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On Tools for Completeness of Kleene Algebra with Hypotheses^{*}

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Abstract. In the literature on Kleene algebra, a number of variants have been proposed which impose additional structure specified by a theory, such as Kleene algebra with tests (KAT) and the recent Kleene algebra with observations (KAO), or make specific assumptions about certain constants, as for instance in NetKAT. Many of these variants fit within the unifying perspective offered by *Kleene algebra with hypotheses*, which comes with a canonical language model constructed from a given set of hypotheses. For the case of KAT, this model corresponds to the familiar interpretation of expressions as languages of guarded strings.

A relevant question therefore is whether Kleene algebra together with a given set of hypotheses is complete with respect to its canonical language model. In this paper, we revisit, combine and extend existing results on this question to obtain tools for proving completeness in a modular way. We showcase these tools by reproving completeness of KAT and KAO, and prove completeness of a new variant of KAT where the collection of tests only forms a distributive lattice.

Keywords: Kleene algebra · Completeness · Reduction

1 Introduction

Kleene algebras (KA) [17,8] are algebraic structures involving an iteration operation, Kleene star, corresponding to reflexive-transitive closure in relational models and to language iteration in language models. Its axioms are complete w.r.t. relational models and language models [18,29,3], and the resulting equational theory is decidable via automata algorithms (in fact, PSPACE-complete [30]).

These structures were later extended in order to deal with common programming constructs. For instance, Kleene algebras with tests (KAT) [22], which combine Kleene algebra and Boolean algebra, make it possible to represent the control flow of while programs. Kleene star is used for while loops, and Boolean

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tests are used for the conditions of such loops, as well as the conditions in if-then-else statements. Again, the axioms of KAT are complete w.r.t. appropriate classes of models, and its equational theory remains in PSPACE. Proving so is non-trivial: Kozen’s proof reduces completeness of KAT to completeness of KA, via a direct syntactic transformation on terms.

Another extension is Concurrent Kleene algebra (CKA) [13], where a binary operator for parallelism is added. The resulting theory is characterised by languages of pomsets rather than languages of words, and is EXPSpace-complete [6]. Trying to have both tests and concurrency turned out to be non-trivial, and called for yet another notion: Kleene algebras with observations (KAO) [15], which are again complete w.r.t. appropriate models, and decidable.

When used in the context of program verification, e.g., in a proof assistant, such structures make it possible to write algebraic proofs of correctness, and to mechanise some of the steps: when two expressions e and f representing two programs happen to be provably equivalent in KA, KAT, or KAO, one does not need to provide a proof, one can simply call a certified decision procedure [4,28,31]. However, this is often not enough [26,1,12]: most of the time, the expressions e and f are provably equal only under certain assumptions on their constituents. For instance, to prove that $(a + b)^*$ and a^*b^* are equal, one may have to use that in the considered instance, we have $ba = ab$. In other words, one would like to prove equations under some assumptions, to have algorithms for the Horn theory of Kleene algebra and its extensions rather than just their equational theories.

Unfortunately, those Horn theories are typically undecidable [21,24], even with rather restricted forms of hypotheses (e.g., commutation of two letters, as in the above example). Nevertheless, important and useful classes of hypotheses can be ‘eliminated’, by reducing to the plain and decidable case of the equational theory. This is for instance the case of *Hoare hypotheses* [23], of the shape $e = 0$, which make it possible to encode Hoare triples for partial correctness in KAT.

In some cases, one wants to exploit hypotheses about specific constituents (e.g. a and b in the above example). In other situations, one wants to exploit assumptions on the whole structure. For instance, in commutative Kleene algebra [33,8,5], one assumes that the product is commutative everywhere.

Many of these extensions of Kleene algebra (KAT, KAO, commutative KA, specific hypotheses) fit into the generic framework of Kleene algebra with hypotheses [10], providing in each case a canonical model in terms of closed languages.

We show that we recover standard models in this way, and we provide tools to establish completeness and decidability of such extensions, in a modular way. The key notion is that of *reduction* from one set of hypotheses to another. We summarise existing reductions and we provide a toolbox for combining those reductions together. We use this toolbox in order to obtain new and modular proofs of completeness for KAT and KAO, as well as for the fragment of KAT where tests are only assumed to form a distributive lattice.

Note however that there are Kleene algebra extensions like *action algebras* [32] or *action lattices* [20], which do not seem to fit into the framework of Kleene

algebra with hypotheses: it is not clear how to interpret the additional operations as letters with additional structure.

2 Kleene algebra, hypotheses, closures

A *Kleene algebra* [8,19] is a tuple $(K, +, \cdot, *, 0, 1)$ such that $(K, +, \cdot, 0, 1)$ is an idempotent semiring, and $*$ is a unary operator on K such that for all $x, y \in K$ the following axioms are satisfied:

$$1 + x \cdot x^* \leq x^* \quad x + y \cdot z \leq z \Rightarrow y^* \cdot x \leq z \quad x + y \cdot z \leq y \Rightarrow x \cdot z^* \leq y$$

There, as later in the paper, we write $x \leq y$ as a shorthand for $x + y = y$. Given the idempotent semiring axioms, \leq is a partial order in every Kleene algebra, and all operations are monotone w.r.t. that order.

We let e, f range over regular expressions over an alphabet Σ , defined by:

$$e, f ::= e + f \mid e \cdot f \mid e^* \mid 0 \mid 1 \mid a \in \Sigma$$

We write $\mathcal{T}(\Sigma)$ for the set of such expressions, or simply \mathcal{T} when the alphabet is clear from the context. Given alphabets Σ and Γ , a function $h: \Sigma \rightarrow \mathcal{T}(\Gamma)$ extends uniquely into a homomorphism $h: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Gamma)$, which we refer to as the *homomorphism generated by h* . As usual, every regular expression e gives rise to a language $\llbracket e \rrbracket \in \mathcal{P}(\Sigma^*)$. Given two regular expressions, we moreover write $\text{KA} \vdash e = f$ when $e = f$ is derivable from the axioms of Kleene algebra. (Equivalently, when the equation $e = f$ holds universally, in all Kleene algebras.)

The central theorem of Kleene algebra is the following:

Theorem 2.1 (Soundness and Completeness of KA [18,29,3]). *For all $e, f \in \mathcal{T}$, we have $\text{KA} \vdash e = f$ if and only if $\llbracket e \rrbracket = \llbracket f \rrbracket$.*

As a consequence, the equational theory of Kleene algebras is decidable.

Our goal is to extend this result to the case where we have additional hypotheses on some of the letters of the alphabet, or axioms restricting the behaviour of certain operations. Those are represented by sets of *inequations*, i.e., pairs (e, f) of regular expressions written $e \leq f$ for the sake of clarity. Given a set H of such inequations, we write $\text{KA}_H \vdash e \leq f$ when the inequation $e \leq f$ is derivable from the axioms of Kleene algebra and the hypotheses in H (similarly for equations). By extension, we write $\text{KA}_H \vdash H'$ when $\text{KA}_H \vdash e \leq f$ for all $e \leq f$ in H' .

Note that we consider letters of the alphabet as constants rather than variables. In particular, while we have $\text{KA}_{ba \leq ab} \vdash (a+b)^* \leq a^*b^*$, we do not have $\text{KA}_{ba \leq ab} \vdash (a+c)^* \leq a^*c^*$. Formally, we use a notion of derivation where there is no substitution rule, and where we have all instances of Kleene algebra axioms as axioms. When we want to consider hypotheses that are universally valid, it suffices to use all their instances. For example, to define commutative Kleene algebra, we simply use the infinite set $\{ef \leq fe \mid e, f \in \mathcal{T}\}$.

We associate a canonical language model to KA with a set of hypotheses H , defined by *closure* under H . For $u, v \in \Sigma^*$ and $L \subseteq \Sigma^*$, let $uLv \triangleq \{uxv \mid x \in L\}$.

Definition 2.2 (H-closure). *Let H be a set of hypotheses and $L \subseteq \Sigma^*$ a language. The H -closure of L , denoted as $\text{cl}_H(L)$, is the smallest language containing L s.t. for all $e \leq f \in H$ and $u, v \in \Sigma^*$, if $u \llbracket f \rrbracket v \subseteq \text{cl}_H(L)$, then $u \llbracket e \rrbracket v \subseteq \text{cl}_H(L)$.*

Fixpoint theory makes it possible to characterise the H -closure of a language L as the least (pre)fixpoint of the function $\text{st}'_{H,L}(X) = \text{st}_H(X) \cup L$, where

$$\text{st}_H(X) = \bigcup \{u \llbracket e \rrbracket v \mid e \leq f \in H, u, v \in \Sigma^*, u \llbracket f \rrbracket v \subseteq X\}.$$

This least fixpoint can be characterised more explicitly by transfinite iteration: we have $\text{cl}_H(L) = \bigcup_{\alpha} \text{st}_H^{\alpha}(L)$ where $\text{st}_H^{\alpha+1}(L) = \text{st}_H(\text{st}_H^{\alpha}(L))$ for every ordinal α , and $\text{st}_H^{\lambda}(L) = \bigcup_{\alpha < \lambda} \text{st}_H^{\alpha}(L)$ for every limit ordinal λ .

This notion of closure gives a *closed interpretation* of regular expressions, $\text{cl}_H(\llbracket - \rrbracket)$, for which KA_H is sound:

Theorem 2.3 ([10, Theorem 2]). *If $\text{KA}_H \vdash e = f$, then $\text{cl}_H(\llbracket e \rrbracket) = \text{cl}_H(\llbracket f \rrbracket)$.*

In the sequel, we shall prove the converse implication, completeness, for specific choices of H : we say that KA_H is *complete* if for all expressions e, f :

$$\text{cl}_H(\llbracket e \rrbracket) = \text{cl}_H(\llbracket f \rrbracket) \text{ implies } \text{KA}_H \vdash e = f.$$

We could hope that completeness always holds, notably because the notion of closure is invariant under inter-derivability of the considered hypotheses, as a consequence of the following lemma:

Lemma 2.4 ([16, Lemma 4.10]). *Let H and H' be sets of hypotheses such that $\text{KA}_H \vdash H'$. Then $\text{cl}_{H'} \subseteq \text{cl}_H$.*

Unfortunately, there are concrete instances for which KA_H is known not to be complete. For instance, there is a finitely presented monoid (thus a finite set H_0 of equations) such that $\{(e, f) \mid \text{cl}_{H_0}(\llbracket e \rrbracket) = \text{cl}_{H_0}(\llbracket f \rrbracket)\}$ is not r.e. [25, Theorem 1]. Since derivability in KA_H is r.e. as soon as H is, KA_{H_0} cannot be complete.

Before turning to techniques for proving completeness, let us describe the closed interpretation of regular expressions for two specific choices of hypotheses.

Let us consider first *commutative Kleene algebra*, obtained as explained in the Introduction using the set $\{ef \leq fe \mid e, f \in \mathcal{T}(\Sigma)\}$. Under Kleene algebra axioms, this set is equiderivable with its restriction to letters, $C = \{ab \leq ba \mid a, b \in \Sigma\}$ (a consequence of [1, Lemma 4.4]).

The associated closure can be characterised as follows:

$$\text{cl}_C(L) = \{w \in \Sigma^* \mid \exists v \in L. |w|_x = |v|_x \text{ for all } x \in \Sigma\}$$

where $|w|_x$ denotes the number of occurrences of x in w . Thus, $w \in \text{cl}_C(L)$ if it is a permutation of some word in L .

This semantics matches precisely the one used in [8] for commutative Kleene algebra: there, a function $\llbracket - \rrbracket_C : \mathcal{T}(\Sigma) \rightarrow \mathcal{P}(\mathbb{N}^{\Sigma})$ interprets regular expressions as subsets of \mathbb{N}^{Σ} , whose elements are thought of as ‘‘commutative words’’: these

assign to each letter the number of occurrences, but there is no order of letters. Let $q: \mathcal{P}(\Sigma^*) \rightarrow \mathcal{P}(\mathbb{N}^\Sigma)$, $q(L) = \{\lambda x. |w|_x \mid w \in L\}$; this map computes the *Parikh image* of a given language L , that is, the set of multisets representing occurrences of letters in words in L . Then this semantics is characterised by $\llbracket - \rrbracket_c = q \circ \llbracket - \rrbracket$.

One may observe that $\llbracket - \rrbracket_c = q(\text{cl}_C(\llbracket - \rrbracket))$, since cl_C only adds words to a language which have the same number of occurrences of each letter as some word which is already there. Conversely, we have $\text{cl}_C(\llbracket - \rrbracket) = q'(\llbracket - \rrbracket_c)$, where $q': \mathcal{P}(\mathbb{N}^\Sigma) \rightarrow \mathcal{P}(\Sigma^*)$, $q'(L) = \{w \mid p \in L, \forall x \in \Sigma, |w|_x = p(x)\}$. As a consequence, we have $\llbracket e \rrbracket_c = \llbracket f \rrbracket_c$ if and only if $\text{cl}_C(\llbracket e \rrbracket) = \text{cl}_C(\llbracket f \rrbracket)$.

From there, we can easily deduce from the completeness result in [8, Chapter 11, Theorem 4], attributed to Pilling (see also [5]), that KA_C is complete.

Let us now consider a single hypothesis: $D = \{ab \leq 0\}$ for some letters a and b . The D -closure of a language L consists of those words that either belong to L , or contain ab as a subword. As a consequence, we have $\text{cl}_D(\llbracket e \rrbracket) = \text{cl}_D(\llbracket f \rrbracket)$ if and only if $\llbracket e \rrbracket$ and $\llbracket f \rrbracket$ agree on all words not containing the pattern ab .

In this example, we can easily obtain decidability and completeness of KA_D . Indeed, consider the function $r: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Sigma)$, $r(e) = e + \Sigma^*ab\Sigma^*$. For all e , we have $\text{KA}_D \vdash e = r(e)$, and $\text{cl}_D(\llbracket e \rrbracket) = \llbracket r(e) \rrbracket$. As a consequence, we have

$$\begin{array}{ll}
 & \text{cl}_D(\llbracket e \rrbracket) = \text{cl}_D(\llbracket f \rrbracket) \\
 \Leftrightarrow & \llbracket r(e) \rrbracket = \llbracket r(f) \rrbracket \\
 \Leftrightarrow & \text{KA} \vdash r(e) = r(f) \quad (\text{Theorem 2.1}) \\
 \Rightarrow & \text{KA}_D \vdash e = f
 \end{array}$$

The first step above establishes decidability of the closed semantics; the following ones reduce the problem of completeness for KA_D to that for KA alone, which is known to hold. By soundness (Theorem 2.3), the last line implies the first one, so that these conditions are all equivalent.

This second example exploits and illustrates a simple instance of the framework we design in the sequel to prove completeness of various sets of hypotheses.

3 Reductions

As illustrated above, the overall strategy is to reduce completeness of KA_H , for a given set of hypotheses H , to completeness of Kleene algebra. The core idea is to provide a map r from expressions to expressions, which incorporates the hypotheses H in the sense that $\llbracket r(e) \rrbracket = \text{cl}_H(\llbracket e \rrbracket)$, and such that $r(e)$ is provably equivalent to e under the hypotheses H . This idea leads to the unifying notion of *reduction*, developed in [25,10,16].

Definition 3.1 (Reduction). *Assume $\Gamma \subseteq \Sigma$ and let H, H' be sets of hypotheses over Σ and Γ respectively. We say that H reduces to H' if $\text{KA}_H \vdash H'$ and there exists a map $r: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Gamma)$ such that for all $e \in \mathcal{T}(\Sigma)$,*

1. $\text{KA}_H \vdash e = r(e)$, and

$$2. \text{cl}_H(\llbracket e \rrbracket) \cap \Gamma^* = \text{cl}_{H'}(\llbracket r(e) \rrbracket).$$

We often refer to such a witnessing map r itself as a reduction. Generalising the above example, we obtain the key property of reductions:

Theorem 3.2. *Suppose H reduces to H' . If $\text{KA}_{H'}$ is complete, then so is KA_H .*

Proof. Let r be the map for the reduction from H to H' . For all $e, f \in \mathcal{T}(\Sigma)$,

$$\begin{aligned} & \text{cl}_H(\llbracket e \rrbracket) = \text{cl}_H(\llbracket f \rrbracket) \\ \Rightarrow & \text{cl}_{H'}(\llbracket r(e) \rrbracket) = \text{cl}_{H'}(\llbracket r(f) \rrbracket) && (r \text{ a reduction (item 2)}) \\ \Rightarrow & \text{KA}_{H'} \vdash r(e) = r(f) && (\text{completeness of } H') \\ \Rightarrow & \text{KA}_H \vdash r(e) = r(f) && (\text{KA}_H \vdash H') \\ \Rightarrow & \text{KA}_H \vdash e = f && (r \text{ a reduction (item 1)}) \quad \square \end{aligned}$$

An important case is when $H' = \emptyset$: given a reduction from H to \emptyset , Theorem 3.2 gives completeness of KA_H , by completeness of KA . Such reductions are what we ultimately aim for. However, in the examples later in this paper, these reductions are composed of smaller ones, which do make use of intermediate hypotheses. Section 3.2 contains general techniques for combining reductions.

While we focus on completeness in this paper, note that reductions can also be used to prove decidability. More precisely, if KA'_H is complete and decidable, and H reduces to H' via a computable reduction r , then KA_H is decidable.

The following result from [16] (cf. Remark 3.5) gives a sufficient condition for the existence of a reduction. This is useful for reductions where the underlying map r is a homomorphism.

Lemma 3.3. *Assume $\Gamma \subseteq \Sigma$ and let H, H' be sets of hypotheses over Σ and Γ respectively, such that $\text{KA}_H \vdash H'$. If there exists a homomorphism $r: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Gamma)$ such that:*

1. *For all $a \in \Gamma$, we have $\text{KA} \vdash a \leq r(a)$.*
2. *For all $a \in \Sigma$, we have $\text{KA}_H \vdash a = r(a)$.*
3. *For all $e \leq f \in H$, we have $\text{KA}_{H'} \vdash r(e) \leq r(f)$.*

then H reduces to H' .

Example 3.4. We consider KA together with a global “top element” \top and the axiom $e \leq \top$. To make this precise in Kleene algebra with hypotheses, we assume an alphabet Σ with $\top \in \Sigma$, and take the set of hypotheses $H_\top = \{e \leq \top \mid e \in \mathcal{T}(\Sigma)\}$. Then $\text{cl}_{H_\top}(L)$ contains those words obtained from a word $w \in L$ by replacing every occurrence of \top in w by arbitrary words in Σ^* .

We claim that H reduces to \emptyset . To this end, define the homomorphism $r: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Sigma)$ by $r(\top) = \Sigma^*$ (where we view Σ as an expression consisting of the sum of its elements) and $r(a) = a$ for $a \in \Sigma$ with $a \neq \top$. Each of the conditions of Lemma 3.3 is now easy to check. Thus r is a reduction, so by Theorem 3.2, KA_{H_\top} is complete.

Note that this implies completeness w.r.t. validity of equations in all (regular) language models, where \top is interpreted as the largest language: indeed, the closed semantics $\text{cl}_{H_\top}(\llbracket - \rrbracket)$ is generated by such a model.

At the end of Section 2, we discussed commutative KA as an instance of Kleene algebra with hypotheses H . While KA_H is complete in that case, there is no reduction from H to \emptyset , as cl_H does not preserve regularity. Indeed, $\text{cl}_H(\llbracket(ab)^*\rrbracket) = \{w \mid |w|_a = |w|_b\}$ which is not regular. The completeness proof in [8,5] is self-contained, and does not rely on completeness of KA.

Remark 3.5. The idea to use two sets of hypotheses in Definition 3.1 is from [16], where reductions are defined slightly differently: the alphabet is fixed (that is, $\Sigma = \Gamma$), and the last condition is instead defined as $\text{cl}_H(\llbracket e \rrbracket) = \text{cl}_H(\llbracket f \rrbracket) \Rightarrow \text{cl}_{H'}(\llbracket r(e) \rrbracket) = \text{cl}_{H'}(\llbracket r(f) \rrbracket)$. An extra notion of *strong* reduction is then introduced, which coincides with our definition if $\Sigma = \Gamma$. By allowing a change of alphabet, we do not need to distinguish reductions and strong reductions. Lemma 3.3 is in [16, Lemma 4.23], adapted here to the case with two alphabets (this is taken care of in *loc. cit.* by assuming $\text{cl}_{H'}$ preserves languages over Γ).

3.1 Basic reductions

The following result collects several sets of hypotheses for which we have reductions to \emptyset . These mostly come from the literature. They form basic building blocks used in the more complex reductions that we present in the examples below.

Lemma 3.6. *Each of the following sets of hypotheses reduce to the empty set (of hypotheses over Σ).*

- (i) $\{u_i \leq w_i \mid i \in I\}$ with $u_i, w_i \in \Sigma^*$ and $|u_i| \leq 1$ for all $i \in I$
- (ii) $\{1 \leq \sum_{a \in S_i} a \mid i \in I\}$ with each $S_i \subseteq \Sigma$ finite
- (iii) $\{e \leq 0\}$ for $e \in \mathcal{T}(\Sigma)$
- (iv) $\{ea \leq a\}$ and $\{ae \leq a\}$ for $a \in \Sigma$, $e \in \mathcal{T}(\Sigma \setminus \{a\})$

Proof. (i) This is [25, Theorem 2]. (The result mentions equations, but in the proof only the relevant inequations are used.)

(ii) This is [10, Proposition 6].

(iii) This is basically due to [7], but since it is phrased differently there we include a proof in Appendix B.

(iv) Hypotheses of a similar form as (iv) are studied in the setting of Kleene algebra with tests in [12], we include a proof in Appendix B. \square

Note that Item iii above covers finite sets of hypotheses of the form $\{e_i \leq 0\}_{i \in I}$, as these can be encoded as the single hypothesis $\sum_{i \in I} e_i \leq 0$.

3.2 Compositional reductions

The previous subsection gives reductions to the empty set for single equations. However, in the examples we often start with a *collection* of hypotheses of different shapes, which we wish to reduce to the empty set. Therefore, we now discuss a few techniques for combining reductions.

Throughout this section, for sets of hypotheses H_1, \dots, H_n we often denote the associated closure by cl_i instead of cl_{H_i} , $\text{cl}_{i,j}$ instead of $\text{cl}_{H_i \cup H_j}$ and $\text{cl}_{i\dots j}$ instead of $\text{cl}_{\bigcup_{i \leq k \leq j} H_k}$. Similarly, we write st_i instead of st_{H_i} etc. First, there are the basic observations that reductions compose (Lemma 3.7) and that equiderivable sets of hypotheses always reduce to each other, via the identity (Lemma 3.8).

Lemma 3.7. *Let H_1, H_2 and H_3 be sets of hypotheses. If H_1 reduces to H_2 and H_2 reduces to H_3 then H_1 reduces to H_3 .*

Equiderivable sets of hypotheses always reduce to each other, via the identity:

Lemma 3.8. *Let H_1, H_2 be sets of hypotheses over a common alphabet. If $\text{KA}_{H_1} \vdash H_2$ and $\text{KA}_{H_2} \vdash H_1$ then H_1 and H_2 reduce to each other.*

The following useful lemma allows to combine reductions by union, and is used in many of the examples. Its assumptions allow one to compose the reductions sequentially. A similar lemma is formulated in the setting of bi-Kleene algebra in [14, Lemma 4.46].

Lemma 3.9. *Let H_1, \dots, H_n, H be sets of hypotheses over a common alphabet Σ , with $n \geq 1$. If H_i reduces to H for all i , and $\text{cl}_{1\dots n} = \text{cl}_n \circ \dots \circ \text{cl}_1$, then $\bigcup_{i \leq n} H_i$ reduces to H .*

The next lemma is useful to show the second requirement in Lemma 3.9.

Lemma 3.10. *Let H_1, \dots, H_n be sets of hypotheses, such that $\text{cl}_i \circ \text{cl}_j \subseteq \text{cl}_j \circ \text{cl}_i$ for all i, j with $i < j$. Then $\text{cl}_{1\dots n} = \text{cl}_n \circ \dots \circ \text{cl}_1$.*

Remark 3.11. The condition $\text{cl}_i \circ \text{cl}_j \subseteq \text{cl}_j \circ \text{cl}_i$ in Lemma 3.10 is equivalent to $\text{cl}_{i,j} = \text{cl}_j \circ \text{cl}_i$. With the latter formulation, Lemma 3.10 is stated in the bi-Kleene algebra setting as [14, Lemma 4.50].

We now proceed with several lemmas that help proving $\text{cl}_1 \circ \text{cl}_2 \subseteq \text{cl}_2 \circ \text{cl}_1$. In particular, these allow to use the “one-step closure” st_H from the fixed point characterisation of cl_H (below Definition 2.2), for cl_1 in Lemma 3.12 and additionally for cl_2 in Lemma 3.13, assuming further conditions.

Lemma 3.12. *For all H_1, H_2 , if $\text{st}_1 \circ \text{cl}_2 \subseteq \text{cl}_2 \circ \text{cl}_1$, then $\text{cl}_1 \circ \text{cl}_2 \subseteq \text{cl}_2 \circ \text{cl}_1$.*

Lemma 3.13. *Let H_1, H_2 be sets of hypotheses such that the right-hand sides of inequations in H_1 are all words. If $\text{st}_1 \circ \text{st}_2 \subseteq \text{cl}_2 \circ \text{st}_1^-$ then $\text{cl}_1 \circ \text{cl}_2 \subseteq \text{cl}_2 \circ \text{cl}_1$.*

We conclude this section by returning to hypotheses of the form $e \leq 0$. If $H_1 = \{e \leq 0\}$ for some term e and H_2 is an arbitrary set of hypotheses, then we have $\text{cl}_1 \circ \text{cl}_2 \subseteq \text{cl}_2 \circ \text{cl}_1$. As a consequence, we can strengthen Lemma 3.6(iii) to the following result, which gives a general treatment of hypotheses of the form $e \leq 0$: we can always get rid of finite sets of hypotheses of this form. A similar result, in terms of Horn formulas and in the context of KAT, is shown in [11].

Lemma 3.14. *For any set of hypotheses H and any term e , there is a reduction from $H \cup \{e \leq 0\}$ to H .*

4 Kleene algebra with tests

In this section we apply the machinery from the previous sections to obtain a modular completeness proof for Kleene algebra with tests [27].

A *Kleene algebra with tests (KAT)* is a Kleene algebra X containing a Boolean algebra L such that the meet of L coincides with the product of X , the join of X coincides with the sum of L , the top element of L is the multiplicative identity of X , and the bottom elements of X and L coincide.

Syntactically, we fix two finite sets Σ and Ω of primitive actions and primitive tests. We denote the set of Boolean expressions over alphabet Ω by \mathcal{T}_{BA} :

$$\phi, \psi ::= \phi \vee \psi \mid \phi \wedge \psi \mid \neg\phi \mid \perp \mid \top \mid o \in \Omega$$

We write $\text{BA} \vdash \phi = \phi'$ when this equation is derivable from Boolean algebra axioms [2,9], and similarly for inequations.

We let α, β range over *atoms*: elements of the set $\text{At} \triangleq 2^\Omega$. Those may be seen as valuations for Boolean expressions, or as complete conjunctions of literals: α is implicitly seen as the Boolean formula $\bigwedge_{\alpha(o)=1} o \wedge \bigwedge_{\alpha(o)=0} \neg o$. They form the atoms of the Boolean algebra generated by Ω . We write $\alpha \models \phi$ when ϕ holds under the valuation α . A key property of Boolean algebras is that for all atoms α and formulas ϕ , we have

$$\alpha \models \phi \Leftrightarrow \text{BA} \vdash \alpha \leq \phi \quad \text{and} \quad \text{BA} \vdash \phi = \bigvee_{\alpha \models \phi} \alpha$$

The KAT terms over alphabets Σ and Ω are the regular expressions over the alphabet $\Sigma + \mathcal{T}_{\text{BA}}$: $\mathcal{T}_{\text{KAT}} \triangleq \mathcal{T}(\Sigma + \mathcal{T}_{\text{BA}})$. We write $\text{KAT} \vdash e = f$ when this equation is derivable from the axioms of KAT, and similarly for inequations.

The standard interpretation of KAT associates to each term a language of guarded strings. A *guarded string* is a sequence of the form $\alpha_0 a_0 \alpha_1 a_1 \dots a_{n-1} \alpha_n$ with $a_i \in \Sigma$ for all $i < n$, and $\alpha_i \in \text{At}$ for all $i \leq n$. We write \mathcal{GS} for the set $\text{At} \times (\Sigma \times \text{At})^*$ of such guarded strings. Now, the interpretation $\mathcal{G}: \mathcal{T}(\Sigma + \mathcal{T}_{\text{BA}}) \rightarrow 2^{\mathcal{GS}}$ is defined as the homomorphic extension of the assignment $\mathcal{G}(a) = \{\alpha a \beta \mid \alpha, \beta \in \text{At}\}$ for $a \in \Sigma$ and $\mathcal{G}(\phi) = \{\alpha \mid \alpha \models \phi\}$ for $\phi \in \mathcal{T}_{\text{BA}}$, where for sequential composition of guarded strings the *coalesced product* is used. The coalesced product of guarded strings $u\alpha$ and βv is defined as $u\alpha v$ if $\alpha = \beta$ and undefined otherwise.

Theorem 4.1 ([27, Theorem 8]). *For all $e, f \in \mathcal{T}_{\text{KAT}}$, we have $\text{KAT} \vdash e = f$ iff $\mathcal{G}(e) = \mathcal{G}(f)$.*

We now reprove this result using Kleene algebra with hypotheses. We start by defining the additional axioms of KAT as hypotheses.

Definition 4.2. *We write **bool** for the set of all instances of Boolean algebra axioms over \mathcal{T}_{BA} and **glue** for the following set of hypotheses relating the Boolean algebra connectives to the Kleene algebra ones*

$$\text{glue} = \{\phi \wedge \psi = \phi \cdot \psi, \phi \vee \psi = \phi + \psi \mid \phi, \psi \in \mathcal{T}_{\text{BA}}\} \cup \{\perp = 0, \top = 1\}$$

We then define $\text{kat} = \text{bool} \cup \text{glue}$.

(Note that all these equations are actually understood as double inequations.)

We prove completeness of KA_{kat} in Section 4.2 below, by constructing a suitable reduction. Recall that this means completeness w.r.t. the interpretation $\text{cl}_{\text{kat}}(\llbracket - \rrbracket)$ in terms of closed languages. Before proving completeness of KA_{kat} , we compare it to the classical completeness (Theorem 4.1). First note that KA_{kat} contains the same axioms as Kleene algebra with tests, so that provability in KA_{kat} and KAT coincide: $\text{KA}_{\text{kat}} \vdash e = f$ iff $\text{KAT} \vdash e = f$. Comparing the interpretation $\text{cl}_{\text{kat}}(\llbracket - \rrbracket)$ to the guarded string interpretation \mathcal{G} is slightly more subtle, and is the focus of the next subsection.

4.1 Relation to guarded string interpretation

To relate the guarded string model and the model obtained with closure under kat , we first develop the following lemmas.

The key step consists in characterising the strings that are present in the closure of a language of guarded strings (Lemma 4.3) below. First observe that a guarded string may always be seen as a word over the alphabet $\Sigma + \mathcal{T}_{\text{BA}}$. Conversely, a word over the alphabet $\Sigma + \mathcal{T}_{\text{BA}}$ can always be decomposed as a sequence $\phi_0 a_0 \cdots \phi_{n-1} a_{n-1} \phi_n$ where $a_i \in \Sigma$ for all $i < n$ and each ϕ_i is a possibly empty sequence of Boolean expressions. We let ϕ range over such sequences, and we write $\overline{\phi}$ for the conjunction of the elements of ϕ .

Lemma 4.3. *Let L be a language of guarded strings. We have*

$$\begin{aligned} & \phi_0 a_0 \cdots \phi_{n-1} a_{n-1} \phi_n \in \text{cl}_{\text{kat}}(L) \\ \Leftrightarrow & \quad \forall (\alpha_i)_{i \leq n}, \quad (\forall i \leq n, \alpha_i \models \overline{\phi_i}) \quad \Rightarrow \quad \alpha_0 a_0 \cdots \alpha_{n-1} a_{n-1} \alpha_n \in L \end{aligned}$$

Then we show that the kat -closures of $\llbracket e \rrbracket$ and $\mathcal{G}(e)$ coincide:

Lemma 4.4. *For all KAT expressions e , $\text{cl}_{\text{kat}}(\llbracket e \rrbracket) = \text{cl}_{\text{kat}}(\mathcal{G}(e))$.*

Let \mathcal{GS} be the set of all guarded strings. We also have:

Lemma 4.5. *For all KAT expressions e , $\mathcal{G}(e) = \text{cl}_{\text{kat}}(\llbracket e \rrbracket) \cap \mathcal{GS}$.*

As an immediate consequence of these two lemmas, we can finally relate the guarded strings languages semantics to the kat -closed languages one:

Corollary 4.6. *Let $e, f \in \mathcal{T}_{\text{kat}}$. We have $\mathcal{G}(e) = \mathcal{G}(f) \Leftrightarrow \text{cl}_{\text{kat}}(\llbracket e \rrbracket) = \text{cl}_{\text{kat}}(\llbracket f \rrbracket)$.*

4.2 Completeness

To prove completeness of the closed language model wrt kat , we take the following steps:

1. We reduce the hypotheses in kat to a simpler set of axioms: by putting the Boolean expressions into normal forms via the atoms, we can get rid of the

hypotheses in **bool**. We do not remove the hypotheses in **glue** directly: we transform them into the following hypotheses about atoms:

$$\mathbf{atom} = \{\alpha \cdot \beta \leq 0 \mid \alpha, \beta \in \mathbf{At}, \alpha \neq \beta\} \cup \{\alpha \leq 1 \mid \alpha \in \mathbf{At}\} \cup \{1 \leq \sum_{\alpha \in \mathbf{At}} \alpha\}$$

We thus first show that **kat** reduces to **atom**.

2. Then we use results from Section 3.1 and Section 3.2 to construct a reduction from **atom** to the empty set, and thereby obtain completeness of $\mathbf{KA}_{\mathbf{kat}}$.

Let $r: \mathcal{T}(\Sigma + \mathcal{T}_{\mathbf{BA}}) \rightarrow \mathcal{T}(\Sigma + \mathbf{At})$ be the homomorphism defined by

$$r(x) = \begin{cases} a & x = a \in \Sigma \\ \sum_{\alpha \models \phi} \alpha & x = \phi \in \mathcal{T}_{\mathbf{BA}} \end{cases}$$

We show below that r yields a reduction from **kat** to **atom**, using Lemma 3.3. In the sequel, we use \mathbf{atom}_1 , \mathbf{atom}_2 and \mathbf{atom}_3 , or simply 1, 2, 3, to denote the three families of inequations in **atom**.

Lemma 4.7. *We have $\mathbf{KA}_{\mathbf{atom}} \vdash 1 = \sum_{\alpha \in \mathbf{At}} \alpha$, and for all α , $\mathbf{KA}_{\mathbf{atom}} \vdash \alpha\alpha = \alpha$.*

Lemma 4.8. *For all $e \leq f \in \mathbf{kat}$, we have that $\mathbf{KA}_{\mathbf{atom}} \vdash r(e) \leq r(f)$.*

Lemma 4.9. *The homomorphism r yields a reduction from **kat** to **atom**.*

Proof. We use Lemma 3.3. We first need to show $\mathbf{KA}_{\mathbf{kat}} \vdash \mathbf{atom}$: for $\alpha, \beta \in \mathbf{At}$ with $\alpha \neq \beta$, we have the following derivations in $\mathbf{KA}_{\mathbf{kat}}$

$$\alpha \cdot \beta = \alpha \wedge \beta = \perp = 0 \quad \alpha \leq \top = 1 \quad 1 = \top = \bigvee_{\alpha \models \top} \alpha = \sum_{\alpha \in \mathbf{At}} \alpha$$

Now for $a \in \Sigma + \mathbf{At}$, we have $a = r(a)$ (syntactically): if $a = \mathbf{a} \in \Sigma$, then $r(\mathbf{a}) = \mathbf{a}$; if $a = \alpha \in \mathbf{At}$, then $r(\alpha) = \sum_{\alpha \models \alpha} \alpha = \alpha$. The first condition about r is thus satisfied, and it suffices to verify the second condition about r for $\phi \in \mathcal{T}_{\mathbf{BA}}$. In this case, we have we have $\mathbf{KA}_{\mathbf{kat}} \vdash r(\phi) = \sum_{\alpha \models \phi} \alpha = \bigvee_{\alpha \models \phi} \alpha = \phi$. The third and last condition was proven in Lemma 4.8. \square

Now we must reduce **atom** to the empty set. We can immediately get rid of \mathbf{atom}_1 : by Lemma 3.14, **atom** reduces to $\mathbf{atom}_{2,3}$. For \mathbf{atom}_2 and \mathbf{atom}_3 , we have individual reductions to the empty set via Lemma 3.6(i) and (ii), respectively. We combine those reductions via Lemma 3.9, by showing that their corresponding closures can be organised as follows:

Lemma 4.10. *We have $\mathbf{cl}_{\mathbf{atom}_{2,3}} = \mathbf{cl}_{\mathbf{atom}_3} \circ \mathbf{cl}_{\mathbf{atom}_2}$*

Proof. We simply write 2 and 3 for \mathbf{atom}_2 and \mathbf{atom}_3 . Since the right-hand sides of \mathbf{atom}_2 are words, by Lemma 3.13 and Lemma 3.10, it suffices to prove $\mathbf{st}_2 \circ \mathbf{st}_3 \subseteq \mathbf{cl}_3 \circ \mathbf{st}_2$. We actually prove $\mathbf{st}_2 \circ \mathbf{st}_3 \subseteq \mathbf{st}_3 \circ \mathbf{st}_2$.

Assume $w \in \text{st}_2(\text{st}_3(L))$ for some language L . Hence, $w = u\alpha v$ for some atom α and words u, v such that $uv \in \text{st}_3(L)$. In turn, w must be equal to $u'v'$ for some words u', v' such that for all atoms β , $u'\beta v' \in L$. By symmetry, we may assume $|u| \leq |u'|$, i.e., $u' = uw$, $v = wv'$ for some word w . In this case, we have $uw\beta v' \in L$ for all β , whence $u\alpha w\beta v' \in \text{st}_2(L)$ for all β , whence $u\alpha v = u\alpha wv' \in \text{st}_3(\text{st}_2(L))$, as required. \square

Putting everything together, we finally obtain completeness of KA_{kat} .

Theorem 4.11. *For all $e, f \in \mathcal{T}_{\text{KAT}}$, $\text{cl}_{\text{kat}}(\llbracket e \rrbracket) = \text{cl}_{\text{kat}}(\llbracket f \rrbracket)$ implies $\text{KA}_{\text{kat}} \vdash e = f$.*

Proof. kat reduces to atom (Lemma 4.9), which reduces to $\text{atom}_{2,3}$ by Lemma 3.14. The latter set reduces to the empty set by Lemma 4.10, Lemma 3.6 and Lemma 3.9. Thus kat reduces to the empty set, and we conclude via completeness of Kleene algebra (Theorem 2.1) and Theorem 3.2. \square

5 Kleene Algebra with Observations

A *Kleene algebra with Observations (KAO)* is a Kleene algebra which also contains a Boolean algebra, but the connection between the Boolean algebra and the Kleene algebra is different than for KAT: instead of having the axiom $\phi \wedge \psi = \phi \cdot \psi$ for all $\phi, \psi \in \mathcal{T}_{\text{BA}}$, we only have $\phi \wedge \psi \leq \phi \cdot \psi$ [15]. This system was introduced to allow for concurrency and tests in a Kleene algebra framework, because associating $\phi \cdot \psi$ and $\phi \wedge \psi$ in a concurrent setting is no longer appropriate: $\phi \wedge \psi$ is one event, where we instantaneously test whether both ϕ and ψ are true, while $\phi \cdot \psi$ performs first the test ϕ , and then ψ , and possibly other things can happen between those tests in another parallel thread. Hence, the behaviour of $\phi \wedge \psi$ should be included in $\phi \cdot \psi$, but they are no longer equivalent. (Note that even if we add the axiom $1 = \top$, in which case we have that $\phi \cdot \psi$ is below both ψ and ϕ , this is not enough to collapse $\phi \cdot \psi$ and $\phi \wedge \psi$, because $\phi \cdot \psi$ need not be an element of the Boolean algebra.)

Algebraically this constitutes a small change, and an ad-hoc completeness proof is in [15]. Here we show how to obtain completeness within our framework. We also show how to add the additional and natural axiom $1 = \top$, which is not present in [15], and thereby emphasise the modular aspect of the approach.

Similar to KAT, we add the additional axioms of KAO to KA as hypotheses. The additional axioms of KAO are the axioms of Boolean algebra and the axioms specifying the interaction between the two algebras. The KAO-terms are the same as the KAT-terms: regular expression over the alphabet $\Sigma + \mathcal{T}_{\text{BA}}$.

Definition 5.1. *We define the set of hypotheses $\text{kao} = \text{bool} \cup \text{glue}'$, where*

$$\text{glue}' = \{\phi \wedge \psi \leq \phi \cdot \psi, \phi \vee \psi = \phi + \psi \mid \phi, \psi \in \mathcal{T}_{\text{BA}}\} \cup \{\perp = 0\}$$

We prove completeness with respect to the closed interpretation under hypotheses: $\text{cl}_{\text{kao}}(\llbracket - \rrbracket)$. As shown below, this also implies completeness for the language model presented in [15]. We take similar steps as for KAT:

1. Reduce **kao** to the simpler set of axioms $\mathbf{contr} = \{\alpha \leq \alpha \cdot \alpha \mid \alpha \in \mathbf{At}\}$, where $\mathbf{At} = 2^\Omega$ is the set of atoms, as in Section 4.
2. Use results from Section 3.1 to reduce **contr** to the empty set.

For the first step, we use the same homomorphism r as for KAT.

Lemma 5.2. *For all $e \leq f \in \mathbf{kao}$, we have $\mathbf{KA}_{\mathbf{contr}} \vdash r(e) \leq r(f)$.*

Lemma 5.3. *The homomorphism r yields a reduction from **kao** to **contr**.*

Proof. Like for Lemma 4.9, we use Lemma 3.3. We show $\mathbf{KA}_{\mathbf{kao}} \vdash \mathbf{contr}$: for $\alpha \in \mathbf{At}$, we have $\mathbf{KA}_{\mathbf{kao}} \vdash \alpha = \alpha \wedge \alpha \leq \alpha \cdot \alpha$. The first and second condition about r are obtained like in the KAT case: the glueing equations for \wedge were not necessary there. The third and last condition was proven in Lemma 5.2. \square

Theorem 5.4. *For all $e, f \in \mathcal{T}_{\mathbf{KAT}}$, $\text{cl}_{\mathbf{kao}}(\llbracket e \rrbracket) = \text{cl}_{\mathbf{kao}}(\llbracket f \rrbracket)$ implies $\mathbf{KA}_{\mathbf{kao}} \vdash e = f$.*

Proof. **kao** reduces to **contr** (Lemma 5.3), which reduces to \emptyset by Lemma 3.6(i), as both α and $\alpha \cdot \alpha$ are words and α is a word of length 1. \square

Note that the semantics defined in [15] actually corresponds to $\text{cl}_{\mathbf{contr}}(\llbracket r(-) \rrbracket)$ rather than $\text{cl}_{\mathbf{kao}}(\llbracket - \rrbracket)$. These semantics are nonetheless equivalent, **kao** reducing to **contr** via r (the proof of Theorem 3.2 actually establishes that when H reduces to H' via r and $\mathbf{KA}_{H'}$ is complete, we have $\text{cl}_H(\llbracket e \rrbracket) = \text{cl}_H(\llbracket f \rrbracket)$ iff $\text{cl}_{H'}(\llbracket r(e) \rrbracket) = \text{cl}_{H'}(\llbracket r(f) \rrbracket)$ for all expressions e, f).

Because we set up KAO in a modular way, we can now easily extend it with the extra axiom $\top = 1$. Combining the proofs that r is a reduction from **kat** to **atom** and from **kao** to **contr**, we can easily see that r is also a reduction from $\mathbf{kao} \cup \{\top = 1\}$ to $\mathbf{contr} \cup \mathbf{atom}_2 \cup \mathbf{atom}_3$. To obtain completeness, it thus suffices to explain how to combine the closures w.r.t. **contr**, **atom**₂, and **atom**₃.

Lemma 5.5. *We have $\text{cl}_2 \circ \text{cl}_{\mathbf{contr}} \subseteq \text{cl}_{\mathbf{contr}} \circ \text{cl}_2$ and $\text{cl}_{\mathbf{contr}} \circ \text{cl}_3 \subseteq \text{cl}_3 \circ \text{cl}_{\mathbf{contr}}$*

Proof. See Appendix D, the second inclusion is non-trivial. \square

Theorem 5.6. *For all $e, f \in \mathcal{T}_{\mathbf{KAT}}$, if $\text{cl}_{\mathbf{kao} \cup \{\top = 1\}}(\llbracket e \rrbracket) = \text{cl}_{\mathbf{kao} \cup \{\top = 1\}}(\llbracket f \rrbracket)$ then $\mathbf{KA}_{\mathbf{kao} \cup \{\top = 1\}} \vdash e = f$.*

Proof. Because we know $\text{cl}_{2,3} = \text{cl}_3 \circ \text{cl}_2$ (Lemma 4.10) and hence $\text{cl}_2 \circ \text{cl}_3 \subseteq \text{cl}_3 \circ \text{cl}_2$, we can use Lemma 3.10 with Lemma 5.5 to deduce $\text{cl}_{\mathbf{contr}, 2, 3} = \text{cl}_3 \circ \text{cl}_{\mathbf{contr}} \circ \text{cl}_2$, and obtain completeness via Lemma 3.9 and the fact that **contr**, 2 and 3 all reduce to the empty set. \square

6 Kleene algebra with positive tests

In KAT, tests are assumed to form a Boolean algebra. Here we study the structure obtained by assuming that they only form a distributive lattice. A *Kleene algebra with positive tests (KAPT)* is a Kleene algebra X containing a lattice L such that the meet of L coincides with the product of X , the join of X coincides with the

sum of X , and all elements of L are below the multiplicative identity of X . (We discuss the variant where we have a bounded lattice at the end, see Remark 6.6). Since the product distributes over sums in X , L must be a distributive lattice. Also note that there might be elements of X below 1 that do not belong to L .

As before, we fix two finite sets Σ and Ω of primitive actions and primitive tests. Then we consider regular expressions over the alphabet $\Sigma + \mathcal{T}_{\text{DL}}$, where \mathcal{T}_{DL} is the set of lattice expressions over Ω : expressions built from elements of Ω and two binary connectives \vee and \wedge .

We write dl for the set of all instances of distributive lattice axioms over \mathcal{T}_{DL} [9], and we set $\text{kapt} \triangleq \text{dl} \cup \text{glue}''$ where

$$\text{glue}'' \triangleq \{\phi \wedge \psi = \phi \cdot \psi, \phi \vee \psi = \phi + \psi \mid \phi, \psi \in \mathcal{T}_{\text{DL}}\} \cup \{\phi \leq 1 \mid \phi \in \mathcal{T}_{\text{DL}}\}$$

Like for Boolean algebras, the free distributive lattice over Ω is finite and can be described easily. An *atom* α is a non-empty subset of Ω , and we write At for the set of such atoms as before. However, while an atom $\{a, b\}$ of Boolean algebra was implicitly interpreted as the term $a \wedge b \wedge \neg c$ (when $\Omega = \{a, b, c\}$), the same atom in the context of distributive lattices is implicitly interpreted as the term $a \wedge b$ —there are no negative literals in distributive lattices. Again similarly to the case of Boolean algebras, the key property for atoms in distributive lattices is the following: for all atoms α and formulas ϕ , we have

$$\alpha \models \phi \Leftrightarrow \text{DL} \vdash \alpha \leq \phi \quad \text{and} \quad \text{DL} \vdash \phi = \bigvee_{\alpha \models \phi} \alpha$$

Like for KAT, such a property makes it possible to reduce kapt to the following set of equations on the alphabet $\Sigma + \text{At}$.

$$\text{atom}' \triangleq \{\alpha \cdot \beta = \alpha \cup \beta \mid \alpha, \beta \in \text{At}\} \cup \{\alpha \leq 1 \mid \alpha \in \text{At}\}$$

(Note that in the right-hand side of the first equation, $\alpha \cup \beta$ is a single atom, whose implicit interpretation is $\alpha \wedge \beta$.)

Lemma 6.1. *There is a reduction from kapt to atom' , witnessed by the homomorphism $r: \mathcal{T}(\Sigma + \mathcal{T}_{\text{DL}}) \rightarrow \mathcal{T}(\Sigma + \text{At})$ defined by*

$$r(x) = \begin{cases} a & x = a \in \Sigma \\ \sum_{\alpha \models \phi} \alpha & x = \phi \in \mathcal{T}_{\text{DL}} \end{cases}$$

As a consequence, in order to get decidability and completeness for KAPT (i.e., kapt), it suffices to reduce atom' to the empty set. Let us number the three kinds of inequations that appear in this set:

$$1 = \{\alpha \cup \beta \leq \alpha \cdot \beta \mid \alpha, \beta \in \text{At}\} \quad 2 = \{\alpha \cdot \beta \leq \alpha \cup \beta \mid \alpha, \beta \in \text{At}\} \quad 3 = \{\alpha \leq 1 \mid \alpha \in \text{At}\}$$

Lemma 3.6(i) gives reductions to the empty set for 1 and 3, but so far we have no reduction for 2. We actually do not know if there is a reduction from 2 to the empty set. Instead, we establish a reduction from 2 together with 3 to 3 alone.

Lemma 6.2. *There is a reduction from 2,3 to 3, witnessed by the homomorphism $r: \mathcal{T}(\Sigma + \text{At}) \rightarrow \mathcal{T}(\Sigma + \text{At})$ defined by*

$$r(x) = \begin{cases} a & x = a \in \Sigma \\ \sum \{\alpha_1 \cdot \dots \cdot \alpha_n \mid \alpha = \bigcup_{i \leq n} \alpha_i, i \neq j \Rightarrow \alpha_i \neq \alpha_j\} & x = \alpha \in \text{At} \end{cases}$$

(Note that the above reduction requires 3 in its target, and cannot be extended directly into a reduction from 1,2,3 to 3: $r(\alpha \cup \beta) \leq r(\alpha \cdot \beta)$ cannot be proved in KA_3 —take $\alpha = \{a\}$, $\beta = \{b\}$, then ba is a term in $r(\alpha \cup \beta)$ which is not provably below $ab = r(\alpha \cdot \beta)$.)

Composed with the existing reduction from 3 to the empty set (Lemma 3.6(i)), we thus have a reduction from 2,3 to the empty set. It remains to combine this reduction to the one from 1 to the empty set (Lemma 3.6(i) again). To this end, we would like to use Lemma 3.9, which simply requires us to prove that the closure $\text{cl}_{\text{atom}'} = \text{cl}_{1,2,3}$ is equal either to $\text{cl}_1 \circ \text{cl}_{2,3}$ or to $\text{cl}_{2,3} \circ \text{cl}_1$. Unfortunately, this is not the case. To see this, suppose we have two atomic tests a and b . For the first option, consider the singleton language $\{ab\}$ (a word consisting of two atoms); we have $ba \in \text{cl}_{1,2,3}(\{ab\})$ (because $(a \wedge b) \in \text{cl}_1(\{ab\})$, and then using cl_2) but $ba \notin \text{cl}_1(\text{cl}_{2,3}(\{ab\}))$. For the second option, consider the singleton language $\{a\}$; we have $(a \wedge b) \in \text{cl}_{1,2,3}(\{a\})$, because $ab \in \text{cl}_3(\{a\})$, but $(a \wedge b) \notin \text{cl}_{2,3}(\text{cl}_1(\{a\}))$ because $\text{cl}_1(\{a\})$ is just $\{a\}$, and $\text{cl}_{2,3}$ does not make it possible to forge conjunctions.

In order to circumvent this difficulty, we use a fourth family of equations:

$$4 = \{\alpha \cup \beta \leq \alpha \mid \alpha, \beta \in \text{At}\}$$

These axioms are immediate consequences of 1 and 3. Therefore, 1,2,3 reduces to 1,2,3,4. Moreover they consist of ‘letter-letter’ inequations, which are covered by Lemma 3.6(i): 4 reduces to the empty set. We shall further prove that $\text{cl}_{1,2,3,4} = \text{cl}_4 \circ \text{cl}_3 \circ \text{cl}_2 \circ \text{cl}_1$ and $\text{cl}_{2,3} = \text{cl}_3 \circ \text{cl}_2$, so that Lemma 3.9 applies to obtain a reduction from 1,2,3,4 to the empty set.

Lemma 6.3. *We have the following inclusions of functions:*

- (i) $\text{st}_1 \circ \text{st}_2 \subseteq \text{st}_2 \circ \text{st}_1 + \text{id}$
- (ii) $\text{st}_1 \circ \text{st}_3 \subseteq \text{st}_3 \circ \text{st}_1 + \text{st}_4$
- (iii) $\text{st}_1 \circ \text{st}_4 \subseteq \text{st}_4 \circ \text{st}_1$
- (iv) $\text{st}_2 \circ \text{st}_3 \subseteq \text{st}_3 \circ \text{st}_2 + \text{st}_3 \circ \text{st}_3$
- (v) $\text{st}_2 \circ \text{st}_4 \subseteq \text{st}_4 \circ \text{st}_2 + \text{st}_4 \circ \text{st}_4 \circ \text{st}_2$
- (vi) $\text{st}_3 \circ \text{st}_4 \subseteq \text{st}_4 \circ \text{st}_3$

Lemma 6.4. *We have $\text{cl}_{1,2,3,4} = \text{cl}_4 \circ \text{cl}_3 \circ \text{cl}_2 \circ \text{cl}_1 = \text{cl}_4 \circ \text{cl}_{2,3} \circ \text{cl}_1$.*

Proof. We use Lemma 3.10 and Lemma 3.13 repeatedly, on (combinations of) the inclusions provided by Lemma 6.3. See the proof in Appendix E. \square

Theorem 6.5. KA_{kapt} *reduces to the empty set, and is complete and decidable.*

Proof. kapt reduces to atom' by Lemma 6.1, which in turn reduces to 1, 2, 3, 4 by Lemma 3.8. The latter is composed of three sets of hypotheses, 1, 4, and 2, 3. All three of them reduce to the empty set: the first two by Lemma 3.6(i), and the third one by Lemma 6.2. These three reductions can be composed together by Lemma 3.9 and Lemma 6.4. \square

Remark 6.6. The case of Kleene algebras containing a *bounded* distributive lattice, with extremal elements \perp and \top coinciding with 0 and 1, may be obtained as follows. Allow the empty atom \emptyset in At (interpreted as \top), and add the inequation $5 = \{1 \leq \emptyset\}$ to atom' . Lemma 6.1 extends easily, and we have a reduction from 5 to the empty set (Lemma 3.6(i)). Therefore it suffices to find how to combine cl_5 with the other closures. We have $\text{cl}_{1,2,3,4,5} = \text{cl}_5 \circ \text{cl}_4 \circ \text{cl}_3 \circ \text{cl}_2 \circ \text{cl}_1$ (see Lemma E.1 in Appendix E), so that we can conclude that the equational theory of Kleene algebras with a bounded distributive lattice is complete and decidable.

7 Related work

There is a range of papers on completeness and decidability of Kleene algebra together with specific forms of hypotheses, starting with [7]. The general case of Kleene algebra with hypotheses, and reductions to prove completeness, has been studied recently in [16,10,25]. The current paper combines and extends these results, and thereby aims to provide a comprehensive overview and a showcase of how to apply these techniques to concrete case studies (KAT, KAO and the new theory KAPT). Below, we discuss each of these recent works in more detail.

Kozen and Mamouras [25] define the canonical language model for KA with a set of hypotheses in terms of rewriting systems, as well as reductions and their role in completeness, and provide reductions for equations of the form $1 = w$ and $a = w$ (cf. Lemma 3.6). Their general results cover completeness results which instantiate to KAT and NetKAT. In fact, the assumptions made in their technical development are tailored towards these cases; for instance, their assumption $\alpha\beta \leq \perp$ (in Assumption 2) rules out KAPT. The current paper focuses more on generality and how to construct reductions in a modular way.

Doumane et al. [10] also define reductions, with an emphasis on (un)decidability. In particular, they cover hypotheses of the form $1 \leq \sum_{a \in S} a$ (cf. Lemma 3.6). A first step towards modularity may also be found in [10, Proposition 3].

Kappé et al. [16] study hypotheses on top of bi-Kleene algebra, where the canonical interpretation is based on pomset languages, and ultimately prove completeness of *concurrent Kleene algebra with observations*; many of the results there apply to the word case as well. We follow this paper for the basic definitions and results for the general theory of Kleene algebra with hypotheses, with a small change in the actual definition of a reduction (Remark 3.5). Compositionality in the sense of Section 3.2 is treated in Kappé's PhD thesis [14]. We extend these results with Lemmas 3.12, 3.13, which simplify the work needed to combine hypotheses. Further, we highlight the word case in this paper (as opposed to the pomset languages in concurrent Kleene algebra), by showcasing several examples.

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A Additional lemmas on closures

Here we prove a few lemmas, which are used for the proofs in the remainder of this appendix.

Lemma A.1. *Let H be a set of hypotheses whose right-hand sides are all words. Then st_H is additive: for every family $(L_i)_{i \in I}$ of languages, $\text{st}_H(\bigcup_{i \in I} L_i) = \bigcup_{i \in I} \text{st}_H(L_i)$. Moreover, in that case, $\text{cl}_H(L) = \bigcup_{n \in \mathbb{N}} \text{st}_H^n(L)$.*

Proof. For additivity, the right-to-left containment follows by monotonicity of st_H . For the left-to-right containment, take a word $w \in \text{st}_H(\bigcup_{i \in I} L_i)$. By definition of st_H , there must be $e \leq f \in H$ such that $w \in u \llbracket e \rrbracket v$ for some words $u, v \in \Sigma^*$ such that $u \llbracket f \rrbracket v \subseteq \bigcup_{i \in I} L_i$. Since f is a word, we know that $u \llbracket f \rrbracket v = \{ufv\}$. Hence, there exists an $i \in I$ such that $ufv \in L_i$. It follows that $u \llbracket e \rrbracket v \subseteq \text{st}_H(L_i)$, and thus $w \in u \llbracket e \rrbracket v \subseteq \text{st}_H(L_i) \subseteq \bigcup_{i \in I} \text{st}_H(L_i)$.

That we can stop the iteration at ω follows: we have $\text{st}_H^{\omega+1}(L) \subseteq \text{st}_H^\omega(L)$ by additivity. \square

Lemma A.2. *For all sets of hypotheses H and languages L, K , we have*

$$\text{cl}_H(L) \cdot K \subseteq \text{cl}_H(L \cdot K) \quad L \cdot \text{cl}_H(K) \subseteq \text{cl}_H(L \cdot K) \quad \text{cl}_H(L) \cdot \text{cl}_H(K) \subseteq \text{cl}_H(L \cdot K)$$

Proof. The first two inclusions are symmetrical, and hold because the rules defining cl_H are contextual. For the third inclusion, we use the first two in order to deduce $\text{cl}_H(L) \cdot \text{cl}_H(K) \subseteq \text{cl}_H(L \cdot \text{cl}_H(K)) \subseteq \text{cl}_H(\text{cl}_H(L \cdot K)) = \text{cl}_H(L \cdot K)$. \square

Lemma A.3. *For all sets of hypotheses H and languages L , we have*

$$\text{cl}_H(L)^* \subseteq \text{cl}_H(L^*)$$

Proof. It suffices to show $1 + L \cdot \text{cl}_H(L^*) \subseteq \text{cl}_H(L^*)$, which follows easily from Lemma A.2. \square

B Proofs for Section 3 (Reductions)

For the proof of Lemma 3.3, we need the following definition.

For a map $r: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Gamma)$, define $\dot{r}: \mathcal{P}(\Sigma^*) \rightarrow \mathcal{P}(\Gamma^*)$ by

$$\dot{r}(L) = \bigcup_{w \in L} \llbracket r(w) \rrbracket$$

(Where words w are seen as regular expressions when fed to r .) We have the following property:

Lemma B.1. *Let $r: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Gamma)$ be a homomorphism. Then we have:*

1. $\dot{r}(\llbracket e \rrbracket) = \llbracket r(e) \rrbracket$ for all $e \in \mathcal{T}(\Sigma)$, and
2. $\dot{r}(L \cdot K) = \dot{r}(L) \cdot \dot{r}(K)$ for all $L, K \subseteq \Sigma^*$.

Now we can prove Lemma 3.3.

Lemma 3.3. *Assume $\Gamma \subseteq \Sigma$ and let H, H' be sets of hypotheses over Σ and Γ respectively, such that $\text{KA}_H \vdash H'$. If there exists a homomorphism $r: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Gamma)$ such that:*

1. *For all $a \in \Gamma$, we have $\text{KA} \vdash a \leq r(a)$.*
2. *For all $a \in \Sigma$, we have $\text{KA}_H \vdash a = r(a)$.*
3. *For all $e \leq f \in H$, we have $\text{KA}_{H'} \vdash r(e) \leq r(f)$.*

then H reduces to H' .

Proof. First, we must prove $\text{KA}_H \vdash e = r(e)$ for all $e \in \mathcal{T}(\Sigma)$, which follows by induction on e , using that r is a homomorphism and Item 2.

Similarly, we get $\text{KA} \vdash e \leq r(e)$ for all $e \in \mathcal{T}(\Gamma)$ from Item 1. By using soundness of KA and this latter result with words in Γ^* , we obtain that for all languages $L \subseteq \Gamma^*$,

$$L \subseteq \dot{r}(L) \tag{1}$$

Secondly, we must prove $\text{cl}_H(\llbracket e \rrbracket) \cap \Gamma^* = \text{cl}_{H'}(\llbracket r(e) \rrbracket)$. For the right-to-left inclusion, the right-hand side is contained in Γ^* by definition, and we have

$$\begin{aligned} \text{cl}_{H'}(\llbracket r(e) \rrbracket) &\subseteq \text{cl}_H(\llbracket r(e) \rrbracket) && (\text{KA}_H \vdash H' \text{ and Lemma 2.4}) \\ &= \text{cl}_H(\llbracket e \rrbracket) && (\text{KA}_H \vdash e = r(e) \text{ and soundness (Theorem 2.3)}) \end{aligned}$$

For the left-to-right inclusion, we first prove that for all languages $L \subseteq \Sigma^*$,

$$\dot{r}(\text{cl}_H(L)) \subseteq \text{cl}_{H'}(\dot{r}(L)) \tag{2}$$

We do so by transfinite induction proving that for all ordinals κ ,

$$\dot{r}(\text{st}_H^\kappa(L)) \subseteq \text{cl}_{H'}(\dot{r}(L)) \tag{3}$$

The base case is trivial. The limit case follows since \dot{r} is defined element-wise.

For the induction step, suppose (3) holds for some κ ; we have to prove $\dot{r}(\text{st}_H(\text{st}_H^\kappa(L))) \subseteq \text{cl}_{H'}(\dot{r}(L))$. Let u, v, e, f be such that $e \leq f \in H$ and $u \llbracket f \rrbracket v \subseteq \text{st}_H^\kappa(L)$. Then we need to show that $\dot{r}(u \llbracket e \rrbracket v) \subseteq \text{cl}_{H'}(\dot{r}(L))$. Since, $u \llbracket f \rrbracket v \subseteq \text{st}_H^\kappa(L)$, the induction hypotheses entails $\dot{r}(u \llbracket f \rrbracket v) \subseteq \text{cl}_{H'}(\dot{r}(L))$. We argue as follows:

$$\begin{aligned} \dot{r}(u \llbracket e \rrbracket v) &\subseteq \text{cl}_{H'}(\dot{r}(u \llbracket e \rrbracket v)) \\ &= \text{cl}_{H'}(\llbracket r(u) \cdot r(e) \cdot r(v) \rrbracket) && (\text{Lemma B.1}) \\ &\subseteq \text{cl}_{H'}(\llbracket r(u) \cdot r(f) \cdot r(v) \rrbracket) && (\star) \\ &= \text{cl}_{H'}(\dot{r}(u \llbracket f \rrbracket v)) && (\text{Lemma B.1}) \\ &\subseteq \text{cl}_{H'}(\text{cl}_{H'}(\dot{r}(L))) && (\text{IH}) \\ &= \text{cl}_{H'}(\dot{r}(L)) \end{aligned}$$

Where we deduce the intermediate inclusion (\star) from soundness of $\mathbf{KA}_{H'}$ and a derivation $\mathbf{KA}_{H'} \vdash r(u) \cdot r(e) \cdot r(v) \leq r(u) \cdot r(f) \cdot r(v)$ obtained from monotonicity and Item 3. This concludes the proof of (3), and thus (2).

Putting everything together, we have, for all expressions $e \in \mathcal{T}(\Sigma)$:

$$\begin{aligned} \text{cl}_H(\llbracket e \rrbracket) \cap \Gamma^* &\subseteq \dot{r}(\text{cl}_H(\llbracket e \rrbracket) \cap \Gamma^*) && \text{(by (1))} \\ &\subseteq \dot{r}(\text{cl}_H(\llbracket e \rrbracket)) \\ &\subseteq \text{cl}_{H'}(\dot{r}(\llbracket e \rrbracket)) && \text{(by (2))} \\ &= \text{cl}_{H'}(\llbracket r(e) \rrbracket) && \text{(Lemma B.1)} \end{aligned}$$

which concludes the proof. \square

Lemma 3.6. *Each of the following sets of hypotheses reduce to the empty set (of hypotheses over Σ).*

- (i) $\{u_i \leq w_i \mid i \in I\}$ with $u_i, w_i \in \Sigma^*$ and $|u_i| \leq 1$ for all $i \in I$
- (ii) $\{1 \leq \sum_{a \in S_i} a \mid i \in I\}$ with each $S_i \subseteq \Sigma$ finite
- (iii) $\{e \leq 0\}$ for $e \in \mathcal{T}(\Sigma)$
- (iv) $\{ea \leq a\}$ and $\{ae \leq a\}$ for $a \in \Sigma$, $e \in \mathcal{T}(\Sigma \setminus \{a\})$

Proof. (i) This is [25, Theorem 2]. (The result mentions equations, but in the proof only the relevant inequations are used.)

(ii) This is [10, Proposition 6].

(iii) This is basically due to [7], but since it is phrased differently there we include a proof.

Define $r: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Sigma)$ by $r(f) = f + \Sigma^* \cdot e \cdot \Sigma^*$ for $f \in \mathcal{T}(\Sigma)$, where Σ^* is seen as the expression $(\sum_{a \in \Sigma} a)^*$. We claim r witnesses a reduction. First, we have $\mathbf{KA}_H \vdash f \leq r(f)$ trivially, and $\mathbf{KA}_H \vdash r(f) \leq f$ follows because $e \leq 0 \in H$.

Second, we have to prove that $\llbracket r(f) \rrbracket = \text{cl}_H(\llbracket f \rrbracket)$. It is easy to check that $\text{st}_H \circ \text{st}_H = \text{st}_H$ for this H , so by the iterative characterisation of cl_H we get $\text{cl}_H(L) = L \cup \text{st}_H(L)$ for any $L \subseteq \Sigma^*$. Now,

$$\begin{aligned} \text{cl}_H(L) &= L \cup \text{st}_H(L) \\ &= L \cup \bigcup \{u \llbracket e \rrbracket v \mid u \llbracket 0 \rrbracket v \subseteq L\} \\ &= L \cup \{u \llbracket e \rrbracket v \mid u, v \in \Sigma^*\} \end{aligned}$$

for all $L \subseteq \Sigma^*$. Therefore,

$$\text{cl}_H(\llbracket f \rrbracket) = \llbracket f \rrbracket \cup \{u \llbracket e \rrbracket v \mid u, v \in \Sigma^*\} = \llbracket f + \Sigma^* e \Sigma^* \rrbracket = \llbracket r(f) \rrbracket$$

for all $f \in \mathcal{T}(\Sigma)$.

- (iv) Hypotheses of a similar form are studied in the setting of Kleene algebra with tests in [12]. Consider the case $\{ea \leq a\}$. Define r as the unique homomorphism satisfying $r(a) = e^*a$, and $r(b) = b$ for all $b \in \Sigma$ with $b \neq a$. We use Lemma 3.3 to show that this witnesses a reduction.

The first two conditions are trivial for letters $b \in \Sigma$ with $a \neq b$, since r is the identity on those letters. For a , we have $\text{KA} \vdash a \leq e^*a = r(a)$ easily from the KA axioms. Further, we get $\text{KA}_H \vdash r(a) = e^*a \leq a$ by $ea \leq a$ and the (left) induction axiom for Kleene star in KA . Finally, we have to prove that $\text{KA} \vdash r(ea) \leq r(a)$. We have $r(a) = e^*a$, and since e does not contain a , we have $r(e) = e$. Therefore, $r(ea) = r(e)r(a) = ee^*a$, and the required inequality follows since $\text{KA} \vdash ee^* \leq e^*$. \square

Lemma 3.7. *Let H_1, H_2 and H_3 be sets of hypotheses. If H_1 reduces to H_2 and H_2 reduces to H_3 then H_1 reduces to H_3 .*

Proof. We have $\text{KA}_{H_1} \vdash H_2$ and $\text{KA}_{H_2} \vdash H_3$, and therefore $\text{KA}_{H_1} \vdash H_3$. Let $r_1: \mathcal{T}(\Sigma_1) \rightarrow \mathcal{T}(\Sigma_2)$ and $r_2: \mathcal{T}(\Sigma_2) \rightarrow \mathcal{T}(\Sigma_3)$ be witnesses of the reductions from H_1 to H_2 and from H_2 to H_3 respectively. We show that $r_2 \circ r_1$ is a reduction from H_1 to H_3 .

Since $\text{KA}_{H_1} \vdash e = r_1(e)$ for all $e \in \mathcal{T}(\Sigma_1)$, $\text{KA}_{H_2} \vdash e = r_2(e)$ for all $e \in \mathcal{T}(\Sigma_2)$ and $\text{KA}_{H_1} \vdash H_2$, we get $\text{KA}_{H_1} \vdash e = r_1(e) = r_2(r_1(e))$.

Finally, we have:

$$\text{cl}_1(\llbracket e \rrbracket) \cap \Sigma_3^* = (\text{cl}_1(\llbracket e \rrbracket) \cap \Sigma_2^*) \cap \Sigma_3^* = \text{cl}_2(\llbracket r_1(e) \rrbracket) \cap \Sigma_3^* = \text{cl}_3(\llbracket r_2(r_1(e)) \rrbracket)$$

using that $\Sigma_3 \subseteq \Sigma_2$, r_1 is a reduction, and r_2 is a reduction, respectively. \square

Lemma 3.8. *Let H_1, H_2 be sets of hypotheses over a common alphabet. If $\text{KA}_{H_1} \vdash H_2$ and $\text{KA}_{H_2} \vdash H_1$ then H_1 and H_2 reduce to each other.*

Proof. The identity map on terms fulfills the requirements; in particular, we have $\text{cl}_1(\llbracket e \rrbracket) = \text{cl}_2(\llbracket e \rrbracket)$ for all e by Lemma 2.4. \square

Lemma 3.9. *Let H_1, \dots, H_n, H be sets of hypotheses over a common alphabet Σ , with $n \geq 1$. If H_i reduces to H for all i , and $\text{cl}_{1\dots n} = \text{cl}_n \circ \dots \circ \text{cl}_1$, then $\bigcup_{i \leq n} H_i$ reduces to H .*

Proof. Since $\text{KA}_{H_i} \vdash e \leq f$ for each i and $e \leq f \in H$, and there is at least one such H_i , we have $\text{KA}_{\bigcup_{i \leq n} H_i} \vdash e \leq f$. Now, let r_1, \dots, r_n be witnesses for the reductions from H_1, \dots, H_n to H . We show that $r = r_n \circ \dots \circ r_1$ is a reduction under these assumptions.

Further, it is easy to prove $\text{KA}_{\bigcup_{i \leq k} H_i} \vdash e = r_k \circ \dots \circ r_1(e)$ for all $k \leq n$, by induction on k , using that each r_i is a reduction.

Finally, for the second requirement on $r_k \circ \dots \circ r_1$ to be a reduction, observe that the restriction to the codomain alphabet is void since it coincides with the alphabet in the domain, that is, we have $\text{cl}_i(\llbracket e \rrbracket) = \text{cl}_H(\llbracket r_i(e) \rrbracket)$ for each i . This implies that, for all $k \leq n$, we have $\text{cl}_k \circ \dots \circ \text{cl}_1(\llbracket e \rrbracket) = \text{cl}_H(\llbracket r_k \circ \dots \circ r_1(e) \rrbracket)$

by induction on k . Indeed, the base case ($k = 1$) holds by assumption. For the inductive case, if it holds for some k with $k < n$, then

$$\begin{aligned}
 & \text{cl}_{k+1} \circ \text{cl}_k \circ \dots \circ \text{cl}_1(\llbracket e \rrbracket) \\
 &= \text{cl}_{k+1}(\text{cl}_H(\llbracket r_k \circ \dots \circ r_1(e) \rrbracket)) && \text{(IH)} \\
 &= \text{cl}_{k+1}(\llbracket r_k \circ \dots \circ r_1(e) \rrbracket) \\
 &= \text{cl}_H(\llbracket r_{k+1} \circ r_k \circ \dots \circ r_1(e) \rrbracket) && (r_{k+1} \text{ is a reduction})
 \end{aligned}$$

using in the second step that $\text{cl}_H \subseteq \text{cl}_{k+1}$, since $\text{KA}_{H_{k+1}} \vdash H$ (Lemma 2.4) and thus $\text{cl}_{k+1} \circ \text{cl}_H \subseteq \text{cl}_{k+1} \circ \text{cl}_{k+1} = \text{cl}_{k+1}$ (the converse inclusion $\text{cl}_{k+1} \subseteq \text{cl}_{k+1} \circ \text{cl}_H$ holds by definition of closure). This concludes the inductive proof. As a consequence, we have $\text{cl}_{1\dots n}(\llbracket e \rrbracket) = \text{cl}_n \circ \dots \circ \text{cl}_1(\llbracket e \rrbracket) = \llbracket r_n \circ \dots \circ r_1(e) \rrbracket$ as needed. \square

Lemma 3.10. *Let H_1, \dots, H_n be sets of hypotheses, such that $\text{cl}_i \circ \text{cl}_j \subseteq \text{cl}_j \circ \text{cl}_i$ for all i, j with $i < j$. Then $\text{cl}_{1\dots n} = \text{cl}_n \circ \dots \circ \text{cl}_1$.*

Proof. The right to left inclusion is trivial. For the other inclusion, we use the fixed point characterisation of $\text{cl}_{1\dots n}$: given a language L , it suffices to prove that $\text{cl}_n \circ \dots \circ \text{cl}_1(L)$ is a pre-fixed point of $\text{st}'_{1\dots n, L}$ (Definition 2.2), i.e.,

$$\text{st}_{1\dots n}(\text{cl}_n \circ \dots \circ \text{cl}_1(L)) \cup L \subseteq \text{cl}_n \circ \dots \circ \text{cl}_1(L).$$

The inclusion $L \subseteq \text{cl}_n \circ \dots \circ \text{cl}_1(L)$ holds since each cl_i is a closure. For the other part, we first observe that

$$\begin{aligned}
 & \text{st}_{1\dots n}(\text{cl}_n \circ \dots \circ \text{cl}_1(L)) \\
 &= \text{st}_1(\text{cl}_n \circ \dots \circ \text{cl}_1(L)) \cup \dots \cup \text{st}_n(\text{cl}_n \circ \dots \circ \text{cl}_1(L))
 \end{aligned}$$

By assumption, for all $i < j$ we have $\text{cl}_i \circ \text{cl}_j \subseteq \text{cl}_j \circ \text{cl}_i$. Hence, we can derive

$$\begin{aligned}
 & \text{st}_i \circ \text{cl}_n \circ \dots \circ \text{cl}_1(L) \\
 & \subseteq \text{cl}_i \circ \text{cl}_n \circ \dots \circ \text{cl}_1(L) \\
 & \subseteq \text{cl}_1 \circ \dots \circ \text{cl}_i \circ \text{cl}_i \circ \dots \circ \text{cl}_n(L) \\
 & = \text{cl}_n \circ \dots \circ \text{cl}_1(L)
 \end{aligned}$$

using that $\text{cl}_i \circ \text{cl}_i \subseteq \text{cl}_i$ in the last step. \square

Lemma 3.12. *For all H_1, H_2 , if $\text{st}_1 \circ \text{cl}_2 \subseteq \text{cl}_2 \circ \text{cl}_1$, then $\text{cl}_1 \circ \text{cl}_2 \subseteq \text{cl}_2 \circ \text{cl}_1$.*

Proof. It suffices to prove that

$$\text{st}_1(\text{cl}_2(\text{cl}_1(L))) \cup \text{cl}_2(L) \subseteq \text{cl}_2(\text{cl}_1(L))$$

by the fixed point characterisation of cl_1 . Indeed, we have $\text{st}_1(\text{cl}_2(\text{cl}_1(L))) \subseteq \text{cl}_2(\text{cl}_1(\text{cl}_1(L))) = \text{cl}_2(\text{cl}_1(L))$ using the assumption. Further, we have $\text{cl}_2(L) \subseteq \text{cl}_2(\text{cl}_1(L))$ by definition of closure. \square

Lemma 3.13. *Let H_1, H_2 be sets of hypotheses such that the right-hand sides of inequations in H_1 are all words. If $st_1 \circ st_2 \subseteq cl_2 \circ st_1^-$ then $cl_1 \circ cl_2 \subseteq cl_2 \circ cl_1$.*

Proof. We first prove by transfinite induction on α that for all languages L ,

$$st_1(st_2^\alpha(L)) \subseteq cl_2(st_1^-(L)) \quad (\dagger)$$

For the successor case, $\alpha + 1$, we get

$$\begin{aligned} st_1(st_2^{\alpha+1}(L)) &= st_1(st_2(st_2^\alpha(L))) \\ &\subseteq cl_2(st_1^-(st_2^\alpha(L))) && \text{(assumption)} \\ &\subseteq cl_2(cl_2(st_1^-(L)) \cup st_2^\alpha(L)) && \text{(IH)} \\ &\subseteq cl_2(cl_2(st_1^-(L)) \cup cl_2(L)) \\ &\subseteq cl_2(cl_2(st_1^-(L))) \\ &= cl_2(st_1^-(L)) \end{aligned}$$

For the limit case, λ , we get

$$\begin{aligned} st_1(st_2^\lambda(L)) &= st_1\left(\bigcup_{\alpha < \lambda} st_2^\alpha(L)\right) \\ &= \bigcup_{\alpha < \lambda} st_1(st_2^\alpha(L)) && \text{(by Lemma A.1)} \\ &\subseteq cl_2(st_1^-(L)) && \text{(IH)} \end{aligned}$$

This concludes the proof of (\dagger) . Then, for every language L , we have:

$$\begin{aligned} st_1(cl_2(L)) &= st_1\left(\bigcup_{\alpha} st_2^\alpha(L)\right) \\ &= \bigcup_{\alpha} st_1(st_2^\alpha(L)) && \text{(by Lemma A.1)} \\ &\subseteq cl_2(st_1^-(L)) && \text{(by } (\dagger)) \end{aligned}$$

Hence $st_1 \circ cl_2 \subseteq cl_2 \circ st_1^- \subseteq cl_2 \circ cl_1$, and we conclude with Lemma 3.12. \square

Lemma B.2. *Let $H_1 = \{e \leq 0\}$ for some term e and let H_2 be an arbitrary set of hypotheses. Then $cl_1 \circ cl_2 \subseteq cl_2 \circ cl_1$.*

Proof. By Lemma 3.12, it suffices to show that

$$st_1 \circ cl_2 \subseteq cl_2 \circ cl_1. \quad (4)$$

Observe that st_1 is constant:

$$\begin{aligned} \text{st}_1(L) &= \bigcup \{u \llbracket e \rrbracket v \mid u, v \in \Sigma^*, u \llbracket 0 \rrbracket v \subseteq L\} \\ &= \bigcup \{u \llbracket e \rrbracket v \mid u, v \in \Sigma^*\} \\ &= \{uvw \mid u, v \in \Sigma^*, w \in \llbracket e \rrbracket\} \end{aligned}$$

using in the second step that $u \llbracket 0 \rrbracket v = u\emptyset v = \emptyset$. Thus, for any language L , we have $\text{st}_1(\text{cl}_2(L)) = \text{st}_1(L) \subseteq \text{cl}_1(L) \subseteq \text{cl}_2(\text{cl}_1(L))$, using that st_1 is constant in the first equality. \square

Lemma 3.14. *For any set of hypotheses H and any term e , there is a reduction from $H \cup \{e \leq 0\}$ to H .*

Proof. Let $H_1 = \{e \leq 0\}$. First, $\text{KA}_{H \cup H_1} \vdash H$ is trivial. Now let $r: \mathcal{T}(\Sigma) \rightarrow \mathcal{T}(\Sigma)$ be a reduction from H_1 to \emptyset , which exists by Lemma 3.6. We claim that r is also a reduction from $H \cup \{e \leq 0\}$ to H . For any term f , $\text{KA}_{H \cup H_1} \vdash f = r(f)$ holds since $\text{KA}_{H_1} \vdash f = r(f)$.

For the second requirement on r , we need to show that $\text{cl}_{H \cup H_1}(\llbracket f \rrbracket) = \text{cl}_H(\llbracket r(f) \rrbracket)$; note that the alphabet restriction is dropped as the domain and codomain of r are equal. By Lemma B.2 and Lemma 3.10 we have $\text{cl}_{H \cup H_1} = \text{cl}_H \circ \text{cl}_1$. Hence, $\text{cl}_{H \cup H_1}(\llbracket f \rrbracket) = \text{cl}_H(\text{cl}_1(\llbracket f \rrbracket)) = \text{cl}_H(\llbracket f \rrbracket)$ using that r is a reduction from H_1 to \emptyset . \square

C Proofs for Section 4 (KAT)

Lemma 4.3. *Let L be a language of guarded strings. We have*

$$\begin{aligned} &\phi_0 a_0 \cdots \phi_{n-1} a_{n-1} \phi_n \in \text{cl}_{\text{kat}}(L) \\ \Leftrightarrow &\quad \forall (\alpha_i)_{i \leq n}, \quad (\forall i \leq n, \alpha_i \models \overline{\phi_i}) \Rightarrow \alpha_0 a_0 \cdots \alpha_{n-1} a_{n-1} \alpha_n \in L \end{aligned}$$

Proof. The proof for the left to right direction proceeds by induction on the closure (seen as the least fixed point of $\text{st}'_{\text{kat}, L}$ —Definition 2.2). In the base case, where the word already belongs to L , and thus is already a guarded string, the ϕ_i must all be single atoms. Since $\alpha \models \beta$ iff $\alpha = \beta$ for all atoms α, β , the condition is immediately satisfied.

Otherwise, we have $\phi_0 a_0 \cdots \phi_{n-1} a_{n-1} \phi_n \in u \llbracket e \rrbracket v$ for $u, v \in (\Sigma + \mathcal{T}_{\text{BA}})^*$ with $u \llbracket f \rrbracket v \subseteq \text{cl}_{\text{kat}}(L)$ and $e \leq f \in \text{kat}$, and we know that the induction hypothesis applies to all words in $u \llbracket f \rrbracket v$. We proceed by a case distinction on the hypothesis in $e \leq f \in \text{kat}$.

- $e \leq f \in \text{bool}$: in this case $e = \phi$ and $f = \psi$ for some formulas ϕ, ψ such that $\text{BA} \vdash \phi = \psi$. Therefore, we have $\alpha \models \phi$ iff $\alpha \models \psi$ for all atoms α , so that the induction hypothesis on the unique word of $u\psi v$ is equivalent to the property we have to prove about the unique word of $u\phi v$.

- $e = \phi \wedge \psi$ and $f = \phi \cdot \psi$: there must be j such that $\phi_j = \phi(\phi \wedge \psi)\psi$ for $\phi, \psi \in \mathcal{T}_{\text{BA}}^*$. The word of $u\phi\psi v$ has the same decomposition as $u(\phi \wedge \psi)v$, except that ϕ_j is replaced by $\phi\phi\psi\psi$. Since $\text{BA} \vdash \overline{\phi\phi\psi\psi} = \overline{\phi(\phi \wedge \psi)\psi}$, the induction hypothesis immediately applies.
- The case for $e = \phi \cdot \psi$ and $f = \phi \wedge \psi$ is handled similarly.
- $e = \phi \vee \psi$ and $f = \phi + \psi$: there must be j such that $\phi_j = \phi(\phi \vee \psi)\psi$ for $\phi, \psi \in \mathcal{T}_{\text{BA}}^*$. We have $\llbracket f \rrbracket = \{\phi, \psi\}$ and the two words of $u\llbracket f \rrbracket v$ have the same decomposition as $u(\phi \vee \psi)v$, except that ϕ_j is replaced by $\phi\phi\psi$ or $\phi\psi\psi$. We use the induction hypothesis and the fact that $\alpha \models \phi \vee \psi$ iff $\alpha \models \phi$ or $\alpha \models \psi$.
- $e = \phi + \psi$ and $f = \phi \vee \psi$: there must be j such that w.l.o.g. $\phi_j = \phi(\phi)\psi$ for $\phi, \psi \in \mathcal{T}_{\text{BA}}^*$. We have $\llbracket f \rrbracket = \{\phi \vee \psi\}$ and the word in $u\llbracket f \rrbracket v$ has the same decomposition as $u(\phi)v$, except that ϕ_j is replaced by $\phi(\phi \vee \psi)\psi$. We use the induction hypothesis and the fact that if $\alpha \models \phi$ then $\alpha \models \phi \vee \psi$.
- $e = \perp$ and $f = 0$: there must be j such that \perp belongs to ϕ_j , so that there are no atoms α such that $\alpha \models \phi_j$: the condition is trivially satisfied.
- $e = 0$ and $f = \perp$: in this case $u\llbracket e \rrbracket v$ is empty so this case is trivially satisfied.
- $e = \top$ and $f = 1$: there is j, ϕ, ψ such that $\phi_j = \phi\top\psi$. The word of uv decomposes like $u\top v$, except that ϕ_j is replaced by $\phi\psi$. Since $\text{BA} \vdash \overline{\phi\psi} = \overline{\phi\top\psi}$, the induction hypothesis suffices to conclude.
- The case where $e = 1$ and $f = \top$ is handled similarly.

We now prove the right-to-left implication. From the assumptions we obtain

$$\left[\sum_{\alpha \models \overline{\phi_0}} \alpha \right] a_0 \dots \left[\sum_{\alpha \models \overline{\phi_{n-1}}} \alpha \right] a_{n-1} \left[\sum_{\alpha \models \overline{\phi_n}} \alpha \right] \subseteq L$$

We deduce using the glueing inequations for \vee and \perp that

$$\left(\bigvee_{\alpha \models \overline{\phi_0}} \alpha \right) a_0 \dots \left(\bigvee_{\alpha \models \overline{\phi_{n-1}}} \alpha \right) a_{n-1} \left(\bigvee_{\alpha \models \overline{\phi_n}} \alpha \right) \in \text{cl}_{\text{kat}}(L)$$

Since $\text{BA} \vdash \phi = \bigvee_{\alpha \models \phi} \alpha$ for all ϕ , the inequations from `bool` yield

$$\overline{\phi_0} a_0 \dots \overline{\phi_{n-1}} a_{n-1} \overline{\phi_n} \in \text{cl}_{\text{kat}}(L)$$

We conclude with the glueing inequations for \wedge and \top , which give

$$\phi_0 a_0 \dots \phi_{n-1} a_{n-1} \phi_n \in \text{cl}_{\text{kat}}(L) \quad \square$$

Now we turn to proving Lemma 4.4, yielding Corollary 4.6. We first prove two lemmas about `kat`-closures of guarded string languages.

Lemma C.1. *For all guarded string languages L, K , $\text{cl}_{\text{kat}}(L \cdot K) \subseteq \text{cl}_{\text{kat}}(L \diamond K)$.*

Proof. We show that $L \cdot K \subseteq \text{cl}_{\text{kat}}(L \diamond K)$, which is sufficient. Take $w \in L \cdot K$. Hence, $w = xy$ with $x \in L$ and $y \in K$. Because x and y are guarded strings, we know both of them begin and end in an atom. So we can write $x = x'\alpha$ and $y = \beta y'$ for some $\alpha, \beta \in \text{At}$.

- If $\alpha = \beta$, then $x\alpha y \in L \diamond K$. In that case, we have $xy = x'\alpha\alpha y'$, but then, since x and y are guarded strings and α is the only atom s.t. $\alpha \models \alpha \wedge \alpha$, we get $xy \in \text{cl}_{\text{kat}}(L \diamond K)$ by Lemma 4.3.
- Otherwise, if $\alpha \neq \beta$, then *a priori* $xy \notin L \diamond K$. However, since there are no atoms γ s.t. $\gamma \models \alpha \wedge \beta$, we have $xy \in \text{cl}_{\text{kat}}(L \diamond K)$ by Lemma 4.3. \square

Given a guarded string language L , let us denote by L^\diamond the language L iterated w.r.t. the coalesced product \diamond : $L^\diamond = \bigcup_n L^{\diamond n}$, where $L^{\diamond 0} = \{\alpha \mid \alpha \in \text{At}\}$ and $L^{\diamond n+1} = L \diamond L^{\diamond n}$.

Lemma C.2. *For all guarded string languages L , $\text{cl}_{\text{kat}}(L^*) \subseteq \text{cl}_{\text{kat}}(L^\diamond)$.*

Proof. It suffices to prove $L^* \subseteq \text{cl}_{\text{kat}}(L^\diamond)$, which in turn follows once we prove $\varepsilon \in \text{cl}_{\text{kat}}(L^\diamond)$ and $L\text{cl}_{\text{kat}}(L^\diamond) \subseteq \text{cl}_{\text{kat}}(L^\diamond)$. The former follows from the fact that $\text{KA}_{\text{kat}} \vdash 1 \leq \sum_\alpha \alpha$, since L^\diamond contains $L^{\diamond 0}$ and thus all atoms α . For the latter, we have

$$\begin{aligned} L \cdot \text{cl}_{\text{kat}}(L^\diamond) &\subseteq \text{cl}_{\text{kat}}(L \cdot L^\diamond) && \text{(Lemma A.2)} \\ &\subseteq \text{cl}_{\text{kat}}(L \diamond L^\diamond) && \text{(Lemma C.1)} \\ &\subseteq \text{cl}_{\text{kat}}(L^\diamond) && \square \end{aligned}$$

We can finally prove Lemma 4.4:

Lemma 4.4. *For all KAT expressions e , $\text{cl}_{\text{kat}}(\llbracket e \rrbracket) = \text{cl}_{\text{kat}}(\mathcal{G}(e))$.*

Proof. It suffices to prove that $\llbracket e \rrbracket \subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e))$ and $\mathcal{G}(e) \subseteq \text{cl}_{\text{kat}}(\llbracket e \rrbracket)$. We prove both results via induction on e .

We start with $\llbracket e \rrbracket \subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e))$.

- 0: trivial as $\llbracket 0 \rrbracket = \emptyset$.
- 1: $\llbracket 1 \rrbracket = \{\varepsilon\}$ and $\mathcal{G}(1) = \{\alpha \mid \alpha \in \text{At}\}$, therefore $\varepsilon \in \text{cl}_{\text{kat}}(\mathcal{G}(1))$ by Lemma 4.3.
- $a \in \Sigma$: $\llbracket a \rrbracket = \{a\}$ and $\mathcal{G}(a) = \{\alpha a \beta \mid \alpha, \beta \in \text{At}\}$, therefore $a \in \text{cl}_{\text{kat}}(\mathcal{G}(a))$ by Lemma 4.3.
- $\phi \in \mathcal{T}_{\text{BA}}$: $\llbracket \phi \rrbracket = \{\phi\}$ and $\mathcal{G}(\phi) = \{\alpha \mid \alpha \models \phi\}$, hence $\phi \in \text{cl}_{\text{kat}}(\mathcal{G}(\phi))$ by Lemma 4.3.
- $e + f$: we have

$$\begin{aligned} \llbracket e + f \rrbracket &= \llbracket e \rrbracket \cup \llbracket f \rrbracket \subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e)) \cup \text{cl}_{\text{kat}}(\mathcal{G}(f)) && \text{(IH)} \\ &\subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e) \cup \mathcal{G}(f)) = \text{cl}_{\text{kat}}(\mathcal{G}(e + f)) && (\text{cl}_{\text{kat}} \text{ monotone}) \end{aligned}$$

- $e \cdot f$: we have

$$\begin{aligned} \llbracket e \cdot f \rrbracket &= \llbracket e \rrbracket \cdot \llbracket f \rrbracket \subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e)) \cdot \text{cl}_{\text{kat}}(\mathcal{G}(f)) && \text{(IH)} \\ &\subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e) \cdot \mathcal{G}(f)) && \text{(Lemma A.2)} \\ &\subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e) \diamond \mathcal{G}(f)) = \text{cl}_{\text{kat}}(\mathcal{G}(e \cdot f)) && \text{(Lemma C.1)} \end{aligned}$$

– e^* : we have

$$\begin{aligned} \llbracket e^* \rrbracket &= \llbracket e \rrbracket^* \subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e))^* && \text{(IH)} \\ &\subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e)^*) && \text{(Lemma A.3)} \\ &\subseteq \text{cl}_{\text{kat}}(\mathcal{G}(e)^\diamond) = \text{cl}_{\text{kat}}(\mathcal{G}(e^*)) && \text{(Lemma C.2)} \end{aligned}$$

Next we prove $\mathcal{G}(e) \subseteq \text{cl}_{\text{kat}}(\llbracket e \rrbracket)$:

- 0: trivial as $\mathcal{G}(0) = \emptyset$.
- 1: $\mathcal{G}(1) = \{\alpha \mid \alpha \in \text{At}\}$; for all α , $\text{KA}_{\text{kat}} \vdash \alpha \leq 1$, whence $\alpha \in \text{cl}_{\text{kat}}\{\varepsilon\} = \text{cl}_{\text{kat}}\llbracket 1 \rrbracket$.
- $a \in \Sigma$: $\mathcal{G}(a) = \{\alpha a \beta \mid \alpha, \beta \in \text{At}\}$; like above, for all α, β , $\alpha a \beta \in \text{cl}_{\text{kat}}\{a\} = \text{cl}_{\text{kat}}\llbracket a \rrbracket$.
- $\phi \in \mathcal{T}_{\text{BA}}$: $\mathcal{G}(\phi) = \{\alpha \mid \alpha \models \phi\}$, for all α such that $\alpha \models \phi$, $\text{KA}_{\text{kat}} \vdash \alpha \leq \phi$, whence $\alpha \in \text{cl}_{\text{kat}}\{\phi\} = \text{cl}_{\text{kat}}\llbracket \phi \rrbracket$.
- $e + f$: we have

$$\begin{aligned} \mathcal{G}(e + f) &= \mathcal{G}(e) \cup \mathcal{G}(f) \\ &\subseteq \text{cl}_{\text{kat}}(\llbracket e \rrbracket) \cup \text{cl}_{\text{kat}}(\llbracket f \rrbracket) && \text{(IH)} \\ &\subseteq \text{cl}_{\text{kat}}(\llbracket e \rrbracket \cup \llbracket f \rrbracket) && \text{(cl}_{\text{kat}} \text{ monotone)} \\ &= \text{cl}_{\text{kat}}(\llbracket e + f \rrbracket) \end{aligned}$$

- $e \cdot f$: suppose $w \in \mathcal{G}(e \cdot f)$, i.e., $w = x\alpha y$ for $x\alpha \in \mathcal{G}(e)$ and $\alpha y \in \mathcal{G}(f)$. Via the induction hypothesis we know that $x\alpha \in \text{cl}_{\text{kat}}(\llbracket e \rrbracket)$ and $\alpha y \in \text{cl}_{\text{kat}}(\llbracket f \rrbracket)$. Hence $x\alpha\alpha y \in \text{cl}_{\text{kat}}(\llbracket e \rrbracket) \cdot \text{cl}_{\text{kat}}(\llbracket f \rrbracket) \subseteq \text{cl}_{\text{kat}}(\llbracket e \cdot f \rrbracket)$ via Lemma A.2. As $\text{KA}_{\text{kat}} \vdash \alpha \leq \alpha\alpha$, we get that $x\alpha y \in \text{cl}_{\text{kat}}(\llbracket e \cdot f \rrbracket)$.
- e^* : a similar argument works. \square

Recall that \mathcal{GS} is the set of all guarded strings. We deduce from Lemma 4.3 that closing a guarded string language under kat does not add new guarded strings:

Lemma C.3. *For all guarded string language L , $L = \text{cl}_{\text{kat}}(L) \cap \mathcal{GS}$.*

Proof. The left-to-right direction is trivial, as $L \subseteq \text{cl}_{\text{kat}}(L)$ and all strings in L are guarded. For the right-to-left direction, given a guarded string $\alpha_0 a_0 \dots a_{n-1} \alpha_n$ in $\text{cl}_{\text{kat}}(L)$, we simply use Lemma 4.3 with the sequence of α_i themselves: we have $\alpha_i \models \alpha_i$ for all i , so that $\alpha_0 a_0 \dots a_{n-1} \alpha_n$ actually belongs to L . \square

Lemma 4.5. *For all KAT expressions e , $\mathcal{G}(e) = \text{cl}_{\text{kat}}(\llbracket e \rrbracket) \cap \mathcal{GS}$.*

Proof. Immediate from Lemma C.3 and Lemma 4.4. \square

Lemma 4.7. *We have $\text{KA}_{\text{atom}} \vdash 1 = \sum_{\alpha \in \text{At}} \alpha$, and for all α , $\text{KA}_{\text{atom}} \vdash \alpha\alpha = \alpha$.*

Proof. The first equation comes from (atom.2) and (atom.3) For the second one,

$$\text{KA}_{\text{atom}} \vdash \alpha = \alpha \cdot 1 = \alpha \cdot \sum_{\beta \in \text{At}} \beta = \sum_{\beta \in \text{At}} \alpha \cdot \beta = \alpha \cdot \alpha$$

where we use (atom₁) in the last step. \square

Lemma 4.8. *For all $e \leq f \in \text{kat}$, we have that $\text{KA}_{\text{atom}} \vdash r(e) \leq r(f)$.*

Proof. We have five families of equations to consider, and thus ten families of inequations. We derive equalities directly, so that we only have to consider five cases:

– $\perp = 0$: we derive

$$\text{KA}_{\text{atom}} \vdash r(\perp) = \sum_{\alpha \models \perp} \alpha = 0 = r(0)$$

– $\top = 1$: we derive

$$\text{KA}_{\text{atom}} \vdash r(\top) = \sum_{\alpha \models \top} \alpha = \sum_{\alpha \in \text{At}} \alpha = 1 = r(1)$$

(using Lemma 4.7 for the last but one step)

– $\phi \wedge \psi = \phi \cdot \psi$: we derive

$$\begin{aligned} \text{KA}_{\text{atom}} \vdash r(\phi \cdot \psi) &= r(\phi) \cdot r(\psi) \\ &= \sum_{\alpha \models \phi} \alpha \sum_{\beta \models \psi} \beta \\ &= \sum_{\alpha \models \phi, \beta \models \psi} \alpha \beta \\ &= \sum_{\alpha \models \phi, \alpha \models \psi} \alpha \alpha && \text{(atom}_1\text{)} \\ &= \sum_{\alpha \models \phi, \alpha \models \psi} \alpha && \text{(Lemma 4.7)} \\ &= \sum_{\alpha \models \phi \wedge \psi} \alpha = r(\phi \wedge \psi) \end{aligned}$$

– $\phi \vee \psi = \phi + \psi$: we derive

$$\begin{aligned} \text{KA}_{\text{atom}} \vdash r(\phi + \psi) &= r(\phi) + r(\psi) \\ &= \sum_{\alpha \models \phi} \alpha + \sum_{\alpha \models \psi} \alpha \\ &= \sum_{\alpha \models \phi \text{ or } \alpha \models \psi} \alpha \\ &= \sum_{\alpha \models \phi \vee \psi} \alpha = r(\phi \vee \psi) \end{aligned}$$

- $\phi = \psi$, an instance of a Boolean algebra axiom: in this case, $\alpha \models \phi$ iff $\alpha \models \psi$ for all α , so that $r(\phi)$ and $r(\psi)$ are identical. \square

D Proofs for Section 5 (KAO)

Lemma 5.2. *For all $e \leq f \in \text{kao}$, we have $\text{KA}_{\text{contr}} \vdash r(e) \leq r(f)$.*

Proof. Similar to the proof of Lemma 4.8. First observe that in that proof, we needed the hypotheses of kat only for the $\top = 1$ case, which is not there, and for the $\phi \wedge \psi = \phi \cdot \psi$ case, which is now only an inequation, and which is dealt with as follows:

$$\begin{aligned}
 \text{KA}_{\text{contr}} \vdash r(\phi \wedge \psi) &= \sum_{\alpha \models \phi \wedge \psi} \alpha \leq \sum_{\alpha \models \phi \wedge \psi} \alpha \cdot \alpha && (\text{contr}) \\
 &\leq \sum_{\alpha \models \phi, \beta \models \psi} \alpha \cdot \beta \\
 &= \sum_{\alpha \models \phi} \alpha \cdot \sum_{\beta \models \psi} \beta = r(\phi) \cdot r(\psi) = r(\phi \cdot \psi) \quad \square
 \end{aligned}$$

We now prove the two parts of Lemma 5.5, in Lemma D.3 and Lemma D.1 below.

Lemma D.1. $\text{cl}_2 \circ \text{cl}_{\text{contr}} \subseteq \text{cl}_{\text{contr}} \circ \text{cl}_2$

Proof. By Lemma 3.13 it is sufficient to prove that $\text{st}_2 \circ \text{st}_1 \subseteq \text{cl}_1 \circ \text{st}_2^-$. Take $w \in \text{st}_2 \circ \text{st}_1(L)$. So $w = u\alpha v$ for some atom α and words such that $uv \in \text{st}_1(L)$. This implies that $uv = l\beta m$ for some words l, m and atom β and $l\beta m \in L$. From this we obtain that $uv = l\beta m$. By symmetry, we may assume that $|u| \leq |l|$, i.e., $l = ul'$, $v = l'\beta m$ for some word l' . Then $ul'\beta m \in L$, so $u\alpha l'\beta m \in \text{st}_2^-(L)$ and subsequently $u\alpha l'\beta m = u\alpha v = w \in \text{cl}_{\text{contr}} \circ \text{st}_2^-(L)$. \square

For the commutativity of contr and atom_3 we need an auxiliary lemma.

Lemma D.2. *Let $u, v \in (\Sigma + \mathcal{T}_{\text{BA}})^*$ and $\alpha \in \text{At}$. If $u\alpha^n v \in \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$ for some $n \geq 1$ then $u\alpha v \in \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$.*

Proof. We show that for all ordinals κ , for all natural numbers $n \geq 1$ we have that

$$u\alpha^n v \in \text{st}_3^\kappa \circ \text{cl}_{\text{contr}}(L) \quad \text{implies} \quad u\alpha v \in \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$$

by transfinite induction on κ . For the base case, we just have the inclusion $\text{cl}_{\text{contr}}(L) \subseteq \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$, and the limit case is trivial.

For the inductive step, assume it holds for some ordinal κ and let $u\alpha^n v \in \text{st}_3(\text{st}_3^\kappa(\text{cl}_{\text{contr}}(L)))$. That means $u\alpha^n v = lm$ for some l, m such that $l \ll [\sum_{\alpha \in \text{At}} \alpha] m \subseteq \text{st}_3^\kappa(\text{cl}_{\text{contr}}(L))$. We distinguish three cases.

1. $l = u\alpha^n l'$ and $v = l'm$. In this case we have

$$u\alpha^n l' \left[\sum_{\alpha \in \text{At}} \alpha \right] m = l \left[\sum_{\alpha \in \text{At}} \alpha \right] m \subseteq \text{st}_3^k(\text{cl}_{\text{contr}}(L))$$

By repeatedly applying the induction hypothesis we obtain that

$$u\alpha l' \left[\sum_{\alpha \in \text{At}} \alpha \right] m \subseteq \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$$

Via definition of cl_3 this leads to $\{u\alpha l' m\} = u\alpha l' \llbracket 1 \rrbracket m \subseteq \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$. As $u\alpha l' m = u\alpha v$ we have obtained the required result.

2. $l = u\alpha^p$ and $m = \alpha^q v$ for $p, q \geq 0$ and $p + q = n$. Here we have

$$u\alpha^p \left[\sum_{\alpha \in \text{At}} \alpha \right] \alpha^q v = l \left[\sum_{\alpha \in \text{At}} \alpha \right] m$$

Hence we know that $u\alpha^p \alpha \alpha^q v \in l \llbracket \sum_{\alpha \in \text{At}} \alpha \rrbracket m \subseteq \text{st}_3^k(\text{cl}_{\text{contr}}(L))$. Via the induction hypothesis we obtain that $u\alpha v \in \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$.

3. $u = lu'$ and $m = u'\alpha v$. This case is symmetrical to case 1. \square

Lemma D.3. $\text{cl}_{\text{contr}} \circ \text{cl}_3 \subseteq \text{cl}_3 \circ \text{cl}_{\text{contr}}$

Proof. By Lemma 3.12 it is sufficient to prove that $\text{st}_{\text{contr}} \circ \text{cl}_3 \subseteq \text{cl}_3 \circ \text{cl}_{\text{contr}}$. Even though contr satisfies the conditions of Lemma 3.13, proving the necessary implication needed in that lemma turned out to be more difficult. Hence, we prove the necessary implication of Lemma 3.12 instead. In order to prove $\text{st}_{\text{contr}} \circ \text{cl}_3 \subseteq \text{cl}_3 \circ \text{cl}_{\text{contr}}$ it is sufficient to show that $\text{st}_{\text{contr}} \circ \text{cl}_3 \circ \text{cl}_{\text{contr}}(L) \subseteq \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$. For a word $w \in \text{st}_{\text{contr}} \circ \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$, we know that $w \in u \llbracket \alpha \rrbracket v$ for some $u, v \in (\Sigma + \mathcal{T}_{\text{BA}})^*$ and $u \llbracket \alpha \rrbracket v \subseteq \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$. Via Lemma D.2, we can immediately conclude that $w = u\alpha v \in \text{cl}_3 \circ \text{cl}_{\text{contr}}(L)$. \square

E Proofs for Section 6 (KAPT)

Lemma 6.2. *There is a reduction from 2,3 to 3, witnessed by the homomorphism $r: \mathcal{T}(\Sigma + \text{At}) \rightarrow \mathcal{T}(\Sigma + \text{At})$ defined by*

$$r(x) = \begin{cases} a & x = a \in \Sigma \\ \sum \{ \alpha_1 \cdot \dots \cdot \alpha_n \mid \alpha = \bigcup_{i \leq n} \alpha_i, i \neq j \Rightarrow \alpha_i \neq \alpha_j \} & x = \alpha \in \text{At} \end{cases}$$

Proof. We use Lemma 3.3. For a letter $\mathbf{a} \in \Sigma$, $r(\mathbf{a}) = \mathbf{a}$ so that the first two conditions are trivial; we thus need to prove them only for atoms $\alpha \in \text{At}$. $\text{KA} \vdash \alpha \leq r(\alpha)$ follows by using the singleton sequence α which is a term in the sum $r(\alpha)$. For the second condition, it thus suffices to show $\text{KA}_{2,3} \vdash r(\alpha) \leq \alpha$, i.e., $\text{KA}_{2,3} \vdash \alpha_1 \dots \alpha_n \leq \alpha$ for all sequences $\alpha_1, \dots, \alpha_n$ of pairwise distinct atoms whose union is α . This follows by $n - 1$ successive applications of inequations in 2. It remains to check the last condition of Lemma 3.3; we consider the two kinds of equations separately:

- $\alpha \cdot \beta \leq \alpha \cup \beta$ (2): we have to prove $r(\alpha) \cdot r(\beta) \leq r(\alpha \cup \beta)$ in KA_3 . By distributivity, this amounts to proving $\alpha_1 \dots \alpha_n \beta_1 \dots \beta_m \leq r(\alpha \cup \beta)$ for all sequences of pairwise distinct atoms $\alpha_1, \dots, \alpha_n$ and β_1, \dots, β_m whose unions are α and β , respectively. The sequence $\alpha_1, \dots, \alpha_n, \beta_1, \dots, \beta_m$ almost yields a term in $r(\alpha \cup \beta)$: its union is $\alpha \cup \beta$, but it may contain duplicate entries. We simply remove such duplicates using inequations in 3 ($\gamma \leq 1$).
- $\alpha \leq 1$ (3): we have to show $r(\alpha) \leq 1$ in KA_3 . This follows by repeated applications of inequations in 3: all terms of the sum $r(\alpha)$ are below the identity. \square

Let us recall the five sets of hypotheses used for atom' :

- 1 = $\{\alpha \cup \beta \leq \alpha \cdot \beta \mid \alpha, \beta \in \text{At}\}$,
- 2 = $\{\alpha \cdot \beta \leq \alpha \cup \beta \mid \alpha, \beta \in \text{At}\}$,
- 3 = $\{\alpha \leq 1 \mid \alpha \in \text{At}\}$,
- 4 = $\{\alpha \cup \beta \leq \alpha \mid \alpha, \beta \in \text{At}\}$,
- 5 = $\{1 \leq \emptyset\}$ (for KAPT with \top , with \emptyset being the atom interpreted as \top).

Lemma 6.3. *We have the following inclusions of functions:*

- (i) $\text{st}_1 \circ \text{st}_2 \subseteq \text{st}_2 \circ \text{st}_1 + \text{id}$
- (ii) $\text{st}_1 \circ \text{st}_3 \subseteq \text{st}_3 \circ \text{st}_1 + \text{st}_4$
- (iii) $\text{st}_1 \circ \text{st}_4 \subseteq \text{st}_4 \circ \text{st}_1$
- (iv) $\text{st}_2 \circ \text{st}_3 \subseteq \text{st}_3 \circ \text{st}_2 + \text{st}_3 \circ \text{st}_3$
- (v) $\text{st}_2 \circ \text{st}_4 \subseteq \text{st}_4 \circ \text{st}_2 + \text{st}_4 \circ \text{st}_4 \circ \text{st}_2$
- (vi) $\text{st}_3 \circ \text{st}_4 \subseteq \text{st}_4 \circ \text{st}_3$

Proof. A word in $\text{st}_i(\text{st}_j(L))$ has shape $rs't$ with $rst = uv'w$ such that $uvw \in L$, where u, v, v', w, r, s, s', t are words with $v' \leq v$ an inequation of j and $s' \leq s$ an inequation of i . We depict this situation as follows:

$$rs't \leftarrow_i rst = uv'w \leftarrow_j uvw \in L$$

In each case we have to consider how the equation $rst = uv'w$ was obtained. When $|rs| \leq |u|$ or $|st| \leq |w|$, we immediately get $rs't \in \text{st}_j(\text{st}_i(L))$: the patterns do not overlap and may be triggered in any order. We only consider the other cases in the following.

- $i = 1, j = 2$. There is an overlap when $r = u, s = v', t = w$ in which case the two reductions cancel each other:

$$u(\alpha \cup \beta)w \xleftarrow{1} u\alpha\beta w \xleftarrow{2} u(\alpha \cup \beta)w \in L$$

In this case we simply have $rs't \in \text{id}(L)$.

There is another overlap when $u = r\gamma$, which is solved as follows, in st_2st_1 :

$$\begin{array}{c}
 r(\gamma \cup \alpha)\beta w \xleftarrow{1} r\gamma\alpha\beta w \xleftarrow{2} r\gamma(\alpha \cup \beta)w \in L \\
 \swarrow \quad \searrow \\
 \quad \quad r(\gamma \cup \alpha \cup \beta)w
 \end{array}$$

The remaining overlap is handled symmetrically.

- $i = 1, j = 3$. There is an overlap when $u = r\gamma$, which is solved as follows, using st_4 :

$$\begin{array}{c}
 r(\gamma \cup \alpha)w \xleftarrow{1} r\gamma\alpha w \xleftarrow{3} r\gamma w \in L \\
 \quad \quad \quad \swarrow \quad \searrow \\
 \quad \quad \quad \quad \quad r\gamma w
 \end{array}$$

The other overlap is handled symmetrically.

- $i = 1, j = 4$. There is an overlap when $u = r\gamma$, which is solved as follows, in st_4st_1 :

$$\begin{array}{c}
 r(\gamma \cup \alpha \cup \beta)w \xleftarrow{1} r\gamma(\alpha \cup \beta)w \xleftarrow{4} r\gamma\alpha w \in L \\
 \swarrow \quad \searrow \\
 \quad \quad r(\gamma \cup \alpha)w
 \end{array}$$

The other overlap is handled symmetrically.

- $i = 2, j = 3$. The only overlap is when $u = r$, which is solved as follows, in st_3^2 :

$$\begin{array}{c}
 u\alpha\beta w \xleftarrow{2} u(\alpha \cup \beta)w \xleftarrow{3} uw \in L \\
 \swarrow \quad \searrow \\
 \quad \quad u\alpha w
 \end{array}$$

- $i = 2, j = 4$. The only overlap is when $u = r$, which is solved as follows, in $st_4^2st_2$:

$$\begin{array}{c}
 u\gamma\delta w \xleftarrow{2} u(\alpha \cup \beta)w \xleftarrow{4} u\alpha w \in L \\
 \uparrow \quad \searrow \\
 \quad \quad u\gamma(\alpha \cap \delta)w \xleftarrow{4} u(\alpha \cap \gamma)(\alpha \cap \delta)w
 \end{array}$$

(Assuming $\gamma \cup \delta = \alpha \cup \beta$ to obtain the horizontal st_2 reduction and then justifying the other st_2 reduction.)

- $i = 3, j = 4$. There cannot be an overlap in this case. □

Lemma 6.4. *We have $cl_{1,2,3,4} = cl_4 \circ cl_3 \circ cl_2 \circ cl_1 = cl_4 \circ cl_{2,3} \circ cl_1$.*

Proof. We use Lemma 3.10 and Lemma 3.13 repeatedly, on (combinations of) the inclusions provided by Lemma 6.3. Note that all inequations considered here have a word as right-hand side, so that all functions st_H considered here are additive by Lemma A.1.

The second equation follows from $\text{cl}_{2,3} = \text{cl}_3\text{cl}_2$, which is a consequence of item (iv) (because $\text{st}_3\text{st}_2 + \text{st}_3^2 \subseteq \text{cl}_3\text{st}_2^-$). For the first equation, we have $\text{cl}_{3,4} = \text{cl}_4\text{cl}_3$ as a consequence of item (vi), so that it suffices to prove $\text{cl}_{1,2,3,4} = \text{cl}_{3,4}\text{cl}_2\text{cl}_1$. This follows by using Lemma 3.10 on cl_1 , cl_2 , and $\text{cl}_{3,4}$. Indeed, we have:

- $\text{cl}_1\text{cl}_2 \subseteq \text{cl}_2\text{cl}_1$ as a consequence of item (i) (because $\text{st}_2\text{st}_1 + \text{id} \subseteq \text{cl}_2\text{st}_1^-$);
- $\text{cl}_1\text{cl}_{3,4} \subseteq \text{cl}_{3,4}\text{cl}_1$ as a consequence of items (ii) and (iii) (yielding $\text{st}_1\text{st}_{3,4} \subseteq \text{cl}_{3,4}\text{st}_1^-$ altogether);
- $\text{cl}_2\text{cl}_{3,4} \subseteq \text{cl}_{3,4}\text{cl}_2$ as a consequence of items (iv) and (v) (yielding $\text{st}_2\text{st}_{3,4} \subseteq \text{cl}_{3,4}\text{st}_2^-$ altogether). \square

Lemma E.1. *We have $\text{cl}_{1,2,3,4,5} = \text{cl}_5\text{cl}_4\text{cl}_3\text{cl}_2\text{cl}_1$.*

Proof. Given Lemma 6.4, it suffices to show that cl_5 commutes over the four other closures.

- For $i \in \{2, 3, 4\}$, since the right-hand sides of the corresponding inequations are words of length at most one, we easily get $\text{st}_i\text{st}_5 \subseteq \text{st}_5\text{st}_i$, and hence $\text{cl}_i\text{cl}_5 \subseteq \text{cl}_5\text{cl}_i$ by Lemma 3.13.
- For 1, we use the fact that a word belongs to $\text{cl}_5(L)$ iff it is obtained from a word of L by removing an arbitrary number of occurrences of \emptyset (i.e., the atom denoting \top), and we prove directly $\text{st}_1\text{cl}_5 \subseteq \text{cl}_5\text{cl}_1$. Indeed a word w belongs to $\text{st}_1(\text{cl}_5(L))$ if $w = u(\alpha \cup \beta)v$ for some $u, u', v, v', \alpha, \beta, n$ such that $u'\alpha\emptyset^n\beta v' \in L$ and u (resp. v) is obtained from u' (resp. v') by removing some occurrences of \top . We deduce $u'(\alpha \cup \beta)v' \in \text{cl}_1(L)$ by using st_1 $n + 1$ times, and thus $w \in \text{cl}_5(\text{cl}_1(L))$ by removing the remaining additional occurrences of \emptyset . We finally deduce $\text{cl}_1\text{cl}_5 \subseteq \text{cl}_5\text{cl}_1$ with Lemma 3.12. \square