

# 2-adic parity explorations of the 3n+1 problem. Francis Laclé

## ▶ To cite this version:

Francis Laclé. 2-adic parity explorations of the 3n+1 problem. 2021. hal-03201180v2

# HAL Id: hal-03201180 https://hal.science/hal-03201180v2

Preprint submitted on 25 Apr 2021

**HAL** is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers. L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.



Distributed under a Creative Commons Attribution - ShareAlike 4.0 International License



Clockwise arc diagram of orbits from the parity Collatz function for  $n_0=25$   $^\dagger$ 

# 2-adic parity explorations of the 3n + 1 problem.

Francis Laclé

Oranjestad, Aruba

April 18, 2021

#### Abstract

Starting with a separate 2-adic reformulation of the Collatz function that maps parity orbits, we show that the stopping time k of this reformulation can be alternatively obtained with the help of some novel ideas as well as hybrid-algorithmic techniques.

 $<sup>^{\</sup>dagger}\mathrm{Figure}$  can be zoomed in to read the label on each node when viewing with a PDF viewer.

**Definition 1.** Let  $n \in \mathbb{N}^*$  be the input term and let  $t \in \mathbb{N}$  denote the index of a sequence of n's after iteratively applying the Collatz function  $c : \mathbb{N}^* \to \mathbb{N}^*$ , defined as the assignment

$$c(n_t) := \begin{cases} \frac{n_t}{2} & \text{if } n_t \equiv 0 \pmod{2} \\ 3n_t + 1 & \text{if } n_t \equiv 1 \pmod{2} \end{cases}$$

**Remark 1.** Throughout this paper we will modify  $c(n_t)$  and t. We have opted to retain the same symbols for clarity purposes.

**Definition 2.** Let e stand for even and  $\mathbb{N}_e = \{2n : n \in \mathbb{N}^*\} = \{2, 4, 6, ..., 2n\}$  represent the subset of even positive numbers.

**Definition 3.** Let o stand for odd and  $\mathbb{N}_o = \{2n - 1 : n \in \mathbb{N}^*\} = \{1, 3, 5, ..., 2n - 1\}$  represent the subset of odd positive numbers.

**Corollary 1.** Given the 2-adic valuation of n where  $\nu$  is the largest possible integer such that  $2^{\nu}$  divides elements of  $\mathbb{N}_e$  into positive integers, we find the outcome of this division to be always odd, i.e. any outcome  $n \in \mathbb{N}_o$ .

**Corollary 2.** Following Corollary 1 we can substitute  $\nu$  with the 2-adic valuation function  $\nu_2(n)$  (OEIS A007814 [3]) defined as the sequence

$$\nu_2(n) = (0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, 4, \dots).$$

**Remark 2.** A primitive recursive  $\mathbb{N}^*$ -adic valuation function is given in Appendix 2.

**Corollary 3.** With Corollaries 1 and 2 successive orbits of  $\frac{n_t}{2}$  in  $c(n_t)$  given  $n_t \in \mathbb{N}_e$  can be mapped by reformulating the assignment when  $n_t \equiv 0 \pmod{2}$  as follows

$$c(n_t) := \begin{cases} \frac{n_t}{2^{\nu_2(n_t)}} & \text{if } n_t \equiv 0 \pmod{2} \\ & \\ 3n_t + 1 & \text{if } n_t \equiv 1 \pmod{2} \end{cases}$$

**Corollary 4.** Any outcome of  $3n_t + 1$  when  $n_t \equiv 1 \pmod{2}$  always equals an  $n_{t+1} \in \mathbb{N}_e$ .

*Proof.* Let o be any odd integer.  $3 \cdot o$  results in another odd integer. Given the parity present in the natural order of  $\mathbb{N}$ , adding a 1 to this result will therefore always return an element of  $\mathbb{N}_e$ .

$n_0$	$(c(n_t))_{n_0\in\mathbb{N}_o}^\infty$
1	(2, 1,)
3	(5, 8, 1,)
5	(8,1,)
7	$(11, 17, 26, 13, 20, 5, 8, 1, \ldots)$
9	$(14, 7, 11, 17, 26, 13, 20, 5, 8, 1, \ldots)$
11	$(17, 26, 13, 20, 5, 8, 1, \ldots)$
13	$(20, 5, 8, 1, \ldots)$
15	$(23, 35, 53, 80, 5, 8, 1, \ldots)$
17	$(26, 13, 20, 5, 8, 1, \ldots)$
19	$(29, 44, 11, 17, 26, 13, 20, 5, 8, 1, \ldots)$
21	(32, 1,)
23	$(35, 53, 80, 5, 8, 1, \ldots)$

Table 1: Orbits of the first 12 integers in  $\mathbb{N}_o$  of  $c(n_t)$  as given in Corollary 5.

**Corollary 5.** Following Corollary 4 we can reformulate  $c(n_t)$  by adding one division of 2 when  $n_t \equiv 1 \pmod{2}$  as

$$c(n_t) := \begin{cases} \frac{n_t}{2^{\nu_2(n_t)}} & \text{if } n_t \equiv 0 \pmod{2} \\ \frac{3n_t + 1}{2} & \text{if } n_t \equiv 1 \pmod{2} \end{cases}$$

thereby further accelerating the succession. Orbits of the first 12 odd integers are listed for demonstration purposes in Table 1.

Given Corollary 5 will continue to explore the subset  $n_t \in \mathbb{N}_o$  starting at t = 0.

**Definition 4.** Let  $s(n_t)$  denote an expansion of  $\frac{3n_t+1}{2}$  as follows

$$s(n_t) = \frac{n_t + 2(n_t) + 1}{2}.$$

Observe that after applying  $s(n_t)$  we get an outcome  $n_{t+1}$ , which can be recursively added to and into the expression  $\frac{n_t+2(n_t)+1}{2}$  provided that  $n_t \equiv 1 \pmod{2}^{\dagger}$ .

**Remark 3.** Throughout this paper we will modify both  $s(n_t)$  and the distance between increments of t. We have opted to retain the same symbols for clarity purposes.

<sup>&</sup>lt;sup>†</sup>Working this out on brown paper results in a shape similar to an upside-down Cantor ternary tree.

**Lemma 1.** The parity of  $n_{t+1}$  can alter from  $n_t$  after applying the operation  $s(n_t)$  and this alteration depends on whether the outcome  $n_{t+1} = s(n_t)$  contains one or more factors of 2 given  $n_t \in \mathbb{N}_o$ .

*Proof.* Let  $n_0 \in \mathbb{N}_o$  be the initial input for s. We have

$$s(n_0) = \frac{n_0 + 2(n_0) + 1}{2} = \frac{3(n_0) + 1}{2} = \frac{3}{2}n_0 + \frac{1}{2} = n_1$$

which can be restated as

 $n_1 - \frac{1}{2} = \frac{3(n_0)}{2}$ 2(n\_1) - 1 = 3(n\_0) (multiply both sides by 2) 2(n\_1) = 3(n\_0) + 1 (add 1 to both sides)

However,  $n_1$  could or could not be even depending on the number of factors of 2 contained [2]. To complete the proof of existence it is suffice to show one example of each case. For the first case we consider  $n_0 = 11$ . This gives  $n_1 = 17$ , while for the second case we consider  $n_0 = 5$ , which gives  $n_1 = 16$  that contains four factors of two, or  $\nu_2(16) = 4$ , and we are done.

**Lemma 2.** Given  $\mathbb{N}_o$  there exists at least one initial value of  $n_t$  where the outcome of more than one recursive iteration of  $s(n_t)$  results in another  $n_t \in \mathbb{N}_e$  i.e. where  $\nu_2(n_t) > 1$  for a particular  $n_t$ .

*Proof.* Let  $n_0 \in \mathbb{N}_o$  be the initial value of  $s(n_t)$  and let  $n_1$  be odd, we can insert  $s(n_0)$  given by  $n_1$  into  $s(n_t)$  to obtain

$$s(n_1) = \frac{\frac{n_0 + 2(n_0) + 1}{2} + 2\left(\frac{n_0 + 2(n_0) + 1}{2}\right) + 1}{2} = n_2$$

which results in the dividend of  $s(n_1)$  being even as  $n_1$  is odd, simplified as

$$s(n_1) = \frac{\frac{n_0 + 2(n_0) + 1}{2} + \frac{2(n_0) + 4(n_0) + 2}{2} + 1}{2} = \frac{\frac{3}{2}n_0 + 3(n_0) + \frac{5}{2}}{2} = \frac{\frac{9}{2}n_0 + \frac{5}{2}}{2} = \frac{9}{4}n_0 + \frac{5}{4} = n_2$$

which can be restated as

$$\begin{split} n_2 - \frac{5}{4} &= \frac{9(n_0)}{4} \quad (\text{multiply both sides by 4}) \\ 4(n_2) - 5 &= 9(n_0) \quad (\text{add 5 to both sides}) \\ 4(n_2) &= 9(n_0) + 5. \end{split}$$

Similar to the result of Lemma 1 we have that 9 multiplied by an odd integer is an odd integer and an odd integer added with 5 gives an even integer, therefore if  $n_0$  is odd then  $4(n_2)$  given  $s(n_0) \rightarrow s(n_1)$  must be even.

Here  $n_2$  is even or odd following Lemma 1. For example,  $n_0 = 7$  gives  $n_2 = 17$ , while  $n_0 = 19$  gives  $n_2 = 44$  as  $\nu_2(44) = 2$ .

We now solve for the third iteration. Let  $n_2$  equal an odd integer after two iterations of  $s(n_t)$ . We can insert  $n_2$  in  $s(n_t)$  to obtain

$$s(n_2) = \frac{\frac{\frac{3}{2}n_0 + 3(n_0) + \frac{5}{2}}{2} + 2\left(\frac{\frac{3}{2}n_0 + 3(n_0) + \frac{5}{2}}{2}\right) + 1}{2} = n_3$$

simplified as

$$s(n_2) = \frac{\frac{\frac{3}{2}n_0 + 3(n_0) + \frac{5}{2}}{2} + \frac{3(n_0) + 6(n_0) + 5}{2} + 1}{2} = \frac{\frac{9}{4}n_0 + \frac{9}{2}n_0 + \frac{19}{4}}{2} = \frac{\frac{27}{4}n_0 + \frac{19}{4}}{2} = \frac{27}{8}n_0 + \frac{19}{8} = n_3$$

from which follows that

$$n_3 - \frac{19}{8} = \frac{27(n_0)}{8}$$
 (multiply both sides by 8)  
8(n\_3) - 19 = 27(n\_0) (add 19 to both sides)  
8(n\_3) = 27(n\_0) + 19

As was the case with the previous two results we have that 27 multiplied by an odd integer is an odd integer and an odd integer added with 19 gives an even integer, therefore  $n_0$  must be odd and  $8(n_3)$  must be even.

To complete the proof it is suffice to provide two examples of each case. As a first case we have,  $n_0 = 15$  gives  $n_3 = 53$ , while for the second case we have  $n_0 = 23$  gives  $n_3 = 80$  as  $\nu_2(80) = 4$ .

**Theorem 1.** The coefficients of recursive iterations of  $s(n_t)$  are uniquely defined for some k, with k being the exponent having bases 2 and 3.

*Proof.* Let  $n_t \in \mathbb{N}_o$  denote iterative outcomes of  $s(n_t)$  and let  $k \in \mathbb{N}$  be an iteration index for subsequent outcomes. From Lemma 1 and Lemma 2 we have the following incremental

results up until k = 3

$$n_{0} = \frac{1}{1}n_{0} + \frac{0}{1}$$

$$n_{1} = \frac{3}{2}n_{0} + \frac{1}{2}$$

$$n_{2} = \frac{9}{4}n_{0} + \frac{5}{4}$$

$$n_{3} = \frac{27}{8}n_{0} + \frac{19}{8}$$

Then by induction we claim

$$n_4 = \frac{81}{16}n_0 + \frac{65}{16}$$

for k = 4. Suppose that this is true, let

$$n_k = \frac{a}{b}n_0 + \frac{z}{b}$$

define all terms where a is the dividend of the first term, b be the divisor of both terms respectively, and z be the dividend of the right term. Then, the claim is valid if we can derive a correspondence with each equation for every k together with valid examples of  $c(n_t)$ . For a it is obvious that these are repeated multiplications of base 3, as such  $a = 3^k$ . For b it is similarly obvious that these are repeated multiplications of base 2, as such  $b = 2^k$ , giving

$$n_k = \frac{3^k}{2^k} n_0 + \frac{z}{2^k}.$$

With z representing  $(0, 1, 5, 19, ..., \mathbb{N})$  we claim this to be the enumeration of the dividends of both terms up until k - 1 for k > 0. For  $1 \le k \le 4$  we have

$$z_{1} = 1 = 1 + 0 \qquad | k = 1$$

$$z_{2} = 5 = 1 + 0 + 3 + 1 \qquad | k = 2$$

$$z_{3} = 19 = 1 + 0 + 3 + 1 + 9 + 5 \qquad | k = 3$$

$$z_{4} = 65 = 1 + 0 + 3 + 1 + 9 + 5 + 27 + 19 | k = 4$$

with  $z_0 = 0$ . This can be rewritten as

$$z_0 = 0 \qquad \qquad | k = 0$$

$$z_1 = 1 = 3^0 + z_0 \qquad | k = 1$$

$$z_2 = 5 = 3^0 + z_0 + 3^1 + z_1 \qquad \qquad | k = 2$$

$$z_{3} = 19 = 3^{0} + z_{0} + 3^{1} + z_{1} + 3^{2} + z_{2} \qquad | k = 3$$
  
$$z_{4} = 65 = 3^{0} + z_{0} + 3^{1} + z_{1} + 3^{2} + z_{2} + 3^{3} + z_{3} | k = 4$$

and further expanded as

$$z_{0} = 0$$

$$z_{1} = 1 = 3^{0} + 0$$

$$z_{2} = 5 = 3^{0} + 0 + 3^{1} + 3^{0} + 0$$

$$z_{3} = 19 = 3^{0} + 0 + 3^{1} + 3^{0} + 0 + 3^{2} + 3^{0} + 0 + 3^{1} + 3^{0} + 0$$

$$z_{4} = 65 = 3^{0} + 0 + 3^{1} + 3^{0} + 0 + 3^{2} + 3^{0} + 0 + 3^{1} + 3^{0} + 0 + 3^{3} + 3^{0} + 0 + 3^{1} + 3^{0}$$

The number of times that  $3^0$  or 1 gets included in the enumerative series is always by an additional factor of 2, and so we can simplify to

$$\begin{aligned} z_0 &= 0\\ z_1 &= 1 = 2^0 + 0\\ z_2 &= 5 = 2^1 + 0 + 3^1 + 0\\ z_3 &= 19 = 2^2 + 0 + 3^1 + 0 + 3^2 + 0 + 3^1 + 0\\ z_4 &= 65 = 2^3 + 0 + 3^1 + 0 + 3^2 + 0 + 3^1 + 0 + 3^3 + 0 + 3^1 + 0 + 3^2 + 0 + 3^1 + 0. \end{aligned}$$

Rearranging and cancelling the zeros we get

$$z_1 = 1 = 2^0$$
  

$$z_2 = 5 = 2^1 + 3^1$$
  

$$z_3 = 19 = 2^2 + 3^1 + 3^1 + 3^2$$
  

$$z_4 = 65 = 2^3 + 3^1 + 3^1 + 3^1 + 3^1 + 3^2 + 3^2 + 3^3.$$

which we claim to have a closed form defined as

$$z = a - b = 3^k - 2^k$$
.

Let us rewrite outcomes of  $z_k$  as binomial integers, we have

$$z_1 = 1 + 0$$
  
 $z_2 = 2 + 3$   
 $z_3 = 4 + 15$   
 $z_4 = 8 + 57.$ 

For each  $z_k$  we can expand it equally so the right term is equal to  $3^k$  while the left term is

equal to  $-2^k$  such as

$$z_1 = (1 - 3) + (0 + 3)$$
$$z_2 = (2 - 6) + (3 + 6)$$
$$z_3 = (4 - 12) + (15 + 12)$$

If this holds then the additional integers in  $z_4$  should equate 24 as the assertion here is that every pair of additional integers contains a factor of  $2^{k-1}$  times 3. Indeed

$$z_4 = (8 - 24) + (57 + 24) = 65$$

as claimed.

We can now reformulate  $s(n_t)$  explicitly. We assume values of  $n_t \in \mathbb{N}_o$  such that  $s(n_t)$  iterates in k+1 steps with initial value k=0 up until an outcome that contains factors of  $2^{n \in \mathbb{N}^*}$ . We have

$$s(n_k) = \frac{3^k(n_0) + 3^k - 2^k}{2^k} = \frac{3^k(n_0) + 3^k}{2^k} - 1$$

up until

 $\max\{k : s(n_k) \equiv 0 \pmod{2}\}.$ 

As an example, let us evaluate  $s(n_k)$  for the  $c(n_t)$  function given in Corollary 5 with  $n_0 = 31$  to be the integer that returns an even integer after five consecutive iterations. We have

$$47 \rightarrow 71 \rightarrow 107 \rightarrow 161 \rightarrow 242$$

from which we can work out as

$$s(31_5) = \frac{3^5(31) + 3^5 - 2^5}{2^5} = \frac{243(31) + 243 - 32}{32} = \frac{7533 + 243 - 32}{32} = \frac{7744}{32} = 242.$$

**Remark 4.** It follows from the definition of  $s(n_k)$  that its domain can be expanded to include all  $n \in \mathbb{N}^*$  as initial values  $n_0$  as  $s(n_0) = n_0$ .

**Corollary 6.** Following Corollary 5, we can make  $s(n_k)$  part of  $c(n_t)$  such that

$$c(n_t) := \begin{cases} \frac{n_t}{2^{\nu_2(n_t)}} & \text{if } n_t \equiv 0 \pmod{2} \\ s(n_{t,k}) & \text{if } n_{t,k} \equiv 1 \pmod{2} \end{cases}$$

The following orbits of initial values of  $n_{0,k} = 1$  till  $n_{0,k} = 12$  for  $n \in \mathbb{N}^*$  with presets k = 0 are listed for demonstration purposes in Table 2.

$n_0$	$(c(n_t))_{n\in\mathbb{N}^*}^\infty$	$\max\{k: s(n_{t,k}) \equiv 0 \pmod{2}\}$
1	(1, 2, 1,)	(1, 0,)
2	(1,)	(0,)
3	(3, 5, 8, 1,)	(2, 0,)
4	(1,)	(0,)
5	(5, 8, 1,)	(1, 0,)
6	(3, 5, 8, 1,)	(2, 0,)
7	$(7, 11, 17, 26, 13, 20, 5, 8, 1, \ldots)$	(3, 2, 2, 0,)
8	(1,)	(0,)
9	(9, 17, 7, 11, 17, 26, 13, 20, 5, 8, 1,)	(5, 2, 2, 0,)
10	(5, 8, 1,)	(1, 0,)
11	$(11, 17, 26, 13, 20, 5, 8, 1, \ldots)$	$(2, 2, 2, 0, \ldots)$
12	$(3, 5, 8, 1, \ldots)$	(2, 0,)

Table 2: Orbits of  $c(n_t)$  for initial values of  $n_0 = 1$  till  $n_0 = 12$  for  $n \in \mathbb{N}^*$ .

**Theorem 2.** From Theorem 1 we derive a substitution for k of s(n) where the number of k-iterations required to reach an  $n \in \mathbb{N}_e$  from n = 0 is the 2-adic valuation function  $\nu_2(n)$  with an offset of 1, i.e.  $k = v_2(n+1)$ .

*Proof.* Let  $n_0$  be an  $n_k \in \mathbb{N}^*$ . We have  $s(n_0) = n_0$  as all k = 0 exponents result in 1 with an even dividend and a divisor of one. Consider  $n_0 = 1$  and k = 1, evaluated as

$$s(1_1) = \frac{3(1) + 3 - 2}{2} = \frac{3(1) + 1}{2}.$$

Observe that  $\frac{3(1)+1}{2}$  is the expression given in the equivalence relation of Corollary 4 for cases where  $n \equiv 1 \pmod{2}$ . Let us denote this expression in terms of even and odd variables e and o, and the dividend and divisor as superscript dd and dr respectively. From the rules of modular arithmetic we have

$$\frac{e^{dd}}{e^{dr}} = \frac{o(o) + o}{e} = \frac{o + o}{e}$$

Outcomes of  $n_0 \equiv 1 \pmod{2}$  for  $e^{dr} > 1$  result in a  $\frac{e^{dd}}{e^{dr}}$  that returns an integer when

$$e^{dd} = e^{dr} \text{ if } e^{dr} > 1 \qquad \text{or}$$
$$e^{dd} > e^{dr} \text{ if } e^{dd} = 2^p(e^{dr}) \text{ for } p > 0.$$

Let us now consider the case of  $n \in \mathbb{N}_e$ . We have

$$\frac{o^{dd}}{e^{dr}} = \frac{o(e) + o}{e} = \frac{e + o}{e}.$$

Outcomes of  $n_0 \equiv 0 \pmod{2}$  are therefore rational and will never return an integer, except for the odd case when we have  $2^0$  or 1 as the divisor.

For the odd case  $n_0 = 1$  recall that

$$s(1_1) = \frac{3(1) + 3 - 2}{2} = \frac{3(1) + 1}{2} = \frac{4}{2}$$

continuing the iteration for k = 2 we have

$$s(1_2) = \frac{9(1) + 9 - 4}{4} = \frac{9(1) + 5}{4} = \frac{14}{4}$$

While both iterations produce  $\frac{e^{dd}}{e^{dr}}$  we observe that solely the first iteration,  $s(1_1)$ , satisfies the constraint  $e^{dd} = 2^p(e^{dr})$  thereby returning the even integer 2. We further observe that  $e^{dd}$  of  $s(1_2)$  contains one less a factor of two than  $e^{dr}$ , namely a factor of  $2^1$  against a factor of  $2^2$ .

Let us consider the subsequent case of  $n_0 = 3$ , of which by definition of  $s(n_k)$  has a maximum of  $0 \le k \le 2$ . We have

$$s(3_1) = \frac{3(3) + 3 - 2}{2} = \frac{3(3) + 1}{2} = \frac{10}{2}$$

and

$$s(3_2) = \frac{9(3) + 9 - 4}{4} = \frac{9(3) + 5}{4} = \frac{32}{4}$$

and finally

$$s(3_3) = \frac{27(3) + 27 - 8}{8} = \frac{27(3) + 19}{8} = \frac{100}{8}$$

The above shows that  $e^{dd}$  of  $s(3_1)$  contains a factor of  $2^1$  against  $2^1$  resulting in the odd integer 5. Continuing the iteration we have  $s(3_2)$  which contains a factor of  $2^5$  against a factor of  $2^2$ , satisfying the  $e^{dd} = 2^p(e^{dr})$  constraint. Finally, we have  $s(3_3)$  which does not satisfy the constraint as it contains a factor of  $2^2$  against a factor of  $2^3$ . Similarly for  $s(5_k)$  $s(7_k)$  we have

$$s(5_1) = \frac{3(5) + 3 - 2}{2} = \frac{3(5) + 1}{2} = \frac{16}{2}$$

and

$$s(7_3) = \frac{27(7) + 27 - 8}{8} = \frac{27(7) + 19}{8} = \frac{208}{8}.$$

Due to the general rational relation between the  $e^{dd}$  and  $e^{dr}$ , an  $n_k \in \mathbb{N}_e$  will result when the rational constraint  $e^{dd} = 2^p(e^{dr})$  is satisfied. To summarise we have for the four examples

$$\max\{k = 1 \text{ for } s(1_k) \in \mathbb{N}_e\} \text{ when } s(1_k) = \frac{2^2}{2^1} \\ \max\{k = 2 \text{ for } s(3_k) \in \mathbb{N}_e\} \text{ when } s(3_k) = \frac{2^5}{2^2} \\ \max\{k = 1 \text{ for } s(5_k) \in \mathbb{N}_e\} \text{ when } s(5_k) = \frac{2^3}{2^1} \\ \max\{k = 3 \text{ for } s(7_k) \in \mathbb{N}_e\} \text{ when } s(7_k) = \frac{2^4}{2^3}.$$

From the above we observe that each maximum of k is in fact the 2-adic valuation of  $n_0 + 1 \in \mathbb{N}_e$  from  $n_0$ .

Recall our formulation of  $s(n_k)$  from Theorem 1

$$s(n_k) = \frac{3^k(n_0) + 3^k - 2^k}{2^k} = \frac{3^k(n_0) + 3^k}{2^k} - 1$$

through the negative exponents rule we get

$$s(n_k) = 2^{-k} \left( 3^k(n_0) + 3^k - 2^k \right) = 2^{-k} 3^k(n_0) + 2^{-k} 3^k - 1$$

expressing the above formulation as a rational gives

$$s(n_k) = \frac{1}{2^k} \cdot \frac{3^k}{1} \cdot \frac{n_0}{1} + \frac{1}{2^k} \cdot \frac{3^k}{1} - \frac{2^k}{2^k}$$

We take the first step of k-iterations starting at  $s(7_0)$  as the example to further derive the proof. Evaluating the dividend for each k gives

- 3(7) + 3 2 = 22 (divided by 2 gives the odd integer 11)
- 9(7) + 9 4 = 68 (divided by 4 gives the odd integer 17)
- 27(7) + 27 19 = 208 (divided by 8 gives the *even* integer 26)

We can expand each outcome of k as follows

$$\frac{1}{2} \cdot \frac{3}{1} \cdot \frac{7}{1} + \frac{1}{2} \cdot \frac{3}{1} - \frac{1}{1} = \frac{3}{2} \cdot \frac{7}{1} + \frac{3}{2} - \frac{1}{1} = \frac{21}{2} + \frac{3}{2} - \frac{1}{1} = \frac{24}{2} - \frac{2}{2} = \frac{22}{2} = \frac{11}{1}$$

$$\frac{1}{4} \cdot \frac{9}{1} \cdot \frac{7}{1} + \frac{1}{4} \cdot \frac{9}{1} - \frac{1}{1} = \frac{9}{4} \cdot \frac{7}{1} + \frac{9}{4} - \frac{1}{1} = \frac{63}{4} + \frac{9}{4} - \frac{1}{1} = \frac{72}{4} - \frac{4}{4} = \frac{68}{4} = \frac{34}{2} = \frac{17}{1}$$

$$\frac{1}{8} \cdot \frac{27}{1} \cdot \frac{7}{1} + \frac{1}{8} \cdot \frac{27}{1} - \frac{1}{1} = \frac{27}{8} \cdot \frac{7}{1} + \frac{27}{8} - \frac{1}{1} = \frac{189}{8} + \frac{27}{8} - \frac{1}{1} = \frac{216}{8} - \frac{8}{8} = \frac{208}{8} = \frac{104}{4} = \frac{52}{2} = \frac{26}{1}$$

From the above expansion it can be observed on an elementary level that each k-iteration adds factors of three to the first two left terms, starting with a multiplication

$$3^k \cdot n_0$$

followed by an addition

$$3^k \cdot n_0 + 3^k$$

and a subtraction of  $-2^k$  for the right dividend term, to finally culminate with a division by  $2^k$ . From this it follows that multiples of 3 with odd numbers result in odd numbers as o(o) = o. Then, adding these multiples with powers of 3 result in even numbers as o(o) + o = e. When we consider the formulation of  $s(n_k)$  without  $-\frac{2^k}{2^k} = -1$  we get

$$\frac{3^k(n_0) + 3^k}{2^k}$$

which is equivalent to

$$\frac{3^k(n_0+1)}{2^k}$$

This implies that

$$3^k \cdot o + 3^k = 3^k \cdot (o+1)$$

If o is odd then (o + 1) must be even. If (o + 1) is even then it must contain at least one factor of 2. Illustrated with natural numbers 7 and 8 as

$$7 = 2 \cdot 2 \cdot 2 - 1$$
$$8 = 2 \cdot 2 \cdot 2$$

This leaves us with the following two simplified formulations of  $s(n_{t,k})$ 

$$s(n_{t,k}) = \frac{3^k(n_t+1) - 2^k}{2^k} = \frac{3^k(n_t+1)}{2^k} - 1 \mid n_t \in \mathbb{N}_c$$

and observe for  $s(7_k)$  the following

$$s(7_0) = \frac{8}{1} - \frac{1}{1} = \frac{2 \cdot 2 \cdot 2 - 1}{1} = 2 \cdot 2 \cdot 2 - 1 \in \mathbb{N}_o$$
  

$$s(7_1) = \frac{24}{2} - \frac{2}{2} = \frac{3 \cdot 2 \cdot 2 \cdot 2 - 2}{2} = 3 \cdot 2 \cdot 2 - 1 \in \mathbb{N}_o$$
  

$$s(7_2) = \frac{72}{4} - \frac{4}{4} = \frac{3 \cdot 3 \cdot 2 \cdot 2 \cdot 2 - 4}{4} = 3 \cdot 3 \cdot 2 - 1 \in \mathbb{N}_o$$
  

$$s(7_3) = \frac{216}{8} - \frac{8}{8} = \frac{3 \cdot 3 \cdot 3 \cdot 2 \cdot 2 \cdot 2 - 8}{8} = 3 \cdot 3 \cdot 3 - 1 \in \mathbb{N}_e$$

For the case of  $s(7_k)$  we thus have k = 3 as it takes three k-iterations to factor out all 2's from 7+1 such that the outcome is again an even integer. More specifically, a subtraction

by  $-\frac{2^k}{2^k}$  in  $s(7_3)$  where its outcome returns an even integer. What remains as a constant term for every k is  $-\frac{2^k}{2^k}$ . This subtraction retains  $s(n_{t,k})$  to have an even outcome up until the same number of factors of 2's have been iteratively added as the initial  $(n_t + 1)$ , such that divisions by  $2^k$  can factor out enough of the even 2's to return back to an outcome in the form of o - 1 = e, thus ending the iterative process of  $s(n_{t,k})$ . This process coincides with the outcome of  $\nu_2(7+1) = 3$  as  $2^3 = 8 = 7 + 1$ . Hence, through this equivalence we can look ahead at the number of k-iterations needed for  $n_{t,0} \in \mathbb{N}_o$  to reach an  $n_{t,k} \in \mathbb{N}_e$ , defined as

$$k = \nu_2(n+1)$$

With the demonstration in Theorem 1 and consequently Theorem 2, we exhibit the recursive divisions by factors of two and multiplications by factors of three in elementary form, using  $s(7_{t,k})$  as the exemplar, while highlighting noteworthy terms in bold. We have, with initial values t = 0 and k = 0

$$\begin{split} s(7_{0,0}) &= \frac{8}{1} - \frac{1}{1} = \frac{2 \cdot 2 \cdot 2 - 1}{1} = \mathbf{2} \cdot \mathbf{2} \cdot \mathbf{2} - 1 \\ s(7_{0,1}) &= \frac{24}{2} - \frac{2}{2} = \frac{22}{2} = \frac{11}{1} = \frac{3 \cdot 2 \cdot 2 - 2}{2} = 3 \cdot \mathbf{2} \cdot \mathbf{2} - 1 \\ s(7_{0,2}) &= \frac{72}{4} - \frac{4}{4} = \frac{68}{4} = \frac{34}{2} = \frac{17}{1} = \frac{3 \cdot 3 \cdot 2 \cdot 2 \cdot 2 - 4}{4} = 3 \cdot 3 \cdot \mathbf{2} - 1 \\ s(7_{0,3}) &= \frac{216}{8} - \frac{8}{8} = \frac{208}{8} = \frac{104}{4} = \frac{52}{2} = \frac{26}{1} = \frac{3 \cdot 3 \cdot 3 \cdot 2 \cdot 2 \cdot 2 - 8}{8} = 3 \cdot 3 \cdot 3 - 1 \\ \frac{26}{2} &= \frac{13}{1} = \frac{3 \cdot 3 \cdot 3 - 1}{2} = \frac{2 \cdot 2 (2 \cdot 2 \cdot 2 - 1) - 2}{2} = \mathbf{2} (\mathbf{2} \cdot \mathbf{2} \cdot \mathbf{2} - 1) - 1 \\ s(13_{1,1}) &= \frac{42}{2} - \frac{2}{2} = \frac{40}{2} = \frac{20}{1} = \frac{3 \cdot 2 (2 \cdot 2 \cdot 2 - 1) - 2}{2} = 3 (\mathbf{2} \cdot \mathbf{2} \cdot \mathbf{2} - 1) - 1 \\ \frac{20}{2} &= \frac{10}{1} = \frac{3 (2 \cdot 2 \cdot 2 - 1) - 1}{2} = \frac{3 (2 \cdot 2 \cdot 2 - 1) - 2}{2} = 3 (\mathbf{2} \cdot \mathbf{2}) - 1 - 1 \\ \frac{10}{2} &= \frac{5}{1} = \frac{3 (2 \cdot 2) - 1 - 1}{2} = (3 \cdot \mathbf{2}) - 1 \\ s(5_{2,1}) &= \frac{18}{2} - \frac{2}{2} = \frac{16}{2} = \frac{8}{1} = \frac{(3 \cdot 3 \cdot 2) - 2}{2} = (3 \cdot 3) - 1 \\ \frac{8}{2} &= \frac{4}{1} = \frac{(3 \cdot 3) - 1}{2} = \frac{2 \cdot 2 \cdot 2}{2} = \mathbf{2} \cdot \mathbf{2} \\ \frac{4}{2} &= \frac{2}{1} = \frac{2 \cdot 2}{2} = \mathbf{2} \\ \frac{2}{2} &= \frac{1}{1} = 1 \end{split}$$

with 1 entering the trivial cycle

$$s(1_{3,1}) = \frac{6}{2} - \frac{2}{2} = \frac{4}{2} = \frac{2}{1} = \frac{2 \cdot 2}{2} = \mathbf{2}$$
$$\frac{2}{2} = \frac{1}{1} = 1$$
$$\vdots$$

$n_0$	$n_1 = s(n_0)$	$n_2 = \frac{n_1}{2^{\nu_2(n_1)}}$	$\left  \begin{array}{c} n_3 = s(n_2) \end{array} \right $	$n_4 = \frac{n_3}{2^{\nu_2(n_3)}}$	$n_5 = s(n_4)$	$n_6 = \frac{n_5}{2^{\nu_2(n_5)}}$	$\left  \begin{array}{c} n_6 = s(n_5) \end{array} \right $	$n_7 = \frac{n_6}{2^{\nu_2(n_6)}}$
1	2	1	2	1	2	1	2	1
2	2	1	2	1	2	1	2	1
3	8	1	2	1	2	1	2	1
4	4	1	2	1	2	1	2	1
5	8	1	2	1	2	1	2	1
6	6	3	8	1	2	1	2	1
7	26	13	20	5	8	1	2	1
8	8	1	2	1	2	1	2	1
9	14	7	26	13	20	5	8	1
10	10	5	8	1	2	1	2	1
11	26	13	20	5	8	1	2	1
12	12	3	8	1	2	1	2	1

Table 3: Outcomes of the parity function  $c(n_t)$  as defined in Corollary 7 for  $1 \le n_0 \le 12 \land 1 \le t \le 7$ .

**Corollary 7.** Following previous theorems together with Corollary 6, the 2-adic valuation property for  $n \in \mathbb{N}^*$  can be added to  $s(n_{t,k})$  such that k becomes a function of  $\nu_2(n_t + 1)$  giving us the reformulation

$$s(n_t) = \frac{3^{\nu_2(n_t+1)}(n_0) + 3^{\nu_2(n_t+1)}}{2^{\nu_2(n_t+1)}} - 1 = \frac{3^{\nu_2(n_t+1)}(n_0+1)}{2^{\nu_2(n_t+1)}} - 1$$

Combining the above with  $\frac{n_t}{2^{v_2(n_t)}}$  effectively reformulates  $c(n_t)$  as a parity assignment of (even-)odd-even-odd and so forth orbits  $\forall n_t \in \mathbb{N}^*$ , defined as

$$c(n_t) := \begin{cases} \begin{array}{c} \frac{n_t}{2^{\nu_2(n_t)}} & \textit{if } n_t \equiv 0 \pmod{2} \\ \\ s(n_t) & \textit{if } n_t \equiv 1 \pmod{2} \end{array} \end{cases}$$

demonstrated in Table 3 for the following 12 outcomes of initial values of  $n_0 = 1$  till  $n_0 = 12$  for  $n_t \in \mathbb{N}^*$ .

**Corollary 8.** We follow Corollary 7 with one approach to describe the stopping time t from  $n_0$  up until  $n_t | c(n_t) = 1$  as a summation series that stems from the parity-alternating formulation of  $c(n_t)$ . Let t be the iteration index after every operation on n, we have

$$t = \begin{cases} \nu_2(n_0) + 2(\nu_2(n_{0\mid 1} + 1)) + \dots + \nu_2(n_t) & \text{if } n_t \neq 2^p \\ \nu_2(n_0) & \text{if } n_t = 2^p \end{cases}$$

where  $p \in \mathbb{N}^*$ . A 2-factor of k is included as each iteration of k is an even-multiple of

 $\frac{3n+1}{2}$ . Substituting i for  $\nu_2(n_t)$  and k for  $\nu_2(n_t+1)$ , we can simplify to

$$t = \begin{cases} i_0 + 2k_1 + \dots + i_t & \text{if } n_t \neq 2^p \\ i_0 & \text{if } n_t = 2^p \end{cases}$$

Considering Theorem 2 where  $s(n_t)$  returns an  $n \in \mathbb{N}_e$  we note that the above ends with an odd number of terms. 16 examples are shown in Appendix 3 for illustration purposes.

**Definition 5.** We introduce the clockwise arc diagram A as a visual aid that illustrates a couple of orbital properties of  $c(n_t)$  from Corollary 7 through conjectures 1, 2 and 3. The horizontal graph in the centre represents the number line  $n \in \mathbb{N}^*$ . Arcs on the top represent increments of  $n_t \to n_{t+1}$  from left to right by  $s(n_t)$ . Arcs on the bottom represent decrements of  $n_t \to n_{t+1}$  from right to left by  $\frac{n_t}{2^{v_2(n_t)}}$ <sup>†</sup>. A few examples are illustrated in Figures 1 and 2<sup>†</sup>.

**Remark 5.** In terms of graph theory, we can restate the main conjecture as follows. Up until reaching a vertex  $n_t = 1$  for a particular  $t \in T$  with the exception of  $n_0 = 1$ , a Hamiltonian or Eulerian path exists in  $A \forall n \in \mathbb{N}^*$ .

**Conjecture 1.** The summation of incremental differences minus the summation of decremental differences is equal to  $n_0 - 1 \forall \mathbb{N}^*$ .

Let  $\operatorname{inc}(n_{t-1}) = s(n_{t-1}) - n_{t-1} | t > 0$  and let  $\operatorname{dec}(n_t) = n_t - \frac{n_t}{2^{\nu_2(n_t)}} | t > 0$  while iterating  $c(n_t)$ . Through Figures 3 and 4 we can visualise the difference between  $\operatorname{inc}(n_{t-1})$  and  $\operatorname{dec}(n_t)$  to be

$$\left(\sum_{t=0}^{t \mid T \in \mathbb{N}_e} \operatorname{\mathbf{dec}}(n_t)\right) - \left(\sum_{t=1}^{t \mid T \in \mathbb{N}_o} \operatorname{\mathbf{inc}}(n_{t-1})\right) = n_0 - 1, \ \forall \ n_0 \in \mathbb{N}$$

which can be interpreted as the incremental step from 1 to  $n_0$  when a particular  $n_0$ -orbit converges back to 1.

**Conjecture 2.** Following conjecture 1 it can also be conjectured that

$$\left(\sum_{t=0}^{t \mid T \in \mathbb{N}_e} \operatorname{dec}(n_t)\right) \ge \left(\sum_{t=1}^{t \mid T \in \mathbb{N}_o} \operatorname{inc}(n_{t-1})\right) \ \forall \ n_0 \in \mathbb{N}^*$$

when a particular  $n_0$ -orbit converges back to 1.

Through A we can visually validate orbits of  $c(n_t)$  as defined in Corollary 7, alternating from  $+\mathbb{N}^* \to -\mathbb{N}^*$  or vice versa from  $-\mathbb{N}^* \to +\mathbb{N}^*$ . The following single-step iteration of initial values of  $n_0 = 1$  till  $n_0 = 12$  for  $n \in \mathbb{N}^*$  are listed for demonstration purposes in Table 4.

<sup>&</sup>lt;sup>†</sup>Source code for drawing the clock-wise arc diagrams can be found at https://github.com/flacle/. <sup>†</sup>Figures can be zoomed in to read the label on each node when viewing with a PDF viewer.



(a)  $n_1 = 26$  is its supremum element.



(b)  $n_3 = 26$  is its supremum element.

Figure 1: A for the parity  $c(n_t)$  function of orbits  $n_0 = 7$  and  $n_0 = 9$ .



(a)  $n_1 = 80$  is its supremum element.



(b)  $n_3 = 152$  is its supremum element.

Figure 2: A for the parity  $c(n_t)$  function for orbits  $n_0 = 15$  and  $n_0 = 39$ .



(a)  $n_1 = 2$  is its supremum element and the trivial cycle exemplar.



(b)  $n_3 = 8$  is its supremum element.

Figure 3: A for the parity  $c(n_t)$  function of orbits  $n_0 = 1$  and  $n_0 = 6$ .



(a)  $n_1 = 26$  is its supremum element.



(b)  $n_0 = 20$  is its supremum element.

Figure 4: A for the parity  $c(n_t)$  function for orbits  $n_0 = 17$  and  $n_0 = 20$ .

$n_0$	$c(n_0)$
1	+1
2	-1
3	+5
4	-3
5	+3
6	-3
7	+19
8	-7
9	+5
10	-5
11	+15
12	-9

Table 4: Alternating operations executed by  $c(n_t)$  as defined in Corollary 7, for initial values of  $n_0 = 1$  till  $n_0 = 12$ .

**Theorem 3.** We can a reformulate  $c(n_t)$  without the classical piecewise notation through the inclusion of divergent subsequences.

*Proof.* We perceive regularities within pairs of  $n_0$  tuples that are equally distanced from one another, listed in Table 4. For example  $\langle 1, 2 \rangle$ ,  $\langle 5, 6 \rangle$ ,  $\langle 9, 10 \rangle$ , and so on. For the first example we perceive an increase of +1 followed by a decrease of -1. For the second we have +3 followed by -3, similarly for  $n_0 = 9$  we have +5 and -5 and so on, with a distance of  $2^1$ . Suppose that the same rationale can be applied for pairs where  $\nu_2(n_0 + 1) > 1$ . Then, distances between pairs should also exhibit regularities. Indeed, this will be proved by construction to inherit  $z = 3^k - 2^k$  as defined in Theorem 1.

First, let the  $n_0$  terms  $\forall n_0 \in \mathbb{N}^*$  be expanded, illustrated for initial values of  $n_0 = 1$  till  $n_0 = 26$  in Table 5, and forming the sequence

 $n_1 = (2, 1, 8, 1, 8, 3, 26, 1, 14, 5, 26, 3, 20, 7, 80, 1, 26, 9, 44, 5, 32, 11, 80, 3, 38, 13)$ 

for  $n_0 = 1$  till  $n_0 = 26$ . This can be further expanded to highlight the recursion within each expression as shown in Table 6. Continuing this approach, expressions are reformulated

$n_0$	$n_0 + n_1$	$n_0$	$n_0 + n_1$
1	1 + 1	14	14 - 7
2	2 - 1	15	15 + 15 + 50
3	3 + 3 + 2	16	16 - 15
4	4 - 3	17	17 + 9
5	5 + 3	18	18 - 9
6	6 - 3	19	19 + 15 + 10
7	7 + 7 + 12	20	20 - 15
8	8 - 7	21	21 + 11
9	9 + 5	22	22 - 11
10	10 - 5	23	23 + 21 + 36
11	11 + 9 + 6	24	24 - 21
12	12 - 9	25	25 + 13
13	13 + 7	26	26 - 13

Table 5: Expanded  $n_0$  terms. For example, 7 + 19 = 14 + 7 = 7 + 7 + 12.

1			1
$n_0$	$n_0 + n_1$	$n_0$	$n_0 + n_1$
1	1 + 1	14	14 - 7
2	2 - 1	15	$15+15\cdot1+50\cdot1$
3	$3 + 3 \cdot 1 + 2 \cdot 1$	16	$16 - 15 \cdot 1$
4	$4 - 3 \cdot 1$	17	17 + 9
5	5 + 3	18	18 - 9
6	6 - 3	19	$19 + 3 \cdot 5 + 2 \cdot 5$
7	$7 + 7 \cdot 1 + 12 \cdot 1$	20	$20 - 3 \cdot 5$
8	$8 - 7 \cdot 1$	21	21 + 11
9	9 + 5	22	22 - 11
10	10 - 5	23	$23 + 7 \cdot 3 + 12 \cdot 3$
11	$11 + 3 \cdot 3 + 2 \cdot 3$	24	$24 - 7 \cdot 3$
12	$12 - 3 \cdot 3$	25	25 + 13
13	13 + 7	26	26 - 13

Table 6: Demonstration of recurrence within each expression.

$n_0$	$n_0 + n_1$	$n_0$	$  n_0 + n_1$
1	$1 + (2^{1} - 1) \cdot 1 + (3^{1} - 2^{1} - (2^{1} - 1)) \cdot 1$	14	$14 - 1 \cdot 7$
2	$2 - 1 \cdot 1$	15	$15 + (2^4 - 1) \cdot 1 + (3^4 - 2^4 - (2^4 - 1)) \cdot 1$
3	$3 + (2^2 - 1) \cdot 1 + (3^2 - 2^2 - (2^2 - 1)) \cdot 1$	16	$16 - 15 \cdot 1$
4	$4 - 3 \cdot 1$	17	$17 + (2^1 - 1) \cdot 9 + (3^1 - 2^1 - (2^1 - 1)) \cdot 9$
5	$5 + (2^1 - 1) \cdot 3 + (3^1 - 2^1 - (2^1 - 1)) \cdot 3$	18	18 - 9
6	$6 - 1 \cdot 3$	19	$19 + (2^2 - 1) \cdot 5 + (3^2 - 2^2 - (2^2 - 1)) \cdot 5$
7	$7 + (2^3 - 1) \cdot 1 + (3^3 - 2^3 - (2^3 - 1)) \cdot 1$	20	$20 - 3 \cdot 5$
8	$8-7\cdot 1$	21	$21 + (2^1 - 1) \cdot 11 + (3^1 - 2^1 - (2^1 - 1)) \cdot 11$
9	$9 + (2^1 - 1) \cdot 5 + (3^1 - 2^1 - (2^1 - 1)) \cdot 5$	22	22 - 11
10	$10 - 1 \cdot 5$	23	$23 + (2^3 - 1) \cdot 3 + (3^3 - 2^3 - (2^3 - 1)) \cdot 3$
11	$11 + (2^2 - 1) \cdot 3 + (3^2 - 2^2 - (2^2 - 1)) \cdot 3$	24	$24 - 7 \cdot 3$
12	$12 - 3 \cdot 3$	25	$25 + (2^{1} - 1) \cdot 13 + (3^{1} - 2^{1} - (2^{1} - 1)) \cdot 13$
13	$13 + (2^{1} - 1) \cdot 7 + (3^{1} - 2^{1} - (2^{1} - 1)) \cdot 7$	26	26 - 13

Table 7: Reformulation as expressions that contain power operations for every  $n_0 \in \mathbb{N}_o$ .

with power operations such as

$$3 + 3 \cdot 1 + 2 \cdot 1 = 3 + (2^{2} - 1) \cdot 1 + (3^{2} - 2^{2} - 3) \cdot 1 = 3 + (2^{2} - 1) \cdot 1 + (3^{2} - 2^{2} - (2^{2} - 1)) \cdot 1$$
$$7 + 7 \cdot 1 + 12 \cdot 1 = 7 + (2^{3} - 1) \cdot 1 + (3^{3} - 2^{3} - 7) \cdot 1 = 7 + (2^{3} - 1) \cdot 1 + (3^{3} - 2^{3} - (2^{3} - 1)) \cdot 1$$
$$15 + 15 \cdot 1 + 50 \cdot 1 = 15 + (2^{4} - 1) \cdot 1 + (3^{4} - 2^{4} - 15) \cdot 1 = 15 + (2^{4} - 1) \cdot 1 + (3^{4} - 2^{4} - (2^{4} - 1)) \cdot 1$$

and so on, listed in Table 7 for every  $n_0 \in \mathbb{N}_o$  up until  $n_0 = 26$ . We further observe that a similar approach can be derived for each For  $n_0 \in \mathbb{N}_e$  subsequence giving

$$2 - (2^{1} - 1) \cdot 1,$$
  

$$4 - (2^{2} - 1) \cdot 1,$$
  

$$8 - (2^{3} - 1) \cdot 1,$$

and so on. Table 8 combines both approaches, completing the expansion for  $n_0 \in \mathbb{N}^*$  up until  $n_0 = 26$ .

As such,  $s(n_t)$  from Theorem 2 applies an expansion of two terms during increments and an expansion of one term during decrements. Given this expanded form we can symbolically derive a separate piecewise function for iterates of  $c(n_t)$  namely

$$F(n_t, p, d) = \begin{cases} -d \cdot (2^p - 1) & \text{if } n_t \equiv 0 \pmod{2} \\ +d \cdot (2^p - 1) + d \cdot (3^p - 2^p - (2^p - 1)) & \text{if } n_t \equiv 1 \pmod{2} \end{cases}$$

$n_0$	$n_0 + n_1$	$n_0$	$  n_0 + n_1$
1	$1 + (2^1 - 1) \cdot 1 + (3^1 - 2^1 - (2^1 - 1)) \cdot 1$	14	$14 - (2^1 - 1) \cdot 7$
2	$2 - (2^1 - 1) \cdot 1$	15	$15 + (2^4 - 1) \cdot 1 + (3^4 - 2^4 - (2^4 - 1)) \cdot 1$
3	$3 + (2^2 - 1) \cdot 1 + (3^2 - 2^2 - (2^2 - 1)) \cdot 1$	16	$16 - (2^4 - 1) \cdot 1$
4	$4 - (2^2 - 1) \cdot 1$	17	$17 + (2^1 - 1) \cdot 9 + (3^1 - 2^1 - (2^1 - 1)) \cdot 9$
5	$5 + (2^1 - 1) \cdot 3 + (3^1 - 2^1 - (2^1 - 1)) \cdot 3$	18	$18 - (2^1 - 1) \cdot 9$
6	$6 - (2^1 - 1) \cdot 3$	19	$19 + (2^2 - 1) \cdot 5 + (3^2 - 2^2 - (2^2 - 1)) \cdot 5$
7	$7 + (2^3 - 1) \cdot 1 + (3^3 - 2^3 - (2^3 - 1)) \cdot 1$	20	$20 - (2^2 - 1) \cdot 5$
8	$8 - (2^3 - 1) \cdot 1$	21	$21 + (2^1 - 1) \cdot 11 + (3^1 - 2^1 - (2^1 - 1)) \cdot 11$
9	$9 + (2^{1} - 1) \cdot 5 + (3^{1} - 2^{1} - (2^{1} - 1)) \cdot 5$	22	$22 - (2^1 - 1) \cdot 11$
10	$10 - (2^1 - 1) \cdot 5$	23	$23 + (2^3 - 1) \cdot 3 + (3^3 - 2^3 - (2^3 - 1)) \cdot 3$
11	$11 + (2^2 - 1) \cdot 3 + (3^2 - 2^2 - (2^2 - 1)) \cdot 3$	24	$24 - (2^3 - 1) \cdot 3$
12	$12 - (2^2 - 1) \cdot 3$	25	$25 + (2^{1} - 1) \cdot 13 + (3^{1} - 2^{1} - (2^{1} - 1)) \cdot 13$
13	$13 + (2^{1} - 1) \cdot 7 + (3^{1} - 2^{1} - (2^{1} - 1)) \cdot 7$	26	$26 - (2^1 - 1) \cdot 13$

Table 8: Combination of both approaches  $n_0 \in \mathbb{N}^*$ .

where p > 0 denotes power and d > 0 the divergence index for F, i.e. the first occurrence of an increment of d in n starting at 1. The additional  $2^p - 1 \mid n_t \equiv 1 \pmod{2}$  can be factored out to obtain

$$F(n_t, p, d) = (n_t \equiv 1) d \cdot (3^p - 2^p) - (n_t \equiv 0) d \cdot (2^p - 1) \pmod{2}$$

that emphasises a modular symmetry between increments and decrements from the parityalternating formulation of  $c(n_t)$ . As an example, the orbit for  $n_t = 7$  converges to 1 as follows

$$7 + F(7,3,1) + F(26,1,13) + F(13,1,7) + F(20,2,5) + F(5,1,3) + F(8,3,1) = 1.$$

Table 9 shows divergence subsequences for  $1 \le p \le 4$  for  $n \in \mathbb{N}^*$  up until n = 26, while Figure 5 combines subsequences into one chart up until n = 100 for illustration purposes.

Suppose  $F(n_t, p, d)$  can be merged back into the metric space of  $c(n_t)$ , then it should be possible to make both p and d functions of  $n_t$ . For p we observe that it satisfies the sequence

$$(1, 1, 2, 2, 1, 1, 3, 3, 1, 1, 2, 2, 1, 1, 4, 4, ...)$$
.

which is OEIS A050603 [4]<sup>†</sup>. For p as a function of  $n_t$  we thus obtain

$$p(n_t) = \nu_2(n_t + 1) + \nu_2(n_t) = \nu_2(n_t^2 + n_t) = \nu_2(n_t + (n_t \equiv 1)) \pmod{2}$$

<sup>&</sup>lt;sup>†</sup>Also OEIS A136480 [5] minus one.

n	F(n,1,n)	F(n,2,n)	F(n,3,n)	F(n,4,n)
1	1	5	19	65
2	-2	-6	-14	-30
3	3	15	57	195
4	-4	-12	-28	-60
5	5	25	95	325
6	-6	-18	-42	-90
7	7	35	133	455
8	-8	-24	-56	-120
9	9	45	171	585
10	-10	-30	-70	-150
11	11	55	209	715
12	-12	-36	-84	-180
13	13	65	247	845
14	-14	-42	-98	-210
15	15	75	285	975
16	-16	-48	-112	-240
17	17	85	323	1105
18	-18	-54	-126	-270
19	19	95	361	1235
20	-20	-60	-140	-300
21	21	105	399	1365
22	-22	-66	-154	-330
23	23	115	437	1495
24	-23	-72	-168	-360
25	25	125	475	1625
26	-26	-78	-182	-390

Table 9: Divergent subsequences for powers  $1 \le p \le 4$  for  $n \in \mathbb{N}^*$  up until n = 26.



Figure 5: Chart showing the metric spaces of  $1 \le p \le 4$  for  $1 \le n \le 100$ .

for the exponential p-terms of F where  $p(n_t) = p(n_t + 1)$  for all consecutive pairs of  $\langle n_t, n_t + 1 \rangle \in \mathbb{N}^* | n_t \in \mathbb{N}_o \land n_t + 1 \in \mathbb{N}_e$ . As proven in Theorem 2, the equivalence  $k = \nu_2(n_t + 1)$  within  $s(n_{t,k})$  is the lookahead power when  $n_t \in \mathbb{N}_o$  while  $\nu_2(n_t)$  is the power that denotes the maximum number of factors of two contained in every  $n_t \in \mathbb{N}_e$ , as such  $p(n_t)$  combines both powers for all  $n_t \in \mathbb{N}^*$ , equating the sequence above.

For the divergence index d, we also define a function that satisfies the following sequence

$$(1, 1, 1, 1, 3, 3, 1, 1, 5, 5, 3, 3, 7, 7, 1, 1, ...)$$

for  $n_t \in \mathbb{N}^*$ . We observe this to be similar to the "change point" sequence as given in OEIS A163575 [6] albeit obtained with the ceiling instead of the flooring function. This yields the following

$$d^{\lceil}(n_t) = \left\lceil \frac{n_t}{2^{p(n_t)}} \right\rceil = \frac{n_t + (n_t \equiv 1)}{2^{p(n_t)}} \pmod{2}$$

where  $\lceil$  denotes the ceiling variety. Following Corollary 5,  $d^{\lceil}(n_t) \mid n_t \in \mathbb{N}_e$  is  $\frac{n_t}{2^{\nu_2(n_t)}}$  that is equivalent to  $c(n_t) \mid n_t \in \mathbb{N}_e$ . While for  $n_t \in \mathbb{N}_o$  it is  $\frac{n_t+1}{2^{\nu_2(n_t+1)}}$  that is equivalent to the denominator in  $s(n_t)$  from Theorem 2 and the equivalence

$$\frac{3^{\nu_2(n_t+1)}(n_t)+3^{\nu_2(n_t+1)}}{2^{\nu_2(n_t+1)}}=\frac{3^{\nu_2(n_t+1)}(n_t+1)}{2^{\nu_2(n_t+1)}}.$$

Both  $p(n_t)$  and  $d^{\lceil}(n_t)$  exhibit cycles for occurrences due to Corollary 5. For  $p(n_t)$  its cycles are of equal 2-adic valuation distances, where the metric is defined as the following function

$$D(n) = n + 2^{p(n)+1}$$

that would reach the same subsequence of p(n), mapping a metric to the next occurrence of a subsequence, i.e. to a proceeding pair  $\langle n, n+1 \rangle$  equal with respect to p(n). Thus arriving at the next pair that contains the same 2-adic valuation outcome as the current pair. E.g.

$$D(7) = 23, D(9) = 13, D(11) = 19$$
  
 $p(7) = 3, p(23) = 3, p(9) = 1, p(13) = 1, p(11) = 2, p(19) = 2$ 

for natural odd numbers, or

$$D(8) = 24, D(10) = 14, D(12) = 20$$
  
 $p(8) = 3, p(24) = 3, p(10) = 1, p(14) = 1, p(12) = 2, p(20) = 2$ 

for natural even numbers. For  $d^{\lceil}(n_t)$  we designate it the label of a divergence index as it embeds a counter for D(n) with a fixed metric of 2 starting at 1, i.e. for the first encounter of a pair in the sequence of p(n). E.g.

$$p(1) = 1, \ p(5) = 1, \ p(9) = 1, \ p(13) = 1$$
  
 $d^{\lceil}(1) = 1, \ d^{\lceil}(5) = 3, \ d^{\lceil}(9) = 5, \ d^{\lceil}(13) = 7$ 

at the onset of  $2^1 - 1$ , or

$$p(3) = 2, \ p(11) = 2, \ p(19) = 2, \ p(27) = 2$$
  
 $d^{\lceil}(3) = 1, \ d^{\lceil}(11) = 3, \ d^{\lceil}(19) = 5, \ d^{\lceil}(27) = 7$ 

at the onset of  $2^2 - 1$ .

Finally, we can combine and simplify it all together with  $n_t$ , which gives

$$c_1(n_t) := n_t + (n_t \equiv 1) d^{\lceil}(n_t) \left(3^{p(n_t)} - 2^{p(n_t)}\right) - (n_t \equiv 0) d^{\lceil}(n_t) \left(2^{p(n_t)} - 1\right) \pmod{2}$$

that embeds orbitals  $\forall n_t \in \mathbb{N}^*$ . E.g., for  $n_0 = 7$  we have the iterations

$$\underbrace{c_1 \left( c_1 \left$$



Figure 6: Subsequences  $1 \le p(n_t) \le 5$  of F contained within  $c(n_t)$  for  $1 \le n_t \le 100$ .

Figure 6 shows divergence subsequences  $F(n_t, p, d)$  when mapped back into  $c_1(n_t)$  together with an identity sequence  $n_t = n_t$  (1-slope). Figure 7 shows the same result but excluding the addition term  $n_t$  that is given in  $c_1(n_t)$ , thereby isolating iterative operations, and including a constant sequence  $n_t = 0$  (0-slope).

**Remark 6.** Subscript 1 applied on  $c(n_t)$  as  $c_1(n_t)$  denotes single-step iterations in the parity-alternating formulation of  $c(n_t)$ , also seen in A. E.g.,  $c_1(7) = 26$  followed by  $c_1(26) = 13$ .

**Corollary 9.** Given  $c_1(n_t)$  where  $n_{t-1} \in \mathbb{N}_o$  and  $n_t \in \mathbb{N}_e$  for t > 0, we illustrate the embedding of seemingly continuous 2-adic subspaces that exclude 0, or  $\mathbb{Z}_2 \setminus 0$  in Table 10 where rows represent  $3^{p(n_{t-1})} - 2^{p(n_{t-1})}$ , columns represent  $d(n_{t-1}) \in \mathbb{N}_o$ , and each cell the 2-adic valuation of odd operations within  $c_1(n_t)$  namely  $\nu_2 \left(n_{t-1} + d(n_{t-1}) \left(3^{p(n_{t-1})} - 2^{p(n_{t-1})}\right)\right)$ . Followed by Table 11 that lists the same results but without the  $\nu_2$  operation. Both tables are limited to  $p(n_{t-1}) \leq 17$  and  $d(n_{t-1}) \leq 33$  for illustrative purposes.

**Lemma 3.** Through Corollary 9 we derive a reformulation of the orbit-transforming function  $c_1(n_t)$  as defined in Theorem 3 with a correspondence to Corollary 7.

*Proof.* Let us first consider  $c_1(n_t) \forall n_t \in \mathbb{N}_o$ . Table 11 illustrates ever-expanding "jumps" within  $\mathbb{N}^*$ . To illustrate this further, indices are listed in Table 12 for  $1 \leq n \leq 32$  and for

$3^{p(n_{l-1})} - 2^{p(n_{l-1})}$	1	3	Ð	- (	9	11	13	15	11	19	21	23	25	27	29	31	- 33
$3^1 - 2^1$	1	3	1	2	1	5	1	2	1	3	1	2	1	4	1	2	1
$3^2 - 2^2$	3	1	2	1	4	1	2	1	3	1	2	1	5	1	2	1	3
$3^3 - 2^3$	1	4	1	2	1	3	1	2	1	9	1	2	1	3	1	2	1
$3^4 - 2^4$	4	1	2	1	3	1	2	1	5	1	2	1	3	1	2	1	4
$3^5 - 2^5$	1	3	1	2	1	4	1	2	1	3	1	2	1	5	1	2	1
$3^6 - 2^6$	3	1	2	1	5	1	2	1	3	1	2	1	4	1	2	1	3
$3^7 - 2^7$	1	5	1	2	1	3	1	2	1	4	1	2	1	3	1	2	1
$3^8 - 2^8$	5	1	2	1	3	1	2	1	4	1	2	1	3	1	2	1	6
$3^9 - 2^9$	1	3	1	2	1	6	1	2	1	3	1	2	1	4	1	2	1
$3^{10} - 2^{10}$	3	1	2	1	4	1	2	1	3	1	2	1	7	1	2	1	3
$3^{11} - 2^{11}$	1	4	1	2	1	3	1	2	1	5	1	2	1	3	1	2	1
$3^{12} - 2^{12}$	4	1	2	1	3	1	2	1	8	1	2	1	3	1	2	1	4
$3^{13} - 2^{13}$	1	3	1	2	1	4	1	2	1	3	1	2	1	6	1	2	1
$3^{14} - 2^{14}$	3	1	2	1	6	1	2	1	3	1	2	1	4	1	2	1	3
$3^{15} - 2^{15}$	1	6	1	2	1	3	1	2	1	4	1	2	1	3	1	2	1
$3^{16} - 2^{16}$	6	1	2	1	3	1	2	1	4	1	2	1	3	1	2	1	5
$3^{17} - 2^{17}$	1	3	1	2	1	5	1	2	1	3	1	2	1	4	1	2	1

 $\begin{vmatrix} 3^{p(n_{t-1})} - 2^{p(n_{t-1})} & | & 1 & | & 3 & | & 5 & | & 7 & | & 9 & | & 11 & | & 13 & | & 15 & | & 17 & | & 19 & | & 21 & | & 23 & | & 25 & | & 27 & | & 29 & | & 31 & | & 33 & | \\ \end{vmatrix}$ 

Table 10: Outcomes of  $\nu_2\left(n_{t-1} + d(n_{t-1})\left(3^{p(n_{t-1})} - 2^{p(n_{t-1})}\right)\right)$  for  $1 \le p(n_{t-1}) \le 17$  and  $1 \le d(n_{t-1}) \le 33$ .

$3^{p(n_{t-1})} - 2^{p(n_{t-1})}$	1	3	5	7	9	11	13	15	17	19	21	23	25	27	29	31	33
$3^1 - 2^1$	2	8	14	20	26	32	38	44	50	56	62	68	74	80	86	92	98
$3^2 - 2^2$	8	26	44	62	80	98	116	134	152	170	188	206	224	242	260	278	296
$3^3 - 2^3$	26	80	134	188	242	296	350	404	458	512	566	620	674	728	782	836	890
$3^4 - 2^4$	80	242	404	566	728	890	1052	1214	1376	1538	1700	1862	2024	2186	2348	2510	2672
$3^5 - 2^5$	242	728	1214	1700	2186	2672	3158	3644	4130	4616	5102	5588	6074	6560	7046	7532	8018
$3^6 - 2^6$	728	2186	3644	5102	6560	8018	9476	10934	12392	13850	15308	16766	18224	19682	21140	22598	24056
$3^7 - 2^7$	2186	6560	10934	15308	19682	24056	28430	32804	37178	41552	45926	50300	54674	59048	63422	67796	72170
$3^8 - 2^8$	6560	19682	32804	45926	59048	72170	85292	98414	111536	124658	137780	150902	164024	177146	190268	203390	216512
$3^9 - 2^9$	19682	59048	98414	137780	177146	216512	255878	295244	334610	373976	413342	452708	492074	531440	570806	610172	649538
$3^{10} - 2^{10}$	59048	177146	295244	413342	531440	649538	767636	885734	1003832	1121930	1240028	1358126	1476224	1594322	1712420	1830518	1948616
$3^{11} - 2^{11}$	177146	531440	885734	1240028	1594322	1948616	2302910	2657204	3011498	3365792	3720086	4074380	4428674	4782968	5137262	5491556	5845850
$3^{12} - 2^{12}$	531440	1594322	2657204	3720086	4782968	5845850	6908732	7971614	9034496	10097378	11160260	12223142	13286024	14348906	15411788	16474670	17537552
$3^{13} - 2^{13}$	1594322	4782968	7971614	11160260	14348906	17537552	20726198	23914844	27103490	30292136	33480782	36669428	39858074	43046720	46235366	49424012	52612658
$3^{14} - 2^{14}$	4782968	14348906	23914844	33480782	43046720	52612658	62178596	71744534	81310472	90876410	100442348	110008286	119574224	129140162	138706100	148272038	157837976
$3^{15} - 2^{15}$	14348906	43046720	71744534	100442348	129140162	157837976	186535790	215233604	243931418	272629232	301327046	330024860	358722674	387420488	416118302	444816116	473513930
$3^{16} - 2^{16}$	43046720	129140162	215233604	301327046	387420488	473513930	559607372	645700814	731794256	817887698	903981140	990074582	1076168024	1162261466	1248354908	1334448350	1420541792
$3^{17} - 2^{17}$	129140162	387420488	645700814	903981140	1162261466	1420541792	1678822118	1937102444	2195382770	2453663096	2711943422	2970223748	3228504074	3486784400	3745064726	4003345052	4261625378

Table 11: Outcomes of  $n_{t-1} + d(n_{t-1}) \left( 3^{p(n_{t-1})} - 2^{p(n_{t-1})} \right)$  for  $1 \le p(n_{t-1}) \le 17$  and  $1 \le d(n_{t-1}) \le 33$ . Equal colours emphasise the distance expansions that take place on  $\mathbb{Z}_2$  from lower powers  $p(n_{t-1})$ .



Figure 7: Subsequences  $1 \le p(n_t) \le 5$  of F contained within  $c(n_t)$  but now without the addition term  $n_t$ , or  $(n \equiv 1) d^{\lceil}(n) (3^{p(n-2^{p(n)})} - (n \equiv 0) d^{\lceil}(n) (2^{p(n-1)}) \pmod{2}).$ 

$3^{p(n_{t-1})} - 2^{p(n_{t-1})}$	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32
$3^1 - 2^1$	-	1	-	-	-	-	-	2	-	-	-	-	-	3	-	-	-	-	-	4	-	-	-	-	-	5	-	-	_	-	-	6
$3^2 - 2^2$	-	-	-	-	-	-	-	1	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	2	-	-	-	_	-	-
$3^3 - 2^3$	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	1	-	-		_	-	-

Table 12: Indices per power expressed in  $c_1(n_t)|n_t \in \mathbb{N}_o$  for  $1 \le p \le 3$ , illustrating ever-expanding "jumps" in  $\mathbb{N}^*$ .

the exponential order of  $1 \le p \le 3$ . Suppose that the expansion holds  $\forall n_t \in \mathbb{N}_o$ , then there should be a correspondence with  $s(n_t)$  following previous conclusions brought forth in this paper. For p = 1 we have

$$1 \to 2, 2 \to 8, 3 \to 14, 4 \to 20, 5 \to 26, 6 \to 32, \dots$$

with a metric of 6 starting at table column 2. For p = 2 we have

$$1 \rightarrow 8, 2 \rightarrow 26, \dots$$

with a metric of 18 starting at table column 8, and finally we have p = 3 with a metric of 54 starting at table column 26. For p = 1 we observe that the metric of 6 can be applied on  $\mathbb{N}^*$  with  $3(3^0 \cdot n) - 1|n \in \mathbb{N}_o$ . For example  $3(3^0 \cdot 1) - 1 = 2$ ,  $3(3^0 \cdot 3) - 1 = 8$ , and so on. For p = 2 we similarly can apply the metric of 18 on  $\mathbb{N}^*$  with  $3(3^1 \cdot n) - 1|n \in \mathbb{N}_o$ . For example  $3(3^1 \cdot 1) - 1 = 8$ ,  $3(3^0 \cdot 3) - 1 = 26$ , and so on. Considering the increase in power

to p = 3 we expect  $3(3^2n) - 1 \mid n \in \mathbb{N}_o$  to hold. This is true because

$$3(3^2n) - 1 = 3^3n - 1 = 3^pn - 1 \mid n \in \mathbb{N}_o$$

Considering that the range of  $d^{\lceil}(n_t)$  is equal to  $\{2n-1 : n \in \mathbb{N}\}$ , we can substitute n and p with their functional counterparts to obtain

$$3^{p(n)}d^{\dagger}(n) - 1 \mid n \in \mathbb{N}_o.$$

For cases  $n \in \mathbb{N}_e$  of  $c_1(n_t)$  we derive this to be  $d^{\lceil}(n)$  as follows. From  $c_1(n_t)$  as defined in Theorem 3 we have the following expression

$$n - d^{\lceil}(n) \left(2^{p(n)} - 1\right)$$

for which we can flatten to obtain

$$n - \frac{n\left(2^{p(n)} - 1\right)}{2^{p(n)}} = n - \frac{n\left(2^{\nu_2(n)} - 1\right)}{2^{\nu_2(n)}} = n - \frac{2^{\nu_2(n)}n - n}{2^{\nu_2(n)}} = n - \frac{n}{2^{\nu_2(n)}}$$

resulting in  $d^{\lceil}(n_t) \mid n_t \in \mathbb{N}_e$  as a function of  $\frac{n_t}{2^{\nu_2(n_t)}}$  that corresponds with Corollary 7. Both equations can now be combined with an addition operator to obtain

$$c_1(n_t) := (n_t \equiv 1) d^{\lceil}(n_t) 3^{p(n_t)} - (n_t \equiv 1) + (n_t \equiv 0) d^{\lceil}(n_t) \pmod{2}$$

To complete the proof we demonstrate the correspondence with  $s(n_t)$  from Corollary 7 for  $n_t \in \mathbb{N}_o$ . Flattening out gives

$$\left[\frac{n_t}{2^{\nu_2(n_t+1)}}\right]3^{\nu_2(n_t+1)} - 1 = \frac{3^{\nu_2(n_t+1)}(n_t+1)}{2^{\nu_2(n_t+1)}} - 1 = s(n_t) \mid n_t \in \mathbb{N}_o$$

as claimed.

**Corollary 10.** The derivation of  $c_1$  from Lemma 3 can be compacted further.

*Proof.* Considering sequence

$$n_1 = (2, 1, 8, 1, 8, 3, 26, 1, 14, 5, 26, 3, 20, 7, 80, 1, \dots)$$

resulting from  $c_1(n_0)$ , we have for  $n_t \in \mathbb{N}_o$ 

$$d^{\lceil}(n_t)3^{p(n_t)} - 1 = d^{\lceil}(n_t)3^{\nu_2(n_t) + \nu_2(n_t+1)} - 1$$

Further considering that  $3^{\nu_2(n_t)} = 1 \mid n_t \in \mathbb{N}_o$  and  $3^{\nu_2(n_t+1)} = 1 \mid n_t \in \mathbb{N}_e$ , then

$$c_1(n_t) := d^{\lceil}(n_t) 3^{\nu_2(n_t+1)} - (n_t \equiv 1) \pmod{2}$$

for the parity-alternating formulation of c in single-step iterations.

**Lemma 4.** Given Lemma 3 and Corollary 10, we can derive  $c_2(n_t)$  by wrapping and embedding the second iteration within  $c_1(n_t)$ .

Proof. If  $n_0 \in \mathbb{N}_o$  then  $c_1(n_0) = n_1 \in \mathbb{N}_e$ . Vice versa, if  $n_0 \in \mathbb{N}_e$  then  $c_1(n_0) = n_1 \in \mathbb{N}_o$ . As such, to iterate in double-steps we can wrap  $d^{\lceil}(n_t) \mid n_{t-1} \in \mathbb{N}_o$  and, following Lemma 3, embed  $d^{\lceil}(n_t)$  within itself and within  $p(n_t)$  when  $n_{t-1} \in \mathbb{N}_e$ . This yields

$$c_2(n_t) := (n_t \equiv 1) d_1^{\lceil} \left( d_1^{\lceil}(n_t) 3^{p_1(n_t)} - 1 \right) + (n_t \equiv 0) d_1^{\lceil} \left( d_1^{\lceil}(n_t) \right) 3^{p_1(d_1^{\lceil}(n_t))} - (n_t \equiv 0) \pmod{2}$$

that results for example in  $n_0 = 7 \rightarrow n_1 = 13$  crossing over 26, or  $n_0 = 6 \rightarrow n_1 = 8$  crossing under 3 through double-step iterations.

**Remark 7.** Subscript 2 applied as  $c_2(n_t)$  denotes double-step iterations in the parityalternating formulation of  $c(n_t)$ , also seen in A. For example,  $c_2(7) = 13$  followed by  $c_2(13) = 5$ . Similarly, subscript 1 applied on d and p as  $p_1(n_t)$  and  $d_1^{\lceil}(n_t)$ , denote the power variant of the original definition  $p(n_t)$  and the divergent index variant of the original definition  $d^{\lceil}(n_t)$  respectively.

Let

$$(n_t \equiv 0) d_1^{\lceil} \left( d_1^{\lceil}(n_t) \right) 3^{\nu_2 \left( d_1^{\lceil}(n_t) + 1 \right)} - (n_t \equiv 0) \pmod{2}$$

define the operation  $c_2(n_0) = n_2 \mid n_2 \in \mathbb{N}_e$ . First, we define a function for  $\nu_2\left(d_1^{\lceil}(n_t) + 1\right)$  as

$$p_{2}(n_{t}) = \nu_{2} \left( d_{1}^{\lceil}(n_{t}) + 1 \right)$$
$$= \nu_{2} \left( \frac{n_{t}}{2^{\nu_{2}(n_{t})}} + 1 \right)$$
$$= \nu_{2} \left( \frac{n_{t} + 2^{\nu_{2}(n_{t})}}{2^{\nu_{2}(n_{t})}} \right)$$
$$= \nu_{2}(n_{t} + 2^{\nu_{2}(n_{t})}) - \nu_{2}(2^{\nu_{2}(n_{t})})$$
$$= \nu_{2}(n_{t} + 2^{\nu_{2}(n_{t})}) - \nu_{2}(n_{t})$$

where the last two right-hand expressions follow from an identity extension for  $\mathbb{Q}$  as given in [1] (p. 2), and forming the sequence OEIS A089309 [7]

$$(1, 1, 2, 1, 1, 2, 3, 1, 1, 1, 2, 2, 1, 3, 4, 1, ...)$$
.

Intuitively, we recognise  $p_2$  as the 2-adic valuation of n from  $\mathbb{N}_o \to \mathbb{N}_e \ \forall n \in \mathbb{N}^*$ . Furthermore,  $p_2(n) = p_1(n) \mid n \in \mathbb{N}_o$  because

$$\nu_2 (n+1) - 0 = \nu_2 (n+1) + 0$$

For  $d_1^{\lceil} \left( d_1^{\rceil}(n_t) \right)$  we can define the function  $d_2^{\rceil}(n_t)$  as

$$d_2^{\lceil}(n_t) = d_1^{\lceil} \left( d_1^{\lceil}(n_t) \right)$$

and reformulate the embedding by first flattening, followed by a formulation for  $n_t \in \mathbb{N}_e$ then by  $n_t \in \mathbb{N}_o$ , to finally combine both simplifications. To start, let

$$\begin{split} d_{2}^{\lceil}(n_{t}) &= d_{1}^{\lceil} \left( d_{1}^{\lceil}(n_{t}) \right) \\ &= d_{1}^{\lceil} \left( \frac{n_{t} + (n_{t} \equiv 1)}{2^{\nu_{2}(n_{t}) + \nu_{2}(n_{t}+1)}} \right) \\ &= \frac{n_{t} + (n_{t} \equiv 1)}{\frac{2^{\nu_{2}(n_{t}) + \nu_{2}(n_{t}+1)}}{2^{\nu_{2}(n_{t}) + \nu_{2}(n_{t}+1)}} + \nu_{2} \left( \frac{n_{t} + (n_{t} \equiv 1)}{2^{\nu_{2}(n_{t}) + \nu_{2}(n_{t}+1)}} + 1 \right) + \nu_{2} \left( \frac{n_{t} + (n_{t} \equiv 1)}{2^{\nu_{2}(n_{t}) + \nu_{2}(n_{t}+1)}} + 1 \right) + \frac{\left( \frac{n_{t} + (n_{t} \equiv 1)}{2^{\nu_{2}(n_{t}) + \nu_{2}(n_{t}+1)}} \equiv 1 \right)}{2^{\nu_{2} \left( \frac{n_{t} + (n_{t} \equiv 1)}{2^{\nu_{2}(n_{t}) + \nu_{2}(n_{t}+1)}} \right) + \nu_{2} \left( \frac{n_{t} + (n_{t} \equiv 1)}{2^{\nu_{2}(n_{t}) + \nu_{2}(n_{t}+1)}} + 1 \right)} \end{split}$$

(mod 2). For the case of  $n_t \in \mathbb{N}_e$  we can cancel out all 0's and obtain

$$\frac{\frac{n_t}{2^{\nu_2(n_t)}}}{2^{\nu_2}\left(\frac{n_t}{2^{\nu_2(n_t)}}+1\right)} + \frac{1}{2^{\nu_2}\left(\frac{n_t}{2^{\nu_2(n_t)}}+1\right)}$$

that through the identity given in  $p_2(n_t)$  can be simplified to

$$\frac{n_t + 2^{\nu_2(n_t)}}{2^{\nu_2(n_t + 2^{\nu_2(n_t)})}}$$

forming the sequence

$$(1, 1, 1, 1, 3, 1, 1, 1, 5, 3, 3, 1, 7, 1, 1, 1, ...).$$

Similarly, for the case of  $n_t \in \mathbb{N}_o$ , we can cancel out all 0's and obtain

$$\frac{\frac{n_t+1}{2^{\nu_2(n_t+1)}}}{2^{\nu_2}\left(\frac{n_t+1}{2^{\nu_2(n_t+1)}}+1\right)} + \frac{1}{2^{\nu_2}\left(\frac{n_t+1}{2^{\nu_2(n_t+1)}}+1\right)}$$

which is simplified to

$$\frac{n_t + 1 + 2^{\nu_2(n_t+1)}}{2^{\nu_2\left(n_t + 1 + 2^{\nu_2(n_t+1)}\right)}}$$

that forms the sequence

$$(1, 1, 1, 3, 1, 1, 1, 5, 3, 3, 1, 7, 1, 1, 1, 9, ...)$$
.

Combining both together we arrive at

$$d_2^{\lceil}(n_t) = \frac{n_t + (n_t \equiv 1) + 2^{\nu_2(n_t + (n_t \equiv 1))}}{2^{\nu_2(n_t + (n_t \equiv 1) + 2^{\nu_2(n_t + (n_t \equiv 1))})}} \pmod{2}$$

that forms the sequence

$$(1, 1, 1, 1, 1, 1, 1, 1, 3, 3, 1, 1, 1, 1, 1, 1, ...)$$
.

For the case of  $c_2(n_0) = n_2 \mid n_2 \in \mathbb{N}_o$  we have

$$(n_t \equiv 1) d_1^{\lceil} \left( d_1^{\lceil}(n_t) 3^{p_1(n_t)} - 1 \right) = (n_t \equiv 1) d_1^{\lceil}(s(n_t)) \pmod{2}.$$

Flattening this out gives

$$\frac{\frac{3^{\nu_2(n_t+1)}(n_t+1)}{2^{\nu_2(n_t+1)}}-1}{2}$$

that through the negative exponentiation rule and  $p_2(n_t)$  identity we can expand as

$$3^{\nu_{2}(n_{t}+1)}2^{-\nu_{2}(n_{t}+1)-\nu_{2}\left(3^{\nu_{2}(n_{t}+1)}(n_{t}+1)-2^{\nu_{2}(n_{t}+1)}\right)+\nu_{2}\left(2^{\nu_{2}(n_{t}+1)}\right)}(n_{t}+1)$$
  
-2<sup>-\nu\_{2}\left(3^{\nu\_{2}(n\_{t}+1)}(n\_{t}+1)-2^{\nu\_{2}(n\_{t}+1)}\right)+\nu\_{2}\left(2^{\nu\_{2}(n\_{t}+1)}\right)}.</sup>

Because  $2^{\nu_2(2^{\nu_2(n_t+1)})} = 2^{\nu_2(n_t+1)}$  we can simplify to

$$\frac{3^{\nu_2(n_t+1)}(n_t+1)}{2^{\nu_2(3^{\nu_2(n_t+1)}(n_t+1)-2^{\nu_2(n_t+1)})}} - \frac{1}{2^{\nu_2(3^{\nu_2(n_t+1)}(n_t+1)-2^{\nu_2(n_t+1)})-\nu_2(n_t+1)}}$$

and finally combine to obtain

$$\frac{3^{\nu_2(n_t+1)}(n_t+1)-2^{\nu_2(n_t+1)}}{2^{\nu_2(3^{\nu_2(n_t+1)}(n_t+1)-2^{\nu_2(n_t+1)})}}.$$

Adding it all together we arrive at

$$c_2(n_t) := (n_t \equiv 1) \left( \frac{3^{\nu_2(n_t+1)} (n_t+1) - 2^{\nu_2(n_t+1)}}{2^{\nu_2(3^{\nu_2(n_t+1)}(n_t+1) - 2^{\nu_2(n_t+1)})}} \right) + (n_t \equiv 0) \left( d_2^{\lceil}(n_t) 3^{p_2(n_t)} - 1 \right) \pmod{2}$$

which gives the following sequence as its result

$$c_2(n_0) = (1, 2, 1, 2, 1, 8, 13, 2, 7, 8, 13, 8, 5, 26, 5, 2, ...).$$

$n_0$	$A_2$ sequences given $c_2(n_t)$	$n_0$	$A_2$ sequences given $c_2(n_t)$
1	θ	14	$\ominus \to \ominus \to \ominus \to \ominus \to \ominus$
2	θ	15	$\ominus \to \ominus \to \ominus$
3	$\ominus \rightarrow \ominus$	16	$\ominus \rightarrow \ominus$
4	$\ominus \to \ominus$	17	$\ominus \to \ominus \to \ominus \to \ominus$
5	$\ominus \rightarrow \ominus$	18	$\bigcirc \rightarrow \bigcirc \rightarrow \bigcirc \rightarrow \bigcirc \rightarrow \bigcirc \rightarrow \bigcirc \rightarrow \bigcirc$
6	$\ominus \to \ominus \to \ominus$	19	$\bigcirc \rightarrow \bigcirc \rightarrow \bigcirc \rightarrow \bigcirc \rightarrow \bigcirc$
7	$\ominus \to \ominus \to \ominus \to \ominus$	20	$\ominus \to \ominus \to \ominus$
8	$\ominus \rightarrow \ominus$	21	$\ominus \rightarrow \ominus$
9	$ \ominus \rightarrow \ominus \rightarrow \ominus \rightarrow \ominus \rightarrow \ominus $	22	$\bigcirc \rightarrow \bigcirc \rightarrow \bigcirc \rightarrow \bigcirc \rightarrow \bigcirc$
10	$\ominus \to \ominus \to \ominus$	23	$\ominus \to \ominus \to \ominus$
11	$\ominus \to \ominus \to \ominus \to \ominus$	24	$\bigcirc \to \ominus \to \ominus$
12	$\ominus \to \ominus \to \ominus$	25	$\bigcirc \rightarrow \oslash \rightarrow \ominus \rightarrow \ominus \rightarrow \ominus \rightarrow \ominus$
13	$\ominus \to \ominus \to \ominus$	26	$\bigcirc \to \bigcirc \to \bigcirc \to \bigcirc$

Table 13: Sequences of  $A_2$  given  $c_2(n_t)$  for  $1 \le n_0 \le 26$ .

**Conjecture 3.** Following Lemma 4 we conjecture nor more nor less than five sub-arc diagrams as part of A for double-step iterations denoted as  $A_2$  and symbolised as  $\bigcirc$  a **positive overpass**: an increment crossing the number line and remaining over  $n_{t-2}$ ,

 $\ominus$  a **positive underpass**: an increment crossing the number line and under  $n_{t-2}$ ,

 $\ominus$  a **negative overpass**: a decrement crossing the number line and remaining over  $n_{t-2}$ ,

 $\ominus$  a **negative underpass**: a decrement crossing the number line and under  $n_{t-2}$ , and

 $\bigcirc$  that symbolises a **cycle** circling back to  $n_{t-2}$ .

**Remark 8.** Increments or decrements in  $A_2$  also express the clock or anti-clockwise winding direction. In other words, for  $A_2$  under this conjecture, sequences with  $n_0 \in \mathbb{N}_o$  solely contain positive passes, while sequences with  $n_0 \in \mathbb{N}_e$  solely contain negative passes. Thus, "remaining over" or "under" in the previous statement takes this direction into account.

 $A_2$  sequences are listed in Table 13 for  $1 \le n_0 \le 26$  for illustration purposes. The orbital length of  $A_2$  of  $c_2(n_t)$  is equal to OEIS A160541 [8] minus one for  $n_0 > 2$ .

**Corollary 11.** Also following Lemma 4 we derive a rational function for  $c_1(n_t)$ .

*Proof.* When  $n_0$  is subtracted from outcomes of  $c_1(n_0)$  we isolate the operation of singlestep iteration and observe linear correspondences between elements (visible Figure 6). For

$n_0$	$\frac{ax+gf}{b} = y$	$n_0$	$\frac{ax+gf}{b} = y$
1	$\frac{2x+1\cdot 2}{4} = y$	14	$\frac{-2x+0\cdot 4}{4} = y$
2	$\frac{-2x+0\cdot 4}{4} = y$	15	$\frac{130x+65\cdot 2}{32} = y$
3	$\frac{10x+5\cdot 2}{8} = y$	16	$\frac{-30x+0\cdot32}{32} = y$
4	$\frac{-6x+0\cdot 8}{8} = y$	17	$\frac{2x+1\cdot 2}{4} = y$
5	$\frac{2x+1\cdot 2}{4} = y$	18	$\frac{-2x+0\cdot 4}{4} = y$
6	$\frac{-2x+0\cdot 4}{4} = y$	19	$\frac{10x+5\cdot 2}{8}$
7	$\frac{38x+19\cdot 2}{16} = y$	20	$\frac{-6x+0\cdot 8}{8} = y$
8	$\frac{-14x+0\cdot16}{16} = y$	21	$\frac{2x+1\cdot 2}{4} = y$
9	$\frac{2x+1\cdot 2}{4} = y$	22	$\frac{-2x+0\cdot 4}{4} = y$
10	$\frac{-2x+0\cdot 4}{4} = y$	23	$\frac{38x+19\cdot 2}{16} = y$
11	$\frac{10x+5\cdot 2}{8} = y$	24	$\frac{-14x+0\cdot16}{16} = y$
12	$\frac{-6x+0\cdot 8}{8} = y$	25	$\frac{2x+1\cdot 2}{4} = y$
13	$\frac{2x+1\cdot 2}{4} = y$	26	$\frac{-2x+0\cdot 4}{4} = y$

Table 14: Linear fractions within element pairs as seen in Figure 6, for  $1 \le n_0 \le 26 | c_1(n_t)$ .

each correspondence we can identify a linear fraction in the form of

$$\frac{ax+gf}{b}$$

with outcomes listed in Table 14. As  $n_0 \in \mathbb{N}_o$  of  $c_1(n_t)$  results in an  $n_t \in \mathbb{N}_e$  we can remove factors of two for each as listed in Table 15. For f we identify this in reduced form to be

$$f(n_0) = 2^{\nu_2(n_0)}$$

or  $f(n_0) = 2^{\nu_2(n_0)+1}$  in unreduced form. For g we have

$$g(n_0) = 3^{\nu_2(n_0+1)} - 2^{\nu_2(n_0+1)}$$

where  $g(n_0) = 0 | n_0 \in \mathbb{N}_e$ . For b we observe these to be powers of two forming the sequence

$$(2, 2, 4, 4, 2, 2, 8, 8, 2, 2, 4, 4, 2, 2, 16, 16, \ldots)$$

or

$$\left(2^1,2^1,2^2,2^2,2^1,2^1,2^3,2^3,2^1,2^1,2^2,2^2,2^1,2^1,2^4,2^4,\ldots\right)$$

where we denote the exponent as  $\delta_1(n_0)$  and in fact

$$\delta_1(n_0) = p_1(n_0).$$

$n_0$	$\frac{ax+gf}{b} = y$	$n_0$	$\frac{ax+gf}{b} = y$
1	$\frac{1x+1\cdot 1}{2} = y$	14	$\frac{-1x+0\cdot 2}{2} = y$
2	$\frac{-1x+0\cdot 2}{2} = y$	15	$\frac{65x + 65 \cdot 1}{16} = y$
3	$\frac{5x+5\cdot 1}{4} = y$	16	$\frac{-15x+0\cdot16}{16} = y$
4	$\frac{-3x+0\cdot 4}{4} = y$	17	$\frac{1x+1\cdot 1}{2} = y$
5	$\frac{1x+1\cdot 1}{2} = y$	18	$\frac{-1x+0\cdot 2}{2} = y$
6	$\frac{-1x+0\cdot 2}{2} = y$	19	$\frac{5x+5\cdot 1}{4}$
7	$\frac{19x+19\cdot 1}{8} = y$	20	$\frac{-3x+0\cdot 4}{4} = y$
8	$\frac{-7x+0\cdot 8}{8} = y$	21	$\frac{1x+1\cdot 1}{2} = y$
9	$\frac{1x+1\cdot 1}{2} = y$	22	$\frac{-1x+0\cdot 2}{2} = y$
10	$\frac{-1x+0\cdot 2}{2} = y$	23	$\frac{19x+19\cdot 1}{8} = y$
11	$\frac{5x+5\cdot 1}{4} = y$	24	$\frac{-7x+0\cdot 8}{8} = y$
12	$\frac{-3x+0\cdot 4}{4} = y$	25	$\frac{1x+1\cdot 1}{2} = y$
13	$\frac{1x+1\cdot 1}{2} = y$	26	$\frac{-1x+0\cdot 2}{2} = y$

Table 15: Simplified linear fractions within element pairs of  $c_1(n_t)$  as seen in Figure 6.

Finally for a we identify this to be

$$a(n_0) = 3^{\nu_2(n_0+1)} - 2^{\delta_1(n_0)}.$$

Bringing it all together we arrive at

$$\frac{\left(3^{\nu_2(n+1)}-2^{\delta_1(n)}\right)n+\left(3^{\nu_2(n+1)}-2^{\nu_2(n+1)}\right)2^{\nu_2(n)}}{2^{\delta_1(n)}}$$

where removing the  $2^{\delta_1(n_t)}$  term from  $a(n_t)$  gives

$$c_1(n_t) := \frac{3^{\nu_2(n_t+1)}n + \left(3^{\nu_2(n_t+1)} - 2^{\nu_2(n_t+1)}\right)2^{\nu_2(n_t)}}{2^{\delta_1(n_t)}}.$$

as claimed.

**Corollary 12.** Following Theorem 3, Lemma's 3, and 4, and Corollaries 10 and 11 we now derive a rational function for  $c_2(n_t) \forall n_t \in \mathbb{N}^*$ .

*Proof.* Similar to Corollary 11, we start by subtracting  $n_0$  from  $c_2(n_0)$  to isolate the operation of double-step iterations. Outcomes are listed in Table 16 and Figure 8. Similarly evident from Figure 8 are apparent 2-tuple outcomes, e.g.  $\langle 9, -2 \rangle$  with  $\langle 10, -2 \rangle$  or  $\langle 17, -4 \rangle$ 

$n_0$	$c_2(n_0) - n_0$	$n_0$	$c_2(n_0) - n_0$
1	1 - 1 = 0	14	26 - 14 = 12
2	2 - 2 = 0	15	5 - 15 = -10
3	1 - 3 = -2	16	2 - 16 = -14
4	2 - 4 = -2	17	13 - 17 = -4
5	1 - 5 = -4	18	14 - 18 = -4
6	8 - 6 = 2	19	11 - 19 = -8
7	13 - 7 = 6	20	8 - 20 = -12
8	2 - 8 = -6	21	1 - 21 = -20
9	7 - 9 = -2	22	26 - 22 = 4
10	8 - 10 = -2	23	5 - 23 = -18
11	13 - 11 = 2	24	8 - 24 = -16
12	8 - 12 = -4	25	19 - 25 = -6
13	5 - 13 = -8	26	20 - 26 = -6

Table 16: Subtraction of  $n_0$  terms from  $c_2(n_0)$  as given in Lemma 3 for  $1 \le n_0 \le 26$ .



Figure 8: Outcomes of  $c_2(n_0) - n_0$  for  $1 \le n_0 \le 100$ .

$n_0$	$\frac{a}{b}x + \frac{g}{h} = y$	$  n_0$	$\frac{a}{b}x + \frac{g}{h} = y$
1	$\frac{-2}{8}x + \frac{1}{4} = y$	14	$\frac{22}{32}x + \frac{19}{8} = y$
2	$\frac{-2}{8}x + \frac{1}{2} = y$	15	$\frac{-350}{512}x + \frac{65}{256} = y$
3	$\frac{-46}{64}x + \frac{5}{32} = y$	16	$\frac{-58}{64}x + \frac{1}{2} = y$
4	$\frac{-10}{16}x + \frac{1}{2} = y$	17	$\frac{-2}{8}x + \frac{1}{4} = y$
5	$\frac{-26}{32}x + \frac{1}{16} = y$	18	$\frac{-2}{8}x + \frac{1}{2} = y$
6	$\frac{2}{16}x + \frac{5}{4} = y$	19	$\frac{-14}{32}x + \frac{5}{16}$
7	$\frac{22}{32}x + \frac{19}{16} = y$	20	$\frac{-10}{16}x + \frac{1}{2} = y$
8	$\frac{-52}{64}x + \frac{1}{2} = y$	21	$\frac{-122}{128}x + \frac{1}{64} = y$
9	$\frac{-2}{8}x + \frac{1}{4} = y$	22	$\frac{2}{16}x + \frac{5}{4} = y$
10	$\frac{-2}{8}x + \frac{1}{2} = y$	23	$\frac{-202}{256}x + \frac{19}{128} = y$
11	$\frac{2}{16}x + \frac{5}{8} = y$	24	$\frac{-46}{64}x + \frac{5}{4} = y$
12	$\frac{-14}{32}x + \frac{5}{4} = y$	25	$\frac{-2}{8}x + \frac{1}{4} = y$
13	$\frac{-10}{16}x + \frac{1}{8} = y$	26	$\frac{-2}{8}x + \frac{1}{2} = y$

Table 17: Linear fractions within element pairs as seen in Figure 8, for  $1 \le n_0 \le 26 \mid c_2(n_t)$ .

with  $\langle 18, -4 \rangle$  with linear correspondences between outcomes of the first element and consequently of the second element of every pair. Solving for each correspondence symbolically we arrive at a set of partial fractions, listed in Table 17. These can be further simplified and combined in the form of

$$\frac{a \cdot n_0 + g \cdot f}{b}$$

, listed in Table 18. For f we observe this to remain  $f(n_0) = 2^{\nu_2(n_0)}$ .

For g we have the sequence

$$(1, 1, 5, 1, 1, 5, 19, 1, 1, 1, 5, 5, 1, 19, 65, ...)$$

which we identify to be  $g(n_0) = 3^{p_2(n_0)} - 2^{p_2(n_0)}$  with  $p_2(n_0)$  as derived in Lemma 4. For b we observe these to be powers of two forming the sequence

$$(4, 4, 32, 8, 16, 8, 16, 16, 4, 4, 8, 16, 8, 16, 256, 32, ...)$$

or

$$\left(2^2,2^2,2^5,2^3,2^4,2^3,2^4,2^4,2^2,2^2,2^3,2^4,2^3,2^4,2^8,2^5,\ldots\right)$$

$n_0$	$\frac{ax+gf}{b} = y$	$ n_0 $	$\frac{ax+gf}{b} = y$
1	$\frac{-1x+1\cdot 1}{4} = y$	14	$\frac{11x + 19 \cdot 2}{16} = y$
2	$\frac{-1x+1\cdot 2}{4} = y$	15	$\frac{-175x+65\cdot 1}{256} = y$
3	$\frac{-23x+5\cdot 1}{32} = y$	16	$\frac{-29x+1\cdot 16}{32} = y$
4	$\frac{-5x+1\cdot 4}{8} = y$	17	$\frac{-1x+1\cdot 1}{4} = y$
5	$\frac{-13x+1\cdot 1}{16} = y$	18	$\frac{-1x+1\cdot 2}{4} = y$
6	$\frac{1x+5\cdot 2}{8} = y$	19	$\frac{-7x+5\cdot 1}{16}$
7	$\frac{11x+19\cdot 1}{16} = y$	20	$\frac{-5x+1\cdot 4}{8} = y$
8	$\frac{-13x+1\cdot 8}{16} = y$	21	$\frac{-61x+1\cdot 1}{64} = y$
9	$\frac{-1x+1\cdot 1}{4} = y$	22	$\frac{1x+1\cdot 2}{8} = y$
10	$\frac{-1x+1\cdot 2}{4} = y$	23	$\frac{-101x+19\cdot 1}{128} = y$
11	$\frac{1x+5\cdot 1}{8} = y$	24	$\frac{-23x+5\cdot 8}{32} = y$
12	$\frac{-7x+5\cdot4}{16} = y$	25	$\frac{-1x+1\cdot 1}{4} = y$
13	$\frac{-5x+1\cdot 1}{8} = y$	26	$\frac{-1x+1\cdot 2}{4} = y$

Table 18: Simplified linear fractions within element pairs as seen in Figure 8, for  $1 \le n_0 \le 26 \mid c_2(n_t)$ .

where the exponent we have identified as the function

$$\delta_2(n_0) = p_1(c_1(n_0)) + p_1(n_0)$$
  
=  $\nu_2 \left( c_1(n_0)^2 + c_1(n_0) \right) + \nu_2 \left( n_0^2 + n_0 \right)$   
=  $\nu_2 \left( c_1(n_0) + \left( c_1(n_0) \equiv 1 \right) \right) + \nu_2 \left( n_0 + (n_0 \equiv 1) \right) \pmod{2}$ 

embedding  $c_1(n_t)$  as a function of  $p_1$  and where  $n_t + 2^{\delta_2(n_t)+1}$  represents the  $\Delta x$  metric to reach a pair  $\langle n_t, n_t + 1 \rangle$  of equal 2-adic valuation in  $\mathbb{N}^* \mid n_t \in \mathbb{N}_o \land n_t + 1 \in \mathbb{N}_e$  during double-step iterations. Removing  $\nu_2$  from  $\nu_2(c_1(n_0) + (c_1(n_0) \equiv 1)) \pmod{2}$  in  $p_1(c_1(n_0))$ yields the following sequence

$$(2, 2, 8, 2, 8, 4, 26, 2, 14, 6, 26, 4, 20, 8, 80, 2, ...)$$

while removing  $\nu_2$  from  $\nu_2 \left( c_1(n_0)^2 + c_1(n_0) \right)$  in  $p_1(c_1(n_0))$  yields the following sequence

(6, 2, 72, 2, 72, 12, 702, 2, 210, 30, 702, 12, 420, 56, 6480, 2, ...).

Because of the following p-adic property

$$\nu_{p}\left(xy\right) = \nu_{p}\left(x\right) + \nu_{p}\left(y\right)$$

where  $x \land y \in \mathbb{N}$  we have that

$$\nu_p (x+1) + \nu_p (y) = \nu_p ((x+1)y) = \nu_p ((xy+y))$$

As such, the result of the 2-adic valuation function holds even though the two integer sequences above are not equal.

For  $\delta_2(n_0)$  as the exponent of powers of two we thus complete the denominator as

$$b(n_0) = 2^{\delta_2(n_0)}.$$

Finally for a we identify this to be

$$a(n_0) = 3^{p_2(n_0)} - 2^{\delta_2(n_0)}.$$

**Remark 9.** Subscript 2 applied on  $\delta(n_0)$  as  $\delta_2(n_0)$  denotes the denominator-exponent variety of double-step iterations in the parity-alternating formulation of  $c(n_t)$ .

Adding it all together results in

$$\frac{\left(3^{p_2(n)}-2^{\delta_2(n)}\right)n+\left(3^{p_2(n)}-2^{p_2(n)}\right)2^{\nu_2(n)}}{2^{\delta_2(n)}}$$

Subtracting  $2^{\delta_2(n)}$  term from  $a(n_0)$  allows for the reformulation with  $n_t$  to obtain

$$c_2(n_t) := \frac{3^{p_2(n_t)} n_t + \left(3^{p_2(n_t)} - 2^{p_2(n_t)}\right) 2^{\nu_2(n_t)}}{2^{\delta_2(n_t)}}$$

which allows us to apply double-steps iteratively. For example for  $n_0 = 7$  we have

$$\underbrace{c_2\left(c_2\left(c_2(7)\right)\right)}_{3 \ times} = 1$$

that maps

$$7 \rightarrow 13 \rightarrow 5 \rightarrow 1$$

as claimed.

**Corollary 13.** We can remove a factor of  $2^{\nu_2(n_t)}$  from  $\frac{a \cdot n_0 + g \cdot f}{b}$  as given in Corollary 11 and Corollary 12 thereby reducing the number of unknowns from three to two fully.

*Proof.* Let

$$c_1(n_t) := \frac{3^{\nu_2(n_t+1)}n + \left(3^{\nu_2(n_t+1)} - 2^{\nu_2(n_t+1)}\right)2^{\nu_2(n_t)}}{2^{\delta_1(n_t)}}$$

We can include  $\frac{n_t}{2^{\nu_2(n_t)}}$  (sequence OEIS A000265 [9]) in  $c_1(n_t)$  as

$$c_1(n_t) := \frac{3^{\nu_2(n_t+1)} \frac{n_t}{2^{\nu_2(n_t)}} + 3^{\nu_2(n_t+1)} - 2^{\nu_2(n_t+1)}}{2^{\nu_2(n_t+1)}}$$

where the denominator is the result of  $2^{\delta_1(n_t)-\nu_2(n_t)}$ . Furthermore, from the simplifications that were identified in Theorem 2 we also obtain

$$c_1(n_t) := \frac{3^{\nu_2(n_t+1)} \left(\frac{n_t}{2^{\nu_2(n_t)}} + 1\right)}{2^{\nu_2(n_t+1)}} - 1$$

Similarly, for  $c_2(n_t)$  we also include  $\frac{n_t}{2^{\nu_2(n_t)}}$  and obtain

$$c_2(n_t) := \frac{3^{p_2(n_t)} \frac{n_t}{2^{\nu_2(n_t)}} + 3^{p_2(n_t)} - 2^{p_2(n_t)}}{2^{p_1(c_1(n_t)) + \nu_2(n_t+1)}}$$

where the denominator is the result of  $2^{\delta_2(n_t)-\nu_2(n_t)}$ . Here, we also have the identity

$$p_1(c_1(n_t)) + \nu_2(n_t+1) = \nu_2(c_1(n_t)) + \nu_2(c_1(n_t)+1) + \nu_2(n_t+1) = \nu_2(c_1(n_t)) + p_2(n_t)$$

with the last term being equal to the function  $p_2(n)$  from Lemma 4 because

$$(p_1(c_1(n_t)) + \nu_2(n_t + 1)) = (\nu_2(c_1(n_t)) + 0 + \nu_2(n_t + 1)) \mid (n_t \equiv 1) \pmod{2}$$
$$= \left(\nu_2(c_1(n_t)) + \nu_2\left(\frac{n_t}{2^0} + 1\right)\right) \mid (n_t \equiv 1) \pmod{2}$$

and

$$(p_1(c_1(n_t)) + \nu_2(n_t + 1)) = (0 + \nu_2(c_1(n_t) + 1) + 0) \mid (n_t \equiv 0) \pmod{2}$$
$$= \nu_2 \left(\frac{n_t}{2^{\nu_2(n_t)}} + 1\right) \mid (n_t \equiv 0) \pmod{2}.$$

As such, we can similarly simplify  $c_2(n_t)$  to

$$c_2(n_t) := \frac{3^{p_2(n_t)} \left(\frac{n_t}{2^{\nu_2(n_t)}} + 1\right)}{2^{\nu_2(c_1(n_t)) + p_2(n_t)}} - \frac{1}{2^{\nu_2(c_1(n_t))}}.$$

**Lemma 5.** Following previous Corollaries, a rational function for  $c_3(n_t)$  exists, effectively mapping  $n_0 \rightarrow n_3$ .

*Proof.* By induction, suppose that a rational function for  $c_3(n_t)$  can be deduced from the previously derived rational functions  $c_2(n_t)$  and  $c_1(n_t)$  with three unknowns. We consider once again

$$\frac{a(n)\cdot n + g(n)\cdot f(n)}{b(n)} + n.$$

As b(n) is present both in the numerator (as part of a(n)) and the denominator we will deduce b(n) first. Considering the result from previous Corollaries we define  $\delta_3(n)$  as

$$\delta_3(n) = p_1(c_2(n)) + p_1(c_1(n)) + p_1(n)$$

that embeds previous c-step functions and where

$$b(n) = 2^{\delta_3(n)}.$$

Then we can assert

$$D_3(n) = n + 2^{\delta_3(n) + 1}$$

that would reach the same subsequence of  $\delta_3(n)$ , mapping a metric to the next occurrence of a subsequence, i.e. to a proceeding pair equal with respect to  $\delta_3(n)$ . Thus arriving at another pair that contains the same 2-adic valuation outcome as the current pair. E.g.

$$D_3(7) = 71, \ D_3(9) = 73, \ D_3(11) = 43$$
  
$$\delta_3(7) = 5, \ \delta_3(71) = 5, \ \delta_3(9) = 5, \ \delta_3(73) = 5, \ \delta_3(11) = 4, \ \delta_3(43) = 4$$

for natural odd numbers, or

$$D_3(8) = 72, \ D_3(10) = 74, \ D_3(12) = 268$$
  
$$\delta_3(8) = 5, \ \delta_3(72) = 5, \ \delta_3(10) = 5, \ \delta_3(74) = 5, \ \delta_3(12) = 7, \\ \delta_3(268) = 7$$

for natural even numbers. Indeed, for  $\delta_3(n)$  we have the sequence

$$\delta_3(n) = (3, 3, 6, 4, 5, 6, 5, 5, 5, 5, 4, 7, 4, 5, 9, 6, \dots)$$

or

$$2^{\delta_3(n)} = \left(2^3, 2^3, 2^6, 2^4, 2^5, 2^6, 2^5, 2^5, 2^5, 2^5, 2^5, 2^4, 2^7, 2^4, 2^5, 2^9, 2^6, \ldots\right)$$

while for  $D_3(n)$  the sequence

$$D_3(n) = (17, 18, 131, 36, 69, 134, 71, 72, 73, 74, 43, 268, 45, 78, 1039, 144, \dots).$$

For f(n) we claim this to remain

$$f(n) = 2^{\nu_2(n)}$$

equal to previous c-step functions. For the exponent in the numerator, which we shall now denote as r that is contained within a(n) and g(n), we claim the following. If

$$r_1(n_t) = \nu_2 (n_t + 1) | c_1(n_t)$$
  

$$r_2(n_t) = p_2(n_t) = \nu_2 \left(\frac{n_t}{2^{\nu_2(n_t)}} + 1\right) | c_2(n_t)$$

then

$$r_3(n_t) = \nu_2 \left( c_2(n_t) + 1 \right) + \nu_2 \left( \frac{n_t}{2^{\nu_2(n_t)}} + 1 \right) \mid c_3(n_t)$$
$$= \nu_2 \left( c_2(n_t) + 1 \right) + r_2(n_t) \mid c_3(n_t)$$

resulting in the sequence

$$r_3(n_t) = (2, 1, 3, 1, 2, 2, 4, 1, 4, 1, 3, 2, 2, 3, 5, 1, \dots).$$

It's worth highlighting that this is a continuation that recurrently follows from

$$r_2(n_t) = p_2(n_t) = \nu_2 \left( c_1(n_t) + 1 \right) + r_1(n_t) \mid c_2(n_t)$$

because

.

$$\nu_2 \left( c_1(n_t) + 1 \right) + \nu_2(n_t + 1) = 0 + \nu_2(n_t + 1) \mid (n_t \equiv 1) \pmod{2}$$

and

$$\nu_2 \left( c_1(n_t) + 1 \right) + \nu_2(n_t + 1) = \nu_2 \left( \frac{n_t}{2^{\nu_2(n_t)}} + 1 \right) + 0 \mid (n_t \equiv 0) \pmod{2}.$$

Finalising a(n) we thus have

$$a(n_0) = 3^{r_3(n_0)} - 2^{\delta_3(n_0)}.$$

For g(n) we experimentally obtain

$$g(n_0) = \left( \left( 3^{\nu_2(c_2(n_0)+1)} - 2^{\nu_2(c_2(n_0)+1)} \right) \left( 2^{\nu_2(c_1(n_0))} \right) - \left( 3^{\nu_2(c_2(n_0)+1)} - 2^{\nu_2(c_2(n_0)+1)} \right) \right) 2^{\nu_2(n_0+1)} \\ + 3^{r_3(n_0)} - 2^{r_3(n_0)} \\ = \left( \left( 3^{r_3(n_0)-r_2(n_0)} - 2^{r_3(n_0)-r_2(n_0)} \right) \left( 2^{\nu_2(c_1(n_0))} \right) - \left( 3^{r_3(n_0)-r_2(n_0)} - 2^{r_3(n_0)-r_2(n_0)} \right) \right) 2^{\nu_2(n_0+1)} \\ + 3^{r_3(n_0)} - 2^{r_3(n_0)} \\ = \left( 3^{r_3(n_0)-r_2(n_0)} - 2^{r_3(n_0)-r_2(n_0)} \right) \left( 2^{\nu_2(c_1(n_0))} - 1 \right) 2^{\nu_2(n_0+1)} + 3^{r_3(n_0)} - 2^{r_3(n_0)} \right)$$

Bringing it all together in non-simplified or original form we have

$$\frac{\left(3^{r_3(n)}2-2^{\delta_3(n)+1}\right)n}{+\left(\left(3^{r_3(n)-r_2(n)}-2^{r_3(n)-r_2(n)}\right)\left(2^{\nu_2(c_1(n))}-1\right)2^{\nu_2(n+1)}+3^{r_3(n)}-2^{r_3(n)}\right)2^{\nu_2(n)+1}}{2^{\delta_3(n)+1}}$$

$n_0$	$\frac{a}{b}x + \frac{g}{h} = y$	$n_0$	$\frac{a}{b}x + \frac{g}{h} = y$
1	$\frac{2x+7\cdot 2}{16} = y$	14	$\frac{-10x+19\cdot 4}{64} = y$
2	$\frac{-10x+1\cdot4}{16} = y$	15	$\frac{-538x + 451 \cdot 2}{1024} = y$
3	$\frac{-74x+47\cdot 2}{128} = y$	16	$\frac{-122x+1\cdot 32}{128} = y$
4	$\frac{-26x+1\cdot 8}{32} = y$	17	$\frac{2x+7\cdot 2}{16} = y$
5	$\frac{-46x+19\cdot 2}{64} = y$	18	$\frac{-10x+1\cdot 4}{16} = y$
6	$\frac{-110x+5\cdot 4}{128} = y$	19	$\frac{34x+125\cdot 2}{128}$
7	$\frac{98x + 73 \cdot 2}{64} = y$	20	$\frac{-122x+1\cdot 8}{128} = y$
8	$\frac{-58x+1\cdot 16}{64} = y$	21	$\frac{-238x+67\cdot 2}{256} = y$
9	$\frac{98x + 103 \cdot 2}{64} = y$	22	$\frac{-14x+5\cdot4}{32} = y$
10	$\frac{-58x+1\cdot4}{64} = y$	23	$\frac{-350x+185\cdot 2}{512} = y$
11	$\frac{22x+23\cdot 2}{32} = y$	24	$\frac{-494x+5\cdot 16}{512} = y$
12	$\frac{-238x+5\cdot 8}{256} = y$	25	$\frac{22x+29\cdot 2}{32} = y$
13	$\frac{-14x+11\cdot 2}{32} = y$	26	$\frac{-26x+1\cdot 4}{32} = y$

Table 19: Non-simplified or original linear fractions within element pairs for  $1 \le n_0 \le 26 \mid c_3(n_t)$ .



Figure 9: Outcomes of  $c_3(n_0) - n_0$  for  $1 \le n_0 \le 100$ .

of which its results are listed in Table 19 and Figure 9. From here we can reformulate to  $c_3(n_t)$  by subtracting  $2^{\delta_3(n)+1}$  from a(n). Additionally, removing a factor of  $2^{\nu_2(n_t)+1}$  from all terms to obtain

$$c_{3}(n_{t}) := \frac{3^{r_{3}(n_{t})} \frac{n_{t}}{2^{\nu_{2}(n_{t})}}}{+ \left(3^{r_{3}(n_{t}) - r_{2}(n_{t})} - 2^{r_{3}(n_{t}) - r_{2}(n_{t})}\right) \left(2^{\nu_{2}(c_{1}(n_{t}))} - 1\right) 2^{\nu_{2}(n_{t}+1)} + 3^{r_{3}(n_{t})} - 2^{r_{3}(n_{t})}}{2^{\delta_{3}(n_{t}) - \nu_{2}(n_{t})}},$$

which can be reduced to

$$c_{3}(n_{t}) := \frac{3^{r_{3}(n_{t})} \left(\frac{n_{t}}{2^{\nu_{2}(n_{t})}} + 1\right)}{+ \left(3^{r_{3}(n_{t}) - r_{2}(n_{t})} - 2^{r_{3}(n_{t}) - r_{2}(n_{t})}\right) \left(2^{\nu_{2}(c_{1}(n_{t}))} - 1\right) 2^{\nu_{2}(n_{t}+1)} - 2^{r_{3}(n_{t})}}{2^{\delta_{3}(n_{t}) - \nu_{2}(n_{t})}}$$

From Corollary 13 we also have the identity

$$2^{\delta_3(n_t)-\nu_2(n_t)} = 2^{\nu_2(c_2(n_t))+\nu_2(c_1(n_t))+r_3(n_t)}$$

thereby allowing the further reduction of  $c_3(n_t)$  to

$$c_{3}(n_{t}) := \frac{3^{r_{3}(n_{t})} \left(\frac{n_{t}}{2^{\nu_{2}(n_{t})}} + 1\right)}{2^{\nu_{2}(c_{2}(n_{t})) + \nu_{2}(c_{1}(n_{t})) + r_{3}(n_{t})}} - \frac{2^{\nu_{2}(c_{2}(n_{t}) + 1)} - \left(3^{\nu_{2}(c_{2}(n_{t}) + 1)} - 2^{\nu_{2}(c_{2}(n_{t}) + 1)}\right) \left(2^{\nu_{2}(c_{1}(n_{t}))} - 1\right)}{2^{p_{1}(c_{2}(n_{t})) + \nu_{2}(c_{1}(n_{t}))}}$$

yielding for  $n_0 \rightarrow n_3$  the sequence

$$c_3(n_0) = (2, 1, 2, 1, 2, 1, 20, 1, 26, 1, 20, 1, 8, 13, 8, 1, ...)$$

For example for  $n_0 = 7$  we have

$$\underbrace{c_3\left(c_3(7)\right)}_{2 \ times} = 1$$

that maps

$$7 \rightarrow 20 \rightarrow 1$$

as claimed.

**Corollary 14.** As  $g(n_t)$  gets more complex to deduce for every iteration jump, we propose  $a_k(n_t)$  and  $b_k(n_t)$ , where k denotes jumps in iteration steps, to first obtain successive results of  $g(n_t)$  computationally, denoted as  $g_k(n)$  with t = 0.

gir(mit)	1+	21	34	44	~ ~	60	24			10+	11+	124	13+	144	150	164	17)	184	194	204	25+	224	23+	24	254	20	231	284	2%	204	33+	324	334
H.(ma)	1	0	2	0	1	0	19	0	1	0	2		1	0	63	- 0	1	0	3	0	1	0	19	0	1	0	3	0	1	0	211		1
git(mit)	1	1		1	1	5	19	1	1	1	5	5		13	65	1		1	5		1	5	19	5	1	1		13	1	65	211		1
8/144	7	1	17	1	19	5	73	1	163	1	23	5	11	19	450	1	7	1	125	1	62	3	185	3	29	1	2903	19	29	65	692		7
gu(mit)	7	-1	17	7	19	47	73	7	165	13	10		11	73	454	7	T	203	125	19	62	23	185	17	29	11	2903	73	5	101	607	7	7
$H_1(m_0)$	32	7	207	7	121	17	347	7	373	19	133	42	161	73	5119	7	53	202	503	19	432	23	2003	17	541	11	9221	73	211	101	2552	7	143
ga(mit)	22	37	205	37	121	20	347	37	373	121	133	20	161	347	5449	10	13	373	503	121	437	133	2003	205	381	161	9221	347	211	5449	2352	37	143
10(144)	175	37	1319	37	629	209	3049	37	1631	125	1423	207	995	347	32755	32	671	373	2533	121	2310	133	16001	207	2255	161	93229	347	1145	5429	2080	37	2367
ge(mit)	175	125	1315	175	625	1329	2015	175	1631	623	1423	1319	995	300	32735	175	671	1635	2533	625	2316	1423	16001	1315	2253	985	93229	300	1145	32731	2050	175	2367
$B_0(m_{\rm H})$	781	175	6005	175	2881	1329	17159	175	13985	623	8365	1329	5033	300	263729	175	2012	1432	22983	629	11281	1423	80771	1319	10861	995	2828179	300	11627	32731	2140(2)3	175	9729
pic(m)	281	751	6005	793	2001	6005	17159	781	13985	201	8365	4005	2033	17438	263729	261	2012	13085	22983	2881	11281	8365	80771	6005	10861	5833	2625179	17438	11627	163229	21400203	793	9729
RL(m)	3347	791	2(2)7	791	12729	6865	80145	781	72923	201	41.079	6005	23291	17439	253330	251	20375	13085	132465	2881	38227	8365	373385	6005	39119	5833	2010159	17439	67629	163729	66429185	781	45431
gii(m)	3347	3367	26297	3367	12729	2001	8115	3367	72933	12738	41479	26207	23291	85145	223333	3367	26235	72023	132465	12729	38227	41479	373385	2(2)7	39115	23294	22(12(1))	85145	67620	753331	66429185	3367	45431
100.00	14192	3387	111389	3367	52001	21217	396507	3387	347141	12739	189973	21207	102641	80143	130608	3362	93893	72023	67/2009	12729	216217	41479	1644443	2(2)7	356581	23294	720423040	80143	332329	753331	2953341347	3387	209027
gia(mi)	14192	14197	111369	14197	54901	11130	396537	14197	3010	5986	180073	111380	102641	38000	130663	14192	33853	317111	67200	54001	214217	189971	1641413	111365	336381	162641	720823365	38000	332329	2333569	2953541547	14197	200037
100(16)	38975	14197	443239	14197	225339	111399	2683809	14197	1562211	32000	832963	111389	43605	380002	14120041	14290	412734	347141	3072373	54901	903790	189923	7030481	111389	2718009	102641	21439399711	380002	1326345	23336569	18252882563	14197	2215/687
gia(mi)	38075	26075	443239	28962	225325	265239	2683809	26/62	1562211	229333	x32963	265239	43693	1682809	14129041	24912	412751	1263711	3072373	225525	\$\$3755	832363	7033481	46229	2718079	13005	21128399711	1682809	1526345	14120011	182526/0563	58965	2315687
102(16)	242010	28075	2920005	38903	100241	00239	71.0029	38903	6294285	229339	3544245	265239	1841273	1683909	38132218	24912	1762541	1260711	13411423	225339	3780964	832063	29480051	443239	12326541	130005	196223306549	1683909	6676287	14120011	12020251191397	38963	10502-365
gin(mi)	212065	212011	2920005	212261	555441	1320005	71.28329	312441	676265	103344	3541265	1123005	1841273	7146079	38137218	313365	1762541	4794285	13411423	100141	3780964	3541765	29480051	2120005	12348541	1441273	196323396525	7146379	6626287	59137249	12020251191367	313161	10502.365
seal(mail)	3005.27	2420(1	2807347	212261	280966.9	1120005	29834345	212461	28771463	100264	1003099	1123005	2(20)71	710079	21029411	212360	2384225	6794285	22011455	100241	15532187	3041793	121994283	2920005	53822839	1441273	133249996566031	710079	284117169	59137249	1133007116453797	242361	49071311
git(mi)	3805.27	386627	2807147	585527	2606659	THOTIET	29834345	188327	28771463	360653	14828599	TK07247	3120171	29834345	31029611	180327	7384775	26771-013	32011465	3609659	155327287	12(2)(309	121994285	2507147	53822K29	7420071	131269996056031	29834345	28417140	244320611	1133307116453797	\$85027	49071311
(m) (m)	4017132	300027	3296(409	585027	13891481	TRUTIET	123057167	198327	12960821	200003	61263013	707267	31251321	29834345	1005997289	100027	30542333	24771463	234143339	3000009	63398777	12(2)(3)9	200206483	2857167	228077381	2420971	12314428260028391	29834345	118865839	244320611	3495796833129953	\$95027	215822797
gii(mi)	4017132	4017137	32960109	4047137	13891481	31960509	123057267	4017157	12960021	15890481	61263013	22360203	31251321	12302402	1005597288	4647132	30542933	113868621	238143319	13891481	63368777	61263013	5002054×3	32960109	228577381	31254521	123164282820028391	12302402	118803839	2002997269	3495706683149953	4007137	215822797
104(14)	9(24)775	4917137	129431159	2(17137	64453659	21960009	503390129	0(17137	200631194	15896481	200497903	21960209	127308093	12302407	2017/122010	4647132	125183231	113968021	182805113	13891481	217275210	61263(63	2012415361	3296(039	954147399	31254521	37238753771858929	12302407	490635545	2002997289	10627637298793187	2007137	945903847
140.00	91243775	9(24)275	129431759	14213773	64452959	129431259	583390129	16213772	20001191	61132638	200407963	129434738	127308095	503390129	2079732010	202123773	125143231	293922296	182865413	64453559	2012153290	250800903	2012475361	129431759	904147389	12730695	37238153171858829	503390129	490633545	4009733691	10(278)729(79)187	142421773	915903827
ma(ma)	62014343	91243775	322322003	16313773	260463821	129431259	2047041299	16213773	2018343485	64432638	2021129645	129434738	306144713	503390129	14332368368	20213773	309767421	210(22210	2022334063	64453659	1040360665	253807993	X250900731	129431759	393(211)21	12730695	112614161222119111	503390129	2008777547	40C9TX3694	3212022342030873	14240773	3621453385
and real	60011010	60011011	322322000	60014041	2010/3811	322522986	2017011299	60014341	2918343465	200203511	2021129665	322522000	206144713	2007001299	14231268303	60314041	309797,011	2014313165	2022339063	260463821	1010/06060	1021129365	\$259909731	322322000	2006211021	206044713	112914161222119111	2007003299	2009777547	16531168369	221002182630873	60014341	3621431300
mc(m)	26362247	60014041	2101130007	60014341	202062670	322522965	X2X06875425	60014341	8202514168	203203411	4137129329	322322000	208330001	2042041299	66752372296	60314345	2066173173	2014313145	1436256245	210463821	4194322947	1021129145	3330963755	322322000	14133791159	346144713	312154043293127629	2042041299	\$173826289	16534168369	206326762603352621	60014341	13738327.891
and real	26362247	26302.07	2101130107	263653287	202920109	2106430927	X2896875420	26362247	\$202514168	1003020379	4127129329	2106434927	2040300001	X200800243	667X2372296	263652387	2066172175	×202111203	1630393.00	1029820309	6196322967	1132129319	31309061753	2331130107	16233791159	2040.300001	3121540K3293127K29	\$28960245	\$173826289	96750324294	201321712403352101	20303287	15758327.391
mar(ma)	10090292917	26362247	8160791129	20303287	4223223761	2104430027	33430757227	263622447	332974250914	1043626979	14706233233	2166434927	\$233390001	X2X86607345	240064209609	263653467	K320003173	×202314303	66265913639	10296269079	26K7K542137	4137128319	13110(72)723	2101130107	65490923661	2040.305/01	1000031823021811223	X200600245	331111363209	66750324294	42133456806411124281	263633247	64432904307
su(nu)	1009292917	1009292917	8.260791129	1658392967	42232227%1	800010129	33430757227	1658392967	2019123001	122122236	INTERIMENT	800091129	\$233290001	33455757227	20006209009	1000302917	8320003173	33290429900	66263413639	1223222791	26878342137	14796335253	131108729723	8260791129	6549072961	FERTENDET	1000031823921811423	33455757227	31111363429	2006/03/9609	1223456506411124291	1658330967	64432904.307
$\pi_{k}(m_{0})$	4201100013	1000202017	33972337979	1658392967	1406243329	X20076229	134727060029	1653392167	133002140071	42222222316	67298032343	820091229	33800130995	33455757227	13X2017730775	1008302917	33(27)(4111)	332102276966	202334917633	4223222761	67¥1340000	14706335253	543845142641	\$160791129	263132391719	\$233200014	200779120076031620X3	33435737227	11303300245	20006539609	2154271282606912651	1658332967	2120130331417
au(nu)		1211100075		1221020022	16962633279	33872317679	131727060029	1231929323	133902149071	1696/2635379	67298032923	33972317979		130727093003	10420177300771	12011020075		132902309071	20233912053		\$2\$13,210010	67298932923	2439453162641	33972337979	20110391119	13400130995		131727093043	11303200245	2000007783271	2194171291409912401	1011920073	2(20)3(3)3(7)
mu(m)	1700029925	4201100675	131220602000	12101201073	00073772021	33972317979	2424129963629	4314129373	120205-00085	101002022279	2706 21291565	33972317979	120700131353	134727093043	4340714710689	12011020075	10240202014	112002200071	107742403903	1070202027	272540463125	1729032923	2172291211911	1002107629	1070400091301	12400130905	210420714312800000.01	131727093043	12020629297	20420277932271	213200341127920658881	101092022	106113099330

Table 20: Outcomes of  $g_k(n_0)$  for  $1 \le n \le 33$  and for  $1 \le k \le 33$  (zoom in for inspection).

$c_k(n_0)$	$c_k(1_0)$	$c_k(2_0)$	$c_k(3_0)$	$c_k(4_0)$	$c_k(5_0)$	$c_k(6_0)$	$c_k(7_0)$	$c_k(8_0)$	$c_k(9_0)$	$c_k(10_0)$	$c_k(11_0)$	$c_k(12_0)$	$c_k(13_0)$	$c_k(14_0)$	$c_k(15_0)$	$c_k(16_0)$	$c_k(17_0)$	$c_k(18_0)$	$c_k(19_0)$	$c_k(20_0)$	$c_k(21_0)$	$c_k(22_0)$	$c_k(23_0)$	$c_k(24_0)$	$c_k(25_0)$	$c_k(26_0)$	$c_k(27_0)$	$c_k(28_0)$	$c_k(29_0)$	$c_k(30_0)$	$c_k(31_0)$	$c_k(32_0)$	$c_k(33_0)$
$c_1(n_0)$	2	1	8	1	8	3	26	1	14	5	26	3	20	7	80	1	26	9	44	5	32	11	80	3	38	13	62	7	44	15	242	1	50
$c_2(n_0)$	1	2	1	2	1	8	13	2	7	8	13	8	5	26	5	2	13	14	11	8	1	26	5	8	19	20	31	26	11	80	121	2	25
$c_3(n_0)$	2	1	2	1	2	1	20	1	26	1	20	1	8	13	8	1	20	7	26	1	2	13	8	1	- 44	5	242	13	26	5	182	1	38
$c_4(n_0)$	1	2	1	2	1	2	5	2	13	2	5	2	1	20	1	2	5	26	13	2	1	20	1	2	11	8	121	20	13	8	91	2	19
$c_5(n_0)$	2	1	2	1	2	1	8	1	20	1	8	1	2	5	2	1	8	13	20	1	2	5	2	1	26	1	182	5	20	1	206	1	44
$c_6(n_0)$	1	2	1	2	1	2	1	2	5	2	1	2	1	8	1	2	1	20	5	2	1	8	1	2	13	2	91	8	5	2	103	2	11
$c_7(n_0)$	2	1	2	1	2	1	2	1	8	1	2	1	2	1	2	1	2	5	8	1	2	1	2	1	20	1	206	1	8	1	350	1	26
$c_8(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	8	1	2	1	2	1	2	5	2	103	2	1	2	175	2	13
$c_{9}(n_{0})$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	8	1	350	1	2	1	890	1	20
$c_{10}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	175	2	1	2	445	2	5
$c_{11}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	890	1	2	1	668	1	8
$c_{12}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	445	2	1	2	167	2	1
$c_{13}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	668	1	2	1	566	1	2
$c_{14}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	167	2	1	2	283	2	1
$c_{15}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	566	1	2	1	638	1	2
$c_{16}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	283	2	1	2	319	2	1
$c_{17}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	638	1	2	1	3644	1	2
$c_{18}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	319	2	1	2	911	2	1
$c_{19}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	3644	1	2	1	4616	1	2
$c_{20}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	911	2	1	2	577	2	1
$c_{21}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	4616	1	2	1	866	1	2
$c_{22}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	577	2	1	2	433	2	1
$c_{23}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	866	1	2	1	650	1	2
$c_{24}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	433	2	1	2	325	2	1
$c_{25}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	650	1	2	1	488	1	2
$c_{26}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	325	2	1	2	61	2	1
$c_{27}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	488	1	2	1	92	1	2
$c_{28}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	61	2	1	2	23	2	1
$c_{29}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	92	1	2	1	80	1	2
$c_{30}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	23	2	1	2	5	2	1
$c_{31}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	80	1	2	1	8	1	2
$c_{32}(n_0)$	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	5	2	1	2	1	2	1
$c_{33}(n_0)$	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	2	1	8	1	2	1	2	1	2

Table 21: Outcomes of  $c_k(n_0)$  for  $1 \le n \le 33$  and for  $1 \le k \le 33$  (zoom in for inspection).

*Proof.* Deductively speaking, if k defines jump iterations over executions of  $c(n_t)$  as  $c_k(n_t)$  then

$$g_k(n_t) = \left(c_k(n_t) - a_k(n_t)\frac{n_t}{2^{\nu_2(n_t)}}\right)b_k(n_t) \mid t = 0.$$

The computational verification of  $g_k(n_0)$  is given in Appendix 4 with some of its results listed in Table 20 along with Table 21 that lists corresponding integers of  $c_k(n_0)$ .

**Remark 10.** Noteworthy from Table 20 are the appearance of integers that are at most  $\nu_2(n) = 1$  between differences of odd elements of k and k + 2 for every  $n_t$ .

**Remark 11.** Also, in this particular instance we subtract  $2^{\delta_k(n_t)}$  from  $a_k(n_t)$  and have  $f(n_t) = 1$  by removing a factor of two from both the numerator and the denominator.

**Definition 6.** We introduce the notion of a *zipper function* for  $c(n_t)$  denoted with C. Let

$$c(n_t) := \frac{a(n_t) \cdot \frac{n_t}{2^{\nu_2(n_t)}} + g(n_t)}{b(n_t)}$$

and let  $k \in \mathbb{N}$  denote the iteration index for every number of k-iteration steps. Then, we can substitute as

$$C_k(n) = \frac{a_k(n) \cdot \frac{n}{2^{\nu_2(n)}} + g_k(n)}{b_k(n)}$$

that contains the expansion

$$C_{\langle 1,...,k\rangle}(n) = \frac{3^{r_{\langle 1,...,k\rangle(n)}} \cdot \frac{n}{2^{\nu_2(n)}} + g_{\langle 1,...,k\rangle}(n)}{2^{\delta_{\langle 1,...,k\rangle}(n)}}$$

up until k = 1 that forms the basis with  $c_0(n) = n$ . The zipper function thus symbolically encodes continuous inclusions of previous iterations from k...1 where outcomes of n trend to converge around the trivial parity cycle (2;1) giving the visual impression of closing a zipper for increments k. As a result we highlight the example of  $C_7(n)$  that gets defined as

$$C_7(n) = \frac{3^{\nu_2(c_6(n)+1)+\nu_2(c_5(n)+1)+\nu_2(c_4(n)+1)+\dots+\nu_2(c_0(n)+1)} \cdot \frac{n}{2^{\nu_2(n)}} + g_{\langle 1,\dots,k \rangle}(n)}{2^{p_1(c_6(n))+p_1(c_5(n))+p_1(c_4(n))+\dots+p_1(c_0(n))}}$$

which through the results of computational verification as listed in Table 21 yields the sequence

$$C_7(n) = (2, 1, 2, 1, 2, 1, 2, 1, 8, 1, 2, 1, 2, 1, 2, 1, ...).$$

**Remark 12.**  $g_k(n)$  encapsulates all terms as previously stated and is by Corollary 14 the only term that is (at this moment) reliant on the outcome of  $c_k(n)$ .

**Theorem 4.** Following Definition 6 we can induce and validate identifications of subsequent results of  $g_k(n) | k > 3$  computationally.

*Proof.* Let

$$g_{1}(n) = \left(3^{\nu_{2}(I_{0})} - 2^{\nu_{2}(I_{0})}\right)$$
  

$$g_{2}(n) = \left(3^{\nu_{2}(I_{1}I_{0})} - 2^{\nu_{2}(I_{1}I_{0})}\right)$$
  

$$g_{3}(n) = \left(3^{\nu_{2}(I_{2}I_{1}I_{0})} - 2^{\nu_{2}(I_{2}I_{1}I_{0})}\right)$$
  

$$+ 2^{\nu_{2}(I_{0})} \left(3^{\nu_{2}(I_{2})} - 2^{\nu_{2}(I_{2})}\right) \left(2^{\nu_{2}(V_{1})} - 1\right).$$

The *p*-adic property for multiplication as stated in Corollary 12 is used here as a way to shorten the exponents of each identified result from a series to a product. We also further denote  $I_{k-1,\dots,1}$  as the variety for iterates  $(c_{k-1,\dots,1}(n_t) + 1)$  and  $V_{k-1,\dots,1}$  as the variety for

iterates  $(c_{k-1,\dots,1}(n_t)+0)$ . For  $4 \leq k \leq 6$  of  $g_k(n)$  we have experimentally identified the following results

$$\begin{split} g_4(n) &= \left(3^{\nu_2(I_3I_2I_1I_0)} - 2^{\nu_2(I_3I_2I_1I_0)}\right) \\ &+ 2^{\nu_2(I_1I_0)} \left(3^{\nu_2(I_3I_2)} - 2^{\nu_2(I_3I_2)}\right) \left(2^{\nu_2(V_2V_1)} - 1\right) \\ g_5(n) &= \left(3^{\nu_2(I_4I_3I_2I_1I_0)} - 2^{\nu_2(I_4I_3I_2I_1I_0)}\right) \\ &+ 2^{\nu_2(I_2I_1I_0)} \left(3^{\nu_2(I_4)} - 2^{\nu_2(I_4)}\right) \left(2^{\nu_2(V_3V_2V_1)} - 1\right) \\ &+ 3^{\nu_2(I_4)}2^{\nu_2(I_1I_0)} \left(3^{\nu_2(I_3I_2)} - 2^{\nu_2(I_3I_2)}\right) \left(2^{\nu_2(V_2V_1)} - 1\right) \\ g_6(n) &= \left(3^{\nu_2(I_5I_4I_3I_2I_1I_0)} - 2^{\nu_2(I_5I_4I_3I_2I_1I_0)}\right) \\ &+ 2^{\nu_2(I_3I_2I_1I_0)} \left(3^{\nu_2(I_5I_4)} - 2^{\nu_2(I_5I_4)}\right) \left(2^{\nu_2(V_4V_3V_2V_1)} - 1\right) \\ &+ 3^{\nu_2(I_5I_4)}2^{\nu_2(I_1I_0)} \left(3^{\nu_2(I_3I_2)} - 2^{\nu_2(I_3I_2)}\right) \left(2^{\nu_2(V_2V_1)} - 1\right) . \end{split}$$

Then, by induction we claim

$$g_{7}(n) = \left(3^{\nu_{2}(I_{6}I_{5}I_{4}I_{3}I_{2}I_{1}I_{0})} - 2^{\nu_{2}(I_{6}I_{5}I_{4}I_{3}I_{2}I_{1}I_{0})}\right)$$
  
+  $2^{\nu_{2}(I_{4}I_{3}I_{2}I_{1}I_{0})} \left(3^{\nu_{2}(I_{6})} - 2^{\nu_{2}(I_{6})}\right) \left(2^{\nu_{2}(V_{5}V_{4}V_{3}V_{2}V_{1})} - 1\right)$   
+  $3^{\nu_{2}(I_{6}I_{5}I_{4})}2^{\nu_{2}(I_{1}I_{0})} \left(3^{\nu_{2}(I_{3}I_{2})} - 2^{\nu_{2}(I_{3}I_{2})}\right) \left(2^{\nu_{2}(V_{2}V_{1})} - 1\right)$   
+  $3^{\nu_{2}(I_{6})}2^{\nu_{2}(I_{4}II_{3}I_{2}I_{1})} \left(3^{\nu_{2}(I_{5}I_{4})} - 2^{\nu_{2}(I_{5}I_{4})}\right) \left(2^{\nu_{2}(V_{4}V_{3}V_{2}V_{1})} - 1\right)$ 

resulting in the sequence

 $g_7(n) = (175, 37, 1319, 37, 619, 269, 3089, 37, 1631, 121, 1423, 269, 995, 347, 32731, 37, \dots)$ 

**Conjecture 4.** We conjecture the following three constraints  $\forall n \in \mathbb{N}^*$  from Theorem 4. First, an upper bound for the total number or series of expressions for each k of g, given as the function

$$B(k) = \left\lfloor \frac{k+1}{2} \right\rfloor \mid k > 0$$

that calculates a growth of one extra expression for every other  $g_k$  by embedding the last expression of  $g_{k-2}$  multiplied with additive factors of  $3^{\nu_2(I_{k-1})}$ . The second is a sequence of  $I_{k-1}$ 's in the second expression that coincides with subsequences of OEIS A004523 [10].

For example, ((2,3), (4), (4,5), (6), ...). The third constraint is a combination of at most four types of expressions, which are  $(3^{\nu_2(\langle I \rangle)} - 2^{\nu_2(\langle I \rangle)})$ ,  $2^{\nu_2(\langle I \rangle)}$ ,  $(2^{\nu_2(\langle V \rangle)} - 1)$ , and lastly  $3^{\nu_2(\langle I \rangle)}$ . All three constraints are demonstrated by Appendix 5.

## Appendix

#### 1 Binary opposite modulo

For definition in Appendix 2 we employ a binary opposite modulo, defined as

$$\mathbf{bomod}(x,y) = \begin{cases} 1 & \text{if } x - y \lfloor \frac{x}{y} \rfloor = 0 \\ 0 & \text{if } x - y \lfloor \frac{x}{y} \rfloor > 0 \end{cases}$$

using floored division. For example if y = 2 then **bomod**(2, 2) = 1, conversely **bomod**(1, 2) = 0 for  $y \in \mathbb{N}^*$ . This provides versatility in testing membership of integers.

**Remark 13.** This concept was developed to assist in computing p-adic valuations on some graphing calculators such as Desmos to accelerate the research brought forth herein.

### 2 Primitive recursive $\mathbb{N}^*$ -adic valuation function

Given Appendix 1 we can define a primitive recursive function of the N<sup>\*</sup>-adic valuation function  $\nu_n(x)$  described in Corollary 2 other than the common algorithmic approach that uses a **while-loop** as

$$\nu_y(x) = \sum_{i=1}^{\lfloor \log_y(x) \rfloor} \mathbf{bomod}(x, y^i)$$

for  $x > 0 \land y > 0$  with the accompanying code for the Python 3 language

```
from mpmath import mp
mp.dps = 300 # adjust precision as needed
# Binary opposite modulo function, returns 1 if 0 or 0 if 1
def bomod(x, y):
    if mp.fsub(x, mp.fmul(y, mp.floor(mp.fdiv(x, y)))) > 0:
        return 0
    else:
        return 1
# General n-adic valuation function in primitive form (a = n-adic number, e.g. 2 prime)
def v(n, a):
    k = 0
    # in Python, end of range function is exclusive, hence + 1
    for i in range(1, int(mp.floor(mp.fadd(mp.log(n, a), 1)))):
        k = k + bomod(n, mp.power(a, i))
    return k
```

# 3 One approach to derive t following Corollary 7

$$\begin{split} n_0 &= 1, t = 0 + 2(1) + 1 = 3 \\ n_0 &= 2, t = 1 \\ n_0 &= 3, t = 0 + 2(2) + 3 = 7 \\ n_0 &= 4, t = 2 \\ n_0 &= 5, t = 0 + 2(1) + 3 = 5 \\ n_0 &= 6, t = 1 + 2(2) + 3 = 8 \\ n_0 &= 7, t = 0 + 2(3) + 1 + 2(1) + 2 + 2(1) + 3 = 16 \\ n_0 &= 8, t = 3 \\ n_0 &= 9, t = 0 + 2(1) + 1 + 2(3) + 1 + 2(1) + 2 + 2(1) + 3 = 19 \\ n_0 &= 10, t = 1 + 2(1) + 3 = 6 \\ n_0 &= 11, t = 0 + 2(2) + 1 + 2(1) + 2 + 2(1) + 3 = 14 \\ n_0 &= 12, t = 2 + 2(2) + 3 = 9 \\ n_0 &= 13, t = 0 + 2(1) + 2 + 2(1) + 3 = 9 \\ n_0 &= 14, t = 1 + 2(3) + 1 + 2(1) + 2 + 2(1) + 3 = 17 \\ n_0 &= 15, t = 0 + 2(4) + 1 + 2(1) + 6 = 17 \\ n_0 &= 16, t = 4 \end{split}$$

## 4 Computational verification of $g_k(n_t)$

```
from mpmath import mp
mp.dps = 300 #adjust precision as needed
import functools
# general algorithmic way to write a n—adic valuation function
@functools.lru_cache(maxsize=None)
def v(n, a):
   i = 0
    while mp.fmod(n, a) == 0:
       n = mp.fdiv(n, a)
       i = i + 1
    return i
# p_1(n) as defined in the paper
@functools.lru_cache(maxsize=None)
def p(n):
    return v(mp.fadd(n, mp.fmod(n, 2)), 2)
# parity alternating Collatz function
def c(n):
   nm = mp.fadd(n, mp.fmod(n, 2))
    d = mp.fdiv(nm, mp.power(2, v(nm, 2)))
    three = mp.power(3, v(mp.fadd(n, 1), 2))
    return int(mp.fsub(mp.fmul(d, three), mp.fmod(n, 2)))
# demonstration of the zipper function (to compute g(n))
def gFromZipper(nMax, k):
   nt = []
   b = []
   r = []
   a = []
   g = []
    # integers of c(n)
    cint = []
    for n in range(1, nMax+1):
        # temp store for results of c(n)
       nk = []
       nk.append(n)
        delta = v(mp.fadd(n, 1), 2)
```

```
rk = delta
        # temp store to recall last result of c(n)
        ck = n
        for ki in range(1, k):
            # iterate c(n), a.k.a. c_1, c_2, etc.
            ck = c(nk[ki-1])
            nk.append(ck)
            delta = mp.fadd(delta, p(ck))
            rk = mp.fadd(v(mp.fadd(ck, 1), 2), rk)
        bn = mp.power(2, delta)
        b.append(int(bn))
        r.append(int(rk))
        ak = mp.power(3, rk)
        a.append(int(ak))
        xdiv = mp.fdiv(n, mp.power(2, v(n, 2)))
        akx = mp.fmul(ak, xdiv)
        akxb = mp.fdiv(akx, bn)
        gkb = mp.fsub(c(ck), akxb)
        cint.append(c(ck))
        gk = mp.fmul(gkb, bn)
        g.append(int(gk))
    #return(a)
    #return(b)
    return(g)
    #return(cint)
# test range
for n in range(1, 33 + 1):
    derived = gFromZipper(33, n)
```

print(derived)

## **5** Computational validation of $g_k(n_t)$

```
# see Appendix 3 for c, p, v and gFromZipper functions
def ckMul(ck, s, e):
    exponent = 1
    for k in range(s, e):
        exponent = mp.fmul(exponent, ck[k])
    return exponent
def express1(ck0dd, s, e):
    exponent = ckMul(ckOdd, s, e)
    return mp.fsub(mp.power(3, v(exponent, 2)), mp.power(2, v(exponent, 2)))
def express2(ck0dd, s, e):
    exponent = ckMul(ckOdd, s, e)
    return mp.power(2, v(exponent, 2))
def express3(ckEven, s, e):
    exponent = ckMul(ckEven, s, e)
    return mp.fsub(mp.power(2, v(exponent, 2)), 1)
def express4(ck0dd, s, e):
    exponent = ckMul(ckOdd, s, e)
    return mp.power(3, v(exponent, 2))
def t1t2t3(ck0dd, ckEven, e):
    t1 = express2(ck0dd, 0, e - 2)
    if (mp.fmod(e, 2) == 0):
        t2 = express1(ck0dd, e - 2, e)
    else:
        t2 = express1(ck0dd, e - 1, e)
    # start at k = 1 as we can remove v_2(c_0) from all
    t3 = express3(ckEven, 1, e - 1)
    return mp.fprod([t1, t2, t3])
def gkV(nMax, nMin, kMax):
    # storage for sum of k—expressions for every n
    nExpressionSums = []
    # go through n—columns (we start at n = 1)
    for n in range(nMin, nMax + 1):
        ck0dd = []
```

```
ckEven = []
        ckOdd.append(int(mp.fadd(n, 1))) # precompute all c_k iterates
        ckEven.append(n)
        for ck in range(0, kMax - 1):
            ckEven.append(int(c(ckEven[ck])))
            ckOdd.append(int(mp.fadd(ckEven[ck + 1], 1)))
        kExpresss = [] # storage of expressions per k
        # max number of expressions for every k
        maxExpressions = int(mp.floor(mp.fdiv(mp.fadd(kMax, 1), 2)))
        # keep track of self—expressions
        expressRemainder = maxExpressions
        for express in range(maxExpressions):
            if (express == 0):
                kExpresss.append(express1(ckOdd, 0, kMax))
                expressRemainder = expressRemainder — 1
            elif (express == 1):
                kExpresss.append(t1t2t3(ckOdd, ckEven, kMax))
                expressRemainder = expressRemainder — 1
            else: # express >= 2
                # for each additional expression we look back
                for iter in range(expressRemainder, expressRemainder + 1):
                    if (mp.fmod(kMax, 2) == 0):
                        lookB4 = expressRemainder * 2
                        t4 = express4(ck0dd, kMax - lookB4, kMax)
                        t4 = mp.fmul(t4, t1t2t3(ckOdd, ckEven, kMax - lookB4))
                        expressRemainder = expressRemainder - 1
                    else:
                        lookB4 = expressRemainder * 2 - 1
                        t4 = express4(ckOdd, kMax - lookB4, kMax)
                        t4 = mp.fmul(t4, t1t2t3(ckOdd, ckEven, kMax - lookB4))
                        expressRemainder = expressRemainder - 1
                    kExpresss.append(t4)
        nExpressionSums.append(int(mp.fsum(kExpresss)))
   return(nExpressionSums)
for k in range(1, 33 + 1):
   obtained = gFromZipper(33, k)
   validated = gkV(33, 1, k)
   difference = list(map(int.__sub__, obtained, validated))
   print(difference) # expectation is a matrix with only zeros
```

## 6 Zipper function $C_k(n)$ demonstration

```
mp.dps = 10000 # adjust precision as needed
# Zipper function (construction)
# n: number in N^*
# k: k—step or number of rows
def cZipper(n, k):
    nk = []
    # here we only iterate once
    if k == 1:
        return c(n)
    # temp store for results of c(n)
    nk.append(n)
    delta = v(mp.fadd(n, 1), 2)
    rk = delta
    # temp store to recall last result of c(n)
    ck = n
    gk = n
    for ki in range(1, k):
        # iterate c(n), a.k.a. c_1, c_2, etc.
        ck = c(nk[ki - 1])
        nk.append(ck)
        delta = mp.fadd(delta, p(ck))
        rk = mp.fadd(v(mp.fadd(ck, 1), 2), rk)
        gk = gkV(n, n, k)[0]
    bn = mp.power(2, delta)
    ak = mp.power(3, rk)
    xdiv = mp.fdiv(n, mp.power(2, v(n, 2)))
    akx = mp.fmul(ak, xdiv)
    akxg = mp.fadd(akx, gk)
    ck = mp.fdiv(akxg, bn)
    return(int(ck))
```

print(cZipper(999999999999999999999999, 314))

## Postface

I would first of all like to extend my apologies as I, the author, does not have a formal education/training in pure mathematics but has instead a degree in computer science. Hence, the information presented in this paper could be described differently through more formal/advanced mathematical concepts. Unfortunately, time is a luxury in these uncertain times and the exploration would not have been possible if the author were to undertake a formal education.

This information should be best considered as meta extracts from an external observer that took a submarine into the number theory ocean, to find a rock that is lurking at the depth of the sea floor, and start chipping it away to see what lies inside. Communicating these ideas are meant to catch the attention of the experts and professionals to help validate them and determine whether it regresses or progresses our understanding of this truth seeking problem.

The author wishes to thank especially his amazing wife, brilliant kids, extended family, and friends for their consistent support throughout this pattern seeking quest, followed by Quanta Magazine with Prof. Terrence Tao for the introduction to the  $\sim 84$  years old conjecture of Lothar Collatz. Last but not least, the vibrant communities that maintain the Encyclopaedia of Integer Sequences, Wikipedia, Math StackExchange, Python, and the folks at Desmos for making this useful graphing calculator free.

#### References

- Amdeberhan, T.; Manna, D.; Moll, V.H. The 2-adic valuations of Stirling numbers, 2007.
- [2] Pólya, G.; Tarjan, R.E.; Woods, D.R. Notes on Introductory Combinatorics, Springer, pp. 21–22, ISBN 9780817649524 2009.
- [3] Sloane, N.J.A. An On-LineVersion of the Encyclopedia of Integer Sequences, https://oeis.org/A007814 2021.
- [4] Sloane, N.J.A. An On-LineVersion of the Encyclopedia of Integer Sequences, https://oeis.org/A050603 2021.
- [5] Sloane, N.J.A. An On-LineVersion of the Encyclopedia of Integer Sequences, https://oeis.org/A136480 2021.
- [6] Sloane, N.J.A. An On-LineVersion of the Encyclopedia of Integer Sequences, https://oeis.org/A163575 2021.
- [7] Sloane, N.J.A. An On-Line Version of the Encyclopedia of Integer Sequences, https://oeis.org/A089309 2021.
- [8] Sloane, N.J.A. An On-Line Version of the Encyclopedia of Integer Sequences, https://oeis.org/A160541 2021.
- [9] Sloane, N.J.A. An On-Line Version of the Encyclopedia of Integer Sequences, https://oeis.org/A000265 2021.
- [10] Sloane, N.J.A. An On-Line Version of the Encyclopedia of Integer Sequences, https://oeis.org/A004523 2021.