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A linear perspective on cut-elimination for non-wellfounded sequent calculi with least and greatest fixed points (extended version)

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Abstract. This paper establishes cut-elimination for μLL^∞ , μLK^∞ and μLJ^∞ , that are non-wellfounded sequent calculi with least and greatest fixed-points, by expanding on prior works by Santocanale and Fortier [20] as well as Baelde et al [4,3]. The paper studies a fixed-point encoding of LL exponentials in order to deduce those cut-elimination results from that of μMALL^∞ . Cut-elimination for μLK^∞ and μLJ^∞ is obtained by developing appropriate linear decorations for those logics.

Keywords: LL , μ -calculus, non-wellfounded proofs, cut elimination

1 Introduction

On the non-wellfounded proof-theory of fixed-point logics. In the context of logics with induction and coinduction (such as logics with inductive definitions *à la* Martin L of [6,9,10,25], or variants of the μ -calculus [22,23,11]), the need for a (co)inductive invariant (in the form of the Park’s rule for induction) is replaced by the ability to pursue the proof infinitely, admitting non-wellfounded branches, when considering non-wellfounded and circular proofs (also called cyclic, or regular proofs, since the proof tree is a regular tree, with finitely many distinct subtrees). In such frameworks, sequent proofs may be finitely branching but non-wellfounded derivation trees and infinite branches shall satisfy some validity condition. (Otherwise one could derive any judgement, see Figure 1(a).) Various validity conditions have been considered in the literature [3].

The non-wellfounded and circular proof-theory of fixed-points attracted a growing attention first motivated by proof-search [18,17,1,16,29,7,8] and more recently by a Curry-Howard perspective, studying the dynamics of the cut-elimination in those logics [4,20,30] where formulas correspond to (co)inductive types. Notice also that when interested in the computational content of proofs, we will not focus solely on the regular fragment as we expect, for instance, that we can write a regular program that computes a non-ultimately periodic stream.

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$$\begin{array}{c}
\vdots \\
\frac{}{\vdash \Gamma, \mu X.X} (\mu) \quad \frac{}{\vdash \nu X.X, \Delta} (\nu) \\
\frac{}{\vdash \Gamma, \mu X.X} (\mu) \quad \frac{}{\vdash \nu X.X, \Delta} (\nu) \\
\hline
\vdash \Gamma, \Delta \quad (\text{Cut})
\end{array}
\quad (a)$$

$$\begin{array}{c}
\vdash \Gamma, C \quad \vdash C^\perp, \Delta, D \quad \vdash D^\perp, \Sigma \\
\hline
\vdash \Gamma, \Delta, \Sigma \quad \text{mcut}(\perp, \perp)
\end{array}
\quad (b)$$

Fig. 1. (a) Example of an invalid circular pre-proof. (b) Schema of the multicut rule.

Cut-elimination and LL. When studying the structure of proofs and their cut-elimination properties, LL, Girard’s Linear Logic [21], is a logic of choice: the careful treatment of structural rules gives access to a lot of information and a fine-grained control over cut-reduction. The constrained use of structural rules indeed renders the cut-elimination theorem more informative than in LJ and of course LK. Interestingly it provided a positive feedback on the understanding of LJ and LK: by decorating intuitionistic and classical proofs with enough exponential modalities ($!$, $?$), they can become LL proofs and one can therefore refine the original cut-elimination relations [21,13]. This approach impacted the understanding of evaluation strategies of programming languages such as call-by-name and call-by-value notably. Another way to view this is by noting that in, say, LK the additive and multiplicative presentation of conjunction (resp. disjunction) can be shown to be interderivable thanks to structural rules. This fails in LL and it is the reason why LL has well-established additive – $\oplus, \&, \top, 0$ – (resp. multiplicative $\wp, \otimes, \perp, 1$) *fragments*. It is the role of the exponential fragment to relate the additive and multiplicative worlds, by mean of the fundamental equivalence: $!A \otimes !B \dashv\vdash !(A \& B)$ (and its dual, $?A \wp ?B \dashv\vdash ?(A \oplus B)$). The exponential modalities are precisely introduced where structural rules are needed to restore the equivalence between the additive and multiplicative conjunctions; in categorical models of LL [27], this principle is referred to as Seely isomorphisms.

Cut-elimination for non-wellfounded proofs. Proving cut-elimination results for non-wellfounded proofs in the presence of least and greatest fixed-points requires to use reasoning techniques coping with the non-inductive structure of the considered formulas (fixed-points formulas regenerate) and proof objects (which are non-wellfounded). For instance, Santocanale and Fortier [20] proved cut-elimination for the regular fragment of non-wellfounded proofs of purely additive linear logic with fixed points, μ ALL while Baelde *et al.* [4] proved cut-elimination for non-wellfounded proofs with additive and multiplicative connectives, μ MALL. In both cases, the proof relies on a generalization of the cut-rule (which abstracts a portion of a proof tree constituted only of cut inferences, the *multicut* rule, see Figure 1(b)) and on a reasoning by contradiction to prove that one can eliminate cuts at the limit of an infinite cut reduction sequence, while preserving the validity condition. Baelde *et al.* [4,3] use a so-called “locative” approach by modelling sequents as sets of formulas paired with addresses which determines uniquely the formula occurrence in a sequent and makes explicit the ancestor relation used to trace the progress along branches. Moreover, the cut-elimination

proof proceeds by a rather complex semantical, indirect, argument relying on a soundness theorem.

In a slightly different direction, Das and Pous [15] proved a cut-elimination result for Kleene algebras and their variants. This can be viewed as a non-commutative version of intuitionistic MALL with a particular form of inductive construction, Kleene’s star. Kuperberg et al [24] and more specifically Pinault’s PhD thesis [28] as well as Das [14] examine non-wellfounded versions of System T based on [15], exploring the computational content of non-wellfounded proofs.

Neither Santocanale and Fortier’s [30,20], nor Baelde et al. [4,3] works captured full linear logic: the exponentials are missing and the proofs cannot deal with them in a simple way. Indeed, the proof for μALL strongly relies on the assumption the sequents are pairs of formulas $(A \vdash B)$ while in μMALL , the locative approach taken by Baelde et al. is not well-suited to work with structural rules: the extension of the proof would be possible though highly technical. In contrast, our motto in the present work is to work with traditional sequents as lists of formulas and to exploit the (co)inductive nature of LL exponentials.

On the (co)inductive nature of exponential modalities in linear logic. The original works by Baelde and Miller on fixed-points in linear logic [5,2] focus on μMALL only and present an encoding of the exponential modalities of LL using least and greatest fixed points. Indeed, the $?$ and $!$ modalities have an infinitary character which is well-known from the early days of linear logic (see Section V.5 of Girard’s seminal paper [21]) and which is in fact respectively inductive for $?$ and coinductive for $!$; let us discuss it briefly here.

One can decide to contract a $?$ -statement any *finite* number of times before it is ultimately weakened or derelicted. It is therefore natural to represent $?A$ with formula $?^{\bullet}A = \mu X.A \oplus (\perp \oplus (X \wp X))$: A allows for dereliction, \perp for weakening and $X \wp X$ will regenerate, by unfolding, two copies of $?^{\bullet}A$, making the contraction derivable. The \oplus and μ connectives respectively provide the ability to choose either of those three inferences and to repeat finitely this process.

On the other hand, a $!$ -formula is a formula which, during cut-elimination, shall maintain a proper interaction with *any number* of contractions, weakenings or derelictions: a proof concluded with a promotion shall be able to react to any number of duplications or erasure before the promotion actually interact with a dereliction to open the *exponential box*: from that follows the coinductive character of $!A$ modelled as $!^{\bullet}A = \nu X.A \& (1 \& (X \otimes X))$.

As informally described above and formally established by Baelde and Miller [5], the exponential rules can be derived in the finitary sequent calculus μMALL : to any LL provable sequent can be associated a provable μMALL sequent via the above translations of the exponentials. However, until now one can hardly say more about this embedding for two deep reasons: (i) the fundamental Seely isomorphisms which relate the additive and multiplicative versions of conjunction (resp. disjunction) are still derivable through this encoding but they are no more isomorphisms and (ii) it is unknown whether the converse of the above provability preservation result holds: whether the μMALL provability of the translation of an LL sequent s entails the LL provability of s itself. A contribution of the

Several appendices are provided for the reviewers, at their discretion.

It details some definitions and provides proofs which, most often, could not be included in the body of the paper due to lack of space. A companion technical report is available online [31]. References to the Appendices are presented in such margin paragraphs.

present paper is to put to work Baelde and Miller’s encoding, showing that, in the case of non-wellfounded proofs, its structure is faithful enough to extract information of the cut-reduction behaviour of the logic.

Contributions and organization of the paper. The main result of this paper is a cut-elimination theorem for μLL^∞ , the non-wellfounded sequent calculus for linear logic extended with least and greatest fixed points. Our proof proceeds by encoding LL exponentials in μMALL^∞ and studying μLL^∞ cut-reduction sequences through their simulation in μMALL^∞ which may be a *transfinite* sequence. In Section 2, we introduce our logics, μMALL^∞ , μLL^∞ , μLK^∞ and μLJ^∞ , altogether with their non-wellfounded proofs and validity conditions. We adapt μMALL^∞ cut-elimination theorem [4] to our setting where sequents are lists and prove a compression lemma for μMALL^∞ transfinite cut-reduction sequences. Section 3 constitutes the core of our paper: we define μLL^∞ cut-reduction rules, study the encoding of exponentials in μMALL^∞ and show that μLL^∞ cut-reduction steps can be simulated in μMALL^∞ , before proving μLL^∞ cut-elimination theorem. We prove in Section 4, as corollaries, cut-elimination for μLK^∞ and μLJ^∞ , the non-wellfounded sequent-calculi for classical and intuitionistic logic. While our result for μLL^∞ shows that any fair cut-reduction sequence produces a cut-free valid proof, our two other cut-elimination results are truly (infinitary) weak-normalization results. We finally conclude in Section 5 with perspectives. A major advantage of our approach is that μMALL^∞ cut-elimination proof and, to some extent, the validity conditions, are regarded as black box, simplifying the presentation of the proof and making it reusable wrt. other validity condition or μMALL^∞ proof techniques. An additional by-product of our approach, to the theory of linear logic, is to illustrate the fact that Seely isomorphisms are not needed to reach a cut-free proof.

2 Non-wellfounded proofs: μMALL^∞ , μLL^∞ , μLK^∞ , μLJ^∞ .

2.1 μ -signatures and formulas.

Definition 1 (μ -signature). A μ -signature is a set \mathcal{C} of pairs (c, p) of a connective symbol c and a tuple p of elements of $\{+, -\}$. The arity of c , $\text{ar}(c)$, is the length of p , while the elements of p indicate the mono/antitonicity of the connective in the given component. The empty tuple will be denoted as $()$.¹

Example 2 (μ -signature associated with μMALL , μLL , μLK , μLJ). The μ -signature associated with μMALL , μLL , μLK , μLJ , considered in this paper are:

- μMALL signature: $\mathcal{C}_{\mu\text{MALL}} = \{\wp, \otimes, \oplus, \&\} \times \{(+, +)\} \cup \{0, 1, \top, \perp\} \times \{()\}$;
- one-sided μLL signature: $\mathcal{C}_{\mu\text{LL}_1} = \mathcal{C}_{\mu\text{MALL}} \cup \{!, ?\} \times \{(+)\}$;
- two-sided μLL signature: $\mathcal{C}_{\mu\text{LL}_2} = \mathcal{C}_{\mu\text{LL}_1} \cup \{(-\circ, (-, +)), (-^\perp, (-))\}$;

¹ μ -signature can be enriched to consider quantifiers but we restrict to the propositional case here.

- μLK signature: $\mathcal{C}_{\mu\text{LK}} = \{\wedge, \vee\} \times \{(+, +)\} \cup \{(\Rightarrow, (-, +))\} \cup \{\top, \text{F}\} \times \{()\}$;
- μLJ signature: $\mathcal{C}_{\mu\text{LJ}} = \mathcal{C}_{\mu\text{LK}}$.

Definition 3 (Pre-formulas). Given a μ -signature \mathcal{C} , a countable set \mathcal{V} of fixed-point variables and a set of atomic formulas \mathcal{A} , the set of **pre-formulas** over \mathcal{S} is defined as the least set $\mathcal{F}_{\mathcal{S}}$ such that: (α) $\mathcal{A} \cup \mathcal{V} \subseteq \mathcal{F}_{\mathcal{S}}$; (β) for every c of arity n in \mathcal{C} and $F_1, \dots, F_n \in \mathcal{F}_{\mathcal{S}}$, $c(F_1, \dots, F_n) \in \mathcal{F}_{\mathcal{S}}$; (γ) for every $X \in \mathcal{V}$ and pre-formula $F \in \mathcal{F}_{\mathcal{S}}$, $\mu X.F \in \mathcal{F}_{\mathcal{S}}$ and $\nu X.F \in \mathcal{F}_{\mathcal{S}}$.

Definition 4 (Positive and negative occurrences of a variable). Given a μ -signature \mathcal{C} and a fixed-point variable $X \in \mathcal{V}$, one defines by induction on pre-formulas the fact, for X , to occur positively (resp. negatively) in a pre-formula : (α) X occurs positively in X ; (β) X occurs positively (resp. negatively) in $c(F_1, \dots, F_n)$, for $(c, p) \in \mathcal{C}$, if there is some $1 \leq i \leq n$ such that X occurs positively (resp. negatively) in F_i and $p_i = +$ or there is some $1 \leq i \leq n$ such that X occurs negatively (resp. positively) in F_i and $p_i = -$; (γ) X occurs positively (resp. negatively) in $\sigma Y.F$, for $\sigma \in \{\mu, \nu\}$, if $Y \neq X$ and X occurs positively (resp. negatively) in F .

Definition 5 (μ -formula). A μ -formula F over a signature \mathcal{S} is a pre-formula containing no free fixed-point variable and such that for any sub-pre-formula of F of the form $\sigma X.G$, all occurrences of X in G are positive.

Definition 6. One-sided μLL formulas are those formulas defined over the signature $\mathcal{C}_{\mu\text{LL}_1}$ together with a set of atomic formulas $\{a, a^\perp \mid a \in \mathcal{A}\}$ for a countable set \mathcal{A} . Negation $(_)^\perp$ is the involution on pre-formulas defined by:

$$(a^\perp)^\perp = a; \perp^\perp = 1; \top^\perp = 0; (F \wp G)^\perp = F^\perp \otimes G^\perp; (F \oplus G)^\perp = F^\perp \& G^\perp; (? F)^\perp = ! F^\perp; X^\perp = X; (\nu X.F)^\perp = \mu X.F^\perp.$$

Definition 7 (μ -Fischer-Ladner subformulas). Given a μ -signature \mathcal{C} and a μ -formula F , $FL(F)$ is the least set of formulas such that:

- $F \in FL(F)$;
- $c(F_1, \dots, F_n) \in FL(F) \Rightarrow F_1, \dots, F_n \in FL(F)$ for $c \in \mathcal{C}$;
- $\sigma X.B \in FL(F) \Rightarrow B[\sigma X.B/X] \in FL(F)$ for $\sigma \in \{\mu, \nu\}$.

Example 8. Let us consider $F = \nu X.((a \wp a^\perp) \otimes (!X \otimes \mu Y.X))$. $FL(F)$ is the set $\{F, (a \wp a^\perp) \otimes (!F \otimes \mu Y.F), a \wp a^\perp, a, a^\perp, !F \otimes \mu Y.F, !F, \mu Y.F\}$.

The finiteness of $FL(F)$ makes it an adequate notion of subformula:

Proposition 9. For any μ -signature \mathcal{S} and μ -formula F , $FL(F)$ is finite.

Appendix B.1 provides more details.

See proof in Doumane's PhD thesis [18].

2.2 μMALL^∞ , μLL^∞ , μLK^∞ & μLJ^∞ inference rules.

Now, we define the inference rules associated with the above μ -signatures.

$$\begin{array}{c}
\text{(a)} \quad \frac{}{\vdash F, F^\perp} \text{ (Ax)} \quad \frac{\vdash \Gamma, F \quad \vdash F^\perp, \Delta}{\vdash \Gamma, \Delta} \text{ (Cut)} \quad \frac{\vdash \Gamma, G, F, \Delta}{\vdash \Gamma, F, G, \Delta} \text{ (X)} \\
\frac{\vdash F, G, \Gamma}{\vdash F \otimes G, \Gamma} \text{ (\otimes)} \quad \frac{\vdash F, \Gamma \quad \vdash G, \Delta}{\vdash F \otimes G, \Gamma, \Delta} \text{ (\otimes)} \quad \frac{\vdash \Gamma}{\vdash \perp, \Gamma} \text{ (\perp)} \quad \frac{}{\vdash \top} \text{ (\top)} \\
\frac{\vdash F, \Gamma \quad \vdash G, \Gamma}{\vdash F \otimes G, \Gamma} \text{ (\&)} \quad \frac{\vdash A_i, \Gamma}{\vdash A_1 \oplus A_2, \Gamma} \text{ (\oplus^i)} \quad \frac{}{\vdash \top, \Gamma} \text{ (\top)} \quad \text{(no rule for 0)} \\
\frac{\vdash G[\nu X.G/X], \Gamma}{\vdash \nu X.G, \Gamma} \text{ (\nu)} \quad \frac{\vdash F[\mu X.F/X], \Gamma}{\vdash \mu X.F, \Gamma} \text{ (\mu)} \\
\hline
\text{(b)} \quad \frac{\vdash F, \Gamma}{\vdash ?F, \Gamma} \text{ (?d)} \quad \frac{\vdash F, ?\Gamma}{\vdash !F, ?\Gamma} \text{ (!p)} \quad \frac{\vdash \Gamma}{\vdash ?F, \Gamma} \text{ (?w)} \quad \frac{\vdash ?F, ?\Gamma, \Gamma}{\vdash ?F, \Gamma} \text{ (?c)}
\end{array}$$

Fig. 2. (a) μMALL^∞ Inferences; (b) μLL^∞ Exponential Inferences

Definition 10 (Sequents and inferences). A sequent $s = \Gamma \vdash \Delta$ over a μ -signature \mathcal{S} is a pair of finite lists Γ, Δ of \mathcal{S} -formulas: Γ is the **antecedent** and Δ the **succedent**. An inference rule r , usually presented by a schema, is the data of a **conclusion sequent**, **premise sequents**, together with an **ancestor relation** relating formulas of the conclusion with formulas of the premises. A rule has a subset of distinguished **principal formulas** of the conclusion.

Appendix B.2 provides more details.

Convention 1 In the following, the ancestor relation will be depicted as colored lines joining related formulas. The **principal** formulas of an inference are the formulas which are explicitly spelled out in the conclusion sequent of an inference, not described via a context meta-variable. A formula occurrence of an inference is said to be **active** if it is principal or related to a principal formula by the ancestor relation. We will freely use the derived rules obtained by **pre- and post-composition with the exchange rule**, adapting the ancestry relation accordingly. Finally, for one-sided sequent calculi with an involutive negation \cdot^\perp , we may write $\Gamma \vdash \Delta$ for sequents $\vdash \Gamma^\perp, \Delta$ to clarify the computational behaviour of our examples (keeping the rule names unchanged).

Definition 11 ($\mu\text{MALL}^\infty, \mu\text{LL}^\infty, \mu\text{LK}^\infty, \mu\text{LJ}^\infty$). μMALL^∞ inferences are given in Figure 2. Those for one-sided μLL^∞ in Fig. 2(a) and 2(b). Those for μLK^∞ in Fig. 3. Those for μLJ^∞ by considering only inference from Fig. 3 where the succedent of both premises and conclusion sequents are singletons.

Appendix B.3 provides more details on μLJ^∞ .

In the above sequent calculi, every inference but the cut satisfies the subformula property wrt. FL-subformulae. The 2-sided μLL^∞ sequent calculus, over $\mathcal{C}_{\mu\text{LL}_2}$, is defined as usual and not recalled here for space constraints.

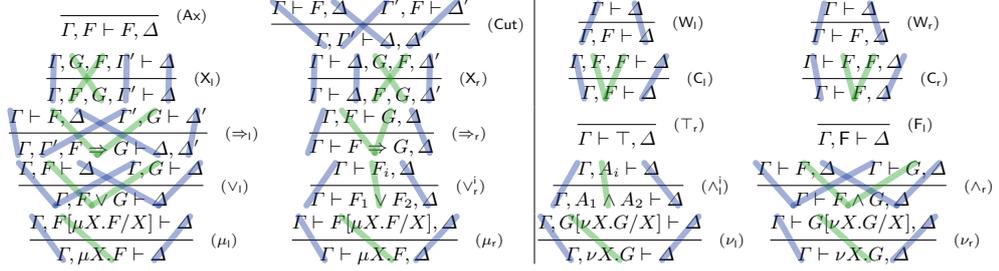


Fig. 3. μLK^∞ Two-sided Inferences

2.3 Pre-proofs and validity conditions.

Definition 12 (Pre-proofs). *The set $\mathcal{P}_{\mathcal{S}, \mathcal{I}}$ of \mathcal{I} -pre-proofs associated to some of the above μ -signatures \mathcal{S} and sets of inferences \mathcal{I} is the set of **finite or infinite trees** whose nodes are correctly labelled with inferences and sequents.*

See details in appendix B.4.

Pre-proofs are equipped with a metric structure as follows: we define a **distance** $d : \mathcal{P}_{\mathcal{S}, \mathcal{I}} \times \mathcal{P}_{\mathcal{S}, \mathcal{I}} \rightarrow \mathbb{R}$ as: $d(\pi, \pi') = 0$ if $\pi = \pi'$ and $d(\pi, \pi') = 2^{-k}$ where k is the length of the shortest position where π and π' differ otherwise.

Example 13. Consider μLJ formulas $N = \mu X. \top \vee X$ and $S = \nu X. N \wedge X$. They represent nats and streams of nats. The μLJ^∞ derivations of Figure 4 respectively represent natural numbers, successor function, $n :: n + 1 :: n + 2 :: \dots$, the double functions and the function that builds a stream enumerating the natural numbers from its input: the cut-elimination process considered below will ensure that cutting π_k with π_{enum} will infinitarily reduce to π_{from}^k . Figure 5 shows other examples of μLL^∞ pre-proofs, discussed with the validity condition.

The back-edge arrow to a lower sequent is a fixed-point definition of the proof object: the subproof rooted in the source is equal to the proof rooted in the target. This fixed-point definition trivially has a unique solution.

In the following, we assume given a μ -signature \mathcal{S} and a sequent calculus \mathcal{S} for this signature and we shall define the valid \mathcal{S} -proofs as a subset of \mathcal{S} -pre-proofs, by introduction a **thread-based validity condition**.

More details provided in appendix B.4.

Definition 14 (Thread and validity). *Given a pre-proof π and an infinite branch $\beta = (s_i)_{i \in \omega}$ in π , a **thread** for β is an infinite sequence θ of formula occurrences such that $\forall i \in \omega$, θ_i is a formula occurrence of s_i and θ_i and θ_{i+1} are ancestor of each other. θ is said to **support** β .*

A formula F is **recurring** in a thread θ of β if there are infinitely many i such that θ_i is an occurrence of F .

A thread θ is **valid** if it contains infinitely often the principal formula (occurrence) of a ν or μ rule and if the set of recurring formulas of θ has a least element (for the usual subformula ordering) which is (i) a ν formula when the least element occurs in the succedents or (ii) a μ formula if it occurs in the antecedents. A pre-proof is **valid** if all its infinite branches have a suffix supported by a valid thread.

$$\begin{array}{c}
\pi_0 = \frac{\overline{\vdash \top} \quad (\top)}{\vdash \top \vee N} \quad (\vee_l) \quad \pi_{k+1} = \frac{\overline{\pi_k}}{\vdash \top \vee N} \quad (\vee_r^2) \quad \pi_{\text{succ}} = \frac{\overline{N \vdash N} \quad (\text{Ax})}{N \vdash \top \vee N} \quad (\vee_r^2) \quad \pi_{\text{from}}^n = \frac{\overline{\pi_n} \quad \overline{\pi_{\text{from}}^{n+1}}}{N \wedge S} \quad (\wedge_r) \\
\vdash N \\
\pi_{\text{double}} = \frac{\overline{\pi_0}}{\top \vdash N} \quad (\top_l) \quad \frac{\overline{N \vdash N} \quad (\vee_r^2) \quad \overline{N \vdash \top \vee N} \quad (\mu_r) \quad \overline{N \vdash N} \quad (\vee_r^2) \quad \overline{N \vdash \top \vee N} \quad (\mu_r) \quad \overline{N \vdash N} \quad (\vee_l)}{\top \vee N \vdash N} \quad (\mu_l) \quad \pi_{\text{enum}} = \frac{\overline{N \vdash N} \quad (\text{Ax}) \quad \overline{\pi_{\text{succ}}} \quad \overline{N \vdash S}}{N \vdash N \wedge S} \quad (\wedge_r) \quad \frac{\overline{N \vdash S} \quad (\text{Cut})}{N \vdash S} \quad (\nu) \quad \frac{\overline{N \vdash S}}{\vdash N \Rightarrow S} \quad (\Rightarrow_r) \\
N \vdash N
\end{array}$$

Fig. 4. Examples of μLJ^∞ pre-proofs.

Example 15 ((Non-)valid pre-proofs). Consider the pre-proof in Fig. 5(a), with $F = \nu X.((a\wp a^\perp) \otimes (!X \otimes \mu Y.X))$ and $G = \mu Y.F$. The rightmost branch is supported by the green thread for which the least recurring formula is F , a ν -formula. All other branches are valid: this pre-proof is valid. Consider now the same pre-proof but with $F = \nu X.((a\wp a^\perp) \otimes (!X \otimes G))$ and $G = \mu Y.\nu X.((a\wp a^\perp) \otimes (!X \otimes Y))$. G is now a subformula of F and G , a μ -formula, and becomes the least recurring formula of all threads along the right-most infinite branch. This branch is invalid: the pre-proof is not a proof. Examples of μLL^∞ invalid pre-proofs are given in Fig. 1(a),5(b-c). In Fig. 4, π_{double} has a left thread on N while $\pi_{\text{from}}^n, \pi_{\text{enum}}$ have right threads on S : they are valid.

2.4 Non-locative μMALL^∞ cut-elimination theorem

The validity condition defines a subset of pre-proofs, ensuring good properties for those non-wellfounded derivations that satisfy the validity condition. In this paper, we will mainly be interested in cut-elimination theorem, which was proved for μMALL^∞ [4] and that we review in this subsection. In [4], a somehow stronger result than cut-elimination is proved: infinitary strong normalization with respect to the class of *fair* reduction sequences.

The only new result developed in this subsection is the lifting of the occurrence-based cut-elimination result of [4] to our setting system, for which we first introduce the multicut inference and review the main multicut-reduction steps for μMALL^∞ before defining fair reductions. The cut-elimination results of [4,20] do not rewrite cuts, *per se*, but subtrees of cuts in the form of an abstraction called **multicut** which is a variable arity inference defined as follows:

See details in appendix B.5.

Definition 16. *The multicut inference is given by the data of (i) a conclusion sequent s , (ii) a non-empty list of premises $(s_1, \dots, s_n), n \geq 1$, (iii) an ancestor relation ι which is an injective map from the conclusion formulas to the premise formulas and relates identical formulas and additionally (iv) a cut-connectedness relation $\perp\!\!\!\perp$ which is a total, symmetric, binary relation among the formula occurrences of the premises which are not ancestor of a conclusion*

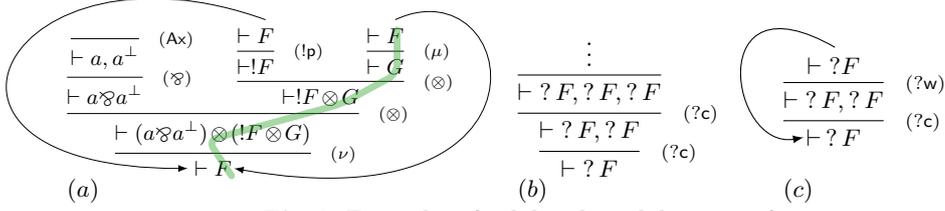


Fig. 5. Examples of valid and invalid pre-proofs.

formula, which relates dual formulas² and which satisfies a connectedness and acyclicity condition (see [3,4]). The multicut inference has no principal formula.

We write this multicut rule as:
$$\frac{s_1 \quad \dots \quad s_n}{s} \text{mcut}(\perp, \perp).$$

In the following, we only consider μMALL^∞ pre-proofs with specific multicuts:

Definition 17 (μMALL_m^∞). μMALL_m^∞ (pre)proofs are those (pre)proofs built from μMALL^∞ inferences and the multicut, such that (i) any branch contains at most one multicut and (ii) any occurrence of a cut is above a multicut inference.

In the following, we shall always assume, even without mentioning it, that we consider proofs in μMALL_m^∞ (as well as μLL_m^∞ , μLJ_m^∞ , μLK_m^∞). We need the following definition (from [4]), identifying the premises of an mcut which are cut-connected to a given formula occurrence:

Definition 18 (Restriction of a mcut-context). Consider an occurrence of a mcut $\frac{s_1 \quad \dots \quad s_n}{s} \text{mcut}(\perp, \perp)$ and assume s_i to be $\vdash F_1, \dots, F_k$. We define \mathcal{C}_{F_j} , $1 \leq j \leq k$, to be the least set of sequent occurrences contained in $\{s_1, \dots, s_n\}$ such that:

- (i) If $\exists k, l$ such that $(k, l) \perp\!\!\!\perp (i, j)$, then $s_k \in \mathcal{C}_{F_j}$;
 - (ii) for any $k, k' \neq i$, if $s_k \in \mathcal{C}_{F_j}$ and $\exists l, l'$ such that $(k, l) \perp\!\!\!\perp (k', l')$, then $s_{k'} \in \mathcal{C}_{F_j}$.
- We define $\mathcal{C}_\emptyset = \emptyset$ and $\mathcal{C}_{F, \Gamma} = \mathcal{C}_F \cup \mathcal{C}_\Gamma$.

When relating μLL^∞ and μMALL^∞ mcut-sequences below, we shall consider not only finite sequence nor ω -indexed sequences but also transfinite sequences. Those are sequences of triples of a proof, a redex and the position of the redex in the proof tree. A position p has a **depth** $\text{dpth}(p)$ which is its length.

Definition 19 (mcut-reduction rules, transfinite sequences). μMALL^∞ mcut-reduction sequences are directly adapted from [3,4]. Given an ordinal λ , a **transfinite reduction sequence** of length λ , or λTRS , is a λ -indexed sequence $(\pi_i, r_i, p_i)_{i \in \lambda}$ such that $\pi_i \xrightarrow{p_i} \pi_{i+1}$, for any i such that $i+1 \in \lambda$, where the reduction occurs at position p_i reducing mcut-redex r_i .

Definition 20 (Weak and strong convergence). A (transfinite) mcut reduction sequence $(\pi_i, r_i, p_i)_{i \in \alpha}$ is **weakly converging** if for any limit ordinal $\beta \in \alpha$, $\lim(\pi_i)_{i \in \beta} = \pi_\beta$. $(\pi_i, r_i, p_i)_{i \in \alpha}$ is **strongly converging** if it is weakly converging and moreover for any limit ordinal $\beta \in \alpha$, $\lim(\text{dpth}(p_i))_{i \in \beta} = +\infty$.

² When working with two-sided sequents, $\perp\!\!\!\perp$ will relate identical formulas, one in a succedent, the other in an antecedent.

See details in appendix B.6.

See details in appendix A.

Remark 21. The cut-reduction rules preserve the property that every branch of a proof has at most one multicut inference: μMALL_m^∞ is closed by cut-reduction.

See details in appendix B.8.

A μMALL_m^∞ pre-proof π may contain multiple cut-redexes: $\pi \xrightarrow{r_1^{p_1}} \pi_1$ and $\pi \xrightarrow{r_2^{p_2}} \pi_2$. As usual, a notion of residual associates to (r_1, p_1) , a set of redexes of π_2 , $(r_1, p_1)/(r_2, p_2)$ which is generalized to reduction sequences: $(r_1, p_1)/\sigma$.

Definition 22 (Fair reduction sequences). A reduction sequence $(\pi_i, r_i, p_i)_{i \in \omega}$ is **fair** if for all $i \in \omega$ and r, p such that $\pi_i \xrightarrow{r, p} \pi'$ there is some $j \geq i$ such that π_j does not contain a residual of (r, p) anymore.

See proof in appendix B.7.

Theorem 23. Every fair mcut-reduction sequence of μMALL^∞ valid proofs of $\vdash \Gamma$ (strongly) converges to a cut-free valid proof of $\vdash \Gamma$.

2.5 Compressing transfinite μMALL^∞ cut-reduction sequences

In the previous paragraph, we introduced not only ω -indexed sequences, but transfinite μMALL^∞ cut-reduction sequences as we shall need reduction beyond ω when simulating μLL^∞ cut-elimination in μMALL^∞ . We shall now prove that a class of transfinite μMALL^∞ mcut-reduction sequences can be compressed to ωTRS . This result can be viewed as adapting to our setting the compression lemma from infinitary rewriting [32], even though we require more on the structure of the compressed sequences as it will be useful to establish μLL^∞ cut-elimination.

See details on infinitary rewriting in appendix A.

Definition 24 (Depth-increasing). A μMALL^∞ cut reduction sequence $\sigma = (\pi_i, r_i, p_i)_{i \in \omega}$ is **depth-increasing** if $(\text{dpth}(p_i))_{i \in \omega}$ is (weakly) increasing.

Definition 25 (Reordering). An mcut reduction sequence $\sigma = (\pi_i, r_i, p_i)_{i \in \alpha}$ is a reordering of $\sigma' = (\pi'_i, r'_i, p'_i)_{i \in \beta}$ if there is a bijection o between α and β such that for any $i \in \alpha$, $(r'_{o(i)}, p'_{o(i)}) = (r_i, p_i)$.

Proposition 26 (Compression lemma). Let $\sigma = (\pi_i, r_i, p_i)_{i \in \alpha}$ be a strongly converging μMALL^∞ transfinite cut-reduction sequence. There exists a μMALL^∞ cut-reduction sequence $\text{Comp}(\sigma) = (\pi'_i, r'_i, p'_i)_{i \in \beta}$ which is a reordering of σ , depth-increasing, strongly converging with the same limit as σ and such that $\beta = \alpha$ if α is finite and $\beta = \omega$ otherwise.

See proof in App B.9.

3 Cut-elimination theorem for μLL^∞

The aim of this section is to prove the following theorem:

Theorem 27. For any valid μLL^∞ proof π , fair μLL^∞ mcut-sequences from π converge to cut-free μLL^∞ proofs.

The idea of the proof and outline of the present section are as follows:

1. We shall first define the cut-reduction rules for μLL^∞ by extending μMALL^∞ *multicut*-reduction with rules for reducing exponential cuts.
 2. We then encode exponentials with fixed-points and translate μLL^∞ sequents (resp. pre-proofs) into μMALL^∞ , preserving validity both ways.
 3. We will then simulate μLL^∞ reductions in μMALL^∞ : a single μLL^∞ step may require an infinite, or even transfinite, μMALL^∞ mcut-reduction sequence.
 4. Finally, we will study the simulation of fair μLL^∞ cut-reduction sequences. Even though the simulation of μLL^∞ sequences builds transfinite sequences, we shall see that one can associate a(n almost) fair μMALL^∞ mcut-reduction sequence to any fair μLL^∞ mcut-reduction sequence, and conclude.
- The next four subsections will closely follow the above pathway.

3.1 Cut-elimination rules for μLL^∞

μLL^∞ mcut-reduction is defined by extending μMALL^∞ multicut-reduction with the steps given in Figure 6. The reduction rules for the exponentials assume a condition on the premisses of the multi-cut rule: all the proofs (hereditarily) cut-connected to some distinguished formula must have promotions as last inferences.

More details can be found in appendix C.1.

Definition 28 ((!p)-ready contexts). *A subset of the subproofs of a multicut is said to be (!p)-ready if all its elements are concluded with an (!p) rule. $\mathcal{C}^!$ will denote a (!p)-ready context and $\mathcal{C}_\Gamma^!$ a context restriction which is (!p)-ready.*

Remark 29. The condition for triggering the exponential key reductions (?w)/(!p) and (?c)/(!p) as well as the (!p)-commutation rule is expressed in terms of (!p)-readiness: for every ?-formula ?G in the context of a promotion which shall either commute or cut-reduce with a ?-rule, we require that $\mathcal{C}_{?G}$ is (!p)-ready.

3.2 Embedding μLL^∞ in μMALL^∞

To extend the cut-elimination result from μMALL^∞ to μLL^∞ , we encode the exponential connectives using fixed points as follows, following Baelde [2]:

Definition 30. $?^\bullet(F) = \mu X.F \oplus (\perp \oplus (X \wp X)); !^\bullet(F) = \nu X.F \& (1 \& (X \otimes X))$

This straightforwardly induces an embedding of μLL^∞ into μMALL^∞ :

Definition 31 (Embedding of μLL^∞ sequents into μMALL^∞).

$$\begin{array}{l|l} (a)^\bullet = a & \text{if } a \text{ is an atom} \\ (u)^\bullet = u & \text{if } u \in \{1, \perp, \top, 0\} \\ (A \star B)^\bullet = (A)^\bullet \star (B)^\bullet & \text{if } \star \in \{\&, \oplus, \wp, \otimes\} \end{array} \left| \begin{array}{l} (\sigma X.F)^\bullet = \sigma X.(F)^\bullet \\ (?F)^\bullet = ?^\bullet(F^\bullet) \\ (!F)^\bullet = !^\bullet(F^\bullet) \end{array} \right. , \sigma \in \{\mu, \nu\}$$

Definition 32 (μMALL^∞ derivability of the exponentials). *μLL^∞ exponential rules can be encoded in μMALL^∞ as shown in Figure 7. We denote the derivable rules by ?d $^\bullet$, ?c $^\bullet$, ?w $^\bullet$ and !p $^\bullet$ respectively. (!p $^\bullet$ uses a circular proof.)*

Proposition 33 (Preservation of validity). *π is a valid μLL^∞ proof of $\vdash \Gamma$ iff π^\bullet is a valid μMALL^∞ proof of $\vdash \Gamma^\bullet$.*

$$\begin{array}{c}
\frac{\frac{\mathcal{C} \quad \frac{\vdash \Delta, F}{\vdash \Delta, ?F} \text{ (?d)}}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash \Delta, F}{\vdash \Sigma, F} \text{ mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma, ?F} \text{ (?d)} \\
\frac{\mathcal{C} \quad \frac{\vdash \Delta, ?F, ?F}{\vdash \Delta, ?F} \text{ (?c)}}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash \Delta, ?F, ?F}{\vdash \Sigma, ?F, ?F} \text{ mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma, ?F} \text{ (?c)} \\
\frac{\mathcal{C} \quad \frac{\vdash \Delta}{\vdash \Delta, ?F} \text{ (?w)}}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash \Delta}{\vdash \Sigma} \text{ mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma, ?F} \text{ (?w)} \\
\frac{\frac{\vdash F, ?\Gamma}{\vdash !F, ?\Gamma} \text{ (!p)} \quad \mathcal{C}^!}{\vdash !F, ?\Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\vdash !F, ?\Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\frac{\vdash F, ?\Gamma}{\vdash !F, ?\Sigma} \text{ (!p)} \quad \mathcal{C}^!}{\vdash !F, ?\Sigma} \text{ (!p)} \\
\frac{\frac{\frac{\vdash F, \Gamma}{\vdash ?F, \Gamma} \text{ (?d)} \quad \frac{\vdash F^\perp, ?\Delta}{\vdash !F^\perp, ?\Delta} \text{ (!p)}}{\vdash \Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash F, \Gamma \quad \vdash F^\perp, ?\Delta}{\vdash \Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp')}}{\vdash \Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp')
\end{array}$$

where $?F \perp\!\!\!\perp !F^\perp$ and $\perp\!\!\!\perp'$ coincides with $\perp\!\!\!\perp$ except for $F \perp\!\!\!\perp' F^\perp$.

$$\frac{\mathcal{C}_\Gamma \quad \mathcal{C}_{?F}^! \quad \frac{\frac{\vdash ?F, ?F, \Gamma}{\vdash ?F, \Gamma} \text{ (?c)}}{\vdash \Gamma', ?\Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Gamma', ?\Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\mathcal{C}_\Gamma \quad \mathcal{C}_{?F}^! \quad \mathcal{C}_{?F}^! \quad \frac{\vdash ?F, ?F, \Gamma}{\vdash \Gamma', ?\Sigma, ?\Sigma} \text{ mcut}(\iota', \perp\!\!\!\perp')}}{\frac{\vdash \Gamma', ?\Sigma}{\vdash \Gamma', ?\Sigma} \text{ (?c)}^*} \text{ mcut}(\iota', \perp\!\!\!\perp')$$

where $\mathcal{C}_{?F}^! \neq \emptyset$, $\perp\!\!\!\perp'$ corresponds to $\perp\!\!\!\perp$ on \mathcal{C}_Γ and is a “duplication” of $\perp\!\!\!\perp$ on $\mathcal{C}_{?F}^!$ and each copy of $?F$ is in $\perp\!\!\!\perp'$ -relation with the corresponding copy of $!F^\perp$ in $\mathcal{C}_{?F}^!$.

$$\frac{\mathcal{C}_\Gamma \quad \mathcal{C}_{?F}^! \quad \frac{\frac{\vdash \Gamma}{\vdash ?F, \Gamma} \text{ (?w)}}{\vdash \Gamma', ?\Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Gamma', ?\Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\mathcal{C}_\Gamma \quad \frac{\vdash \Gamma}{\vdash \Gamma'} \text{ mcut}(\iota', \perp\!\!\!\perp')}}{\frac{\vdash \Gamma', ?\Sigma}{\vdash \Gamma', ?\Sigma} \text{ (?w)}^*} \text{ mcut}(\iota', \perp\!\!\!\perp')$$

where $\mathcal{C}_{?F}^! \neq \emptyset$ and $\perp\!\!\!\perp'$ corresponds to the restriction of $\perp\!\!\!\perp$ on $\mathcal{C}_\Gamma, \Gamma$.

Fig. 6. μLL^∞ mcut-reduction rules.

Proof (Proof sketch). We simply relate the infinite branches in both pre-proofs. Assuming that π is valid, consider the special case of an infinite branch β of π^\bullet that, when entering the encoding of a promotion, follows the left-most premise of the ($\&$) rule. To such an infinite branch it is easy to associate an infinite branch b of π . b is valid and supported by a thread t with least formula $\nu X.F$. $(\nu X.F)^\bullet$ is the least recurring formula in the thread θ associated with t in β : β is valid.

The detailed proof can be found in Appendix C.2.

3.3 Simulation of μLL^∞ cut-elimination steps

Now we have to show that μLL^∞ cut-elimination steps can be simulated by the previous encoding. *E.g.*, the commutation rule for dereliction is simulated by a

<p>Dereliction :</p> $\frac{\frac{\vdash F, \Delta}{\vdash F \oplus (\perp \oplus (?^{\bullet} F \wp ?^{\bullet} F)), \Delta} \text{ } (\oplus^1)}{\vdash ?^{\bullet} F, \Delta} \text{ } (\mu)$	<p>Contraction :</p> $\frac{\frac{\frac{\frac{\vdash ?^{\bullet} F, ?^{\bullet} F \Delta}{\vdash ?^{\bullet} F \wp ?^{\bullet} F, \Delta} \text{ } (\wp)}{\vdash \perp \oplus (?^{\bullet} F \wp ?^{\bullet} F), \Delta} \text{ } (\oplus^2)}{\vdash F \oplus (\perp \oplus (?^{\bullet} F \wp ?^{\bullet} F)), \Delta} \text{ } (\oplus^2)}{\vdash ?^{\bullet} F, \Delta} \text{ } (\mu)$	<p>Weakening :</p> $\frac{\frac{\frac{\vdash \Delta}{\vdash \perp, \Delta} \text{ } (\perp)}{\vdash \perp \oplus (?^{\bullet} F \wp ?^{\bullet} F), \Delta} \text{ } (\oplus^1)}{\vdash F \oplus (\perp \oplus (?^{\bullet} F \wp ?^{\bullet} F)), \Delta} \text{ } (\oplus^2)}{\vdash ?^{\bullet} F, \Delta} \text{ } (\mu)$
<p>Promotion :</p> $\frac{\frac{\frac{\frac{\vdash 1}{\vdash 1, ?^{\bullet} \Delta} \text{ } (1)}{\vdash F, ?^{\bullet} \Delta} \text{ } (?w^{\bullet})^*}{\vdash !^{\bullet} F, ?^{\bullet} \Delta} \text{ } (\nu), (\&), (\&)}{\frac{\frac{\frac{\frac{\vdash !^{\bullet} F, ?^{\bullet} \Delta}{\vdash !^{\bullet} F \otimes !^{\bullet} F, ?^{\bullet} \Delta, ?^{\bullet} \Delta} \text{ } (\otimes)}{\vdash !^{\bullet} F \otimes !^{\bullet} F, ?^{\bullet} \Delta} \text{ } (?c^{\bullet})^*}{\vdash !^{\bullet} F, ?^{\bullet} \Delta} \text{ } (\nu), (\&), (\&)}{\vdash !^{\bullet} F, ?^{\bullet} \Delta} \text{ } (\nu), (\&), (\&)} \text{ } (\nu), (\&), (\&)$		

Fig. 7. μMALL^{∞} encoding of the exponential inferences.

$(\mu)/(\text{Cut})$ commutation followed by a $(\oplus)/(\text{Cut})$ commutation as follows:

$$\frac{\frac{\frac{\vdash F, G, \Gamma}{\vdash ?^{\bullet} F, G, \Gamma} \text{ } (?d^{\bullet})}{\vdash ?^{\bullet} F, \Gamma, \Delta} \text{ } (\text{Cut}) \quad \vdash G^{\perp}, \Delta}{\vdash ?^{\bullet} F, \Gamma, \Delta} \text{ } (\text{Cut})} \longrightarrow^2 \frac{\frac{\vdash F, G, \Gamma \quad \vdash G^{\perp}, \Delta}{\vdash F, \Gamma, \Delta} \text{ } (\text{Cut})}{\vdash ?^{\bullet} F, \Gamma, \Delta} \text{ } (?d^{\bullet})$$

The challenge is to show that the simulation of reductions also holds (i) for the reductions involving $[\!|p\rangle]$ as well as (ii) for reductions occurring *above* a promotion rule (aka. in a box) since the encoding of $[\!|p\rangle]$ uses an infinite, circular derivation. In the promotion commutation case for instance, we have:

$$\frac{\frac{\frac{\vdash F, ?^{\bullet} G, ?^{\bullet} \Gamma}{\vdash !^{\bullet} F, ?^{\bullet} G, ?^{\bullet} \Gamma} \text{ } (!p^{\bullet})}{\vdash !^{\bullet} F, ?^{\bullet} \Gamma, ?^{\bullet} \Delta} \text{ } (\text{Cut}) \quad \frac{\frac{\vdash G^{\perp}, ?^{\bullet} \Delta}{\vdash !^{\bullet} G^{\perp}, ?^{\bullet} \Delta} \text{ } (!p^{\bullet})}{\vdash !^{\bullet} F, ?^{\bullet} \Gamma, ?^{\bullet} \Delta} \text{ } (!p^{\bullet})}}{\vdash !^{\bullet} F, ?^{\bullet} \Gamma, ?^{\bullet} \Delta} \text{ } (\text{Cut})} \longrightarrow^{\omega} \frac{\frac{\frac{\vdash F, ?^{\bullet} G, ?^{\bullet} \Gamma}{\vdash !^{\bullet} F, ?^{\bullet} G, ?^{\bullet} \Gamma} \text{ } (!p^{\bullet})}{\vdash !^{\bullet} F, ?^{\bullet} \Gamma, ?^{\bullet} \Delta} \text{ } (!p^{\bullet})} \quad \frac{\frac{\vdash G^{\perp}, ?^{\bullet} \Delta}{\vdash !^{\bullet} G^{\perp}, ?^{\bullet} \Delta} \text{ } (!p^{\bullet})}{\vdash !^{\bullet} F, ?^{\bullet} \Gamma, ?^{\bullet} \Delta} \text{ } (!p^{\bullet})}}{\vdash !^{\bullet} F, ?^{\bullet} \Gamma, ?^{\bullet} \Delta} \text{ } (\text{Cut})} \text{ } (!p^{\bullet})$$

Proposition 34. *Each μLL^{∞} mcut-reduction r can be simulated in μMALL^{∞} by a (possibly infinite) sequence of mcut-reductions, denoted r^{\bullet} .*

See proof in App C.3.

Remark 35. Conversely, one can wonder whether a possible reduction in π^{\bullet} necessarily comes from the simulation of a reduction step in π . It is *almost* the case except when the reduction in π^{\bullet} comes from exponential cuts requiring a $(!p)$ -ready context (*ie.* $(!p)$ commutation as well as $(?w)/(!p)$ and $(?c)/(!p)$ key cases, see above): in those cases indeed, if the context is “partially ready” – meaning that some, but not all, the required premises are promoted – a prefix of the sequence simulating the reduction step can indeed be performed, before being stuck. As consequence – and we shall exploit it in the next section when proving μLL^{∞} cut-elimination – the simulation of a fair reduction sequence is not necessarily fair, *but only as long as the above cases are involved*:

Proposition 36. *There exists a fair reduction ρ from some μLL^{∞} (pre-)proof π such that ρ^{\bullet} is an ω -indexed unfair μMALL^{∞} cut-reduction sequence*

See proof in App C.4.

3.4 Proof of μLL^∞ cut-elimination theorem

μLL^∞ cut-elimination theorem follows from the following two lemmas:

See proof in
App C.5.

Lemma 37. *Let π be a μLL^∞ -proof of $\vdash \Gamma$ and $\sigma = (\pi_i, r_i, p_i)_{i \in \omega}$ a fair μLL^∞ cut-reduction sequence from π . σ converges to a cut-free μLL^∞ -pre-proof of $\vdash \Gamma$.*

See proof in
App C.5.

Lemma 38. *Let π be a μLL^∞ pre-proof of $\vdash \Gamma$ and let us consider a cut-reduction sequence $\sigma = (\pi_i, r_i, p_i)_{i \in \omega}$ in μLL^∞ from π that converges to a cut-free μLL^∞ pre-proof π' . σ^\bullet is a strongly converging (possibly transfinite) sequence.*

See details in
App C.5.

Proof (Sketch for Thm. 27). Let π be a μLL^∞ -proof of $\vdash \Gamma$ and $\sigma = (\pi_i, r_i, p_i)_{i \in \omega}$ be a fair μLL^∞ mcut-reduction sequence from π . Consider the associated (transfinite) μMALL^∞ mcut-reduction sequence σ^\bullet from π^\bullet obtained by simulation. By Lemma 37, σ converges (strongly) to a cut-free μLL^∞ pre-proof π' .

Let us prove that π' is valid. By Lemma 38, σ^\bullet is a *transfinite* mcut-reduction sequence from π^\bullet strongly converging to π'^\bullet . By Prop. 26, σ^\bullet can be compressed into $\rho = (\pi'_i, r'_i, p'_i)_{i \in \omega}$ an ω -indexed depth-increasing μMALL^∞ mcut-reduction sequence which converges to π'^\bullet and contains the same reductions as σ^\bullet . By Proposition 36, ρ may not be fair: this prevents us from concluding directly by Proposition 33 but we can still conclude. Let us consider ρ_f a fair reduction sequence obtained from ρ by reducing those redexes which cause the lack of fairness of ρ and let us consider the limit of ρ_f, π_f . To any infinite branch β of π'^\bullet , one can associate a branch β_f of π_f : it coincides with β except when the next inference of β_f is on a $(!F)^\bullet$ (in a sequent, say, $\vdash (!F)^\bullet, ?^\bullet \Delta^\bullet$ which is not principal along β). In that case, we expand β_f by following the unique premise of the (ν) rule, the second premise of the first $(\&)$ rule and the first premise of the second $(\&)$ rule, reaching $\vdash 1, ?^\bullet \Delta^\bullet$, in which case we know that the 1 is not principal (and never will be) and we follow back β . β_f has exactly the same threads as β : finite threads may only be extended *finitely* on occurrences of $(!F)^\bullet$. Since ρ_f is fair, β_f is valid and so is β .

We can then conclude that π'^\bullet is cut-free and valid and, using preservation of validity (Proposition 33), that π' is a valid cut-free μLL^∞ -proof. \square

Infinitary cut-elimination for μLL^∞ two-sided sequent calculus is an easy corollary of Theorem 27. Indeed, fair cut-reduction sequences in two-sided μLL^∞ are mapped to fair reduction sequences in one-sided μLL^∞ from which follows:

See proof in
App C.5.

Corollary 39. *Fair 2-sided μLL^∞ valid mcut-reduction sequences eliminate cuts.*

4 Cut-elimination theorem for μLK^∞ and μLJ^∞

Cut-elimination theorems for both μLK^∞ and μLJ^∞ can be established as corollaries of Theorem 27. For lack of space, we directly go to our results and postpone to future work a detailed study of the generalizations to non-wellfounded sequent calculi of the linear embeddings of LK and LJ into LL developed since Girard seminal paper. We shall comment on those translations in the conclusion.

The translation preserves validity, following from $[X]^j = !X$, by induction.

Lemma 46. *The following hold:*

- For any μLJ formulas A, B , $\sigma \in \{\mu, \nu\}$, $[A[\sigma X.B/X]]^j = [A]^j[\sigma X.[B]^j/X]$.
- For any μLJ formula A , $[A]^j$ is a μLL formula.
- If π is a μLJ^∞ proof of $\Gamma \vdash F$, then $[\pi]^j$ is a μLL^∞ proof of $[\Gamma]^j \vdash [F]^j$.

On μLL^∞ proofs, the skeletons of the previous section can be reused: $\text{Sk}(\cdot)$ transports valid μLL^∞ proof to valid μLJ^∞ proofs. Moreover μLL^∞ proofs are closed by μLL^∞ cut-reductions from which we deduce, as for μLK^∞ , that:

Theorem 47. *μLJ^∞ enjoys cut-elimination.*

5 Conclusion

In the present paper, we established several cut-elimination results for non-wellfounded proof systems for logics with least and greatest fixed-points expanding on previous works [4,20]: (i) for μMALL^∞ with sequents as lists in contrast sequents as sets of locative occurrences [4], (ii) for the 1-sided and 2-sided sequent calculi of μLL^∞ , (iii) for μLK^∞ and (iv) for μLJ^∞ . We also established additional results from a compression lemma for μMALL^∞ strongly converging cut-reduction sequences to linear embeddings of μLK^∞ and μLJ^∞ into μLL^∞ .

On the meaning and expressiveness of tree-exponential modalities. The proof of our main result proceeds by encoding LL exponentials in μMALL^∞ following an encoding first considered by Baelde and Miller, and studying μLL^∞ cut-reduction sequences through their simulation in μMALL^∞ , which was first conjectured in Doumane’s thesis [18]. We think that the present paper does not only demonstrate the usefulness of the encoding but that it also suggests new questions. Indeed, this encoding has interesting features:

- this “rigid” tree-like exponential does not exhibit the Seely isomorphism but, even though those isomorphisms are common in axiomatizations of categorical models of linear logic, it is not necessary to have them as isomorphisms to build a denotational model of linear logic (that is, which quotients proofs up to cut-equivalence): the present work is actually an example of this fact. (They are crucial, though, to encode the λ -calculus in linear logic, as additional equations are needed, which are realized by Seely isos.)
- These exponentials allow for a realization of a somehow non-uniform promotion: indeed, while a proof of $\vdash !^\bullet F, ?^\bullet \Gamma$ has to provide a proof of $\vdash F, ?^\bullet \Gamma$, the circular definition of the promotion is not the only possible definition: one can consider as well promotions that would provide a distinct value each time a box is opened (*e.g.* a proof of $\vdash !^\bullet \mu X.1 \oplus X$ may provide distinct integers depending on how structural rules managed the resource).

This tree-like exponential is being investigated with Ehrhard and Jafarrahmani.

See discussion
& example in
appendix E.

Benefiting from advances in infinitary rewriting. Our cut-elimination proof by encoding μLL^∞ into μMALL^∞ relies on a simulation of reductions sequences which makes use of transfinite reductions sequences and compression results. Those techniques are inspired and adapted from the literature on infinitary rewriting. We plan to make clearer the connection between non-wellfounded proof theory and infinitary rewriting in the future, even though in the present state it was not possible to readily apply results from infinitary rewriting such as the compression lemma which we has to reprove in our setting [32]. Moreover, we did not make use of coinductive formulations of infinitary rewriting [19]. That is another direction for future work: currently, we do not know how to use those formulations of infinitary rewriting because the sequences we consider by simulation are not given as (strongly) converging sequences. We plan to reconsider this and benefit from the coinductive approach to infinite reduction sequences.

On linear translations for fixed-point logics and non-wellfounded proofs. We obtained a cut-elimination theorem for μLK^∞ and μLJ^∞ thanks to linear translations which deserve some comments. While the linear translation used for μLJ^∞ is standard (it is a call-by-value translation dating back to Girard’s seminal paper), the treatment of classical logic was more complex. Indeed, usual linear translation for classical logic introduce, at places, cuts. Due to the sensitivity of the straight-thread validity condition with respect to the presence of cuts in cycles, we could not use those translations. However, we plan to investigate whether a more standard translation can be used in the specific case of bouncing validity [3].

A treatment of cut-elimination which is agnostic to validity conditions. Last but not least, a major advantage of our approach is that μMALL^∞ cut-elimination proof and, to some extent, the validity conditions, are regarded as black boxes, simplifying the presentation of the proof and making it reusable wrt. other validity condition or μMALL^∞ proof techniques. The proof seems to be reusable easily with bouncing validity for instance (even though setting up an adequate definition of bouncing validity for μLL^∞ is quite tricky). A fragment which seems promising and that we wish to investigate in the near future, is μMELL^∞ equipped with bouncing validity [3].

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A Appendix on infinitary rewriting

In the present section, we review some standard notions of infinitary rewriting which shall be useful in the remainder of the paper and are standard from the literature, see for instance TERESE book [32] for a general reference for infinitary rewriting covering (more than) what we will need.

Definition 48 (Non-wellfounded terms). *Given a signature Σ , that is a set of pairs (f, a) of a function symbol f and an arity $a \in \omega$, we write as usual $\text{Term}(\Sigma)$ for the set of finite terms over Σ . The set of **non-wellfounded terms**, or **infinite terms** over Σ , written $\text{Term}^\infty(\Sigma)$, is given as the set of Σ -labelled infinite trees given by a set P of (finite) sequences of natural numbers, the **positions of the tree**, and a function $\ell : P \mapsto \Sigma$, the **labelling function** satisfying the following conditions:*

- P is prefix closed: if $p \cdot k \in P$, then $p \in P$;
- if $p \cdot k \in P$ and $k' \leq k$, then $p \cdot k' \in P$;
- for any $p \in P$, if a is the arity of $\ell(p)$, then $p \cdot k \in P$ if and only if $0 \leq i < a$.

If $t \in \text{Term}^\infty(\Sigma)$, we denote by $\text{Pos}(t)$ the set of **positions of t** and if $p \in \text{Pos}(t)$, the **subterm of t at p** , written $t_{@p}$, is the subtree of t rooted in p .

Two positions p, p' of a term are **incomparable** if none is a prefix of the other. The length of a position is defined as usual

We define a **distance** between terms $d : \text{Term}^\infty(\Sigma) \times \text{Term}^\infty(\Sigma) \rightarrow \mathbb{R}$ as: $d(s, t) = 0$ if $s = t$ and $d(s, t) = 2^{-k}$ where k is the length of the shortest position $p \in \text{Pos}(s) \cap \text{Pos}(t)$ where s and t differ otherwise.

Definition 49 (Infinitary rewriting rules, infinitary rewriting step). *An **infinitary rewriting rule** (∞RR) is a pair (t, s) with $t \in \text{Term}(\Sigma)$ and $s \in \text{Term}^\infty(\Sigma)$ such that t is not reduced to a variable and $\text{Var}(s) \subseteq \text{Var}(t)$.*

*If r is an ∞RR , $t, s \in \text{Term}^\infty(\Sigma)$ and $p \in \text{Pos}(t)$, we call an **infinitary rewriting step** (∞RS) from t to s along r at p , and write $t \xrightarrow[r]{p} s$, if $r = (t', s')$ and there exists a substitution $\sigma : \text{Var}(t') \mapsto \text{Term}^\infty(\Sigma)$ such that $t_{@p} = \sigma(t')$, $s_{@p} = \sigma(s')$.*

*If p is a position and $R = t \xrightarrow[r]{p} s$ is an ∞RS , the **depth** of p (resp. R), written $\text{dpth}(p)$, is the length of p*

Definition 50 (Transfinite reduction sequence). *Given a signature Σ , a set of infinitary rewriting rules \mathcal{R} and an ordinal λ , a **transfinite reduction sequence** of length λ , or λTRS , is a λ -indexed sequence $(t_i, r_i, p_i)_{i \in \lambda}$ such that:*

- $\forall i \in \lambda, t_i \in \text{Term}^\infty(\Sigma), r_i \in \mathcal{R}, p_i \in \text{Pos}(t_i)$;
- $t_i \xrightarrow[r_i]{p_i} t_{i+1}$, for any i such that $i + 1 \in \lambda$.

Definition 51 (Weak and strong convergence). *Let $\rho = (t_i, r_i, p_i)_{i \in \alpha}$ be a αTRS . ρ is **weakly converging** if for every limit ordinal $\lambda \leq \alpha$, $(d(t_i, t_\lambda))_{i \in \lambda}$ tends to 0. ρ is **strongly converging** if it is weakly converging and if for every limit ordinal $\lambda \leq \alpha$, $(\text{dpth}(p_i))_{i \in \lambda}$ tends to infinity.*

Example 52. Let us consider $\pi = \frac{\frac{\vdots}{\vdash \Gamma, \mu X.X} \quad (\mu)}{\vdash \Gamma, \mu X.X} \quad (\mu) \quad \frac{\frac{\vdots}{\vdash \nu X.X, \Delta} \quad (\nu)}{\vdash \nu X.X, \Delta} \quad (\nu)}{\vdash \Gamma, \Delta} \quad (\text{Cut})$. With

the cut-reduction system to be considered in this paper, $\pi \xrightarrow{\epsilon}_{(\text{princ}, \{\mu X.X, \nu X.X\})} \pi$, therefore there is a weakly converging reduction sequence from π but not strongly converging reduction.

Remark 53. It is easy to see that in a strongly converging reduction sequence, there can only be finitely many reduction steps occurring at each depth; therefore every strongly converging reduction sequence is indexed by a countable ordinal.

We end this section with a property that will be essential in our construction, the compression property:

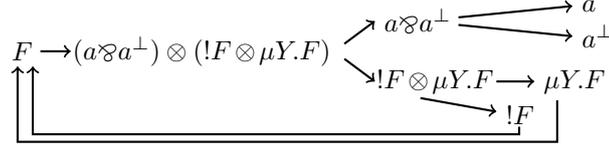
Definition 54 (Compression property). *A infinitary term rewriting system is said to have the **compression property** (or is said to be **CP**) when for every transfinite strongly converging reduction sequence from s converging to t , there is a strongly converging reduction sequence indexed by ω or a finite ordinal from s converging to t .*

Proposition 55 (Compression Lemma [32]). *Every left-linear infinitary term rewriting system is CP.*

B Appendix on the non-wellfounded proof systems: μMALL^∞ , μLL^∞ , μLK^∞ , μLJ^∞ .

B.1 μ -Formulas and subformulas (details on Section 2.1)

One can represent the Fischer-Ladner closure as a graph, following Example 8.



B.2 Inference rules

Definition 56 (\mathcal{S} -sequents and inferences). A sequent $s = \Gamma \vdash \Delta$ over a μ -signature \mathcal{S} is a pair of finite lists Γ, Δ of \mathcal{S} -formulas. (As usual, Γ is called the antecedent and Δ the succedent.)

An inference rule $r = (s, (s_1, \dots, s_n), \text{Princ}, \alpha, l)$ is the data of

- a conclusion \mathcal{S} -sequent s , also noted $\text{Conc}(r)$;
- a tuple of premise \mathcal{S} -sequents s_1, \dots, s_n , also noted $\text{Prem}(r)$;
- together with an ancestor relation $\alpha \subseteq [1, \text{Card}(s)] \times \{(i, j), 1 \leq i \leq n, 1 \leq j \leq \text{Card}(s_i)\}$;
- a subset $\text{Princ} \subseteq \{1, \dots, \text{Card}(s)\}$ of cardinality at most 2, identifying the so-called principal formulas of the (conclusion of the) inference, also noted $\text{Princ}(r)$;
- and a label l , the rule name, also noted $\text{RName}(r)$.

We set a number of conventions that will facilitate the description and manipulation of sequent derivations.

Convention 2 In the following, inference rules will be described, as is usual, as inference schema thanks to both formula meta-variables (ranging over F, G, \dots) and context meta-variables (ranging over Γ, Δ, \dots) while the ancestor relation will be depicted as colored lines joining related formulas. When identical context meta-variables are related, we mean that the part of the ancestor relation such described relates the i th formula of the bottom context with the i th formula of the topmost context: that will allow us to avoid the very explicit description of inference rules given above, while allowing to reconstruct it at will.

Convention 3 In the sequent calculi defined in the following, we follow the usual convention that the principal formulas of an inference are all those formulas which are explicitly spelled out in the conclusion sequent and not described via a context meta-variable.

A formula occurrence of an inference is said to be active if it is principal or if its a related to a principal formula of the inference by the ancestor relation.

Convention 4 In general, all the considered sequent calculi will have the exchange rule and we shall make the assumption that the set of inference rules to be closed by pre-composition and post-composition of the exchange rule. By that, we mean that we will freely use the derived rules which are obtained by adding any number of exchange rule above and below each inference, adapting the ancestry relation of course.

Convention 5 When considering a proof system with one-sided sequents and an involutive negation \cdot^\perp , we shall sometimes write $\Gamma \vdash \Delta$ for sequents $\vdash \Gamma^\perp, \Delta$ in order to make the computational behaviour of our examples clearer. The labelling of our inferences will be unchanged (note that the index associated to formulas in a sequent remain unchanged by this operation).

B.3 μLJ^∞ sequent calculus

We provide details on Definition 11:

Definition 57 (μLJ^∞ inferences). The inference rules of μLJ^∞ sequent calculus, defined with respect to the μ -signature $\mathcal{C}_{\mu\text{LJ}}$, is given in Figure 8.

$$\begin{array}{c}
\frac{}{\Gamma, F \vdash F} \text{ (Ax)} \quad \frac{\Gamma \vdash F \quad \Delta, F \vdash H}{\Gamma, \Delta \vdash H} \text{ (Cut)} \quad \left| \quad \frac{\Gamma, G, F, \Delta \vdash H}{\Gamma, F, G, \Delta \vdash H} \text{ (X}_1\text{)} \quad \frac{\Gamma, F, F \vdash H}{\Gamma, F \vdash H} \text{ (C}_1\text{)} \\
\frac{\Gamma, F \vdash H \quad \Gamma, G \vdash H}{\Gamma, F \vee G \vdash H} \text{ (v}_1\text{)} \quad \frac{\Gamma \vdash F_i}{\Gamma \vdash F_1 \vee F_2} \text{ (v}_2\text{)} \quad \frac{\Gamma, F_i \vdash H}{\Gamma, F_1 \wedge F_2 \vdash H} \text{ (\wedge}_1\text{)} \quad \frac{\Gamma \vdash F \quad \Gamma \vdash G}{\Gamma \vdash F \wedge G} \text{ (\wedge}_r\text{)} \\
\frac{}{\Gamma \vdash \top} \text{ (tr)} \quad \frac{}{\Gamma, F \vdash H} \text{ (f}_1\text{)} \quad \frac{\Gamma \vdash F \quad \Delta, G \vdash H}{\Gamma, \Delta, F \Rightarrow G \vdash H} \text{ (\Rightarrow}_1\text{)} \quad \frac{\Gamma, F \vdash G}{\Gamma \vdash F \Rightarrow G} \text{ (\Rightarrow}_r\text{)} \\
\frac{\Gamma, F[\mu X.F/X] \vdash H}{\Gamma, \mu X.F \vdash H} \text{ (\mu)} \quad \frac{\Gamma \vdash F[\mu X.F/X]}{\Gamma \vdash \mu X.F} \text{ (\mu}_r\text{)} \quad \frac{\Gamma, G[\nu X.G/X] \vdash H}{\Gamma, \nu X.G \vdash H} \text{ (\nu)} \quad \frac{\Gamma \vdash G[\nu X.G/X]}{\Gamma \vdash \nu X.G} \text{ (\nu}_r\text{)}
\end{array}$$

Fig. 8. μLJ^∞ Inferences

B.4 Pre-proofs and validity (Appendix for Section 2.3)

We provide some details on Definition 12:

Definition 58 (Pre-proofs). Given one of the above considered μ -signatures and sets of inferences, the corresponding set of pre-proofs is the set of finite and infinite trees whose nodes are labelled with inference rules and which are locally correct in the following sense: for every node n of a pre-proof, the arity n coincides with the number of premises of the corresponding inference and the

i th premise of n matches the conclusion sequent of the root of the i th subtree of n . Such a tree is called a pre-proof of sequent s if s is the conclusion sequent of the root of the pre-proof.

We provide more details here about the precise treatment of infinite branches and threads given in Section 2.3.

Definition 59 (Infinite branch). Given a pre-proof π , the set of **infinite branches** of π , $\infty\text{Br}(\pi)$ is defined as $\infty\text{Br}(\pi) = \{(a_i)_{i \in \omega} / \pi = \frac{\pi_1 \cdots \pi_k}{r} \ \& \ 1 \leq a_0 \leq k \ \& \ (a_{i+1})_{i \in \omega} \in \infty\text{Br}(\pi_{a_0})\}$.

It is useful to consider also **enriched branches**: $\infty\text{EBr}(\pi) = \{(a_i, r_i)_{i \in \omega} / \pi = \frac{\pi_1 \cdots \pi_k}{r} \ \& \ 1 \leq a_0 \leq k \ \& \ (a_{i+1}, r_{i+1})_{i \in \omega} \in \infty\text{EBr}(\pi_{a_0})\}$.

Enriched branches are convenient to define threads:

Definition 60 (Thread). Let π be a pre-proof and $\beta = (a_i, r_i)_{i \in \omega} \in \infty\text{EBr}(\pi)$. We define

$$\infty\text{Th}(\beta) = \{(t_i)_{i \in \omega} / r_0 = (s, (s_1, \dots, s_k), p, \alpha, l), 1 \leq t_0 \leq \text{Card}(s) \ \& \ \alpha(t_0, (a_0, t_1)) \ \& \ (t_{i+1})_{i \in \omega} \in \infty\text{Th}((a_{i+1}, r_{i+1})_{i \in \omega})\}.$$

It is useful to consider also enriched threads:

$$\infty\text{ETH}((a_i, r_i)_{i \in \omega}) = \{(a_i, r_i, t_i)_{i \in \omega} / r_0 = (s, (s_1, \dots, s_k), p, \alpha, l), 1 \leq t_0 \leq \text{Card}(s) \ \& \ \alpha(t_0, (a_0, t_1)) \ \& \ (a_{i+1}, r_{i+1}, t_{i+1})_{i \in \omega} \in \infty\text{ETH}((a_{i+1}, r_{i+1})_{i \in \omega})\}.$$

Definition 61 (Recurring formula). A formula F is recurring in an enriched thread $(a_i, r_i, t_i)_{i \in \omega}$ if $\{i \in \omega / \text{Conc}(r_i)(t_i) = F\}$ is infinite.

Definition 62 (Valid thread). A thread $\tau = (a_i, r_i, t_i)_{i \in \omega} \in \infty\text{ETH}(\beta)$ is said to be a valid thread for β if for any $i \in \mathbb{N}$, there exists $j \geq i$ such that t_j is principal in r_j (and r_j is not an exchange rule) and the set of recurring formulas of τ has a least elements for the subformula ordering which is ν when it occurs in the succedents of sequents or a μ if it occurs in the antecedent of sequents.

Definition 63 (Valid pre-proof). A pre-proof is valid if all its infinite branches have suffixes which are supported by valid threads.

B.5 Details on the multicut rule (Definition 16)

Definition 64 (Multicut inference rule (Expanded definition)). The multicut inference is given by the data of

- a conclusion sequent s ;
- non-empty tuple (s_1, \dots, s_n) of premises, $n \geq 1$;
- an ancestor relation ι which is an injective map from the conclusion formulas to the premise formulas, that is from $\mathbf{C} = [1, \text{Card}(s)]$ to $\mathbf{Pr} = \{(j, k), 1 \leq j \leq n, 1 \leq k \leq \text{Card}(s_j)\}$, that is, for any $1 \leq i \leq \text{Card}(s)$ there exists a unique $(j, k) \in \mathbf{Pr}$ such that $\iota(i, (j, k))$ and a given pair (j, k) is in relation with at most one conclusion formula.

Moreover, if $\iota(i, (j, k))$, then $s(i) = s_j(k)$: the ancestor relation relates identical formulas.

- the rule has no principal formula (as the cut rule);
- the rule name is $\text{mcut}(\iota, \perp\!\!\!\perp)$, where $\perp\!\!\!\perp$ is a symmetric binary relation over $\text{Pr} \setminus \iota(\mathbf{C})$ defined as follows:
 - $\perp\!\!\!\perp$ is total: for any $t \in \text{Pr} \setminus \iota(\mathbf{C})$, there exists some $t' \in \text{Pr} \setminus \iota(\mathbf{C})$ such that $t \perp\!\!\!\perp t'$;
 - $(i, j) \perp\!\!\!\perp (k, l) \implies s_i(j) = (s_k(l))^\perp$ (the multicut relation relates dual formulas);
 - $\forall 1 \leq i, k \leq n$, if $(i, j) \perp\!\!\!\perp (k, l)$ and $(i, j') \perp\!\!\!\perp (k, l')$, then $j = j', l = l'$ (two sequents are related in at most one way by the multicut relation.);
 - if $n_1, \dots, n_k \in \text{Pr} \setminus \iota(\mathbf{C})$, $n_i \perp\!\!\!\perp n_{i+1}$ and $n_k = n_1$ then there is $1 \leq i, j \leq k$ such that $n_i = n_{j+1}$ and $n_j = n_{i+1}$ (the only cycles come from the symmetry of the relation);
 - for any $1 \leq i, j \leq n$, there exists k, l and $n_1, \dots, n_p \in \text{Pr} \setminus \iota(\mathbf{C})$ such that $n_1 = (i, k)$, $n_p = (j, l)$ and $n_q \perp\!\!\!\perp n_{q+1}$ for $1 \leq q < p$ (the relation induce a connected relation over premise sequents).

We write this multicut rule as:

$$\frac{s_1 \quad \dots \quad s_n}{s} \text{mcut}(\iota, \perp\!\!\!\perp)$$

B.6 μMALL^∞ cut-reductions and cut-elimination

μMALL^∞ cut-reduction rules

Definition 65. External reductions are defined in fig. 9.

Definition 66. Internal reductions are the principal reductions given in ?? together with the following two reductions:

- the merge ($\text{mcut}/(\text{Cut})$) reduction

$$\frac{\frac{\frac{\frac{\vdash \Delta, F \quad \vdash \Gamma, F^\perp}{\vdash \Delta, \Gamma} (\text{Cut})}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\frac{\frac{\vdash \Delta, F \quad \vdash \Gamma, F^\perp}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}$$

where $\perp\!\!\!\perp'$ extends $\perp\!\!\!\perp$ with $F \perp\!\!\!\perp' F^\perp$ and $r = (\text{merge}, \{F, F^\perp\})$.

- the axiom reduction ($\text{mcut}/(\text{Ax})$) defined as follows, with F, F' and F'' being three occurrences of the same formula, that we distinguish with primes instead of explicitly describing their positions in sequents.

$$\frac{\frac{\frac{\frac{\vdash F, F'^\perp}{\vdash \Sigma} (\text{Ax})}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\frac{\vdash F'', \Gamma}{\vdash \Sigma} \text{mcut}(\iota', \perp\!\!\!\perp')}{\vdash \Sigma} \text{mcut}(\iota', \perp\!\!\!\perp')}$$

where $r = (\text{CutAx}, \{F, F'^\perp\})$, $F'^\perp \perp\!\!\!\perp F''$ and $\iota', \perp\!\!\!\perp'$ are defined as follows:

- for all $G \in \Sigma$, if $\iota(G) = F$ then $\iota'(G) = F''$, otherwise $\iota'(G) = \iota(G)$;
- $\perp\!\!\!\perp' = \perp\!\!\!\perp \cup \{\{F'', G\} \mid \{F, G\} \in \perp\!\!\!\perp\}$.

$$\begin{array}{c}
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F \quad \vdash \Gamma, G}{\vdash \Delta, \Gamma, F \otimes G} (\otimes)}{\vdash \Sigma_{\Delta}, \Sigma_{\Gamma}, F \otimes G} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma_{\Delta}, \Sigma_{\Gamma}, F \otimes G} \text{mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\frac{\mathcal{C}_{\Delta} \quad \vdash \Delta, F}{\vdash \Sigma_{\Delta}, F} \text{mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma_{\Delta}, \Sigma_{\Gamma}, F \otimes G} \quad \frac{\mathcal{C}_{\Gamma} \quad \vdash \Gamma, G}{\vdash \Sigma_{\Gamma}, G} \text{mcut}(\iota'', \perp\!\!\!\perp)}{\vdash \Sigma_{\Delta}, \Sigma_{\Gamma}, F \otimes G} (\otimes) \\
\\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F, G}{\vdash \Delta, F \wp G} (\wp)}{\vdash \Sigma, F \wp G} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, F \wp G} \text{mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\frac{\mathcal{C} \quad \vdash \Delta, F, G}{\vdash \Sigma, F, G} \text{mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma, F \wp G} (\wp) \\
\\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F \quad \vdash \Delta, G}{\vdash \Delta, F \& G} (\&)}{\vdash \Sigma, F \& G} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, F \& G} \text{mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\frac{\mathcal{C} \quad \vdash \Delta, F}{\vdash \Sigma, F} \text{mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma, F \& G} \quad \frac{\mathcal{C} \quad \vdash \Delta, G}{\vdash \Sigma, G} \text{mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma, F \& G} (\&) \\
\\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F_i}{\vdash \Delta, F_1 \oplus F_2} (\oplus^i)}{\vdash \Sigma, F_1 \oplus F_2} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, F_1 \oplus F_2} \text{mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\frac{\mathcal{C} \quad \vdash \Delta, F_i}{\vdash \Sigma, F_i} \text{mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma, F_1 \oplus F_2} (\oplus^i) \\
\\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F[\sigma X.F/X]}{\vdash \Delta, \sigma X.F} (\sigma)}{\vdash \Sigma, \sigma X.F} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, \sigma X.F} \text{mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\frac{\mathcal{C} \quad \vdash \Delta, F[\sigma X.F/X]}{\vdash \Sigma, F[\sigma X.F/X]} \text{mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma, \sigma X.F} (\sigma) \\
\\
\frac{\mathcal{C} \quad \frac{\overline{\vdash \Delta, \top} (\top)}{\vdash \Sigma, \top} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, \top} \text{mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\overline{\vdash \Sigma, \top} (\top)}{\vdash \Sigma, \top} (\top) \\
\\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta}{\vdash \Delta, \perp} (\perp)}{\vdash \Sigma, \perp} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma, \perp} \text{mcut}(\iota, \perp\!\!\!\perp)} \xrightarrow{r} \frac{\frac{\mathcal{C} \quad \vdash \Delta}{\vdash \Sigma} \text{mcut}(\iota', \perp\!\!\!\perp)}{\vdash \Sigma, \perp} (\perp) \quad \frac{\overline{\vdash 1} (1)}{\vdash 1} \text{mcut}(\iota, \perp\!\!\!\perp) \xrightarrow{r} \overline{\vdash 1} (1)
\end{array}$$

In all the above rules, one requires that ι' and ι'' are the ancestry relation induced by ι after the possible modification of the premises when taking restrictions of the premises as in the \otimes commutation and relate the subformulas of the commuted formula occurrence that are not in the premise and conclusion of the multicut.

Fig. 9. External reduction rules, where $r = (\text{ext}, F)$ and F is the principal occurrence.

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F \quad \vdash \Gamma, G}{\vdash \Delta, \Gamma, F \otimes G} (\otimes) \quad \frac{\vdash \Theta, F^\perp, G^\perp}{\vdash \Theta, F^\perp \wp G^\perp} (\wp)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash \Delta, F \quad \vdash \Gamma, G \quad \vdash \Theta, F^\perp, G^\perp}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)'}{\vdash \Sigma}$$

where $F \otimes G \perp\!\!\!\perp F^\perp \wp G^\perp$ and $\perp\!\!\!\perp'$ coincides with $\perp\!\!\!\perp$ except for $F \perp\!\!\!\perp' F^\perp$ and $G \perp\!\!\!\perp' G^\perp$

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F_1 \quad \vdash \Delta, F_2}{\vdash \Delta, F_1 \& F_2} (\&) \quad \frac{\vdash \Gamma, F_i^\perp}{\vdash \Gamma, F_1^\perp \oplus F_2^\perp} (\oplus^i)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash \Delta, F_i \quad \vdash \Gamma, F_i^\perp}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)'}{\vdash \Sigma}$$

If $\{F_1 \& F_2, F_1^\perp \oplus F_2^\perp\} \in \perp\!\!\!\perp$, where: $\perp\!\!\!\perp' \triangleq \perp\!\!\!\perp \cup \{\{F_i, F_i^\perp\}\}$.

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F[\mu X.F/X]}{\vdash \Delta, \mu X.F} (\mu) \quad \frac{\vdash \Gamma, F^\perp[\nu X.F^\perp/X]}{\vdash \Gamma, \nu X.F^\perp} (\nu)}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\vdash \Sigma} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash \Delta, F[\mu X.F/X] \quad \vdash \Gamma, F^\perp[\nu X.F^\perp/X]}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)'}{\vdash \Sigma}$$

where $\mu X.F \perp\!\!\!\perp \nu X.F^\perp$ and $\perp\!\!\!\perp'$ coincides with $\perp\!\!\!\perp$ except for $F[\mu X.F/X] \perp\!\!\!\perp' F^\perp[\nu X.F^\perp/X]$

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash \Gamma}{\vdash \Gamma, \perp} (\perp) \quad \frac{}{\vdash 1} (1)}{\Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\Sigma} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash \Gamma}{\Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\Sigma} \quad \text{where } \perp \perp\!\!\!\perp 1$$

Fig. 10. Principal reductions, where $r = (\text{princ}, \{F, F^\perp\})$ with $\{F, F^\perp\}$ the principal occurrences that have been reduced for the μMALL rules.

B.7 Adapting Baelde *et al.* μMALL^∞ cut-elimination proof to sequents as lists (Proof of Theorem 23).

We here sketch how to transfer the cut-elimination result for μMALL^∞ formulated as in Baelde et al [4,3] where sequents are sets of occurrences (ie of pairs of a formula and of an address) and the system of the present paper where sequents are ordered lists of formulas, with the exchange rule.

Let us refer to proofs with sequents as sets of occurrences as μMALL_o^∞ , proofs with sequents as ordered lists as μMALL_l^∞ for clarity.

- It is straightforward to lift the μMALL_l^∞ proof system considered in the present paper to a proof system in which sequents are *ordered lists of occurrences* (following the methodology described in [4] that we do not recall here.) Call this system $\mu\text{MALL}_{ol}^\infty$.
- One can define two forgetful maps:
 - $(\cdot)^{\text{occ}} : \mu\text{MALL}_{ol}^\infty \longrightarrow \mu\text{MALL}_l^\infty$
 - $(\cdot)^{\text{list}} : \mu\text{MALL}_{ol}^\infty \longrightarrow \mu\text{MALL}_o^\infty$

that respectively erase the addresses or the ordering of the formulas (together with the exchange rules, accordingly).

- Both of those maps are validity preserving and are onto.
- It is straightforward that:
 - If $\pi, \pi' \in \mu\text{MALL}_{ol}^\infty$ such that $\pi \xrightarrow{\text{cut}}^{\mu\text{MALL}_{ol}^\infty} \pi'$, then $(\pi)^{\text{occ}} \xrightarrow{\text{cut}}^{\mu\text{MALL}_{ol}^\infty} (\pi')^{\text{occ}}$ and if $\pi, \pi' \in \mu\text{MALL}_{ol}^\infty$ such that $\pi \xrightarrow{\text{cut}}^{\mu\text{MALL}_l^\infty} \pi'$, then either $(\pi)^{\text{list}} = (\pi')^{\text{list}}$ or $(\pi)^{\text{list}} \xrightarrow{\text{cut}}^{\mu\text{MALL}_o^\infty} (\pi')^{\text{list}}$. (Indeed, in the case of an exchange commutation for instance, the two proofs are translated to the same proof.)
- It is straightforward as well that there exists no infinite cut-reduction sequence made of only exchange commutations.
- Now, consider some proof π in μMALL_l^∞ and consider an antecedent of π by $(\cdot)^{\text{occ}}$ and consider a fair cut-reduction sequence σ from π . Clearly, σ can be lifted to a sequence σ_o in $\mu\text{MALL}_{ol}^\infty$, preserving fairness as reduction steps exactly match.
- Finally, $(\sigma_o)^{\text{list}}$ is a fair reduction sequence in μMALL_o^∞ , which is therefore productive and has a valid cut-free proof as a limit, from which we conclude that σ_o and finally σ has a cut-free and valid limit.

B.8 Residuals

Definition 67 (Residual). *Let π be a μMALL pre-proof such that $\pi \xrightarrow{r_1^{p_1}} \pi_1$ and $\pi \xrightarrow{r_2^{p_2}} \pi_2$. We define the notion of residual of (r_1, p_1) after (r_2, p_2) , written $(r_1, p_1)/(r_2, p_2)$ as follows:*

- if p_1, p_2 are incomparable, then $(r_1, p_1)/(r_2, p_2) = \{(r_1, p_1)\}$

- if $p_1 = p_2 = p$ then $(r_1, p_1)/(r_2, p_2) =$

$$\left\{ \begin{array}{ll} \emptyset & \text{if } r_1 \text{ and } r_2 \text{ are the same reduction step} \\ \emptyset & \text{if } r_1 \text{ and } r_2 \text{ are both (Cut)/(Ax) key cases involving two distinct axioms whi} \\ \{(r_1, p_1)\} & \text{if } r_2 \text{ is a principal reduction} \\ \{(r_1, p_1 \cdot 1), (r_1, p_1 \cdot 2)\} & \text{if } r_2 \text{ is a } (\&) \text{ commutation} \\ \{(r_1, p_1 \cdot 0)\} & \text{if } r_2 \text{ is a } (\oplus), (\&S) \text{ or } (\sigma) \text{ commutation} \\ \{(r_1, p_1 \cdot 1)\} & \text{if } r_2 = (\text{ext}, (\otimes)) \text{ and } r_1 \text{ is cut-connected to the } \otimes \text{ left context} \\ \{(r_1, p_1 \cdot 2)\} & \text{if } r_2 = (\text{ext}, (\otimes)) \text{ and } r_1 \text{ is cut-connected to the } \otimes \text{ right context} \end{array} \right.$$

B.9 Details on the compression of μMALL^∞ cut-reduction sequence: Proof of Proposition 26

The following lemma is a trivial observation:

Lemma 68 (Cut-reductions at incomparable positions commute). *Let π a μMALL^∞ pre-proof such that $\pi \xrightarrow{r_1^{p_1}} \pi_1$, $\pi \xrightarrow{r_2^{p_2}} \pi_2$ with p_1, p_2 incomparable. Then there exists π_3 such that $\pi_1 \xrightarrow{r_2^{p_2}} \pi_3$ and $\pi_2 \xrightarrow{r_1^{p_1}} \pi_3$.*

Lemma 69. *Let $\sigma = (\pi_i, r_i, p_i)_{i \in \omega+1}$ be a strongly converging μMALL^∞ cut-reduction sequence, with $\pi_\omega \xrightarrow{r_\omega^{p_\omega}} \pi_{\omega+1}$. There exists a strongly converging μMALL^∞ cut-reduction sequence $\text{Comp}(\sigma) = (\pi'_i, r'_i, p'_i)_{i \in \omega}$ which converges to $\pi_{\omega+1}$. More precisely, there exists some j such that*

- for any $i < j$, $(\pi'_i, r'_i, p'_i) = (\pi_i, r_i, p_i)$
- $r'_j = r_\omega$, $p'_j = p_\omega$
- and $r'_{i+1} = r_i$, $p'_{i+1} = p_i$ for any $i \geq j$ and
- $\text{Comp}(\sigma)$ strongly converges to $\pi_{\omega+1}$

Proof. The lemma is a direct consequence of Lemma 68 since, by strong convergence, there exists some j such that for any $i \geq j$, p_i is longer than p_ω , therefore the reduction r_ω occurring at p_ω in π_ω can already be performed in π_j since the root of the subproof at p_ω will never be modified in $(\pi_i, r_i, p_i)_{j \leq i \in \omega}$, therefore there must be a multicut at p_ω in π_i for every $j \leq i \in \omega$ and since this is there is at most one multicut per branch, all p_i for $j \leq i \in \omega$ are incomparable with p_ω .

We can now prove that μMALL^∞ cut-reduction is CP:

Proposition 70 (Compression lemma). *Let $\sigma = (\pi_i, r_i, p_i)_{i \in \alpha}$ be a strongly converging μMALL^∞ transfinite cut-reduction sequence. There exists a μMALL^∞ cut-reduction sequence $\text{Comp}(\sigma) = (\pi'_i, r'_i, p'_i)_{i \in \beta}$ which is a reordering of σ , depth-increasing, strongly converging with the same limit as σ and such that if α is finite, then $\beta = \alpha$ and $\beta = \omega$ otherwise.*

Proof. By transfinite induction on the length of the reduction σ that we want to compress and make depth-increasing.

- If α is a successor ordinal $\alpha = \lambda + 1$, the property follows directly from Lemma 68: applying induction hypothesis to the prefix of length λ results in a reduction $\sigma' = (\pi'_i, r'_i, p'_i)_{i \in \beta'}$ which is depth-increasing and such that every redex fired in the prefix of σ of length λ is still present. In particular, σ' strongly converges to π_α .

σ' being depth-increasing, there exists some j such that for any $i < j$, $\text{dpth}(p'_i) \leq \text{dpth}(p_\alpha)$ and for any $j \leq i \in \beta'$, $\text{dpth}(p'_i) > \text{dpth}(p_\alpha)$, therefore the reduction r_α occurring at p_α in π_α can already be performed in π'_j since the root of the subproof at p_α will never be modified in $(\pi'_i, r'_i, p'_i)_{j \leq i \in \beta'}$, therefore there must be a multicut at p_α in π'_i for every $j \leq i \in \beta'$. Since there is at most one multicut per branch, all p'_i for $j \leq i \in \beta'$ are incomparable with p_α . In particular, $\pi'_j \xrightarrow{r_\alpha} \pi''_j$ and for any $j \leq i \in \beta'$ $\pi''_i \xrightarrow{r'_i} \pi''_{i+1}$ by using the commutation property.

In particular $\sigma'' = (\pi''_i, r''_i, p''_i)_{j \leq i \in \beta'}$ strongly converges to $\pi_{\alpha+1}$. Considering the reduction sequence $(\pi'_i, r'_i, p'_i)_{i < j} \cdot (\pi'_j, p_\alpha, r_\alpha) \cdot \sigma''$ one get a depth-increasing reduction sequence indexed by $\beta = 1 + \beta'$ (due to the additional reduction step $(\pi'_j, p_\alpha, r_\alpha)$ inserted in the reduction sequence) with is such that every redex of σ is still present. Finally, notice that α was a finite ordinal, $\beta = 1 + \beta' = 1 + \lambda = \lambda + 1 = \alpha$ as expected and if α is infinite, $\beta' = \omega$ and $\beta = 1 + \omega = \omega$.

- If α is a limit ordinal.

Let us first observe the following fact: if d is the minimal depth of a reduction step in σ , by strong convergence, there is a finite number n_d of reduction steps of σ at depth d and there is an ordinal $\alpha' < \alpha$ such that for any $\alpha' \in i \in \alpha$, $\text{dpth}(p_i) > d$. By applying the transfinite induction hypothesis to $\sigma_{\alpha'} = (\pi_j, r_j, p_j)_{j \in \alpha'}$, one gets a depth-increasing reduction sequence $\sigma'_{\beta'} = (\pi'_i, r'_i, p'_i)_{i \in \beta'}$ with the same limit and containing the same reduction steps as $\sigma_{\alpha'}$ and which is indexed by $\min(\alpha', \omega)$. In particular, $\sigma'_{\beta'} = \sigma'_d \cdot \sigma''_d$ with all reduction steps in σ'_d being at depth d and all reduction steps in σ''_d being at depth at least $d + 1$. Therefore $\sigma''_d \cdot (\pi_i, r_i, p_i)_{\alpha' \in i \in \alpha}$ is a strongly converging sequence of length at most α and having all its reduction steps at depth at least $d + 1$. Then one builds by induction on $i \in \omega$, $(k_i, \sigma_i, \sigma'_i)_{i \in \omega}$ such that:

1. for any $i \in \omega$, $\sigma_i = (\pi'_j, r'_j, p'_j)_{j \in \alpha_i}$ with $\alpha_i \leq \alpha$ and for any $j \in \alpha_i$, $\text{dpth}(p'_j) \geq i$;
2. for any $i \in \omega$, $k_i \leq k_{i+1}$, $\sigma'_i = (\pi'_j, r'_j, p'_j)_{k_i \leq j < k_{i+1}}$ and if $k_i \leq j < k_{i+1}$, $\text{dpth}(p'_j) = i$;
3. for any $i \in \omega$, $(\pi'_j, r'_j, p'_j)_{j < k_i} \cdot \sigma_i$ is a strongly converging μMALL^∞ cut reduction sequence that contains the same reduction steps as σ up to reordering.

Initialization. We set $\sigma_0 = \sigma$ and $k_0 = 0$. By the above observation with $d = 0$, we get an $\alpha'_0 \leq \alpha$ we build σ'_0 , we set $k_1 = k_0 + n_0$ and set $\sigma_1 = \sigma''_0 \cdot (\pi'_j, r'_j, p'_j)_{\alpha'_0 \in j \in \alpha}$. Conditions 1–3 are satisfied for $i = 0$.

Heredity. Assuming $(k_i)_{i \leq l}, (\sigma_i)_{i \leq l}, (\sigma'_i)_{i < l}$ are built, satisfying the conditions, one proceeds as above taking σ_l for σ and $l + 1$ for d and setting $k_{l+1} = k_l + n_l$ and σ'_l being provided as before and σ_{l+1} as the concatenation of σ''_l and $(\pi'_j, r'_j, p'_j)_{\alpha'_l \in j \in \alpha}$, satisfying the required conditions.

Finally, we set $\sigma' = (\pi'_j, r'_j, p'_j)_{j \in \omega}$ which is trivially depth increasing, strongly converging, and which is a reordering of the reduction steps of σ : any reduction step of σ' comes from a reduction in σ and for any reduction step (r_i, p_i) in σ , it occurs by construction in $\sigma'_{\text{dpth}(p_i)}$ and therefore in σ' .

C Appendix on Section 3

C.1 Details on μLL^∞ cut-reduction rules

We detail the rules of μLL^∞ cut elimination.

Definition 71. *Cut commutations or External reductions are defined in fig. 11.*

Definition 72. *Key cut-reductions, or internal reductions, are the principal reductions given in fig. 12 together with the following two reductions:*

- the merge (mcut)/(Cut) reduction

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F \quad \vdash \Gamma, F^\perp}{\vdash \Delta, \Gamma} \text{ (Cut)}}{\vdash \Sigma} \text{ mcut}(\iota, \perp)}{\vdash \Sigma} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash \Delta, F \quad \vdash \Gamma, F^\perp}{\vdash \Sigma} \text{ mcut}(\iota, \perp')}{\vdash \Sigma}$$

where \perp' extends \perp with $F \perp' F^\perp$ and $r = (\text{merge}, \{F, F^\perp\})$.

- the axiom reduction (mcut)/(Ax) defined as follows, with F, F' and F'' being three occurrences of the same formula, that we distinguish with primes instead of explicitly describing their positions in sequents.

$$\frac{\mathcal{C} \quad \frac{\frac{}{\vdash F, F'^\perp} \text{ (Ax)}}{\vdash \Sigma} \quad \vdash F'', \Gamma \text{ mcut}(\iota, \perp)}{\vdash \Sigma} \xrightarrow{r} \frac{\mathcal{C} \quad \frac{\vdash F'', \Gamma}{\vdash \Sigma} \text{ mcut}(\iota', \perp')}{\vdash \Sigma}$$

where $r = (\text{CutAx}, \{F, F'^\perp\})$, $F'^\perp \perp F''$ and ι', \perp' are defined as follows:

- for all $G \in \Sigma$, if $\iota(G) = F$ then $\iota'(G) = F''$, otherwise $\iota'(G) = \iota(G)$;
- $\perp' = \perp \cup \{\{F'', G\} \mid \{F, G\} \in \perp\}$.

C.2 Proof of Proposition 33

Proposition 73 (Preservation of validity). *π is a valid μLL^∞ proof of $\vdash \Gamma$ iff π^\bullet is a valid μMALL^∞ proof of $\vdash \Gamma^\bullet$.*

Proof. It is simply a matter of relating the infinite branches in both pre-proofs: Since every infinite branch in π gives rise to *at least* an infinite branch of π^\bullet , validity of π^\bullet entails that of π . Conversely, let us prove that all infinite branches of π^\bullet are valid. Let us consider three types of infinite branches:

- Let us consider an infinite branch β of π^\bullet such that a suffix of β loops around the backedges of the encoding of some promotion introducing $!F$ in π . In that case β is valid as it is inhabited by a thread having as least recurring formula $(!F)^\bullet$ (indeed, in such a case, the only recurring formulas of the thread are $(!F)^\bullet, F^\bullet \& (1 \& (!F)^\bullet \otimes (!F)^\bullet), 1 \& (!F)^\bullet \otimes (!F)^\bullet$ and $(!F)^\bullet \otimes (!F)^\bullet$)

$$\begin{array}{c}
\frac{\frac{\mathcal{C} \quad \frac{\vdash \Delta, F}{\vdash \Delta, ?F} \text{ (?d)}}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp)}{\vdash \Delta, ?F, ?F} \text{ (?c)} \quad \longrightarrow_r \quad \frac{\mathcal{C} \quad \frac{\vdash \Delta, F}{\vdash \Sigma, F} \text{ mcut}(\iota', \perp)}{\vdash \Sigma, ?F} \text{ (?d)} \\
\frac{\mathcal{C} \quad \frac{\vdash \Delta, ?F}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp)}{\vdash \Delta, ?F, ?F} \text{ (?c)} \quad \longrightarrow_r \quad \frac{\mathcal{C} \quad \frac{\vdash \Delta, F}{\vdash \Sigma, ?F, ?F} \text{ mcut}(\iota', \perp)}{\vdash \Sigma, ?F} \text{ (?c)} \\
\frac{\mathcal{C} \quad \frac{\vdash \Delta, ?F}{\vdash \Sigma, ?F} \text{ mcut}(\iota, \perp)}{\vdash \Delta, ?F} \text{ (?w)} \quad \longrightarrow_r \quad \frac{\mathcal{C} \quad \frac{\vdash \Delta}{\vdash \Sigma} \text{ mcut}(\iota', \perp)}{\vdash \Sigma, ?F} \text{ (?w)} \\
\frac{\frac{\vdash F, ?\Gamma}{\vdash !F, ?\Gamma} \text{ (!p)} \quad \mathcal{C}^!}{\vdash !F, ?\Sigma} \text{ mcut}(\iota, \perp)}{\vdash \Delta, F \quad \vdash \Gamma, G} \text{ (?p)} \quad \longrightarrow_r \quad \frac{\frac{\mathcal{C}_\Delta \quad \frac{\vdash \Delta, F}{\vdash \Sigma_\Delta, F} \text{ mcut}(\iota', \perp)}{\vdash \Sigma_\Delta, \Sigma_\Gamma, F \otimes G} \text{ (?p)} \quad \frac{\mathcal{C}_\Gamma \quad \frac{\vdash \Gamma, G}{\vdash \Sigma_\Gamma, G} \text{ mcut}(\iota'', \perp)}{\vdash \Sigma_\Delta, \Sigma_\Gamma, F \otimes G} \text{ (?p)}}{\vdash \Sigma_\Delta, \Sigma_\Gamma, F \otimes G} \text{ (?p)} \\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F, G}{\vdash \Delta, F \wp G} \text{ (?8)}}{\vdash \Sigma, F \wp G} \text{ mcut}(\iota, \perp)}{\vdash \Delta, F, G} \text{ (?8)} \quad \longrightarrow_r \quad \frac{\mathcal{C} \quad \frac{\vdash \Delta, F, G}{\vdash \Sigma, F, G} \text{ mcut}(\iota', \perp)}{\vdash \Sigma, F \wp G} \text{ (?8)} \\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F \quad \vdash \Delta, G}{\vdash \Delta, F \& G} \text{ (?&)}}{\vdash \Sigma, F \& G} \text{ mcut}(\iota, \perp)}{\vdash \Delta, F} \text{ (?&)} \quad \longrightarrow_r \quad \frac{\frac{\mathcal{C} \quad \frac{\vdash \Delta, F}{\vdash \Sigma, F} \text{ mcut}(\iota', \perp)}{\vdash \Sigma, F \& G} \text{ (?&)} \quad \frac{\mathcal{C} \quad \frac{\vdash \Delta, G}{\vdash \Sigma, G} \text{ mcut}(\iota', \perp)}{\vdash \Sigma, F \& G} \text{ (?&)}}{\vdash \Sigma, F \& G} \text{ (?&)} \\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F_i}{\vdash \Delta, F_1 \oplus F_2} \text{ (?}\oplus\text{)}}{\vdash \Sigma, F_1 \oplus F_2} \text{ mcut}(\iota, \perp)}{\vdash \Delta, F_i} \text{ (?}\oplus\text{)} \quad \longrightarrow_r \quad \frac{\mathcal{C} \quad \frac{\vdash \Delta, F_i}{\vdash \Sigma, F_i} \text{ mcut}(\iota', \perp)}{\vdash \Sigma, F_1 \oplus F_2} \text{ (?}\oplus\text{)} \\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F[\sigma X.F/X]}{\vdash \Delta, \sigma X.F} \text{ (?}\sigma\text{)}}{\vdash \Sigma, \sigma X.F} \text{ mcut}(\iota, \perp)}{\vdash \Delta, F[\sigma X.F/X]} \text{ (?}\sigma\text{)} \quad \longrightarrow_r \quad \frac{\mathcal{C} \quad \frac{\vdash \Delta, F[\sigma X.F/X]}{\vdash \Sigma, F[\sigma X.F/X]} \text{ mcut}(\iota', \perp)}{\vdash \Sigma, \sigma X.F} \text{ (?}\sigma\text{)} \\
\frac{\mathcal{C} \quad \frac{\overline{\vdash \Delta, \top} \text{ (?}\top\text{)}}{\vdash \Sigma, \top} \text{ mcut}(\iota, \perp)}{\vdash \Delta, \top} \text{ (?}\top\text{)} \quad \longrightarrow_r \quad \overline{\vdash \Sigma, \top} \text{ (?}\top\text{)} \\
\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta}{\vdash \Delta, \perp} \text{ (?}\perp\text{)}}{\vdash \Sigma, \perp} \text{ mcut}(\iota, \perp)}{\vdash \Delta, \perp} \text{ (?}\perp\text{)} \quad \longrightarrow_r \quad \frac{\mathcal{C} \quad \frac{\vdash \Delta}{\vdash \Sigma} \text{ mcut}(\iota', \perp)}{\vdash \Sigma, \perp} \text{ (?}\perp\text{)} \quad \frac{\overline{\vdash 1} \text{ (?}1\text{)}}{\vdash 1} \text{ mcut}(\iota, \perp)}{\vdash 1} \quad \longrightarrow_r \quad \overline{\vdash 1} \text{ (?}1\text{)}
\end{array}$$

In the promotion commutation case, all premisses in $\mathcal{C}^!$ are assumed to end with promotion rules.

In all the above rules, one requires that ι' and ι'' are the ancestry relation induced by ι after the possible modification of the premisses when taking restrictions of the premisses as in the \otimes commutation and relate the subformulas of the commuted formula occurrence that are not in the premise and conclusion of the multicut. (For instance, for the contraction commutation, each occurrence of $?F$ is related to the occurrence in the corresponding position in the premise.)

Fig. 11. External reduction rules, where $r = (\text{ext}, F)$ and F is the principal occurrence.

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash F, \Gamma}{\vdash ?F, \Gamma} \text{ (?d)} \quad \frac{\vdash F^\perp, ?\Delta}{\vdash !F^\perp, ?\Delta} \text{ (!p)}}{\vdash \Sigma} \quad \text{mcut}(\iota, \perp\!\!\!\perp)} \longrightarrow_r \frac{\mathcal{C} \quad \frac{\vdash F, \Gamma \quad \vdash F^\perp, ?\Delta}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)'}{\vdash \Sigma}$$

where $?F \perp\!\!\!\perp !F^\perp$ and $\perp\!\!\!\perp'$ coincides with $\perp\!\!\!\perp$ except for $F \perp\!\!\!\perp' F^\perp$.

$$\frac{\mathcal{C}_\Gamma \quad \mathcal{C}_{?F}^! \quad \frac{\frac{\vdash ?F, ?F, \Gamma}{\vdash ?F, \Gamma} \text{ (?c)}}{\vdash \Gamma', ?\Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)} \longrightarrow_r \frac{\mathcal{C}_\Gamma \quad \mathcal{C}_{?F}^! \quad \mathcal{C}_{?F}^! \quad \frac{\vdash ?F, ?F, \Gamma}{\vdash \Gamma', ?\Sigma, ?\Sigma} \text{mcut}(\iota', \perp\!\!\!\perp')}{\frac{\vdash \Gamma', ?\Sigma}{\vdash \Gamma', ?\Sigma} \text{ (?c) }^*}$$

where $\mathcal{C}_{?F}^! \neq \emptyset$, $\perp\!\!\!\perp'$ corresponds to $\perp\!\!\!\perp$ on \mathcal{C}_Γ and is a “duplication” of $\perp\!\!\!\perp$ on $\mathcal{C}_{?F}^!$ and each copy of $?F$ is in $\perp\!\!\!\perp'$ -relation with the corresponding copy of $!F^\perp$ in $\mathcal{C}_{?F}^!$.

$$\frac{\mathcal{C}_\Gamma \quad \mathcal{C}_{?F}^! \quad \frac{\frac{\vdash \Gamma}{\vdash ?F, \Gamma} \text{ (?w)}}{\vdash \Gamma', ?\Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)} \longrightarrow_r \frac{\mathcal{C}_\Gamma \quad \frac{\vdash \Gamma}{\vdash \Gamma'} \text{mcut}(\iota', \perp\!\!\!\perp')}{\vdash \Gamma', ?\Sigma} \text{ (?w) }^*$$

where $\mathcal{C}_{?F}^! \neq \emptyset$ and $\perp\!\!\!\perp'$ corresponds to the restriction of $\perp\!\!\!\perp$ on \mathcal{C}_Γ , Γ .

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F \quad \vdash \Gamma, G}{\vdash \Delta, \Gamma, F \otimes G} \text{ (?\otimes)} \quad \frac{\vdash \Theta, F^\perp, G^\perp}{\vdash \Theta, F^\perp \wp G^\perp} \text{ (?\wp)}}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)} \longrightarrow_r \frac{\mathcal{C} \quad \frac{\vdash \Delta, F \quad \vdash \Gamma, G \quad \vdash \Theta, F^\perp, G^\perp}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)'}{\vdash \Sigma}$$

where $F \otimes G \perp\!\!\!\perp F^\perp \wp G^\perp$ and $\perp\!\!\!\perp'$ coincides with $\perp\!\!\!\perp$ except for $F \perp\!\!\!\perp' F^\perp$ and $G \perp\!\!\!\perp' G^\perp$

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F_1 \quad \vdash \Delta, F_2}{\vdash \Delta, F_1 \& F_2} \text{ (?\&)} \quad \frac{\vdash \Gamma, F_i^\perp}{\vdash \Gamma, F_1^\perp \oplus F_2^\perp} \text{ (?\oplus)}}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)} \longrightarrow_r \frac{\mathcal{C} \quad \frac{\vdash \Delta, F_i \quad \vdash \Gamma, F_i^\perp}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)'}{\vdash \Sigma}$$

If $\{F_1 \& F_2, F_1^\perp \oplus F_2^\perp\} \in \perp\!\!\!\perp$, where: $\perp\!\!\!\perp' \triangleq \perp\!\!\!\perp \cup \{\{F_i, F_i^\perp\}\}$.

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash \Delta, F[\mu X.F/X]}{\vdash \Delta, \mu X.F} \text{ (?\mu)} \quad \frac{\vdash \Gamma, F^\perp[\nu X.F^\perp/X]}{\vdash \Gamma, \nu X.F^\perp} \text{ (?\nu)}}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)} \longrightarrow_r \frac{\mathcal{C} \quad \frac{\vdash \Delta, F[\mu X.F/X] \quad \vdash \Gamma, F^\perp[\nu X.F^\perp/X]}{\vdash \Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)'}{\vdash \Sigma}$$

where $\mu X.F \perp\!\!\!\perp \nu X.F^\perp$ and $\perp\!\!\!\perp'$ coincides with $\perp\!\!\!\perp$ except for $F[\mu X.F/X] \perp\!\!\!\perp' F^\perp[\nu X.F^\perp/X]$

$$\frac{\mathcal{C} \quad \frac{\frac{\vdash \Gamma}{\vdash \Gamma, \perp} \text{ (?1)} \quad \frac{}{\vdash 1} \text{ (?1)}}{\Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)} \longrightarrow_r \frac{\mathcal{C} \quad \frac{\vdash \Gamma}{\Sigma} \text{mcut}(\iota, \perp\!\!\!\perp)}{\Sigma} \quad \text{where } \perp \perp\!\!\!\perp 1$$

Fig. 12. Principal reductions, where $r = (\text{princ}, \{F, F'^\perp\})$ with $\{F, F'^\perp\}$ the principal occurrences that have been reduced for the μ MALL rules and $r \in \{(\text{princ}, \{F, F'^\perp\}), (?d), (\text{princ}, \{F, F'^\perp\}), (?c), (\text{princ}, \{F, F'^\perp\}), (?w))\}$ for the exponential case.

In the principal reduction involving contraction and weakening, the restrictions $\mathcal{C}_{?F}$ are required to contain only proofs ending with a promotion rule.

- Otherwise, let us first consider the special case of an infinite branch that, each time it enters the encoding of a promotion, takes the left-most premise of the ($\&$) rule. To such an infinite branch it is easy to associate an infinite branch b of π which is valid, therefore supported by a thread the least formula of which is some $\nu X.F$. Then one observes that $(\nu X.F)^\bullet$ is recurring in β as is the least such formula: β is valid.
- In the last case, on each promotion, one shall first backedge a finite number of times on the backedges of the promotion before moving towards the encoding of the premise of the promotion. Note that during those loop phases, the subformulas of the context $(? \Delta)^\bullet$ are preserved as well as F^\bullet . This allows to find, as in the previous case, a ν -formula that is the least recurring formula along some thread.

C.3 Simulation of μLL^∞ cut-reduction: proof of Proposition 34

Proposition 74 (Simulation of μLL^∞ cut-elimination steps). *Each μLL^∞ mcut-reduction r can be simulated in μMALL^∞ by a (possibly infinite) sequence of mcut-reductions, denoted r^\bullet .*

Proof. We first assume that the position of the reduction is not above some (!p) rule. In that case, the simulation of μMALL^∞ reduction rules is trivial.

The simulation of exponential cut-reductions are shown in Figures 13 & 14 in the simplified case of cuts. The simulation for the general mcut case can easily be inferred: it does not modify the length of the simulated sequence for commutation rules nor in the (?d)/(!p) key reduction, while the simulation of contraction and weakening key reductions shall be adapted to take into account the length of the (!p)-ready context of the mcut rule.

In the case the position of the reduction step is above a certain number of (!p), the redex is in fact encoded as an infinity of redexes, which induce, in each case, a simulation of length ω done as a depth-increasing reduction, which is always possible.

C.4 Proof of Proposition 36

Proposition 75 (Non-preservation of fairness by simulation). *There exists a μLL^∞ (pre-)proof π and a fair reduction ρ from π such that ρ^\bullet is an ω -indexed μMALL^∞ cut-reduction sequence from π^\bullet which is not fair.*

Proof. Consider pre-proofs $\pi_k, k \in \omega$:

$$\pi_k = \frac{\frac{\frac{\vdots}{\vdash \nu Y.Y} \text{ } (\nu)}{\vdash ? \mu X.X, \nu Y.Y} \text{ } (?w)}{\frac{\frac{\frac{\frac{\vdots}{\vdash \nu X.X, ? B} \text{ } (\nu)}{\vdash ! \nu X.X, ? B} \text{ } (!p)}{\vdash ! B^\perp, \sigma Z.Z} \text{ } (\sigma)}{\vdash \nu Y.Y, \sigma Z.Z} \text{ } (\sigma)^k} \text{ } \text{mcut}(\iota, \perp\!\!\!\perp)}$$

$$\begin{array}{c}
\frac{\frac{\pi}{\vdash F, \Gamma}}{\vdash ?^\bullet F, \Gamma} \quad \frac{\frac{\pi'}{\vdash F^\perp, ?^\bullet \Delta}}{\vdash !^\bullet F^\perp, ?^\bullet \Delta} \quad \frac{(\text{!p}^\bullet)}{\longrightarrow^2} \quad \frac{\frac{\pi}{\vdash F, \Gamma} \quad \frac{\pi'}{\vdash F^\perp, ?^\bullet \Delta}}{\vdash \Gamma, ?^\bullet \Delta} \quad (\text{Cut}) \\
\frac{\frac{\pi}{\vdash ?^\bullet F, \Gamma}}{\vdash \Gamma, ?^\bullet \Delta} \quad \frac{\frac{\pi'}{\vdash F^\perp, ?^\bullet \Delta}}{\vdash !^\bullet F^\perp, ?^\bullet \Delta} \quad \frac{(\text{!p}^\bullet)}{\longrightarrow^{4int, 4 \times \# \Delta ext}} \quad \frac{\frac{\pi}{\vdash ?^\bullet F, ?^\bullet F, \Gamma} \quad \frac{\pi'}{\vdash !^\bullet F^\perp, ?^\bullet \Delta} \quad \frac{\pi'}{\vdash !^\bullet F^\perp, ?^\bullet \Delta}}{\vdash \Gamma, ?^\bullet \Delta, ?^\bullet \Delta} \quad (\text{Cut}) \quad (\text{?c}^\bullet)^* \quad (\text{mcut}) \\
\frac{\frac{\pi}{\vdash \Gamma}}{\vdash ?^\bullet F, \Gamma} \quad \frac{\frac{\pi'}{\vdash F^\perp, ?^\bullet \Delta}}{\vdash !^\bullet F^\perp, ?^\bullet \Delta} \quad \frac{(\text{?w}^\bullet)}{\longrightarrow^{3int, 3 \times \# \Delta ext}} \quad \frac{\frac{\pi}{\vdash \Gamma}}{\vdash \Gamma, ?^\bullet \Delta} \quad (\text{?w}^\bullet)^* \quad (\text{Cut})
\end{array}$$

Fig. 14. Simulation of μLL^∞ key-(m)cut rules

By the fairness assumption, that implies that there is a suffix τ of σ (in fact $(\pi_i, r_i, p_i)_{i \geq k}$) containing infinitely many cut-reduction steps occurring at the root of the subproof of π_i rooted in p ($\pi_i @ p$), for $i \geq k$. We want to view $\tau^p = (\pi_i @ p)_{i \geq k}$ as a reduction sequence. As such, this is not a μLL^∞ cut-reduction sequence as there are steps of τ reducing outside of the $\pi_i @ p$: in that case, $\pi_i @ p = \pi_{i+1} @ p$.

Let $I = \{i \geq k \mid p_i = p\}$ and let κ be an enumeration of the infinite set I . Since for any $j \geq k$ such that $j \notin I$, $\pi_j @ p = \pi_{j+1} @ p$, the sequence $\rho = (\pi_{\kappa(i)} @ p, r_{\kappa(i)}, \epsilon)_{i \in \omega}$ is a cut-reduction sequence which is a fair μLL^∞ cut-reduction sequence and which consists only of key-cut-reduction such that the multicuts remain at the root.

Consider the μMALL^∞ encoding of ρ , ρ^\bullet . One can first remark that it is an ω -indexed sequence since no reduction occurs above a promotion in ρ and since there is no commutation rule. ρ^\bullet is therefore unproductive, from which we want to draw a contradiction with μMALL^∞ cut-elimination theorem.

Is ρ^\bullet a fair reduction sequence? Proposition 36 shows that it is not necessarily the case: there may be available reduction steps which are not fired along ρ^\bullet (in fact they can only be $(\mu)/(\nu)$, $(\oplus)/(\&)$ key cases as they come only from the condition on the application of key cases for weakening and contraction) and which have residual along the whole sequence ρ^\bullet . Let σ be a reduction obtained from ρ^\bullet by firing those redexes causing the lack of fairness of ρ^\bullet : it is easy to notice that σ is fair. But remark that σ contains exactly the same commutation rules as ρ^\bullet , that is none, contradicting μMALL^∞ cut-elimination theorem.

Proof of Lemma 38

Lemma 77. *Let π be a μLL^∞ pre-proof of $\vdash \Gamma$ and let us consider a cut-reduction sequence $\sigma = (\pi_i, r_i, p_i)_{i \in \omega}$ in μLL^∞ initiated with π that converges to a cut-free μLL^∞ pre-proof π' . σ^\bullet is a strongly converging (possibly transfinite) sequence.*

Proof. Let λ be the indexing ordinal of σ^\bullet : $\sigma^\bullet = (\pi'_i)_{i \in \lambda}$. We want to prove that σ^\bullet is strongly converging. This comes from two facts:

- first, because π'^\bullet – the encoding of the limit of σ – is cut-free, the depth of the redexes tend to infinity when tending to the limit ordinal λ indexing σ^\bullet as σ is itself strongly converging.
- Second, because only !-commutations as well as reductions above a promotion can create infinite reduction sequence. Let λ' be a limit ordinal strictly below λ : the depths of the cuts reduced in π'_i tend to infinity as i tends to λ' as it corresponds either to the simulation of a ! commutation, which is strongly converging or to a reduction in an exponential box, which is also simulated by a strongly converging sequence.

Details on the cut-elimination proof, page 14

Theorem 78. *Fair μLL^∞ mcut-reduction sequences converge to cut-free μLL^∞ proofs.*

Proof. Let π be a μLL^∞ -proof of $\vdash \Gamma$ and let us consider a fair cut-reduction sequence $\sigma = (\pi_i, r_i, p_i)_{i \in \omega}$ in μLL^∞ initiated with π . We want to prove that σ has a limit, say π' , which is a cut-free μLL^∞ -proof of $\vdash \Gamma$.

Consider the associated (potentially transfinite) μMALL^∞ cut-reduction sequence σ^\bullet initiated in $\pi^\bullet \vdash \Gamma^\bullet$ obtained by simulating each cut-reduction step of σ . By Lemma 37, we know that σ converges to a cut-free μLL^∞ pre-proof. From this, one can conclude that σ is productive and therefore that σ *strongly converges* to some μLL^∞ cut-free pre-proof π' .

We are left with proving that the limit of σ, π' , is actually valid. By Lemma 38, σ^\bullet is a *transfinite* cut-reduction sequence from π^\bullet which *strongly converges* to π'^\bullet . Therefore, the Compression Lemma applies to π'^\bullet : there exists $\rho = (\pi'_i, r'_i, p'_i)_{i \in \omega}$ an ω -indexed μMALL^∞ cut-reduction sequence converging to π'^\bullet .

In order to apply μMALL^∞ cut-elimination theorem, we would need to ensure that ρ is fair, from which validity of π'^\bullet would follow and that of π' by Proposition 33. However, we know by Proposition 36 that it is not necessary the case. We can still conclude by considering what happens when ρ is made fair (in a reduction sequence called ρ_f). By remark 35, we know that the loss of fairness of ρ can only have one reason: the lack of (!p)-readiness, in a multicut $\text{mcut}(\iota, \perp\!\!\!\perp)$, of the cut-context associated to the context of a (!p) rule the principal formula of which (i) either is ι -related to a formula in the conclusion of the (mcut), therefore blocking a (!p) commutation rule, or (ii) is $\perp\!\!\!\perp$ -related with a ?F which is conclusion of a (?w) or of a (?c) rule, therefore blocking a key cut-reduction step of type (?w)/(!p) or (?c)/(!p). Indeed, in those two cases, the μMALL^∞ -encoding of the proof can perform partially the simulation of the reduction that is blocked.

We shall analyze more precisely what happen in this case to prove that to any infinite branch β of π'^{\bullet} , one can associate an infinite branch β_f of the limit, say π_f , of ρ_f (ρ_f being fair by assumption, it has a limit which is valid) containing the same infinite threads as β .

Before proceeding, let us notice that the content of remark 35 does not entail that the only reduction steps added to ρ in order to obtain ρ_f are the partial reduction of (the simulation of) the blocked reductions. Indeed, let us consider the case of a blocked (!p)-commutation:

$$\frac{\frac{!(\pi)}{\vdash !F, ?\Gamma} \text{ (!p)} \quad \mathcal{C}_{? \Gamma}}{\vdash !F, \Sigma} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\text{with } !(\pi) = \frac{\pi}{\vdash !F, ?\Gamma} \text{ (!p)}, \text{ assuming that } \mathcal{C}_{? \Gamma} \text{ is not (!p)-ready.}}$$

$$\text{The } \mu\text{MALL}^\infty \text{ encoding of the above proof is: } \frac{!(\pi^\bullet)}{\vdash !F^\bullet, \Sigma^\bullet} \text{ } (\mathcal{C}_{? \Gamma}^\bullet) \text{ mcut}(\iota, \perp\!\!\!\perp)$$

and since $!(\pi^\bullet)$ is concluded with $(\&)$ and (ν) rules on $!F^\bullet$, the corresponding commutation step can be performed (and will be performed in a fair reduction) resulting in

$$\frac{\frac{\mathcal{C}_{? \Gamma}^\bullet \quad \pi^\bullet}{\vdash F^\bullet, \Sigma^\bullet} \text{ mcut}(\iota, \perp\!\!\!\perp) \quad \frac{\frac{\overline{\vdash 1} \text{ (1)}}{\vdash 1, ?\Gamma^\bullet} \text{ (?w)}^{\bullet+}}{\vdash 1, \Sigma^\bullet} \text{ mcut}(\iota, \perp\!\!\!\perp) \quad \frac{\frac{!(\pi^\bullet) \quad !(\pi^\bullet)}{\vdash !F^\bullet \otimes !F^\bullet, ?\Gamma^\bullet, ?\Gamma^\bullet} \text{ } (\otimes)}{\vdash !F^\bullet \otimes !F^\bullet, ?\Gamma^\bullet} \text{ (?c)}^{\bullet+}}{\vdash !F^\bullet \otimes !F^\bullet, \Sigma^\bullet} \text{ mcut}(\iota, \perp\!\!\!\perp)}{\vdash !F^\bullet, \Sigma^\bullet} \text{ } (\nu) \text{ } (\&) \text{ } (\&)$$

In the left-most premise, the subproof π^\bullet is cut with $\mathcal{C}_{? \Gamma}^\bullet$ which could unlock productive new reduction steps.

On the other hand, considering the middle premise, one notices that this cannot happen as the (1) inference is guarded by at least one (encoding of) weakening (since Γ is non-empty, otherwise $\mathcal{C}_{? \Gamma}$ is (!p)-ready) which will not interact with some premise of $\mathcal{C}_{? \Gamma}^\bullet$ by non-readiness.

Similarly, analyzing the cases of blocked key cut-steps of type (?w)/(!p) or (?c)/(!p), one realizes that ρ_f could contain some additional $(\mu)/(\nu)$, $(\oplus^2)/(\&)$, $(\oplus^1)/(\&)$ steps or some additional $(\mu)/(\nu)$, $(\oplus^2)/(\&)$, $(\oplus^1)/(\&)$ steps respectively. In particular, the (encoding of the) premise of the blocked structural rule cannot produce part of the cut-free proof π_f as they are blocked respectively by $(?w)^\bullet$ and by $(?c)^\bullet$ respectively.

This remark is the key to conclude the proof: indeed, let β be an infinite branch of π'^{\bullet} and let us consider ρ_f a fair reduction sequence obtained from ρ by reducing those redexes which have residuals along the whole of ρ causing the lack of fairness of ρ (and doing that hereditarily) and let us consider the limit π_f of this sequence.

The above analysis of the blocking situations allows us to define a branch β_f of π_f as follows: β_f agrees with β as long as possible. When the next inference in π_f differ from that of β , that means that it comes from the commutations of inferences used in the simulation of a (!p) which is blocked.

In that case one requires β_f to follow the unique premise of the (ν) rule, the second premise of the $(\&)$ rule and finally the first premise of the following $(\&)$ rule, ending in sequent $\vdash 1, ?\Delta$ as described below:

$$\begin{array}{c}
\vdots \\
\hline
\vdash 1, ?\Delta^\bullet \quad \vdash (!F)^\bullet \otimes (!F)^\bullet, ?\Delta^\bullet \\
\hline
\vdash F^\bullet, ?\Delta^\bullet \quad \vdash 1 \& ((!F)^\bullet \otimes (!F)^\bullet), ?\Delta^\bullet \quad (\&) \\
\hline
\vdash F^\bullet \& (1 \& ((!F)^\bullet \otimes (!F)^\bullet)), ?\Delta^\bullet \quad (\&) \\
\hline
\vdash (!F)^\bullet, ?\Delta^\bullet \quad (\nu) \\
\hline
\vdots
\end{array}$$

and we follow back the branch according to β as before.

It is easy to notice that in β_f we have exactly the same (infinite) threads as in β (some finite thread may possibly have been extended *finitely* on occurrences of $(!F)^\bullet$): indeed, the only inferences on which β_f differs from β are of the type pictured in green above and therefore end in a 1 formula. Therefore β is valid if, and only if, β_f is valid, which is the case since ρ_f is fair.

We can then conclude that π'^\bullet is cut-free and valid and, using preservation of validity, that π' is a valid cut-free μLL^∞ -proof.

Proof of Corollary 39

Corollary 79. *Fair two-sided μLL^∞ mcut-reduction sequences converge to cut-free two-sided μLL^∞ proofs.*

Proof. A fair reduction sequence in two-sided μLL^∞ can easily be seen to map to a fair reduction sequence in one-sided μLL^∞ from which the result follows.

A two-sided sequent $\Gamma \vdash \Delta$ is mapped to $\vdash \Gamma^\perp, \Delta$ with $A \multimap B$ interpreted as $A^\perp \wp B$. Each right rule is mapped to the same rule and each left rule is mapped to the same rule for the de Morgan dual connective. In such a way, any two-sided μLL^∞ pre-proof π can be mapped to a one-sided μLL^∞ pre-proof π^{mono} preserving validity: π is valid if, and only if, π^{mono} is valid.

From the above remark, any fair reduction cut-reduction sequence $(\pi_i)_{i \in \omega}$ in two-sided μLL^∞ can be mapped to a fair one-sided cut-reduction sequence $(\pi_i^{\text{mono}})_{i \in \omega}$ which has a limit Π which is valid. It follows that $(\pi_i)_{i \in \omega}$ is productive and that its limit π satisfies $\pi^{\text{mono}} = \Pi$ from which validity follows.

D Appendix on Section 4

D.1 Details on skeletons

To any μLL^∞ formulas and μLL^∞ proofs, one can associate their classical skeletons, that is corresponding μLK^∞ formulas and proofs, after erasing of the linear information of the sequent and proof:

Definition 80 (μLK^∞ -Skeleton). *The μLK^∞ -skeleton of a μLL^∞ formula is defined inductively as follows ($\sigma \in \{\mu, \nu\}$):*

$$\begin{array}{l|l|l} \text{Sk}(A \otimes B) = \text{Sk}(A) \wedge \text{Sk}(B) & \text{Sk}(A \wp B) = \text{Sk}(A) \vee \text{Sk}(B) & \text{Sk}(!A) = \text{Sk}(A) \\ \text{Sk}(A \& B) = \text{Sk}(A) \wedge \text{Sk}(B) & \text{Sk}(A \oplus B) = \text{Sk}(A) \vee \text{Sk}(B) & \text{Sk}(?A) = \text{Sk}(A) \\ \text{Sk}(1) = \top & \text{Sk}(\perp) = \text{F} & \text{Sk}(a) = a \\ \text{Sk}(\top) = \top & \text{Sk}(0) = \text{F} & \text{Sk}(a^\perp) = \neg a \\ \text{Sk}(A \multimap B) = \text{Sk}(A) \Rightarrow \text{Sk}(B) & \text{Sk}(\sigma X.A) = \sigma X.\text{Sk}(A) & \text{Sk}(X) = X \end{array}$$

Let π be a two-sided μLL^∞ pre-proof of $\Gamma \vdash \Delta$. $\text{Sk}(\pi)$ is the μLK^∞ pre-proof of $\text{Sk}(\Gamma) \vdash \text{Sk}(\Delta)$ obtained by the following corecursive process by case analysis on the last rule r of π :

- if $r = (\text{Ax})$, then $\text{Sk}(\pi)$ is an axiom;
- if $r = (\text{Cut})$, then $\text{Sk}(\pi)$ ends with a cut of premises the skeletons of the premises of π ;
- if r is a fixed-point rules, then $\text{Sk}(\pi)$ ends with the same rule of μLK^∞ with, for premise, the skeleton of the premise of π ;
- if r is a multiplicative or additive rule, then $\text{Sk}(\pi)$ ends with the corresponding rule for the associated connective with, for premises, the skeletons of the premises of π ;
- if r is a structural rules for an exponential formula or an exchange rule, then $\text{Sk}(\pi)$ ends with the corresponding structural rules with, for premise, the skeleton of the premise of π ;
- if r is a promotion or dereliction, then $\text{Sk}(\pi)$ is the skeleton of the premise of π of $\text{Sk}(\pi)$.

D.2 Proof of Proposition 41

Proposition 81. *For any valid μLL^∞ proof π of s , $\text{Sk}(\pi)$ is a valid μLK^∞ proof of $\text{Sk}(s)$.*

Proof. It is easy to check that $\text{Sk}(\pi)$ is a μLK^∞ pre-proof (as soon as one admit both additive and multiplicative inference for classical conjunction and disjunction):

- even though the corecursive call, in the promotion and dereliction cases are non-guarded, the procedure is productive as there can only be finitely many such inferences
- $\text{Sk}(\pi)$ is trivially locally correct and of conclusion $\text{Sk}(s)$.

Every infinite branch of $\text{Sk}(\pi)$ is valid since all fixed-points unfoldings are preserved in the skeleton and therefore every infinite thread of a branch of π with least recurring formula F may be mapped to a thread of minimal formula $\text{Sk}(F)$

D.3 A linear decoration of μLK^∞ into μLL^∞

We consider the following embedding of μLK^∞ (and μLK^∞ sequents) into two-sided μLL^∞ :

$$\begin{aligned}
[a]^k &= !a \\
[X]^k &= !X \\
[A \Rightarrow B]^k &= !((?[A]^k) \multimap ([B]^k)) \\
[A \wedge B]^k &= !((?[A]^k) \& ([B]^k)) \\
[A \vee B]^k &= !((?[A]^k) \oplus ([B]^k)) \\
[\sigma X.A]^k &= !\sigma X.?[A]^k \quad \sigma \in \{\mu, \nu\} \\
[\top]^k &= !\top \\
[\text{F}]^k &= !0 \\
[\Gamma \vdash \Delta]^k &= [\Gamma]^k \vdash ?[\Delta]^k
\end{aligned}$$

One easily checks that the translation is consistent with positivity conditions of μLK and that:

Lemma 82. *For any μLK formula A , $[A]^k$ is a μLL formula.*

Moreover, setting $[X]^k = !X$ entails easily that:

Lemma 83. *For any μLK formula A, B , $\sigma \in \{\mu, \nu\}$,*

$$[A[\sigma X.B/X]]^k = [A]^k[\sigma X.?[B]^k/X].$$

Proof. This simply follows from the fact that $[X]^k = !X$, by induction on A .

Definition 84 (Translation of μLK^∞ proofs in μLL^∞ proofs). *In Figures 15 and 16, each μLK^∞ inference is translated in a μLL^∞ proof pattern relating the translation of the premisses sequents to the translation of the conclusion sequent.*

A non-wellfounded μLK^∞ pre-proof π of $\Gamma \vdash \Delta$ can therefore be translated into a non-wellfounded pre-proof $[\pi]^k$ of $[\Gamma \vdash \Delta]^k$ of μLL^∞ .

In the above definition and the corresponding figures, we do not make explicit the ancestor relations as it is unambiguous.

Remark 85. Notice that the translation preserve the ancestor relation in the following sense: for each inference r of μLK^∞ , if the i th formula of conclusion s is in ancestor relation with the j th formula of the k th premise s_k , then in $[r]^k$, the i th formula of $[s]^k$ is in relation, for the transitive closure of μLL^∞ ancestor relation, with the j th formula of the k th open sequent $[s_k]^k$.

Proposition 86. *Let be a non-wellfounded μLK^∞ pre-proof π of Γ . π is a valid μLK^∞ proof if, and only if, $[\pi]^k$ is a valid μLL^∞ proof.*

Proof. Indeed, the infinite branches of π are in one-to-one correspondence with the infinite branches of $[\pi]^k$, the thread are preserved thanks to the previous remark and the translation of formula ensures that a thread is valid in π if, and only if, it is valid in $[\pi]^k$.

$$\begin{aligned}
\left[\frac{}{\Gamma, F \vdash F, \Delta} \text{(Ax)} \right]^k &= \frac{\frac{}{[F]^k \vdash [F]^k} \text{(Ax)}}{[F]^k \vdash ?[F]^k} \text{(?d}_r)}{[\Gamma]^k, [F]^k \vdash ?[F]^k, ?[\Delta]^k} \text{(!w}_i)^*, \text{(?w}_r)^* \\
\left[\frac{\Gamma \vdash \Delta, F \quad \Gamma', F \vdash \Delta'}{\Gamma, \Gamma' \vdash \Delta, \Delta'} \text{(Cut)} \right]^k &= \frac{[\Gamma]^k \vdash ?[\Delta]^k, ?[F]^k \quad \frac{[\Gamma']^k, [F]^k \vdash ?[\Delta']^k}{[\Gamma']^k, ?[F]^k \vdash ?[\Delta']^k} \text{(?p}_i)}{[\Gamma, \Gamma']^k \vdash ?[\Delta, \Delta']^k} \text{(Cut)} \\
\left[\frac{\Gamma \vdash \Delta}{\Gamma \vdash F, \Delta} \text{(W}_r) \right]^k &= \frac{[\Gamma]^k \vdash ?[\Delta]^k}{[\Gamma]^k \vdash ?[F]^k, ?[\Delta]^k} \text{(?w}_r) \\
\left[\frac{\Gamma \vdash \Delta}{\Gamma, F \vdash \Delta} \text{(W}_i) \right]^k &= \frac{[\Gamma]^k \vdash ?[\Delta]^k}{[\Gamma]^k, [F]^k \vdash ?[\Delta]^k} \text{(!w}_i) \\
&\quad \text{(Since } [F]^k \text{ is a !-formula.)} \\
\left[\frac{\Gamma \vdash F, F, \Delta}{\Gamma \vdash F, \Delta} \text{(C}_r) \right]^k &= \frac{[\Gamma]^k \vdash ?[F, F, \Delta]^k}{[\Gamma]^k \vdash ?[F, \Delta]^k} \text{(?c}_r) \\
\left[\frac{\Gamma, F, F \vdash \Delta}{\Gamma, F \vdash \Delta} \text{(C}_i) \right]^k &= \frac{[\Gamma, F, F]^k \vdash ?[\Delta]^k}{[\Gamma, F]^k \vdash ?[\Delta]^k} \text{(!c}_i) \\
&\quad \text{(Since } [F]^k \text{ is a !-formula.)} \\
\left[\frac{\Gamma \vdash \Delta, G, F, \Delta'}{\Gamma \vdash \Delta, F, G, \Delta'} \text{(X}_r) \right]^k &= \frac{[\Gamma]^k \vdash ?[\Delta, G, F, \Delta']^k}{[\Gamma]^k \vdash ?[\Delta, F, G, \Delta']^k} \text{(X}_r) \\
\left[\frac{\Gamma, G, F, \Gamma' \vdash \Delta}{\Gamma, F, G, \Gamma' \vdash \Delta} \text{(X}_i) \right]^k &= \frac{[\Gamma, G, F, \Gamma']^k \vdash ?[\Delta]^k}{[\Gamma, F, G, \Gamma']^k \vdash ?[\Delta]^k} \text{(X}_i)
\end{aligned}$$

Fig. 15. Translation of μLK^∞ inferences into μLL^∞ : identity and structural rules

$$\begin{aligned}
\left[\frac{\Gamma \vdash F, \Delta \quad \Gamma \vdash G, \Delta}{\Gamma \vdash F \wedge G, \Delta} (\wedge_r) \right]^k &= \frac{\frac{[T]^k \vdash ?[F]^k, ?[\Delta]^k \quad [T]^k \vdash ?[G]^k, ?[\Delta]^k}{[T]^k \vdash !(?[F]^k \& ?[G]^k), ?[\Delta]^k} (\&r), (!p_r)}{[T]^k \vdash ?[F \wedge G]^k, ?[\Delta]^k} (?d_r)} \\
\left[\frac{\Gamma, F_i \vdash \Delta}{\Gamma, F_1 \wedge F_2 \vdash \Delta} (\wedge_i) \right]^k &= \frac{\frac{[T]^k, [F_i]^k \vdash ?[\Delta]^k}{[T]^k, ?[F_i]^k \vdash ?[\Delta]^k} (?p_i)}{[T]^k, ?[F_1]^k \& ?[F_2]^k \vdash ?[\Delta]^k} (\&i)}{[T]^k, [F_1 \wedge F_2]^k \vdash ?[\Delta]^k} (!d_i)} \\
\left[\frac{\Gamma \vdash F_i, \Delta}{\Gamma \vdash F_1 \vee F_2, \Delta} (\vee_i) \right]^k &= \frac{\frac{[T]^k \vdash ?[F_i]^k, ?[\Delta]^k}{[T]^k \vdash !(?[F_1]^k \oplus ?[F_2]^k), ?[\Delta]^k} (\oplus_i), (!p_r)}{[T]^k \vdash ?[F_1 \vee F_2]^k, ?[\Delta]^k} (?d_r)} \\
\left[\frac{\Gamma, F \vdash \Delta \quad \Gamma, G \vdash \Delta}{\Gamma, F \vee G \vdash \Delta} (\vee_i) \right]^k &= \frac{\frac{\frac{[T]^k, [F]^k \vdash ?[\Delta]^k}{[T]^k, ?[F]^k \vdash ?[\Delta]^k} (?p_i) \quad \frac{[T]^k, [G]^k \vdash ?[\Delta]^k}{[T]^k, ?[G]^k \vdash ?[\Delta]^k} (?p_i)}{[T]^k, ?[F]^k \oplus ?[G]^k \vdash ?[\Delta]^k} (\oplus_i)}{[T]^k, [F \vee G]^k \vdash ?[\Delta]^k} (!d_i)} \\
\left[\frac{}{\Gamma \vdash \top, \Delta} (\top_r) \right]^k &= \frac{\frac{}{[T]^k \vdash \top, ?[\Delta]^k} (\top_r)}{[T]^k \vdash !\top, ?[\Delta]^k} (!p_r)}{[T]^k \vdash ?[\top]^k, ?[\Delta]^k} (?d_r)} \\
\left[\frac{}{\Gamma, \mathbf{F} \vdash \Delta} (\mathbf{F}_i) \right]^k &= \frac{\frac{}{[T]^k, 0 \vdash ?[\Delta]^k} (0_i)}{[T]^k, [\mathbf{F}]^k \vdash ?[\Delta]^k} (!d_i)} \\
\left[\frac{\Gamma, F \vdash G, \Delta}{\Gamma \vdash F \Rightarrow G, \Delta} (\Rightarrow_r) \right]^k &= \frac{\frac{\frac{[T]^k, [F]^k \vdash ?[G]^k, ?[\Delta]^k}{[T]^k, ?[F]^k \vdash ?[G]^k, ?[\Delta]^k} (?p_i)}{[T]^k \vdash !(?[F]^k \multimap ?[G]^k), ?[\Delta]^k} (\multimap_r), (!p_r)}{[T]^k \vdash ?[F \Rightarrow G]^k, ?[\Delta]^k} (?d_r)} \\
\left[\frac{\Gamma \vdash F, \Delta \quad \Gamma', G \vdash \Delta'}{\Gamma, \Gamma', F \Rightarrow G \vdash \Delta, \Delta'} (\Rightarrow_i) \right]^k &= \frac{\frac{\frac{[T]^k \vdash ?[F]^k, ?[\Delta]^k \quad \frac{[T]'^k, [G]^k \vdash ?[\Delta']^k}{[T]'^k, ?[G]^k \vdash ?[\Delta']^k} (?p_i)}{[T], [T]'^k, ?[F]^k \multimap ?[G]^k \vdash ?[\Delta]^k, ?[\Delta']^k} (\multimap_i)}{[T], \Gamma', F \Rightarrow G \vdash ?[\Delta, \Delta']^k} (!d_i)} \\
\left[\frac{\Gamma \vdash G[\sigma X.G/X], \Delta}{\Gamma \vdash \sigma X.G, \Delta} (\sigma_r) \right]^k &= \frac{\frac{[T]^k \vdash ?[G]^k[\sigma X.?[G]^k/X], ?[\Delta]^k}{[T]^k \vdash \sigma X.?[G]^k, ?[\Delta]^k} \text{By Lemma 83}}{[T]^k \vdash !\sigma X.?[G]^k, ?[\Delta]^k} (!p_r)}{45 \quad \frac{[T]^k \vdash ?[\sigma X.G]^k, ?[\Delta]^k}{} (?d_r)} \\
\left[\frac{\Gamma, G[\sigma X.G/X] \vdash \Delta}{\Gamma, \sigma X.G \vdash \Delta} (\sigma_i) \right]^k &= \frac{\frac{[T]^k, [G[\sigma X.G/X]]^k \vdash ?[\Delta]^k}{[T]^k, ?([G]^k[\sigma X.?[G]^k/X]) \vdash ?[\Delta]^k} (?p_i)}{[T]^k, \sigma X.?[G]^k \vdash ?[\Delta]^k} (\sigma_i)}{[T]^k, [\sigma X.G]^k \vdash ?[\Delta]^k} (!d_i)} \text{By Lemma 83}
\end{aligned}$$

Fig. 16. Translation of μLK^∞ inferences into μLL^∞ : logical rules

D.4 Proof of Theorem 44

Theorem 87. μLK^∞ eliminates cuts.

Proof. Let $\pi \in \mu\text{LK}^\infty$ be a proof of $\Gamma \vdash \Delta$. π^d is a valid proof of $(\Gamma \vdash \Delta)^d$ so that by cut-elimination, there is a fair reduction sequence $\sigma = (II_i)_{i \in \omega}$ with $II_0 = \pi^d$ which converges to a cut-free μLL^∞ proof II' . For any $i \in \omega$, one has either $\text{Sk}(II_i) = \text{Sk}(II_{i+1})$ or $\text{Sk}(II_i) \longrightarrow_{\mu\text{LK}^\infty} \text{Sk}(II_{i+1})$ so that one can extract from σ a strongly converging $\longrightarrow_{\mu\text{LK}^\infty}$ -reduction sequence which converges to $\text{Sk}(II')$, a cut-free valid μLK^∞ proof of $\Gamma \vdash \Delta$.

E On Seely (iso)morphisms.

E.1 From multiplicatives to additives and back

In LL, there are two conjunctions and two disjunctions, the additives and multiplicatives, which are clearly distinguished and have no interderivability properties.

On the other hand, in LK, the additive and multiplicative presentations of conjunction and disjunction are interderivable thanks to structural rules. One has:

$$\frac{\frac{\overline{A \vdash A} \quad (\text{Ax})}{A, B \vdash A} \quad (\text{W}_l) \quad \frac{\overline{B \vdash B} \quad (\text{Ax})}{A, B \vdash B} \quad (\text{W}_l)}{A, B \vdash A \wedge^a B} \quad (\wedge^a_r)}{A \wedge^m B \vdash A \wedge^a B} \quad (\wedge^m_l)$$

$$\frac{\frac{\overline{A \vdash A} \quad (\text{Ax})}{A \wedge^a B \vdash A} \quad (\wedge^{a2}_l) \quad \frac{\overline{B \vdash B} \quad (\text{Ax})}{A \wedge^a B \vdash B} \quad (\wedge^{a1}_l)}{A \wedge^a B, A \wedge^a B \vdash A \wedge^m B} \quad (\wedge^m_r)}{A \wedge^a B \vdash A \wedge^m B} \quad (\text{C}_l)$$

A, B are weakened on the left, $A \wedge^a B$ is contracted on the left.

In LL, we do not have free structural rules, but only thanks to exponentials, so we need to decorate formulas with exponentials where structural rules are needed, leading to:

E.2 $!A \otimes !B \dashv\vdash !(A \& B)$

$$\pi_S = \frac{\frac{\frac{\overline{A \vdash A} \quad (\text{Ax})}{!A \vdash A} \quad (?d) \quad \frac{\overline{B \vdash B} \quad (\text{Ax})}{!B \vdash B} \quad (?d)}{!A, !B \vdash A} \quad (?w) \quad \frac{\overline{B \vdash B} \quad (\text{Ax})}{!B \vdash B} \quad (?d)}{!A, !B \vdash B} \quad (?w)}{!A, !B \vdash A \& B} \quad (\&)}{!A, !B \vdash !(A \& B)} \quad (!p)}{!A \otimes !B \vdash !(A \& B)} \quad (!\otimes)$$

$$\pi'_S = \frac{\frac{\frac{\overline{A \vdash A} \quad (\text{Ax})}{A \& B \vdash A} \quad (\oplus^1) \quad \frac{\overline{B \vdash B} \quad (\text{Ax})}{A \& B \vdash B} \quad (\oplus^2)}{!(A \& B) \vdash A} \quad (?d) \quad \frac{\overline{B \vdash B} \quad (\text{Ax})}{A \& B \vdash B} \quad (?d)}{!(A \& B) \vdash !A} \quad (!p) \quad \frac{\overline{B \vdash B} \quad (\text{Ax})}{A \& B \vdash B} \quad (?d)}{!(A \& B) \vdash !B} \quad (!p)}{!(A \& B), !(A \& B) \vdash !A \otimes !B} \quad (\otimes)}{!(A \& B) \vdash !A \otimes !B} \quad (?c)$$

