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Projective toric codes

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Any integral convex polytope P in \mathbb{R}^N provides a N -dimensional toric variety \mathbf{X}_P and an ample divisor D_P on this variety. This paper gives an explicit construction of the algebraic geometric error-correcting code on \mathbf{X}_P , obtained by evaluating global section of the line bundle corresponding to D_P on every rational point of \mathbf{X}_P . This work presents an extension of toric codes analogous to the one of Reed-Muller codes into projective ones, by evaluating on the whole variety instead of considering only points with non-zero coordinates. The dimension of the code is given in terms of the number of integral points in the polytope P and an algorithmic technique to get a lower bound on the minimum distance is described.

Keywords: Error-correcting codes, Toric variety, Polytope, Algebraic geometric codes, Gröbner bases.

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Introduced by J.P. Hansen [14], toric codes consist in evaluating Laurent monomials $x_1^{m_1} x_2^{m_2} \dots x_n^{m_n}$ at points $(x_1, \dots, x_n) \in (\mathbb{F}_q^*)^n$ where $m = (m_1, \dots, m_n)$ runs through the integral points of a given lattice polytope $P \subset \mathbb{R}^n$. This simple definition hides the algebraic geometric nature of these codes. A lattice polytope P defines a toric variety \mathbf{X}_P , which contains a dense torus $\mathbb{T}^n \simeq (\overline{\mathbb{F}_q^*})^n$. The toric code associated to P , denoted here by C_P , is the code obtained by evaluating of the functions in the Riemann-Roch space of a divisor D_P on \mathbf{X}_P at the rational points of \mathbb{T}^n .

Theoretically, algebraic geometric codes have excellent parameters but their practical use requires the computation of Riemann-Roch spaces. Arithmetic [16, 30] and geometric [9, 10, 13, 19] algorithms have been developed and refined to compute the Riemann-Roch space of a divisor D on a curve C . However, the literature is undoubtedly sparser on higher dimensional varieties. Toric varieties present the extremely convenient property of having an explicit description of the set of their divisors and their Riemann-Spaces in combinatorial terms, in any dimension. This explains why the parameters of toric codes have been broadly investigated, on sur-

faces [14, 20, 26] and in higher dimension [21, 25].

In this paper, we construct a linear code from the polytope P by evaluating the elements of the Riemann Roch space $L(D_P)$ at the \mathbb{F}_q -points of the whole variety \mathbf{X}_P , not only of the dense torus, as in toric codes. In the present work, we call the code defined from the polytope P the *projective toric code* associated to P and we denote it by PC_P . This terminology has been chosen because the difference between a toric code C_P and its projective version PC_P is comparable to the one between a Reed-Muller code and the projective Reed-Muller code of the same degree [18]. In the latter case, we evaluate *homogeneous* polynomials in $n + 1$ variables instead of polynomials in n variables. Actually, toric varieties, like projective spaces, come with a homogenization process which, given a field k , turns Laurent polynomials of $k[t_1^\pm, \dots, t_n^\pm]$ into polynomials of the Cox ring $k[X_1, \dots, X_r]$ ($r > n$). The parameters of PC_P compared to those of C_P behave the same way as for (projective) Reed-Muller codes : for q large enough, the dimension stays the same but the length and the minimum distance are increased.

This work presents a general framework of the study of algebraic codes on toric varieties, in which notably (weighted) projective Reed-Muller codes fit (see [18, 27] for parameters). As several toric codes are champion codes [2, 3, 22], extending them by evaluating outside the big torus is likely to provide new champion codes (see Sec. 6.2). Besides, this extension of toric codes will probably find applications in information theory. For instance, J.P. Hansen built a *secret sharing scheme with strong multiplication* based on toric codes [15] in which, as usual, the number of participants is bounded by the length of the code. Projective toric codes may provide similar secret sharing schemes with more participants and analogous techniques are likely to give the parameters of these schemes.

When studying classical toric codes, the rational points and the global sections have a simple and explicit description. Here arises the problem of handling the \mathbb{F}_q -points of the abstract variety \mathbf{X}_P . Instead of embedding the toric variety into \mathbf{X}_P into a possibly large projective space [4], we focus on the case where the variety \mathbf{X}_P is *simplicial* and thus can be represented as a *good geometric quotient* of an open affine subset of \mathbb{A}^r under the action of a group G , where r is the number of facets of the polytope P . This standpoint proved to be effective for the study of Goppa codes on Hirzebruch surfaces [24]. This not only provides a good grasp on the points of \mathbf{X}_P but it also expresses the global sections in $L(D_P)$ as polynomials of $\mathbb{F}_q[X_1, \dots, X_r]$. To evaluate them at the rational points \mathbf{X}_P , we determine which orbits of r -tuples correspond to \mathbb{F}_q -points, choose a representative among each of these orbits and finally evaluate naively polynomials at these representatives. Note that a generator matrix can then be constructed thanks to Hermite Normal forms without any knowledge of toric varieties (see Section 2.4).

The combinatorial properties of P echo the geometric properties of the variety \mathbf{X}_P and are related to the parameters of the toric codes C_P and PC_P . Integral points of the polytope P correspond to monomials forming a basis of the space of

global sections $L(D_P)$. Regarding the dimension, there is a correspondance between a basis of the toric code C_P and the lattice points of P modulo $q-1$ [25]. Using the decomposition of \mathbf{X}_P as a disjoint union of tori, one for each face of P , we prove here that reducing the lattice points of P modulo $q-1$ *face by face* provides a basis of PC_P (Th. 3.5).

To bound the minimum distance of PC_P from below, we aim to bound from above the number of zeroes in \mathbf{X}_P of a global section $f \in L(D_P)$, as usual in the context of algebraic-geometric code. Here, we use a footprint bound technique [1, 8] generalizing the one set for codes on Hirzebruch surfaces [24]. We use Gröbner basis theory to express a bound on the minimum distance in terms of integral points in a polytope (Th. 5.2), without suffering from the exponential growth in the number of variables of complexity of the actual computation of a Gröbner basis.

Altogether, constructing the projective toric code PC_P and estimating its parameters heavily relies on the determination of lattice points of some polytopes. For a polytope of dimension N of volume V with s vertices, this can be done in time $\tilde{O}\left(\left(s^{\lceil \frac{N}{2} \rceil} + V\right) \log \delta\right)$ where δ is the maximum modulus of the coordinates of the vertices of P [28, Prop. 3.5]. This encourages a thorough study of (projective) toric codes associated to polygons before investigating in higher dimension.

1. Toric variety from a polytope

This section sums up different definitions and properties of toric varieties defined by polytopes, from the reference book of D. Cox, J. B. Little, and H. K. Schenck [6].

1.1. Normal fan of P

Let $P \subset \mathbb{R}^N$ be a full dimensional convex lattice polytope. Throughout this paper, all polytopes are assumed *convex*. We write $Q \preceq P$ if Q is a face of P . For each $k \in \{0, \dots, N\}$, the set of k -dimensional faces of P is denoted by $P(k)$. The faces of dimension $N-1$ are called *facets* and those of dimension 0 *vertices*.

Definition 1.1 (Primitive normal vector). For each facet F of P , denote by u_F the shortest inner normal vector of F with integer coordinates. We call u_F the *primitive normal vector* of the facet F .

For every face $Q \preceq P$, we define the cone $\sigma_Q = \text{Cone}(u_F \mid Q \preceq F) \subset \mathbb{R}^N$. The union of these cones forms the *normal fan* of P defined as $\Sigma := \bigcup_{Q \preceq P} \sigma_Q$. [6, Th. 2.3.2]

Such a fan defines an abstract toric variety \mathbf{X}_P , by gluing some affine spaces corresponding to these cones [6, Prop. 3.1.6].

Fact 1.2. The toric varieties associated to two polytopes P and P' are isomorphic if and only if their normal fans are equal up to a unimodular transformation, *i.e.* multiplication by a matrix of $\text{GL}_N(\mathbb{Z})$.

1.2. Simplicial toric variety as quotient

Let \mathbb{F}_q be a finite field with q elements, where q is a prime power. **From now on, we assume that both of the following hypotheses hold.**

(H1) The polytope P is simple: each vertex belongs to exactly N facets.

For each vertex $v \in P(0)$, let us define the N -square matrix $A(v)$ whose rows are the primitive normal vectors u_F (Def. 1.1) of the N facets F containing v .

(H2) The determinant of $A(v)$ is coprime with the characteristic of the field \mathbb{F}_q .

Remark 1.3.

- (1) (H1) is only meaningful when $N \geq 3$. Every convex polygon is simple. In dimension 3, cubes and tetrahedra are simple but square-based pyramids are not, as their top vertex belongs to 4 faces.
- (2) Note that (H2) does not depend on the order of the rows of $A(v)$. Moreover, it only matters when the variety given by the polytope P is not smooth, otherwise all these determinants are equal to ± 1 .

These hypotheses ensure that the toric variety \mathbf{X}_P is *simplicial* and enable us to consider it as the quotient of an open affine subset modulo a group action. More precisely, we see \mathbf{X}_P in the affine space \mathbb{A}^r , where r is the number of facets of P . In the polynomial ring,

$$R := \mathbb{F}_q[X_F \mid F \in P(N-1)], \quad (1.1)$$

we consider the *irrelevant ideal* $B := \langle \prod_{F \ni v} X_F \mid v \in P(0) \rangle$. We denote by $\mathcal{Z} \subset \mathbb{A}^r$ its set of zeroes.

Theorem 1.4. [6, Prop. 5.1.9] *The variety \mathbf{X}_P is isomorphic to the quotient of $\mathbb{A}^r \setminus \mathcal{Z}$ under the action of the algebraic group*

$$G = \left\{ (t_F)_{F \in P(N-1)} \in (\overline{\mathbb{F}_q}^*)^r : \forall j \in \{1, \dots, N\}, \prod_{F \in P(N-1)} t_F^{u_F^{(j)}} = 1 \right\} \quad (1.2)$$

where $u_F^{(j)}$ is the j -th coordinate of the vector $u_F \in \mathbb{Z}^N$, with the following action:

$$\begin{cases} (\mathbb{A}^r \setminus \mathcal{Z}) \times G \rightarrow \mathbb{A}^r \setminus \mathcal{Z} \\ (\mathbf{x}, \mathbf{t}) \mapsto \mathbf{t} \cdot \mathbf{x} := (t_F x_F)_{F \in P(N-1)}. \end{cases}$$

Remark 1.5. Proposition 5.1.9 [6] only holds in characteristic zero. It follows from the reductivity of the group G [12], which is guaranteed here by (H2).

In this presentation of \mathbf{X}_P , rational points correspond to Frobenius-invariant G -orbits, *i.e.* elements $\mathbf{x} \in \mathbb{A}^r \setminus \mathcal{Z}$ such that there exists $\mathbf{t} \in G$ satisfying $\mathbf{x}^q = \mathbf{t} \cdot \mathbf{x}$.

1.3. Cox ring and polynomials

Beside an elegant description of the abstract variety \mathbf{X}_P , the quotient representation provides a polynomial coordinate ring to the variety \mathbf{X}_P as the ring R (Eq. (1.1)) endowed with a grading by the class group of \mathbf{X}_P .

1.3.1. Class group of \mathbf{X}_P

The *class group* of the variety \mathbf{X}_P , denoted by $\text{Cl}(\mathbf{X}_P)$, is the set of its Weil divisors modulo linear equivalence. We denote by $[D] \in \text{Cl}(\mathbf{X}_P)$ the class of a Weil divisor D on \mathbf{X}_P . The r facets $F \in P(N-1)$ give r torus-invariant prime divisors D_F on \mathbf{X}_P . Their classes $[D_F]$ generate $\text{Cl}(\mathbf{X}_P)$ as \mathbb{Z} -module [6, Th. 4.1.3]. More precisely,

$$\text{Cl}(\mathbf{X}_P) \simeq \mathbb{Z}^{r-N} \oplus \bigoplus_{i=1}^N \mathbb{Z}/d_i\mathbb{Z}$$

where $1 \leq d_1 \mid d_2 \mid \cdots \mid d_N$ are the *invariant factors* of the $r \times N$ -matrix A_P whose rows are the primitive normal vectors u_F . If $d_i = 0$, then $\mathbb{Z}/d_i\mathbb{Z} = \{0\}$.

1.3.2. Cox ring

In the polynomial ring $R = \mathbb{F}_q[x_F \mid F \in P(N-1)]$, we define the *degree* of a monomial $M = \prod_{F \in P(N-1)} X_F^{\alpha_F}$ as the class $[\sum \alpha_F D_F] \in \text{Cl}(\mathbf{X}_P)$. Denote by R_α the \mathbb{F}_q -vector space spanned by monomials of degree $\alpha \in \text{Cl}(\mathbf{X}_P)$. Then

$$R = \bigoplus_{\alpha \in \text{Cl}(\mathbf{X}_P)} R_\alpha.$$

Fact 1.6. [6, Th 5.4.1] The component R_α is isomorphic to the Riemann-Roch space $L(D) = H^0(\mathbf{X}_P, \mathcal{O}(D))$ of any divisor such that $[D] = \alpha$.

Riemann-Roch spaces come with a handy depiction thanks to a correspondence between effective divisors and polytopes. For a divisor $D = \sum \alpha_F D_F$, we define a polytope P_D as follows:

$$P_D = \{m \in \mathbb{R}^N \mid \forall F \in P(N-1), \langle m, u_F \rangle \geq -\alpha_F\}. \quad (1.3)$$

For $m \in P_D$, set $\chi^{\langle m, P_D \rangle}$ the monomial in R defined by

$$\chi^{\langle m, P_D \rangle} = \prod X_F^{\langle m, u_F \rangle + \alpha_F}. \quad (1.4)$$

Then $L(D) = \text{Span} \{\chi^{\langle m, P_D \rangle}, m \in P_D \cap \mathbb{Z}^N\}$.

2. Explicit construction of the projective toric code associated to a polytope

2.1. Definition of PC_P

Definition 2.1. Let P be a lattice polytope and \mathbf{X}_P its corresponding toric variety. Call D the divisor such that $P = P_D$ (1.3). Choose a set \mathcal{P} of representatives of \mathbb{F}_q -rational points of \mathbf{X}_P as points of \mathbb{A}^r . We define the *projective toric code* associated

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to P , which we denote by PC_P , as the image of the evaluation map

$$\text{ev}_P : \begin{cases} L(D) \rightarrow \mathbb{F}_q^n \\ \chi^{\langle m, P \rangle} \mapsto (\chi^{\langle m, P \rangle}(\mathbf{x}))_{\mathbf{x} \in P} \end{cases}$$

where $n := \#X_P(\mathbb{F}_q)$. A formula for n in terms of q and the number of faces of P is given in (2.10).

Remark 2.2. One can easily check that, given two r -tuples \mathbf{x}, \mathbf{x}' such that there exists $\mathbf{t} \in G$ satisfying $\mathbf{x}' = \mathbf{t} \cdot \mathbf{x}$, we have

$$\chi^{\langle m, P \rangle}(\mathbf{x}') = \prod_{F \in P(N-1)} t_F^{\alpha_F} \cdot \chi^{\langle m, P \rangle}(\mathbf{x}).$$

Therefore, taking a different set of representatives of the \mathbb{F}_q -points of \mathbf{X}_P gives a Hamming-equivalent code.

Why *projective* toric codes? The classical toric code [14] associated to a polytope P is defined by

$$\mathcal{C}_P = \text{Span}\{(\chi^m(\mathbf{t}) \mid \mathbf{t} \in (\mathbb{F}_q^*)^N), m \in P \cap \mathbb{Z}^N\}, \text{ where } \chi^m(t_1, \dots, t_N) = \prod_{i=1}^N t_i^{m_i}. \quad (2.1)$$

It is the evaluation of some regular functions χ^m on the dense torus \mathbb{T}^N , called *characters*. To get from toric codes to projective ones, we *homogenize* these characters into the monomials $\chi^{\langle m, P \rangle}$ (1.4) in the Cox ring R , which we evaluate at every \mathbb{F}_q -rational point of \mathbf{X}_P , not only at points on the torus. This process is actually the same that turns a Reed-Muller code into a projective one, hence the terminology *projective* toric codes chosen here.

2.2. (Rational) points of \mathbf{X}_P

The main issue arising from Def. 2.1 when handling the code PC_P is the determination of a set \mathcal{P} of representatives of the \mathbb{F}_q -rational points of \mathbf{X}_P in its quotient representation. For instance, in the projective space \mathbb{P}^N , we are used to represent a point by a $(N+1)$ -tuple with its far-left (or far-right) non-zero coordinate equal to 1. We aim to generalize this to other simplicial toric varieties by choosing representatives in a *normalized form*, *i.e.* as tuples with some determined zero coordinates and as many coordinates as possible equal to 1. Let us first fix some notations.

Notation 2.3. Let \mathbb{K} be a field. Let A be a matrix of size $\ell \times m$ with integer entries. For any ℓ -tuple \mathbf{t} with coordinates in \mathbb{K} , we write \mathbf{t}^A the m -tuple whose j -th coordinate of \mathbf{t}^A is equal to $\prod_{i=1}^{\ell} t_i^{a_{ij}}$.

For any couple of matrices A, B with compatible sizes, we have $(\mathbf{t}^A)^B = \mathbf{t}^{AB}$.

This notation and Eq. (1.2) gives a new definition of the group G :

Fact 2.4. Let A_P be the $r \times N$ -matrix whose rows are the primitive normal vectors u_F . The group G is the kernel of the homomorphism defined by

$$\begin{cases} (\overline{\mathbb{F}_q^*})^r \rightarrow (\overline{\mathbb{F}_q^*})^N \\ \mathbf{t} \mapsto \mathbf{t}^{A_P} \end{cases}.$$

While handling \mathbf{X}_P as a quotient space under the action of G , we will be often led to exhibit the existence of preimages under some maps $\mathbf{t} \mapsto \mathbf{t}^A$ for some square matrices A by using the following lemma.

Lemma 2.5. *Let \mathbb{K} be a field. Let A be a square integer matrix of size ℓ . If the determinant of A is coprime to the characteristic of the field \mathbb{K} , then the map*

$$\begin{cases} \overline{\mathbb{F}_q}^\ell \rightarrow \overline{\mathbb{F}_q}^\ell \\ \mathbf{t} \mapsto \mathbf{t}^A \end{cases}$$

is surjective.

Proof. The inverse of A is equal to the transpose of its cofactor matrix, whose entries are integers, divided by its determinant. Therefore, finding a solution of $\mathbf{t}^A = \mathbf{s}$ is possible if we are able to find roots of order $\det(A)$ of any element of $\overline{\mathbb{F}_q}$, which is possible if the determinant of A and the characteristic of $\overline{\mathbb{F}_q}$ are coprime. \square

2.2.1. Local coordinates

The toric variety \mathbf{X}_P is covered by $\#P(0)$ affine toric charts corresponding to the vertices of the polytope P : given a vertex $v \in P(0)$, define the affine open set of \mathbf{X}_P

$$\mathcal{U}_v := \{x_F \neq 0 \mid F \not\ni v\} \subset \mathbb{A}^r \setminus \mathcal{Z}. \quad (2.2)$$

Lemma 2.6. *The G -orbit of a geometric point of \mathcal{U}_v contains an r -tuple $(x_F)_{F \in P(N-1)}$ such that $x_F = 1$ for $F \not\ni v$.*

Proof. Fix $\mathbf{x} = (x_F) \in \mathcal{U}_v$. Let us prove there exists $\mathbf{t} \in G$ such that $t_F = x_F^{-1}$ if $F \not\ni v$. By Eq. (1.2), we can find such a r -tuple $\mathbf{t} = (t_F)_{F \in P(N-1)}$ in G if there exists a N -tuple $(t_F)_{F \ni v}$ such that

$$\forall j \in \{1, \dots, N\}, \prod_{F \ni v} t_F^{u_F^{(j)}} = \prod_{F \not\ni v} x_F^{u_F^{(j)}}. \quad (2.3)$$

Take $A(v)$ the N -square matrix whose rows are the vectors u_F of facets $F \ni v$ (see (H1)). We can solve Eq. (2.3) by finding a N -tuple $\tilde{\mathbf{t}}$ satisfying $\tilde{\mathbf{t}}^{A(v)} = \mathbf{y}$ where \mathbf{y} is the vector formed by the N right-handsides of (2.3), which is possible by Lem. 2.5 and (H2). \square

2.2.2. Normalized forms

The action of the big torus \mathbb{T}^N splits the toric variety \mathbf{X}_P into a disjoint union of finitely many \mathbb{T}^N -orbits, each orbit itself being an algebraic torus:

$$\mathbf{X}_P = \bigsqcup_{Q \preceq P} \mathbb{T}_Q, \tag{2.4}$$

where $\mathbb{T}_Q \simeq (\overline{\mathbb{K}}^*)^{\dim Q}$ [6, Th. 3.2.6]. The torus \mathbb{T}_Q is well-described in the affine set $\mathbb{A}^r \setminus \mathcal{Z}$: it is the set of points whose coordinates x_F are zero only for the facets F containing Q . It is included in any affine chart \mathcal{U}_v (see Eq. (2.2)) associated to a vertex v of Q and it is G -invariant.

We shall give the points of \mathbf{X}_P by determining those on the tori \mathbb{T}_Q as r -tuples in $\mathbb{A}^r \setminus \mathcal{Z}$ up to the action of G .

Definition 2.7. Let $Q \preceq P$ be a face of the polytope P and $v \in P(0)$ a vertex of Q . A v -normalized representative of a point of \mathbb{T}_Q is a r -tuple $\mathbf{x} \in \mathcal{U}_v$ such that $x_F = 1$ for $F \not\ni v$ and $x_F = 0$ for $F \supseteq Q$.

We will call \mathbf{x} a normalized representative of a point p if Q and/or v need not to be specified.

From this definition, some points may have several v -normalized representatives, as illustrated in Ex. 2.8. But we will describe the normalized tuples that represent the same point by using Fact 2.4.

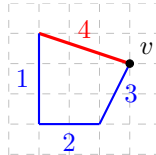


Fig. 1: Polygon $P = \text{Conv}((0, 0), (2, 0), (3, 2), (0, 3))$.

Example 2.8. The four inner normal vectors of the quadrilateral P of Fig. 1 are $u_{F_1} = (1, 0)$, $u_{F_2} = (0, 1)$, $u_{F_3} = (-2, 1)$ and $u_{F_4} = (-1, -3)$. The determinants of the matrices $A(v)$ defined in (H2) are 1, 2, 7 and 3. Let us work in a finite field \mathbb{F}_q with characteristic different from 2, 3 and 7.

Call v the vertex $(3, 2)$. A v -normalized representative of a point p on \mathbb{T}_{F_4} has the form $(1, 1, x, 0)$, since v does not lie on the edges F_1 and F_2 .

Take $t \in \overline{\mathbb{F}}_q$ such that $t^7 = 1$ and $t^3 \neq 1$, which is possible by assumption on the characteristic. Then $\mathbf{t} = (1, 1, t^3, t)$ belongs to the group

$$G = \{(t_1, t_2, t_3, t_4) \in (\overline{\mathbb{K}}^*)^4 \mid t_1 = t_3^2 t_4 \text{ and } t_2 t_3 = t_4^3\},$$

which implies that $(1, 1, x, 0)$ and $(1, 1, t^3 x, 0)$ are both some v -normalized representatives of a same point of \mathbb{T}_{F_4} .

2.2.3. Orbits and Hermite normal form

Notation 2.9. Fix a k -dimensional face Q of P and a vertex v of Q . Since the polytope P is simple (H1), we can number the facets F_1, \dots, F_r of P so that

- F_1, \dots, F_N are exactly the N facets containing the vertex v ,
- $Q = \bigcap_{i=1}^{N-k} F_i$.

Set $A(v, Q)$ the matrix whose rows are the primitive normal vectors u_{F_i} for $i \in \{1, \dots, N\}$. Let $H(v, Q)$ be the *lower triangular Hermite Normal Form* (HNF) of $A(v, Q)$: there exists $T(v, Q) \in \text{GL}_N(\mathbb{Z})$ that satisfies

$$A(v, Q)T(v, Q) = H(v, Q). \quad (2.5)$$

Let us write the matrix $H(v, Q)$ by blocks as follows:

$$H(v, Q) = \begin{pmatrix} L_1(v, Q) & 0 \\ B(v, Q) & L_2(v, Q) \end{pmatrix} \quad (2.6)$$

where $L_1(v, Q)$ and $L_2(v, Q)$ are lower triangular square matrices of size $N - k$ and k respectively.

Lemma 2.10. *Using Notations 2.9 and 2.3, two r -tuples \mathbf{x} and \mathbf{x}' are some v -normalized forms of the same point on \mathbb{T}_Q if and only if*

$$(x_{F_{N-k+1}}, \dots, x_{F_N})^{L_2(v, Q)} = (x'_{F_{N-k+1}}, \dots, x'_{F_N})^{L_2(v, Q)}.$$

Proof.

Two tuples \mathbf{x} and \mathbf{x}' are some v -normalized forms of the same point on \mathbb{T}_Q if and only if there exists $\mathbf{t} \in G$ such that $\mathbf{t} \cdot \mathbf{x} = \mathbf{x}'$. This is equivalent to $(t_F)_{F \in P(N-1)} \in G$, with $t_F = 1$ for $F \not\ni v$ and $t_F = \frac{x'_F}{x_F}$ for $F \ni v$ such that $F \not\supseteq Q$. There is no condition on the other coordinates t_F since $x_F = x'_F = 0$ for $F \supseteq Q$.

Such a tuple \mathbf{t} belongs to G if and only if

$$\left(t_{F_1}, \dots, t_{F_{N-k}}, \frac{x'_{F_{N-k+1}}}{x_{F_{N-k+1}}}, \dots, \frac{x'_{F_N}}{x_{F_N}}, 1, \dots, 1 \right)^{A_P} = \mathbf{1}_r$$

where A_P is the $r \times N$ matrix whose rows are the u_{F_i} , ordered in the same way as the facets. The last coordinates of \mathbf{t} being ones, this is equivalent to

$$\left(t_{F_1}, \dots, t_{F_{N-k}}, \frac{x'_{F_{N-k+1}}}{x_{F_{N-k+1}}}, \dots, \frac{x'_{F_N}}{x_{F_N}} \right)^{A(v, Q)} = \mathbf{1}_N.$$

The vector $\mathbf{1}_N$ being invariant under powering by $T(v, Q)$, this is tantamount to

$$\left(t_{F_1}, \dots, t_{F_{N-k}}, \frac{x'_{F_{N-k+1}}}{x_{F_{N-k+1}}}, \dots, \frac{x'_{F_N}}{x_{F_N}} \right)^{H(v, Q)} = \mathbf{1}_N, \quad (2.7)$$

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where $H(v, Q)$ is the HNF of $A(v, Q)$ (Eq. (2.5)).

It is thus enough to prove that there exists $(t_{F_1}, \dots, t_{F_{N-k}}) \in (\overline{\mathbb{F}_q^*})^{N-k}$ satisfying (2.7) if and only if

$$\left(\frac{x'_{F_{N-k+1}}}{x_{F_{N-k+1}}}, \dots, \frac{x'_{F_N}}{x_{F_N}} \right)^{L_2(v, Q)} = \mathbf{1}_k, \quad (2.8)$$

where the matrix $L_2(v, Q)$ appears in the bloc form of the HNF $H(v, Q)$ (Eq. (2.6)).

The "if" part is clear by definition of $L(v, Q)$ (2.6). *Conversely*, if we assume that Eq. (2.8) holds, computing $(t_{F_1}, \dots, t_{F_{N-k}}) \in (\overline{\mathbb{F}_q^*})^{N-k}$ satisfying Eq. (2.7) essentially consists in extracting roots whose degrees are the diagonal entries of $H(v, Q)$, which are coprime with the characteristic (H2) since $\det A(v, Q) = \pm \det H(v, Q)$. \square

The rational points of \mathbb{T}_Q viewed in the affine chart \mathcal{U}_v are easily characterized.

Lemma 2.11. *Assume that \mathbf{x} is a v -normalized representative of a point p on \mathbb{T}_Q . Then p is \mathbb{F}_q -rational if and only if $(x_{F_{N-k+1}}, \dots, x_{F_N})^{L_2(v, Q)} \in (\mathbb{F}_q^*)^k$.*

Proof. The point p is rational if and only if \mathbf{x} and \mathbf{x}^q are both some of its v -normalized representatives. By Lem. 2.10, this is equivalent to

$$\left(x_{F_{N-k+1}}^{q-1}, \dots, x_{F_N}^{q-1} \right)^{L_2(v, Q)} = \mathbf{1}_k.$$

The proof is concluded by noticing that

$$\left(x_{F_{N-k+1}}^{q-1}, \dots, x_{F_N}^{q-1} \right)^{L_2(v, Q)} = \left((x_{F_{N-k+1}}, \dots, x_{F_N})^{L_2(v, Q)} \right)^{(q-1)I_k}$$

and that $(y_1, \dots, y_k)^{(q-1)I_k} = \mathbf{1}_k$ if and only if $(y_1, \dots, y_k) \in (\mathbb{F}_q^*)^k$. \square

2.3. Piecewise toric code

To construct the evaluation code, we take the naive evaluation of monomials $\chi^{(m, P)}$ at the r -tuples representing the rational points of \mathbf{X}_P , described in Lem. 2.11.

Remark 2.12. Such r -tuples are likely to have coordinates in an extension of \mathbb{F}_q , as their computation may require root extractions. At first glance, one may doubt that the evaluation of a monomial at a rational point lies in \mathbb{F}_q . But, as fairly expected, this is the case as stated in Cor. 2.17.

Given a proper face $Q \preceq P$, let us have a look at the evaluation of monomials $\chi^{(m, P)}$ at the points on \mathbb{T}_Q . They have some zero coordinates, which makes some monomials vanish.

Lemma 2.13. *Let \mathbf{x} be a representative of a point on \mathbb{T}_Q . Then a monomial $\chi^{(m, P)}$ has a non zero evaluation at \mathbf{x} if and only if m belongs to the face Q .*

Proof. Since $x_F = 0$ if and only if $F \supsetneq Q$, the evaluation of $\chi^{\langle m, P \rangle} = \prod X_F^{\langle m, u_F \rangle + a_F}$ is non-zero if and only if $\langle m, u_F \rangle + a_F = 0$ for $F \supseteq Q$, which exactly means that $m \in Q$. \square

The previous lemma and the expression of \mathbf{X}_P as union of tori associated to faces of the polytope P (see Eq. (2.4)) suggest some links between the projective toric codes associated to the polytope P and the classical toric codes C_Q associated to the faces Q of P . The rest of this paragraph aims to prove of the following result.

Proposition 2.14. *For any $Q \in P(k)$, the puncturing of the code PC_P at coordinates corresponding to points outside of \mathbb{T}_Q is monomially equivalent to the toric code C_Q .*

We recall that two codes C_1 and C_2 of length n and dimension k are said to be *monomially equivalent* if, given a generator matrix G_1 of C_1 , there is an invertible diagonal matrix Δ and a permutation matrix Π of size n such that $G_2 = G_1 \Delta \Pi$ is a generator matrix for C_2 .

To prove Prop. 2.14, we first display a more precise expression of the evaluation of a monomial $\chi^{\langle m, P \rangle}$ at a point $p \in \mathbb{T}_Q$ by *straightening* the face Q : we want a unimodular transformation that maps $Q \in P(k)$ into a subspace of dimension k spanned by k vectors of the *canonical basis*. For instance, straightening an edge $Q \in P(1)$ of a polygon makes it horizontal or vertical. We prove that the unimodular transformation associated to the HNF of $A(v, Q)$ (see Eq (2.5)) provides such a map.

Lemma 2.15 (Straightening a face). *Let $Q \in P(k)$ and take a vertex v of Q . Then, the matrix $T(v, Q)$, defined in (2.5), satisfies:*

$$\forall m \in Q, \exists! (\tilde{m}_1, \dots, \tilde{m}_k) \in \mathbb{Z}^k \text{ such that } (m - v)(T(v, Q)^\top)^{-1} = \sum_{i=1}^k \tilde{m}_i e_{N-k+i},$$

where (e_1, \dots, e_N) is the canonical basis of \mathbb{R}^N .

Proof. The matrix $T(v, Q)$ is unimodular: its columns form \mathbb{Z} -basis of \mathbb{Z}^N . Then, for any $m \in Q$, there exists a unique N -tuple $\tilde{m} = (\tilde{m}_1, \dots, \tilde{m}_N) \in \mathbb{Z}^N$ such that

$$m - v = \sum_{j=1}^N \tilde{m}_j e_j T(v, Q)^\top.$$

We want to prove that the $N - k$ first coordinates of \tilde{m} are zero. For $i \in \{1, \dots, N\}$,

$$\langle m - v, u_{F_i} \rangle = \sum_{j=1}^N \tilde{m}_j e_j T(v, Q)^\top (e_i A(v, Q))^\top = \sum_{j=1}^N \tilde{m}_j H(v, Q)_{i,j}.$$

The matrix $H(v, Q)$ being lower triangular, the entry $H(v, Q)_{i,j}$ is zero for $j > i$. In addition, the lattice points m and v belong to Q , then for $i \in \{1, \dots, N - k\}$,

$$0 = \langle m - v, u_{F_i} \rangle = \sum_{j=1}^i \tilde{m}_j H(v, Q)_{i,j}.$$

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Using that $H(v, Q)_{i,i} \neq 0$, a simple induction on i proves that $\tilde{m}_i = 0$ for $i \in \{1, \dots, N - k\}$. Shifting the tuple \tilde{m} gives the expected expression. \square

Proposition 2.16. *Let us take $Q \in P(k)$, a lattice element $m \in Q \cap \mathbb{Z}^N$, a vertex v of Q and \mathbf{x} a v -normalized representative of a point on \mathbb{T}_Q . Then, using Notation 2.9 of Section 2.2.3, there exist $\tilde{m}_1, \dots, \tilde{m}_k \in \mathbb{Z}$ such that*

$$\chi^{\langle m, P \rangle}(\mathbf{x}) = \prod_{j=1}^k y_j^{\tilde{m}_j} \text{ where } (y_1, \dots, y_k) = (x_{F_{N-k+1}}, \dots, x_{F_N})^{L_2(v, Q)} \in (\overline{\mathbb{F}_q}^*)^k. \quad (2.9)$$

Proof. As $m \in Q$, we have $\langle m, u_{F_i} \rangle + a_{F_i} = 0$ for $i \in \{1, \dots, N - k\}$. The element \mathbf{x} is v -normalized, which means that $x_{F_i} = 1$ for $i > N$. Then

$$\chi^{\langle m, P \rangle}(\mathbf{x}) = \prod_{i=N-k}^N x_{F_i}^{\langle m, u_{F_i} \rangle + a_{F_i}}.$$

Since $v \in F_i$ for every $i \in \{1, \dots, N\}$, we have $\langle v, u_{F_i} \rangle = -a_{F_i}$. Therefore, for every $i \in \{1, \dots, k\}$, Lem. 2.15 implies that

$$\langle m, u_{F_{N-k+i}} \rangle + a_{F_{N-k+i}} = \langle m - v, u_{F_{N-k+i}} \rangle = \sum_{j=1}^k \tilde{m}_j (H(v, Q))_{N-k+i, j}.$$

Moreover, we have $(H(v, Q))_{N-k+i, j} = (L_2(v, Q))_{i, j}$ since $L_2(v, Q)$ is the lower right block of $H(v, Q)$ (2.6). Then $(H(v, Q))_{N-k+i, j} = (L_2(v, Q))_{i, j}$ and

$$\chi^{\langle m, P \rangle}(\mathbf{x}) = \prod_{j=1}^k \prod_{i=1}^k x_{F_{N-k+i}}^{\tilde{m}_j (L_2(v, Q))_{i, j}} = \prod_{j=1}^k y_j^{\tilde{m}_j}. \quad \square$$

Corollary 2.17. *The evaluation of a monomial $\chi^{\langle m, P \rangle}$ at a normalized representative of a \mathbb{F}_q -rational point of \mathbb{T}_Q belongs to \mathbb{F}_q .*

Proof. If $m \notin Q$, then $\chi^{\langle m, P \rangle}(x) = 0$ by Lem. 2.13. Otherwise, Lem. 2.11 states that a v -normalized representative of an \mathbb{F}_q -rational point on \mathbb{T}_Q satisfies $\mathbf{y} = (x_{F_{N-k+1}}, \dots, x_{F_N})^{L_2(v, Q)} \in (\mathbb{F}_q^*)^k$. In this case, the quantity given in Prop. 2.16 clearly lies in \mathbb{F}_q . \square

The expression in Prop. 2.16 has exactly the form encountered in toric codes (See Def. 2.1). Before precisely stating a comparison between the restriction of the code PC_P on points of \mathbb{T}_Q and the toric code C_Q , let us recall a result about toric codes made by J. Little and R. Schwarz [21].

We say that integral polytopes P_1 and P_2 in \mathbb{R}^N are *lattice equivalent* if there exists an affine transformation ϕ defined by $\phi(m) = mT + \lambda$ with $T \in \text{GL}_N(\mathbb{Z})$

and $\lambda \in \mathbb{Z}^N$ such that $\phi(P_1) = P_2$. Not only lattice equivalent polytopes define isomorphic toric varieties (Fact 1.2), but they also give equivalent codes.

Theorem 2.18. [21, Th. 3.3] *If two polytopes P_1 and P_2 are lattice equivalent, then the toric codes \mathcal{C}_{P_1} and \mathcal{C}_{P_2} are monomially equivalent.*

Note that the converse is false (see [23] for instance). We have now gathered all the ingredients to prove Prop. 2.14.

Proof of Prop. 2.14. Fix a vertex $v \in Q$. By Prop. 2.16, puncturing PC_P outside \mathbb{T}_Q gives the toric code $\mathcal{C}_{\phi(Q)}$ where ϕ is the invertible affine map defined by $\phi(m) = (m - v)(T(v, Q)^\top)^{-1}$. By Th. 2.18, it is thus monomially equivalent to \mathcal{C}_Q . \square

2.4. A "generator matrix" of the code PC_P

The purpose of this section is to give a matrix $M(P)$ whose rows span the code PC_P . In fact, this matrix represents the evaluation map of Def. 2.1 in the basis formed by the monomials $\chi^{(m, P)}$. It has $\#(P \cap \mathbb{Z}^N)$ rows and n columns, where

$$n := \#\mathbf{X}_P(\mathbb{F}_q) = \sum_{k=0}^N \#P(k) \cdot (q-1)^k \quad (2.10)$$

can be computed from Eq. (2.4), using that $\mathbb{T}_Q(\mathbb{F}_q) \simeq (\mathbb{F}_q^*)^{\dim(Q)}$. Beware that the matrix $M(P)$ may not have full rank : in this case, it is not strictly speaking a generator matrix of the code.

The depiction of $M(P)$ relies on Prop. 2.16. Before giving a protocol to construct $M(P)$ for any polytope, let us focus on the simpler case when P is a polygon.

2.4.1. Example for polygons ($N = 2$)

We number the facets (which are edges) F_1, \dots, F_r and the vertices v_1, \dots, v_r so that $v_1 \in F_1 \cap F_r$ and $v_i \in F_{i-1} \cap F_i$ for $i \geq 2$. By Prop. 2.16, the restriction of the code along an edge is a Reed-Solomon code. Then we can sort the rational points of \mathbf{X}_P and the lattice points of P so that the matrix of the linear map ev_P in the monomial basis has the form given in Fig. 2, where the grey blocks are Vandermonde-type matrices. For $i \in \{1, \dots, r\}$, we set

$$V_i = \begin{pmatrix} \alpha_1 & \alpha_2 & \dots & \alpha_{q-1} \\ \alpha_1^2 & \alpha_2^2 & \dots & \alpha_{q-1}^2 \\ \vdots & \vdots & & \vdots \\ \alpha_1^{l_i} & \alpha_2^{l_i} & \dots & \alpha_{q-1}^{l_i} \end{pmatrix}$$

where $\mathbb{F}_q^* = \{\alpha_1, \alpha_2, \dots, \alpha_{q-1}\}$ and $l_i := \#(F_i^\circ \cap \mathbb{Z})$ is the number of lattice points that are not vertices on the i -th edge.

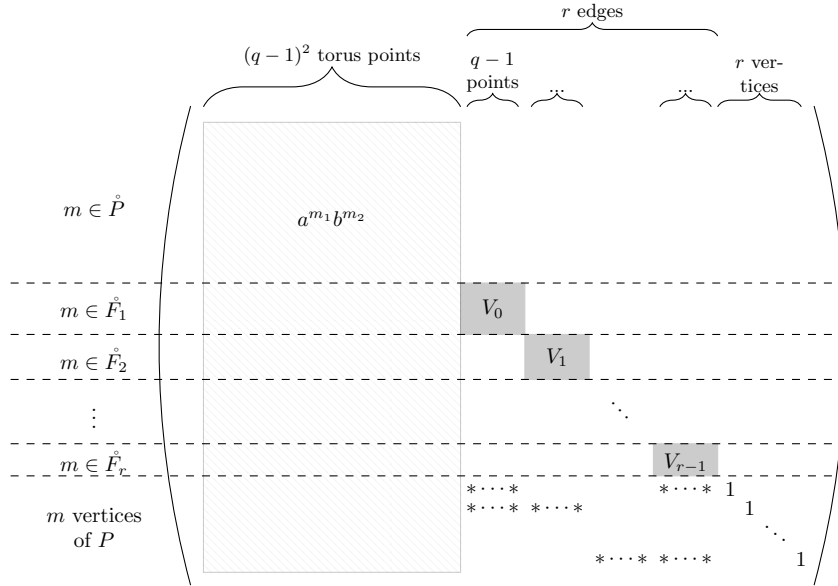


Fig. 2: Matrix of the evaluation map associated to a polygon P ($N = 2$)

2.5. General case

To construct the matrix $M(P)$, we benefit from the piecewise toric structure of the code PC_P , displayed in Prop. 2.14. For each face Q of P , we need an affine transformation ϕ that straightens Q , which will give the evaluation on the torus \mathbb{T}_Q a convenient form. From Sec. 2.2.3, if v is a vertex of Q , the HNF of $A(v, Q)$ provides such a map ϕ . Instead of running through faces Q and choosing a vertex $v \in Q$ to compute $A(v, Q)$, one idea to save some unnecessary computations of HNFs is to choose flags of faces of P , denoted by $\mathcal{F} = (Q_0 \prec Q_1 \prec \dots \prec Q_{N-1} \prec P)$, with $\dim Q_j = j$ ($i = 1, \dots, b$) so that we have one matrix $A_{\mathcal{F}}$ satisfying $A_{\mathcal{F}} = A(Q_0, Q_j)$ for all $j \in \{1, \dots, N\}$. Then the transition matrix associated to the Hermitite normal form of $A_{\mathcal{F}}$ straighten all the faces Q_j at once.

- **Step 0:** Order the points of $P \cap \mathbb{Z}^N$. See Remark 2.19 for suggestion of an order that gives the resulting matrix $M(P)$ a triangular by blocks form.
- **Step 1:** Find b flags of faces of P , denoted by $\mathcal{F}_i = (Q_{i,0} \prec Q_{i,1} \prec \dots \prec Q_{i,N-1} \prec P)$ with $\dim Q_{i,j} = j$ ($i = 1, \dots, b$) so that every face $Q \preceq P$ belongs to one of these flags. We can choose $b = \max_j \#P(j)$.
- **Step 2:** Fix $i \in \{1, \dots, b\}$. For every $j \in \{0, \dots, N - 1\}$, call $E_{i,j}$ the underlying vector space of $Q_{i,j}$, centered in the point $Q_{i,0}$. We want to

determine an invertible integer affine map ϕ_i such that

$$\phi_i(E_{i,j}) = \text{Span}(e_1, \dots, e_j). \quad (\star)$$

Set A_i the N -square matrix whose rows are the normal vectors u_{F_j} of the N facets containing the vertex $Q_{i,0}$, ordered so that $Q_{i,j} = \bigcap_{\ell=0}^{N-j} F_\ell$ for every $j \in \{0, \dots, N\}$. Let H_i be the HNF of A_i and $T_i \in \text{GL}_N(\mathbb{Z})$ the unimodular matrix such that $A_i T_i = H_i$. Denote by T'_i the matrix obtained by reversing the order of the columns of T_i . Then the affine map $\phi_i(m) = T'_i(m - Q_{i,0})$ satisfies (\star) by a similar^a argument as in the proof of Lem. 2.15.

- **Step 3:** Now fix a face $Q \in P(k)$. Then there exists $i \in \{1, \dots, b\}$ such that $Q = Q_{i,k}$.

For every $m \in P \cap \mathbb{Z}^N$, compute the vector $w(m, Q) \in (\mathbb{F}_q)^{(q-1)^k}$ defined by

$$w(m, Q) = \begin{cases} \left(\prod_{j=1}^k x_j^{\phi_i(m)^{(j)}} \mid (x_1, \dots, x_k) \in (\mathbb{F}_q^*)^k \right) & \text{if } m \in Q, \\ (0, \dots, 0) & \text{if } m \notin Q, \end{cases} \quad (2.11)$$

where $\phi_i(m)^{(j)}$ is the j -th coordinate of the vector $\phi_i(m)$. By construction of ϕ_i , if $m \in Q$, then $\phi_i(m)^{(j)} = 0$ for $j \in \{k+1, \dots, N\}$.

Form the matrix $M(P, Q)$, of size $\#(P \cap \mathbb{Z}^N) \times (q-1)^k$ by stacking the vectors $w(m, Q)$ on top of each other, in the order chosen at Step 0.

- **Step 4:** Form the matrix $M(P)$ by putting side by side the matrices $M(P, Q)$ constructed in the previous step.

Remark 2.19. Similarly to the case $N = 2$ (See Fig. 2), we can order the lattice points and the faces of P to make the matrix $M(P)$ triangular by blocks.

First put the block $M(P, P)$, followed by the blocks $M(P, F_i)$ corresponding to facets, then those corresponding to faces of dimension $N - 2$ and so on, making the dimension of the faces decrease:

$$M(P) = \begin{pmatrix} M(P, P) & M(P, F_i) & \dots & M(P, Q_i^k) & \dots & M(P, v_i) \end{pmatrix}$$

$\underbrace{\hspace{10em}}$
big torus

$\underbrace{\hspace{10em}}$
facets

\dots

$\underbrace{\hspace{10em}}$
 k -dim.
faces

\dots

$\underbrace{\hspace{10em}}$
vertices

In the same way, we order the lattice points of P at Step 0 starting by points in the interior of P , then in the interior of facets and so on, ending with the vertices.

This way, given a face $Q \preceq P$, the first rows of $M(P, Q)$ correspond to evaluations of $\chi\langle m, P \rangle$ on \mathbb{T}_Q for lattice points m inside faces $Q' \neq Q$ with $\dim Q' \geq \dim Q$. These rows are thus filled with zeros, by definition of $w(m, Q)$ (Eq. (2.11)).

^aWe just reverse the columns of T_i for (\star) to hold. Otherwise, we would have $\phi_i(E_{i,j}) = \text{Span}(e_{N-j+1}, \dots, e_N)$ as in Lem. 2.15.

3. Dimension

As for classical toric codes, the computation of the dimension of a projective toric code PC_P relies on the reduction modulo $q-1$ of the integral points of the polytope P .

3.1. Reduction modulo $q-1$ and dimension

Definition 3.1. Let us define an equivalence relation on \mathbb{Z}^N as follows: for two elements $(u, v) \in (\mathbb{Z}^N)^2$, we write $u \sim v$ if $u - v \in (q-1)\mathbb{Z}^N$.

Let U be a subset of \mathbb{Z}^N . We will call a *reduction of U* , denoted by \bar{U} , a set of representatives of U under the equivalence relation \sim , that is to say a subset of U with the following property:

$$\forall u \in U, \exists! \bar{u} \in \bar{U} \text{ s.t. } u \sim \bar{u}.$$

If P is a lattice polytope in \mathbb{R}^N , we denote by \bar{P} a reduction of the lattice points of P , i.e. $\bar{P} = P \cap \mathbb{Z}^N$.

Reducing the lattice points of the polytope P modulo $q-1$ provides a basis of a classical toric code [25].

Theorem 3.2. [25] *Let P be a polytope and its associated toric code C_P . Take \bar{P} a reduction of P modulo $q-1$. The kernel of the evaluation map is the \mathbb{F}_q -vector space spanned by*

$$\{\chi^m - \chi^{\bar{m}}, m \in (P \cap \mathbb{Z}^n), \bar{m} \in \bar{P} \text{ such that } m \neq \bar{m}\}.$$

A basis of the code C_P is

$$\{(\chi^{\bar{m}}(\mathbf{t}), \mathbf{t} \in \mathbb{T}^N) \mid \bar{m} \in \bar{P}\}$$

and therefore the dimension of C_P is equal to $\#\bar{P}$.

We benefit from the form of the matrix $M(P)$, constructed in Sec. 2.5, to use this result on classical toric codes in order to get the dimension of projective ones. We are led to reduce the polytope P face by face.

Definition 3.3. For a polytope P , we define the interior of P as the set of points of P that do not belong to any proper face of P , which we denote by

$$P^\circ = \{m \in P \text{ such that } \forall Q \preceq P \text{ with } Q \neq P, m \notin Q\}.$$

Definition 3.4. Given a lattice polytope $P \subset \mathbb{Z}^N$, we define an equivalence relation \sim_P on the set of its lattice points $P \cap \mathbb{Z}^N$ by

$$m \sim_P m' \Leftrightarrow \exists Q \preceq P \text{ such that } m, m' \in Q^\circ \text{ and } m - m' \in (q-1)\mathbb{Z}^N.$$

We call a *projective reduction of P* any set of representatives of elements of $P \cap \mathbb{Z}^N$ modulo \sim_P .

We will denote a projective reduction of P by $\text{Red}(P)$ to emphasize the difference between the relation defined in Def. 3.1 in the study of classical toric codes and in the one of Def. 3.4 that we use here. However, given reductions of the interior of every face $Q \preceq P$, we get a projective reduction of P by setting $\text{Red}(P) = \bigcup_{Q \preceq P} \overline{Q^\circ}$.

Theorem 3.5. *Let P be a polytope and PC_P its associated toric code. Let $\text{Red}(P)$ be a projective reduction of P . Then*

$$\ker \text{ev}_P = \text{Span}\{\chi^m - \chi^{\overline{m}} : m \in (P \cap \mathbb{Z}^n), \overline{m} \in \text{Red}(P) \text{ such that } m \neq \overline{m}\}$$

and the set

$$\left\{ \text{ev}_P(\chi^{\langle \overline{m}, P \rangle}) \mid \overline{m} \in \text{Red}(P) \right\}$$

forms a basis of the code PC_P . The dimension of the code PC_P is equal to $\dim \text{PC}_P = \#\text{Red}(P)$.

Proof. The dimension of PC_P is the rank of the matrix $M(P)$, defined in Section 2.5. By Remark 2.19, it is the sum of the ranks of the matrices $\tilde{M}(P, Q)$. By Prop. 2.14, the rank of $\tilde{M}(P, Q)$ is the dimension of the toric code C_{Q° , which is equal to $\#\overline{Q^\circ}$, by Prop. 3.2. \square

Remark 3.6. Beware that Theorem 3.5 does not provide the dimension of *any* algebraic geometric code associated to a divisor D on \mathbf{X}_P . The result holds if and only if the polytope associated to D has the same normal fan than P , which is equivalent for D to be ample on \mathbf{X}_P [6][Th. 6.2.1]. If this is not the case, the kernel may not be spanned by binomials (see [24] for example on Hirzebruch surfaces).

Example 3.7. Let a, b and η be three positive integers. Consider the quadrilateral P_η whose vertices are $(0, 0)$, $(a, 0)$, (a, b) and $(0, b + \eta a)$ (See Fig. 3). The trapezoids $(P_\eta)_{\eta \in \mathbb{N}}$ describe a family of toric surfaces parametrized by the integer η , called Hirzebruch surfaces (See Fig. 2 [14]).

The black dots are the lattice points of $\overline{Q^\circ}$ for every face Q of P . Figures 3a and 3b respectively display points of $\overline{P^\circ}$ and those of the interior of edges of P . Figure 3c shows the final reduction face by face of P , when adding the vertices to the two previous steps.

The number of black dots is thus equal to the dimension of the code over \mathbb{F}_7 , whose explicit formula depending on (a, b, η) is given by [24, Th. 2.4.1].

4. Gröbner basis theory in the Cox Ring

To compute a lower bound for the minimum distance, we use a technique similar to the one used in the study of Hirzebruch surfaces [24]: we estimate the dimension of a quotient of vectors spaces. We shall take advantage of the graded structure of the Cox ring and use Gröbner basis theory, which may be powerful in this framework as it provides a basis of the quotient of a ring modulo an ideal. Let us first recall some facts about Gröbner bases (See [5] for details).

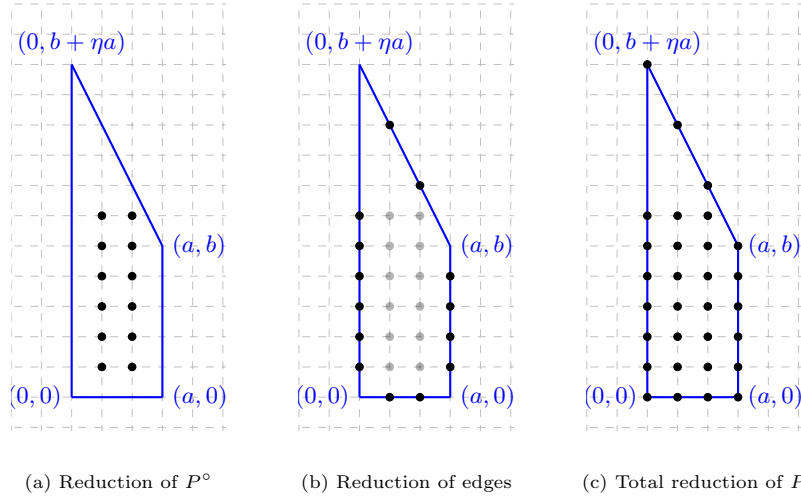


Fig. 3: Reduction face by face of a polygon associated to a Hirzebruch surface \mathcal{H}_2 for $q = 7$

Let R be a polynomial ring. Denote by \mathbb{M} the set of monomials of R . A *monomial order* is a total order on \mathbb{M} , denoted by $<$ that is compatible with the multiplication and such that $1 < M$ for any $M \in \mathbb{M}$. For every polynomial $f \in R$, we call the *leading monomial* of f and denote by $\text{LM}(f)$ the greatest monomial for this ordering that appears in f .

Let I be an ideal of a polynomial ring R over a field^b, endowed with a monomial order $<$. The *monomial ideal* $\text{LM}(I) \subset R$ associated to I is the ideal generated by the leading monomials $\text{LM}(f)$ of all the polynomials $f \in I$. A finite subset G of an ideal I is a *Gröbner basis* of the ideal I if $\text{LM}(I) = \langle \text{LM}(g) \mid g \in G \rangle$.

The pleasing property of Gröbner bases that will be used to estimate the minimum distance of the code is the following.

Proposition 4.1. [29, Prop. 1.1] *Let I be an ideal of a polynomial ring R over a field \mathbb{K} with Gröbner basis G . Then, setting π as the canonical projection of R onto R/I , the set*

$$\{\pi(M), M \in \mathbb{M} \text{ s.t. } \forall g \in G, \text{LM}(g) \nmid M\}$$

is a basis of R/I as a \mathbb{K} -vector space.

4.1. Application in the Cox ring

Let us set a monomial order on the Cox Ring of \mathbf{X}_P by using its grading by $\text{Cl}(\mathbf{X}_P)$ and the correspondence between lattice points in P_D and element of $R_{[D]}$ (See Par.

^bWe assume R to be a polynomial ring over a *field* to avoid the distinction between leading term and leading monomial.

1.3.2).

Notation 4.2. Fix a total order $<$ on $\text{Cl}(\mathbf{X}_P)$ such that $[0] < [D]$ for any non-zero effective divisor D . Fix an order on \mathbb{Z}^N that is compatible with the addition, which we also denote by $<$. From these, we define a total order $<$ on R as follows.

Given D_1 and D_2 two divisors on \mathbf{X}_P and m_1, m_2 two tuples in \mathbb{Z}^N such that $m_i \in P_{D_i}$, we set $\chi^{\langle m_1, D_1 \rangle} < \chi^{\langle m_2, D_2 \rangle}$ if one of the following holds:

- (1) $[D_1] < [D_2]$ in $\text{Cl}(\mathbf{X}_P)$,
- (2) $[D_1] = [D_2]$ (then $\exists m \in \mathbb{Z}^n$, such that $P_{D_1} = m + P_{D_2}$) and $m_1 < m + m_2$.

In the Cox ring R , we define I as the homogeneous vanishing ideal in the subvariety consisting of the \mathbb{F}_q -rational points of \mathbf{X}_P . As the Cox ring is graded by the class group of the variety, the homogeneous component of I of "degree" $\alpha \in \text{Cl}(\mathbf{X}_P)$ is isomorphic to the kernel of ev_D , for any $D \in \alpha$.

Provided a monomial order on R , Prop. 4.1 gives a basis of R/I as \mathbb{F}_q -vector space. Taking the elements of degree $[D]$ gives a basis of the homogeneous component $R_{[D]}/I_{[D]}$. Since

$$R_{[D]}/I_{[D]} \simeq L(D)/\ker \text{ev}_P,$$

Theorem 3.5 states that a projective reduction of P provides a basis of this same vector space. We would like to make these two bases coincide by choosing a projective reduction of P that takes into account the monomial order we consider.

Definition 4.3. Let $<$ a total order on \mathbb{Z}^N . Let P be a lattice polytope and $m \in P \cap \mathbb{Z}^N$. Define $\text{Red}_{<}(m, P)$ as the minimum lattice point of P under the order $<$ such that $\text{Red}_{<}(m, P) \sim_P m$. By extension, we define the *projective reduction of P with respect to the order $<$* :

$$\text{Red}_{<}(P) = \{\text{Red}(m, P), m \in P \cap \mathbb{Z}^N\}.$$

Example 4.4. The points displayed in Fig. 3 are in fact the elements of $\text{Red}_{<}(P)$ for the lexicographic order.

Lemma 4.5. *Given an order $<$ on the Cox ring as in Not. 4.2, there exists a Gröbner basis \mathcal{G} of the ideal I such that*

$$\{\chi^{\langle m, P \rangle}, m \in \text{Red}_{<}(P)\} = \{M \in \mathbb{M} \cap R_{[D]} \text{ such that } \forall g \in \mathcal{G}, \text{LM}(g) \nmid M\}.$$

Proof. By Theorem 3.5, the set

$$\mathcal{B}_D := \{\chi^{\langle m, P \rangle} - \chi^{\langle \text{Red}_{<}(m, P), P \rangle}, m \in P \text{ such that } m \neq \text{Red}_{<}(m, P)\}$$

forms a a basis of $I_{[D]} \subset I$, where D is the divisor associated to P . Take a Gröbner basis \mathcal{G} of I that contains \mathcal{B}_D . It is always possible: a Gröbner basis of I to which we add elements of I remains a Gröbner basis of I .

The set of monomials $\{\chi^{\langle m, P \rangle}, m \in \text{Red}_{<}(P)\}$ contains exactly the monomials that are not divisible by any $\text{LM}(g)$ for any $g \in \mathcal{B}_D$. Then, we have

$$\{\chi^{\langle m, P \rangle}, m \in \text{Red}_{<}(P)\} \supseteq \{M \in \mathbb{M} \cap R_{[D]} \text{ such that } \forall g \in \mathcal{G}, \text{LM}(g) \nmid M\}.$$

Both of these sets being some bases of the vector space $R_{[D]}/I_{[D]}$, they are thus equal. \square

The result brought by the Gröbner basis theory involves divisibility relation in the polynomial ring we consider, here the Cox ring of \mathbf{X}_P . The following lemma takes advantage of the correspondence between monomials in this Cox ring and lattice points of polytopes to translate divisibility of monomials into a relation between the corresponding points in \mathbb{Z}^N .

Lemma 4.6. *Take two divisors D_1 and D_2 of \mathbf{X}_P and their associated polytopes $P_{D_1} = P_1$ and $P_{D_2} = P_2$ (defined in Eq. (1.3)). Let P_3 be the polytope corresponding to the divisor $D_2 - D_1$. Then, in the Cox ring of \mathbf{X}_P , for every $m_1 \in P_1 \cap \mathbb{Z}^N$ and $m' \in P_2 \cap \mathbb{Z}^N$,*

$$\chi^{\langle m_1, P_1 \rangle} \mid \chi^{\langle m_2, P_2 \rangle} \Leftrightarrow m_2 - m_1 \in P_3$$

Proof. Write $D_i = \sum_F a_{i,F} D_F$ ($i = 1, 2$). The monomial $\chi^{\langle m_1, P_1 \rangle}$ divides $\chi^{\langle m_2, P_2 \rangle}$ if and only if for every $F \in P(N-1)$, we have

$$\langle m_1, u_F \rangle + a_{1,F} \leq \langle m_2, u_F \rangle + a_{2,F} \Leftrightarrow \langle m_2 - m_1, u_F \rangle \geq -(a_{2,F} - a_{1,F}).$$

This exactly means that $m_2 - m_1 \in P_3$ by (1.3). \square

5. Minimum distance

The dimension of the code PC_P has been related to the lattice points of the polytope P we hope for a similar kind of result for the minimum distance. Here another polytope comes into play.

Definition 5.1. Let P be a lattice polytope. A lattice polytope P' is said to be *P -surjective* if it has the same normal fan as P , contains P and $\text{PC}_{P'} = \mathbb{F}_q^{\#\mathbf{X}_P(\mathbb{F}_q)}$.

Two polytopes with the same normal fan are said normally equivalent and define the same toric variety (Fact 1.2). Note that the set of polytopes which are normally equivalent to a given simple polytope P modulo translation is well characterized [7, Lem. 2.1].

Asking a P -surjective polytope to contain P enables us to embed $L(D_P)$ into $L(D_{P'})$. Finally, Prop. 3.5 entails that the reduction modulo $q-1$ of the interior lattice points of each k -dimensional faces of P' has $(q-1)^k$ elements.

One can always search for a P -surjective polytope as a dilatation λP of P . Indeed, for any n -dimensional lattice polytope, the dilated polytope $((n+1)P)$ has

at least one interior point. Then some multiple of P has the expected number of interior points in each of its faces.

Given a P -surjective polytope and an order on \mathbb{Z}^N , one can deduce a lower bound of the minimum distance, as stated by the following theorem.

Theorem 5.2. *Let P and P_{surj} be two lattice polytopes which are respectively associated to the divisors D and D_{surj} on \mathbf{X}_P . Assume that P_{surj} is P -surjective. Denote by P' the polytope associated to the divisor $D_{\text{surj}} - D$.*

Let $<$ be an order on \mathbb{Z}^N that is compatible with the addition. Then the minimum distance of the code PC_P satisfies

$$d(\text{PC}_P) \geq \min_{m \in \text{Red}_{<}(P)} \#((m + P') \cap \text{Red}_{<}(P_{\text{surj}})).$$

Remark 5.3. For each couple $(<, P')$ of an order on \mathbb{Z}^N and a P -surjective polytope, Th. 5.2 provides a lower bound for the minimum distance of PC_P . Computing $\text{Red}_{<}(P_{\text{surj}})$ is easier when P_{surj} is not too big but finding a small P -surjective polytope may not be trivial, as illustrated in Paragraph 6.1. Experimentally the lower bound does not seem to depend on the choice of P_{surj} . It is thus more reasonable to change the ordering. We shall avoid computing several times the reduction of P and P_{surj} by computing the "classes" modulo $q - 1$ once and for all and then taking the minimum elements with respect to different orderings. The bound highly depends on the order $<$, as illustrated by the following example.

Example 5.4. Take $P = \text{Conv}((0, 0), (1, 4), (4, 1))$. The classical toric code C_P , studied by D. Joyner [17], has parameters [49, 11, 28] over \mathbb{F}_8 . The smallest multiple of P that is P -surjective is $P_{\text{surj}} = 8P$. The lower bound for the minimum distance given by the previous theorem is equal to 30 for the lexicographic order and 33 for the reverse lexicographic order. Let us consider the order on \mathbb{Z}^2 defined by

$$(x, y) < (x', y') \text{ if } y < y' \text{ or } (y = y' \text{ and } x > x').$$

For this order, the bound is trivial since $(1, 3) + 7P$ does not meet $\text{Red}_{<}(8P)$. The actual parameters of the code are [73, 11, 35] (using MAGMA).

The proof of Th. 5.2 is split in two parts. Lem. 5.5 gives a lower bound of $d(\text{PC}_P)$ in terms of the dimension of a quotient of vector spaces. Lem. 5.6 benefits from results of the previous section about Gröbner bases to handle this dimension as number of points in the polytope P_{surj} .

Lemma 5.5. *Call D (resp. D_{surj}) the divisor on \mathbf{X}_P such that $P = P_D$ (1.3) (resp. $P_{\text{surj}} = P_{D_{\text{surj}}}$). For any $f \in L(D) \setminus \ker \text{ev}_P$, we define*

$$\widetilde{n}_f := \dim \left(L(D_{\text{surj}}) /_{\ker \text{ev}_{P_{\text{surj}}}} + f \cdot L(D_{\text{surj}} - D) \right).$$

Then the minimum distance $d(\text{PC}_P)$ satisfies $n - d(\text{PC}_P) \leq \max_{\substack{f \in L(D) \\ f \notin \ker \text{ev}_P}} \widetilde{n}_f$.

Proof. As usual in the framework of algebraic geometric codes, we can write that

$$d(\text{PC}_P) = n - \max_{\substack{f \in L(D) \\ f \notin \ker \text{ev}_P}} n_f$$

where $n_f = \#\mathcal{Z}(f)(\mathbb{F}_q)$ is the number of \mathbb{F}_q -points in the zero set $\mathcal{Z}(f)$ of f in \mathbf{X}_P .

Fix a polynomial $f \in L(D) \setminus \ker \text{ev}_P$. The map

$$\text{ev}_{P_{\text{surj}},f} : \begin{cases} L(D_{\text{surj}}) \rightarrow \mathbb{F}_q^{n_f} \\ g \mapsto (g(p))_{p \in \mathcal{Z}(f)(\mathbb{F}_q)} \end{cases},$$

is the composition of the evaluation $\text{ev}_{P_{\text{surj}}}$ and the projection onto the zero points of f . Then, it is surjective and

$$n_f = \dim \left(L(D_{\text{surj}}) / \ker \text{ev}_{P_{\text{surj}},f} \right).$$

The kernel of $\text{ev}_{P_{\text{surj}},f}$, contains the kernel of the map $\text{ev}_{P_{\text{surj}}}$ and the multiples of F , which implies that $n_f \leq \widetilde{n}_f$ and concludes the proof. \square

We want to handle the quantity \widetilde{n}_f for $f \in L(D) \setminus \ker \text{ev}_P$. Thanks to Th. 3.5, we are able to pick a monomial basis $\mathcal{B}_{D_{\text{surj}}}$ of $L(D_{\text{surj}})$ modulo $\ker \text{ev}_{P_{\text{surj}}}$ by reducing lattice points of P_{surj} modulo $q-1$ while taking into account the monomial order. We now pick some of these monomials to form a linearly independent family in $L(D_{\text{surj}})$ modulo $(\ker \text{ev}_{P_{\text{surj}}} + f \cdot L(D_{\text{surj}} - D))$, which provides an easy bound on $n - \widetilde{n}_f$.

Lemma 5.6. *Using notations of Theorem 5.2 and Lem. 5.5, let us denote by m_f the lattice point of P such that $\text{LM}(F) = \chi^{\langle m_f, P \rangle}$ for $f \in L(D)$. Then*

$$n - \widetilde{n}_f \geq \#\{m \in \text{Red}_{<}(P_{\text{surj}}) \text{ such that } m - m_f \in P'\}.$$

Proof.

As the polynomial $f \in L(D)$ is a homogeneous element, the sum of the vanishing ideal I of $\mathbf{X}(\mathbb{F}_q)$ and the ideal generated by f in R is also homogeneous. Let $\widehat{\mathcal{G}}$ be a Gröbner basis of the ideal $I + \langle f \rangle$ that contains the Gröbner basis \mathcal{G} of I provided by Lem. 4.5 and the polynomial f . Using Prop. 4.1 and restricting to the component of degree $[D_{\text{surj}}]$ of R/I , we can deduce that the set

$$\mathcal{M}_f := \{M \in \mathbb{M} \cap R_{[D_{\text{surj}}]} \text{ such that } \forall h \in \widehat{\mathcal{G}}, \text{LM}(h) \nmid M\}$$

is a basis of $L(D_{\text{surj}})$ modulo $\ker \text{ev}_{P_{\text{surj}}} + f \cdot L(D_{\text{surj}} - D)$ of cardinality \widetilde{n}_f . Since $\mathcal{G} \subset \widehat{\mathcal{G}}$ and $f \in \widehat{\mathcal{G}}$,

$$\mathcal{M}_f \subset \{M \in \mathbb{M} \cap R_{[D_{\text{surj}}]} \text{ such that } \forall g \in \mathcal{G}, \text{LM}(g) \nmid M \text{ and } \text{LM}(f) \nmid M\}.$$

Applying Lem. 4.5, we get

$$\mathcal{M}_f \subset \{\chi^{\langle m, P_{\text{surj}} \rangle}, m \in \text{Red}_{<}(P_{\text{surj}}) \text{ such that } \chi^{\langle m_f, P \rangle} \nmid \chi^{\langle m, P_{\text{surj}} \rangle}\}.$$

Since $n = \#\text{Red}_<(P_{\text{surj}})$, we have

$$n - \tilde{n}_f \geq \#\{m \in \text{Red}_<(P_{\text{surj}}) \text{ such that } \chi^{\langle m_f, P \rangle} \mid \chi^{\langle m, P_{\text{surj}} \rangle}\},$$

which, thanks to Lem. 4.6, gives the expected result. \square

Combining the two previous lemmas, we are now able to prove Th. 5.2.

Proof of Th. 5.2. By Lem. 5.5,

$$d(\text{PC}_P) \geq \min_{\substack{f \in L(D) \\ f \notin \ker \text{ev}_P}} n - \tilde{n}_f.$$

Lem. 5.6 gives a lower bound of $n - \tilde{n}_f$ in terms of the lattice point $m_f \in P$ such that $\text{LM}(f) = \chi^{\langle m_f, P \rangle}$. However, by Lem. 4.5, we can assume that $m_f \in \text{Red}_<(P)$, as the associated monomials in $L(D)$ form a basis of $L(D)$ modulo $\ker \text{ev}_P$. This proves that it is enough to take the minimum over the lattice points of $\text{Red}_<(P)$ \square

6. Examples

6.1. A toy example

Let us work out a toy example to illustrate the results on the dimension and the minimum distance in combinatorial terms. Let P be the triangle of vertices $(0, 0)$, $(1, 0)$ and $(-2, 3)$. We want to determine the parameters of the code PC_P over \mathbb{F}_4 . The dimension of this code is equal to the number of lattice points of P , since P is entirely contained in the square $[0, q-2]^2$.

For the minimum distance, we look for a P -surjective polygon in the form $P_{\text{surj}} = \lambda P$ with $\lambda \in \mathbb{N}$. Note that $\lambda \geq 4$ so that P_{surj} has at least $q-1 = 3$ interior lattice points of each of its edges.

Let us detail the projective reduction modulo $q-1$ of $4P$ (See Fig. 4a). Each edge contains at least 3 points in distinct classes modulo $q-1$ in its interior. Empty circles represent such points that are minimal for the lexicographic order. However, the reduction of the interior of $4P$ contains only 8 classes, whose representatives are depicted by diamonds. Points belonging to the same class are linked by a path in Fig. 4a. A P -surjective polygon is expected to have $(q-1)^2 = 9$ classes for its interior lattice points. Therefore, the polygon $4P$ is not P -surjective and we need to choose λ larger. One can easily see that $\lambda = 5$ suits.

Now we can bound the minimum distance: make the projective reduction of P and $P_{\text{surj}} = 5P$ with respect to an ordering $<$ on \mathbb{Z}^2 - here the lexicographic order - (see Fig. 4b) to compute the minimum of Th. 5.2. It is reached for $m \in \{(1, 0), (0, 1)\}$: the cardinality of $(P_{\text{surj}} - P + m) \cap \text{Red}_<(P_{\text{surj}})$ equals 8. A computation with MAGMA ensures that it is the actual minimum distance.

Choosing $P_{\text{surj}} = \lambda P$ with $\lambda \geq 5$ leads to the same lower bound but the reduction of P_{surj} become heavier as λ grows.

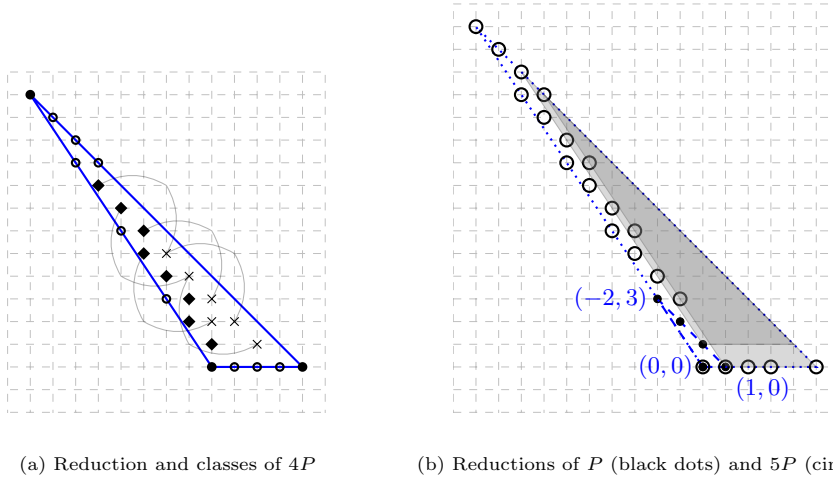


Fig. 4: Reductions of P , $4P$ and $5P$ for $P = \text{Conv}((0, 0), (1, 0), (-2, 3))$ on \mathbb{F}_4 with respect to the lexicographic order

6.2. *Towards new champion codes*

G. Brown and A. M. Kasprzyk [2] systematically investigated toric codes and their generalized versions associated to points in small polygons. This way, they exhibit large families of good codes, achieving and sometimes beating the best-known parameters [11]. Given a champion toric code C_P , even if its projective version PC_P is unlikely to be a champion code itself, it may indicate how to extend C_P while keeping good parameters.

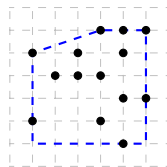


Fig. 5: A polygon containing the points defining a champion generalized toric code [49, 14, 26] over \mathbb{F}_8 [2]

Example 6.1. Take U the set formed by the black dots in Fig. 5. It defines a champion generalizing toric code [49, 14, 26] over \mathbb{F}_8 [2]. Its convex hull does not define a simplicial toric variety over \mathbb{F}_8 since it does not fulfill (H2). However, it is contained in the blue dotted polygon P that defines a simplicial toric surface X_P over \mathbb{F}_8 and a [87, 14, 34] code PC_P . With MAGMA, evaluating the monomials

corresponding to the lattice points in U at torus points but also at two other points gives a $[51, 14, 27]$ code, and adding again two other points produces a $[53, 14, 28]$ code. Both of these codes have best-known parameters [11]. We have to add three other evaluation points to improve the minimum distance by 1 once more, whereas a $[54, 14, 29]$ code is already referenced.

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