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SUBSTITUTIVE SYSTEMS AND A FINITARY VERSION OF COBHAM'S THEOREM

JAKUB BYSZEWSKI, JAKUB KONIECZNY, AND ELŻBIETA KRAWCZYK

Abstract. We study substitutive systems generated by nonprimitive substitutions and show that transitive subsystems of substitutive systems are substitutive. As an application we obtain a complete characterisation of the sets of words that can appear as common factors of two automatic sequences defined over multiplicatively independent bases. This generalises the famous theorem of Cobham.

Introduction

Let \mathcal{A} be a finite alphabet, let \mathcal{A}^* be the set of finite words over \mathcal{A} and let \mathcal{A}^ω be the set of sequences $(a_n)_{n \geq 0}$ with values in \mathcal{A} . We say that a dynamical system $X \subseteq \mathcal{A}^\omega$ is *substitutive* if it arises as the orbit closure of a substitutive sequence. Such systems were extensively studied in the context of primitive substitutions [28, 29, 33], necessarily restricting such studies to minimal systems. There is also a close relationship between substitutive systems and D0L-systems, studied e.g. in [26]. In the recent years there has been growing interest in the study of nonminimal substitutive systems, e.g. with connection to Bratteli diagrams [7] and tiling spaces [31]. Nevertheless, it seems that treatments of substitutive systems arising from nonprimitive substitutions are still scarce. In particular, the following basic question seems not to have been studied: Is every transitive subsystem of a substitutive system substitutive? In other words, if X is a substitutive system and x is a sequence in X , is there a substitutive sequence y such that x and y have the same set of factors? The same question can be posed for k -automatic systems (for the precise definitions of substitutive and k -automatic systems see Section 1). Note that substitutive systems may contain uncountably many points, while the number of substitutive sequences is countable, and so most sequences in a substitutive system will often not be substitutive.

The aim of this paper is twofold. First, we study general substitutive systems and provide a positive answer to the above question. Second, we apply this result to obtain a finitary version of the classical theorem of Cobham, answering a question posed by Shallit¹ (see also the discussion in [32]).

We focus our study on noninvertible substitutive systems, but we briefly present analogous results for invertible systems as well. Noninvertible substitutive systems have a considerably more complicated and interesting dynamical structure than the invertible ones. For instance, it follows from [31] that the number of subsystems of an invertible substitutive system is finite (see Remark 2.14), while noninvertible substitutive systems can have infinitely many subsystems (see e.g. Example 2.5).

Throughout the article, we only consider substitutions that are growing (i.e. substitutions $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ such that the length of the words $\varphi^n(a)$ tends to infinity for all letters $a \in \mathcal{A}$). We do not know what happens when this assumption is removed, though we would not be surprised if it could be shown that Theorem A below continues to hold.

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¹Jeffrey Shallit, journal entry, 2 February 2018 (private communication)

Our first main result yields the following description of transitive subsystems of substitutive systems.

Theorem A. *Every transitive subsystem of a substitutive system is substitutive. Every transitive subsystem of a k -automatic system is k -automatic.*

In fact, we obtain a much more precise description of substitutive (resp., k -automatic) sequences generating such subsystems. We present here a simplified version of the result in the noninvertible case (for more details see Proposition 2.6 and the proof of Proposition 2.8).

Theorem B. *Let x be a substitutive sequence produced by a substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$, and let X be the orbit closure of x . There exists a power $\tau = \varphi^m$ of φ and a finite set of words $W \subset \mathcal{A}^*$ such that every transitive subsystem $Y \subset X$ can be generated by a sequence $y \in Y$ that is a suffix of a biinfinite sequence of the form*

$$\dots \tau^2(v)\tau(v)vabw\tau(w)\tau^2(w)\dots \quad (1)$$

for some $v \in W$, $w \in W \setminus \{\epsilon\}$, and $a, b \in \mathcal{A} \cup \{\epsilon\}$.

Substitutive sequences of such a form have been considered before in specific contexts. Let x be a substitutive sequence over an alphabet \mathcal{A} . A sequence z in the orbit closure of x is called extremal if it is lexicographically minimal with respect to some total order on the alphabet \mathcal{A} . In [16] it was shown that (under some additional assumptions) all extremal sequences are substitutive. We note here a curious observation that all extremal sequences in (purely) substitutive systems considered in [16] are of the form (1) ([16, Lemma 9]).

In the second part of the article we restrict our attention to automatic sequences. One of the most fundamental results about automatic sequences is Cobham's theorem, which gives a strong relation between k -automaticity of a sequence and the chosen base k . Recall that two integers $k, l \geq 2$ are called *multiplicatively independent* if they are not both powers of the same integer. Cobham's theorem states that a sequence is simultaneously automatic with respect to two multiplicatively independent bases if and only if it is ultimately periodic [14]. This result has sparked a lot of research and has been generalised to a variety of different settings. An extension of Cobham's theorem to the class of substitutive sequences was obtained by Durand in 2011 [17].

A considerable effort went also into strengthening Cobham's original theorem. Let x and y be two automatic sequences defined over multiplicatively independent bases. In [22], Fagnot showed that for the claim of Cobham's theorem to hold for x and y it is sufficient that they contain the same factors; that is, if the languages $\mathcal{L}(x)$ and $\mathcal{L}(y)$ coincide, then both x and y are ultimately periodic. In [11] the first- and second-named authors showed that the claim of Cobham's theorem holds if the sequences x and y agree on a set of upper density 1. In the spirit of Shallit's question, Mol, Rampersad, Shallit and Stipulanti obtained in [32] an explicit bound on the length of a common prefix of x and y that depends on the number of states in the automata generating x and y . They further asked for a characterisation of the set $\mathcal{L}(x) \cap \mathcal{L}(y)$ of common factors of x and y . Since all ultimately periodic sequences are k -automatic for all $k \geq 2$, it is clear that we cannot hope for a bound on the length of common factors of x and y . We might hope, however, that the set of common factors exhibits some simple periodic-like structure.

In this paper, we show the following finitary version of Cobham's theorem, which provides a complete characterisation of the sets of words that can appear as common factors of two automatic sequences defined over multiplicatively independent bases. In particular, this set can always be described by a finite amount of data.

Theorem C. *Let $k, l \geq 2$ be multiplicatively independent integers, let \mathcal{A} be an alphabet, and let $U \subset \mathcal{A}^*$. The following conditions are equivalent:*

- (i) *there exist a k -automatic sequence x and an l -automatic sequence y such that U is the set of common factors of x and y ;*
- (ii) *the set U is a finite union of sets of the form $\mathcal{L}({}^\omega v u w^\omega)$, where u, v, w are (possibly empty) words over \mathcal{A} and ${}^\omega v u w^\omega = \dots u v u w u v u w \dots$.*

Note that Cobham's theorem follows immediately from Theorem C. One of the crucial ingredients in the proof of Theorem C is Theorem A applied to k -automatic systems. Indeed, as a simple application of Theorem A and Fagnot's result we can already obtain the following generalisation: if $\mathcal{L}(z) \subset \mathcal{L}(x) \cap \mathcal{L}(y)$ for some k -automatic sequence x and l -automatic sequence y , then z is ultimately periodic (see Corollary 3.3).

Cobham's original result can be understood in the framework of recognisability of subsets of integers in base- k numeration systems. In this context, Cobham's theorem has been seen to hold in many nonstandard numeration systems over the integers as well as in the higher dimensional setting over \mathbf{N}^d , the latter due to Semenov [35]. For a comprehensive treatment of these developments and the interplay between substitutions, numeration systems and logic, see the surveys [10] and [18].

In the setting of recognisable subsets of \mathbf{R}^d , an analogue of Cobham's theorem for integer-based numeration systems was obtained by Boigelot, Brusten and Leroux [8, 9] with recognisability being defined with respect to (weak) Büchi automata. The one-dimensional case was obtained independently by Adamczewski and Bell [2], although it was framed in a different language inspired by the kernel-based definition of automaticity. The two approaches were thoroughly linked in [12], providing further connections with (graph directed) iterated function systems and Cobham-like theorems for iterated function systems obtained in [23] and [20]. For more about these developments we refer to [12] and references therein.

In another direction, Cobham's theorem proved amenable to various algebraic extensions, in part thanks to the characterisation of automaticity in terms of algebraicity of power series over \mathbf{F}_p obtained by Christol [13]. Cobham's theorem has been successfully generalised to the class of regular sequences [6], quasi-automatic functions [1] (introduced by Kedlaya in [25] in order to give a description of the algebraic closure of the field $\mathbf{F}_q(t)$) and Mahler functions (over fields of characteristic 0) [3]. All these proofs made use of Cobham's original theorem. A much simpler proof of Cobham's theorem for Mahler functions that does not rely on the original result of Cobham has been obtained by Schäfke and Singer [34].

It would be interesting to see if the finitary version of Cobham's theorem can be generalised to any of these wider settings. It is easy to see that we cannot hope for a straightforward extension to the class of substitutive sequences since we can construct two non-ultimately periodic substitutive sequences x and y over multiplicatively independent bases such that $\mathcal{L}(x) \subset \mathcal{L}(y)$ (see Remark 3.5 for more details). It is, however, reasonable to hope that some extension of Theorem C holds in a higher dimensional setting and can lead to a generalisation of the Cobham–Semenov theorem. A possible approach could involve extending Theorem A to automatic systems over \mathbf{N}^d .

We briefly discuss the contents of the paper. In the first section, we recap some basic facts about (topological) dynamical systems and substitutions, and introduce the class of substitutive systems that we will work with. In Definition 1.7 we introduce the notion of an idempotent substitution and show that for every substitution φ some power φ^n is idempotent. Idempotency gathers all the technical assumptions that we need from the substitution in order to carry out the proofs in Section 2. In Lemma 1.2 we show that certain sequences that will turn out to be closely linked with sequences generating transitive subsystems of substitutive (resp., k -automatic) systems are indeed substitutive (resp., k -automatic).

The second section is devoted to the classification of minimal and transitive subsystems of substitutive systems, and contains the proofs of Theorems A and B. The analogues of Theorems A and B for invertible substitutive systems are presented at the end of the section.

The proof of Theorem C is given in the third section. The main ingredients in the proof are Theorem A and Proposition 3.6, which describes occurrences of cyclic factors in automatic sequences. We also discuss the problem of effective computability of the set of common factors, that is, existence of an algorithm that, given as input two automatic sequences $x, y \in \mathcal{A}^\omega$ defined over multiplicatively independent bases, returns a finite set of triples of words $(v, u, w) \in (\mathcal{A}^*)^3$ that describe the set of common factors $\mathcal{L}(x) \cap \mathcal{L}(y)$ in the sense of Theorem C. Our proof of Theorem C uses the compactness of the space \mathcal{A}^ω and is not effective. We believe that the question of whether Theorem C admits an effective proof is interesting and worthy of further study.

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1. Preliminaries

In this section we recall some classical definitions and state a few preliminary lemmas.

Symbolic dynamics. A (topological) dynamical system is a compact metric space X together with a continuous map $T: X \rightarrow X$. We denote by T^n the n -th iterate of T and by $\mathcal{O}(x) = \{T^n(x) \mid n \geq 0\}$ the orbit of a point $x \in X$. A point $x \in X$ is *periodic* if $T^k(x) = x$ for some $k \geq 1$. A point $x \in X$ is *ultimately periodic* if there exists $m \geq 0$ such that $T^m(x)$ is periodic. A *subsystem* of X is a closed subset of X that is invariant under the map T . A system X is called *minimal* if $X \neq \emptyset$ and if X has no subsystems other than \emptyset and X ; equivalently, a system $X \neq \emptyset$ is minimal if the orbit of every point is dense in X [19, Ex. 4.2.1.a]. A system X is called *transitive* if it has a point with a dense orbit. An easy application of Zorn's lemma shows that every dynamical system has a minimal subsystem [19, Ex. 4.2.1.c]. We say that a dynamical system (Y, S) is a (topological) *factor* of the system (X, T) if there exists a continuous surjective map $\pi: X \rightarrow Y$ such that $\pi \circ T = S \circ \pi$. Such a map π is called a *factor map*. We will need the following simple fact.

Lemma 1.1. *Let X and Y be dynamical systems and let $\pi: X \rightarrow Y$ be a factor map. Let Y' be a minimal subsystem of Y . Then there exists a minimal subsystem X' of X such that $\pi(X') = Y'$.*

Proof. Let $X'' \subset X$ be the preimage of Y' by the map π . Clearly, X'' is a subsystem of X . Let X' be some minimal subsystem of X'' . Then $\pi(X')$ is a subsystem of Y' , and since Y' is minimal, $\pi(X') = Y'$. \square

In this paper we are interested in dynamical systems coming from substitutive sequences. Let \mathcal{A} be a finite set (called an *alphabet*). We denote by \mathcal{A}^* the set of finite words over \mathcal{A} . This is a monoid under concatenation. The empty word is denoted by ϵ . We say that a word w is a *factor* of a word v or that w *appears* in v if $v = ywz$ for some words y and z . A word w is a *prefix* of a word v if $v = wz$ for some word z . We similarly define a *suffix*. For a word w we denote by $|w|$ the length of w .

We denote by \mathcal{A}^ω the set of sequences over \mathcal{A} . For a sequence x and integers $i \leq j$ we write $x_{[i, j]}$ for the word $x_i x_{i+1} \cdots x_{j-1}$ and $x_{[i, \infty)}$ for the sequence $x_i x_{i+1} \cdots$. (In particular, $x_{[i, i)} = \epsilon$.) The notions of concatenation, factor, prefix and suffix are used for words and sequences as long as they make obvious sense. For a word $w \neq \epsilon$ we denote by w^ω the sequence $w^\omega = www \cdots$, and we put $\epsilon^\omega = \epsilon$. While we will always use the notation x_n for the n -th term of a sequence $x = (x_n)_{n \geq 0} \in \mathcal{A}^\omega$, we regard words themselves as not indexed.

The set \mathcal{A}^ω with the product topology (where we use discrete topology on each copy of \mathcal{A}) is a compact metrisable space. We define the shift map $T: \mathcal{A}^\omega \rightarrow \mathcal{A}^\omega$ by $T((x_n)_n) = (x_{n+1})_n$. The space \mathcal{A}^ω together with the shift map T is a dynamical system. We refer to subsystems X of \mathcal{A}^ω as *subshifts*. By $\mathcal{L}(X)$ we denote the *language* of the subshift X , i.e. the set of all finite words that appear in some $x \in X$. A subshift X is uniquely determined by its language since

$$X = \{x \in \mathcal{A}^\omega \mid \text{all factors of } x \text{ are in } \mathcal{L}(X)\}.$$

We also use $\mathcal{L}(y)$ to denote the set of factors of a word or a sequence y . By a slight abuse of terminology we say that a sequence of words w_n converges to the sequence x if $|w_n| \rightarrow \infty$ and for every $m \geq 0$ the prefixes of x and w_n of length m agree for sufficiently large n .

We will occasionally also work with backwards infinite sequences in ${}^\omega\mathcal{A}$ and biinfinite sequences in ${}^\omega\mathcal{A}^\omega$. The definitions of factor, prefix, suffix and language generalise to these cases in a straightforward manner. We always regard sequences $(a_n)_n$ in \mathcal{A}^ω as indexed by $n \in \{0, 1, \dots\}$, backwards infinite sequences in ${}^\omega\mathcal{A}$ as indexed by $n \in \{\dots, -2, -1\}$, and biinfinite sequences in ${}^\omega\mathcal{A}^\omega$ as indexed by $n \in \{\dots, -2, -1, 0, 1, \dots\}$.

Substitutive sequences. Let \mathcal{A} be an alphabet. A *substitution* is a map $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ that assigns to each letter a some finite word w in \mathcal{A}^* . We only consider substitutions that are *growing*, i.e. $|\varphi^n(a)| \rightarrow \infty$ as $n \rightarrow \infty$ for each $a \in \mathcal{A}$, and throughout the paper the term ‘substitution’ is used for a growing substitution. A substitution φ is called *primitive* if there exists an integer $n \geq 1$ such that for any $a, b \in \mathcal{A}$ the letter a appears in $\varphi^n(b)$. A letter $a \in \mathcal{A}$ is *prolongable* if a is the initial letter of $\varphi(a)$. A letter $a \in \mathcal{A}$ is *backwards prolongable* if a is the final letter of $\varphi(a)$. If a is prolongable (resp., backwards prolongable), then the sequence $\varphi^n(a)$ converges to a sequence in \mathcal{A}^ω (resp., in ${}^\omega\mathcal{A}$) that we denote by $\varphi^\omega(a)$ (resp., ${}^\omega\varphi(a)$). A *coding* is an arbitrary map $\pi: \mathcal{A} \rightarrow \mathcal{B}$ between alphabets \mathcal{A} and \mathcal{B} . A surjective coding π naturally extends to a factor map $\pi: \mathcal{A}^\omega \rightarrow \mathcal{B}^\omega$ between dynamical systems.

A substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ induces a natural map $\varphi: \mathcal{A}^\omega \rightarrow \mathcal{A}^\omega$, denoted by the same letter. We say that a sequence x is *purely substitutive* if it is a fixed point of some substitution φ , i.e. $\varphi(x) = x$. In this case we also say that the sequence x is *produced* by the substitution φ . Sequences produced by a substitution φ are exactly of the form $\varphi^\omega(a)$ for a prolongable letter a . A *substitutive* sequence is the image of a purely substitutive sequence under a coding.

We say that a substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ is of *constant length* k if $|\varphi(a)| = k$ for each $a \in \mathcal{A}$. A fixed point of a substitution of constant length k is called a *purely k -automatic sequence*. A *k -automatic* sequence is the image of a purely k -automatic sequence under a coding. The classes of substitutive and k -automatic sequences are invariant under changing finitely many terms of a sequence and under the forward and backward shift operations [5, Cor. 6.8.5 and Thm. 7.6.1 & 7.6.3]. We also mention here the trivial case of Cobham’s theorem, which says that for any integer $t \geq 1$ the classes of k -automatic and k^t -automatic sequences coincide [5, Theorem 6.6.3].

The term *automatic* has its origin in theoretical computer science. Informally speaking, automata (or more precisely finite deterministic k -automata with output) are simple finite computational devices that compute the n -th term of a sequence from the base- k digits of n . For more details, see [5]. The famous theorem of Cobham asserts that the description of k -automatic sequences in terms of substitutions of constant length k and in terms of k -automata are equivalent [5, Thm. 6.3.2], [15]. In this paper, we will only work with the former definition. Note that the assumption that a substitution is growing is trivially satisfied when the substitution is of constant length $k \geq 2$.

Yet another definition of an automatic sequence can be given in terms of kernels. Let $x = (x_n)_{n \geq 0}$ be a sequence over an alphabet \mathcal{A} . Let $\mathbf{N} = \{0, 1, \dots\}$ denote the nonnegative integers. The *k -kernel* of x is defined as the family of sequences

$$K_k(x) = \{(x_{k^m n + r})_{n \geq 0} \mid m, r \in \mathbf{N}, 0 \leq r < k^m\}.$$

A theorem of Cobham asserts that a sequence is k -automatic if and only if its k -kernel is finite (see [15] or [5, Thm. 6.6.2]).

Later we will need the following result.

Lemma 1.2. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution and let $w \in \mathcal{A}^*$ be nonempty. Consider the sequence $x = w\varphi(w)\varphi^2(w)\cdots$.*

- (i) *The sequence x is substitutive.*
- (ii) *If φ is of constant length k , then x is k -automatic.*

Proof. Part (i) follows from [16, Lemma 5]; since the argument is short, we include it for completeness. Let \spadesuit be a letter not belonging to the alphabet \mathcal{A} and consider the substitution $\tau: \mathcal{A} \cup \{\spadesuit\} \rightarrow (\mathcal{A} \cup \{\spadesuit\})^*$ given by $\tau(\spadesuit) = \spadesuit w$, $\tau(a) = \varphi(a)$ for $a \in \mathcal{A}$. The sequence $\tau^\omega(\spadesuit)$ takes the form $\spadesuit x$ and is clearly substitutive. Therefore, x is substitutive as well.

We now assume that φ is of constant length k . We will prove part (ii). We first reduce to the case where $k \geq 3$ by replacing w by $w\varphi(w)$ and φ by φ^2 . Next, we reduce to the case when $|w| = 1$. The substitution φ induces in a natural way a substitution φ' on the alphabet $\mathcal{A}^{|w|}$ consisting of all the words over \mathcal{A} of length $|w|$. Let a' denote the word w considered as an element of $\mathcal{A}^{|w|}$ and let $x' = a'\varphi'(a')\varphi'^2(a')\cdots$. The sequence x is then the image of x' by the morphism $\mathcal{A}^{|w|} \rightarrow \mathcal{A}^*$ that sends each letter from $\mathcal{A}^{|w|}$ to the corresponding word in \mathcal{A}^* . Since images of k -automatic sequences by morphisms of constant length are k -automatic [5, Cor. 6.8.3], we get the desired reduction.

We now assume that $|w|=1$. Let $v \in \mathcal{A}^*$ be any word of length $k-2$. Let \spadesuit be a letter not belonging to the alphabet \mathcal{A} and consider the substitution $\tau: \mathcal{A} \cup \{\spadesuit\} \rightarrow (\mathcal{A} \cup \{\spadesuit\})^*$ of constant length k given by $\tau(\spadesuit) = v\spadesuit w$, $\tau(a) = \varphi(a)$ for $a \in \mathcal{A}$. The substitution τ is a so-called $(k-2, 1)$ -morphism (see [5, 6.7]), and in the terminology of [5, 5.3] the two-sided sequence

$$y = \cdots \varphi^2(v)\varphi(v)v\spadesuit w\varphi(w)\varphi^2(w)\cdots$$

is $(k-2, 1)$ -automatic by [5, Thm. 6.7.2]. It follows from [5, Theorem 5.3.4] that $y_{[1,\infty]} = w\varphi(w)\varphi^2(w)\cdots$ is then k -automatic. \square

Remark 1.3. The proof of Lemma 1.2.(ii) above uses the notion of (k, l) -automatic sequence. It is possible, albeit somewhat lengthy and tedious, to prove this result using only the notion of k -automatic sequence and showing directly that the k -kernel of x is finite.

Substitutive systems. Let \mathcal{A} be an alphabet. A system $X \subseteq \mathcal{A}^\omega$ is called *purely substitutive* (resp., *substitutive*, *k -automatic*) if it arises as the orbit closure of a purely substitutive (resp., substitutive, k -automatic) sequence. Note that any such system is automatically transitive.

There is a more general notion of systems arising from substitutions. Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution and let X_φ denote the dynamical system generated by φ , i.e.

$$X_\varphi = \{z \in \mathcal{A}^\omega \mid \text{every factor of } z \text{ appears in } \varphi^n(a) \text{ for some } n \geq 0 \text{ and } a \in \mathcal{A}\}.$$

The system X_φ does not have to be transitive; consider, e.g. the substitution $\varphi: \{0, 1\} \rightarrow \{0, 1\}^*$ given by $\varphi(0) = 00$, $\varphi(1) = 11$ for which $X_\varphi = \{0^\omega, 1^\omega\}$. It is clear that every substitutive system has the form $\pi(X_\varphi)$ for some substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ and coding $\pi: \mathcal{A} \rightarrow \mathcal{B}$. It is also well-known that if X_φ is minimal, then it is substitutive [33, section 5.2]. Proposition 2.8 below shows more generally that a system X_φ is substitutive if and only if it is transitive.

Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution. For $b \in \mathcal{A}$ we denote by \mathcal{A}_b the set of all letters $a \in \mathcal{A}$ appearing in $\varphi^n(b)$ for some $n \geq 0$. For $a, b \in \mathcal{A}$, we write $b \geq a$ if $a \in \mathcal{A}_b$. Note that the relation \geq is only a preorder on the set \mathcal{A} . If $b \geq a \geq b$ we write $a \sim b$. We also write $b > a$ if $b \geq a$ and $b \not\sim a$. The relation \sim is an equivalence relation and $>$ is a strict partial order on \mathcal{A} . These relations clearly depend on the substitution φ , but we will nevertheless write $b > a$ or $b \geq a$ and

call a letter minimal, maximal or equivalent to another letter when the substitution is clear from the context. For $b \in \mathcal{A}$ let $X_{\varphi,b}$ denote the subsystem of X_φ generated by b , i.e.

$$X_{\varphi,b} = \{z \in X \mid \text{every factor of } z \text{ appears in } \varphi^n(b) \text{ for some } n \geq 0\}.$$

Thus, $X_{\varphi,b}$ is equal to the system $X_{\varphi'}$, where φ' is the substitution φ restricted to the alphabet \mathcal{A}_b . We will often write X_b instead of $X_{\varphi,b}$ when the substitution φ is clear from the context. Note that if $a \geq b$, then $X_a \supseteq X_b$. Note also that $X_\varphi = \bigcup_{a \in \mathcal{A}} X_a$ and $\varphi(X_b) \subset X_b$ for each $b \in \mathcal{A}$.

The following lemma is a variant of [16, Lemma 6] (cf. [7, Prop. 5.10] for a version for two-sided dynamical systems). We recall that x_n denotes the n -th term of a sequence $x = (x_n)_{n \geq 0} \in \mathcal{A}^\omega$.

Lemma 1.4. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution and let $x \in X_\varphi$. There exist $y \in X_\varphi$ and a proper prefix u of $\varphi(y_0)$ such that $ux = \varphi(y)$.*

Proof. Let v_n be the prefix of x of length $n \geq 1$. Since $x \in X_\varphi$, there exist $b \in \mathcal{A}$ and $k \geq 1$ such that v_n is a factor of $\varphi^k(b)$. Removing from $\varphi^{k-1}(b)$ the longest prefix whose image by φ does not intersect v_n , we obtain a suffix w_n of $\varphi^{k-1}(b)$ with initial letter a_n and a proper prefix u_n of $\varphi(a_n)$ such that $u_n x$ and $\varphi(w_n)$ agree on the first n positions.

Since there are only finitely many possibilities for u_n and \mathcal{A}^ω is compact, there exist $u \in \mathcal{A}^*$, $y \in X_\varphi$ and an increasing sequence $(k_m)_{m \geq 0}$ of positive integers such that $u_{k_m} = u$ for all $m \geq 0$ and w_{k_m} converge to y as $m \rightarrow \infty$. By construction, u is then a proper prefix of $\varphi(y_0)$ and we have $ux = \varphi(y)$. \square

The following lemma is well-known for primitive substitutions (see, e.g. [33, Prop. 5.4]); we will however need the claim under the weaker assumption of transitivity.

Lemma 1.5. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution. If X_φ is transitive, then $X_\varphi = X_{\varphi^n}$ for any $n \geq 1$.*

Proof. Let $n \geq 1$ and for $0 \leq i \leq n-1$ write

$$X_i = \{x \in X_\varphi \mid \text{every factor of } x \text{ appears in } \varphi^k(b) \text{ for some } b \in \mathcal{A} \text{ and } k \equiv i \pmod{n}\}.$$

Note that $X_0 = X_{\varphi^n}$ and $X_\varphi = \bigcup_{i=0}^{n-1} X_i$. Since X_φ is transitive, $X_\varphi = X_j$ for some $0 \leq j \leq n-1$. Let $x \in X_\varphi$. By Lemma 1.4 there exist $y \in X_\varphi = X_j$ and a proper prefix u of $\varphi(y_0)$ such that $ux = \varphi(y)$. Since every factor of y is a factor of $\varphi^k(b)$ for some $b \in \mathcal{A}$ and $k \equiv j \pmod{n}$, every factor of x is a factor of $\varphi^{k+1}(b)$ and hence x lies in X_{j+1} (where $X_n = X_0$). Thus $X_\varphi = X_{j+1}$. Repeating the argument, we get that $X_\varphi = X_i$ for all $0 \leq i \leq n-1$. In particular, $X_\varphi = X_0 = X_{\varphi^n}$. \square

Remark 1.6. The result above is not necessarily true without the assumption of transitivity. For an example, consider the substitution $\varphi: \{0, 1, 2\} \rightarrow \{0, 1, 2\}^*$ given by $\varphi(0) = 12$, $\varphi(1) = 22$, $\varphi(2) = 11$. Then

$$X_\varphi = \{1^n 2^\omega \mid n \geq 0\} \cup \{2^n 1^\omega \mid n \geq 0\} \quad \text{and} \quad X_{\varphi^2} = \{2^n 1^\omega \mid n \geq 0\} \cup \{2^\omega\}.$$

Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution. We say that a letter $a \in \mathcal{A}$ is *ample* if a appears in $\varphi^n(a)$ for some $n \geq 1$, and *very ample* if a appears at least twice in $\varphi^n(a)$ for some $n \geq 1$. Note that for all $n \geq 1$ the sets of ample and very ample letters with respect to substitutions φ and φ^n are the same. We denote the set of all ample letters by \mathcal{A}' .

We note two easy properties of ampleness. First, a letter equivalent to an ample letter is itself ample. Second, for any ample letter a the word $\varphi(a)$ contains at least one letter equivalent to a . For $a \in \mathcal{A}'$ let $\lambda_\varphi(a)$ denote the letter b which is equivalent to a and which occurs in $\varphi(a)$ at the last position among all letters equivalent to a . This gives rise to a map $\lambda_\varphi: \mathcal{A}' \rightarrow \mathcal{A}'$. Note that $\lambda_\varphi^n = \lambda_{\varphi^n}$ for all $n \geq 1$.

Let S be a set and let $\psi: S \rightarrow S$ be a map. We say that ψ is *idempotent* if $\psi^2 = \psi$. Note that if x is a substitutive sequence produced by a substitution φ , then x is also produced by the substitution φ^n for any integer $n \geq 1$. Similarly, the notions of k -automatic and k^n -automatic

sequences coincide [5, Thm. 6.6.4]. For these reasons, we may freely replace φ by φ^n , which will often have nicer properties. In Definition 1.7 we gather all the technical properties of the substitution that we intend to obtain in this manner. We will often need only some of these properties, but for simplicity we will not attempt to always state the precise minimal hypotheses.

Definition 1.7. A substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ is called *idempotent* if it satisfies the following conditions:

- (i) for all $a \in \mathcal{A}$ and $n \geq 1$ the set of letters appearing in $\varphi(a)$ is the same as the set of letters appearing in $\varphi^n(a)$;
- (ii) for all $a \in \mathcal{A}$ and $n \geq 1$ the set of letters appearing at least twice in $\varphi(a)$ is the same as the set of letters appearing at least twice in $\varphi^n(a)$;
- (iii) for all $a \in \mathcal{A}$ the initial letter of $\varphi(a)$ is prolongable;
- (iv) the map $\lambda_\varphi: \mathcal{A}' \rightarrow \mathcal{A}'$ is idempotent.

Note that if $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ is an idempotent substitution, then for each $b \in \mathcal{A}$, \mathcal{A}_b consists exactly of b and the letters appearing in $\varphi(b)$. Furthermore, a letter b is ample if and only if b appears in $\varphi(b)$ and it is very ample if and only if it appears at least twice in $\varphi(b)$.

Lemma 1.8. *Let S be a finite set, and let $\psi: S \rightarrow S$ be a map. There exists an integer $m \geq 1$ such that ψ^m is idempotent.*

Proof. Take, for example, $m = |S|!$. □

Lemma 1.9. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution. There exists an integer $m \geq 1$ such that the substitution φ^m is idempotent.*

Proof. Note first that properties (i), (ii), (iii) and (iv) in Definition 1.7 are preserved after replacing φ by its iterate. We will first choose m so that properties (i) and (ii) hold. For $a \in \mathcal{A}$ and a substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$, let $S_\varphi(a)$ denote the set of letters appearing in $\varphi(a)$. Consider the map $\psi: 2^\mathcal{A} \rightarrow 2^\mathcal{A}$ that sends a subset A of \mathcal{A} to the set $\bigcup_{a \in A} S_\varphi(a)$. Note that $\psi^n(\{a\}) = S_{\varphi^n}(a)$ for all $n \geq 1$ and $a \in \mathcal{A}$. By Lemma 1.8 there exists an integer $m \geq 1$ such that $\psi^m = \psi^{2m}$. This implies that $S_{\varphi^m}(a) = S_{\varphi^{2m}}(a)$ for all $a \in \mathcal{A}$ and $n \geq 1$, and hence the substitution φ^m satisfies property (i). We then obtain property (ii) by repeating the reasoning above with the set $2^\mathcal{A}$ replaced by the set $3^\mathcal{A}$, which includes the information on whether a letter appears in a word at least twice, exactly once or not at all. Thus, we may assume that properties (i) and (ii) hold.

To prove the remaining properties, let $\alpha_\varphi(a)$ denote the initial letter of $\varphi(a)$ for $a \in \mathcal{A}$. Note that $\alpha_\varphi^n = \alpha_{\varphi^n}$ for all $n \geq 1$. By Lemma 1.8, there exists $m' \geq 1$ such that $\alpha_{\varphi^{m'}} = \alpha_{\varphi^{2m'}}$, and hence $\varphi^{m'}$ satisfies property (iii). A similar reasoning applied to the map $\lambda_\varphi: \mathcal{A}' \rightarrow \mathcal{A}'$ proves that some iterate of $\varphi^{m'}$ satisfies the remaining property (iv). □

2. Subsystems of substitutive systems

This section studies subsystems of substitutive systems. Recall that every substitutive system is a topological factor of a purely substitutive system X (the factor map being given by a coding), and that we may assume that the substitutive sequence generating X is produced by an idempotent substitution (see Lemma 1.9). Our main result is Theorem 2.1 below, which says that transitive subsystems of substitutive systems are still substitutive. The proof of this result will occupy the whole section. We will first prove the statement for purely substitutive sequences and obtain the general result by an easy reduction. Along the way, we will obtain a more detailed description of all transitive subsystems.

Theorem 2.1. *Every transitive subsystem of a substitutive system is substitutive. Every transitive subsystem of a k -automatic system is k -automatic.*

Minimal subsystems of substitutive systems. We start by investigating minimal subsystems of (purely) substitutive systems. Actually, we work in a slightly more general context of systems of the form X_φ with φ idempotent. We will show that in this case all minimal subsystems arise as X_b for some minimal letter b . In particular, every minimal subsystem of a substitutive system is substitutive and every minimal subsystem of a k -automatic system is k -automatic.

Proposition 2.2. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be an idempotent substitution. Let Y be a subsystem of X_φ . Then Y is minimal if and only if $Y = X_b$ for some minimal letter $b \in \mathcal{A}$.*

Proof. First we show that every system X_b with b minimal is minimal. If $b \in \mathcal{A}$ is a minimal letter, then the substitution $\varphi|_{\mathcal{A}_b}: \mathcal{A}_b \rightarrow \mathcal{A}_b^*$ is primitive. Since every primitive substitution gives rise to a minimal system [33, Prop. 5.5], X_b is minimal.

Now assume that Y is a minimal subsystem of X_φ . Fix an integer $m \geq 1$. Choosing a sufficiently long word $w \in \mathcal{L}(Y)$, we can find a letter $a \in \mathcal{A}$ such that $\varphi^m(a)$ appears in w , and hence $\varphi^m(a) \in \mathcal{L}(Y)$. Since \mathcal{A} is finite, there is some letter $a \in \mathcal{A}$ such that $\varphi^m(a) \in \mathcal{L}(Y)$ for infinitely many m . Since φ is idempotent, the set of letters appearing in $\varphi^l(a)$ is independent of $l \geq 1$, and hence some minimal letter b appears in $\varphi^l(a)$ for all $l \geq 1$. It follows that $\varphi^n(b) \in \mathcal{L}(Y)$ for all $n \geq 0$, and hence $X_b \subset Y$. By minimality of Y , we have $X_b = Y$. \square

Corollary 2.3. *Let X be a substitutive system. The number of minimal subsystems of X is finite.*

Proof. If X is purely substitutive, then it is of the form $X = X_\varphi$ for some idempotent substitution φ , and the claim follows from Proposition 2.2. In the general case, write X as a topological factor of a purely substitutive system and use Lemma 1.1. \square

Remark 2.4. A very special case of Proposition 2.2 was proven in a different language in [11, Lemma 2.3] by the first-named and the second-named author for constant length substitutions and one-point subsystems (and with a slightly weaker notion of idempotency).

Transitive subsystems of substitutive systems. Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution. If a is a (not necessarily minimal) letter in \mathcal{A} , then X_a is a subsystem of X_φ . It would be tempting to conjecture that all transitive subsystems of X_φ are of this form. The following examples show that this is not the case.

Example 2.5.

(i) Let $\mathcal{A} = \{0, 1, 2, 3\}$ and let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be the substitution given by

$$\varphi(0) = 12, \quad \varphi(1) = 11, \quad \varphi(2) = 23, \quad \varphi(3) = 32.$$

Let y denote the biinfinite sequence $y = \omega_\varphi(1)\varphi^\omega(2)$. For an integer n consider the suffix $y_{[n, \infty)} = y_n y_{n+1} \dots$ of y . This is just the Thue–Morse sequence on the alphabet $\{2, 3\}$ with n first symbols removed if $n \geq 0$ or preceded by $1^{|n|}$ if $n < 0$, and it lies in X_φ since every factor of y is a factor of some $\varphi^n(12) = \varphi^{n+1}(0)$. Consider the subsystems $Y_n = \overline{\mathcal{O}(y_{[n, \infty)})} \subset X_\varphi$. For $n \geq 0$, the system Y_n is just the Thue–Morse system (since it is minimal), while for $n < 0$ it is equal to the Thue–Morse system with $|n|$ extra points adjoined. Hence, for $n < 0$ the systems Y_n are pairwise distinct, and are different from each X_b for $b \in \mathcal{A}$. In fact, $X_\varphi = \bigcup_{n \leq 0} Y_n \cup \{1^\omega\}$ and $X_0 = X_\varphi$, $X_1 = \{1^\omega\}$ and $X_2 = X_3 = Y_0$ (cf. Corollary 2.7 below and note that X_2 is minimal).

(ii) Let $\mathcal{A} = \{0, 1, 2, 3\}$ and let $\tau: \mathcal{A} \rightarrow \mathcal{A}^*$ be the substitution given by

$$\tau(0) = 01023, \quad \tau(1) = 12, \quad \tau(2) = 22, \quad \tau(3) = 33.$$

Write $v = 01$ and $w = 23$. Let z denote the biinfinite sequence (indexed so that the 0 below occurs at the 0-th position)

$$z = \dots \tau^2(v)\tau(v)v0w\tau(w)\tau^2(w)\dots$$

For an integer n consider the suffix $z_{[n,\infty)} = z_n z_{n+1} \cdots$ of z . Every factor of $z_{[n,\infty)}$ is a factor of some $\tau^m(0)$, and hence $z_{[n,\infty)}$ lies in X_τ . Consider the subsystems $Z_n = \overline{\mathcal{O}(z_{[n,\infty)})} \subset X_\tau$. It is easy to see that $Z_n = \mathcal{O}(z_{[n,\infty)}) \cup \{3^k 2^\omega \mid k \geq 0\} \cup \{2^k 3^\omega \mid k \geq 0\}$, and hence the systems Z_n are pairwise distinct and different from each X_b for $b \in \mathcal{A}$.

Now assume that $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ is an idempotent substitution. The next proposition characterises points $y \in X_\varphi$ such that $\overline{\mathcal{O}(y)}$ is not equal to any X_b for $b \in \mathcal{A}$. We show that all such points are substitutive. Note that this is by no means obvious. In fact, substitutive systems have often continuum many points (e.g. the Thue–Morse system), while the number of substitutive sequences over a given alphabet is only countable.

Proposition 2.6. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be an idempotent substitution. Let $y \in X_\varphi$ and let Y be the orbit closure of y . Then at least one of the following conditions holds:*

- (a.i) *there exists a letter a in $\mathcal{L}(X_\varphi)$ such that $y \in X_a$;*
- (a.ii) *there exist a backwards prolongable letter a and a prolongable letter c such that $ac \in \mathcal{L}(X_\varphi)$ and y is a suffix of ${}^\omega\varphi(a)\varphi^\omega(c)$.*

Assume moreover that Y is different from each X_b for $b \in \mathcal{A}$. Then at least one of the following conditions holds:

- (b.i) *there exists a letter a such that $\varphi(a) = v_a a w_a$ for some words v_a and w_a such that $w_a \neq \epsilon$, w_a contains only letters b such that $b < a$, and y is a suffix of*

$$\cdots \varphi^2(v_a)\varphi(v_a)v_a a w_a \varphi(w_a)\varphi^2(w_a)\cdots;$$

- (b.ii) *the sequence y satisfies condition (a.ii).*

Proof. Since $X_\varphi = \bigcup_{b \in \mathcal{A}} X_b$ and Y is transitive, we see that Y is contained in X_b for some $b \in \mathcal{A}$. We choose b to be minimal among such letters.

Let $y^0 = y$. Using Lemma 1.4, we inductively construct for $i \geq 0$ letters a_i , sequences $y^i \in X_b$ with initial letters a_i , and proper prefixes u_i of $\varphi(a_{i+1})$ such that

$$\varphi(y^{i+1}) = u_i y^i.$$

Note that since a_i appears in $\varphi(a_{i+1})$, we have $a_{i+1} \geq a_i$, and hence letters a_i become equivalent for sufficiently large i . We will now show that y satisfies one of the properties (a.i) and (a.ii). We consider two cases.

Case I (for infinitely many i the length of u_i is strictly smaller than $|\varphi(a_{i+1})| - 1$). In this case for infinitely many i the prefix of y^i of length 2 is a factor of $\varphi(a_{i+1})$. Let a be a letter that occurs infinitely many times among a_i . Since a is the initial letter of some y^i , it lies in $\mathcal{L}(X_\varphi)$. Since φ is growing and $\varphi(y^{i+1}) = u_i y^i$, it follows that every prefix of y is a factor of $\varphi^i(a)$ for some $i \geq 0$. In particular, $y \in X_a$ and hence property (a.i) holds.

Case II (for all sufficiently large i the length of u_i is equal to $|\varphi(a_{i+1})| - 1$). Let i_0 be such that we have $|u_i| = |\varphi(a_{i+1})| - 1$ and $a_{i+1} \sim a_i$ for all $i \geq i_0$. Take $i \geq i_0$. Then a_i is the final letter of $\varphi(a_{i+1})$, which implies that a_{i+1} is ample and $\lambda(a_{i+1}) = a_i$. Since the map λ is idempotent, we have $a_{i+1} = a_i$ for all $i \geq i_0$. Denote this letter by a , and note that it is backwards prolongable.

For $i \geq i_0$ the sequence $T(y^i)$ is the image of $T(y^{i+1})$ by φ . Iterating this for $i \geq i_0$, we see that for each $n \geq 0$ and $d = T(y^{i_0+n})_0$ the word $\varphi^n(d)$ is a prefix of $T(y^{i_0})$. Choose some letter d that arises in this manner for infinitely many n , and put $c = \varphi(d)_0$. Since φ is idempotent, c is prolongable, and the assumption on d shows that $T(y^{i_0}) = \varphi^\omega(c)$. This shows that $y^{i_0} = a\varphi^\omega(c)$ and in particular $ac \in \mathcal{L}(X_\varphi)$. Since y is a suffix of $\varphi^{i_0}(y^{i_0})$, it is also a suffix of ${}^\omega\varphi(a)\varphi^\omega(c)$ and property (a.ii) holds. This ends the proof of the first assertion.

Now assume that Y is different from each X_c for $c \in \mathcal{A}$. To show the second claim, we only need to treat Case I. We will show that in this case y satisfies property (b.i). As in the reasoning

above, let a be a letter that occurs infinitely many times among a_i , and recall that $a \in \mathcal{L}(X_\varphi)$ and $y \in X_a$. By the choice of b it follows that $a \sim b$ and b is ample.

We claim that for sufficiently large i the sequence y^i contains no letters equivalent to a at non-initial positions. Indeed, if y^i contains a letter $c \sim a$ at a non-initial position, then y contains $\varphi^i(c)$. If this happened for infinitely many i , the word $\varphi^n(a)$ would appear in y for each $n \geq 0$, contradicting the assumption that Y is a proper subset of X_b .

Since for sufficiently large i the letters a_i are all equivalent (and hence ample) and since the sequence y^i contains no letters equivalent to b at non-initial positions, we have $\lambda(a_{i+1}) = a_i$. Since the map λ is idempotent, the sequence a_i is eventually constant with value a . It follows that $\lambda(a) = a$ and $\varphi(a) = v_a a w_a$ with w_a nonempty (since we are in Case I) and containing only letters $c < a$.

Choose $i_0 \geq 0$ so that for $i \geq i_0$ we have $a_i = a$ and the sequence y^i contains no letters equivalent to a at non-initial positions. Since $\varphi(y^{i+1}) = u_i y^i$, we must have

$$y^{i_0} = a w_a \varphi(w_a) \varphi^2(w_a) \cdots .$$

Hence y is a suffix of

$$\varphi^{i_0}(y^{i_0}) = \varphi^{i_0-1}(v_a) \cdots \varphi^2(v_a) \varphi(v_a) v_a a w_a \varphi(w_a) \varphi^2(w_a) \cdots . \quad \square$$

Corollary 2.7. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be an idempotent substitution. Let $y \in X_\varphi$ and let Y be the orbit closure of y . Assume that Y is different from each X_b for $b \in \mathcal{A}$. Then y is a substitutive sequence. If furthermore φ is a substitution of constant length k , then y is k -automatic.*

Proof. This follows immediately from Proposition 2.6, Lemma 1.2 and the fact that substitutive (resp., k -automatic) sequences are closed under backward and forward shifting. \square

The next proposition characterises systems X_φ that are transitive.

Proposition 2.8. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution. Let $n \geq 1$ be such that φ^n is idempotent. The following conditions are equivalent:*

- (i) X_φ is transitive,
- (ii) $X_\varphi = X_{\varphi^n, b}$ for some letter $b \in \mathcal{A}$ that is either prolongable under φ^n or very ample.

Moreover, if X_φ is transitive, then it is substitutive. Furthermore, if φ is a substitution of constant length k , then X_φ is k -automatic.

Proof. Note first that under either of the assumptions (i) and (ii) we have $X_\varphi = X_{\varphi^n}$ (in the former case by Lemma 1.5, in the latter case it is obvious). Hence, we may assume that φ itself is idempotent and $n = 1$.

We first show that (ii) implies both (i) and the final claim. If b is prolongable, then $X_\varphi = X_b$ is the orbit closure of $\varphi^\omega(b)$, from which all the remaining claims follow easily. Suppose that this is not the case. Then b appears at least twice in $\varphi(b)$ and we can write $\varphi(b) = vbw$, where v, w are words such that b appears in w . The word $x_n = w\varphi(w) \cdots \varphi^n(w)$ is a suffix of $\varphi^{n+1}(b)$ and hence the sequence $x = w\varphi(w)\varphi^2(w) \cdots$ lies in X_b . Since $\varphi^n(b)$ is a factor of x for all $n \geq 0$, the orbit of x is dense in X_b . The sequence x is substitutive by Lemma 1.2. Furthermore, if φ is a substitution of constant length k , then x is k -automatic.

It remains to prove that (i) implies (ii). Assume that X_φ is transitive and let y be a point in X_φ with a dense orbit. Since $X_\varphi = \bigcup_{b \in \mathcal{A}} X_b$, there exists $b \in \mathcal{A}$ such that $X_\varphi = X_b$; pick minimal b with this property. Suppose that b is not prolongable and not very ample. It means that b appears in $\varphi(b)$ at most once, and at a non-initial position.

Case I (b appears in $\varphi(b)$ exactly once, and at a non-initial and non-final position). Write $\varphi(b) = vbw$ for nontrivial words $v, w \in \mathcal{A}^*$. Every word

$$\varphi^n(b) = \varphi^{n-1}(v) \cdots \varphi(v) vbw \varphi(w) \cdots \varphi^{n-1}(w)$$

contains exactly one occurrence of b , and hence every point in X_b contains at most one occurrence of b . On the other hand, every suffix of the biinfinite sequence $\cdots \varphi^2(v)\varphi(v)vbw\varphi(w)\varphi^2(w)\cdots$ lies in X_b , and hence X_b contains infinitely many points in which b appears. It follows that X_b is not transitive.

Case II (either b does not appear in $\varphi(b)$ or appears only at the final position). In this case $b \notin \mathcal{L}(X_b)$ and for all $a \in \mathcal{A}_b$ different from b we have $a < b$. Applying Proposition 2.6 to the system X_b , we see that either $y \in X_a$ for some $a < b$ or y is a suffix of ${}^\omega\varphi(a)\varphi^\omega(c)$ for some backwards prolongable letter a and prolongable letter c such that $ac \in \mathcal{L}(X_b)$. The first case implies that $X_\varphi = X_a$ and contradicts the choice of b . In the second case $ac \in \mathcal{L}(X_b)$ implies that all suffixes of ${}^\omega\varphi(a)\varphi^\omega(c)$ lie in X_b . Since the orbit of y is dense in X_b , for each $n \geq 1$ the sequence $\varphi^n(a)\varphi^\omega(c)$ has arbitrarily long prefixes in common with some forward shift of y . Since $\varphi^n(a)$ is a suffix of $\varphi^{n+1}(a)$ for each $n \geq 0$ and y has the form $y = u\varphi^\omega(c)$ for some finite word u , all $\varphi^n(a)\varphi^\omega(c)$ are in fact factors of $\varphi^\omega(c)$. Letting n tend to infinity, we conclude that all suffixes of ${}^\omega\varphi(a)\varphi^\omega(c)$ lie in X_c . In particular, $y \in X_c$, which again contradicts the choice of b . This ends the proof. \square

Proof of Theorem 2.1. The claim for transitive subsystems of systems of the form $X = X_\varphi$ for an idempotent substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ follows immediately from Corollary 2.7 and Proposition 2.8.

In general, if Y is a transitive subsystem of a substitutive system X , we consider X as a topological factor $X = \pi(X_\varphi)$ of some X_φ for an idempotent substitution φ and a coding π . Choose $y \in Y$ such that $Y = \overline{\mathcal{O}(y)}$ and let $z \in X_\varphi$ be such that $\pi(z) = y$. Put $Z = \overline{\mathcal{O}(z)}$. By compactness we have $\pi(Z) = Y$. Since Z is a transitive subsystem of X_φ , it is substitutive, and hence so is the system $Y = \pi(Z)$. A similar argument proves the claim concerning k -automatic systems. \square

Two-sided substitutive shifts and their subsystems. We close this section with the remark that the results formulated above have their analogues for two-sided shifts. For a substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ we let $X_\varphi^{\mathbf{Z}}$ denote the two-sided dynamical system generated by φ , i.e.

$$X_\varphi^{\mathbf{Z}} = \{z \in {}^\omega\mathcal{A}^\omega \mid \text{every factor of } z \text{ appears in } \varphi^n(a) \text{ for some } n \geq 0 \text{ and } a \in \mathcal{A}\}.$$

For a letter a , the system $X_a^{\mathbf{Z}}$ is defined accordingly. A sequence $y = (y_n)_n \in {}^\omega\mathcal{A}^\omega$ is *substitutive* if both $(y_n)_{n \geq 0}$ and $(y_n)_{n < 0}$ are substitutive as one-sided sequences. This is obviously the same as saying that all (one-sided) prefixes and suffixes of y are substitutive. Let T denote the shift map on ${}^\omega\mathcal{A}^\omega$. For two-sided systems we consider the two-sided orbit $\mathcal{O}^{\mathbf{Z}}(y) = \{T^n(y) \mid n \in \mathbf{Z}\}$ of a point y . A *two-sided substitutive system* is the (two-sided) orbit closure of a two-sided substitutive sequence. We define a two-sided k -automatic sequence and a two-sided k -automatic system in the same way. The main results for two-sided shifts are the same or simpler as for the one-sided ones. The proofs are mutatis mutandis the same, and we present them in a briefer manner. The most notable difference between two-sided and one-sided shifts is that in the two-sided case every substitutive system has only finitely many subsystems.

Theorem 2.9. *Every transitive subsystem of a two-sided substitutive system is substitutive. Every transitive subsystem of a two-sided k -automatic system is k -automatic.*

To prove this result, we first state three lemmas, which are analogous to the previously described results for one-sided systems.

Lemma 2.10. *Every two-sided substitutive system X arises as the image $X = \pi(X_\varphi^{\mathbf{Z}})$ of a transitive system $X_\varphi^{\mathbf{Z}}$ generated by a substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ via a coding $\pi: \mathcal{A} \rightarrow \mathcal{B}$. If X is k -automatic, we may choose φ to be of constant length k .*

Proof. Let X be a two-sided substitutive system arising as the orbit closure of a sequence $y = (y_n)_n$. Since $(y_n)_{n \geq 0}$ and $(y_n)_{n < 0}$ are one-sided substitutive, we may find substitutions $\varphi_1: \mathcal{A}_1 \rightarrow \mathcal{A}_1^*$ and $\varphi_2: \mathcal{A}_2 \rightarrow \mathcal{A}_2^*$, codings $\pi_1: \mathcal{A}_1 \rightarrow \mathcal{B}$ and $\pi_2: \mathcal{A}_2 \rightarrow \mathcal{B}$, a prolongable letter $a_1 \in \mathcal{A}_1$ and a backwards

prolongable letter $a_2 \in \mathcal{A}_2$ such that $y = \pi_2(\omega\varphi_2(a_2))\pi_1(\varphi_1^\omega(a_1))$. We may assume that \mathcal{A}_1 and \mathcal{A}_2 are disjoint. Define a new alphabet $\mathcal{A} = \mathcal{A}_1 \cup \mathcal{A}_2 \cup \{\spadesuit\}$ with a new symbol $\spadesuit \notin \mathcal{A}_1 \cup \mathcal{A}_2$. Glue φ_i and π_i to maps $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ and $\pi: \mathcal{A} \rightarrow \mathcal{B}$ by putting $\varphi|_{\mathcal{A}_i} = \varphi_i$, $\pi|_{\mathcal{A}_i} = \pi_i$, and $\varphi(\spadesuit) = a_2w$, where w is a prefix of $\varphi^\omega(a_1)$ chosen to be of arbitrary length in the substitutive case and of length $k-1$ in the k -automatic case. It is easy to see that $X_\varphi^{\mathbf{Z}}$ is a transitive system generated by the sequence $\omega\varphi_2(a_2)\varphi_1^\omega(a_1)$ and that $X = \pi(X_\varphi^{\mathbf{Z}})$. \square

The remaining two lemmas are two-sided analogues of Lemmas 1.4 and 1.5. The former of these lemmas is proven in [7, Prop. 5.10].

Lemma 2.11. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution and let $x \in X_\varphi^{\mathbf{Z}}$. There exists $y \in X_\varphi$ such that $x = T^l(\varphi(y))$ for some l with $0 \leq l < |\varphi(y_0)|$.*

Lemma 2.12. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be a substitution. If $X_\varphi^{\mathbf{Z}}$ is transitive, then $X_\varphi^{\mathbf{Z}} = X_{\varphi^n}^{\mathbf{Z}}$ for any $n \geq 1$.*

The proofs of these lemmas are analogous to those of Lemmas 1.4 and 1.5. The main result is derived via essentially the same reasoning as before from Proposition 2.13 and 2.15 below.

Proposition 2.13. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be an idempotent substitution. Let $y \in X_\varphi^{\mathbf{Z}}$ and let*

$$Y = \{z \in X_\varphi^{\mathbf{Z}} \mid \text{every factor of } z \text{ appears in } y\}$$

be the orbit closure of y . Then at least one of the following conditions holds:

- (i) *there exists $a \in \mathcal{L}(Y)$ such that $Y = X_a^{\mathbf{Z}}$;*
- (ii) *there exist a backwards prolongable letter a and a prolongable letter c such that $ac \in \mathcal{L}(y)$ and y is a shift of $\omega\varphi(a)\varphi^\omega(c)$.*

In particular, the number of subsystems of $X_\varphi^{\mathbf{Z}}$ is finite.

Proof. Let $y^0 = y$. Repeating the reasoning from the proof of Proposition 2.6, we construct letters c_i , sequences $y^i \in X_\varphi^{\mathbf{Z}}$ with initial letters c_i and integers $0 \leq l_i < |\varphi(c_{i+1})|$ such that

$$y^i = T^{l_i}\varphi(y^{i+1})$$

for each $i \geq 0$. We consider two cases depending on the asymptotic behaviour of l_i .

Assume first that $l_i = 0$ for all sufficiently large i and put $a_i = y_{-1}^i$. For sufficiently large i , c_i is the initial letter of $\varphi(c_{i+1})$ and a_i is the final letter of $\varphi(a_{i+1})$. Since φ is idempotent, it follows from properties (iii) and (iv) in Definition 1.7 that a_i and c_i are eventually constant, say $a_i = a$ and $c_i = c$ for i sufficiently large. It follows that c is prolongable, a is backwards prolongable, and y is a shift of $\omega\varphi(a)\varphi^\omega(c)$. Similarly, if we assume that $l_i = |\varphi(c_{i+1})| - 1$ for all sufficiently large i , then we may apply the same reasoning with $T(y^i)$ in place of y^i .

Now assume that $l_i \neq 0$ and $l_i \neq |\varphi(c_{i+1})| - 1$ for infinitely many i 's. Let a be a letter that occurs infinitely many times among c_i . Then $a \in \mathcal{L}(Y)$ and we can find arbitrarily large j such that $\varphi^j(a) = y_{[n_j, m_j]}$ where $n_j \rightarrow -\infty$ and $m_j \rightarrow +\infty$ as $j \rightarrow \infty$. It follows that $Y = X_a^{\mathbf{Z}}$.

It follows immediately from our claim that $X_\varphi^{\mathbf{Z}}$ has only finitely many transitive subsystems, and hence finitely many subsystems. \square

Remark 2.14. Two-sided substitutive systems have been considered by Maloney and Rust, mostly under different assumptions on the substitution, namely that it is recognisable and tame (for the definition, see [31, Definition 2.4]). Note that all growing substitutions are tame, but not all growing substitutions are recognisable. Under these assumptions, the finiteness of the number of subsystems of $X_\varphi^{\mathbf{Z}}$ follows from [31, Lemma 5.13]. The authors work with the tiling space Ω_φ (see [31, Section 1.3]) associated with a substitution φ and prove that the number of closed unions of path components of Ω_φ is finite. Since there is a bijective correspondence between subsystems of $X_\varphi^{\mathbf{Z}}$ and closed unions

of path components of Ω_φ , the claim follows. Finiteness of the number of minimal subsystems of two-sided substitutive systems if the substitution is either aperiodic or growing has also been observed by Bezuglyi–Kwiatkowski–Medynets [7, Prop. 5.6 and Remark 5.7].

For simplicity we only state the following result for idempotent substitutions, but the more general analogue can readily be derived in the same way as in Proposition 2.8.

Proposition 2.15. *Let $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ be an idempotent substitution. The following conditions are equivalent:*

- (i) $X_\varphi^{\mathbf{Z}}$ is transitive;
- (ii) one of the following conditions holds:
 - (a) $X_\varphi^{\mathbf{Z}} = X_b^{\mathbf{Z}}$ for some letter $b \in \mathcal{A}$ that is either very ample or ample but neither prolongable nor backwards prolongable;
 - (b) $X_\varphi^{\mathbf{Z}} = X_a^{\mathbf{Z}} \cup X_c^{\mathbf{Z}} \cup \mathcal{O}^{\mathbf{Z}}(\omega_\varphi(a)\varphi^\omega(c))$ for some backwards prolongable letter a and prolongable letter c .

Moreover, if $X_\varphi^{\mathbf{Z}}$ is transitive, then it is substitutive. Furthermore, if φ is a substitution of constant length k , then $X_\varphi^{\mathbf{Z}}$ is k -automatic.

Proof. Suppose (ii) holds. Consider first the case (ii).(a). If $\varphi(b) = vbw$ with $v, w \in \mathcal{A}^*$ nonempty, then $X_b^{\mathbf{Z}}$ is the orbit closure of the point

$$y = \cdots \varphi^2(v)\varphi(v)vbw\varphi(w)\varphi^2(w)\cdots,$$

which is substitutive by Lemma 1.2. Otherwise, by (ii).(a) we have $\varphi(b) = bvb$ for some $v \in \mathcal{A}^*$, and $\varphi^2(b) = bvb\varphi(v)bvb$. Thus, the previous property is satisfied for φ^2 and the claim follows from the equality $X_\varphi^{\mathbf{Z}} = X_{\varphi^2}^{\mathbf{Z}}$, which holds since $\varphi^n(b)$ is a prefix of $\varphi^{n+1}(b)$ for all $n \geq 0$. Finally, in case (ii).(b) the system $X_\varphi^{\mathbf{Z}}$ is the orbit closure of $y = \omega_\varphi(a)\varphi^\omega(c)$.

It remains to show that (i) implies (ii). Assume that $X_\varphi^{\mathbf{Z}}$ is transitive and let y be a point in $X_\varphi^{\mathbf{Z}}$ with a dense orbit. Since $X_\varphi^{\mathbf{Z}} = \bigcup_{b \in \mathcal{A}} X_b^{\mathbf{Z}}$, there exists $b \in \mathcal{A}$ such that $X_\varphi^{\mathbf{Z}} = X_b^{\mathbf{Z}}$; pick minimal b with this property. Suppose that b does not satisfy the conditions in (ii).(a), and so b appears at most once in $\varphi(b)$, either at the initial or final position. Then $b \notin \mathcal{L}(X_b^{\mathbf{Z}})$ and $a < b$ for all $a \in \mathcal{A} \setminus \{b\}$. In particular, $X_b^{\mathbf{Z}} \neq X_a^{\mathbf{Z}}$ for all $a \in \mathcal{L}(X_b^{\mathbf{Z}})$, so Proposition 2.13 implies that there exists $ac \in \mathcal{L}(X_b^{\mathbf{Z}})$ such that y is up to a shift equal to $\omega_\varphi(a)\varphi^\omega(c)$. Hence, case (ii).(b) holds. \square

Proof of Theorem 2.9. The proof is analogous to the proof of Theorem 2.1, replacing the use of Propositions 2.6 and 2.8 by Propositions 2.13 and 2.15. \square

3. Finitary version of Cobham's theorem

In this section, we prove a finitary generalisation of Cobham's classical theorem. Recall that integers $k, l \geq 2$ are *multiplicatively independent* if they are not both powers of the same integer; equivalently, $\log k / \log l \in \mathbf{R} \setminus \mathbf{Q}$. Let $k, l \geq 2$ be multiplicatively independent integers. The celebrated theorem of Cobham classifies sequences that are simultaneously k -automatic and l -automatic: these are precisely the sequences that are ultimately periodic. The following theorem provides a complete characterisation of sets of words that can occur as common factors of automatic sequences defined over multiplicatively independent bases.

Theorem 3.1. *Let $k, l \geq 2$ be multiplicatively independent integers, let \mathcal{A} be an alphabet, and let $U \subset \mathcal{A}^*$. The following conditions are equivalent:*

- (i) *there exist a k -automatic sequence x and an l -automatic sequence y such that U is the set of common factors of x and y ;*

(ii) the set U is a finite union of sets of the form $\mathcal{L}({}^\omega v u w {}^\omega)$, where u, v, w are (possibly empty) words over \mathcal{A} .

In the proof, we will use the following strengthening of Cobham's theorem due to Isabelle Fagnot [22]. Note that, conversely, Cobham's theorem (or indeed its generalisation by Fagnot) follows immediately from Theorem 3.1.

Theorem 3.2 (Fagnot). *Let $k, l \geq 2$ be multiplicatively independent integers. Let x be a k -automatic sequence and let y be an l -automatic sequence. If $\mathcal{L}(x) = \mathcal{L}(y)$, then both sequences x and y are ultimately periodic.*

From the result above and Theorem 2.1 we get the following corollary.

Corollary 3.3. *Let $k, l \geq 2$ be multiplicatively independent integers and let \mathcal{A} be an alphabet. Let $X, Y \subset \mathcal{A}^\omega$ be subsystems such that X is k -automatic and Y is l -automatic. If a sequence z belongs to both X and Y , then it is ultimately periodic.*

Proof. Let Z be the orbit closure of z . Then Z is a transitive subsystem of both systems X and Y . By Theorem 2.1, there exist a k -automatic sequence $x \in X$ and an l -automatic sequence $y \in Y$ such that Z is the orbit closure of x and y , and hence the sequences x and y have the same language. By Theorem 3.2, the system Z is finite, and hence the sequence z is ultimately periodic. \square

The problem of describing common factors of automatic sequences was considered in [32]. The authors obtained, among other things, an upper bound on the length of a common prefix of aperiodic automatic sequences defined over multiplicatively independent bases in terms of the number of states of the automata generating the sequences. They further asked about the structure of the set of common factors of automatic sequences defined over multiplicatively independent bases. Since every ultimately periodic sequence is k -automatic for all integers $k \geq 2$, it is clear that we can get common factors of the form vu^n for some words v, u and arbitrarily large n . The following example shows that common factors can be somewhat more complicated.

Example 3.4. Let $\mathcal{A} = \{0, 1, 2\}$. Consider the 3-automatic sequence $x = \varphi^\omega(0)$ produced by the substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ given by

$$\varphi(0) = 012, \varphi(1) = 111, \varphi(2) = 222$$

and the 4-automatic sequence $y = \tau^\omega(0)$ produced by the substitution $\tau: \mathcal{A} \rightarrow \mathcal{A}^*$ given by

$$\tau(0) = 0121, \tau(1) = 1111, \tau(2) = 2222.$$

Then

$$x = 0121^3 2^3 1^9 2^9 1^{27} 2^{27} \dots$$

and hence

$$X_\varphi = \mathcal{O}(x) \cup \{2^n 1^\omega \mid n \geq 0\} \cup \{1^n 2^\omega \mid n \geq 0\}.$$

Similarly,

$$y = 0121^5 2^4 1^{20} 2^{16} 1^{80} 2^{64} \dots$$

and hence

$$X_\tau = \mathcal{O}(y) \cup \{2^n 1^\omega \mid n \geq 0\} \cup \{1^n 2^\omega \mid n \geq 0\}.$$

The common factors of x and y are exactly the words in $\mathcal{L}({}^\omega 12^\omega) \cup \mathcal{L}({}^\omega 21^\omega) \cup \mathcal{L}(0121^3)$.

We will use Corollary 3.3 to show that common factors of automatic sequences defined over multiplicatively independent bases are all of the form suggested by the example above. We need to introduce some additional terminology. Let \mathcal{A} be an alphabet and let x be a sequence over \mathcal{A} . Let X denote the orbit closure of x . We say that a factor u of x is *cyclic* if u is nonempty and the periodic sequence u^ω lies in X . We say that u is *primitive* if it is not of the form $u = v^n$ for

some $v \in \mathcal{A}^*$ and $n \geq 2$. Since the orbit closure of a periodic sequence is minimal, it follows from Corollary 2.3 that the set of primitive cyclic factors of a substitutive sequence is finite. We say that a common factor of sequences $x, y \in \mathcal{A}^\omega$ is *cyclic* if it is cyclic as a factor of x and as a factor of y .

Remark 3.5. We cannot hope for a straightforward generalisation of Theorem 3.1 to the class of substitutive sequences. Recall that with every substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ we can associate a matrix $M = (m_{ij})_{i,j \in \mathcal{A}}$, where m_{ij} is the number of occurrences of the letter i in the word $\varphi(j)$. By the Frobenius–Perron theorem M always has a dominant eigenvalue $\lambda > 0$. The eigenvalue λ plays the role of a base for a substitutive sequence x produced by φ (for details see, e.g. [18]), which allows us to define the class of λ -substitutive systems. The reason why Theorem C fails in this setting is that transitive subsystems of λ -substitutive systems need not be λ -substitutive. Consider the following example. Let $\mathcal{A} = \{0, 1\}$ and $\mathcal{B} = \{0, 1, 2, 3\}$, let $x = \varphi^\omega(0)$ be the sequence produced by the substitution $\varphi: \mathcal{A} \rightarrow \mathcal{A}^*$ given by

$$\varphi(0) = 01, \quad \varphi(1) = 10,$$

and let $y = \tau^\omega(2)$ be the sequence produced by the substitution $\tau: \mathcal{B} \rightarrow \mathcal{B}^*$ given by

$$\tau(0) = 01, \quad \tau(1) = 10, \quad \tau(2) = 203, \quad \tau(3) = 3233.$$

The systems $X = \overline{\mathcal{O}(x)}$ and $Y = \overline{\mathcal{O}(y)}$ are 2-substitutive and $(2 + \sqrt{2})$ -substitutive, respectively, and the set of common factors of x and y consists precisely of the factors of x .

Occurrences of cyclic factors in automatic sequences. Let x be a k -automatic sequence. To proceed with the proof of Theorem 3.1, we first need to understand the structure of sets of the form

$$S_x = \{n \geq 0 \mid vu^n w \text{ is a factor of } x\}$$

for fixed words v, w and u . This is only interesting if u is a cyclic factor of x , since otherwise S_x is finite. If v is a suffix of some power of u or w is a prefix of some power of u , then the set S_x is easy to determine, and either consists of all integers, or is finite and consists of all integers smaller than some constant. Assume conversely that v is not a suffix of any power of u and w is not a prefix of any power of u . We will show that the set S_x is up to a finite set a finite union of translates of geometric progressions, and deduce that for automatic sequences x and y defined over multiplicatively independent bases the set $S_x \cap S_y$ is finite.

The problem above was also considered by Fagnot in [22, Proposition 8] in the special case when the sequence x takes values in $\{0, 1\}$ and $v = w = 1, u = 0$. This result was used to show that if x is a k -automatic sequence, y is an l -automatic sequence, $k, l \geq 2$ are multiplicatively independent, and $\mathcal{L}(x) \subset \mathcal{L}(y)$, then either x contains only finitely many 1's or 1's occur in x with bounded gaps (see [22, Corollaire 10]). This is the crucial step in her proof of Theorem 3.2. It would be interesting to see if the general statement below could be reduced to the special case considered by Fagnot, but we found no such reduction. We give a proof of the general result that uses similar ideas as the one of Fagnot but seems quite different in details, and we deduce a much stronger finiteness result in Corollary 3.8. We also discuss the question of effectiveness.

Proposition 3.6. *Let $k \geq 2$ be an integer. Let x be a k -automatic sequence over an alphabet \mathcal{A} . Let u, v, w be nonempty words over \mathcal{A} . Assume that v is not a suffix of u^n and w is not a prefix of u^n for any integer n . Let $S = \{n \geq 0 \mid vu^n w \text{ is a factor of } x\}$. The set S is a finite union of sets of the form $\{ak^{mn} + b \mid n \geq 0\}$ for some $a, b \in \mathbf{Q}$ and $m \in \mathbf{Z}$ with $a, b, m \geq 0, a + b \in \mathbf{Z}$ and $(k^m - 1)a \in \mathbf{Z}$.*

Proof. We begin the proof with a few reductions.

Step I (reduction to the case when $|u| = |v| = |w|$). First, we show that it is enough to prove the claim under the additional assumption that $|v| \leq |u|$ and $|w| \leq |u|$. Let $j > 0$ be an integer

such that $\max(|v|, |w|) \leq |u|^{j/2}$. Write

$$S_i = \{n \in S \mid n \equiv i \pmod{j}\}.$$

Then $S = \bigcup_{i=0}^{j-1} S_i$, and hence it is enough to show the claim separately for each of the sets S_i . For $0 \leq i < j$, we put $u' = u^j$, $v' = vu^{\lfloor (i+1)/2 \rfloor}$ and $w' = u^{\lfloor i/2 \rfloor}w$. We then have $\max(|v'|, |w'|) \leq |u'|$ and

$$S_i = \{n = jm + i \mid v'(u')^m w' \text{ is a factor of } x\}.$$

In order to further obtain $|v'| = |w'| = |u'|$, we consider all possible prolongations of the words v' and w' to words v'' and w'' of length $|u'|$ and such that v' is a suffix of v'' and w' is a prefix of w'' . Every element of S_i lies in one of the sets

$$\{n = jm + i \mid v''(u')^m w'' \text{ is a factor of } x\}$$

for some choice of v'' and w'' , except for the values $n \in S_i$ corresponding to factors $v'(u')^m w'$ which occur only at starting positions $< |u'| - |v'|$. Since there are only finitely many such values of m (at most one for each starting position), we may assume that $|u| = |v| = |w|$.

Step II (reduction to the case when $|u| = |v| = |w| = 1$). Write ℓ for the common length of u , v and w . We will now show that we may assume that $\ell = 1$ by changing the alphabet. For $0 \leq i < \ell$ let

$$\tilde{S}_i = \{n \in S \mid vu^n w = x_{[m, m+(n+2)\ell]} \text{ for some } m \equiv i \pmod{\ell}\}$$

denote the set of all integers $n \in S$ such that the factor $vu^n w$ occurs in x at a position $m \equiv i \pmod{\ell}$. Clearly, S is the union of the sets \tilde{S}_i .

Let \mathcal{A}^ℓ denote the set of words of length ℓ over \mathcal{A} . Identifying words u, v, w with letters $u', v', w' \in \mathcal{A}^\ell$ and the sequence x with the corresponding sequence $x' \in (\mathcal{A}^\ell)^\omega$, we see that the set \tilde{S}_0 is equal to the set of all integers n such that $v'(u')^n w'$ is a factor of x' . The same reasoning applied to the sequence $T^i(x)$ instead of x shows that the set \tilde{S}_i is equal to the set of all integers n such that $v'(u')^n w'$ is a factor of the k -automatic sequence $(T^i(x))'$. This allows us to assume that u, v, w are single letters.

Step III (restating the problem in terms of purely automatic sequences). Write the sequence x as the image of a purely k -automatic sequence y produced by a substitution $\varphi: \mathcal{B} \rightarrow \mathcal{B}^*$ of constant length k under a coding $\pi: \mathcal{B} \rightarrow \mathcal{A}$. Let T, C and D denote the preimages of the letters u, v and w under the coding π , respectively. Note that $C \cap T = D \cap T = \emptyset$. The set S can be expressed in terms of the sequence y as

$$S = S(C, D, T) = \{n \geq 0 \mid cwd \text{ is a factor of } y \text{ for some } c \in C, d \in D \text{ and } w \in T^* \text{ with } |w| = n\}.$$

We will prove that sets $S = S(C, D, T)$ satisfy the claim for all purely k -automatic sequences y over an alphabet \mathcal{B} and subsets $T, C, D \subset \mathcal{B}^*$ with $C \cap T = D \cap T = \emptyset$. Dividing S into a finite union, we may further assume that the sets C and D consist of single letters $c, d \in \mathcal{B} \setminus T$, and we write $S(c, d, T)$ for $S(\{c\}, \{d\}, T)$.

Step IV (constructing a recurrence for the set S). For $m \geq 1$ we let $\varphi^{-m}(T)$ denote the set of letters $a \in \mathcal{A}$ such that $\varphi^m(a) \in T^*$. We write $\mathcal{A}_T = \mathcal{A} \setminus \varphi^{-1}(T)$ for the set of letters $a \in \mathcal{A}$ such that $\varphi(a) \notin T^*$. For a letter $a \in \mathcal{A}_T$ we denote by $\alpha(a)$ the first letter in $\varphi(a)$ that is not in T and by $\omega(a)$ the last letter in $\varphi(a)$ that is not in T . We replace the substitution φ by its power in order to get the property $\varphi^{-1}(T) = \varphi^{-2}(T)$ (this is possible by Lemma 1.8).

With every pair $(a, a') \in \mathcal{A}_T^2$ we associate an integer $q(a, a')$ in the following way. Write $\varphi(a) = v\omega(a)w$ and $\varphi(a') = w'\alpha(a')v'$ for some $w, w' \in T^*$ and $v, v' \in \mathcal{A}^*$. Let q denote the length of w and let q' denote the length of w' . Put $q(a, a') = |w| + |w'|$. Consider the sets

$$\Omega_c = \{a \in \mathcal{A}_T \mid \omega(a) = c\} \quad \text{and} \quad A_d = \{a \in \mathcal{A}_T \mid \alpha(a) = d\}.$$

We let $E \subset S$ denote the set of integers $n \in S$ such that $n < k - 1$. We claim that

$$S(c, d, T) = \bigcup_{(c', d') \in \Omega_c \times A_d} (kS(c', d', \varphi^{-1}(T)) + q(c', d')) \cup E. \quad (2)$$

Let $n \geq k - 1$. By definition, n lies in S if and only if there exists a word $w \in T^*$ with $|w| = n$ such that cwd is a factor of y . Since $\varphi(y) = y$ and $|cwd| \geq k + 1$, this happens if and only if there exist a pair $(c', d') \in \Omega_c \times A_d$ and a word $w' \in (\varphi^{-1}(T))^*$ such that cwd is a factor of $\varphi(c'w'd')$. It follows that $n \in S$ if and only if $n \in kS(c', d', \varphi^{-1}(T)) + q(c', d')$ for some $(c', d') \in \Omega_c \times A_d$, with the ‘if’ claim not requiring the assumption that $n \geq k - 1$. This proves (2).

Observe that (2) implies that it is enough to prove the claim for each of the sets $S(c', d', \varphi^{-1}(T))$. By the assumption on φ , we have that $\varphi^{-1}(\varphi^{-1}(T)) = \varphi^{-1}(T)$, and hence it is enough to show the claim for sets of the form $S = S(c, d, T)$ under the additional assumption that $\varphi^{-1}(T) = T$. It follows that $\alpha(\mathcal{A}_T) \subset \mathcal{A}_T$ and $\omega(\mathcal{A}_T) \subset \mathcal{A}_T$. Writing temporarily α_φ and ω_φ for the maps $\alpha, \omega: \mathcal{A}_T \rightarrow \mathcal{A}_T$ defined with respect to the substitution φ , we note that $\alpha_\varphi^n = \alpha_{\varphi^n}$ and $\omega_\varphi^n = \omega_{\varphi^n}$ for all $n \geq 1$. Another application of Lemma 1.8 shows that after replacing φ by an appropriate power we get that the maps α and ω are idempotent, which we henceforth assume.

If the set S is finite, the claim is obvious, so assume that S is infinite. By (2), the sets Ω_c and A_d are nonempty, and hence (since ω and α are idempotent) we have $\omega(c) = c$, $\alpha(d) = d$.

We now consider a family of recurrence sequences. For an element $r \in S$, consider the sequence $(n_t^r)_{t \geq 0}$ given by the formula

$$\begin{aligned} n_0^r &= r, \\ n_t^r &= kn_{t-1}^r + q(c, d), \quad t \geq 1. \end{aligned}$$

By (2) it is clear that $n_t^r \in S$ for all $r \in S$ and $t \geq 0$. We claim that

$$S = \{n_t^r \mid t \geq 0, r \in S, 0 \leq r < k^2 - 1\}. \quad (3)$$

This will end the proof of the claim, since the recurrence sequence $(n_t^r)_{t \geq 0}$ has the closed form $n_t^r = ak^t + b$ with $a = r + q(c, d)/(k - 1)$ and $b = -q(c, d)/(k - 1)$ satisfying $a + b \in \mathbf{Z}$ and $(k - 1)a \in \mathbf{Z}$. (Note that in the process of the proof we have replaced the substitution φ by its iterate, which has the effect of replacing the original k by its power.)

Step V (proving the formula (3)). We have already remarked that all n_t^r are elements of S . For the converse claim (with the extra statement that one can take $r < k^2 - 1$), we will inductively apply (2), which takes a simpler form since $\varphi^{-1}(T) = T$.

Choose $m \in S$ and write $c_0 = c$, $d_0 = d$, $m_0 = m$. If $m \geq k - 1$, then by (2) we may write $m_0 = km_1 + q(c_1, d_1)$ for some $c_1, d_1 \in \mathcal{A}_T$ with $\omega(c_1) = c$, $\alpha(d_1) = d$ and $m_1 \in S(c_1, d_1, T)$. If $m_1 \geq k - 1$, we may repeat this procedure. In this way, we inductively construct sequences $(c_i)_{0 \leq i \leq s}$, $(d_i)_{0 \leq i \leq s}$ and $(m_i)_{0 \leq i \leq s}$ with $c_i, d_i \in \mathcal{A}_T$, $\omega(c_{i+1}) = c_i$, $\alpha(d_{i+1}) = d_i$ and $m_i \in S(c_i, d_i, T)$ with $m_i = km_{i+1} + q(c_{i+1}, d_{i+1})$. Furthermore, we have $m_i \geq k - 1$ for $i < s$ and $m_i < k - 1$ for $i = s$.

Since the maps ω and α are idempotent, the conditions on (c_i) and (d_i) imply that $c_i = c$ and $d_i = d$ for $0 \leq i < s$ (but not necessarily for $i = s$). This shows that $m_i = km_{i+1} + q(c, d)$ for $0 \leq i < s$, and hence $m = n_t^r$ for $t = s - 1$ and $r = m_{s-1}$. Since $m_{s-1} = km_s + q(c_s, d_s)$, $m_s \leq k - 2$ and $q(c_s, d_s) \leq 2k - 2$, we get that $r = m_{s-1} < k^2 - 1$, which ends the proof of the claim. \square

Remark 3.7. Note that the proof of Proposition 3.6 is effective, in the sense that given a k -automatic sequence and words u, v, w satisfying the conditions of the proposition, one may explicitly determine the set $S = \{n \geq 0 \mid vu^n w \text{ is a factor of } x\}$ as a finite union of translates of geometric progressions and a finite set. In fact, even if u, v, w fail to satisfy the assumptions, i.e. either v is a suffix of some power of u or w is a prefix of some power of u , we may still determine S as either an explicit finite set or as all of \mathbf{N} . Indeed, the reasoning in Steps I and II in the proof of Proposition

3.6 allows us to assume that u, v and w are single letters. Fix u and vary v and w . If $v, w \neq u$, we already know how to find the corresponding set S . The remaining cases reduce to this one since every factor $vu^n w$ of x is either a subfactor of some factor $\tilde{v}u^m\tilde{w}$ of x for some letters \tilde{v}, \tilde{w} with $\tilde{v}, \tilde{w} \neq u$ or else arises as a factor of some prefix of x of the form $u^m w$ or finally we have that x is ultimately periodic with suffix u^ω , in which case it is easy to find S .

Corollary 3.8. *Let $k, l \geq 2$ be multiplicatively independent integers and let \mathcal{A} be an alphabet. Let x be a k -automatic sequence over \mathcal{A} and let y be an l -automatic sequence over \mathcal{A} . Let u, v, w be words over \mathcal{A} . Assume that v is not a suffix of u^n and w is not a prefix of u^n for any integer n . Then the word $vu^n w$ is a common factor of x and y only for finitely many n .*

Proof. Let $S_x = \{n \in \mathbf{N} \mid vu^n w \text{ is a factor of } x\}$ and $S_y = \{n \in \mathbf{N} \mid vu^n w \text{ is a factor of } y\}$. By Proposition 3.6, S_x is a finite union of sets of the form $\{ak^{mn} + b \mid n \geq 0\}$ for some $a, b \in \mathbf{Q}$ and $m \geq 0$. Similarly, S_y is a finite union of sets of the form $\{al^{mn} + b \mid n \geq 0\}$ for some $a, b \in \mathbf{Q}$ and $m \geq 0$. In order to prove that the set $S_x \cap S_y$ is finite, it suffices to note that for any multiplicatively independent integers k, l and rational numbers $a, b, c \in \mathbf{Q}$ with a, b, c not all equal to zero the exponential diophantine equation

$$ak^n + bl^m = c$$

has only finitely many integer solutions $m, n \in \mathbf{N}$. This follows, e.g. from the finiteness of the number of solutions of S -unit equations due to Mahler [30] (see also [21, Ch. 4] or [27, p. 28] for a more general, but very convenient, statement). \square

Proof of the main result. We are now ready to prove Theorem 3.1. We begin with a lemma.

Lemma 3.9. *Let u, \tilde{u}, v be words over an alphabet \mathcal{A} , and let $n, m \geq 0$ be integers. Assume that u and \tilde{u} are primitive and that \tilde{u}^m is a suffix of $u^n v$ with $m|\tilde{u}| \geq |v| + |u| + |\tilde{u}| - \gcd(|u|, |\tilde{u}|)$. Then u and \tilde{u} are cyclic shifts of each other and for any $q \geq 0$ we have $u^n v \tilde{u}^q = u^{n+q} v$.*

Proof. Let \tilde{v} denote the word \tilde{u}^m with the suffix v removed. It follows from the Fine–Wilf theorem [5, Thm. 1.5.6] applied to the (backwards infinite) periodic sequences ${}^\omega u$ and ${}^\omega \tilde{u} \tilde{v}$ that ${}^\omega u = {}^\omega \tilde{u} \tilde{v}$, and hence ${}^\omega uv = {}^\omega \tilde{u}$. Since u and \tilde{u} are primitive, they are cyclic shifts of each other. Both the words $u^n v \tilde{u}^q$ and $u^{n+q} v$ are suffixes of ${}^\omega uv = {}^\omega \tilde{u}$ of the same length, and hence are equal. \square

Proof of Theorem 3.1. We first prove that (ii) implies (i). Write U in the form

$$U = \bigcup_{i=1}^p \mathcal{L}({}^\omega v_i u_i w_i^\omega), \quad u_i, v_i, w_i \in \mathcal{A}^*.$$

Replacing k and l with their powers, we may assume that $k, l \geq \max(p+2, 3)$. Choose two symbols \clubsuit, \spadesuit outside of \mathcal{A} . We will construct a k -automatic sequence $x \in (\mathcal{A} \cup \{\clubsuit\})^\omega$ whose set of factors not containing \clubsuit coincides with U . Consider the sequences

$$c^i = \begin{cases} u_i w_i^\omega & \text{if } w_i \neq \epsilon, \\ u_i \clubsuit^\omega & \text{if } w_i = \epsilon. \end{cases} \quad d^i = \begin{cases} {}^\omega v_i & \text{if } v_i \neq \epsilon, \\ {}^\omega \clubsuit & \text{if } v_i = \epsilon. \end{cases}$$

Since these sequences are ultimately periodic, they are both k - and l -automatic. Define the sequence $x \in (\mathcal{A} \cup \{\clubsuit\})^\omega$ by the formula

$$x_n = \begin{cases} c_n^i & \text{if } i \cdot k^t \leq n < i \cdot k^t + k^{t-1}, t \geq 1, 1 \leq i \leq p, \\ d_{n-i \cdot k^t}^i & \text{if } i \cdot k^t - k^{t-1} \leq n < i \cdot k^t, t \geq 1, 1 \leq i \leq p, \\ \clubsuit & \text{otherwise.} \end{cases}$$

(Recall that we always regard sequences in \mathcal{A}^ω as indexed with $0, 1, \dots$, and sequences in ${}^\omega \mathcal{A}$ as indexed with $\dots, -2, 1$.) Using either the characterisation of automaticity in terms of kernels or

in terms of finite automata, it is elementary (albeit tedious) to verify that x is k -automatic. It is immediate that the set of factors of x not containing the symbol \clubsuit coincides with U . Replacing k by l and \clubsuit by \spadesuit , we define an l -automatic sequence $y \in (\mathcal{A} \cup \{\spadesuit\})^\omega$ whose set of factors not containing \spadesuit coincides with U . It follows that the set of common factors of x and y is exactly U . This ends the proof that (ii) implies (i).

For the proof that (i) implies (ii), let x be a k -automatic sequence and let y be an l -automatic sequence. For simplicity, in the rest of the proof we will refer to common factors of x and y simply as common factors. It follows from Corollary 2.3 that there are only finitely many primitive cyclic common factors (in fact, both x and y have only finitely many primitive cyclic factors). Let ℓ denote the maximal length of such a factor. We write common factors t in the form

$$t = v_0 u_1^{n_1} v_1 u_2^{n_2} \cdots v_{s-1} u_s^{n_s} v_s \quad (4)$$

for some integer $s \geq 0$, integers $n_i \geq 0$ and words u_i, v_i satisfying the following properties:

- (i) the words u_i are primitive cyclic common factors,
- (ii) the words v_i have length $|v_i| \leq \ell$,
- (iii) the integer s , called the *rank* of the representation, is the smallest possible,
- (iv) given the choice of s , the sequence of integers (n_1, \dots, n_s) , called the *sequence of exponents*, is lexicographically maximal.

We will refer to (4) as a *representation* of a common factor t .

Note that if t' is a prefix of t , then the rank of t' is at most equal to the rank of t . We will prove that common factors have bounded rank. To this end, we first prove the following claim.

Claim 1: For each $i \in \{2, \dots, s-1\}$ if $n_i > 0$, then there exist

- (i) a suffix y of $v_0 u_1^{n_1} \cdots u_{i-1}^{n_{i-1}} v_{i-1}$ of length $|y| \leq 4\ell$ such that y is not a suffix of u_i^n for any integer $n \geq 0$; and
- (ii) a prefix z of $v_i u_{i+1}^{n_{i+1}} \cdots u_s^{n_s} v_s$ of length $|z| \leq 4\ell$ such that z is not a prefix of u_i^n for any integer $n \geq 0$.

Proof of Claim 1: We only prove (i), the proof of (ii) being analogous. Write

$$v_0 u_1^{n_1} \cdots u_{i-1}^{n_{i-1}} v_{i-1} = w' w,$$

where $w, w' \in \mathcal{A}^*$, w is a suffix of u_i^n for some $n \geq 0$, and w is chosen as long as possible. Let $m \geq 0$ be the largest integer such that u_i^m is a suffix of w , and write $w = w'' u_i^m$. Note that $|w''| < \ell$.

Consider the following cases (which cover all possibilities):

- (i) If $w' = \epsilon$, then t admits the representation $t = w'' u_i^{m+n_i} v_i \cdots v_{s-1} u_s^{n_s} v_s$, which is of rank $s - i + 1 < s$. This is a contradiction.
- (ii) If $|u_i^m| > |u_{i-1}^{n_{i-1}} v_{i-1}|$, we claim that t also admits a representation of smaller rank. In fact, the word $w' w''$ is a prefix of $v_0 u_1^{n_1} \cdots u_{i-2}^{n_{i-2}} v_{i-2}$, and hence has a representation of rank at most $i - 2$. Concatenating it with $u_i^{m+n_i} v_i \cdots v_{s-1} u_s^{n_s} v_s$, we obtain a representation of t of rank $\leq s - 1$. This is a contradiction.
- (iii) If $|u_i^m| \leq |u_{i-1}^{n_{i-1}} v_{i-1}|$ and $|w| \geq 4\ell$, then $|u_i^m| > 3\ell$, and hence by Lemma 3.9 we may write $u_{i-1}^{n_{i-1}} v_{i-1} u_i^{n_i} = u_{i-1}^r v_{i-1}$ with $r = n_{i-1} + n_i$. Replacing in the representation of t the word $u_{i-1}^{n_{i-1}} v_{i-1} u_i^{n_i} v_i$ by $u_{i-1}^r v_{i-1} u_i^0 v_i$, we obtain a representation whose sequence of exponents is $(n_1, \dots, n_{i-2}, r, 0, n_{i+1}, \dots, n_s)$, and hence (since $n_i > 0$) is lexicographically larger than $(n_1, \dots, n_{i-2}, n_{i-1}, n_i, n_{i+1}, \dots, n_s)$. This is a contradiction.
- (iv) If $|w| < 4\ell$ and $w' \neq \epsilon$, then the suffix y of $w' w$ of length $|w| + 1$ satisfies the claim. \square

Claim 2: There is a constant C such that for any common factor t the values of n_2, \dots, n_{s-1} in any representation (4) of t are bounded by C .

Proof of Claim 2: Since there are only finitely many primitive cyclic common factors and finitely many words of length $\leq 4\ell$, Corollary 3.8 and Claim 1 show that the values of n_2, \dots, n_{s-1} are bounded by a constant independent of t . \square

Claim 3: The rank of common factors is bounded.

Proof of Claim 3: Let t be a common factor with representation $t = v_0 u_1^{n_1} v_1 u_2^{n_2} \cdots v_{s-1} u_s^{n_s} v_s$. We claim that any cyclic factor of t (not necessarily primitive) can occur at positions intersecting at most four of the u_i 's. Suppose this is not the case and write such a factor in the form \tilde{u}^n for some integer $n \geq 0$ and primitive cyclic common factor \tilde{u} . Consider in the representation of t the shortest factor \tilde{w} consisting of a concatenation of u_i 's and v_i 's and containing \tilde{u}^n . Replacing \tilde{w} by a word of the form $v' \tilde{u}^n v''$ with $|v'|, |v''| \leq \ell$, we obtain a representation of smaller rank, which gives a contradiction.

Now suppose that the rank of common factors is unbounded. Then the words

$$w(t) := u_2^{n_2} v_2 \cdots v_{s-2} u_{s-1}^{n_{s-1}}$$

can be arbitrarily long, and hence by compactness of \mathcal{A}^ω there exists a sequence z with arbitrarily long prefixes of the form $w(t)$ for some common factors t . Let X and Y denote the orbit closures of x and y , respectively. Then $z \in X \cap Y$, and hence by Corollary 3.3 it is ultimately periodic. However, by Claim 2 any cyclic factor of $w(t)$ has length at most $(3C + 4)\ell$, which is a contradiction. \square

Claim 4: There exists a constant C' such that any common factor t can be written in the form $t = v' u^n v w^m v''$ for integers $n, m \geq 0$, primitive cyclic common factors u, w and words v, v', v'' of length at most C' .

Proof of Claim 4: Follows immediately from Claims 3 and 4. \square

Due to Claim 4, in order to prove our result, it is sufficient to study common factors t of the form $t = v' u^n v w^m v''$ for fixed words u, w, v, v', v'' . Call a set S of common factors special if it takes the form

$$S = \{t = v' u^n v w^m v'' \mid t \text{ is a common factor, } n, m \in \mathbf{N}\}$$

for some words $u, w, v, v', v'' \in \mathcal{A}^*$. If we may further take $w = v'' = \epsilon$, we call the set S degenerate. By Claim 4, the set of all common factors is a finite union of special sets. We write $\mathcal{L}(S)$ for the set of all factors of words in S . We will prove that for any special set S the set $\mathcal{L}(S)$ is a finite union of sets of the form $\mathcal{L}({}^\omega \tilde{u} \tilde{v} \tilde{w} {}^\omega)$ with $\tilde{u}, \tilde{v}, \tilde{w} \in \mathcal{A}^*$. This will conclude the proof of the theorem.

We first prove the claim for a degenerate special set

$$S = \{t = v' u^n v \mid t \text{ is a common factor, } n \in \mathbf{N}\}.$$

If the set S is finite, then it is certainly of the desired form. By Corollary 3.8 this is the case if v' is not a suffix of any power of u and v is not a prefix of any power of u . If on the other hand v' is a suffix of some power of u , and S is infinite, then $\mathcal{L}(S)$ is equal to $\mathcal{L}({}^\omega u v)$. A similar reasoning proves the claim if v is a prefix of some power of u .

Consider now the case of a general special set

$$S = \{t = v' u^n v w^m v'' \mid t \text{ is a common factor, } n, m \in \mathbf{N}\}.$$

If S does not contain factors of the form $t = v' u^n v w^m v''$ for arbitrarily large values of both n and m , then S can be rewritten as a finite union of degenerate special sets, and the claim follows. Suppose that S contains factors t corresponding to arbitrarily large values of both n and m . If either $v w^m$ is a prefix of some power of u for arbitrarily large m or $u^n v$ is a suffix of some power of w for arbitrarily large n , then we may again rewrite S as a finite union of degenerate special sets. Finally, if neither is $v w^m$ a prefix of some power of u for sufficiently large m nor is $u^n v$ a suffix of some power of w for sufficiently large n , then we conclude from Corollary 3.8 that v' is a suffix of some power of u and v'' is a prefix of some power of w . In this case the set $\mathcal{L}(S)$ is equal to $\mathcal{L}({}^\omega u v w {}^\omega)$, which finishes the proof. \square

Remark 3.10. It is an interesting question whether Theorem 3.1 can be made effective, i.e. whether there is an algorithm which, given a k -automatic sequence x and an l -automatic sequence y , produces words $u_i, v_i, w_i, 1 \leq i \leq p$, such that the set U of common factors of x and y is equal to

$$U = \bigcup_{i=1}^p \mathcal{L}(\omega v_i u_i w_i^\omega).$$

The only place in the proof of Theorem 3.1 where it is not clear if the proof is effective is the bound on the rank of common factors (Claim 3), which uses a compactness argument. Let us briefly comment on how to make other parts of the proof effective.

First of all, we can determine all primitive cyclic factors of an automatic sequence x . In fact, write x as the image of a fixed point of a substitution under a coding. Replace the substitution by an idempotent one using Lemma 1.9, which is effective. Proposition 2.2 together with Lemma 1.1 describe all minimal subsystems of $\overline{\mathcal{O}(x)}$ as closures of orbits of explicitly given automatic sequences. To find the cyclic factors of x , we need to determine which of these automatic sequences are periodic, for which a decision procedure was given by Honkala [24] (see also [4] for a simpler approach). Another crucial ingredient of the proof is Corollary 3.8, which uses the S -unit equation. Here, solutions can be effectively bounded using Baker’s method (for a comprehensive discussion, see [21]). In particular, the constant C in the proof can be effectively computed. Finally, given words v, v', v'', u, w we may effectively determine all common factors of x and y of the form $t = v'u^n v w^m v''$ using Remark 3.7 and an effective version of Corollary 3.8. Thus, in order to make the proof fully effective, we need to find a computable bound on the rank of common factors or—equivalently—the constant C' .

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