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► To cite this version:

Fedor Fomin, Daniel Lokshtanov, Saket Saurabh, Dimitrios M. Thilikos. Bidimensionality and Kernels. SIAM Journal on Computing, 2020, 49 (6), pp.1397-1422. 10.1137/16M1080264 . hal-03094544

HAL Id: hal-03094544

<https://hal.science/hal-03094544>

Submitted on 4 Jan 2021

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Bidimensionality and Kernels^{1,2}

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Abstract

Bidimensionality Theory was introduced by [E. D. DEMAINE, F. V. FOMIN, M. HAJIAGHAYI, AND D. M. THILIKOS. *Subexponential parameterized algorithms on graphs of bounded genus and H -minor-free graphs*, J. ACM, 52 (2005), pp. 866–893] as a tool to obtain sub-exponential time parameterized algorithms on H -minor-free graphs. In [E. D. DEMAINE AND M. HAJIAGHAYI, *Bidimensionality: new connections between FPT algorithms and PTASs*, in Proceedings of the 16th Annual ACM-SIAM Symposium on Discrete Algorithms (SODA), SIAM, 2005, pp. 590–601] this theory was extended in order to obtain polynomial time approximation schemes (PTASs) for bidimensional problems. In this work, we establish a third meta-algorithmic direction for bidimensionality theory by relating it to the existence of linear kernels for parameterized problems. In particular, we prove that every minor (respectively contraction) bidimensional problem that satisfies a separation property and is expressible in Countable Monadic Second Order Logic (CMSO), admits a linear kernel for classes of graphs that exclude a fixed graph (respectively an apex graph) H as a minor. Our results imply that a multitude of bidimensional problems admit linear kernels on the corresponding graph classes. For most of these problems no polynomial kernels on H -minor-free graphs were known prior to our work.

Keywords: Kernelization, Parameterized algorithms, Treewidth, Bidimensionality

¹Part of the results of this paper have appeared in [42].

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⁴Supported by the European Research Council under the European Union’s Seventh Framework Programme (FP/2007-2013) / ERC Grant Agreements no. 267959 and by the Norwegian Research Council via MULTIVAL project.

⁵Supported by the Research Council of Norway and the French Ministry of Europe and Foreign Affairs, via the Franco-Norwegian project PHC AURORA 2019.

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⁸Supported by the European Research Council under the European Union’s Seventh Framework Programme (FP/2007-2013) / ERC Grant Agreements no. 306992.

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¹⁰Supported by projects DEMOGRAPH (ANR-16-CE40-0028) and ESIGMA (ANR-17-CE23-0010).

1 Introduction

Bidimensionality theory was introduced by Demaine et al. in [26]. This theory is build on cornerstone theorems from Graph Minors Theory of Robertson and Seymour [64] and initially it was developed to unify and extend subexponential fixed-parameter algorithms for NP-hard graph problems to a broad range of graphs including planar graphs, map graphs, bounded-genus graphs and graphs excluding any fixed graph as a minor [24–26, 28, 29] (see also [7, 40, 41, 49] for other graph classes). Roughly speaking, the problem is bidimensional if the solution value for the problem on a $k \times k$ -grid is $\Omega(k^2)$, and contraction/removal of edges does not increase solution value. Many natural problems are bidimensional, including DOMINATING SET, FEEDBACK VERTEX SET, EDGE DOMINATING SET, VERTEX COVER, r -DOMINATING SET, CONNECTED DOMINATING SET, CYCLE PACKING, CONNECTED VERTEX COVER, GRAPH METRIC TSP, and many others.

The second application of bidimensionality was given by Demaine and Hajiaghayi in [27], where it has been shown that bidimensionality is a useful theory not only in the design of fast fixed-parameter algorithms but also in the design of fast PTASs. Demaine and Hajiaghayi established a link between parameterized and approximation algorithms by proving that every bidimensional problem satisfying some simple separation properties has a PTAS on planar graphs and other classes of sparse graphs. See also [40, 41] for further development of the applications of bidimensionality in the theory of EPTASs. We refer to the surveys [29, 31, 66] for further information on bidimensionality and its applications, as well as the book [22].

In this work we give the third application of bidimensionality, that is *kernelization*. Kernelization can be seen as the strategy of analyzing preprocessing or data reduction heuristics from a parameterized complexity perspective. Parameterized complexity introduced by Downey and Fellows is basically a two-dimensional generalization of “P vs. NP” where, in addition to the overall input size n , one studies the effects on computational complexity of a secondary measurement that captures additional relevant information. This additional information can be the solution size or the quantification of some structural restriction on the input, such as the treewidth or the genus of the input graph. The secondary information is quantified by a positive integer k and is called the *parameter*. Parameterization can be deployed in many different ways; for general background on the theory see [22, 33, 36, 62].

A parameterized problem with a parameter k is said to admit a *polynomial kernel* if there is a polynomial time algorithm (the degree of polynomial is independent of k), called a *kernelization* algorithm, that reduces the input instance down to an instance with size bounded by a polynomial $p(k)$ in k , while preserving the answer. Kernelization has been extensively studied in parameterized complexity, resulting in polynomial kernels for a variety of problems. Notable examples of known kernels are a $2k$ -sized vertex kernel for VERTEX COVER [17], a $355k$ vertex-kernel for DOMINATING SET on planar graphs [2], which later was improved to a $67k$ vertex-kernel [16], or an $\mathcal{O}(k^2)$ kernel for FEEDBACK VERTEX SET [67] parameterized by the solution size. One of the most intensively studied directions in kernelization is the study of problems on planar graphs and other classes of sparse graphs. This study was initiated by Alber et al. [3] who gave the first linear-sized kernel for

the DOMINATING SET problem on planar graphs. The work of Alber et al. [3] triggered an explosion of papers on kernelization, and kernels of linear sizes were obtained for a variety of parameterized problems on planar graphs including CONNECTED VERTEX COVER, MINIMUM EDGE DOMINATING SET, MAXIMUM TRIANGLE PACKING, EFFICIENT EDGE DOMINATING SET, INDUCED MATCHING, FULL-DEGREE SPANNING TREE, FEEDBACK VERTEX SET, CYCLE PACKING, and CONNECTED DOMINATING SET [3, 11, 12, 16, 51, 53, 59, 61]. We refer to the surveys [50, 56, 60] as well as the recent textbook [22] for a detailed treatment of the area of kernelization. Since most of the problems known to have polynomial kernels on planar graphs are bidimensional, the existence of links between bidimensionality and kernelization was conjectured and left as an open problem in [26].

In this work we show that every bidimensional problem with a simple separation property, which is a weaker property than the one required in the framework of Demaine and Hajiaghayi for PTASs [27] and which is expressible in the language of Counting Monadic Second Order Logic (CMSO) (we postpone these definitions till the next section) has a linear kernel on planar and even much more general classes of graphs. In this paper all the problems are parameterized by the solution size. Our main result is the following meta-algorithmic result.

Theorem 1.1. *Every CMSO-definable linear-separable minor-bidimensional problem Π admits a linear kernel on graphs excluding some fixed graph as a minor. Every CMSO-definable linear-separable contraction-bidimensional problem Π admits a linear kernel on graphs excluding some fixed apex graph as a minor.*

Theorem 1.1 implies the existence of linear kernels for many parameterized problems on apex-minor-free or minor-free graphs. For example, it implies that TREEWIDTH- η -MODULATOR, which is to decide whether an input graph can be turned into a graph of treewidth at most η , by removing at most k vertices, admits a kernel of size $\mathcal{O}(k)$ in H -minor-free graphs. Other applications of this theorem are linear kernels for r -DOMINATING SET, CONNECTED DOMINATING SET, CONNECTED VERTEX COVER, INDEPENDENT SET, or r -SCATTERED SET on apex-minor-free graphs and for CYCLE PACKING and FEEDBACK VERTEX SET on minor-free graphs, as well as for many other packing and covering problems. For many of these problems these are the first polynomial kernels on such classes of graphs.

High level overview of the main proof ideas. Our approach is built on the work of Bodlaender et al. [10] who proved the first meta-theorems on kernelization. The results in [10] imply that every parameterized problem that has finite integer index and satisfies an additional surface-dependent property, called a quasi-coverability property, has a linear kernel on graphs of bounded genus.

The kernelization framework in [10] is based on the following idea. Suppose that every yes-instance of a given parameterized problem admits a protrusion decomposition. In other words, suppose that the vertex set of an input graph G can be partitioned in sets R_0, R_1, \dots, R_ℓ , where $|R_0|$ and ℓ are of size linear in parameter k , R_0 separates R_i and R_j , for $1 \leq i < j \leq \ell$, and every set R_i , $i \in \{1, \dots, \ell\}$, is a protrusion, i.e. induces a graph of constant treewidth with a constant

number of neighbours in R_0 . Then the kernelization algorithm uses only one reduction rule, which is based on finite integer index properties of the problem in question, and replaces a protrusion with a protrusion of constant size.

In this paper we use exactly the same reduction rule as the one used in the kernelization algorithm given in [10] for obtaining kernels on planar graphs and graphs of bounded genus. The novel technical contribution of this paper is twofold. First, we introduce a new way the kernel sizes for bidimensional problems are analyzed. The analysis of kernel sizes in [10] requires “topological” decompositions of the given graph, in the sense that the partitioning of the graph into regions with small border, or protrusions, strongly depends on the embedding of the graph into a surface. Then topological properties of the embedding are used to prove the existence of a protrusion decomposition. While such an approach works well when we have a topological embedding it seems difficult to extend it to graphs excluding some fixed graph as a minor. Instead of taking the topological approach, we apply bidimensionality and suitable variants of the Excluded Grid Theorem [30, 37]. Roughly speaking, we show that bidimensionality and separability implies the existence of a protrusion decomposition. This makes our arguments not only much more general but also considerably simpler than the analysis in [10]. Our second technical contribution is the proof that every CMSO-definable separable problem has a finite integer index. Pipelined with the framework from [10], these results imply the proof of Theorem 1.1.

The remaining part of the paper is organized as follows. In Section 2, we provide definitions and notations used in the remaining part of the paper. Section 3 is devoted to combinatorial lemmata on separation properties of bidimensional problems. Based on these combinatorial lemmata, we prove a novel decomposition theorem (Theorem 3.11), which is the first main technical contribution of this work. In Section 4, we prove the second main technical contribution of the paper (Theorem 4.4) about the finite integer index of separable CMSO-optimization problems. In Section 5, we prove the main result of the paper about linear kernels. In the concluding Section 6 we discuss the connection between the separable-bidimensional property of a problem and the quasi-coverable property, which was used in the meta-theorem from [10].

Relevant results Let us provide a brief overview of the relevant results appeared since 2010, when the conference version of this paper was published. The properties of SQGM and SQGC graph classes were used to design approximation, FPT, counting, and kernelization algorithms on various graph classes in [7, 41, 41, 49, 55]. The issues of constructiveness of the kernelization algorithms provided in this paper is discussed by Garnero et al. in [45, 46]. Extension of some of our results to graphs excluding a topological minor is given by Kim et al. [54] (see also [55] for recent applications to counting problems). For DOMINATING SET or CONNECTED DOMINATING SET, linear kernels obtained in this paper for apex-minor free were extended to much more general classes: H -topological-minor-free graphs [43]. For DOMINATING SET linear kernels were obtained even for more general classes of graphs like graphs of bounded expansion [34], see also [35] for even more general results. Finally, see [48] for kernelization results when excluding graphs under other partial ordering relations, different than minors.

2 Preliminaries

In this section we give various definitions used in the paper. We use \mathbb{N} to denote the set of all non-negative integers and \mathbb{Z} to denote the set of all integers.

Concepts from Graph Theory Let G be a graph. We use the notation $V(G)$ and $E(G)$ for the vertex set and the edge set of G respectively. We say that a graph H is a *subgraph* of G if $V(H) \subseteq V(G)$ and $E(H) \subseteq E(G)$. Given a set $S \subseteq V(G)$ we denote by $G[S]$ the subgraph G' of G where $V(G') = S$ and $E(G') = \{xy \in E(G) \mid \{x, y\} \subseteq S\}$ and we call G' the *subgraph of G induced by S* or we simply say that G' is an induced subgraph of G .

For every $S \subseteq V$, we denote by $G - S$ the graph obtained from G by removing the vertices of S , i.e. $G - S = G[V(G) \setminus S]$. For vertex $v \in V(G)$ we also use $G - v$ for $G - \{v\}$. For a set $S \subseteq V(G)$, we define $N_G(S)$ to be the *open neighborhood* of S in G , which is the set of vertices from $V(G) \setminus S$ adjacent to vertices of S . The *closed neighborhood* of S is $N_G[S] := N(S) \cup S$. Given a set $S \subseteq V(G)$, we denote by $\partial_G(S)$ the set of all vertices in S that are adjacent in G with vertices not in S . Thus $N_G(S) = \partial_G(V(G) \setminus S)$.

Treewidth. A *tree decomposition* of a graph G is a pair $\mathcal{T} = (T, \{X_i\}_{i \in V(T)})$, where T is a tree whose every node i is assigned a vertex subset $X_i \subseteq V(G)$, called a *bag*, such that the following three conditions hold:

- (T1) $\bigcup_{i \in V(T)} X_i = V(G)$. In other words, every vertex of G is in at least one bag.
- (T2) For every $uv \in E(G)$, there exists a node i of T such that bag X_i contains both u and v .
- (T3) For every $u \in V(G)$, the set $T_u = \{i \in V(T) \mid u \in X_i\}$ (i.e., the set of nodes whose corresponding bags contain u) induces a subtree of T .

The *width* of a tree decomposition $\mathcal{T} = (T, \{X_i\}_{i \in V(T)})$ equals $\max_{i \in V(T)} |X_i| - 1$, that is, the maximum bag size, minus one. The *treewidth* of a graph G , denoted by $\text{tw}(G)$, is the minimum possible width of a tree decomposition of G . To distinguish between the vertices of the decomposition tree T and the vertices of the graph G , we will refer to the vertices of T as *nodes*. Treewidth is can be seen as a measure of the topological resemblance of a graph to the structure of a tree. It has been used by Robertson and Seymour, in [63], as a cornerstone parameter of Graph Minors and its trace as a parameter goes back to the early 70s [8, 47, 52].

Separators and separations Let G be a graph, $Q \subseteq V(G)$, and let $A_1, A_2 \subseteq V(G)$ such that $A_1 \cup A_2 = V(G)$. We say that the pair (A_1, A_2) is a *separation* of G if there is no edge with one endpoint in $A_1 \setminus A_2$ and the other in $A_2 \setminus A_1$. The *order* of a separation (A_1, A_2) is $|A_1 \cap A_2|$. For a vertex subset $Q \subseteq V(G)$ we say that a separation (A_1, A_2) is a *2/3-balanced separation* of (G, Q) if each of the parts $A_1 \setminus A_2$ and $A_2 \setminus A_1$ contains at most $\frac{2}{3}|Q|$ vertices of Q . Balanced separators have been extensively studied in the context of graph algorithms (see e.g. [4, 5, 57, 58]).

The following separation property of graphs of small treewidth is well known, see e.g. [9, 22].

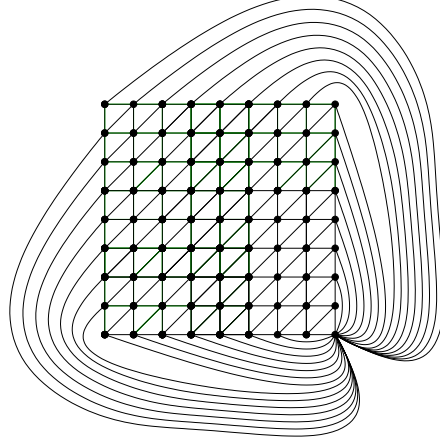


Figure 1: Graph Γ_9 .

Proposition 2.1. *Let G be a graph and let $S \subseteq V(G)$. There is a $2/3$ -balanced separation (A_1, A_2) of (G, S) of order at most $\text{tw}(G) + 1$.*

The following is an easy fact about treewidth, see e.g. [9].

Proposition 2.2. *The treewidth of a graph is the maximum treewidth of its connected components.*

Minors and contractions Given an edge $e = xy$ of a graph G , the graph G/e is obtained from G by contracting the edge e , that is, the endpoints x and y are replaced by a new vertex $v_{x,y}$ which is adjacent to the old neighbors of x and y (except from x and y). A graph H obtained by a sequence of edge-contractions is said to be a *contraction* of G . We denote it by $H \leq_c G$. A graph H is a *minor* of a graph G if H is the contraction of some subgraph of G and we denote it by $H \leq_m G$. We say that a graph G is *H -minor-free* when it does not contain H as a minor. We also say that a graph class \mathcal{G} is *H -minor-free* (or, excludes H as a minor) when all its members are H -minor-free. A graph G is an *apex graph* if there exists a vertex v such that $G - v$ is planar. A graph class \mathcal{G} is *apex-minor-free* if there exists an apex graph H such that \mathcal{G} is H -minor free.

A graph class \mathcal{G} is said to be *subgraph-closed* (*minor-closed*/*contraction-closed*) if every subgraph (minor/contraction) of a graph in \mathcal{G} also belongs to \mathcal{G} .

Grids and triangulated grids. Given a $k \in \mathbb{N}$, we denote by \boxplus_k the $(k \times k)$ -grid that is the graph with vertex set $\{(x, y) \mid x, y \in \{1, \dots, k\}\}$ and where two different vertices (x, y) and (x', y') are adjacent if and only if $|x - x'| + |y - y'| = 1$. Notice that \boxplus_k has exactly k^2 vertices.

For $k \in \mathbb{N}$, the graph Γ_k is obtained from the grid \boxplus_k by adding, for all $1 \leq x, y \leq k - 1$, the edge with endpoints $(x+1, y)$ and $(x, y+1)$ and additionally making vertex (k, k) adjacent to all the other vertices (x, y) with $x \in \{1, k\}$ or $y \in \{1, k\}$, i.e., to the whole perimetric border of \boxplus_k . Graph Γ_9 is shown in Fig. 1. The graph Γ_k has been defined in [37] in the context of bidimensionality theory.

We also need the following result of Robertson and Seymour [63].

Proposition 2.3. *For every $k \geq 0$, $\text{tw}(\boxplus_k) = k$.*

Parameterized graph problems. In general, a parameterized graph problem Π can be seen as a subset of $\Sigma^* \times \mathbb{N}$ where, in each instance (x, k) of Π , x encodes a graph and k is the parameter. In this paper we use an extension of this definition used in [10] that permits the parameter k to be negative with the additional constraint that either all pairs with non-positive value of the parameter are in Π or that no such pair is in Π . Formally, a parameterized problem Π is a subset of $\Sigma^* \times \mathbb{Z}$ where for all $(x_1, k_1), (x_2, k_2) \in \Sigma^* \times \mathbb{Z}$ with $k_1, k_2 < 0$ it holds that $(x_1, k_1) \in \Pi$ if and only if $(x_2, k_2) \in \Pi$. This extended definition encompasses the traditional one and is being adopted for technical reasons (see Subsection 4.3). In an instance of a parameterized problem (x, k) , the integer k is called the parameter.

The notion of *kernelization* is due to Downey and Fellows [32]. Kernelization is formally defined as follows.

Definition 2.4 (Kernelization). A *kernelization algorithm*, or simply a *kernel*, for a parameterized problem Π is an algorithm \mathcal{A} that, given an instance (x, k) of Π , works in polynomial-time and returns an equivalent instance (x', k') of Π . Moreover, there exists a computable function $g(\cdot)$ such that whenever (x', k') is the output for an instance (x, k) , then it holds that $|x'| + k' \leq g(k)$. If the upper bound $g(\cdot)$ is a polynomial (linear) function of the parameter, then we say that Π admits a *polynomial (linear) kernel*.

We often abuse the notation and call the output of a kernelization algorithm, the “reduced” equivalent instance, also a kernel.

Bidimensionality Bidimensionality theory was introduced by Demaine et al. in [26]. Here we closely follow the presentation of bidimensionality from the book of Cygan et al. [22].

We will restrict our attention to vertex or edge subset problems. A *vertex subset problem* Π is a parameterized problem where input is a graph G and an integer k , the parameter is k . An instance (G, k) is a yes-instance if and only if there exists a set $S \subseteq V(G)$ such that $|S| \leq k$ for minimization problems (or $|S| \geq k$ for maximization problems) so that a predicate $\phi(G, S)$ is true. Here ϕ can be any computable function which takes as input a graph G and set $S \subseteq V(G)$ and outputs **true** or **false**. In the rest of this paper, we always assume that a problem Π is generated by the choice of such a predicate ϕ and the choice of whether it is a minimization or a maximization problem. We will refer to ϕ as the *feasibility function* of Π .

The interpretation is that ϕ defines the space of *feasible solutions* S for a graph G by returning whether S is feasible for G . For an example, for the DOMINATING SET problem we have that $\phi(G, S) = \mathbf{true}$ if and only if $N[S] = V(G)$. *Edge subset problems* are defined similarly, with the only difference being that S is a subset of $E(G)$ rather than $V(G)$.

Let us remark that there are many vertex/edge subset problems which, at a first glance, do not look as if they could be captured by this definition. An example is the CYCLE PACKING problem.

Here the input is a graph G and integer k , and the task is to determine whether G contains k pairwise vertex-disjoint cycles C_1, C_2, \dots, C_k . This is a vertex subset problem because G has k vertex-disjoint cycles if and only if there exists a set $S \subseteq V(G)$ of size at least k and $\phi(G, S)$ is true, where $\phi(G, S)$ is defined as follows.

$\phi(G, S) = \exists$ subgraph G' of G such that

- each connected component of G' is a cycle,
- and each connected component of G' contains exactly one vertex of S .

This definition may seem a bit silly, since checking whether $\phi(G, S)$ is true for a given graph G and set S is NP-complete. In fact this problem is considered as a more difficult problem than CYCLE PACKING. Nevertheless, this definition shows that CYCLE PACKING is a vertex subset problem, which will allow us to give a linear kernel for CYCLE PACKING on minor free graphs (see also [46] for explicit bounds on linear kernels for packing problems).

For any vertex or edge subset minimization problem Π we have that $(G, k) \in \Pi$ implies that $(G, k') \in \Pi$ for all $k' \geq k$. Similarly, for a vertex or edge subset maximization problem we have that $(G, k) \in \Pi$ implies that $(G, k') \in \Pi$ for all $k' \leq k$. Thus the notion of “optimality” is well defined for vertex and edge subset problems.

Definition 2.5 (Optimum solution). For a vertex/edge subset minimization problem Π , we define

$$OPT_{\Pi}(G) = \min \{k \mid (G, k) \in \Pi\}.$$

If no k such that $(G, k) \in \Pi$ exists, $OPT_{\Pi}(G)$ returns $+\infty$. For a vertex/edge subset maximization problem Π ,

$$OPT_{\Pi}(G) = \max \{k \mid (G, k) \in \Pi\}.$$

If no k such that $(G, k) \in \Pi$ exists, $OPT_{\Pi}(G)$ returns $-\infty$. We say that a vertex (edge) set S is an *optimum solution* of Π for G if $\phi(G, S) = \mathbf{true}$ and $|S| = OPT_{\Pi}(G)$. We define $SOL_{\Pi}(G)$ to be a function that given as an input a graph G returns an optimum solution S and returns **null** if no such set S exists (if many optimum solutions exist, we pick one arbitrarily).

For many problems it holds that contracting an edge can not increase the size of the optimal solution. We will say that such problems are contraction closed. Formally we have the following definition.

Definition 2.6 (Contraction-closed problem). A vertex/edge subset problem Π is *contraction-closed* if for any G and $uv \in E(G)$, $OPT_{\Pi}(G/uv) \leq OPT_{\Pi}(G)$.

If contracting edges, deleting edges and deleting vertices can not increase the size of the optimal solution, we say that the problem is minor-closed.

Definition 2.7 (Minor-closed problem). A vertex/edge subset problem Π is *minor-closed* if for any G , edge $uv \in E(G)$ and vertex $w \in V(G)$, $OPT_{\Pi}(G/uv) \leq OPT_{\Pi}(G)$, $OPT_{\Pi}(G \setminus uv) \leq OPT_{\Pi}(G)$, and $OPT_{\Pi}(G - w) \leq OPT_{\Pi}(G)$.

The following slight modification of OPT_{Π} makes possible to avoid dealing with asymptotic inequalities in the rest of the paper.

Definition 2.8 (Modified OPT). We also define $OPT_{\Pi}^*(G) = \max\{OPT_{\Pi}(G), 1\}$.

We are now ready to give the definition of bidimensional problems.

Definition 2.9 (Bidimensional problem). A vertex/edge subset problem Π is

- *contraction-bidimensional*: if it is contraction-closed and there exists a positive real constant β such that for every $k \in \mathbb{N}$, $OPT_{\Pi}^*(\Gamma_k) \geq \beta \cdot k^2$.
- *minor-bidimensional*: if it is minor-closed and there exists a positive real constant β such that for every $k \in \mathbb{N}$, $OPT_{\Pi}^*(\boxplus_k) \geq \beta \cdot k^2$.

It is usually quite easy to determine whether a problem is contraction (or minor)-bidimensional. Take for an example INDEPENDENT SET. Contracting an edge may never increase the size of the maximum independent set, so the problem is contraction closed. Furthermore, Γ_k is a planar graph and because of the Four Color theorem, it contains an independent set of size at least $\lfloor \frac{|V(\Gamma_k)|}{4} \rfloor$. Thus INDEPENDENT SET is contraction-bidimensional. On the other hand, deleting edges may increase the size of a maximum size independent set in G . Thus INDEPENDENT SET is not minor-bidimensional.

We would also like to comment why we define bidimensionality by making use of OPT_{Π}^* and not OPT_{Π} . The reason is that for most interesting problems like VERTEX COVER or FEEDBACK VERTEX SET, the values of $OPT_{\Pi}(\boxplus_k)$ on very small grids (like \boxplus_1 for VERTEX COVER and \boxplus_2 for FEEDBACK VERTEX SET) is zero. The definition of OPT_{Π}^* takes care of such degenerate situations.

An alternative way to define bidimensionality, by making use of OPT_{Π} , would be to ask for the existence of a positive β such that that $OPT_{\Pi}(\boxplus_k)$ (or $OPT_{\Pi}(\Gamma_k)$) is *asymptotically bigger* than $\beta \cdot k^2$, i.e., there exists an integer $k_0 > 0$ such that for every $k \geq k_0$, $OPT_{\Pi}(\boxplus_k) \geq \beta \cdot k^2$. This new definition is equivalent to the one of Definition 2.9. To see this, in the non-trivial direction, we set $k_1 = \max\{k_0, \sqrt{1/\beta}\}$ which implies that $\forall k \geq k_1$ $OPT_{\Pi}(\boxplus_k) = OPT_{\Pi}^*(\boxplus_k) \geq \beta \cdot k^2$. We now set $\beta' = \min\{\beta, 1/\sqrt{k_1}\}$ and observe that if $0 \leq k \leq k_1$, then $OPT_{\Pi}^*(\boxplus_k) \geq 1 \geq \beta' \cdot k^2$, while if $k \geq k_1$, then $OPT_{\Pi}^*(\boxplus_k) = OPT_{\Pi}(\boxplus_k) \geq \beta' \cdot k^2$, as required.

In this paper we adopted the definition of Bidimensionality that uses OPT_{Π}^* because this makes our proofs easier to present.

Counting Monadic Second Order Logic The syntax of Monadic Second Order Logic (MSO) of graphs includes the logical connectives $\vee, \wedge, \neg, \Leftrightarrow, \Rightarrow$, variables for vertices, edges, sets of vertices, and sets of edges, the quantifiers \forall, \exists that can be applied to these variables, and the following five binary relations:

1. $u \in U$ where u is a vertex variable and U is a vertex set variable;
2. $d \in D$ where d is an edge variable and D is an edge set variable;

3. **inc**(d, u), where d is an edge variable, u is a vertex variable, and the interpretation is that the edge d is incident with the vertex u ;
4. **adj**(u, v), where u and v are vertex variables and the interpretation is that u and v are adjacent;
5. equality of variables representing vertices, edges, sets of vertices, and sets of edges.

In addition to the usual features of monadic second-order logic, if we have atomic sentences testing whether the cardinality of a set is equal to q modulo r , where q and r are integers such that $0 \leq q < r$ and $r \geq 2$, then this extension of the MSO is called *counting monadic second-order logic*. Thus CMSO is MSO enriched with the following atomic sentence for a set S :

$$\mathbf{card}_{q,r}(S) = \mathbf{true} \text{ if and only if } |S| \equiv q \pmod{r}.$$

For a detailed introduction on CMSO, see [6, 20, 21].

We consider CMSO sentences evaluated either on graphs or on annotated graphs. In this paper by annotated graph we mean a pair (G, S) , where G is a graph and S is either a vertex or edge subset of G . We remark that a graph can be annotated with more than one set and all of the considerations in this paper could be generalized in this sense. For simplicity, we omit this in this paper. A class of graphs \mathcal{G} is *CMSO-definable* if there is a CMSO sentence on graphs ϕ such that $G \in \mathcal{G}$ if and only if $G \models \phi$. Similarly a predicate ϕ on annotated graphs is *CMSO-definable* if there is a CMSO sentence ε on annotated graphs such that $\phi(G, S) = \mathbf{true}$ if and only if $(G, S) \models \varepsilon$.

A vertex/edge subset minimization (or maximization) problem with feasibility function ϕ is a MIN-CMSO problem (or MAX-CMSO problem) if ϕ is CMSO-definable.

Subquadratic grid minor/contraction property. In general, it is known that there exists a constant c such that any graph G which excludes a \boxplus_k as a minor has treewidth at most $O(k^c)$. The exact value of c remains unknown, but it is more than 2 (see e.g., [65]) and at most 20 (see [15, 18, 19]). We will restrict our attention to graph classes on which $c < 2$. In particular, we say that a graph class \mathcal{G} has the *subquadratic grid minor property* (SQGM property, for short) if there exist constants $\lambda > 0$ and $1 \leq c < 2$ such that every graph $G \in \mathcal{G}$ which excludes \boxplus_k as a minor has treewidth at most λk^c .

Problems that are contraction-closed but not minor closed are considered on more restricted classes of graphs. We say that a graph class \mathcal{G} has the *subquadratic graph contraction property* (SQGC property for short) if there exist constants $\lambda > 0$ and $1 \leq c < 2$ such that any *connected* graph $G \in \mathcal{G}$ which excludes a Γ_k as a contraction has treewidth at most λk^c .

Observation 2.10. *Every graph class \mathcal{G} with the SQGC property has the SQGM property.*

Proof. Suppose that the graph class \mathcal{G} has the SQGC property. This means that there is a $c \in [1, 2)$ such that if G is a connected graph in \mathcal{G} that cannot be contracted to Γ_k , then $\text{tw}(G) \leq \lambda k^c$. Suppose that G is any graph in \mathcal{G} excluding \boxplus_k as a minor. Clearly all connected components of G exclude

\boxplus_k as a minor and, as \boxplus_k is a minor of Γ_k (in fact it is a subgraph), the connected components of G also exclude Γ_k as a contraction. This implies that all connected components of G have treewidth at most λk^c , therefore G has treewidth at most λk^c , as required. \square

The following proposition follows directly from the linearity of excluded grid-minor in H -minor-free graphs proven by Demaine and Hajiaghayi [30] and its analog for contraction-minors from [37].

Proposition 2.11. *For every graph H , H -minor-free graph class \mathcal{G} has the SQGM property with $c = 1$. If H is an apex graph, then \mathcal{G} has the SQGC property with $c = 1$.*

Problem restrictions. We say that a parameterized problem Π is a problem on the graph class \mathcal{G} if every yes-instance (G, k) of Π satisfies $G \in \mathcal{G}$. The *restriction* of a parameterized problem Π to a graph class \mathcal{G} , denoted by $\Pi \upharpoonright \mathcal{G}$ is defined as follows.

$$\Pi \upharpoonright \mathcal{G} = \{(G, k) \mid (G, k) \in \Pi \text{ and } G \in \mathcal{G}\}.$$

For a parameterized problem Π (on general graphs) we will refer to the restriction of Π to \mathcal{G} by “ Π on \mathcal{G} ”.

3 Decomposing into protrusions

In this section we give the main technical combinatorial contribution of this work establishing a protrusion decomposition theorem for linearly-separable bidimensional problems. The proof of the theorem is done in two steps. First we show that every graph class with the SQGM or the SQGC property admit a treewidth-modulator of size linear in the parameter of a linearly-separable (minor- or contraction-) bidimensional problem (Subsection 3.2). Then in Subsection 3.3, we show that graph classes with SQGM and SQGC properties and having linear treewidth-modulators, can be decomposed into protrusions.

3.1 Parameter-treewidth bounds

The following lemmata establish *parameter-treewidth bounds*, that is tight relationships between the size of the optimal solution and the treewidth of the input graph. This relationship was first observed by Demaine et al. [26]. The bound for contraction-bidimensional problems presented here is essentially identical to the one presented in Fomin et al. [37]. We re-prove the lemmata here because of slight differences in definitions.

Lemma 3.1. *For any minor-bidimensional problem Π on a graph class \mathcal{G} with the SQGM property, there exist constants $0 < \alpha$, $\frac{1}{2} \leq \mu < 1$ such that for any graph $G \in \mathcal{G}$, $\text{tw}(G) \leq \alpha \cdot (\text{OPT}_\Pi^*(G))^\mu$.*

Proof. Let $\lambda > 0$ and $1 \leq c < 2$ be the constants from the definition of the SQGM property, that is any graph $G \in \mathcal{G}$ which excludes a \boxplus_k as a minor has treewidth at most λk^c . Because of the minor-bidimensionality of Π , there is $\beta > 0$ such that $\text{OPT}_\Pi^*(\boxplus_k) \geq \beta k^2$ for every $k \in \mathbb{N}$. Consider now a graph $G \in \mathcal{G}$.

Let k be the maximum integer such that G contains \boxplus_k as a minor. This means that G excludes \boxplus_{k+1} as a minor, therefore $\text{tw}(G) \leq \lambda(k+1)^c$. Rearranging terms yields that $k \geq (\frac{\text{tw}(G)}{\lambda})^{1/c} - 1$. Since Π is minor-closed, it follows that $\text{OPT}_\Pi(G) \geq \text{OPT}_\Pi(\boxplus_k)$, therefore

$$\text{OPT}_\Pi^*(G) \geq \text{OPT}_\Pi^*(\boxplus_k) \geq \beta k^2 \geq \beta \left(\left(\frac{\text{tw}(G)}{\lambda} \right)^{1/c} - 1 \right)^2,$$

hence

$$\text{tw}(G) \leq \lambda \left(\left(\frac{\text{OPT}_\Pi^*(G)}{\beta} \right)^{\frac{1}{2}} + 1 \right)^c.$$

Recall that $\text{OPT}_\Pi^*(G) \geq 1$. We set $x = \left(\frac{\text{OPT}_\Pi^*(G)}{\beta} \right)^{\frac{1}{2}}$ and, as $x \geq \frac{1}{\sqrt{\beta}}$, we obtain that $x + 1 \leq (1 + \sqrt{\beta})x$. This implies that

$$\text{tw}(G) \leq \lambda (1 + \sqrt{\beta}) \left(\frac{\text{OPT}_\Pi^*(G)}{\beta} \right)^{\frac{1}{2}})^c = \lambda \cdot \left(1 + \frac{1}{\sqrt{\beta}} \right)^c \cdot (\text{OPT}_\Pi^*(G))^{\frac{c}{2}}.$$

We can now set $\alpha = \lambda \cdot \left(1 + \frac{1}{\sqrt{\beta}} \right)^c$, $\mu = c/2$ and observe that $\text{tw}(G) \leq \alpha \cdot (\text{OPT}_\Pi^*(G))^\mu$.

Since $c < 2$ in the definition of the SQGM property, we obtain that $\mu < 1$ and the statement of the lemma follows. \square

Lemma 3.2. *For any contraction-bidimensional problem Π on a graph class \mathcal{G} with the SQGC property, there exist constants $0 < \alpha$, $\frac{1}{2} \leq \mu < 1$, such that for any connected graph $G \in \mathcal{G}$, $\text{tw}(G) \leq \alpha \cdot (\text{OPT}_\Pi^*(G))^\mu$.*

The proof of the Lemma 3.2 is almost identical to the proof of Lemma 3.1. The differences between Lemmata 3.1 and 3.2 are as follows. Lemma 3.1 is for minor-closed problems, on graph classes \mathcal{G} with the SQGM property and works for *every* $G \in \mathcal{G}$. Lemma 3.2 is for contraction-closed problems, on graph classes \mathcal{G} with the SQGC property and works for *connected* graphs $G \in \mathcal{G}$. The connectivity requirement here is necessary: For SQGC property we require that there exist constants $\lambda > 0$ and $1 \leq c < 2$ such that any *connected* graph $G \in \mathcal{G}$ which excludes a Γ_k as a contraction has treewidth at most λk^c . Moreover, it is possible to provide an example of a contraction-bidimensional problem Π such that Lemma 3.2 does not hold for Π and disconnected graphs, see [22, Exercise 7.42]. However, as we will see soon, if in addition to contraction-bidimensionality the problem is separable, then the connectivity condition is not necessary anymore.

3.2 Separability and treewidth modulators

We now restrict our attention to problems Π that are somewhat well-behaved in the sense that whenever we have a small separator in the graph that splits the graph in two parts L and R , the intersection $|X \cap L|$ of L with any optimal solution X to the entire graph is a good estimate of $\text{OPT}_\Pi(G[L])$. This restriction allows us to prove decomposition theorems which are very useful for giving kernels. Similar decomposition theorems may also be used to give approximation schemes, see [27, 40].

Definition 3.3 (Separability). Let $f : \mathbb{N} \rightarrow \mathbb{N}$ be a function. We say that a vertex subset problem Π is *f-separable* if for every graph G , every optimum solution S of Π for G , and every subset $L \subseteq V(G)$

$$|S \cap L| - f(t) \leq \text{OPT}_\Pi(G[L]) \leq |S \cap L| + f(t)$$

where $t = |\partial_G(L)|$. In the case where Π is an edge subset problem the same definition as above applies with the difference that we agree to interpret $S \cap L$ by $S \cap E(G[L])$. The problem Π is called *separable* if there exists a function f such that Π is *f-separable*. Π is called *linear-separable* if function f is linear. That is, there exists a constant σ such that Π is $\sigma \cdot t$ -separable.

There are many problems that are contraction (or minor)-bidimensional, linear-separable including r -DOMINATING SET, CONNECTED DOMINATING SET, CONNECTED VERTEX COVER, VERTEX COVER, INDEPENDENT SET, FEEDBACK VERTEX SET, and r -SCATTERED SET, CYCLE PACKING as well as many packing and covering problems. Another important generic problem which is minor-bidimensional and linear-separable, is the TREEWIDTH- η -MODULATOR problem defined below. We refer to [27] and [41, Section 4] for definitions of these problems and the proofs that they are contraction-bidimensional and linear-separable.

Just as one example, let us consider CYCLE PACKING. It is clearly minor-bidimensional. It is easy to see that CYCLE PACKING is a minor-closed problem. Since grid \boxplus_k contains $\Omega(k^2)$ vertex-disjoint cycles, CYCLE PACKING is a minor-bidimensional problem. For linear-separability, we can view CYCLE PACKING as a problem of finding a maximum vertex set X such that there is a subgraph H of G , such that every connected component of H contains exactly one vertex of X and is a cycle. Observe that by deleting t vertices from graph G , we cannot hit more than t vertex-disjoint cycles from the solution. Hence the linear-separability of CYCLE PACKING follows.

A nice feature of separable contraction-bidimensional problems is that it is possible to extend Lemma 3.2 to disconnected graphs.

Lemma 3.4. *For any contraction-bidimensional separable problem Π on a graph class \mathcal{G} with the SQGC property, there exist constants $0 < \alpha$ and $\frac{1}{2} \leq \mu < 1$ such that for any graph $G \in \mathcal{G}$, $\text{tw}(G) \leq \alpha \cdot (\text{OPT}_\Pi^*(G))^\mu$.*

Proof. Let S be any optimal solution for Π . Since Π is separable there exists a function $f : \mathbb{N} \rightarrow \mathbb{N}$ such that for every graph G and every connected component C of G , it holds that

$$\text{OPT}_\Pi(G[C]) \leq |S \cap C| + f(|\partial_G(C)|) \leq |S| + f(0) = \text{OPT}_\Pi(G) + f(0).$$

We set $c = f(0)$. The above implies that $\text{OPT}_\Pi^*(G[C]) \leq \text{OPT}_\Pi^*(G) + c$. By Lemma 3.2, there exist constants $\alpha' > 0$ and $\frac{1}{2} \leq \mu < 1$ such that $\text{tw}(G[C]) \leq \alpha' \cdot (\text{OPT}_\Pi^*(C))^\mu$ for every connected component C . By Proposition 2.2, the treewidth of G is at most the maximum of the treewidth of its connected components. Since $\text{OPT}_\Pi^*(G) \geq 1$, we have that

$$\text{tw}(G) \leq \alpha' \cdot (\text{OPT}_\Pi^*(G) + c)^\mu \leq \alpha' \cdot ((c+1) \cdot \text{OPT}_\Pi^*(G))^\mu = \alpha' \cdot (c+1)^\mu \cdot (\text{OPT}_\Pi^*(G))^\mu.$$

The lemma follows if we set $\alpha = \alpha' \cdot (c+1)^\mu$. □

We will now consider a sequence of “canonical” bidimensional problems, with one problem for every $\eta \geq 0$. We say that a set $S \subseteq V(G)$ is a *treewidth- η -modulator* if $\text{tw}(G - S) \leq \eta$. We define the following problem.

TREEWIDTH- η -MODULATOR

Instance: A graph G , and integer $k \geq 0$.

Parameter: k .

Problem: Decide whether there exists a treewidth- η -modulator of G of size at most k .

Clearly, $\Pi = \text{TREEWIDTH-}\eta\text{-MODULATOR}$ is a minimization problem and it is easy to see that it is also minor-closed. By Proposition 2.3, every $(\eta + 1) \times (\eta + 1)$ subgrid of \boxplus_k , for $k \geq 1$, must contain at least one vertex of any solution, therefore

$$OPT_{\Pi}^*(\boxplus_k) \geq \lfloor \frac{k^2}{(\eta + 1)^2} \rfloor.$$

Thus the problem is minor-bidimensional.

To see that the problem is separable, consider any graph G , and vertex subset $L \subseteq V(G)$ with $\partial(L) = t$. For any treewidth- η -modulator S of G we have that $(S \cap L) \cup \partial(L)$ is a treewidth- η -modulator of $G[L]$. This proves $OPT_{\Pi}(G[L]) \leq |S \cap L| + t$. On the other hand, for any optimal treewidth- η -modulator S of G and treewidth- η -modulator S_L of $G[L]$, we have that $|S \cap L| \leq |S_L| + t$, since otherwise $(S \setminus L) \cup S_L \cup \partial(L)$ is a treewidth- η -modulator of G of size strictly smaller than $|S|$, contradicting the optimality of S . It follows that $|S \cap L| - t \leq OPT_{\Pi}(G[L])$. This shows that the TREEWIDTH- η -MODULATOR problem is minor-bidimensional and separable.

The fact that the TREEWIDTH- η -MODULATOR problem is minor-bidimensional, together with Lemma 3.1 yields the following observation.

Observation 3.5. *For every graph class \mathcal{G} with the SQGM property and every $\eta \geq 0$ there exist constants $\alpha > 0$ and $\frac{1}{2} \leq \mu < 1$ such that every graph $G \in \mathcal{G}$ with a non-empty treewidth- η -modulator S has treewidth at most $\alpha|S|^{\mu}$.*

Proof. Let $\Pi = \text{TREEWIDTH-}\eta\text{-MODULATOR}$, let $G \in \mathcal{G}$ and let S be a treewidth- η -modulator of G . As Π is minor-bidimensional and \mathcal{G} has the SQGM property, by Lemma 3.1, there exists $0 < \alpha$, $\frac{1}{2} \leq \mu < 1$ such that $G \in \mathcal{G}$, $\text{tw}(G) \leq \alpha \cdot (OPT_{\Pi}^*(G))^{\mu}$. If $OPT_{\Pi}(G) = 0$, then $OPT_{\Pi}^*(G) = 1$ which implies that $\text{tw}(G) \leq \alpha \leq \alpha|S|^{\mu}$ (recall that $|S| \geq 1$). If $OPT_{\Pi}(G) \geq 1$, then $OPT_{\Pi}(G) = OPT_{\Pi}^*(G)$, therefore $\text{tw}(G) \leq \alpha(OPT_{\Pi}(G))^{\mu} \leq \alpha|S|^{\mu}$. \square

Next we show that the TREEWIDTH- η -MODULATOR problems are canonical bidimensional problems in the following sense.

Lemma 3.6. *For any real $\varepsilon > 0$ and minor-bidimensional linear-separable problem Π on graph class \mathcal{G} with the SQGM property, there exists an integer $\eta \geq 0$ such that every $G \in \mathcal{G}$ has a treewidth- η -modulator S of size at most $\varepsilon \cdot OPT_{\Pi}(G)$.*

Proof. Let σ be a constant such that Π is $(\sigma \cdot t)$ -separable. Let $\alpha' > 0$ and $\frac{1}{2} \leq \mu < 1$ be the constants from Lemma 3.1. In particular $\text{tw}(G) \leq \alpha' \cdot (\text{OPT}_\Pi(G))^\mu$ if $\text{OPT}_\Pi(G) > 0$ and $\text{tw}(G) \leq \alpha'$ if $\text{OPT}_\Pi(G) = 0$. Set $\alpha = \max\{\alpha', 1\}$. Furthermore, if $\sigma < 1$ then Π is t -separable, and so we may assume without loss of generality that $\sigma \geq 1$.

We now define a few constants. The reason these constants are defined the way they are will become clear during the course of the proof. Finally we set η based on α, σ, μ and ε .

- Set $\rho = \frac{1^\mu + 2^\mu - 3^\mu}{3^\mu}$ and note that $\rho > 0$.
- set $\gamma = 4\alpha\sigma$,
- set $\delta = \frac{\gamma(2\varepsilon+1)}{\rho}$,
- set $k_0 = (3 + 3\gamma)^{\frac{1}{1-\mu}} + 3 \cdot (\frac{\delta}{\varepsilon})^{\frac{1}{1-\mu}}$, (notice that $k_0 > 3$) and
- set $\eta = \alpha \cdot k_0^\mu$.

We prove, by induction on k , the following stronger statement.

Claim: For any integer $k \geq \frac{1}{3}k_0$, every graph $G \in \mathcal{G}$ such that $\text{OPT}_\Pi(G) \leq k$ has a treewidth- η -modulator of size at most $\varepsilon k - \delta k^\mu$.

Proof of Claim: In the base case we consider any k such that $\frac{1}{3}k_0 \leq k \leq k_0$. By Lemma 3.1, any graph $G \in \mathcal{G}$ such that $\text{OPT}_\Pi(G) \leq k$ has treewidth at most $\alpha \cdot k^\mu \leq \eta$ (recall that $k \geq \frac{1}{3}k_0 \geq 1$). Thus G has a treewidth- η -modulator of size 0. Also by the definition of k_0 , we have that $k_0 \geq 3 \cdot (\frac{\delta}{\varepsilon})^{\frac{1}{1-\mu}} = (\frac{3\delta}{\varepsilon \cdot 3^\mu})^{\frac{1}{1-\mu}}$. Hence $k_0^{1-\mu} \geq (\frac{3\delta}{\varepsilon \cdot 3^\mu})$. This yields that $\frac{3^\mu}{\delta} k_0^{1-\mu} \geq \frac{3}{\varepsilon}$, therefore $\delta \frac{1}{3^\mu} \cdot k_0^{\mu-1} \leq \frac{\varepsilon}{3}$, and finally, $\delta \frac{1}{3^\mu} \cdot k_0^\mu \leq \frac{\varepsilon}{3} k_0$. Consequently,

$$0 \leq \frac{\varepsilon \cdot k_0}{3} - \delta \left(\frac{k_0}{3} \right)^\mu. \quad (1)$$

By (1), we have that the size of treewidth- η -modulator of G , which is 0, satisfies the following

$$0 \leq \varepsilon \frac{1}{3} k_0 - \delta \left(\frac{1}{3} k_0 \right)^\mu \leq \varepsilon k - \delta k^\mu.$$

In the last inequality we used the fact that for any $\frac{1}{2} \leq \mu < 1$, ε and δ , the function $\varepsilon k - \delta k^\mu$ is monotonically increasing from the first point where it becomes positive. This fact may easily be verified by differentiation. This completes the proof of the base case.

For the inductive step, let $k > k_0$ and suppose that the statement is true for all values below k . We prove the statement for k . Consider a graph $G \in \mathcal{G}$ such that $\text{OPT}_\Pi(G) \leq k$. By Lemma 3.1, the treewidth of G is at most $\text{tw}(G) \leq \alpha \cdot k^\mu$. By Proposition 2.1 applied to $(G, \text{SOL}_\Pi(G))$, there is a $2/3$ -balanced separation (A_1, A_2) of $(G, \text{SOL}_\Pi(G))$ of order at most $\text{tw}(G) + 1 \leq \alpha \cdot k^\mu + 1$. Let $L = A_1 \setminus A_2$, $S = A_1 \cap A_2$ and $R = A_2 \setminus A_1$. Note that there are no edges from L to R . Since (A_1, A_2) is a $2/3$ -balanced separation it follows that there exists a real $\frac{1}{3} \leq a \leq \frac{2}{3}$ such that $|L \cap \text{SOL}_\Pi(G)| \leq a|\text{SOL}_\Pi(G)|$ and $|R \cap \text{SOL}_\Pi(G)| \leq (1-a)|\text{SOL}_\Pi(G)|$.

Consider now the graph $G[L \cup S]$. Since L has no neighbors in R (in G) and Π is $(\sigma \cdot t)$ -separable, it follows that

$$\begin{aligned} OPT_{\Pi}(G[L \cup S]) &\leq |SOL_{\Pi}(G) \cap (L \cup S)| + \sigma|S| \\ &\leq ak + (\alpha k^{\mu} + 1) + \sigma(\alpha k^{\mu} + 1) \\ &\leq ak + (\alpha k^{\mu} + 1)(\sigma + 1) \leq ak + \gamma k^{\mu}. \end{aligned}$$

Here the last inequality follows from the assumption that $k \geq k_0 \geq 1$ and the choice of γ .

We claim that for $k \geq k_0$,

$$ak + \gamma k^{\mu} \leq k - 1. \quad (2)$$

By the choice of k_0 and the fact that $\mu > 0$, we have

$$\begin{aligned} k_0 &\geq (3 + 3\gamma)^{\frac{1}{1-\mu}} \Rightarrow k_0 \geq \left(\frac{3}{k_0^{\mu}} + 3\gamma\right)^{\frac{1}{1-\mu}} \Rightarrow \\ k_0^{1-\mu} &\geq \frac{3}{k_0^{\mu}} + 3\gamma \Rightarrow \frac{k_0}{3} \geq 1 + \gamma k_0^{\mu} \Rightarrow \\ &\frac{2}{3}k_0 + \gamma k_0^{\mu} \leq k_0 - 1 \end{aligned}$$

Because for $k \geq k_0$ the function $\frac{2}{3}k + \gamma k^{\mu} - k + 1$ decreases, which is easily verified by differentiation. Hence $\frac{2}{3}k + \gamma k^{\mu} \leq k - 1$ and because $a \leq \frac{2}{3}$, (2) follows.

Further $ak + \gamma k^{\mu} \geq \frac{1}{3}k_0$ since $a \geq \frac{1}{3}$. Thus we may apply the induction hypothesis to $G[L \cup S]$ and obtain a treewidth- η -modulator Z_L of $G[L \cup S]$, such that

$$\begin{aligned} |Z_L| &\leq \varepsilon(ak + \gamma k^{\mu}) - \delta(ak + \gamma k^{\mu})^{\mu} \\ &\leq \varepsilon(ak + \gamma k^{\mu}) - \delta k^{\mu} a^{\mu}. \end{aligned}$$

An identical argument applied to $G[R \cup S]$ yields a treewidth- η -modulator Z_R of $G[R \cup S]$, such that

$$|Z_R| \leq \varepsilon((1-a)k + \gamma k^{\mu}) - \delta k^{\mu}(1-a)^{\mu}.$$

We now make a treewidth- η -modulator Z of G as follows. Let $Z = Z_L \cup S \cup Z_R$. The set Z is a treewidth- η -modulator of G because every connected component of $G - Z$ is a subset of L or R , and Z_L and Z_R are treewidth- η -modulators for $G[L \cup S]$ and $G[R \cup S]$ respectively. Finally we bound the size of Z .

$$\begin{aligned} |Z| &\leq |Z_L| + |Z_R| + |S| \\ &\leq \varepsilon(ak + \gamma k^{\mu}) - \delta k^{\mu} a^{\mu} + \varepsilon((1-a)k + \gamma k^{\mu}) - \delta k^{\mu}(1-a)^{\mu} + \gamma k^{\mu} \\ &= \varepsilon k - \delta k^{\mu}((1-a)^{\mu} + a^{\mu}) + k^{\mu}\gamma(2\varepsilon + 1) \\ &\leq \varepsilon k - \delta k^{\mu} + k^{\mu}(\gamma(2\varepsilon + 1) - \delta\rho) \\ &= \varepsilon k - \delta k^{\mu} \end{aligned}$$

In the transition from the third to the fourth line we used that $(1 - a)^\mu + a^\mu - 1 \geq \rho$ for any a between $\frac{1}{3}$ and $\frac{2}{3}$. The claim follows.

To conclude the proof, we observe that the statement of the lemma follows from the above claim. If $OPT_\Pi(G) \leq k_0$ then $\text{tw}(G) \leq \alpha \cdot k_0^\mu = \eta$ and the empty set is a treewidth- η -modulator of G of size $0 \leq \varepsilon \cdot OPT_\Pi(G)$. If $OPT_\Pi(G) > k_0$ then G has a treewidth- η -modulator of size at most $\varepsilon \cdot OPT_\Pi(G) - \delta(OPT_\Pi(G))^\mu \leq \varepsilon \cdot OPT_\Pi(G)$. This completes the proof. \square

An identical argument yields an analogous lemma for contraction-bidimensional problems. The only difference is that we now use Lemma 3.4 instead of Lemma 3.1.

Lemma 3.7. *For any real $\varepsilon > 0$ and contraction-bidimensional linear-separable problem Π on graph class \mathcal{G} with the SQGC property, there exists an integer $\eta \geq 0$ such that any graph $G \in \mathcal{G}$ has a treewidth- η -modulator S of size at most $\varepsilon \cdot OPT_\Pi^*(G)$.*

3.3 Protrusion decomposition

The notions of *protrusion* and *protrusion decomposition* were introduced in [10].

Definition 3.8. [t -**protrusion**] For a graph G , a set $X \subseteq V(G)$ is a t -*protrusion* of G if $|\partial(X)| \leq t$ and $\text{tw}(G[X]) \leq t$.

Definition 3.9. [(α, r) -**Protrusion decomposition**] An (α, r) -*protrusion decomposition* of a graph G is a sequence $\mathcal{P} = \langle R_0, R_1, \dots, R_\ell \rangle$ of pairwise disjoint subsets of $V(G)$ such that

- $\bigcup_{i \in \{1, \dots, \ell\}} R_i = V(G)$,
- $\max\{\ell, |R_0|\} \leq \alpha$,
- each $R_i^+ = N_G[R_i]$, $i \in \{1, \dots, \ell\}$, is an r -protrusion of G , and
- for every $i \in \{1, \dots, \ell\}$, $N_G(R_i) \subseteq R_0$.

We call the sets R_i^+ , $i \in \{1, \dots, \ell\}$, the *protrusions* of \mathcal{P} and the set R_0 the *core* of \mathcal{P} .

Next we prove that the existence of a treewidth- η -modulator S implies the existence of a protrusion decomposition of G into r -protrusions with a core not much bigger than S , and r only depending on η (and the graph class \mathcal{G} we are working with). The inductive proof of the following lemma is similar to the proof of Lemma 3.6 about treewidth- η -modulator, however this time we use induction to construct the required protrusion decomposition.

Lemma 3.10. *Let \mathcal{G} be a graph class with the SQGM property. For any real $\varepsilon > 0$ and positive integer η , there exists an integer r such that if $G \in \mathcal{G}$ has a non-empty treewidth- η -modulator S , then G has $((1 + \varepsilon)|S|, r)$ -protrusion decomposition \mathcal{P} with S contained in the core of \mathcal{P} .*

Proof. Let $\alpha' > 0$ and $\frac{1}{2} \leq \mu < 1$ be the constants from Observation 3.5, in particular any graph in \mathcal{G} with a non-empty treewidth- η -modulator of size at most k has treewidth at most $\alpha'k^\mu$. Set $\alpha = \max\{\alpha', 1\}$. We define a series of constants. As in the beginning of the proof of Lemma 3.6 the purpose of these constants will become clear during the course of the proof.

- We set $\rho = \frac{1^\mu + 2^\mu - 3^\mu}{3^\mu}$ and note that $\rho > 0$.
- We define $\delta = \frac{(1+\varepsilon)4\alpha}{\rho}$, and
- $k_0 = (3 + 6\alpha)^{\frac{1}{1-\mu}} + 3 \cdot \left(\frac{\delta}{\varepsilon}\right)^{\frac{1}{1-\mu}}$.
- Finally, we set $r = \max\{k_0, \alpha k_0^\mu, \eta\}$

We first proof, using induction, the following claim:

Claim: for every $k \geq \frac{1}{3}k_0$, if a graph $G \in \mathcal{G}$ has a non-empty treewidth- η -modulator S of size at most k , then G has $((1+\varepsilon)k - \delta k^\mu, r)$ -protrusion decomposition with S contained in the core of this protrusion decomposition.

Proof of claim: In the base case we consider any k such that $\frac{1}{3}k_0 \leq k \leq k_0$. By Observation 3.5, any graph that has a non-empty treewidth- η -modulator S of size at most k has treewidth at most $\alpha'k^\mu \leq \alpha k_0^\mu \leq r$. Consider the protrusion-decomposition $\mathcal{P} = \langle S, V(G) \setminus S \rangle$, i.e., \mathcal{P} has core S and only one protrusion, namely $V(G) \setminus S$. Since $k \leq k_0 \leq r$, this is a (k, r) -protrusion decomposition. To complete the proof of the base case we need to show that $k \leq (1+\varepsilon)k - \delta k^\mu$. As in Lemma 3.6, by the choice of k_0 , we have that $0 \leq \varepsilon \frac{1}{3}k_0 - \delta \left(\frac{1}{3}k_0\right)^\mu$ (see (1)). Hence

$$0 \leq \varepsilon \frac{1}{3}k_0 - \delta \left(\frac{1}{3}k_0\right)^\mu \leq \varepsilon k - \delta k^\mu \Rightarrow k \leq (1+\varepsilon)k - \delta k^\mu.$$

In the last inequality we used that for any $\frac{1}{2} \leq \mu < 1$, $\varepsilon > 0$ and $\delta > 0$ the function $\varepsilon k - \delta k^\mu$ is monotonically increasing from the first point where it becomes positive. Thus the base case follows.

For the inductive step we let $k > k_0$ and assume that the statement holds for all values less than k . To prove the statement for k , let us consider a graph $G \in \mathcal{G}$ and a non-empty treewidth- η -modulator S of G of size at most k . If $|S| < k$ then by the induction hypothesis, G has a $((1+\varepsilon)(k-1) - \delta(k-1)^\mu, r)$ -protrusion decomposition. Since the function $(1+\varepsilon)k - \delta k^\mu$ is non-decreasing whenever it is non-negative and $k-1 \geq k_0$ it follows that this is also a $((1+\varepsilon)k - \delta k^\mu, r)$ -protrusion decomposition. Thus we may assume that $|S| = k$.

By Observation 3.5, the treewidth of G is at most $\text{tw}(G) \leq \alpha k^\mu$. We apply Lemma 2.1 to (G, S) and obtain a $2/3$ -balanced separation (A_1, A_2) of (G, S) of order at most $\text{tw}(G) + 1 \leq \alpha \cdot k^\mu + 1$. Let $L = A_1 \setminus A_2$, $X = A_1 \cap A_2$ and $R = A_2 \setminus A_1$. As $|X| \leq \text{tw}(G) + 1$ we obtain that $|X| \leq \alpha \cdot k^\mu + 1 \leq 2\alpha k^\mu$ (here we use the fact that $k \geq \frac{k_0}{3} \geq 1$). Since (A_1, A_2) is a $2/3$ -balanced separation, it follows that there exists a real $\frac{1}{3} \leq a \leq \frac{2}{3}$ such that $|L \cap S| \leq a|S|$ and $|R \cap S| \leq (1-a)|S|$.

Consider now the graph $G[L \cup X]$. Let $S_L = (S \cap L) \cup X$. Clearly S_L is a treewidth- η -modulator of $G[L \cup X]$ and $|S_L| = |S \cap L| + |X| \leq a|S| + |X| \leq ak + 2\alpha k^\mu$. In order to proceed with the induction, we have to verify that for $k > k_0$,

$$ak + 2\alpha k^\mu \leq k - 1 \quad (3)$$

The proof of (3) is very similar to the proof of (2) from Lemma 3.6. By the choice of k_0 , we have

$$\begin{aligned} k_0 &\geq (3 + 6\alpha)^{\frac{1}{1-\mu}} \Rightarrow k_0 \geq \left(\frac{3}{k_0^\mu} + 6\alpha\right)^{\frac{1}{1-\mu}} \Rightarrow \\ k_0^{1-\mu} &\geq \frac{3}{k_0^\mu} + 6\alpha \Rightarrow \frac{k_0}{3} \geq 1 + 2\alpha k_0^\mu \Rightarrow \\ \frac{2}{3}k_0 + 2\alpha k_0^\mu &\leq k_0 - 1 \end{aligned} \quad (4)$$

Now, because of (4), it is easy to verify by differentiation that the inequality $\frac{2}{3}k + 2\alpha k^\mu \leq k - 1$ holds for every $k > k_0$. Thus (3) follows. Notice also that S_L is non-empty because

$$|S_L| \geq |S \setminus R| = |S| - |S \cap R| \geq |S| - (1-a)|S| = a|S| \geq \frac{1}{3}k \geq \frac{1}{3}k_0 \geq 1,$$

and therefore we may apply the induction hypothesis to $G[L \cup X]$ with S_L as treewidth- η -modulator. We obtain a $((1+\varepsilon)|S_L| - \delta|S_L|^\mu, r)$ -protrusion decomposition \mathcal{P}_L of $G[L \cup X]$ with core containing S_L .

We now consider the graph $G[R \cup X]$ and define $S_R = (S \cap R) \cup X$. Working symmetrically to the case of S_L it is possible to deduce that $|S_R| \leq (1-a)k + 2\alpha k^\mu$, $(1-a)k + 2\alpha k^\mu \leq k - 1$, and $|S_R| \geq 1$, therefore we can apply the induction hypothesis to $G[R \cup X]$ and have a $((1+\varepsilon)|S_R| - \delta|S_R|^\mu, r)$ -protrusion decomposition \mathcal{P}_R of $G[R \cup X]$ with core containing S_R .

From the protrusion decompositions \mathcal{P}_L and \mathcal{P}_R we construct a protrusion decomposition \mathcal{P} of G . The core of \mathcal{P} is the union of the cores of \mathcal{P}_L and \mathcal{P}_R . The set of protrusions of \mathcal{P} is the union of the set of protrusions of \mathcal{P}_L and \mathcal{P}_R respectively. Since the cores of \mathcal{P}_L and \mathcal{P}_R contain X , the protrusions of \mathcal{P}_L and \mathcal{P}_R are also protrusions in G . Therefore \mathcal{P} is a $((1+\varepsilon)(|S_L| + |S_R|) - \delta(|S_L|^\mu + |S_R|^\mu), r)$ -protrusion decomposition of G containing S in its core. Thus, to finish the proof of the claim it is sufficient to show the following

$$(1+\varepsilon)(|S_L| + |S_R|) - \delta(|S_L|^\mu + |S_R|^\mu) \leq (1+\varepsilon)k - \delta k^\mu. \quad (5)$$

We now proceed with the proof of (5). Since $|S_L| \leq ak + 2\alpha k^\mu$ and $|S_R| \leq (1-a)k + 2\alpha k^\mu$ it follows that

$$(1+\varepsilon)(|S_L| + |S_R|) \leq (1+\varepsilon)(k + 4\alpha k^\mu). \quad (6)$$

Recall that $S_L = S \setminus (R \cap S)$, therefore $|S_L| = |S| - |R \cap S| \geq |S| - (1-a)|S| = a|S| = ak$. Working analogusly on S_R , we can show that $|S_R| \geq (1-a)k$. The two last inequalities imply that $|S_L|^\mu + |S_R|^\mu \geq (ak)^\mu + ((1-a)k)^\mu$, therefore

$$\delta(|S_L|^\mu + |S_R|^\mu) \geq \delta k^\mu(a^\mu + (1-a)^\mu). \quad (7)$$

Departing from (6) and (7) we prove 5 as follows.

$$\begin{aligned}
& (1 + \varepsilon)(|S_L| + |S_R|) - \delta(|S_L|^\mu + |S_R|^\mu) \\
& \leq (1 + \varepsilon)(k + 4\alpha k^\mu) - \delta k^\mu(a^\mu + (1 - a)^\mu) \\
& = (1 + \varepsilon)k - \delta k^\mu + k^\mu(4\alpha(1 + \varepsilon) - \delta(a^\mu + (1 - a)^\mu - 1)) \\
& \leq (1 + \varepsilon)k - \delta k^\mu + k^\mu(4\alpha(1 + \varepsilon) - \rho\delta) \\
& = (1 + \varepsilon)k - \delta k^\mu
\end{aligned}$$

In the transition from the third to the fourth line we used that $(1 - a)^\mu + a^\mu - 1 \geq \rho$ for any a between $\frac{1}{3}$ and $\frac{2}{3}$. This completes the proof of the claim.

We have now proved that for any $k \geq \frac{1}{3}k_0$ and for every graph $G \in \mathcal{G}$, if G has a non-empty treewidth- η -modulator S of size at most k , then G has an $((1 + \varepsilon)k - \delta k^\mu, r)$ -protrusion decomposition, which is also a $((1 + \varepsilon)k, r)$ -protrusion decomposition containing S in its core.

In the remaining case, where G has a non-empty treewidth- η -modulator S and $|S| \leq \frac{k_0}{3}$ then the protrusion decomposition $\mathcal{P} = \langle S, V(G) \setminus S \rangle$, where S is the core and $V(G) \setminus S$ is the unique protrusion, is an $(|S|, \eta)$ -protrusion decomposition. Since $|S| \leq (1 + \varepsilon)|S|$ and $\eta \leq r$ this completes the proof of the lemma. \square

Lemmata 3.6, 3.7, and 3.10 along with Observation 2.10, imply the following theorem, which is the first main technical contribution of this paper.

Theorem 3.11. *Let \mathcal{G} be a graph class with the SQGM (resp. SQGC) property and Π be a minor-bidimensional (resp. contraction-bidimensional) linear-separable problem. There exists a constant c such that every graph $G \in \mathcal{G}$, admits a $(c \cdot k, c)$ -protrusion decomposition, where $k = OPT_{\Pi \cap \mathcal{G}}(G)$.*

Proof. Let Π be a minor-bidimensional linear-separable problem and $k = OPT_{\Pi \cap \mathcal{G}}(G)$. We apply Lemma 3.6 for $\varepsilon = 1$, and we deduce that there exists an integer $\eta \geq 0$ such that graph G has a treewidth- η -modulator S of size at most k . If $|S| > 0$, then we can apply Lemma 3.10, for $\varepsilon = 1$ and obtain that there exists an integer r such that G admits a $(2 \cdot k, r)$ -protrusion decomposition. If $S = \emptyset$, then $\text{tw}(G) \leq \eta$ and $\langle S, V(G) \rangle$ is trivially an $(2 \cdot k, \eta)$ -protrusion decomposition of G . In any case, by setting $c = \max\{2, r, \eta\}$, we have that G admits a $(c \cdot k, c)$ -protrusion decomposition, as required.

The proof for the case where Π is contraction-bidimensional is the same as above with the difference that we now apply Lemma 3.7 instead of Lemma 3.6 and, before we apply Lemma 3.10, we use Observation 2.10 in order to show that \mathcal{G} has the SQGM property. \square

4 Finite index and finite integer index

In this section we prove the second main technical contribution of the paper: “CMSO + separability \Rightarrow finite integer index”.

4.1 Definitions on boundaried graphs

Boundaried graphs. A *boundaried graph* $\mathbf{G} = (G, B, \lambda)$ is a triple consisting of a graph G , a set $B \subseteq V(G)$ of distinguished vertices and an injective labelling λ from B to the set $\mathbb{Z}_{>0}$. The set B is called the *boundary* of \mathbf{G} and the vertices in B are called *boundary vertices* or *terminals*. Also, the graph G is the *underlying graph* of \mathbf{G} . Given a boundaried graph $\mathbf{G} = (G, B, \lambda)$, we define its *label set* by $\Lambda(\mathbf{G}) = \{\lambda(v) \mid v \in B\}$. If $\mathbf{G} = (G, B, \lambda)$ is a boundaried graph and $S \subseteq V(G)$, then the pair $\hat{\mathbf{G}} = (\mathbf{G}, S)$ is an *annotated boundaried graph* and we say that G is the *underlying graph* of $\hat{\mathbf{G}}$, B is its *boundary* (labeled by λ), and S is its *annotated set*.

Given a finite set $I \subseteq \mathbb{Z}_{>0}$, we define \mathcal{F}_I (resp. $\hat{\mathcal{F}}_I$) as the class of all boundaried graphs (resp. annotated boundaried graphs) whose label set is I . Similarly, we define $\mathcal{F}_{\subseteq I} = \bigcup_{I' \subseteq I} \mathcal{F}_{I'}$ (resp. $\hat{\mathcal{F}}_{\subseteq I} = \bigcup_{I' \subseteq I} \hat{\mathcal{F}}_{I'}$). We also denote by \mathcal{F} (resp. $\hat{\mathcal{F}}$) the class of all boundaried graphs, i.e., $\mathcal{F} = \mathcal{F}_{\subseteq \mathbb{Z}_{>0}}$ (resp. $\hat{\mathcal{F}} = \hat{\mathcal{F}}_{\subseteq \mathbb{Z}_{>0}}$). Finally we say that a boundaried graph \mathbf{G} is a *t-boundaried graph* if $\Lambda(\mathbf{G}) \subseteq \{1, \dots, t\}$.

Let \mathcal{G} be a class of (not boundaried) graphs. We say that a boundaried graph \mathbf{G} *belongs to* \mathcal{G} if the underlying graph of \mathbf{G} belongs to \mathcal{G} . We also use $V(\mathbf{G})$ to denote the vertex set of the underlying graph of \mathbf{G} .

The gluing operation. Let $\mathbf{G}_1 = (G_1, B_1, \lambda_1)$ and $\mathbf{G}_2 = (G_2, B_2, \lambda_2)$ be two boundaried graphs. We denote by $\mathbf{G}_1 \oplus \mathbf{G}_2$ the graph (not boundaried) obtained by taking the disjoint union of G_1 and G_2 and identifying equally-labeled vertices of the boundaries of G_1 and G_2 . The gluing operation maintains edges of both graphs that are glued, i.e., in $\mathbf{G}_1 \oplus \mathbf{G}_2$ there is an edge between two labeled vertices if there is either an edge between them in G_1 or in G_2 , (in case of multi-graphs, multiplicities of edges are summed up in the new graph).

Let $G = \mathbf{G}_1 \oplus \mathbf{G}_2$ where $\mathbf{G}_1 = (G_1, B_1, \lambda_1)$ and $\mathbf{G}_2 = (G_2, B_2, \lambda_2)$ are boundaried graphs. We define the *glued set* of \mathbf{G}_i as the set $B_i^* = \lambda_i^{-1}(\Lambda(\mathbf{G}_1) \cap \Lambda(\mathbf{G}_2))$, $i = 1, 2$. For a vertex $v \in V(\mathbf{G}_1)$ we define its *heir* $\text{heir}(v)$ in G as follows: if $v \notin B_1^*$ then $\text{heir}(v) = v$, otherwise $\text{heir}(v)$ is the result of the identification of v with an equally labeled vertex in \mathbf{G}_2 . The *heir* of a vertex in G_2 is defined symmetrically. The common boundary of \mathbf{G}_1 and \mathbf{G}_2 in G is equal to $\text{heir}(B_1^*) = \text{heir}(B_2^*)$ where the evaluation of *heir* on vertex sets is defined in the obvious way.

Let now $\hat{\mathbf{G}}_1 = (\mathbf{G}_1, S_1)$ and $\hat{\mathbf{G}}_2 = (\mathbf{G}_2, S_2)$ be two annotated boundaried graphs. We define $\hat{\mathbf{G}}_1 \oplus \hat{\mathbf{G}}_2$ as the annotated graph $\hat{G} = (\mathbf{G}_1 \oplus \mathbf{G}_2, \text{heir}(S_1) \cup \text{heir}(S_2))$.

4.2 Finite Index

Let ϕ be a predicate on annotated graphs. In other words ϕ is a function that takes as input a graph G and vertex set $S \subseteq V(G)$, and outputs true or false. We define a *canonical equivalence relation* \equiv_ϕ on boundaried annotated graphs as follows. For two annotated boundaried graphs $\hat{\mathbf{G}}_1 = (\mathbf{G}_1, S_1)$ and $\hat{\mathbf{G}}_2 = (\mathbf{G}_2, S_2)$, we say that $\hat{\mathbf{G}}_1 \equiv_\phi \hat{\mathbf{G}}_2$ if $\Lambda(\mathbf{G}_1) = \Lambda(\mathbf{G}_2)$ and for every annotated boundaried graph $\hat{\mathbf{G}} = (\mathbf{G}, S)$ we have that

$$\phi(\hat{\mathbf{G}}_1 \oplus \hat{\mathbf{G}}) = \text{true} \Leftrightarrow \phi(\hat{\mathbf{G}}_2 \oplus \hat{\mathbf{G}}) = \text{true}$$

It is easy to verify that \equiv_ϕ is an equivalence relation. We say that ϕ is *finite state* if, for every finite $I \subseteq \mathbb{Z}_{>0}$, the equivalence relation \equiv_ϕ has a finite number of equivalence classes when restricted to $\hat{\mathcal{F}}_I$. A formal proof of the following fact can be found in [10, 21].

Proposition 4.1. *For every CMSO-definable predicate ϕ on annotated graphs, ϕ has finite state.*

Given a CMSO-definable predicate ϕ on annotated graphs and an $I \subseteq \mathbb{Z}_{>0}$, we use the notation $\hat{\mathcal{R}}_{\phi, I}$ for a set containing one minimum-size representative from each of the equivalence classes of \equiv_ϕ when restricted to $\hat{\mathcal{F}}_I$.

4.3 Finite Integer Index

Definition 4.2. [Canonical equivalence on boundaried graphs.] Let Π be a parameterized graph problem whose instances are pairs of the form (G, k) . Given two boundaried graphs $\mathbf{G}_1, \mathbf{G}_2 \in \mathcal{F}$, we say that $\mathbf{G}_1 \equiv_\Pi \mathbf{G}_2$ if $\Lambda(\mathbf{G}_1) = \Lambda(\mathbf{G}_2)$ and there exists a *transposition constant* $c \in \mathbb{Z}$ such that

$$\forall(\mathbf{F}, k) \in \mathcal{F} \times \mathbb{Z} \quad (\mathbf{G}_1 \oplus \mathbf{F}, k) \in \Pi \Leftrightarrow (\mathbf{G}_2 \oplus \mathbf{F}, k + c) \in \Pi.$$

Note that the relation \equiv_Π is an equivalence relation. Observe that c could be negative in the above definition. This is the reason we gave the definition of parameterized problems to include negative parameters also.

Notice that two boundaried graphs with different label sets belong to different equivalence classes of \equiv_Π . We are now in position to give the following definition.

Definition 4.3. [Finite Integer Index] A parameterized graph problem Π whose instances are pairs of the form (G, k) has *Finite Integer Index* (or simply has *FII*), if and only if for every finite $I \subseteq \mathbb{Z}^+$, the number of equivalence classes of \equiv_Π that are subsets of \mathcal{F}_I is finite.

The notion of FII first appeared in the works of [13, 23] and is similar to the notion of *finite state* [1, 14, 20].

4.4 A condition for proving FII

Theorem 4.4. *If Π is a separable CMSO-minimization or CMSO-maximization problem, then Π has FII.*

Proof. We prove the statement for vertex subset CMSO-minimization problems. The proofs for CMSO-maximization problems and the edge subset variants are identical. Let ε be the CMSO-predicate defining Π , that is $(G, k) \in \Pi$ if and only if there exists a vertex set S of size at most k such that $(G, S) \models \varepsilon$. Let I be a finite subset of $\mathbb{Z}_{>0}$ and let $\hat{\mathcal{R}}_{\varepsilon, I}$ be a set of minimum-size representatives of \equiv_ε when restricted to $\hat{\mathcal{F}}_I$, the set of annotated boundaried graphs with label set I . By Proposition 4.1, the set $\hat{\mathcal{R}}_{\varepsilon, I}$ is finite. Let f be a function such that Π is f -separable.

Given an annotated boundaried graph $\mathbf{G} \in \hat{\mathcal{F}}_I$, we define the function $\xi_{\mathbf{G}} : \hat{\mathcal{R}}_{\varepsilon, I} \rightarrow \mathbb{Z}_{\geq 0} \cup \{\perp\}$ such that if $(\mathbf{G}^*, S^*) \in \hat{\mathcal{R}}_{\varepsilon, I}$, then

$$\xi_{\mathbf{G}}(\mathbf{G}^*, S^*) = \min\{|S| \mid (\mathbf{G}, S) \equiv_{\varepsilon} (\mathbf{G}^*, S^*)\}. \quad (8)$$

Here $\xi_{\mathbf{G}}(\mathbf{G}^*, S^*) = \perp$ if no set S such that $(\mathbf{G}, S) \equiv_{\varepsilon} (\mathbf{G}^*, S^*)$ exists. We may think of $\xi_{\mathbf{G}}$ as a partial function where \perp means that the function is left undefined. We also define function $\chi_{\mathbf{G}}(\mathbf{G}^*, S^*) : \hat{\mathcal{R}}_{\varepsilon, I} \rightarrow \mathbb{Z} \cup \{\perp\}$ as follows

$$\chi_{\mathbf{G}}(\mathbf{G}^*, S^*) = \begin{cases} \xi_{\mathbf{G}}(\mathbf{G}^*, S^*) - \text{OPT}_{\Pi}(G), & \text{if } \xi_{\mathbf{G}}(\mathbf{G}^*, S^*) \in [\text{OPT}_{\Pi}(G) - f(|I|), \text{OPT}_{\Pi}(G) + f(|I|)] \\ \perp, & \text{otherwise.} \end{cases}$$

Thus $\chi_{\mathbf{G}}$ outputs \perp if $\xi_{\mathbf{G}} = \perp$ or when $\xi_{\mathbf{G}}(\mathbf{G}^*, S^*) \notin [\text{OPT}_{\Pi}(G) - f(|I|), \text{OPT}_{\Pi}(G) + f(|I|)]$. Clearly $\chi_{\mathbf{G}}(\mathbf{G}^*, S^*) \in [-f(|I|), f(|I|)] \cup \{\perp\}$ for all choices of $(\mathbf{G}^*, S^*) \in \hat{\mathcal{R}}_{\varepsilon, I}$. This means that $\chi_{\mathbf{G}}$ may take at most $2 \cdot f(|I|) + 2$ different values, and hence there are at most $(2 \cdot f(|I|) + 2)^{|\hat{\mathcal{R}}_{\varepsilon, I}|}$ different functions $\chi_{\mathbf{G}}$.

Now, we define an equivalence relation \sim on \mathcal{F}_I such that, given $\mathbf{G}_1, \mathbf{G}_2 \in \mathcal{F}_I$,

$$\mathbf{G}_1 \sim \mathbf{G}_2 \text{ if and only if } \chi_{\mathbf{G}_1} = \chi_{\mathbf{G}_2}$$

and keep in mind that \sim has a finite number of equivalence classes. We will prove that \sim is a refinement of \equiv_{Π} . For this we prove that if $\mathbf{G}_1 \sim \mathbf{G}_2$, then $\mathbf{G}_1 \equiv_{\Pi} \mathbf{G}_2$.

We assume that $\mathbf{G}_1 \sim \mathbf{G}_2$, where $\mathbf{G}_1 = (G_1, B_1)$ and $\mathbf{G}_2 = (G_2, B_2)$. We also set $c = \text{OPT}_{\Pi}(G_2) - \text{OPT}_{\Pi}(G_1)$ (notice that c depends only on \mathbf{G}_1 and \mathbf{G}_2). In what follows, we prove that, for every $(\mathbf{F}, k) \in \mathcal{F}_I \times \mathbb{Z}$, $(\mathbf{G}_1 \oplus \mathbf{F}, k) \in \Pi$ implies that $(\mathbf{G}_2 \oplus \mathbf{F}, k + c) \in \Pi$ (the direction $(\mathbf{G}_1 \oplus \mathbf{F}, k) \notin \Pi \Rightarrow (\mathbf{G}_2 \oplus \mathbf{F}, k + c) \notin \Pi$ is symmetric).

We consider some $(\mathbf{F}, k) \in \mathcal{F}_I \times \mathbb{Z}$ and we denote $H = \mathbf{G}_1 \oplus \mathbf{F}$. Let S be an optimal solution of H . That is, $|S| = \text{OPT}_{\Pi}(H)$ and $(H, S) \models \varepsilon$. The fact that $(H, k) \in \Pi$ means that the size of S is at most k . We define $S_1 = S \cap V(\mathbf{G}_1)$ and $S^* = S \setminus S_1$ and we know that $(\mathbf{G}_1 \oplus \mathbf{F}, S_1 \cup S^*) \models \varepsilon$. Notice that $|S_1| + |S^*| = |S_1 \cup S^*| \leq k$.

Claim: There exists no $S'_1 \subseteq V(\mathbf{G}_1)$ such that $|S'_1| < |S_1|$ and $(\mathbf{G}_1, S'_1) \equiv_{\varepsilon} (\mathbf{G}_1, S_1)$.

Proof of Claim: Suppose in contrary that such an S'_1 exists. Since $(\mathbf{G}_1, S'_1) \equiv_{\varepsilon} (\mathbf{G}_1, S_1)$, we have that $(\mathbf{G}_1 \oplus \mathbf{F}, S'_1 \cup S^*) \models \varepsilon$. However, this implies that $|S'_1 \cup S^*| = |S'_1| + |S^*| < |S_1| + |S^*| = \text{OPT}_{\Pi}(H)$, which is a contradiction. This concludes the claim. \square

Let $(\mathbf{G}', S') \in \hat{\mathcal{R}}_{\varepsilon, I}$ such that $(\mathbf{G}', S') \equiv_{\varepsilon} (\mathbf{G}_1, S_1)$. By the above claim and (8), we have that $\xi_{\mathbf{G}_1}(\mathbf{G}', S') = |S_1|$. Let $A_1 = \text{heir}(V(\mathbf{G}_1))$ and $A_2 = \text{heir}(V(\mathbf{F}))$ (here by the *heir* of a vertex set we mean the union of all vertex heir it contains). Here, function *heir* is defined with respect to the operation $H = \mathbf{G}_1 \oplus \mathbf{F}$ and certainly (A_1, A_2) is a separation of H of order $|I|$.

Recall that $S_1 = S \cap A_1$ and $|S| = \text{OPT}_{\Pi}(H)$. The f -separability of Π implies that

$$|S \cap A_1| \in [\text{OPT}_{\Pi}(H[A_1]) - f(|I|), \text{OPT}_{\Pi}(H[A_1]) + f(|I|)],$$

and therefore,

$$\xi_{\mathbf{G}_1}(\mathbf{G}', S') \in [OPT_{\Pi}(G_1) - f(|I|), OPT_{\Pi}(G_1) + f(|I|)].$$

Then $\chi_{\mathbf{G}_1}(\mathbf{G}', S') \neq \perp$ and thus $\chi_{\mathbf{G}_1}(\mathbf{G}', S') = |S_1| - OPT_{\Pi}(G_1)$. As $\mathbf{G}_1 \sim \mathbf{G}_2$, this means that $\chi_{\mathbf{G}_2}(\mathbf{G}', S') = |S_1| - OPT_{\Pi}(\mathbf{G}_1)$. Since $\chi_{\mathbf{G}_2}(\mathbf{G}', S') \neq \perp$ it follows that $\chi_{\mathbf{G}_2}(\mathbf{G}', S') = \xi_{\mathbf{G}_2}(\mathbf{G}', S') - OPT_{\Pi}(G_2)$. We conclude that

$$\xi_{\mathbf{G}_2}(\mathbf{G}', S') = |S_1| + c. \quad (9)$$

By the definition of ξ and (9), there exists a set S_2 of size $|S_1| + c$ such that $(\mathbf{G}_2, S_2) \equiv_{\varepsilon} (\mathbf{G}', S')$. As $(\mathbf{G}_1, S_1) \equiv_{\varepsilon} (\mathbf{G}_2, S_2)$, the fact that $(\mathbf{G}_1 \oplus \mathbf{F}, S_1 \cup S^*) \models \varepsilon$, implies that $(\mathbf{G}_2 \oplus \mathbf{F}, S_2 \cup S^*) \models \varepsilon$. But then $|S_2 \cup S^*| = |S_2| + |S^*| = |S_1| + c + |S^*| \leq k + c$. This means that $(\mathbf{G}_2 \oplus \mathbf{F}, k + c) \in \Pi$. Hence $\mathbf{G}_1 \equiv_{\Pi} \mathbf{G}_2$.

Thus for every pair of boundaried graphs \mathbf{G}_1 and \mathbf{G}_2 , condition $\mathbf{G}_1 \sim \mathbf{G}_2$, yields that $\mathbf{G}_1 \equiv_{\Pi} \mathbf{G}_2$. Since \sim has a finite number of equivalence classes, we conclude that Π has FII. \square

We remark that the proof of Theorem 4.4 is similar in spirit to the proof of [10, Lemma 7.3] with a few key differences. Being separable is a looser constraint for CMSO-optimization problems than being strongly monotone (see [10] for the definition.) In particular, separability only puts restrictions on how an *optimum* solution can interact with both sides of the separation. On the other hand, strong monotonicity puts constraints on how *any* solution can interact with the two sides. Therefore, it is often easier to verify that a CMSO-optimization problem is separable than that it is strongly monotone. However, strong monotonicity allows us to conclude even stronger properties than FII as observed in [39].

5 Proof of the main theorem: Putting things together

Now everything prepared to pipeline the results about protrusion decomposition and FII with the framework from [10]. We start from the following definitions.

Definition 5.1. $[(f, a)$ -**protrusion replacement family**] Let Π be a parameterized graph problem, let $f : \mathbb{Z}^+ \rightarrow \mathbb{Z}^+$ be a non-decreasing function and let $a \in \mathbb{Z}^+$. An (f, a) -*protrusion replacement family* for Π is a collection $\mathcal{A} = \{A_i \mid i \geq 0\}$ of algorithms, such that algorithm A_i receives as input a pair (I, X) , where

- I is an instance of Π whose graph and parameter are G and $k \in \mathbb{Z}$,
- X is an i -protrusion of G with at least $f(i) \cdot k^a$ vertices,

and outputs an equivalent instance I^* such that, if G^* and k^* are the graph and the parameter of I^* , then $|V(G^*)| < |V(G)|$ and $k^* \leq k$.

The following two properties for a parameterized graph problem Π were defined in [10].

A [Protrusion replacement:] There exists an (f, a) -protrusion replacement family \mathcal{A} for Π , for some function $f : \mathbb{Z}^+ \rightarrow \mathbb{Z}^+$ and some $a \in \mathbb{Z}^+$.

B [Protrusion decomposition:] There exists a constant c such that, if G and $k \in \mathbb{Z}^+$ are the graph and the parameter of a yes-instance of Π then G admits a $(c \cdot \text{OPT}_\Pi(G), c)$ -protrusion decomposition.

Our kernelization result is based on the following master theorem from [10].

Theorem 5.2. *If a parameterized graph problem Π has property **A** for some nonnegative constant a and property **B** for some constant c , then Π admits a kernel of size $O(k^{a+1})$.*

The following lemma proven in [10], shows that every parameterized problem that has FII admits $(f, 0)$ -protrusion replacement families.

Lemma 5.3. *Every parameterized graph problem Π that has FII has the protrusion replacement property **A** for $a = 0$.*

Proof of Theorem 1.1. Let Π be a CMSO-definable linear-separable minor-bidimensional problem on an H -minor-free graph class \mathcal{G} for some fixed graph H . By Proposition 2.11, \mathcal{G} has the SQGM property.

By Theorem 4.4, every separable CMSO-minimization or CMSO-maximization problem Π has FII. Thus by Lemma 5.3, Π has the protrusion replacement property **A** for $a = 0$. By Theorem 3.11, Π has the protrusion decomposition property **B**. By Theorem 5.2, Π admits a linear kernel.

The proof for contraction-bidimensional problems is almost identical. \square

6 Conclusion

We conclude with a discussion on the relation between the new and previous kernelization meta-theorems, the running times of kernelization algorithms and an open question.

Bidimensionality vs. quasi-coverability. Let us finally note how the techniques developed to prove the kernelization meta-theorem in this paper, can be used to refine the results from [10].

Let r be a non-negative integer. We say that a parameterized graph problem Π has the *radial r -coverability property* if all yes-instances of Π encode graphs embeddable in some surface of Euler genus at most r and there exist such an embedding of G and a set $S \subseteq V(G)$ such that $|S| \leq r \cdot k$ and $\mathbf{R}_G^r(S) = V(G)$ (here we denote by $\mathbf{R}_G^r(S)$ the set of all vertices of G that are within radial distance¹ at most r from some vertex in S). A parameterized graph problem Π has the *radial r -quasi-coverability property* if all yes-instances of Π encode graphs embeddable in some surface of Euler genus at most r and there exist such an embedding and a set $S \subseteq V(G)$ such that $|S| \leq r \cdot k$

¹The *radial distance* between two vertices x, y in a surface-embedded graph is one less than the minimum size of a sequence of alternating vertices and faces, starting on x and finishing on y , where if an vertex v and a face f appear consecutively in the sequence, then v is a vertex of the boundary of f .

and $\text{tw}(G - \mathbf{R}_G^r(S)) \leq r$. Every problem Π that has the radial r -coverability property is radially r -quasi-coverable. The converse is not necessarily true.

The main two meta-theorems from [10] say that quasi-coverable problems with finite integer index admit linear kernels on graphs of bounded genus and that coverable CMSO-definable problems admit polynomial kernels on graphs of bounded genus. By making use of SQGC property of graphs of bounded genus, it is not difficult to prove that if a set S is such that $\text{tw}(G - \mathbf{R}_G^r(S)) \leq r$, then S is also a treewidth- η -modulator for $\eta = \mathcal{O}(r)$. The following lemma allows to refine the results from [10].

Lemma 6.1. *If Π is a quasi-coverable problem on graphs of bounded genus, then there exists a contraction-bidimensional separable problem Π^* such that if $(G, k) \in \Pi$, then $\text{OPT}_{\Pi^*}(G) = \mathcal{O}(k)$.*

Proof. The lemma follows from the following observation. Let Π be an r -quasi-coverable problem. Consider the following problem Π^* : for a graph G of genus g , pair $(G, k) \in \Pi^*$ if and only if there is a subset of vertices $S \subseteq V(G)$ such that $|S| \leq r \cdot k$ and $\text{tw}(G - \mathbf{R}_G^r(S)) \leq r$. Because Π is r -quasi-coverable, we have that if $(G, k) \in \Pi$, then $\text{OPT}_{\Pi^*}(G) \leq k$. As in the case with TREEWIDTH- η -MODULATOR, it is easy to see that Π^* is contraction-bidimensional and separable. \square

By Lemma 3.10, there is a protrusion decomposition for Π^* and thus there is a protrusion decomposition for Π . This implies one of the main result of [10]: Every quasi-coverable problem with finite integer index admits a linear kernel on graphs of bounded genus. This also can be used to show that every quasi-coverable CMSO-definable problem admits a polynomial kernel on graphs of bounded genus. Thus the results of this paper subsume all the linear and polynomial kernels on graphs of bounded genus from [10]. We refer to [10] for the list of these problems.

Running time. In Theorem 1.1 we do not specify the running time of our kernelization algorithms. The running time of such algorithms depends on the running time hidden in the lemma from [10] (Lemma 5.2), which in turn, depends on how fast one can identify and replace protrusions in a graph. By applying the fast “protrusion replacer” from [38], see also [44, Chapter 16], it is possible to achieve kernelization algorithms in Theorem 1.1 which run in time linear in the input size.

Open question. We conclude with the following open question. For separable contraction-bidimensional CMSO-definable problems, the technique developed in this paper yields the existence of a linear kernel on H -minor-free graphs only when H is an apex graph. An interesting open question is to identify general logic and combinatorial conditions for contraction-bidimensional problems which yield a polynomial kernelization on H -minor-free graphs. For some contraction-bidimensional problems, like DOMINATING SET or CONNECTED DOMINATING SET, linear kernels are known to exist for H -minor-free graphs and H -topological-minor-free graphs [43], and even for graphs of bounded expansion [34], see also [35] but we are still far from a general understanding of the complexity of kernelization for such problems beyond apex-minor-free graphs.

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