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Article

# Energy Efficient Hardware and Improved Cluster-Tree Topology for Lifetime Prolongation in ZigBee Sensor Networks

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**Abstract:** In wireless sensor networks, building energy-efficient systems is one of the major challenges. In such networks, nodes are usually supplied by low power and small batteries. Many factors are involved in the energy consumption, and this issue may be considered as a cross-layer problem, from the hardware architecture to the application layer. This paper aims at presenting a hybrid solution for sensor networks based on two main aspects. The first one is the hardware architecture, where we present a prototype of a sensor node we designed. This node proved its efficiency in terms of energy consumption. The second aspect is related to the topology construction and presents a new topology control algorithm based on graph computing. Thus, our system consists of a real indoor application for temperature and humidity monitoring, applicable to home automation or industrial monitoring. We performed the experiments using a set of sensor nodes deployed over a building and proved the efficiency of the system in terms of energy consumption, network lifetime and data delivery.

**Keywords:** sensor architecture; ZigBee network; topology; energy consumption

## 1. Introduction

Today, we assist the great emergence of wireless sensor networks. These kinds of networks are self-organized and consist in general of a large amount of autonomous sensor nodes with low resources transmitting sensed data to the base stations [1]. A sensor network can be easily deployed over large areas, but once deployed, it can be very difficult or even impossible to access. That is why most of these networks intend to optimize the use of the limited energy contained in each node. Energy optimization can be obtained by an efficient hardware architecture, efficient data routing and aggregation, energy harvesting and many other ways. Many application fields tend to take advantage of the technology of wireless sensor networks. The principal fields are related to temperature and humidity monitoring, surveillance, rescue operations and home automation. Most of the applications that require large-scale networks, such as the detection of forest fires, use high range transmission standards like IEEE 802.11b. Other applications, such as home automation that require less transmission range, use other standards, such as IEEE 802.15.4 related to ZigBee specifications. This standard is normally used for home automation and industrial monitoring applications. In fact, many works have demonstrated that conventional wireless standards such as Bluetooth or “WLAN” [2,3] are not suitable to be used

in this type of network. It has also been shown that the best standard taking into consideration the constraints of indoor networks is IEEE 802.15.4 [4]. Furthermore, the ZigBee network is has the least, compared with other wireless networks, SNR [5,6], which makes it the best standard for indoor applications, where noise has a strong presence, especially in industrial manufacturing.

It is also important to note that several works have been conducted to improve the energy efficiency of sensor networks by enhancing routing protocols, topology control and data aggregation, and the results are confirmed using simulation tools such as Ns-2, Omnet, MATLAB, etc. [7,8]. However, building energy-efficient hardware has not yet gained much interest from the research community. Consequently, most of the commercialized sensor nodes have a high consumption level, energy still being an unresolved problem. In addition, the commercialized nodes do not give any possibility to improve the performances related to the physical and MAC layers, namely the registers of transmission and reception, the size of packets, and others.

Thus, the first challenge that we tackle in this paper is to design a prototype of the sensor node that reduces the energy consumption while maintaining a high level of performance. As a matter of fact, the idea was to use a combination of a ZigBee radio transceiver provided by Microchip [9] and an "Arduino" microcontroller [10]. Our designed node showed satisfying results in terms of energy saving and the reliability of data processing. Secondly, we focus on topology construction to enhance the performances of the sensor network. We developed an improved cluster-tree topology based on the minimum spanning tree (MST) that we called minimum spanning cluster-tree (MSCT).

The remainder of the paper is organized as follows; Section 2 presents some concepts related to our contribution, namely ZigBee specifications and graph theory concepts. Section 3 describes our hardware architecture based on the ZigBee transceiver and Arduino microcontroller, as well as an overview of our system. In Section 4, we treat our developed topology based on minimum spanning trees. Section 5 evaluates the performances of the system by presenting our experimental scenario and results. Finally, Section 6 concludes the paper.

## 2. Background

In this section, we present some concepts related to the paper's contributions, where we introduce graph theory related to the construction of the minimum spanning tree topology, as well as the particularities of ZigBee networks.

### 2.1. Minimum Spanning Tree

A minimum spanning tree (MST) of an undirected graph is a graph that spans all the nodes as vertices and contains no cycles [11]. It can also be presented as a sub-graph that contains all the vertices where the sum of the weights of edges is no larger than the total weight of any other spanning tree from the graph [12]. The concepts of the spanning tree may be applied to construct the topology. Thus, several algorithms for routing and topology control based on MST have been developed. These algorithms use a few methods that allow the construction of the MST-based topology, such as the "relax" algorithm [13], the "adjust weight" algorithm, the "adjust tree" algorithm [14], the "STPtree-based" algorithm [15] and others. In general, spanning tree algorithms have shown energy efficiency for different extreme sensor network conditions. The Kruskal algorithm [16] starts by computing the cost between all nodes and continues until all the possible pairs have calculated there costs. Then, it sorts all edges in increasing order and considers all of these edges for inclusion in the tree. The tree initially has all vertices, but no edges; an edge is included in the tree if and only if its addition does not create a cycle in the already constructed tree.

## 2.2. ZigBee Network

IEEE 802.15.4 is a low-power and low-cost standard for wireless communications, destined for personal area networks (PAN). ZigBee networks are basically composed of one or many PAN coordinators, which monitor the entire network by dealing with the ordinary nodes that request association with the PAN managed by the coordinator in order to participate in the network's operations [17]. Two layers are specified in the IEEE 802.15.4 standard: the physical and MAC layers.

The physical layer supports three frequency bands [17]: a 2450-MHz band (16 channels), a 915-MHz band (10 channels) and an 868-MHz band (1 channel). Concerning the MAC layer, the 802.15.4 standard defines two types of nodes: reduced function devices (RFD) and full function devices (FFD) [17]. The FFD devices are equipped with all MAC layer functions. ZigBee also standardizes the higher layers of the protocol stack [18]. The network layer has the task of organizing routing over a multi-hop network, while the application layer intends to provide a framework for distributed application development and communication.

ZigBee specifies three principal devices: the coordinator, routers and end devices; where routers and coordinator must be FFD. From these devices, three types of topologies are defined: star, mesh and cluster tree. The star topology [19] consists only of end-devices sending packets to the coordinator. The mesh topology [20] is more complex, but robust and resilient to faults. Routers maintain a routing table and employ a route discovery algorithm to construct and update data structures on the path nodes. The routing table is generally consulted by a node before any transmission to decide about the next hop.

The cluster-tree [21] is a simple hierarchical topology based on a relationship between the father and children, as shown in Figure 1. In this topology, several clusters are present and able to communicate using multi-hop routes. They are also controlled by a PAN coordinator [22,23]. The cluster topology performs as follows: after the network's construction based on association-request and association-response, the parents of the clusters form a tree structure and act as intermediate routers. For better network control, the maximum number of children and the maximum depth of the tree that each router may have need to be fixed.

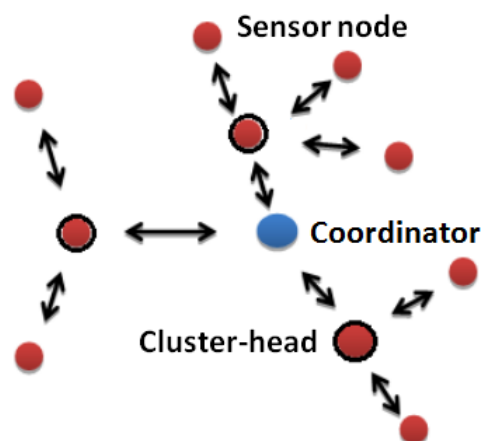


Figure 1. Cluster-tree topology.

In the next section, we present the architecture of our sensor nodes based on ZigBee specifications over the Arduino platform. The system was designed and programmed in order to follow the characteristics cited above.

### 3. Hardware and System

#### 3.1. Hardware Architecture

In this section, we present the hardware architecture of our designed sensor node. Basically, a node includes three main components: sensor, processor and radio frequency (RF) transceiver.

The sensor is a device that produces data measurements in order to monitor the physical environment. Sensed data could be temperature, pressure, humidity, voltage, etc. The sensor gives an analog signal based on the observed phenomenon converted to digital information by the AD converters and sent to the processor for further processing. In our system, we opt for the “Aosong DHT22” temperature-humidity sensor.

The processor is generally integrated with the microcontroller. It performs tasks, processes data and coordinates the functionality of all components connected to the sensor board. In our system, we used the “Atmel Sam3x8e ARM Cortex-M3” micro-processor embedded on an “Arduino Due” microcontroller [24]. A similar system using the Xbee mote, also based on Arduino, has been presented in [25] for a railway control case. Another advantage of working with the Arduino microcontroller is that it gives an easy programming interface for all layers, while many other commercial nodes are fixed for the physical and MAC layers.

The radio frequency (RF) transceiver connects the sensor node to the network. Generally, the RF module operates in three different modes: transmit, receive and sleep. The transmit mode consumes more energy. The energy consumed during the other modes may vary according to the used medium access algorithm. In our system, we used the “Microship MRF24J40MA” ZigBee transceiver [9] for routers and end-devices and the “Microship MRF24J40MD” for the coordinator due to its high transmission range.

Tables 1 and 2 present comparisons between the characteristics of some sensor nodes, including our designed sensor node, based respectively on the CPU, memory and data rate in Table 1 and power consumption, range and sensitivity in Table 2.

This comparison shows that our hardware gives an alternative solution concerning the transmission range-energy consumption problem, by reducing considerably the energy consumption while maintaining a high transmission range compared with the other sensor nodes.

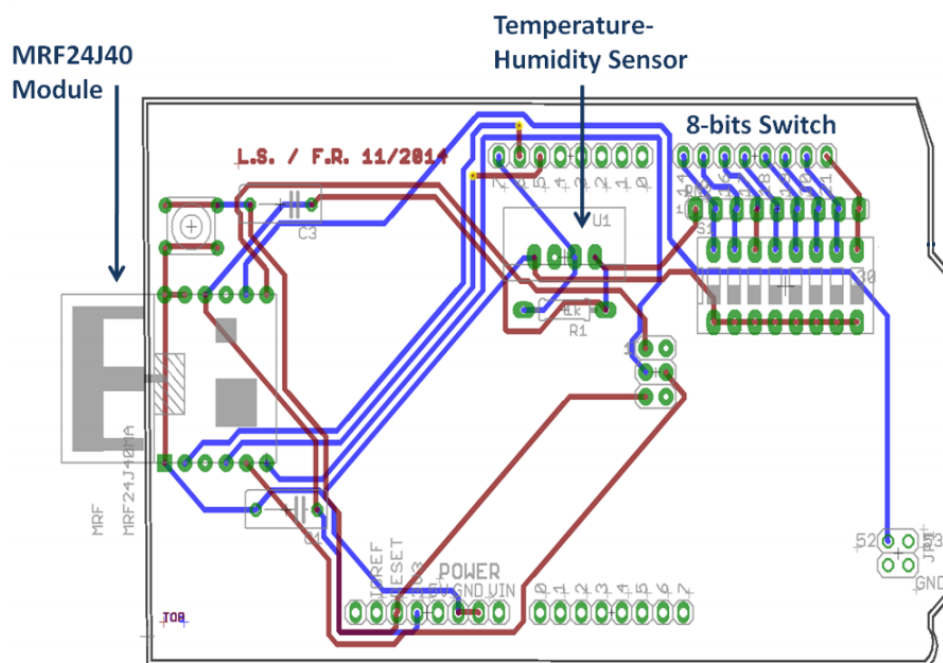
**Table 1.** Sensor Nodes’ Comparison 1: hardware specifications.

Name of Sensor Node	Micro Controller	RF Module	Flash Memory	Program Memory	Frequency Range	Data Rate (kbps)
Texas Instruments Mote	ARM Cortex-M3	SmartRF05EB	512 K	8 K RAM	868–870 MHz	250
Microchip	PIC24	MRF24J40	128 K	4 K RAM	2.405–2.48 GHz	250
TelosB	Atmega128	TI CC2420	48 K	10 K RAM	2.4 GHz	250
Memsic	ATmega128L	TI Chipcon CC2420	128 K	4 K RAM	2.4 GHz	250
Tiny Node	TI MSP430	Smart RF 125 MTD	512 K	8 K RAM	2.4–2.4835 GHz	152.3
Lotus	NXPLPC1758	RF231 Atmel	512 K	64 K SRAM	2.4 GHz	250
Sun SPOT	ARM 920T	TI CC2420	512 K	128 K + 4 K + 4 K	2.4 GHz	250
Our Designed Node	Atmel SAM3X8E ARM	MRF24J40	512 K	64 K RAM	2.4 GHz	250

**Table 2.** Sensor Nodes Comparison 2: energy and power.

Name of Sensor Nodes	Transmitter Power	Average Consumption	Average Indoor Range	Average Outdoor Range	Receive Sensitivity
Texas Instruments Mote	−24 dBm–0 dBm	24–30 mA	108 ft	173 ft	−90 dBm (min), −94 dBm
Microchip	−24 dBm–0 dBm	19–27 mA	110 ft	183 ft	−94 dBm
TelosB	−28 dBm–0 dBm	27–36 mA	82 ft	102 ft	−90 dBm (min), −94 dBm
Memsic	−26 dBm–0 dBm	25–37 mA	89 ft	103 ft	−90 dBm (min), −94 dBm
Tiny Node	max 12 dBm	43–62 mA	121 ft	210 ft	−106 dBm
Lotus	3 dBm	34–52 mA	51 ft	82 ft	−101 dBm
Sun SPOT	−24 dBm–0 dBm	23–46 mA	112 ft	169 ft	−90 dBm (min)
Our Designed Node	−24 dBm–0 dBm	17–25 mA	117 ft	153 ft	−90 dBm (min), −92 dBm

The circuit design of the node (excluding the Arduino board) is modeled in Figure 2, and a prototype of it is shown in Figure 3.



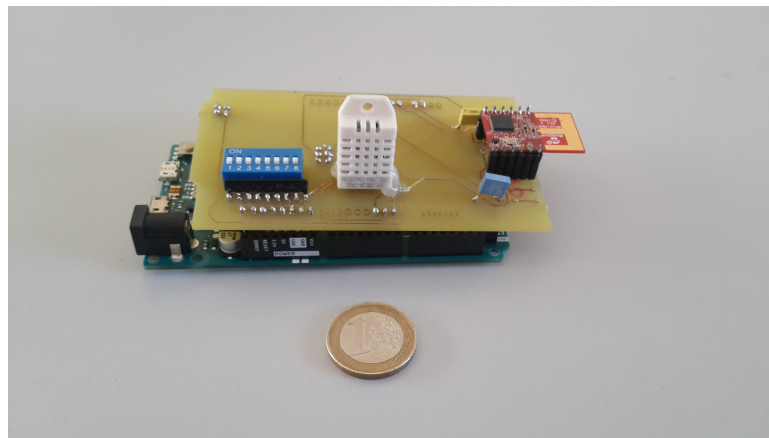
**Figure 2.** Circuit design of the sensor node.

We also note that, besides the three main components cited above, the sensor node includes other hardware components: one eight-bit, so the node provides a MAC address to nodes before the logical address assignment, and two capacitors. The circuit board, including all electronic components, is finally connected to the Arduino microcontroller.

In order to highlight the added value of our sensor device in terms of energy consumption and the performance of the network, it is imperative to note that the used microcontroller and ZigBee transceiver are, in addition to the low power consumption, open to be configured in all network layers, contrary to the commercial sensor nodes described in Tables 1 and 2, where it is very difficult to access or configure the operations performed in the physical and MAC layers. However, the ZigBee specifications only define three topologies inside the MAC layer: star, mesh and cluster-tree. Since the



goal of this work is to develop and test a new topology, it was necessary to design a new sensor node prototype that must be fully open to configuration.



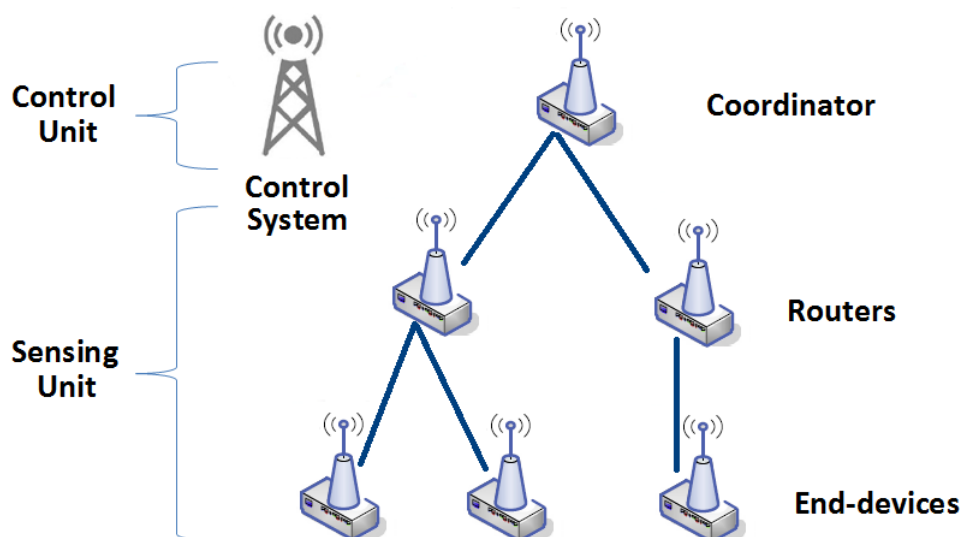
**Figure 3.** The sensor node prototype.

The development of this new device has also allowed performing some improvements related to the physical layer and transceiver registers, such as reducing the size of packets to minimize the computational delay and save energy.

Thus, the hardware has permitted obtaining a new platform that represents a cross-layer solution, where users can define the specifications of the network based on the type and application of the network.

### 3.2. System Overview

In this section, we present the network system based on our designed sensor nodes. The system is based on the IEEE 802.15.4 standard described in the previous section to allow low range communications and energy consumption between the different nodes. It includes a set of sensor nodes organized as a network. The coordinator is connected to a control system acting as the base station. An illustration of this system is shown in Figure 4. The software architecture of the network includes two units as follows.



**Figure 4.** System overview.

### 3.2.1. Sensing Unit

This unit includes an important number of nodes that act simultaneously as sensors and routers. In this system, end-devices are considered as reduced function devices (RFD) and are programmed to activate the sensor, read the physical measurement and transmit it toward the base station. On the other hand, the routers are considered as full function devices (FFD). They are programmed, in addition to the task performed by end-nodes, to receive data from other nodes and route these toward the control unit.

### 3.2.2. Control Unit

The control unit is composed of a coordinator node connected to the control system. The coordinator has to organize the network and its communications. It has also the capability to add or remove any node from the network, provide addresses and create/control the topology. The algorithms used by coordinator, like the topology construction algorithm, have to be sent by the control system to the coordinator.

## 4. MSCT Topology

In this section, we present our improved cluster-tree topology based on the minimum spanning tree. We call our topology the minimum spanning cluster-tree (MSCT). Basically, as discussed in the previous sections, the typical cluster tree is based on the association-request and association-response sent between the node and the coordinator; the coordinator then decides which position has to be attributed to the requesting sensor node. This position depends more on the timing of association-request; the network is created progressively with its expansion. In the MSCT topology, the process of network creation is totally different. We explain in this section how the topology is created in our system.

### 4.1. Indoor Propagation Models

Since our proposed algorithm for topology control is intended for ZigBee networks, the deployment environment is considered to be inside buildings. The indoor networks have some particularities to take into consideration, especially the presence of walls.

We take into consideration a model treating the placement of walls that dictates the signal path inside a building. Generally, in an indoor environment, the average path loss  $PL(d)$  for a transmitter and receiver over a distance  $d$  is given in [26] by:

$$PL(db) = PL(d_0) + 10n\log(d/d_0) + X_\sigma \quad (1)$$

where  $PL(d_0)$  is the path loss for 1 m,  $n$  the path loss exponent, and  $X_\sigma$  is a zero-mean Gaussian distributed random variable with standard deviation  $\sigma$ .

This model has been generalized in [27] to a multi-wall environment. This consists of adding a further attenuation term due to losses introduced by the walls penetrated by the direct path between the receiver and transmitter:

$$L(db) = PL(db) + M_\omega \quad (2)$$

where  $M_\omega$  represent the multi-wall component expressed as:

$$M_\omega = l_c + \sum_{i=1}^I k_{wi}l_i \quad (3)$$

where  $k_{wi}$  is the number of penetrated walls,  $l_c$  is constant and  $l_i$  depends on wall size as:  $l_1$  for a light wall (0–20 cm),  $l_2$  for a medium wall (20–40 cm) and  $l_d$  for a door.



In this paper, we consider that the cost of the transmission of a packet is based on the distance between transmitter and receiver. This cost can be considered as the weight of connection between the two nodes. In the case of free space transmission, the weight would simply be the Euclidean distance between the nodes. However, in the indoor environment, the presence of obstacles and path loss must be taken into consideration. Thus, based on the model above, we conclude this equation that calculates the weight of the connection between two nodes in indoor conditions:

$$\text{Weight} = D_E * (1 - L(db) \times \frac{P_r}{P}) \quad (4)$$

where  $D_E$  represents the Euclidean distance between the transmitter and the receiver and  $P_r$  and  $P$  represent respectively the received power at the destination and the power from the source.

#### 4.2. MSCT Creation

We suppose a graph  $G = (V, E)$  representing the topology of the static sensor network, where  $V$  represents the sensor nodes and  $E$  the link weight between two connected nodes. In theory, the link weight in the sensor network represents the Euclidean distance between two nodes. This distance is not representative in the case of an indoor environment. Therefore, our algorithm has to take into consideration some obstacles for transmission, mainly the walls, and does not consider other kinds of obstacles. Thus, in Equation (4), we presented a formula to compute the weight of connection from the Euclidean distance that considers the path loss. To construct the MST, we used Algorithm 1 based on Kruskal to sort the weight of edges in an increasing order. We then considered each sorted edge for inclusion in the tree if it did not create a cycle. Following this algorithm, we obtain a topology with the minimum cost of communication.

---

#### Algorithm 1 Minimum Spanning Tree Construction Algorithm

---

**Require:** Graph  $(G, w)$  with  $N$  nodes

**Ensure:** Minimum Spanning Tree  $T$

```

1:  $T \leftarrow \emptyset$ 
2: for all Vertex do
3:   Make-set( $V$ )
4: end for
5: Sort the edges of the graph in increasing order by weight  $W$ 
6: for  $i = 1$  to  $N$  do
7:   Select the next smallest weight edge
8:   if No cycle is created then
9:      $T \leftarrow \text{edge}(i)$ 
10:  end if
11: end for
12: return  $T$ 

```

---

Figures 5 and 6 show an illustration of the minimum spanning tree topology construction and how costs are reduced. Let us consider the topology (graph) in Figure 5 with eight nodes (vertices) that can be connected by several edges. The numbers on the edges represent the weight of connection that may be the distance between nodes and is considered as the cost of transmission in our study. Algorithm 1 sorts the weights of edges in an increasing manner and starts adding them to the tree. The edges that generate cycles will be excluded from the tree.

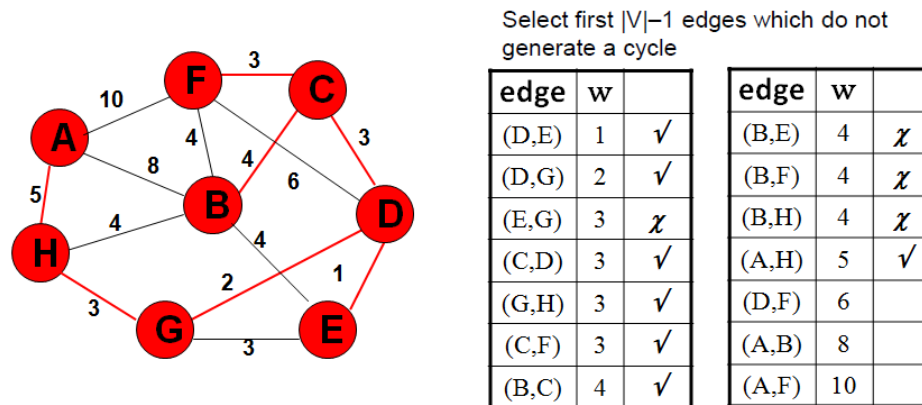


Figure 5. MST construction illustration: Part 1.

At the end of this process, the topology obtained is a spanning tree where the total weight is minimal, as illustrated in Figure 6. As the weight represents the transmission cost, mainly related to the energy dissipation, we can suppose that building such a topology may reduce considerably the consumption of energy in the network.

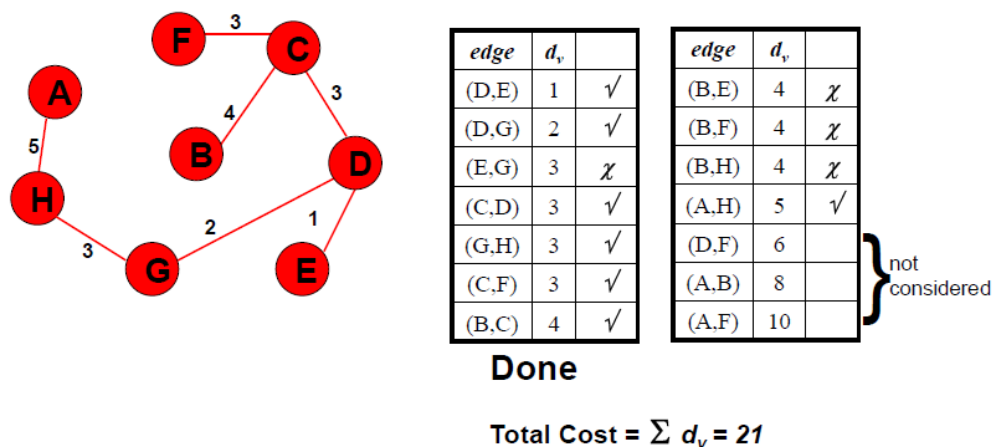


Figure 6. MST construction illustration: Part 2.

### 4.3. MSCT Topology Control

In this section, we explain how the MSCT topology based on the algorithm above is implemented in a real ZigBee sensor network. When the network is created, the only node to be connected is the coordinator. In order to add nodes, the network proceeds as follows:

- When a node sends an association-request, the coordinator adds the node to the network (assigning the PAN id), but not to the topology (no address assigned yet). Thus, during a time  $\tau$ , the coordinator does not create the topology, but only receives association-requests from nodes to store them in a table.
- All nodes added in the network calculate the weight of connections with their neighbors and send this information to the coordinator.
- After  $\tau$  time has passed, the coordinator constructs a map containing all nodes that joined the network, as well as their neighborhood tables. Thus, it creates the topology following Algorithm 1 steps.

- After the construction of the topology, the coordinator defines a period of time  $\tau'$  where the association-requests are received and stored in order to face any change of topology. This process allows the saving of energy by reducing the computational cost.

This control process allows the reduction of the energy consumption in the network because of:

- The algorithms for the topology construction and control are executed at the coordinator node. The energy of the sensor nodes is saved for sensing, transmitting and routing tasks.
- The total cost of transmissions in the network is minimal.
- Topology maintenance is executed periodically, contrary to cluster-tree or mesh, where it is calculated after any change in topology.

To better illustrate the topology control process, Algorithm 2 shows the operations performed by the coordinator. To simplify this algorithm, we assume that all association-requests are accepted by the coordinator.

---

#### Algorithm 2 Network Operations at the Coordinator

---

```

1: Coordinator joins the network
2: while  $t < \tau$  do
3:   Coordinator receives requests
4: end while
5: Store nodes positions and links in Nodes Table
6: Use Algorithm 1 to construct MSCT
7: while  $t - \tau < \tau'$  do
8:   Coordinator receives requests and data
9:   if Data Received then
10:    Send data to user's computer
11:   else if Request Received then
12:    Store request in Standby Table
13:   else if A node leaves the topology then
14:    Remove node from Nodes Table
15:    Exit While
16:   end if
17:   Update  $\tau'$  value
18: end while
19: if Standby Table contains nodes then
20:   Join Standby Table to Nodes Table
21:   Go to line 6
22: end if
23: if Node removed from topology then
24:   Restart algorithm
25: end if

```

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## 5. Performance Evaluation

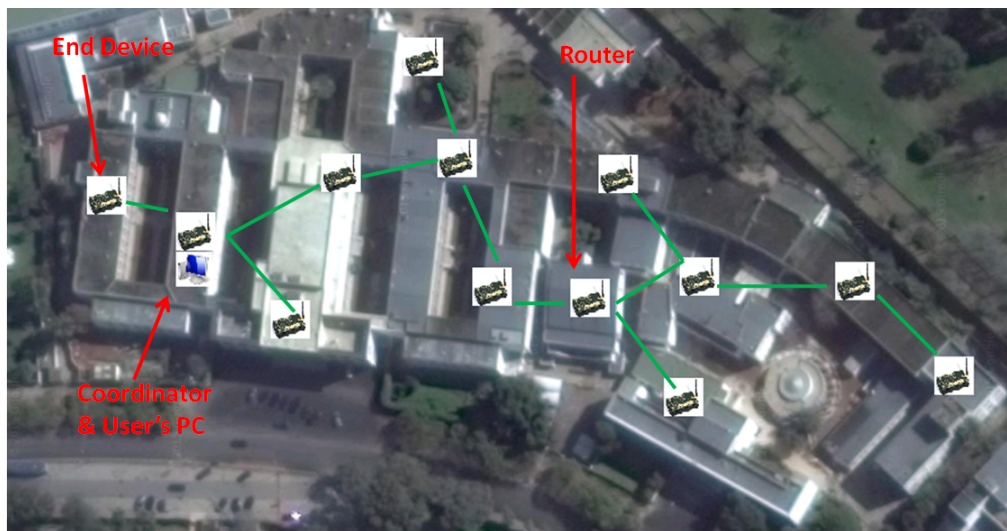
In this section, we present the experimental tests we performed to evaluate the performances of the system based on our designed sensor nodes.

### 5.1. Experiments' Scenario and Parameters

To evaluate the performances of our system in the same conditions as most of the ZigBee systems, we conducted the tests in an indoor environment. In these experiments, the network is composed of 13 ZigBee sensor nodes. We used only FFD devices that can act as routers or end-devices. This is due to the fact that the topology can change at any time, and end-devices may become routers in a new topology. The scenario is to deploy our network inside the building at the University of Rabat.

The system is programmed in such a way that all sensor nodes measure data every 2 s and transmit these toward the coordinator.

The end-device nodes sense and transmit data to the parent node; meanwhile, the routers can receive data from neighbors and transmit these toward the coordinator using the links defined in the topology. Figure 7 shows the deployment of nodes over the building and also shows the links between nodes in one of the topologies constructed during this experiment.



**Figure 7.** Sensor network deployment and topology over the building of the university.

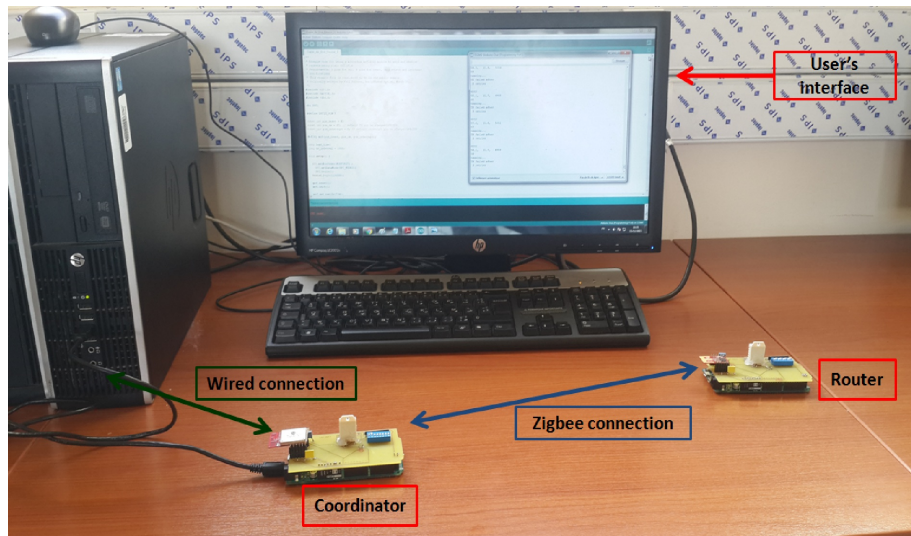
The detailed parameters of the experiments and their corresponding values are summarized in Table 3.

**Table 3.** Experimental parameters. FFD, full function devices; RFD, reduced function devices.

Experiment Parameters	Values
Deployment distribution	Unique
Deployment area	200 m × 40 m
Number of FFDs	13
Number of RFDs	0
Number of nodes	13
Initial energy	1 (J)
Number of coordinators	1
Maximum indoor distance between adjacent nodes	24 m
Sensing interval	2 s
Packets size	127 bits
Time $\tau$	2 min
Time $\tau'$	4 min

The coordinator sends received data from the network to a computer that represents the control system. Figure 8 shows the coordinator area. The connection between coordinator and the control system computer is wired. This can be explained by two reasons:

- (i) The wireless connection will consume much energy because all of the network’s data will be sent through this connection.
- (ii) To ensure that collected data coming from nodes arrive safely at the control system.



**Figure 8.** Architecture of the coordinator area.

### 5.2. Results and Discussion

Concerning the main results, we focus on the energy consumption of sensor nodes to test the performances of the system on it from two aspects: the hardware and the topology.

In order to evaluate the performances of our designed hardware, we compare its performances with a famous commercial sensor node: the MicaZ sensor node of Memsic (presented in Tables 1 and 2).

The MicaZ sensor node (Figure 9) is a sensor node based on a Texas Instruments microcontroller and respects ZigBee specifications. We choose to compare the performances of our hardware with MicaZ nodes because of the nearly similar features of both nodes in term of frequency, communication range and transmitting power (see Tables 1 and 2).



**Figure 9.** Memsic MicaZ sensor node.

The first result regards the discharge process. We performed the test of this metric using our designed sensor node. As illustrated in Figure 10, we compare the different topologies in terms of the battery's voltage discharge in router nodes according to time. The initial voltage for all router nodes is approximately 2 V. It is clear that the discharge process using MSCT is longer than the mesh and cluster-tree topologies. We notice also that this process is accentuated for the mesh topology when

the voltage level is below 1.5 V. In this case, the node is considered out of resources. Note that this value is lower in the other topologies: 1, 3 V for the cluster-tree and 0, 9 V for the MSCT. This first result shows that the energy consumption is minimal and the discharge process is slower by using the MSCT topology.

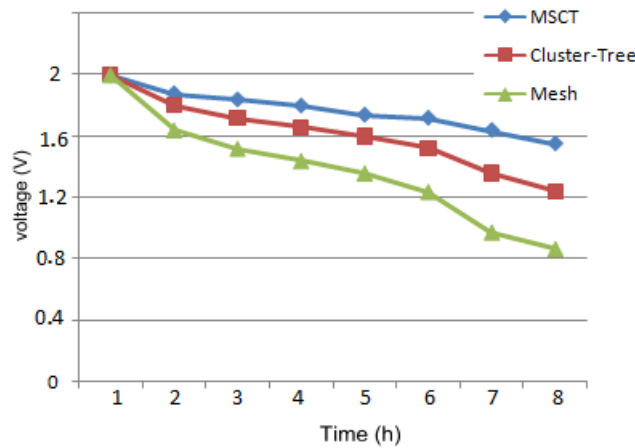


Figure 10. Average voltage of the designed sensor node batteries according to time and topology.

Figure 11 presents the second result where we compare the average energy consumptions for the two nodes. This metric includes the average consumption of all components in the Tx (transmitting) and Rx (receiving) modes, and values are measured in milliamperes (mA). We plot this result for all types of devices; coordinator, routers and end-devices. Logically, the coordinator and the routers consume more energy than end-devices. We can deduce that the MSCT topology shows better results in terms of energy consumption compared with the other typical topologies of ZigBee. Note that in Figure 11, the acronyms “C-T” and “M” represent respectively cluster-tree and mesh topology.

This result is very important since it shows also how energy consumption is reduced by our hardware device compared with the Memsic nodes. This can be explained by the fact that the devices used in our prototype have low power consumption. Furthermore, the size of packets has been reduced at the physical layer, which has positively impacted the computational costs, as well as the energy consumption of the nodes.

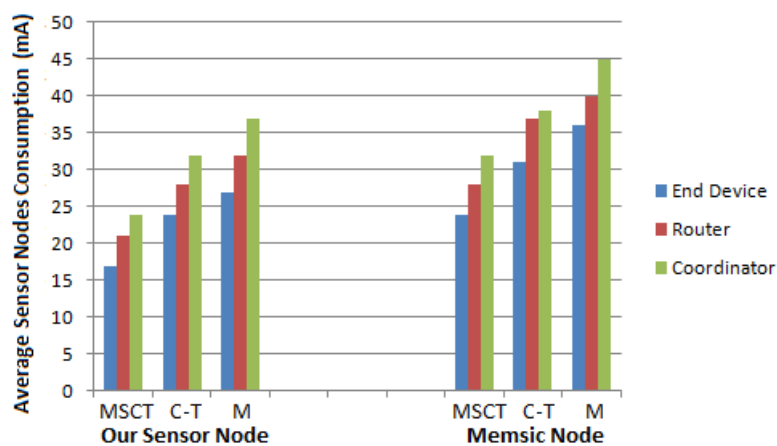


Figure 11. Average of the power consumption of nodes according to hardware and topology. MSCT, minimum spanning cluster-tree; C-T, cluster-tree; M, mesh.



The next result shown in Figure 12 concerns the network’s lifetime, considered as the main result. We represent the lifetime with two metrics. The first one is the time until the first node dies (FND). The FND time is considered as a stability period for the network since there is a node that leaves the network during it. The second one is the total network’s lifetime that represents the time until there is no sensor node in the transmission range of the coordinator; this time is called the network’s lifetime (NL).

We can clearly deduce from this result that our designed sensor node, enhanced by the MSCT topology, considerably improve the network’s lifetime compared with Memsic nodes by reducing the energy consumption.

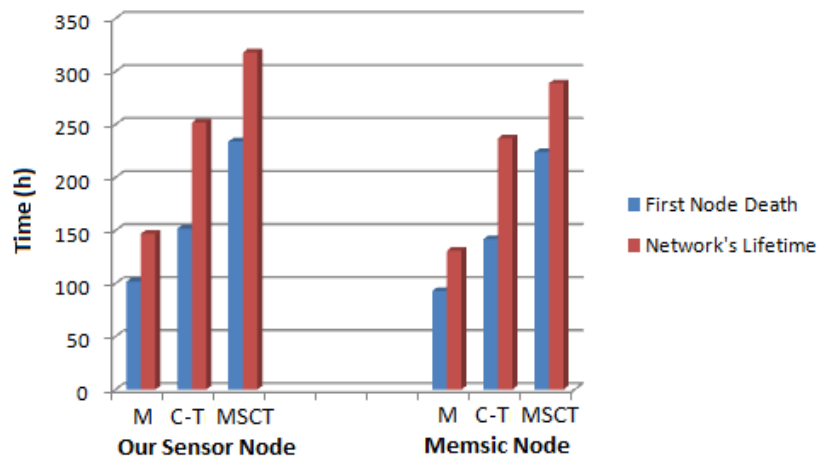


Figure 12. First node dead and the network’s lifetime according to hardware and topology.

Finally, in order to evaluate the communication performances of the hardware, we measure the delivery of packets to the coordinator. We illustrate in Figure 13 the packet loss ratio according to hardware and topology.

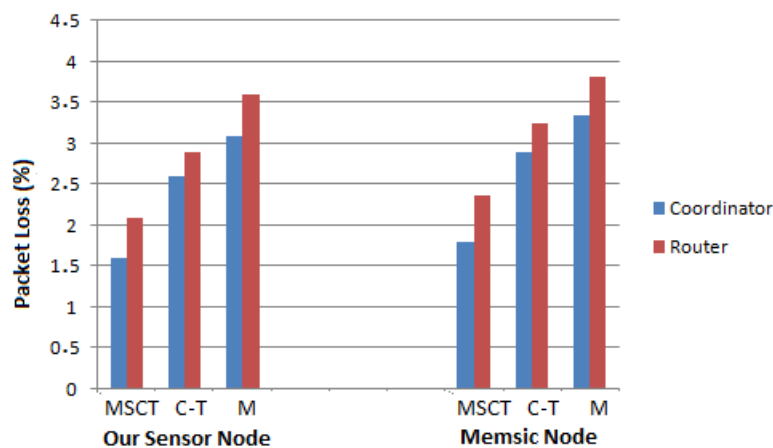


Figure 13. Packet loss ratio according to hardware and topology.

We first deduce that the packet loss ratio is reduced by our designed sensor nodes compared with the Memsic nodes. This is explained by the efficiency of data processing offered by the Arduino platform and by the reduction of the packet size, which leads to the reduction of the dropping of packets in the router nodes.



We also note from Figure 13 that the MSCT topology reduces the packet loss ratio; this can be explained mainly by the minimal cost of transmissions and also because most of the parent nodes have a maximum of two children.

In summary, we deduce from the above experiments that the developed system significantly improves the energy consumption and the network's lifetime. The combination of effective hardware architecture and MSCT topology has led to the minimization of the communications cost and the increase of the network's performances.

## 6. Conclusions

Wireless sensor networks have been widely exploited in the last few decades in many application fields due to their promising features. However, most of the sensor nodes are equipped with limited batteries, which makes power a serious issue. Thus, energy saving and optimization becomes one of the major concerns in the wireless sensor networks' hardware design and software programming.

In this paper, we developed the energy issue of an indoor sensing system based on the ZigBee network and presented as a solution our designed sensor nodes based on Arduino microcontrollers and ZigBee radio transceivers. The device showed good performances in terms of packet delivery, while maintaining a low power consumption. These results were enhanced by an energy-efficient cluster-tree topology called MSCT that reduced the transmissions cost of nodes by computing minimum spanning trees taking into consideration indoor conditions. The combination of the designed hardware and the developed topology showed very satisfactory results in terms of energy consumption and the network's lifetime.

**Author Contributions:** In this work, Yassin El Hillali proposed the hardware architecture for the sensor node, and participated to the implementation of the physical and MAC layers in the sensor nodes. Atika Menhaj-Rivenq supervised the team work at Valenciennes and participated in the design of the hardware prototype. Mourad Ouadou integrated the circuit components to build the sensor node prototype, implemented the first network layers, and developed the proposed algorithms and implemented them into the system. Ouadoudi Zytoune participated in the implementation of the algorithm. Finally, Driss Aboutajdine supervised the working team in Rabat and validated results.

**Conflicts of Interest:** The authors declare no conflict of interest.

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