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# Convexity preserving deformations of digital sets: Characterization of removable and insertable pixels

Lama Tarsissi<sup>a,b,\*</sup>, Yukiko Kenmochi<sup>c</sup>, Pascal Romon<sup>d</sup>, David Coeurjolly<sup>e</sup>,  
Jean-Pierre Borel<sup>f</sup>

<sup>a</sup>*LIGM, Univ Gustave Eiffel, CNRS, ESIEE Paris, F-77454 Marne-la-Vallée, France*

<sup>b</sup>*Univ Sorbonne Abu-Dhabi, United Arab of Emirates*

<sup>c</sup>*GREYC, CNRS – Univ. Caen Normandie – ENSICAEN, F-14000, Caen, France*

<sup>d</sup>*LAMA, Univ Gustave Eiffel, UPEC, CNRS, F-77454 Marne-la-Vallée, France*

<sup>e</sup>*Univ Lyon, CNRS, LIRIS, France*

<sup>f</sup>*Mathis-XLim, Univ. Limoges - CNRS, F-87060 Limoges, France*

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## Abstract

In this paper, we are interested in digital convexity. This notion is applied in several domains like image processing and discrete tomography. We choose to study the inflation and deflation of digital convex sets while maintaining the convexity property. Knowing that any digital convex set can be read and identified by its boundary word, we use the combinatorics on words perspective instead of a purely geometric approach. In this context, we characterize the points that can be added or removed over the digital convex sets without losing its convexity. Some algorithms are given at the end of each section with examples on each process.

*Keywords:* Digital convexity, Christoffel words, Lyndon words, Digital set inflation, Digital set deflation

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## 1. Introduction

In this paper, we provide characterizations of convexity-preserving removable and insertable points of digital convex sets. In  $\mathbb{Z}^2$ , even a simple transformation (such as rotation) of a digital convex set can cause the loss of the digital convexity property. We aim to provide the discrete-equivalent of infinitesimal transform, from which we can derive convexity-preserving set deformations. In order to study such atomic deformations that preserve digital convexity, we need to investigate deflation and inflation processes [29]. The theory of combinatorics on words provides useful tools and techniques for our investigation [21]. Relying

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\*Corresponding author

*Email addresses:* lama.tarsissi@esiee.fr (Lama Tarsissi),  
yukiko.kenmochi@unicaen.fr (Yukiko Kenmochi), pascal.romon@univ-eiffel.fr (Pascal Romon), david.coeurjolly@liris.cnrs.fr (David Coeurjolly),  
jean-pierre.borel@unilim.fr (Jean-Pierre Borel)

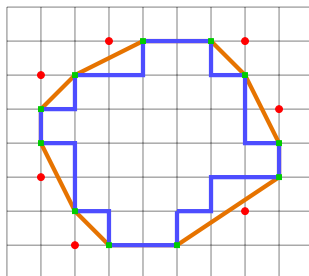


Figure 1: The green points are the removable points, which are the vertices of the convex hull of a digital convex set  $C$ . The red points are the insertable points, which are the exterior Bézout point of each segment of the convex hull.

on this theory, we are able to characterize the removable points  $x$  of a digital convex set  $C$ , such that  $C \setminus \{x\}$  is still digitally convex. Similarly, we provide the characterization of the insertable points  $x \in \overline{C}$  (the complement of  $C$  in  $\mathbb{Z}^2$ ) such that  $C \cup \{x\}$  is still digitally convex. These two actions can only be applied at some specific points on the digital convex set if we want to preserve digital convexity. As a matter of fact, deflation is easy; finding the correct point to remove from a digital set is simple from a geometrical point of view. On the other hand, inflation is more involved. Propositions and theorems for inflation and deflation are naturally associated with algorithms in the mathematical sense. In [30], we provide implementable algorithm, containing all the necessary details.

The plan of the paper is the following. Section 2 provides an overview on digital convexity, from a combinatorics on words perspective, and gives the basic notations needed to understand our results. In Section 3, we characterize removable points and prove that such point can be any simple point that is a corner of the convex hull of  $C$ . As said before, finding a characterization of insertable points is not an easy task. We provide necessary and sufficient conditions to determine candidate points in Section 4. The necessary condition is based on a result from [13]: the authors proved that adding the closest outer point of a segment maintains its digital convexity. Based on this result, we provide in this paper the characterization of all insertable points for the whole digital convex set. We provide two main results for the sufficient condition. The first one is the general case and leads to propagation after inflation. The second results imposes a strong constraint on the sufficient condition. Figure 1 shows an example of a digitally convex set with some removable and insertable points. For both procedures, and after each iteration, we must consider an update for the segments of the convex hull. In this paper, we discuss all the possible cases for this update. One of them is presented in Figure 2. The last section is left for the conclusions and perspectives.

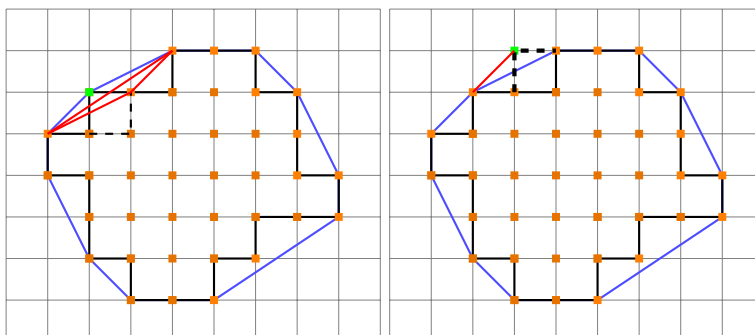


Figure 2: The first picture shows an example of deflation. The green point  $x$  is the point to be removed from a digital convex set (orange points)  $C$ . The second picture shows an example of inflation. The green point  $x$  is the new point to be added. The blue and red segments in both pictures are respectively  $\text{Conv}(C)$  and its modifications after each process.

## 2. Digital convexity and combinatorics on words

In this section, we first give the definition of digital convex sets based on convex hull [18], which are also called H-convex sets [15]. We then show that their boundaries can be expressed by words represented with the Freeman chain code [17], called a boundary word. After recalling basic notions of combinatorics on words, we present the characterization of digital convexity along boundary words using those notions [8].

### 2.1. Definitions of digital convex sets

In  $\mathbb{R}^2$ , a subset  $R$  is said to be convex if for any pair of points  $x, y \in R$ , every point on the straight line segment joining  $x$  and  $y$  is also within  $R$ . This notion, however, cannot be straightforwardly applied to subsets in  $\mathbb{Z}^2$ . In order to tackle this issue, various notions of convexity of a set  $S$  of  $\mathbb{Z}^2$  have been proposed. They are categorized into four approaches; the first approach is based on no existence of triplet of collinear points  $(p, q, r)$  such that  $p, r \in S$  and  $q \in \bar{S}$  [24]. The second one is based on existence of a convex set  $R \subset \mathbb{R}^2$  such that the digitization of  $R$  is  $S$  [28]. Kim proposed another definition, which can be characterized by verifying if the digitization of the convex hull of  $S$  is equal to  $S$ , and also showed its equivalence with the two former [18]; the definition based on convex hull is later called H-convex by Eckhardt [15]. The fourth approach is based on the notion of digital line segments [19], such that  $S$  is digitally convex if the digital line segment joining any pair of points of  $S$  belongs to  $S$ . Note that this last notion guarantees the connectivity of  $S$  contrary to the other three. Under the connectivity assumption, it is then shown that the above different concepts coincide [15].

The definition of connectivity on  $\mathbb{Z}^2$  is based on the notion of neighborhood. Given a point  $p \in \mathbb{Z}^2$ , the neighborhood of  $p$  is defined by  $\mathcal{N}(p) = \{q \in \mathbb{Z}^2 : \|p - q\|_1 \leq 1\}$ , also called 4-neighborhood. We say that a point  $q$  is 4-adjacent to  $p$  if  $q \in \mathcal{N}(p) \setminus \{p\}$ . From the reflexive-transitive closure of this adjacency

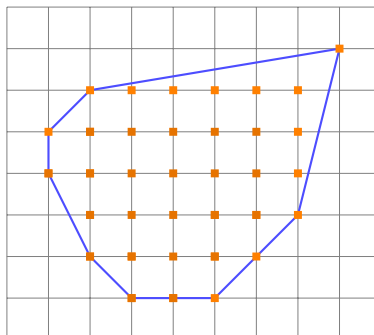


Figure 3: An example of a digital set  $C$  that is digitally convex but not 4-connected. The orange points represents the elements  $x \in C$  while the blue line represents  $\text{Conv}(C)$ .

relation on a finite subset  $X \subset \mathbb{Z}^2$ , we derive the 4-connectivity relation on  $X$ , which is an equivalence relation. If there is exactly one equivalence class for this relation, then we say that  $X$  is 4-connected. In this article, we consider finite and 4-connected sets of  $\mathbb{Z}^2$ . As this 4-connectivity assumption yields the coincidence of the above different definitions of digital convexity [15]<sup>1</sup>, we use the one based on convex hull [15, 18], which is defined for a finite subset  $Y \subset \mathbb{R}^2$  as:

$$\text{Conv}(Y) := \{x \in \mathbb{R}^2 \mid x = \sum_{y \in Y} \lambda_y y \wedge \sum_{y \in Y} \lambda_y = 1 \wedge \forall y \in Y, \lambda_y \geq 0\}.$$

**Definition 1.** (*Digital convexity [18]*) A finite subset  $S$  of  $\mathbb{Z}^2$  is digitally convex if  $\text{Conv}(S) \cap \mathbb{Z}^2 = S$ .

From this definition we have the following remark, which is contrary to the implication of connectivity in the concept of convexity in  $\mathbb{R}^2$  (see Figure 3 for an example).

**Remark 1.** *Digital convexity does not imply connectivity in  $\mathbb{Z}^2$ .*

The above definition of digital convexity is also used in [12], whose aim is verifying if a given finite 4-connected set  $S$  of  $\mathbb{Z}^2$  is digitally convex. Their approach focuses rather on the boundary of  $S$ , on which maximal line segments and their arithmetic properties are analyzed [12]. On the other hand, properties of the boundary of  $S$  based on combinatorics on words are also studied [8].

The problems treated in this article are different from the ones from [8, 12], but are related. Our problem is stated as follows; given a finite, 4-connected and digitally convex set  $C$  of  $\mathbb{Z}^2$  and a point  $p$  of  $C$  (resp. the complement  $\bar{C}$ ),

<sup>1</sup>In [15], 8-connectivity, to which the 4-connectivity leads, is assumed to show the coincidence of the different concepts of digital convexity.

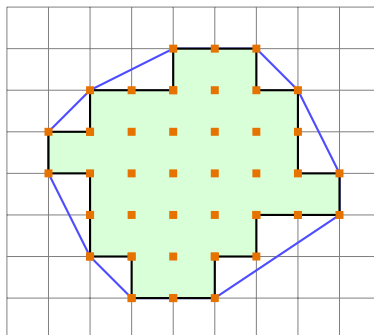


Figure 4: A digital convex set  $C$  represented by the orange points, and its boundary path  $Bd(C)$  represented by the ordered orange points linked by black line segments. The boundary word  $W(C)$  is then given by 10100100110100010100 while the blue polygonal line represents the convex hull  $Conv(C)$ .

we would like to verify if  $C \setminus \{p\}$  (resp.  $C \cup \{p\}$ ) is still digitally convex (and 4-connected). In order to answer such questions, we use the boundary properties based on combinatorics on words that are presented in [8].

## 2.2. Boundary words of digitally convex sets

Let  $C \subset \mathbb{Z}^2$  be a finite, 4-connected digitally convex set. The border points of  $C$  can be tracked by a classical border following algorithm (for example, see [1] for “left-hand-on-wall” border following), which generates a 4-connected sequence of the border points of  $C$ . Note that the sequence can include dead-ends and thus sometimes turnaround sub-sequences. Such a sequence is also called the boundary path of  $C$ , denoted by  $Bd(C)$ , and represented by a word obtained by Freeman chain code [17], denoted by  $W(Bd(C))$  or simply  $W(C)$  and called the boundary word of  $C$  [16]. Boundary words are thus defined over an alphabet of four letters  $0, 1, \bar{0}, \bar{1}$ , which are associated to the right, up, left and down steps, respectively. See Figure 4 for an example.

It is then observed that the boundary word  $W(C)$  of any digital convex set  $C$  can be factorized into four sub-words, such that each sub-word is a binary word, i.e., it contains only two letters. For such a factorization, we first enclose the boundary path  $Bd(C)$  by its bounding box, and cut  $Bd(C)$  into four parts at the four intersections with the bounding box; for example, on the left side of the bounding box, we consider the lowest intersection point  $W$  as a cutting position. Similarly, we can find the other three cutting positions, denoted by  $N$ ,  $E$  and  $S$ , as seen in Figure 5. Starting from  $W$  in the clockwise direction and ending at  $N$ , we obtain the  $WN$ -path as a part of  $Bd(C)$ , which is associated to the  $WN$ -word of  $W(C)$ . Similarly we obtain  $NE$ -,  $ES$ -, and  $SW$ -paths and their associated words. Figure 5 illustrates the factorization of boundary word and shows that each word is a binary word. This factorization allows us to treat the four parts of the boundary of any digital convex set independently, and to introduce the notion of digital convexity adapted to each part [8].

**Definition 2 ([8]).** A word  $w$  is said WN-convex, if it codes the WN-word of the boundary word of some finite, 4-connected and digitally convex set of  $\mathbb{Z}^2$ .

Similarly, we can define WE-, ES-, SW-convex words.

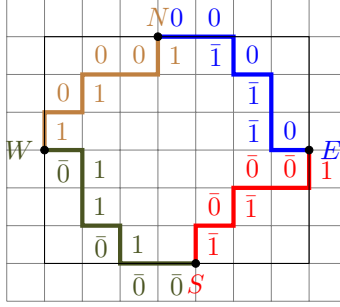


Figure 5: The four parts WN, NE, ES and SW of the boundary of a digital convex set  $C$  are represented in four different colors. The word that codes the WN-path is  $w = 101001$ .

Our aim is to deform digital convex sets with preserving their digital convexity and our approach is based on combinatorics on words. In order to present the important characterization of digital convexity by combinatorics on words [8], we first recall the necessary notions of combinatorics on words in the following.

### 2.3. Basic notations of words

We first present some terminologies of words that can be found in [21]. An *alphabet*  $A$  is a nonempty finite set of symbols called letters; in this article, we have four letters  $0, 1, \bar{0}, \bar{1}$  as mentioned above. A word  $w$  is a sequence of concatenated letters from  $A$ . The empty word  $\epsilon$  is a sequence of zero symbols.  $A^*$  denotes the set of all finite words over  $A$ . We let  $|w|$  represent the length of a word  $w$ , while  $|w|_a$  represents the number of occurrences of  $a$  in  $w$ . For all  $a \in A$ , we thus have  $|w| = \sum_{a \in A} |w|_a$ . The  $n$ -times concatenation of  $w$  is denoted by  $w^n$ . The sub-word  $s$  of  $w$  from position  $k$  to position  $l$  is denoted by:  $s = w[k : l]$ .

A word is said *primitive* if it is not the power of a nonempty word. We say that  $w$  and  $w'$  are *conjugate*, denoted by  $w \equiv w'$ , if there exist two factors  $u, v$  in both  $w$  and  $w'$ , such that  $w = uv$  and  $w' = vu$ . The *reversal* of a word  $w = a_1 a_2 \dots a_n$  is  $\tilde{w} = a_n \dots a_2 a_1$  where each  $a_i$  is a letter. If  $\tilde{w} = w$ , the word is called a *palindrome*. In this paper, we use the total *lexicographic order*, denoted by  $<$ .

We recall that the boundary word over  $\{0, 1, \bar{0}, \bar{1}\}$  can be divided into 4 parts where each of them belongs to a different binary alphabet. For each of these parts, we give the lexicographic order between the two letters in such a way we preserve the decreasing order of the Lyndon factorization (see below for the definition). In this case, the slope of each factor will be calculated depending on which part it belongs to. Table 1 shows this information with respect to each part of the boundary word.

Words	Alphabet	Order	Slopes
$WN$	$\{0, 1\}$	$0 < 1$	$\frac{ w _1}{ w _0}$
$NE$	$\{0, \bar{1}\}$	$\bar{1} < 0$	$\frac{ w _0}{ w _{\bar{1}}}$
$ES$	$\{\bar{0}, \bar{1}\}$	$\bar{0} < \bar{1}$	$\frac{ w _{\bar{1}}}{ w _{\bar{0}}}$
$SW$	$\{\bar{0}, 1\}$	$1 < \bar{0}$	$\frac{ w _{\bar{0}}}{ w _1}$

Table 1: The alphabet, the lexicographic order and the slope's calculation used in each of the 4 parts of the boundary word

In the following, we introduce the two families of words needed in this article, Lyndon and Christoffel words.

#### 2.4. Lyndon words

We introduce the lexicographic order over the words in order to talk about the Lyndon family that was introduced by R. C. Lyndon [22].

**Property 1 ([21]).** *Let  $u$  and  $v$  be two words such that  $u > v$ . this implies that:*

$$\exists r \in A^* u = vr \text{ or } \exists m, l, e \in A^* u = m1e \text{ and } v = m0l$$

**Definition 3 ([22]).** *A word  $w$  is Lyndon if it is the smallest among its conjugates using its lexicographic order.*

We then have the following unique factorization of any word, which is called *Lyndon factorization*, introduced by Lyndon, Chen and Fox in 1958.

**Theorem 1 ([27]).** *Every non-empty word  $w$  admits a unique factorization as a lexicographical decreasing sequence of Lyndon words,  $w = \ell_1^{n_1} \ell_2^{n_2} \dots \ell_k^{n_k}$ , where every  $n_i \in \mathbb{N}$  and every  $\ell_i$  is a primitive Lyndon word such that  $\ell_1 > \ell_2 > \dots > \ell_k$ .*

A linear-time algorithm to compute the Lyndon factorization was proposed by Duval [14], while there also exists a  $O(\log n)$ -time parallel algorithm proposed by Apostolico and Crochemore [2], where  $n$  is the length of the word.

**Definition 4.** *Let  $C \subset \mathbb{Z}^2$  be a finite 4-connected, digitally convex set. The points on  $Bd(C)$  that separate different Lyndon factors of  $W(C)$  are called Lyndon points of  $C$ .*

Geometrically, Lyndon points correspond to the vertices of the convex hull of  $C$ . They are used in the next section and you can see them in Figure 1 as the green points in  $C$ .



### 2.5. Christoffel words

Christoffel words are regarded as the discretization of line segments with rational slopes.

Geometrically, a Christoffel word is determined by encoding with Freeman chain code the discretization of a line segment of rational slope [3]. More precisely, for any two non-negative co-prime integers  $a$  and  $b$ , the discretization of a line segment of rational slope  $\frac{a}{b}$  is the closest digital path below the line segment such that no integer point exists between the path and the line segment. This digital path is called the Christoffel path. If  $a$  and  $b$  are positive co-prime numbers, a Christoffel word  $w$  of slope  $\frac{a}{b}$ , denoted by  $\mathcal{C}(\frac{a}{b})$ , is a sequence of  $a+b$  letters chosen from the binary alphabet  $\{0, 1\}$ . The choice of letters is not random, but it is obtained by assigning the letter 0 (resp. 1), to each increasing (resp. decreasing) step in the sequence of all the multiples of  $a$  modulo  $(a+b)$  as given below.

**Definition 5 ([9, 27]).** *Given a pair of non-negative co-prime integers  $a$  and  $b$ , the Christoffel word  $w$  of slope  $\frac{a}{b}$  is the sequence of  $a+b$  letters of  $\{0, 1\}^*$ , such that the  $i$ -th letter of  $w$  is given by:*

$$\forall i \in \{1, \dots, n\} w[i] := \begin{cases} 0 & \text{if } r_{i-1} < r_i, \\ 1 & \text{otherwise} \end{cases}$$

where the remainder  $r_i$  is defined by

$$r_i = ia \bmod (a+b)$$

for  $0 \leq i \leq a+b$ .

The remainder sequence  $r_i$  and the  $i$ -th letter of  $\mathcal{C}(\frac{5}{8})$  are shown in Table 2.

$i$	0	1	2	3	4	5	6	7	8	9	10	11	12	13
$r_i$	0	5	10	2	7	12	4	9	1	6	11	3	8	0
$w[i]$		0	0	1	0	0	1	0	1	0	0	1	0	1

Table 2: The Christoffel word  $w$  of slope  $\frac{5}{8}$  with the remainder sequence  $(r_i)_{0 \leq i \leq 13}$ .

It is also noticed that each Christoffel word starts with a horizontal step, *i.e.* 0, and ends with a vertical step, *i.e.* 1, while the central part is a palindrome.

**Property 2 ([4]).** *Let  $\mathcal{C}(\frac{a}{b})$  be the Christoffel word of slope  $\frac{a}{b}$  with  $a$  and  $b$  two positive co-prime numbers, we can write  $\mathcal{C}(\frac{a}{b}) = 0w'1$  where  $w'$  is a palindrome.*

As the slope of a Christoffel word is exactly the number of occurrences of the letter 1 over the number of occurrences of the letter 0, the slope of each Christoffel word is defined as follow.

**Definition 6 ([6]).** Given a non-empty primitive Christoffel word  $w \in A^*$ , the morphism  $\rho(w)$  from  $A^*$  to  $\mathbb{Q}^+ \cup \{\infty\}$  is defined by

$$\rho(w) := \begin{cases} \infty & \text{if } w = 1 \\ \frac{|w|_1}{|w|_0} & \text{otherwise.} \end{cases}$$

**Property 3 ([6]).** For any two non-empty Christoffel words  $u$  and  $v$ , we have

$$u > v \Leftrightarrow \rho(u) > \rho(v).$$

For this family of words, Borel and Laubie introduced the *standard factorization* [6, 7], which allows to write any Christoffel word as the concatenation of two other Christoffel words in a unique way as follows.

**Theorem 2 ([6, 7]).** Any Christoffel word  $w$  of length greater or equal to 1 can be written in a unique way such that  $w = uv$  where  $u$  and  $v$  are both primitive Christoffel words. The couple  $(u, v)$  is called the *standard factorization* of  $w$ .

Geometrically, this factorization can be seen as the decomposition of  $w$  into  $u$  and  $v$  at the *closest* point of the path corresponding to  $w$  with respect to the line segment between the origin to  $(|w|_0, |w|_1)$ . It is equivalent to say that this factorization corresponds exactly to the position where  $r_i = 1$ . This closest point of  $w$  is denoted by  $\mathcal{cl}(w)$ . Figure 6 illustrates the standard factorization of the Christoffel word in Table 2, that is (00100101, 00101).

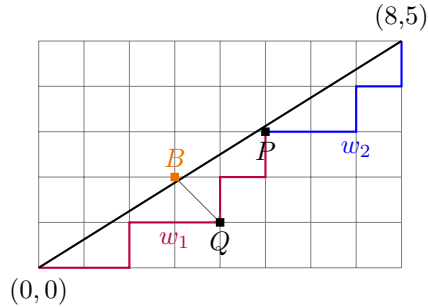


Figure 6: The standard factorization of the Christoffel word  $w$  of slope  $\frac{5}{8}$ .  $P$  corresponds to the closest point  $\mathcal{cl}(w)$  while  $Q$  corresponds to the furthest point  $\mathcal{fu}(w)$ .  $P$  is also called the interior Bézout point, while  $B$  is the exterior Bézout point. We have:  $\mathcal{C}(\frac{5}{8}) = \mathcal{C}(\frac{3}{5})\mathcal{C}(\frac{2}{5})$ .

A direct application of Theorem 2 and Property 2 induces the *palindromic factorization*, which allows to write any Christoffel word as a concatenation of two palindromes [6, 27].

**Property 4 ([6, 10, 27]).** Any primitive Christoffel word  $w$  of length strictly greater than 1 can be written in a unique way as  $w = p_1p_2$  where  $p_1$  and  $p_2$  are palindromes.

This property comes from the fact that the central part of each factor of the couple  $(u, v)$ , which forms the standard factorization of  $w$ , is a palindrome (*i.e.*  $u$  can be written as  $0u_11$  with  $u_1$  a palindrome). Namely, we have:

$$w = uv = 0u_110v_11 = 0p1 = 0v_101u_11 = p_1p_2$$

where  $p, p_1$  and  $p_2$  are palindromes (Property 2 and Theorem 2).

**Corollary 1 ([6]).** *Let  $w$  be a Christoffel word of length greater than 1. If its palindromic factorization is given by  $w = p_1p_2$ , then  $\rho(p_1) < \rho(w) < \rho(p_2)$ .*

**Corollary 2 ([13]).** *Let  $w$  be a Christoffel word of length greater than 1. Then the palindromic factorization decomposes  $w$  into  $p_1$  and  $p_2$  exactly at the furthest point of the Christoffel path of  $w$  with respect to the line segment from the origin to the grid point  $(|w|_0, |w|_1)$ , which is unique respect to  $w$ .*

This furthest point is denoted  $\mathfrak{f}u(w)$ . Note that the position of the grid point  $\mathfrak{f}u(w)$  over  $w$  is  $r_i = |w| - 1$  as seen in Figure 6. The two types of points, the closest point  $\mathfrak{c}l(w)$  and the furthest point  $\mathfrak{f}u(w)$ , on the Christoffel path are highlighted due to particular interest for the rest of the paper;  $\mathfrak{f}u(w)$  is used in Section 4 for inflation of digital convex sets.

In order to link Lyndon and Christoffel words, we need the notion of  $k$ -balanced words [23]; a word  $w \in \{0, 1\}^*$  is  $k$ -balanced if and only if for every pairs of sub-words  $s, t$  of  $w$ , we have:

$$|s| = |t| \implies ||s|_1 - |t|_1| \leq k.$$

It is shown that Christoffel words are 1-balanced [6, 7], and the following theorem connects the two families of words.

**Theorem 3 ([23]).** *A word  $w$  is a Christoffel word if and only if it is a 1-balanced Lyndon word.*

## 2.6. Digital convexity interpreted by combinatorics on words

The authors in [8] gave a characterization for the boundary word  $W(C)$  using the notions of combinatorics on words, in particular Lyndon and Christoffel words.

**Theorem 4 ([8]).** *A word of  $\{0, 1\}^*$  is WN-convex if and only if its Lyndon factorization is unique, and its factors are all primitive Christoffel words.*

Thanks to Theorem 3, if each factor of Lyndon factorization is in addition 1-balanced, we can say that the factors are also Christoffel words. From Theorem 4, we can also conclude that the slopes of the Christoffel factors are in a decreasing order.

**Corollary 3.** *If  $w = w_1^{n_1}w_2^{n_2}\dots w_k^{n_k}$  be the Lyndon factorization of the boundary word of a certain digital convex set  $C$  then  $\rho(w_1) > \rho(w_2) > \dots > \rho(w_k)$ .*

The following example shows a WN-path and the factorization of its boundary word.

**Example 1.** Let us consider a WN-convex word  $w$  whose Lyndon factorization is:

$$w = (1)(011101111)(0111)(01)(001)(00010001001)(0001).$$

The slopes  $\rho$  of the factors are decreasing:  $\frac{1}{0}, \frac{7}{2}, \frac{3}{1}, \frac{1}{1}, \frac{1}{2}, \frac{3}{8}, \frac{1}{3}$ .

In [8], we can find a linear-time algorithm over the word length that checks the WN-convexity of a path encoded by a binary alphabet. Then, Theorem 4 allows us to detect the exact position of the vertices of the convex hull of  $Bd(C)$  over the boundary word  $w$ .

**Property 5 ([8]).** Given a digital convex 4-connected set  $C$ , the vertices of the convex hull of  $C$  or its boundary  $Bd(C)$  corresponds to the Lyndon points of  $Bd(C)$ .

Figure 7 shows the result of Theorem 4 and Property 5 for the digitally convex set with boundary word  $W(C) = 10100100\bar{1}0\bar{1}\bar{1}0\bar{1}0\bar{0}\bar{1}0\bar{1}0\bar{0}\bar{1}0\bar{1}1\bar{1}0$ .

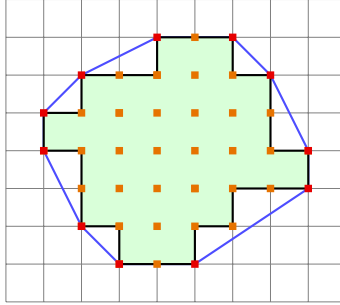


Figure 7: Let us consider the digital convex set  $C$  of the orange and red points, whose boundary  $Bd(C)$  is drawn as the integer points on the black rectilinear polygonal line. The Lyndon points of  $Bd(C)$  are colored in red. Each Lyndon point corresponds to the end of each distinct factor of the Lyndon factorization of the boundary word  $W(C) = (1)(01)(001)(0)^2(\bar{1}0)(\bar{1}\bar{1}0)(\bar{1})(\bar{0}\bar{0}\bar{1}\bar{0}\bar{1})(\bar{0})^2(\bar{1}\bar{0})(1\bar{1}\bar{0})$ . We have  $\rho(1) > \rho(01) > \rho(001) > \rho(0)$  in WN-path.

In the next sections, we show how to deflate and inflate a digital convex 4-connected set using the above tools of combinatorics on words. In such a context, we characterize removable points for deflation process in Section 3 and similarly identify insertable points for inflation in Section 4.

### 3. Deflation of digital convex sets

In this section, we first define *removable points* that allow to deflate a digital convex 4-connected set  $C$  while preserving its convexity. We then characterize such points in order to make the list of all removable points using concepts

of digital topology and combinatorics on words. They are defined along the boundary path  $Bd(C)$  via the boundary word  $W(C)$  (see Section 2.2). Applying iteratively such point-wise deflation operation may be required in various applications of computer imagery. This implies that we need for each step to choose a removable point among the list of removable points and then to update the list, which is made by the new Lyndon factorization of  $W(C)$ . We show that this update can be made locally. In practice, we need to choose one removable point in the list by using some priority. In this article, however, we do not discuss how to define such priority and simply assume that we have such a priority a priori [29]. We focus on characterization of removable points and update of the list of removable points, which are critical issues in the deflation algorithm.

### 3.1. Removable points

We first give a definition of removable points, which are used for deflating a 4-connected digital convex set  $C$  while preserving its convexity.

**Definition 7.** *Given a digitally convex, 4-connected set  $C$ , a point  $x \in C$  is removable if  $C \setminus \{x\}$  is still 4-connected and digitally convex.*

First of all, let us consider the preservation of connectivity. In digital topology, given a subset  $X \subset \mathbb{Z}^2$ , a point  $x \in X$  is said to be *simple* if deleting it from  $X$  preserves the topological characteristics of  $X$ , i.e. the number of connected components of both  $X$  and its complement in  $\mathbb{Z}^2$  [11, 20]. As we consider  $X$  to be 4-connected and digitally convex here, preserving the topological characteristics of  $X$  is equivalent to preserving the connectivity of  $X$ ;  $X$  can have no hole due to the digital convexity. Therefore, removable points must be simple if we would like to preserve the connectivity of  $C$ . Note that digital convexity does not imply connectivity (see Remark 1). The definition of simple points relies on the notion of connected components, which is a global characteristic such that the whole object must be taken into account. However, it is well known that simple points can be characterized locally [11, 25], for example, using the connectivity number defined in the 8-neighborhood [5].

Given a digital convex 4-connected set  $C$ , we recall that the boundary path  $Bd(C)$  is decomposed into four parts  $Bd_{WN}(C), Bd_{NE}(C), Bd_{ES}(C)$  and  $Bd_{SW}(C)$  whose boundary words are  $W_{WN}(C), W_{NE}(C), W_{ES}(C)$  and  $W_{SW}(C)$ , respectively. Let us consider removing a point on  $W_{WN}(C)$ , except its extremities, with preserving the convexity. Then we observe that if boundary points are removable, then the points and their neighboring points back and forth in the path  $Bd_{WN}(C)$  may form the sub-word 10 in  $W_{WN}(C)$ . Removing a point from  $C$  means switching a factor of the form 10 into 01 in  $W_{WN}(C)$  (resp.  $\bar{1}0, 0\bar{1}, \bar{1}0$  into  $0\bar{1}, \bar{1}0, 0\bar{1}$  in  $W_{NE}(C), W_{ES}(C), W_{SW}(C)$ ) as seen in Figure 8. For the extremity points, we usually use the same concept. A particular study will be considered further in this paragraph. This *switch* operator can be defined over the alphabet  $\{0, 1, \bar{0}, \bar{1}\}$  so that two consecutive letters are exchanged at a position  $k$  as follows.

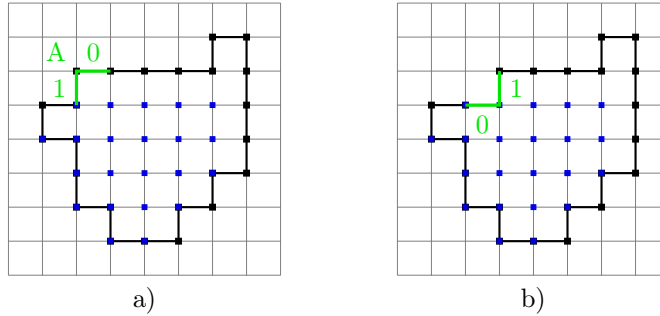


Figure 8: a) A is a point in the  $WN$ -path, where its neighboring points back and forth in  $Bd_{WN}(C)$  form the sub-word 10. b) Removing the point A corresponds to switching the sub-word 10 into 01, i.e applying the switch operator at the position of the point A

Note that any point all over  $Bd(C)$ , whose neighboring points back and forth in  $Bd(C)$  form a sub-word of the same letter can be removed. But, in this case we lose the convexity as we can see in Figure 9.

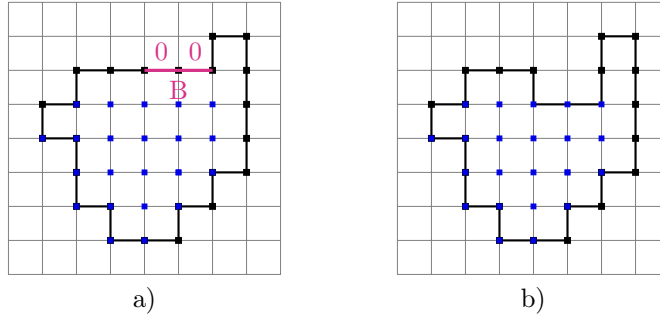


Figure 9: a) B is a point in the  $WN$ -path, where its neighboring points back and forth in  $Bd_{WN}(C)$  form the sub-word 00. b) The digital convexity is lost when removing the point B

**Definition 8 ([29]).** Given a word  $w = a_1a_2 \dots a_n$ , the switch operator at position  $k$ ,  $k < n$ , on  $w$  is defined by:

$$switch_k(w) := a_1a_2 \dots a_{k+1}a_k \dots a_n.$$

We should mention that sometimes same grid points appear twice in the boundary path  $Bd(C)$  so that the boundary word  $W(C)$  contains two consecutive letters of opposite directions. They are always positioned at the junction of two of four decomposed paths,  $Bd_{WN}(C)$ ,  $Bd_{NE}(C)$ ,  $Bd_{ES}(C)$  and  $Bd_{SW}(C)$ , such that the two consecutive letters are in different paths. We can also remove such a grid point, so that we replace the sub-word consisting of two letters of opposite directions by  $\epsilon$ , as seen in Figure 10.

Over  $W(C)$ , we can find several points that form with their neighboring points either sub-word  $10, 0\bar{1}, \bar{1}0, \bar{0}1, 0\bar{0}, 1\bar{1}, \bar{0}\bar{0}$  or  $\bar{1}\bar{1}$ . According to Definition

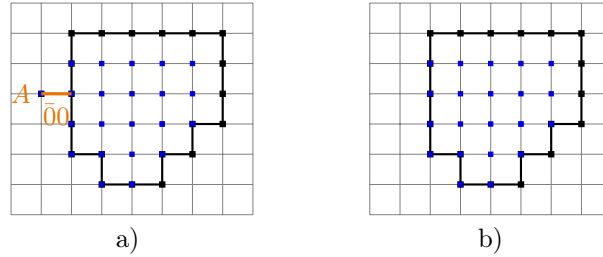


Figure 10: a) The digital set  $C$  whose boundary word is:  $W(C) = 01100000\bar{1}\bar{1}\bar{1}\bar{0}\bar{1}\bar{0}\bar{1}\bar{0}\bar{0}\bar{1}\bar{0}\bar{1}\bar{1}\bar{0}$ . The point  $A \in C$  creates the sub-word  $\bar{0}\bar{0}$  in  $W(C)$ . b) Removing  $A$  means replacing this factor by  $\epsilon$ .

4, they correspond to Lyndon points; we recall that Lyndon points are geometrically the vertices of  $\text{Conv}(C)$ , each of which is at the end of each factor of the Lyndon factorization of  $W(C)$  according to Property 5. However, not all of Lyndon points are removable. Indeed, the switch operation on such factors may lead to losing the connectivity. To avoid this problem, Theorem 5 gives the characterization of removable points for any  $C$ . In order to prove this theorem, we recall that the Christoffel words, which are the one balanced Lyndon words of the Lyndon factorization of  $W(C)$  (Theorem 4), have the following form:  $0u1, \bar{1}k0, \bar{0}\bar{\ell}\bar{1}$  and  $1m\bar{0}$ , where  $u, k, \ell$  and  $m$  are palindromes in  $W_{WN}(C), W_{NE}(C), W_{ES}(C)$  and  $W_{SW}(C)$  (Property 2).

**Theorem 5.** *Given a digital convex 4-connected set  $C$  of  $\mathbb{Z}^2$ , a point  $x \in C$  is removable if and only if  $x$  is a simple point with respect to  $C$  and a Lyndon point of the boundary  $\text{Bd}(C)$ .*

*Proof:* Let us consider the boundary word  $W(C)$  of  $C$ , which is decomposed by the Lyndon factorization such that  $W(C) = \ell_1^{n_1} \ell_2^{n_2} \dots \ell_s^{n_s}$ . Since  $C$  is digitally convex, each  $\ell_i$ ,  $1 \leq i \leq s$ , is a Christoffel word (Theorems 3 and 4). We give the proof only for the binary sub-word  $W_{WN}(C)$  of  $W(C)$  as follows; the similar proofs are found to the three other sub-words  $W_{NE}(C), W_{ES}(C)$  and  $W_{SW}(C)$ . As mentioned above, removing a point from the boundary means applying the switch operator at a Lyndon point, so that the corresponding sub-word  $10$  is replaced by  $01$ , or  $\bar{0}\bar{0}, \bar{1}\bar{1}$  by  $\epsilon$ . The simplicity of a point  $x$  guarantees the 4-connectivity of  $C \setminus \{x\}$ . If a point is not a Lyndon point, i.e. the boundary sub-word belongs to one of the Lyndon factors  $\ell_i$  of  $W_{WN}(C)$ , it cannot be removed as we lose  $WN$ -convexity. Let us consider any boundary sub-word  $10$  such that  $1$  appears at the end of one of the Lyndon factors in  $W_{WN}(C)$  and  $0$  appears at the beginning of the consecutive Lyndon factor. If these two Lyndon factors are identical, then switching the pair  $10$  makes us lose the  $WN$ -convexity. Now, let us focus on sub-words  $10$  that are obtained by two consecutive distinct Lyndon factors  $\ell_i \ell_{i+1}$  of  $W_{WN}(C)$ . By Property 2, we can write  $\ell_i = 0u1$  and  $\ell_{i+1} = 0v1$  where  $u$  and  $v$  are both palindromes, so that  $\ell_i \ell_{i+1} = 0u10v1$ . By applying the switch operator on  $\ell_i \ell_{i+1}$  at position  $|\ell_i|$ , we obtain:  $\text{switch}_k(\ell_i \ell_{i+1}) = 0u01v1$ , which is also a Christoffel word if  $u01v$  is

a palindrome, as seen in Property 4; otherwise, we get several other Christoffel words. This effect will be studied and detailed in the next subsection where we give the updates effect on the sub-word after removing a point (see also the proof of Theorem 1 in [29]). Finally, when the corresponding sub-word of the point to remove is  $0\bar{0}$  (resp.  $1\bar{1}$ ), this means that we are at the intersection between the  $SW$  and  $WN$  paths (resp.  $WN$  and  $NE$  paths). Replacing this sub-word by  $\epsilon$  means that the position of the point  $W$  will change and the  $Bd_{WN}(C)$  will be longer and starts from the first letter after the last  $\bar{0}$  in  $Bd_{SW}(C)$ .  $\square$

Property 5 provides the geometrical interpretation of Theorem 5: any vertex of the convex hull of  $C$  corresponds to a Lyndon point on  $Bd(C)$ . We note that detecting all the removable points for  $C$  is performed in linear complexity with respect to  $|W(C)|$  as it is done via Lyndon factorization of  $W(C)$ .

### 3.2. Updating the removable points

Given a digital convex 4-connected set  $C$  of  $\mathbb{Z}^2$ , the first iteration of our deflation process of  $C$  is started by making the list of all the possible removable points with respect to  $C$ . Once a candidate has been chosen among them, the boundary word  $W(C)$  is modified and the Lyndon factorization is made again. This means that new candidates can arise or disappear from the list of removable points.

**Example 2.** Let  $w_1 = C(\frac{2}{5})C(\frac{1}{3})$ ,  $w_2 = C(\frac{3}{7})C(\frac{1}{3})$  and  $w_3 = C(\frac{14}{19})C(\frac{4}{17})$  be three different factors of the boundary word  $W(C)$  of a digital convex 4-connected set  $C$ . As proved before, the removable points with respect to  $C$  are simple and Lyndon points. From the definition, the Lyndon points of each  $w_i$  are positioned at the joint between consecutive distinct Christoffel words. If we apply the switch operator on each  $w_i$  at the Lyndon point, we obtain:  $switch_7(w_1) = C(\frac{3}{8})$ ;  $switch_{10}(w_2) = (C(\frac{2}{5}))^2$  and  $switch_{33}(w_3) = C(\frac{11}{15})C(\frac{2}{3})C(\frac{1}{2})C(\frac{1}{4})^4$ . These examples show that Lyndon points can disappear or newly appear due to the switch operation (see Figure 11).

The above deflation step is iterated in general, and we remark that several removable point candidates exist at each iteration step. Various heuristics can be considered to choose one of the removable points [29]. It is also important to know what happen after each iteration because of the update of the Lyndon factorization. Property 6 shows all the possible effects that arise after removing a point. We can also see that the update of removable points is made locally, which means that the Lyndon factorization before the factor  $u$  and after the consecutive factor  $v$  is unchanged if the Lyndon point is removed at the joint of  $u$  and  $v$ .

**Property 6 ([29]).** Let the binary word  $w = uv$  such that  $u$  and  $v$  are two Christoffel words with  $\rho(u) > \rho(v)$ . By applying the switch operator to  $w$  at position  $|u|$ , we obtain  $switch_{|u|}(w) = m_1^{\ell_1} \dots m_k^{\ell_k}$ ,  $k \geq 1$  such that  $\rho(u) > \rho(m_1) > \dots > \rho(m_k)$ .



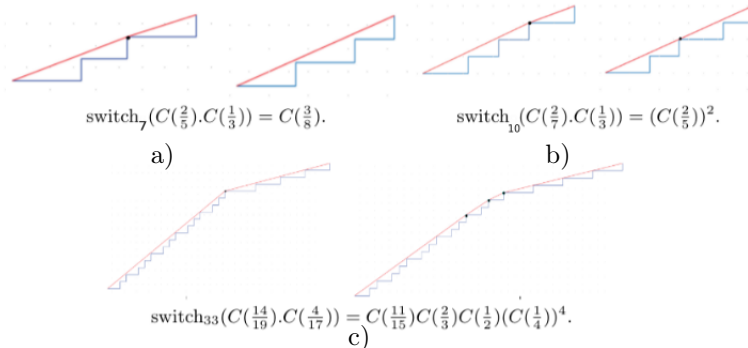


Figure 11: Before and after the pixelwise deflation of (a)  $w_1$ , (b)  $w_2$  and (c)  $w_3$  and their Lyndon factorization (left and right); a) no removable point remains after the update; b) a removable point arises at a different position; c) several new removable points appear.

### 3.3. Algorithm

We give a small recap about all the previous results considered in a combinatorics on words point of view. In order to deflate a digital convex 4-connected set  $C$ , we must start the procedure by considering the boundary word  $w$  of  $C$ . By Property 5, it is sufficient to apply the Lyndon factorization on  $w$  in order to get the positions of all the Lyndon points. At each step of the deflation procedure, several points are considered as good candidates. In order to know which point to be removable, we need to be sure that this point is also a simple point thanks to Theorem 5. We then choose a certain heuristic in order to stock these candidates in a priority queue. Once a point is selected from the queue, we apply the switch operator at its position, given in Definition 8. Some local updates must be considered over the Lyndon factorization. This will certainly affect the priority queue of removable points, which requires an update too since the Lyndon points may vary. This procedure must be iterated  $k$  times if we would like to remove  $k$  points from  $C$  while preserving digital convexity. Algorithm 1 represents such an iterative pointwise deflation over a 4-connected digital convex set  $C$ .

Figure 12 shows the deflation procedure applied on a digital convex 4-connected set using the heuristic of area-change, which minimize or maximize the area difference between the convex hull of  $C$  and  $C \setminus \{p\}$  (see [29] for more details).

## 4. Inflation of digital convex sets

In this section, we study the pixelwise inflation of any digital convex 4-connected set  $C$  with preservation of digital convexity, which is more complex than the deflation process as seen in the previous section. We first give a necessary condition and then a sufficient condition for characterizing insertable

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**Algorithm 1** Point-wise deflation

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**Input:** a digital convex 4-connected set  $C$ , a number of removing points  $k$

**Output:** a sequence of removed points  $R$

- 1: Compute the Lyndon Factorization  $F$  of the boundary word of  $C$
  - 2: Insert all the Lyndon points in a priority queue  $L$
  - 3: **while**  $k \geq 0$  **do**
  - 4:   Pull the highest-priority simple point  $p$  from  $L$  and add  $p$  to  $R$
  - 5:   Let  $u$  and  $v$  be the two distinct Christoffel words of  $F$  around  $p$
  - 6:   Compute  $w = \text{switch}_{|u|}(uv)$
  - 7:   Compute the Lyndon Factorization of  $w$  and update  $F$  and  $L$
  - 8:   Decrement  $k$
  - 9: **end while**
  - 10: **return**  $R$
- 

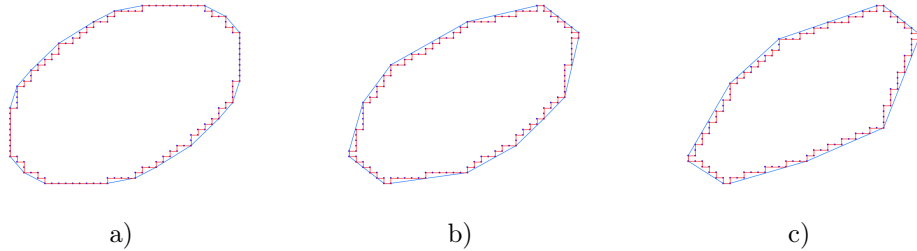


Figure 12: The deflation process of a digital convex 4-connected set (a) represented after 150 and 250 iterations respectively in (b) and (c) using the heuristic of area-change.

points for such inflation. We will show that such insertable points cannot be at any arbitrary place around  $C$ . Here the key process, if such inflation is applied iteratively, is also the update of Lyndon factorization over the boundary word  $W(C)$ . The question raised is that the induced factorization can be made locally and efficiently.

Here, we focus on furthest points, each of which is uniquely defined for every Lyndon factor of the boundary word  $W(C)$ , with length greater than 1 (see Corollary 2). Lyndon factors of length 1 have no furthest point, and this case will be treated differently later on in this section. In order to inflate  $C$ , we need to determine the *insertable point*, which are necessarily furthest points indeed. In order to verify the insertability of a furthest point, we have to take in consideration the local and global effects on  $W(C)$  after adding this point, which requires the Lyndon factorization update of  $W(C)$ . In fact, we might lose the digital convexity after adding a furthest point if we don't choose the right one. We characterize insertable points with some iterative local digital convexity verification. In order to avoid such an expensive iterative verification, we also propose some strong constraints to locally characterize a subset of insertable points.

#### 4.1. Definition of insertable points

We define insertable points, which are used for pixelwise inflation of a digital convex 4-connected set  $C$  without loosing its digital convexity.

**Definition 9.** *Given a digital convex 4-connected set  $C$ , a point  $x \in \bar{C}$  is insertable if  $C \cup \{x\}$  is still 4-connected and digitally convex.*

In this section and unlike Section 3, there is no need to verify if an insertable point refers to a simple point or not. We only need to verify the digital convexity as the simple 4-connectedness of  $C \cup \{x\}$  is kept. We recall that the Lyndon factorization of the boundary word  $W(C)$  is of the following form:  $W(C) = \ell_1^{n_1} \dots \ell_s^{n_s}$  where all  $\ell_i$  are made of primitive Christoffel words of length greater or equal to 1, with the decreasing slope order, and every  $n_i$  is a positive integer. If  $n_i > 1$ , this signifies that a Christoffel word  $\ell_i$  is repeated  $n_i$  times in  $W(C)$ . To find insertable points, we need to consider the following two cases for each Lyndon factor  $\ell_i$  of  $W(C)$ :  $|\ell_i| > 1$  and  $|\ell_i| = 1$ .

#### 4.2. Link between insertable points and furthest points

In this part, we show the link between insertable points and furthest points associated to primitive Christoffel words. Let us call the diagonal opposite point in  $\bar{C}$  of each furthest point in  $C$  *the closest outer point*<sup>2</sup>. We show that for each primitive Christoffel word of length strictly greater than 1, we can find a unique position in the boundary word  $W(C)$  where a point is added at its closest outer point. We first present the definition of the *split operator* applied on primitive Christoffel words, proposed in [13]. This operator helps us to see the local modification over the boundary word  $W(C)$ , which can influence the Lyndon factorization of  $W(C)$  globally sometimes.

##### 4.2.1. Necessary condition for insertable points

Adding a point to a digital convex 4-connected set  $C$  at neither its vertical nor horizontal parts of the bounding box of  $C$  correlates, in a viewpoint of combinatorics on words, to applying the switch operator, given in Definition 8, to a sub-word of form 01 (resp.  $\bar{1}0$ ,  $0\bar{1}$ ,  $1\bar{0}$ ) over the boundary word  $W_{WN}(C)$  (resp.  $W_{NE}(C)$ ,  $W_{ES}(C)$ ,  $W_{SW}(C)$ ); on the vertical and horizontal sides, the letters  $0\bar{1}0$  (resp.  $10\bar{1}$ ,  $0\bar{1}0$ ,  $\bar{1}0\bar{1}$ ) replaces 1 (resp. 0,  $\bar{1}$ ,  $\bar{0}$ ) instead. We aim at inflating  $C$  without loosing its digital convexity. This means that after adding a point, the updated Lyndon factorization of  $W(C)$  must remain made of factors of Christoffel words whose slopes are in a decreasing order, as seen in Corollary 3. In general, this condition is not satisfied if we add a point randomly. We will show that any insertable point must be a closest outer point, which is the diagonally opposite point of a furthest point. Note that the reverse is not always true.

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<sup>2</sup>This point corresponds exactly to the exterior Bézout point of a line segment in  $\mathbb{Z}^2$ .

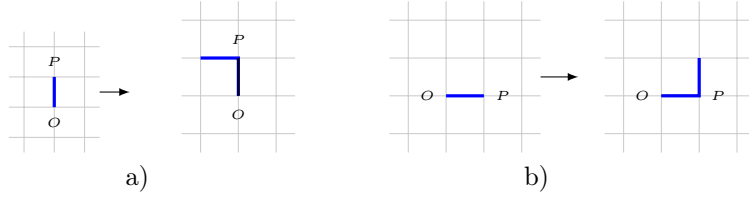


Figure 13: Split operator applied to Christoffel paths of slopes (a)  $\infty$  so that  $split(1) = (1\bar{0}, 0)$ , and (b)  $0$  such that  $split(0) = (01, \bar{1})$ .

**Definition 10 ([13]).** *The split operator applied on a primitive Christoffel word  $w$  decomposes  $w$  into two Christoffel words  $w^+$  and  $w^-$ , defined as  $split(w) := (w^+, w^-)$  such that:*

- when  $|w| \geq 2$ ,

$$(w^+, w^-) := (w'[1 : k], w'[k + 1 : |w|])$$

where  $w' = switch_k(w)$ ; with  $k$  the position of the furthest point of  $w$ ,

- when  $|w| = 1$ ,

$$(w^+, w^-) := \begin{cases} (1\bar{0}, 0) & \text{if } w = 1, \\ (01, \bar{1}) & \text{if } w = 0, \\ (\bar{1}0, \bar{0}) & \text{if } w = \bar{1}, \\ (\bar{0}\bar{1}, 1) & \text{if } w = \bar{0}, \end{cases}$$

In Figure 13, we show examples of the case that the length of a Christoffel word is equal to 1. Note that, if the Christoffel word  $w$  is not primitive, i.e  $w = \ell_i^{n_i}$  with  $n_i > 1$ , we can apply the split operator on any of these  $\ell_i$ .

Corollary 2 shows that the split operator applied at the furthest point position of a Christoffel word of length greater than 1 gives two other Christoffel words. These two words are of decreasing order with respect to their slopes. Lemma 1 also shows that the concatenation of these two new factors gives the same result for the switch operator, defined previously, when applied at the same position. This can be seen in Example 3 and illustrated in Figure 14.

**Lemma 1 ([13]).** *Given a Christoffel word  $w$  with  $|w| > 1$  such that  $w = uv$ , and  $u$  and  $v$  are the factors of the standard factorization of  $w$ . Then we have*

$$split(w) = (v, u)$$

where  $\rho(v) > \rho(u)$ .

From Theorem 2, we know that,  $u$  and  $v$ , the two factors of the standard factorization of  $w$ , are both primitive Christoffel words. Lemma 1 then shows that the result of the split operator exactly consists of the two primitive Christoffel words  $v$  and  $u$ , which are in the reverse order of the standard factorization.

**Example 3.** (Example of the split operator) Let  $w = 00100100101$  be the Christoffel word of slope  $\frac{4}{7}$  with its standard factorization  $w = w^-w^+$  where  $w^- = \mathcal{C}(\frac{1}{2})$ ,  $w^+ = \mathcal{C}(\frac{3}{5})$  and  $\rho(w^-) < \rho(w^+)$ . By applying the split operator on  $w$ , we obtain

$$\text{split}(w) = (w^+, w^-) = (00100101, 001).$$

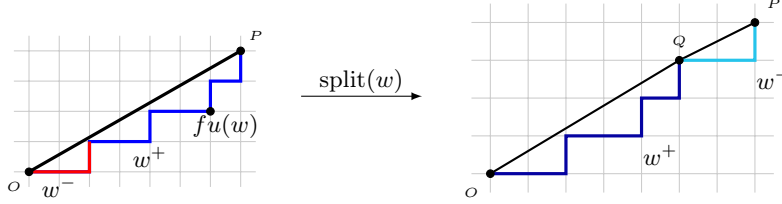


Figure 14: Christoffel word (left) whose standard factorization is  $w = w^-w^+$ , and the result of the split operator (right),  $\text{split}(w) = w^+w^- = \text{switch}_8(w)$ .

Property 4 ensures the uniqueness of the furthest point for each Christoffel word. Based on that, it is shown in [13] that any Christoffel word of length greater than 1 can only be split at this position. In other words, if we need to add a point from  $\bar{C}$  to  $C$  in order to inflate a certain digital line segment of the boundary of  $C$ , we must choose a closest outer point associated to the digital line segment.

So far, we show how to split one Christoffel word with or without multiplicity; Lemma 1 ensures that the order of the slopes for the new Christoffel words, obtained after splitting, is still decreasing. The question that arises next is the following: will this order be preserved also all around the boundary word?. In other words, does this operation affect the order of the slopes around a chosen Christoffel word to be split? These questions are answered in the following part. They help us to give the characterization of insertable points.

#### 4.2.2. Example of closest outer points that are not insertable

Let us consider a digital convex 4-connected set  $C$  and its boundary word whose Lyndon factorization is give by  $W(C) = \ell_1^{n_1} \dots \ell_s^{n_s}$ . Each factor  $\ell_i^{n_i}$  represents one of the polygonal line segments of  $\text{Conv}(C)$  and each  $\ell_i$  is a primitive Christoffel word. As we have seen, in order to add a point around  $C$ , we choose one of the factors of  $W(C)$  together with its closest outer point. This must be done by respecting the conditions given in Definition 10 and Lemma 1. Two examples are illustrated in Figure 15

During the inflation process, we can also get the case where we can completely loose the convexity property. In this case, we know that this closest outer point, can not be chosen as an insertable point. This is illustrated in Example 4.

**Example 4.** (Example of a closest outer point that is not insertable) Let  $w_1 = \mathcal{C}(\frac{30}{41})$  and  $w_2 = \mathcal{C}(\frac{5}{7})$  be two consecutive Christoffel words on the boundary

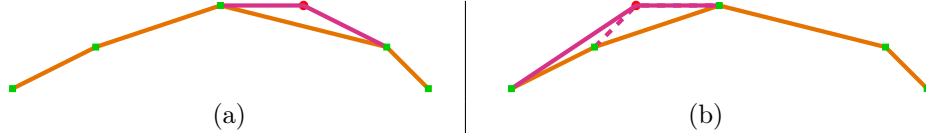


Figure 15: Figure (a) shows that the inflation at this position maintains the convexity. Figure (b) shows that there is an additional step of propagation to verify the convexity on the left side of the segment.

word of a certain digital convex 4-connected set. From Theorem 4, we have  $\rho(w_1) > \rho(w_2)$ . If we apply the split operator on  $w_2$ , i.e. add the closest outer point of  $w_2$ , we loose the digital convexity while it is not the case if we apply it on  $w_1$ :

$$\begin{aligned} \text{split}(w_1)w_2 &= \left( \mathcal{C}\left(\frac{11}{15}\right), \mathcal{C}\left(\frac{19}{26}\right) \right) \mathcal{C}\left(\frac{5}{7}\right) \\ w_1 \text{split}(w_2) &= \mathcal{C}\left(\frac{30}{41}\right) \left( \mathcal{C}\left(\frac{3}{4}\right), \mathcal{C}\left(\frac{2}{3}\right) \right). \end{aligned}$$

This example indicates that a closest outer point is not always insertable. In the following, we study sufficient conditions for insertable points in details.

#### 4.2.3. Characterization of insertable points and Lyndon factorization update

As seen before, not all closest outer points correspond to insertable points. In order to verify such insertability, we need to update, after adding a closest outer point, the Lyndon factorization over the boundary word  $W(C)$ . As mentioned before, we might need to do some propagation for certain cases. This propagation can be made on the right and/or left side of the Christoffel word where we split, and can reach the beginning or the end of the sub-word  $W_{WN}(C)$ ,  $W_{NE}(C)$ ,  $W_{ES}(C)$  or  $W_{SW}(C)$ , in the worst case. We give the characterization of an insertable point in Theorem 6. The proof will be given at the end of this section since we need to show some notions before. Definition 11 introduces two types of insertability verification for grid points, one on the left side and the other one on the right side.

**Definition 11.** Given a digital convex 4-connected set  $C$ , let us consider the boundary word  $W(C)$  and its Lyndon factorization  $W(C) = \ell_1^{n_1} \dots \ell_m^{n_m}$ . Let  $x$  be the closest upper point in  $\overline{C}$  of the  $j$ -th Lyndon factor  $\ell_i$  for  $i \in [1, m]$  and  $j \in [1, n_i]$  in  $W(C)$  such that  $\text{split}(\ell_i) = (\ell_i^+, \ell_i^-)$ , we say that:

- $x$  is insertable on the left if  $\exists k \in \mathbb{Z}^*$ ,  $\ell_{i-k-1} \geq L_k$  such that for every  $h \leq k$ ,  $L_h$  is recursively defined by

$$L_h := \begin{cases} \ell_i^{j-1} \ell_i^+ & \text{for } h = 0 \\ \ell_{i-h}^{n_{i-h}} L_{h-1} & \text{for } h \geq 1 \text{ if } \exists m_{h-1} \in \mathbb{Z}^+, \ell_{i-h} = \ell_{i-h-1} L_{h-1}^{m_{h-1}} \end{cases}$$

- $x$  is insertable on the right if  $\exists k \in \mathbb{Z}^*, \ell_{i+k+1} \leq R_k$  such that for every  $h \leq k$ ,  $R_h$  is recursively defined by

$$R_h := \begin{cases} \ell_i^- \ell_i^{n_i-j} & \text{for } h = 0 \\ R_{h-1} \ell_{i+h}^{n_{i+h}} & \text{for } h \geq 1 \text{ if } \exists m_{h-1} \in \mathbb{Z}^+, \ell_{i+h} = R_{h-1}^{m_{h-1}} \ell_{i+h-1} \end{cases}$$

Theorem 6 is one of the main results of this paper. It provides the characterization of the insertability of such a closest upper point  $x \in \overline{C}$ .

**Theorem 6.** *Given a digital convex 4-connected set  $C$ , let  $x$  be the closest upper point in  $\overline{C}$  with respect to one of the Lyndon factors of the boundary word  $W(C)$ . Then,  $x$  is insertable if and only if  $x$  is insertable on both the left and right sides.*

Theorem 6 characterizes insertable points. Using this characterization, we can obtain the positions of all the insertable points to inflate a digital convex set by preserving its digital convexity. Once one of these eligible candidates is chosen and added, we update  $W(C)$ , which also updates the list of insertable points for the next step of inflation. Theorem 6 and Definition 11 show all the possible cases we can face after adding an insertable point. In order to prove this theorem, we first need to define the following morphism that maps the set of Christoffel words to the set of binary words.

**Definition 12 ([7]).** *Given an ordered pair of Christoffel words over  $\{0,1\}^*$   $B = (\mathcal{C}(\frac{a}{b}), \mathcal{C}(\frac{c}{d}))$ , we define the Christoffel morphism  $\Theta_B$  from the set of Christoffel words to  $A^*$  such that  $\Theta_B(0) = \mathcal{C}(\frac{a}{b})$  and  $\Theta_B(1) = \mathcal{C}(\frac{c}{d})$ .*

Being a morphism,  $\Theta_B(uv) = \Theta_B(u)\Theta_B(v)$ , and it is ordered as we can see in the following property.

**Property 7 ([7]).** *If  $\mathcal{C}(\frac{a}{b}) < \mathcal{C}(\frac{c}{d})$ , then the Christoffel morphism  $\Theta_B$  with  $B = (\mathcal{C}(\frac{a}{b}), \mathcal{C}(\frac{c}{d}))$  is an increasing morphism. In other words, for any two Christoffel words  $w_1$  and  $w_2$  such that  $w_1 < w_2$ , we have:  $\Theta_B(w_1) < \Theta_B(w_2)$ .*

**Example 5.** *Let  $B = (\mathcal{C}(\frac{3}{5}), \mathcal{C}(\frac{2}{3}))$ ,  $w_1 = \mathcal{C}(\frac{3}{4})$  and  $w_2 = \mathcal{C}(\frac{3}{2})$ . From the lexicographic order, we have that  $\mathcal{C}(\frac{3}{5}) < \mathcal{C}(\frac{2}{3})$  and  $w_1 < w_2$ . By applying Definition 12 and Property 7, we get the following words:*

$$\begin{aligned} \Theta_B(w_1) &= \mathcal{C}\left(\frac{3}{5}\right)^2 \mathcal{C}\left(\frac{2}{3}\right) \mathcal{C}\left(\frac{3}{5}\right) \mathcal{C}\left(\frac{2}{3}\right) \mathcal{C}\left(\frac{3}{5}\right) \mathcal{C}\left(\frac{2}{3}\right) \\ \Theta_B(w_2) &= \mathcal{C}\left(\frac{3}{5}\right) \mathcal{C}\left(\frac{2}{3}\right) \mathcal{C}\left(\frac{3}{5}\right) \mathcal{C}\left(\frac{2}{3}\right)^2 \end{aligned}$$

so that  $\Theta_B(w_1) < \Theta_B(w_2)$ .

Lemma 2 studies the effect of the split operator of a certain Christoffel word with multiplicity higher than 1.

**Lemma 2.** *Let us consider a  $WN$ -convex word  $w$  and its Lyndon factorization  $w = \ell_1^{n_1} \dots \ell_s^{n_s}$ . For any  $j \in [1, s]$  such that  $|\ell_j| > 1$ , if we split the  $i$ -th element of  $\ell_j^{n_j}$  where  $1 \leq i \leq n_j$ , we obtain  $split(\ell_j) = (\ell_j^+, \ell_j^-)$  and the factors  $\ell_j^{i-1} \ell_j^+$  and  $\ell_j^- \ell_j^{n_j-i}$  are two decreasing Christoffel words, i.e  $\ell_j^{i-1} \ell_j^+ > \ell_j^- \ell_j^{n_j-i}$ .*

*Proof:* First, we prove that the two factors  $\ell_j^{i-1} \ell_j^+$  and  $\ell_j^- \ell_j^{n_j-i}$  are Christoffel words. This is obtained by applying the Christoffel morphism  $\Theta_A$  over  $A = (\ell_j^-, \ell_j^+)$  on the words  $(01)^{i-1}1$  and  $0(01)^{n_j-i}$  respectively. Second, we must prove that  $\ell_j^{i-1} \ell_j^+ > \ell_j^- \ell_j^{n_j-i}$ . This inequality comes from the fact that  $\Theta_A$  is an increasing morphism, as seen in Property 7; since  $\ell_j^- < \ell_j^+$  and  $0(01)^{n_j-i} < (01)^{i-1}1$ , we have  $\Theta_A((01)^{i-1}1) < \Theta_A(0(01)^{n_j-i})$ .  $\square$

When the length of a Christoffel word is equal to one, the split operator gives two Christoffel words with different binary alphabets such as  $split(1) = (1\bar{0}, 0)$ . These two new Christoffel words can not be comparable between each other as the two Christoffel words do not belong to the same part of the four parts of the boundary path (ex.  $WN$ -path). When they belong to different parts of the boundary word of a digital convex 4-connected set  $C$ , the right insertability of the associated point is verified in the initial binary boundary sub-word while the left insertability is verified in the previous one. In the case of  $split(1) = (1\bar{0}, 0)$ , the right insertability verification is made in  $W_{WN}(C)$  while the left one is made in  $W_{SW}(C)$ .

To simplify the proof of Theorem 6, we treat only the right insertability as the left one can be proved similarly. We consider all the possible situations that can arise after applying the split operation.

**Proof of Theorem 6:**

Let us consider the boundary word  $W(C)$  and its Lyndon factorization  $W(C) = \ell_1^{n_1} \dots \ell_m^{n_m}$ . Let  $x$  be the closest upper point in  $\bar{C}$  of the  $j$ -th Lyndon factor  $\ell_i$  for  $i \in [1, m]$  and  $j \in [1, n_i]$  in  $W(C)$  such that  $split(\ell_i) = (\ell_i^+, \ell_i^-)$ . From Lemma 2, we know that  $\ell_i^- \ell_i^{n_i-1}$  is a Christoffel word and that  $\ell_i^+ > \ell_i^- \ell_i^{n_i-1}$ . Let us consider that  $R_{i+1}$  is in  $W_{WN}(C)$ . To check the  $WN$ -convexity, we must compare  $R_0 = \ell_i^- \ell_i^{n_i-1}$  to  $\ell_{i+1}$ . We get the following cases:

1.  $\ell_{i+1} < R_0$ , so no digital further verification is needed and the digital convexity is maintained.
2.  $\ell_{i+1} = R_0$ , the convexity is also preserved in this case since the line segments discretized by  $R_0$  and  $\ell_{i+1}^{n_{i+1}}$  will be aligned, and the multiplicity of  $\ell_{i+1}$  is increased by 1 in the factorization of  $w$ .
3. If  $\ell_{i+1} = R_0^{m_i} \ell_i$ , then  $R_1 = R_0 \ell_{i+1}^{n_{i+1}}$  is set and  $R_1$  is now compared with  $\ell_{i+2}^{n_{i+2}}$  using the same reasoning. This propagation is kept until the right insertability verification is done with one of the cases 1, 2 and 4 or until the end of  $W_{WN}(C)$ . Hence the sequence  $R_k$  is constructed.
4. If  $\ell_{i+1} > R_0$  and  $\ell_{i+1} \neq R_0^{m_i} \ell_i$ , then we lose the decreasing order, namely the  $WN$ -convexity. Hence, the point  $x$  is not insertable.



A similar proof can be given for the left insertability. Note that if the length of  $\ell_i$  is equal to one, the left insertability should be verified in the previous binary boundary sub-word.  $\square$

Note that the propagation test is limited by the extremity of the  $WN$  side. In Figure 16, we show an example of the propagation following Theorem 6.

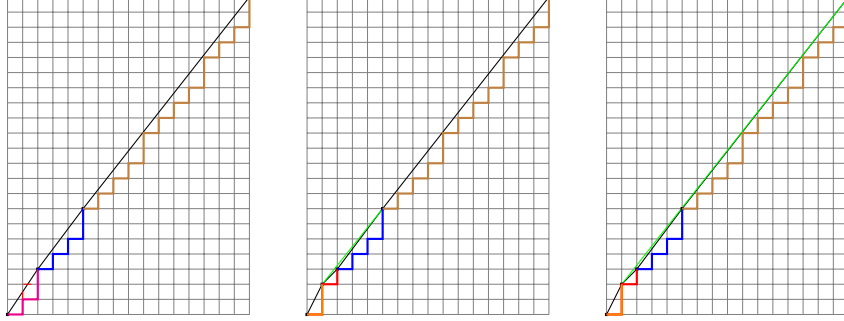


Figure 16: Procedure of insertability verification on the right with propagation: the closest upper point  $(1, 2)$  of  $\ell_i$  ( $\ell_i$  in pink) is inserted (left); as  $\ell_{i+1} \leq R_0$  ( $\ell_{i+1}$  in blue and  $R_0 = \ell^-$  in red) is not satisfied but we have  $\ell_{i+1} = R_0 \ell_i$ , we obtain  $R_1 = R_0 \ell_{i+1}$  ( $R_1$  in green (=red+blue)) (center); as  $\ell_{i+2} \leq R_1$  ( $\ell_{i+2}$  in brown) is not satisfied but we have  $\ell_{i+2} = R_1^2 \ell_{i+1}$ , we obtain  $R_2 = R_1 \ell_{i+2}$  ( $R_2$  in green (=red+blue+brown)) (right).

With Theorem 6, we have the characterization of the insertable pixels over  $W(C)$ . We give in Example 6, 7 and 8, several numerical cases showing the inflation process.

**Example 6.** Let us consider a  $WN$ -convex word  $w = (\mathcal{C}(\frac{3}{4}))^4 \mathcal{C}(\frac{3}{5}) \mathcal{C}(\frac{1}{2})$  and the split of the second factor. As  $\text{split}(\mathcal{C}(\frac{3}{5})) = (\mathcal{C}(\frac{2}{3}), \mathcal{C}(\frac{1}{2}))$ , we obtain the new  $WN$ -convex word  $w'$ ;

$$w' = \left( \mathcal{C}\left(\frac{3}{4}\right) \right)^4 \mathcal{C}\left(\frac{2}{3}\right) \left( \mathcal{C}\left(\frac{1}{2}\right) \right)^2.$$

**Example 7.** Let us consider a  $WN$ -convex word  $w = (\mathcal{C}(\frac{3}{4}))^4 (\mathcal{C}(\frac{3}{5}))^3 \mathcal{C}(\frac{1}{2})$  and the split of the second word of the second factor  $\ell_2$ . As  $\text{split}(\mathcal{C}(\frac{3}{5})) = (\mathcal{C}(\frac{2}{3}), \mathcal{C}(\frac{1}{2}))$  and  $\ell_3 = \ell_2^-$ , we obtain:

$$w' = \left( \mathcal{C}\left(\frac{3}{4}\right) \right)^4 \mathcal{C}\left(\frac{3}{5}\right) \mathcal{C}\left(\frac{2}{3}\right) \mathcal{C}\left(\frac{1}{2}\right) \mathcal{C}\left(\frac{3}{5}\right) \mathcal{C}\left(\frac{1}{2}\right)$$

whose Lyndon factorization is  $(\mathcal{C}(\frac{3}{4}))^4 \mathcal{C}(\frac{5}{8}) \mathcal{C}(\frac{4}{7}) \mathcal{C}(\frac{1}{2})$

**Example 8.** Let us consider a  $WN$ -convex word  $w = (\mathcal{C}(\frac{3}{4}))^4 (\mathcal{C}(\frac{3}{5}))^5 \mathcal{C}(\frac{24}{41})$ , so that  $\ell_3 = (\ell_2^- \ell_2^2)^3 \ell_2$ , and consider splitting the third word of the second factor

$\mathcal{C}(\frac{3}{5})$ . We get:

$$w' = \left( \mathcal{C}\left(\frac{3}{4}\right) \right)^4 \mathcal{C}\left(\frac{2}{3}\right) \mathcal{C}\left(\frac{13}{22}\right) \mathcal{C}\left(\frac{24}{41}\right)$$

whose Lyndon factorization is  $(\mathcal{C}(\frac{3}{4}))^4 \mathcal{C}(\frac{8}{13}) \mathcal{C}(\frac{31}{53})$ .

#### 4.3. Strong sufficient condition to insertable points

Given a digital convex 4-connected set  $C$ , till now, for each closest upper point of  $C$ , we must check the conditions given in Theorem 6 in order to know if it is an insertable point or not. However, checking if a point is insertable, using Theorem 6 may require processing the whole boundary word  $W(C)$  due to the propagation process for the left and right insertability verification (see Definition 11). In this section, we consider some extra local constraint on a chosen Lyndon factor, which ensures that the associated closest upper point is always insertable without the propagation process. For this aim, the authors in [7] restricted the study within the case where all the factors of the Lyndon factorization of the boundary word are primitive. They choose the closest upper point of the segment that is correlated to a primitive Christoffel word of maximal length with respect to previous and next primitive Christoffel words. With this constraint, removing the corresponding closest upper point of the locally longest Christoffel word always preserves the digital convexity. This result was not proved and neither generalized in the case where the factors of the boundary word are not primitive. We provide here the general result with a full study for all the possible cases and updates.

In fact, we would like to show that splitting a locally longest primitive Christoffel word of  $W(C)$  guarantees an inflation at any step, while preserving the digital convexity. In other terms, the closest upper points of all the primitive Christoffel words of local maximal length correspond to insertable points. Before reaching this theorem, we prove first that if we have three consecutive decreasing Christoffel words such that the first one is longer than the neighbors, then its split preserves this decreasing order in the local part.

From Definition 12 and Proposition 7, we recall that the Christoffel morphism  $\Theta_B$  induces an increasing bijection between the set of Christoffel words and itself. Based on this we can get the following Corollary.

**Corollary 4.** *Any Christoffel word  $C$  such that  $\mathcal{C}(\frac{a}{b}) < C < \mathcal{C}(\frac{c}{d})$  verifies  $|C| \geq |\mathcal{C}(\frac{a}{b})| + |\mathcal{C}(\frac{c}{d})| > \max(|\mathcal{C}(\frac{a}{b})|, |\mathcal{C}(\frac{c}{d})|)$ .*

*Proof:* Let  $B = (\mathcal{C}(\frac{a}{b}), \mathcal{C}(\frac{c}{d}))$ , and  $\Theta = \Theta_B$  be the Christoffel morphism defined by  $\Theta(0) = \mathcal{C}(\frac{a}{b})$  and  $\Theta(1) = \mathcal{C}(\frac{c}{d})$ . If  $\mathcal{C}(\frac{a}{b}) < C < \mathcal{C}(\frac{c}{d})$  then  $C = \Theta(U)$ , where  $U$  is a Christoffel word other than 0 or 1. Hence,  $U$  contains one letter 0 and one letter 1. Therefore the image contains two disjoint factors  $\mathcal{C}(\frac{a}{b})$  and  $\mathcal{C}(\frac{c}{d})$ .  $\square$

This shows also that it is true for any length  $\ell$ . Corollary 4 allows us to give an algorithm that inflates a digital convex set  $C$  while preserving its digital

convexity. We start by defining the notion of maximal primitive length of a Christoffel with respect to the previous and next one.

**Definition 13.** Let  $W(C) = \ell_1^{n_1} \dots \ell_s^{n_s}$  be the Lyndon factorization of a digital convex set  $C$ , and let  $\ell_j$ ,  $1 \leq j \leq s$ , be one of the Christoffel words, seen in a cyclic way at  $j = 1$  and  $j = s$ . We say that  $\ell_j$  has a local maximal primitive length if  $|\ell_j| > \max(|\ell_{j-1}|, |\ell_{j+1}|)$ .

By inflating the Christoffel word  $\ell_j$ , with a local maximal primitive on the closest upper point, we preserve the digital convexity, thanks to Theorem 7.

**Theorem 7.** Given a digital convex 4-connected set  $C$ ,  $x \in \overline{C}$  is an insertable point if and only if  $x$  is the closest upper point of a local maximal primitive length Christoffel of the Lyndon factorization of the boundary word  $W(C)$ .

*Proof:* The direct implication is trivial while the second implication is deduced from Corollary 4.  $\square$

This means that adding this constraint assures that this particular closest upper point is an insertable one.

After inflating  $C$  with respect to this strong constraint, we only can face one out of four cases, when we update  $W(C)$ . They are mentioned in Lemma 3.

**Lemma 3.** Let  $\ell_j$  be one of the Christoffel words of local maximal primitive length in the boundary word,  $W(C) = \ell_1^{n_1} \dots \ell_s^{n_s}$ , of a 4-connected digital convex set  $C$ . By applying the split operator on the  $i$ -th  $\ell_j$  for any  $1 \leq i \leq n_j$  such that  $\text{split}(\ell_j) = (\ell_j^+, \ell_j^-)$ , the Lyndon factorization can be updated by the following local replacement:

1. if  $(i > 1 \text{ or } \ell_{j-1} > \ell_j^+)$  and  $(i < n_j \text{ or } \ell_{j+1} < \ell_j^-)$ ,

$$\ell_1^{n_1} \ell_2^{n_2} \dots \ell_{j-1}^{n_{j-1}} (\ell_j^{i-1} \ell_j^+) (\ell_j^- \ell_j^{n_j-i}) \ell_{j+1}^{n_{j+1}} \dots \ell_k^{n_k}$$

2. if  $(i = 1 \text{ and } \ell_{j-1} = \ell_j^+)$  and  $(i < n_j \text{ or } \ell_{j+1} < \ell_j^-)$ ,

$$\ell_1^{n_1} \ell_2^{n_2} \dots \ell_{j-1}^{n_{j-1}+1} (\ell_j^- \ell_j^{n_j-i}) \ell_{j+1}^{n_{j+1}} \dots \ell_k^{n_k}$$

3. if  $(i > 1 \text{ or } \ell_{j-1} > \ell_j^+)$  and  $(i = n_j \text{ and } \ell_{j+1} = \ell_j^-)$ :

$$\ell_1^{n_1} \ell_2^{n_2} \dots \ell_{j-1}^{n_{j-1}} (\ell_j^{i-1} \ell_j^+) \ell_{j+1}^{n_{j+1}+1} \dots \ell_k^{n_k} .$$

4. if  $(i = 1 \text{ and } \ell_{j-1} = \ell_j^+)$  and  $(i = n_j \text{ and } \ell_{j+1} = \ell_j^-)$ :

$$\ell_1^{n_1} \ell_2^{n_2} \dots \ell_{j-1}^{n_{j-1}+1} \ell_{j+1}^{n_{j+1}+1} \dots \ell_k^{n_k} .$$

*Proof:* The proof of this lemma relies on the following two points:

1. Showing that  $u = \ell_j^{i-1} \ell_j^+$  and  $v = \ell_j^- \ell_j^{n_j-i}$  are Christoffel words. This follows from Lemma 2 by taking the base  $B = (\ell_j^-, \ell_j^+)$ . We obtain  $u = \Theta_B((01)^{i-1}1)$  and  $v = \Theta_B(0(01)^{n_j-i})$ .
2. Proving the following inequalities:  $\ell_{j-1} > \ell_j^{i-1} \ell_j^+ > \ell_j^- \ell_j^{n_j-i} > \ell_{j+1}$ .
  - The inequality in the middle comes from the fact that  $\Theta_B$  defined earlier is increasing and  $0(01)^{n_j-i} < (01)^{i-1}1$ .
  - If the last inequality is not correct, we have:  $\ell_j^- \leq \ell_j^- \ell_j^{n_j-i} \leq \ell_{j+1} < \ell_j$ . Then  $\ell_j$  is a Christoffel word in the  $B$  and the condition doesn't allow the equality  $\ell_{j+1} = \ell_j$ ; in this case it has to be longer than  $\ell_j$ , contradicting to the condition that  $\ell_j$  is longer than  $\ell_{j+1}$ .
  - The first inequality is treated in a symmetric way as the previous one.

□

From Lemma 3, we can remark that in all the four cases, the digital convexity is preserved. The inflated segment of the  $\text{conv}(C)$  is replaced by two others. For the second case,  $\ell_j^+$  is the same as  $\ell_j^{n_j-1}$  and a concatenation from the left side arises. Similarly and by symmetry, for the third case,  $\ell_j^-$  is the same as  $\ell_j^{n_j+1}$  and a concatenation from the right side arises. For the last case,  $\ell_j^+$  (resp.  $\ell_j^-$ ) is the same as  $\ell_j^{n_j-1}$  (resp.  $\ell_j^{n_j+1}$ ), and concatenations from both sides arise. In other words, we loose one of the segments of  $\text{conv}(C)$ . For all the cases, we can note that the propagation doesn't exceed the neighboring Christoffel words  $\ell_{j+1}^{n_j+1}$  and  $\ell_{j-1}^{n_j-1}$ .

**Corollary 5.** *Splitting the Christoffel word of the Lyndon factorization which has the local maximal primitive length limits the propagation and bounds it only by the previous and next factor.*

#### 4.4. Algorithm

We give now a general algorithm to inflate  $C$  based on the previous results: for the general case or the case with strong constraint. In Algorithm 2, we determine the list of all the possible insertable points obtained either by applying Theorem 6 or Theorem 7. For each iteration, we choose the point with the highest priority queue and we apply the necessary updates on the Lyndon factorization. We recall that for the general case, this update can lead to a propagation, while for the case with strong constraint, this update is local. Two different and more detailed algorithms, for each of the cases separately, will be given in a future work. By using this approach for the inflation, if we keep choosing the side with primitive local maximal length at each iteration, we remark that the horizontal and vertical segments of  $\text{conv}(C)$  will not be chosen. In fact the discretization of these segments are either factors of the form  $0^p$ ,  $1^q$ ,  $\bar{0}^r$  or  $\bar{1}^s$  for certain integer numbers  $p, q, r$ , or  $s$ . This means that their primitive

length is always equal to 1. Hence, at any step of the inflation they can not be designated. We see in this case, that the inflation will happen at the beginning on the  $WN, NE, ES$  and  $SW$  sides without considering the horizontal and vertical sides which leads to an octagonal shape as seen in Figure 17. Once we are at this step, all the remaining factors are of length 2 and 1. The inflation will continue until we reach the form:  $1^k 0^l \bar{1}^k \bar{0}^l$  for certain integers  $l, k$ , which is our rectangle bounded box.

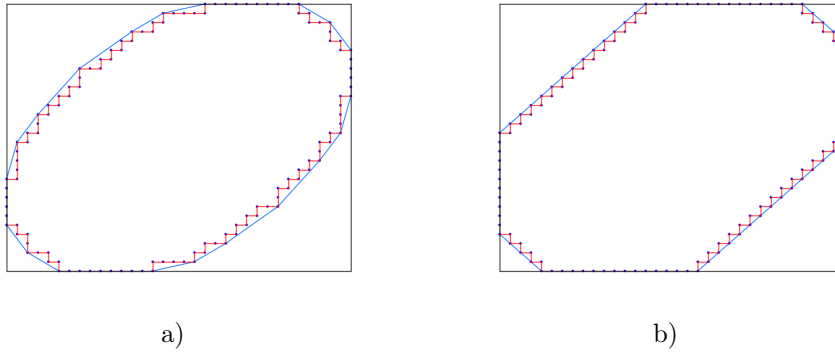


Figure 17: Figure b) shows the inflation of the digital convex set represented in a) by applying the algorithm with the stronger constraint based on local maximal primitive length of a Lyndon factor.

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**Algorithm 2** Point-wise inflation

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**Input:** a digital convex 4-connected set  $C$ , a number of adding points  $k$

**Output:** a sequence of inserted points  $I$

- 1: Compute the Lyndon Factorization  $F$  of the boundary word  $w$  of  $C$
  - 2: Insert all insertable points in a priority queue  $P$
  - 3: **while**  $k \geq 0$  **do**
  - 4: Pull the highest-priority insertable point  $p$  from  $P$  and add  $p$  to  $I$
  - 5: Let  $\ell_j$  be the associated primitive Christoffel word of  $p$  in  $F$ ;
  - 6: Compute  $split(\ell_j)$  and update  $F$  and  $P$  following Lemma 2 or 3
  - 7: Decrement  $k$
  - 8: **end while**
  - 9: **return**  $I$
- 

## 5. Conclusion

We have proposed a combinatorics-on-words study of the points that can be chosen to inflate and deflate a 4-connected digital convex set  $C$  while preserving its digital convexity property. The approach relies on Christoffel and Lyndon factorizations of the boundary word of  $C$  that is represented on a four letter

alphabet. Some update procedures are to be done on these factors, in order to maintain the Lyndon factorizations while adding/removing specific points. For both operations, we have characterized for  $C$  the set of points that can be inserted or removed, while maintaining the convexity. For the deflation process, the updates and modifications are local. In contrast, for the inflation process, the updates, using the general procedure, can be global. In worst case, these updates do not go past the side where the inflation is applied. Adding the strong condition on the choice of the insertable points corresponding to the local maximal primitive Christoffel word, the updates during this procedure become local.

In this work, we have focused on the characterization and theoretical properties of geometrical set operations on boundary word factorizations. The algorithmic details and optimization can be found in [30]. Another question that can arise is to determine if there exists an optimal heuristic for deflating a digital convex set. The choice of the heuristic is crucial, since for each choice we get a different convergent shape. Another stimulating perspective would be to apply these algorithms on non-convex shapes by studying the locally convex boundary using combinatorics on words [12, 26].

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