



ISKO-Maghreb



جامعة بجاية
Tasdawit n Bgayet
Université de Béjaïa



ISKO MAGHREB 2018

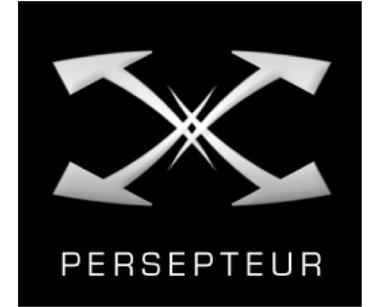
The concept of a Pseudo-polygon and its application to Big Data and IoT

Ahcène Bounceur

November 25th, 2018



ANR Project: PERSEPTEUR



- 3D Smart City and IoT WSN simulator (SCI-WSN)
- 3D Signal propagation and interferences
- Physical and spatial modeling of the radio channel
- Mobility and dynamicity
- Visualization and optimization of the radio frequency space
- Designing reliable networks (QoS, test and Security)

Partners



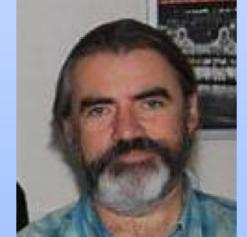
Laurent Clavier (IEMN)



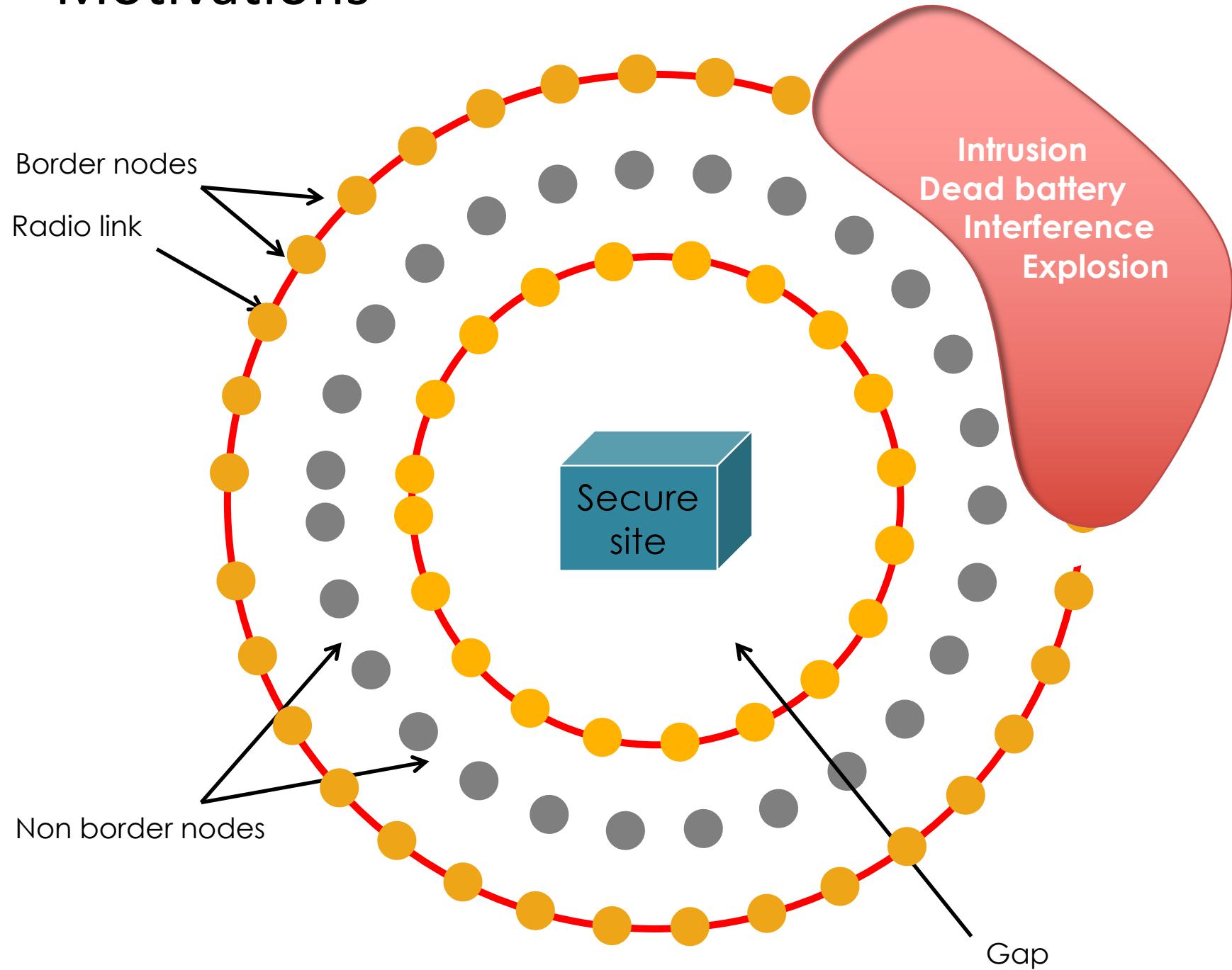
Pierre Combeau (Xlim)



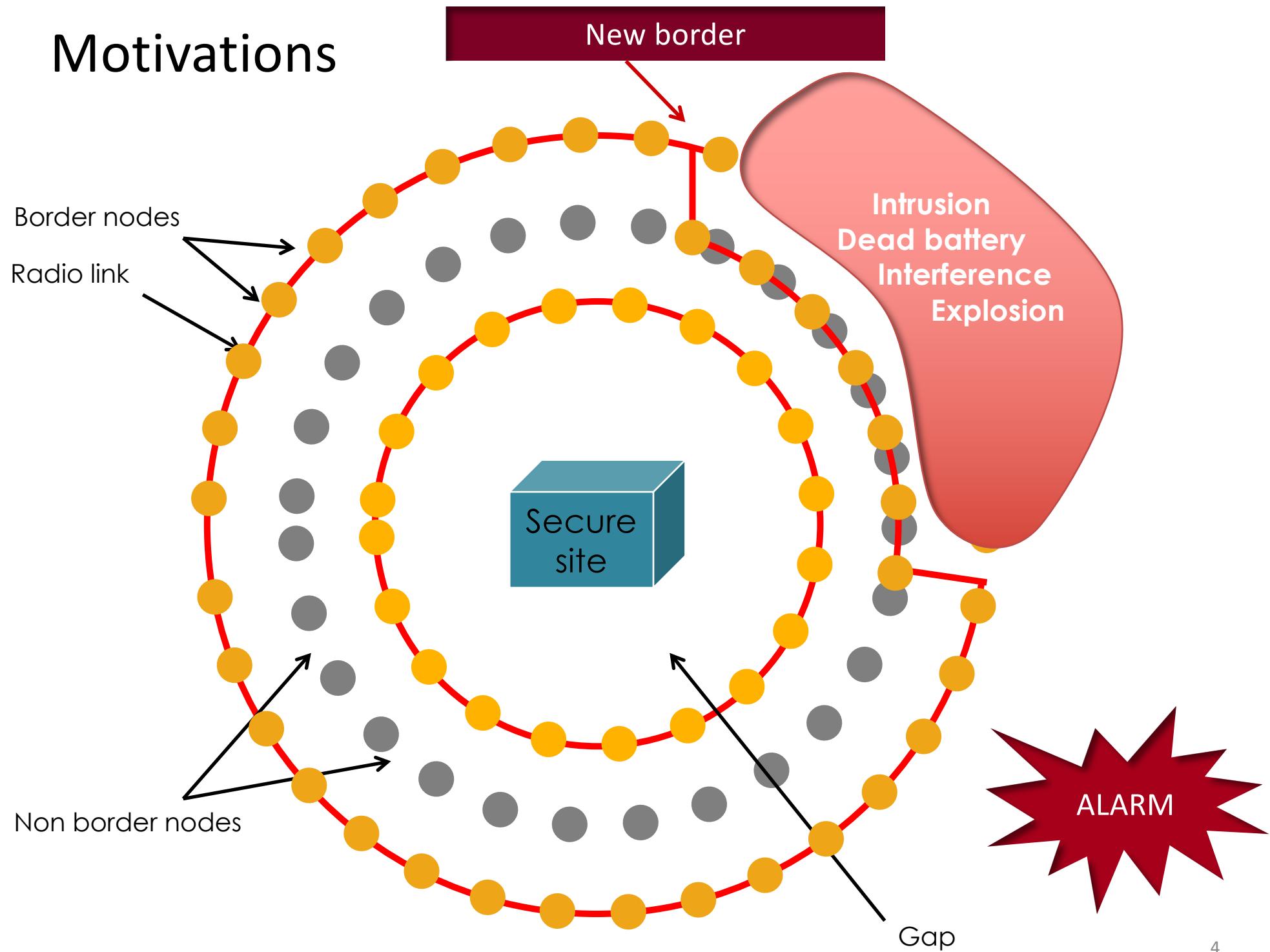
Olivier Marc (Virtualys)
Virtual Vision & Process



Motivations



Motivations



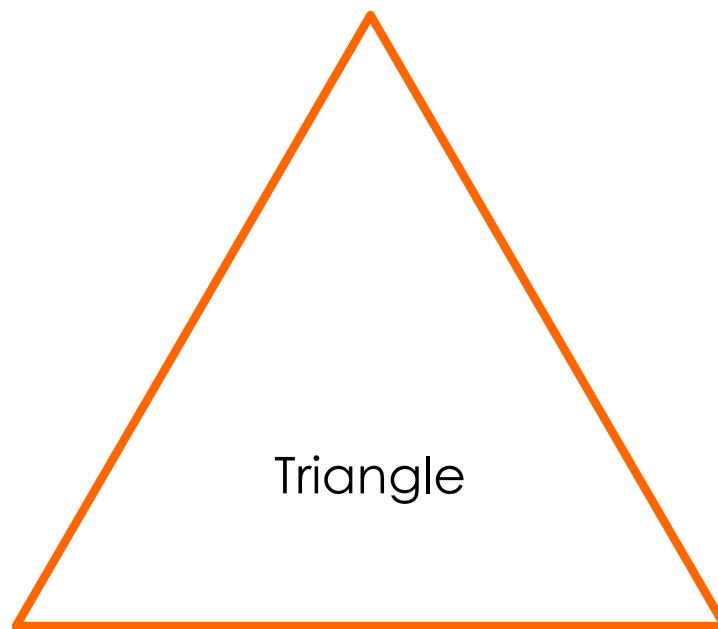
Outline

- Part 1: Polygons and Pseudo-Polygons
- Part 2: CupCarbon for IoT Simulation
- Part 3: Polygon Hulls and Leader Election
- Part 4: Some applications
- Part 5: Conclusion and Future work

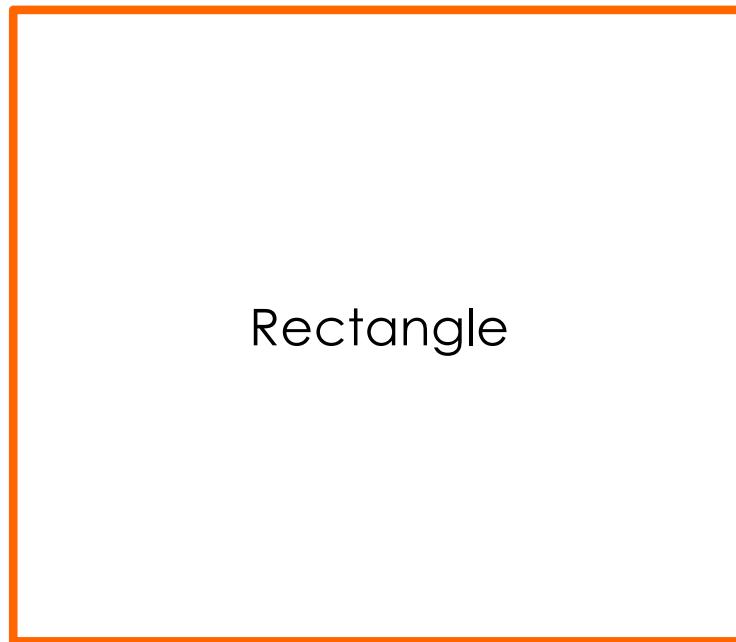
1

POLYGONS AND PSEUDO-POLYGONS

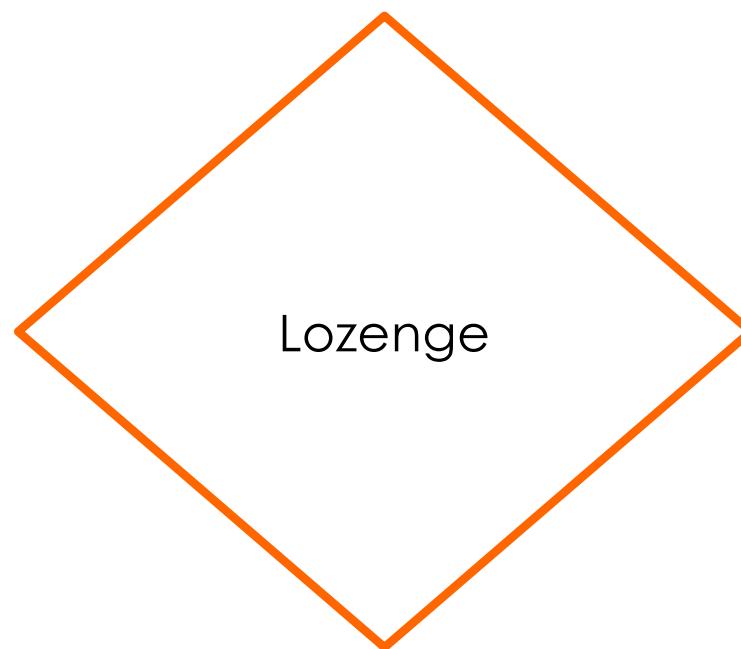
What is a Polygon?



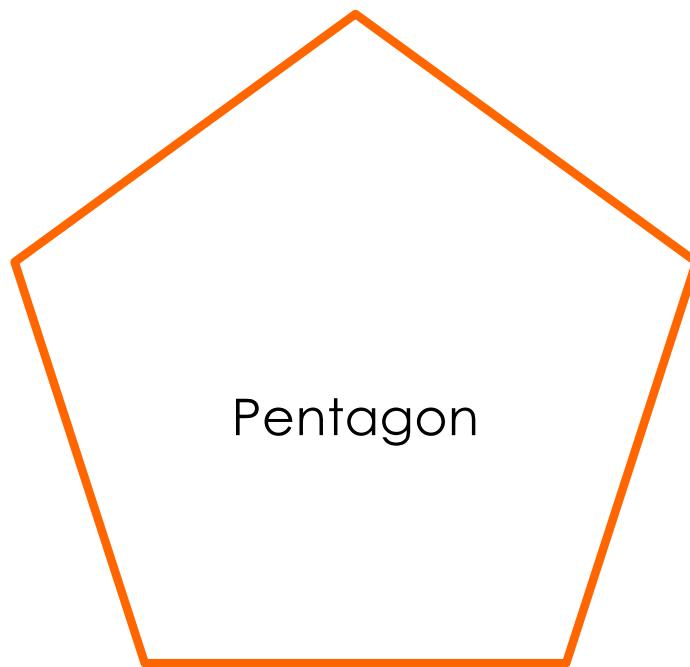
What is a Polygon?



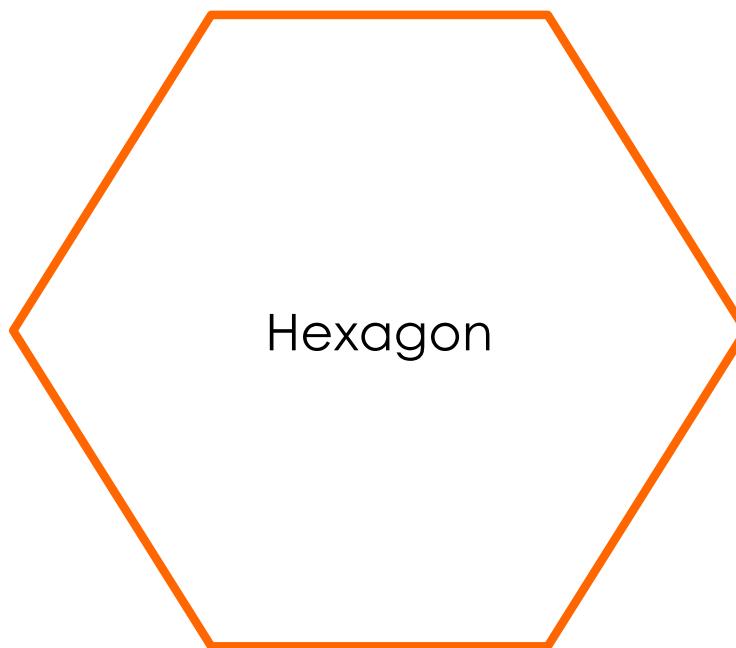
What is a Polygon?



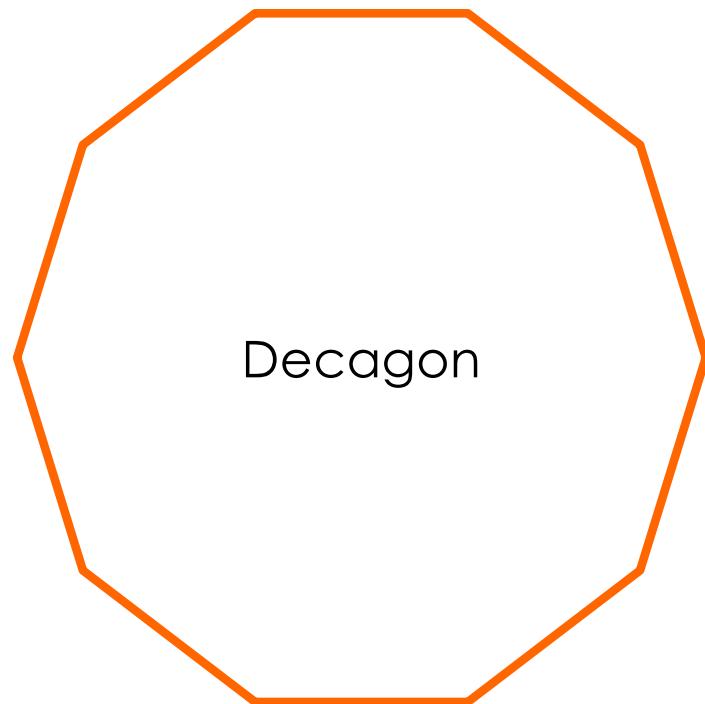
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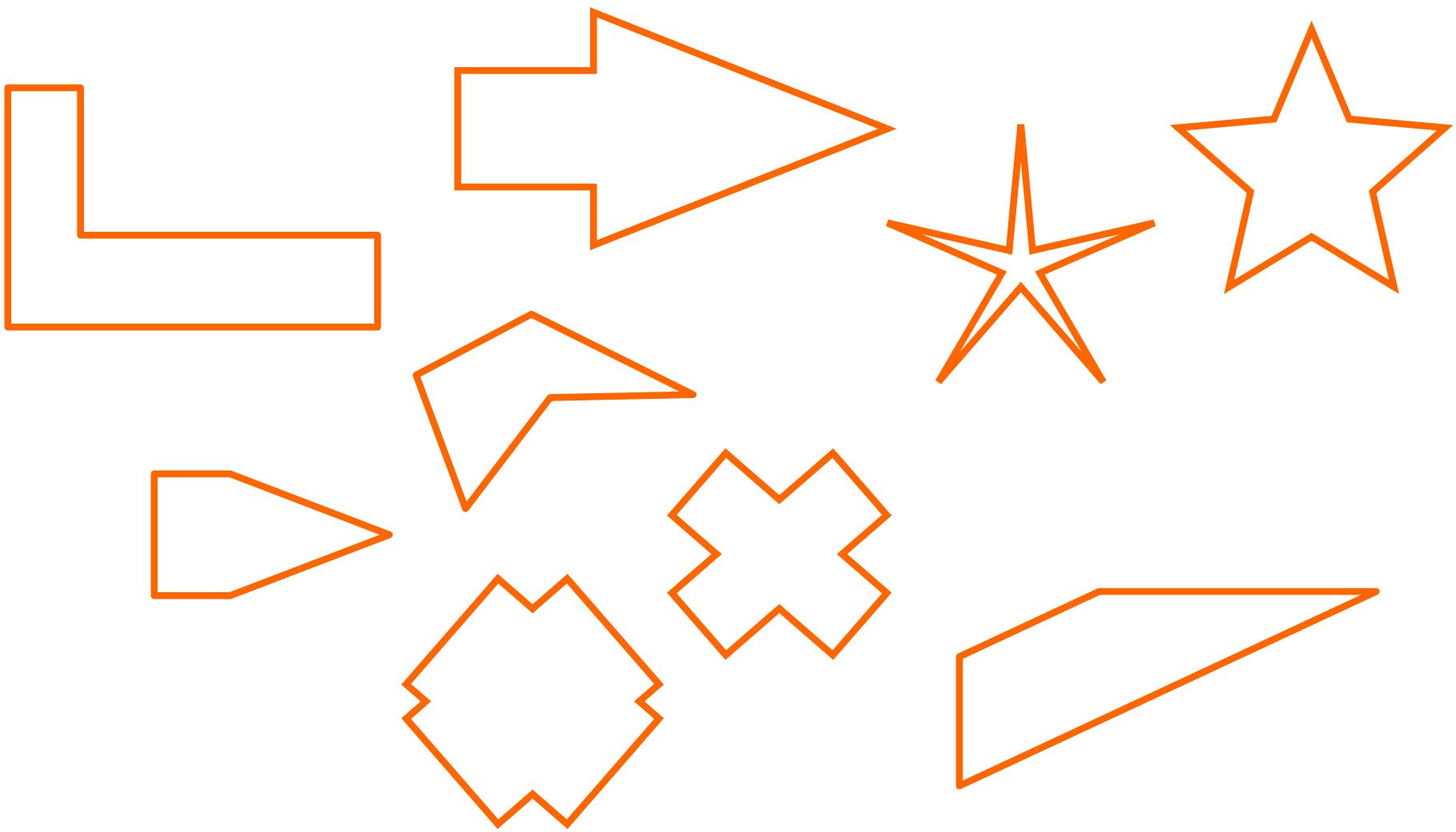
What is a Polygon?



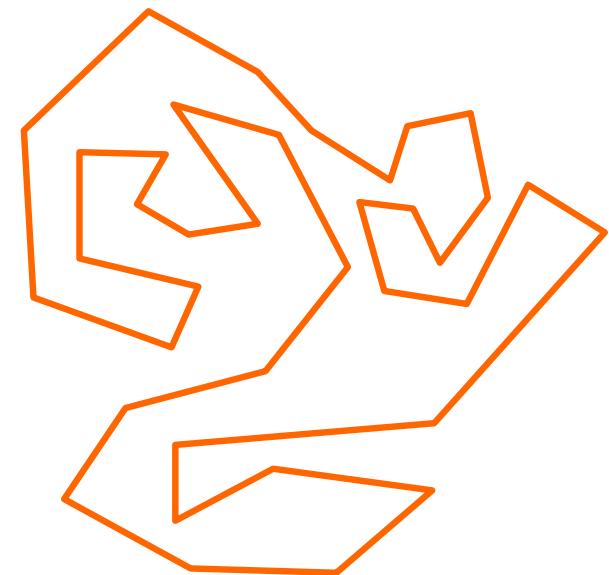
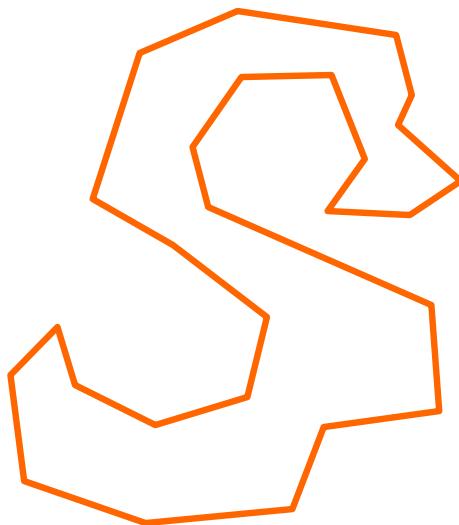
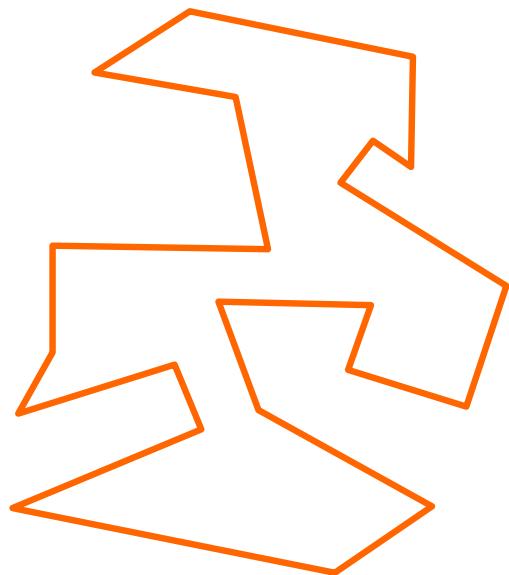
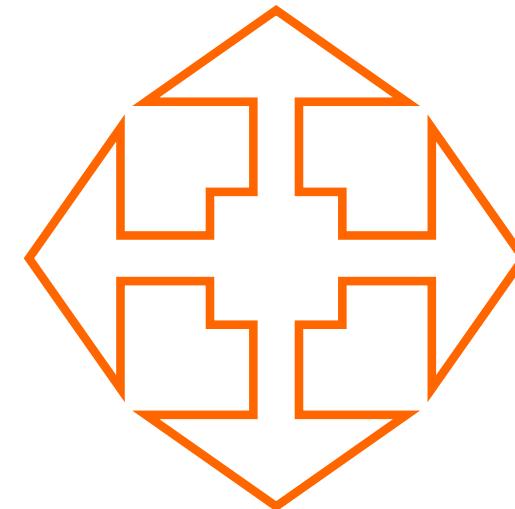
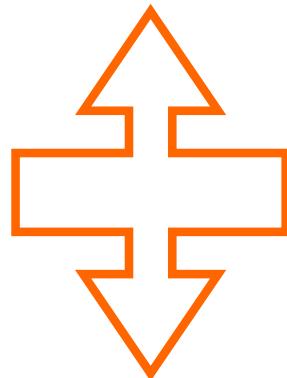
What is a Polygon?



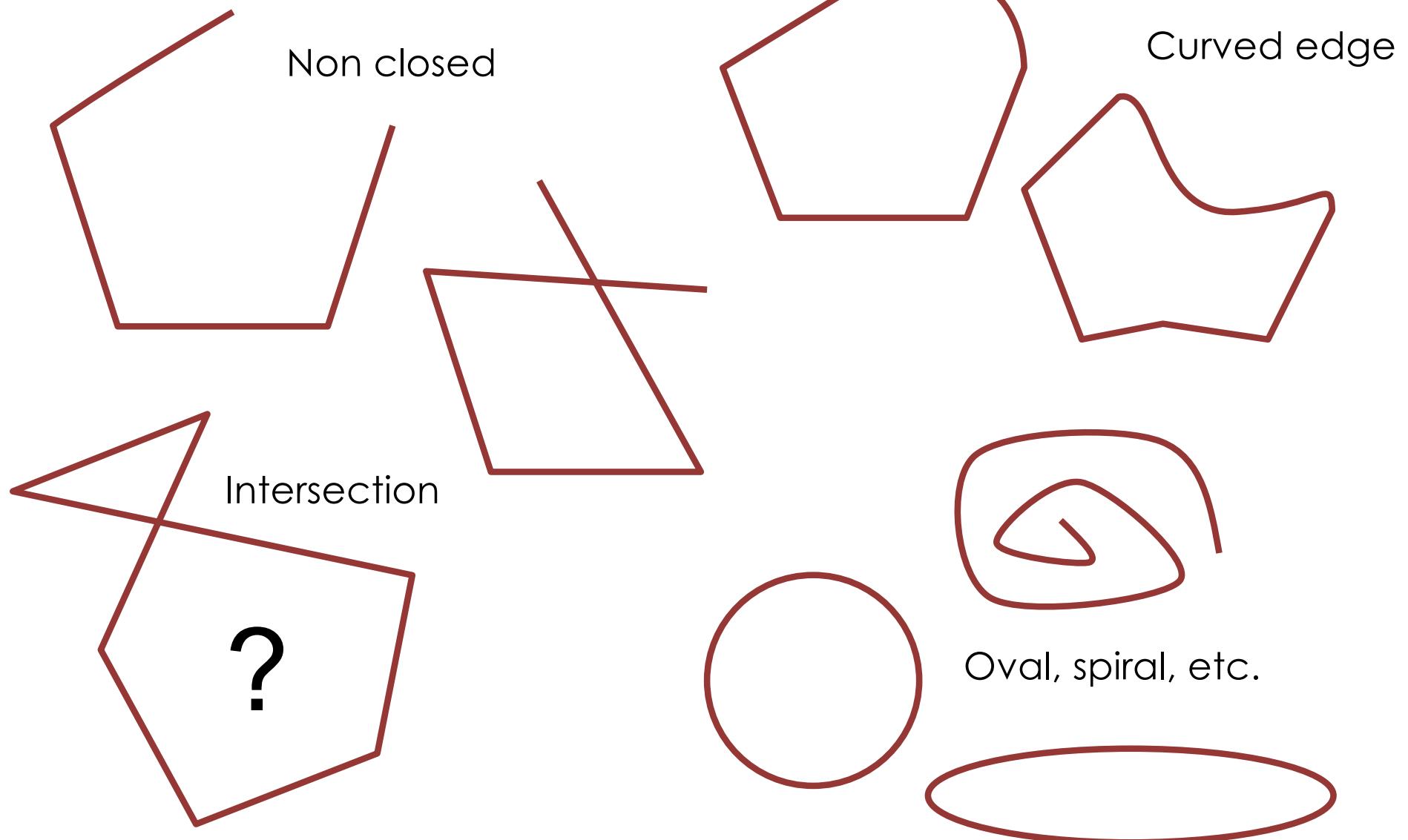
What is a Polygon?



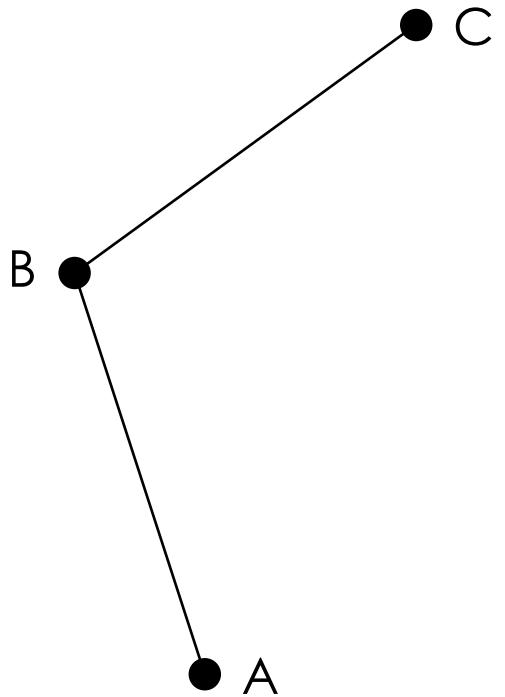
What is a Polygon?



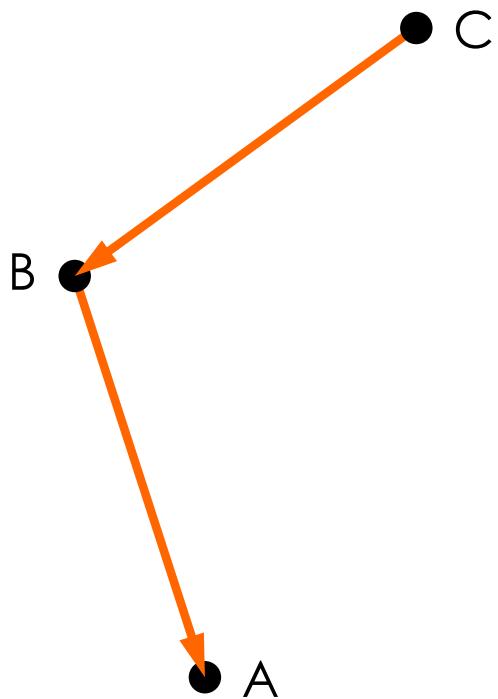
What is a NON-Polygon?



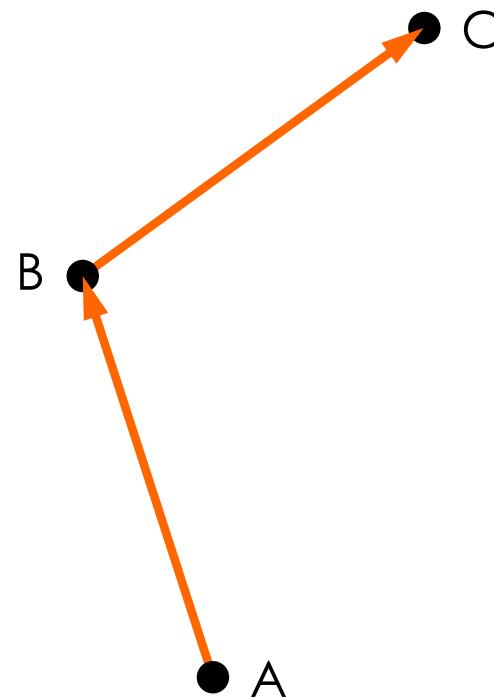
A 2 HOP VISIT



A 2 HOP VISIT

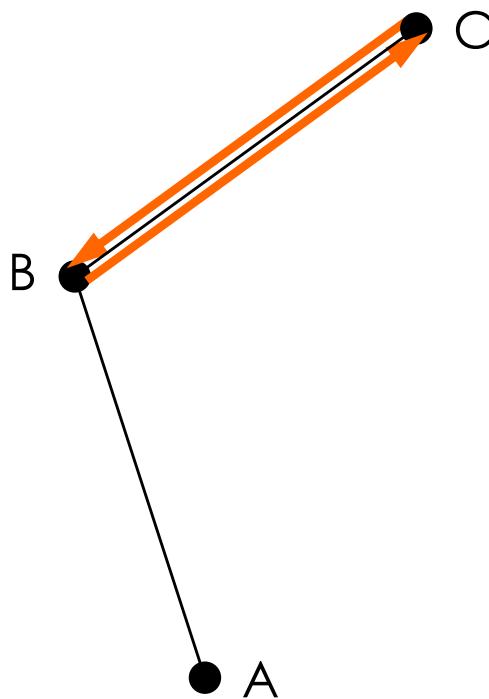


$C \rightarrow B \rightarrow A$

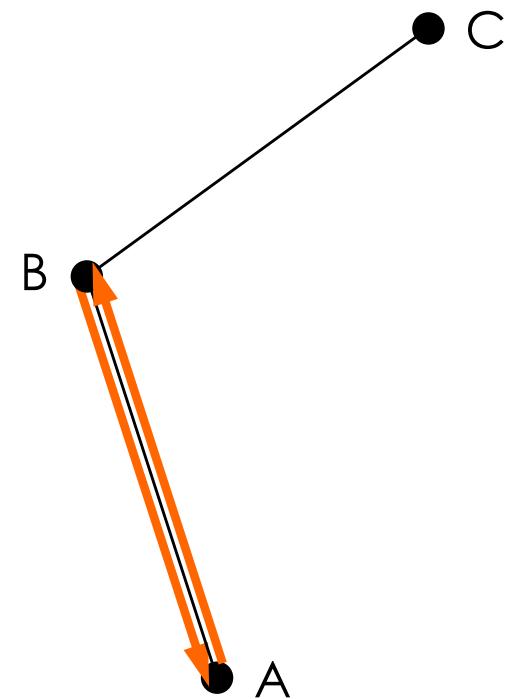


$A \rightarrow B \rightarrow C$

A 2 HOP VISIT

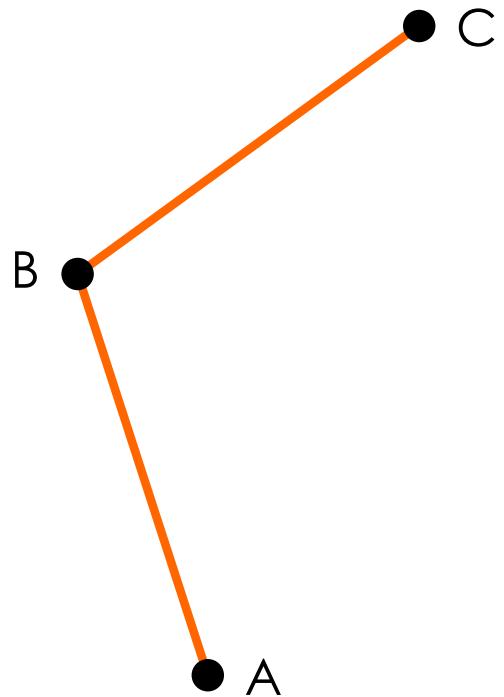


$B \rightarrow C \rightarrow B$
or
 $C \rightarrow B \rightarrow C$

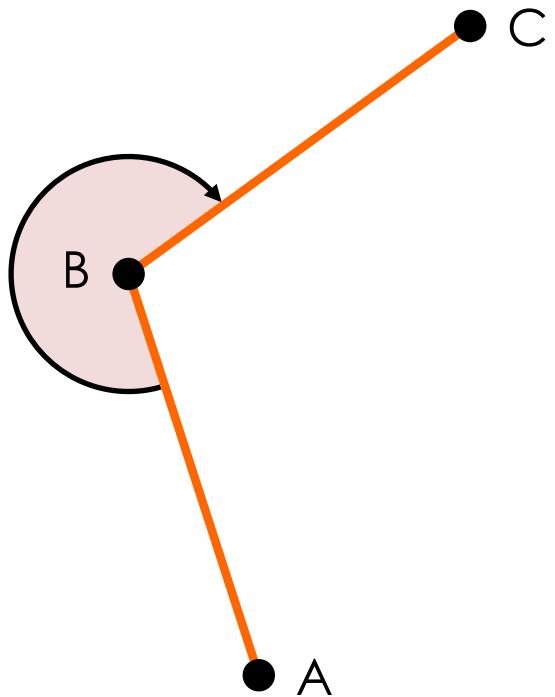


$B \rightarrow A \rightarrow B$
or
 $A \rightarrow B \rightarrow A$

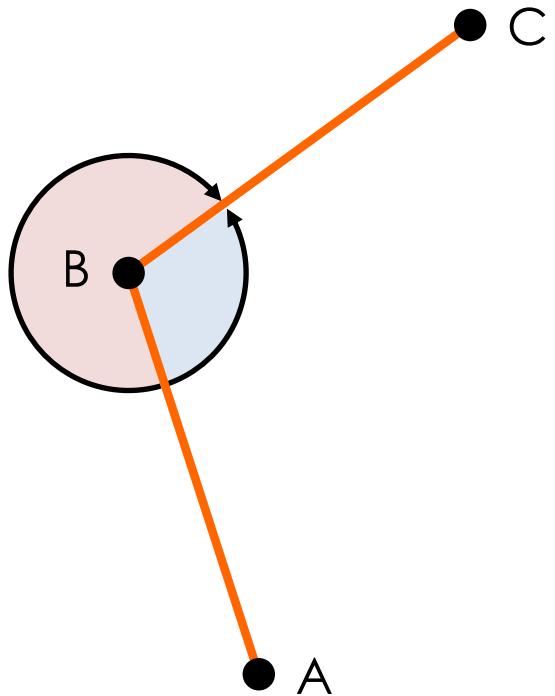
An Angle



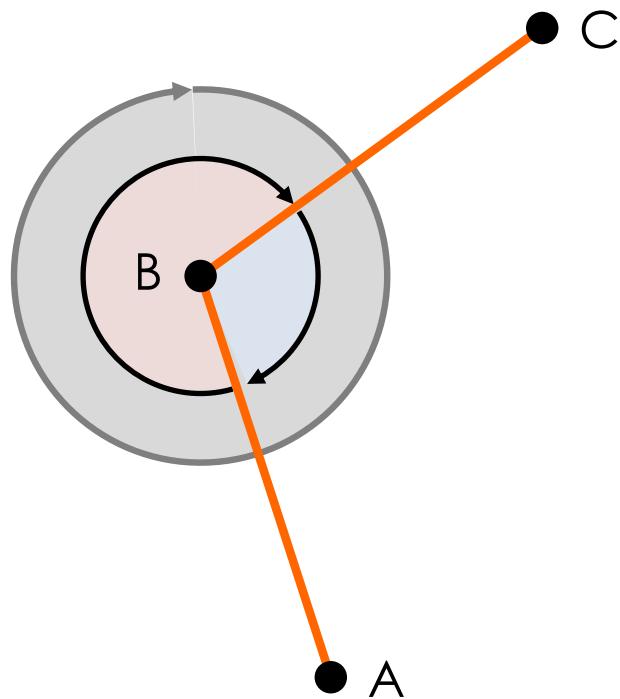
An Angle



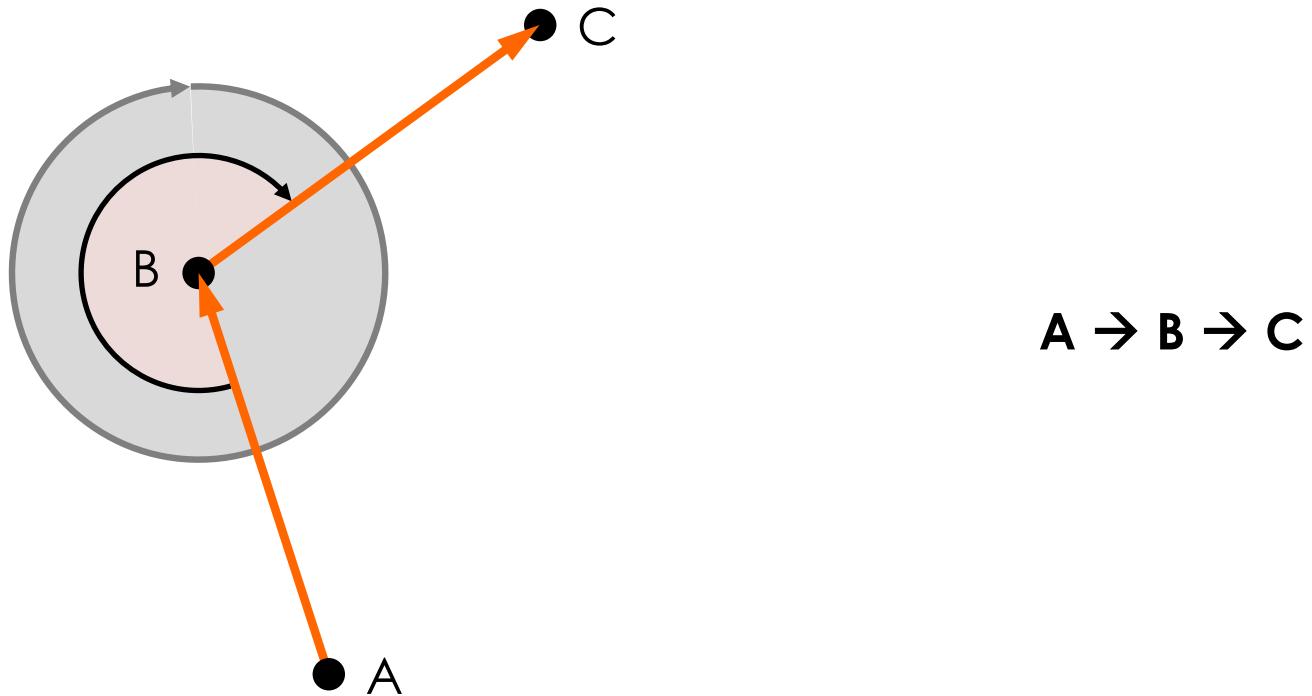
An Angle



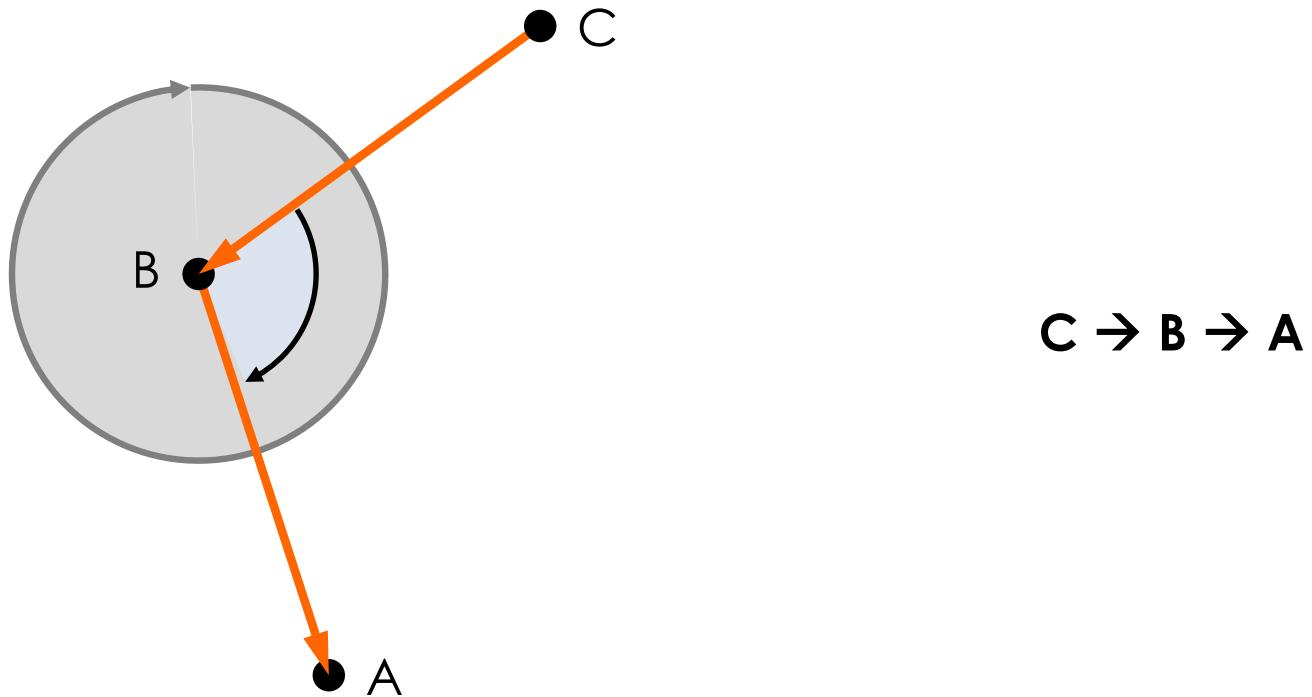
A POLAR Angle



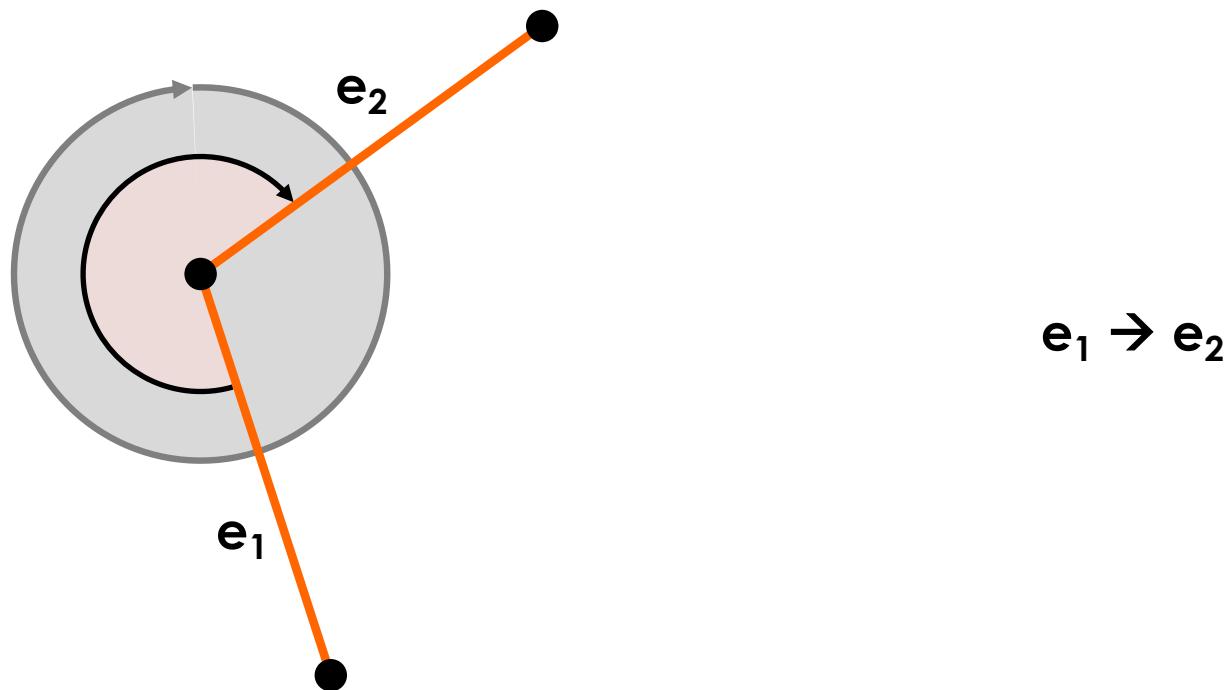
A 2 HOP VISIT and a POLAR Angle



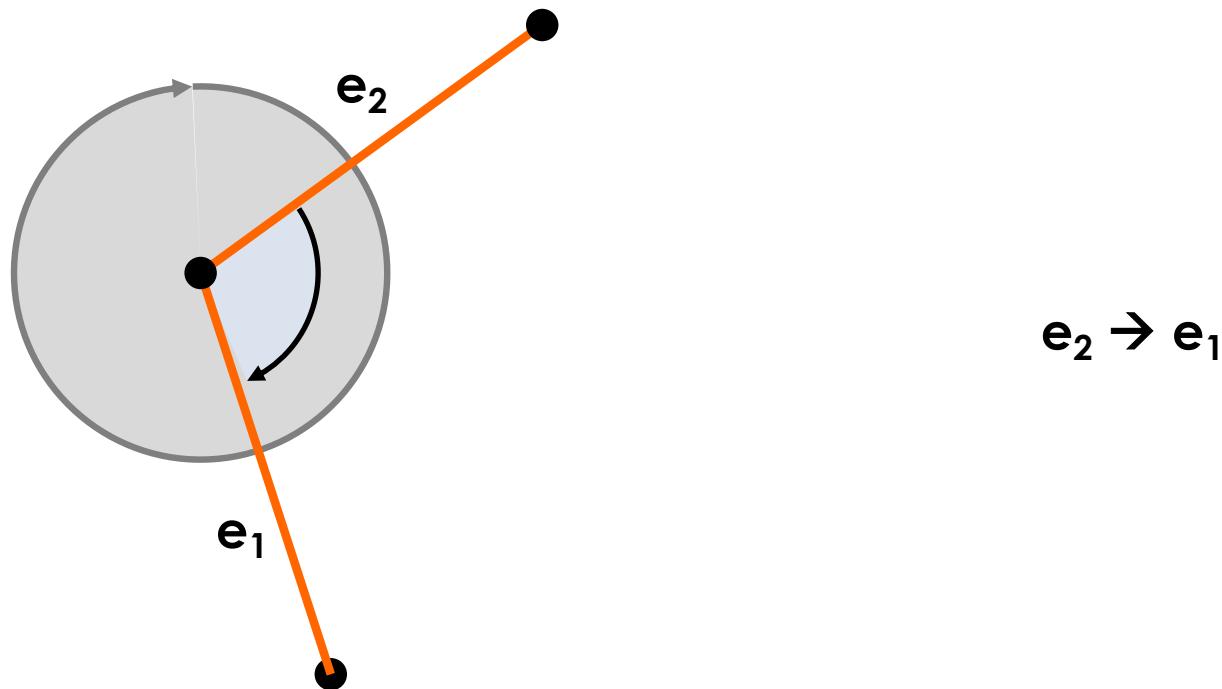
A 2 HOP VISIT and a POLAR Angle



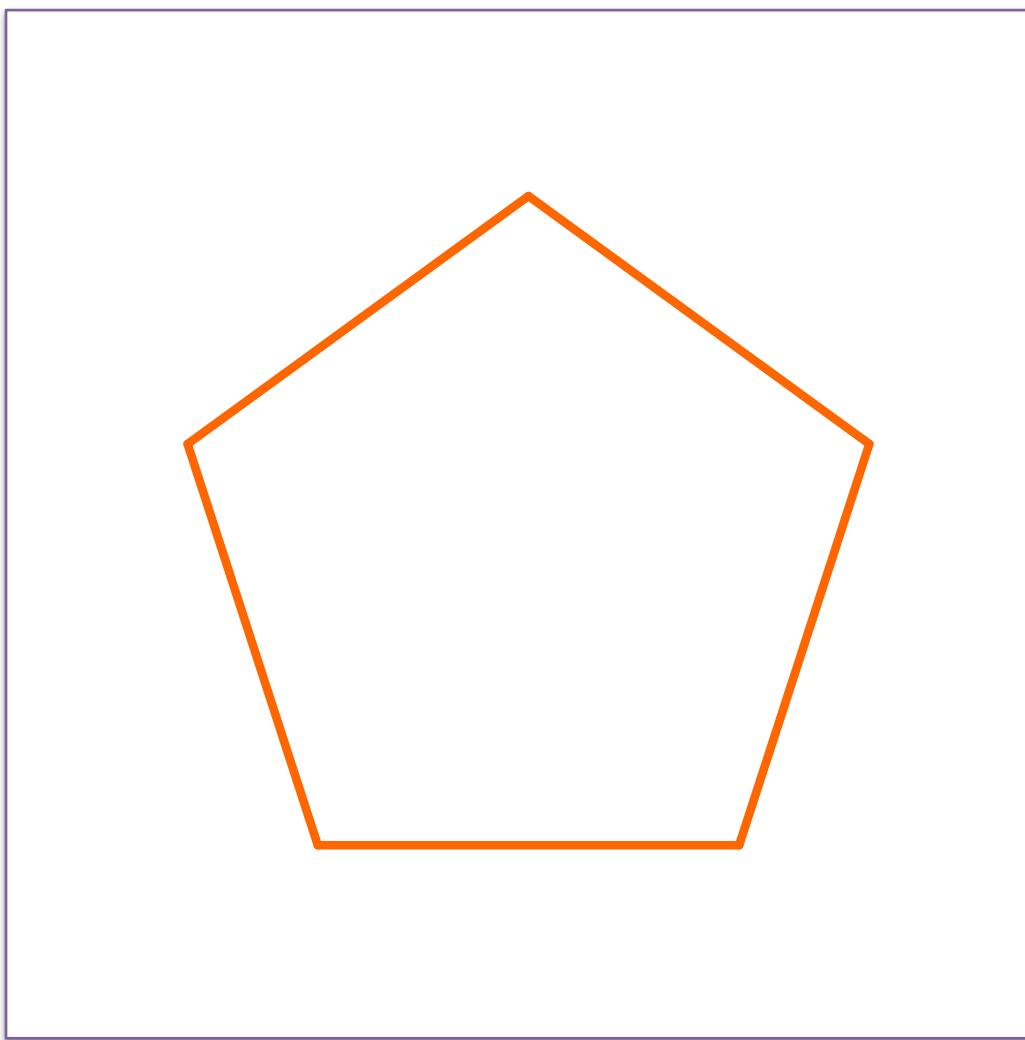
EDGES and a POLAR Angle



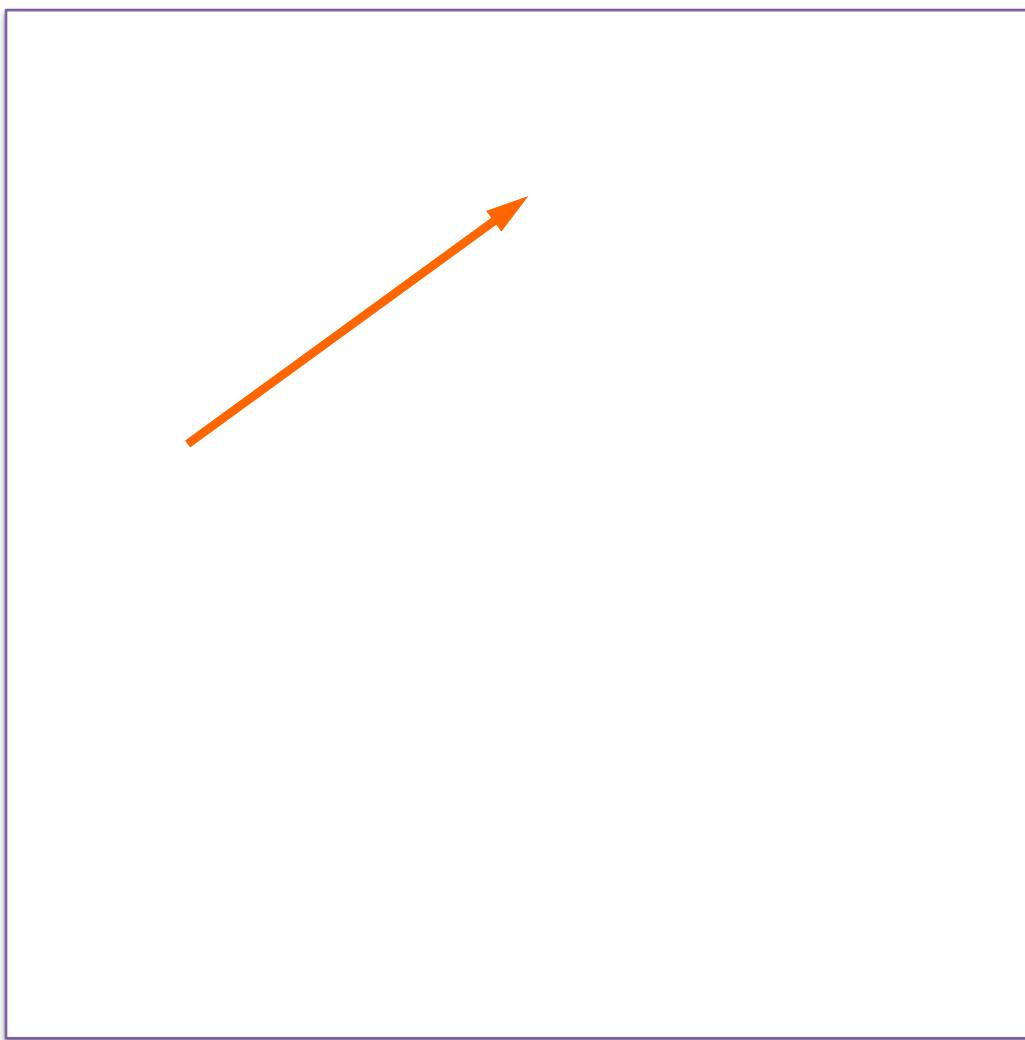
EDGES and a POLAR Angle



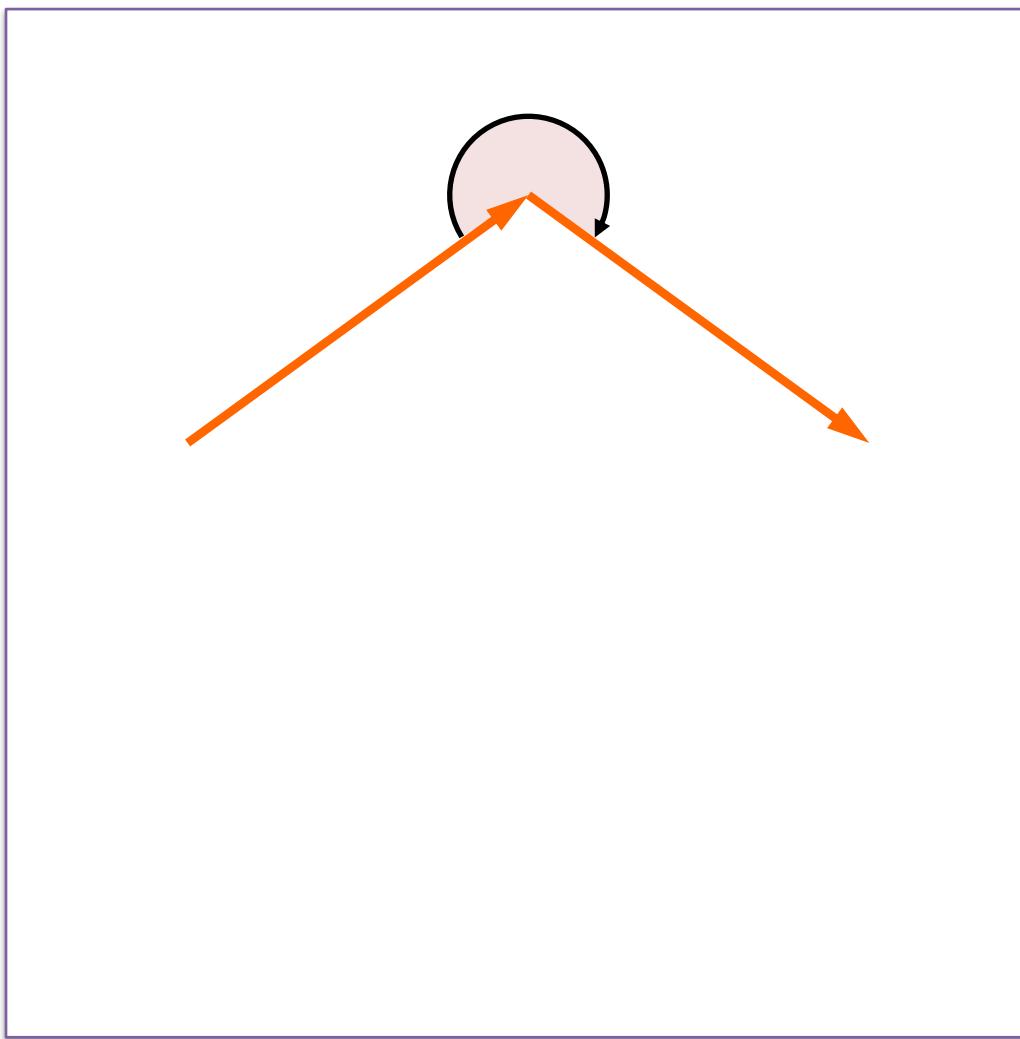
A Polygon



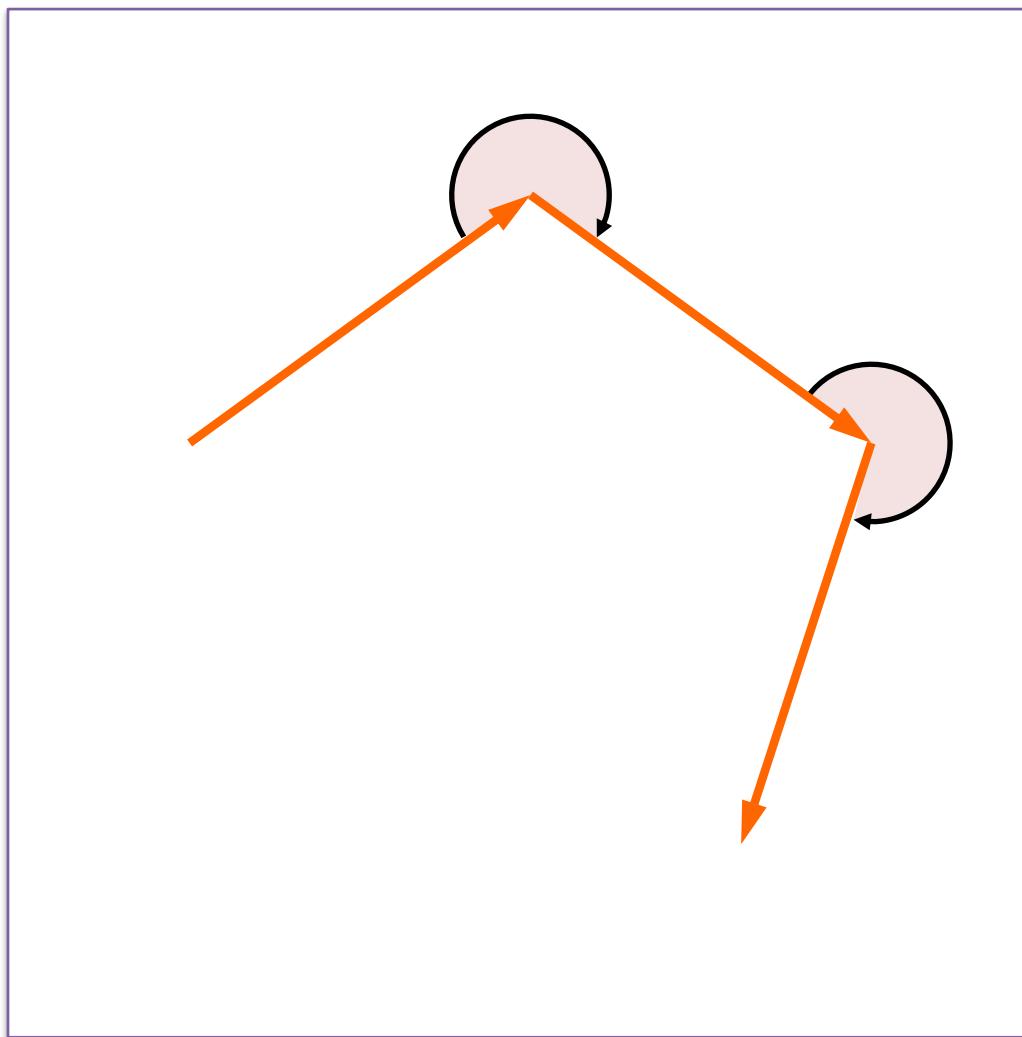
A Polygon



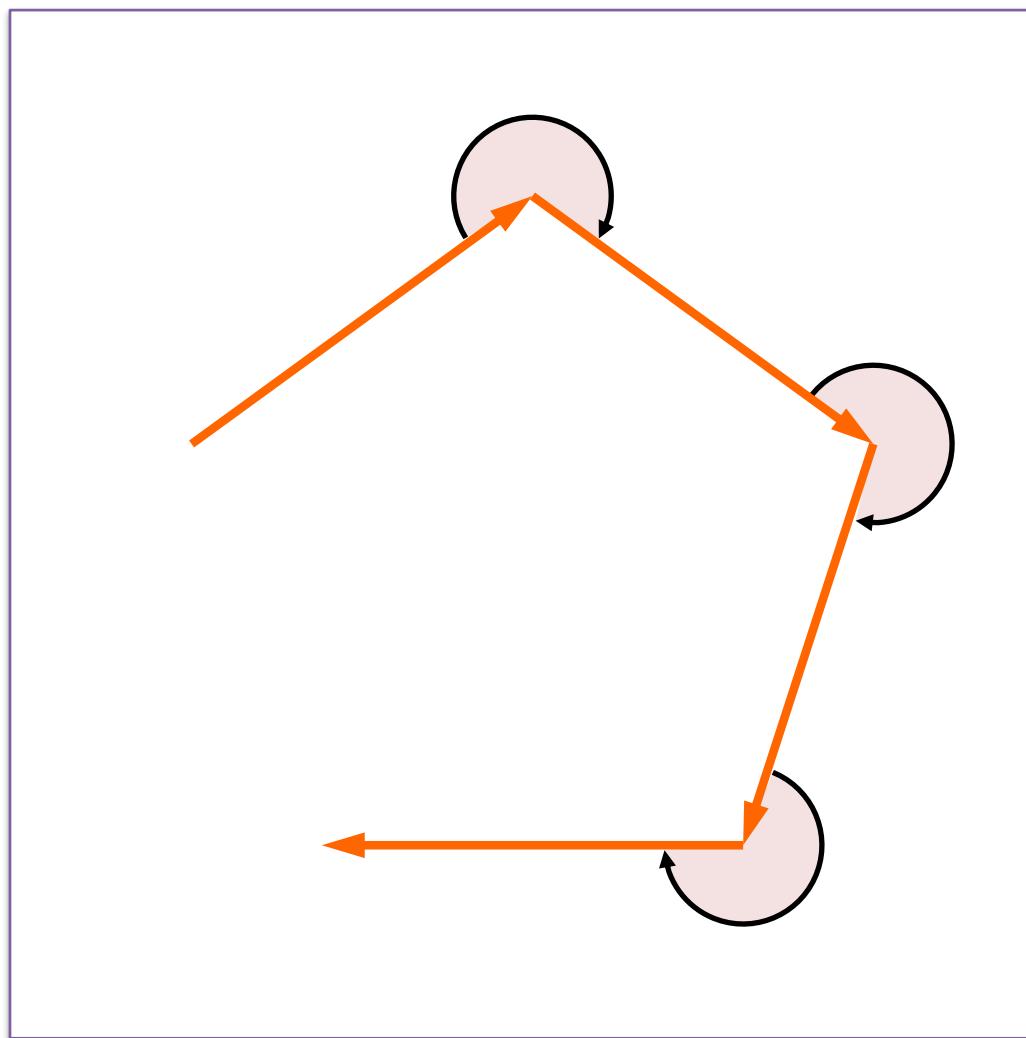
A Polygon



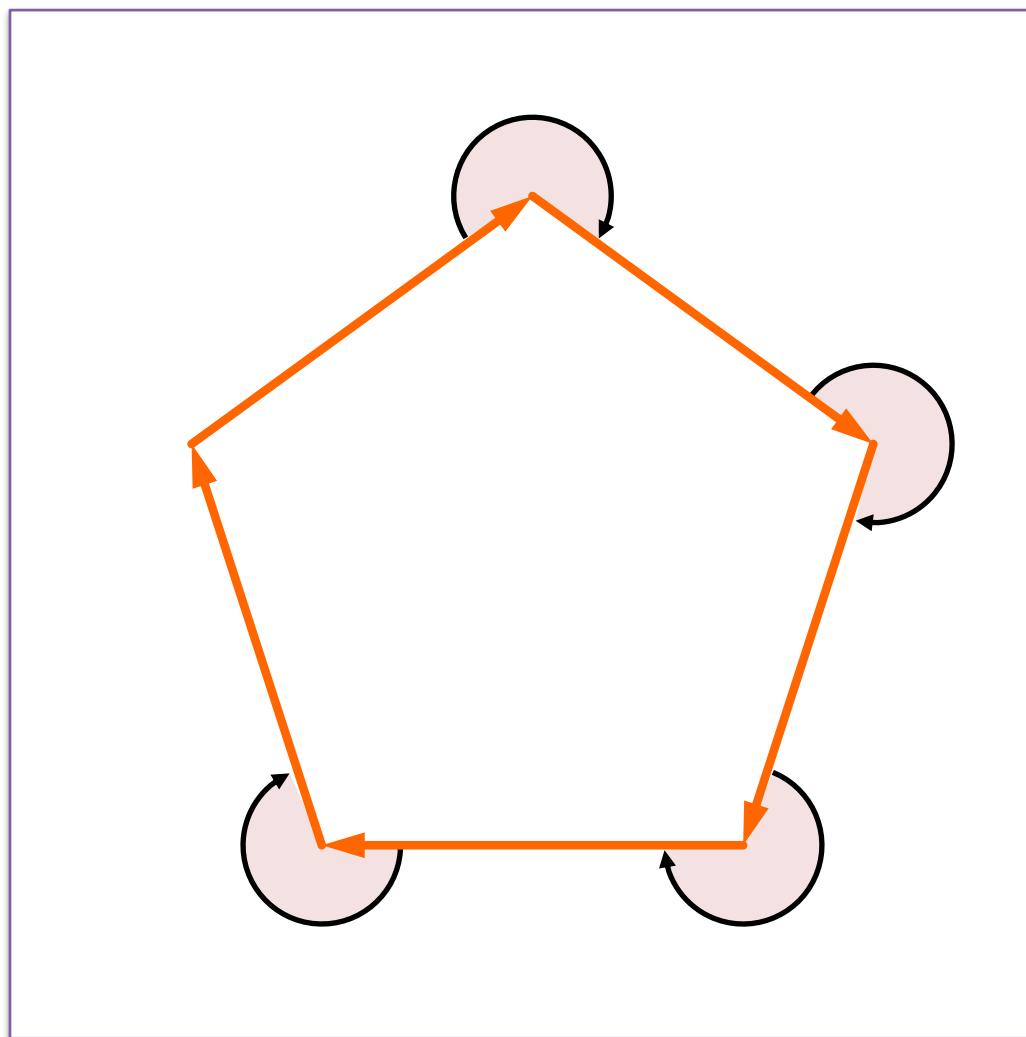
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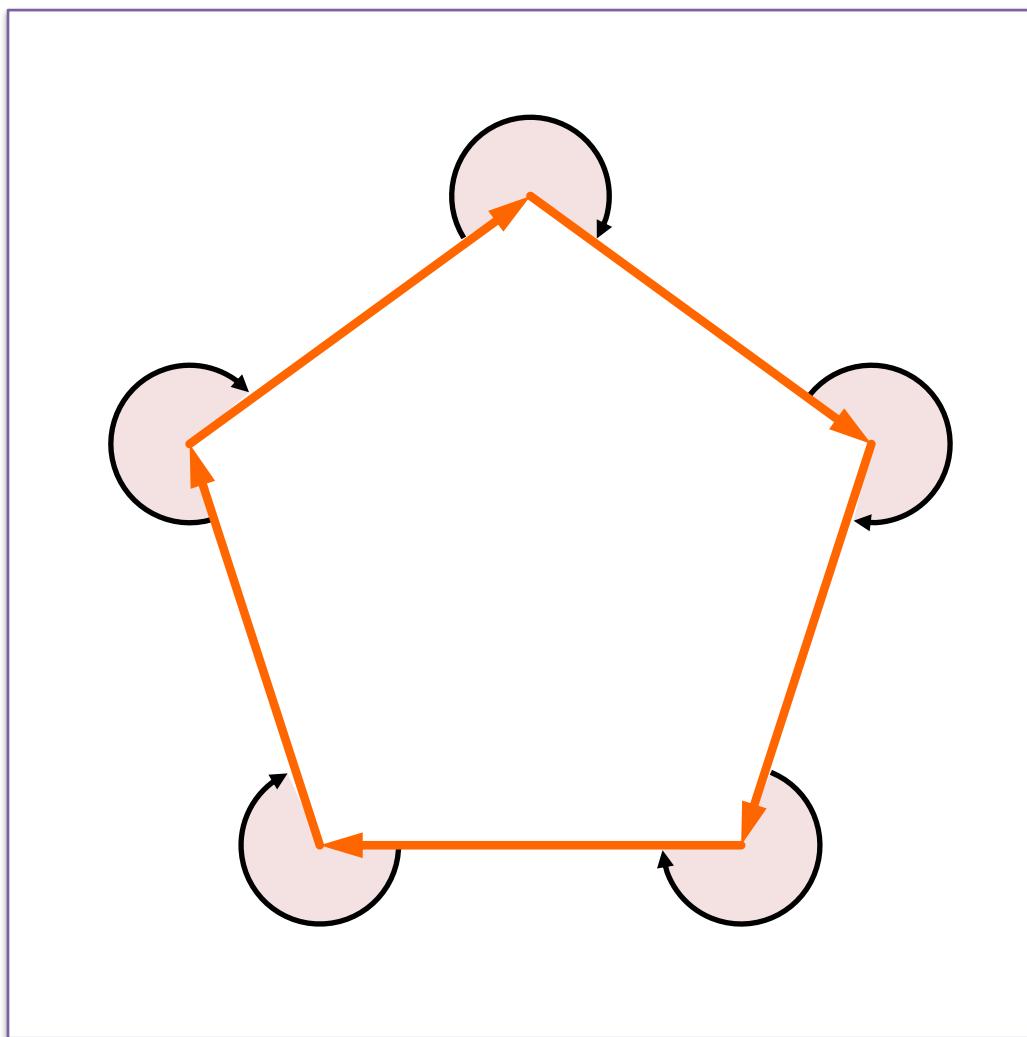
A Polygon



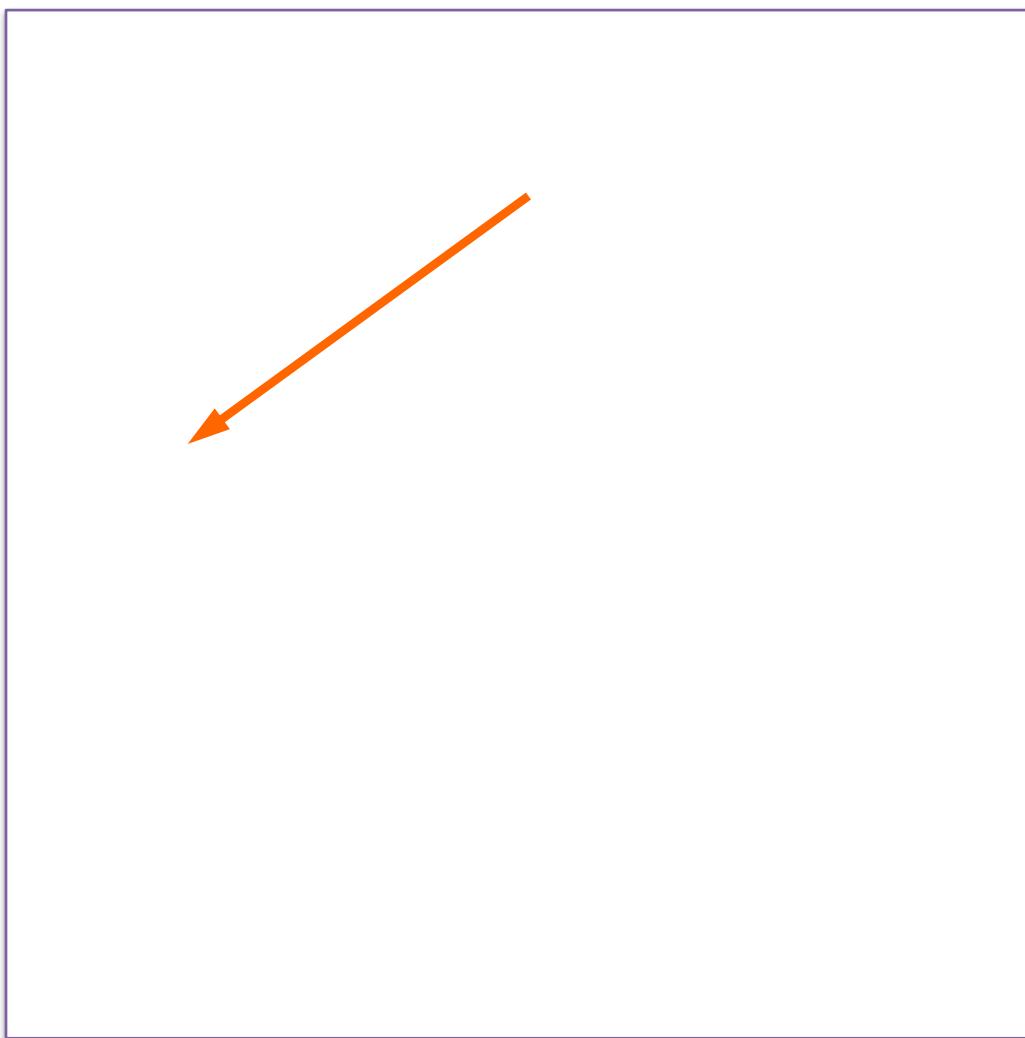
A Polygon



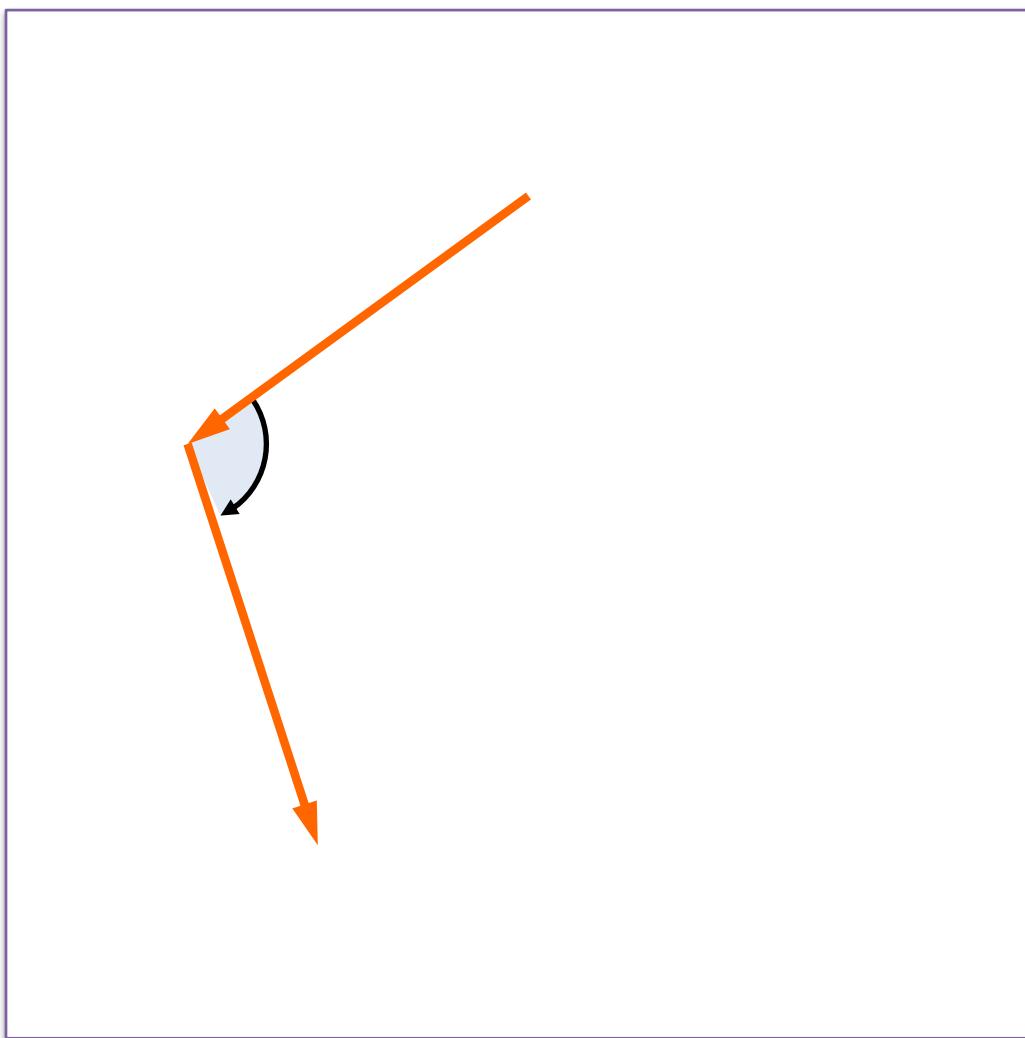
A Polygon



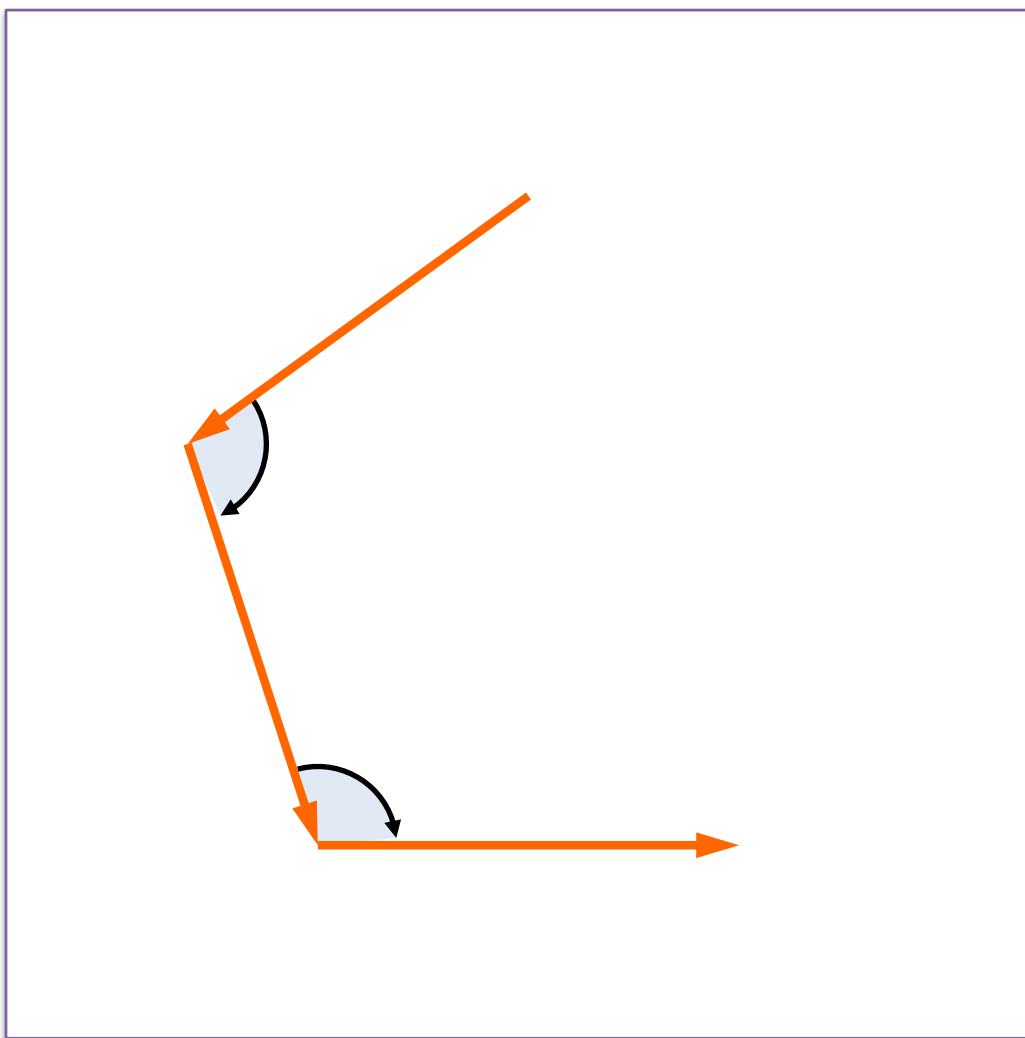
A Polygon



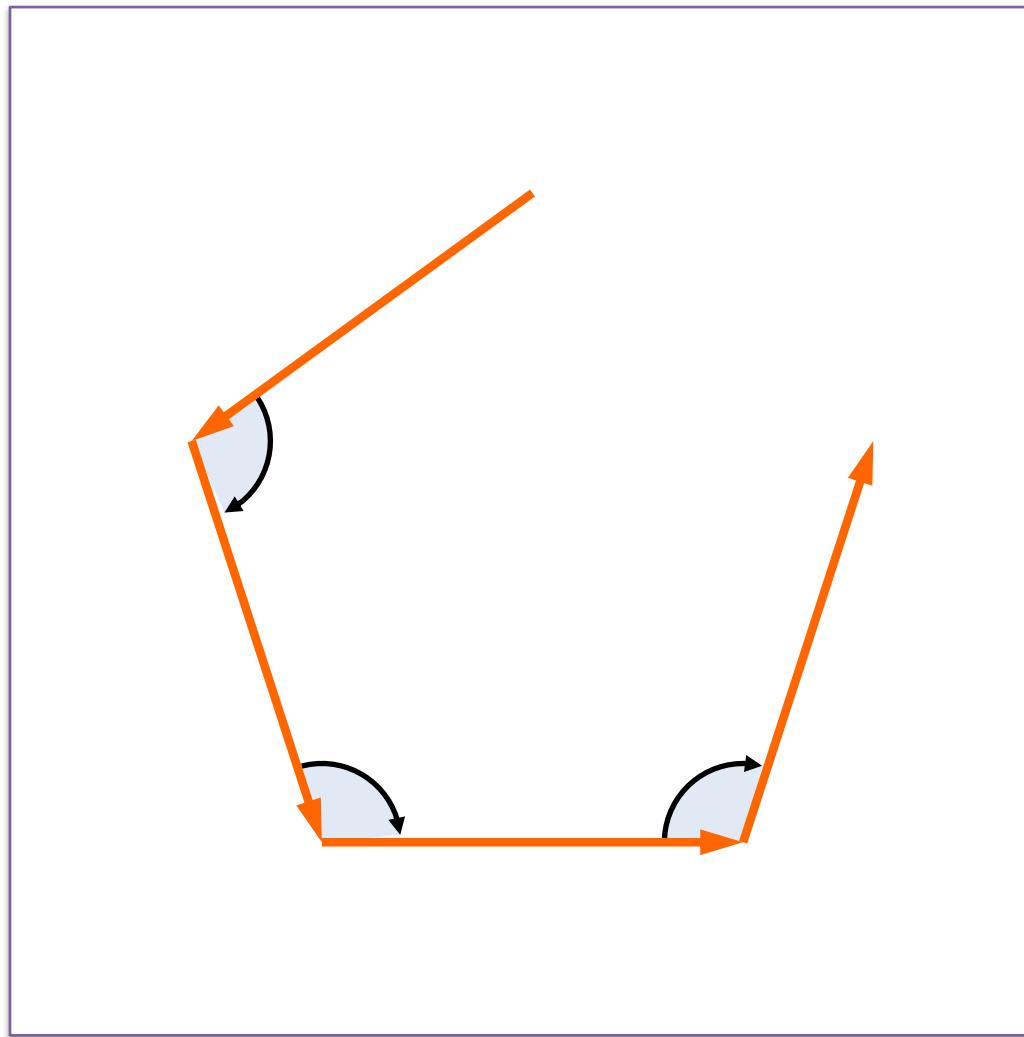
A Polygon



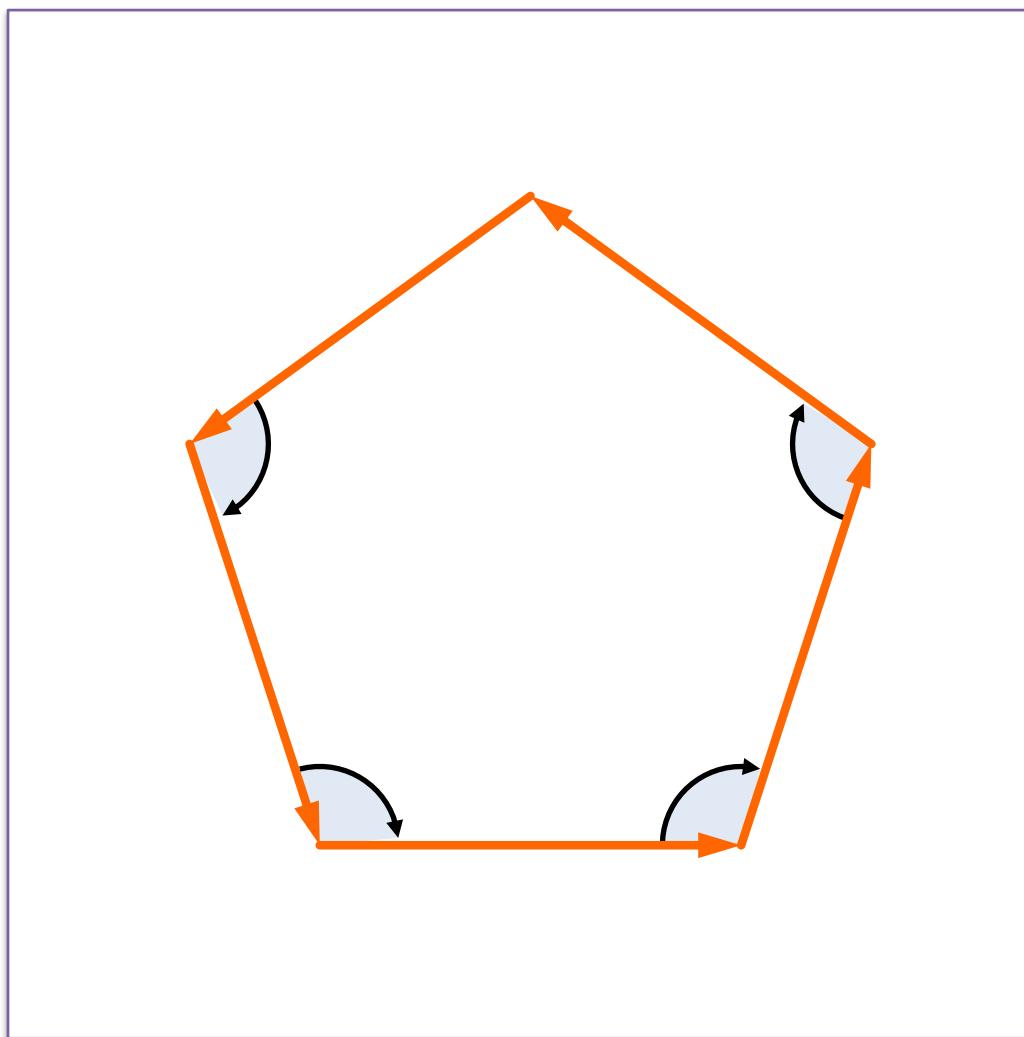
A Polygon



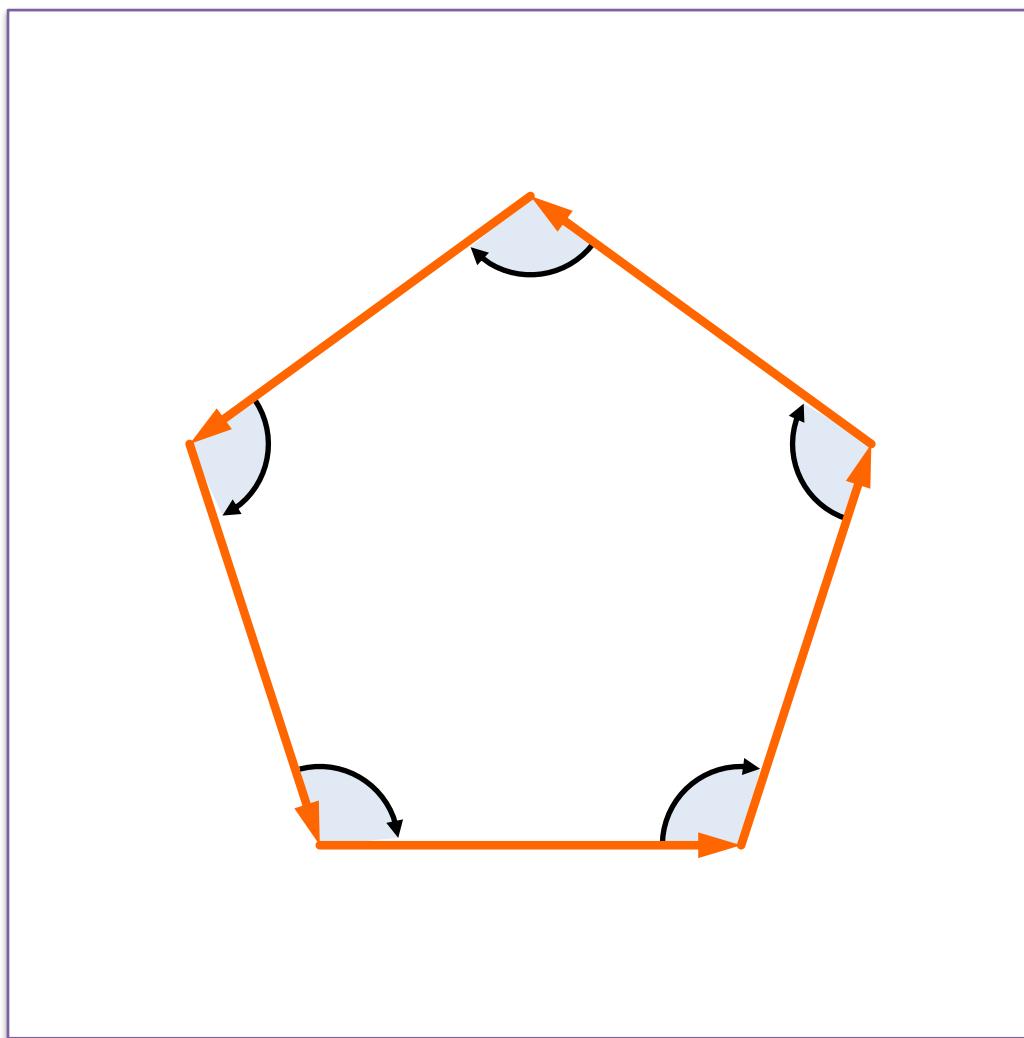
A Polygon



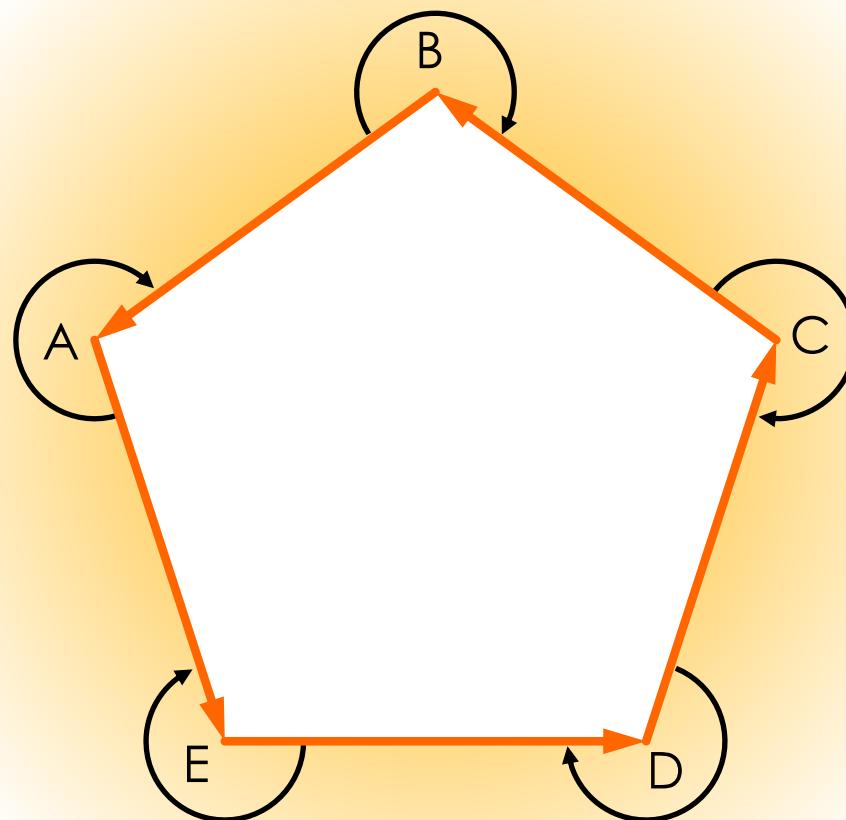
A Polygon



A Polygon

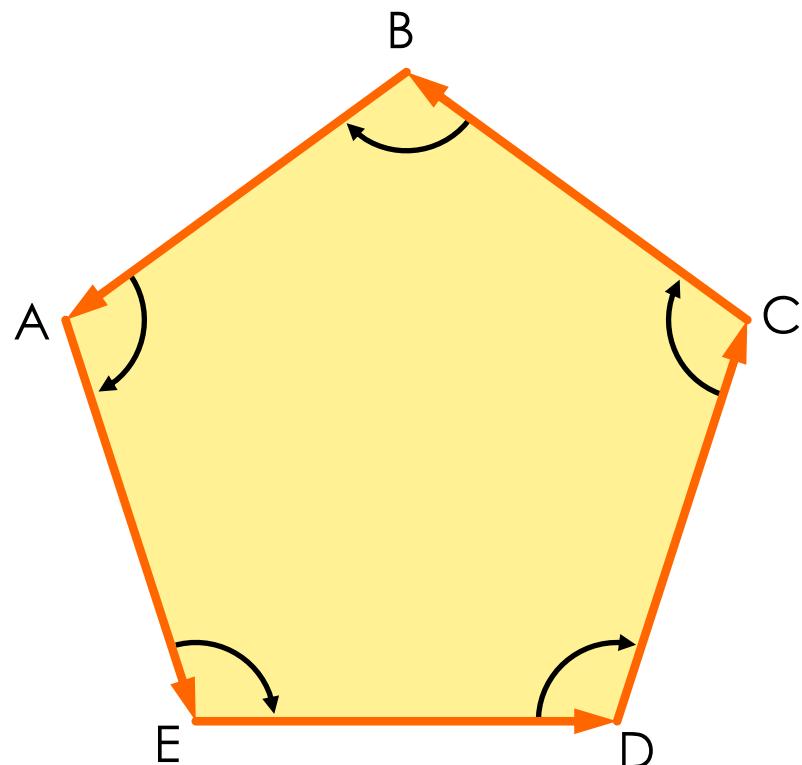


EXTERIOR Polygon



A → B → C → D → E → A

INTERIOR Polygon



$A \rightarrow E \rightarrow D \rightarrow C \rightarrow B \rightarrow A$

Conclusion

A POLYGON is defined by its
VISITED ANGLES

Definition

A **POLYGON**

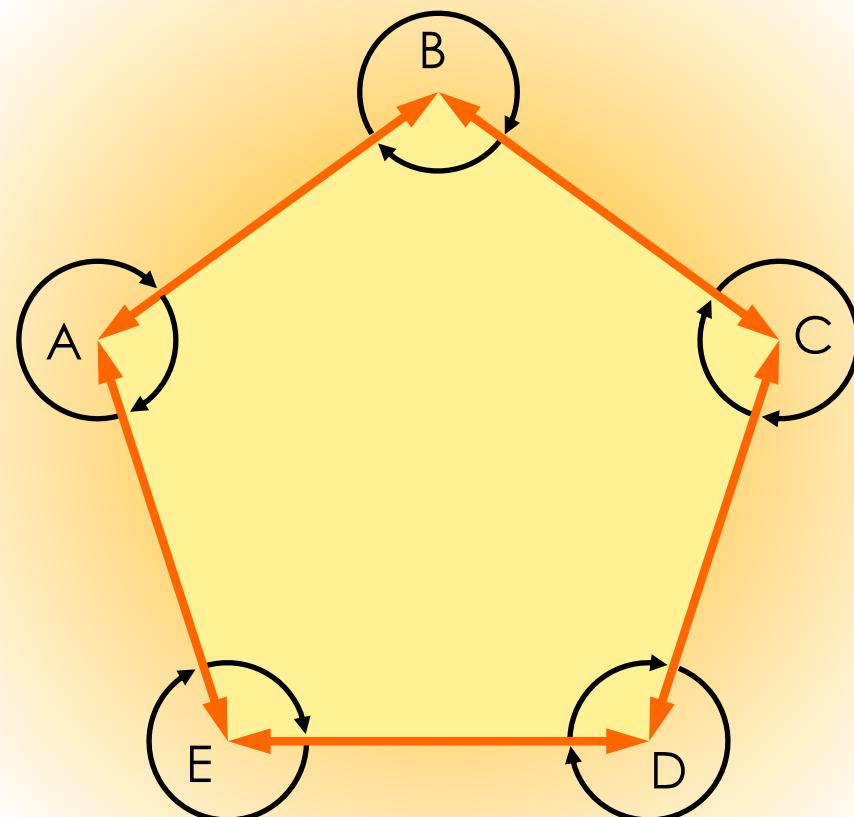
POLus = many

GONia = angles

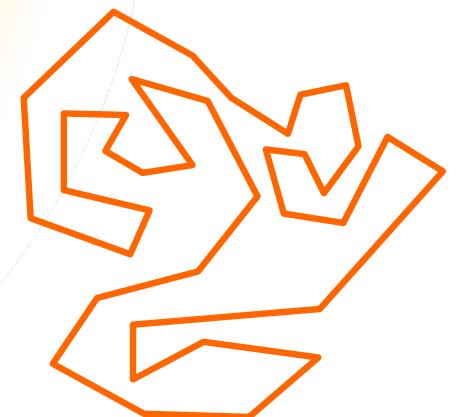
Question

**How to determine that a
polygon is **INTERIOR** or
EXTERIOR?**

INTERIOR or EXTERIOR Polygon?

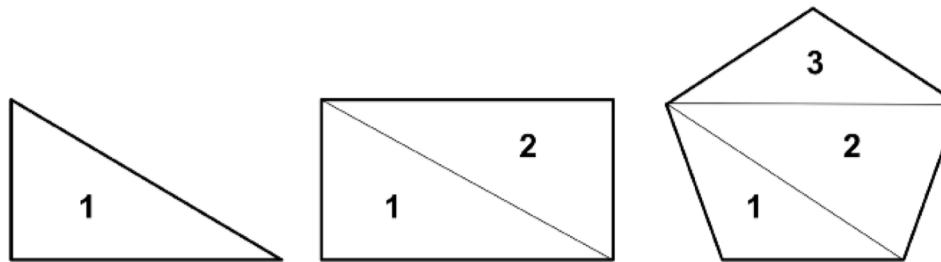


$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow A$
 $A \rightarrow E \rightarrow D \rightarrow C \rightarrow B \rightarrow A$



INTERIOR ANGLES

There are $(n-2)$ triangles in a polygon with n vertices



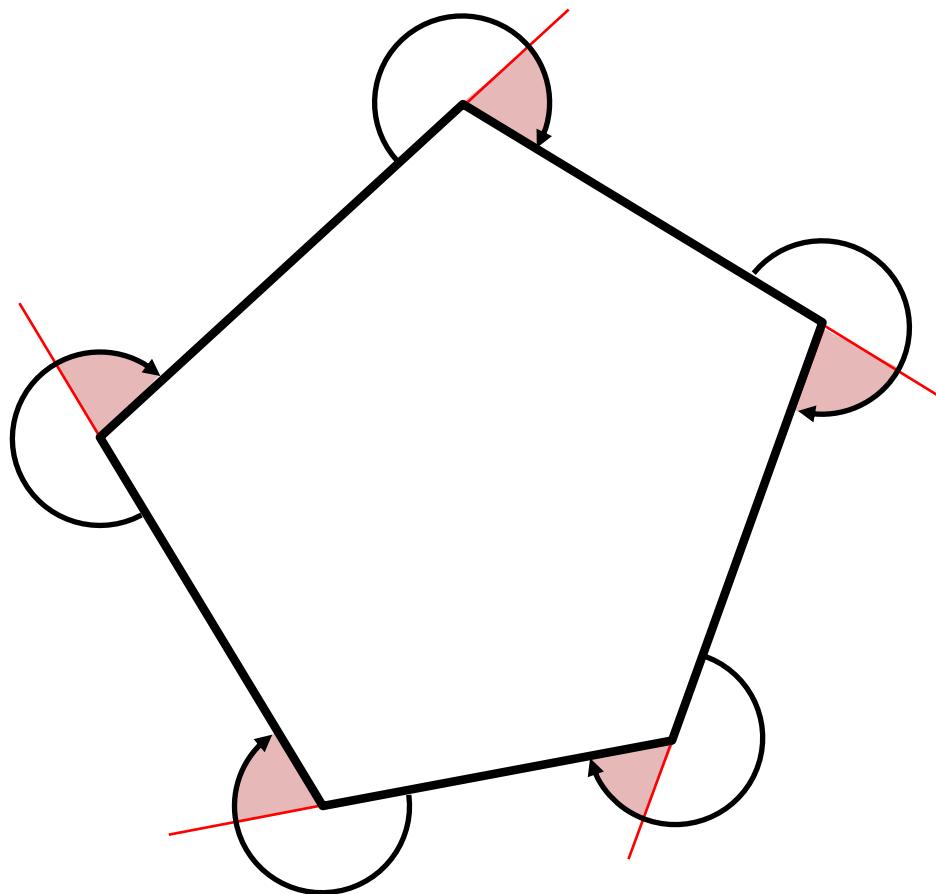
The sum of the interior angles of a triangle is equal to 180°

**The sum of the interior angles of a polygon is
equal to**

$$(n - 2) \times 180^\circ$$

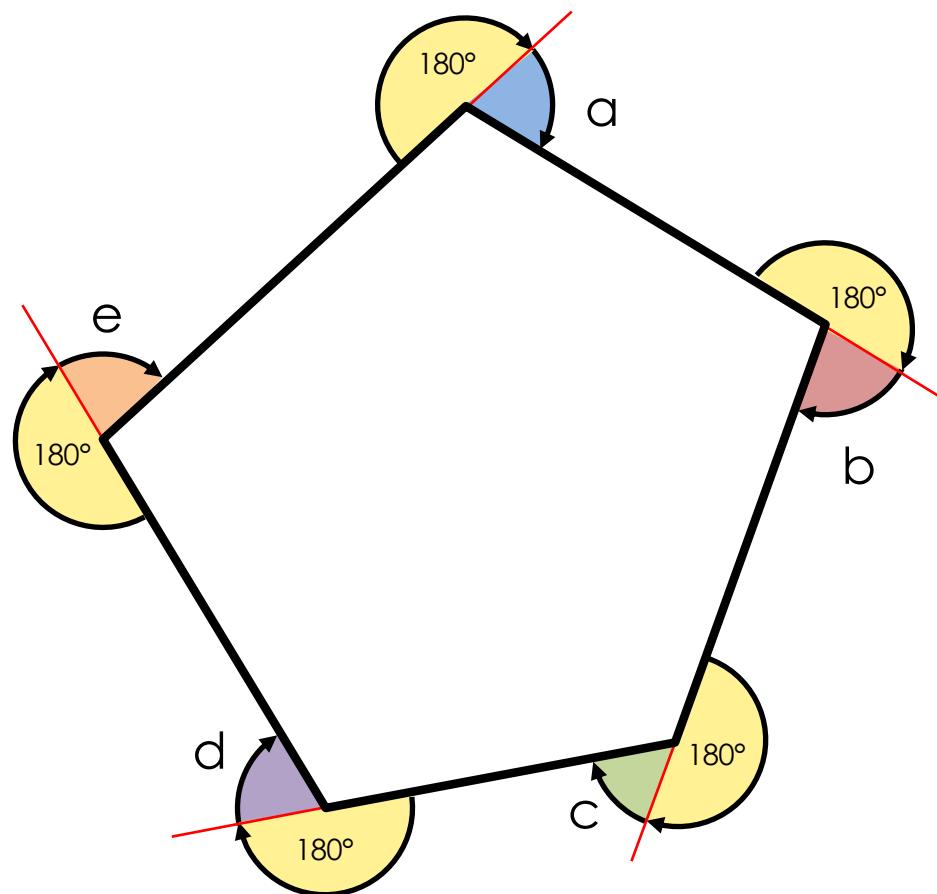
EXTERIOR ANGLES

The sum of the exterior (angles)



EXTERIOR ANGLES

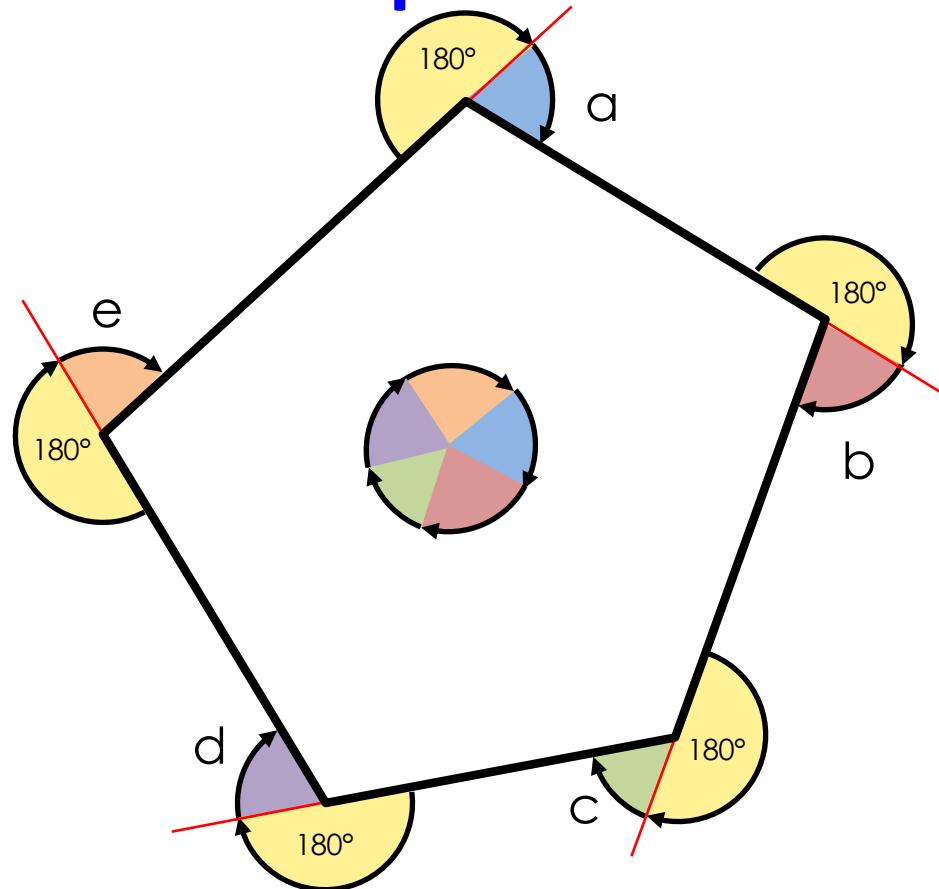
The sum of each (exterior angle – 180) of a polygon is equal to



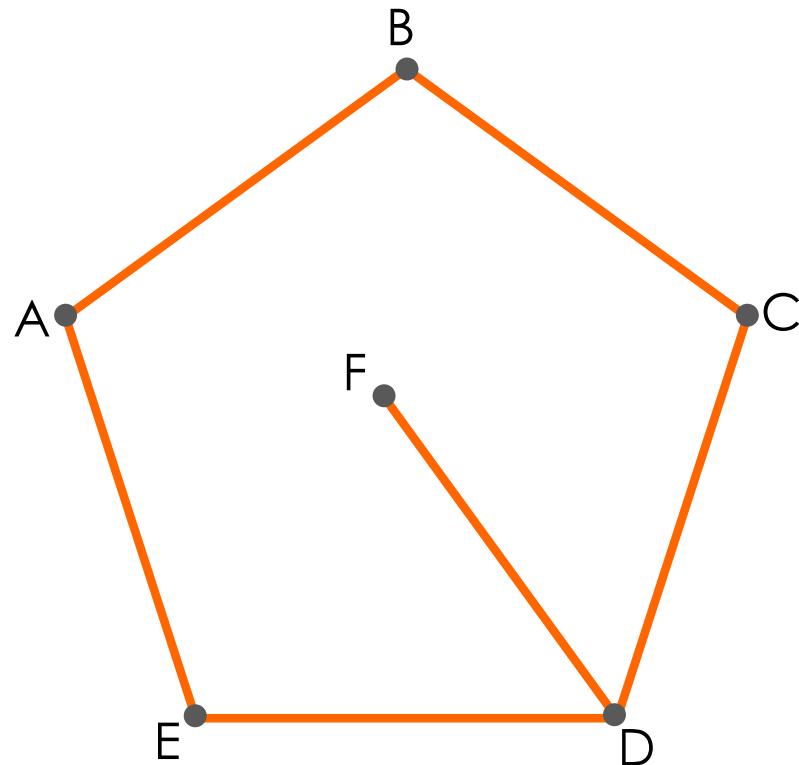
EXTERIOR ANGLES

The sum of each (exterior angle – 180) of a polygon

is equal to 360°

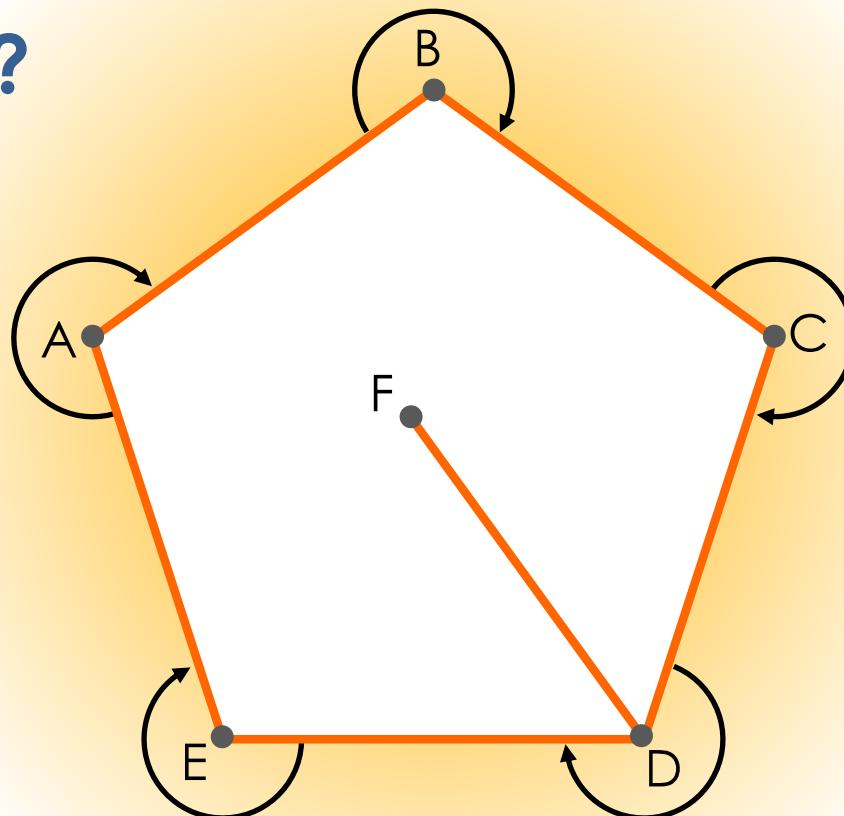


And what about this shape?



What about this shape?

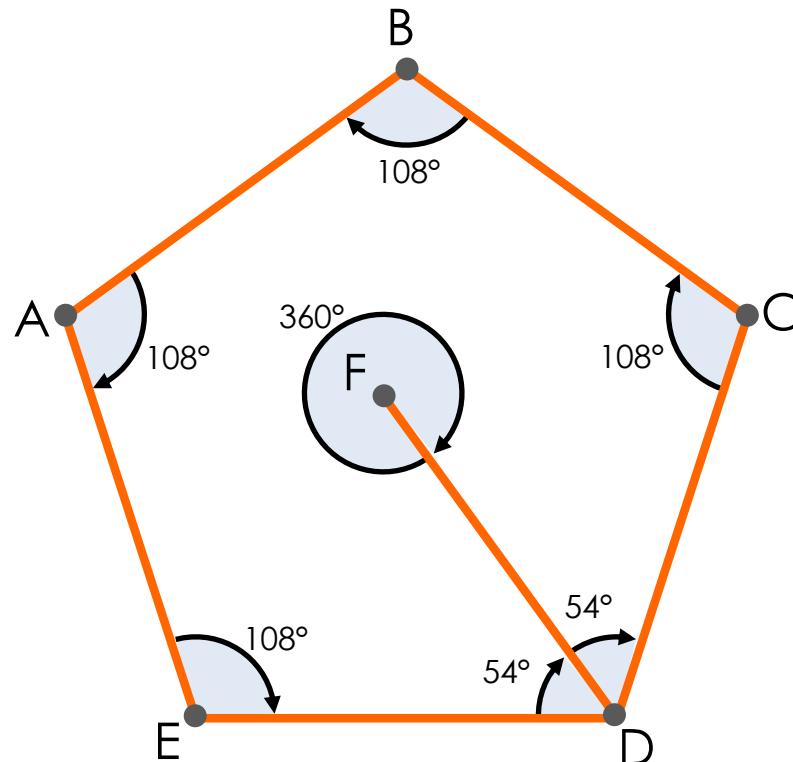
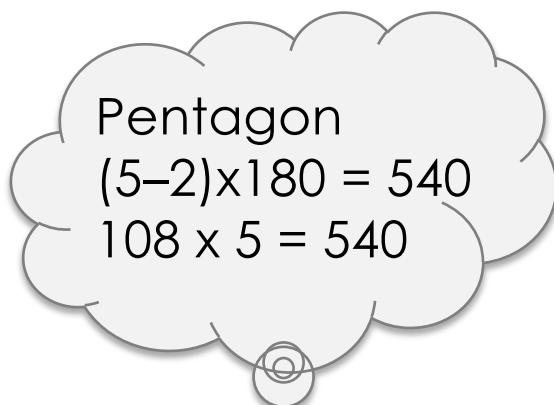
Is it EXTERIOR ?



EXTERIOR ✓

What about this shape?

Is it INTERIOR ?



$$(6-2) \times 180 = 720$$
$$108 \times 5 + 360 = 900$$

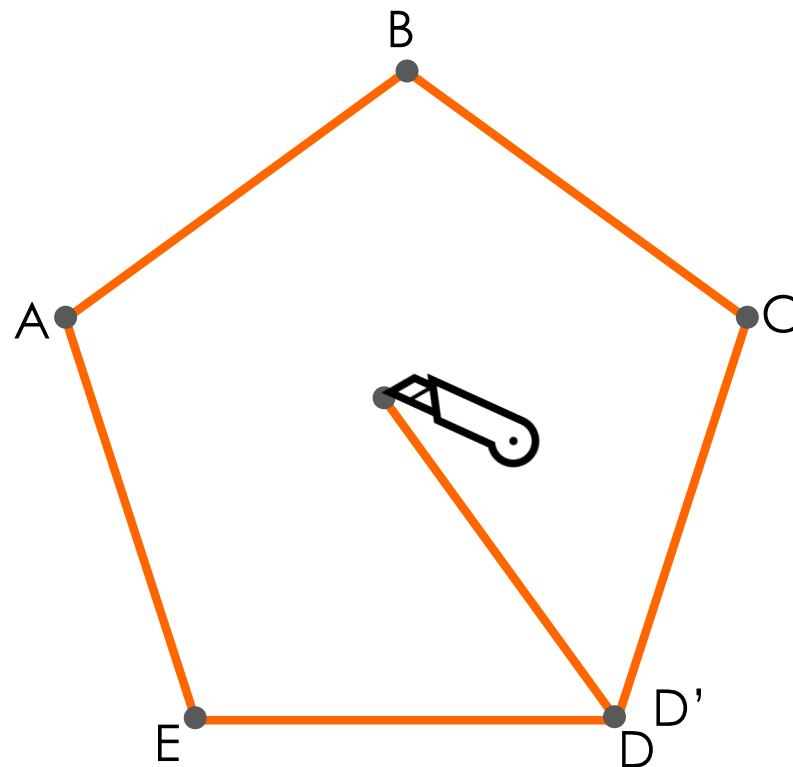
NON-INTERIOR X

BUT ...



What is a **TORN** Polygon

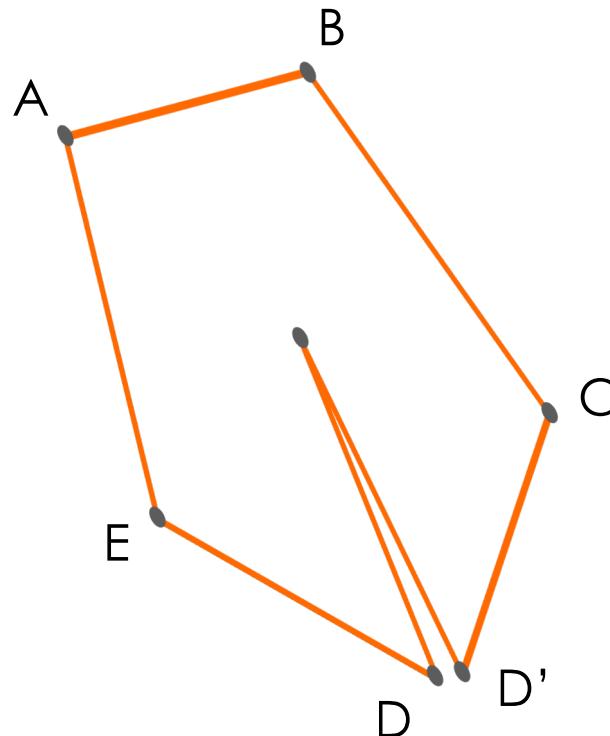
Is it INTERIOR ?



2 overlapped vertices

What is a **TORN** Polygon

Is it **INTERIOR** ?



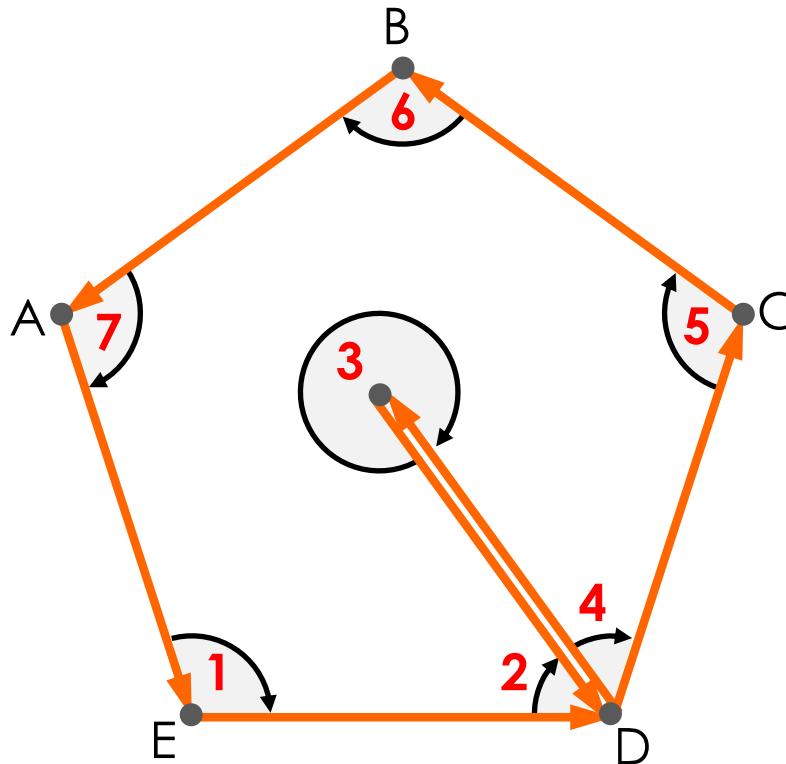
In this case, this shape looks like a polygon with **7** vertices

$$(7-2) \times 180 = 900$$
$$108 \times 5 + 360 = 900$$

INTERIOR ✓

What about visits & angles?

Is it INTERIOR ?



$$n=7$$

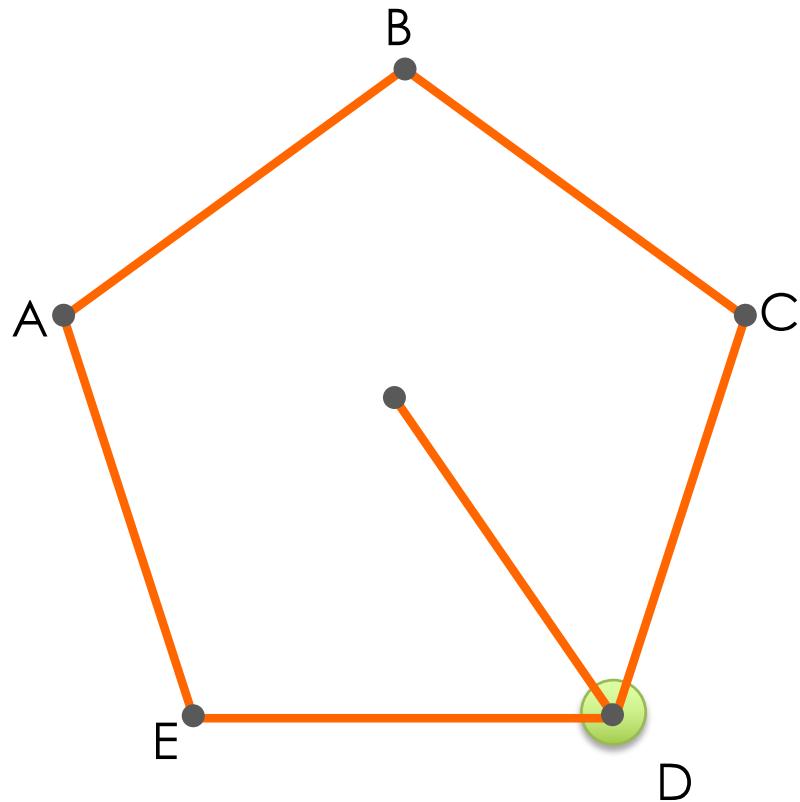
$$(7-2) \times 180 = 900$$

$$108 \times 5 + 360 = 900$$

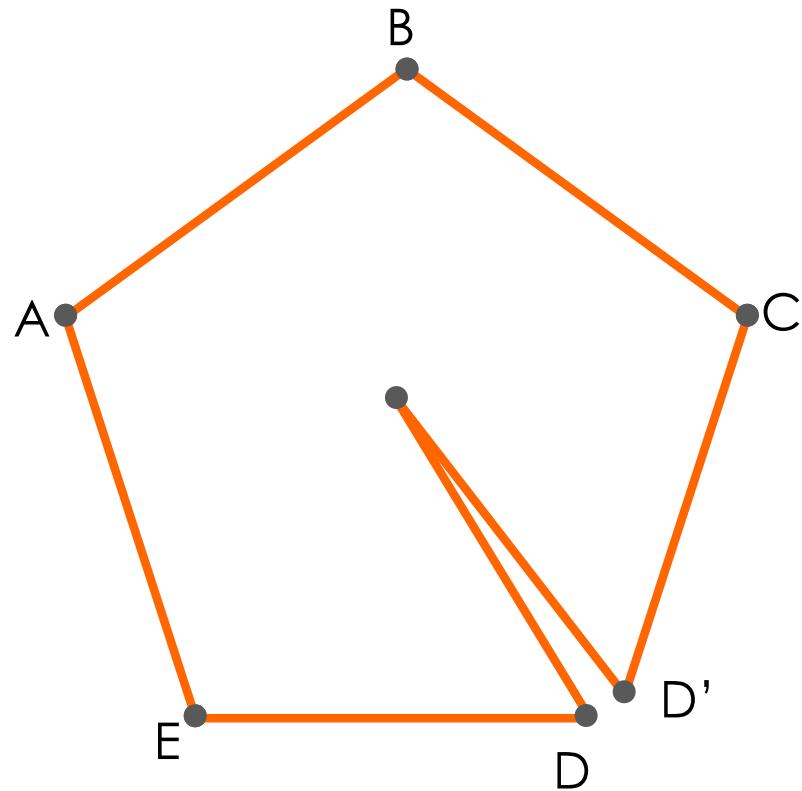
Let us consider **n** as the
number of formed angles
instead just the number of
vertices

INTERIOR ✓

A PSEUDO-POLYGON



A PSEUDO-POLYGON

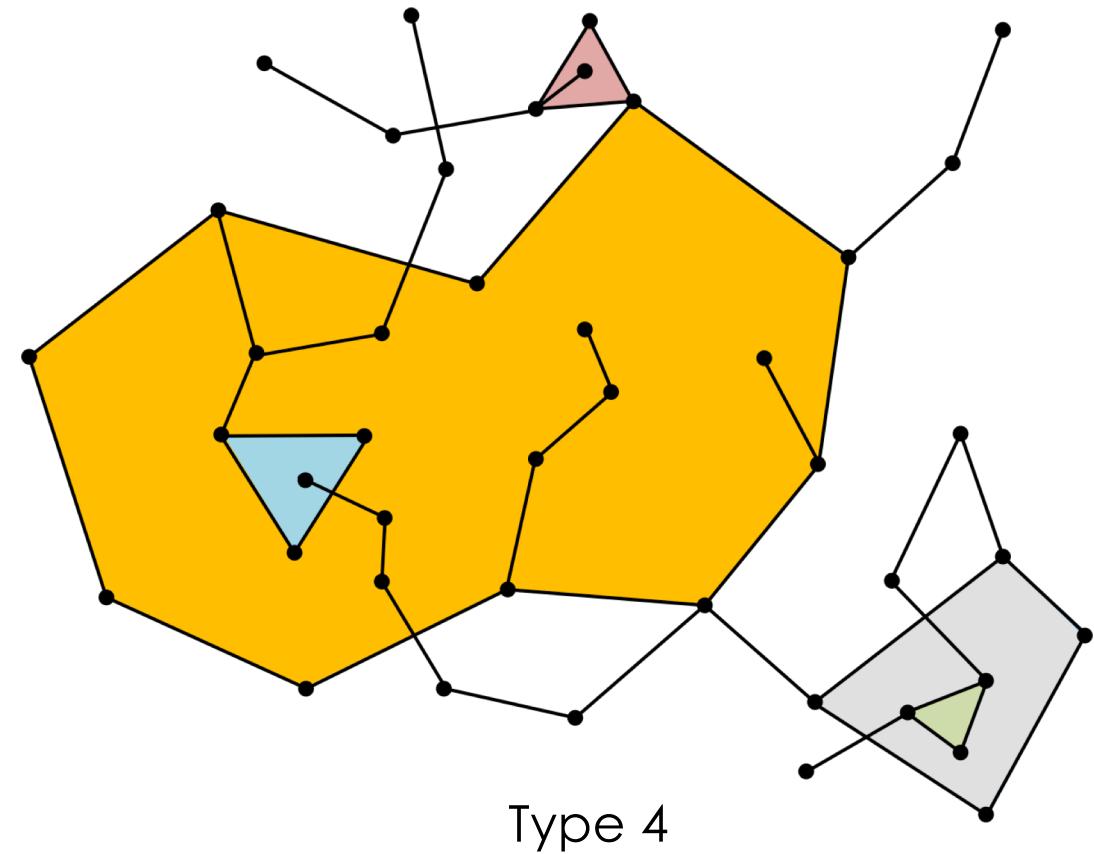
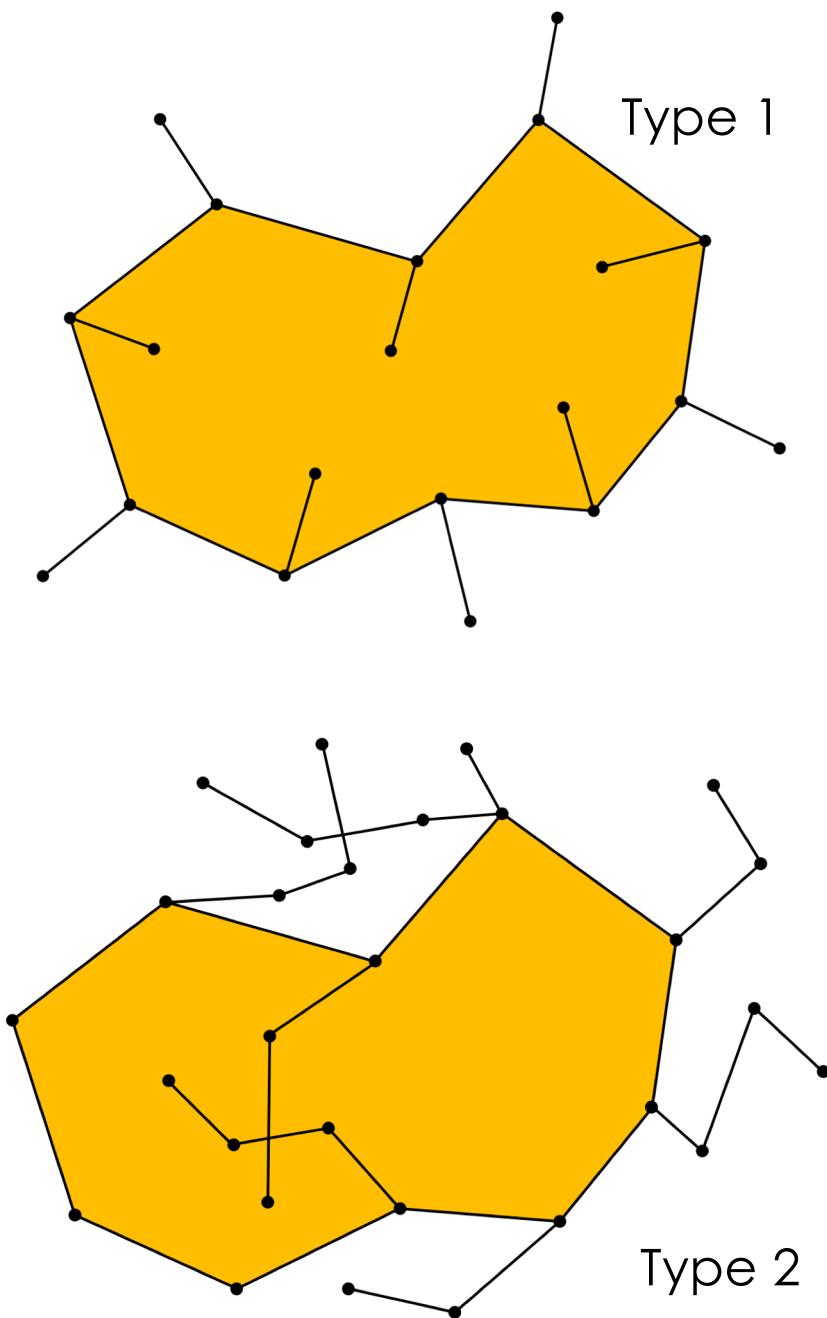


The equivalent Torn Polygon

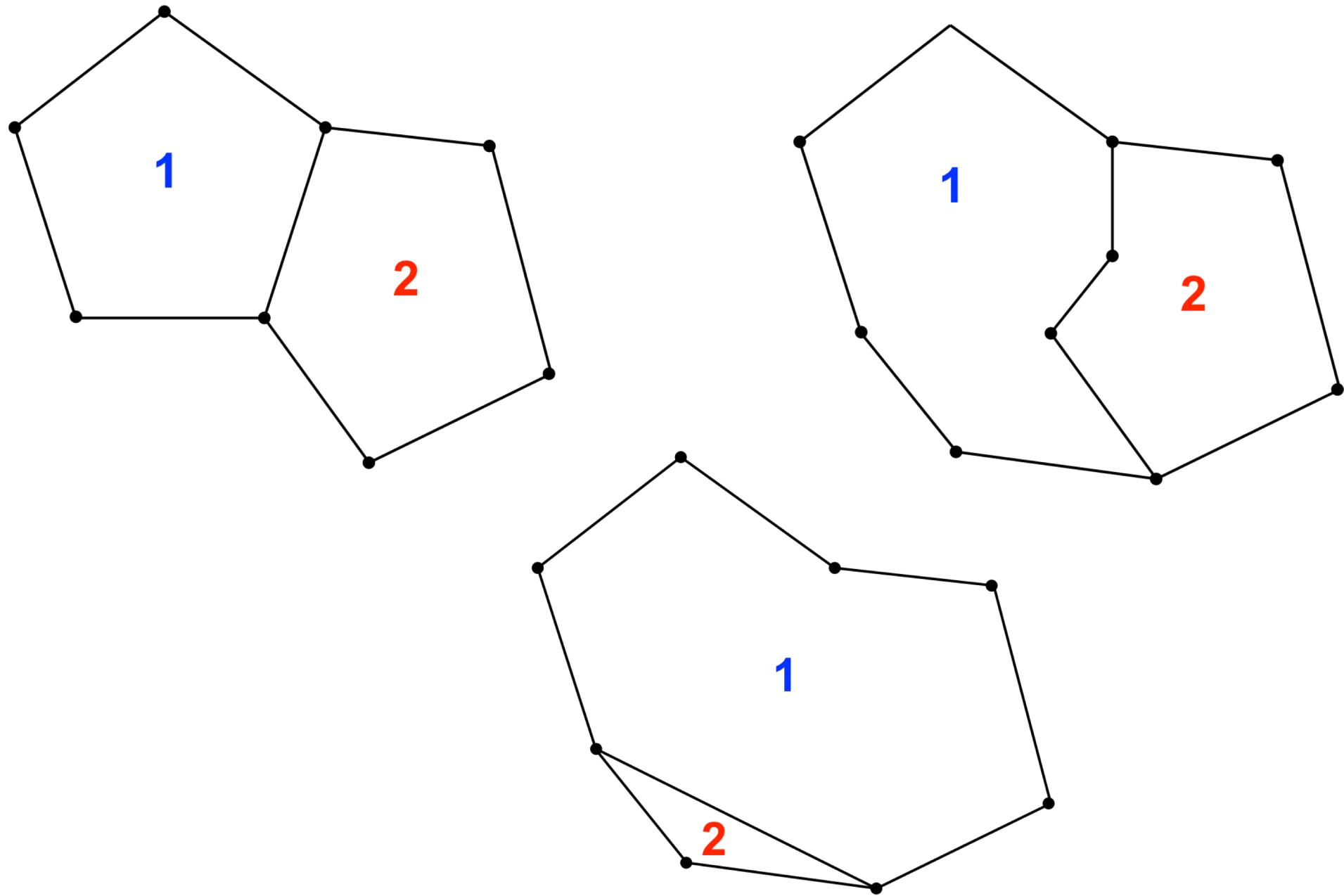
A PSEUDO-POLYGON

A pseudo-polygon P is a set of simple polygons in which each couple of polygons is connected by at most one chain in such a way that none of the polygons is subdivided into two adjacent polygons.

Some Pseudo Polygons



A NON-PSEUDO-POLYGON



2



CUPCARBON FOR IOT SIMULATION



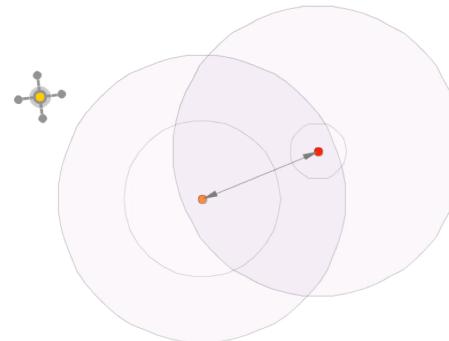
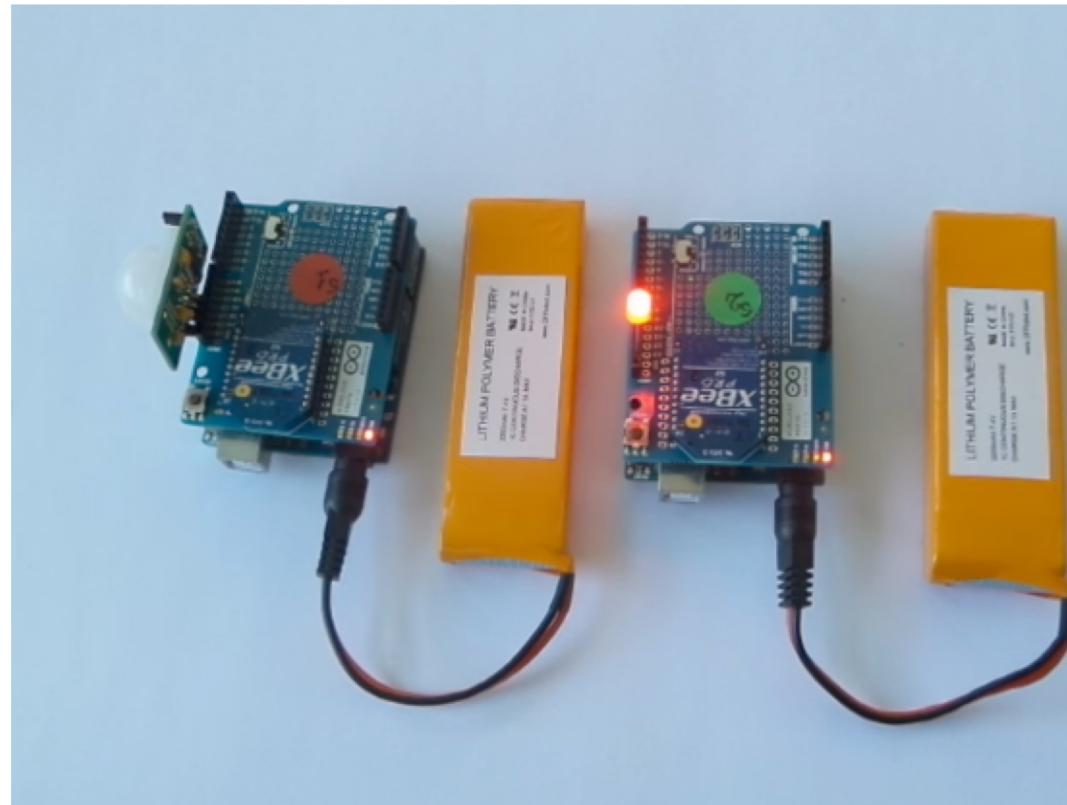
SenScript: Example

S1

```
loop
dreadsensor x
if ($x==1)
  send A
else
  send B
endif
stop
```

S2

```
loop
wait
read x
if ($x==A)
  led 13 1
else
  led 13 0
endif
```



NEW: CupCarbon/SenScript

CupCarbon 3.0.0 (U-One) [dlpcn2.cup]

ROUGE Simulation Time: 51,7251 s

Devices & Objects
Device Parameters
Radio Parameters
Marker Parameters
Simulation Parameters

SenScript

Simulation time: 100.0 s
Speed: 10 ms
Arrow speed: 50 ms

Run Simulation
Stop Simulation
Apply

Symmetrical links
Results
Mobility / Events
ACK
Probability: 1.0

Visibility
Log
Clock Drift
Show

CupCarbon

SEN
SCRIPT

Connexions Nodes Radios Sensors All Details Distances Links Node Arrows

Line Endings Unicode (UTF-8) SenScript Syntax Style

Get Info

dlpcn2_sp.csc

```
1 tab t 10 1
2 set idx -1
3 atget id cid
4 getpos2 cx cy
5 set once 1
6 atnd n
7 set i 0
8 set first 39
9 set nid_ref 0
10 loop
11 if ($cid==$first && ($once==1))
12   set color 2
13   set xmin $cx
14   set ymin $cy
15   minus px $cx 1
16   set py $cy
17   set once 0
18   set type SN
19   set id $first
20   set idx -1
21   data p 0 SN $px $py $xmin
22 else
23   wait
24   read p
25   rdata $p id type
26 end

27 if ($type==AC)
28   data p $cid CS $cx $cy
```

Lines: 92 Chars: 1915 Location: 1915 Line: 92

2 KB

Simulation time: 64 N = 34 RT = 13,6140 s

Presentation of CupCarbon

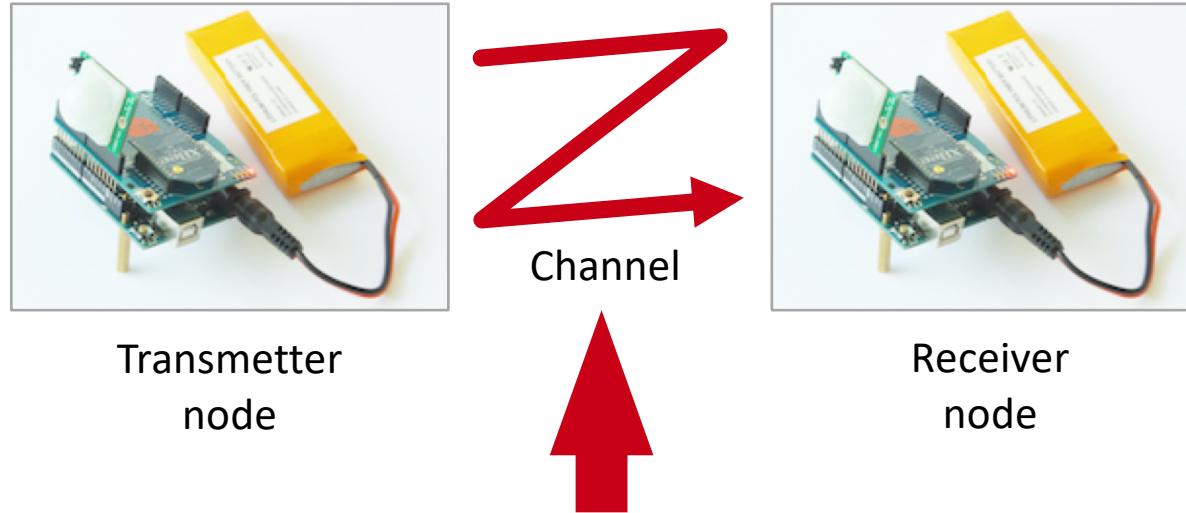
CupCarbon is a WSN simulator dedicated to Smart Cities and Internet of Things

- Its objective is to **design, visualize, debug, simulate and validate distributed algorithms,**
- It can be used for educational and **scientific projects,**
- It offers an **ergonomic** and easy to use interface using the **OpenStreetMap (OSM)** framework to deploy sensors directly on the map,

Presentation of CupCarbon

- It includes a script called **SenScript** which allows to **program** and to **configure** each sensor node individually,
- It is possible to **generate codes** for **hardware platforms** such as Arduino/XBee,
- CupCarbon offers the possibility to **simulate many algorithms** and scenarios in **several steps**,
- The **energy consumption** can be **modelled**, calculated and displayed as a function of the simulated time,

Modelling the radio propagation channel



Existing WSN simulators:

- NS-2/3
- OMNET
- TOSSIM
- OPNET
- ...

→ Classical models used in WSN

- Free space model
- Two rays ground model
- Log normal model
- ...

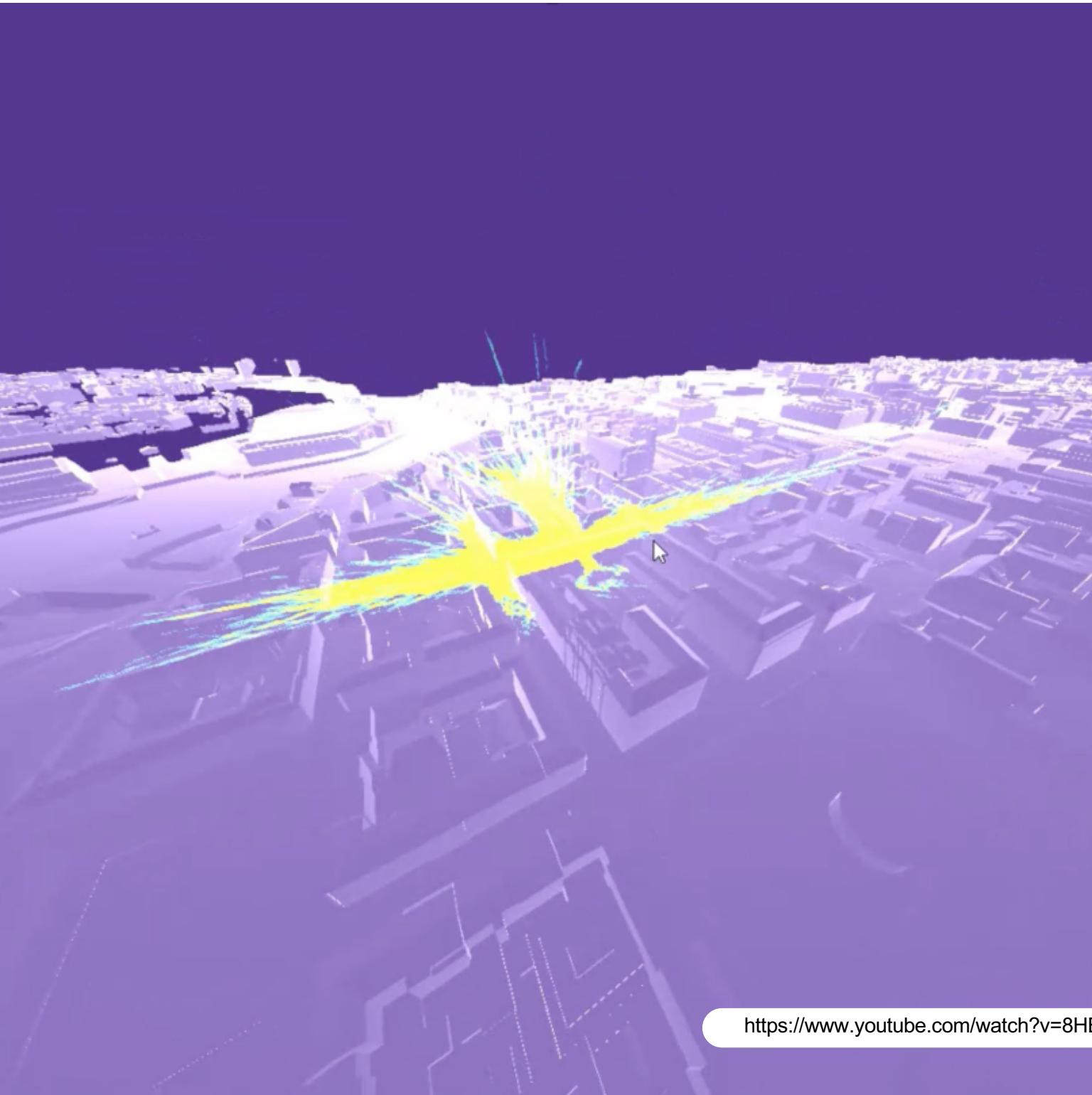
+ → Very simple and fast models ;
- → Not Site specific ;
- → Does not take into account multipath propagation phenomenon.

Modelling the radio propagation channel : Visibility tree



Modelling the radio propagation channel : Visibility tree



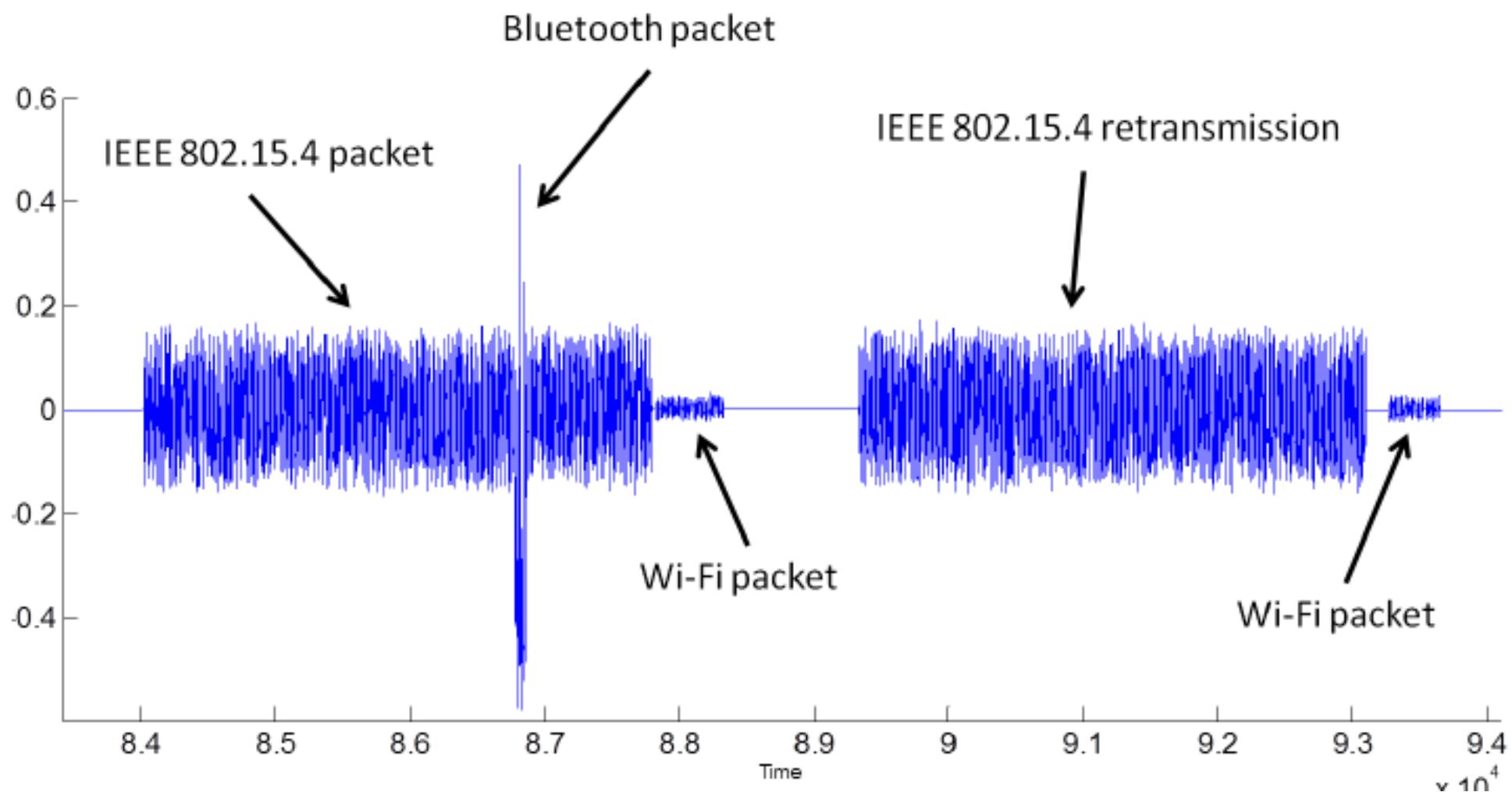


<https://www.youtube.com/watch?v=8HBhjVNIZgl>

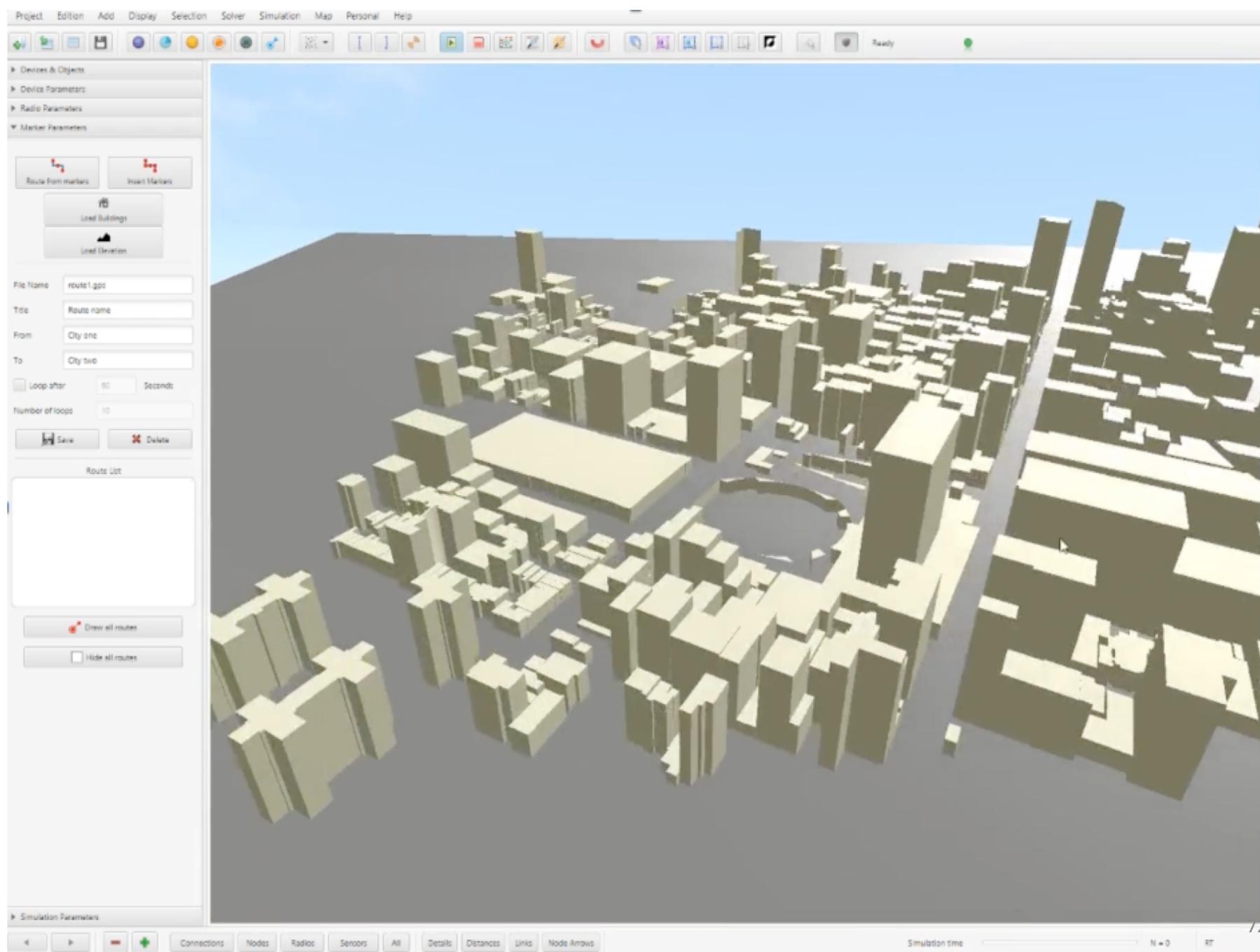


Interference

- Interference: noise characterization



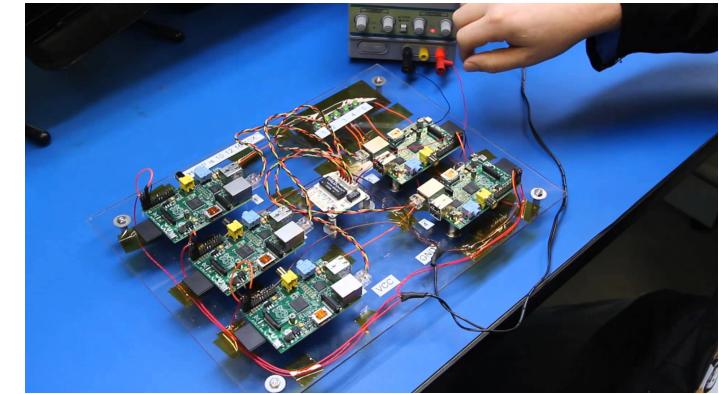
CupCarbon 3D (Demo)



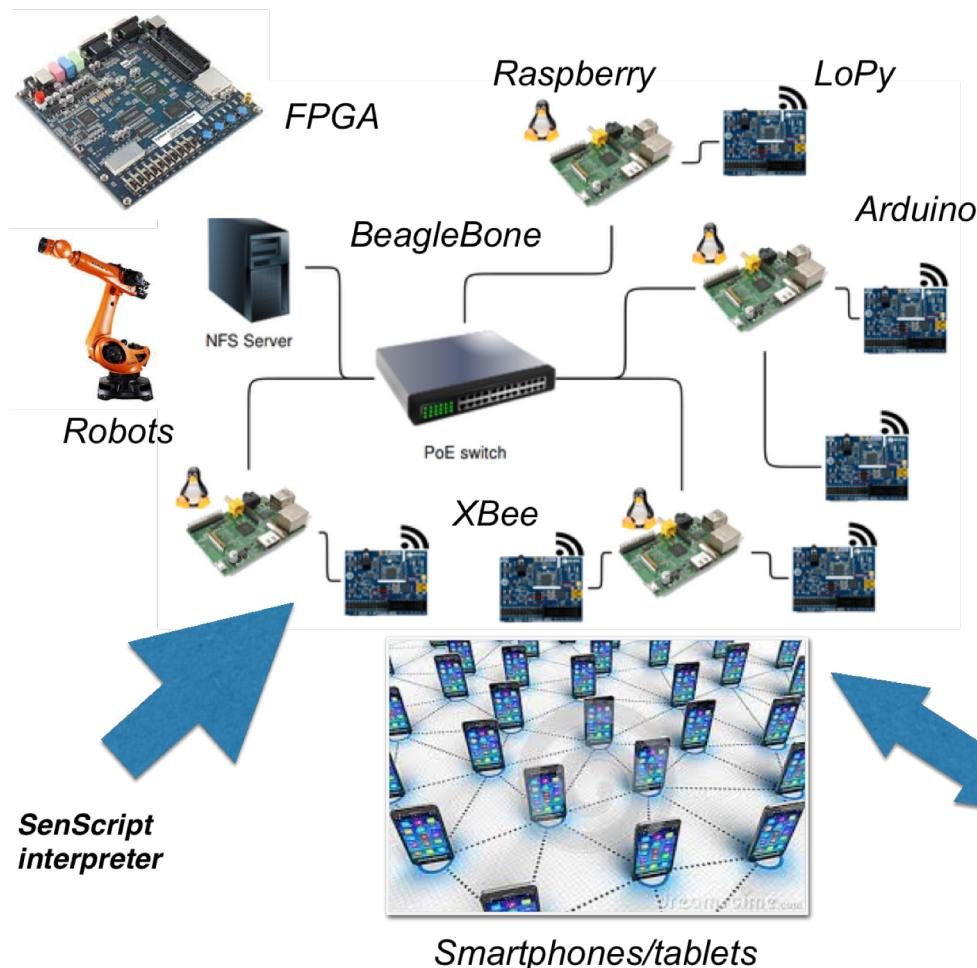
From simulation to real network



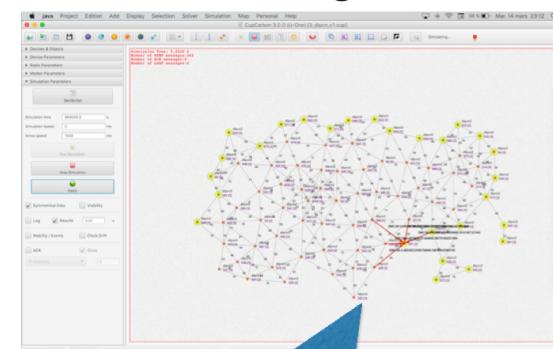
CupCarbonLab



- CupCarbon-Lab: IoT Emulator
- Cybersecurity

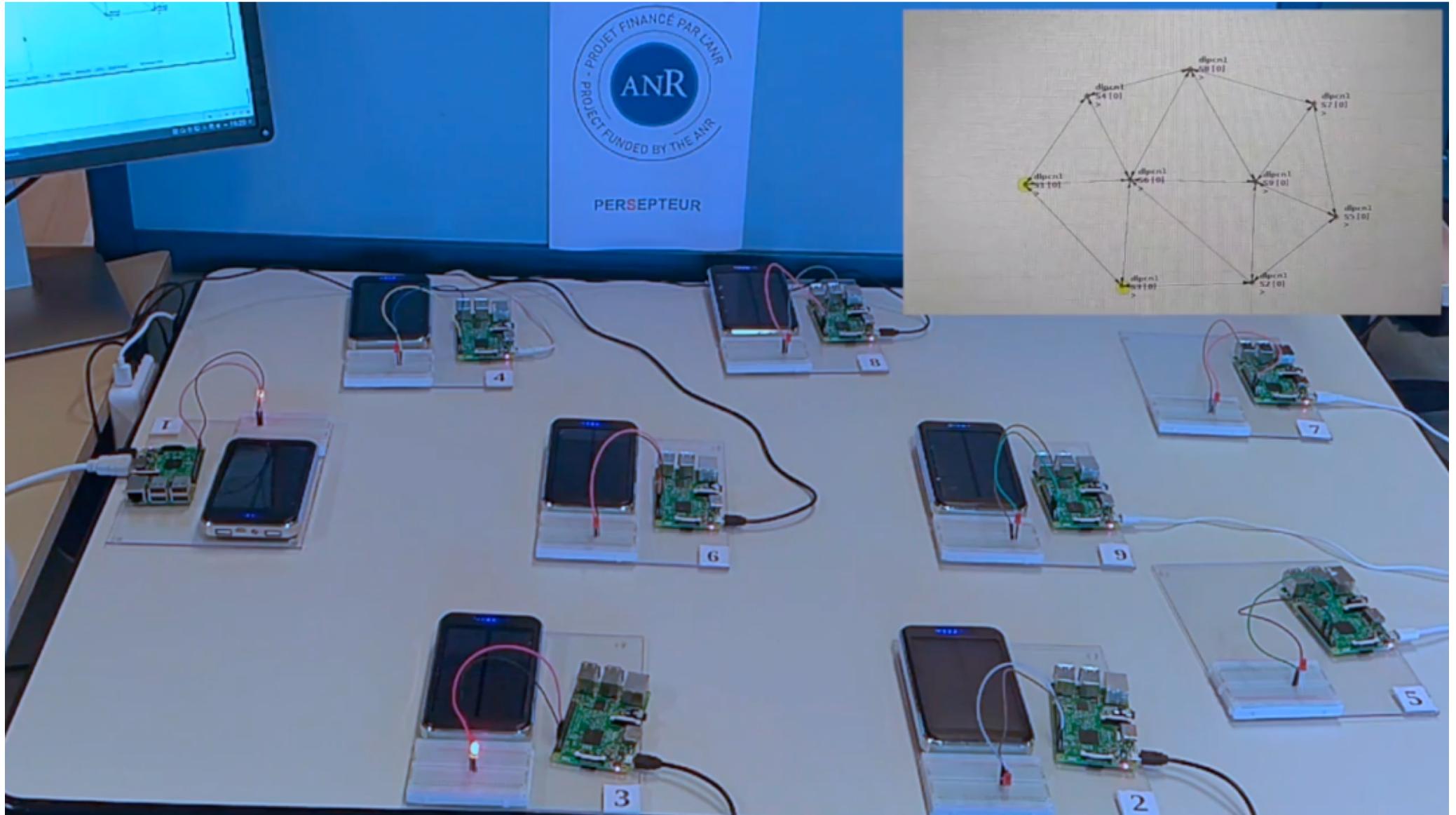


Simulation
Programming
Code injection
Monitoring

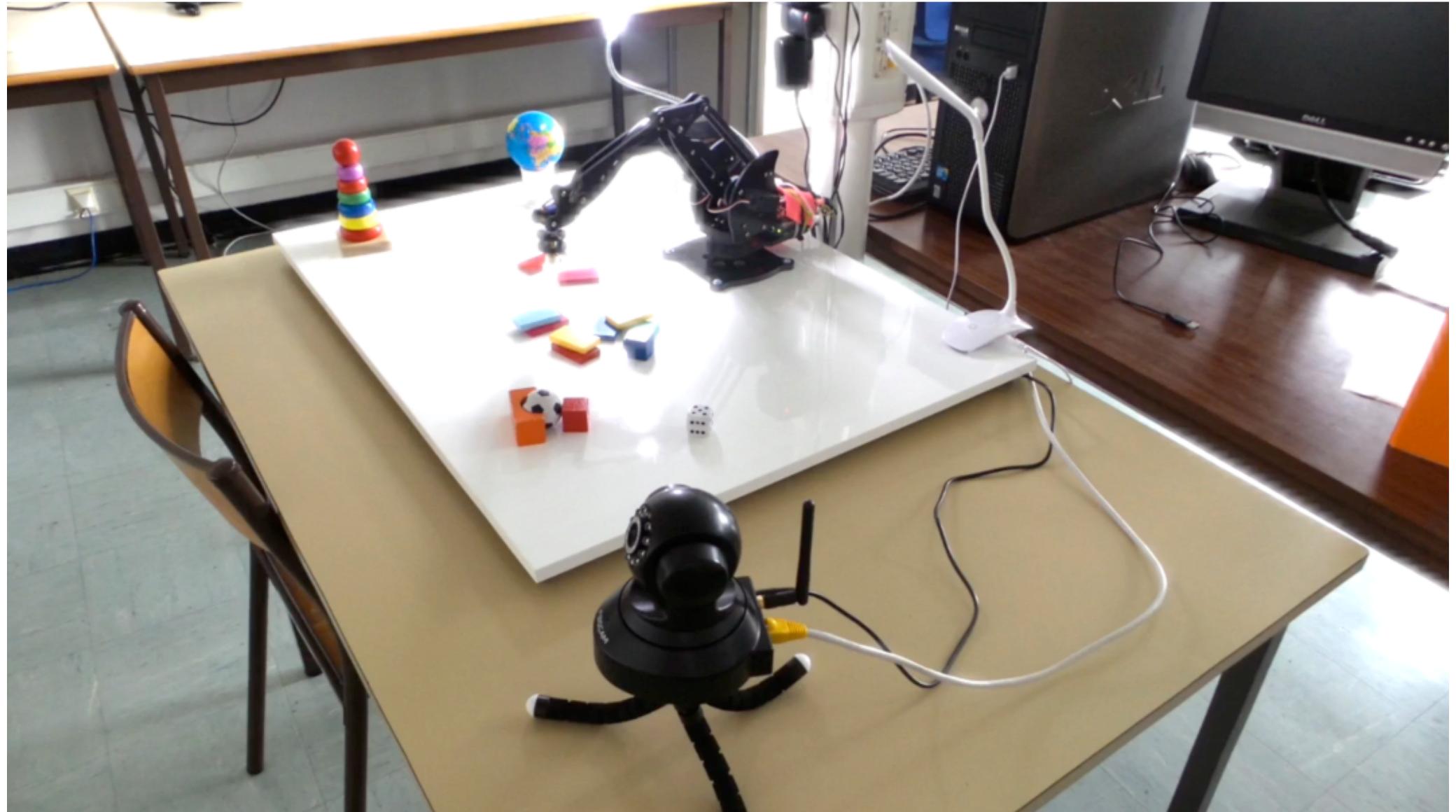


Web Service

BCS (By Card) Simulation – Direct Network Reconfiguration



BCS (By Card) Simulation – Remote Monitoring IoRGame Platform



CupCarbon : Quelques stats

450 visites / mois



| Pays | Utilisateurs | % Utilisateurs |
|----------------------|--------------|----------------|
| 1. 🇮🇳 India | 2 550 | 24,21 % |
| 2. 🇫🇷 France | 1 054 | 10,01 % |
| 3. 🇺🇸 United States | 969 | 9,20 % |
| 4. 🇬🇧 United Kingdom | 402 | 3,82 % |
| 5. 🇧🇷 Brazil | 379 | 3,60 % |
| 6. 🇵🇰 Pakistan | 289 | 2,74 % |
| 7. 🇮🇷 Iran | 279 | 2,65 % |
| 8. 🇮🇹 Italy | 251 | 2,38 % |
| 9. 🇩🇿 Algeria | 229 | 2,17 % |
| 10. 🇰🇷 South Korea | 220 | 2,09 % |

CupCarbon : Quelques stats

150 Téléchargements / mois



| Pays | Utilisateurs | % Utilisateurs |
|----------------------|--------------|----------------|
| 1. 🇮🇳 India | 923 | 27,60 % |
| 2. 🇺🇸 United States | 233 | 6,97 % |
| 3. 🇬🇧 United Kingdom | 176 | 5,26 % |
| 4. 🇫🇷 France | 155 | 4,64 % |
| 5. 🇵🇰 Pakistan | 116 | 3,47 % |
| 6. 🇻🇪 Ecuador | 102 | 3,05 % |
| 7. 🇧🇷 Brazil | 101 | 3,02 % |
| 8. 🇮🇹 Italy | 101 | 3,02 % |
| 9. 🇮🇷 Iran | 94 | 2,81 % |
| 10. 🇩🇿 Algeria | 85 | 2,54 % |

CupCarbon : Quelques stats

100 Téléchargements du guide utilisateurs / mois



| Pays | Utilisateurs | % Utilisateurs |
|----------------------|--------------|----------------|
| 1. 🇮🇳 India | 570 | 24,52 % |
| 2. 🇺🇸 United States | 221 | 9,51 % |
| 3. 🇫🇷 France | 174 | 7,48 % |
| 4. 🇧🇷 Brazil | 90 | 3,87 % |
| 5. 🇬🇧 United Kingdom | 90 | 3,87 % |
| 6. 🇵🇰 Pakistan | 71 | 3,05 % |
| 7. 🇪🇸 Spain | 64 | 2,75 % |
| 8. 🇩🇿 Algeria | 58 | 2,49 % |
| 9. 🇻🇪 Ecuador | 56 | 2,41 % |
| 10. 🇴️ Iran | 53 | 2,28 % |

CupCarbon : Quelques stats



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Evaluation of CupCarbon Network Simulator for Wireless Sensor Networks

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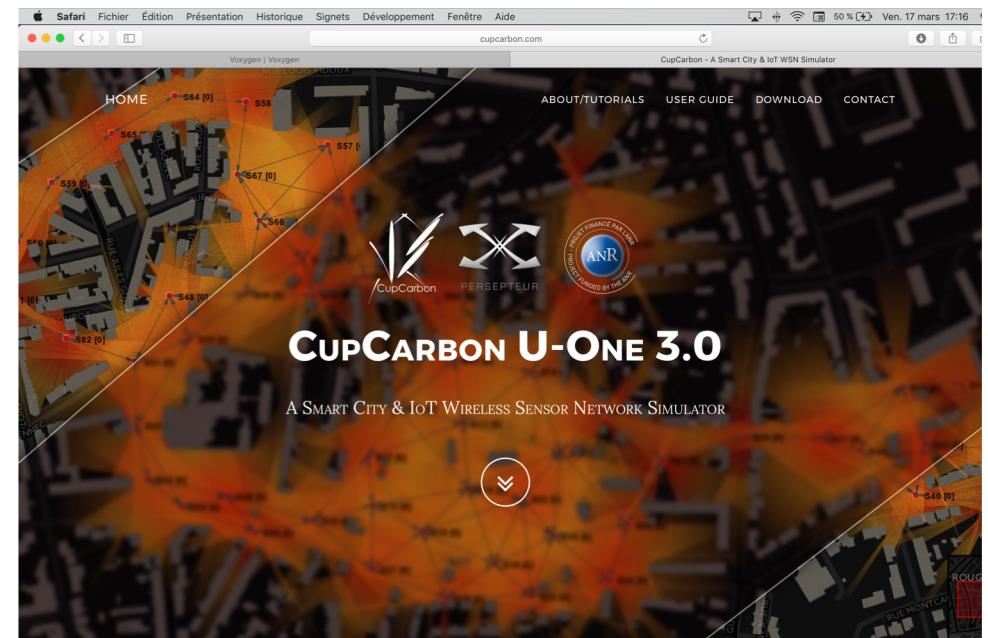
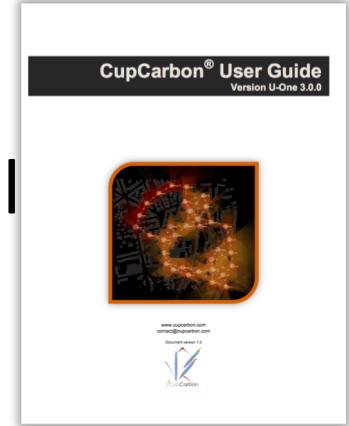
DOI: 10.5296/npa.v10i2.13201

URL: <https://doi.org/10.5296/npa.v10i2.13201>

CupCarbon web site

www.cupcarbon.com

- Presentation of the simulator (video)
- Source code (Git Hub) and Jar file (executable)
- User Guide
- 17 (Basic) + 5 (Advanced) tutorials
- 2 tutorials (video)
- Newsletter
- Forum
- Facebook page



3

POLYGON HULLS & LEADER ELECTION

- 3.1. *LPCN Algorithm*
- 3.2. *D-LPCN Algorithm*
- 3.3. *Leader Election*

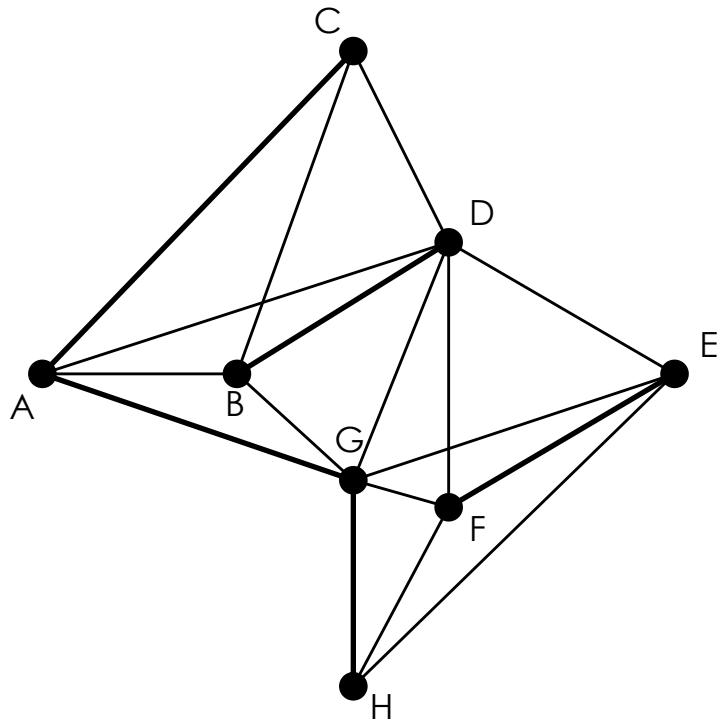
Ahcène Bounceur, Trouver l'Enveloppe Polygonale d'un Réseau de Capteurs Sans Fil : Théorie et Pratique, séminaire, GDR-RO, CEA, Paris, May 2015.

Ahcène Bounceur, Madani Bezoui, Reinhardt Euler and Marc Sevaux, Finding the Boundary Nodes of a Wireless Sensor Network Without Conditions on the Starting Node, Invited talk, EURO 2016, Poznan.

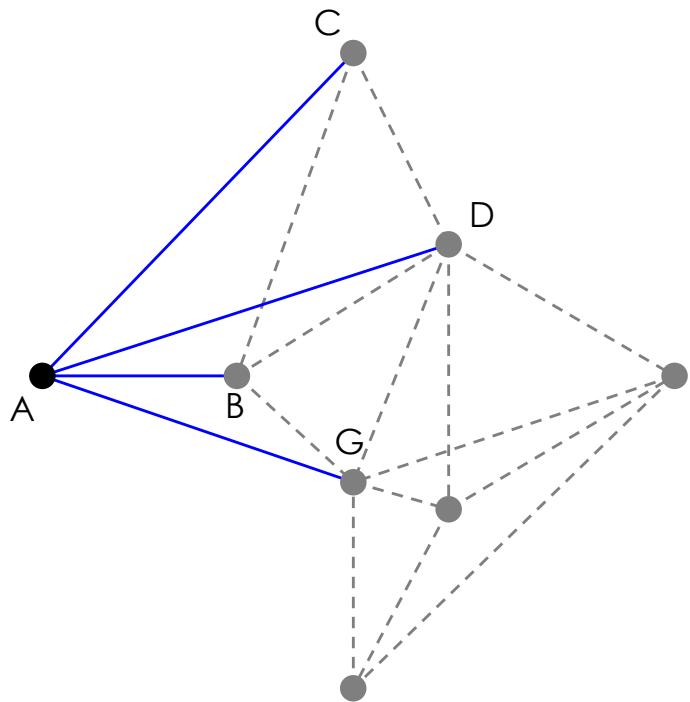
3.1

LPCN ALGORITHM

LPCN Algorithm

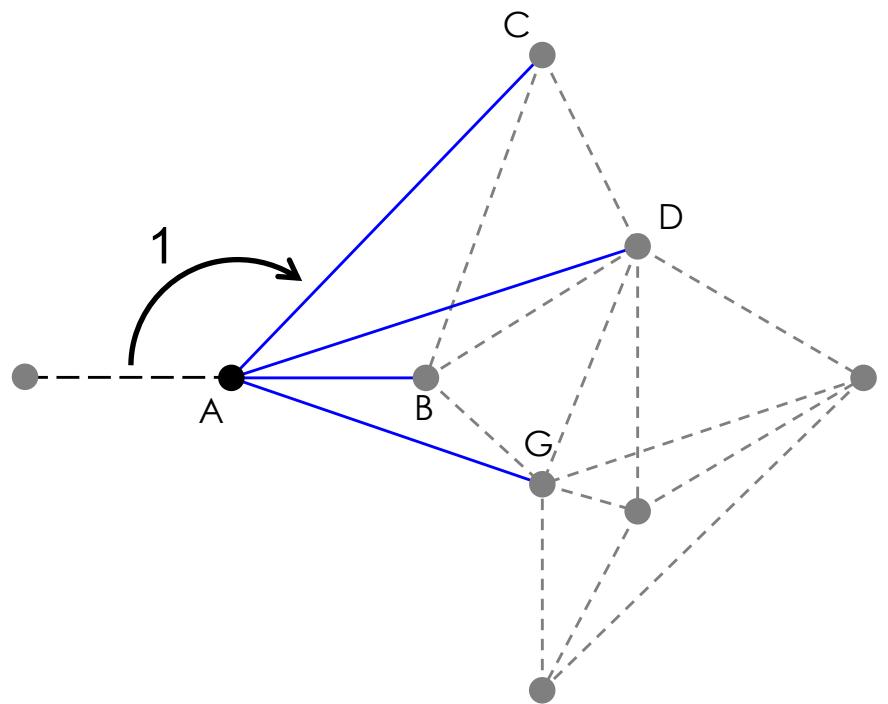


LPCN Algorithm



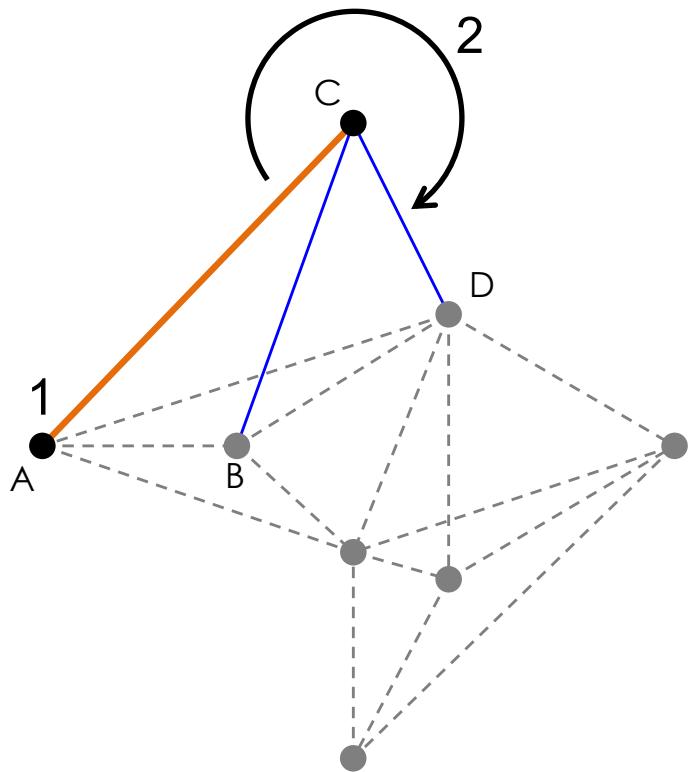
[A]

LPCN Algorithm



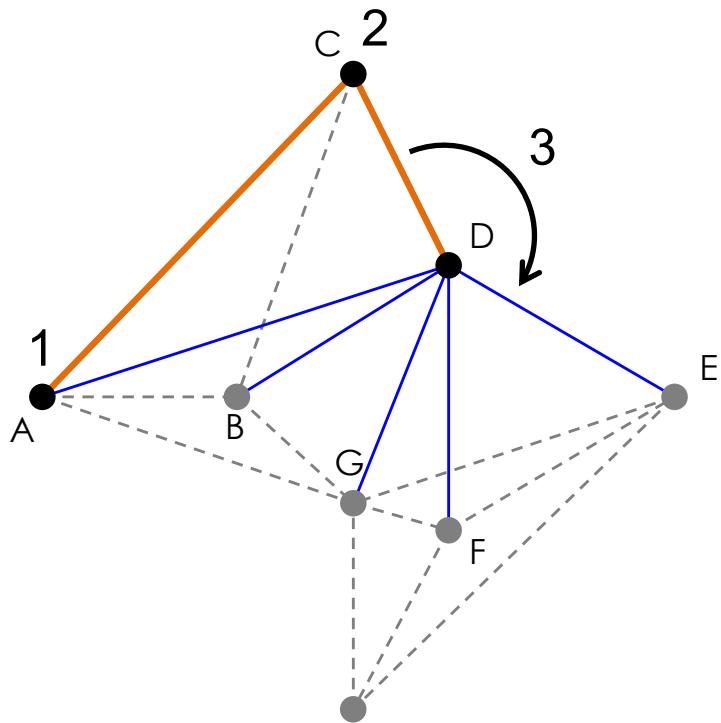
[A]

LPCN Algorithm



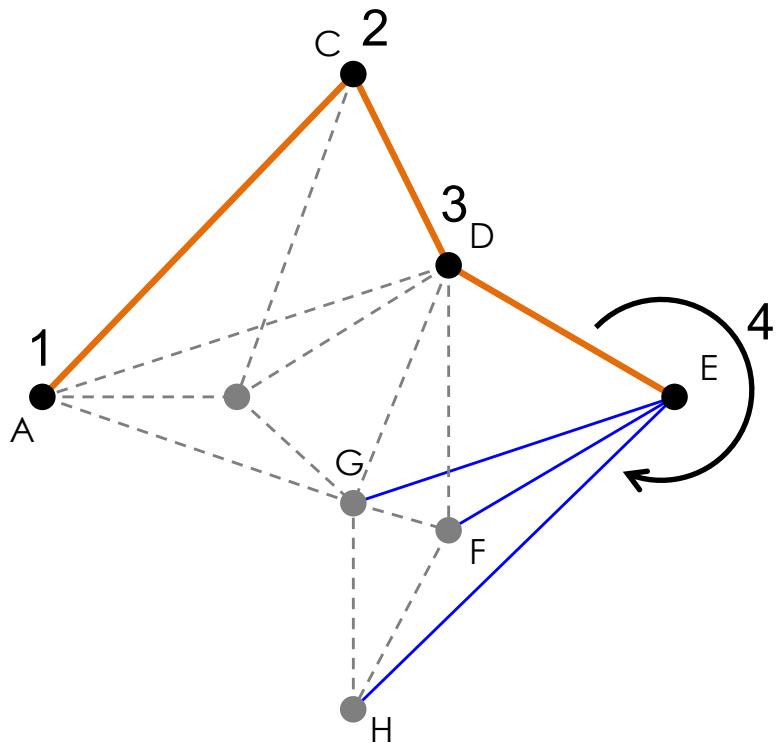
[A, C]

LPCN Algorithm



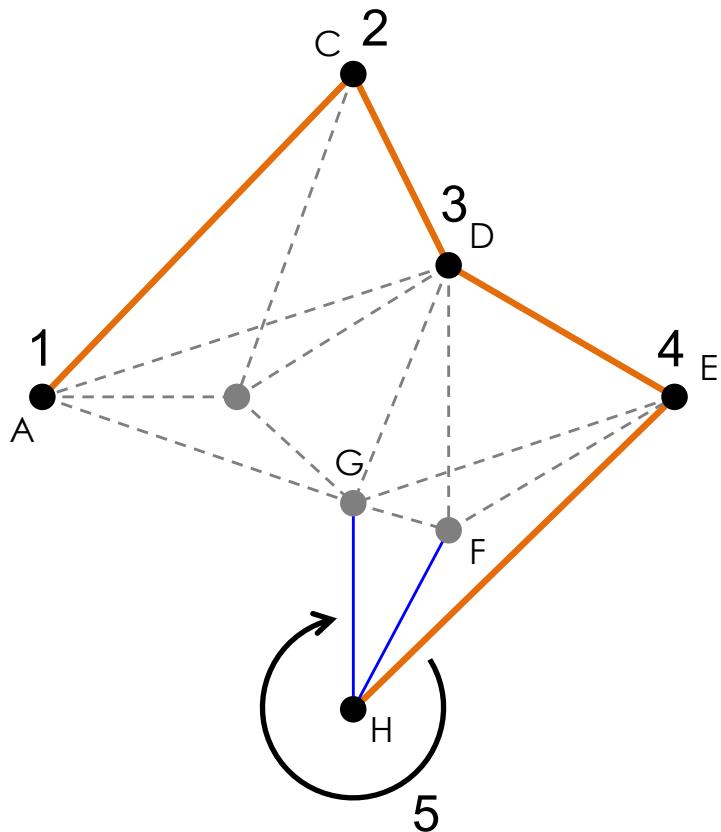
[A, C, D]

LPCN Algorithm



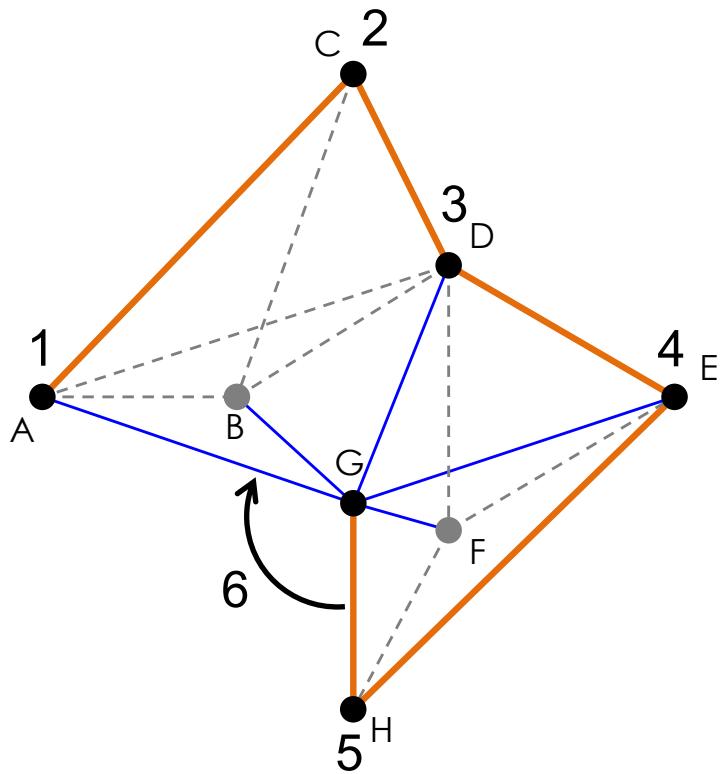
[A, C, D, E]

LPCN Algorithm



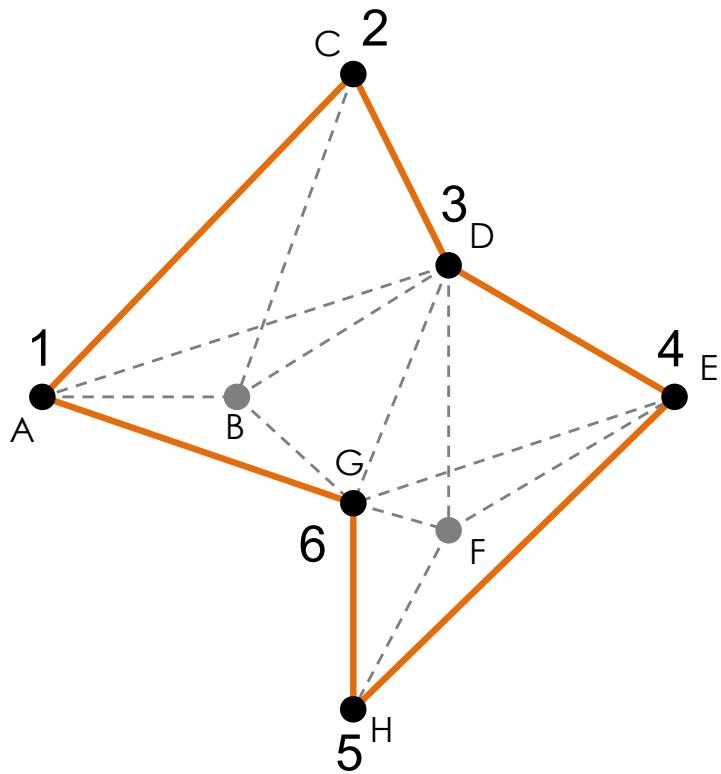
[A, C, D, E, H]

LPCN Algorithm



[A, C, D, E, H, G]

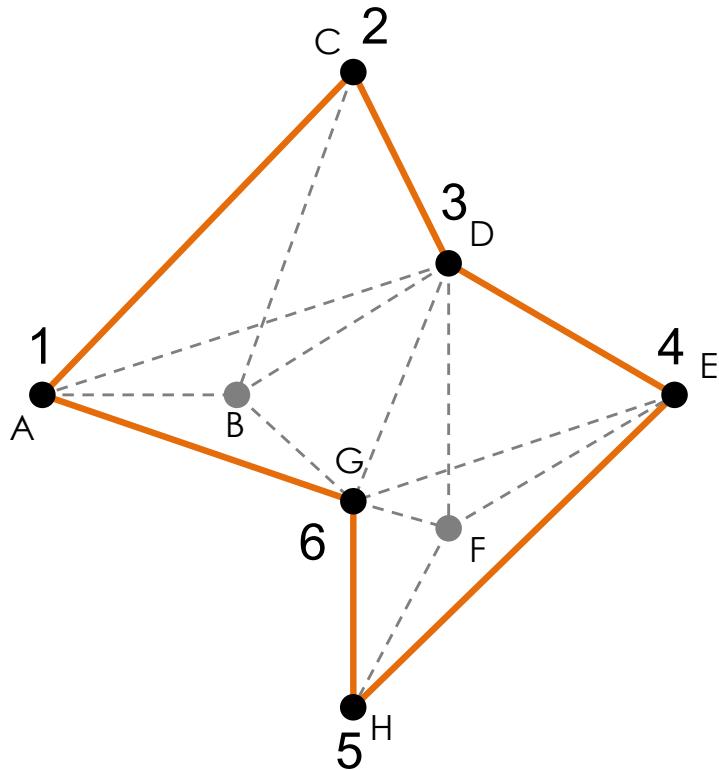
LPCN Algorithm



[A, C, D, E, H, G, A]

LPCN Algorithm

[A, C, D, E, H, G, A]



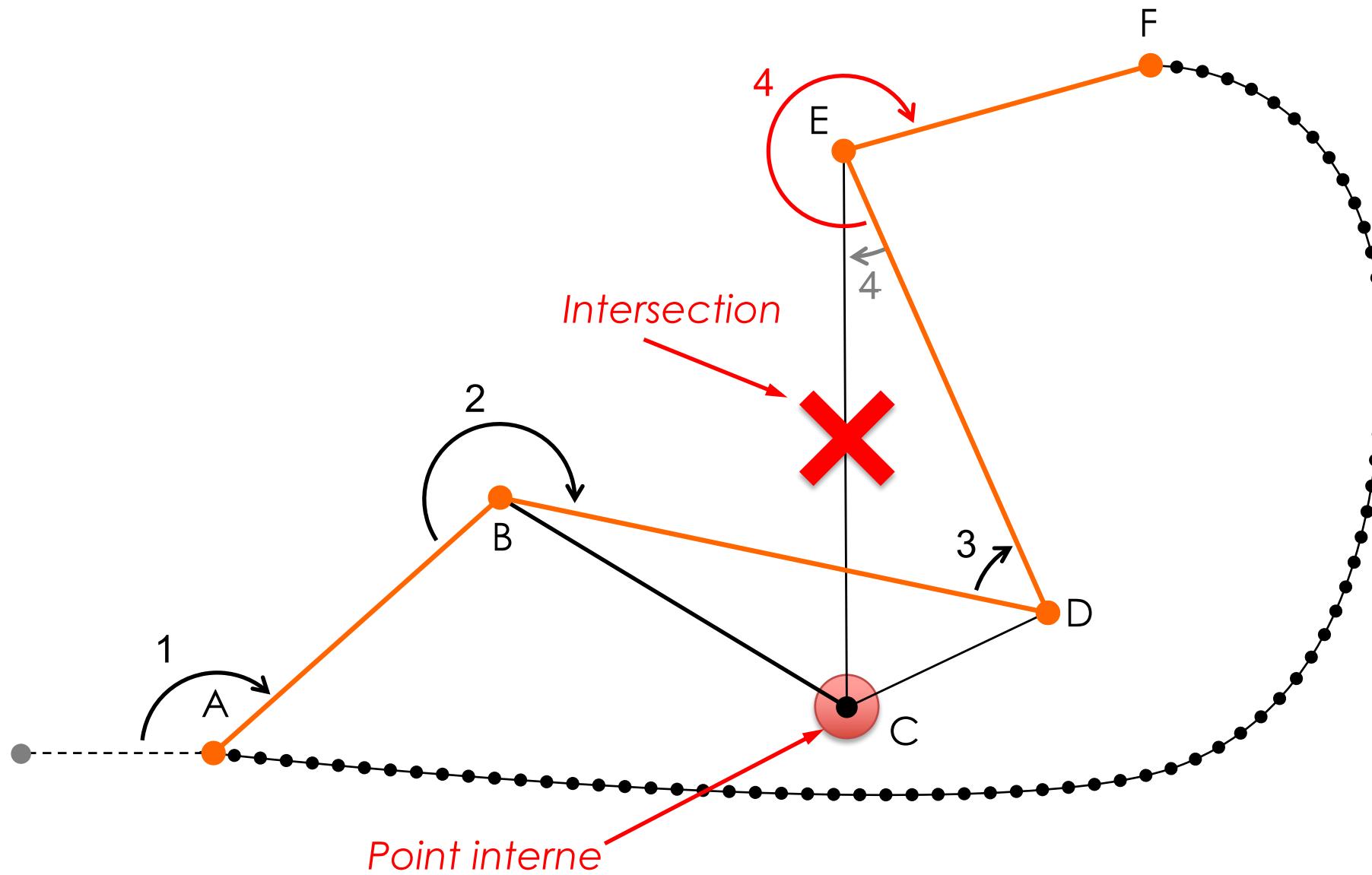
Version 1

```
2:  $P_c \leftarrow$  point having the minimum x-coordinate
3:  $\mathbb{B}_V \leftarrow [P_c]$ 
4:  $P_f \leftarrow P_c$ 
5:  $P_p \leftarrow$  fictive point situated on the left of  $P_f$ 
6: repeat
7:      $P_v = \operatorname{argmin}_{P_j \in N(P_c)} \{pa(P_p, P_c, P_j)\}$ 
8:      $\mathbb{B}_V \leftarrow [\mathbb{B}_V, P_v]$ 
9:      $P_p \leftarrow P_c$ 
10:     $P_c \leftarrow P_v$ 
11:   until  $P_v = P_f$ 
12:   return  $\mathbb{B}_V$ 
13: end procedure
```

LPCN Algorithm

- Version 1
 - Complexity: **O(kh)**
 - **k**: Maximum degree of the graph
 - **h**: Number of the border nodes

LPCN Algorithm: with a Boat graph



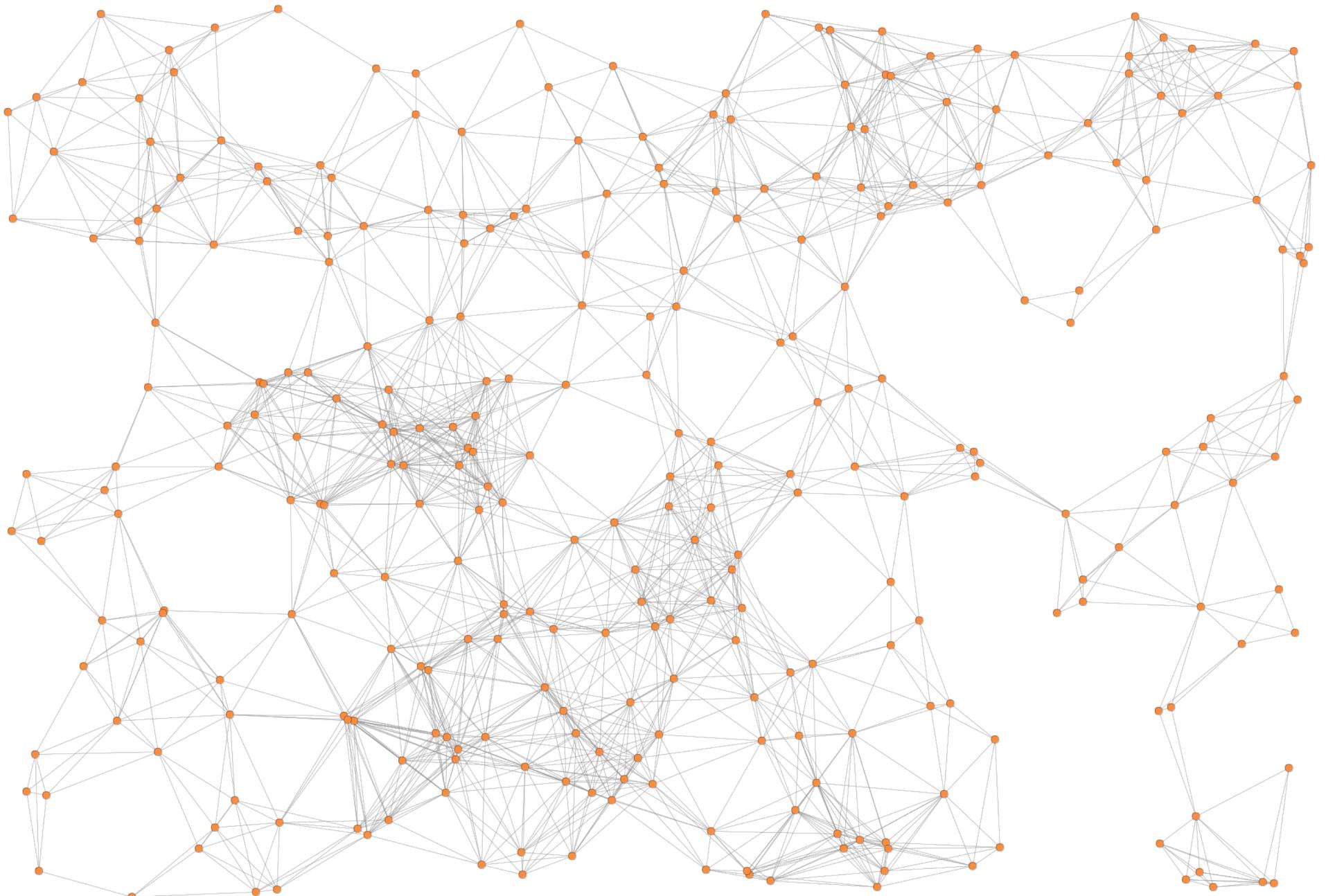
LPCN Algorithm: Version 2

```
1: procedure LPCN( $V, E$ )
2:    $P_0 \leftarrow P_c \leftarrow$  point having the minimum  $x$ -coordinate
3:    $P_p \leftarrow$  fictive point situated in the left of  $P_c$ 
4:    $\mathbb{B}_V \leftarrow \{P_c\}$ 
5:    $\mathbb{B}_E \leftarrow \emptyset$ 
6:    $once \leftarrow true$ 
7:   repeat
8:      $\mathbb{A} \leftarrow \{P \in N(P_c) / \mathbb{B}_E \cap \{\{P_c, P\}\} = \emptyset\}$ 
9:      $P_{min} \leftarrow \underset{P \in \mathbb{A}}{\operatorname{argmin}} \{\varphi(P_p, P_c, P)\}$ 
10:     $\mathbb{B}_V \leftarrow \mathbb{B}_V \cup \{P_{min}\}$ 
11:     $\mathbb{B}_E \leftarrow \mathbb{B}_E \cup \{\{P_c, P_{min}\}\}$ 
12:     $P_p \leftarrow P_c$ 
13:     $P_c \leftarrow P_{min}$ 
14:    if ( $once = true$ ) then
15:       $once \leftarrow false$ 
16:       $P_{first} \leftarrow P_{min}$ 
17:    end if
18:    until ( $(P_c = P_0)$  and ( $P_{min} = P_{first}$ ))
19:    return  $\mathbb{B}_V, \mathbb{B}_E$ 
20: end procedure
```

LPCN Algorithm

- Version 1
 - Complexity: **O(kh)**
 - **k**: Maximum degree of the graph
 - **h**: Number of the border nodes
- Version 2
 - Complexity: **O(kh²)**
 - **k**: Maximum degree of the graph
 - **h(h+1)/2**: Number of the border nodes by considering the intersection detection

LPCN Algorithm: Demo 1



LPCN Algorithm: Demo 2



3.2

D-LPCN ALGORITHM

D-LPCN Algorithm: Messages

- 3 types of messages :

AC : Ask for Coordinates

CS : Coordinate Sending

SN : Select Node

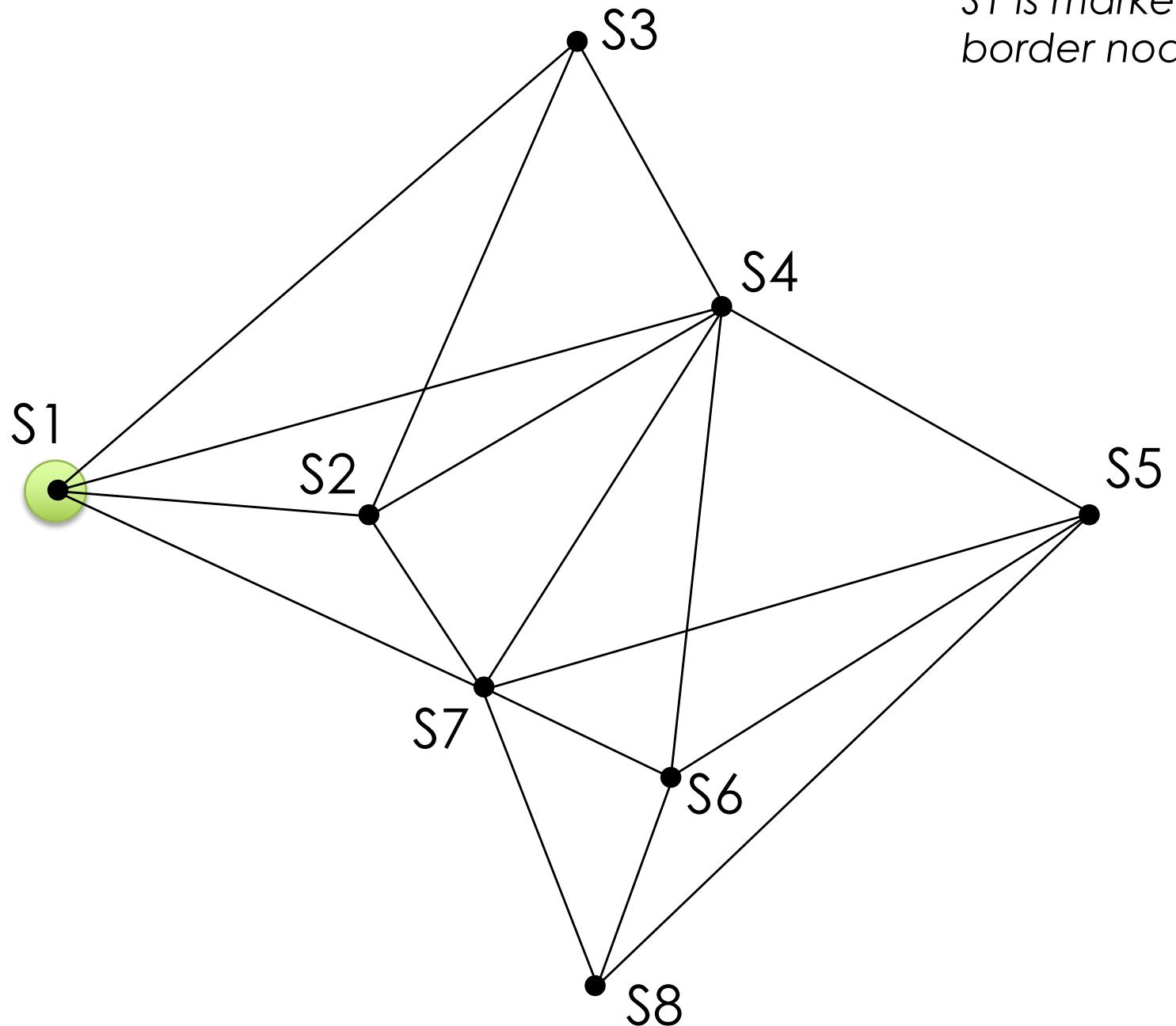
- Messages format:

Message AC : **id | AC**

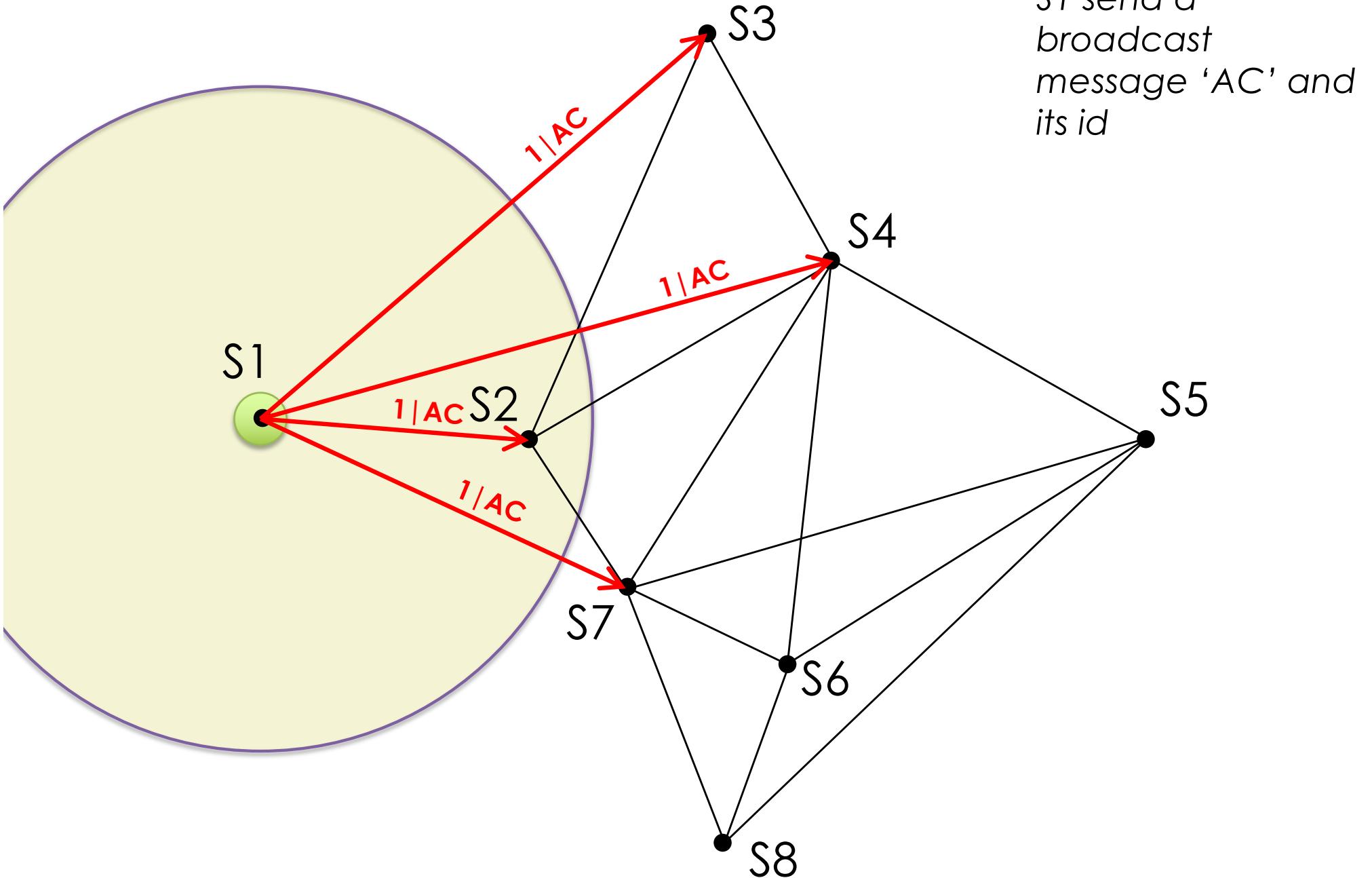
Message CS : **id | CS | x,y**

Message SN : **id | SN | x,y**

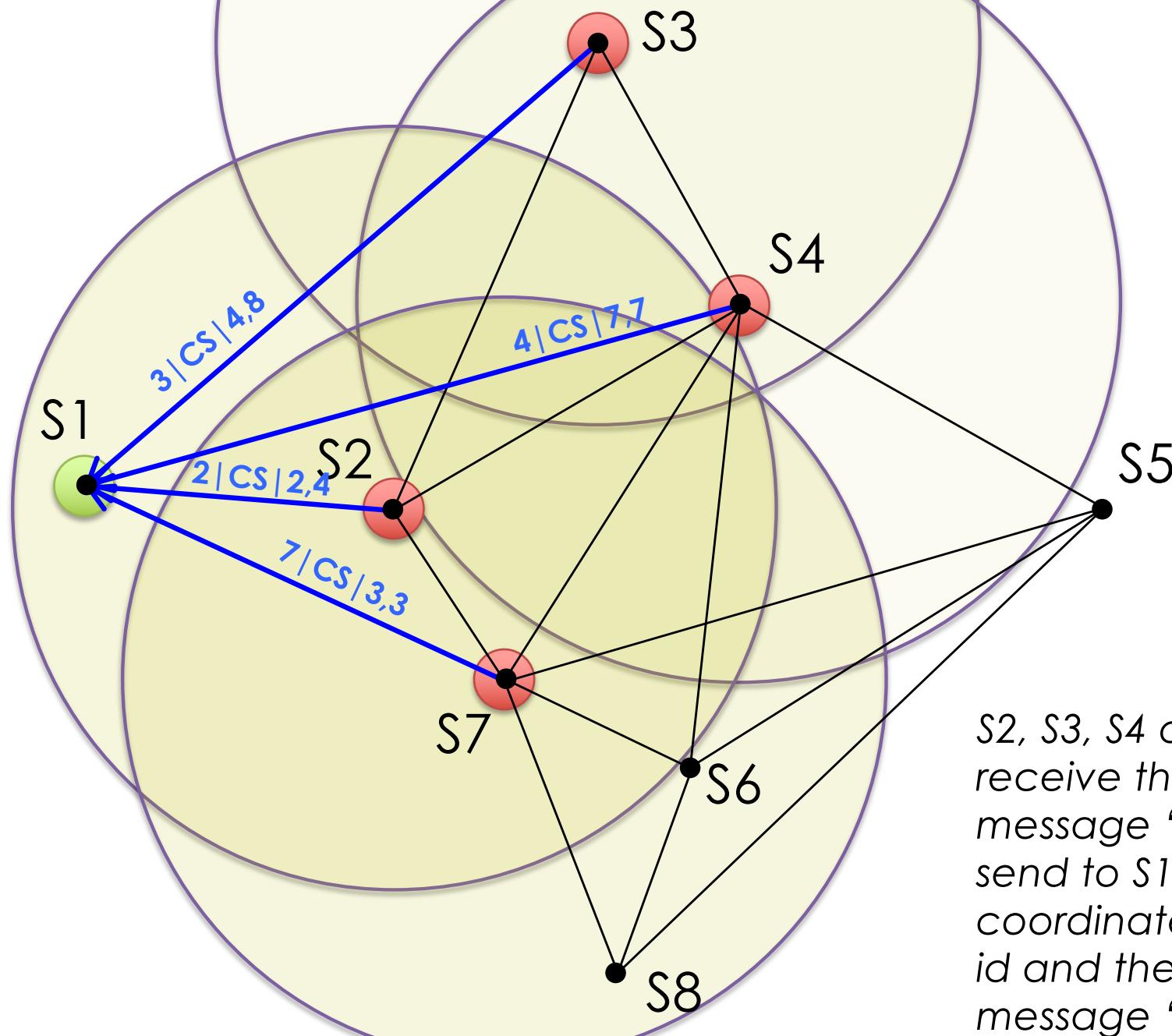
D-LPCN Algorithm



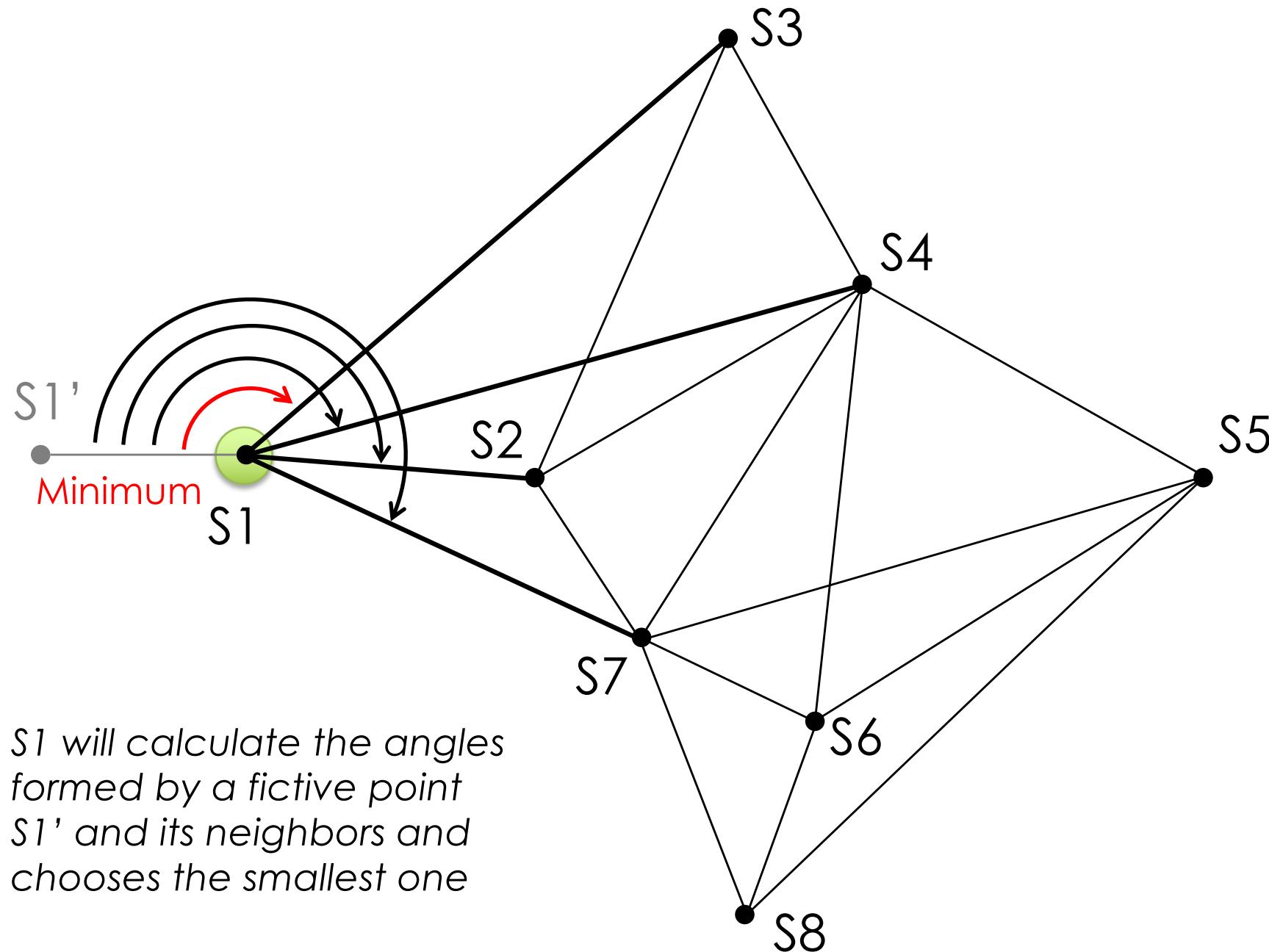
D-LPCN Algorithm



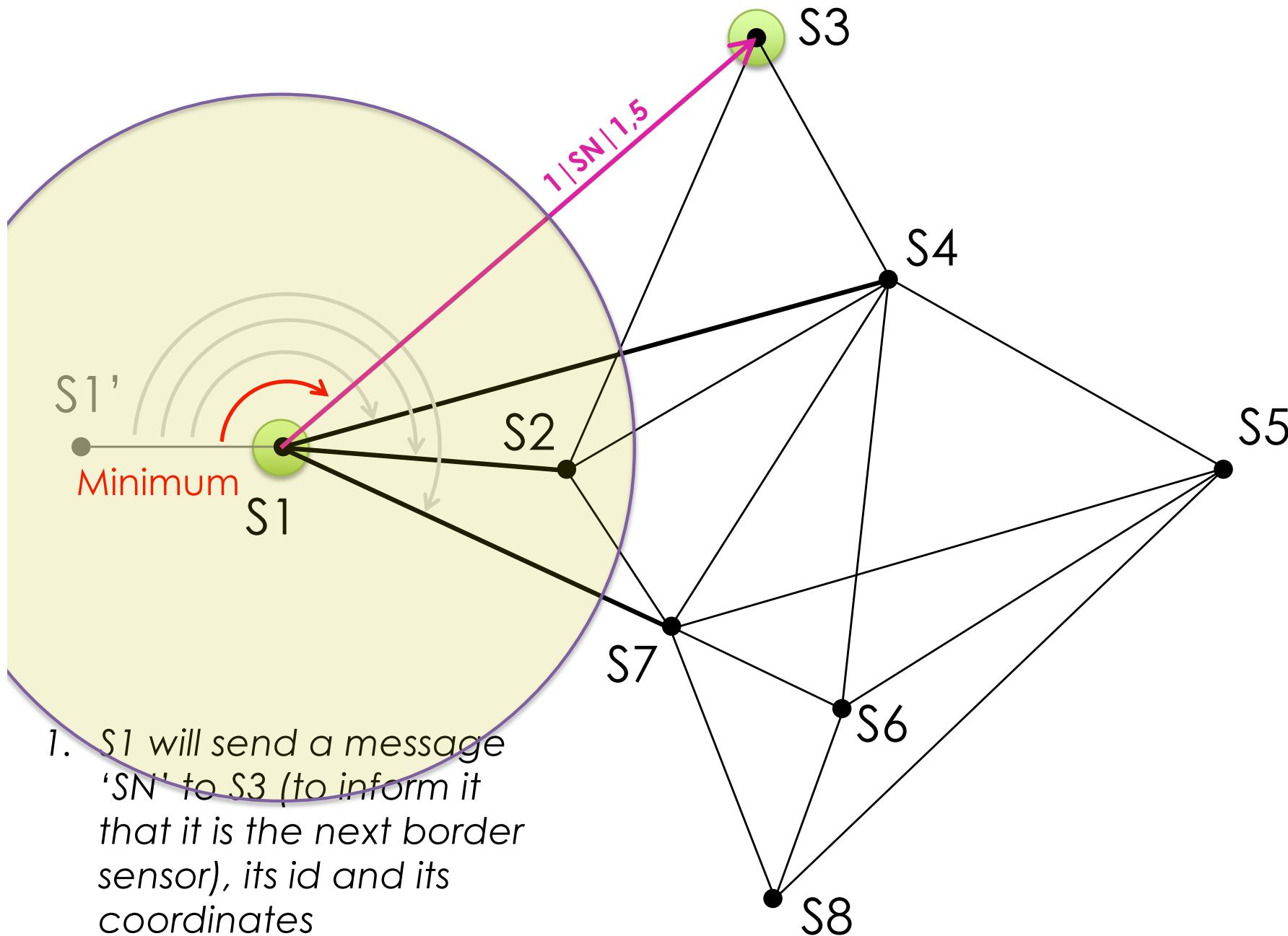
D-LPCN Algorithm



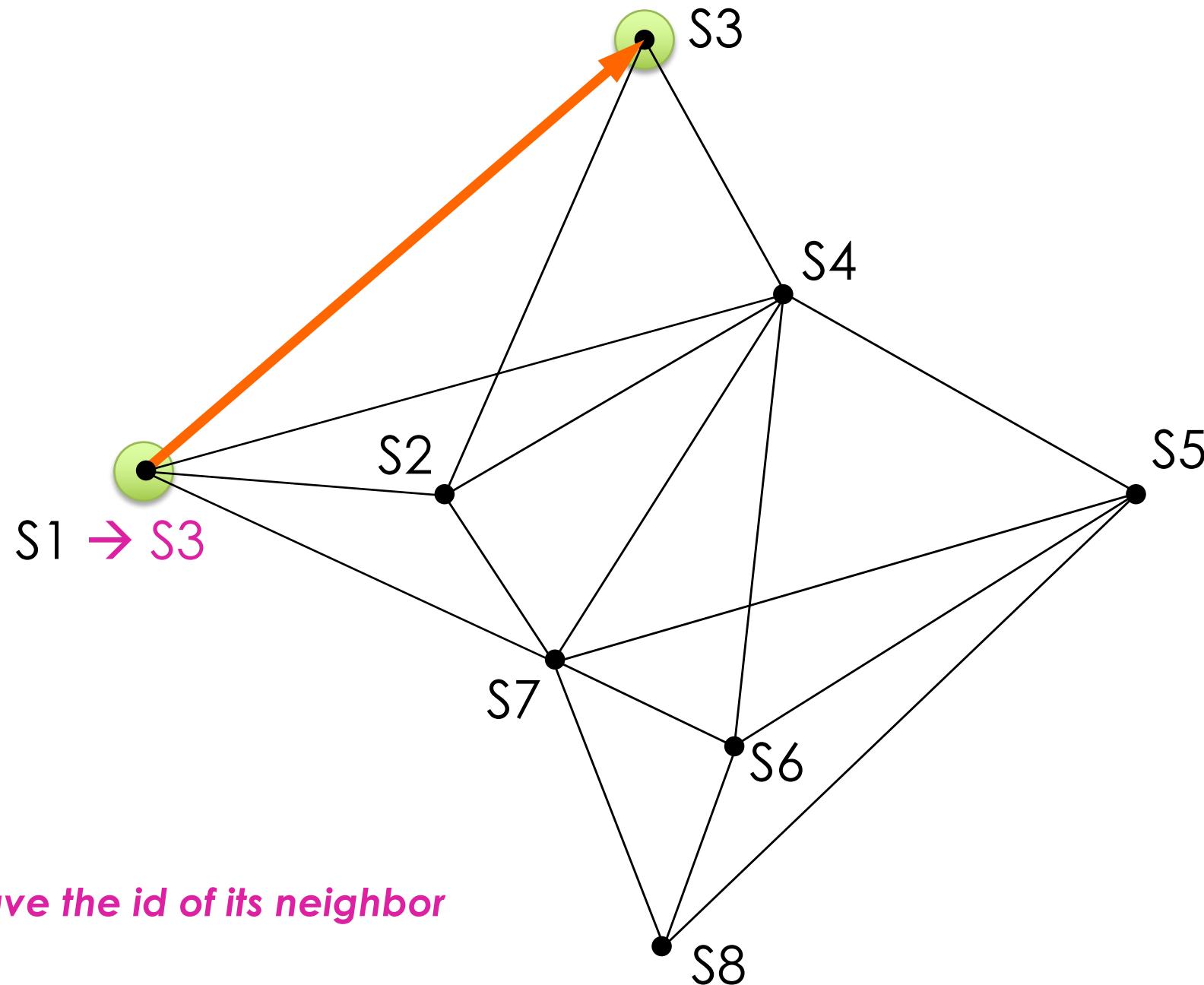
D-LPCN Algorithm



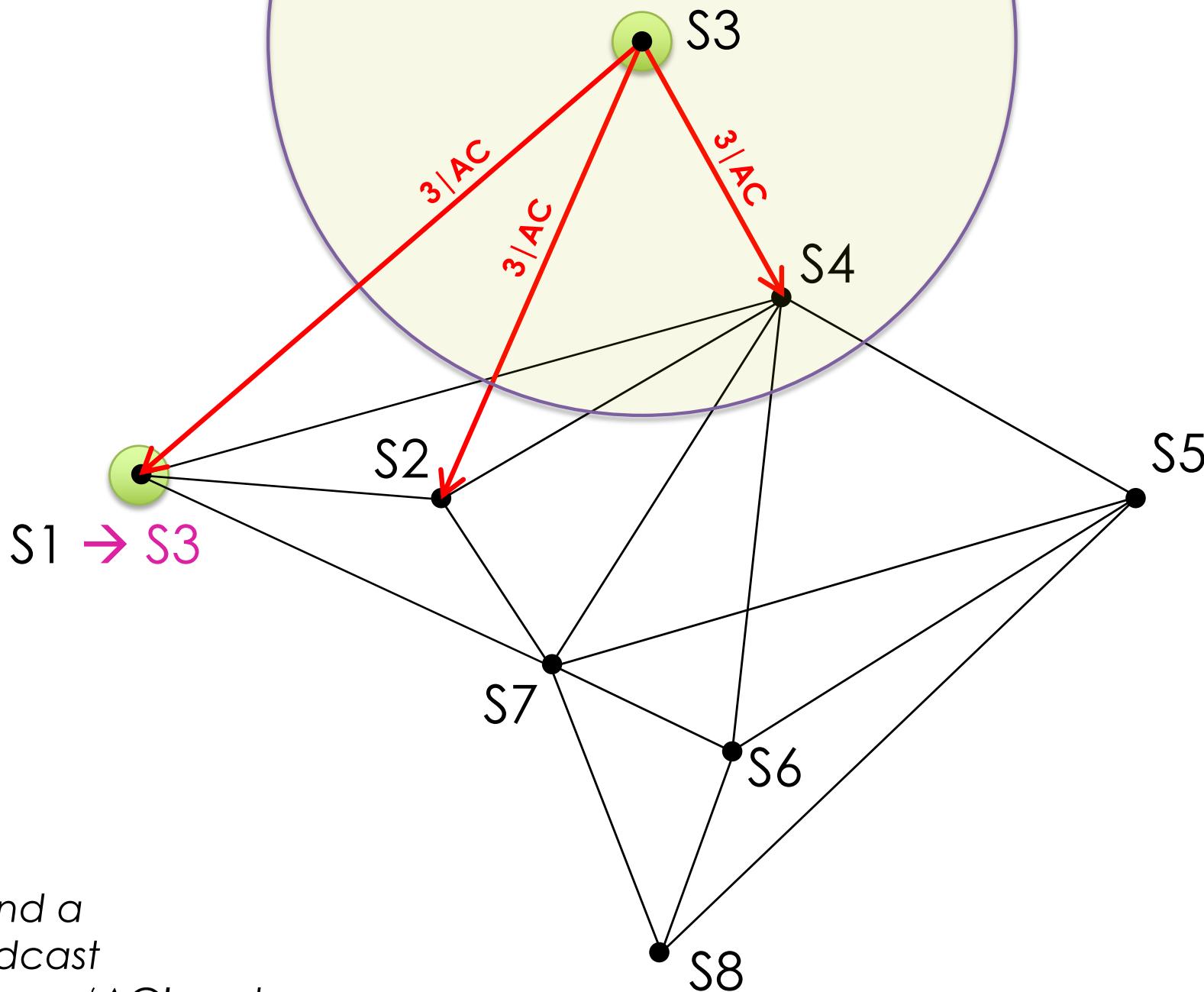
D-LPCN Algorithm



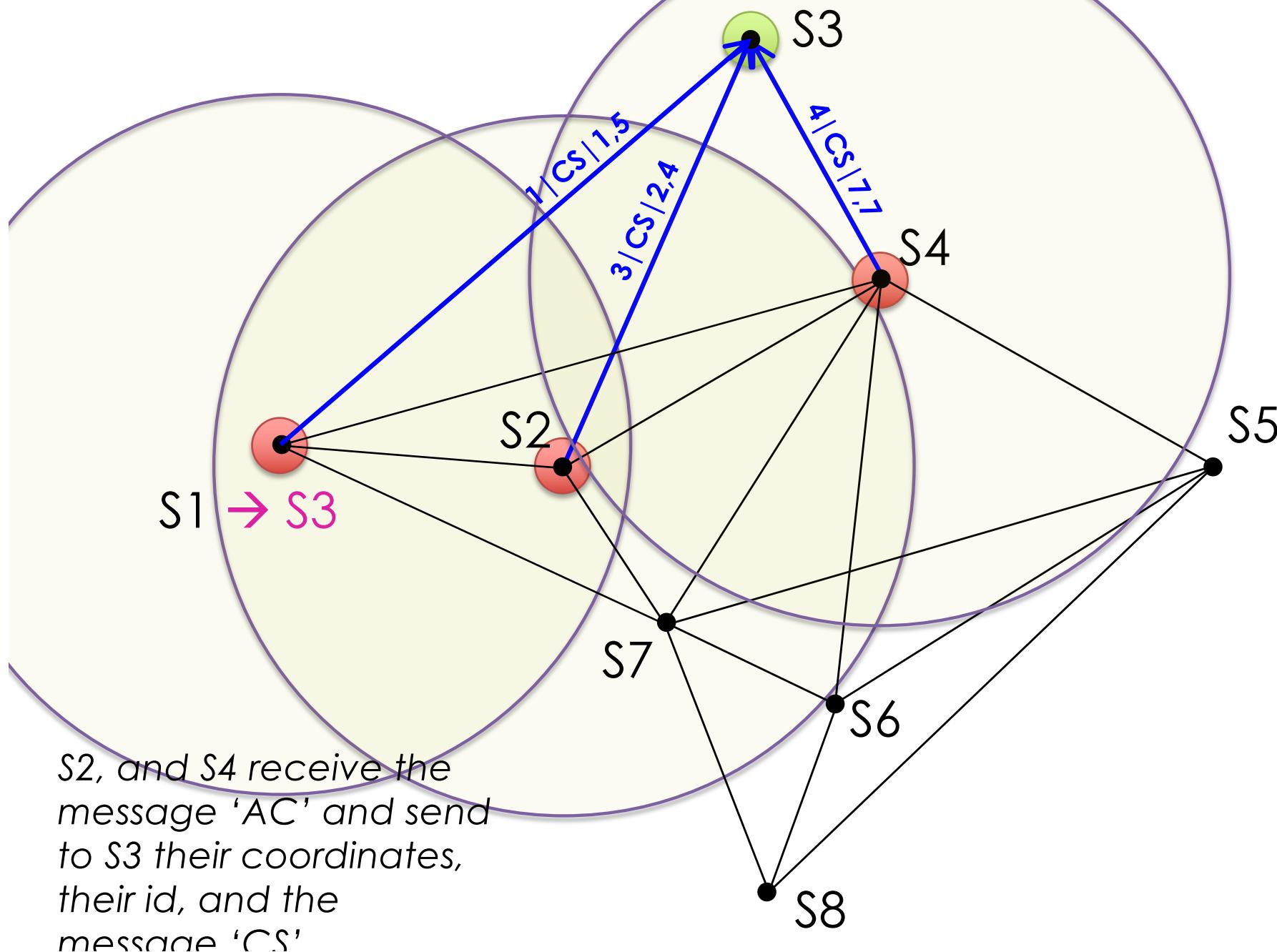
D-LPCN Algorithm



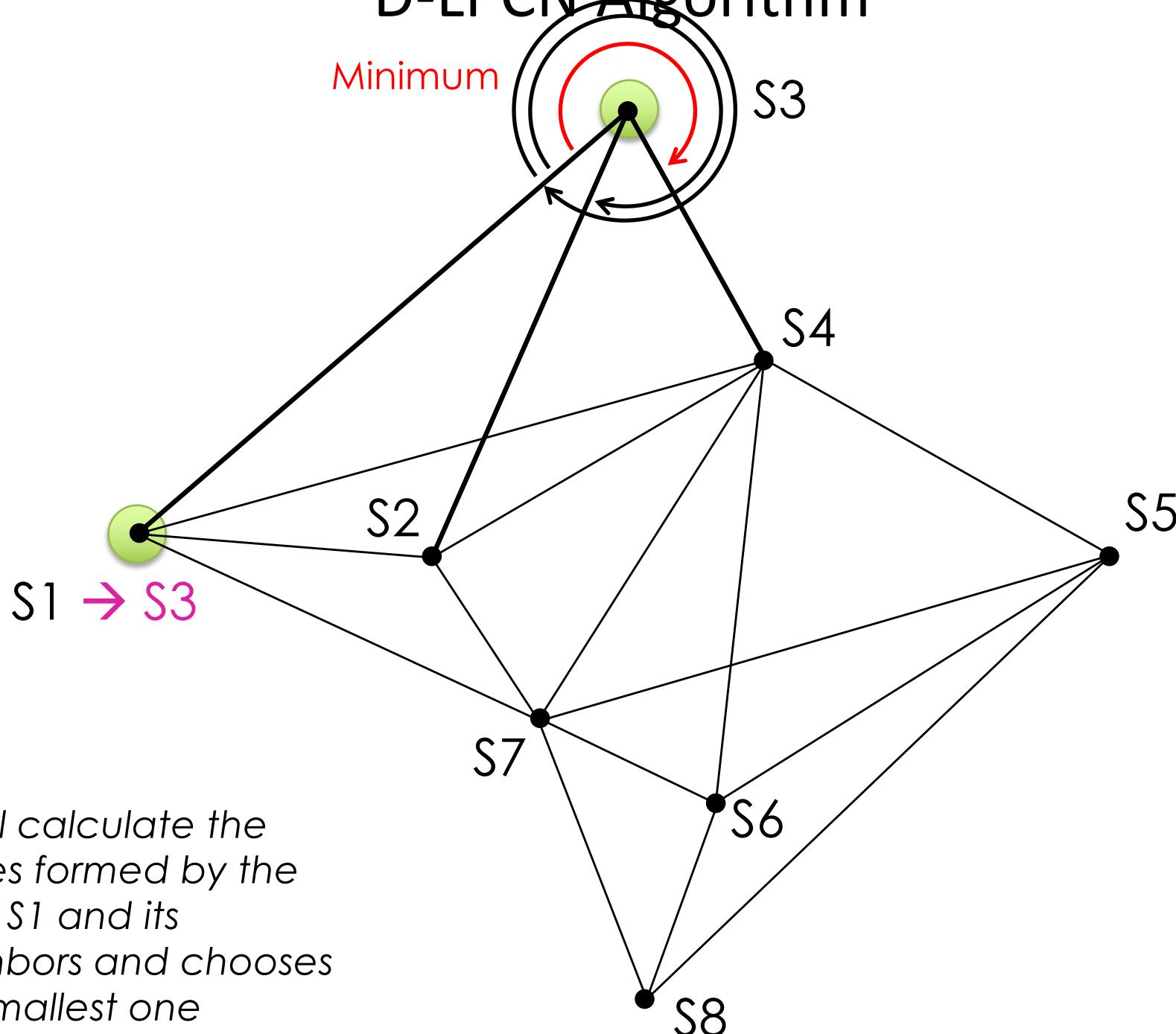
D-LPCN Algorithm



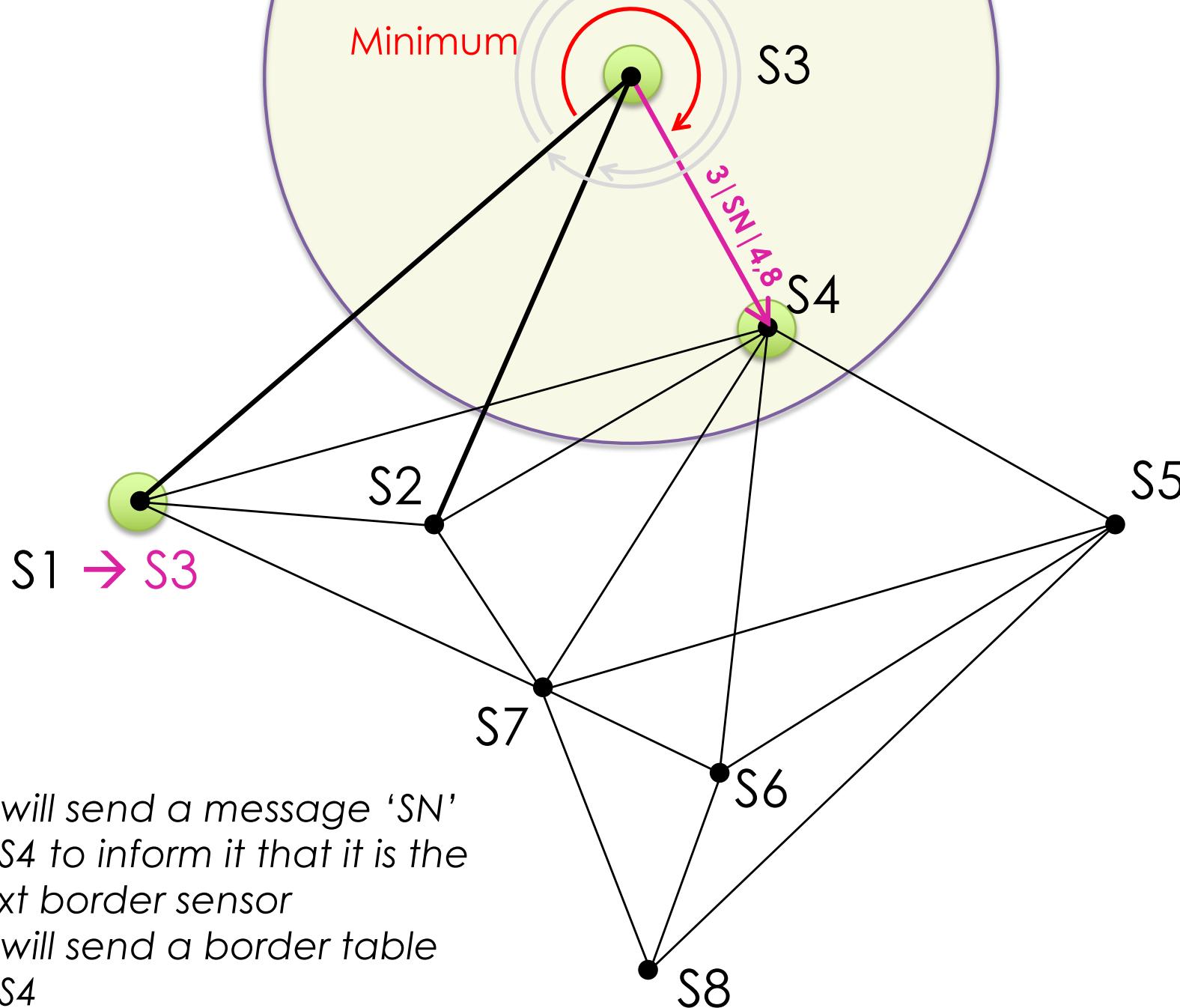
D-LPCN Algorithm



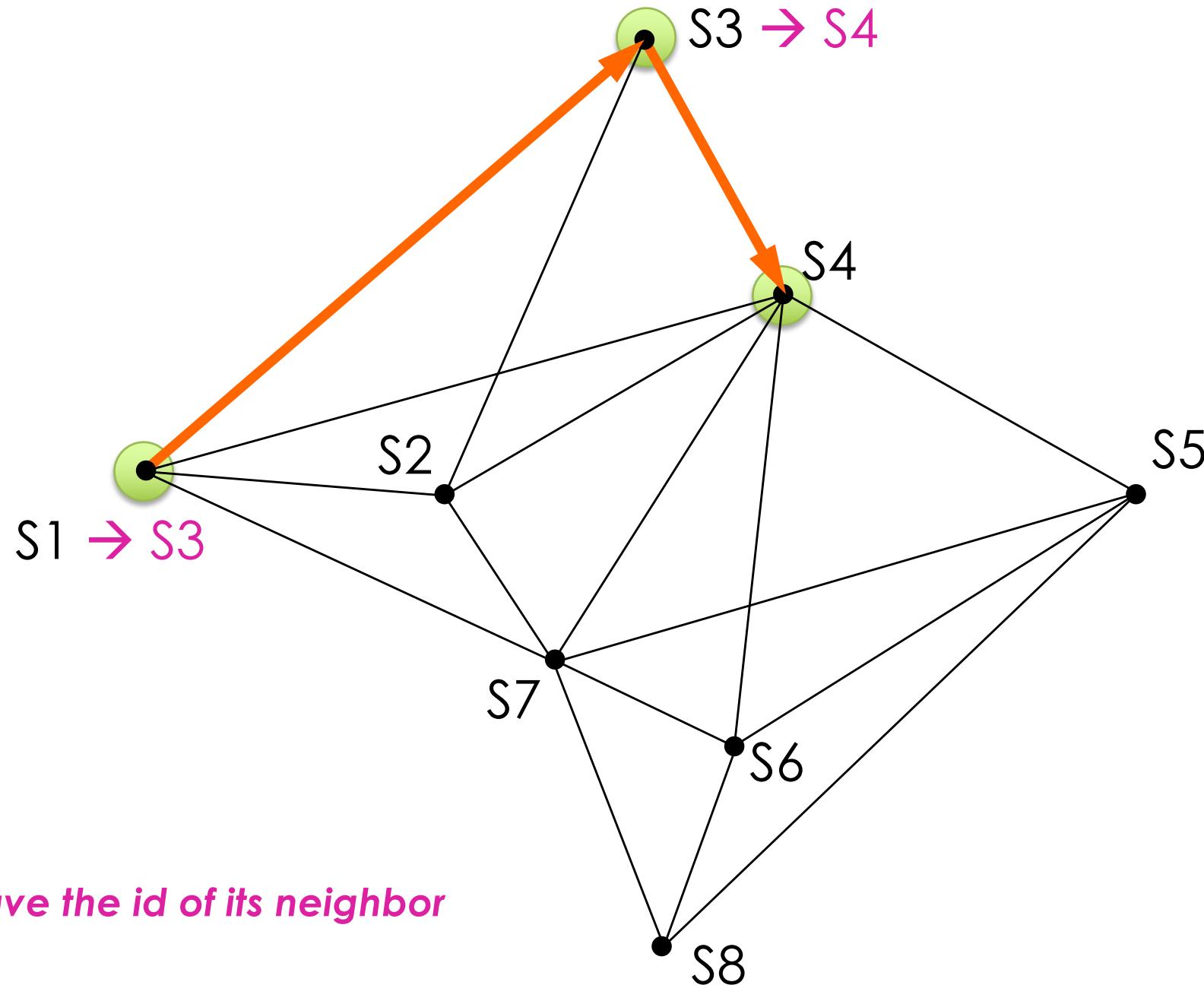
D-LPCN Algorithm



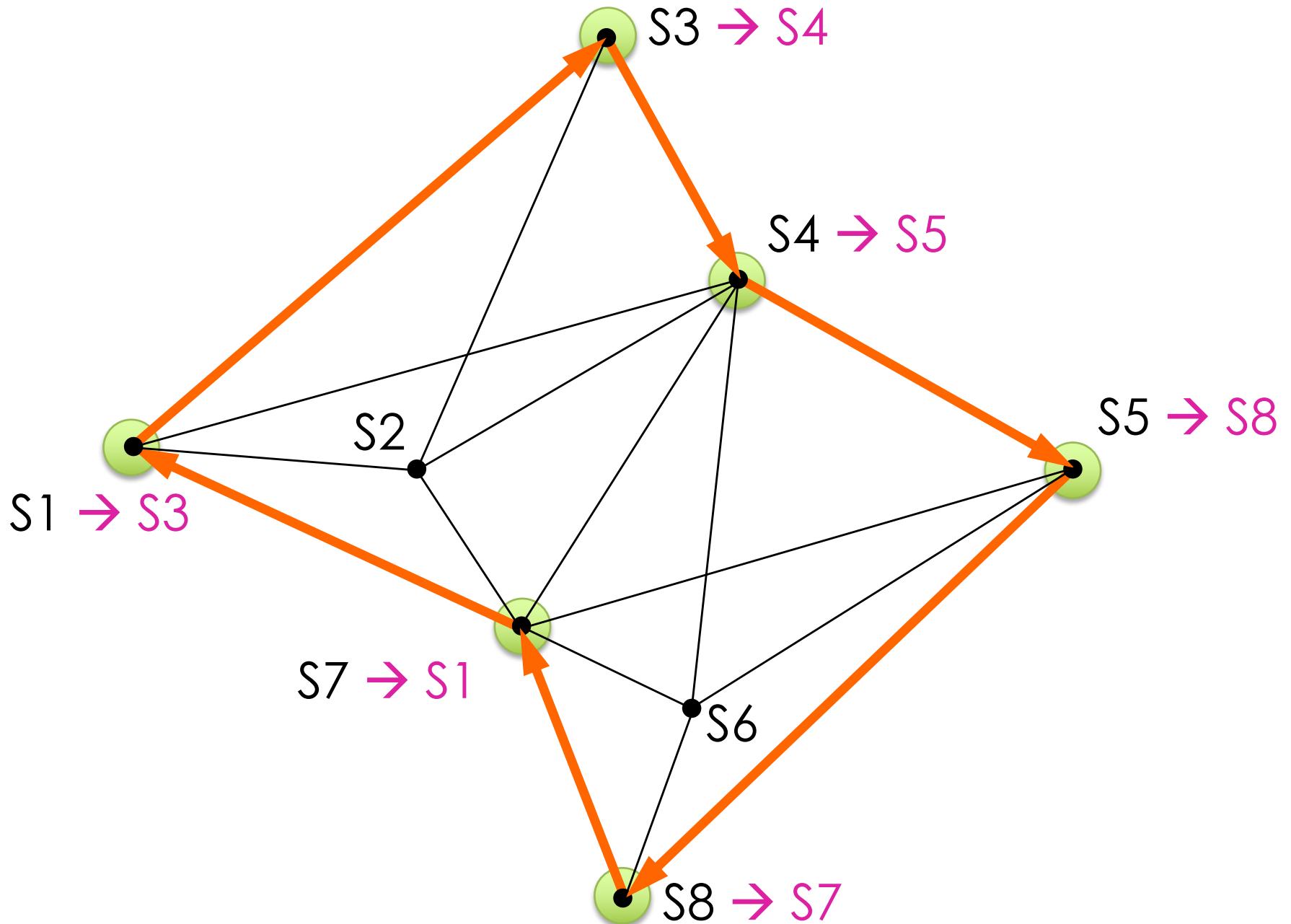
D-LPCN Algorithm



D-LPCN Algorithm



D-LPCN Algorithm



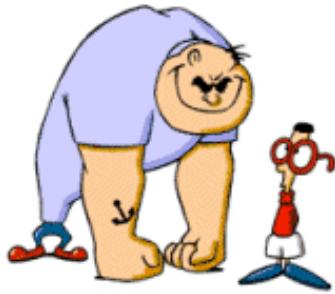
D-LPCN Algorithm: Demo (TelosB/Arduino-XBee)



3.3

How to determine the starting node?

LEADER ELECTION



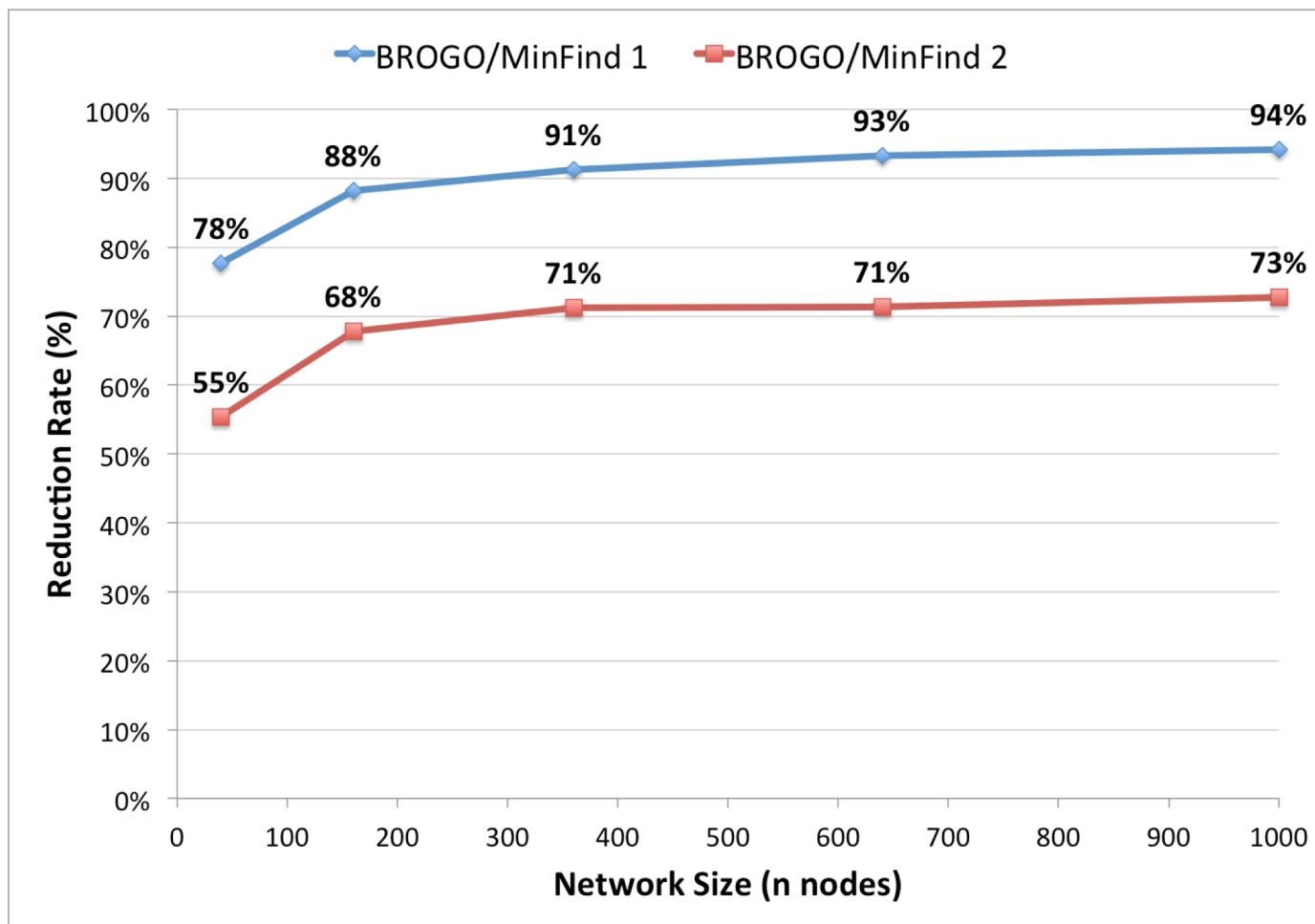
DoTRO Algorithm



Comparison with the MinFind algorithm

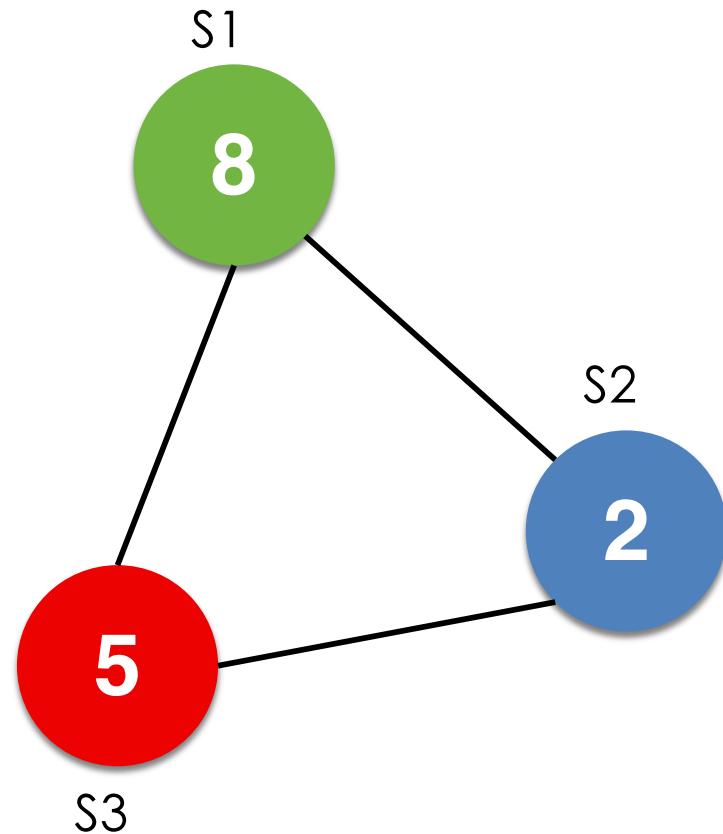
same density of 10 nodes/hm
n varies from 0 to 1000 sensor nodes

MinFind 1: xmin
MinFind 2: random



WBS: Wait-Before-Starting

Values



```
atget id id
mult v $id 1000
set once 0
loop
wait $v
read x
if(($x==\) && ($once==0))
    set once 1
    mark 1
    send A
    stop
end
if(($x==A) && ($once==0))
    set once 1
    send A
    stop
end
```

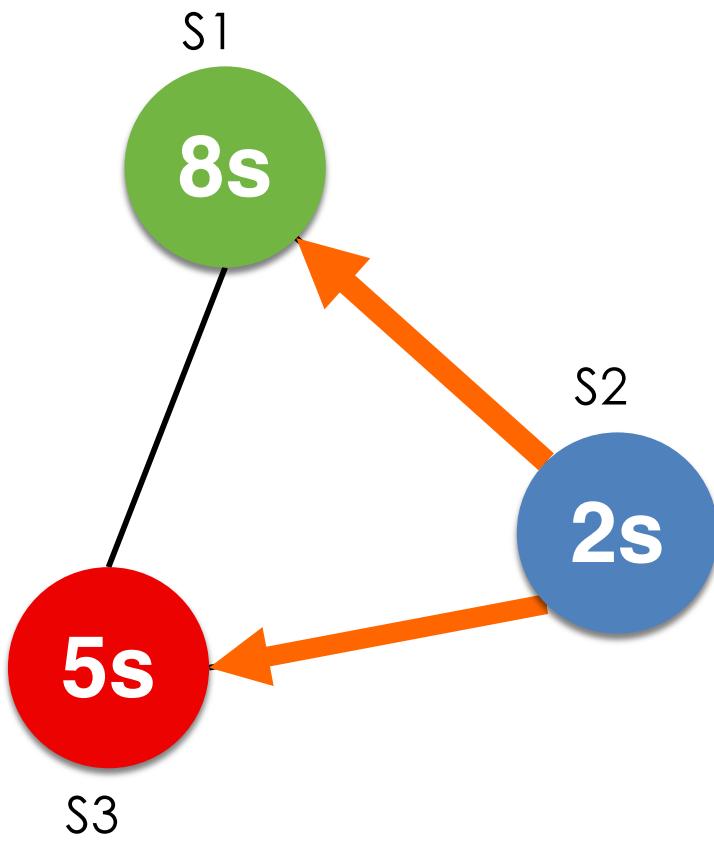
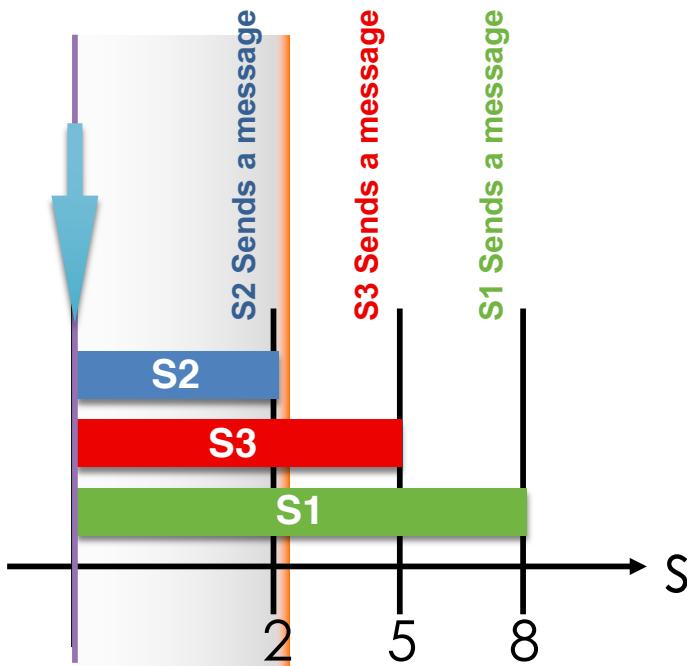
WBS: Wait-Before-Starting

Values

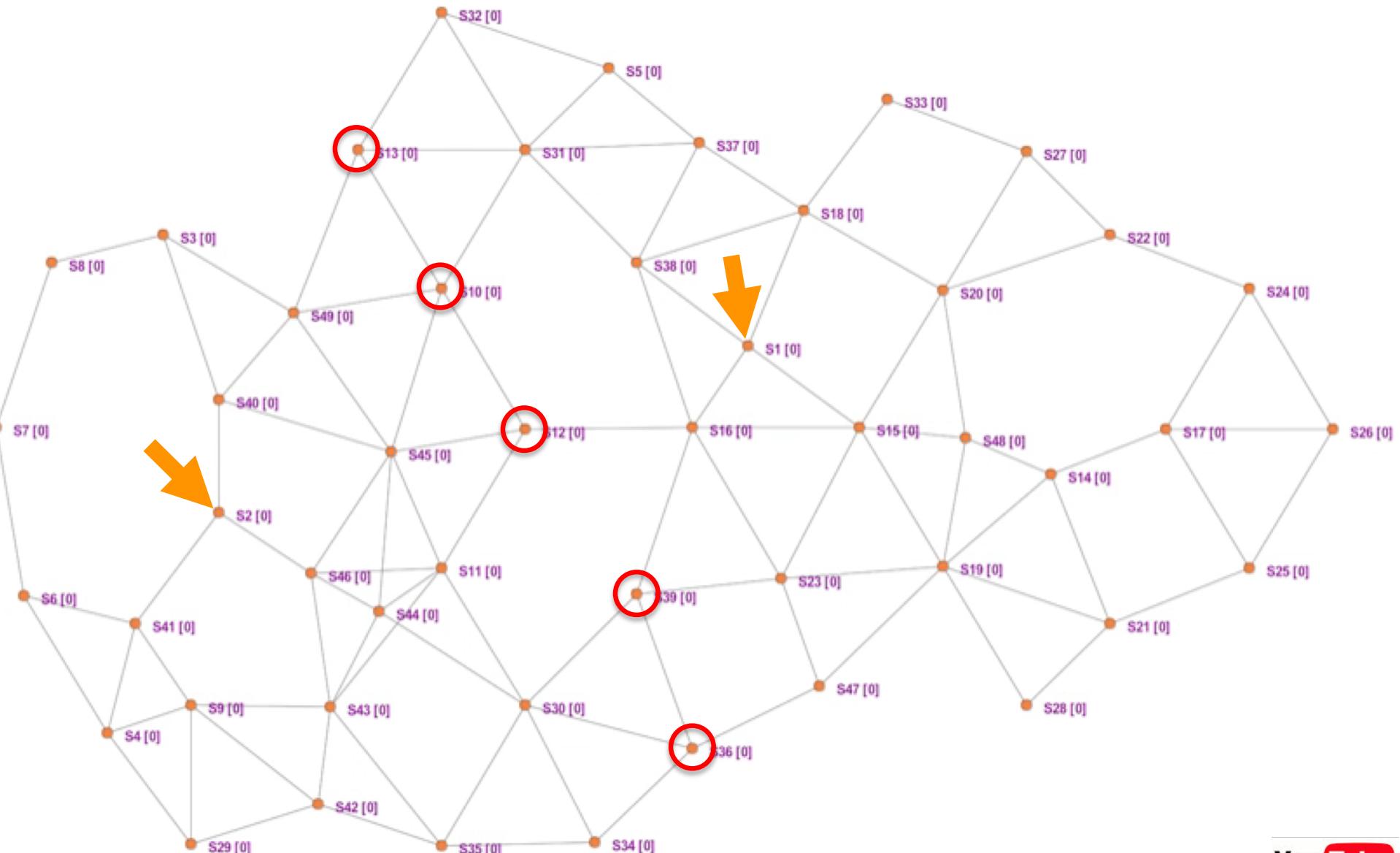
Waiting time unit = 1s

Waiting times

After 2s?



RBROGO & D-LPCN: disconnected network



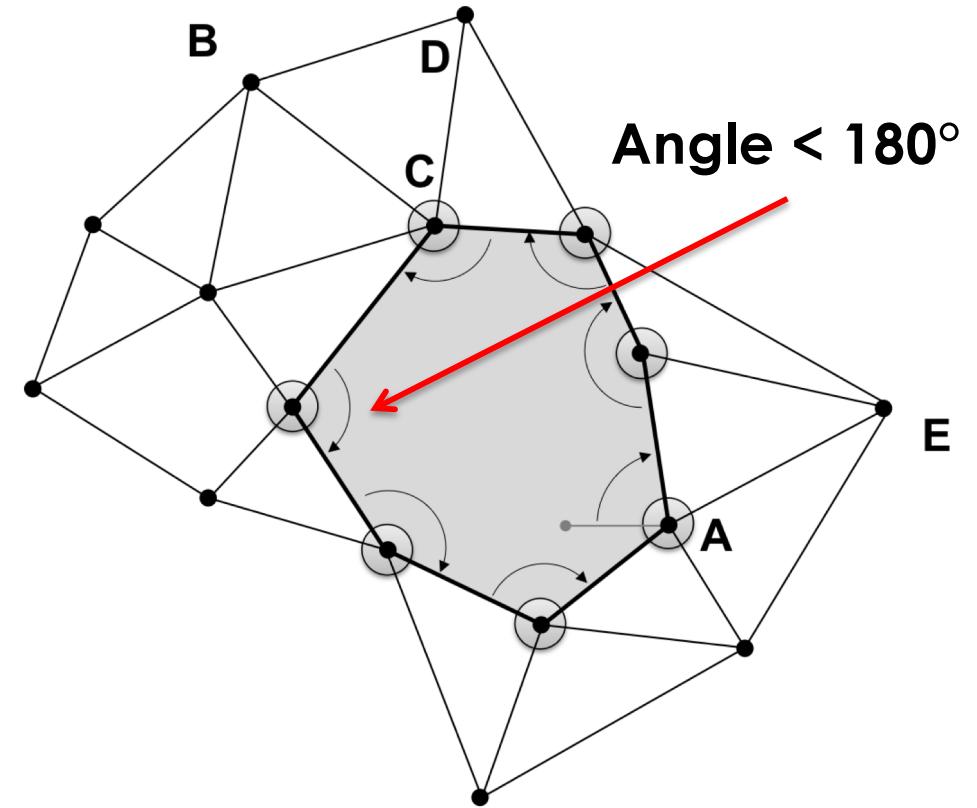
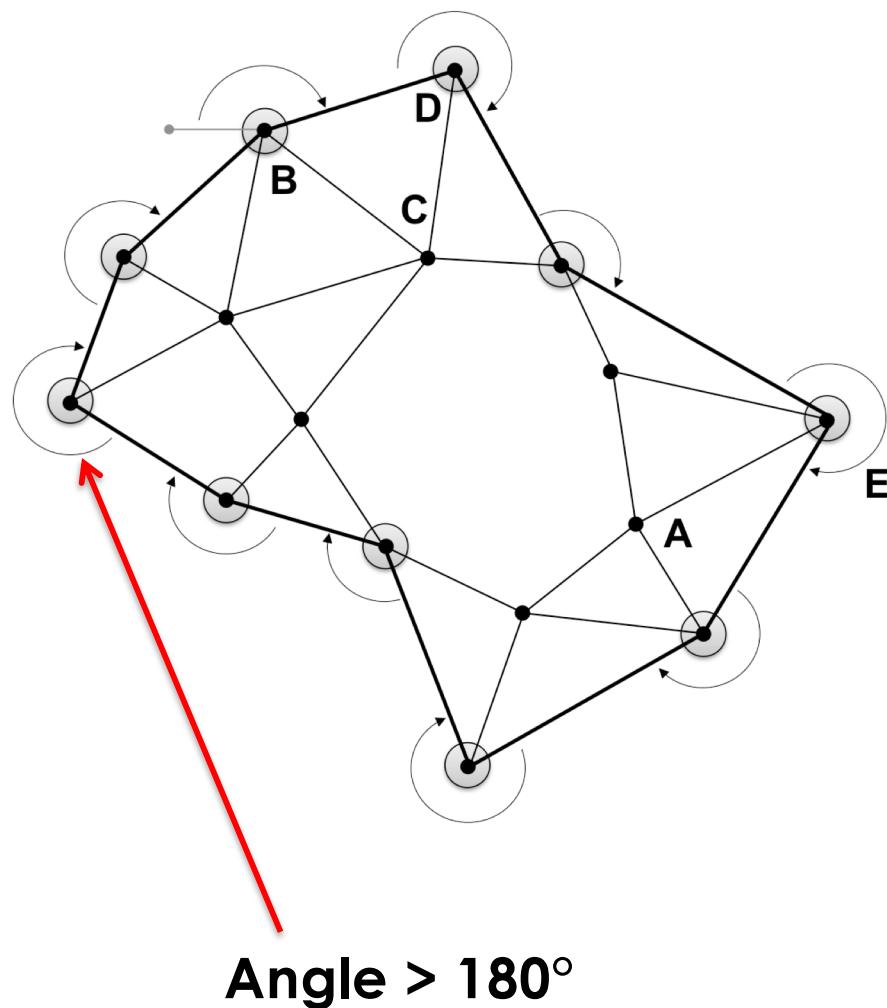
○ Faulty node

→ Root

<https://www.youtube.com/watch?v=iGStMtjeSHg>



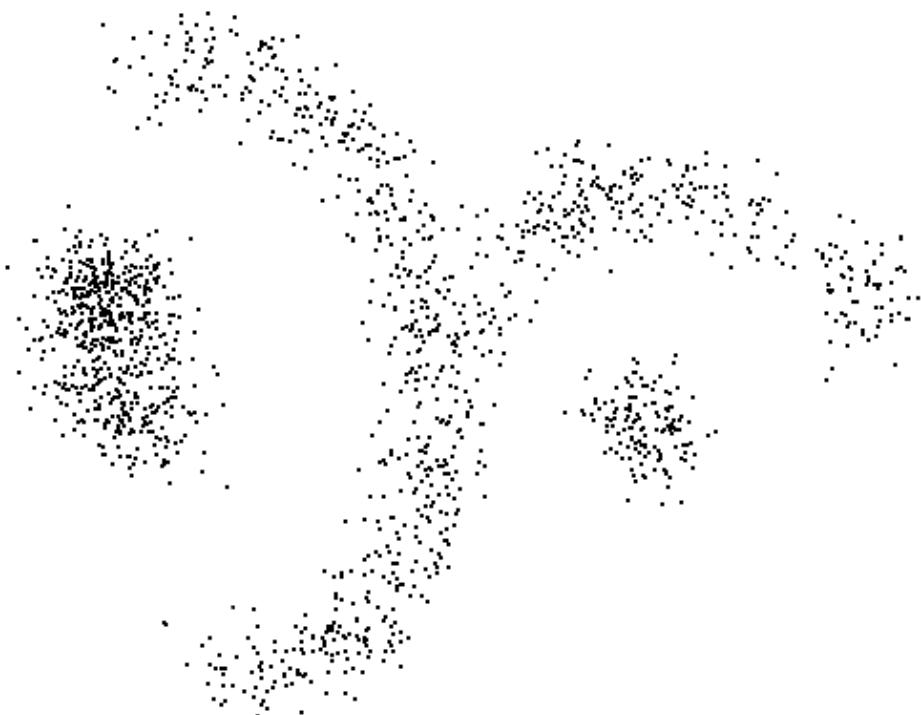
D-LPCN Algorithm: Gap detection



4

QUELQUES APPLICATIONS

Big Data: Finding Clusters



Big Data: Finding Clusters

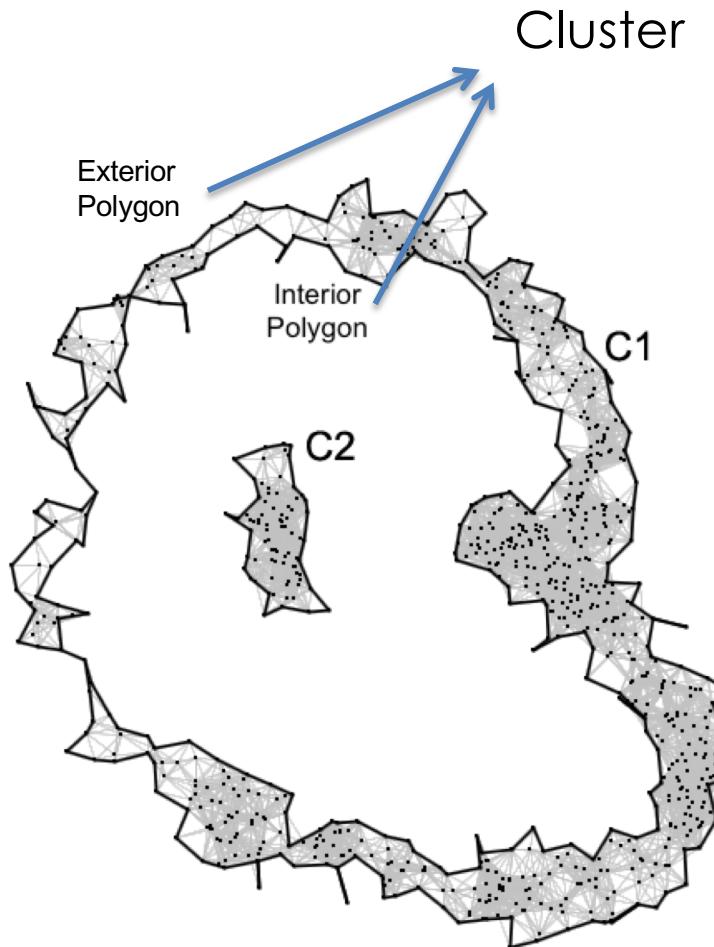


Image Processing: Contour drawing

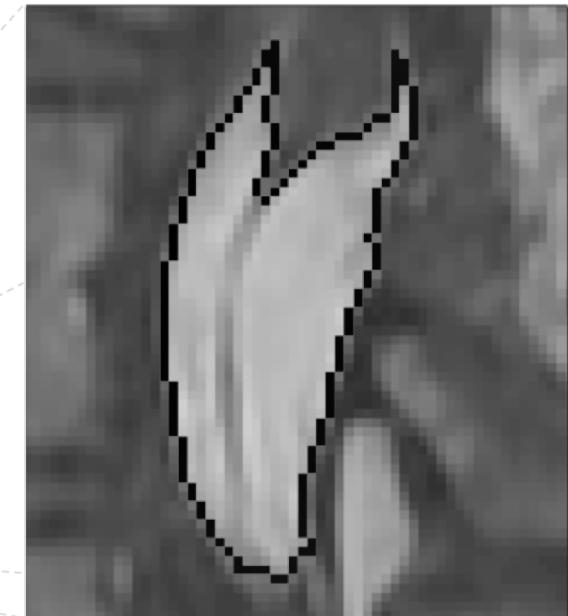
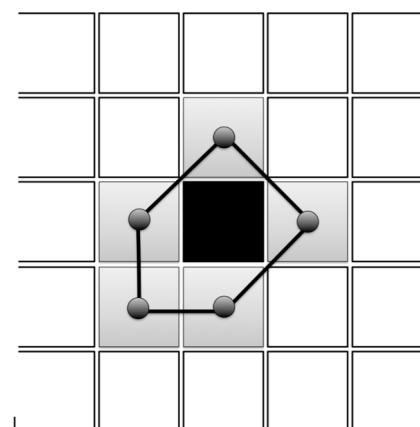
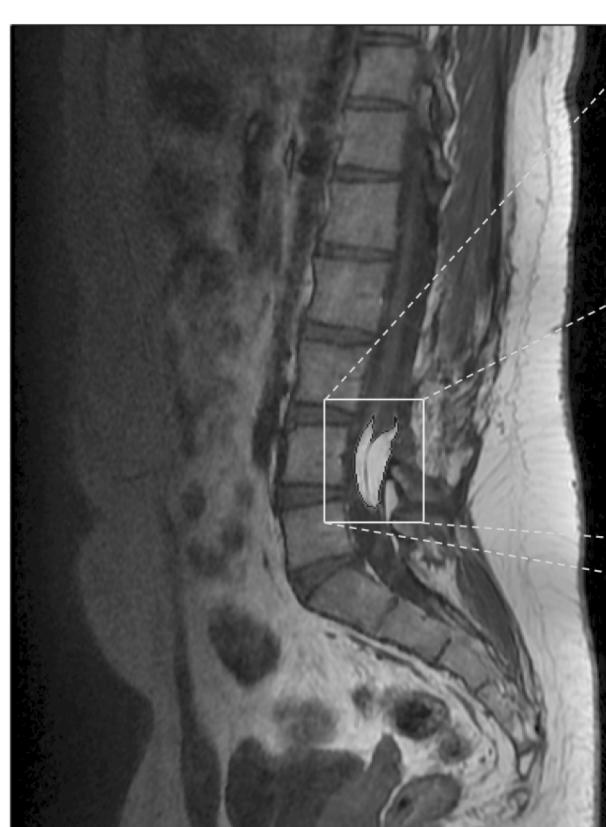
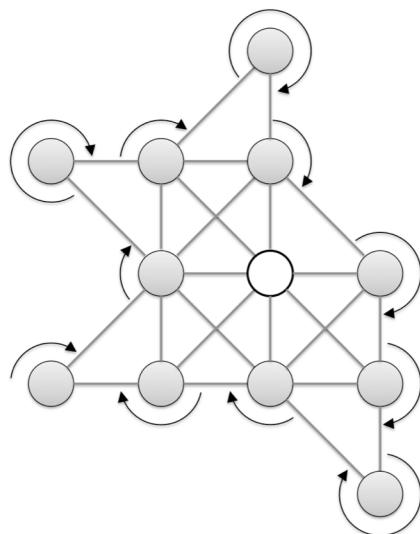
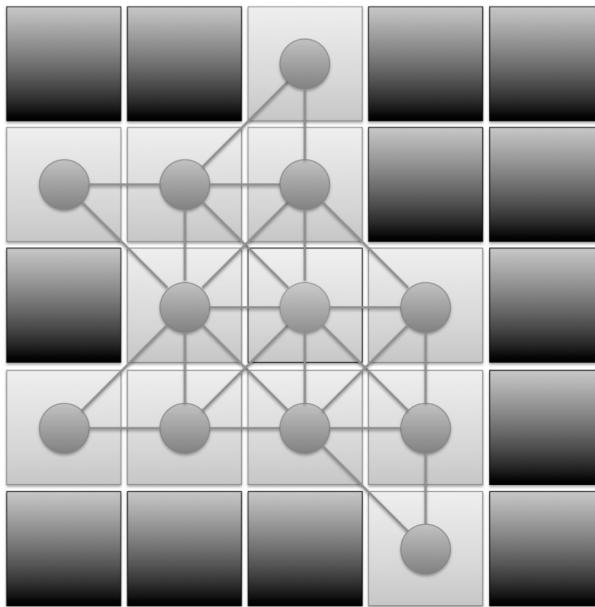


Image Processing: Countour drawing



5

CONCLUSION & FUTURE WORK

Conclusion

- We have presented:
 - A new concept of a Pseudo polygon
 - A new platform for the simulation of WSNs dedicated to Smart-cities and IoT
 - A new distributed algorithms allowing to determine the polygonal hull and the Gaps of Euclidean graphs
- We are developing other new algorithms:
 - What about you?
 - Just use CupCarbon ...

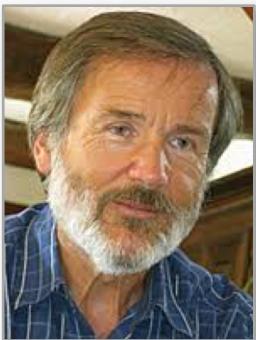
Future Work

- Questions:
 1. Can we consider that a graph is composed of a set of pseudo-polygons?
 2. Can we use pseudo-polygons to secure Data?
 3. Can we use pseudo-polygons to extract knowledge and in digital humanities?
 4. Can we ...

Boundaries and Hulls of Euclidean Graphs: From Theory to Practice



Madani Bezoui



Reinhardt Euler





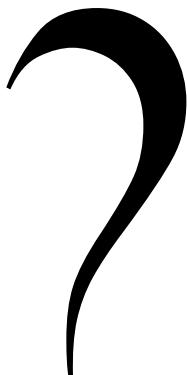
Thank You
for your attention



ISKO-Maghreb

Ahcène Bounceur

Ahcene.Bounceur@univ-brest.fr



www.cupcarbon.com

