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► To cite this version:

Marina Groshaus, Leandro Montero. Structural properties of biclique graphs and the distance formula. The Australasian Journal of Combinatorics, 2021, 81 (2), pp.301-318. hal-03338962

HAL Id: hal-03338962

<https://hal.science/hal-03338962>

Submitted on 9 Sep 2021

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Structural properties of biclique graphs and the distance formula

MARINA GROSHAUS*

UTFPR-DAINF

Brazil

`marinagroshaus@yahoo.es`

LEANDRO MONTERO

KLaIM team, L@bisen, AIDE Lab., Yncrea Ouest

33 Q, Chemin du Champ de Manœuvres, 44470 Carquefou

France

`lpmontero@gmail.com`

Abstract

A biclique is a maximal induced complete bipartite subgraph of G . The biclique graph of a graph G , denoted by $KB(G)$, is the intersection graph of the family of all bicliques of G . In this work we study some structural properties of biclique graphs which are necessary conditions for a graph to be a biclique graph. In particular, we prove that for biclique graphs that are neither a K_3 nor a diamond, the number of vertices of degree 2 is less than half the number of vertices in the graph. Also, we present forbidden structures. For this, we introduce a natural definition of the distance between bicliques in a graph. We give a formula that relates the distance between bicliques in a graph G and the distance between their respective vertices in $KB(G)$. Using these results, we can prove not only this new necessary condition involving the degree, but also that some graphs are not biclique graphs. For example, we show that the crown is the smallest graph that is not a biclique graph, although the known necessary condition for biclique graphs holds, answering an open problem about biclique graphs. Finally, we present some interesting related conjectures and open problems.

* Partially supported by CNPq (428941/2016-8). CONICET / Universidad de Buenos Aires, Argentina.

1 Introduction

Intersection graphs of certain special subgraphs of a general graph have been studied extensively. We can mention line graphs (intersection graphs of the edges of a graph), interval graphs (intersection graphs of subpaths of a path), and in particular, clique graphs (intersection graphs of the cliques of a graph) [5, 6, 10, 11, 12, 21, 23].

The *clique graph* of G , denoted by $K(G)$, is the intersection graph of the family of all cliques of G . Clique graphs were introduced by Hamelink in [16] and characterized in [31]. It was proved in [1] that the clique graph recognition problem is NP-Complete.

Bicliques have been studied in many contexts. Depending on the context and the author, bicliques are defined in different ways: induced or not induced subgraphs, maximal or not, etc. [2, 9, 19, 24, 26, 30, 32]. All of them are rather natural and clearly justified. In our work, we consider bicliques as being maximal induced complete bipartite subgraphs.

Bicliques have applications in various fields, for example, biology: protein-protein interaction networks [7], social networks: web community discovery [20], genetics [3], medicine [25], information theory [15]. More applications (including some of these) can be found in [22].

The *biclique graph* of a graph G , denoted by $KB(G)$, is the intersection graph of the family of all bicliques of G . It was defined and characterized in [14]. However, no polynomial time algorithm is known for recognizing biclique graphs. Biclique graphs have been studied over the last few years. See [8, 13, 28] for some examples of recent articles on the subject.

In this work we study structural properties of biclique graphs. For this, we introduce the concept of the distance between bicliques in a graph. Previous related work in the context of cliques can be found in [4, 17, 18, 27, 29].

In [14], a necessary condition for a graph to be a biclique graph was given. It was an open problem whether this condition was sufficient. Using the distances formula, we give a different proof for this necessary condition¹. Moreover, we prove that this necessary condition is not sufficient, that is, we present some structural properties and forbidden structures which allow us to identify graphs that verify the condition but they are not biclique graphs. Finally, we present a new necessary condition of biclique graphs: given a biclique graph that is not a K_3 or a diamond, we prove that the number of vertices of degree 2 is strictly less than half the number of vertices in the graph. Consequently, we give some forbidden structures.

We hope that these tools give some light to the main open problem in the context of biclique graphs, that is, the recognition of this class. In the appendix, we present a complete list of all biclique graphs up to 6 vertices.

This work is organized as follows. In Section 2 the notation is given. Section 3

¹We remark that, although the original proof is short, the new proof shows how to use this formula to solve problems in biclique graphs.

contains some preliminary and general properties. In Section 4 we present the relation between distances in graphs and biclique graphs plus some structural properties of biclique graphs. Section 5 contains a bound on the number of vertices of degree 2 for biclique graphs and forbidden structures. In the last two sections we present some open and interesting related problems, and we end with a concluding one.

2 Preliminaries

Throughout the paper we restrict our attention to undirected simple graphs. Let $G = (V, E)$ be a graph with vertex set $V(G)$ and edge set $E(G)$, and let $n = |V(G)|$ and $m = |E(G)|$. A *subgraph* G' of G is a graph $G' = (V', E')$, where $V' \subseteq V$ and $E' \subseteq E$ such that all endpoints of the edges of E' are in V' . When E' has all the edges of E whose endpoints belong to the vertex subset V' , we say that V' *induces* the subgraph $G' = (V', E')$, that is, G' is an *induced subgraph* of G . A graph $G = (V, E)$ is *bipartite* when there exist sets U and W such that $V = U \cup W$, $U \cap W = \emptyset$, $U \neq \emptyset$, $W \neq \emptyset$ and $E \subseteq U \times W$. Say that G is a *complete graph* when every possible edge belongs to E . A complete graph on n vertices is denoted K_n . A bipartite graph is *complete bipartite* when every vertex of the first set is connected to every vertex of the second set. A complete bipartite graph on p vertices in one set and q vertices in the other is denoted $K_{p,q}$. A *clique* of G is a maximal complete induced subgraph, while a *biclique* is a maximal induced complete bipartite subgraph of G . The *open neighborhood* of a vertex $v \in V(G)$, denoted $N(v)$, is the set of vertices adjacent to v . The *closed neighborhood* of a vertex $v \in V(G)$, denoted $N[v]$, is the set $N(v) \cup \{v\}$. The *degree* of a vertex v , denoted by $d(v)$, is defined as $d(v) = |N(v)|$. A vertex $v \in V(G)$ is *simplicial* if $N[v]$ is a clique. A vertex $v \in V(G)$ is *universal* if it is adjacent to all of the other vertices in $V(G)$. A *path* of k vertices, denoted by P_k , is a sequence of vertices $v_1 v_2 \dots v_k \in V(G)$ such that $v_i \neq v_j$ for all $1 \leq i \neq j \leq k$ and v_i is adjacent to v_{i+1} for all $1 \leq i \leq k-1$. A graph is *connected* if there exists a path between each pair of vertices. The distance between two vertices $v, w \in G$ is defined as the number of edges in a shortest path between them and is denoted by $d_G(v, w)$. Whenever no confusion arises, we will simply write $d(v, w)$ instead of $d_G(v, w)$. We assume that all the graphs of this paper are connected.

A *diamond* is a complete graph with 4 vertices minus an edge. A *gem* is an induced path with 4 vertices plus a universal vertex.

Given a family of sets \mathcal{H} , the *intersection graph* of \mathcal{H} is a graph that has the members of \mathcal{H} as vertices and there is an edge between two sets $E, F \in \mathcal{H}$ when $E \cap F \neq \emptyset$. A graph G is an *intersection graph* if there exists a family of sets \mathcal{H} such that G is the intersection graph of \mathcal{H} .

3 General properties

In this section we present some properties of the biclique graph related to connectivity. First we recall the theorem in [14] that gives a necessary condition for a graph

to be a biclique graph.

Theorem 3.1 ([14]). *Let G be a graph such that $G = KB(H)$, for some graph H . Then every induced P_3 of G is contained in an induced diamond or an induced gem of G as shown in Figure 1.*

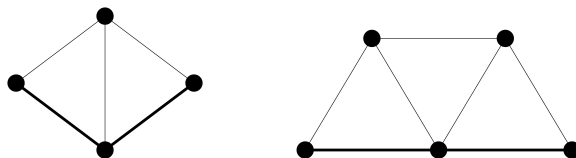


Figure 1: Induced P_3 in bold edges contained in a diamond and in a gem respectively.

In Section 4 we give a different proof of Theorem 3.1. One question that arises from Theorem 3.1 is: Given a graph G such that every induced P_3 is contained in a diamond or in a gem; is $G = KB(H)$ for some graph H ? In Section 4 we show that the answer is “No”, by proving a result that allows us to construct graphs that have every induced P_3 in a diamond or in a gem although they are not biclique graphs.

Next we show the connectivity relation between G and $KB(G)$.

Proposition 3.2. *Let G be a graph. G is connected if and only if $KB(G)$ is connected.*

Proof. Suppose G is connected. Let B and B' be bicliques of G . If B intersects B' then their corresponding vertices in $KB(G)$ are adjacent. If they do not intersect, as G is connected, there is a path between each vertex of B and B' . Let $b \in B$, $b' \in B'$ such that $d(b, b') = \min\{d(v, w) \mid v \in B, w \in B'\}$. Let $k = d(b, b')$. Clearly $k > 0$, so take a path $P = bv_1 \dots v_{k-1}b'$ of length k between b and b' . Now, each triple of consecutive vertices of P is contained in a different biclique since both endpoints of each triple are not adjacent. Finally taking the bicliques that contain the following triples $\{b, v_1, v_2\}, \{v_2, v_3, v_4\}, \dots, \{v_{k-2}, v_{k-1}, b'\}$ for k even and $\{b, v_1, v_2\}, \{v_2, v_3, v_4\}, \dots, \{v_{k-3}, v_{k-2}, v_{k-1}\}, \{v_{k-2}, v_{k-1}, b'\}$ for k odd (note that for $k = 1$, the edge bb' is in a biclique that intersects both B and B'), we have that each biclique only intersects with the previous and the following one. Therefore, their corresponding vertices in $KB(G)$ form a path between the vertices corresponding to B and B' . Hence $KB(G)$ is connected.

The converse is clear. □

The following result is a direct consequence of Theorem 3.1.

Lemma 3.3. *Let G be a connected graph such that $G = KB(H)$, for some graph H . If G has at least 3 vertices, then $d(v) \geq 2$ for all $v \in V(G)$. Moreover, G is 2-connected.*

Proof. If there is a vertex with degree 1, since G has at least 3 vertices, that vertex would be an extreme of a P_3 not contained in an induced diamond nor an induced gem, that is, a contradiction by Theorem 3.1. Finally, suppose G is not 2-connected and let v be a vertex such that $G - \{v\}$ is disconnected. Taking two vertices v_1, v_2 in different connected components of $G - \{v\}$ that are adjacent to v , we obtain a P_3 not contained in an induced diamond nor an induced gem, and again, a contradiction by Theorem 3.1. \square

4 Distances in G and $KB(G)$

In this section we define the distance between bicliques in a graph. Also, we study the relation between the distance of bicliques in a graph G and the distance between their respective vertices in $KB(G)$.

Definition 4.1. Let G be a graph and let B, B' be bicliques of G . We define the distance between B and B' as $d(B, B') = \min\{d(b, b') \mid b \in B, b' \in B'\}$.

The next formula states the relationship between the distances of G and $KB(G)$. This result is useful to show that the condition of Theorem 3.1 is not sufficient.

Lemma 4.2. Let G be a graph and let B, B' be two different bicliques of G . Then $d_{KB(G)}(B, B') = \lfloor \frac{d_G(B, B') + 1}{2} \rfloor + 1$.

Proof. Let v_0, v_k be vertices of G such that $v_0 \in B$, $v_k \in B'$ and $d(v_0, v_k) = d_G(B, B') = k$. If $k = 0$ then B and B' intersect in G . So they are adjacent as vertices in $KB(G)$. Therefore $d_{KB(G)}(B, B') = \lfloor \frac{0+1}{2} \rfloor + 1 = 1$. Suppose now that $k > 0$. Let $P_1 = v_0 v_1 \dots v_k$ be a path in G between B and B' of length k . Take $B_i \in V(KB(G))$ such that $\{v_i, v_{i+1}, v_{i+2}\} \subseteq B_i$ in G , $i = 0, \dots, k-2$. These bicliques B_i of G exist since $v_i v_{i+2} \notin E(G)$ for $i = 0, \dots, k-2$, otherwise there would be a path of length less than k between B and B' .

Then $BB_0 B_2 B_4 \dots B_{2j} \dots B'$ is a path in $KB(G)$ between B and B' of length $\lfloor \frac{k+1}{2} \rfloor + 1$ and therefore, as $d_G(B, B') = k$, we have that $d_{KB(G)}(B, B') \leq \lfloor \frac{d_G(B, B') + 1}{2} \rfloor + 1$. This situation can be observed in Figure 2.

Now let $P_2 = B_0 B_1 \dots B_s$ be a path of minimum length in $KB(G)$ between $B = B_0$ and $B' = B_s$ (Fig. 3). Then $d_{KB(G)}(B, B') = s > 0$. Let $v_{2i} \in B_i \cap B_{i+1}$ be vertices of $V(G)$ for $i = 0, \dots, s-1$. Thus, we obtain the vertices $v_0, v_2, \dots, v_{2s-2}$ of G . Now, for $i = 1, \dots, s-1$, either v_{2i-2} is adjacent to v_{2i} or v_{2i-2} is not adjacent to v_{2i} . If they are not adjacent then there exists one vertex v_{2i-1} adjacent to both since v_{2i-2} and v_{2i} belong to the biclique B_i of G . Then, the longest path between v_0 and v_{2s-2} occurs when these two consecutive vertices are not adjacent. In this situation, adding the vertex adjacent to both between each pair, we have an induced path $v_0 v_1 \dots v_{2s-2}$ in G between B and B' of length $2s - 2$. Then

$$d_G(B, B') \leq 2s - 2,$$

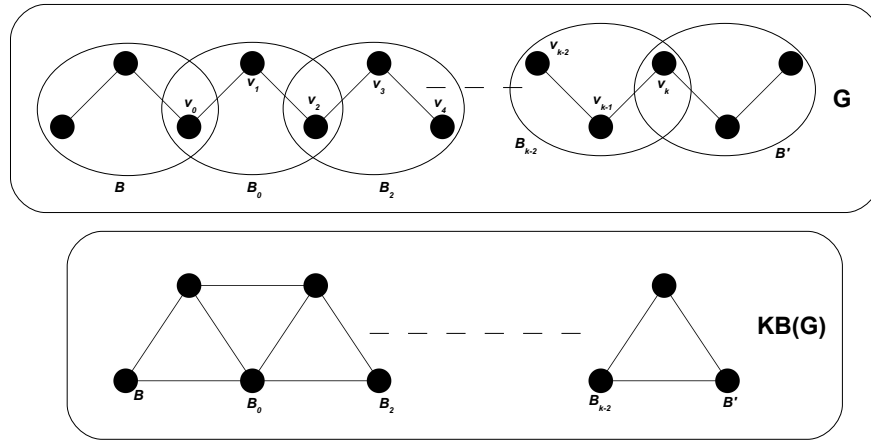


Figure 2: First inequality.

$$\frac{d_G(B, B') + 2}{2} \leq s = d_{KB(G)}(B, B'),$$

$$\frac{d_G(B, B')}{2} + 1 \leq d_{KB(G)}(B, B').$$

Finally, since $d_G(B, B')$ and $d_{KB(G)}(B, B')$ are integers, it follows that

$$\left\lfloor \frac{d_G(B, B') + 1}{2} \right\rfloor + 1 \leq d_{KB(G)}(B, B').$$

Combining both inequalities we obtain the desired result. \square

Now, based on the distance between two bicliques of a graph G , we can ensure the existence of other bicliques “between them”. That is, if the distance between the bicliques B and B' of G is k , then there exist other bicliques at distance at most $k - 1$ to each of B and B' . This result will be very useful for proving not only Theorem 3.1 but also that the condition of the theorem is not sufficient.

Theorem 4.3. *Let G be a graph and let B, B' be bicliques of G such that $d_G(B, B') = k > 0$. Then there exist at least $k + 1$ bicliques in G such that they are at distance at most $k - 1$ from both B and B' .*

Proof. We will prove the theorem in two parts, first when $d_G(B, B') = 1$ and last when $d_G(B, B') = k > 1$.

Suppose first that $d_G(B, B') = 1$. Let $v \in B, w \in B'$ be adjacent vertices and let $x \in B, y \in B'$ be vertices such that $vx, wy \in E(G)$. In Figure 4 we show all possible

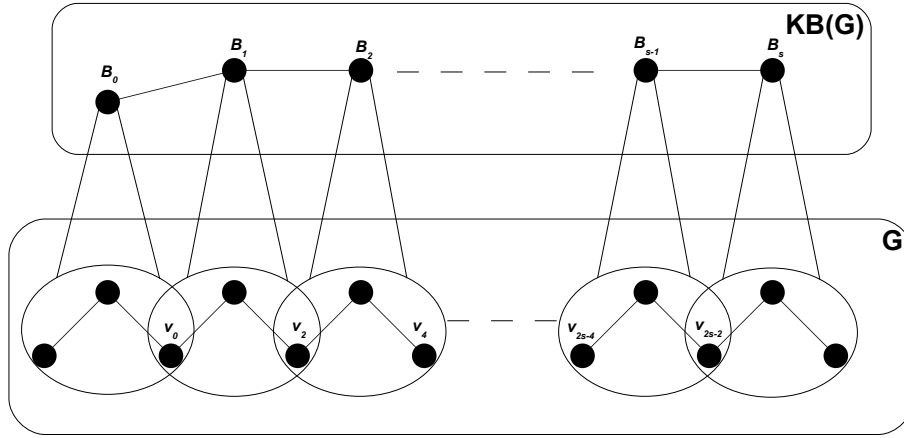
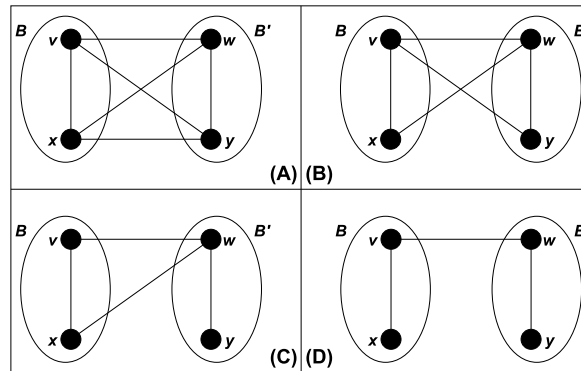


Figure 3: Second inequality.

configurations (up to symmetry) according to the adjacencies between vertices. We count the number of different bicliques in each case.

Figure 4: Different cases for bicliques B and B' at distance 1.

- (A) $\{x, y\}, \{v, y\}, \{v, w\}, \{x, w\}$ are contained in four different bicliques.
- (B) $\{x, v, y\}, \{x, w, y\}$ and $\{v, w\}$ are contained in three different bicliques.
- (C) $\{x, w, y\}$ and $\{v, w, y\}$ are contained in two different bicliques.
- (D) $\{x, v, w\}$ and $\{v, w, y\}$ are contained in two different bicliques.

Note that if $xy \in E(G)$ and $xw, vy \notin E(G)$, then $\{x, v, w, y\}$ is contained in a biclique. Now, as the biclique containing $\{v, w, x, y\}$ should be different from B , there exists a vertex $z \in B$ such that either $zv, zw \in E(G)$ (or $zx, zy \in E(G)$) or $zv \in E(G)$, $zy \notin E(G)$ (or $zx \in E(G)$, $zw \notin E(G)$). By symmetry, as the biclique containing $\{v, w, x, y\}$ is different from B' , there should also exist a vertex $u \in B'$ with similar adjacencies. In these situations, we can see that we fall in one of the

previous cases **(B)**, **(C)** or **(D)**, therefore taking into account the biclique containing $\{v, w, x, y\}$, we obtain at least three different bicliques.

In all cases there are at least two different bicliques that intersect B and B' , that is, they are at distance $k - 1 = 0$ to each of them. We remark that we exactly count in each case the minimum number of bicliques that can exist between B and B' as we will use that later in the paper.

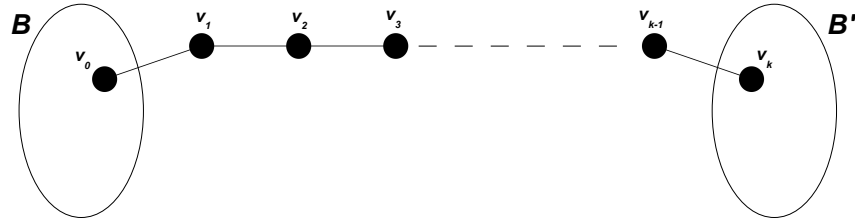


Figure 5: Bicliques B and B' are at distance k .

Suppose last that $d_G(B, B') = k > 1$ and let $v_0 v_1 \dots v_k$ be a shortest path between B and B' such that $v_0 \in B$, $v_k \in B'$. Then, each triple $\{v_i, v_{i+1}, v_{i+2}\}$ is contained in a different biclique of G for $i = 0, \dots, k-2$. Therefore we obtain $k-1$ bicliques that are at distance at most $k-1$ to each of B and B' . We obtain the two remaining bicliques as follows. Let $x \in B$, $y \in B'$ such that $xv_0, yv_k \in E(G)$. If $xv_1 \notin E(G)$ (respectively $yv_{k-1} \notin E(G)$) there is a biclique containing $\{x, v_0, v_1\}$ (respectively $\{y, v_k, v_{k-1}\}$). We call this biclique sharing an edge with B (respectively B') a *special biclique*. Otherwise, if $xv_1 \in E(G)$ (respectively $yv_{k-1} \in E(G)$), then $\{x, v_1, v_2\}$ (respectively $\{y, v_{k-1}, v_{k-2}\}$) is contained in a biclique. \square

As an immediate result of Theorem 4.3 for $k = 1$ and $k = 2$ we obtain the following two corollaries.

Corollary 4.4. *Let G be a graph, let B, B' be bicliques of G such that $d_G(B, B') = 1$, let e be an edge with one endpoint in B and the other in B' , and let B_2 be a biclique containing e . Then there exists a biclique B_1 which intersects B, B' and B_2 . Moreover, the vertices $b, b', b_1, b_2 \in KB(G)$ corresponding to the bicliques B, B', B_1, B_2 in G respectively, induce a diamond in $KB(G)$.*

Corollary 4.5. *Let G be a graph, and let B, B' be bicliques of G such that $d_G(B, B') = 2$. Consider an induced P_3 with its extremes in B and B' , and let B_2 be a biclique containing that P_3 . Then there exist two different bicliques B_1, B_3 “between” B and B' , such that, calling b, b', b_1, b_2, b_3 the vertices in $KB(G)$ corresponding to the bicliques B, B', B_1, B_2, B_3 in G respectively, we either have*

- both B_1 and B_3 are special bicliques, and thus $\{b, b', b_1, b_2, b_3\}$ induces a gem in $KB(G)$, or
- at least one of B_1 and B_3 is not special (suppose B_1 is not), and thus $\{b, b', b_1, b_2\}$ induces a diamond in $KB(G)$.

The proof of Theorem 3.1 follows from Theorem 4.3, and Corollaries 4.4 and 4.5.

Proof of Theorem 3.1. Let bb_2b' be an induced P_3 in G and let B , B_2 and B' be the bicliques of H associated to the vertices b , b_2 and b' of G . The biclique B_2 contains either an edge with one endpoint in B and the other in B' , or B_2 contains a P_3 with its extremes in B and B' . In the first case, by Corollary 4.4, we obtain that bb_2b' is contained in an induced diamond of G . In the second case, by Corollary 4.5, we obtain that bb_2b' is contained in either an induced diamond or an induced gem of G . \square

Now we will show that although in the *crown* (Fig. 6) every induced P_3 is contained in an induced diamond, it is not a biclique graph. This result is a counterexample of the question of the sufficiency of the property of Theorem 3.1.

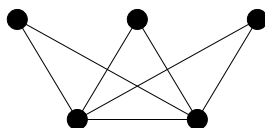


Figure 6: The *crown* is not a biclique graph but has every P_3 in a diamond.

Indeed, we will prove a more general result that implies not only that the crown graph is not a biclique graph but neither are many other graphs. We remark that the crown is the smallest graph that verifies the condition of Theorem 3.1 but is not a biclique graph.

Proposition 4.6. *Let $G = KB(H)$ for some graph H , where G is not isomorphic to the diamond. Then there do not exist $v_1, v_2 \in V(G)$ such that $N(v_1) = N(v_2)$ and their neighbors induce a K_2 .*

Proof. Suppose that there exist $v_1, v_2 \in V(G)$ such that $N(v_1) = N(v_2)$ and their neighbors induce a K_2 . Then $d_G(v_1, v_2) = 2$, and therefore, if B is the biclique of H that corresponds to v_1 and B' is the biclique of H that corresponds to v_2 , by Lemma 4.2, $d_H(B, B') = 2$ or $d_H(B, B') = 1$. We will analyze each case.

Case $d_H(B, B') = 2$.

Now, H must contain a subgraph as depicted in Figure 7. We will show that we arrive at a contradiction. Suppose first that one of both dotted edges does not exist, say vy . Then $\{v, x, y\}$ is contained in a biclique that does not intersect B' . This is a contradiction since $N(v_1) = N(v_2)$ in G . Suppose next that both dotted edges vy and vy' exist. In this case we also arrive at a contradiction since in H there are at least four bicliques that intersect with B and B' . We obtain one for each choice of one element in $\{x, y\}$ in B , one element in $\{x', y'\}$ in B' and v . Therefore $N(v_1) = N(v_2)$ does not induce a K_2 , which is a contradiction.

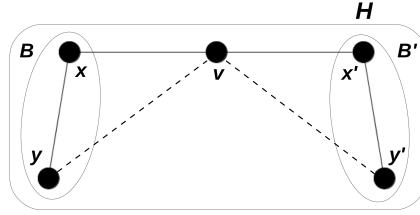


Figure 7: Graph H when $d_H(B, B') = 2$.

Case $d_H(B, B') = 1$.

In this case, by Theorem 4.3, there exist at least two bicliques B_1, B_2 in H that intersect both B and B' . Clearly there must be exactly two, or otherwise $N(v_1) = N(v_2)$ would not induce a K_2 . Now, by the first part of the proof of Theorem 4.3, only in the cases (C) and (D) is it possible that there are exactly two bicliques intersecting both B and B' in H . Figure 8 shows both possible options. The labels of Theorem 4.3 were as given in Figure 4.

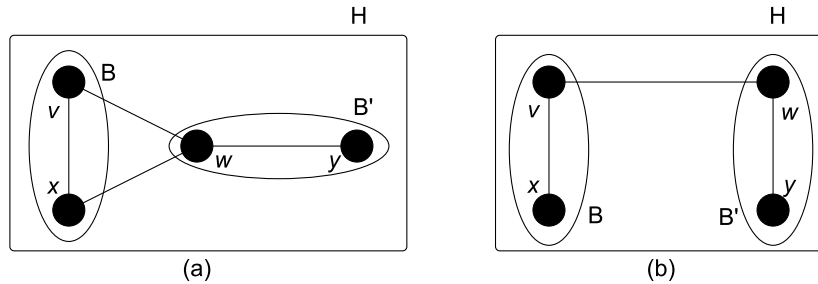


Figure 8: Only options for H with two bicliques that intersect B and B' .

– Case (a) (follows case (C) of the proof of Theorem 4.3):

As we can observe in Figure 8a, if B_1 is the biclique containing $\{v, w, y\}$ and B_2 is the biclique containing $\{x, w, y\}$, H has four bicliques such that they induce a diamond in G . As G is not isomorphic to the diamond, there must exist another biclique in H that intersects neither B nor B' . Let B_3 be another biclique such that $V(B_3) \subseteq V(H) \setminus (B \cup B')$. Suppose now that there is a vertex $u \in B_3$ such that $u \notin (B_1 \cup B_2)$. Since H is connected, we can choose B_3 such that there is a vertex $u' \in (B \cup B' \cup B_1 \cup B_2)$ adjacent to u . Clearly, if $u' \in (B \cup B')$, we obtain a contradiction because the edge uu' is contained in a biclique different to B_1, B_2 that intersects B or B' . Otherwise, since the edge $wy \in (B_1 \cap B_2)$ and $u' \in (B_1 \cup B_2)$, we have that u' is adjacent to w or y , and therefore we obtain that $\{u, u', w\}$ or $\{u, u', y\}$ is contained in another biclique intersecting B' respectively, which is a contradiction. We can conclude that $V(H) \setminus (B \cup B')$ is contained in $B_1 \cup B_2$.

We show now that there are at least three bicliques that intersect B or B' , i.e., v_1 or v_2 has an open neighborhood bigger than two vertices, which would be

a contradiction. For this, let ab be an edge in the biclique B_3 . Because $V(B_3)$ is contained in $B_1 \cup B_2$, both a and b are adjacent to at least one vertex in $\{w, y\}$.

If a is adjacent to both w and y , then the edges ab in B_3 and wy in B' determine case **(A)** or **(B)** of the proof of Theorem 4.3, which is a contradiction. Therefore a is adjacent to just one vertex in $\{w, y\}$. Similarly, the same can be said about b .

If a and b are adjacent to different vertices in $\{w, y\}$, then these four vertices induce a C_4 and hence there are three bicliques between B_3 and B' like in the proof of Theorem 4.3, a contradiction.

If a and b are both adjacent to y , consider the bicliques containing $\{w, y, a\}$ and $\{w, y, b\}$. They cannot be the same biclique, so one is B_1 and the other is B_2 . Suppose without loss of generality that the first biclique is B_1 and the second biclique is B_2 . Recalling that B_1 is the biclique containing $\{v, w, y\}$ and B_2 is the biclique containing $\{x, w, y\}$, one can conclude that a is adjacent to v and that b is adjacent to x . Then, if we consider the edge vx in B and the edge ab in B_3 , we have that $\{a, b, v, x\}$ either induces a C_4 , a diamond (case **(B)**) or a K_4 (case **(A)**). Thus there are at least three bicliques between B and B_3 , a contradiction.

Finally, if both a and b are adjacent to w , this time suppose without loss of generality that B_1 is the biclique containing $\{a, w\}$ and that B_2 is the biclique containing $\{b, w\}$. If v is not adjacent to b , then any biclique containing $\{v, w, b\}$ is different from B_1 and B_2 , so there are three bicliques that intersect B' , a contradiction. Consequently, v must be adjacent to b . Similarly, x is adjacent to a . Therefore when we consider the edges vx in B and ab in B_3 , we see that $\{a, b, v, x\}$ induces a C_4 , diamond or K_4 , so there are at least three bicliques between B and B_3 , a contradiction.

– Case **(b)** (follows case **(D)** of the proof of Theorem 4.3):

As we can observe in Figure 8b, if B_1 is the biclique containing $\{x, v, w\}$ and B_2 is the biclique containing $\{v, w, y\}$, similarly to the case **(a)**, H has four bicliques that induce a diamond in G . Then there must exist another biclique B_3 in H different from these four which does not intersect B or B' .

One can show that $V(H) \setminus (B \cup B')$ is contained in $B_1 \cup B_2$ with exactly the same argument as the previous case, using now the edge $vw \in (B_1 \cap B_2)$ rather than the edge $wy \in (B_1 \cap B_2)$ as we did in case **(a)**.

Now we show for this case that there are at least three bicliques that intersect B or B' . Let B_3 be a fifth biclique in H not intersecting B or B' and let ab be an edge in B_3 . Since $V(B_3)$ is contained in $B_1 \cup B_2$, both a and b are adjacent to at least one vertex in $\{v, w\}$.

If a is adjacent to both v and w , then the bicliques containing edges av and aw are different from B_1 and B_2 , which is a contradiction. Therefore a is adjacent to just one vertex in $\{v, w\}$ and the same can be said about b .

If a and b are adjacent to different vertices in $\{v, w\}$, suppose without loss of generality that a is adjacent to v and b is adjacent to w ; then these four vertices induce a C_4 that is contained in a biclique \tilde{B} . If \tilde{B} is different from B_1 and B_2 , we obtain a contradiction. Otherwise \tilde{B} should be B_1 or B_2 . Without loss of generality,

suppose that $\tilde{B} = B_1$; therefore x must be adjacent to b . Considering the edge ab in B_3 and xv in B , $\{a, b, x, v\}$ induces a C_4 and hence there are three bicliques between B_3 and B as in the proof of Theorem 4.3, a contradiction.

Finally, if a and b are both adjacent to v (symmetric if both are adjacent to w), we have that either $\{x, v, a\}$ and $\{x, v, b\}$ are contained in bicliques different from B_2 which is a contradiction, or $\{a, b, x, v\}$ induces a diamond or a K_4 and therefore there are at least three bicliques between B and B_3 which is also a contradiction.

As no more cases are left, there do not exist $v_1, v_2 \in V(G)$ such that $N(v_1) = N(v_2)$ with their neighbors inducing a K_2 ; this completes the proof. \square

Figure 9 shows some examples of graphs that are not biclique graphs, where every P_3 is included in a diamond.

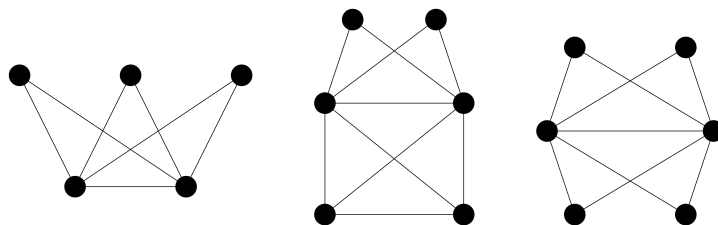


Figure 9: Graphs that are not biclique graphs by Proposition 4.6.

5 Vertices of degree 2 in biclique graphs

In this section we give a strong property for biclique graphs that have an induced P_3 contained in a gem and not in a diamond. Also, we show some forbidden structures. These properties give more tools to recognize graphs that are not biclique graphs.

The next result implies that the *Hajós graph*, the *rising sun* and the X_1 graph (see Fig. 10) are not biclique graphs by giving a forbidden structural property.

Proposition 5.1. *Let $G = KB(H)$ for some graph H , and let bb_2b' be an induced P_3 such that b, b' do not belong to any induced diamond. Let b_1, b_3 be the vertices of G corresponding to special bicliques in H such that $\{b, b', b_1, b_2, b_3\}$ induces a gem (Corollary 4.5). If \tilde{b} is a vertex such that $\tilde{b} \notin \{b, b', b_1, b_2, b_3\}$ and \tilde{b} does not belong to an induced diamond with b , then \tilde{b} is not adjacent to b_1 .*

Proof. Recall that since b, b' do not belong to any induced diamond and are at distance 2, by Corollary 4.5, such vertices b_1, b_3 exist. Let $B_1, B', B_2, B_3, \tilde{B}$ be the bicliques of H corresponding to the vertices $b, b', b_1, b_2, b_3, \tilde{b}$ of G . Let xy be an edge that belongs to $B \cap B_1$. Note that this edge exists since B_1 is a special biclique. By contradiction, if \tilde{B} and B_1 intersect (i.e., \tilde{b} is adjacent to b_1 in G), then either

x is adjacent to some vertex z of \tilde{B} , or the vertex y is adjacent to some vertex z of \tilde{B} . In either case, there is an edge between a vertex of B and a vertex of \tilde{B} , which is a contradiction by Corollary 4.4, as \tilde{b} and b would belong to an induced diamond in G . \square

Corollary 5.2. *The Hajós graph, the rising sun and the X_1 graph are not biclique graphs (Fig. 10).*

Moreover, these three graphs give forbidden structures for biclique graphs.

Corollary 5.3. *Let G be a graph which contains the Hajós graph, the rising sun or the X_1 graph as an induced subgraph where the vertices of degree 2 in the subgraph are also of degree 2 in G . Then G is not a biclique graph.*

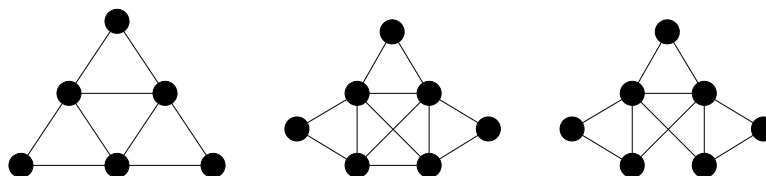


Figure 10: The Hajós graph, the rising sun and the X_1 graph.

Next we present the theorem that gives an upper bound on the number of vertices of degree 2 in a biclique graph.

Theorem 5.4. *Let $G = KB(H)$ for some graph H , where G is not isomorphic to K_3 or the diamond, and $|V(G)| = n$. Then the number of vertices of degree 2 in G is strictly less than $n/2$.*

Proof. Let $V_2 = \{v \in G : d(v) = 2\}$. First we show that there are no edges between the vertices of V_2 . Suppose by contradiction that $v_i, v_j \in V_2$ are adjacent. If they have a common neighbor, say w , since G is not isomorphic to K_3 , both $\{v_i, w, w'\}$ and $\{v_j, w, w'\}$, where w' is any other neighbor of w , induce a P_3 which is not contained in a diamond or a gem. Otherwise, if the vertex w is adjacent to v_i and not adjacent to v_j , then $\{w, v_i, v_j\}$ induces a P_3 that is not included in a diamond or a gem since no other vertex is adjacent to v_i . In both cases we arrive at a contradiction by Theorem 3.1.

Next we show that for each $v_i \in V_2$, there exists a vertex $w_i \in N(v_i)$ such that $w_i \notin N(v_j)$, for all $v_j \in V_2$, $j \neq i$. Clearly, $w_i \notin V_2$. By contradiction, suppose that there is a vertex $v_i \in V_2$ such that each of its two neighbors, say v', v'' , are adjacent to some other vertex of V_2 . Note again that by Theorem 3.1, v' and v'' are adjacent. Consider the following cases:

- $N(v_i) = N(v_j) = \{v', v''\}$, for some $v_j \in V_2$. Since G is not isomorphic to the diamond, by Proposition 4.6 we arrive at a contradiction.

- $v' \in N(v_j)$ and $v'' \in N(v_k)$, for some $v_j, v_k \in V_2$, $j \neq k$. Let v''' be the other vertex adjacent to v_j . Now, as $\{v_i, v', v_j\}$ induces a P_3 , then by Theorem 3.1, $\{v_i, v', v_j, v'', v'''\}$ induces a gem (the only one) containing that P_3 , since v', v''' and v'', v''' should be adjacent. Finally, if the vertices $v_i, v', v_j, v'', v''', v_k$ are respectively called $b, b_2, b', b_1, b_3, \tilde{b}$, then by Proposition 5.1 we obtain a contradiction since v_k cannot be adjacent to v'' . Note that depending on the other neighbor v'''' of v_k and their adjacencies, the *Hajós graph* ($v'''' = v'''$), the *rising sun* (v'''' adjacent to v' and v''') or the X_1 graph (v'''' adjacent to v' and not adjacent to v''') appears.

Therefore, for each vertex of degree 2, we can associate a unique neighbor of degree greater than 2 that is not adjacent to any other vertex of degree 2. Thus, as there is at least one neighbor of a vertex of V_2 that is not associated to any vertex, and has degree greater than 2, it follows that $|V_2| < n/2$. \square

As an application of Theorem 5.4, in Figure 9, the first and third graphs have more vertices of degree 2 than vertices of degree greater than 2. Therefore they are not biclique graphs. Another two examples are shown in Figure 11.

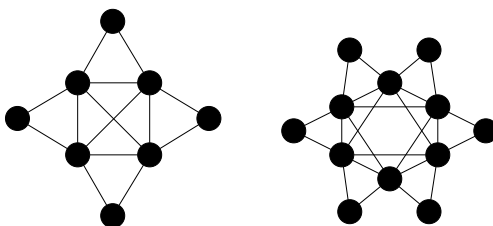


Figure 11: Graphs that are not biclique graphs by Theorem 5.4.

As another application of Proposition 5.1, we obtain the following result. For this, we need one more definition. A family of sets \mathcal{A} is *Helly* when every subfamily of pairwise intersecting subsets has a non-empty intersection.

Proposition 5.5. *Let G be a biclique graph and let $A = \{N[v] : v \in G \text{ and } d(v) = 2\}$. Then A is Helly.*

Proof. Let, to the contrary, A' be a minimal non-Helly subfamily of A . Now, since A' is minimal and each $N[v_i] \in A'$ induces a K_3 (as by Theorem 3.1, neighbors of v_i should be adjacent), we have $|A'| = 3$. Moreover, as A' is non-Helly, it induces the *Hajós graph* where the vertices of degree 2 in the subgraph are also of degree 2 in G . Therefore by Corollary 5.3, G is not a biclique graph, which is a contradiction. We conclude that A is a Helly family. \square

6 Open problems

In this section, we present some conjectures. We look for proofs or counterexamples.

We propose first the following conjecture that generalizes Proposition 5.5.

Conjecture 6.1. *Let G be a biclique graph and let $A = \{N[v] : v \in G \text{ and } v \text{ is simplicial}\}$. Then A is Helly.*

Proposition 4.6 can be extended, leading to the following conjecture.

Conjecture 6.2. *Let $G = KB(H)$ for some graph H , where G is not isomorphic to the diamond. Then there do not exist $v_1, v_2, \dots, v_i \in V(G)$ such that $N(v_1) = N(v_2) = \dots = N(v_i)$ and their neighbors are a subgraph of K_i for $i \geq 2$.*

7 Conclusions

In this work we have given a formula for the distances between vertices in the biclique graph $KB(G)$ using the distances between bicliques of G . This is a useful tool for proving structural properties in bicliques graphs. In particular, it allows us to give a different proof for the necessary condition for a graph to be a biclique graph to that given in [14]. Also, it is used to answer (negatively) the question about the condition being sufficient or not.

Finally, we have given an upper bound on the number of vertices of degree 2 for biclique graphs. Also we gave some forbidden structures that are useful to recognize graphs which are not biclique graphs.

Appendix

In this section we present a complete list of biclique graphs up to 6 vertices. (See Figure 12.) For those graphs that verify the necessary condition of biclique graphs, we have checked by computer whether they are biclique graphs using the characterization given in [14].

Acknowledgements

The authors are deeply grateful to the reviewers for their useful and very detailed comments.

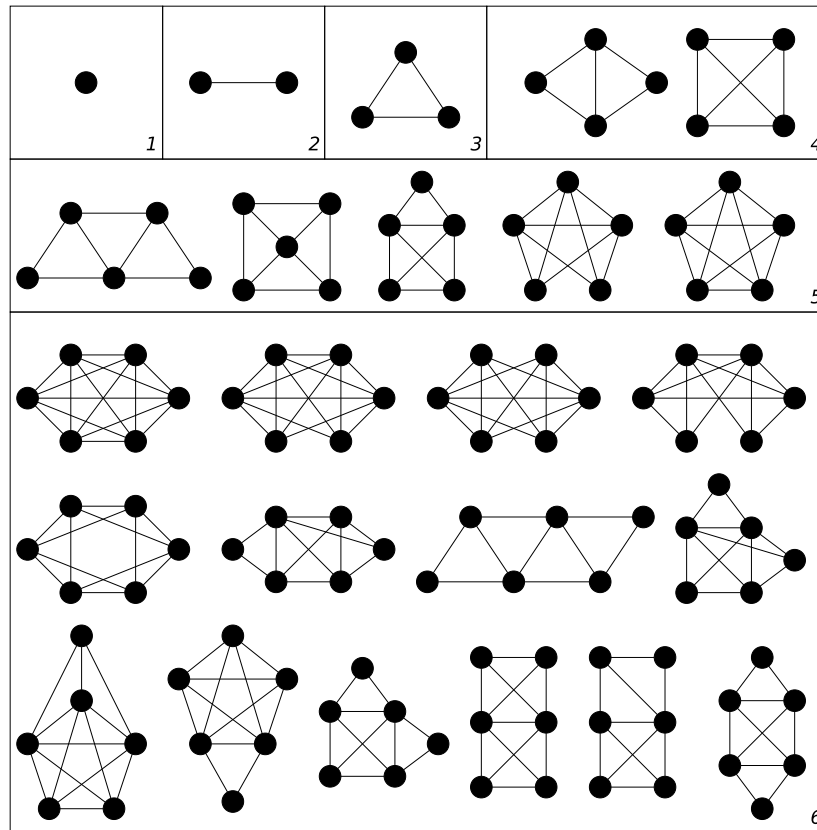


Figure 12: Biclique graphs up to 6 vertices.

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(Received 4 Sep 2020; revised 22 Apr 2021, 21 July 2021)