

#### Efficient Storage of Images onto DNA Using Vector Quantization

Melpomeni Dimopoulou, Marc Antonini

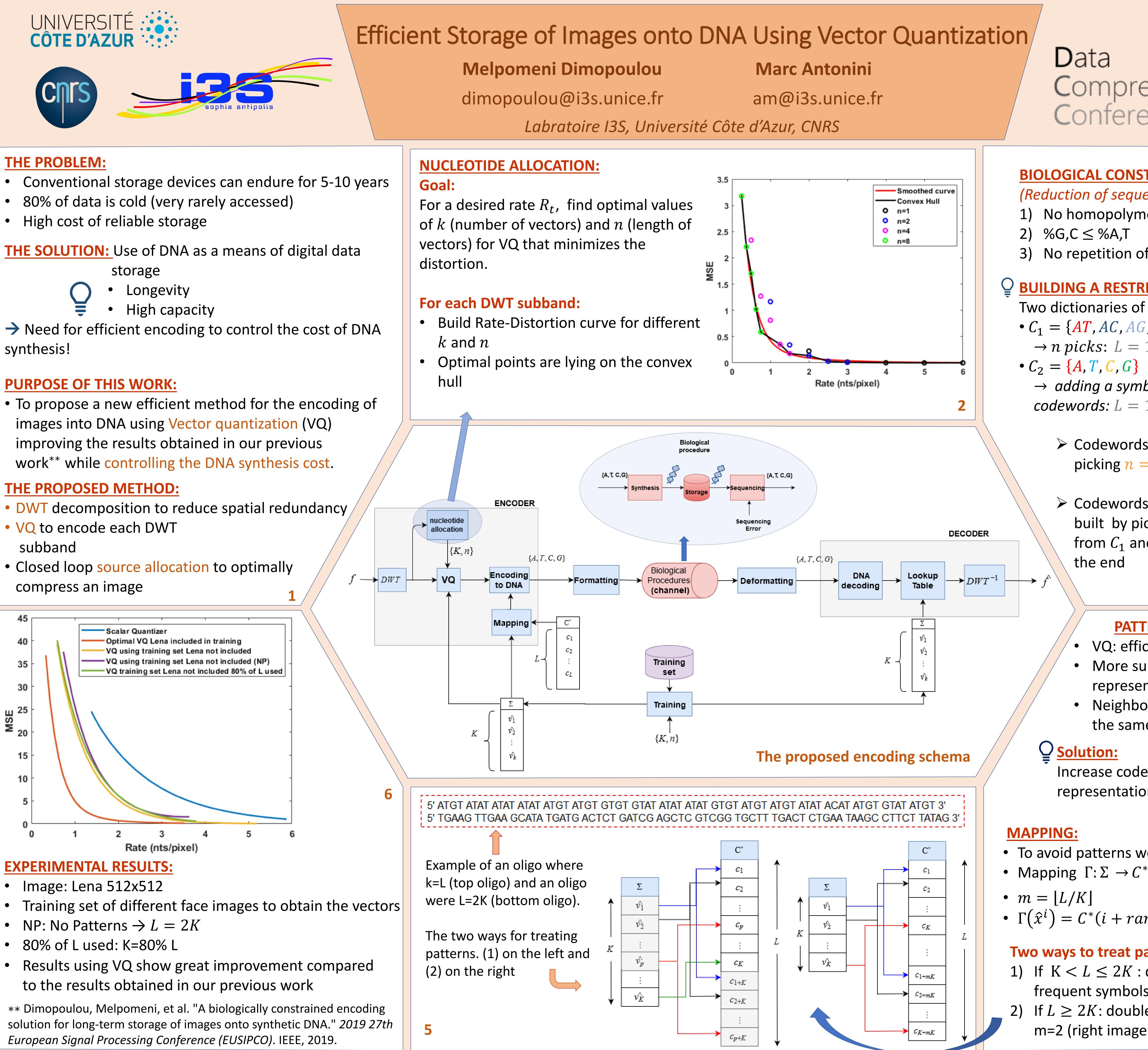
#### ► To cite this version:

Melpomeni Dimopoulou, Marc Antonini. Efficient Storage of Images onto DNA Using Vector Quantization. Data Compression Conference (DCC) 2020, Mar 2020, Utah, United States. 10.1109/DCC47342.2020.00085. hal-02549709

#### HAL Id: hal-02549709 https://hal.science/hal-02549709

Submitted on 21 Apr 2020  $\,$ 

**HAL** is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers. L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.



## **THE PROBLEM:**

- 80% of data is cold (very rarely accessed)
- High cost of reliable storage

**THE SOLUTION:** Use of DNA as a means of digital data

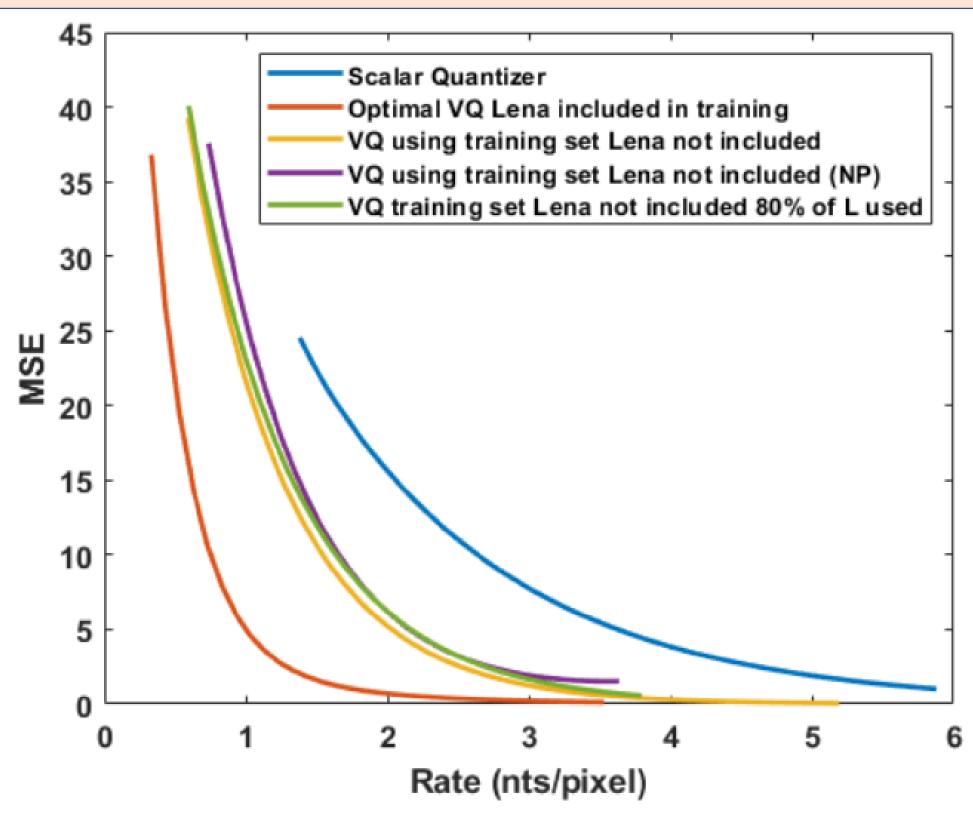
synthesis!

# **PURPOSE OF THIS WORK:**

images into DNA using Vector quantization (VQ) improving the results obtained in our previous work<sup>\*\*</sup> while controlling the DNA synthesis cost.

# **THE PROPOSED METHOD:**

- DWT decomposition to reduce spatial redundancy
- VQ to encode each DWT subband
- Closed loop source allocation to optimally compress an image



### **EXPERIMENTAL RESULTS:**

- Image: Lena 512x512
- NP: No Patterns  $\rightarrow L = 2K$
- 80% of L used: K=80% L

*European Signal Processing Conference (EUSIPCO)*. IEEE, 2019.

Compression Conterence



**BIOLOGICAL CONSTRAINTS ON THE ENCODING:** (Reduction of sequencing error) 1) No homopolymers 3) No repetition of short patterns

**BUILDING A RESTRICTED QUATERNARY CODE:** 

Two dictionaries of pair elements: •  $C_1 = \{AT, AC, AG, TA, TC, TG, CA, CT, GA, GT\}$  $\rightarrow n \ picks: L = 10^n \ codewords \ of \ length \ l=n*2$  $\rightarrow$  adding a symbol from  $C_2$  at the end of codewords:  $L = 10^n * 4$ 

> $\succ$  Codewords of an <u>even length</u> *l* are built by picking n = l/2 pair-elements from  $C_1$

 $\succ$  Codewords of odd length *l* codewords are built by picking n = (l - 1)/2 pair-elements from  $C_1$  and adding a pair element from  $C_2$  at the end

## **PATTERN REPETITIONS:**

• VQ: efficient for compression

- More subband coefficients will be
- represented by the same vector
- Neighboring coefficients will be encoded to
- the same codeword  $\rightarrow$  pattern repetitions!

Increase code size L to allow double representation



• To avoid patterns we need:  $L \ge K$ 

•  $\Gamma(\hat{x}^i) = C^*(i + rand(0, m - 1) * K)$ 

## Two ways to treat pattern repetition:

1) If  $K < L \leq 2K$ : double representation of most frequent symbols , m=1 (left image) 2) If  $L \ge 2K$ : double representation of every word, m=2 (right image)