SSEP can achieve almost the same AIR as VEP [6] and DEP [7] with BICM, but it cannot reach the maximum data rate and is stuck at 120 bits/channel use while VEP and DEP achieve it. The gain of AIR between the proposed algorithm and SEP / ORD SEP is visible starting from code rates R = 1/3 and higher for BICM, and from rate 1/2 and higher for BICM-ID.

Fig. 4 shows the Signal-to-Noise Ratio (SNR) needed to achieve a Frame Error Rate (FER) of 10⁻² given a computational complexity for each algorithm through turbo-iteration. The SU-MIMO scenarios studied are a 32×32 MIMO with a 16 QAM, shown with "+" markers, and a 10×10 MIMO with an 8 PSK, shown with "o" markers. The SNR is defined as $\frac{E_s}{N_0}$ with $E_s = N_t P_s$ and P_s the energy of one QAM symbol, the constellation power is normalized to 1. 5G compliant LDPC of length $K_b = 3840$ and of rate 1/2 and 2/3 are used. The LDPC decoder makes five inner iterations per turbo iteration, and the internal messages of the decoder are kept and reused as initialization messages for the next turbo iteration. A maximum of 9 turbo iterations (i.e. ten decoder calls) brings the overall LDPC iterations to 50. The y-axis measures the detection complexity for a given number of turbo iterations (each point of a curve is associated with a turbo iteration), starting from T=0 with the lowest and least complex point, to T=9 with the highest and most complex point. All detectors use only L=1 auto-iteration per turbo-iteration.

With the 10×10 MIMO scenario with R=1/2, the proposed algorithm, QRD SSEP, has a better performance-to-complexity trade-off than VEP/DEP for a fixed number of turbo-iterations as its performance is similar to VEP/DEP. Compared to SEP and QRD SEP, it has a far better performance-to-complexity trade-off as it can achieve better performance with fewer turbo iterations. When using an LDPC

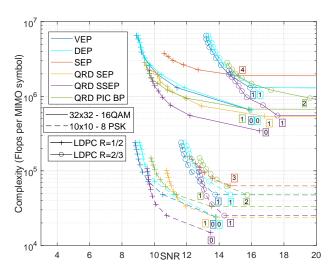


Fig. 4. Performance vs. complexity in flops per MIMO symbol for two ST-BICM-ID SU-MIMO scenarios. Numbers show the first turbo iteration number T per detector, which allows reaching the target FER of 10^{-2} .

code of R=2/3, the same behavior occurs between QRD SSEP and VEP/DEP, while SEP and QRD SEP are not able at all to achieve a FER of 10^{-2} , so there are no curves representing them.

For the 32×32 MIMO with a 16 QAM scenario with R=1/2, QRD SSEP is always less complex than VEP/DEP. It achieves a FER of 10^{-2} with an SNR slightly higher than its competitor. With a few turbo iterations, QRD SSEP performs significantly better than QRD SEP and SEP, and performs slightly better with nine turbo iterations. Like for the previous scenario, with R=2/3, SEP and QRD SEP cannot achieve a FER of 10^{-2} , so there are no curves for them. VEP/DEP can perform better than QRD SSEP, but with a few turbo iteration QRD SSEP performs better for a fixed complexity.

Fig. 5 shows the performance of the algorithms for a BICM-ID 4×8 MU-MIMO with a 64QAM constellation. Only the first, second, and ninth turbo iterations are presented. As in SU-MIMO, SEP performs poorly without many turbo-iterations (Turbo 0 and 1), while QRD SEP performs better. But it is still far from the VEP performance. The proposed SSEP QRD achieves the same performance as VEP for each turbo-iteration, but with a smaller complexity, as shown in Tab. I. The proposed SSEP QRD SIC outperforms VEP with a smaller complexity with a few turbo iterations (Turbo 0 and 1). With nine turbo-iterations, SSEP QRD SIC is still the best-performing MU-MIMO detector, but VEP is very close, as well as SSEP QRD. SEP and SEP QRD are not so far, but they still lack behind by 0.25-0.5dB.

V. CONCLUSION

In this article, we proposed two new algorithms, SSEP QRD for SU-MIMO and MU-MIMO and SSEP QRD SIC for MU-MIMO, that can achieve the same or better performance with respect to VEP and DEP. They both have similar low complexity, lower than VEP/DEP, since they do not require a costly matrix inversion. This makes them great candidates for efficient MIMO and mMIMO detection. This proposal is an additional step toward filling the gap between vector and scalar approaches of EP while keeping a low complexity.

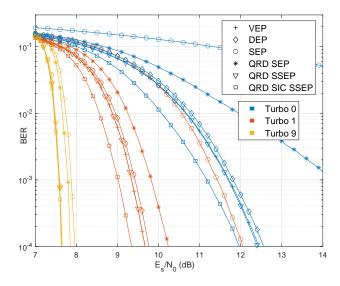


Fig. 5. Performance of BICM-ID MU-MIMO detection with 4×8 MIMO 64QAM on a Rayleigh channel using a LDPC ($K_b = 3840, R = 1/2$).

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