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► **To cite this version:**

Jérémie Chalopin, Victor Chepoi, Fionn Mc Inerney, Sébastien Ratel, Yann Vaxès. Sample Compression Schemes for Balls in Graphs. [Research Report] CISPA Helmholtz Center for Information Security, Saarbrücken, Germany; Aix-Marseille Université, France. 2022. hal-03705798v1

HAL Id: hal-03705798

<https://hal.science/hal-03705798v1>

Submitted on 27 Jun 2022 (v1), last revised 25 Mar 2024 (v2)

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SAMPLE COMPRESSION SCHEMES FOR BALLS IN GRAPHS

JÉRÉMIE CHALOPIN, VICTOR CHEPOI, FIONN MC INERNEY, SÉBASTIEN RATEL,
AND YANN VAXÈS

ABSTRACT. One of the open problems in machine learning is whether any set-family of VC-dimension d admits a sample compression scheme of size $O(d)$. In this paper, we study this problem for balls in graphs. For balls of arbitrary radius r , we design proper sample compression schemes of size 2 for trees, of size 3 for cycles, of size 4 for interval graphs, of size 6 for trees of cycles, and of size 22 for cube-free median graphs. For balls of a given radius, we design proper labeled sample compression schemes of size 2 for trees and of size 4 for interval graphs. We also design approximate sample compression schemes of size 2 for balls of δ -hyperbolic graphs.

Keywords: VC-dimension, Sample Compression Schemes, Proper Sample Compression Schemes

1. INTRODUCTION

Sample compression schemes were introduced by Littlestone and Warmuth [22], and have been vastly studied in the literature due to their importance in computational machine learning. Roughly, a sample compression scheme consists of a compressor α and a reconstructor β , and the aim is to compress data as much as possible, such that data coherent with the original data can be reconstructed from the compressed data. For balls in graphs, sample compression schemes of size k can be defined as follows. Given a ball $B = B_r(x)$ of a graph $G = (V, E)$, a realizable sample for B is a signed subset $X = (X^+, X^-)$ of V such that X^+ is included in B , and X^- is disjoint from B . Given a realizable sample X , X is compressed to a subsample $\alpha(X) \subseteq X$ of size at most k . The reconstructor β takes $\alpha(X)$ as an input and returns $\beta(\alpha(X))$, a subset B' of vertices of G that is consistent with X , *i.e.*, X^+ is included in B' , and X^- is disjoint from B' . If B' is always a ball of G , then the compression scheme is proper. If $X^+ = B$ and $X^- = V \setminus B$, then $\beta(\alpha(X))$ must coincide with B . Note that a proper sample compression scheme of size k for the family of all balls of G yields a sample compression scheme of size k for any subfamily of balls (*e.g.*, for balls of a fixed radius r), but this scheme is no longer proper. Sample compression schemes are labeled if β knows the labels of the elements of $\alpha(X)$, and are unlabeled otherwise (abbreviated LSCS and USCS, resp.). The Vapnik-Chervonenkis dimension (VC-dimension) of a set system was introduced by Vapnik and Chervonenkis [27] as a complexity measure of set systems. VC-dimension is central in PAC-learning, and is important in combinatorics and discrete geometry. Floyd and Warmuth [17] asked whether any set-family of VC-dimension d has a sample compression scheme of size $O(d)$. This remains one of the oldest open problems in machine learning.

In this paper, we consider the family of balls in graphs, which is as general as the sample compression conjecture. Indeed, the sample compression conjecture for set families in general is equivalent to the same conjecture restricted to the family of balls of radius 1 on split graphs in which samples only contain vertices in the clique, and the centers of the unit balls are in the stable set. Balls in graphs also constitute an important topic in graph theory, and moreover, their VC-dimension has often been considered in the literature (see, *e.g.*, [4, 7, 9, 15, 26]).

The VC-dimension of the balls of radius r of a graph not containing K_{n+1} as a minor is at most n [9]. This result was extended to arbitrary balls in [7]. Hence, the VC-dimension of balls of planar graphs is at most 4 (2 for trees and 3 for trees of cycles), and the VC-dimension of balls of a chordal graph G is at most its clique number $\omega(G)$. The VC-dimension of balls of interval

This work has been supported by the ANR project DISTANCIA (ANR-17-CE40-0015) and the European Research Council (ERC) consolidator grant No. 725978 SYSTEMATICGRAPH.

graphs was shown to be at most 2 in [15]. Finally, the VC-dimension of balls of cube-free median graphs is unknown, but we can prove that it is at least 4.

Our results. In this paper, we design proper sample compression schemes of small size for the family of balls of a graph G . We investigate this problem for different graph classes. For trees, we exhibit proper USCS of size 2 for all balls, and proper LSCS of size 2 for balls of equal radius. We also give proper USCS of size 3 for all balls in cycles. For trees of cycles, we exhibit proper LSCS of size 6 for all balls. Then, we design proper LSCS of size 22 for all balls of cube-free median graphs. We also construct proper LSCS of size 4 for all balls of interval graphs. This is followed by some partial results for split graphs and planar graphs. Mainly, for any split graph G , we give proper LSCS of size $\omega(G)$ for all balls of G , and we design proper LSCS of size 4 for balls of radius 1 of planar graphs. Finally, we define (ρ, μ) -approximate proper sample compression schemes, and design $(2\delta, 3\delta)$ -approximate LSCS of size 2 for δ -hyperbolic graphs.

Related work. Floyd and Warmuth [17] proved that, for any concept class of VC-dimension d , any LSCS has size at least $\frac{d}{5}$, and that, for some maximum classes of VC-dimension d , they have size at least d . Pálvölgyi and Tardos [25] proved that some concept classes of VC-dimension 2 do not admit USCS of size at most 2. On the positive side, it was shown by Moran and Yehudayoff [24] that LSCS of size $O(2^d)$ exist (their schemes are not proper). For particular concept classes, better results are known. Floyd and Warmuth [17] designed LSCS of size d for regions in arrangements of central hyperplanes in \mathbb{R}^d . Ben-David and Litman [5] obtained USCS of size d for regions in arrangements of affine hyperplanes in \mathbb{R}^d . Helmbold, Sloan, and Warmuth [19] (implicitly) constructed USCS of size d for intersection-closed concept classes. Moran and Warmuth [23] designed proper LSCS of size d for ample classes. Chalopin et al. [8] designed USCS of size d for maximum families. They also combinatorially characterized USCS for ample classes via the existence of *unique sink orientations* of their graphs. However, the existence of such orientations is open. Chepoi, Knauer, and Philibert [12] extended the result of [23], and designed proper LSCS of size d for concept classes defined by Complexes of Oriented Matroids (COMs). COMs were introduced in [3] as a natural common generalization of ample classes and Oriented Matroids [6].

2. DEFINITIONS

Concept classes and sample compression schemes. Let V be a non-empty finite set. Let $\mathcal{C} \subseteq 2^V$ be a family of subsets (also called a *concept class*) of V . The *VC-dimension* $\text{VC-dim}(\mathcal{C})$ of \mathcal{C} is the size of a largest set $Y \subseteq V$ *shattered* by \mathcal{C} , *i.e.*, such that $\{C \cap Y : C \in \mathcal{C}\} = 2^Y$. In machine learning, a (*labeled*) *sample* is a set $X = \{(x_1, y_1), \dots, (x_m, y_m)\}$, where $x_i \in V$ and $y_i \in \{-1, +1\}$. To X is associated the unlabeled sample $\underline{X} = \{x_1, \dots, x_m\}$. A sample X is *realizable by a concept* C if $y_i = +1$ if $x_i \in C$, and $y_i = -1$ if $x_i \notin C$. A sample X is *realizable by a concept class* \mathcal{C} if X is realizable by some $C \in \mathcal{C}$.

We adopt the language of sign maps and sign vectors from [6]. Let \mathcal{L} be a *set of sign vectors*, *i.e.*, maps from V to $\{\pm 1, 0\} := \{-1, 0, +1\}$. The elements of \mathcal{L} are also called *covectors*. For $X \in \mathcal{L}$, let $X^+ := \{v \in V : X_v = +1\}$ and $X^- := \{v \in V : X_v = -1\}$. $\underline{X} = X^- \cup X^+$ is called the *support* of X , and its complement $X^0 := V \setminus \underline{X} = \{v \in V : X_v = 0\}$ the *zero set* of X . Since $X^0 = V \setminus (X^- \cup X^+)$, we will view any sample X as $X^- \cup X^+$. Let \preceq be the product ordering on $\{\pm 1, 0\}^V$ relative to the ordering of signs with $0 \preceq -1$ and $0 \preceq +1$. Any concept class $\mathcal{C} \subseteq 2^V$ can be viewed as a set of sign vectors of $\{\pm 1\}^V$: for any $C \in \mathcal{C}$ we consider the sign vector $X(C)$, where $X_v(C) = +1$ if $v \in C$ and $X_v(C) = -1$ if $v \notin C$. For simplicity, we will consider \mathcal{C} as a family of sets and as a set of $\{\pm 1\}$ -vectors. We now define sample compression schemes. This way of presenting them seems novel. From the definition, it follows that a sample X is just a $\{\pm 1, 0\}$ -sign vector. Given a concept class $\mathcal{C} \subseteq 2^V$ and $C \in \mathcal{C}$, the set of samples realizable by C consists of all covectors $X \in \{\pm 1, 0\}^V$ such that $X \preceq C$. We denote by $\downarrow \mathcal{C}$ the set of all samples realizable by \mathcal{C} .

A *proper labeled sample compression scheme* (*proper LSCS*) of size k for a concept class $\mathcal{C} \subseteq \{\pm 1\}^V$ is defined by a *compressor* $\alpha : \{\pm 1, 0\}^V \rightarrow \{\pm 1, 0\}^V$ and a *reconstructor* $\beta : \{\pm 1, 0\}^V \rightarrow$

\mathcal{C} such that, for any realizable sample $X \in \downarrow \mathcal{C}$, $\alpha(X) \preceq X \preceq \beta(\alpha(X))$ and $|\underline{\alpha}(X)| \leq k$, where \preceq is the order between sign vectors defined above, and $\underline{\alpha}(X)$ is the support of the subsample of the sign vector X . Hence, $\alpha(X)$ is a signed vector with a support of size at most k such that $\alpha(X) \preceq X$, and $\beta(\alpha(X))$ is a concept C of \mathcal{C} viewed as a sign vector. It suffices to define the map α only on $\downarrow \mathcal{C}$, and the map β only on $\text{Im}(\alpha) := \alpha(\downarrow \mathcal{C})$. The condition $X \preceq \beta(\alpha(X))$ is equivalent to the condition $\beta(\alpha(X))|_{\underline{X}} = X$, which means that the restriction of the concept $\beta(\alpha(X))$ to the support of X coincides with the sign vector X . *Proper unlabeled sample compression schemes (proper USCS)* are defined analogously, only that $\alpha(X)$ is not a signed vector, but a subset of size at most k of the support of X . Formally, a proper unlabeled sample compression scheme of size k for a concept class $\mathcal{C} \subseteq 2^V$ is defined by the functions $\alpha : \{\pm 1, 0\}^V \rightarrow 2^V$ and $\beta : 2^V \rightarrow \mathcal{C}$ such that, for any realizable sample $X \in \downarrow \mathcal{C}$, $\alpha(X) \subseteq \underline{X}$, $X \preceq \beta(\alpha(X))$, and $|\alpha(X)| \leq k$. For graphs, any preprocessing on the input graph G , such as a labeling or an embedding of G , is permitted and known to both the compressor and the reconstructor. As in, *e.g.*, [22, 24], information, like representing the support as a vector with coordinates, is also permitted, and when we use such information, we refer to α and β as vectors rather than maps. Lastly, in our schemes, the reconstructor returns the empty set when $X^+ = \emptyset$, and thus, one may consider that our schemes are not proper. We note that in all of the LSCS for the family of balls of arbitrary radius we exhibit in this paper, we could simply put an ordering on the vertices of the graph $G = (V, E)$, and put a single vertex $z \in X^-$ into $\alpha(X)$, and return the ball $B_0(z')$, where $z' \in V$ is the successor of z in the ordering and $B_0(z') \cap X = \emptyset = X^+$. However, to not overcomplicate our schemes for such degenerate cases, we make use of the empty set.

Graphs. Every graph $G = (V, E)$ in this paper is simple and connected. The *distance* $d(u, v) := d_G(u, v)$ between two vertices u and v of a graph G is the length of a (u, v) -shortest path. The *interval* $I(u, v)$ is the set of vertices contained in (u, v) -shortest paths. A set S is *gated* if, for any vertex $x \in V$, there is a vertex $x' \in S$ (the *gate* of x , with $x' = x$ if $x \in S$) such that $x' \in I(x, y)$ for any $y \in S$. A *median* of a triplet u, v, w is any vertex in $I(u, v) \cap I(v, w) \cap I(w, u)$. A graph G is *median* [1] if any triplet of vertices u, v, w has a unique median. For any vertex $x \in V$ and any integer $r \geq 0$, the *ball of radius r centered at x* is the set $B_r(x) := \{v \in V : d(v, x) \leq r\}$. The unit ball $B_1(x)$ is usually denoted by $N[x]$ and called the *closed neighborhood of x* . The *sphere* of radius r centered at x is the set $S_r(x) = \{z \in V : d(z, x) = r\}$. Let also $\text{c}B_r(u) = V \setminus B_r(u)$. Two balls $B_{r_1}(x)$ and $B_{r_2}(y)$ are *distinct* if $B_{r_1}(x)$ and $B_{r_2}(y)$ are distinct as sets. We denote by $\mathcal{B}(G)$ the set of all distinct balls of G , and by $\mathcal{B}_r(G)$ the set of all distinct balls of radius r of G . For a subset $Y \subseteq V$, we call $\text{diam}(Y) = \max\{d(u, v) : u, v \in Y\}$ the *diameter* of Y , and we call any pair $u, v \in Y$ such that $d(u, v) = \text{diam}(Y)$ a *diametral pair* of Y .

3. TREES

3.1. Proper USCS for $\mathcal{B}(T)$ for metric trees. It is well known that balls in trees are gated, the family of balls has VC-dimension 2, and trees are median graphs. For simplicity, we first consider metric trees T , *i.e.*, each edge of T is homeomorphic to the segment $[0, 1]$ (combinatorial trees are treated later).

First, let X be a realizable sample for the family $\mathcal{B}(T)$. Define $\alpha(X)$ to be the sign map such that $\alpha^+(X)$ is any diametral pair of the set X^+ and $\alpha^-(X) := \emptyset$ if $|X^+| \geq 2$, and $\alpha(X) = X^+$ otherwise. For each pair of vertices $\{u, v\}$ of T , let $\beta(\{u, v\})$ be any ball B of T having u, v as a diametral pair. The center of B is in the middle of the (u, v) -path, and the radius of B is $\frac{1}{2}d(u, v)$. Also, let $\beta(\{u\})$ be $B_0(u)$ and $\beta(\emptyset)$ be the empty ball.

Proposition 1. *For any tree $T = (V, E)$, the pair (α, β) of maps defines a proper unlabeled sample compression scheme of size 2 for $\mathcal{B}(T)$.*

Proof. Let X be a realizable sample for a ball $B_{r'}(y)$. If $|X^+| < 2$, then $\alpha(X) = X^+$ and $\beta(\alpha(X^+)) = X^+$, so assume that $|X^+| \geq 2$. Let $\{u, v\}$ be a diametral pair of X^+ , and let $B_r(x)$ be the ball of T returned as $\beta(\{u, v\})$. Thus, $r = \frac{1}{2}d(u, v)$. We assert that $B_r(x)$ is consistent with X . Otherwise, either there exists $w' \in B_r(x) \cap X^-$ or $w'' \in X^+ \setminus B_r(x)$. First, let there exist $w' \in B_r(x) \cap X^-$. Suppose, without loss of generality, that the median t of the triplet u, v, w'

belongs to $I(x, v)$. Since $w' \in B_r(x)$, $d(x, w') \leq r = d(x, v)$, and therefore, $d(t, w') \leq d(t, v) \leq r$. Note also that $d(t, w') \leq d(t, u)$ because $d(t, w') \leq r$ and $d(t, u) \geq r$. Let z be the median of the triplet u, v, y . If $z \in I(u, t)$, then since $d(t, w') \leq d(t, v)$ and $v \in B_{r'}(y)$, we conclude that $w' \in B_{r'}(y)$, contrary to the assumption that $w' \in X^-$. Thus, $z \in I(v, t)$. Since t is the median of u, w', y , and $d(t, w') \leq d(t, u)$, we conclude that $d(y, w') \leq d(y, u) \leq r'$, contrary to the assumption that $w' \in X^-$. Now, let there exist $w'' \in X^+ \setminus B_r(x)$. Again, let t be the median of u, v, w'' , and suppose that $t \in I(x, v)$. Since $w'' \notin B_r(x)$, we conclude that $d(t, w'') > d(t, v)$. This implies that $d(u, w'') > d(u, v)$. Since $w'' \in X^+$, this contradicts that $\{u, v\}$ is a diametral pair of X^+ . \square

3.2. Proper LSCS for $\mathcal{B}(T)$ for combinatorial trees. We present a proper labeled sample compression scheme for the family of balls of a fixed (combinatorial) tree.

Consider a realizable sample X for the family $\mathcal{B}(T)$ such that $|X^+| \geq 2$. Consider a diametral pair $\{u^+, v^+\}$ of X^+ , and let $r = \left\lceil \frac{d(u^+, v^+)}{2} \right\rceil$. Let x be the unique vertex in $I(u^+, v^+)$ such that $d(u^+, x) = \left\lfloor \frac{d(u^+, v^+)}{2} \right\rfloor$ and $d(x, v^+) = r = \left\lceil \frac{d(u^+, v^+)}{2} \right\rceil$, and y be the unique vertex in $I(u^+, v^+)$ such that $d(u^+, x) = r = \left\lceil \frac{d(u^+, v^+)}{2} \right\rceil$ and $d(x, v^+) = \left\lfloor \frac{d(u^+, v^+)}{2} \right\rfloor$. Note that if $d(u^+, v^+) = 2r$, then $x = y$.

Claim 2. *The balls $B_r(x)$ and $B_r(y)$ satisfy the following properties:*

- (1) $X^+ \subseteq B_r(x) \cap B_r(y)$;
- (2) $B_r(x) \cap X^- = \emptyset$ or $B_r(y) \cap X^- = \emptyset$.

Consequently, one of the balls (possibly both) $B_r(x)$ or $B_r(y)$ realizes X .

Proof. The proof is similar to the proof of Proposition 1. Consider $w \in X^+$, and let w' be the median of u, v, w . Without loss of generality, assume that $w' \in I(u^+, x)$. Then, $d(v^+, w) = d(v^+, x) + d(x, w) = d(v^+, y) + d(y, w)$. Since u^+, v^+ is a diametral pair of X^+ , we have $d(v^+, w) \leq d(u^+, v^+) = d(u^+, x) + d(x, v^+) = d(u^+, y) + d(y, v^+)$. Consequently, $d(x, w) \leq d(x, u^+) \leq r$ and $d(y, w) \leq d(y, u^+) \leq r$. Therefore, $X^+ \subseteq B_r(x)$ and $X^+ \subseteq B_r(y)$.

Consider a ball of center z and of radius r' that realizes X . Without loss of generality, assume that the median of z, u^+, v^+ belongs to $I(u^+, x)$. Observe that $d(z, v^+) = d(z, x) + d(x, v^+)$, and thus, $r' \geq d(z, v^+) = d(z, x) + r$. Consequently, for any $w \in B_r(x)$, $d(w, z) \leq d(w, x) + d(x, z) \leq r + d(x, z) \leq r'$, and thus, $B_r(x) \subseteq B_{r'}(z)$. Consequently, $B_r(x) \cap X^- = \emptyset$. \square

If $d(u^+, v^+) = 2r$ is even, then $x = y$ and $B_r(x)$ realizes X by Claim 2. If $d(u^+, v^+) = 2r - 1$, by Claim 2, $B_r(x)$ or $B_r(y)$ realizes X . Observe that $B_r(x)$ and $B_r(y)$ are the only balls of radius r in $\mathcal{B}(T)$ containing u^+ and v^+ . If $B_r(x)$ is the only ball of radius r realizing X , then there exists $w \in X^-$ such that $w \in B_r(y) \setminus B_r(x)$. Since x and y are adjacent, necessarily, $d(w, x) = d(w, y) + 1 = r + 1$. Observe that this implies that y is the median of u, v, w , that $d(y, w) = d(y, u^+) = d(y, v^+) + 1 = r$, and that x is the unique vertex of $I(u^+, w)$ such that $d(u^+, x) = r - 1$.

Proposition 3. *For any tree T , there exists a proper labeled sample compression scheme of size 2 for $\mathcal{B}(T)$.*

Proof. We first define the compressor α . Consider a realizable sample X for the family $\mathcal{B}(T)$. If $|X^+| \leq 1$, then let $\alpha^+(X) := X^+$ and $\alpha^-(X) := \emptyset$. Suppose now that $|X^+| \geq 2$, and consider the vertices u^+, v^+, x, y defined as above, and let $r = \left\lceil \frac{d(u^+, v^+)}{2} \right\rceil$. If $B_r(x)$ and $B_r(y)$ both realize X , then $\alpha^+(X) := \{u^+, v^+\}$ and $\alpha^-(X) := \emptyset$. Otherwise, by Claim 2, either $B_r(x)$ or $B_r(y)$ realizes X , say the first (the other case is symmetric). Consequently, there exists $w \in X^-$ such that $d(w, x) = d(w, y) + 1 = r + 1$. In this case, let $\alpha^+(X) := \{u^+\}$, and $\alpha^-(X) := \{w\}$. Note that, in this case, $d(u^+, w) = 2r$.

We now define the reconstructor β . Consider a sample $Y \in \text{Im}(\alpha)$. If $Y^+ = Y^- = \emptyset$, then $\beta(Y)$ is the empty ball. If $Y^+ = \{y\}$ and $Y^- = \emptyset$, then $\beta(Y) = B_0(y)$. If $Y^+ = \{u, v\}$ and $Y^- = \emptyset$, then let $r = \left\lceil \frac{d(u, v)}{2} \right\rceil$, and let $\beta(Y)$ be any ball of radius r containing u and v . Suppose

now that $Y^+ = \{u\}$ and $Y^- = \{w\}$. Since $Y \in \text{Im}(\alpha)$, we can assume that $d(u, w) = 2r$. Let x be the unique vertex in $I(u, w)$ such that $d(u, x) = r - 1$, and let $\beta(Y) = B_r(x)$.

We claim that (α, β) is a proper labeled sample compression scheme for $\mathcal{B}(T)$. Consider a sample X . If $|X^+| \leq 1$, then $\alpha(X) = X^+$ and $\beta(\alpha(X^+)) = X^+$. Suppose now that $|X^+| \geq 2$. If $\alpha^+(X) = \{u^+, v^+\}$, then u^+, v^+ is a diametral pair of X^+ , and any ball of T of radius $r = \left\lceil \frac{d(u^+, v^+)}{2} \right\rceil$ containing u^+ and v^+ realizes X . Consequently, $\beta(\alpha(X))$ realizes X in this case. Suppose now that $\alpha^+(X) = \{u^+\}$ and $\alpha^-(X) = \{w\}$. Then, in this case, the ball $B_r(x)$ realizes X , where $r = \frac{d(u^+, w)}{2}$ and x is the unique vertex in $I(u^+, w)$ that is at distance $r - 1$ from u^+ . Since $\beta(\alpha(X)) = B_r(x)$, we are done. \square

3.3. Proper LSCS for $\mathcal{B}_r(T)$ for combinatorial trees. Now, let X be a realizable sample for the family $\mathcal{B}_r(T)$. The sample compression scheme for $\mathcal{B}(T)$ cannot be applied to $\mathcal{B}_r(T)$ since β may return a ball B whose radius can be much smaller than r . Keeping the center of B and increasing its radius until r may result in a ball which is no longer compatible with X^- . This leads us to design a different technique that encodes the center of a ball realizing the input sample. Let $X = X^+ \cup X^-$ be a realizable sample for $\mathcal{B}_r(T)$. Let $\{u^+, v^+\}$ be a diametral pair of X^+ (if $X^+ = \emptyset$, then u^+ and v^+ are not defined). Similarly to the proof of Proposition 1, we obtain the following result:

Lemma 4. *Any ball $B_r(x)$ of T containing $\{u^+, v^+\}$ also contains X^+ .*

If X^- is located far away from X^+ , then any r -ball realizing X^+ also realizes X . The next lemma provides the conditions under which this choice of the r -ball is no longer true:

Lemma 5. *Either any r -ball containing X^+ is disjoint from X^- , or there exists a ball $B_r(x)$ realizing X and a vertex $s \in X^-$ such that $d(x, s) = r + 1$.*

Proof. By Lemma 4, we can suppose that $X^- \neq \emptyset$. Let $B_r(y)$ be an r -ball containing X^+ and a vertex $z \in X^-$. Among all r -balls realizing X , pick a ball $B_r(x)$ minimizing $d(x, z)$. If $d(x, z) = r + 1$, then we are done by taking $s = z$. So, let $d(x, z) > r + 1$. Let x' be the neighbor of x in $I(x, z)$. By the choice of x , the ball $B_r(x')$ does not realize X . Since $d(x, z) > r + 1$, $z \notin B_r(x')$. If $B_r(x')$ contains a vertex $z' \in X^-$, then $d(x, z') = r + 1$, and we are done by taking $s = z'$. Thus, there exists $w \in X^+ \setminus B_r(x')$. Since $w \in B_r(x)$, we deduce that $d(x', w) = r + 1$ and $d(x, w) = r$. Consequently, $d(w, z) = d(w, x) + d(x, x') + d(x', z) = r + 1 + d(x', z) > 2r + 1$. This contradicts the assumption that $B_r(y)$ contains $X^+ \cup \{z\}$. \square

In view of Lemma 5, in certain cases, the center x of a ball $B_r(x)$ realizing X is located on the sphere $S_{r+1}(s)$ centered at some vertex s of X^- . If s is given, then we have to encode the position of x on $S_{r+1}(s)$ using only the vertices of X . Denote by ℓ_s the labeling of the sphere $S_{r+1}(s)$ obtained by performing a Depth-First Search (DFS) of T with root s . This labeling assigns increasing labels (starting from 0) to the vertices of $S_{r+1}(s)$, *i.e.*, the first vertex of $S_{r+1}(s)$ reached by the DFS is labeled 0, the second vertex of $S_{r+1}(s)$ reached by the DFS is labeled 1, and so on. Let $n_s = |S_{r+1}(s)| - 1$ (we also set $n_s = 0$ if $S_{r+1}(s) = \emptyset$). Consider the labels $\{0, 1, \dots, n_s\}$ in clockwise circular order and let us call any sequence of the form $\{i, i + 1, \dots, j - 1, j\}$ or $\{j, j + 1, \dots, n_s, 0, 1, \dots, i - 1, i\}$ a *circular interval*.

Lemma 6. *For any $v \in V(T)$, $S_{r+1}(s) \cap B_r(v)$ and $S_{r+1}(s) \cap \text{c}B_r(v)$ are circular intervals.*

Proof. Pick $x, y \in S_{r+1}(s) \cap B_r(v)$ with $\ell_s(x) = i < \ell_s(y) = j$, and $z \in S_{r+1}(s)$ with $i < \ell_s(z) < j$. We assert that $z \in B_r(v)$. Let a be the median of s, x, y . Then, a is the lowest common ancestor of x, y in T rooted at s . According to the DFS, for $z \in S_{r+1}(s)$, we have $i < \ell_s(z) < j$ if and only if a is an ancestor of z , *i.e.*, a belongs to a path between s and z . Since $x, y, z \in S_{r+1}(s)$, we have $d(a, x) = d(a, z) = d(a, y)$ for any such z . Let b be the closest to v vertex in the subtree T' of T spanned by the vertices s, x, y . If $b \in I(s, a)$, then $a \in I(v, x) \cap I(v, y) \cap I(v, z)$, and therefore, $d(v, z) = d(v, a) + d(a, z) = d(v, a) + d(v, x) = d(v, x) \leq r$, yielding $z \in B_r(v)$. Now, suppose that $b \in I(a, x)$ (the case $b \in I(x, y)$ is similar). By the triangle inequality and since $d(a, z) = d(a, y)$, we obtain $d(v, z) \leq d(v, a) + d(a, z) = d(v, b) + d(b, a) + d(a, y) = d(v, y) \leq r$,

and hence, $z \in B_r(v)$ in this case too. Finally, $S_{r+1}(s) \cap \text{c}B_r(v)$ is an interval as the complement of $S_{r+1}(s) \cap B_r(v)$. \square

If $S_{r+1}(s) \cap B_r(v)$ is a non-empty proper subset of $S_{r+1}(s)$, we denote by $\phi_s^+(v)$ the last vertex of the circular interval $S_{r+1}(s) \cap B_r(v)$, and by $\phi_s^-(v)$ the last vertex of the circular interval $S_{r+1}(s) \cap \text{c}B_r(v)$. Let $B_r(x)$ and $s \in X^-$ be as in Lemma 5: $B_r(x)$ realizes the sample X and $d(x, s) = r + 1$. To locate x on the sphere $S_{r+1}(s)$, we identify x by another vertex $t \in X$ such that $\phi_s^+(t) = x$ if $t \in X^+$ or $\phi_s^-(t) = x$ if $t \in X^-$ holds. We call $t \in X$ the *center designator* of the vertex s . Now, we strengthen the assertion of Lemma 5.

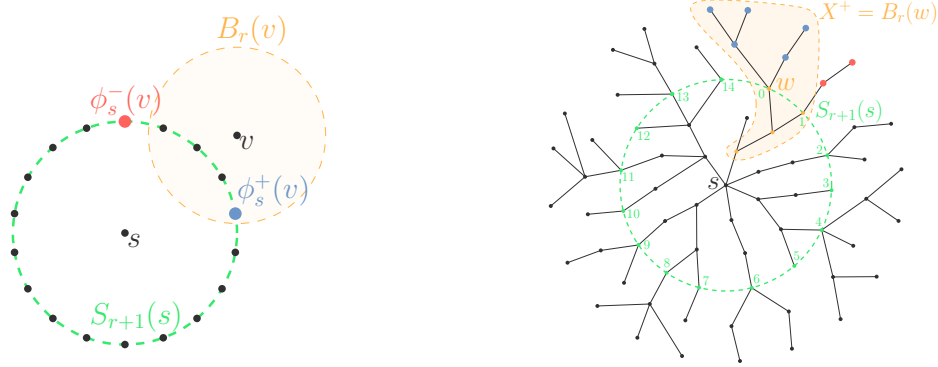


FIGURE 1. On the left is an illustration of the vertices $\phi_s^+(v)$ and $\phi_s^-(v)$ associated to a vertex $v \in V(T)$ with respect to a vertex $s \in X^-$. On the right is an example of the possible center designators of a vertex $s \in X^-$. The vertices $t \in X^+$ such that the r -ball centered at $\phi_s^+(t)$ realizes X are in blue. The vertices $t \in X^-$ such that the r -ball centered at $\phi_s^-(t)$ realizes X are in red.

Lemma 7. *One of the following conditions holds:*

- (1) *any r -ball realizing X^+ also realizes X ;*
- (2) *there exists a vertex $s \in X^-$ such that, for all $x \in S_{r+1}(s)$, the ball $B_r(x)$ realizes X ;*
- (3) *there exists a vertex $s \in X^-$ having a center designator $t \in X$.*

Proof. By Lemma 5, either any r -ball realizing X^+ also realizes X (and (1) holds) or there exists a ball $B_r(w)$ realizing X , and a vertex $s \in X^-$ with $d(s, w) = r + 1$. In the second case, if all the vertices of $S_{r+1}(s)$ are centers of r -balls realizing X , then (2) is satisfied. Thus, we can suppose that $S_{r+1}(s)$ contains a vertex whose r -ball does not realize X . In this case, we assert that s admits a center designator, *i.e.*, s satisfies (3). Indeed, moving clockwise along $S_{r+1}(s)$ according to the DFS-order ℓ_s , we will find two consecutive vertices x and y such that $B_r(x)$ realizes the sample X and $B_r(y)$ does not realize X . The latter implies that either there exists $t' \in X^+ \setminus B_r(y)$ or there exists $t'' \in X^- \cap B_r(y)$. If t' exists, since $t' \in B_r(x)$ and $t' \notin B_r(y)$, we conclude that $x \in B_r(t')$ and $y \notin B_r(t')$. Hence, x is the last vertex of the circular interval $S_{r+1}(s) \cap B_r(t')$, and thus, $x = \phi_s^+(t')$, yielding that t' is a center designator of s . On the other hand, if t'' exists, since $t'' \notin B_r(x)$ and $t'' \in B_r(y)$, we conclude that $x \notin B_r(t'')$ and $y \in B_r(t'')$. Thus, x is the last vertex of the circular interval $S_{r+1}(s) \cap \text{c}B_r(t'')$, and thus, $x = \phi_s^-(t'')$, yielding that t'' is a center designator of s . \square

We now use the previous results to define the compressor $\alpha(X)$ for any realizable sample X . Let $\{u^+, v^+\}$ be a diametral pair of X^+ . We define $\alpha(X)$ according to Lemma 7:

- (1) if any r -ball realizing X^+ also realizes X , then set $\alpha^-(X) := \emptyset$ and $\alpha^+(X) := (u^+, v^+)$;
- (2) if there exists a vertex $s \in X^-$ such that, for all $x \in S_{r+1}(s)$, the ball $B_r(x)$ realizes X , then set $\alpha^-(X) := (s)$ and $\alpha^+(X) := \emptyset$;
- (3) if there exists $s \in X^-$ having a center designator $t \in X$, then set $\alpha^-(X) := (s, t)$ and $\alpha^+(X) := \emptyset$ if $t \in X^-$, and set $\alpha^-(X) := (s)$ and $\alpha^+(X) := (t)$, otherwise.

In Case (3), when $|\alpha^-(X)| = 2$, we suppose that the second vertex of $\alpha^-(X)$ is the center designator of the first vertex. The reconstructor β takes any sign vector Y with a support of size at most 2 (from $\text{Im}(\alpha) = \alpha(\downarrow \mathcal{C})$) and returns a ball $B_r(y)$ defined in the following way:

- (1) if $Y^- = \emptyset$, then $\beta(Y)$ is any ball $B_r(y)$ containing Y^+ ;
- (2) if $Y^+ = \emptyset$ and $Y^- = (s)$, then $\beta(Y)$ is any ball $B_r(y)$ centered at a vertex y of $S_{r+1}(s)$;
- (3') if $Y^+ = \emptyset$ and $Y^- = (s, t)$, then $\beta(Y)$ is the ball $B_r(y)$ centered at the vertex $y = \phi_s^-(t)$;
- (3'') if $Y^+ = (t)$ and $Y^- = (s)$, then $\beta(Y)$ is the ball $B_r(y)$ centered at the vertex $y = \phi_s^+(t)$.

From the correspondence of the cases in Lemma 7 and the definitions of α and β , we get:

Proposition 8. *For any tree $T = (V, E)$ and any fixed radius r , the pair (α, β) of vectors defines a proper labeled sample compression scheme of size 2 for $\mathcal{B}_r(T)$.*

Proof. Let $Y = \alpha(X)$ and $B_r(y) = \beta(Y)$. Cases (1) and (2) in the definition of β correspond to Cases (1) and (2) in the definition of α , and to Cases (1) and (2) in Lemma 7. In these cases, $B_r(y)$ realizes X . Cases (3') and (3'') in the definition of β correspond to Case (3) in the definition of α and to Case (3) in Lemma 7. Thus, in Cases (3') and (3''), t is the center designator of s . By Lemma 7 and its proof, $B_r(y)$ realizes X . \square

3.4. Proper LSCS without information for $\mathcal{B}_r(T)$ for combinatorial trees. Next, we show how to design proper labeled sample compression schemes without information for balls of fixed radius in trees. We consider the case where the vertices of Y are not distinguishable by their order, *i.e.*, $\alpha(X)$ is not a vector. The only case where this causes problems is when $\alpha(X) := \{t, s\} =: Y$ and $t \in X^-$ is the center designator of $s \in X^-$, since it is not known which of the vertices in Y is the center designator. To get past this problem, we use an arbitrary labeling ℓ' of the vertices of T , that assigns a distinct integer in $1, \dots, n$ to each of the vertices of T , *i.e.*, for any two distinct vertices $u, v \in V(T)$, we have that $\ell'(u) \neq \ell'(v)$, and we increase the size of the support to at most 6 in all cases. Namely, we replace item (3'') in the encoding with information by the following lines (where t denotes a center designator of s):

- (1) if there exists $s, t \in X^-$ such that $\ell'(s) < \ell'(t)$, or if $|X| = 2$, then set $\alpha^-(X) := \{s, t\}$ and $\alpha^+(X) := \emptyset$;
- (2) if $|X^+| \geq 1$, then set $\alpha^-(X) := \{s, t\}$ and $\alpha^+(X) := \{w\}$, where w is an arbitrary vertex of X^+ ;
- (3) at this point, $|X^+| = 0$ and $\ell'(t) < \ell'(s)$ for any center designator $t \in X^-$ of $s \in X^-$. If there exist $s, t, p \in X^-$ such that $\ell'(t) < \ell'(s) < \ell'(p)$, or if $X^- = \{s, t, p\}$, then set $\alpha^-(X) := \{s, t, p\}$;
- (4) from now on, $\ell'(s) = \max\{\ell'(v) : v \in X^-\}$. If there exists $s, t, p, q \in X^-$ such that $\ell'(t) < \ell'(p) < \ell'(q) < \ell'(s)$, or if $X^- = \{s, t, p, q\}$, then set $\alpha^-(X) := \{s, t, p, q\}$;
- (5) if there exists $s, t, p, q, w \in X^-$ such that $\ell'(w) < \ell'(q) < \ell'(t) < \ell'(p) < \ell'(s)$, or if $X^- = \{s, t, p, q, w\}$, then set $\alpha^-(X) := \{s, t, p, q, w\}$, where q and w are two arbitrary vertices of X^- distinct from s, t, p , and each other;
- (6) if $|X^-| \geq 6$, then set $\alpha^-(X) := \{s, t, p, q, w, z\}$, where p, q, w , and z are four arbitrary vertices of X^- distinct from s, t , and each other.

The reconstructor β takes any sign vector Y from $\text{Im}(\alpha) := \alpha(\downarrow \mathcal{C})$ (with a support of size at most 6) and returns a ball B_r of radius r of T as described in the following lines:

- (1) if $Y^+ = \emptyset$ and $Y^- = \{s\}$, then $\beta(Y)$ is any ball $B_r(y)$ centered at a vertex y of $S_{r+1}(s)$;
- (2) if $|Y^+| \geq 1$ and $Y^- = \emptyset$, then $\beta(Y)$ is any ball $B_r(y)$ containing Y ;
- (3) if $Y^+ = \{t\}$ and $Y^- = \{s\}$, then $\beta(Y)$ is the ball $B_r(y)$ centered at the vertex $y = \phi_s^+(t)$;
- (4) if $|Y| = |Y^-| = 2$, then let $s, t \in Y$ be such that $\ell'(s) < \ell'(t)$. Let $y := \phi_s^-(t)$. If $B_r(y) \cap Y = \emptyset$, then $\beta(Y) := B_r(y)$. Otherwise, $\beta(Y)$ is any ball $B_r(y)$ avoiding Y^- ;
- (5) if $|Y^-| = 2$ and $|Y^+| = 1$, then let $s, t \in Y$ be such that $\ell'(s) > \ell'(t)$. Then, $\beta(Y)$ is the ball $B_r(y)$ centered at the vertex $y = \phi_s^-(t)$;
- (6) if $|Y| = |Y^-| = 3$, then let $s, t, p \in Y$ be such that $\ell'(t) < \ell'(s) < \ell'(p)$. Let $y := \phi_s^-(t)$. If $B_r(y) \cap Y = \emptyset$, set $\beta(Y) := B_r(y)$. Otherwise, $\beta(Y)$ is any ball $B_r(y)$ avoiding Y^- ;

- (7) if $|Y| = |Y^-| = 4$, then let $t \in Y$ be such that $\ell'(t) = \min\{\ell'(y) : y \in Y\}$, and let $s \in Y$ be such that $\ell'(s) = \max\{\ell'(y) : y \in Y\}$. Let $y := \phi_s^-(t)$. If $B_r(y) \cap Y = \emptyset$, then $\beta(Y) := B_r(y)$. Otherwise, $\beta(Y)$ is any ball $B_r(y)$ avoiding Y^- ;
- (8) if $|Y| = |Y^-| = 5$, then let $s, t, p, q, w \in Y$ be such that $\ell'(w) < \ell'(q) < \ell'(t) < \ell'(p) < \ell'(s)$. Let $y := \phi_s^-(t)$. If $B_r(y) \cap Y = \emptyset$, then $\beta(Y) := B_r(y)$. Otherwise, $\beta(Y)$ is any ball $B_r(y)$ avoiding Y^- ;
- (9) if $|Y| = |Y^-| = 6$, then let $s \in Y$ be such that $\ell'(s) = \max\{\ell'(y) : y \in Y\}$, and let $t \in Y$ be such that $\forall y \in Y, y \neq s, \ell'(y) < \ell'(t)$. Then, $\beta(Y)$ be the ball $B_r(y)$ centered at the vertex $y = \phi_s^-(t)$.

The following Proposition 9 is now easy to obtain since in line (4) ((6), resp.), either t is the center designator of s , or $|X^+| = 0$ and $|X^-| = 2$ ($|X^-| = 3$, resp.). Moreover, in line (7) ((8), resp.), either t is the center designator of s , or $|X^+| = 0$ and $|X^-| = 4$ ($|X^-| = 5$, resp.).

Proposition 9. *For any tree $T = (V, E)$ and any radius r , the pair (α, β) of maps defines a proper labeled sample compression scheme of size 6 for $\mathcal{B}_r(T)$.*

4. CYCLES

In this subsection, we consider the problem of constructing sample compression schemes for cycles. This section serves as a warm-up for the section in which we study sample compression schemes for trees of cycles. In particular, this section is here for completeness, and it shows that the main difficulty when considering sample compression schemes for trees of cycles does not come from the case of a single cycle, just as it does not come from a tree (in fact, as will be discussed in Remark 18, it comes from a spider which is a single cycle with paths of different lengths emanating from it). For these reasons, and since we only study the case of the family of balls of arbitrary radius in trees of cycles, we also only study the case of the family of balls of arbitrary radius in cycles. As in the case of trees, in what follows, instead of a graphic cycle, we consider a solid circle C (homeomorphic to the sphere \mathbb{S}^1), endowed with the circle-distance. Also, note that the family of balls of a cycle has VC-dimension at most 3.

Given a cycle C , we define a proper unlabeled sample compression scheme of size 3 for the set $\mathcal{B}(C)$ of balls of C as follows. We first define the compressor function $\alpha : \{\pm 1, 0\}^V \rightarrow 2^V$. For any realizable sample X for $\mathcal{B}(C)$,

- (1) if $|X^+| \leq 1$, then set $\alpha(X) := X^+$;
- (2) if there exist distinct vertices $u, v \in X^+$ such that $X^+ \subseteq I(u, v)$ and $X^- \cap I(u, v) = \emptyset$, then we set $\alpha(X) := \{u, v\}$;
- (3) if $X^- = \emptyset$ and, for all $u, v \in X^+$, we have $X^+ \setminus I(u, v) \neq \emptyset$, then we choose $u, v, w \in X^+$ such that $d(u, v) + d(v, w) + d(u, w)$ is maximum and we set $\alpha(x) = \{u, v, w\}$;
- (4) if there exist distinct vertices $u, v \in X^+$ and a shortest path P between u and v on C such that $X^+ \cap P = \{u, v\}$ and there exists $w \in P \cap X^-$, then set $\alpha(X) := \{u, v, w\}$.

Observe that if $|X^+| = 2$, then either we are in Case (2) or in Case (4). In Case (4), if u, v is not a diametral pair of C , then we can replace the path P by $I(u, v)$.

We now describe the reconstructor function $\beta : 2^V \rightarrow \mathcal{B}$. Let $Y := \alpha(X)$.

- (1) if $Y = \emptyset$, then $\beta(Y)$ is the empty ball;
- (2) if $Y = \{u\}$, then $\beta(Y) = B_0(u)$;
- (3) if $Y = \{u, v\}$, then $\beta(Y) = I(u, v)$ (note that if u, v is a diametral pair of C , $I(u, v) = B_{d(u,v)}(u)$ and that $I(u, v)$ is a ball of radius $d(u, v)/2$ otherwise);
- (4) if $Y = \{u, v, w\}$ and $d(v, w) \leq d(u, w) \leq d(u, v)$, we distinguish two cases:
 - (a) if $w \in I(u, v)$, then $\beta(Y)$ is the unique ball B of C containing u, v and not w such that u, v is a diametral pair of B ;
 - (b) any ball of C covering C otherwise.

Proposition 10. *For any cycle $C = (V, E)$, the pair (α, β) of maps defines a proper unlabeled sample compression scheme of size 3 for $\mathcal{B}(C)$.*

Proof. We just need to show that, for any realizable sample X for $\mathcal{B}(C)$, $\beta(\alpha(X))$ realizes X . If $X^+ = \emptyset$, then $\beta(\alpha(X)) = \emptyset = X^+$. If $X^+ = \{u\}$, then $\beta(\alpha(X)) = \{u\} = X^+$. If there exist

distinct vertices $u, v \in X^+$ such that $X^+ \subseteq I(u, v)$ and $X^- \cap I(u, v) = \emptyset$, then $\beta(\alpha(X)) = I(u, v)$ realizes X .

If $X^- = \emptyset$ and, for all $u, v \in X^+$, we have $X^+ \setminus I(u, v) \neq \emptyset$, then $\alpha(X) = \{u, v, w\}$ and $d(u, v) + d(v, w) + d(u, w)$ is maximum among all triplets of vertices in X^+ . Consequently, $u \notin I(v, w)$, $v \notin I(u, w)$, and $w \notin I(u, v)$, and thus, $\beta(\alpha(X)) = V$ realizes X .

Finally, if $X^- \neq \emptyset$ and X^+ is not contained in some interval $I(u, v)$, then there exist $u, v \in X^+$ such that when we consider the two simple (u, v) -paths P_1 and P_2 in C , we have $X^+ \subseteq P_1$, $X^- \subseteq P_2$, and $|P_1| \geq |P_2|$. In this case, $\alpha(X) = \{u, v, w\}$ for some $w \in P_2$. Consequently, $w \in I(u, v)$ and $\beta(\alpha(X)) = P_1$ realizes X . \square

5. TREES OF CYCLES

A *tree of cycles* (or *cactus*) is a graph in which each *block* (2-connected component) is a cycle or an edge. In this section, we give a proper labeled sample compression scheme of size 6 for balls of trees of cycles. Let G be a tree of cycles. For a vertex v of G that is not a cut vertex, let $C(v)$ be the unique cycle containing v . If v is a cut vertex or a degree-one vertex, then set $C(v) = \{v\}$. Let $T(G)$ be the tree whose vertices are the cut vertices and the blocks of G , and where a cut vertex v is adjacent to a block B of G if and only if $v \in B$. For any two vertices u, v of G , let $C(u, v)$ denote the union of all cycles and/or edges on the unique path of $T(G)$ between $C(u)$ and $C(v)$. Note that $C(u, v)$ is a path of cycles, and that $C(u, v)$ is gated. Let X be a realizable sample for $\mathcal{B}(G)$, and $\{u^+, v^+\}$ a diametral pair of X^+ . The next lemma shows that the center of a ball realizing X can always be found in $C(u^+, v^+)$.

Lemma 11. *Let $B_r(x)$ be a ball realizing X , x' be the gate of x in $C(u^+, v^+)$, and $r' = r - d(x, x')$. Then, the ball $B_{r'}(x')$ also realizes X .*

Proof. Let $B' := B_{r'}(x')$. From the definition of x' and r' , it immediately follows that $B' \subseteq B_r(x) = B$, implying that $X^- \cap B' = \emptyset$. To prove the inclusion $X^+ \subseteq B'$, suppose by way of contradiction that there exists $z \in X^+ \setminus B'$. First, suppose that $x' \in I(z, u^+) \cup I(z, v^+)$, say $x' \in I(z, u^+)$. Then, $d(u^+, v^+) \leq d(u^+, x') + d(x', v^+) \leq d(u^+, x') + r' < d(u^+, x') + d(x', z) = d(u^+, z)$, and thus, $\{u^+, v^+\}$ is not a diametral pair of X^+ , a contradiction. Consequently, $x' \notin I(z, u^+) \cup I(z, v^+)$. In particular, $x' \neq u^+, v^+$. This shows that, in the graph $G \setminus \{x'\}$, the three vertices u^+, v^+ , and z belong to the same connected component and x' is a cut vertex separating the vertex x from u^+, v^+, z . Since $z \notin B_{r'}(x')$, we have $d(x, z) = d(x, x') + d(x', z) > d(x, x') + r' = r$, contradicting the assumption that $z \in X^+ = X \cap B_r(x)$. \square

In what follows, let $B_r(x)$ be a ball realizing X with x in $C(u^+, v^+)$ (it exists by Lemma 11). Let C be a cycle of $C(u^+, v^+)$ containing x . The main idea is to encode a region of $C(u^+, v^+)$ where the center x of $B_r(x)$ is located (this region may be C), the center, and the radius of $B_r(x)$ by a few vertices of X . The diametral pair $\{u^+, v^+\}$ is in $\alpha(X)$. If X contains a vertex $w \neq u^+, v^+$ whose gate in $C(u^+, v^+)$ is in C , then C is easily detected by including w in $\alpha(X)$. In this case, it remains to find the position of x in C and to compute the radius r . This is done by using 2 or 3 vertices of X . Otherwise, if the gates in $C(u^+, v^+)$ of all vertices $w \in X \setminus \{u^+, v^+\}$ are outside C , then we show that $B_r(x)$ is determined by 4 vertices in X .

The partitioning of X . For a vertex $y \in C(u^+, v^+)$, set $r_y := \max\{d(y, u^+), d(y, v^+)\}$ and $r_y^* := \max\{d(y, w) : w \in X^+\}$. Clearly, $B_{r_y^*}(y)$ is the smallest ball centered at y containing X^+ . For any vertex z of G , we denote by z' its gate in $C(u^+, v^+)$. Let u^* and v^* be the gates of u^+ and v^+ in C . We partition X and X^- as follows. Let X_u (X_u^- , resp.) consist of all $w \in X$ ($w \in X^-$, resp.) whose gate w' in $C(u^+, v^+)$ belongs to $C(u^+, u^*)$. The sets X_v and X_v^- are defined analogously. Let X_C (X_C^- , resp.) consist of all the vertices $w \in X$ ($w \in X^-$, resp.) whose gates w' in $C(u^+, v^+)$ belong to the cycle C . Note that some of these sets can be empty and that $X_u^- \subseteq X_u$, $X_v^- \subseteq X_v$, and $X_C^- \subseteq X_C$. Let u_0 be the cut vertex of $C(u^+, v^+)$ farthest from u^* , and such that, for any vertex $w \in X_u$, its gate $w' \in C(u^+, v^+)$ is not in $C(u_0, u^*)$. Analogously, we define the cut vertex v_0 with respect to v^* and X_v . If $u^* = u^+$ ($v^* = v^+$, resp.), then set $u_0 = u^* = u^+$ ($v_0 = v^* = v^+$, resp.).

First, suppose that $X_C = \emptyset$. Let w_1 be a vertex of X_u closest to u_0 , and z_1 a vertex of X_u^- closest to x . Note that w_1 always exists as u^+ is in X_u , and that z_1 exists if and only if X_u^- is non-empty. Similarly, we define the vertices w_2 and z_2 with respect to X_v and X_v^- . See Fig. 2 for an illustration. The next lemmas show how to compute $B_r(x)$ in this case.

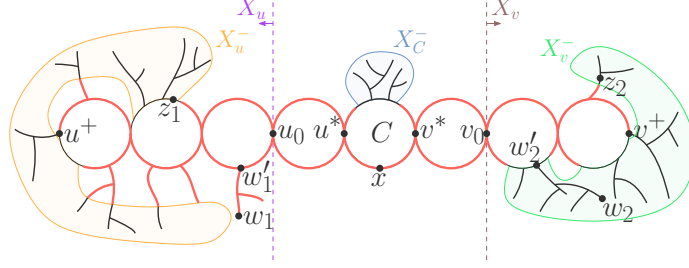


FIGURE 2. The vertices and sets used in the proper labeled sample compression scheme for trees of cycles. The ball $B_r(x)$ is represented in red. The cycles outside $C(u^+, v^+)$ are represented as paths.

Lemma 12. *For $y \in C(u_0, v_0)$, if there exists a vertex $w \in X^+ \setminus B_{r_y}(y)$, then $w' \in C(y)$. Consequently, if $X_C = \emptyset$, then, for any $y \in C(u_0, v_0)$, we have $X^+ \subset B_{r_y}(y)$ and $r_y = r_y^*$.*

Proof. Suppose that the gate w' of w in $C(u^+, v^+)$ does not belong to $C(y)$. By the definition of u_0 and v_0 , it follows that w' either belongs to $C(u_0, u^+)$ or to $C(v_0, v^+)$, say $w' \in C(v_0, v^+)$. Then, v_0 separates u^+ and y from v^+ and w . Consequently,

$$\begin{aligned} d(y, w) &= d(y, v_0) + d(v_0, w) & \text{and} & \quad d(y, v^+) = d(y, v_0) + d(v_0, v^+), \text{ and} \\ d(u^+, w) &= d(u^+, v_0) + d(v_0, w) & \text{and} & \quad d(u^+, v^+) = d(u^+, v_0) + d(v_0, v^+). \end{aligned}$$

From the first two equalities, we obtain $d(v_0, w) > d(v_0, v^+)$. From the last two equalities, it follows that $d(u^+, w) > d(u^+, v^+)$, contradicting that $\{u^+, v^+\}$ is a diametral pair of X^+ . \square

Lemma 13. *If $X^- \neq \emptyset$ and $X_C = \emptyset$, then $B_{r_y^*}(y) \cap X^- = \emptyset$ for any vertex $y \in C(u_0, v_0)$ such that $B_{r_y^*}(y) \cap \{z_1, z_2\} = \emptyset$.*

Proof. By Lemma 12, we have $r_y^* = r_y$. Suppose by way of contradiction that there exists a vertex $z \in B_{r_y}(y) \cap X^-$. Since $X_C = \emptyset$, z belongs to one of the sets X_u^- or X_v^- , say $z \in X_v^-$. Then, z_2 exists and the vertex v_0 separates x and y from z_2 and z . Consequently,

$$\begin{aligned} d(y, z_2) &= d(y, v_0) + d(v_0, z_2) & \text{and} & \quad d(y, z) = d(y, v_0) + d(v_0, z), \text{ and} \\ d(x, z_2) &= d(x, v_0) + d(v_0, z_2) & \text{and} & \quad d(x, z) = d(x, v_0) + d(v_0, z). \end{aligned}$$

Since $z_2 \notin B_{r_y}(y)$ and $z \in B_{r_y}(y)$, we conclude that $d(v_0, z_2) > d(v_0, z)$. From the last two equalities, $d(x, z_2) > d(x, z)$, contrary to the choice of z_2 as a vertex of X_v^- closest to x . \square

Now, suppose that $X_C \neq \emptyset$. By the definition of r_x^* , $B_{r_x^*}(x)$ also realizes X . Let w be a vertex of X whose gate w' in $C(u^+, v^+)$ is in C . If, for every $y \in C$, $B_{r_y^*}(y)$ realizes X , then $B_{r_{w'}}^*(w')$ realizes X , and, in this case, let $s \in X^+$ be such that $d(w', s) = r_{w'}^*$. Otherwise, we can find two adjacent vertices x and y of C such that $B_{r_x^*}(x)$ realizes X , but $B_{r_y^*}(y)$ does not. This implies that there is a vertex $z \in X^-$ with $z \in B_{r_y^*}(y) \setminus B_{r_x^*}(x)$. In this case, let $s, t \in X^+$ be such that $r_y^* = d(y, s)$ and $r_x^* = d(x, t)$, with $t = s$ whenever $r_y^* = r_x^* + 1$. Let s', t' , and z' be the respective gates of s, t , and z in C . If $s = t$ ($s \neq t$, resp.), then let P' be the path of C between s' and z' (t' , resp.) containing the edge xy . See Fig. 3 for an illustration.

Lemma 14. *For adjacent vertices $x, y \in C$, and the corresponding vertices $z \in X^-$ and $s \in X^+$, one of the following conditions holds:*

- (1) $r_y^* = r_x^* + 1$, $d(x, z) = d(y, z) + 1$, and $d(x, s) = d(y, s) - 1$;
- (2) $r_y^* = r_x^* + 1$, $d(x, z) = d(y, z)$, and $d(x, s) = d(y, s) - 1$;
- (3) $r_y^* = r_x^*$, $d(x, z) = d(y, z) + 1$, and $d(x, s) = d(y, s)$;

(4) $r_y^* = r_x^*$, $d(x, z) = d(y, z) + 1$, and $d(x, s) = d(y, s) - 1$.

Proof. Since x and y are adjacent, $|r_x^* - r_y^*| \leq 1$. If $r_x^* = r_y^* + 1$, then z also belongs to the ball $B_{r_x^*}(x)$, contrary to our choice of x . This shows that $r_y^* \in \{r_x^*, r_x^* + 1\}$, establishing the first equality in each of the four cases.

Now, notice that $x \notin I(y, z)$. Otherwise, since $r_x^* \geq r_y^* - 1$, we obtain that $z \in B_{r_x^*}(x)$, a contradiction. Since x and y are adjacent, we have either $d(x, z) = d(y, z) + 1$ (i.e., $y \in I(x, z)$) or $d(x, z) = d(y, z)$.

If $r_y^* = r_x^* + 1$, we immediately obtain that $d(x, s) = d(y, s) - 1$ (i.e., $x \in I(y, s)$). Now, suppose that $r_y^* = r_x^*$. Since $d(y, s) = r_y^*$, we conclude that either $d(x, s) = d(y, s) - 1$ (i.e., $x \in I(y, s)$) or $d(x, s) = d(y, s)$. This establishes the last equality in each of the four cases.

It remains to prove that, if $r_y^* = r_x^* + 1$, then we are either in Case (1) or (2), and if $r_y^* = r_x^*$, then we are either in Case (3) or (4). Indeed, if $r_y^* = r_x^* + 1$, then $d(y, s) = d(x, s) + 1 = r_x^* + 1$, and we are in Case (1) or (2). On the other hand, if $r_x^* = r_y^*$, then $d(x, z) = d(y, z) + 1$ since $z \notin B_{r_x^*}(x)$, and we are in Case (3) or (4). \square

Without the knowledge of r_x^* and r_y^* , the relationships between $d(x, z)$ and $d(y, z)$, and between $d(x, s)$ and $d(y, s)$ do not allow us to distinguish between the cases (1) and (4). This can be done by additionally using the vertex $t \in X^+$ defined above. Indeed, in Case (1) we have $t = s$, while in Case (4) we have $t \neq s$ and $d(x, t) = d(y, t) + 1$. We continue with the following simple lemma for paths (where, for each edge xy in the path, x is to the left of y):

Lemma 15. *Let Q be a graph which is a path with end-vertices $a \neq b$, and let d' be its distance function. Then, Q contains a unique edge x_0y_0 such that $d'(x_0, b) - d'(x_0, a) \in \{1, 2\}$.*

Proof. Let $\|Q\|$ be the solid path obtained from Q by replacing each edge of Q by a solid segment of unit length ($\|Q\|$ is the line-segment of length $d'(a, b)$). Denote the distance function on $\|Q\|$ by d' as well. Consider the function $f(q) := d'(q, b) - d'(q, a)$ defined on $\|Q\|$. Since the distances are measured on the segment $\|Q\|$, the function f is continuous and takes integer values on the vertices of Q and on the middles of the edges. Moreover, while moving on $\|Q\|$ from b to a , the function f is strictly increasing. By the intermediate value theorem, $\|Q\|$ contains a unique point q_0 such that $f(q_0) = 0$.

First, suppose that q_0 is an interior point of the edge x_0y_0 , where y_0 is closer to b than x_0 is. Then, q_0 is the middle of this edge, and $d'(x_0, b) = d'(y_0, b) + 1$ and $d'(x_0, a) = d'(y_0, a) - 1$. Since $f(q_0) = d'(q_0, b) - d'(q_0, a) = 0$, then $f(x_0) = 1$, and the edge x_0y_0 is such that $d'(x_0, b) - d'(x_0, a) = 1$. Pick any other edge x_1y_1 of Q . If x_1y_1 is located between a and x_0 , then $d'(x_1, a) \leq d'(x_0, a) - 1$ and $d'(x_1, b) \geq d'(x_0, b) + 1$, leading to $d'(x_1, b) - d'(x_1, a) \geq 3$. If x_1y_1 is located between b and y_0 , then $d'(x_1, a) \geq d'(x_0, a) + 1$ and $d'(x_1, b) \leq d'(x_0, b) - 1$, leading to $d'(x_1, b) - d'(x_1, a) \leq -1$.

Now, suppose that q_0 is a vertex of Q . Since $a \neq b$, we have that $q_0 \notin \{a, b\}$. Then, set $y_0 = q_0$, and let x_0 be the neighbor of y_0 located between a and y_0 . Since $d'(y_0, b) = d'(y_0, a)$, we obtain that $d'(x_0, b) - d'(x_0, a) = 2$. Similarly to the previous case, we can show that, for all other edges x_1y_1 of Q , we have that $d'(x_1, b) - d'(x_1, a)$ is either at least 4 or at most -2 . \square

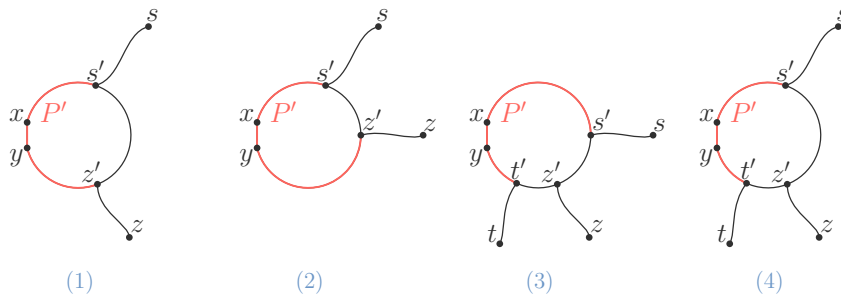


FIGURE 3. Definition and positioning of s , t , and z in the four cases of Lemma 14.

We use Lemma 15 to find adjacent vertices x_0 and y_0 of C and an integer r_x^* that satisfy a condition of Lemma 14. Let P' be the path between z' and s' (or between s' and t') containing the edge xy as defined above. Let P be the path of G obtained by joining the shortest (s', s) - and (z, z') -paths ((s, s') - and (t', t) -paths, resp.) to P' . Let d' be the distance function on P .

Lemma 16. *Let P be the (s, z) -path or (s, t) -path of G defined above. Let $x_0 y_0$ be the unique edge of P satisfying the conclusion of Lemma 15. Then, $x_0 = x$ and $y_0 = y$. Moreover,*

- (1) *if P is an (s, z) -path, $d'(x_0, z) = d(x_0, z)$, and $d'(y_0, s) = d(y_0, s)$, then $r_x^* = d(y_0, s) - 1$;*
- (2) *if P is an (s, z) -path, $d'(x_0, z) = d(x_0, z) + 1$, and $d'(y_0, s) = d(y_0, s)$, then $r_x^* = d(y_0, s) - 1$;*
- (3) *if P is an (s, z) -path, $d'(x_0, z) = d(x_0, z)$, and $d'(y_0, s) = d(y_0, s) + 1$, then $r_x^* = d(y_0, s)$;*
- (4) *if P is an (s, t) -path, then $r_x^* = d(x_0, t)$.*

Proof. By its definition, the edge xy satisfies one of the four conditions of Lemma 14. First, suppose that xy satisfies Condition (1). Then, $d(x, s) = r_y^* - 1$. Since $z \in B_{r_y^*}(y) \setminus B_{r_x^*}(x)$ and $r_y^* = r_x^* + 1$, we conclude that $r_y^* \leq d(x, z) \leq r_y^* + 1$. Since $d'(x, z) = d(x, z)$ and $d'(x, s) = d(x, s)$, we conclude that $d'(x, z) - d'(x, s) = d(x, z) - d(x, s) \in \{1, 2\}$.

Now, suppose that xy satisfies condition (2). Then, $r_y^* = d(y, z)$ because $z \in B_{r_y^*}(y) \setminus B_{r_x^*}(x)$ and $d(x, z) = d(y, z) \leq r_y^* = r_x^* + 1$. Since $d'(x, z) = d(x, z) + 1$ and $d'(x, s) = d(x, s)$, we obtain that

$$\begin{aligned} d'(x, z) - d'(x, s) &= d(x, z) + 1 - d(x, s) \\ &= d(y, z) - (d(y, s) - 1) + 1 \\ &= r_y^* - r_x^* + 2 = 2. \end{aligned}$$

Now, suppose that xy satisfies Condition (3). Then, $d'(y, s) = d(y, s) + 1 = d(x, s) + 1$, $d'(x, z) = d(x, z)$, and $d'(x, s) = d(x, s)$. As in the previous case, we can show that $r_y^* = d(y, z)$. Since $r_x^* = r_y^*$ and $d(x, z) = d(y, z) + 1 = r_y^* + 1$, we conclude that

$$\begin{aligned} d'(x, z) - d'(x, s) &= d(x, z) - d(x, s) \\ &= r_x^* + 1 - r_x^* = 1. \end{aligned}$$

Since in Cases (1)–(3) of Lemma 14, we have $d'(x, z) - d'(x, s) \in \{1, 2\}$, then the edge xy satisfies the conclusion of Lemma 15. Therefore, $x = x_0$ and $y = y_0$. We can easily check that in the first two cases, $r_x^* = d(y_0, s) - 1$, and in the third case, $r_x^* = d(y_0, s)$.

Finally, if xy satisfies Condition (4) of Lemma 14, then $d'(x, z) = d(x, z)$ and $d'(x, s) = d(x, s)$, and, as in Case (1), we conclude that $x = x_0$ and $y = y_0$. In this case, the path P is an (s, t) -path, and therefore, $r_x^* = d(x_0, t)$ by the definition of t . \square

The compressor $\alpha(X)$. The compressor $\alpha(X)$ is a vector with six coordinates, which are grouped into three pairs: $\alpha(X) := (\alpha_1(X), \alpha_2(X), \alpha_3(X))$. The pair $\alpha_1(X) \subseteq X^+$ is a diametral pair (u^+, v^+) of X^+ , $\alpha_2(X)$ is used to specify the region of $C(u^+, v^+)$ where the center of the target ball is located, and the pair $\alpha_3(X)$ is used to compute the radius of this ball. We use the symbol $*$ to indicate that the respective coordinate of $\alpha(X)$ is empty.

We continue with the definitions of $\alpha_2(X)$ and $\alpha_3(X)$. First, suppose that $X_C = \emptyset$, *i.e.*, $X_u \cup X_v = X$ and $X_u^- \cup X_v^- = X^-$. Then, set $\alpha_2(X) := (w_1, w_2)$ and $\alpha_3(X) := (z_1, z_2)$. Now, suppose that $X_C \neq \emptyset$. Let w be a vertex of X whose gate w' in $C(u^+, v^+)$ belongs to C . If $B_{r_x^*}(x)$ realizes X for any vertex x of C , then set $\alpha_2(X) := (w, *)$ and $\alpha_3(X) := (s, *)$, where $s \in X^+$ is such that $d(w', s) = r_{w'}^*$. Otherwise, we pick an edge xy of C such that $B_{r_x^*}(x)$ realizes X and $B_{r_y^*}(y)$ does not realize X . Let s' , t' , and z' be the respective gates in C of the vertices s , t , and z as defined previously. If $s = t$, then the path P is defined by the vertices s and z , and set $\alpha_3(X) := (s, z)$. Otherwise, the path P is defined by the vertices s and t , and set $\alpha_3(X) := (s, t)$. Moreover, set $\alpha_2(X) := (*, w)$ if the edge xy belongs to the path from s' to z' (from s' to t' , resp.) in the clockwise traversal of C , and $\alpha_2(X) := (w, *)$ otherwise. Formally, the compressor function α is defined in the following way:

- (C1) if $X^- = \emptyset$, set $\alpha_1(X) = \alpha_2(X) = \alpha_3(X) := (*, *)$;
- (C2) otherwise, if $|X^+| = 0$, set $\alpha_1(X) = \alpha_2(X) := (*, *)$ and $\alpha_3(X) := (z, *)$, where z is an arbitrary vertex of X^- ;

- (C3) otherwise, if $X^+ = \{u\}$, set $\alpha_1(X) := (u, *)$, $\alpha_2(X) := (*, *)$, and $\alpha_3(X) := (z, *)$, where z is an arbitrary vertex of X^- ;
- (C4) otherwise, if $|X^+| \geq 2$ and $X_C = \emptyset$, set $\alpha_1(X) := (u^+, v^+)$, $\alpha_2(X) := (w_1, w_2)$, and
 - (C4i) if the vertex z_2 does not exist, then set $\alpha_3(X) := (z_1, *)$;
 - (C4ii) if the vertex z_1 does not exist, then set $\alpha_3(X) := (*, z_2)$;
 - (C4iii) if the vertices z_1 and z_2 exist, set $\alpha_3(X) := (z_1, z_2)$;
- (C5) otherwise ($|X^+| \geq 2$ and $X_C \neq \emptyset$), and
 - (C5i) if, for any vertex $y \in C$, the ball $B_{r_y^*}(y)$ realizes X , then set $\alpha_1(X) := (u^+, v^+)$, $\alpha_2(X) := (w, *)$, and $\alpha_3(X) := (s, *)$, where $s \in X^+$ is such that $d(w', s) = r_{w'}^*$;
 - (C5ii) otherwise, if s and z are given, and the edge xy belongs to the clockwise (s', z') -path of C , then set $\alpha_2(X) := (*, w)$ and $\alpha_3(X) := (s, z)$;
 - (C5iii) otherwise, if s and z are given, and the edge xy belongs to the counterclockwise (s', z') -path of C , then set $\alpha_2(X) := (w, *)$ and $\alpha_3(X) := (s, z)$;
 - (C5iv) otherwise, if s and t are given, and the edge xy belongs to the clockwise (s', t') -path of C , then set $\alpha_2(X) := (*, w)$ and $\alpha_3(X) := (s, t)$;
 - (C5v) otherwise, if s and t are given, and the edge xy belongs to the counterclockwise (s', t') -path of C , then set $\alpha_2(X) := (w, *)$ and $\alpha_3(X) := (s, t)$.

The reconstructor $\beta(X)$. Let Y be a vector on six coordinates grouped into three pairs Y_1 , Y_2 , and Y_3 . If $Y_1 = (y_1, y_2)$, then, for any vertex t of G , we denote by t' its gate in $C(y_1, y_2)$. For any vertex y of $C(y_1, y_2)$, we also set $r_y := \max\{d(y, y_1), d(y, y_2)\}$. The reconstructor β takes Y and returns a ball $B_r(y)$ of G defined in the following way:

- (R1) if $Y = ((*, *), (*, *), (*, *))$, then $\beta(Y)$ is any ball that contains the vertex set of G ;
- (R2) if $Y = ((*, *), (*, *), (y_5, *))$, then $\beta(Y)$ is the empty set;
- (R3) if $Y = ((y_1, *), (*, *), (y_5, *))$, then $\beta(Y)$ is the ball $B_0(y_1)$;
- (R4) if $Y_1 = (y_1, y_2)$ and $Y_2 = (y_3, y_4)$, then let u_0 be the cut vertex of $C(y_3')$ between y_3' and y_2 , and v_0 be the cut vertex of $C(y_4')$ between y_4' and y_1 . Then, $\beta(Y)$ is any ball $B_{r_y}(y)$ centered at $y \in C(u_0, v_0)$ such that $B_{r_y}(y)$ contains no vertex of Y_3 .
- (R5i) if $Y = ((y_1, y_2), (y_3, *), (y_5, *))$, then $\beta(Y)$ is the ball $B_r(y_3')$ of radius $r = d(y_3', y_5)$;
- (R5ii) if $Y = ((y_1, y_2), (*, y_4), (y_5, y_6))$ and $(y_5, y_6) \in X^+ \times X^-$, let xy be the edge of the (y_5', y_6') -path in the clockwise traversal of the cycle $C(y_4')$ such that $|d'(x, y_6) - d'(x, y_5)| \in \{1, 2\}$ and y is closer to y_6 than x is. Let $\beta(Y)$ be the ball $B_r(x)$, where $r = d(y, y_5)$ if $d'(y, y_5) = d(y, y_5) + 1$, and $r = d(y, y_5) - 1$ otherwise;
- (R5iii) if $Y = ((y_1, y_2), (y_3, *), (y_5, y_6))$ and $(y_5, y_6) \in X^+ \times X^-$, let xy be the edge of the (y_5', y_6') -path in the counterclockwise traversal of $C(y_3')$ such that $|d'(x, y_6) - d'(x, y_5)| \in \{1, 2\}$ and y is closer to y_6 than x is. Let $\beta(Y)$ be the ball $B_r(x)$, where $r = d(y, y_5)$ if $d'(y, y_5) = d(y, y_5) + 1$, and $r = d(y, y_5) - 1$ otherwise;
- (R5iv) if $Y = ((y_1, y_2), (*, y_4), (y_5, y_6))$ and $(y_5, y_6) \in X^+ \times X^+$, let xy be the edge of the (y_5', y_6') -path in the clockwise traversal of the cycle $C(y_4')$ such that $|d'(x, y_6) - d'(x, y_5)| \in \{1, 2\}$ and y is closer to y_6 than x is. Let $\beta(Y)$ be the ball $B_r(x)$, where $r = d(x, y_6)$;
- (R5v) if $Y = ((y_1, y_2), (y_3, *), (y_5, y_6))$ and $(y_5, y_6) \in X^+ \times X^+$, let xy be the edge of the (y_5', y_6') -path in the counterclockwise traversal of the cycle $C(y_3')$ such that $|d'(x, y_6) - d'(x, y_5)| \in \{1, 2\}$ and y is closer to y_6 than x is. Let $\beta(Y)$ be the ball $B_r(x)$, where $r = d(x, y_6)$;

Proposition 17. *For any tree of cycles G , the pair (α, β) of vectors defines a proper labeled sample compression scheme of size 6 for $\mathcal{B}(G)$.*

Proof. Let X be a realizable sample for \mathcal{B} . Let $Y = \alpha(X)$ and $B_r(x^*) = \beta(Y)$. We will prove case by case that the ball $B_r(x^*)$ realizes the sample X , *i.e.*, that $X^+ \subseteq B_r(x^*)$ and $X^- \cap B_r(x^*) = \emptyset$. One can easily see that the cases (Rk) and their subcases in the definition of β correspond to the cases (Ck) and their subcases in the definition of α : namely, the vector Y in Case (Rk) has the same specified coordinates as the vector $\alpha(X)$ in Case (Ck).

In Case (R1), we have $Y = ((*, *), (*, *), (*, *))$. Since $Y = \alpha(X)$, this implies that $|X^-| = 0$, which corresponds to Case (C1). Consequently, the ball covering G is compatible with X . In Case (R2), $Y = ((*, *), (*, *), (y_5, *))$. Since $|Y^+| = 0$ and $|Y^-| \neq 0$, this implies that $|X^+| = 0$

and $|X^-| \neq 0$, which corresponds to Case (C2). In this case, the empty set is compatible with X . In Case (R3), $Y = ((y_1, *), (*, *), (y_5, *))$. The fact that $|Y_1| = 1$ implies that $|X^+| = 1$. Thus, the ball $B_0(y_1)$ is compatible with X .

Consider now Case (R4). In this case, $Y_1 = (y_1, y_2)$ and $Y_2 = (y_3, y_4)$. Since $Y = \alpha(X)$, this implies that X satisfies the conditions of Case (C4), *i.e.*, $|X^+| \geq 2$ and $X_C = \emptyset$. Therefore, $Y_1 = (y_1, y_2) = (u^+, v^+) = \alpha_1(X)$, $Y_2 = (y_3, y_4) = (w_1, w_2) = \alpha_2(X)$, and Y_3 (containing one or two vertices) coincides with $\alpha_3(X)$ (containing one or two vertices z_1, z_2 as in subcases (C4i)–(C4iii)). The ball $B_{r_y}(y)$ returned by Case (R4) is centered at $y \in C(u_0, v_0)$, contains Y_1 , and is disjoint from Y_3 . Since the target ball $B_r(x)$ has its center on $C(u_0, v_0)$ and is compatible with $X \supseteq Y_1 \cup Y_3$, the ball $B_{r_y}(y)$ is well-defined. By Lemma 12, $B_{r_y}(y)$ contains X^+ . By Lemma 13, $B_{r_y}(y)$ is disjoint from X^- . Consequently, $B_{r_y}(y)$ is compatible with X .

In Case (R5), we have $Y_2 = (*, y_4)$ or $Y_2 = (y_3, *)$. This distinguishes the Case (R5) from the Cases (R1)–(R4). If $Y_3 = (y_5, *)$, then we are in Case (R5i), which corresponds to Case (C5i). In this case, from the definition of y'_3 and y_5 , we conclude that $r_{y'_3}^* = d(y'_3, y_5)$, and thus, the ball $\beta(Y) = B_{r_{y'_3}^*}(y'_3)$ realizes X . Each subcase (R5ii)–(R5v) corresponds to the respective subcase (C5ii)–(C5v), and its analysis is based on Lemma 16. If $Y_3 = (y_5, y_6) \in X^+ \times X^+$, then we are in Case (R5iv) or (R5v), and also in Case (4) of Lemma 16. Therefore, $\beta(Y) = B_r(x)$ with $r = d(x, y_6)$ realizes X . If $Y_3 = (y_5, y_6) \in X^+ \times X^-$, then we are in Case (R5ii) or (R5iii), and in Cases (1)–(3) of Lemma 16. If $d'(y, y_5) = d(y, y_5) + 1$, then we are in Case (3) of Lemma 16. Therefore, by Lemma 16, $r_x^* = d(y, y_5) =: r$, and the ball $\beta(Y) = B_r(x)$ realizes X . Otherwise, $d'(y, y_5) = d(y, y_5)$, and we are in Case (1) or (2) of Lemma 16. In these cases, if we set $r = d(y, y_5) - 1$, then the ball $B_r(x)$ realizes X . \square

Remark 18. The most technically involved case of the previous result is the case $X_C \neq \emptyset$. In fact, this case corresponds to proper labeled sample compression schemes in *spiders*, *i.e.*, in graphs consisting of a single cycle C and paths of different lengths emanating from this cycle. Due to this case, $\alpha(X)$ in our result is not a signed map but a signed vector of size 6. Thus, in this case, we need extra information compared to the initial definition of proper labeled sample compression schemes. The VC-dimension of the family of balls in a spider and in a tree of cycles is 3. We wonder *whether the family of balls in spiders admits a proper labeled sample compression scheme without any information that is of (a) size 3 or (b) constant size.*

6. CUBE-FREE MEDIAN GRAPHS

The *dimension* $\dim(G)$ of a median graph G is the largest dimension of a hypercube of G . A *cube-free median graph* is a median graph of dimension 2, *i.e.*, a median graph not containing 3-cubes as isometric subgraphs. For references about median graphs, see [1]. For cube-free median graphs, see [2, 10, 11, 13]. We use the fact that intervals of median graphs are gated. We describe a proper LSCS of size 22 for balls of cube-free median graphs.

Let G be a cube-free median graph. Let X be a realizable sample for $\mathcal{B}(G)$, and $\{u^+, v^+\}$ a diametral pair of X^+ . The next lemma shows that the center of a ball realizing X can always be found in $I(u^+, v^+)$ (this result does not hold for all median graphs):

Lemma 19. *If x' is the gate of x in the interval $I(u^+, v^+)$, and $r' = r - d(x, x')$, then X is a realizable sample for $B_{r'}(x')$, *i.e.*, $X^+ \subseteq B_{r'}(x')$ and $X^- \cap B_{r'}(x') = \emptyset$.*

Proof. Let $B' := B_{r'}(x')$. From the definitions of x' and r' , it immediately follows that $B' \subseteq B_r(x) = B$, implying that $X^- \cap B' = \emptyset$. Notice also that $d(x', u^+) \leq r'$ and $d(x', v^+) \leq r'$. To prove the inclusion $X^+ \subseteq B'$, suppose by way of contradiction that there exists $z \in X^+ \setminus B'$. Let z' be the gate of z in $I(u^+, v^+)$. First, suppose that $z' \in I(x, z)$. Since $x' \in I(x, z')$, the vertices z' and x' belong to a common shortest path between z and x . Consequently,

$$d(x, z) = d(x, x') + d(x', z') + d(z', z).$$

Since $d(x, z) \leq r$, we conclude that $d(x', z) \leq r - d(x, x') = r'$. The same conclusion holds if $x' \in I(x, z)$. This implies that the median of x, z , and z' is a vertex $z'' \neq z'$.

Now, suppose that the vertices x' and z' belong to a common shortest path between u^+ and v^+ . Say $z' \in I(u^+, x')$ and $x' \in I(z', v^+)$. Since $d(x', u^+) \leq r'$ and $d(x', z) > r'$, we deduce that $d(x', u^+) < d(x', z)$. Since $z' \in I(x', u^+) \cap I(x', z)$, we deduce that $d(z, z') > d(u^+, z')$. Since $z' \in I(v^+, u^+) \cap I(v^+, z)$, we conclude that $d(v^+, z) > d(v^+, u^+)$, contradicting the choice of the pair $\{u^+, v^+\}$ as a diametral pair of X^+ .

Finally, suppose that x' and z' do not belong to a common shortest path between u^+ and v^+ . This is equivalent to the assertion that the median u' of the triplet $\{u^+, x', z'\}$ is different from x' and z' , and that the median v' of the triplet $\{v^+, x', z'\}$ is different from x' and z' . Notice that u' and v' belong to a common shortest path between u^+ and v^+ . Let p be a neighbor of z' in $I(z', u')$ and let q be a neighbor of z' in $I(z', v')$. If $p = q$, then this vertex is the median of (z', u', v') and of (z', u^+, v^+) , contrary to the assumption that $z' \in I(u^+, v^+)$. It follows that $p \neq q$. Let s be a neighbor of z' in $I(z', z'')$. Since z' is the gate of z in $I(u^+, v^+)$, s does not belong to $I(u^+, v^+)$, and is thus distinct from p and q . From the definitions of the vertices x' , z' , and z'' , we conclude that the vertices p , q , and s belong to the interval $I(z', x)$ and therefore belong to a 3-cube of G . This contradicts the fact that G is a cube-free median graph. \square

By [13], $I(u^+, v^+)$ of a cube-free median graph has an isometric embedding in the square grid \mathbb{Z}^2 . We denote by (z_a, z_b) the coordinates in \mathbb{Z}^2 of a vertex $z \in I(u, v)$. We consider isometric embeddings of $I(u, v)$ in \mathbb{Z}^2 for which $u = (0, 0)$ and $v = (v_a, v_b)$ with $v_a \geq 0$ and $v_b \geq 0$. We fix a canonical isometric embedding, which can be used both by the compressor and the reconstructor. Finally, we use the same notation for the vertices and their images under the embedding, and we denote by \mathbf{I} the interval $I(u^+, v^+)$ embedded in \mathbb{Z}^2 . As usual, for a vertex $z \in V$, we denote by z' its gate in the interval $I(u^+, v^+)$.

The compressor $\alpha(X)$. The compressor $\alpha(X)$ is a vector with 22 coordinates grouped into four parts $\alpha(X) := (\alpha_1(X), \alpha_2(X), \alpha_3(X), \alpha_4(X))$. The part $\alpha_1(X) \subseteq X^+$ consists of a diametral pair (u^+, v^+) of X^+ . The part $\alpha_2(X) \subseteq X$ has size 4, and is used to specify a region $\mathbf{R} \subseteq \mathbf{I} = I(u^+, v^+)$ such that the gates in $I(u^+, v^+)$ of all the vertices of X are located outside or on the boundary of \mathbf{R} . Moreover, \mathbf{R} contains the center x of the target ball $B_r(x)$. The parts $\alpha_3(X) \subseteq X^+$ and $\alpha_4(X) \subseteq X^-$ each have size 8 and are used to locate the center and the radius of a ball $B_{r''}(y)$ realizing X . Now, we formally define $\alpha_i(X)$, $i = 1, \dots, 4$. Let $X_1 := \{w \in X : w'_b \geq x_b\}$, $X_2 := \{w \in X : w'_a \geq x_a\}$, $X_3 := \{w \in X : w'_b \leq x_b\}$, and $X_4 := \{w \in X : w'_a \leq x_a\}$. Since $I(u^+, v^+)$ is gated, $X = \cup_{i=1}^4 X_i$. Denote by X'_i , $i = 1, \dots, 4$, the gates of the vertices of X_i in $I(u^+, v^+)$. Set $\alpha_2(X) := (w_1, w_2, w_3, w_4) \in X^4$, where:

- w_1 is a vertex of X_1 whose gate w'_1 has the smallest ordinate among the vertices of X'_1 ;
- w_2 is a vertex of X_2 whose gate w'_2 has the smallest abscissa among the vertices of X'_2 ;
- w_3 is a vertex of X_3 whose gate w'_3 has the largest ordinate among the vertices of X'_3 ;
- w_4 is a vertex of X_4 whose gate w'_4 has the largest abscissa among the vertices of X'_4 ;

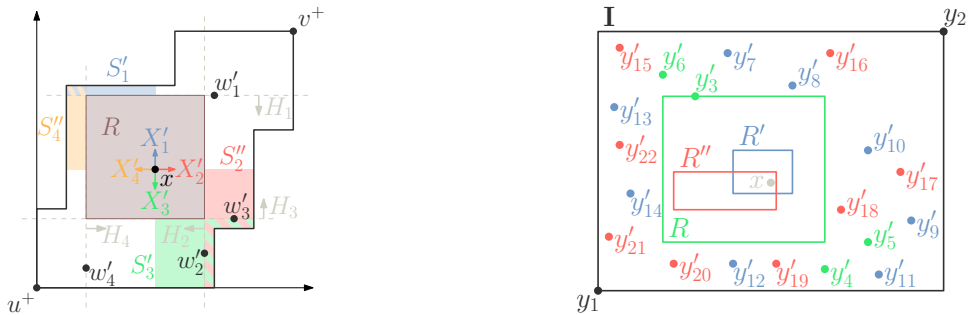


FIGURE 4. On the left, the region \mathbf{R} and the halfstrips $\mathbf{S}'_1(x)$, $\mathbf{S}''_2(x)$, $\mathbf{S}'_3(x)$, and $\mathbf{S}''_4(x)$. On the right, the regions \mathbf{R} , \mathbf{R}' , and \mathbf{R}'' computed from $\alpha(X)$. Steps 1-4 of the reconstruction correspond to the black, green, blue, and red parts of the figure. The target center x is given in gray.

For a vertex $w = (w_a, w_b) \in \mathbb{Z}^2$, we consider the four coordinate halfplanes $\mathbf{H}_{\leq w_a} := \{t : t_a \leq w_a\}$, $\mathbf{H}_{\geq w_a}$, $\mathbf{H}_{\leq w_b}$, and $\mathbf{H}_{\geq w_b}$. Let \mathbf{R} be the set of vertices of \mathbf{I} that belong to the intersection of the halfplanes $\mathbf{H}_1 := \mathbf{H}_{\leq w_{1b}}$, $\mathbf{H}_2 := \mathbf{H}_{\leq w_{2a}}$, $\mathbf{H}_3 := \mathbf{H}_{\geq w_{3b}}$, and $\mathbf{H}_4 := \mathbf{H}_{\geq w_{4a}}$. If a vertex w_i does not exist, then the corresponding halfplane \mathbf{H}_i is not defined. From the definition, the inside of \mathbf{R} does not contain gates of vertices of X . We denote by \mathbf{S}_i , $i = 1, \dots, 4$, the intersection of \mathbf{I} with the closure of the complementary halfspace of \mathbf{H}_i . We call \mathbf{S}_i , $i = 1, \dots, 4$, a *strip of \mathbf{I}* . Consequently, the interval \mathbf{I} is covered by the region \mathbf{R} , two *horizontal* strips \mathbf{S}_1 and \mathbf{S}_3 , and two *vertical* strips \mathbf{S}_2 and \mathbf{S}_4 . Using this notation, we can redefine X_i as the sets of all the vertices of X whose gate in \mathbf{I} belongs to the strip \mathbf{S}_i . Consequently, $X_i' \subseteq \mathbf{S}_i$. Furthermore, for a vertex $z \in \mathbb{Z}^2$, each strip \mathbf{S}_i is partitioned into two strips $\mathbf{S}_i'(z)$ and $\mathbf{S}_i''(z)$ by the vertical or horizontal line passing via z . The labeling of the strips is done in the clockwise order around z , see Fig. 4 (left). Let $\alpha_3(X) := (s_1, t_1, s_2, t_2, s_3, t_3, s_4, t_4)$, where

- s_1 is a vertex of X^+ furthest from x , whose gate s_1' belongs to $\mathbf{S}_1'(x)$, and t_1 is a vertex of X^+ such that its gate t_1' belongs to $\mathbf{S}_1''(x)$ and the abscissa of t_1' is closest to x_a ;
- s_2 is a vertex of X^+ furthest from x , whose gate s_2' belongs to $\mathbf{S}_2'(x)$, and t_2 is a vertex of X^+ such that its gate t_2' belongs to $\mathbf{S}_2''(x)$ and the ordinate of t_2' is closest to x_b ;
- s_3 is a vertex of X^+ furthest from x , whose gate s_3' belongs to $\mathbf{S}_3'(x)$, and t_3 is a vertex of X^+ such that its gate t_3' belongs to $\mathbf{S}_3''(x)$ and the abscissa of t_3' is closest to x_a ;
- s_4 is a vertex of X^+ furthest from x , whose gate s_4' belongs to $\mathbf{S}_4'(x)$, and t_4 is a vertex of X^+ such that its gate t_4' belongs to $\mathbf{S}_4''(x)$ and the ordinate of t_4' is closest to x_b .

Let $\alpha_4(X) := (p_1, q_1, p_2, q_2, p_3, q_3, p_4, q_4)$, where p_i is a vertex of X^- closest to x , whose gate p_i' belongs to $\mathbf{S}_i'(x)$, and q_i is a vertex of X^- closest to x , whose gate q_i' belongs to $\mathbf{S}_i''(x)$. If any of the vertices of the four groups is not defined, then its corresponding coordinate in $\alpha(X)$ is set to $*$.

The reconstructor $\beta(Y)$. Let Y be a vector of 22 coordinates corresponding to a realizable sample and grouped into four parts $Y_1 := (y_1, y_2)$, $Y_2 := (y_3, y_4, y_5, y_6)$, $Y_3 := (y_7, y_8, y_9, y_{10}, y_{11}, y_{12}, y_{13}, y_{14})$, and $Y_4 := (y_{15}, y_{16}, y_{17}, y_{18}, y_{19}, y_{20}, y_{21}, y_{22})$. The reconstructor $\beta(Y)$ returns a ball $B_{r''}(y)$ by performing the following steps (see Fig. 4 (right)):

- (1) Using Y_1 , canonically isometrically embed $I(y_1, y_2)$ into \mathbb{Z}^2 as \mathbf{I} .
- (2) Using Y_2 , compute the gates y_i' of y_i in \mathbf{I} and compute the region \mathbf{R} as the intersection of the halfplanes $\mathbf{H}_{\leq y_{1b}}$, $\mathbf{H}_{\leq y_{2a}}$, $\mathbf{H}_{\geq y_{3b}}$, and $\mathbf{H}_{\geq y_{4a}}$ with \mathbf{I} .
- (3) Using Y_3 , compute the set $\mathbf{R}' \subseteq \mathbf{R}$ of all $y = (y_a, y_b) \in \mathbf{R}$ such that the gates $y_7', y_8', y_9', y_{10}', y_{11}', y_{12}', y_{13}', y_{14}'$ belong to $\mathbf{S}_1'(y)$, $\mathbf{S}_1''(y)$, $\mathbf{S}_2'(y)$, $\mathbf{S}_2''(y)$, $\mathbf{S}_3'(y)$, $\mathbf{S}_3''(y)$, $\mathbf{S}_4'(y)$, $\mathbf{S}_4''(y)$, resp. For each $y \in \mathbf{R}'$, let r_y' be the *smallest* radius such that $Y_1 \cup Y_3 \subseteq B_{r_y'}(y)$.
- (4) Using Y_4 , compute the region $\mathbf{R}'' \subseteq \mathbf{R}$ consisting of all the vertices $y \in \mathbf{R}$ such that the gates $y_{15}', y_{16}', \dots, y_{21}', y_{22}'$ belong to the strips $\mathbf{S}_1'(y)$, $\mathbf{S}_1''(y)$, \dots , $\mathbf{S}_4'(y)$, $\mathbf{S}_4''(y)$, respectively. For each $y \in \mathbf{R}''$, let r_y'' be the *largest* radius such that $B_{r_y''}(y) \cap Y_4 = \emptyset$.
- (5) Let $\mathbf{R}_0 := \{y \in \mathbf{R}' \cap \mathbf{R}'' : r_y'' \geq r_y'\}$ and return as $\beta(Y)$ any ball $B_{r_y''}(y)$ with $y \in \mathbf{R}_0$.

Proposition 20. *For any cube-free median graph G , the pair (α, β) of vectors defines a proper labeled sample compression scheme of size 22 for $\mathcal{B} = \mathcal{B}(G)$.*

Proof. Let X be a realizable sample for \mathcal{B} , and let $B_r(x)$ be a ball realizing X . Let $Y := \alpha(X)$. We suppose that $|X^+| \geq 2$, as otherwise we are trivially done. By the definition of the maps α and β , we deduce that x belongs to the regions \mathbf{R}' and \mathbf{R}'' defined in Steps 3 and 4. Also, $r_x' \leq r \leq r_x''$, and therefore, $x \in \mathbf{R}_0$. This implies that $\mathbf{R}_0 \neq \emptyset$ and the map $\beta(Y)$ is well-defined. Denote by \mathbf{S}_1 , \mathbf{S}_2 , \mathbf{S}_3 , and \mathbf{S}_4 the strips defined by the gates y_3' , y_4' , y_5' , and y_6' of the vertices of Y_2 . Since $Y_2 = \alpha_2(X)$, all the vertices of X have their gates in the union $\bigcup_{i=1}^4 \mathbf{S}_i$ of the four strips. Now, suppose that $\beta(Y) = B_{r_y''}(y)$.

Claim 21. *For any $y \in \mathbf{R}'$, $X^+ \subseteq B_{r_y'}(y)$.*

Proof. Pick any $z \in X^+$. Assume, without loss of generality, that its gate z' in \mathbf{I} belongs to the strip \mathbf{S}_1 .

Case 1. $z' \in \mathbf{S}_1''(y)$.

In this case, $d(u^+, z) = d(u^+, z') + d(z', z)$ and $d(u^+, v^+) = d(u^+, z') + d(z', v^+)$. Also, $d(y, z) = d(y, z') + d(z', z)$ and $d(y, v^+) = d(y, z') + d(z', v^+)$. If $z \notin B_{r_y'}(y)$, then since $v^+ \in B_{r_y'}(y)$, we conclude that $d(z', z) > d(z', v^+)$, yielding $d(u^+, z) > d(u^+, v^+)$, a contradiction with the choice of $\{u^+, v^+\}$ as a diametral pair of X^+ .

Case 2. $z' \in \mathbf{S}_1'(y)$.

Subcase 2.1. $z' \in \mathbf{S}_1''(x)$.

In this case, the vertex $y_8 \in \mathbf{S}_1''(y)$ exists and $y'_{8a} \leq z'_a$. Since $z \in \mathbf{S}_1''(x)$, we have $y_a \geq z'_a$. On the other hand, by the construction of \mathbf{R}' , $y_a \leq y'_{8a}$. Consequently, $y_a \leq y'_{8a} \leq z'_a \leq y_a$, showing that $y_a = y'_{8a} = z'_a$. Consequently, $z' \in \mathbf{S}_1''(y)$, which is the condition of Case 1.

Subcase 2.2. $z' \in \mathbf{S}_1'(x)$.

In this case, the vertex y_7 exists. By its definition, y_7 belongs to $\mathbf{S}_1'(x)$, and by the construction of \mathbf{R}' , y_7 belongs to $\mathbf{S}_1'(y)$. If $y_a \leq x_a$, then let t be the vertex on the upper side of \mathbf{R} such that $t_a = y_a$. Analogously, if $y_a \geq x_a$, then let t be the vertex on the upper side of \mathbf{R} such that $t_a = x_a$. Notice that, in both cases, we have $t \in I(x, z) \cap I(x, y_7)$ and $t \in I(y, z) \cap I(y, y_7)$. From the definition of the vertex y_7 , we have $d(x, y_7) \geq d(x, z)$, yielding $d(t, y_7) \geq d(t, z)$. As a result, we obtain that $z \in B_{r_y'}(y)$. \square

Claim 22. For any $y \in \mathbf{R}''$, $X^- \cap B_{r_y''}(y) = \emptyset$.

Proof. Let z be any vertex of X^- . Suppose, without loss of generality, that its gate z' belongs to $\mathbf{S}_1'(x)$. This implies that the vertex y_{15} exists. By the definition of y_{15} and the construction of \mathbf{R}'' , we have $y_{15} \in \mathbf{S}_1'(x) \cap \mathbf{S}_1'(y)$. We distinguish two cases.

Case 1. $z' \in \mathbf{S}_1'(y)$.

If $y_a \leq x_a$, then let t be the vertex on the upper side of \mathbf{R} such that $t_a = y_a$. Analogously, if $y_a \geq x_a$, then let t be the vertex on the upper side of \mathbf{R} such that $t_a = x_a$. In both cases we have $t \in I(y, z) \cap I(y, y_{15})$ and $t \in I(x, z) \cap I(x, y_{15})$. Since y_{15} is the vertex of X^- closest to x whose gate is in $\mathbf{S}_1'(x)$, we must have $d(x, y_{15}) \leq d(x, z)$. This implies that $d(t, y_{15}) \leq d(t, z)$. Consequently, $d(y, y_{15}) \leq d(y, z)$, and thus, $z \notin B_{r_y''}(y)$.

Case 2. $z' \in \mathbf{S}_1''(y)$.

In this case, the vertices z' and y'_{15} are to the left of x , y'_{15} is to the left of y , and z' is to the right of y . Let t be the vertex on the upper side of \mathbf{R} such that $t_a = z'_a$. Then, $t \in I(x, y_{15}) \cap I(x, z)$. From the choice of y_{15} and this inclusion, we deduce that $d(t, y_{15}) \leq d(t, z)$. Let s be the vertex on the upper side of \mathbf{R} such that $s_a = y_a$. Notice that $t \in I(s, z)$ and $s \in I(t, y_{15})$. Since $d(t, y_{15}) \leq d(t, z)$, from the previous inclusion, we obtain $d(s, y_{15}) \leq d(t, y_{15}) \leq d(t, z) \leq d(s, z)$. Since $s \in I(y, y_{15}) \cap I(y, z)$, we conclude that $d(y, y_{15}) \leq d(y, z)$. \square

Since $r_y'' \geq r_y'$ for any vertex $y \in \mathbf{R}_0 = \mathbf{R}' \cap \mathbf{R}''$, by Claims 21 and 22, we deduce that the ball $B_{r_y''}(y)$ realizes the sample X , concluding the proof. \square

7. INTERVAL GRAPHS

For any interval graph $G = (V, E)$, we construct proper LSCS of size 4 for $\mathcal{B}(G)$ and $\mathcal{B}_r(G)$. We consider a representation of G by a set of segments $J_v, v \in V$ of \mathbb{R} with pairwise distinct ends. For any $u \in V$, we denote by $J_u = [s_u, e_u]$ its segment, where s_u is the start of J_u , and e_u is the end of J_u , i.e., $s_u \leq e_u$. We use the following property of interval graphs:

Lemma 23. If $u, v \in B_r(x)$, $s_u, s_z < s_v$, and $e_u < e_v, e_z$, then $z \in B_r(x)$.

Proof. Since $s_z < s_v$ and $e_u < e_z$, if J_u and J_v intersect, then J_z covers the segment $[s_v, e_u]$, and otherwise, J_z intersects $[e_u, s_v]$. Let P be a path obtained from a shortest (x, u) -path of G by removing u , and Q be a path obtained from a shortest (x, v) -path by removing v . The union J_S of all segments of $S := P \cup \{x\} \cup Q$ intersects J_u and J_v . If J_u and J_v intersect, then J_z covers $[s_v, e_u]$, and thus, intersects J_S . Otherwise, J_S covers $[e_u, s_v]$, and J_z intersects $[e_u, s_v]$. In both cases, J_z and J_S intersect, whence a segment of S intersects J_z . Since all segments of S are at distance at most $r - 1$ from x , $z \in B_r(x)$. \square

Let X be a realizable sample for $\mathcal{B}(G)$. A *farthest pair* of X^+ is a pair $\{u^+, v^+\}$ such that u^+ is the vertex in X^+ whose interval J_{u^+} ends farthest to the left, and v^+ is the vertex in X^+ whose interval J_{v^+} begins farthest to the right, *i.e.*, for any $w \in X^+$, we have $e_{u^+} < e_w$ and $s_w < s_{v^+}$. If $u^+ \neq v^+$, then $[e_{u^+}, s_{v^+}] \cap J_w \neq \emptyset$ for any $w \in X^+$. If $u^+ = v^+$, then $J_{u^+} \subseteq J_w$ for any $w \in X^+$. A vertex p^- of X^- is a *left-bounder* if there is a ball $B_r(x)$ realizing X such that $e_{p^-} < s_x$ and, for all $p \in X^-$ with $e_p < s_x$, it holds that $e_p \leq e_{p^-}$. Analogously, a vertex q^- of X^- is a *right-bounder* if there is a ball $B_r(x)$ realizing X such that $e_x < s_{q^-}$ and, for all $q \in X^-$ with $e_x < s_q$, it holds that $s_{q^-} \leq s_q$. If p^- is a left-bounder and q^- is a right-bounder, then $\{p^-, q^-\}$ is a *bounding pair* of X^- . The farthest pair $\{u^+, v^+\}$ of X^+ and the bounding pair $\{p^-, q^-\}$ of X^- have the following properties:

Lemma 24. *If $u^+, v^+ \in B_r(x)$ and $r > 0$, then $X^+ \subseteq B_r(x)$.*

Proof. Pick any $w \in X^+ \setminus \{u^+, v^+\}$. By the definition of u^+ and v^+ , we have $s_w < s_{v^+}$ and $e_{u^+} < e_w$. If $u^+ \neq v^+$, then $s_{u^+} < s_{v^+}$ and $e_{u^+} < e_{v^+}$, and so, $s_{u^+}, s_w < s_{v^+}$ and $e_{u^+} < e_{v^+}, e_w$. By Lemma 23, $w \in B_r(x)$. Now, let $u^+ = v^+$. Then, $J_{u^+} \subset J_w$, and thus, any segment intersecting J_{u^+} also intersects J_w . Consequently, w is included in any ball of G of radius $r > 0$ containing u^+ , and, in particular, $w \in B_r(x)$. \square

Lemma 25. *If $e_{p^-} < s_x$ and $p^- \notin B_r(x)$, then, for all $z \in X^-$ with $e_z < e_{p^-}$, $z \notin B_r(x)$. Also, if $e_x < s_{q^-}$ and $q^- \notin B_r(x)$, then, for all $w \in X^-$ with $s_{q^-} < s_w$, $w \notin B_r(x)$.*

Proof. For the first statement, towards a contradiction, suppose that $e_{p^-} < s_x$ and $p^- \notin B_r(x)$, but there exists $z \in X^-$ such that $e_z < e_{p^-}$ and $z \in B_r(x)$. Then, $s_z, s_{p^-} < s_x$ since $s_z \leq e_z < e_{p^-} < s_x$, and $e_z < e_x, e_{p^-}$ as $e_z < e_{p^-} < s_x \leq e_x$. By Lemma 23, $p^- \in B_r(x)$, a contradiction. For the second statement, suppose by way of contradiction that $e_x < s_{q^-}$ and $q^- \notin B_r(x)$, but there exists $w \in X^-$ such that $s_{q^-} < s_w$ and $w \in B_r(x)$. Then, $s_x, s_{q^-} < s_w$ since $s_x \leq e_x < s_{q^-} < s_w$, and $e_x < e_w, e_{q^-}$ as $e_x < s_{q^-} < s_w \leq e_w$. By Lemma 23, $q^- \in B_r(x)$, a contradiction. \square

The compressor $\alpha(X)$. The compressor $\alpha(X)$ of X is a vector with four coordinates grouped into two pairs: $\alpha(X) := (\alpha_1(X), \alpha_2(X))$. The pair $\alpha_1(X)$ is a farthest pair $\{u^+, v^+\}$ of X^+ and the pair $\alpha_2(X)$ is a bounding pair $\{p^-, q^-\}$ of X^- . We use the symbol $*$ to indicate that the respective coordinate of $\alpha(X)$ is empty. We define $\alpha(X)$ as follows:

- (C1) if $X^+ = \emptyset$, then set $\alpha_1(X) = \alpha_2(X) := (*, *)$;
- (C2) if $X^+ = \{x\}$, then set $\alpha_1(X) := (x, *)$ and $\alpha_2(X) := (*, *)$;
- (C3) if $|X^+| \geq 2$, then set $\alpha_1(X) := (u^+, v^+)$ if $u^+ \neq v^+$ and $\alpha_1(X) := (*, v^+)$ if $u^+ = v^+$;
- (C3i) if $X^- = \emptyset$, then set $\alpha_2(X) := (*, *)$;
- (C3ii) if there exists a bounding pair of X^- , then set $\alpha_2(X) := (p^-, q^-)$;
- (C3iii) if there exists a left-bounder, but not a right-bounder of X^- , then set $\alpha_2(X) := (p^-, *)$;
- (C3iv) if exists a right-bounder, but not a left-bounder vertex of X^- , then set $\alpha_2(X) := (*, q^-)$.

The reconstructor $\beta(Y)$. The reconstructor β takes any signed vector Y on four coordinates grouped into two pairs Y_1 and Y_2 from $\text{Im}(\alpha)$, and returns a ball $\beta(Y)$ defined as follows:

- (R1) if $Y_1 = Y_2 = (*, *)$, then $\beta(Y)$ is the empty ball;
- (R2) if $Y_1 = (y_1, *)$ and $Y_2 = (*, *)$, then $\beta(Y)$ is the ball of radius 0 centered at y_1 ;
- (R3) if $Y_1 = (y_1, y_2)$ or $Y_1 = (*, y_2)$, then $\beta(Y)$ is any ball $B_r(x)$ of radius $r \geq 1$ containing Y_1 , not intersecting Y_2 , and such that:
 - (R3i) if $Y_2 = (*, *)$, then no condition;
 - (R3ii) if $Y_2 = (y_3, y_4)$, then $e_{y_3} < s_x$ and $e_x < s_{y_4}$;
 - (R3iii) if $Y_2 = (y_3, *)$, then $e_{y_3} < s_x$;
 - (R3iv) if $Y_2 = (*, y_4)$, then $e_x < s_{y_4}$.

Now, let X be a realizable sample for $\mathcal{B}_r(G)$. If $|X^+| \geq 2$ or $r \geq 1$, then we define α and β as above, since, in these cases, we do not specify the radius of the ball realizing X in α , nor the radius of the ball returned by β . So, we can exhibit a proper LSCS of size 4 for $\mathcal{B}_r(G)$ if we can deal with the case $|X^+| \leq 1$. The only difference is that if $|X^+| \leq 1$, then we set $\alpha_2(X)$ as in Case (C3), but we set $\alpha_1(X) := (*, *)$ when $X^+ = \emptyset$, and $\alpha_1(X) := (*, x)$ when $X^+ = \{x\}$. Now, let $r = 0$. If $|X^+| = 0$ and there is a ball $B_0(y)$ such that $y \notin X^-$ and $e_y < e_z$ for any $z \in V, z \neq y$, then $\alpha(X) := ((*, *), (*, *))$. Otherwise, if $|X^+| = 0$, there is a ball $B_0(y)$ such that $y \notin X^-$, $w' \in X^-$, $e_{w'} < e_y$, and, for all $w \in V$ with $e_w < e_y$, we have $e_w \leq e_{w'}$. In this case, $\alpha(X) := ((*, *), (w', *))$. If $X^+ = \{x\}$, set $\alpha(X) := ((x, *), (*, *))$. Given any signed vector Y on four coordinates, β returns a ball $\beta(Y)$ defined as follows. If $Y = ((*, *), (*, *))$, then $\beta(Y)$ is the ball $B_0(x)$ such that $e_x < e_z$ for any $z \in V \setminus \{x\}$. If $Y = ((*, *), (y_3, *))$, then $\beta(Y)$ is the ball $B_0(x)$ such that $e_{y_3} < e_x$, and, for all $w \in V$ with $e_w < e_x$, it holds that $e_w \leq e_{y_3}$. Lastly, if $Y = ((x, *), (*, *))$, then $\beta(Y)$ is the ball $B_0(x)$.

Proposition 26. *For any interval graph $G = (V, E)$, the pair (α, β) of vectors defines a proper labeled sample compression scheme of size 4 for $\mathcal{B}(G)$ and $\mathcal{B}_r(G)$.*

Proof. Let X be a realizable sample for $\mathcal{B}(G)$ (the case of $\mathcal{B}_r(G)$ is similar), $Y = \alpha(X)$, and $B_r(x) = \beta(Y)$. The cases (Rk) and their subcases in the definition of β correspond to the cases (Ck) and their subcases in the definition of α . The correctness is trivial if $k = 1, 2$. Now, let $k = 3$. Since Y_1 always contains a farthest pair of X^+ and the returned ball $B_r(x)$ contains Y_1 and $r \geq 1$, by Lemma 24, $X^+ \subseteq B_r(x)$. Furthermore, in Case (C3), any ball realizing X must have a radius $r \geq 1$ since $|X^+| \geq 2$. Now, we prove that $X^- \cap B_r(x) = \emptyset$. This is trivial in subcase (R3i) since $X^- = \emptyset$. In the remaining subcases of (R3), $X^- \cap B_r(x) = \emptyset$ follows from the definition of the corresponding subcase of case (C3) and Lemma 25. \square

8. SPLIT GRAPHS

In this section, we study proper labeled sample compression schemes for split graphs. First, recall that the sample compression conjecture for ball of radius 1 of split graphs is as hard as the general sample compression conjecture. Indeed, it is well known that, to any concept class \mathcal{C} of a set S , one can associate a split graph G having S as a clique, and a vertex v_C for each concept C of \mathcal{C} . Then, each ball $B_1(v_C)$ coincides with $C \cup \{v_C\}$. One can easily check that the VC-dimension of the family of balls $\{B_1(v_C) : C \in \mathcal{C}\}$ is the same as the VC-dimension of \mathcal{C} . A sample compression scheme for \mathcal{C} corresponds to a sample compression scheme for $\{B_1(v_C) : C \in \mathcal{C}\}$, where each sample X has its support in S .

Let $G = (V, E)$ be a split graph, and let $S = \{w_1, \dots, w_{\omega(G)}\}$ be the clique and I be the independent set in the partition of its vertices, where $\omega(G)$ is the clique number of G (the number of vertices in a maximum clique in G). For the family $\mathcal{B}(G)$, we construct proper labeled sample compression schemes of size $k = \max\{2, \omega(G)\}$ with the assumption that $\alpha(X)$ is a vector on k coordinates. Note that the family of balls of a split graph has VC-dimension at most k . We first define the compressor α .

The compressor $\alpha(X)$. We use $*$ to indicate that the respective coordinate of $\alpha(X)$ is empty. For any realizable sample X for $\mathcal{B}(G)$,

- (C1) if $X^+ = \emptyset$, then set $\alpha(X) := (*, \dots, *)$;
- (C2) otherwise, if $X^- = \emptyset$, then set $\alpha(X) := (u, *, \dots, *)$, where u is some vertex in X^+ ;
- (C3) otherwise, if $X^+ = \{u\}$, then set $\alpha(X) := (*, u, *, \dots, *)$;
- (C4) otherwise, if there exists $u \in X^+ \cap I$ such that $B_1(u) \cap X = X^+$, then set $\alpha(X) := (u, v, *, \dots, *)$, where $v \in X^+ \cap N(u)$;
- (C5) otherwise, $k = \omega(G)$, and, if there exists $u \in I \setminus X^+$ such that $B_1(u) \cap X = X^+$, then, for each $1 \leq i \leq k$, set the i^{th} coordinate of $\alpha(X)$ to w_i if $w_i \in X$, and $*$ otherwise;
- (C6) otherwise, if there exists $u \in S$ such that $B_1(u) \cap X = X^+$, then, for each $w_i \in S$,
 - (a) if there exists $y \in N(w_i) \cap X^-$ such that $y \notin \alpha(X)$, then set y as the i^{th} coordinate of $\alpha(X)$;

- (b) otherwise, if there exists $z \in X^+ \setminus N(w_i)$ and $z \notin \alpha(X)$, then set z as the i^{th} coordinate of $\alpha(X)$;
 - (c) otherwise, set $*$ as the i^{th} coordinate of $\alpha(X)$.
- (C7) otherwise, there exists $u \in I$ such that $B_2(u) \cap X = X^+$, and, for each $w_i \in S$,
- (a) if there exists $y \in N(w_i) \cap X^-$ such that $y \notin \alpha(X)$, then set y as the i^{th} coordinate of $\alpha(X)$;
 - (b) otherwise, if $w_i \in N(u)$ and there exists $z \in X^+ \cap N(w_i)$ such that $z \in I$ and $z \notin \alpha(X)$, then set z as the i^{th} coordinate of $\alpha(X)$;
 - (c) otherwise, set $*$ as the i^{th} coordinate of $\alpha(X)$.

The reconstructor $\beta(Y)$. The reconstructor β takes any sign vector Y on k coordinates and returns a ball defined in the following way:

- (R1) if $Y = (*, \dots, *)$, then $\beta(Y)$ is the empty ball;
- (R2) if $Y = (u, *, \dots, *)$ for some vertex $u \in Y^+$, then $\beta(Y)$ is any ball of G covering G ;
- (R3) if $Y = (*, u, *, \dots, *)$ for some vertex $u \in Y^+$, then $\beta(Y) = B_0(u)$;
- (R4) if $Y = (u, v, *, \dots, *)$ for two vertices $u, v \in Y^+$ such that $u \in I$ and $v \in S$, then $\beta(Y) = B_1(u)$;
- (R5) if there exist two vertices $x, y \in Y^+$ such that $x, y \in S$, then $\beta(Y)$ is any ball of radius 1 centered at a vertex $u \in I$ such that $B_1(u) \cap Y = Y^+$;
- (R6) if there exists $x \in Y^-$ and, for all $y_j \in Y^+$ such that y_j is the j^{th} coordinate of Y , $y_j \notin N(w_j)$, then $\beta(Y)$ is any ball of radius 1 centered at a vertex $w_t \in S$ such that the t^{th} coordinate of Y is $*$ and $B_1(w_t) \cap Y = Y^+$;
- (R7) otherwise, $\beta(Y)$ is any ball B of radius 2 centered at a vertex $u \in I$ such that $B \cap Y^- = \emptyset$, and, for all $y_t \in Y^+$ such that y_t is the t^{th} coordinate of Y , $w_t \in N(u)$.

Proposition 27. *For any split graph G , the pair (α, β) of vectors defines a proper labeled sample compression scheme of size $\omega(G)$ for $\mathcal{B}(G)$.*

Proof. Let X be a realizable sample for $\mathcal{B}(G)$. We will prove that the ball $\beta(Y)$ realizes the sample X , *i.e.*, that $\beta(Y) \cap X = X^+$. For each $1 \leq x \leq 7$, the Case (Rx) in the definition of β corresponds to the Case (Cx) in the definition of α , and the correctness follows as long as no two cases (Ci) and (Cj) produce the same type of vector. First, note that, in (C5), $\alpha(X)$ contains at least two vertices of X^+ that are in S since $|X^+| \geq 2$ and $B_1(u)$ realizes X for some vertex $u \in I$ such that $u \notin X^+$. Also, by definition, in (C6) and (C7), $\alpha(X)$ does not contain any vertex of X^+ that is in S , and contains at least one vertex of X^- that is in I since $|X^-| \geq 1$ and any ball realizing X contains all of S in both of these cases. Lastly, in (C7), since we are not in any of the other cases, there must be at least one vertex in $X^+ \cap I$, and thus, for some $1 \leq j \leq k$, the j^{th} coordinate of $\alpha(X)$ contains a vertex $y \in X^+ \cap N(w_j)$, while this is never the case in (C6). Hence, no two cases produce the same vector, and the correctness follows. \square

9. PLANAR GRAPHS

In this section, we study proper labeled sample compression schemes for planar graphs. There seems to be an inherent difficulty when dealing with planar graphs, and thus, we only give partial results in this section. Let G be a planar graph and fix a planar embedding of G . For the family $\mathcal{B}_1(G)$, we construct proper labeled sample compression schemes of size 4. Note that the family of balls of a planar graph has VC-dimension at most 4 [7, 9]. For a vertex x , we denote by $N(x)$ its open neighborhood, *i.e.*, $N(x) = B_1(x) \setminus \{x\}$. We begin with the next simple lemma, which is a direct consequence of the fact that a planar graph does not admit $K_{3,3}$ as a minor. Let X be a realizable sample. A *potential center* for X is a vertex v of G such that $X^+ \subseteq N(v)$ and $v \notin X^+$.

Lemma 28. *If $|X^+| \geq 3$ and, for any vertex $v \in X^+$, it holds that $N(v) \cap X^- \neq \emptyset$, then X either has one or two potential centers.*

Proof. Since X is realizable, there exists $u \in V$ such that $u \notin X^+$, $X^+ \subseteq N(u)$, and $X^- \cap N(u) = \emptyset$, and thus, X has at least one potential center (namely u). If X has three or more potential centers, then since $|X^+| \geq 3$, we obtain a forbidden $K_{3,3}$ (see Fig. 5 for an illustration). \square

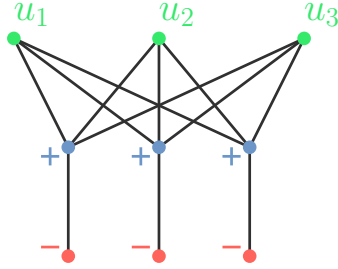


FIGURE 5. Illustration of the forbidden $K_{3,3}$ obtained in the proof of Lemma 28 when there are three potential centers u_1 , u_2 , and u_3 .

We can now describe our proper labeled sample compression scheme of size 4. The compressor $\alpha(X)$ is a vector on four coordinates. We use the symbol $*$ to indicate that the respective coordinate of $\alpha(X)$ is empty. If $|X^+| = 0$, then $\alpha(X) := (*, *, *, *)$, and $\beta((*, *, *, *))$ is the empty ball. If the center of a ball realizing X is contained in X^+ , say u , then $\alpha(X) := (u, *, *, *)$, and $\beta((u, *, *, *))$ is $B_1(u)$. Otherwise, if $|X^+| \geq 3$, then the compressor $\alpha(X)$ is defined as follows. First, let $\alpha(X)$ be equal to $(v_1, v_2, v_3, *)$, where v_1, v_2, v_3 are three arbitrary vertices of X^+ . If X^- contains a vertex z adjacent to a potential center of X , then change the last coordinate of $\alpha(X)$ to z , *i.e.*, set $\alpha(X) := (v_1, v_2, v_3, z)$. If Y is a signed vector with $|Y^+| = 3$, then the reconstructor $\beta(Y)$ is defined as follows. If $|Y^-| = 0$, then $\beta(Y)$ is $B_1(u)$ for any vertex u adjacent to the three vertices of Y^+ . If $|Y^-| = 1$, then $\beta(Y)$ is $B_1(u)$ for any vertex u adjacent to the three vertices of Y^+ , and not adjacent to the unique vertex of Y^- . The correctness immediately follows from Lemma 28.

Now, suppose that $|X^+| \leq 2$ and that the center of a ball realizing X is not contained in X^+ . In this case, the number of potential centers is potentially unbounded, and thus, we need a technique different from the one described above. We first deal with the case $|X^+| = 2$. For any two vertices $u, v \in V$, let $\ell_{u,v}$ be the labeling that assigns labels 0 through $k = |N(u) \cap N(v)| - 1$ to the vertices of $N(u) \cap N(v)$. In particular, let $\ell_{u,v}$ have the property that, for all $0 \leq i < j \leq k$, if $w_i, w_j \in N(u) \cap N(v)$ are such that $\ell_{u,v}(w_i) = i$ and $\ell_{u,v}(w_j) = j$, then, for all $i < y < j$, the vertex w_y with label y according to $\ell_{u,v}$ is contained in the interior region of the cycle u, w_i, v, w_j, u (see Fig. 6 for an illustration). Such a labeling always exists since the subgraph induced by the vertices u, v , and w_1, \dots, w_k is planar and contains $K_{2,k}$ as a subgraph.

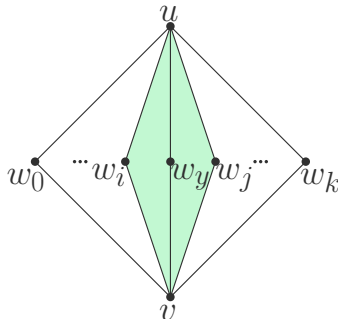


FIGURE 6. Illustration of the labeling $\ell_{u,v}$ described in the case of planar graphs when $X^+ = \{u, v\}$, $W = \{w_1, \dots, w_k\}$ is the set of the potential centers of X , where, for all $0 \leq i \leq k$, w_i has label i according to $\ell_{u,v}$, and $i < y < j$.

Now, let X be a realizable sample with $X^+ = \{u, v\}$, and let $W = \{w_0, \dots, w_k\}$ be the set of its potential centers, where, for all $0 \leq i \leq k$, w_i has label i according to $\ell_{u,v}$. Since X is

realizable, $k \geq 0$, and in what follows, all indices are considered modulo $k + 1$. Then, as a consequence of the property of $\ell_{u,v}$ described in the paragraph above, we get the following:

Lemma 29. *For every vertex $x \in X^-$, one of the following holds:*

- (1) $B_1(x) \cap W = \emptyset$;
- (2) $B_1(x) \cap W = \{w_j\}$, for some $0 \leq j \leq k$;
- (3) $B_1(x) \cap W = \{w_j, w_{j+1}\}$, for some $0 \leq j \leq k$;
- (4) $B_1(x) \cap W = \{w_j, w_{j+1}, w_{j+2}\}$, for some $0 \leq j \leq k$.

The compressor $\alpha(X)$ is defined as follows. First, let $\alpha(X) = (u, v, *, *)$ and note that $\alpha^+(X) := X^+$. If X^- does not contain a vertex that is a potential center nor a vertex adjacent to a potential center, then change nothing, *i.e.*, $\alpha^-(X) := \emptyset$. Otherwise, there exists some $0 \leq s \leq k$ such that $B_1(w_s) \cap X = X^+$, and some vertex $t \in X^-$ such that $w_{s-1} \in B_1(t)$, and so, change the third coordinate of $\alpha(X)$ to t , and note that $\alpha^-(X) = (t)$. If Y is a signed map with $Y^+ = (u, v)$, then the reconstructor $\beta(Y)$ is defined as follows. If $|Y^-| = 0$, then $\beta(Y) = N(w)$ for any vertex w adjacent to both the vertices of Y^+ . If $Y^- = (t)$, then $\beta(Y)$ is $N(w_s)$, where the vertices w_s and w_{s-1} are labeled s and $s - 1$, respectively, according to $\ell_{u,v}$, and $w_{s-1} \in B_1(t)$. The correctness follows from Lemma 29.

Now, let us finish with the case $|X^+| = 1$. For any vertex $u \in V$, let ℓ_u be the labeling that assigns labels 0 through $|N(u)| - 1$ to the vertices of $N(u)$ according to the circular order they appear in topologically around u in the clockwise direction, and starting at any vertex of $N(u)$. For a subset S of vertices of $N(u)$, the *predecessor* (*successor*, *resp.*) of a vertex $x \in S$, is the first vertex of S that appears in the circular order topologically around u in the counterclockwise (clockwise, *resp.*) direction. Now, let X be a realizable sample with $X^+ = \{u\}$. Since $B_1(u)$ does not realize X (which means that $N(u) \cap X^- \neq \emptyset$), any ball of radius 1 realizing X is centered at a neighbor of X . Consider the set $N(u) = \{w_0, \dots, w_k\}$ of the neighbors of u , where, for all $0 \leq i \leq k$, w_i has label i according to ℓ_u . Since u has a neighbor in X^- , there exists some $0 \leq s \leq k$ such that $B_1(w_s) \cap X = X^+$, and some vertex $t \in X^-$ such that $w_{s-1} \in B_1(t)$ (again, here and further, the indices are considered modulo $k + 1$). Let $W = \{w_s \in N(u) : w_{s-1} \in B_1(t) \text{ and } w_s \notin B_1(t)\}$. Our labeled sample compression scheme will return a ball of radius 1 centered at a vertex of W . Note that in the embedding of G , if u and t have at least two common neighbors, there exists $0 \leq p, q \leq k$ such that all vertices of $N(u) \cap B_1(t)$ distinct from w_p and w_q are in the interior region of the cycle u, w_p, t, w_q (where t is omitted if it coincides with w_p or w_q) and appear between p and q in the order ℓ_u . If $N(u) \cap B_1(t) = \{w_i\}$, we let $p = q = i$. Note that, for any $w_r, w_{r'} \in N(u) \cap B_1(t)$ such that w_r is the predecessor of $w_{r'}$ in $N(u) \cap B_1(t)$, there is at most one vertex w_s in W with $r < s < r'$. In particular, there is at most one vertex $w_s \in W$ that appears between w_q and w_p in the circular order ℓ_u .

To define $\alpha(X)$, we consider the first element w_s of W appearing after w_p (according to the order ℓ_u) such that $B_1(w_s) \cap X = X^+$. If $w_r \in B_1(t)$ for all $p \leq r \leq s - 1$ (*i.e.*, w_s is the first element of W appearing after w_p in $N(u)$), we let $\alpha^-(X) := (t)$ and $\alpha^+(X) := (u)$. Note that if u and t have a unique common neighbor, we are in this case. Otherwise, let $w_{s'}$ be the predecessor of w_s in W . Observe that $w_{s'}$ belongs to the interior region of the cycle $u, w_{s'-1}, t, w_{s-1}$ (where t is omitted if it coincides with $w_{s'-1}$ or w_{s-1}). Since $B_1(w_{s'})$ does not realize X , there exists $z \in B_1(w_{s'}) \cap X^-$. In this case, we let $\alpha^+(X) := (u)$ and $\alpha^-(X) := (t, z)$ (where we remember the order between t and z).

Claim 30. *Either $B_1(z) \cap W = \{w_{s'}\}$, or $B_1(z) \cap W = \{w_{s'}, w_{s''}\}$, where $w_{s''}$ is the predecessor of $w_{s'}$ in W .*

Proof. Since z is adjacent to $w_{s'}$, and $w_{s'}$ belongs to the interior region of the cycle $u, w_{s'-1}, t, w_{s-1}$, either z also belongs to the interior region of the cycle $u, w_{s'-1}, t, w_{s-1}$, or it coincides with one of the vertices $u, t, w_{s'}$ or w_{s-1} . Since $z \in X^-$ and $u \in X^+$, we have $z \neq u$. Since $w_{s'} \in B_1(z) \setminus B_1(t)$, we have $z \neq t$.

If z belongs to the interior region R of the cycle $u, w_{s'-1}, t, w_{s-1}$, then $w_{s'}$ is the only neighbor of z in W , since the only vertex of W in R is $w_{s'}$. If z coincides with $w_{s'-1}$ and $p < s' - 1 < s - 1$,

let $w_{r'}$ be the predecessor of $w_{s'-1}$ in $N(u) \cap B_1(t)$. Note that $p \leq r < s' - 1$. Then, z is in the interior of the region of the cycle $u, w_{r'}, t, w_{s-1}$. In this case, since the only vertices of W in this region are $w_{s'}$ and potentially $w_{s''}$ (if $r = s'' - 1$), we are done. If z coincides with $w_{s'-1} = w_p$, then z is not adjacent to any vertex lying in the interior region of u, w_{s-1}, t, w_q . The only vertices of W that are not lying in this region are $w_{s'}$ and w_{q+1} if it exists. Since w_{q+1} is the predecessor of $w_{s'}$ in this case, we are done. Suppose now that z coincides with w_{s-1} . If $s - 1 < q$, then z belongs to the interior cycle of $u, w_{s'-1}, t, w_r$, where w_r is the successor of w_{s-1} in $N(u) \cap B_1(t)$. Since the only vertices of W lying in this region are $w_{s'}$ and w_s , and since $z \notin B_1(w_s)$, $w_{s'}$ is the unique neighbor of z in W . Finally, suppose that $z = w_{s-1} = w_q$. In this case, z is not adjacent to any vertex lying in the interior region of the cycle $u, w_p, t, w_{s'-1}$. Since $w_{s'}$ and w_s are the only vertices of W that are not lying in this region, and since z is not adjacent to w_s , we are done. \square

We now define the reconstructor $\beta(Y)$ when $Y^+ = (u)$ and $|Y^-| \geq 1$. Consider the set $N(u) = \{w_0, \dots, w_k\}$ of the neighbors of u , where, for all $0 \leq i \leq k$, w_i has label i according to ℓ_u . If $|Y^-| = 1$, we let $Y^- = (t)$, and if $|Y^-| = 2$, we let $Y^- = (t, z)$. If $|N(u) \cap B_1(t)| \geq 2$, there exists $0 \leq p, q \leq k$ such that all the vertices of $N(u) \cap B_1(t)$ are in the interior region of the cycle u, w_p, t, w_q , and appear between p and q in the order ℓ_u . If $N(u) \cap B_1(t) = \{w_i\}$, we let $p = q = i$. Let $W = \{w_s \in N(u) : w_{s-1} \in B_1(t) \text{ and } w_s \notin B_1(t)\}$. If $Y^- = (t)$, let w_s be the first element of W appearing after w_p in the order ℓ_u , and let $\beta(Y) = B_1(w_s)$. Otherwise, $Y^- = (t, z)$, and let w' be the last vertex in W such that $w' \in B_1(z)$. In this case, $\beta(Y) = B_1(w)$, where w is the successor of w' in W .

Proposition 31. *For any planar graph G , the pair (α, β) of vectors defines a proper labeled sample compression scheme of size 4 for $\mathcal{B}_1(G)$.*

Proof. Let X be a realizable sample for $\mathcal{B}_1(G)$. We prove that the ball $\beta(Y)$ realizes the sample X , i.e., that $\beta(Y) \cap X = X^+$. When $Y = (*, *, *, *)$, then $|X^+| = 0$, and so, the empty set returned is compatible with X . When $|Y^+| = 1$ and $|Y^-| = 0$, this implies that the ball of radius 1 centered at the unique vertex of Y^+ is compatible with X , and so, $\beta(Y)$ is compatible with X in this case. When $|Y^+| = 3$, the correctness follows from Lemma 28, and when $|Y^+| = 2$, the correctness follows from Lemma 29. When $Y^+ = (u)$ and $Y^- = (t)$ (or $Y^- = (t, z)$), then we define W , w_p , and w_q as before. If $Y^- = (t)$, the ball of radius one centered at the first vertex in W appearing after w_p (in the order ℓ_u) realizes X , and it is the ball returned by $\beta(Y)$. If $Y^- = (t, z)$, then by the definition of α , there exists $w \in W$ such that $B_1(w) \cap X = X^+$ and z is adjacent to the predecessor of w in W . By Claim 30, w is necessarily the successor in W of the last vertex $w' \in W \cap B_1(z)$. Since $\beta(Y)$ returns $B_1(w)$ in this case, we are done. \square

10. HYPERBOLIC GRAPHS

A (ρ, μ) -approximate proper labeled sample compression scheme of size k for the family of balls $\mathcal{B}(G)$ of a graph G compresses any realizable sample X to a subsample $\alpha(X)$ of support of size k , such that $\beta(\alpha(X))$ is a ball $B_r(x)$ such that $X^+ \subseteq B_{r+\rho}(x)$ and $X^- \cap B_{r-\mu}(x) = \emptyset$. Let (V, d) be a metric space and $w \in V$. Let $\delta \geq 0$. A metric space (X, d) is δ -hyperbolic [18] if, for any four points u, v, x, y of X , the two larger of the sums $d(u, v) + d(x, y)$, $d(u, x) + d(v, y)$, and $d(u, y) + d(v, x)$, differ by at most $2\delta \geq 0$. Next, we show that δ -hyperbolic graphs admit a $(2\delta, 3\delta)$ -approximate labeled sample compression scheme of size 2.

An interval $I(u, v)$ of a graph is ν -thin if $d(x, y) \leq \nu$ for any two points $x, y \in I(u, v)$ with $d(u, x) = d(u, y)$ and $d(v, x) = d(v, y)$. Intervals of δ -hyperbolic graphs are 2δ -thin. A metric space (X, d) is *injective* if, whenever X is isometric to a subspace Z of a metric space (Y, d') , there is a map $f : Y \rightarrow Z$ such that $f(z) = z$ for any $z \in Z$ and $d'(f(x), f(y)) \leq d(x, y)$ for any $x, y \in Y$. By a construction of Isbell [16] (rediscovered by Dress [14]), any metric space (V, d) has an *injective hull* $E(V)$, i.e., the smallest injective metric space into which (V, d) isometrically embeds. Lang [21] proved that the injective hull of a δ -hyperbolic space is δ -hyperbolic. It was shown in [14] that the injective hull $T := T(u, v, y, w)$ of a metric space on 4 points u, v, y, w is a rectangle $R := R(u', v', y', w')$ with four attached tips uu', vv', yy', ww' (one or several tips

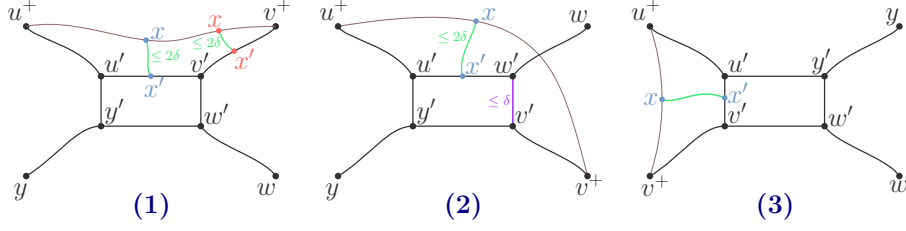


FIGURE 7. Cases 1-3 of Proposition 32.

may reduce to a single point or R may reduce to a segment or a single point). The smallest side of R is exactly the hyperbolicity of the quadruplet u, v, y, w .

Let X be a realizable sample of $\mathcal{B}(G)$ and $\{u^+, v^+\}$ be a diametral pair of X^+ . Let $B_{r^*}(y)$ be a ball of smallest radius such that $X^+ \subseteq B_{r^*}(y)$ and $X^- \cap B_{r^*}(y) = \emptyset$. We set $\alpha(X) := \emptyset$ if $X^+ = \emptyset$, $\alpha(X) := X^+$ if $|X^+| = 1$, and $\alpha(X) := \{u^+, v^+\}$ if $|X^+| \geq 2$. Given a subset Y of size at most 2, the reconstructor returns $\beta(Y) = \emptyset$ if $Y = \emptyset$, $\beta(Y) = B_0(y)$ if $Y = \{y\}$, and $\beta(Y) = B_{d(y_1, y_2)/2}(x)$ if $Y = \{y_1, y_2\}$, where x is the middle of a (y_1, y_2) -geodesic.

Proposition 32. *For any δ -hyperbolic graph $G = (V, E)$, the pair (α, β) defines a $(2\delta, 3\delta)$ -approximate proper labeled sample compression scheme of size 2 for $\mathcal{B}(G)$.*

Proof. We first show that $X^+ \subseteq B_{r+2\delta}(x)$, where $r = d(u^+, v^+)/2$ and x is a middle of a (u^+, v^+) -geodesic. Pick any $w \in X^+$. Since u^+, v^+ is a diametral pair of X^+ , $d(u^+, w) \leq 2r$ and $d(v^+, w) \leq 2r$. We also have $d(u^+, v^+) = 2r$ and $d(x, u^+) = d(x, v^+) = r$. Thus, the three distance sums have the form $d(u^+, w) + d(x, v^+) \leq 3r$, $d(v^+, w) + d(x, u^+) \leq 3r$, and $d(u^+, v^+) + d(x, w) = 2r + d(x, w)$. By the definition of δ -hyperbolicity, we conclude that either $d(x, w) \leq r$ (if $d(u^+, v^+) + d(x, w)$ is at most $3r$) or $d(x, w) \leq r + 2\delta$ (if $d(u^+, v^+) + d(x, w)$ is the largest sum). Hence, $w \in B_{r+2\delta}(x)$. We now show that $X^- \cap B_{r-3\delta}(x) = \emptyset$. Pick $w \in X^-$ and consider the injective hull T of the points $\{u^+, v^+, y, w\}$. T is a rectangle R with four tips (see Fig. 7) and is a subspace of the injective hull $E(V)$. Since $w \in X^-$, $w \notin B_{r^*}(y)$. Since $u^+, v^+ \in B_{r^*}(y)$, we deduce that $d(y, w) > d(y, u^+)$ and $d(y, w) > d(y, v^+)$. Let x' be a point of $I(u^+, v^+) \cap T$ such that $d(u^+, x') = d(u^+, x) = r$ and $d(v^+, x') = d(v^+, x) = r$. Since the injective hull T is δ -hyperbolic, its intervals are 2δ -thin, and thus, $d(x, x') \leq 2\delta$.

Case 1. u^+, v^+, y , and w are as in Fig. 7(1). First, suppose that x' belongs to the tip between u^+ and u' or to the segment between u' and v' . Since y' and w' belong to a common geodesic from y to w and from y to v^+ , and since $v^+ \in B_{r^*}(y)$ and $w \notin B_{r^*}(y)$, we deduce that $d(w, w') > d(w', v^+) \geq d(v', v^+)$. Consequently, $d(v', w) > d(v', v^+)$. If x' is located on the tip between u^+ and u' or on the segment between u' and v' , then, since $r = d(x', v^+) = d(x', v') + d(v', v^+)$ and $d(x', w) = d(x', v') + d(v', w)$, we obtain that $w \notin B_r(x')$. Since $d(x, x') \leq 2\delta$, $w \notin B_{r-2\delta}(x)$. If x' belongs to the tip between v' and v^+ , then $r = d(x', v^+) \leq d(v', v^+) \leq d(v', w)$, whence $w \notin B_r(x')$ and $w \notin B_{r-2\delta}(x)$.

Case 2. u^+ and v^+ , and y and w are opposite in T as in Fig. 7(2). Consider x' to be on the boundary of T containing the vertices u', w' , and v' . Since $v^+ \in B_{r^*}(y)$ and $w \notin B_{r^*}(y)$, then $d(v', w') + d(w', w) > d(v', v^+)$. Note also that $d(v', w') \leq \delta$, and thus, $d(w, v') > d(v', v^+) - \delta$. Independently of the location of x' on the boundary of T , $w \notin B_{r-\delta}(x')$. Thus, $w \notin B_{r-3\delta}(x)$.

Case 3. u^+, v^+, y , and w are as in Fig. 7(3). Since w' belongs to a geodesic between y and w and between y and v^+ , and $w \notin B_{r^*}(y)$, $v^+ \notin B_{r^*}(y)$, we deduce that $d(w', w) > d(w', v') + d(v', v^+) \geq d(v', v^+)$. Independently of the location of x' , we obtain that $w \notin B_{r-2\delta}(x)$. \square

11. PERSPECTIVES

A direction of interest is to design proper sample compression schemes for balls of radius r in trees of cycles or cube-free median graphs. Designing sample compression schemes of size $O(d)$ for balls in general median graphs G of dimension d is also open, as well as whether the VC-dimension of $\mathcal{B}(G)$ is $O(d)$ or not. For general median graphs, it no longer holds that the

interval between a diametral pair of X^+ contains a center of a ball realizing X . However, one can show that X^+ contains $2d$ vertices whose convex hull contains such a center. This convex hull can be d -dimensional and it is unclear how to encode the center in this region.

Other open questions are to design proper sample compression schemes of constant size for balls of planar graphs and of size $O(\omega(G))$ for balls of a chordal graph G . We showed that the former is possible for balls of radius 1 in Section 9, and that the latter is possible for split graphs in Section 8. Finding proper sample compression schemes of size $O(\omega(G))$ for $\mathcal{B}(G)$ is also interesting for other classes of graphs from metric graph theory: bridged graphs (generalizing chordal graphs) and Helly graphs; for their definitions and characterizations, see [1].

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AIX-MARSEILLE UNIVERSITÉ, UNIVERSITÉ DE TOULON, CNRS, LIS, MARSEILLE, FRANCE
Email address: `jeremie.chalopin@lis-lab.fr`

AIX-MARSEILLE UNIVERSITÉ, UNIVERSITÉ DE TOULON, CNRS, LIS, MARSEILLE, FRANCE
Email address: `victor.chepoi@lis-lab.fr`

CISPA HELMHOLTZ CENTER FOR INFORMATION SECURITY, SAARBRÜCKEN, GERMANY
Email address: `fmcinern@gmail.com`

AIX-MARSEILLE UNIVERSITÉ, UNIVERSITÉ DE TOULON, CNRS, LIS, MARSEILLE, FRANCE
Email address: `ratel.seb@gmail.com`

AIX-MARSEILLE UNIVERSITÉ, UNIVERSITÉ DE TOULON, CNRS, LIS, MARSEILLE, FRANCE
Email address: `yann.vaxes@lis-lab.fr`