



# A Matheuristic for the Commodity Constrained Split Delivery VRP

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## A Matheuristic for the Commodity Constrained Split Delivery VRP

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## Introduction

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The Commodity Constrained Split Delivery VRP  
Context and related literature

# INTRODUCTION

## Introduction

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# EXAMPLE [Archetti et al., 2016]

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# Separate routing strategy

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## The Commodity Constrained Split Delivery VRP

Context and related literature

# Commodity constrained split delivery mixed routing strategy

## Context and related literature

The C-SDVRP:

- was first introduced in [Archetti et al., 2016]
- exact algorithms : Branch-and Cut algorithm [Archetti et al., 2016]; Branch-price-and-cut algorithms [Archetti et al., 2015],[Gschwind et al., 2019]
- heuristics: VRP based heuristic with node replications [Archetti et al., 2016]; ALNS [Gu et al., 2019]
- has an application in the short and local fresh food supply chain [Gu et al., 2022]

## AN EXTENDED FORMULATION

## Notation

- $\mathcal{R}$  set of **feasible routes**
- $C_r$ : **cost of route**  $r \in \mathcal{R}$
- $\mathcal{C}$  set of customers.
- $\mathcal{K}$  set of commodities.
- $R_{jk} \geq 0$  demand of customer  $j \in \mathcal{C}$  for commodity  $k \in \mathcal{K}$ .
- $a_{jk}^r$ : **binary parameter** = 1 if route  $r \in \mathcal{R}$  delivers  $k \in \mathcal{K}_j$  to customer  $j \in \mathcal{N}$ , = 0 otherwise
- For all routes  $r \in \mathcal{R}$  a **binary variable**

$$\lambda_r = \begin{cases} 1 & \text{if route } r \text{ is in the solution,} \\ 0 & \text{otherwise.} \end{cases}$$

## Formulation

$$\min \sum_{r \in \mathcal{R}} C_r \lambda_r \quad (1)$$

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$$\lambda_r \in \{0, 1\} \quad \forall r \in \mathcal{R} \quad (4)$$

## A MATHEURISTIC

## General setting

- We price  $\lambda_r$  variables,  $r \in \mathcal{R}_o$ ,  $o \in \mathcal{D}$
- The pricing problem  $\rho_o := \min\{\bar{C}_r \mid r \in \mathcal{R}_o\}$ ,  $o \in \mathcal{D}$  can be formulated as an **Elementary Shortest Path Problem with Resource Constraints (ESPPRC)**  
[Archetti et al., 2015]
- We tackle the ESPPRC by a **label setting dynamic programming technique** [Feillet et al., 2004]

## Classical features

Our matheuristic incorporates

- the **ng-path** relaxation [Baldacci et al., 2011] and the **bidirectional labeling search** [Righini and Salani, 2006]
- **heuristic approaches** to solve the ESPPRC
- automatic dual pricing smoothing **stabilization** [Pessoa et al., 2020]
- **capacity cuts** separation [Ralphs et al., 2003]
- separation of valid inequalities for the set covering polytope and the number partitioning polytope

## Heuristic Column Generator (HCG)

A three phase procedure:

- **First phase: define a subgraph** by using a clustering technique on the set of customers
- **Second phase: compute a lower bound** on the optimal value of the pricing problem by solving an ESPPRC on a modified graph
- **Third phase: determine all the negative reduced cost feasible routes** by solving the ESPPRC in **acyclic graphs**

## Second phase

## Third phase

# Valid Inequalities

## Valid Inequalities

- **Set covering polytope** [Balas and Ng, 1989]:

- for all customer  $j \in \mathcal{C}$  and all subsets  $\hat{\mathcal{K}} \subseteq \mathcal{K}_j$

$\mathcal{R}_j^{\hat{\mathcal{K}}}$ : set of routes delivering to  $j$  all the commodities in  $\hat{\mathcal{K}}$

$\bar{\mathcal{R}}_j^{\hat{\mathcal{K}}}$ : set of routes delivering to  $j$  some (but not all) commodities in  $\hat{\mathcal{K}}$

$$2 \sum_{r \in \mathcal{R}_j^{\hat{\mathcal{K}}}} \lambda_r + \sum_{r \in \bar{\mathcal{R}}_j^{\hat{\mathcal{K}}}} \lambda_r \geq 2, \quad \forall j \in \mathcal{C}, \forall \hat{\mathcal{K}} \subseteq \mathcal{K}_j$$

## Valid Inequalities

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- **Number partitioning problem polytope**:

- for all customer  $j \in \mathcal{C}$  and all  $k = 1, \dots, |\mathcal{K}_j|$

$x_k^j$ : number of routes which deliver  $k$  commodities to  $j$

$$\sum_{k=1}^{|\mathcal{K}_j|} kx_k^j \leq |\mathcal{K}_j|, \quad \forall j \in \mathcal{C}$$

## COMPUTATIONAL RESULTS

## Instances

Instances from [Archetti et al., 2016] built from the R101 and C101 Solomon instances for the VRP with Time Windows.

- $|\mathcal{C}| \in \{15, 20, 40, 60, 80, 100\}$
- $|\mathcal{K}| \in \{2, 3\}$

Each commodity required with a probability  $p$  of 0.6 or 1  
Quantity of each commodity within  $[1, 100]$  or  $[40, 60]$ .

## Summary of results on small sized instances.

instances					Matheuristic vs ALNS			
	<i>n</i>	<i>M</i>	<i>p</i>	<i>avg.n<sub>cc</sub></i>	<i>nblns</i>	<i>avg.gap</i>	<i>avg.time</i>	<i>%time.improv.</i>
C101	15	2	0.6/1	22/30	16	0.36	0.36	-96.4
C101	15	3	0.6/1	28/45	16	0.02	0.22	-97.6
R101	15	2	0.6/1	22/30	16	0.17	3.62	-81.3
R101	15	3	0.6/1	28/45	16	0.16	1.20	-93.7

## Preliminary results on large sized instances.

instances					Matheuristic vs ALNS			
	<i>n</i>	<i>M</i>	<i>p</i>	avg. <i>n<sub>cc</sub></i>	<i>nblns</i>	avg.gap	avg.time	%time.improv.
C101	100	2	0.6	136.4	10	0.67	74.45	-71.7
C101	100	3	1	300	5	0.86	633.5	-44.5
R101	100	2	0.6	136.4	10	0.72	271.3	-24.0
R101	100	3	1	300	5	-0.28	82.6	-90.8

## FUTURE WORK

## Future work

- Perform extensive **computational experiments**
- Develop a Branch&Price&Cut algorithm

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THANK YOU!