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# AUTOMATIC DETECTION OF REPEATED OBJECTS IN IMAGES 

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#### Abstract

The definition of an "object" through the presentation of several of its instances is certainly one of the most efficient ways for humans and machines to learn. An object can be "learned" from a single image, just because it is repeating. In this paper, we explore a three step algorithm to detect repeated objects in images. Starting from a graph of auto-correspondences inside an image, we first extract subgraphs composed of repetitions of unbreakable pieces of objects, that we call atoms. Then, these graphs of atoms are grouped into initial propositions of object instances. Finally, geometry inconsistencies are filtered out to end up with the final repeated object. The meaningfulness of object repetitions is measured by their Number of False Alarms (NFA), which provides a natural order among repeated objects in images; a very low NFA being a strong proof of existence of the discovered object. Source codes are available at https: //rdguez-mariano.github.io/pages/autosim.


Index Terms- image comparison, SIFT, autosimilarities, symmetry detection, RANSAC, NFA.

## 1. INTRODUCTION

Everyday images often contain repeated objects, such as roof tiles, windows on buildings or chairs in a classroom. Humans not only identify these repetitions but also extract meaningful information from them. Sometimes, repetitions follow regularity patterns (e.g. directional, grids, symmetries, etc), but in other cases, no apparent structure is perceived.

Reliably detecting repetitions is challenging. Most approaches of the literature require strong assumptions. For example, in [1] a 2 D repetition grid is required. In [2], and similar to [3], authors assume that large repetitive structures repeat along the horizontal direction in facade images. More specific relationships for repeated elements have also been investigated in [4-7]. Yet another active field imposing patterns to repetitions is symmetry detection [8-13].

Some state-of-the-art matching methods are compatible with the detection of repeated structures [14-19]. Others sacrifice this feature in order to be distinctive in practical applications [20]. Ideally, the best suited matching method for repetitions should return fully connected groups of similar key points whenever repeated objects are present in images; also, it should not connect to unrelated objects.

This paper proposes a method free from any a priori geometric assumption on the repetition patterns. Figure 1 describes the main steps of our proposal. Our input data is a graph of autocorrespondences in the considered image, obtained by applying a reformulation of SIFT [21] presented in [19] (based on the $a$ contrario procedure [22]) between the image and itself. Our bottom up method starts by identifying sub-graphs corresponding to repetitions of small object pieces (called atoms), then groups these atoms to form repeated objects. An atom in our approach is an image patch


Fig. 1: Proposed steps for automatic detection of repeated objects. (a): Creating graphs of repeated atoms and filtering out mutations (black $\times$ ); (b): Grouping of key points for graphs of equal cardinality; (c): A repeated object is identified if geometry consistency is assessed among groups.
(a local element), equipped with a key point (the center's coordinates) and a local description. Repetitions are identified regardless of any a priori pattern, but the presence of several atoms in an object is required to validate the geometric transformation between instances of the object, and therefore identify this object. At this point, two atoms may be connected by a path of strong similarity scores, while not corresponding to the exact same physical part of an object (see Figure 2). We call these, mutations. Unfortunately, these mutations often occur between two spatially close atoms, and if not treated, they may harm the detection process. To avoid this situation, we treat these mutations by splitting the sub-graphs containing overlapping atoms.

## 2. GRAPHS OF AUTO-CORRESPONDENCES IN IMAGES

In this section we introduce the mathematical objects that will be useful in the paper. Let $V=\left\{v_{1}, \cdots, v_{n}\right\}$ be a set of atoms in an image $\mathbf{u}$, and let $w_{i j}$ be the similarity score between the local descriptors from atoms $v_{i}$ and $v_{j}$. We define $G=(V, E)$, the nondirected graph with set of nodes $V$ and set of edges $E$. Edges represent the detected auto-correspondences between atoms. An edge between nodes $v_{i}$ and $v_{j}$, denoted by $v_{i} \leftrightarrow v_{j}$, is weighted by $w_{i j}$.

Definition 1. Let $\tilde{G}_{1}=\left(\tilde{V}_{1}, \tilde{E}_{1}\right)$ and $\tilde{G}_{2}=\left(\tilde{V}_{2}, \tilde{E}_{2}\right)$ be two subgraphs of $G=(V, E)$. We define the merging operation between two graphs, $\tilde{G}_{1}$ and $\tilde{G}_{2}$, as the graph

$$
\tilde{G}_{1} \bigoplus \tilde{G}_{2}:=\left(\tilde{V}_{1} \bigcup \tilde{V}_{2}, \tilde{E}_{1} \bigcup \tilde{E}_{2}\right) \subset G
$$

Definition 2. Let $A \subset V$, we say $A$ is a connected component if and only if any two nodes in $A$ are connected by a path of positive edges in $E$ and there is no path of positive edges connecting $A$ and $A^{c}$.

(a) (Ideal behavior) A fully connected component (4 nodes, 6 edges) from the similarity graph, presenting one and only one atom.

(b) (Undesired behavior) A connected component (7 nodes, 7 edges) from the similarity graph, presenting three atom mutations; see the three pairs of overlapping atoms (i.e. spatially close atoms).

Fig. 2: SIFT connected components.

Definition 3. Let $\tilde{G}=(\tilde{V}, \tilde{E})$ be a sub-graph of $G=(V, E)$. We denote

$$
\text { Int }(\tilde{G}):=\{m \in E \mid m \text { connects two nodes of } \tilde{V}\}=\tilde{E}
$$

and
$B d(\tilde{G}):=\left\{m \in E \mid m\right.$ connects a node of $\tilde{V}$ to a node of $\left.\tilde{V}^{c}\right\}$,
the interior and boundary sets of $\tilde{G}$, respectively.
Remark 1. Both interior and boundary sets of a graph defined above are composed of edges, not nodes.

Definition 4. We say that $\mathcal{G}$ is a partition of $G=(V, E)$ if and only if

$$
\tilde{G} \in \mathcal{G} \Rightarrow \tilde{G} \subset G ; \quad \bigcup_{\tilde{G} \in \mathcal{G}} \operatorname{Int}(\tilde{G})=E
$$

and

$$
\tilde{G}_{1}, \tilde{G}_{2} \in \mathcal{G} \Rightarrow \operatorname{Int}\left(\tilde{G}_{1}\right) \bigcap \operatorname{Int}\left(\tilde{G}_{2}\right)=\emptyset .
$$

Definition 5. Let $G=(V, E)$ be a graph. Let $\{x \leftrightarrow y\}=m \in E$. We define the singleton graph as: $S(m):=(\{x, y\}, m) \subset G$.

Definition 6. Let $\mathcal{G}$ be a partition of $G$. We denote by $\mathcal{G}_{m}, m \in$ Int $(G)$, the only element $\tilde{G} \in \mathcal{G}$ such that $m \in \operatorname{Int}(\tilde{G})$.

## 3. RETRIEVING ATOMS IN IMAGES

Small local elements, called atoms, are the unbreakable pieces that compose objects. Ideally, a perfect matching method should return fully connected components in the similarity graph whenever atom repetitions are observed (see Figure 2a); in those cases, any node can act as the class representative for the atom repetition. Unfortunately, it is often the case with existing matching methods that more than one atom repetition are present in a connected component. This is due to mutations, see Figure 2b, which are often observed with spatially close atoms. Also, atom repetitions do not always appear as fully connected components, due to missed matches. Nevertheless, the intuition supports that all 'copies' of an atom should belong to a unique connected component.

If not treated properly, mutations from spatially close atoms are a nuisance to the detection process, for they may point to a bigger
(and misleading) number of repetitions, see Figure 2b. In order to avoid the aforementioned nuisance, we encourage the splitting of connected components harboring overlapping atoms. Other types of mutation not harming the inferred number of repetitions are more easily handled as geometry consistency will be imposed in the next steps.

In this section we want to identify atoms to the interior of some similarity graph. Each interior should be free from mutations, so it can be considered to harbor several instances of only one atom. For that, a pre-processing of the raw similarity graph is required. Usually, clustering algorithms, related to the minimization of a functional (i.e. RatioCut, Ncut, etc), are used for this purpose. However, the main drawback of best-performing clustering algorithms is the added time in computations. Furthermore, it is not straightforward for classic clustering algorithms to add a free-from-mutation constraint to the functional. With this in mind, we define a simple functional, well suited to our problem, and propose a heuristic to optimize it.

Let $\mathcal{G}=\left\{G_{1}, \cdots, G_{k}\right\}$ be a partition of graphs. Starting from the thinnest partition $\mathcal{G}$ (i.e. composed of isolated matches), the merging procedure will be guided by the following functional:

$$
\begin{equation*}
f(\mathcal{G}):=\sum_{G \in \mathcal{G}}\left[f_{I}(G)+f_{O}(G)\right] \tag{1}
\end{equation*}
$$

where

$$
f_{I}(G)=\frac{1}{\sum_{e \in E} w_{e}},
$$

and

$$
f_{O}(G)=\left\{\begin{aligned}
0, & G \text { has no overlapping atoms } \\
\infty, & G \text { has overlapping atoms }
\end{aligned}\right.
$$

The functional $f$ will encourage the merging of graphs resulting in higher cumulative similarities and no overlapping atoms. Algorithm 1 proposes an heuristic to minimize the functional appearing in Equation 1. It ensures that sub-graphs in the resulting partition will not exhibit the undesired behavior of Figure 2b. In order to minimize the proposed functional, Algorithm 1 keeps track of the interior and boundary sets after each merge with the following formulas:

$$
\begin{aligned}
& \operatorname{Int}\left(\tilde{G}_{1} \oplus \tilde{G}_{2}\right)=\operatorname{Int}\left(\tilde{G}_{1}\right) \cup \operatorname{Int}\left(\tilde{G}_{2}\right), \\
& \operatorname{Bd}\left(\tilde{G}_{1} \oplus \tilde{G}_{2}\right)=\operatorname{Bd}\left(\tilde{G}_{1}\right) \cup \operatorname{Bd}\left(\tilde{G}_{2}\right) \backslash \operatorname{Int}\left(\tilde{G}_{1} \oplus \tilde{G}_{2}\right) .
\end{aligned}
$$

```
Algorithm 1 Greedy heuristic to minimize Equation 1
input: M - set of Matches. It includes similarity measurements.
start:
M.SORT() \(\triangleright\) Order matches with respect to similarity measurements
\(\mathcal{G}=\bigcup_{m \in M} S(m) \quad \triangleright\) Initial partition of singleton-graphs
\(f^{*}=\infty \quad \triangleright f^{*}\) serves to identify when \(\mathcal{G}\) is not being modified
while \(f(\mathcal{G})<f^{*}\) do
    \(f^{*}=f(\mathcal{G})\)
    foreach \(m \in M\) do
                                    \(\triangleright\) Matches with higher similarity scores are analyzed first
        \(m_{1}^{*} \in \arg \min _{m_{1} \in \operatorname{Bd}\left(\mathcal{G}_{m}\right)} f\left(\mathcal{G}_{m} \oplus \mathcal{G}_{m_{1}}\right)-f\left(\mathcal{G}_{m}\right)-f\left(\mathcal{G}_{m_{1}}\right)\)
        if \(f\left(\mathcal{G}_{m} \bigoplus \mathcal{G}_{m_{1}^{*}}\right)<f\left(\mathcal{G}_{m}\right)+f\left(\mathcal{G}_{m_{1}^{*}}\right)\) then
            \(\mathcal{G}=\mathcal{G}_{m} \bigoplus \mathcal{G}_{m_{1}^{*}} \cup \mathcal{G} \backslash\left\{\mathcal{G}_{m}, \mathcal{G}_{m_{1}^{*}}\right\}\)
                                    \(\triangleright\) Update the graph partition
```

return: $\mathcal{G}$

## 4. FROM ATOMS TO OBJECTS

In the previous section we have addressed the problem of identifying repetitions of unbreakable pieces of objects, atoms. These atom repetitions are provided by Algorithm 1 as a partition of graphs, $\mathcal{G}$. In this section, we pass from repeated atoms to the detection of repeated objects.

### 4.1. Pre-labeling object instances

Intuitively, an object atom is supposed to be present in all instances of the object in question, i.e., the number of object instances equals the number of instances of any of its atoms. Let $\Upsilon_{N_{G}}$ be the set of all graphs in $\mathcal{G}$ with fixed cardinality $N_{G}$, representing all atoms that repeat $N_{G}$ times in the image. For a fixed number of instances $N_{G}$, we define the most meaningful object repetition as the one that possesses the more atoms in $\Upsilon_{N_{G}}$. We assume that no image zone should belong to two instances of an object at the same time.

In this section, we look for a function $L$ that assigns to each node of a graph in $\Upsilon_{N_{G}}$, a tentative instance of the most meaningful $N_{G^{-}}$ times-repeated object. Let us define such a function, called labeling function, as

$$
\begin{array}{cccc}
L: & \Upsilon_{N_{G}} \times\left\{1, \cdots, N_{G}\right\} & \rightarrow & \left\{1, \cdots, N_{G}\right\}  \tag{2}\\
(G, i) & \mapsto & c
\end{array}
$$

where the $i$-th node of $G$ is assigned to the object instance $c$. We say that $G$ is unconsensual if and only if $\exists i, j, i \neq j$, such that $L(G, i)=L(G, j)$. We also define the consensual score as:

$$
\begin{equation*}
L_{\text {score }}(L):=\sum_{G \in \Upsilon_{N_{G}}} \prod_{c=1}^{N_{G}}|\{i: L(G, i)=c\}| \tag{3}
\end{equation*}
$$

where the inner expression in the sum equals zero if two or more nodes of a graph are assigned to the same cluster by $L$; and one otherwise.

In practice, the labeling function can be provided by one of the following clustering algorithms: Spectral Clustering on KNN graphs, Gaussian Mixture Models clustering, or even the simple KNN clustering algorithm. We argue that the set of all nodes is not difficult to cluster, as our graphs are based on information provided by a distinctive matching method. Indeed, a simple KNN algorithm already provides fair enough results, see Figure 3.

Algorithm 2 presents how a generic clustering method is used to create labeling functions $L$ in the presence of graph outliers. In the following we describe this process for a simple clustering based on Voronoi regions. First, we select one graph $G$ at random from $\Upsilon_{N_{G}}$ and propose it as center of each $N_{G}$ cluster. We then count the number of graphs from $\Upsilon_{N_{G}}$ whose nodes belong to one and only one Voronoi region. Finally, we iterate this process and keep the clusters (with centers at $G$ ) providing the biggest number of graphs in consensus with them. As depicted in Figure 3, each node from each graph is then labeled with a reference to the cluster to which it belongs. As one can imagine, these labels might not yet represent repetitions of objects due to spurious graphs. They will be filtered out in the next section.

For simplicity, we assume that all unconsensual graphs (0labeled by Algorithm 2) are removed from $\Upsilon_{N_{G}}$ and that all nodes are rearranged so as to

$$
L(G, i)=i, \quad \forall i \in\left\{1, \cdots, N_{G}\right\}
$$

Algorithm 2 Generic pre-labeling of nodes
input: $\Upsilon_{N_{G}}$ - All graphs of matches having $N_{G}$ nodes; $N_{\text {iters }}-$
Number of iterations; FINDClusters - Clustering method.
start:
for $i \in\left\{1, \cdots, N_{\text {iters }}\right\}$ do
Select $G$ at random from $\Upsilon_{N_{G}}$.
Define $L_{i}$ with the help of FindClusters, initialized at $G$.
$i^{\star}=\arg \max _{i} L_{\text {score }}\left(L_{i}\right)$.
Set $L_{i^{\star}}$ labels of non consensual graphs to 0 .
$\triangleright$ The 0 label acts as an unconsensual flag.
return: $L_{i \star}$.


Fig. 3: Pre-labeling graphs nodes for $\Upsilon_{6}$. Labeling provided by spectral clustering, gaussian mixture models and KNN. Random colors indicates appartenance to clusters. Black nodes belong to a non consensual graph, as there are two nodes belonging to the same cluster (the top right cluster).

### 4.2. Geometry consistency

The standard RANSAC (RAndom SAmple Consensus) algorithm [23] computes the parameters fitting a mathematical model from observed data in the presence of outliers. Numerous improvements have been proposed in the literature for RANSAC, see [24-28], but the core idea behind them remains.

In the case of homography estimation, the classic RANSAC algorithm returns the homography $\eta_{j}$ computed in iteration $j$ having the largest consensus of inliers among all iterations. The $j$-iteration can be briefly described in two steps:

1. (Model fitting) Randomly select $s$ matches $\left(x_{i} \leftrightarrow y_{i}\right)_{i=1, \ldots, s}$ from the set of all matches ( $M_{T}$ ) and compute the homography $\eta_{j}$ that yields the best fit.
2. (Model consensus) Count the number of matches from $M_{T}$ that are within a distance threshold of $\sigma$ (i.e. counting inliers).
We propose to detect a repeated object from pre-labeled graphs by finding graphs in consensus with respect to multiple homographies. These homographies describe the transformations of a query object to all other present instances of it. The labeling function $L$ in Equation 2 assigns to each node (i.e. atom) from a fixed graph, a label. These labels indicate an initial proposition of object instances, but without geometric validation. We describe in Algorithm 3 how to identify homography consistent copies of an object from graphs in $\Upsilon_{N_{G}}$ and a pre-labeling function. For simplicity, we select the query object to be among nodes assigned to the first label $(c=1)$ and all other labels relate to different instances of this object. The homography fitting function HFITTING $N_{s}$ of Algorithm 3 can be based on raw matches $\left(N_{s}=4\right)$ or on the affine information $\left(N_{s}=2\right)$. Notice that $\alpha_{\max }=(\infty, \infty, \infty, \infty)$ implies the classic score for homography consensus in Equation 4.
```
Algorithm 3 REP-RANSAC
input: \(\Upsilon_{N_{G}}\) - All graphs of matches having \(N_{G}\) nodes;
HFITTING \(N_{s}\) - a homography model estimator from \(N_{s}\) matches;
\(N_{\text {iters }}\) - Number of iterations; \(\alpha_{\max }\) - vector of maximal affine dis-
cordance; \(\sigma\) - maximal distance error threshold; \(N_{s}\) - Number of
samples needed for Homography fitting.
start:
for \(i \in\left\{1, \cdots, N_{\text {iters }}\right\}\) do
    Randomly select \(G_{1}, \cdots, G_{N_{s}}\) from \(\Upsilon_{N_{G}}\)
    for \(c=2, \cdots, N_{G}\) do
        \(\eta_{c}^{i}:=\operatorname{HFITTING}_{N_{s}}\left(\mathbf{x}_{G_{1}}^{1} \leftrightarrow \mathbf{x}_{G_{1}}^{c}, \cdots, \mathbf{x}_{G_{N_{s}}}^{1} \leftrightarrow \mathbf{x}_{G_{N_{s}}}^{c}\right)\)
        \(\mathbb{A}_{c}:=\mathbb{1}_{\left\{\alpha\left(A_{E}\left(\mathbf{x}_{G}^{1} \leftrightarrow \mathbf{x}_{G}^{c}\right), A_{\eta_{c}^{i}}\left(\mathbf{x}_{G}^{1}\right)\right)<\alpha_{\max }\right\}}\)
                            \(\triangleright\) where \(\mathbf{x}_{G}^{c}\) correspond to
``` the \(G\) node pre-labeled to cluster \(c, \alpha\) is the affine discordance measure appearing in [28], \(\alpha\left(A_{E}\left(\mathbf{x}_{G}^{1} \leftrightarrow \mathbf{x}_{G}^{c}\right)\right.\) are the estimated affine maps from \(\mathbf{x}_{G}^{1}\) to \(\mathbf{x}_{G}^{c}\), and \(A_{\eta_{c}^{i}}\left(\mathbf{x}_{G}^{1}\right)\) are the first order Taylor approximations of \(\eta_{c}^{i}\) around \(\mathbf{x}_{G}^{1}\).
\[
\begin{equation*}
S_{i}\left(\left\{\eta_{c}^{i}\right\}_{c}\right)=\sum_{G \in \Upsilon_{N_{G}}} \prod_{c=2}^{N_{G}} \mathbb{1}_{\left\{\left\|\eta_{c}^{i}\left(\mathbf{x}_{G}^{1}\right)-\mathbf{x}_{G}^{c}\right\|_{L^{2}} \leq \sigma\right\}} \times \mathbb{A}_{c} \tag{4}
\end{equation*}
\]
\(i^{\star}=\arg \max _{i} S_{i}\)
return: \(\left\{\eta_{c}^{i^{\star}}\right\}_{c \in\left\{2, \cdots, N_{G}\right\}}\)

\subsection*{4.3. Detecting meaningful objects across cardinalities}

Up until now and for each set \(\Upsilon_{N_{G}}\), a most significant repeated object is proposed and validated. In this section, we order these propositions by their Number of False Alarms (NFA), based on the a-contrario procedure [22]. Notice that we do not use the NFA to automatically set up parameters as done in [24]; instead we use it to measure meaningfulness across cardinalities.

For each possible cardinality \(N_{G}\), we define the associated NFA of its most meaningful object (computed in Algorithm 3) as,
\[
\begin{equation*}
N F A\left(\left\{\eta_{c}\right\}_{c}\right)=\left(n-N_{s}\right)\binom{n}{k}\binom{k}{N_{s}} P(k, n) \tag{5}
\end{equation*}
\]
where \(n=\left|\Upsilon_{N_{G}}\right|, k=S\left(\left\{\eta_{c}\right\}_{c}\right)\) (the function \(S\) appears in Equation 4) and the probability of a random graph to be an inlier of the model is expressed as
\[
P(k, n)=\mathbb{P}\left(\left\{\|\mathbf{X}\|_{L^{2}} \leq \sigma\right\} \bigcap\left\{\mathbf{A}<\alpha_{\max }\right\}\right)^{\left(N_{G}-1\right)\left(k-N_{s}\right)}
\]

Indeed, if there are \(k\) graphs inliers, potentially all \(N_{s}\) out of them yield the correct configuration. Each one has a probability
\[
\mathbb{P}\left(\left\{\|\mathbf{X}\|_{L^{2}} \leq \sigma\right\} \bigcap\left\{\mathbf{A}<\alpha_{\max }\right\}\right)^{\left(N_{G}-1\right)}
\]
of being inlier according to the background model, see Equation 4. Through the assumption of independence among graphs, the probability that all \(k\) of them are inliers is the above quantity powered by \(k-N_{s}\), since \(N_{s}\) ones were used for estimation are automatically inliers (they have error 0 ). The number \(k\) of inliers is usually not known in advance, so all values of \(k\) are tested (from \(N_{s}+1\) to \(n\) ), which explains the factor \(\left(n-N_{s}\right)\) in Equation 5.

Equation 5 allow us to assign a natural order between all detected objects by Algorithm 3. Figure 4 shows these ordered detections across cardinalities. Clearly, a strong NFA (i.e. \(\log\) NFA \(\ll 0\) )


Fig. 4: In descending order from left to right, the three most significant detections. Each instance of an object is randomly colored.


Fig. 5: Most significant repeated objects with weak NFAs.
supports the evidence of a repeated object. On the other hand, weak NFA signatures (i.e. \(\log\) NFA \(\approx 0\) ) shows poor evidence of repeated objects. If for a fixed image all \(\log\) NFA \(>0\), then a no detection is reported. Examples of most significant repeated objects with weak NFAs can be found in Figure 5.

\section*{5. CONCLUSIONS}

We have presented a methodology to identify object repetitions based on graphs derived from multiple SIFT matches. The SIFT method can be replaced by any other local matching method. The underlying idea is simple: if several local descriptions of a repeated object stand out (for the similarity score), then a coherence between their associated graphs points out repetitions of a same object. The pipeline for detecting object repetitions starts by analysing each detected similarity graph, associated with the repetition of a local atom. Then, an initial proposition of the most significant repeated object is deduced from graphs of equal cardinality, and its geometric inconsistencies filtered out to provide a final proposition of "repeated object". The validity of the object's repetitions is determined by computing an NFA. Clearly, in order to keep detecting more repeated objects, a simple workaround would consist in removing all geometrically consistent graphs and iterate the pipeline over the remaining graphs till no object repetition is detected. This will be explored in the future.

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