# Intersection Cuts for Mixed-Integer Signomial Sets 

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## 1 Introduction

Intersection cuts $[3,10]$ have been tackled in recent years to tighten the polyhedral outer approximation of certain non-convex sets such as lattice set [1], bi-level set [5], outer product set [2] and quadratic set [9].
In this work, we study the intersection cuts for signomial set and mixed-integer-signomial set. Signomial Programs (SPs) and Mixed-Integer Signomial Programs (MISPs) [6, 8] contain signomial terms, the intersection cuts can tighten their Linear Programming (LP) relaxations.

Given an $n$-dimensional multi-index $\alpha=\left(\alpha_{1}, \alpha_{2}, \ldots, \alpha_{n}\right) \in \mathbf{R}^{n}$, a signomial term is defined as $x^{\alpha}:=\prod_{i=1}^{n} x_{i}^{\alpha_{i}}$. We assume the variables being positive, thus allowing the multi-index $\alpha$ to have negative components. We define the signomial set as follows :

$$
\mathcal{S}_{s}:=\left\{(x, y) \in \mathbf{R}_{+}^{n} \times \mathbf{R}: y(\cdot) x^{\alpha}\right\},
$$

where $\cdot$ can be $\leq,=$ or $\geq$.
We define the mixed-integer counterpart of the signomial set. Let $[\underline{x}, \bar{x}]$ be the box constraint on $x$ such that $\underline{x}>0$ implicitly imposes the positiveness of $x$. W.l.o.g., we assume the first $n_{1} \leq n$ entries of $x$ to be integer, and let $\mathcal{D}=\left(\mathbf{Z}_{+}^{n_{1}} \times \mathbf{R}_{+}^{n-n_{1}}\right) \cap[\underline{x}, \bar{x}]$ be the bounded mixedinteger domain. Similarly, the mixed-integer-signomial set is defined as :

$$
\mathcal{S}_{\text {mis }}:=\left\{(x, y) \in \mathcal{D} \times \mathbf{R}: y(\cdot) x^{\alpha}\right\} .
$$

For a non-convex set $\mathcal{S}$, a convex set $\mathcal{C}$ is $\mathcal{S}$-free if $\operatorname{int} \mathcal{C} \cap \mathcal{S}=\emptyset$. The intersection cut framework requires two ingredients to generate a valid inequality to separate $\bar{z}$ :
i) a translated simplicial conic relaxation $\mathcal{R}$ with vertex $\bar{z}$ such that $\mathcal{S} \subset \mathcal{R}$ and $\bar{z} \notin \mathcal{S}$;
ii) a $\mathcal{S}$-free set $\mathcal{C}$ such that $\bar{z} \in \operatorname{int} \mathcal{C}$.

The spatial-branch-and-bound (sBB) algorithm [7] is a general technique to find the global solution of non-convex MINLPs. The second ingredient of intersection cuts, i.e., recession conic relaxation, can be retrieved from the LP-based sBB algorithm. If two $\mathcal{S}$-free sets $\mathcal{C}_{1}, \mathcal{C}_{2}$ are such that $\mathcal{C}_{1} \subset \mathcal{C}_{2}$, the intersection cut derived from $\mathcal{C}_{2}$ dominates the one derived from $\mathcal{C}_{1}$ [4]. Therefore, the maximal $\mathcal{S}$-free sets are ideal and crucial to generate strong intersection cuts.
The set $\mathcal{S}$ has a difference-of-concave (DC) representation if it can be expressed as

$$
\mathcal{S}:=\left\{z \in \mathbf{R}^{p}: f_{1}(z)-f_{2}(z) \leq 0\right\},
$$

where $f_{1}$ and $f_{2}$ are concave functions.
[10] proposed the linearization technique to derive $\mathcal{S}$-free sets in DC-representations. The linearization technique constructs the following $\mathcal{S}$-free set :

$$
\mathcal{C}:=\left\{z \in \mathbf{R}^{p}: f_{1}(z)-\left(f_{2}(\tilde{z})+\nabla f_{2}(\tilde{z}) \cdot(z-\tilde{z})\right) \geq 0\right\},
$$

where $f_{2}(\tilde{z})+\nabla f_{2}(\tilde{z}) \cdot(z-\tilde{z})$ is the linearization (over-estimator) of $f_{2}$ at $\tilde{z}$.
Indeed, a non-convex set $\mathcal{S}$ may have various DC-representations. An ideal DC-representation should yield maximal $\mathcal{S}$-free sets using the linearization technique.

## 2 Our results

We show that the signomial set can be transformed into the following normalized DCrepresentation :

$$
\begin{equation*}
\mathcal{S}_{s}=\left\{(u, v) \in \mathbf{R}_{+}^{h} \times \mathbf{R}_{+}^{l}: u^{\beta}-v^{\gamma} \leq 0\right\}, \tag{1}
\end{equation*}
$$

where $\beta \in \mathbf{R}_{+}^{h}$ and $\gamma \in \mathbf{R}_{+}^{l}$ are $h-$ and $l-$ dimensional multi-indices such that $\max \left\{\|\beta\|_{1},\|\gamma\|_{1}\right\}=$ 1. Because $u^{\beta}, v^{\gamma}$ are concave, the representation yields a $\mathcal{S}_{s}-$ free set using the linearization technique.
For any non-convex set $\mathcal{S}$, we give a sufficient condition to guarantee the maximality of a $\mathcal{S}$-free set; then, we show that the normalized DC-representation is ideal in the sense that it yields the maximal $\mathcal{S}_{s}-$ free set in the non-negative orthant. Some non-ideal DC-representations of $\mathcal{S}_{s}$ are given as well.
We propose a method to enlarge the $\mathcal{S}_{s}$-free set into the mixed-integer-signomial-free ( $\mathcal{S}_{\text {mis }}-$ free) set, where both nonlinear constraints and bounded/integer variables are featured. A byproduct is a strong convex relaxation of the following mixed-integer power cone :

$$
\begin{equation*}
\left\{(x, t) \in \mathcal{D} \times \mathbf{R}: x^{\alpha} \geq t\right\} \tag{2}
\end{equation*}
$$

where $\alpha \in \mathbf{R}_{+}^{p},\|\alpha\|_{1} \leq 1$, and $\mathcal{D} \subset \mathbf{R}_{+}^{p}$ is a bounded mixed-integer domain. The relaxation is stronger than the continuous relaxation $\left\{(x, t) \in \mathbf{R}_{+}^{p} \times \mathbf{R}: x^{\alpha} \geq t\right\}$.

The intersection cut separation problem is reduced to univariate root-finding problems. Such root-finding problems can be solved numerically in quadratic convergence rate by Newton-like algorithms.

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