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Zero-Error Coding for Computing with Encoder Side-Information

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Abstract—We study the zero-error source coding problem in which an encoder with Side Information (SI) $g(Y)$ transmits source symbols X to a decoder. The decoder has SI Y and wants to recover $f(X, Y)$ where f, g are deterministic. We exhibit a condition on the source distribution and g that we call “pairwise shared side information”, such that the optimal rate has a single-letter expression. This condition is satisfied if every pair of source symbols “share” at least one SI symbol for all output of g . It has a practical interpretation, as Y models a request made by the encoder on an image X , and $g(Y)$ corresponds to the type of request. It also has a graph-theoretical interpretation: under “pairwise shared side information” the characteristic graph can be written as a disjoint union of OR products. In the case where the source distribution is full-support, we provide an analytic expression for the optimal rate. We develop an example under “pairwise shared side information”, and we show that the optimal coding scheme outperforms several strategies from the literature.

I. INTRODUCTION

Consider the source coding scenario depicted in Figure 1 where two correlated sequences (X^n, Y^n) of discrete i.i.d. random source symbols are drawn with a distribution $P_{X,Y}$. The encoder knows X^n , has $(g(Y_t))_{t \leq n}$ as side information, and transmits information to the decoder through a perfect channel of capacity R . The decoder has the side information Y^n and wants to reconstruct $(f(X_t, Y_t))_{t \leq n}$, where f, g are deterministic. What is the minimal R such that $(f(X_t, Y_t))_{t \leq n}$ can be retrieved by the decoder with probability of error 0?

This coding problem appears in video compression [1], [2], where X^n models a set of images known at the encoder. The decoder does not always want to retrieve each image, but has instead a sequence Y^n of particular requests for each image, e.g. detection: cat, dog, car, bike; or scene recognition: street/city/mountain, etc... The encoder does not know the

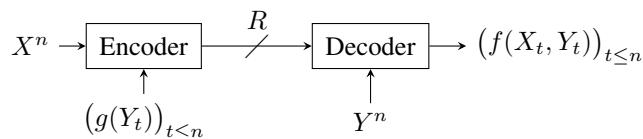


Fig. 1: Zero-error coding for computing with side information at the encoder.

decoder’s exact request but has prior information about it (e.g. type of request), which is modeled by $(g(Y_t))_{t \leq n}$.

The problem of Figure 1 relates to the “restricted inputs” zero-error problem of Alon and Orlitsky [3], as it is obtained as a special case by taking g constant and $f(X, Y) = X$. The optimal rate in “restricted inputs” is given by asymptotic chromatic entropies of graph products. Koulgi et al. show in [4] that this optimal rate is equal to the complementary graph entropy, introduced in [5] by Körner and Longo. No single-letter expression for these quantities is known. In [6], Marton shows that these quantities are closely related to the Shannon capacity of a graph (see [7]), which is a wide open problem.

The similar “unrestricted inputs” zero-error setting of [3] has a single-letter formula for the optimal rate, but its zero-error constraint is much stronger than “restricted inputs” as (X^n, Y^n) can take values out of the support of $P_{X,Y}^n$.

Now the scheme of Figure 1 has been studied with different coding constraints than zero-error “restricted inputs”, and the optimal rate has been characterized in each case: the lossless case by Orlitsky and Roche in [8], the lossy case by Yamamoto in [9], and the zero-error “unrestricted inputs” case by Shayevitz in [10]. These results can only be used as bounds here: the zero-error “restricted inputs” problem depicted in Figure 1 does not have a characterization of the optimal rate.

Numerous extensions of this problem have been studied recently. The distributed context, for instance, has an additional encoder which encodes Y before transmitting it to the decoder. Achievability schemes have been proposed for this setting by Krithivasan and Pradhan in [11] using abelian groups; by Basu et al. in [12] using hypergraphs for the case with maximum distortion criterion; and by Malak and Médard in [13] using hyperplane separations for the continuous lossless case.

Another related context is the network setting, where the function of source random variables from source nodes has to be retrieved at the sink node of a given network. For tree networks, the feasible rate region is characterized by Feizi and Médard in [14] for networks of depth one; and by Sefidgaran and Tchamkerten in [15] under a Markov source distribution hypothesis. In [16], Ravi and Dey consider a bidirectional relay with zero-error “unrestricted inputs” and characterize the rate region for a specific class of functions. In [17], Guang et al. study zero-error function computation on acyclic networks

with limited capacities, and give an inner bound based on network cut-sets. For both distributed and network settings, the “restricted inputs” zero-error problem remains open.

In this paper, we formulate an hypothesis on $P_{X,Y}$ and g that we call “pairwise shared side information” that allows us to derive a single-letter characterization of the optimal rate. This hypothesis is satisfied if every pair of source symbols “share” at least one side information symbol for all output of g . It has graph-theoretic interpretations, as the single-letter formula stems from the particular structure of the characteristic graph of the problem: a disjoint union of OR products. Moreover, this result is of practical interest as it covers the cases with $P_{X,Y}$ full-support, without any assumption on f, g .

In Section II, we present formally the problem of Figure 1. In Section III, we build the characteristic graphs and give an asymptotic formula for the general case. In Section IV, we make the “pairwise shared side information” assumption and give a single-letter characterization of the optimal rate, along with a refinement for $P_{X,Y}$ full-support. We prove it in Section VI and illustrate it in Section V with an example.

II. PROBLEM STATEMENT

We denote sequences by $x^n = (x_1, \dots, x_n)$. The set of probability distributions over \mathcal{X} is denoted by $\Delta(\mathcal{X})$. The distribution of X is denoted by $P_X \in \Delta(\mathcal{X})$, its support is denoted by $\text{supp } P_X$. Given the sequence length $n \in \mathbb{N}^*$, we denote by $\Delta_n(\mathcal{X}) \subset \Delta(\mathcal{X})$ the set of empirical distributions of sequences from \mathcal{X}^n . We denote by $\{0, 1\}^*$ the set of binary words. The setting of Fig. 1 is described by:

- Four finite sets $\mathcal{X}, \mathcal{Y}, \mathcal{U}, \mathcal{Z}$, a couple of random variables $(X, Y) \in \mathcal{X} \times \mathcal{Y}$ drawn with the distribution $P_{X,Y}$ (with P_X and P_Y full-support), and deterministic functions

$$f : \mathcal{X} \times \mathcal{Y} \rightarrow \mathcal{U}, \quad (1)$$

$$g : \mathcal{Y} \rightarrow \mathcal{Z}. \quad (2)$$

For n iterated source uses, we denote by (X^n, Y^n) the sequence of n independent copies of (X, Y) , with probability distribution $P_{X,Y}^n \in \Delta(\mathcal{X}^n \times \mathcal{Y}^n)$.

- The encoder observes the realizations of $X^n, (g(Y_t))_{t \leq n}$ and sends information to the decoder over a noiseless channel of capacity $R \geq 0$.
- The decoder has to recover $(f(X_t, Y_t))_{t \leq n}$ based on the encoder message and the side information Y^n .

Definition II.1 (Zero-error source code, achievable rates)

Given $n \in \mathbb{N}^*$, a (n, R_n) -zero-error source code consists of an encoding function $\phi_e : \mathcal{X}^n \times \mathcal{Z}^n \rightarrow \{0, 1\}^*$ and a decoding function $\phi_d : \mathcal{Y}^n \times \{0, 1\}^* \rightarrow \mathcal{U}^n$ such that:

- 1) the set $\phi_e(\mathcal{X}^n \times \mathcal{Z}^n)$ is prefix-free,
- 2) $R_n = \frac{1}{n} \mathbb{E}[l \circ \phi_e(X^n, (g(Y_t))_{t \leq n})]$, where $l(\cdot)$ denotes the length of a binary word,
- 3) the “restricted inputs” zero-error property is satisfied:

$$\begin{aligned} \forall (x^n, y^n) \in \text{supp } P_{X,Y}^n, \\ \phi_d\left(y^n, \phi_e\left(x^n, (g(y_t))_{t \leq n}\right)\right) = (f(x_t, y_t))_{t \leq n}. \end{aligned} \quad (3)$$

A rate R is achievable if there exists a sequence of (n, R_n) -zero-error source codes such that $\lim_n R_n = R$. The optimal rate is denoted by $R^* = \inf\{R \geq 0 \mid R \text{ is achievable}\}$.

The prefix-free hypothesis guarantees that the decoder knows when the encoder’s message stops. A relaxation of this hypothesis is considered in [3, Theorem 3], without influence on the asymptotic optimal rate. Satisfying all three conditions imply a correct decoding with probability one.

III. GENERAL SETTING

A probabilistic graph G is a tuple $(\mathcal{V}, \mathcal{E}, P_V)$, where \mathcal{V} is the set of vertices, \mathcal{E} is the set of edges, and $P_V \in \Delta(\mathcal{V})$ is an underlying probability distribution on the vertices.

We first build the characteristic graph $G_{[n]}$, which is a probabilistic graph that captures the zero-error encoding constraints on a given number n of source uses. It differs from the graphs used in [10], as we do not need a cartesian representation of these graphs to study the optimal rates. Furthermore, it has a vertex for each possible realization of $(X^n, (g(Y_t))_{t \leq n})$ known at the encoder, instead of \mathcal{X}^n , as in [18] and [3].

Definition III.1 (Characteristic graph $G_{[n]}$) The characteristic graph $G_{[n]}$ is defined by:

- $\mathcal{X}^n \times \mathcal{Z}^n$ as set of vertices with distribution $P_{X,g(Y)}^n$,
- $(x^n, z^n)(x'^n, z'^n)$ are adjacent if $z^n = z'^n$ and there exists $y^n \in g^{-1}(z^n)$ such that:

$$\forall t \leq n, P_{X,Y}(x_t, y_t) P_{X,Y}(x'_t, y_t) > 0, \quad (4)$$

$$\text{and } \exists t \leq n, f(x_t, y_t) \neq f(x'_t, y_t); \quad (5)$$

$$\text{where } g^{-1}(z^n) = \{y^n \in \mathcal{Y}^n \mid (g(y_t))_{t \leq n} = z^n\}.$$

The characteristic graph $G_{[n]}$ is designed with the same core idea as in [18]: (x^n, z^n) and (x'^n, z'^n) are adjacent if there exists a side-information symbol y^n compatible with the observation of the encoder (i.e. $z^n = z'^n$ and $y^n \in g^{-1}(z^n)$), such that $f(x^n, y^n) \neq f(x'^n, y^n)$. In order to prevent erroneous decodings, the encoder must map adjacent pairs of sequences to different codewords; hence the use of graph colorings.

Definition III.2 (Coloring, independent subset) Let $G = (\mathcal{V}, \mathcal{E}, P_V)$ be a probabilistic graph. A subset $\mathcal{S} \subseteq \mathcal{V}$ is independent if $xx' \notin \mathcal{E}$ for all $x, x' \in \mathcal{S}$. Let \mathcal{C} be a finite set (the set of colors), a mapping $c : \mathcal{V} \rightarrow \mathcal{C}$ is a coloring if $c^{-1}(i)$ is an independent subset for all $i \in \mathcal{C}$.

The chromatic entropy of $G_{[n]}$ gives the best rate of n -shot zero-error encoding functions, as in [3].

Definition III.3 (Chromatic entropy H_χ) The chromatic entropy of a probabilistic graph $G = (\mathcal{V}, \mathcal{E}, P_V)$ is defined by

$$H_\chi(G) = \inf \{H(c(V)) \mid c \text{ is a coloring of } G\}. \quad (6)$$

Theorem III.4 (Optimal rate) The optimal rate writes:

$$R^* = \lim_{n \rightarrow \infty} \frac{1}{n} H_\chi(G_{[n]}). \quad (7)$$

Proof Outline. An encoding function ϕ_e is a coloring of $G_{[n]}$ if and only if it satisfies (3) with some decoding function ϕ_d .

Thus the best achievable rate writes

$$R^* = \inf_n \inf_{\phi_e \text{ coloring of } G_{[n]}} H\left(\phi_e\left(X^n, (g(Y_t))_{t \leq n}\right)\right) \quad (8)$$

$$= \lim_{n \rightarrow \infty} \frac{1}{n} H_\chi(G_{[n]}). \quad (9)$$

where (9) comes from Fekete's lemma and (6).

A general single-letter expression for R^* is missing, due to the lack of intrinsic structure of $G_{[n]}$. In Section IV, we introduce a hypothesis that gives structure to $G_{[n]}$ and allows us to derive a single-letter expression for R^* .

IV. PAIRWISE SHARED SIDE INFORMATION

Definition IV.1 *The distribution $P_{X,Y}$ and the function g satisfy the “pairwise shared side information” condition if*

$$\forall z \in \mathcal{Z}, \forall x, x' \in \mathcal{X}, \exists y \in g^{-1}(z), P_{XY}(x, y)P_{XY}(x', y) > 0. \quad (10)$$

This means that for all z output of g , every pair (x, x') “shares” at least one side information symbol $y \in g^{-1}(z)$.

Note that any full-support distribution $P_{X,Y}$ satisfies the “pairwise shared side information” hypothesis.

Definition IV.2 (AND, OR product) *Let $G_1 = (\mathcal{V}_1, \mathcal{E}_1, P_{V_1})$, $G_2 = (\mathcal{V}_2, \mathcal{E}_2, P_{V_2})$ be two probabilistic graphs; their AND (resp. OR) product denoted by $G_1 \wedge G_2$ (resp. $G_1 \vee G_2$) is defined by: $\mathcal{V}_1 \times \mathcal{V}_2$ as set of vertices, $P_{V_1}P_{V_2}$ as probability distribution on the vertices, and $(v_1v_2), (v'_1v'_2)$ are adjacent if*

$$v_1v'_1 \in \mathcal{E}_1 \text{ AND } v_2v'_2 \in \mathcal{E}_2, \quad (11)$$

resp. $(v_1v'_1 \in \mathcal{E}_1 \text{ and } v_1 \neq v'_1)$ OR $(v_2v'_2 \in \mathcal{E}_2 \text{ and } v_2 \neq v'_2)$;

with the convention that all vertices are self-adjacent. We denote by $G_1^{\wedge n}$ (resp. $G_1^{\vee n}$) the n -th AND (resp. OR) power.

For the “restricted inputs” source coding problem in [3], the n -shot characteristic graph is the n -th AND product of the one-shot characteristic graph, and the optimal rate in this problem $\lim_{n \rightarrow \infty} \frac{1}{n} H_\chi(G^{\wedge n})$ does not have a single-letter expression. However, for the “unrestricted inputs” setting there exists such a formula: the Körner graph entropy introduced in [19], which relates to the OR product as shown in Proposition IV.4.

Definition IV.3 (Körner graph entropy H_κ) *For all $G = (\mathcal{V}, \mathcal{E}, P_V)$, let $\Gamma(G)$ be the collection of independent sets of vertices in G . The Körner graph entropy of G is defined by*

$$H_\kappa(G) = \min_{V \in W \in \Gamma(G)} I(W; V), \quad (12)$$

where the minimum is taken over all distributions $P_{W|V} \in \Delta(\mathcal{W})^\mathcal{V}$, with $\mathcal{W} = \Gamma(G)$ and with the constraint that the random vertex V belongs to the random independent set W with probability one, i.e. $V \in W \in \Gamma(G)$ in (12).

Proposition IV.4 (Properties of H_κ) [3, Theorem 5] *For all probabilistic graphs G and G' ,*

$$H_\kappa(G) = \lim_{n \rightarrow \infty} \frac{1}{n} H_\chi(G^{\vee n}), \quad (13)$$

$$H_\kappa(G \vee G') = H_\kappa(G) + H_\kappa(G'). \quad (14)$$

By using a convex combination of Körner graph entropies, we provide a single-letter expression for the optimal rate R^* .

Definition IV.5 (Auxiliary graph G_z^f) *For all $z \in \mathcal{Z}$, we define the auxiliary graph G_z^f by*

- \mathcal{X} as set of vertices with distribution $P_{X|g(Y)=z}$,
- xx' are adjacent if $f(x, y) \neq f(x', y)$ for some $y \in g^{-1}(z) \cap \text{supp } P_{Y|X=x} \cap \text{supp } P_{Y|X=x'}$.

Theorem IV.6 (Pairwise shared side information) *If $P_{X,Y}$ and g satisfy (10), the optimal rate writes:*

$$R^* = \sum_{z \in \mathcal{Z}} P_{g(Y)}(z) H_\kappa(G_z^f). \quad (15)$$

The proof is in Section VI, the keypoint is the particular structure of $G_{[n]}$: a disjoint union of OR products.

Remark IV.7 *The “pairwise shared side information” assumption (10) implies that the adjacency condition (4) is satisfied, which makes $G_{[n]}$ a disjoint union of OR products. Moreover, Körner graph entropies appear in the final expression for R^* , even if $G_{[n]}$ is not an n -th OR power.*

Now consider the case where $P_{X,Y}$ is full-support. This is a sufficient condition to have (10). The optimal rate in this setting is derived from Theorem IV.6, which leads to the analytic expression in Theorem IV.8.

Theorem IV.8 (Optimal rate when $P_{X,Y}$ is full-support)

When $P_{X,Y}$ is full-support, the optimal rate writes:

$$R^* = H(j(X, g(Y))|g(Y)), \quad (16)$$

where the function j returns a word in \mathcal{U}^ , defined by*

$$j : \mathcal{X} \times \mathcal{Z} \rightarrow \mathcal{U}^* \quad (17)$$

$$(x, z) \mapsto (f(x, y'))_{y' \in g^{-1}(z)}.$$

Proof Outline. By Theorem IV.6, $R^* = \sum_{z \in \mathcal{Z}} P_{g(Y)}(z) H_\kappa(G_z^f)$. It can be shown that G_z^f is complete multipartite for all z as $P_{X,Y}$ is full support; and it satisfies $H_\kappa(G_z^f) = H(j(X, g(Y))|g(Y) = z)$.

V. EXAMPLE

In this example, the “pairwise shared side information” assumption is satisfied and R^* is strictly less than a conditional Huffman coding of X knowing $g(Y)$; and also strictly less than the optimal rate without exploiting $g(Y)$ at the encoder.

Consider the probability distribution and function outcomes depicted in Figure 2, with $\mathcal{U} = \{a, b, c\}$, $\mathcal{X} = \{0, \dots, 3\}$, $\mathcal{Y} = \{0, \dots, 7\}$, and $\mathcal{Z} = \{0, 1\}$. Let us show that the “pairwise shared side information” assumption is satisfied. The source

$P_{X,Y}$		Y							
		0	1	2	3	4	5	6	7
		$g(Y) = 0$				$g(Y) = 1$			
X	0	0.1	0.05	*	*	0.05	0.05	*	*
	1	0.1	*	0.05	*	0.05	0.05	0.05	*
	2	0.1	*	*	0.05	*	0.05	*	*
	3	*	0.05	0.05	0.05	*	0.05	*	0.05

$f(\cdot, \cdot)$		Y							
		0	1	2	3	4	5	6	7
		$g(Y) = 0$				$g(Y) = 1$			
X	0	a	b	*	*	b	a	*	*
	1	a	*	b	*	a	a	b	*
	2	b	*	*	c	*	b	*	*
	3	*	c	c	c	*	c	*	c

Fig. 2: An example of $P_{X,Y}$ and g that satisfy (10); along with the outcomes $f(X, Y)$. The elements outside $\text{supp } P_{X,Y}$ are denoted by $*$.

symbols $0, 1, 2 \in \mathcal{X}$ share the SI symbol 0 (resp. 5) when $g(Y) = 0$ (resp. $g(Y) = 1$). The source symbol $3 \in \mathcal{X}$ shares the SI symbols 1, 2, 3 with the source symbols 0, 1, 2, respectively, when $g(Y) = 0$; and the source symbol 3 shares the SI symbol 5 with all other source symbols when $g(Y) = 1$.

Since the ‘‘pairwise shared side information’’ assumption is satisfied, we can use Theorem IV.6; the optimal rate writes

$$R^* = P_{g(Y)=0} H_\kappa(G_0^f) + P_{g(Y)=1} H_\kappa(G_1^f). \quad (18)$$

First we need to determine the probabilistic graphs G_0^f and G_1^f . In G_0^f , the vertex 0 is adjacent to 2 and 3, as $f(0, 0) \neq f(2, 0)$ and $f(0, 1) \neq f(3, 1)$. The vertex 1 is also adjacent to 2 and 3 as $f(1, 0) \neq f(2, 0)$ and $f(1, 2) \neq f(3, 2)$. Furthermore $P_{X|g(Y)=0}$ is uniform, hence $G_0^f = (C_4, \text{Unif}(\mathcal{X}))$ where C_4 is the cycle graph with 4 vertices.

In G_1^f , the vertices 1, 2, 3 are pairwise adjacent as $f(1, 5)$, $f(2, 5)$ and $f(3, 5)$ are pairwise different; and 0 is adjacent to 1, 2 and 3 because of the different function outputs generated by $Y = 4$ and $Y = 5$. Thus, $G_1^f = (K_4, P_{X|g(Y)=1})$ with $P_{X|g(Y)=1} = (\frac{1}{4}, \frac{3}{8}, \frac{1}{8}, \frac{1}{4})$ and K_4 is the complete graph with 4 vertices. An illustration of C_4 and K_4 is given in Figure 3.

Now let us determine $H_\kappa(G_0^f)$ and $H_\kappa(G_1^f)$. On one hand,

$$H_\kappa(G_0^f) = H(V_0) - \max_{V_0 \in W \in \Gamma(G_0^f)} H(V_0|W) \quad (19)$$

$$= 2 - 1 = 1, \quad (20)$$

with $V_0 \sim P_{X|g(Y)=0} = \text{Unif}(\mathcal{X})$; and where $H(V_0|W)$ in (19) is maximized by taking $W = \{0, 1\}$ when $V_0 \in \{0, 1\}$, and $W = \{2, 3\}$ otherwise.

On the other hand,

$$H_\kappa(G_1^f) = \min_{V_1 \in W \in \Gamma(G_1^f)} I(W; V_1) \quad (21)$$

$$= H(V_1) \simeq 1.906, \quad (22)$$



Fig. 3: The graph G depicted here with its underlying probability distribution satisfies $G = (C_4, \text{Unif}(\{1, \dots, 4\})) \sqcup (K_4, \text{Unif}(\{1, \dots, 4\}))$; where C_4 (resp. K_4) is the cycle (resp. complete) graph with 4 vertices.

with $V_1 \sim P_{X|g(Y)=1}$; where (22) follows from $\Gamma(G_1^f) = \{\{0\}, \dots, \{3\}\}$, as G_1^f is complete. Hence $R^* \simeq 1.362$.

The rate that we would obtain by transmitting X knowing $g(Y)$ at both encoder and decoder with a conditional Huffman algorithm writes: $R_{\text{Huff}} = H(X|g(Y)) \simeq 1.962$.

The rate that we would obtain without exploiting $g(Y)$ at the encoder is $R_{\text{No } g} = H(X) \simeq 1.985$, because of the different function outputs generated by $Y = 4$ and $Y = 5$.

Finally, $H(f(X, Y)|Y) \simeq 0.875$.

In this example we have

$$H(X) = R_{\text{No } g} > R_{\text{Huff}} > R^* > H(f(X, Y)|Y). \quad (23)$$

This illustrates the impact of the side information at the encoder in this setting, as we can observe a large gap between the optimal rate R^* and $R_{\text{No } g}$.

VI. PROOF OF THEOREM IV.6

A. Definitions

We will use the disjoint union of probabilistic graphs, which generalizes the existing concept of a disjoint union of graphs without underlying probability distribution [20, Section 1.4]. An example of disjoint union is depicted in Figure 3. We also need to formalize the concept of isomorphic probabilistic graphs, i.e. same structure and underlying distribution.

Definition VI.1 (Disjoint union of probabilistic graphs)

Let $N \in \mathbb{N}^*$; let $G = (\mathcal{V}, \mathcal{E}, P_V)$ and for all $i \leq N$, let $\tilde{G}_i = (\mathcal{V}_i, \mathcal{E}_i, P_{V_i})$. We say that G is the disjoint union of the (\tilde{G}_i) , denoted by $G = \bigsqcup_{i \leq N} \tilde{G}_i$, if the following is satisfied:

- \mathcal{V} is the disjoint union of the sets $(\mathcal{V}_i)_{i \leq N}$, i.e. $\mathcal{V} = \bigcup_i \mathcal{V}_i$ and $\mathcal{V}_i \cap \mathcal{V}_{i'} = \emptyset$ for all $i \neq i'$;
- For all $v, v' \in \mathcal{V}$, let i, i' be the unique indexes such that $v \in \mathcal{V}_i$ and $v' \in \mathcal{V}_{i'}$. Then if $i = i'$, $vv' \in \mathcal{E} \iff vv' \in \mathcal{E}_i$; if $i \neq i'$, $vv' \notin \mathcal{E}$;
- For all $i \leq n$, $P_{V|V \in \mathcal{V}_i} = P_{V_i}$.

Definition VI.2 (Isomorphic probabilistic graphs)

Let $G_1 = (\mathcal{V}_1, \mathcal{E}_1, P_{V_1})$ and $G_2 = (\mathcal{V}_2, \mathcal{E}_2, P_{V_2})$. We say that G_1 is isomorphic to G_2 (denoted by $G_1 \simeq G_2$) if there exists an isomorphism between them, i.e. a bijection $\psi : \mathcal{V}_1 \rightarrow \mathcal{V}_2$ such that:

- For all $v_1, v'_1 \in \mathcal{V}_1$, $v_1 v'_1 \in \mathcal{E}_1 \iff \psi(v_1) \psi(v'_1) \in \mathcal{E}_2$,
- For all $v_1 \in \mathcal{V}_1$, $P_{V_1}(v_1) = P_{V_2}(\psi(v_1))$.

B. Main proof

Let us specify the adjacency condition in $G_{[n]}$ under the assumption (10). Two vertices are adjacent if they satisfy (4) and (5); however (4) is always satisfied under (10). Thus $(x^n, z^n)(x'^n, z'^n)$ are adjacent if $z^n = z'^n$ and

$$\exists y^n \in g^{-1}(z^n), \exists t \leq n, f(x_t, y_t) \neq f(x'_t, y_t). \quad (24)$$

It can be observed that the condition (24) is the adjacency condition of an OR product of adequate graphs; more precisely,

$$G_{[n]} = \bigsqcup_{z^n \in \mathcal{Z}^n} \bigvee_{t \leq n} G_{z_t}^f. \quad (25)$$

Although $G_{[n]}$ cannot be expressed as an n -th OR power, we will show that its chromatic entropy asymptotically coincide with that of an appropriate OR power: we now search for an asymptotic equivalent of $H_\chi(G_{[n]})$.

Definition VI.3 \mathcal{S}_n is the set of colorings of $G_{[n]}$ that can be written as $(x^n, z^n) \mapsto (T_{z^n}, \tilde{c}(x^n, z^n))$ for some mapping $\tilde{c}: \mathcal{X}^n \times \mathcal{Z}^n \rightarrow \tilde{\mathcal{C}}$; where T_{z^n} denotes the type of z^n .

In the following, we define $Z^n \doteq (g(Y_t))_{t \leq n}$. Now we need several Lemmas. Lemma 1 states that the optimal coloring $c(x^n, z^n)$ of $G_{[n]}$ has the type of z^n as a prefix at a negligible rate cost. Lemma 2 is an adapted version for chromatic entropies of the following observation: minimum colorings on each connected component induce a minimum coloring of the whole graph. Lemma 3 gives an asymptotic formula for the minimal entropy of the colorings from \mathcal{S}_n .

Lemma 1 The following asymptotic comparison holds:

$$H_\chi(G_{[n]}) = \inf_{\substack{c \text{ coloring of } G_{[n]} \\ \text{s.t. } c \in \mathcal{S}_n}} H(c(X^n, Z^n)) + O(\log n). \quad (26)$$

Lemma 2 Let $N \in \mathbb{N}^*$; let $G = (\mathcal{V}, \mathcal{E}, P_V)$ and for all $i \leq N$, let $\tilde{G}_i = (\mathcal{V}_i, \mathcal{E}_i, P_{V_i})$ be probabilistic graphs such that $G = \bigsqcup_i \tilde{G}_i$ and $\tilde{G}_1 \simeq \dots \simeq \tilde{G}_N$. Then we have $H_\chi(G) = H_\chi(\tilde{G}_1)$.

Lemma 3 The following asymptotic comparison holds:

$$\inf_{\substack{c \text{ coloring of } G_{[n]} \\ \text{s.t. } c \in \mathcal{S}_n}} H(c(X^n, Z^n)) = n \sum_{z \in \mathcal{Z}} P_{g(Y)}(z) H_\kappa(G_z^f) + o(n). \quad (27)$$

The keypoint of the proof of Lemma 1 is the asymptotically negligible entropy of the prefix T_{Z^n} of the colorings of \mathcal{S}_n .

Lemma 2 is proved using the concavity of the entropy, which implies the following: an optimal coloring colors all the isomorphic connected components the same way.

The proof of Lemma 3 relies on the decomposition $G_{[n]} = \bigsqcup_{Q_n \in \Delta_n(\mathcal{Z})} G_{[n]}^{Q_n}$, where $G_{[n]}^{Q_n}$ is the subgraph induced by the vertices (x^n, z^n) such that the type of z^n is Q_n . We show that $G_{[n]}^{Q_n}$ is a disjoint union of isomorphic graphs whose chromatic entropy is given by Lemma 2 and (14): $|H_\chi(G_{[n]}^{Q_n}) - n \sum_{z \in \mathcal{Z}} Q_n(z) H_\kappa(G_z^f)| \leq n\epsilon_n$. Finally, uniform convergence arguments enable us to conclude.

Now let us combine these results together:

$$R^* = \frac{1}{n} H_\chi(G_{[n]}) + o(1) \quad (28)$$

$$= \frac{1}{n} \inf_{\substack{c \text{ coloring of } G_{[n]} \\ \text{s.t. } c \in \mathcal{S}_n}} H(c(X^n, Z^n)) + o(1) \quad (29)$$

$$= \sum_{z \in \mathcal{Z}} P_{g(Y)}(z) H_\kappa(G_z^f) + o(1), \quad (30)$$

where (28) comes from Theorem III.4, (29) comes from Lemma 1, and (30) comes from Lemma 3. The proof of Theorem IV.6 is complete.

C. Proof of Lemmas 1, 2, 3

Proof of Lemma 1. Let c_n^* be the coloring of $G_{[n]}$ with minimal entropy. Then we have:

$$H_\chi(G_{[n]}) = \inf_{c \text{ coloring of } G_{[n]}} H(c(X^n, Z^n)) \quad (31)$$

$$\leq \inf_{\substack{c \text{ coloring of } G_{[n]} \\ \text{s.t. } c \in \mathcal{S}_n}} H(c(X^n, Z^n)) \quad (32)$$

$$= \inf_{\substack{c: (x^n, z^n) \\ \mapsto (T_{z^n}, \tilde{c}(x^n, z^n))}} H(T_{Z^n}, \tilde{c}(X^n, Z^n)) \quad (33)$$

$$\leq H(T_{Z^n}) + H(c_n^*(X^n, Z^n)) \quad (34)$$

$$= H_\chi(G_{[n]}) + O(\log n), \quad (35)$$

where (33) comes from Definition VI.3; (34) comes from the subadditivity of the entropy, and the fact that $(x^n, z^n) \mapsto (T_{z^n}, c_n^*(x^n, z^n))$ is a coloring of $G_{[n]}$ that belongs to \mathcal{S}_n ; and (35) comes from $H(T_{Z^n}) = O(\log n)$, as $\log |\Delta_n(\mathcal{Z})| = O(\log n)$. The desired equality comes from the bounds $H_\chi(G_{[n]})$ and $H_\chi(G_{[n]}) + O(\log n)$ on (32).

Proof of Lemma 2. Let $(\tilde{G}_i)_{i \leq N}$ be isomorphic probabilistic graphs and G such that $G = \bigsqcup_i \tilde{G}_i$. Let $c_1^*: \mathcal{V}_1 \rightarrow \mathcal{C}$ be the coloring of \tilde{G}_1 with minimal entropy, and let c^* be the coloring of G defined by

$$c^*: \mathcal{V} \rightarrow \mathcal{C} \quad (36)$$

$$v \mapsto c_1^* \circ \psi_{i_v \rightarrow 1}(v), \quad (37)$$

where i_v is the unique integer such that $v \in \mathcal{V}_{i_v}$, and $\psi_{i_v \rightarrow 1}: \mathcal{V}_{i_v} \rightarrow \mathcal{V}_1$ is an isomorphism between \tilde{G}_{i_v} and \tilde{G}_1 . In other words c^* applies the same coloring pattern c_1^* on each connected component of G . We have

$$H_\chi(G) \leq H(c^*(V)) \quad (38)$$

$$= h\left(\sum_{j \leq N} P_{i_v}(j) P_{c^*(V_j)}\right) \quad (39)$$

$$= h\left(\sum_{j \leq N} P_{i_v}(j) P_{c_1^*(V_1)}\right) \quad (40)$$

$$= H(c_1^*(V_1)) \quad (41)$$

$$= H_\chi(\tilde{G}_1), \quad (42)$$

where h denotes the entropy of a distribution; (40) comes from the definition of c^* ; and (42) comes from the definition of c_1^* .

Now let us prove the upper bound on $H_\chi(\tilde{G}_1)$. Let c be a coloring of G , and let $i^* \doteq \operatorname{argmin}_i H(c(V_i))$ (i.e. i^* is the index of the connected component for which the entropy of the coloring induced by c is minimal). We have

$$H(c(V)) = h\left(\sum_{j \leq N} P_{i_V}(j) P_{c(V_j)}\right) \quad (43)$$

$$\geq \sum_{j \leq N} P_{i_V}(j) h(P_{c(V_j)}) \quad (44)$$

$$\geq \sum_{j \leq N} P_{i_V}(j) H(c(V_{i^*})) \quad (45)$$

$$\geq H_\chi(\tilde{G}_{i^*}), \quad (46)$$

$$= H_\chi(\tilde{G}_1), \quad (47)$$

where (44) follows from the concavity of h ; (45) follows from the definition of i^* ; (46) comes from the fact that c induces a coloring of \tilde{G}_{i^*} ; (47) comes from the fact that \tilde{G}_1 and \tilde{G}_{i^*} are isomorphic. Now, we can combine the bounds (42) and (47): for all coloring c of G we have

$$H_\chi(G) \leq H_\chi(\tilde{G}_1) \leq H(c(V)), \quad (48)$$

which yields the desired equality when taking the infimum over c .

Proof of Lemma 3. For all $Q_n \in \Delta_n(\mathcal{Z})$, let

$$G_{[n]}^{Q_n} = \bigsqcup_{\substack{z^n \in \mathcal{Z}^n \\ T_{z^n} = Q_n}} \bigvee_{t \leq n} G_{z_t}^f, \quad (49)$$

with the probability distribution induced by $P_{X,Z}^n$. This graph is formed of the connected components of $G_{[n]}^{Q_n}$ whose corresponding z^n has type Q_n . We need to find an equivalent for $H_\chi(G_{[n]}^{Q_n})$. Since $G_{[n]}^{Q_n}$ is a disjoint union of isomorphic graphs, we can use Lemma 2:

$$H_\chi(G_{[n]}^{Q_n}) = H_\chi\left(\bigvee_{z \in \mathcal{Z}} (G_z^f)^{\vee n Q_n(z)}\right). \quad (50)$$

On one hand,

$$H_\chi\left(\bigvee_{z \in \mathcal{Z}} (G_z^f)^{\vee n Q_n(z)}\right) \geq H_\kappa\left(\bigvee_{z \in \mathcal{Z}} (G_z^f)^{\vee n Q_n(z)}\right) \quad (51)$$

$$= n \sum_{z \in \mathcal{Z}} Q_n(z) H_\kappa(G_z^f), \quad (52)$$

where (51) comes from $H_\kappa \leq H_\chi$ [3, Lemma 14], (52) comes from (14). On the other hand,

$$H_\chi\left(\bigvee_{z \in \mathcal{Z}} (G_z^f)^{\vee n Q_n(z)}\right) \leq \sum_{z \in \mathcal{Z}} Q_n(z) H_\chi((G_z^f)^{\vee n}) \quad (53)$$

$$= n \sum_{z \in \mathcal{Z}} Q_n(z) H_\kappa(G_z^f) + n\epsilon_n, \quad (54)$$

where $\epsilon_n \doteq \max_z \frac{1}{n} H_\chi((G_z^f)^{\vee n}) - H_\kappa(G_z^f)$ is a quantity that does not depend on Q_n and satisfies $\lim_{n \rightarrow \infty} \epsilon_n = 0$; (53)

comes from the subadditivity of H_χ . Combining equations (50), (52) and (54) yields

$$\left| H_\chi(G_{[n]}^{Q_n}) - n \sum_{z \in \mathcal{Z}} Q_n(z) H_\kappa(G_z^f) \right| \leq n\epsilon_n. \quad (55)$$

Now, we have an equivalent for $H_\chi(G_{[n]}^{Q_n})$.

$$\inf_{\substack{c \text{ coloring of } G_{[n]} \\ \text{s.t. } c \in \mathcal{S}_n}} H(c(X^n, Z^n)) \quad (56)$$

$$= \inf_{\substack{c: (x^n, z^n) \\ \mapsto (T_{z^n}, \tilde{c}(x^n, z^n))}} H(\tilde{c}(X^n, Z^n) | T_{Z^n}) + H(T_{Z^n}) \quad (57)$$

$$= \inf_{\substack{c: (x^n, z^n) \\ \mapsto (T_{z^n}, \tilde{c}(x^n, z^n))}} \sum_{Q_n \in \Delta_n(\mathcal{Z})} P_{T_{Z^n}}(Q_n) H(\tilde{c}(X^n, Z^n) | T_{Z^n} = Q_n) + O(\log n) \quad (58)$$

$$= \sum_{Q_n \in \Delta_n(\mathcal{Z})} P_{T_{Z^n}}(Q_n) \inf_{c_{Q_n} \text{ coloring of } G_{[n]}^{Q_n}} H(c_{Q_n}(X^n, Z^n) | T_{Z^n} = Q_n) + O(\log n) \quad (59)$$

$$= \sum_{Q_n \in \Delta_n(\mathcal{Z})} P_{T_{Z^n}}(Q_n) H_\chi(G_{[n]}^{Q_n}) + O(\log n) \quad (60)$$

$$= \sum_{Q_n \in \Delta_n(\mathcal{Z})} P_{T_{Z^n}}(Q_n) \left(n \sum_{z \in \mathcal{Z}} Q_n(z) H_\kappa(G_z^f) \pm n\epsilon_n \right) + O(\log n) \quad (61)$$

$$= n \sum_{Q_n \in \Delta_n(\mathcal{Z})} 2^{-nD(Q_n \| P_{g(Y)}) + o(n)} \left(\sum_{z \in \mathcal{Z}} Q_n(z) H_\kappa(G_z^f) \right) \pm n\epsilon_n + O(\log n) \quad (62)$$

$$= n \sum_{z \in \mathcal{Z}} P_{g(Y)}(z) H_\kappa(G_z^f) + o(n), \quad (63)$$

where (58) comes from $H(T_{Z^n}) = O(\log n)$, as $\log |\Delta_n(\mathcal{Z})| = O(\log n)$; (59) follows from the fact that the entropy of \tilde{c} can be minimized independently on each $G_{[n]}^{Q_n}$; (60) follows from the definition of $G_{[n]}^{Q_n}$; (61) comes from (55); (62) comes from [21, Lemma 2.6] and the fact that ϵ_n does not depend on Q_n .

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