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Mesh Density and Geodesic Tortuosity in Planar Triangular Tesselations Devoted to Fracture Mechanics

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In fracture mechanics, the mesh sensitivity is a key issue. It is particularly true concerning cohesive-volumetric finite element methods in which the crack path and the overall behavior are respectively influenced by the mesh topology and the mesh density. Poisson-Delaunay tessellations parameters, including the edge length distributions, were widely studied in the literature but very few works concern the mesh density and topology in Delaunay type meshes suitable for finite element simulations, which is of crucial interest for practical use.

Starting from previous results concerning Poisson-Delaunay tessellations and studying in details the Lloyd relaxation algorithm, we propose estimates for the probability density functions of the edge length and triangle top angles sets. These estimates depend both on the intensity of the underlying point process and on an efficiency index associated to the global quality of the mesh. The global and local accuracies of these estimates are checked for various standard mesh generators. Finally the mesh density and geodesic tortuosity are estimated for standard random or structured triangular meshes typically used in finite element simulations.

These results provide practical formulas to estimate bias introduced by the mesh density and topology on the results of cohesive-volumetric finite element simulations.

Keywords triangular tessellations, edge length distribution, angle distribution, mesh density, mesh tortuosity, Delaunay mesh

Short nomenclature

Nomenclature

τ denotes the intensity of a point process, defined in [section 2](#).

N_E denotes the number of edges of a mesh.

\mathcal{E} denotes the efficiency index of a mesh, introduced in [section 3.3](#).

$f_G(\tau, x)$ is the probability density function (pdf) of a quantity x for a graph/tessellation G with intensity τ ; this pdf has mean $\mu_G(\tau, x)$ and variance $\sigma_G^2(\tau, x)$.

Z_G is the mesh density of a graph/tessellation G defined as the total length of an edge set divided by the area A of the discretized planar domain.

$T_G(n)$ denotes the tortuosity of a particular continuous path between two nodes of a tessellation G composed by n connected edges.

PD and St stand for Poisson-Delaunay and Standard type meshes respectively.

1 Introduction

The mesh dependency is a crucial problem in fracture mechanics. In continuum damage mechanics as well as in surface fracture methods, softening behaviors lead to ill-posed problems and the field equations which describe the motion of the body lose hyperbolicity as soon as softening occurs ([Lasry et al. 1988](#); [G. Francfort et al. 2008](#)). Many regularization methods have been proposed,

most of them introducing a characteristic length scale (nonlocal formulations, higher gradient theories, extended or enriched finite element methods, etc.) or a characteristic time scale (volume or surface viscosity), see e.g. J.-L. Chaboche et al. (2001); X. Lu et al. (2009) for a review.

For applications we have in mind, we focus on surface damage approaches and in particular on cohesive-volumetric finite element method (F. Perales et al. 2008). It consists in inserting surface cohesive elements (a traction-separation law) between two adjacent bulk elements. These cohesive elements can be inserted on a predefined crack path or between each element of a given mesh. The later is often used when the crack path is not known a priori and has to emerge naturally from the loading path. Despite the success of these models, their use suffers from a twofold mesh dependency.

Firstly, since cracks are forced to propagate along element boundaries, a dependence on the mesh *topology* appears (Xu et al. 1994; M. Tijssens et al. 2000). This difficulty is related to the shortest path problem in tessellations. The local distributions of the edge length and triangle top angles sets have a strong influence on the shortest path topology and have to be carefully analyzed. The computed crack path length is always greater or equal to the theoretical one as a straight path between two nodes of a planar triangular mesh is not always accessible. As a consequence, the computed overall damage energy, proportional to the crack length (resp. surface) in 2D (resp. 3D) may be overestimated. The geodesic tortuosity is introduced in order to quantify this overestimation. It is defined as the ratio between the length of the shortest path between two nodes of the mesh composed of connected edges and the euclidian distance between these two nodes. It is a deterministic variable for structured meshes and a stochastic variable characterized by a mean and a variance for random meshes. Various solutions were proposed in the literature to avoid this dependency on the mesh topology. Among others, the extended finite element method (J. Dolbow et al. 2000), the pin-wheel based meshes (K. D. Papoulia et al. 2006) and K-means meshes (J. Rimoli et al. 2012) are very promising.

Secondly, the cohesive-volumetric finite element approach exhibits a *mesh size* dependency. The latter is related to the characteristic length scale introduced by cohesive zone models, since they are defined with the help of a surface energy and a peak stress. A rigorous micromechanical-based approach has been proposed by N. Blal et al. (2011); N. Blal et al. (2012a); N. Blal et al. (2012b) in order to limit this mesh size sensitivity. In this approach, the cohesive zone models are considered as inclusions distributed in a matrix (the set of the bulk elements). The spatial distribution of these inclusions strictly corresponds to the one of the edges (resp. surfaces) in a 2-D (resp. 3-D) mesh. For sake of simplicity, we focus our attention in the sequel to planar meshes. For any random planar meshes with an equiprobable distribution of the orientation of the edges, the spatial distribution of the inclusions is isotropic in space and in orientation. Since these inclusions have zero thickness, they are also characterized by a surface density, referred in this paper as *mesh density*. The mesh density is defined as the total length of the edge set divided by the area of the discretized domain and is directly related to the intensity of the underlying point process. As it is the case for the geodesic tortuosity, it is a deterministic variable for structured meshes and a stochastic variable characterized by a mean and a variance for random meshes. The main objective of this paper is to estimate this mesh density and the geodesic tortuosity for standard triangular tessellations suitable for finite element simulations.

Many works were devoted to distribution of the edge length and triangle top angles sets in Poisson-Delaunay tessellations since they are of practical interest in network/graph theories or in computational geometry. However the context of finite element simulations is less studied. The finite element tessellations are built to avoid badly shaped elements in order to preserve the quality of the finite element approximation, see e.g. Shewchuk (2002) for a review. Roughly speaking, the standard mesh generators thus lead to less “disordered” tessellations than the Poisson-Delaunay ones. To the best of our knowledge, there exists no estimate for the edge length and triangle top angles distributions neither for the mesh density and the mesh geodesic tortuosity of such a finite element tessellation.

Introducing some notations (section 2) and starting from standard results concerning Poisson-Delaunay tessellations, we show that the Lloyd relaxation algorithm leads to random triangular tessellations with edge length and triangle top angles distributions having the same means as Poisson-Delaunay tessellations but smaller variances (section 3). The “quality” of the resulting

meshes is also studied using a global efficiency index. Estimates are proposed for both edge length and triangle top angles distributions of any random triangular tessellation with arbitrary efficiency index and intensity of the underlying point process (section 4). The local robustness and the validity range of these estimates are analyzed in appendix A and appendix B. Moreover, we show that the edge length and the triangle top angles distributions of a Lloyd relaxed random triangular tessellation, are suitable for standard mesh generators. Finally, the mesh density and geodesic tortuosity are estimated for random tessellations (section 5). Additional results concerning the mesh density of usual structured tessellations are available in appendix C. Practical estimates are exhibited.

2 Notations

A Delaunay triangle is obtained by joining three points of a set of points P if and only if the resulting triangle does not contain any other points of P in the interior of its circumcircle. For a N points generating Poisson point process (intensity $\tau = N/A$, where A is the area covered by the set of points), table 1 holds. For a bounded Delaunay tessellation, the number of faces in a triangulation is not exactly twice the number of points and the edge set intensity and the face set intensity given in table 1 are only approximations in that case. We assume in the sequel that these approximations are accurate for high intensity point processes even for bounded tessellations.

intensities	point set	edge set	face set
Delaunay - G_D	τ	3τ	2τ

Table 1: Densities of the point set, of the edge set and of the face set of a Delaunay Tessellation based on a generating Poisson point process with intensity τ . According to Meijering (1953) and Miles (1970).

3 Poisson-Delaunay (PD) tessellation and Lloyd relaxed random triangular tessellation

3.1 Poisson-Delaunay tessellation: probability density functions of edge lengths and triangle top angles

If the set P corresponds to a Poisson point process with intensity τ , Collins (1968), Miles (1970) and Sibson (1980) show that the probability density functions of both the length l of the Delaunay edge and the top angle α of a Delaunay triangle are:

$$f_{PD}(\tau, l)(x) = \frac{1}{3} \pi x \tau \left(\operatorname{erfc} \left(\frac{1}{2} \sqrt{\pi} x \sqrt{\tau} \right) + x \sqrt{\tau} e^{-\frac{1}{4} \pi x^2 \tau} \right), \quad \text{for } x \geq 0, \quad (1)$$

and

$$f_{PD}(\alpha)(x) = 4 [(\pi - x) \cos(x) + \sin(x)] \frac{\sin(x)}{3\pi}, \quad \text{for } x \in [0, \pi], \quad (2)$$

where erfc is the complementary error function.

The corresponding means and variances read:

$$\mu_{PD}(\tau, l) = \frac{32}{9\pi\sqrt{\tau}} \quad (\text{length mean}), \quad \sigma_{PD}^2(\tau, l) = \frac{405\pi - 32^2}{81\pi^2\tau} \quad (\text{length variance}), \quad (3)$$

and

$$\mu_{PD}(\alpha) = \frac{\pi}{3} \quad (\text{top angle mean}), \quad \sigma_{PD}^2(\alpha) = \frac{\pi^2}{9} - \frac{5}{6} \quad (\text{top angle variance}). \quad (4)$$

These pdf are represented in figure 1 and compared to histograms of a particular PD tessellation.

We underline that both the mean and the variance of the edge lengths depend on τ . However, the triangle top angle distribution does not depend on the intensity of the point process. With the same number of points generated by the point process, the larger the domain, the greater the variance of the edge length distribution.

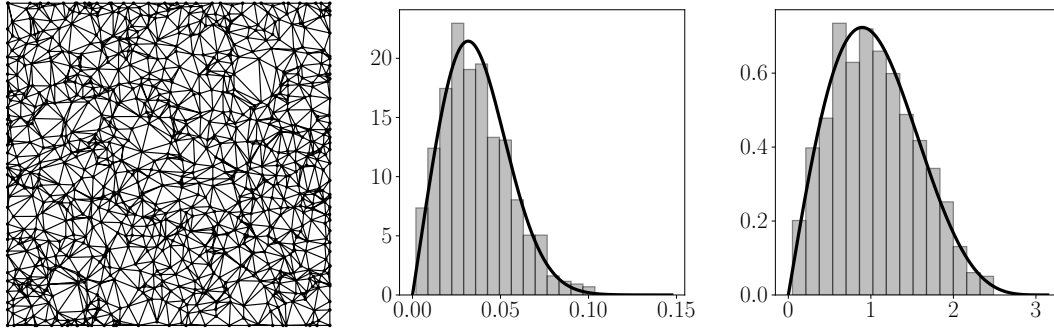


Figure 1: Left: Delaunay tessellation of a Poisson point process ($\tau = 1000$). Middle: Density function of the length l of a typical Delaunay edge. Right: Density function of a top angle α of a typical Delaunay triangle; black lines correspond to [equation \(1\)](#) and [equation \(2\)](#).

3.2 Lloyd relaxation

The Lloyd algorithm ([Lloyd 1982](#)) is an iterative scheme for computing Centroidal Voronoi tessellations. Centroidal Voronoi tessellations are Voronoi tessellations of a bounded geometric domain Ω such that the generating points of the tessellations are also the centroids (mass centers) of the corresponding Voronoi regions. The Lloyd algorithm consists in: i) selecting an initial set of points (seeds) according to a point process, ii) constructing the Voronoi tessellation associated to these seeds, iii) computing the mass centroids of the constructed Voronoi zones and iv) consider these centroids as the seeds of the new Voronoi tessellation. The process is iterated, see [figure 2](#). At convergence, the obtained Voronoi cells compose the so-called Centroidal Voronoi tessellation. The dual of this centroidal Voronoi tessellation is as close as possible to a tessellation of equilateral triangles.

Since Lloyd iterations lead to triangles as close as possible to equilateral ones and knowing that the mean of triangle top angles in a PD tessellation corresponds to the top angle of equilateral triangles, see [equation \(4\)](#), we consider that the triangle top angles mean remains unchanged through the Lloyd iterations. The same goes with the edge length mean following [Chang et al. \(1984\)](#): denoting by the subscript $L(i)$ parameters of the tessellation obtained by the Lloyd algorithm at iteration i :

$$\mu_{L(i)}(\tau, l) = \mu_{PD}(\tau, l) \quad \text{and} \quad \mu_{L(i)}(\alpha) = \mu_{PD}(\alpha). \quad (5)$$

Despite its wide application in many problems, the convergence of the Lloyd algorithm still remains an open issue. Numerous (local and global) convergence properties are presented in [Q. Du et al. \(2006\)](#) and in the references cited therein. The Lloyd algorithm may be viewed as a fixed point iteration of a mapping from a set of distinct seeds to the mass centers. The global convergence is only guaranteed in one-dimensional space for any positive and smooth density function. For multidimensional spaces, no rigorous convergence can be pronounced. Moreover, when the convergence holds, the speed of convergence is very low and can be significantly improved using dedicated algorithms ([Y. Liu et al. 2009](#)).

However, we show in the sequel that Lloyd relaxed tessellations lead to high quality triangles. For a sufficient number of iteration, this quality is similar to the quality of triangular mesh generators devoted to finite element computations. This quality is estimated with a global efficiency index. The associated variances of both edge length and triangle top angle sets are also estimated.

3.3 Efficiency index and variance of edge length and triangle top angle sets in PD tessellations improved by the Lloyd relaxation algorithm

The quality of a planar triangular mesh suitable for finite element calculations can be evaluated by two main ways. The first one concerns a topological measure of this quality: spatial homogeneity of the area of each triangle or of the edge length set, respect of minimal or maximal angles, etc. The second one refers directly to the quality of the finite element approximation. Various authors have thus proposed some interpolation quality measures, see [Shewchuk \(2002\)](#) for example. Since

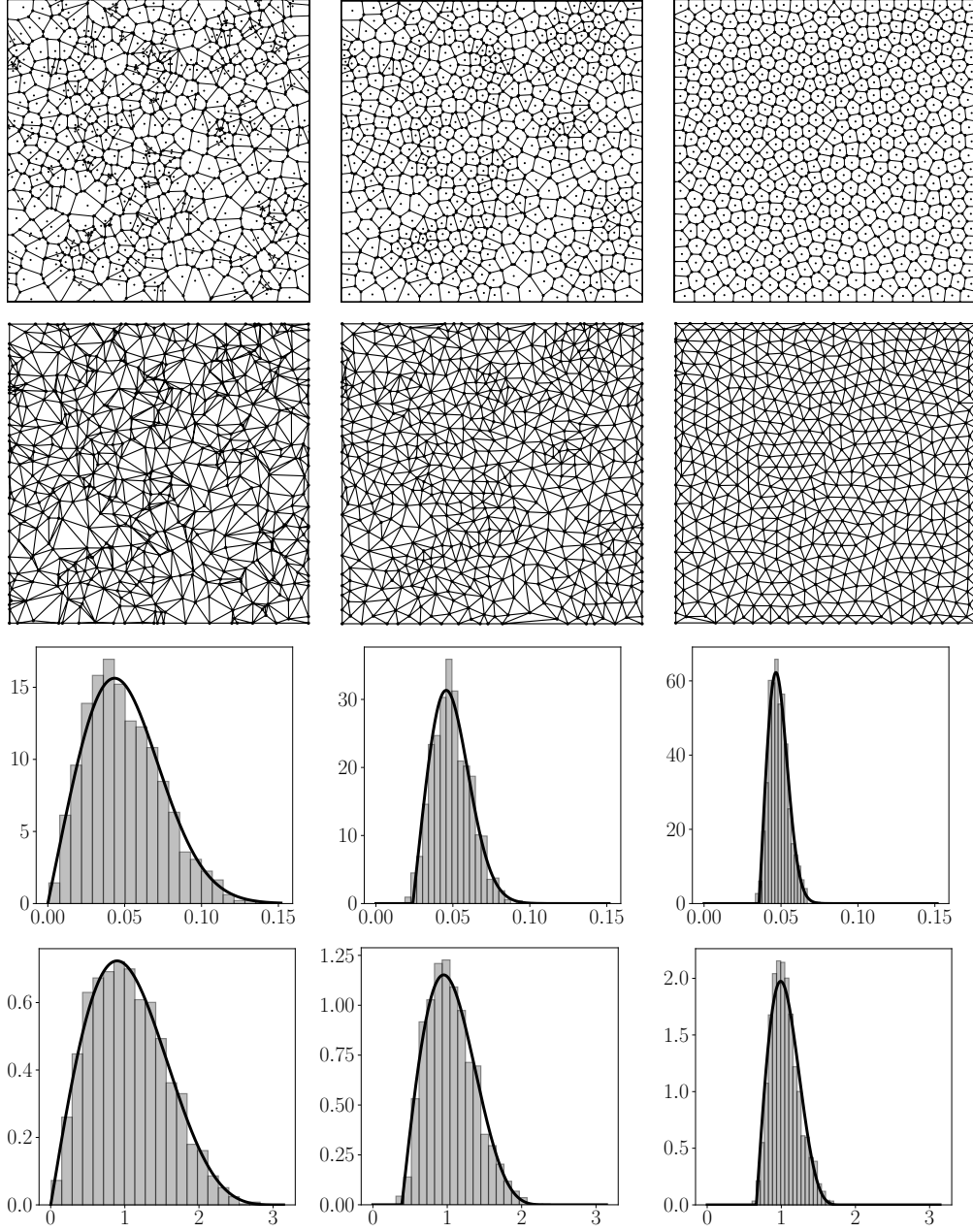


Figure 2: Lloyd relaxation algorithm. Top row: successive Lloyd relaxations of a Voronoi tessellation (intensity $\tau = 500$). Second row: dual graph. Third row: pdf of the edge length in the dual graph (bar charts). Bottom row: pdf of the triangle top angle in the dual graph (bar charts). Black lines over bar charts are estimates [equation \(14\)](#) and [equation \(15\)](#). From left to right various iterations of the Lloyd algorithm: initial tessellation ($\bar{\mathcal{E}} = 0.025$), 2d iteration ($\bar{\mathcal{E}} = 0.40$), 30th iteration ($\bar{\mathcal{E}} = 0.71$). Retained boundary elements may introduce a small bias. Thus edge lengths and summit angles of triangles containing a boundary node were not considered for computing efficiency indexes and establishing bar charts.

this later type of measure depends on the degree of interpolation in the finite element strategy, we focussed our attention on topological indexes.

Following [George et al. \(2000\)](#), the efficiency index of a triangular tessellation reads:

$$\mathcal{E} = e^{\langle \eta \rangle}, \quad \text{with} \quad \begin{cases} \eta = \bar{l} - 1 & \text{if } \bar{l} < 1 \\ \eta = 1/\bar{l} - 1 & \text{if } \bar{l} \geq 1 \end{cases}, \quad (6)$$

where \bar{l} is the adimensional length of an arbitrary edge e_i ($\rho(\mathbf{x})$ is a prescribed size field):

$$\bar{l} = \int_{e_i} \frac{1}{\rho(\mathbf{x})} dl, \quad (7)$$

and $\langle \gamma \rangle$ denotes the average of any quantity $\gamma = \{\gamma_1, \dots, \gamma_{N_E}\}$ defined on the set of edges (according to [table 1](#), $N_E = 3\tau A$ if τ is the intensity of the underlying point process over the area $A = |\Omega|$):

$$\langle \gamma \rangle = \frac{1}{N_E} \sum_{i=1}^{N_E} \gamma_i. \quad (8)$$

The efficiency index \mathcal{E} belongs to $[0, 1]$ and should be as close as possible to 1 in order to have a mesh suitable for finite element simulations.

In the sequel, the size field $\rho(\mathbf{x})$ is chosen uniform: $\rho(\mathbf{x}) = \mu_{PD,l}(\tau)$, $\forall \mathbf{x} \in \Omega$. Propagating the probability density function of the edge length [equation \(1\)](#) through the efficiency index [equation \(6\)](#), one obtains the PD tessellation approximated efficiency index ($\forall \tau$):

$$\mathcal{E}_{PD} \simeq 0.725. \quad (9)$$

In the sequel, efficiency indexes, edge length and triangle top angle variances are computed. In order to avoid bias induced by boundary elements, only triangles whose summit are inner tessellation nodes have been considered for computing these values.

The Lloyd algorithm leads to a triangular tessellation whose triangle elements are as close as possible to a tiling of equilateral triangles. As a consequence, both the edge length and triangle top angle variances diminish with Lloyd iterations (see [figure 3](#) on the left). Thus Lloyd relaxations improves the efficiency index, as the efficiency index increases when the edge length variance decreases (see [figure 3](#) on the right).

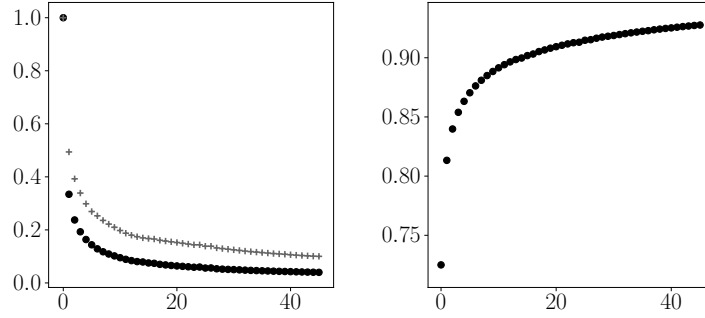


Figure 3: Left: edge length (black marks) and triangle top angle (gray marks) variances for the Lloyd algorithm (normalized by the corresponding variances of the Poisson-Delaunay tessellation respectively $\sigma_{PD}^2(\tau, l)$ and $\sigma_{PD}^2(\alpha)$) as a function of the successive iterations. Right: efficiency index \mathcal{E} of the Lloyd algorithm as a function of the successive iterations ($\tau = 300$).

A semi-empirical relationship can be exhibited between the variances and the efficiency for the Lloyd algorithm. Starting from a PD tessellation and remarking that for a structured tessellation based on equilateral triangles the efficiency index is optimal $\mathcal{E}^+ = 1$ and both the edge length and triangle top angle variances vanish $\sigma(\tau, l)^2 = 0$ and $\sigma(\alpha)^2 = 0$, we can derive semi-empirical relationships (see [figure 4](#)):

$$\frac{\sigma_{L(i)}(\tau, l)^2}{\sigma_{PD}(\tau, l)^2} \simeq \frac{(1 - \bar{\mathcal{E}}_{L(i)})^2}{1 + \bar{\mathcal{E}}_{L(i)}} \quad \text{and} \quad \frac{\sigma_{L(i)}(\alpha)^2}{\sigma_{PD}(\alpha)^2} \simeq \frac{(1 - \bar{\mathcal{E}}_{L(i)})^{3/2}}{(1 + \bar{\mathcal{E}}_{L(i)})^{1/2}}, \quad (10)$$

where $\bar{\mathcal{E}} \in [0, 1]$ denotes the reduced efficiency index of any index \mathcal{E} greater than or equal to \mathcal{E}_{PD} :

$$\bar{\mathcal{E}} = \frac{\mathcal{E} - \mathcal{E}_{PD}}{\mathcal{E}^+ - \mathcal{E}_{PD}} = \frac{\mathcal{E} - \mathcal{E}_{PD}}{1 - \mathcal{E}_{PD}}. \quad (11)$$

The corresponding means and variances of edge length and triangle top angle sets of a Lloyd tessellation read (combining [equation \(3\)](#), [equation \(4\)](#), [equation \(5\)](#) and [equation \(10\)](#)):

$$\mu_{L(i)}(\tau, l) = \frac{32}{9\pi\sqrt{\tau}} \quad \text{and} \quad \sigma_{L(i)}^2(\tau, l) \simeq \frac{(1 - \bar{\mathcal{E}}_{L(i)})^2}{1 + \bar{\mathcal{E}}_{L(i)}} \left(\frac{405\pi - 32^2}{81\pi^2\tau} \right), \quad (12)$$

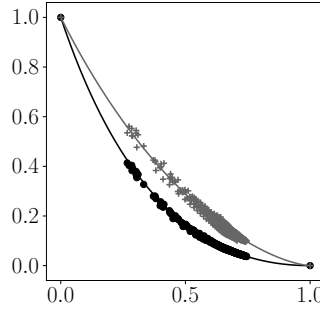


Figure 4: Edge length and triangle top angle variances for the Lloyd algorithm (normalized by the corresponding variances of the Poisson-Delaunay tessellation respectively $\sigma_{PD}^2(\tau, l)$ and $\sigma_{PD}^2(\alpha)$) as a function of the reduced efficiency index $\bar{\mathcal{E}}$. Points: numerical results starting from 10 different PD tessellations ($\tau = 300$). Lines: semi-empirical relationships [equation \(10\)](#).

and

$$\mu_{L(i)}(\alpha) = \frac{\pi}{3} \quad \text{and} \quad \sigma_{L(i)}^2(\alpha) \simeq \frac{(1 - \bar{\mathcal{E}}_{L(i)})^{3/2}}{(1 + \bar{\mathcal{E}}_{L(i)})^{1/2}} \left(\frac{\pi^2}{9} - \frac{5}{6} \right). \quad (13)$$

The next section proposes a probability density functions for both edge length and triangle top angle sets for standard triangular random meshes. These meshes try to avoid as much as possible distorted elements to ensure the accuracy of the finite element method.

4 Edge length and triangle top angle distributions in standard triangular random meshes

4.1 Standard meshing strategies

In the huge literature corresponding to mesh generators, three main approaches can be exhibited for unstructured (or random triangular) meshes (octree methods are not taking into account):

- *Delaunay tessellations.* Starting from a Poisson point process or using iterative George-Frey-like algorithms, this situation leads to $f_{PD}(\tau, l)$.
- *Frontal algorithms.* Using advancing-front techniques, meshes are generated starting from boundaries and adding progressively some “ideal” points as Steiner-like points ([Rebay 1993](#); [R. Lohner 1988](#)). The added points are chosen in order to have a new sequence of elements as close as possible as those of the previous sequence corresponding to the advancing-front (see e.g. [de Oliveira \(1988\)](#) for a review). Far away from the boundaries of the domain $\partial\Omega$, these frontal algorithms lead to semi-structured meshes with hexagonal shape.
- *Ad hoc optimized meshes.* Each mesh generator has its own strategy to construct elements having a *good* shape. Mathematical indicators of this good shape can be local (minimal angle, minimal aspect ratio, etc.) or global as [equation \(6\)](#). As for Delaunay tessellations, these mesh generators lead to random meshes.

Since, we have here interest for random meshes, we consider in the sequel the “ad hoc optimized meshes” and in particular the so-called “MeshAdapt” strategy proposed by Gmsh ([Geuzaine et al. 2009](#)) as a representative strategy of triangular random meshes generators. Starting from a PD tessellation, the “MeshAdapt” strategy consists in four local mesh modifications:

- *Edge Splitting:* long edges ($\bar{l}_{e_i} > 1.4$) are splitted if the two new obtained edges will have a sufficient adimensional size ($\min(\bar{l}_{e_j}, \bar{l}_{e_k}) > 0.7$).
- *Edge Collapsing:* short edges ($\bar{l}_{e_i} < 0.7$) are collapsed.
- *Edge Swapping:* an edge between two triangles is swapped if the efficiency index [equation \(6\)](#) will improved.
- *Triangle re-positioning:* each triangle (vertex) is moved inside the cavity made of all its surrounding triangles in order to maximize the worst element quality ([Freitag et al. 1997](#)).

This strategy leads to triangle-based random meshes with edge length and triangle top angle probability density functions respectively defined by [equation \(14\)](#) and [equation \(15\)](#), see [figure 5](#).

4.2 Edge length and triangle top angle distributions in random triangular tessellations with any efficiency index and intensity

The Lloyd relaxation algorithm depicted in [figure 2](#) and [figure 3](#) shows that when the efficiency index of a random triangular mesh increases, the probability density functions of the edge length and triangle top angle sets seem to be contracted and shifted while globally keeping their shapes. Moreover edge lengths and triangle top angles means are preserved. We assume that this observation is valid at any intensity τ and any efficiency index $\bar{\mathcal{E}}$ and we propose to define the pdf of any random triangular tessellation edge length and triangle top angle sets (referred with the subscript St for Standard mesh generators) as follows:

$$f_{\text{St}}(\bar{\mathcal{E}}, \tau, l)(x) = k(\bar{\mathcal{E}}, l) f_{\text{PD}}(\tau, l) (\tilde{x}(l, \bar{\mathcal{E}}, x)), \quad (14)$$

and

$$f_{\text{St}}(\bar{\mathcal{E}}, \alpha)(x) = k(\bar{\mathcal{E}}, \alpha) f_{\text{PD}}(\alpha) (\tilde{x}(\alpha, \bar{\mathcal{E}}, x)), \quad (15)$$

where $f_{\text{PD}}(\tau, l)$ and $f_{\text{PD}}(\alpha)$ are prolonged to be defined over the whole set of real numbers:

$$\forall x < 0, f_{\text{PD}}(\tau, l)(x) = 0 \quad \text{and} \quad \forall x \in]-\infty; 0[\cup]\pi; +\infty[, f_{\text{PD}}(\alpha)(x) = 0, \quad (16)$$

with

$$\begin{cases} \tilde{x}(l, \bar{\mathcal{E}}, x) = k(\bar{\mathcal{E}}, l)x + p(\bar{\mathcal{E}}, l) \\ \tilde{x}(\alpha, \bar{\mathcal{E}}, x) = k(\bar{\mathcal{E}}, \alpha)x + p(\bar{\mathcal{E}}, \alpha) \end{cases}, \quad (17)$$

the contraction parameters being defined as follows:

$$k(\bar{\mathcal{E}}, l) = \frac{\sigma_{\text{PD}}(\tau, l)}{\sigma_{\text{St}}(\bar{\mathcal{E}}, \tau, l)} = \frac{(1 + \bar{\mathcal{E}})^{1/2}}{1 - \bar{\mathcal{E}}} \quad (18)$$

and

$$k(\bar{\mathcal{E}}, \alpha) = \frac{\sigma_{\text{PD}}(\alpha)}{\sigma_{\text{St}}(\bar{\mathcal{E}}, \alpha)} = \frac{(1 + \bar{\mathcal{E}})^{1/4}}{(1 - \bar{\mathcal{E}})^{3/4}}, \quad (19)$$

and the shift parameters by:

$$p(\bar{\mathcal{E}}, \tau, l) = (1 - k(\bar{\mathcal{E}}, l)) \mu_{\text{PD}}(\tau, l) \quad (20)$$

and

$$p(\bar{\mathcal{E}}, \alpha) = (1 - k(\bar{\mathcal{E}}, \alpha)) \mu_{\text{PD}}(\alpha) = (1 - k(\bar{\mathcal{E}}, \alpha)) \frac{\pi}{3}, \quad (21)$$

where $\bar{\mathcal{E}}$ states again for the reduced efficiency index given in [equation \(11\)](#).

Accuracy of estimates [equation \(14\)](#) and [equation \(15\)](#) is shown in [figure 2](#), [figure 5](#), [figure 6](#) and [figure 7](#).

Since we have interest for standard mesh generators, four new cases are considered in [figure 6](#) and [figure 7](#) covering various standard efficiency indexes and intensities:

1/ a mesh generated with the *frontal* algorithm of Gmsh with low intensity but high efficiency index, 2/ a mesh generated with [Abaqus \(2008\)](#) with very high intensity and high efficiency index, 3/ a mesh obtained with the pre-processor of the [Cast3m \(2011\)](#) finite element code with high intensity and mid efficiency index, 4/ a mesh drawn with the mesh generator Triangle ([Shewchuk 2005](#)) with high intensity and low efficiency index.

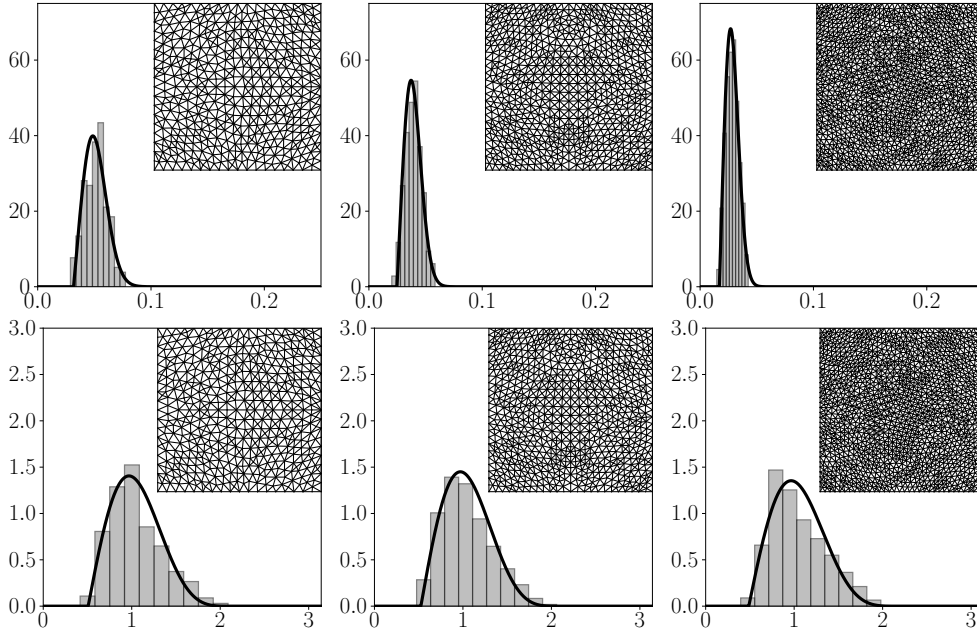


Figure 5: Histograms of the edge length and triangle top angle sets obtained with the Gmsh mesh generator (Geuzaine et al. 2009), with “MeshAdapt” strategy. First line: edge lengths distributions. Second line: triangle top angle distributions. From left to right the meshes with increasing intensity τ , (coarser mesh $\tau = 525$, intermediate mesh $\tau = 831$, finest mesh $\tau = 1611$) black lines refer to the standard mesh estimators for the density functions of both the edge length [equation \(14\)](#) and the triangle top angle [equation \(15\)](#) sets. The efficiency index of the finest mesh is about $\mathcal{E} \approx 0.86$.

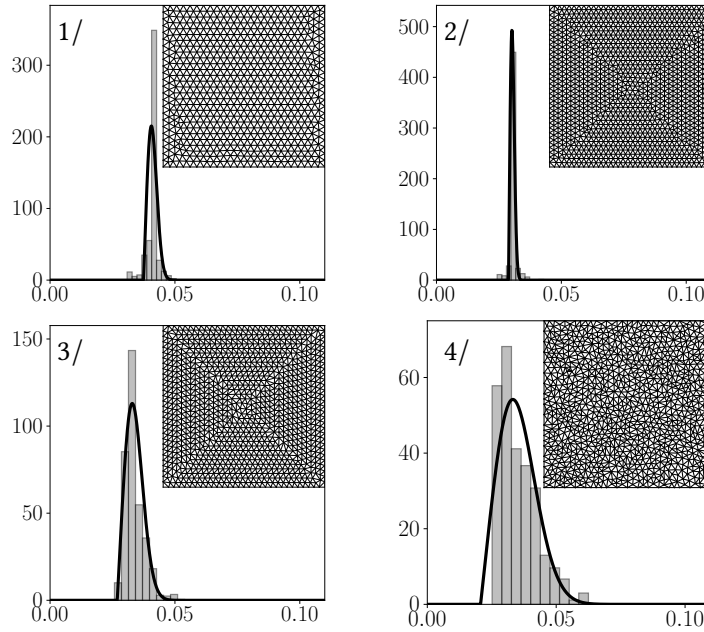


Figure 6: Edge length probability density: accuracy of estimate [equation \(14\)](#) (black line). Meshes obtained with: 1/ the “frontal” algorithm of Gmsh ($\tau = 750$, $\mathcal{E} = 0.976$), 2/ the Abaqus software ($\tau = 1386$, $\mathcal{E} = 0.981$), 3/ the Cast3m software ($\tau = 1102$, $\mathcal{E} = 0.917$), 4/ the Triangle mesh generator ($\tau = 1008$, $\mathcal{E} = 0.863$).

Figure 5, figure 6 and figure 7 show that estimates [equation \(14\)](#) and [equation \(15\)](#) lead to accurate estimations of the edge length and triangle top angle means. However variances estimations are accurate only for tessellations whose efficiency index is not too high (typically equal or at least less than 0.92, see [figure 6](#) and [figure 7](#)). As a conclusion, estimates [equation \(14\)](#) and [equation \(15\)](#) are accurate for standard random meshes with efficiency index between 0.725 and 0.92 (i.e. with reduced efficiency index at least less than 0.71).

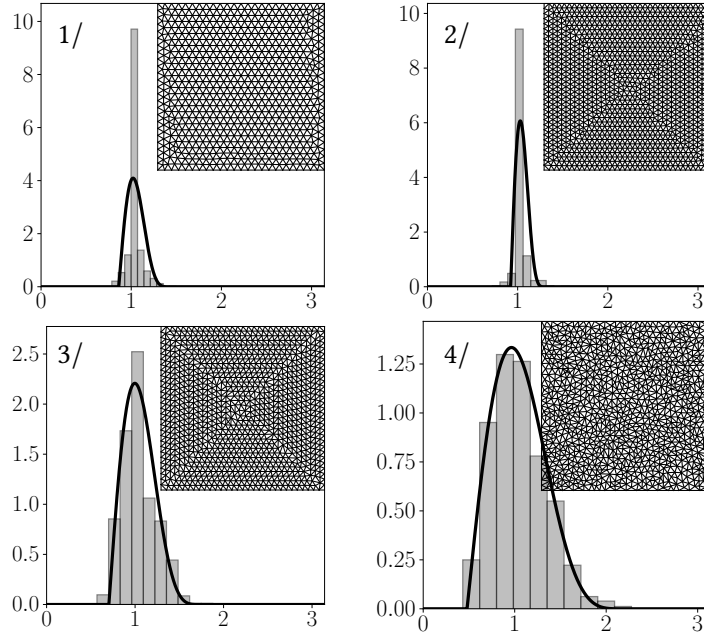


Figure 7: Top angle probability density: accuracy of estimate [equation \(15\)](#) (black line). Meshes obtained with: 1/ the “frontal” algorithm of Gmsh ($\tau = 750$, $\mathcal{E} = 0.976$), 2/ the Abaqus software ($\tau = 1386$, $\mathcal{E} = 0.981$), 3/ the Cast3m software ($\tau = 1102$, $\mathcal{E} = 0.917$), 4/ the Triangle mesh generator ($\tau = 1008$, $\mathcal{E} = 0.863$).

5 Density of random triangular meshes

The mesh density Z is defined as the sum of the edge lengths $\sum_{i=1}^{N_E} l_i$ divided by the area A of the meshes domain Ω :

$$Z = \frac{1}{A} \sum_{i=1}^{N_E} l_i$$

where N_E is the number of edges of the mesh.

For practical purpose, we showed that the pdf of the length l of the typical edge is given by [equation \(14\)](#). Due to the central limit theorem, the sum $\sum_{i=1}^{N_E} l_i$ thus follows a normal distribution with mean $N_E \times \mu_{St}(\tau, l)$ and variance $N_E \times \sigma_{St}^2(\bar{\mathcal{E}}, \tau, l)$ when the number of edges N_E is large enough ([Billingsley 1995](#)). Then the mesh density Z follows a normal distribution with mean $(N_E/A) \times \mu_{St}(\tau, l)$ and variance $(N_E/A^2) \times \sigma_{St}^2(\bar{\mathcal{E}}, \tau, l)$.

Recalling that $N_E = 3\tau A$, one obtains the probability density function $f_{St}(\bar{\mathcal{E}}, \tau, Z)$ of the mesh density as a normal distribution with mean $\mu_{St}(\tau, Z) = 3\tau \mu_{St}(\tau, l)$ and variance $\sigma_{St}^2(\bar{\mathcal{E}}, \tau, Z) = (3\tau/A) \sigma_{St}^2(\bar{\mathcal{E}}, \tau, l)$:

$$f_{St}(\bar{\mathcal{E}}, \tau, Z) = \mathcal{PDF} \left[\mathcal{N} \left(\mu_{St}(\tau, Z), \sigma_{St}^2(\bar{\mathcal{E}}, \tau, Z) \right) \right] \quad (22)$$

with

$$\mu_{St}(\tau, Z) = \frac{32\sqrt{\tau}}{3\pi} \text{ (mean)} \quad \text{and} \quad \sigma_{St}^2(\bar{\mathcal{E}}, \tau, Z) \simeq \frac{(1 - \bar{\mathcal{E}})^2}{1 + \bar{\mathcal{E}}} \left(\frac{405\pi - 32^2}{27\pi^2 A} \right) \text{ (variance)}. \quad (23)$$

It is worth noting that the variance of this mesh density no more depends on the intensity τ , but now depends on the area A of the meshed domain. In particular, σ_Z^2 depends on the unit length used to define the boundary of the domain $\partial\Omega$.

The ratio in percentage between the standard deviation and the mean of a standard mesh density is given by :

$$\frac{\sigma_{St}(\bar{\mathcal{E}}, \tau, Z)}{\mu_{St}(\tau, Z)} \times 100 \simeq \left(0.49 \frac{1 - \bar{\mathcal{E}}}{\sqrt{1 + \bar{\mathcal{E}}}} \right) \frac{1}{\sqrt{N_E}} \times 100. \quad (24)$$

This relation leads to the following result :

$$\forall \mathcal{E} \geq 0.86, \frac{\sigma_{St}(\bar{\mathcal{E}}, \tau, Z)}{\mu_{St}(\tau, Z)} \times 100 \leq \frac{20}{\sqrt{N_E}}. \quad (25)$$

In fracture mechanics, N_E is sufficiently large to consider that the variance of Z is negligible. We then consider for practical purpose that Z can be understood as a scalar value depending exclusively on the intensity of the underlying point-process or equivalently on the edge length mean :

$$Z_{St}(\tau) = \frac{32\sqrt{\tau}}{3\pi} = \frac{1024}{27\pi^2\mu_{St}(\tau, l)}. \quad (26)$$

The accuracy of estimate [equation \(26\)](#) is illustrated in [figure 8](#).

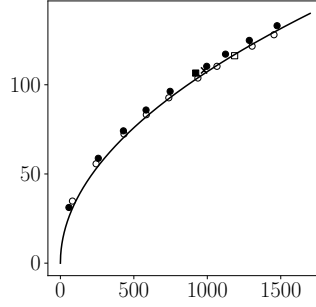


Figure 8: The mesh density Z as a function of the intensity of the point-process τ (line: $\mu_{St}(\tau, Z)$ in [equation \(22\)](#), unitary area). Gmsh results with “MeshAdapt” strategy (points) and with “Frontal” strategy (open circles); Cast3m result (cross); Triangle result (filled square); Abaqus result (open square). Corresponding meshes are given in [figure 6](#) and [figure 7](#).

6 Geodesic tortuosity of random Delaunay meshes

6.1 Shortest path approximation

Let A and B be two nodes of a planar triangular mesh. Among all the paths which are composed of a collection of adjacent edges and whose extremities are nodes A and B , there exists at least one with a minimum length. Let an edge inclination be the angle between this edge and the direction AB , whose value has to be taken between $-\frac{\pi}{2}$ and $\frac{\pi}{2}$.

The orientation with respect to a given direction AB of the edges of a PD mesh is uniformly distributed over $[-\frac{\pi}{2}; \frac{\pi}{2}]$. It is reasonably supposed that this observation remain valid for meshes obtained from a PD tessellation after few Lloyd relaxations (as long as the mesh remain sufficiently unstructured). In the sequel, such meshes will be referred as “unstructured Delaunay-like meshes”. As a consequence, only the distance $d(A, B)$ has to be considered for estimating the shortest path length between A and B in such meshes. Considering a fixed value for the distance $d(A, B)$ is somehow the same as fixing the number of edges n composing the shortest path. For sake of simplicity, the choice is therefore made to caracterise a shortest path only by the number n of edges it contains.

At that point, the edge length and triangle top angle distributions are the only geometrical data on which we can rely on for estimating the shortest path length between two nodes of a given planar triangular mesh. Although, to the best of our knowledge, there is no explicit mathematical expression relying these quantities. The difficulty lies in the fact that a shortest path is a “global” quantity in the way that there is no explicit rule followed individually by each edge. Thus all the edges of the shortest path have to be found at the same time and cannot be chosen independently from each other.

Facing this difficult problem, it has been considered to focus on n -edges paths verifying the following property, hereafter referred as “inclination rule”:

$$\text{Each edge inclination of the path is as close as possible to zero.} \quad (27)$$

The idea is to consider paths as close as possible to the shortest one and whose description is made easier by the inclination rule they fulfill. Thanks to this rule the probability density

function of the tortuosity of such a path - i.e. the ratio between the path length and the length of its projection along the direction AB - can be estimated (see [section 6.4](#)).

[Figure 9](#) represents on its first column examples of a shortest path between two nodes A and B of a PD tessellation and a Lloyd relaxed PD tessellation. A distinction is made between edges following the inclination rule (blue ones) and those who are not (red ones). On the same figure, on the second column, corresponding histograms of the percentage of the path edges which are not following the inclination rule are displayed.

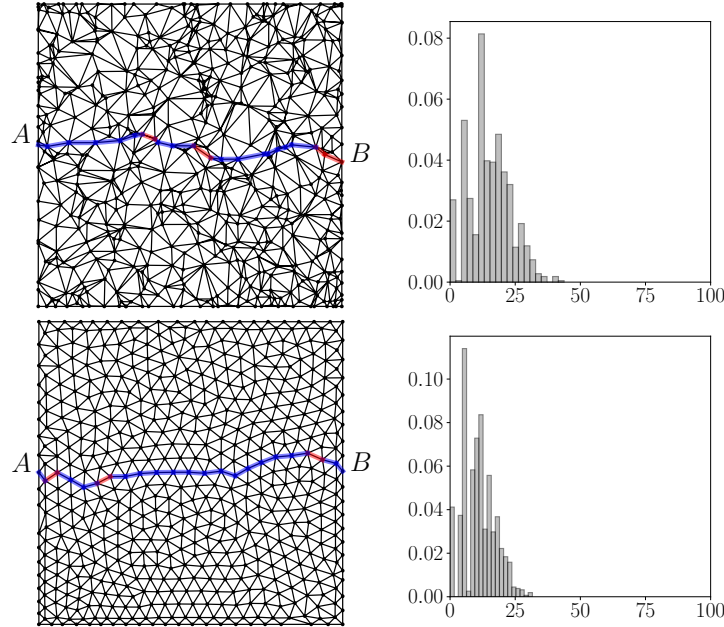


Figure 9: First column: Representation of the shortest path between two nodes A and B in a PD tessellation (top) and a Lloyd relaxed tessellation (bottom) of a Poisson point process ($\tau = 300$). Blue edges follow the inclination rule, red ones do not. On the second column: over 1000 PD tessellations (top) and over 1000 Lloyd relaxed tessellations (bottom) of a Poisson point process ($\tau = 300$), histograms of the percentage of edges of the shortest path which are not following the inclination rule.

Although all the edges are not following the inclination rule, a major part of them does. To consider that the shortest path is following the rule [equation \(27\)](#) is then not optimal but seems acceptable for estimating the shortest path tortuosity.

6.2 Definition of the tortuosity of a path following the inclination rule

Among all edges connected a given node, at least one is associated to a minimum inclination (along direction AB) absolute value. The inclination θ of this particular edge is referred in the sequel as “best inclination” (see [figure 10](#)).

Edges best inclinations of a path following the inclination rule [equation \(27\)](#) are particular events of the random variable θ . For sake of simplicity, we consider in the sequel that these events are statistically independent. The tortuosity of a n -edges path of a given tessellation G following the inclination rule is given by:

$$T_G(n) = \frac{\sum_{i=1}^n l_i}{\sum_{i=1}^n l_i \cos(\theta_i)}, \quad (28)$$

where l_i are the edge lengths of the path and θ_i its edge inclinations.

In the particular case where the path is a straight line, $\forall i \in \llbracket 1; n \rrbracket$, $\theta_i = 0$, the tortuosity is equal to 1. For other case the tortuosity is greater than 1.

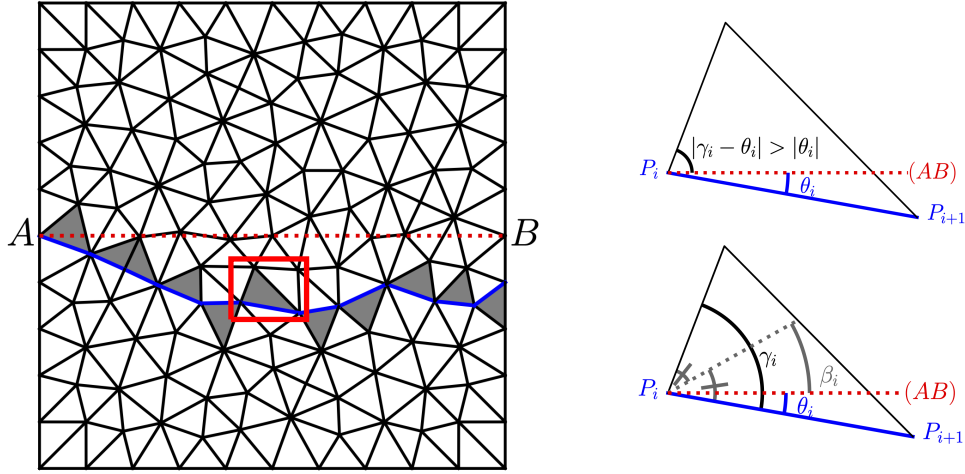


Figure 10: Schematic definition of the incident angle γ , bisector orientation β and best inclination θ .

6.3 Probability density function of the best inclination in Delaunay-type meshes

Except for trivial cases, the half straight line from a node P with direction AB crosses exactly one of the triangle elements surrounding P . We respectively denote as β and γ the bisector orientation with respect to the direction AB and the angle of this triangle at top P , hereafter called the incident angle (see figure 10).

Node P is surrounded by triangles having their own angle at summit P . The orientation with respect to a given direction AB of the edges of an unstructured Delaunay-like mesh may reasonably supposed uniformly distributed over $[-\frac{\pi}{2}; \frac{\pi}{2}]$. As a consequence, the probability that a summit angle corresponds to a particular value of γ is proportional to the value of this angle. In other words, it is more likely that among the surrounding triangles at node P , the triangle crossed by the half straight line from node P with direction AB be the one having the highest angle at summit P . Considering then that the angles at summit P are independent events of the random variable α (as defined in section 3), the following expression for the pdf of γ is obtained:

$$f_{\text{St}}(\bar{\mathcal{E}}, \gamma)(x) = \frac{k(\bar{\mathcal{E}}, \alpha)}{\frac{\pi}{3} - p(\bar{\mathcal{E}}, \alpha)} \times x f_{\text{St}}(\bar{\mathcal{E}}, \alpha)(x), \quad (29)$$

where the prefactor ensures that $\int_0^\pi f_{\text{St}}(\bar{\mathcal{E}}, \gamma)(x) dx = 1$.

By construction, the bisector orientation β verifies:

$$\beta \in [-\frac{\gamma}{2}; \frac{\gamma}{2}]. \quad (30)$$

Since the edges inclination follows a uniform distribution, the conditional density of β knowing $\gamma = \gamma_0$ is derived from equation (30):

$$\begin{aligned} g_{\text{St}}(\bar{\mathcal{E}}, \beta | \gamma = \gamma_0)(x) &= \begin{cases} \frac{1}{\gamma_0} & \text{if } x \in [-\frac{\gamma_0}{2}, \frac{\gamma_0}{2}] \\ 0 & \text{if not} \end{cases} \\ &= \begin{cases} \frac{1}{\gamma_0} & \text{if } |x| \leq \frac{\pi}{2} \text{ and } \gamma_0 \in [2|x|, \pi] \\ 0 & \text{if not} \end{cases}, \end{aligned} \quad (31)$$

Using then the law of total probability, the pdf of β reads:

$$f_{\text{St}}(\bar{\mathcal{E}}, \beta)(x) = \int_0^\pi g_{\text{St}}(\bar{\mathcal{E}}, \beta | \gamma = z)(x) f_{\text{St}}(\bar{\mathcal{E}}, \gamma)(z) dz, \quad (32)$$

which, according to [equation \(29\)](#) and [equation \(31\)](#), could be written as follow:

$$\begin{aligned} f_{\text{St}}(\bar{\mathcal{E}}, \beta)(x) &= \int_{2|x|}^{\pi} \frac{f_{\text{St}}(\bar{\mathcal{E}}, \gamma)(z)}{z} dz = \left(\frac{k(\bar{\mathcal{E}}, \alpha)}{\frac{\pi}{3} - p(\bar{\mathcal{E}}, \alpha)} \right) \int_{2|x|}^{\pi} f_{\text{St}}(\bar{\mathcal{E}}, \alpha)(z) dz \\ &= \left(\frac{k(\bar{\mathcal{E}}, \alpha)}{\frac{\pi}{3} - p(\bar{\mathcal{E}}, \alpha)} \right) (1 - F_{\text{St}}(\bar{\mathcal{E}}, \alpha)(2|x|)), \end{aligned} \quad (33)$$

where $F_{\text{St}}(\bar{\mathcal{E}}, \alpha)$ is the cumulative density function of α given by:

$$\begin{aligned} F_{\text{St}}(\bar{\mathcal{E}}, \alpha)(x) &= \frac{1}{3\pi} \left[\pi + 2\tilde{x}(\alpha, \bar{\mathcal{E}}, x) + (\tilde{x}(\alpha, \bar{\mathcal{E}}, x) - \pi) \cos(2\tilde{x}(\alpha, \bar{\mathcal{E}}, x)) \right. \\ &\quad \left. - \frac{3}{2} \sin(2\tilde{x}(\alpha, \bar{\mathcal{E}}, x)) \right], \quad \text{if } \tilde{x}(\alpha, \bar{\mathcal{E}}, x) \in [0, \pi], \end{aligned} \quad (34)$$

if $\tilde{x}(\alpha, \bar{\mathcal{E}}, x) \in [0, \pi]$, and by

$$F_{\text{St}}(\bar{\mathcal{E}}, \alpha)(x) = 0, \quad \text{if } \tilde{x}(\alpha, \bar{\mathcal{E}}, x) \leq 0 \quad (35)$$

and

$$F_{\text{St}}(\bar{\mathcal{E}}, \alpha)(x) = 1, \quad \text{if } \tilde{x}(\alpha, \bar{\mathcal{E}}, x) \geq \pi, \quad (36)$$

where $\tilde{x}(\alpha, \bar{\mathcal{E}}, x)$ is defined in [equation \(17\)](#).

Regarding [figure 10](#) and denoting by i^{th} edge of a path following the inclination rule [equation \(27\)](#), a relation can be exhibited between θ_i , β_i and γ_i :

$$\text{if } \beta_i \geq 0, \theta_i = \beta_i - \frac{\gamma_i}{2} \leq 0, \quad \text{else, } \theta_i = \beta_i + \frac{\gamma_i}{2} \geq 0, \quad (37)$$

which can be rewritten as follows:

$$\text{if } \theta_i \geq 0, \beta_i = \theta_i - \frac{\gamma_i}{2} \leq 0, \quad \text{else, } \beta_i = \theta_i + \frac{\gamma_i}{2} \geq 0. \quad (38)$$

The symmetry of relations [equation \(37\)](#) and [equation \(38\)](#) implies that β and θ have the same density:

$$f_{\text{St}}(\bar{\mathcal{E}}, \theta)(x) = \left(\frac{k(\bar{\mathcal{E}}, \alpha)}{\frac{\pi}{3} - p(\bar{\mathcal{E}}, \alpha)} \right) (1 - F_{\text{St}}(\bar{\mathcal{E}}, \alpha)(2|x|)). \quad (39)$$

[Figure 11](#) compares the distributions of γ , β and θ for a PD tessellation and a Lloyd relaxed PD tessellation. A very good agreement is obtained between the actual distributions and the different pdf expressions [equation \(29\)](#), [equation \(33\)](#) and [equation \(39\)](#).

6.4 Probability density function of the tortuosity of a path following the inclination rule

Equation [equation \(28\)](#) can be rewritten as follows:

$$T_{\text{St}}(n) = \frac{1}{\sum_{i=1}^n w_i \cos(\theta_i)}, \quad (40)$$

with weight w_i defined by the proportion of the i^{th} edge length in the total length of the path:

$$l_i = w_i \sum_{k=1}^n l_k. \quad (41)$$

It is worth noting that with such a definition, summation of w_i values is equal to 1.

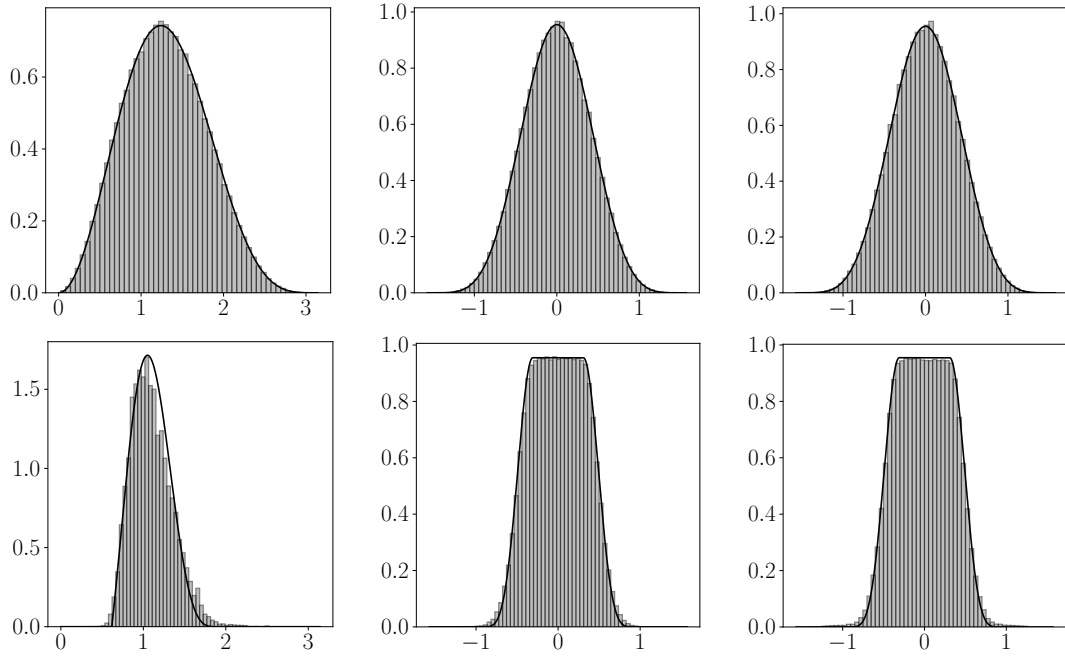


Figure 11: First line: angles distribution of a PD tessellation ($\tau = 500$). Second line: angles distribution of a Lloyd relaxed PD tessellation ($\mathcal{E} = 0.89$, $\tau = 500$). From left to right: γ , β and θ distributions. Black lines: density functions [equation \(29\)](#), [equation \(33\)](#) and [equation \(39\)](#).

[Figure 12](#) represents w_i values versus $\cos(\theta_i)$ values of a PD tessellation (left) and a Lloyd relaxed PD one (right). A correlation can be exhibited between w_i and $\cos(\theta_i)$ since low values of $\cos(\theta_i)$ seem to be associated to low values of w_i . The following hypothesis is therefore considered in the sequel for purpose of simplification:

The expected value of w_i knowing that $\cos(\theta_i) = c_0$ is proportional to c_0 . (42)

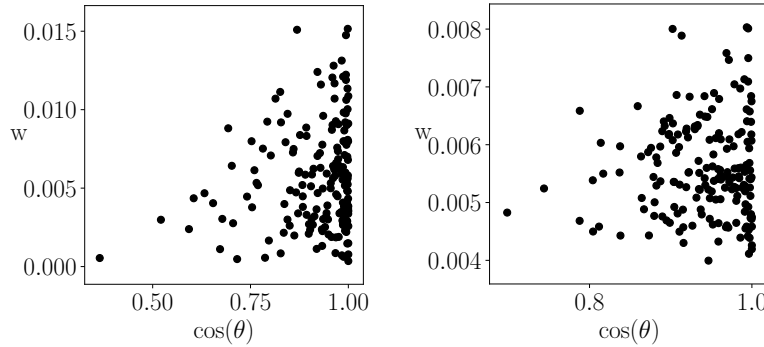


Figure 12: Weights values w_i versus $\cos(\theta_i)$ values of a PD tessellation (left) and a Lloyd relaxed PD one (right)

As a consequence, it is necessary to compute the correlation coefficient (which would differs from zero) between weigths w_i and $\cos(\theta_i)$ sets for the computation of the probabilistic density function of the tortuosity.

Splitting then I , the interval of accessible values for $\cos(\theta_i)$, into m intervals I_k , $k \in \llbracket 1; m \rrbracket$, the sum at denominator of expression [equation \(40\)](#) reads:

$$\sum_{i=1}^n w_i \cos(\theta_i) = \sum_{k=1}^m \left(\sum_{i=1}^{n_k} w_i \cos(\theta_i) \right), \quad (43)$$

where n_k corresponds to the number of $\cos(\theta)$ realizations which are in I_k .

Therefore if n_k is large enough and according to the central limit theorem, the sum can be approximated by:

$$\sum_{i=1}^n w_i \cos(\theta_i) \simeq \sum_{k=1}^m n_k \langle w \cos(\theta) \rangle_k = \sum_{k=1}^m \left(\sum_{i=1}^{n_k} \langle w \cos(\theta) \rangle_k \right), \quad (44)$$

where notation $\langle X \rangle_k$ stands for the expected value of the random variable X inside the interval I_k and θ and w are respectively the random variables associated to the values θ_i and w_i .

If the interval I_k is narrow enough, values $\cos(\theta_i)$ in I_k can be considered as equal to a constant value corresponding to the center c_k of interval I_k . Then $\langle w \cos(\theta) \rangle_k \simeq \langle w \rangle_k c_k$ and estimation [equation \(44\)](#) becomes:

$$\sum_{i=1}^n w_i \cos(\theta_i) \simeq \sum_{k=1}^m \left(\sum_{i=1}^{n_k} \langle w \rangle_k c_k \right). \quad (45)$$

Referring to the hypothesis [equation \(42\)](#), and remembering that the summation of w_i values is equal to 1, values of $\langle w \rangle_k$ can be expressed as follows:

$$\langle w \rangle_k \simeq \frac{c_k}{\sum_{i=1}^m n_i c_i}, \quad (46)$$

where c_i and n_i are respectively the center and the number of elements of the interval I_i .

Therefore the following estimate is accurate when the number of edges n is large enough and considering that hypothesis [equation \(42\)](#) is true:

$$\sum_{i=1}^n w_i \cos(\theta_i) \simeq \sum_{i=1}^n \frac{\cos(\theta_i)^2}{\sum_{k=1}^m \cos(\theta_k)}. \quad (47)$$

This estimate lead to the following approximation for the tortuosity of an n -edges path following the inclination rule [equation \(27\)](#):

$$T_{St}(n) \simeq \frac{\sum_{i=1}^n \cos(\theta_i)}{\sum_{i=1}^n \cos(\theta_i)^2}. \quad (48)$$

In the sequel the numerator and the denominator in expression [equation \(48\)](#) are respectively denoted by X and Y . Using an academic result on the variance of the sum of two correlated variables:

$$\sigma_{X+Y}^2 = \sigma_X^2 + \sigma_Y^2 + 2\text{Cov}(X, Y), \quad (49)$$

the correlation coefficient between variables X and Y can be expressed as follows:

$$\rho_{X,Y} = \frac{\text{Cov}(X, Y)}{\sigma_X \sigma_Y} = \frac{\sigma_{X+Y}^2 - \sigma_X^2 - \sigma_Y^2}{2\sigma_X \sigma_Y}, \quad (50)$$

where variances are estimated by the central limit theorem as follows:

$$\sigma_{X+Y}^2 \simeq n \sigma_{\cos(\theta) + \cos(\theta)^2}^2 \quad \sigma_X^2 \simeq n \sigma_{\cos(\theta)}^2 \quad \text{and} \quad \sigma_Y^2 \simeq n \sigma_{\cos(\theta)^2}^2. \quad (51)$$

Relations [equation \(50\)](#) and [equation \(51\)](#) imply that $\rho_{X,Y}$ does not depends of n . Expression [equation \(39\)](#) allows then to directly estimate $\rho_{X,Y}$ in function of $\tilde{\mathcal{E}}$.

Having estimated $\rho_{X,Y}$ and given the fact that $T_{St}(n)$ is a ratio between two correlated normal random variables, [Hinkley \(1969\)](#) gives the estimation of the probabilistic density function of $T_{St}(n)$:

$$f_{St}(\bar{\mathcal{E}}, T, n)(x) \simeq \frac{\sqrt{1 - \rho_{X,Y}^2}}{\pi \sigma_X \sigma_Y a(x)^2} \exp\left(-\frac{c}{2(1 - \rho_{X,Y}^2)}\right) + \frac{b(x)d(x)}{\sqrt{2\pi} \sigma_X \sigma_Y a(x)^3} \left[\phi\left(\frac{b(x)}{a(x)\sqrt{1 - \rho_{X,Y}^2}}\right) - \phi\left(-\frac{b(x)}{a(x)\sqrt{1 - \rho_{X,Y}^2}}\right) \right], \quad (52)$$

where

$$a(x) = \sqrt{\frac{x^2}{\sigma_X^2} - \frac{2\rho_{X,Y}x}{\sigma_X \sigma_Y} + \frac{1}{\sigma_Y^2}}, \quad b(x) = \frac{\mu_X x}{\sigma_X^2} - \frac{\rho_{X,Y}(\mu_X + \mu_Y x)}{\sigma_X \sigma_Y} + \frac{\mu_Y}{\sigma_Y^2}, \quad (53)$$

$$c = \frac{\mu_X^2}{\sigma_X^2} - \frac{2\rho_{X,Y}\mu_X\mu_Y}{\sigma_X \sigma_Y} + \frac{\mu_Y^2}{\sigma_Y^2}, \quad d(x) = \exp\left(\frac{b(x)^2 - ca(x)^2}{2(1 - \rho_{X,Y}^2)a(x)^2}\right), \quad (54)$$

$$\mu_X = \int_{-\frac{\pi}{2}}^{\frac{\pi}{2}} \cos(x) f_{St}(\bar{\mathcal{E}}, \theta)(x) dx, \quad \mu_Y = \int_{-\frac{\pi}{2}}^{\frac{\pi}{2}} \cos(x)^2 f_{St}(\bar{\mathcal{E}}, \theta)(x) dx, \quad (55)$$

and ϕ is the normal cumulative distribution function:

$$\phi(y) = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^y \exp\left(-\frac{1}{2}u^2\right) du. \quad (56)$$

[Figure 13](#) shows that [equation \(52\)](#) is accurate for PD tessellations while considering n -edges path following the inclination rule [equation \(27\)](#) with $n \geq 20$. For Lloyd relaxed PD tessellation, approximation [equation \(52\)](#) slightly underestimates the actual tortuosity of paths following the inclination rule.

[Figure 14](#) shows that estimate [equation \(52\)](#) leads to an overestimation of the actual geodesic tortuosity which tends to decrease as the quality index increases. As a result, estimate [equation \(52\)](#) appears to be accurate for unstructured Delaunay-like meshes suitable to finite element simulations (see [figure 14](#), right). This statement has been verified by a complementary study carried out on meshes generated by the Gmsh Delaunay algorithm.

[Figure 15](#) shows comparisons between the geodesic tortuosity pdf estimation [equation \(52\)](#) and distributions measured on 1000 Gmsh Delaunay meshes for several path lengths.

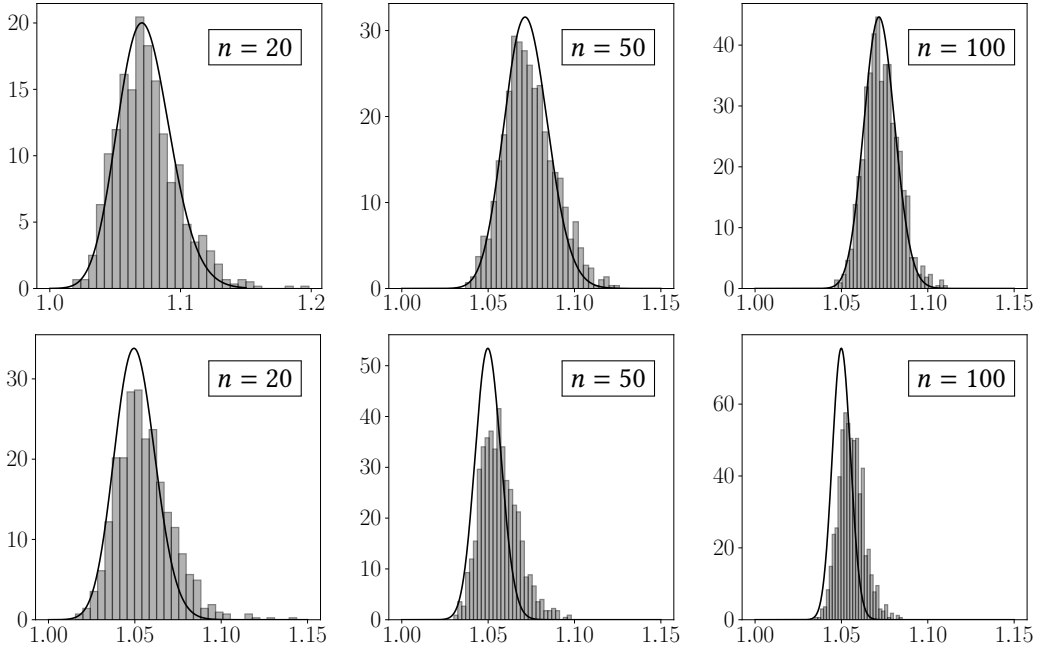


Figure 13: Distribution of the tortuosity of n -edges path following the inclination rule. Top: distribution obtained from a PD tessellation ($\tau = 500$). Bottom: distribution obtained from a Lloyd relaxed PD tessellation ($\mathcal{E} = 0.86$, $\tau = 500$). Left: 20-edges path, middle: 50-edges path and right: 100-edges path. Black line: tortuosity pdf estimation [equation \(52\)](#).

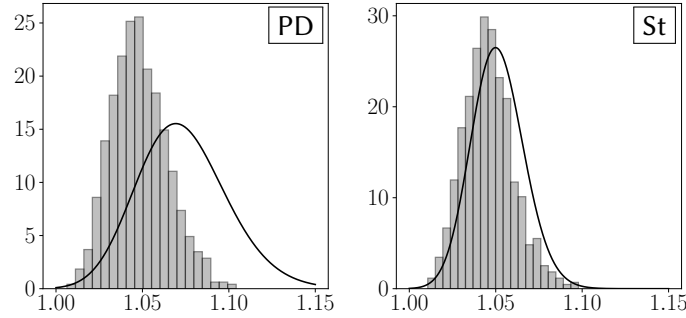


Figure 14: Left: geodesic tortuosity distribution of a PD tessellation ($\tau = 500$). Right: geodesic tortuosity distribution of a Lloyd relaxed PD tessellation ($\mathcal{E} = 0.89$, $\tau = 500$). Black lines: geodesic tortuosity pdf estimates [equation \(52\)](#).

The pdf estimates correctly mimic the geodesic tortuosity distributions of Gmsh Delaunay meshes with a slight overestimation: the average value obtained for Gmsh meshes is about 1.043 while estimate [equation \(52\)](#) gives an average value of 1.051.

6.5 Practical use of the geodesic tortuosity distribution of Delaunay meshes in computational fracture mechanics

In volumetric-cohesive finite-element approach, crack paths are restrained to propagate along elements boundaries (edges). If the localization and/or the orientation of the crack path is not known a priori, an isotropic tessellation is used, typically a Delaunay tessellation.

The fact that a Delaunay tessellation does not necessarily includes the theoretical path followed by the crack leads to the obtention of a more tortuous crack path which may lead to an overestimation of the crack length (or crack surface in 3D).

As the theoretical overall fracture energy W_{th} is equal to the Griffith energy multiply by the crack length times the specimen width (in 2D), the computed overall fracture energy W_{comp} may also be overestimated.

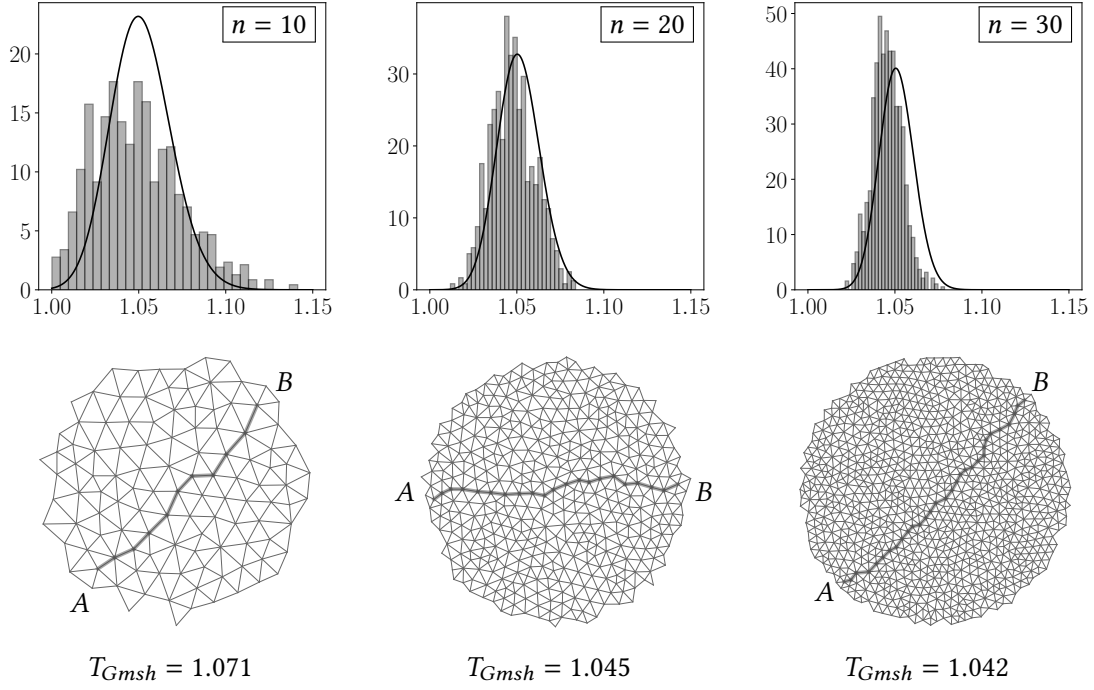


Figure 15: Obtained on Gmsh meshes ($\mathcal{E} = 0.89$, $\tau = 550$): top: distributions of geodesic tortuosities of cracks containing distinct number of edges n . Black lines: geodesic tortuosity pdf estimates [equation \(52\)](#) ; bottom: zoom on particular shortest path obtained on meshes generated by the Gmsh Delaunay algorithm (from left to right : paths containing respectively 10, 20 and 30 edges).

It is reasonably expected that the fracture energy and equivalently the crack length overestimations are of the same order of magnitude as the tortuosity estimate $T_{St}(n)$. Thus, denoting by L_{th} the theoretical length of the crack path and by L_{comp} the computed one, the following relations are obtained:

$$\frac{W_{comp}}{W_{th}}(n) = \frac{L_{comp}}{L_{th}}(n) \simeq T_{St}(n), \quad (57)$$

where n is the number of edges composing the crack path.

From [equation \(52\)](#), confidence intervals of geodesic tortuosity estimate $T_{St}(n)$ as well as its expected value can be derived and estimated for distinct values of n (see [figure 16](#)).

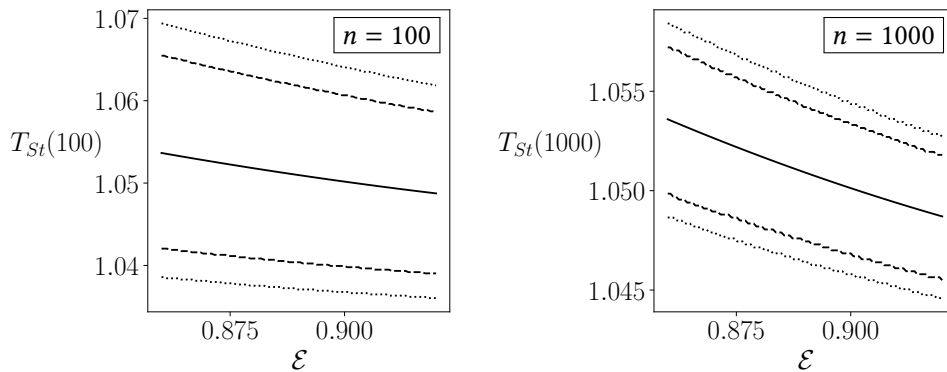


Figure 16: Estimate of the geodesic tortuosity expected value (continuous line) and 95% (dashed line) and 99% (dotted line) confidence intervals in function of the efficiency index. Left: for 100-edges paths. Right: for 1000-edges paths.

[Table 2](#) quantifies the dependency of the geodesic tortuosity estimate T_{St} distribution to the number of edges n composing the crack path.

Number of edges	Mean tortuosity	95% certainty interval	99% certainty interval
20	1.0512	[1.0279; 1.0760]	[1.0212; 1.0844]
50	1.0510	[1.0362; 1.0665]	[1.0317; 1.0716]
100	1.0510	[1.0404; 1.0617]	[1.0371; 1.0653]
1000	1.0509	[1.0474; 1.0542]	[1.0464; 1.0553]

Table 2: Estimates of the geodesic tortuosity expected value and 95% and 99% confidence intervals in function of the number of edges ($\mathcal{E} = 0.89$).

Considering relation [equation \(57\)](#), and [table 2](#), we expect that an overestimation of the fracture energy of about 5% may be obtained while using an unstructured Delaunay-like mesh in a cohesive-volumetric finite elements simulation.

7 Conclusion

The main objective of this paper was to derive accurate estimates for both the mesh density and geodesic tortuosity of planar triangular tessellations suitable for finite element simulations in the context of cohesive-based fracture mechanics. The following results were obtained.

- New estimates for the edge length and the triangle top angle distributions were proposed for Poisson-Delaunay tessellations and Lloyd relaxed ones for arbitrary efficiency indexes and intensities (equations [equation \(14\)](#) and [equation \(15\)](#)).
- The edge length and triangle top angle distributions of random triangular meshes obtained by standard mesh generators (Abaqus, Gmsh, Cast3M and Triangle) have been compared to these estimates depending on the intensity of the underlying point process and its efficiency index. These estimates are accurate for meshes with efficiency index at least less than 0.92. In other words, estimates are no longer accurate for semi-structured meshes (obtained for exemple by the GMSH “Frontal” algorithm and Abaqus). Moreover, these estimates are robust enough to be used at a local scale (see [appendix A](#)).
- Based on these practical estimates, the distribution of the mesh density was derived as a normal law (equation [equation \(22\)](#)). Its mean grows as the square root of the intensity of the point process, while its variance does not depend on this intensity but decreases when the area of the discretized domain increases. The expected value of the mesh density does not depend on the global efficiency index. A dependence to the efficiency index is observed for the variance of the mesh density. However, for standard meshes used in finite elements simulations, the variance of the mesh density appears to be negligible in comparison to its expected value. It is then stated that the mesh density can be viewed as a singular function of the intensity of the underlying point-process (equal to the mesh density expected value). The mesh densities of usual structured triangular meshes are also given in [appendix C](#).
- For unstructured Delaunay-like meshes (whose edge orientations may be considered uniformly distributed), a tortuosity estimate has been derived from an hypothetical “orientation rule” followed by each edge of a shortest path between two nodes. Since the orientation rule is not exactly fulfilled by shortest paths, the estimated tortuosity is slightly overestimated to the actual mesh geodesic tortuosity. However, for application we have in mind, the overestimation is small enough for considering that the geodesic tortuosity estimate [equation \(52\)](#) is accurate.
- Overestimation of both the overall fracture energy and cracks tortuosities computed by a cohesive-volumetric finite elements simulation are expected while using an unstructured Delaunay-like mesh. Using the geodesic tortuosity estimate T_{St} , we deduced that this overestimation is approximatively equal to 5%.

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A Local robustness of the edge length and triangle top angle distributions [equation \(14\)](#) and [equation \(15\)](#)

Since we have in mind fracture mechanics of complex microstructures, involving local heterogeneities, gradients, cavities, etc., we have to check at which scale or for which local intensity τ estimates [equation \(14\)](#) and [equation \(15\)](#) are accurate enough. The accuracy is quantified by the standard deviations of the means $\mu_{St}(\tau, l)$ and $\mu_{St}(\alpha)$, respectively denoted by $\sigma_{\mu_{St}}(\tau, l)$ and $\sigma_{\mu_{St}}(\alpha)$, and by the standard deviations of the variances $\sigma_{St}^2(\tau, l)$ and $\sigma_{St}^2(\alpha)$, respectively denoted by $\sigma_{\sigma_{St}^2}(\tau, l)$ and $\sigma_{\sigma_{St}^2}(\alpha)$, for 100 samples. Each sample corresponds to a subset of a Gmsh random triangular mesh (“MeshAdapt” strategy, efficiency index $\mathcal{E} \simeq 0.86$) consisting of a disk with random center and radius R , see [figure 17](#). The mesh of [figure 17](#) has an efficiency index of $\mathcal{E} \simeq 0.86$.

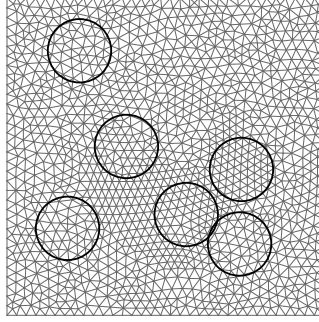


Figure 17: Six circular subset samples with random center and radius $R \simeq 3.3\mu_{St}(\tau, l)$.

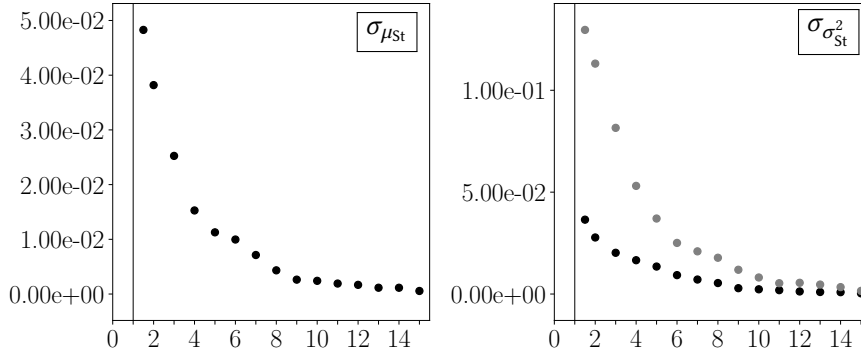


Figure 18: Left: standard deviation of the mean of the edge length divided by its actual mean with respect to $R/\mu_{St}(\tau, l)$. Right: standard deviation of the variance of the edge length (resp. the triangle top angle) divided by its actual variance with respect to $R/\mu_{St}(\tau, l)$ (black dots/resp. gray dots). The standard deviations are computed over 100 samples of circular random subsets belonging to a Gmsh triangular mesh (“MeshAdapt” strategy, efficiency index $\mathcal{E} \simeq 0.86$).

[Figure 18](#) shows that, when the radius R is larger than the expected mean of the edge length, $R \geq \mu_{St}(\tau, l)$, the standard deviation of the mean $\sigma_{\mu_{St}}(\tau, l)$ is lower than about 5% of the actual edge length mean and the standard deviations of the variances $\sigma_{\sigma_{St}^2}(l)$ and $\sigma_{\sigma_{St}^2}(\alpha)$ are respectively lower than 4% and 12% of the corresponding actual variances. In addition, the standard deviation of the mean $\mu_{St}(\alpha)$, which is not represented in [figure 18](#), is lower than the machine epsilon whatever $R \geq \mu_{St}(\tau, l)$. Finally, the local robustness of [equation \(14\)](#) and [equation \(15\)](#) is such that they can be applied at the length scale of a few triangles in an unstructured Delaunay-like mesh.

B Some hints on the range of validity for estimates [equation \(14\)](#) and [equation \(15\)](#)

The proposed estimates [equation \(14\)](#) and [equation \(15\)](#) are strictly devoted to unstructured Delaunay-like meshes. These estimates are not able to deal with:

- *Badly shaped elements.* Efficiency indexes lower than the minimal value \mathcal{E}_{PD} corresponding to PD tessellations (see [equation \(11\)](#)) can not be taken into account by these estimates. They are restricted to “good” enough meshes (i.e. with a efficiency index greater than 0.725).
- *Structured or quasi-structured meshes.* Structured meshes lead to discrete values of the edge length and triangle top angle sets. Their distributions thus correspond to successive Dirac delta functions. Estimates [equation \(14\)](#) and [equation \(15\)](#) are restricted to “sufficiently unstructured” meshes (i.e. with a efficiency index at least less than 0.92).

To underline these two limitations, planar hand-generated meshes are considered as counterexamples. These meshes consist in cross-triangular square patterns stretched in one direction (degradation of the shape) with random perturbation of the inner points (in order to control the regularity), tessellations remaining valid. Three cases are considered.

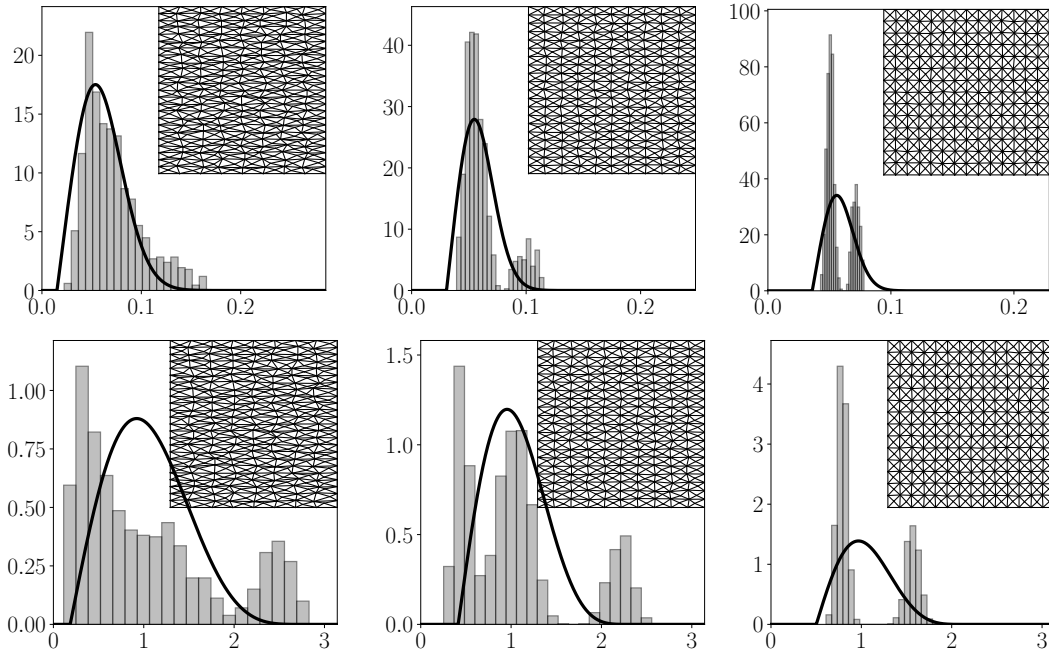


Figure 19: Limit of validity for estimates [equation \(14\)](#) and [equation \(15\)](#) in case of quasi-structured meshes. Pdf of the edge length (first row) and triangle top angle (second row) sets for hand-generated cross-triangular square meshes, stretched in the horizontal direction and randomly perturbed of the inner points. Left: stretching/elongation ratio of 3 with high distortion, efficiency index $\mathcal{E} \approx 0.78$, intensity $\tau = 466$. Middle: stretching ratio of 2 with intermediate distortion, efficiency index $\mathcal{E} \approx 0.84$, intensity $\tau = 472$. Right: stretching ratio of 1 with small distortion, efficiency index $\mathcal{E} \approx 0.87$, intensity $\tau = 458$.

With highly random perturbation, the mesh is no more structured and estimates [equation \(14\)](#) and [equation \(15\)](#) exhibit a rough validity. In addition, these estimates do not accurately predict the tail of the edge length and triangle angle distributions when badly shaped elements are numerous; see [figure 19](#), left. With small random perturbation, the edge length distribution consists in multi-modal normal distribution. Estimates [equation \(14\)](#) and [equation \(15\)](#) are able to predict the mode with higher intensity, but are not able to predict the other modes; [figure 19](#), middle. With even smaller random perturbation, a quasi-structured mesh is obtained and estimates [equation \(14\)](#) and [equation \(15\)](#) fail to predict the quasi-Dirac combs; [figure 19](#), right.

Thus, the use of estimates [equation \(14\)](#) and [equation \(15\)](#) is restricted to standard random triangular tessellations whose efficiency index is greater than or equal to the one of a PD tessellation and at least less than 0.92 (see [section 4.2](#)).

C Complementary results for structured triangular meshes

The estimate of the total length of the edges of a tessellation containing N_P points regularly distributed in an area A has been presented in [Watanabe \(2008\)](#) for various measures. Their results concern some triangular meshes, square meshes and hexagonal meshes. We focus on the mesh density Z of structured triangular meshes usually used in finite element discretizations ([M. Tijssens et al. 2000](#)).

The total length of edges in a structured triangular tessellation based on a hexagonal (H) point process reads (see [table 3](#)):

$$\sum_{i=1}^{N_E} l_i = 3 \sqrt{\frac{2\tau}{\sqrt{3}}} A. \quad (58)$$

The corresponding mesh density is thus

$$Z_H = 3 \sqrt{\frac{2}{\sqrt{3}}} \sqrt{\tau}. \quad (59)$$

Regular points distributed along a square lattice (see [table 3](#)) lead to a structured “Triangle Quadrilateral” (TQ) pattern whose the total length of the edges reads:

$$\sum_{i=1}^{N_E} l_i = (2 + \sqrt{2}) \sqrt{\tau} A, \quad (60)$$

and the mesh density

$$Z_{TQ} = (2 + \sqrt{2}) \sqrt{\tau}. \quad (61)$$

In fracture mechanics, “Cross-Triangle Quadrilateral” (CTQ) patterns are often used (see [table 3](#)). For that mesh topology, the mesh density is deduced from the previous one:

$$Z_{CTQ} = 2 (1 + \sqrt{2}) \sqrt{\tau}. \quad (62)$$

This CTQ pattern can be stretched in one direction and the underlying quadrilateral is not a square but a rectangle (aspect ratio $s \geq 1$). In this case, the mesh density becomes:

$$Z_{CTQs} = \left(2 \sqrt{1 + \frac{1}{s^2}} + 1 + \frac{1}{s} \right) \sqrt{\tau}. \quad (63)$$

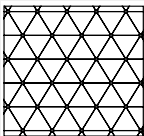
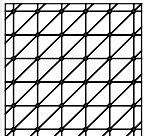
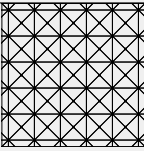
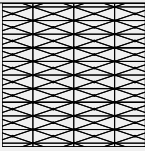
	$Z_H = 3 \sqrt{\frac{2}{\sqrt{3}}} \sqrt{\tau}$		$Z_{TQ} = (2 + \sqrt{2}) \sqrt{\tau}$
	$Z_{CTQ} = 2 (1 + \sqrt{2}) \sqrt{\tau}$		$Z_{CTQs} = \left(2 \sqrt{1 + \frac{1}{s^2}} + 1 + \frac{1}{s} \right) \sqrt{\tau}$

Table 3: Mesh density of the structured triangular meshes usually used in finite element discretizations

[Table 3](#) summarizes the mesh densities for the meshes usually used in finite element discretizations. It is worth noting that these mesh densities for structured triangular meshes evolve as $\sqrt{\tau}$, exactly as the mean of the mesh densities for random triangular meshes do. In particular, the mean $\mu_{St}(\tau, Z)$ of standard random triangular meshes given in [equation \(22\)](#)

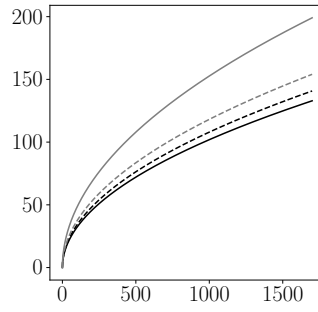


Figure 20: The mesh density Z as a function of the point process intensity τ for different types of structured triangular meshes. Gray line: Cross-Triangle Quadrilateral mesh [equation \(62\)](#). Gray dashed line: stretched Cross-Triangle Quadrilateral mesh [equation \(63\)](#) with stretch factor $s = 2$. Black dashed line: Triangle Quadrilateral mesh [equation \(61\)](#). Black line: hexagonal mesh [equation \(59\)](#).

is very close to the mesh density of TQ tessellations: the prefactor of $\mu_{St}(\tau, Z)$ is about 3.40, while the prefactor of $Z_{TQ}(\tau)$ is about 3.41:

$$Z_{TQ}(\tau) \simeq \mu_{St}(\tau, Z). \quad (64)$$

From the point of view of the mesh density, a Triangle Quadrilateral pattern accurately mimics a standard random triangular mesh (keeping in mind that the distribution of the edge length set is bimodal in TQ tessellations). Moreover, the mesh density Z of stretched CTQ (CTQs) decreases as the stretch factor s increases, see [figure 20](#).