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► **To cite this version:**

Anni Hakanen, Ismael G. Yero. Complexity and equivalency of multiset dimension and ID-colorings. 2023. hal-04026222v2

HAL Id: hal-04026222

<https://hal.science/hal-04026222v2>

Preprint submitted on 21 Nov 2023

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Complexity and equivalency of multiset dimension and ID-colorings

Anni Hakanen ^{*† ‡} Ismael G. Yero ^{§¶}

November 20, 2023

Abstract

This investigation is firstly focused into showing that two metric parameters represent the same object in graph theory. That is, we prove that the multiset resolving sets and the ID-colorings of graphs are the same thing. We also consider some computational and combinatorial problems of the multiset dimension, or equivalently, the ID-number of graphs. We prove that the decision problem concerning finding the multiset dimension of graphs is NP-complete. We consider the multiset dimension of king grids and prove that it is bounded above by 4. We also give a characterization of the strong product graphs with one factor being a complete graph, and whose multiset dimension is not infinite.

Keywords: multiset resolving set; multiset dimension; ID-colorings; ID-number; king grids.

AMS Subj. Class. (2020): 05C12; 05C76.

1 Introduction

The concept of metric dimension in graphs is one of the classical parameters in the area of graph theory. It is understood it has been independently introduced in the decade of 1970 in the two separate works [9] and [23], which were aimed to consider identification properties of vertices in a graph. These identification properties were also connected with the Mastermind game in [4], and problems related to pattern recognition and image processing in [19]. After these two seminal works, the research on the topic remained relatively quiet until the first years of the new century, where the number of articles on the topic exploded. From this point on, several theoretical and applied results have been appearing, and nowadays, the metric dimension of graphs is very well studied. It is not our goal to include a lot of references on this fact, and we simply suggest the interested reader to consult the two recent surveys [17, 24], which have a fairly complete amount of information on metric dimension in graphs and related topics.

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Among the research lines addressed in the investigation with metric dimension in graphs, a remarkable one is that of considering variations of the classical concept that are giving more insight into the main concept, and are indeed of independent interest. The range of variations include generalizations of the concept, particularizations in the style of identification of vertices, union of the metric properties with other graph situations, identification of other elements (edges for instance) of graphs, etc. The survey [17] contains information on several of the most important variations already known.

As it happens, sometimes some variations could relate much between them, and sometimes there could be some that indeed represent the same structure. This is one of the contributions of our work. We describe the equivalence of two metric concepts that, at a first glance, appear to be separate from each other. To this end, we first present some basic terminology and notations that shall be used in our article.

Along our work, all graphs are simple, undirected and connected, unless we will specifically state the contrary. Let G be a graph with vertex set $V(G)$. Two vertices u, v are identified (recognized or determined) by a vertex w if $d_G(u, w) \neq d_G(v, w)$ where $d_G(x, y)$ stands for the distance (in its standard version) between x and y . A set $S \subset V(G)$ is a *resolving set* for G , if every two vertices $u, v \in V(G)$ are identified by a vertex of S . The cardinality of a smallest possible resolving set for G is the *metric dimension* of G , denoted $\dim(G)$. A resolving set of cardinality $\dim(G)$ is called a *metric basis*. These concepts are from [9] and [23], although in the latter work, they had different names. Some recent significant contributions on the metric dimension of graphs are for instance [6, 7, 18, 21, 25].

The metric representation of a vertex x with respect to an ordered set of vertices $S = \{v_1, \dots, v_k\}$ is the vector $r(x|S) = (d_G(x, v_1), \dots, d_G(x, v_k))$. It can be readily observed that a set S of vertices is a resolving set of a graph G if and only if the set of metric representations of vertices of G are pairwise different. This terminology turns out to be very useful while working with this concept.

Throughout our exposition, we are then focused on studying two (indeed one) metric dimension related parameters: the multiset dimension and the ID-number of graphs, which are variants of metric dimension that use multisets of distances instead of vectors, to uniquely identify the vertices of a graph. Our exposition is organized as follows. The next section is dedicated to prove that in fact the multiset dimension and the ID-number of graphs are the same parameter, and based on this, we just follow the terminology of multiset resolving sets and multiset dimension in our work. Section 3 is focused on showing that the decision problem concerning finding the multiset dimension of graphs is NP-complete. Next, in Section 4, we consider the multiset dimension of king grids, namely, the strong product of a path with itself. Section 5 contains a characterization of the strong product graphs with one factor being a complete graph, and whose multiset dimension is not infinite. Finally, we end our work with some questions and open problems that might be of interest for future research.

2 Two equivalent metric concepts

A modified version of the classical metric dimension of graphs was first presented in [22] where the use of “multisets” instead of vectors was initiated while considering metric representations of vertices with respect to a given set. That is, if $u \in V(G)$ and $W = \{w_1, \dots, w_t\}$, then the *multiset representation of u with respect to W* is given by

$$m_G(u|W) = \{\{d_G(u, w_1), \dots, d_G(u, w_t)\}\},$$

where $\{\cdot\}$ represents a multiset. In order to facilitate our exposition, we write $\{\cdot\} + i$ to denote the multiset obtained from $\{\cdot\}$ by adding i to every element of such multiset. The set W is a *multiset resolving set* for G if the collection of multisets $m_G(u|S)$ with $u \in V(G)$ are pairwise distinct. The *multiset dimension* of G , denoted $\dim_{ms}(G)$, represents the cardinality of a smallest possible multiset resolving set of G . Multiset resolving sets do not always exist in a given graph. For those graphs G which do not contain any multiset resolving set, the agreement that $\dim_{ms}(G) = \infty$ was taken in [22]. Some other investigations on the multiset dimension of graphs are [1, 3, 10–12]. It is clear that any multiset resolving set is also a resolving set, since the fact that two multisets are different implies that the vectors with the same elements are also different. This means that for any graph G ,

$$\dim(G) \leq \dim_{ms}(G). \quad (1)$$

On the other hand, the following concepts were defined in [5]. Consider a connected graph G of diameter d and a set of vertices $S \subset V(G)$. Now, for every vertex $x \in V(G)$, the *code* of x with respect to S is the d -vector $\vec{d}(x|S) = (a_1, a_2, \dots, a_d)$ where a_i , with $i \in \{1, \dots, d\}$ represents the number of vertices in S at distance i from x . If all the codes of vertices of G are pairwise different, then S is called an *identification coloring* or *ID-coloring*. Moreover, a graph G that has an ID-coloring is called an ID-graph. In this sense, for any ID-graph G , the cardinality of a smallest ID-coloring is the *ID-number* of G , denoted by $ID(G)$. Other contributions in this direction are [14–16].

We next show that the two parameters defined above are indeed the same. To this end, we may remark that for an ID-graph G of diameter d with an ID-coloring S , any vertex $x \in V(G)$ with code $\vec{d}(x|S) = (a_1, a_2, \dots, a_d)$ satisfies $\sum_{i=1}^d a_i = |S|$ if $x \notin S$ and $\sum_{i=1}^d a_i = |S| - 1$ if $x \in S$.

Theorem 1. *Let G be a graph of diameter d . Then $S \subset V(G)$ is an ID-coloring for G if and only if S is a multiset resolving set for G .*

Proof. The result follows directly from the following fact. Let $S \subset V(G)$ and let $x \in V(G)$. Consider the multiset representation $m_G(x|S) = \{d_G(x, w_1), \dots, d_G(x, w_t)\}$. From $m_G(x|S)$ we construct the vector $\vec{d}(x|S) = (a_1, a_2, \dots, a_d)$ as follows. Each a_i equals the number of elements in $m_G(x|S)$ with value i for any $i \in \{1, \dots, d\}$. On the contrary, if we have the vector $\vec{d}(x|S) = (a_1, a_2, \dots, a_d)$, then the multiset representation $m_G(x|S)$ can be obtained in the following way. First, if $\sum_{i=1}^d a_i = |S| - 1$, then $x \in S$ and we need to add the value zero (0) to $m_G(x|S)$. Otherwise, if $\sum_{i=1}^d a_i = |S|$, then $x \notin S$, and the value zero (0) does not belong to $m_G(x|S)$. Now, in both cases, for every $a_i \in \vec{d}(x|S)$ such that $a_i \neq 0$, we add to $m_G(x|S)$ a_i elements equal to i .

Based on the equivalence that exists between these two representations, the fact that the codes of vertices of G are pairwise different implies that the collection of multisets are pairwise distinct as well, and vice versa. Consequently, it is clear that a given set $S \subset V(G)$ is an ID-coloring of G if and only if S is a multiset resolving set. \square

Based on the equivalence above, we conclude the next consequence.

Corollary 2. *For any graph G , $\dim_{ms}(G) = ID(G)$.*

It is then now clear that graphs defined in [5] as ID-graphs are those ones satisfying that $\dim_{ms}(G) < \infty$ according to the terminology from [22].

3 Complexity results

This section is centered into considering the proof of the NP-completeness for the following decision problem, which in addition allows to conclude that computing the multiset dimension of graphs is NP-hard, and clearly, based on Theorem 1 and Corollary 2, it means that computing the ID-number of graphs is NP-hard as well. Our proof is somewhat inspired by the proof of the NP-hardness of the outer multiset dimension problem presented in [8].

Distinct vertices u and v are called *twins* if they have the same set of neighbors. It is well known that if two vertices are twins, then each (multiset) resolving set contains at least one of them.

MULTISET DIMENSION

Instance: A graph $G = (V, E)$ and an integer k satisfying $1 \leq k \leq |V|$.

Question: Is $\dim_{ms}(G) \leq k$?

Theorem 3. *The MULTISET DIMENSION problem is NP-complete.*

Proof. The problem is clearly in NP. We prove the NP-completeness by a reduction from 3-SAT. Consider an arbitrary input to 3-SAT, that is, a formula F with n variables and m clauses, which does not have a clause containing both the positive and the negative literals of the same variable. Let x_1, x_2, \dots, x_n be the variables, and let C_1, C_2, \dots, C_m be the clauses of F . We next construct a connected graph G based on this formula F . To this end, we use the following gadgets.

For each variable x_i we construct a *variable gadget* as follows (see Figure 1(a)).

- Vertices T_i, F_i are the “true” and “false” ends of the gadget. The gadget is attached to the rest of the graph only through these vertices.
- Vertices $a_i^1, a_i^2, b_i^1, b_i^2$ represent the value of the variable x_i , that is, a_i^1 and a_i^2 will be used to represent that variable x_i is true, and b_i^1 and b_i^2 that it is false. The vertices a_i^1 and b_i^1 are adjacent, and so are the vertices a_i^2 and b_i^2 . Additionally, the vertices a_i^1 and a_i^2 are adjacent to T_i and the vertices b_i^1 and b_i^2 are adjacent to F_i .
- $P_{t_i} = d_i^1 d_i^2 \dots d_i^{t_i}$ is a path such that d_i^1 is adjacent to T_i and F_i , while $d_i^{t_i}$ is adjacent to the two vertices e_i^1 and e_i^2 . Notice that the vertices e_i^1 and e_i^2 are twins, and so, each multiset resolving set of G must contain at least one of them.

For each clause C_j we construct a *clause gadget* as follows (see Figure 1(b)).

- The vertices c_j^1, c_j^2 and c_j^3 form a path. The vertices c_j^1 and c_j^3 will be helpful in determining whether the clause C_j is satisfied.
- $P_{s_j} = f_j^1 f_j^2 \dots f_j^{s_j}$ is a path such that f_j^1 is adjacent to c_j^2 , while $f_j^{s_j}$ is adjacent to the two vertices g_j^1 and g_j^2 . The vertices g_j^1 and g_j^2 are twins, and so, each multiset resolving set of G must contain at least one of them.

The orders of the paths P_{t_i} and P_{s_j} must be pairwise distinct in order for the multiset representations to be different and our reduction to work. Let $t_i = 5(i+1)$ for all $i \in \{1, \dots, n\}$ and $s_j = 5(n+j+1)$ for all $j \in \{1, \dots, m\}$. The sum of all the vertices in the variable and clause gadgets is clearly polynomial in terms of $n+m$.

The variable and clause gadgets are connected in the following way in order to construct our graph G .

- Vertices c_j^1 are adjacent to vertices T_i, F_i for all j and i .

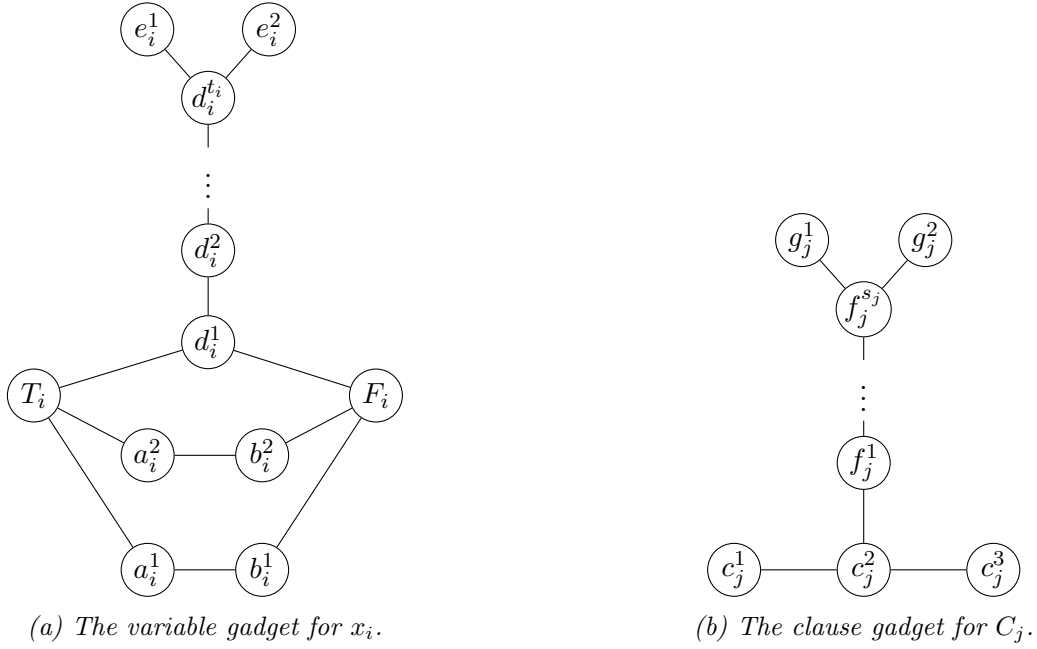


Figure 1: The variables and clause gadgets used in the reduction.

- If a variable x_i does not appear in a clause C_j , then the vertices T_i , F_i are adjacent to c_j^3 .
- If a variable x_i appears as a positive literal in a clause C_j , then the vertex F_i is adjacent to c_j^3 .
- If a variable x_i appears as a negative literal in a clause C_j , then the vertex T_i is adjacent to c_j^3 .

Observe that G is connected and its order is polynomial in terms of the quantity of variables and clauses of the 3-SAT instance. We shall show that F is satisfiable if and only if $\dim_{ms}(G) = 2m + n$. To this end, we proceed with a series of claims that will complete our whole reduction.

Claim 1. *We have $\dim_{ms}(G) \geq 2n + m$.*

Proof of Claim 1. Let S be a multiset basis of G . The vertices e_i^1 and e_i^2 are twins, and thus S contains at least one of them for every $i \in \{1, \dots, n\}$. Similarly, the vertices g_j^1 and g_j^2 are twins and at least one of them is in S for every $j \in \{1, \dots, m\}$. Finally, in order to have distinct multiset representation for the vertices a_i^1 , a_i^2 , b_i^1 and b_i^2 , at least one of these four vertices must be in S . Thus, we have $\dim_{ms}(G) = |S| \geq 2n + m$. \blacklozenge

Claim 2. *If F is satisfiable, then $\dim_{ms}(G) = 2n + m$.*

Proof of Claim 2. Consider a satisfying assignment for F and construct a set S^* containing $2n + m$ vertices as next described.

- For each $i \in \{1, \dots, n\}$, we add the vertex e_i^1 to S^* .
- For each $j \in \{1, \dots, m\}$, we add the vertex g_j^1 to S^* .
- For each variable x_i , if $x_i = \mathbf{true}$, then we add the vertex a_i^1 to S^* , otherwise if $x_i = \mathbf{false}$, then we add the vertex b_i^1 to S^* .

We will show that the set S^* is a multiset resolving set of G . We denote by $S_{x_i}^*$ and $S_{C_j}^*$ the vertices of the set S^* that are *not* in the gadget of x_i and C_j , respectively. We will first express the multiset representations of the vertices in the variable gadgets with the help of the multiset representation of T_i with respect to $S_{x_i}^*$. Since the vertices c_j^1 are adjacent to all T_i and F_i , the distance from T_i to all $T_{i'}$ and $F_{i'}$ for $i' \neq i$ is 2. Now, we have

$$m(T_i|S_{x_i}^*) = \left\{ \begin{array}{ll} 3, & | \text{ for each } i' \in \{1, \dots, n\}, i' \neq i \\ t_{i'} + 3, & | \text{ for each } i' \in \{1, \dots, n\}, i' \neq i \\ s_j + 3, & | \text{ for each } j \in \{1, \dots, m\} \end{array} \right\}.$$

For instance, the distance between the vertex T_i and $e_{i'}^1$ equals $t_{i'} + 3$, because one shortest path between them is $T_i c_1^1 T_{i'} d_{i'}^1 \cdots d_{i'}^{t_{i'}} e_{i'}^1$. Other cases are deduced similarly from the construction of the graph G .

Now, we can write the multiset representations of the vertices in the variable gadget of x_i as follows.

- $m(T_i|S^*) = m(T_i|S_{x_i}^*) \cup \{t_i + 1, y\}$, where $y = 1$ when $a_i^1 \in S^*$ and $y = 2$ when $b_i^1 \in S^*$.
- $m(F_i|S^*) = m(T_i|S_{x_i}^*) \cup \{t_i + 1, y\}$, where $y = 1$ when $b_i^1 \in S^*$ and $y = 2$ when $a_i^1 \in S^*$.
- $m(d_i^h|S^*) = (m(T_i|S_{x_i}^*) + h) \cup \{t_i - h + 1, h + 1\}$ for all $h \in \{1, \dots, t_i\}$.
- $m(e_i^h|S^*) = (m(T_i|S_{x_i}^*) + t_i + 1) \cup \{t_i + 2, y\}$, where $y = 0$ when $e_i^h = e_i^1$ and $y = 2$ when $e_i^h = e_i^2$.
- $m(a_i^h|S^*) = (m(T_i|S_{x_i}^*) + 1) \cup \{t_i + 2, y\}$, where when $a_i^1 \in S^*$, we have $y = 0$ for a_i^1 and $y = 2$ for a_i^2 , and when $b_i^1 \in S^*$, we have $y = 1$ for a_i^1 and $y = 3$ for a_i^2 .
- $m(b_i^h|S^*) = (m(T_i|S_{x_i}^*) + 1) \cup \{t_i + 2, y\}$, where when $a_i^1 \in S^*$, we have $y = 1$ for b_i^1 and $y = 3$ for b_i^2 , and when $b_i^1 \in S^*$, we have $y = 0$ for b_i^1 and $y = 2$ for b_i^2 .

As for the vertices of the clause gadgets, we will express them with an auxiliary representation as well. To that end, observe that

$$m(c_j^1|S_{C_j}^*) = \left\{ \begin{array}{ll} 2, & | \text{ for each } i \in \{1, \dots, n\} \\ t_i + 2, & | \text{ for each } i \in \{1, \dots, n\} \\ s_{j'} + 4, & | \text{ for each } j' \in \{1, \dots, m\}, j' \neq j \end{array} \right\}.$$

We then write the multiset representations of the vertices of the clause gadget of C_j other than c_j^3 as follows.

- $m(c_j^1|S^*) = m(c_j^1|S_{C_j}^*) \cup \{s_j + 2\}$.
- $m(c_j^2|S^*) = (m(c_j^1|S_{C_j}^*) + 1) \cup \{s_j + 1\}$.
- $m(f_j^h|S^*) = (m(c_j^1|S_{C_j}^*) + h + 1) \cup \{s_j - h + 1\}$ for all $h \in \{1, \dots, s_j\}$.
- $m(g_j^h|S^*) = (m(c_j^1|S_{C_j}^*) + s_j + 2) \cup \{y\}$, where $y = 0$ when $g_j^h = g_j^1$ and $y = 2$ when $g_j^h = g_j^2$.

Since the values of t_i and s_j are large and distinct enough (recall that $t_i = 5(i + 1)$ and $s_j = 5(n + j + 1)$), the multiset representations of the vertices in the variable and clause gadgets (other than c_j^3) are pairwise distinct. Indeed, notice that the values $t_{i'} + 3$ and $s_j + 3$ form a pattern to the multiset representations of a vertex of a variable gadget which is easy to distinguish from the corresponding pattern of $t_i + 2$ and $s_{j'} + 4$ of a vertex from a clause gadget. Thus, we readily observe that the multiset representations of the vertices in a variable gadget are distinct from those of the vertices of clause gadgets. These patterns, or the anomalies

present in them, to be more precise, are also the reason why vertices in two different variable gadgets (or clause gadgets) have distinct multiset representations. Indeed, there is a “gap” in this pattern where $t_i + 3$ should be for all vertices of the gadget of x_i . Thus, the multiset representations of the vertices within the same variable or clause gadget are distinct.

Let us then consider the vertices c_j^3 . Similarly to the vertex c_j^1 , the multiset representation of c_j^3 contains $s_j + 2$ once, $t_i + 2$ for each $i \in \{1, \dots, n\}$, and $s_{j'} + 2$ for each $j' \in \{1, \dots, m\}$, $j' \neq j$. However, the distance from c_j^3 to the vertices a_i^1 and b_i^1 is 2 or 3 depending on whether the variable x_i appears in the clause C_j and in which form (positive or negative), and whether $x_i = \mathbf{true}$ or $x_i = \mathbf{false}$ according to the truth assignment. More precisely, if the variable x_i does not appear in the clause C_j , then $d(c_j^3, a_i^1) = d(c_j^3, b_i^1) = 2$. If x_i appears in C_j but the clause C_j is not satisfied due to x_i (note that here C_j can be satisfied but due to the truth value of some other variable $x_{i'}$), then the distance from c_j^3 to whichever of a_i^1 and b_i^1 is in S^* is again 2. However, if C_j is satisfied due to x_i , the distance from c_j^3 to whichever of a_i^1 and b_i^1 is in S^* is 3. (For example, if x_i appears as a positive literal in C_j , the edge $c_j^3 F_i$ is present in G whereas $c_j^3 T_i$ is not. Thus, we have $d(c_j^3, a_i^1) = 3$ and $d(c_j^3, b_i^1) = 2$. Now, if the truth assignment of x_i leads to C_j being satisfied, we have $x_i = \mathbf{true}$ and $a_i^1 \in S^*$. Therefore, the multiset representation of c_j^3 contains 3 due to a_i^1 .)

Since the set S^* is constructed using a truth assignment that satisfies F , there is at least one 3 in the multiset representation of c_j^3 . Thus, the multiset representation of c_j^3 is almost the same as the multiset representation of c_j^1 except that at least one 2 (in $m(c_j^1|S^*)$) is swapped to 3. Thus, c_j^3 and c_j^1 have distinct multiset representations. Furthermore, based on the arguments concerning the multiset representations of the other vertices of G , it is clear that each c_j^3 has a distinct multiset representation compared to all other vertices of G .

Consequently, the set S^* is a multiset resolving set of G , and the claim holds due to Claim 1. \blacklozenge

Claim 3. *If $\dim_{ms}(G) = 2n + m$, then F is satisfiable.*

Proof of Claim 3. Let S be a multiset basis of G . By the arguments in the proof of Claim 1, the set S must contain exactly one of the two vertices e_i^1 or e_i^2 for every $i \in \{1, \dots, n\}$; exactly one of the two vertices g_j^1 or g_j^2 for every $j \in \{1, \dots, m\}$; and exactly one of the vertices a_i^1 , a_i^2 , b_i^1 or b_i^2 for every $i \in \{1, \dots, n\}$. By the same arguments as at the end of the proof of Claim 2, the vertices c_j^1 and c_j^3 must have distinct multiset representations due to some a_i^h or b_i^h . Thus, the truth assignment where $x_i = \mathbf{true}$ if a_i^1 or a_i^2 is in S , and $x_i = \mathbf{false}$ if b_i^1 or b_i^2 is in S for all $i \in \{1, \dots, n\}$ satisfies F . \blacklozenge

This completes the reduction from 3-SAT to the MULTISSET DIMENSION problem. \square

4 The king grid

Based on the NP-completeness reduction made in the proof of Theorem 3, it is then desirable to consider the multiset dimension (or ID-number) of some non-trivial families of graphs. In connection with this, in this section we consider the strong product of a path P_n with itself, also known as the king grid.

The graph $G \boxtimes H$ is the *strong product* of G and H . The vertex set of $G \boxtimes H$ is the set $V(G) \times V(H) = \{(u, v) \mid u \in V(G), v \in V(H)\}$. Two vertices $(g, h), (g', h') \in V(G \boxtimes H)$ are adjacent if $g = g'$ and h is adjacent to h' in H ; or g is adjacent to g' in G and $h = h'$; or g

is adjacent to g' in G and h is adjacent to h' in H . We write $V(P_n) = \{1, \dots, n\}$ so that the vertices of $P_n \boxtimes P_n$ correspond to the coordinates of the \mathbb{Z}^2 lattice. Moreover, for a vertex $(i, j) \in V(P_n \boxtimes P_n)$ and an integer $q \geq 1$, by $D_q(i, j)$ we mean the set of vertices in $V(P_n \boxtimes P_n)$ at distance q from (i, j) . Notice that such set $D_q(i, j)$ represents a kind of (not necessarily whole) “border” of a subgraph of $P_n \boxtimes P_n$ isomorphic to the strong product of two paths. See Figure 2 for two representative examples.

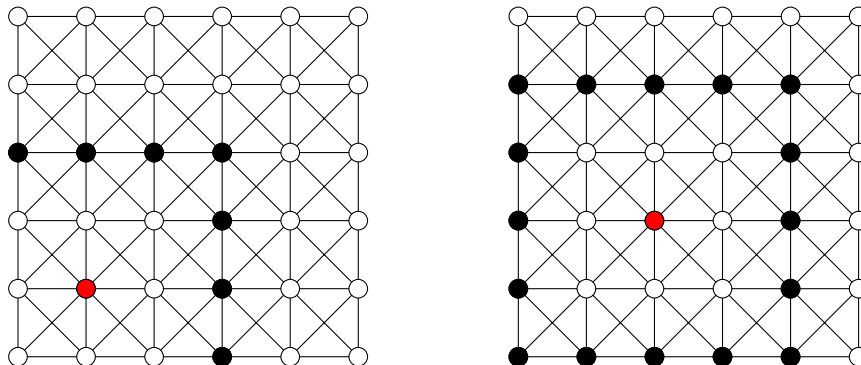


Figure 2: The graph $P_6 \boxtimes P_6$ with the sets $D_2(2, 2)$ and $D_2(3, 3)$, respectively illustrated in black. The vertices $(2, 2)$ and $(3, 3)$ appear in red color.

Remark 4. We have $\dim_{ms}(P_2 \boxtimes P_2) = \dim_{ms}(P_3 \boxtimes P_3) = \infty$, since both of these graphs are of diameter at most 2, and such graphs have no multiset resolving sets [22]. As for the case $n = 4$, the set $\{(1, 1), (2, 1), (4, 1), (1, 3), (2, 3), (2, 4)\}$ is a multiset resolving set of $P_4 \boxtimes P_4$. We have checked with an exhaustive computer search that no smaller multiset resolving sets exist for this graph, and thus $\dim_{ms}(P_4 \boxtimes P_4) = 6$

We begin our exposition of the larger king grids with the two smallest cases, and further on proceed with the general case.

Proposition 5. *We have $\dim_{ms}(P_5 \boxtimes P_5) = \dim_{ms}(P_6 \boxtimes P_6) = 4$.*

Proof. It is known from [22] and [5] that no graph has a multiset resolving set of cardinality 2. Also, the only graph that has a multiset resolving set consisting of only one element is the path graph P_n . Thus, $\dim_{ms}(P_5 \boxtimes P_5) \geq 3$ and $\dim_{ms}(P_6 \boxtimes P_6) \geq 3$. We first prove that $\dim_{ms}(P_5 \boxtimes P_5) \neq 3$ by a simple counting argument. Suppose that S is a multiset resolving set of $P_5 \boxtimes P_5$ such that $|S| = 3$. The maximum number of distinct multiset representations that do not contain 0 is $\binom{6}{3} = 20$. Since $|V(G) \setminus S| = 22$, some vertices of $P_5 \boxtimes P_5$ have the same multiset representations, a contradiction. Thus, $\dim_{ms}(P_5 \boxtimes P_5) > 3$, and the first equality follows since the set $S = \{(1, 1), (2, 1), (5, 1), (1, 5)\}$ is a multiset resolving set of $P_5 \boxtimes P_5$. The sets S along with the multiset representations is illustrated in Figure 3.

On the other hand, observe that the set $S' = \{(2, 1), (2, 2), (6, 1), (1, 6)\}$ is a multiset resolving set of $P_6 \boxtimes P_6$, as shown in Figure 3, throughout the multiset representations of each vertex with respect to S' , being pairwise different. Thus $\dim_{ms}(P_6 \boxtimes P_6) \leq 4$.

Now, in contrast to the case of $P_5 \boxtimes P_5$, to prove that $\dim_{ms}(P_6 \boxtimes P_6) \neq 3$, the counting argument used does not directly work. That is, if we suppose S'' is a multiset resolving set of $P_6 \boxtimes P_6$ such that $|S''| = 3$, then the maximum number of distinct multiset representations that do not contain 0 is $\binom{7}{3} = 35$, and $|V(P_6 \boxtimes P_6) \setminus S''| = 33$. Thus, some extra arguments are

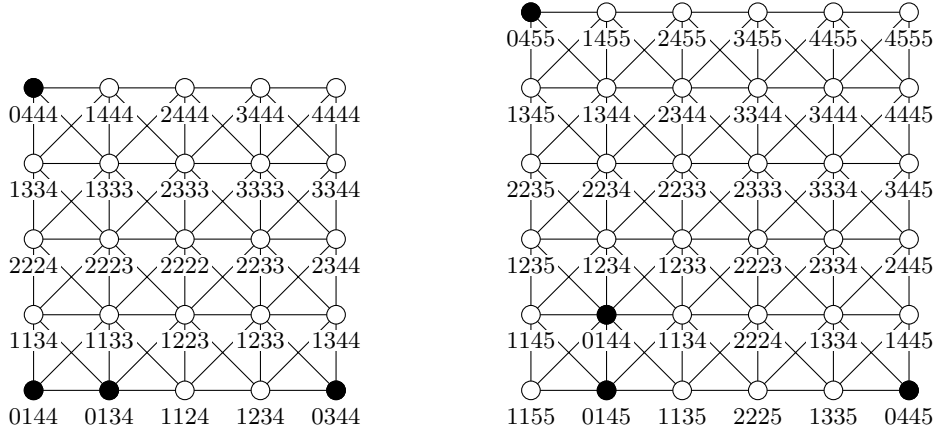


Figure 3: The graphs $P_5 \boxtimes P_5$ and $P_6 \boxtimes P_6$ with the sets S and S' , respectively illustrated in black. The four digits below each vertex are the distances in the multiset representation sorted in ascending order.

required. To this end, assume there is a vertex $(\alpha, \beta) \in V(P_6 \boxtimes P_6)$ such that it has multiset representation $\{\{1, 1, a\}\}$ for some $a \in \{1, \dots, 5\}$. Let $(i, j), (i', j')$ be two neighbors of (α, β) in S'' (note that $(i, j), (i', j')$ are at distance at most two). Hence, the third vertex (i'', j'') of S'' must be in the set $D_a(\alpha, \beta)$, namely, $S'' = \{(i, j), (i', j'), (i'', j'')\}$.

It is now just a matter of checking all the possibilities that can occur between the two vertices $(i, j), (i', j')$ and the third vertex (i'', j'') , to observe that one can always find two vertices that have the same multiset representation with respect to S'' . In order to avoid a lengthy and time consuming case analysis, we have simply checked this by computer. Thus, the multisets representations $\{\{1, 1, a\}\}$ are not possible for every $a \in \{1, \dots, 5\}$ with respect to the set S'' , and there are 5 of them. But then this means we have a total amount of $\binom{7}{3} - 5 = 30$ possible distinct multiset representations that do not contain 0. However, there are $|V(P_6 \boxtimes P_6) \setminus S''| = 33$ vertices, which is a contradiction. Therefore, $\dim_{ms}(P_6 \boxtimes P_6) \neq 3$, and the equality $\dim_{ms}(P_6 \boxtimes P_6) = 4$ follows. \square

From now on, in order to facilitate the exposition, by a *row* or a *column* in $P_n \boxtimes P_n$ we mean the path induced by the vertices $(1, j), (2, j) \dots, (n, j)$ or $(i, 1), (i, 2) \dots, (i, n)$, respectively, for any $i, j \in \{1, \dots, n\}$.

Theorem 6. *If $n \geq 7$ is an integer, then $3 \leq \dim_{ms}(P_n \boxtimes P_n) \leq 4$.*

Proof. The lower bound follows from the fact that any graph different from a path has multiset dimension at least 3, or by using inequality (1), since $\dim(P_n \boxtimes P_n) = 3$ (see [2, 20]). To show the upper bound we use an induction procedure that separately works for odd and even values of n .

Case 1: $n \geq 7$ is odd. Assume that the set $S_{n-2} = \{(1, 1), (2, 1), (n-2, 1), (1, n-2)\}$ is a multiset resolving set of $P_{n-2} \boxtimes P_{n-2}$. This is true for $n = 7$ according to Proposition 5, which shows the base case of the induction process. We will show that the set $S_n = \{(1, 1), (2, 1), (n, 1), (1, n)\}$ is a multiset resolving set of $P_n \boxtimes P_n$.

Denote $F = \{(x, y) \in V(P_n \boxtimes P_n) \mid x \in \{1, n\} \text{ or } y \in \{1, n\}\}$ and $I = V(P_n \boxtimes P_n) \setminus F$ (see Figure 4(a)). Observe that each vertex in F has $n-1$ in its multiset representation with

respect to S_n , whereas the vertices in I do not. Thus, vertices from F and I clearly have multiset representations distinct from one another.

Consider the vertices in I . The graph $P_n \boxtimes P_n$ can be viewed as a graph constructed from $P_{n-2} \boxtimes P_{n-2}$ by adding an additional row or column of vertices to all four sides of the graph. The set S_n can then be obtained by moving the elements of S_{n-2} diagonally away from the middle. Thus, for each $(v_1, v_2) \in I \setminus \{(2, 2), (3, 2)\}$, we have

$$m((v_1, v_2)|S_n) = m((v_1 - 1, v_2 - 1)|S_{n-2}) + 1.$$

Since S_{n-2} is a multiset resolving set of $P_{n-2} \boxtimes P_{n-2}$, all vertices in $I \setminus \{(2, 2), (3, 2)\}$ have distinct multiset representations with respect to S_n in $P_n \boxtimes P_n$. Since $m((2, 2)|S_n) = \{\{1, 1, n-2, n-2\}\}$ and $m((3, 2)|S_n) = \{\{1, 2, n-3, n-2\}\}$, the vertices $(2, 2)$ and $(3, 2)$ have distinct multiset representations with respect to each other. Moreover, the vertex $(2, 2)$ is the only vertex in I that is adjacent to two elements of the set S_2 . The vertex $(3, 2)$ is adjacent to one element of the set S_n , and so are two other vertices in I . However, neither of these vertices has the distance 2 in their multiset representations, which the vertex $(3, 2)$ does have. Thus, all vertices in I have pairwise distinct multiset representations with respect to S_n .

Consider then the vertices in F . The multiset representations of these vertices are as follows.

- For vertices on the top row, i.e. (i, n) where $i \in \{1, \dots, n\}$, we have $m((i, n)|S_n) = \{\{i-1, n-1, n-1, n-1\}\}$.
- For vertices on the right column, excluding the top corner, i.e. (n, i) where $i \in \{1, \dots, n-1\}$, we have $m((n, i)|S_n) = \{\{i-1, n-2, n-1, n-1\}\}$.
- For vertices on the left column, excluding top and bottom corners, i.e. $(1, i)$ where $i \in \{2, \dots, n-1\}$, we have $m((1, i)|S_n) = \{\{i-1, i-1, n-i-2, n-1\}\}$.
- For vertices on the bottom row, excluding both corners, i.e. $(i, 1)$, where $i \in \{2, \dots, n-1\}$, we have $m((i, 1)|S_n) = \{\{i-1, i-2, n-i-2, n-1\}\}$.
- For the bottom left corner $(1, 1)$, we have $m((1, 1)|S_n) = \{\{0, 1, n-1, n-1\}\}$.

Since the vertices in the top row are the only ones to have (at least) three $(n-1)$'s in their multiset representations, the vertices on the top row have distinct representations from the other vertices of F . Moreover, the multiset representations of the vertices on the top row are pairwise distinct. Similarly, from the vertices left to be considered, the vertices in the right column are the only vertices that have exactly two $(n-1)$'s in their representations (other than $(1, 1)$, but that has a distinct multiset representation due to the 0). Thus it is clear that the vertices in the right column have distinct multiset representations with respect to each other and other vertices in F . The vertices in the left column have pairwise distinct representations with respect to each other, and the same holds also for vertices in the bottom row. The only thing left to show is that two vertices, one in the left column and the other in the bottom row, cannot have the same multiset representations.

To that end, suppose to the contrary that $(i, 1)$ and $(1, j)$ have the same multiset representation for some $i, j \in \{2, \dots, n-1\}$. Since $j-1$ appears twice in the multiset representation of $(1, j)$, we must have $i-1 = n-i-2$ or $i-2 = n-i-2$. Since n is odd, we have $i-1 = n-i-2$ and $i-1 = \frac{n-3}{2}$. Now, $m((i, 1)|S_n) = \{\{\frac{n-3}{2}, \frac{n-3}{2}, \frac{n-3}{2} - 1, n-1\}\}$. This implies that $j-1 = \frac{n-3}{2}$, but now $m((1, j)|S_n) = \{\{\frac{n-3}{2}, \frac{n-3}{2}, \frac{n-3}{2}, n-1\}\} \neq m((i, 1)|S_n)$, a contradiction. Thus, all vertices of F have pairwise distinct multiset representations with respect to S_n .

As a consequence, we obtain that S_n is a multiset resolving set as claimed, and so, $\dim_{ms}(P_n \boxtimes P_n) \leq 4$ in this case.

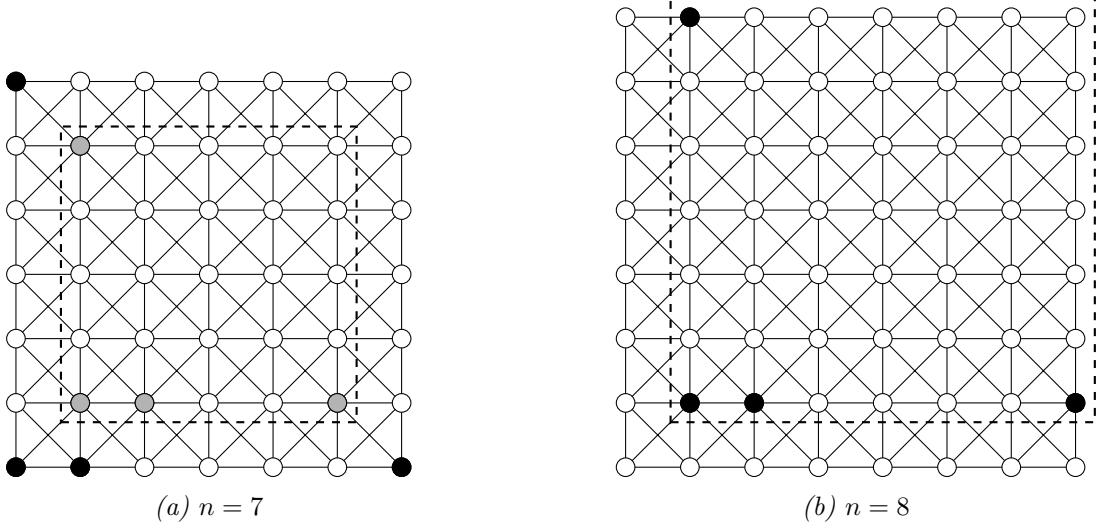


Figure 4: Vertices within the dashed line are the vertices in I , and the vertices outside the dashed line are the vertices in F . The black vertices are the elements of the set S_n .

Case 2: $n \geq 8$ is even. Assume that the set $S_{n-1} = \{(1, 1), (2, 1), (n-1, 1), (1, n-1)\}$ is a multiset resolving set of $P_{n-1} \boxtimes P_{n-1}$. If $n = 8$, then by the Case 1 we know that S_7 is a multiset resolving set of $P_7 \boxtimes P_7$, which shows the base case. We will show that the set $S_n = \{(2, 2), (2, 3), (n, 2), (2, n)\}$ is a multiset resolving set of $P_n \boxtimes P_n$.

We now denote $F = \{(x, y) \in V(P_n \boxtimes P_n) \mid x = 1 \text{ or } y = 1\}$ and $I = V(P_n \boxtimes P_n) \setminus F$ (see Figure 4(b)). The vertices in I clearly have pairwise distinct multiset representations as $m(v|S_n) = m(v|S_{n-1})$ for all $v \in I$, and S_{n-1} is a multiset resolving set of $P_{n-1} \boxtimes P_{n-1}$. Also, all the vertices in F have $n-1$ in their multiset representations with respect to S_n , whereas vertices in I do not. Thus, the multiset representation of a vertex in F is always distinct from that of a vertex of I .

We will show next that the multiset representations of vertices of F are pairwise distinct. To that end, we first consider the vertices that are adjacent to some element of S_n . Their multiset representations are the following:

$$\begin{aligned}
 m((1, 1)|S_n) &= \{\{1, 2, n-1, n-1\}\}, & m((1, n)|S_n) &= \{\{1, n-2, n-2, n-1\}\}, \\
 m((1, 2)|S_n) &= \{\{1, 2, n-2, n-1\}\}, & m((1, n-1)|S_n) &= \{\{1, n-3, n-3, n-1\}\}, \\
 m((1, 3)|S_n) &= \{\{1, 2, n-3, n-1\}\}, & m((n, 1)|S_n) &= \{\{1, n-3, n-2, n-1\}\}, \\
 m((4, 1)|S_n) &= \{\{1, 2, n-4, n-1\}\}, & m((n-1, 1)|S_n) &= \{\{1, n-4, n-3, n-1\}\}, \\
 m((2, 1)|S_n) &= \{\{1, 1, n-2, n-1\}\}, & m((3, 1)|S_n) &= \{\{1, 1, n-3, n-1\}\}.
 \end{aligned}$$

Since $n \geq 8$, we have $n-4 \neq 2$ and all these multiset representations are pairwise distinct. Moreover, since these vertices are the only vertices in F that have the distance 1 in their multiset representations, these multiset representations are distinct from the multiset representation of other vertices in F as well. Let us then consider the rest of the

vertices in F . The vertices $(i, 1)$ where $i \in \{5, \dots, n-2\}$ have multiset representations of the form $\{\{i-2, i-3, n-i-2, n-1\}\}$, and these representations are pairwise distinct. Similarly, the vertices $(1, j)$ where $j \in \{4, \dots, n-2\}$ have multiset representations of the form $\{\{j-2, j-2, n-j-2, n-1\}\}$, and these representations are clearly pairwise distinct. Suppose then that $(i, 1)$ and $(1, j)$ have the same multiset representation for some $i \in \{5, \dots, n-2\}$ and $j \in \{4, \dots, n-2\}$. As in the proof for odd n , the distance $j-2$ appears twice in the multiset representation of $(1, j)$. This implies that $i-2 = n-i-2$ or $i-3 = n-i-2$. As n is even, we have $i-2 = n-i-2$, and thus $i-2 = \frac{n-4}{2}$. Now, $m((i, 1)|S_n) = \{\{\frac{n-4}{2}, \frac{n-4}{2}, \frac{n-4}{2}-1, n-1\}\}$. This implies that $j-2 = \frac{n-4}{2}$. However, now $m((1, j)|S_n) = \{\{\frac{n-4}{2}, \frac{n-4}{2}, \frac{n-4}{2}, n-1\}\} \neq m((i, 1)|S_n)$, a contradiction. Thus, all vertices in F have pairwise distinct multiset representations with respect to S_n .

Therefore, we again conclude that S_n is a multiset resolving set in this situation, and so, $\dim_{ms}(P_n \boxtimes P_n) \leq 4$ follows as well, which completes the proof. \square

5 Strong products involving a complete graph

Several graphs with infinite multiset dimension (or equivalently that are not ID-graphs) are already known from the seminal works [5, 22], in their corresponding terminologies. For instance, it is known from these mentioned works that for a given graph G of diameter two, $\dim_{ms}(G \boxtimes K_n) < \infty$ if and only if G is P_3 . From this result we can identify a lot of interesting and non-trivial families of graphs like for instance the join of two non-complete graphs, the Cartesian or direct product of two complete graphs, and the Kneser graph $K(n, 2)$, among others, having infinite multiset dimension.

We next describe some other graphs with diameter larger than two that have infinite multiset dimension. Specifically, we consider the case of strong product graphs $G \boxtimes H$ when H is a complete graph. We first need the following definition from [13]. A graph G is called a *multiset distance irregular graph* if for any two vertices $x, y \in V(G)$ it follows that $m(x|V(G)) \neq m(y|V(G))$. An example of a multiset distance irregular graph is for instance a tree obtained from a star S_n with leaves v_0, \dots, v_{n-1} ($n \geq 3$) and center x , by subdividing i times the edge xv_i for every $i \in \{0, \dots, n-1\}$.

Theorem 7. *Let G be a graph and let $n \geq 2$ be an integer. We have $\dim_{ms}(G \boxtimes K_n) < \infty$ if and only if $n = 2$ and G is a multiset distance irregular graph.*

Proof. Notice that the n vertices of each copy of K_n are twins, that is, they have the same closed neighborhood. Hence, if S is a multiset resolving set of $G \boxtimes K_n$, then at least $n-1$ vertices from each copy of K_n must be in S . If $n \geq 3$, then from each copy of K_n there are at least two vertices in S . But then, these two vertices have the same multiset representation, which is not possible. Consequently, we deduce that $G \boxtimes K_n$ does not have multiset resolving sets when $n \geq 3$ and so, $\dim_{ms}(G \boxtimes K_n) = \infty$ in this case.

Assume next that $n = 2$. By the same reasons as above, at least one vertex from each copy of K_2 must be in S . Thus $\dim_{ms}(G \boxtimes K_2) \geq |V(G)|$. Now, assume G is multiset distance irregular, and consider a set X of vertices containing one whole copy of G . Since the multisets representations of any two vertices $x, y \in V(G)$ satisfy that $m(x|V(G)) \neq m(y|V(G))$, we deduce that X is a multiset resolving set, and so $\dim_{ms}(G \boxtimes K_n) \leq |V(G)|$, which gives the equality $\dim_{ms}(G \boxtimes K_n) = |V(G)| < \infty$.

Assume now that $\dim_{ms}(G \boxtimes K_n) < \infty$. Clearly, $n = 2$, for otherwise we get a contradiction. Also, we readily see that every multiset resolving set has nonempty intersection with

every copy of K_2 in $G \boxtimes K_2$ and that $\dim_{ms}(G \boxtimes K_2) \geq |V(G)|$. Let X' be a multiset basis of $G \boxtimes K_2$. If X' contains both vertices of one copy of K_2 , then these two vertices have the same multiset representation as they are twins. Thus, X' contains exactly one vertex from each copy of K_2 . If G is not multiset distance irregular, then there are two vertices $x, y \in V(G)$ such that $m(x|V(G)) = m(y|V(G))$. Now, the vertices in the copies of K_2 corresponding to x and y that are not in X' have the same multiset representation, a contradiction. Therefore, G is multiset distance irregular, and the proof is completed. \square

Corollary 8. *For any multiset distance irregular graph G , $\dim_{ms}(G \boxtimes K_2) = |V(G)|$.*

6 Concluding remarks

This work firstly shows that two metric parameters represent the same in graph theory, that is, multiset dimension and ID-colorings are the same. We have also considered some computational and combinatorial problems on this parameter. As a consequence of the study a number of possible future research lines have been detected. We next remark a few that could be of interest from our humble opinion.

- We have proved that finding the multiset dimension of graphs is in general NP-hard. However, not much is known on special classes of graphs. Does finding such parameters remain NP-hard even when restricted to trees, chordal graphs or planar graphs?
- The multiset dimension of the king grid has been bounded above by 4. Is it true that $\dim_{ms}(P_n \boxtimes P_n) = 4$ for any $n \geq 7$? In addition, this result gives a first step into considering the multiset dimension of the strong product graphs in general. Moreover, the study of such parameter in some other related graph products is worthwhile as well.
- We have related the multiset dimension of graphs with multiset distance irregular graphs. Does such graphs play a significant role in some investigations on the multiset dimension of general graphs or at least in some product related structures?
- The notion of identification spectrum of an ID-graph G was introduced in [5], which can be understood as the set of positive integers r for which there is a multiset resolving set of cardinality r . In this sense, since we have proved that $\dim_{ms}(P_n \boxtimes P_n) \leq 4$ (equivalently $P_n \boxtimes P_n$ is an ID-graph) for any $n \geq 7$, we wonder which is the identification spectrum of $P_n \boxtimes P_n$.
- Characterize all the graphs G for which $\dim(G) = \dim_{ms}(G)$.

Acknowledgements

Anni Hakanen has been supported by the Jenny and Antti Wihuri Foundation, and also partially supported by the ANR project GRALMECO (ANR-21-CE48-0004-01) and Academy of Finland grant number 338797. Ismael G. Yero has been partially supported by the Spanish Ministry of Science and Innovation through the grant PID2019-105824GB-I00.

Declaration of interests

The authors declare that they have no known competing financial interests or personal relationships that could have appeared to influence the work reported in this paper.

Data availability

Our manuscript has no associated data.

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