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Complexity and algorithms for constant diameter augmentation problems

Eun Jung Kim*, Martin Milanič†, Jérôme Monnot*, Christophe Picouleau‡

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Abstract

We study the following problem: for given integers d, k and graph G , can we obtain a graph with diameter d via at most k edge deletions? We determine the computational complexity of this and related problems for different values of d .

Keywords: graph diameter, blocker problem, polynomial algorithm, NP-complete.

To the memory of our colleague and friend Jérôme.

1 Introduction

A typical graph modification problem aims to modify a graph parameter π via a small number of operations. The operations involved may be vertex deletion, edge deletion, edge addition, edge contraction, etc. That is, for a fixed graph operation, we ask, given a graph G and two integers k and d , whether G can be transformed into a graph G' by using at most k operations so that $\pi(G') \leq \pi(G) - d$ or $\pi(G') \geq \pi(G) + d$. Typically, $\pi(G)$ is defined as the optimal value of an instance of an optimization problem on graphs. Thus, such problems are called *blocker problems*, as the set of vertices or edges involved in the operations “block” the optimal value of the problem. Identifying the part of the graph responsible for a significant change of the parameter under consideration gives crucial information on the graph.

Blocker problems have been given much attention over the last few years [4, 5, 10, 33]. Graph parameters considered were the matching number [30, 36], the s, t -shortest path, the maximum flow and minimum s, t -cut values [10], the chromatic number [2, 12], the independence number [2, 3, 12], the clique number [12, 28], the domination number [14, 29], and the vertex cover number [3, 24, 29], for instance. Concerning the operations involved to the minimal modification of a graph parameter or a graph property the reader can see [12, 22] for node deletion, [29, 34] for edge deletion, [12, 14, 17, 24, 34] for edge contraction.

Here, we aim to determine the number of edges to delete from a given graph to increase the diameter of the graph by a fixed amount. Our goal is to determine the computational complexity,

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that is, to either find polynomial-time algorithms or prove NP-completeness, for several related problems.

We consider the following decision problems:

EXACT DIAMETER- d AUGMENTATION (d -EDA):

Input: A connected graph G .

Question: Is there a set $F \subseteq E(G)$ such that $G - F$ is a connected graph of diameter d ?

DIAMETER- d AUGMENTATION (d -DA):

Input: A connected graph G .

Question: Is there a set $F \subseteq E(G)$ such that $G - F$ is a connected graph of diameter at least d ?

MINIMUM EXACT DIAMETER- d AUGMENTATION (d -MEDA)

Input: A connected graph G , a positive integer k .

Question: Is there a set $F \subseteq E(G)$ such that $|F| \leq k$ and $G - F$ is a connected graph of diameter d ?

MINIMUM DIAMETER- d AUGMENTATION (d -MDA)

Input: A connected graph G , a positive integer k .

Question: Is there a set $F \subseteq E(G)$ such that $|F| \leq k$ and $G - F$ is a connected graph of diameter at least d ?

In all the above problems we consider d to be a fixed constant independent of the input graph. If d is part of input, then all four problems are NP-complete as they generalize the HAMILTONIAN PATH problem (which is obtained by taking $d = |V(G)| - 1$ and, when necessary, $k = |E(G)|$), see [31].

Our results and organization of the paper. In Section 2 we give the main definitions and notations we use throughout the paper together with some preliminary observations about the decisions problems defined above. In Sections 3 and 4 we give polynomial-time algorithms for the 3-MDA, 3-MEDA, and 3-EDA problems. Section 5 is devoted to NP-completeness results for k -MEDA and k -MDA problems for $k \geq 5$. We show that 5-MEDA is NP-complete for graphs of diameter $d \in \{3, 4\}$, and show the NP-completeness of k -MEDA for all $k \geq 6$, for some different values of the diameter of the input graph. The same results also hold for the k -MDA problems. In Section 6 we give a conclusion and propose some future research problems.

Related work. There is a large literature studying extremal questions related to diameter (see, e.g., [7]). Among the questions most relevant to our study, let us mention the effect on the diameter of deleting t edges from a $(t + 1)$ -edge-connected graph [9], the maximum number of edges in an n -vertex diameter- d -critical graph (that is, a graph with diameter d such that the diameter increases upon deletion of any edge; see [25]) and the minimum number of edges in an n -vertex graph such that deleting any edge results in a graph with diameter d (see [8]). A connection between the 3-MEDA problem and Moore graphs will be established in Section 4. From the algorithmic point of view, related problems extensively studied in the literature include the complementary problem of adding a minimum number of edges to achieve diameter d (see, e.g., [13, 15, 19, 23, 31]) and the length-bounded max flow and min cut problems (see [1, 26]).

2 Notations and preliminaries

We only consider finite, simple, and undirected graphs. We refer to [35] for undefined terminology.

Let $G = (V, E)$ be a graph. For a subset $F \subseteq E$, we let $G - F$ denote the subgraph of G with vertex set V and edge set $E \setminus F$. The subgraph of G induced by a set $S \subseteq V$ is denoted by $G[S]$ and defined as $G[S] = (S, F)$ where $F = \{xy \in E(G) | x, y \in S\}$. The length of a path is its number of edges. Given two vertices u and v , $\text{dist}_G(u, v)$ denotes the minimum length of a path between u and v (note that since G is assumed to be connected $\text{dist}_G(u, v)$ is finite). A u, v -path of minimum length is called a shortest u, v -path. The *diameter* of G is the maximum length of a shortest path, that is, $\text{diam}(G) = \max_{(u,v) \in V^2} \{\text{dist}_G(u, v)\}$. A *diametral path* in a graph G is any shortest path between vertices u and v such that $\text{dist}_G(u, v) = \text{diam}(G)$. A *cut-vertex* in a connected graph G is a vertex $v \in V(G)$ such that $G - v$ is not connected. A graph is *biconnected* if it is connected and has no cut-vertices. A *cut-edge* in a connected graph G is an edge $e \in E(G)$ such that $G - e$ is not connected. A vertex in a graph is *universal* if it is adjacent to all other vertices. The *girth* of a graph G , denoted $g(G)$, is the shortest length of a cycle in G (or ∞ if G is acyclic). A *triangle* in a graph G is a cycle of length three. Given two graphs G and H , we say that G is *H -free* if no induced subgraph of G is isomorphic to H . The *\mathcal{F} -free graphs* where \mathcal{F} is a finite set of graphs are defined analogously.

(Non)monotonicity

It follows directly from the definitions that the d -DA and d -MDA problems satisfy a certain monotonicity, in the sense that for every $d < d'$, every graph G that is a yes instance to the d' -DA problem is also a yes instance to the d -DA problem, and similarly, every pair (G, k) that is a yes instance to the d' -MDA problem is also a yes instance to the d -MDA problem. Unsurprisingly, the analogous monotonicity fails to hold for the d -EDA and d -MEDA problems. It is not difficult to construct graphs G and pairs of positive integers d, d' , and k , with $d < d'$ such that G is a yes instance to d -EDA as well as to d' -EDA, while (G, k) is a yes instance to d' -MEDA but not to d -MEDA. In other words, there is a set of at most k edges such that deleting them from G results in a graph with diameter d' but there is no set of at most k edges whose deletion would produce a graph with diameter d .

Example 2.1. *The graph shown in Fig. 1 has diameter 3, deleting any of the edges $v_1v_6, v_3v_4, v_4v_5, v_5v_6$ results in a graph with diameter 5, while deleting any other edge does not change the diameter. However, a graph with diameter 4 can be obtained by deleting, for example, the two edges v_0v_1 and v_0v_3 .*

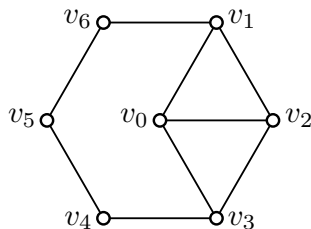


Figure 1: A graph in which the diameter can be increased by two by deleting a single edge, while two edges need to be deleted to increase the diameter by one.

The d -DA problem

We show how the d -DA problem is related to the problem of searching for a path of length d . Then we deduce that it is polynomial-time solvable.

Proposition 2.2. *For a connected graph G and a positive integer d , the following are equivalent.*

1. G is a yes instance to the d -DA problem.
2. G has a connected spanning subgraph of diameter at least d .
3. G has a spanning tree of diameter at least d .
4. G contains a path of length at least d .

Proof. The equivalence between statements 1 and 2 is immediate from the definition of the d -DA problem. Clearly, statement 3 implies statement 2. Moreover, statement 2 implies statement 4, as if G' is a connected spanning subgraph of G of diameter at least d , then any diametral path in G' is a path in G of length at least d . Finally, we argue that statement 4 implies statement 3. Let P be a path in G of length at least d . Then, G has a depth-first traversal starting with the edges of P along the path. The corresponding DFS tree is a spanning tree of G of diameter at least d , implying statement 3. \square

Since for each fixed d , the presence of a path of length at least d can be tested in polynomial time, Proposition 2.2 implies the following.

Corollary 2.3. *For every positive integer d , the d -DA problem is solvable in polynomial time.*

One may wonder whether there are similar characterizations of yes instances for the d -EDA problem as those given in Proposition 2.2, in particular, whether the existence of a connected spanning subgraph of diameter d implies the existence of a spanning tree of diameter d . While this implication clearly fails if d is allowed to be the diameter of G (take for instance any cycle of length at least 4), examples can also be constructed showing that the implication fails if $d < \text{diam}(G)$.

Example 2.4. *Let G be the graph obtained from the 5-cycle by adding to it a new vertex v and connecting it by an edge to exactly two non-adjacent vertices of the 5-cycle. Graph G is of diameter two, it has a connected spanning subgraph of diameter three (for example, the subgraph obtained by deleting one of the two edges incident with v) but it is not difficult to verify that G has no spanning tree of diameter three.*

Complete graphs

In the case when the input graph G is complete, the d -DA, d -EDA, d -MDA, and d -MEDA problems are either trivial or related to a known and solved question in extremal graph theory.

First, since the maximum diameter of an n -vertex connected graph is $n - 1$, the complete graph K_n is a yes instance to the d -DA (or to the d -EDA) problem if and only if $n \geq d + 1$.

Second, the d -MEDA and the d -MDA problems on K_n are equivalent to the problems of deleting the smallest number of edges from K_n so that the resulting graph is of diameter exactly d (resp., at least d). Both of these questions are equivalent to a well-known problem in extremal graph theory asking for the maximum number of edges in an n -vertex graph of diameter d . O. Ore proved that

the maximum number of edges in a graph with n vertices and diameter d is upper bounded by $d + \frac{1}{2}(n-d-1)(n-d+4)$ and this bound is tight [27]. Thus, the answer to either d -MEDA or d -MDA problem given K_n and a positive integer k is yes if and only if $k \geq \binom{n}{2} - d - \frac{1}{2}(n-d-1)(n-d+4)$.

For later use, we make note of the above observations.

Theorem 2.5 (Ore). *For all $d \geq 1$, the d -EDA, d -MDA, and d -MEDA problems are solvable in linear time for complete graphs.*

The d -EDA, d -MDA, and d -MEDA problems for $d \in \{1, 2\}$

We give some observations solving the cases $d \in \{1, 2\}$. First, note that for each of the d -EDA, d -MDA, and d -MEDA problems, we may assume that $\text{diam}(G) < d$ since otherwise the answer is trivial:

- If $\text{diam}(G) > d$, then G is a no instance to the d -EDA problem (since no spanning subgraph of G has diameter exactly d). Similarly, any (G, k) with $k \geq 0$ is a yes instance to the d -MDA problem and a no instance to the d -MEDA problem.
- If $\text{diam}(G) = d$, then G is a yes instance to the d -EDA problems (take $F = \emptyset$) and any (G, k) with $k \geq 0$ is a yes instance to the d -MDA and d -MEDA problems.

Since complete graphs are the only graphs with diameter strictly less than 2, the above observation implies that all three problems are polynomially solvable for $d \in \{1, 2\}$.

3 A polynomial-time algorithm for the 3-MDA problem

We prove the following.

Theorem 3.1. *The 3-MDA problem is solvable in polynomial time.*

The algorithm is based on a polynomial-time algorithm for the following problem for the case $d = 3$ (see [26]).

MINIMUM DISTANCE- d AUGMENTATION (d -MDISA)

Input: A connected graph G , distinct vertices x, y of G , a positive integer k .
Question: Is there a set $F \subseteq E(G)$ such that $|F| \leq k$, $G - F$ is a connected graph and $\text{dist}_{G-F}(x, y) \geq d$?

We note that the inapproximability of two optimization problems related to d -MDISA when d is not constant was established in [21]. Beier et al. showed that for all $d \geq 4$, the optimization version of the d -MDISA problem is NP-hard to approximate within a factor of 1.1377 [1]. On the other hand, Mahjoub and McCormick showed that even a weighted version of the 3-MDISA problem is solvable in polynomial time, by reducing it to a max flow problem [26]. Hence, taking a weight $w_e = 1$ for each edge $e \in E(G)$ in their construction, we obtain the following.

Theorem 3.2 (Mahjoub and McCormick). *The 3-MDISA problem is solvable in polynomial time.*

Proof of Theorem 3.1. Let (G, k) be an instance of the 3-MDA problem. Note that (G, k) is a yes instance if and only if G has a pair of distinct vertices x, y and a set $F \subseteq E(G)$ such that $|F| \leq k$, $G - F$ is connected and $\text{dist}_{G-F}(x, y) \geq 3$. This last condition is equivalent to the existence of a pair $x, y \in V(G)$ of distinct vertices such that (G, x, y, k) is a yes instance to 3-MDISA, which can be solved in polynomial time by Theorem 3.2. \square

4 Polynomial-time algorithms for the 3-EDA and 3-MEDA problems

In this section we show how to solve the 3-EDA and 3-MEDA problems in polynomial time.

The following weight function on the edges of a graph G will be useful. For each edge $e \in E(G)$, we denote by $w_G(e)$ the minimum length of a cycle in G containing e (or ∞ if no such cycle exists). The weight function w_G is well-defined and can be computed in polynomial time; in fact, for every edge $e = uv \in E(G)$, the value of $w_G(e) - 1$ equals the length of a shortest u, v -path in $G - e$ (or ∞ if no such path exists).

Lemma 4.1. *Let G be a graph with $\text{diam}(G) = 2$. Then, for every edge e of G , we have $w_G(e) \in \{3, 4, 5, \infty\}$. Furthermore, the following holds.*

$$\text{diam}(G - e) \begin{cases} \in \{2, 3\} & \text{if } w_G(e) = 3, \\ = 3 & \text{if } w_G(e) = 4, \\ \geq 4 & \text{if } w_G(e) = 5, \\ = \infty & \text{if } w_G(e) = \infty. \end{cases}$$

Proof. To see the first statement, suppose that e is an edge with $w_G(e) = k \geq 6$ and let $C = (v_1, \dots, v_k, v_1)$ with $e = v_1v_2$ be a shortest cycle containing e . Since v_1 and v_4 are non-adjacent and G is of diameter two, v_1 and v_4 have a common neighbor in G , say z . Observe that $\{v_1, v_2, v_3, v_4, z\}$ induce a cycle containing e shorter than C , a contradiction. Hence, we have $w_G(e) \in \{3, 4, 5, \infty\}$ for every edge e of G .

We prove the second statement. The case when $w_G(e) = 3$ is trivial as the distance of any vertex pair increases by at most one by the deletion of e .

If $w_G(e) = 4$, let $C = (v_1, v_2, v_3, v_4, v_1)$ be a cycle of length four with $e = v_1v_4$. We claim that $\text{diam}(G - e) = 3$. Since $w_G(e) > 3$, vertices v_1 and v_4 have no common neighbor in G and hence $\text{dist}_{G-e}(v_1, v_4) = 3$. Consequently $\text{diam}(G - e) \geq 3$. Suppose for a contradiction that $\text{diam}(G - e) \geq 4$. Then there exists a pair of vertices $x, y \in V(G)$ such that $\text{dist}_{G-e}(x, y) \geq 4 > \text{dist}_G(x, y) = 2$, and observe that all shortest x, y -paths in G must use the deleted edge v_1v_4 . Since x and y are at distance two in G , this is only possible if $|\{x, y\} \cap \{v_1, v_4\}| = 1$, say $x = v_1$ and $y \notin \{v_1, v_4\}$. Since $\{v_2, y\} \cap \{v_1, v_4\} = \emptyset$, we infer that no shortest v_2, y -path in G (which are of length at most two) uses the edge v_1v_4 , which implies $\text{dist}_{G-e}(v_2, y) = \text{dist}_G(v_2, y)$. But then $\text{dist}_{G-e}(x, y) \leq \text{dist}_{G-e}(x, v_2) + \text{dist}_{G-e}(v_2, y) = 1 + \text{dist}_G(v_2, y) \leq 3$, a contradiction.

If $w_G(e) = 5$ for $e = uv$, then we know that $\text{dist}_{G-e}(u, v) = 4$. Consequently, $\text{diam}(G - e) \geq 4$. Lastly, an edge with $w_G(e) = \infty$ is a cut-edge, and thus $G - e$ is disconnected. By definition, $\text{diam}(G - e) = \infty$ in this case. \square

Lemma 4.2. *Let G be a graph with at least four vertices with $\text{diam}(G) = 2$ and girth at most four. Then there exists a set $D \subseteq E(G)$ such that $\text{diam}(G - D) = 3$. Furthermore, such a set D of minimum size can be found in polynomial time.*

Proof. If there exists an edge e whose deletion results in a graph of diameter three, then $\{e\}$ is an optimal solution. We can thus make the following.

Assumption 1. For every edge $e \in E(G)$, we have $\text{diam}(G - e) \neq 3$.

Assumption 1 and Lemma 4.1 imply the following.

Claim 1. The girth of G is three and for every edge $e \in E(G)$ such that $w_G(e) = 3$, $\text{diam}(G - e) = 2$.

Furthermore, Assumption 1 has strong implications on the structure of G .

Claim 2. G is biconnected.

Proof of Claim 2. Suppose that v is a cut-vertex and observe that v is adjacent with every other vertex of G . Since the girth of G is three, there exists an edge xy such that $v \notin \{x, y\}$. Then $G - vx$ has a vertex pair at distance three, namely x and an arbitrary vertex from a connected component of $G - v$ that does not contain x , a contradiction with Assumption 1. \square

Claim 3. For every edge e of G , we have $w_G(e) = 3$.

Proof of Claim 3. By Claim 2, all edge weights are finite. By Assumption 1 and Lemma 4.1, all edge weights are either 3 or 5. Since the graph has girth three, there is an edge with weight 3. Suppose for a contradiction that there exists an edge with weight 5. Then, since G is connected, there exists a pair of edges with different weights sharing a vertex, say $e = uv$ and $f = vz$ where $w_G(e) = 3$ and $w_G(f) = 5$. The definition of weight function w implies that $\text{dist}_G(u, z) = 2$ and $u - v - z$ is the unique shortest u, z -path in G . It follows that the distance in $G - uv$ between u and z is at least three. It is also at most three, since the edge uv is contained in a triangle in G . It follows that $G - uv$ has diameter three, contradicting Assumption 1. \square

Claims 1 and 3 imply the following.

Claim 4. For every edge e of G , we have $\text{diam}(G - e) = 2$.

The following notation will be useful in the rest of the proof. Given a four-vertex path $P = (a, b, c, d)$ in G , we denote by G^P the graph $G - (E(G) \cap \{ac, ad, bd\})$.

Claim 5. For every four-vertex path $P = (a, b, c, d)$ in G , we have $\text{diam}(G^P) \leq 3$.

Proof of Claim 5. Suppose for a contradiction that $\text{diam}(G^P) \geq 4$ for some four-vertex path $P = (a, b, c, d)$. Since the graph G^P is connected, it has finite diameter and thus contains a pair x, y of vertices at distance at least 4. Let $G' = G - (E(G) \cap \{ad\})$. Using Claim 4, we infer that $\text{diam}(G') = 2$, which implies that vertices x and y have a common neighbor z in G' . On the other hand, since $\text{dist}_{G^P}(x, y) \geq 4$, at least one of the edges xz and yz is not present in G^P , that is, at least one of the edges xz and yz belongs to $E(G') \setminus E(G^P)$. However, since $E(G') \setminus E(G^P) \subseteq \{ac, bd\}$, no two edges in $E(G') \setminus E(G^P)$ have an endpoint in common; in particular, not both of the edges xz and yz can belong to $E(G') \setminus E(G^P)$. By symmetry, we may assume that $(E(G') \setminus E(G^P)) \cap \{xz, yz\} = \{xz\} = \{ac\}$ (and thus $yz \in E(G^P)$). Since $x = a$ and $z = c$ or vice versa, b is a common neighbor of x and z in G^P . Thus, (x, b, z, y) is an x, y -path of length 3 in G^P , a contradiction with $\text{dist}_{G^P}(x, y) \geq 4$. \square

A four-vertex path $P = (a, b, c, d)$ in G is said to be *irrelevant* if there exists a vertex $v \in V(G)$ such that $N(a) \cap N(d) = \{v\}$ and $w_{G^P}(va) = w_{G^P}(vd) = 5$, and *relevant* otherwise. The importance of relevant paths is given by the following.

Claim 6. *Suppose that $\text{diam}(G - D) = 3$ for some $D \subseteq E(G)$ and let P be a diametral path in $G - D$. Then P is relevant in G .*

Proof of Claim 6. Let $P = (a, b, c, d)$. Suppose for a contradiction that there exists a vertex $v \in V(G)$ such that $N(a) \cap N(d) = \{v\}$ and $w_{G^P}(va) = w_{G^P}(vd) = 5$. Since vertices a and d are at distance 3 in $G - D$, at least one of the edges va and vd belongs to D . By symmetry, we may assume that $va \in D$. But now, the fact that $w_{G^P}(va) = 5$ implies that $\text{dist}_{G^P - va}(v, a) = 4$, yielding $\text{diam}(G - D) \geq \text{diam}(G^P - va) \geq \text{dist}_{G^P - va}(v, a) = 4$, a contradiction. \square

By Claim 6, it suffices to consider relevant paths. However, we do not yet know if G contains an edge set D whose deletion results in a graph with diameter three. We show that this is indeed the case in the next two claims, by first establishing the existence of a relevant path and then showing that, given any relevant path P , it is possible to delete some edges incident to the vertices of P to increase the diameter of G by exactly one.

Claim 7. *G has a relevant path.*

Proof of Claim 7. First, note that G has a four-vertex path. Indeed, by Claim 1 G has a cycle C of length three. Since G has at least four vertices and it is biconnected by Claim 2, at least one vertex of C has a neighbor outside C , and a four-vertex path in G can be easily found.

Suppose for a contradiction that all four-vertex paths in G are irrelevant. Let $P = (a, b, c, d)$ be an arbitrary four-vertex path in G . Since P is irrelevant, there exists a vertex $v \in V(G)$ such that $N(a) \cap N(d) = \{v\}$ and $w_{G^P}(va) = w_{G^P}(vd) = 5$. Note that $vc \notin E(G)$, since otherwise the existence of a 4-cycle (v, c, b, a, v) in G^P would contradict $w_{G^P}(va) = 5$. A symmetric argument shows that $vb \notin E(G)$. Furthermore, since $w_G(va) = 3$ but $w_{G^P}(va) = 5$, we infer that vertices a and d are adjacent in G . Since the above arguments hold for an arbitrary four-vertex path in G , it follows that G does not contain any induced four-vertex path. In particular, since the path (v, a, b, c) is not induced but $vb \notin E(G)$ and $vc \notin E(G)$, we must have $ac \in E(G)$.

The above arguments hold for an arbitrary four-vertex path in G , and thus we showed that for an arbitrary four-vertex path (x, y, z, w) in G , we have $xz \in E(G)$. But now, applying this argument to the path (v, a, b, c) , we obtain that $vb \in E(G)$, a contradiction. \square

Claim 8. *Let $P = (a, b, c, d)$ be a relevant path in G , let $N_G(a) \cap N_G(d) = \{v_1, \dots, v_k\}$ for some $k \in \mathbb{Z}_{\geq 0}$, and let $F = \cup_{1 \leq i \leq k} \{v_i a, v_i d\}$. Then, there exists an edge set $D \subseteq F \cup (E(G) \cap \{ac, ad, bd\})$ with $|D \cap \{v_i a, v_i d\}| \leq 1$ for each $i \in \{1, \dots, k\}$ such that $\text{diam}(G - D) = 3$. Furthermore, such an edge set can be computed in polynomial time.*

Proof of Claim 8. By Claim 5, we have $\text{diam}(G^P) \leq 3$. If $\text{diam}(G^P) = 3$, then we are done (by taking $D = E(G) \cap \{ac, ad, bd\}$). So we assume that $\text{diam}(G^P) = 2$. Note that this implies that $k \geq 1$. Consider the case when $k = 1$. Since P is relevant, the weights $w_{G^P}(v_1 a)$ and $w_{G^P}(v_1 d)$ cannot both be equal to 5. Suppose that both are at least 5. Then $w_{G^P}(v_1 a) = w_{G^P}(v_1 d) = 5$, since (a, b, c, d, v_1, a) is a cycle of length 5 in G^P containing the edges $v_1 a$ and $v_1 d$, a contradiction. We thus have $\min\{w_{G^P}(v_1 a), w_{G^P}(v_1 d)\} \leq 4$, and we choose e_1 as an edge of minimum weight among $\{v_1 a, v_1 d\}$. Then $\text{dist}_{G^P - e_1}(a, d) = 3$ and it follows that $\text{diam}(G^P - e_1) = 3$ by Lemma 4.1.

Now we consider the case $k \geq 2$. Notice that $w_{G^P}(e) \leq 4$ for every edge e of the form $v_i a$ or $v_i d$. If there exists an edge $f \in F$ with $w_{G^P}(f) = 4$, then $D = \{f\} \cup (E(G) \cap \{ac, ad, bd\})$ is a desired edge set by Lemma 4.1. Therefore, we may assume that $w_{G^P}(e) = 3$ for every edge $e \in F$.

Consider an arbitrary edge set $D \subseteq F$ such that $|D \cap \{v_i a, v_i d\}| \leq 1$ for each i and $\text{diam}(G^P - D) = 2$ (for example, $D = \emptyset$). Obviously, there exists a vertex v_j such that D contains neither $v_j a$ nor $v_j d$. We claim that there exists an edge $e_j \in \{v_j a, v_j d\}$ such that $w_{G^P - D}(e_j) = 3$. Recall that v_j and a must have a common neighbor in G^P by the above assumption. If v_j and a have a common neighbor in $V(G) \setminus \{v_1, \dots, v_k\}$ in the graph G^P , then clearly $w_{G^P - D}(v_j a) = 3$. Otherwise v_j has at least one neighbor v_ℓ in G^P . From $|D \cap \{v_\ell a, v_\ell d\}| \leq 1$, at least one of $\{a, v_j, v_\ell\}$ and $\{d, v_j, v_\ell\}$ forms a triangle in $G^P - D$, thus ensuring the existence of $e_j \in \{v_j a, v_j d\}$ such that $w_{G^P - D}(e_j) = 3$.

We now define a sequence of sets D_0, \dots, D_k such that for all $j \in \{0, 1, \dots, k\}$ we have $D_j \subseteq F$, $|D_j \cap \{v_i a, v_i d\}| \leq 1$ for each i , $\text{diam}(G^P - D_j) = 2$ for all $j < k$, and $\text{diam}(G^P - D_k) = 3$. We set $D_0 = \emptyset$ and for each $j = 1, \dots, k$, we set $D_j = D_{j-1} \cup \{e_j\}$ where v_i is a vertex such that D_{j-1} contains neither $v_i a$ nor $v_i d$ and e_j is an edge in $\{v_i a, v_i d\}$ such that $w_{G^P - D_{j-1}}(e_j) = 3$. Note that such an edge exists by the claim proved in the previous paragraph. Clearly, the final set D_k can be computed in polynomial time and satisfies $|D_k \cap \{v_i a, v_i d\}| = 1$ for each i and $\text{diam}(G^P - D_k) \geq \text{dist}_{G^P - D_k}(a, d) = 3$. Thus, $D_k \cup (E(G) \cap \{ac, ad, bd\})$ is a desired set. This proves Claim 8. \square

For a relevant path $Q = (a, b, c, d)$, let us denote by $f(Q)$ the quantity given by the following expression:

$$f(Q) = |E(G) \cap \{ad, bc, bd\}| + |N_G(a) \cap N_G(d)|.$$

By Claim 8, for every relevant path $Q = (a, b, c, d)$ one can compute an edge set D of size at most $f(Q)$ such that $G - D$ has diameter three.

By Claims 7 and 8, we conclude that there exists an edge set D such that $\text{diam}(G - D) = 3$. On the other hand, for every relevant path Q , any solution D which makes Q a diametral path contains at least $f(Q)$ edges. Furthermore, by Claim 6 any path P^* achieving diameter three after removing an optimal solution D^* is relevant, and an output solution of size at most $f(P^*)$ provided by Claim 8 meets the lower bound and thus is an optimal solution (even though it need not be necessarily identical with D^*). These observations bring forth the following algorithm: we iterate over all four-vertex paths $Q = (a, b, c, d)$ in G and compute a solution of size at most $f(Q)$ whenever Q is relevant. By keeping track of a thus-far best solution, we can compute an optimal solution. Clearly, the algorithm runs in polynomial time. \square

Theorem 4.3. *Let G be a graph. Then, G has a spanning subgraph of diameter three if and only if one of the following conditions holds.*

- (i) $\text{diam}(G) = 1$ (that is, G is a complete graph) and G has at least four vertices.
- (ii) $\text{diam}(G) = 2$ and G has girth at most four.
- (iii) $\text{diam}(G) = 3$.

Proof. Suppose first that G has a spanning subgraph G' such that $\text{diam}(G') = 3$. Since $\text{diam}(G') \geq \text{diam}(G)$, we infer that G has diameter at most three. Furthermore, since G' contains a pair of vertices at distance three, G' (and therefore G) has at least four vertices. If G has diameter one or three, then one of conditions (i) and (iii) holds.

Let us assume now that $\text{diam}(G) = 2$. It remains to show that G has girth at most four. Suppose that the girth of G is at least five. By Lemma 4.1, we have $w_G(e) \in \{5, \infty\}$ for every edge e of G . Then Lemma 4.1 implies that G does not contain any spanning subgraph of diameter three, a contradiction. This establishes the forward implication.

We now prove that the conditions are also sufficient. Let G be a graph satisfying one of conditions (i)–(iii). We will show that in each case, G contains a spanning subgraph G' with diameter three. If condition (i) holds, observe that for every edge $e \in E(G)$, the graph $G - e$ has diameter two and girth three, which reduces to the case of (ii). If condition (ii) holds, that is, G has diameter two and girth at most four, then a desired subgraph G' exists by Lemma 4.2. If condition (iii) holds, then we can take $G' = G$. This completes the proof. \square

Theorem 4.3 implies the following.

Corollary 4.4. *The 3-EDA problem is solvable in polynomial time.*

In case of a yes instance, a smallest set of edges whose deletion results in a graph of diameter 3 can also be computed efficiently.

Theorem 4.5. *The 3-MEDA problem is solvable in polynomial time.*

Proof. If G has diameter three, the solution is trivial (delete nothing). If G has at most three vertices or has diameter at least four, then it is a trivial no instance. We may also assume that G is not complete. Indeed, if G is complete, then we can delete any edge to transform G into an equivalent graph G' of diameter two and girth three. Solving the resulting instance G' optimally will yield an optimal solution for G . Clearly, whether G satisfies one of the above conditions (and if so, which one) can be checked in polynomial time. Therefore, we may assume that G has at least four vertices and diameter two. Moreover, we may assume that G has girth at most four, since otherwise G is trivially a no instance by Lemma 4.1. Now, applying Lemma 4.2 we can compute an optimal solution. \square

The result of Theorem 4.3 connects our study to a well-known concept in extremal and algebraic graph theory, the Moore graphs. A connected graph with diameter d is either acyclic or has girth at most $2d+1$. A connected graph G with diameter d and girth exactly $2d+1$ is said to be a *Moore graph* (see [32]). A Moore graph is necessarily regular (that is, all vertices have the same degree), and, apart from odd cycles and complete graphs, all Moore graphs have diameter two. Furthermore, it is known that every Moore graph G with diameter two is k -regular with $|V(G)| = k^2 + 1$ for some $k \in \{2, 3, 7, 57\}$. For each $k \in \{2, 3, 7\}$, there is a unique k -regular Moore graph with diameter two, namely the 5-cycle (for $k = 2$), the (10-vertex) Petersen graph (for $k = 3$), and the (50-vertex) Hoffman-Singleton graph (for $k = 7$); see, e.g., [6]. The existence of a (3250-vertex) 57-regular Moore graph is a famous open problem in graph theory (see [11] for a recent survey). Theorem 4.3 implies the following characterization of graphs with diameter two that are yes / no instances to the 3-EDA problem.

Corollary 4.6. *For every graph G with diameter two, the following holds.*

1. *If G has girth at most four, then G has a spanning subgraph of diameter three.*
2. *If G has girth at least five, then G is either a Moore graph or G is acyclic (in which case G is isomorphic to a star $K_{1,n}$ for some $n \geq 2$); in either case, G does not have any spanning subgraph of diameter three.*

5 NP-completeness results for d -MEDA and d -MDA for $d \geq 5$

We prove that 5-MEDA is NP-complete for graphs with diameter three and four. Then we generalize this result to d -MEDA for all $d \geq 6$. Our approach also proves NP-completeness of d -MDA for all $d \geq 5$.

Theorem 5.1. *The 5-MEDA problem is NP-complete for graphs of diameter 3.*

Proof. Since computing the diameter of a graph can be done in polynomial time, 5-MEDA is in NP.

Our reduction is from the NP-complete VERTEX COVER problem (see [16]). A *vertex cover* in a graph G is a set of vertices intersecting every edge.

VERTEX COVER

Input: A graph $\Gamma = (W, E')$, an integer c .

Question: Does Γ have a vertex cover of size at most c ?

Given an instance $(\Gamma = (W, E'), c)$ of VERTEX COVER we create an instance (G, k) of the 5-MEDA problem such that $G = (V, E)$ is of diameter 3 as follows.

We set $k = |W| + c$.

For each vertex $v \in W$ we create a path $v_1 - v_2 - v_3 - v_4$.

For each edge $e = uv \in E'$ we create two paths $v_2 - e_1 - u_3$ and $v_3 - e_2 - u_2$.

There is a vertex s with the edges sv_1 and sv_2 for every $v \in W$.

There is a vertex t with the edges tv_3 and tv_4 for every $v \in W$.

Finally, there are K^1 and K^2 , two complete graphs each with $k + 1$ vertices. For every vertex $u \in K^1$ and each $v \in W$ there are the edges uv_1 and uv_2 , and there is an edge ue_1 for every $e \in E'$. For every vertex $u \in K^2$ and each $v \in W$ there are the edges uv_3 and uv_4 , and there is an edge ue_2 for every $e \in E'$. There is an edge between any pair of vertices of $K^1 \cup K^2$.

There are no other vertices or edges in G .

Clearly (G, k) can be computed from (Γ, c) in polynomial time. Furthermore, it can be verified that G has diameter 3.

Suppose that (Γ, c) is a yes instance of VERTEX COVER. Let $S \subseteq W$ be a vertex cover in Γ of size $|S| \leq c$. We build $F \subseteq E$ as follows: for each $v \in S$ we put the edges sv_2 and tv_3 in F ; for each $v \in W \setminus S$ we put edge v_2v_3 in F . Thus $|F| \leq k$. We show that $\text{dist}_{G-F}(s, t) = 5$. First, $s - v_1 - z_1 - z_2 - v_4 - t$ where $v \in W$, $z_1 \in K^1$ and $z_2 \in K^2$, is a path of length 5 in $G - F$. Now let P be a shortest s, t -path in $G - F$. Since $\text{dist}_G(s, t) = 3$ the length of P is at least 3. If P has length less than 5 it cannot contain a vertex of $K^1 \cup K^2$. Assume that the length of P is 3. Then $P = s - v_2 - v_3 - t$ for some $v \in W$. If $v \in S$ then $sv_2 \in F$ and P cannot be a path in $G - F$. If $v \notin S$ then $v_2v_3 \in F$ and P cannot be a path in $G - F$. So P has length at least 4. Assume that the length of P is 4. Suppose that the second vertex in P is v_1 for some $v \in W$. Then $P = s - v_1 - v_2 - v_3 - t$, which is impossible since either $v \in S$ and $v_3t \in F$ or $v \notin S$ and $v_2v_3 \in F$. So the second vertex of P is v_2 for some $v \in W$. Therefore $sv_2 \notin F$, which implies that $v \notin S$, and thus $v_2v_3 \in F$. Thus $P = s - v_2 - e_i - w_3 - t$ for some $e = vw \in E'$ and some $i \in \{1, 2\}$. Since $v \notin S$ and S is a vertex cover of Γ , we have $w \in S$. But $w_3t \in F$, a contradiction. We conclude that the length of P is 5. Now, for any pair of vertices $\{x, y\} \neq \{s, t\}$ we have $\text{dist}_{G-F}(x, y) \leq 4$ (passing through the vertices of $K^1 \cup K^2$). Thus any shortest path that is not an s, t -path has a length no more than 4. Hence the diameter of $G - F$ is 5 and (G, k) is a yes instance of 5-MEDA.

Suppose that (G, k) is a yes instance of 5-MEDA. So there exists a set $F \subseteq E$, $|F| \leq k$, such that the diameter of $G - F$ is 5. It can be verified that for any two distinct vertices $x, y \in V(G)$ such that $\{x, y\} \neq \{s, t\}$, there are at least $k + 1$ edge-disjoint x, y -paths of length at most four in G (using vertices of $K^1 \cup K^2$). Therefore, $\text{dist}_{G-F}(x, y) \leq 4$. This implies that a shortest path of length 5 in $G - F$ is an s, t -path. Note that each s, t -path in G with a vertex in $K^1 \cup K^2$ has length at least five. We may thus assume that F does not contain any edge incident to a vertex of $K^1 \cup K^2$.

The only s, t -paths of length 3 in G are of the form $s - v_2 - v_3 - t$ for $v \in W$. They are pairwise edge-disjoint, so at least $|W|$ of their edges are in F , at least one per path. The s, t -paths of length 4 in G are paths $s - v_1 - v_2 - v_3 - t$ or $s - v_2 - v_3 - v_4 - t$ for $v \in W$, and paths $s - v_2 - e_i - w_3 - t$ for $e = vw \in E'$ and $i \in \{1, 2\}$. All the other s, t -paths of G have length more than 4.

Suppose that F contains an edge of the form v_2e_i . This edge cannot be in an s, t -path of length 3. Since the path $s - v_2 - e_i - u_3 - t$ is a unique s, t -path of length 4 that contains this edge, we can remove v_2e_i from F and add instead the edge sv_2 (if sv_2 is not already in F). So from now, by symmetry, we may assume that for every $v \in W$, the edges v_2e_i and e_iv_3 are not in F .

Suppose that F contains an edge of the form v_1v_2 . Since there is exactly one s, t -path of length 4 that contains this edge, the path $s - v_1 - v_2 - v_3 - t$, and the edge v_1v_2 cannot be in an s, t -path of length 3, we can remove v_1v_2 from F and add instead the edge v_2v_3 (if v_2v_3 is not already in F). The same argument can be used with the edge sv_1 . So from now, by symmetry, we may assume that each edge in F is of the form sv_2 , v_2v_3 , or v_3t for some $v \in W$.

Since for every $v \in W$, set F contains at least one of the edges of the path $s - v_1 - v_2 - v_3 - t$, we infer that F contains at least one of the edges v_2v_3 and v_3t , for every $v \in W$. By symmetry, F also contains at least one of the edges sv_2 and v_2v_3 for every $v \in W$. Thus, if for some $v \in W$ we have $v_2v_3 \notin F$, then $sv_2 \in F$ and $v_3t \in F$.

Furthermore, if for some $v \in W$ we have $sv_2 \in F$ and $v_3t \in F$, then we may assume that $v_2v_3 \notin F$. Indeed, if $\{sv_2, v_2v_3, v_3t\} \subseteq F$ for some $v \in W$, then we can replace F with $F \setminus \{v_2v_3\}$: adding the edge v_2v_3 to the graph $G - F$ cannot create an s, t -path of length 4 or less since $sv_2 \in F$ and $v_3t \in F$.

From the above we infer that F has the following structure. For each $v \in W$ either $v_2v_3 \in F$ or $sv_2, v_3t \in F$. Since $|F| \leq k = |W| + c$, at most c pairs sv_2, v_3t are in F . From F we build a set $S \subseteq W$ as follows. We take $v \in S$ if and only if $sv_2, v_3t \in F$. So $|S| \leq c$. For each edge $uv \in E'$ we have $sv_2, v_3t \in F$ or $su_2, u_3t \in F$ else $s - v_2 - e_i - u_3 - t$ for some $i \in \{1, 2\}$ would be an s, t -path in $G - F$ of length 4. Thus S is a vertex cover of Γ and (Γ, c) is a yes instance. \square

Theorem 5.2. *The 5-MEDA problem is NP-complete for graphs of diameter 4.*

Proof. The proof is the same as the proof of Theorem 5.1 except for the following. Instead of two complete subgraphs K^1, K^2 , here there are four pairwise disjoint complete subgraphs K^1, K^2, K^3, K^4 with $k + 1$ vertices each. All the vertices of K^3 are adjacent to all the vertices of $K^1 \cup K^4$. All the vertices of K^4 are adjacent to all the vertices of $K^2 \cup K^3$. The vertices of K^1 and K^2 are linked to the rest of G as previously. The so obtained graph now has diameter 4 (the diametral paths are from $u \in K^3$ to t and from $v \in K^4$ to s). \square

Lemma 5.3. *For any fixed $d \geq 5$ the $(d+1)$ -MEDA problem is NP-complete for graphs of diameter d .*

Proof. Given an instance (G, p) of the 5-MEDA problem such that $G = (V, E)$ is of diameter 4 we create an instance (G', p') of $(d + 1)$ -MEDA problem such that $G' = (V', E')$ is of diameter d as

follows. We take $p = p'$. Let $\{s, t\}$ be a diametral pair of G . Let $d' = d - 4$. We add to G a path with (new) vertices $z_1, z_2, \dots, z_{d'}$ and an edge tz_1 . Let $G = (V, E)$ be a yes instance of 5-MEDA. Let $F \subseteq E, |F| \leq p$ be such that $G - F$ has diameter 5. Then $G' - F$ has diameter $5 + d' = d + 1$. Now, let $G' = (V', E')$ be a yes instance of $(d + 1)$ -MEDA and $F \subseteq E', |F| \leq p' = p$ be such that $G' - F$ has diameter $d + 1$. Since $G' - F$ is connected, $F \subseteq E$ and $G - F$ is connected. It follows from the proof of Theorem 5.1 (which is the basis for the proof of Theorem 5.2) that for every two vertices x, y in G such that $\{x, y\} \neq \{s, t\}$, we have $\text{dist}_{G'-F}(x, y) = \text{dist}_{G-F}(x, y) \leq 4$. This implies that vertices s and $z_{d'}$ form the only diametral pair of G' and thus $G - F$ has diameter $d + 1 - d' = d + 1 - (d - 4) = 5$. \square

Clearly, Theorem 5.2 and Lemma 5.3 have the following consequence.

Corollary 5.4. *For every $d \geq 5$, the d -MEDA problem is NP-complete.*

We now show how to further adapt the above constructions to prove that for any fixed $d \geq 5$ and $q \geq 2$, $(d + q)$ -MEDA is NP-complete for graphs of diameter d . The possible values of q for which we obtain NP-completeness depend on the diameter d of the target graph (or, equivalently, on the diameter of the input graph); see Theorem 5.8 for the precise statement.

Lemma 5.5. *For any fixed $d \geq 4$ and $\delta \geq 1$, the following restriction of the $(d + 1)$ -MEDA problem is NP-complete:*

Input: A connected graph G with diameter d and a positive integer p such that there is no set $F \subseteq E$ with fewer than p edges such that $G - F$ has diameter $d + 1$ and there is no set $F' \subseteq E$ such that $|F'| \leq p + \delta$ and $G - F'$ has diameter at least $d + 2$.

Question: Is there a set $F \subseteq E(G)$ such that $|F| \leq p$ and $G - F$ is a connected graph of diameter $d + 1$?

Proof. We showed in the proof of Theorem 5.1 that the 5-MEDA problem remains NP-complete for instances (G, k) such that $\text{diam}(G) = 3$. The proof was based on a reduction transforming an instance (Γ, c) of VERTEX COVER to an instance (G, k) of the 5-MEDA problem such that $\text{diam}(G) = 3$. Without loss of generality, we may assume that the instance (Γ, c) of VERTEX COVER is an instance obtained from Karp's reduction from 3-SAT to VERTEX COVER [20], in which case c vertices are necessary to cover all the edges. This in turn implies that there is no set $F \subseteq E(G)$ with fewer than k edges such that $G - F$ has diameter 5 (since otherwise one could use the arguments from the proof Theorem 5.1 to infer that Γ has a vertex cover of size less than c). Hence through the proof of Theorem 5.2 the same holds for an instance (G, k) of 5-MEDA when G has diameter 4. From such an instance (G, k) we build (G_δ, p_δ) , another instance of 5-MEDA where G_δ has diameter 4. We take $p_\delta = (\delta + 1)k$. We take $\delta + 1$ copies of the subgraph of G induced by the all vertices v_i and e_j . Here the four cliques K^1, K^2, K^3, K^4 have size $\delta + 1$ times the size they have in the proof of Theorem 5.2. The copies of the vertices v_i, e_j are connected in the same manner as above to the four cliques K^1, K^2, K^3, K^4 . There are two vertices s, t connected to the copies of the vertices v_j as before. Hence the $\delta + 1$ copies of G are pairwise independent in G_δ . Note that G_δ has diameter 4. As noted in the proof of Theorem 5.1, only the edges between the vertices s, t, v_i, e_j can be put inside an edge set F whose removal increases the diameter of the graph by one unit. Thus there is no set $F \subseteq E(G_\delta)$ with fewer than $p_\delta = (\delta + 1)k$ edges such that $G_\delta - F$ has diameter 5. Furthermore, there is no set $F' \subseteq E(G_\delta)$ with at most $p_\delta + \delta$ edges such that $G_\delta - F'$ has diameter more than 5, since the $\delta + 1$ copies of G are pairwise independent. With the same

construction as in the proof of Lemma 5.3 the same conclusion can be derived for instances (G_δ, p_δ) of $(d+1)$ -MEDA such that G_δ has diameter d . \square

Lemma 5.6. *For every fixed d and q , $d \geq 4, q \geq 1$, the $(d+2q+1)$ -MEDA problem is NP-complete for graphs of diameter $d+q$.*

Proof. Given an instance (G, p) of the $(d+1)$ -MEDA problem such that $G = (V, E)$ is of diameter d and satisfies the condition of Lemma 5.5 with $\delta = q$, we create an instance (G', p') of $(d+2q+1)$ -MEDA problem such that $G' = (V', E')$ is of diameter $d+q$ as follows. Let $p' = p+q$. Let $\{s, t\}$ be a diametral pair of G . We add to G a path with (new) vertices z_1, z_2, \dots, z_q and an edge tz_1 . We denote $t = z_0$. For each $i \in \{1, \dots, q\}$ we add a new vertex r_i with the two edges $z_{i-1}r_i, r_iz_i$. Let us denote by T_i the edge set of the triangle containing r_i . Clearly G' has diameter $d+q$. Let (G, p) be a yes instance of $(d+1)$ -MEDA. Let $F \subseteq E, |F| \leq p$ be such that $G - F$ has diameter $d+1$. Let $F' = F \cup \{z_0z_1\} \cup \dots \cup \{z_{q-1}z_q\}$. Then $|F'| = p+q \leq p'$ and $G' - F'$ has diameter $d+1+2q$. Now, let $(G' = (V', E'), p')$ be a yes instance of $(d+2q+1)$ -MEDA and $F' \subseteq E', |F'| \leq p'$ be such that $G' - F'$ has diameter $d+2q+1$. Since (G, p) satisfies the condition of Lemma 5.5, the diameter d of G cannot be increased to $d+k$ where $k \geq 2$, by deleting only $p+q = p'$ edges. The graph $G' - F'$ is connected, so $|F' \cap T_i| \leq 1$ whenever $1 \leq i \leq q$. This implies that $F' \cap T_i = \{z_{i-1}z_i\}$ for all $i \in \{1, \dots, q\}$ and $|F' \cap E| \leq p$. Thus $F = F' \cap E$ is such that $G - F$ has diameter $d+1$. \square

Lemma 5.7. *For every fixed d and q , $d \geq 3, q \geq 1$, with $d+q \geq 5$ the $(d+2q+2)$ -MEDA problem is NP-complete for graphs of diameter $d+q$.*

Proof. From Theorem 5.1 we know that the 5-MEDA problem is NP-complete for graphs of diameter 3. Using similar arguments as in the proof of Lemma 5.3 the $(d+2)$ -MEDA problem is NP-complete for graphs of diameter d . Now with similar arguments as those used in the proofs of Lemmas 5.5 and 5.6 the proof is complete. \square

Using Lemmas 5.3, 5.6, and 5.7 we now derive the following general result.

Theorem 5.8. *For every fixed $d \geq 5$ and q such that $1 \leq q \leq d-1$, the $(d+q)$ -MEDA problem is NP-complete for graphs of diameter d .*

Proof. If $q = 1$, then we use Lemma 5.3.

For $q = 2$, let $d' = d-1$ and $q' = 1$. Then $d' \geq 4$ and applying Lemma 5.6 with d' and q' in place of d and q , respectively, we have that $(d'+3)$ -MEDA is NP-complete for graphs of diameter $d'+1$. Since $d = d'+1$ and $d+2 = d'+3$ the result follows.

For $q \geq 3$, let $d' = d-q+2$ and $q' = q-2$. Since $d \geq q+1$, we have $d' \geq 3$. Since $q \geq 3$ and $d \geq 5$, we also have $q' \geq 1$ and $d'+q' \geq 5$. From Lemma 5.7 we have that $(d'+2q'+2)$ -MEDA is NP-complete for graphs of diameter $d'+q'$. Since $d = d'+q'$ and $d+q = d'+2q'+2$, the result follows. \square

Our approach used in the proofs of Theorems 5.1 and 5.8 (including Lemma 5.5 and its proof) actually shows that analogous results hold also for the MINIMUM DIAMETER- $(d+q)$ AUGMENTATION problems.

Theorem 5.9. *The 5-MDA problem is NP-complete for graphs of diameter 3.*

Theorem 5.10. *For every fixed $d \geq 5$ and q such that $1 \leq q \leq d-1$, the $(d+q)$ -MDA problem is NP-complete for graphs of diameter d .*

6 Conclusion and future works

We determined the computational complexity of several problems related to the blocker problem for the diameter of a graph. In particular, we summarize in Table 1 the complexity results for the $(d+q)$ -MEDA and $(d+q)$ -MDA problems when restricted to graphs of diameter d for various values of $d \geq 1$ and $q \geq 1$.

$d \setminus q$	1	2	3	4	5	6	7
1	P (T 2.5)	P (T 2.5)	P (T 2.5)	P (T 2.5)	P (T 2.5)	P (T 2.5)	P (T 2.5)
2	P (T 4.5)	?	?	?	?	?	?
3	?	NP-c (T 5.1)	?	?	?	?	?
4	NP-c (T 5.2)	?	?	?	?	?	?
5	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	?	?	?
6	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	?	?
7	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	?
8	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)	NP-c (T 5.8)

Table 1: Complexity of the MINIMUM EXACT DIAMETER- $(d+q)$ AUGMENTATION problem when restricted to graphs of diameter d , for various values of $d \geq 1$ and $q \geq 1$. Rows are indexed by the value d of the diameter of the input graph. Columns are indexed by the desired value q of the diameter increase. P denotes that the problem is solvable in polynomial time, NP-c that it is NP-complete, and a question mark that the complexity is open. Next to each entry, we give a reference for the corresponding statement (T stands for Theorem). The same results and open questions hold also for the MINIMUM DIAMETER- $(d+q)$ AUGMENTATION problem when restricted to graphs of diameter d , for various values of $d \geq 1$ and $q \geq 1$ (see Theorems 2.5, 3.1, 5.9 and 5.10).

In conclusion, we list some related open problems and research directions that seem interesting for future consideration. The d -MEDA and d -MDA problems are polynomial for $d = 3$ and NP-complete for all $d \geq 5$. Thus, to obtain a dichotomy, two open problems remain, the complexity of 4-MEDA and 4-MDA. It would be natural to also investigate the d -EDA problem for $d \geq 4$. In the cases of NP-complete constant diameter augmentation problems, it would be interesting to understand restrictions on the input graphs under which the problems become tractable, and perform a more detailed complexity analysis of the problems, for example with respect to parameterized complexity, existence of polynomial kernels, and approximability properties of the corresponding optimization problems.

From the structural point of view, an interesting question related to the EXACT DIAMETER- d AUGMENTATION problems is the following. What are the graphs G such that the set of diameter values of the spanning connected subgraphs of G forms an interval of consecutive integers? This class of graphs generalizes trees and complete graphs. Examples of graphs not in the class include Moore graphs of diameter two (for example, the Petersen graph).

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