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Extending Drawings of Graphs to Arrangements of Pseudolines

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1 Abstract

2 In the recent study of crossing numbers, drawings of graphs that can be extended to an
3 arrangement of pseudolines (pseudolinear drawings) have played an important role as they are a
4 natural combinatorial extension of rectilinear (or straight-line) drawings. A characterization of the
5 pseudolinear drawings of K_n was found recently. We extend this characterization to all graphs, by
6 describing the set of minimal forbidden subdrawings for pseudolinear drawings. Our characterization
7 also leads to a polynomial-time algorithm to recognize pseudolinear drawings and construct the
8 pseudolines when it is possible.

2012 ACM Subject Classification

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Lines 500

9 1 Introduction

10 Since 2004, geometric methods have been used to make impressive progress for determining
11 the crossing number of (certain classes of drawings of) the complete graph K_n . In particular,
12 drawings that extend to straight lines, or, more generally, arrangements of pseudolines, have
13 been central to this work, spurring interest in such drawings for arbitrary graphs, not just
14 complete graphs [2, 5, 6, 7, 12].

In particular, for pseudolinear drawings, it is now known that, for $n \geq 10$, a pseudolinear drawing of K_n has more than

$$H(n) := \frac{1}{4} \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor \left\lfloor \frac{n-2}{2} \right\rfloor \left\lfloor \frac{n-3}{2} \right\rfloor$$

15 crossings [1, 13]. The number $H(n)$ is conjectured by Harary and Hill to be the smallest
16 number of crossings over all topological drawings of K_n ; that is, the crossing number $\text{cr}(K_n)$
17 is conjectured to be $H(n)$.

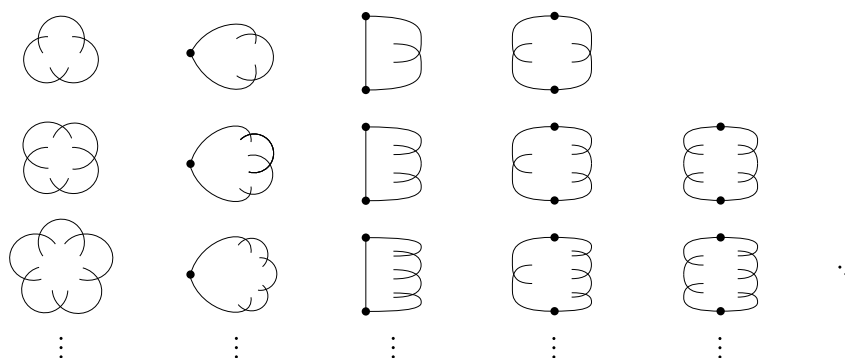
18 A *pseudoline* is the image ℓ of a continuous injection from the real numbers \mathbb{R} to the plane
19 \mathbb{R}^2 such that $\mathbb{R}^2 \setminus \ell$ is not connected. An *arrangement of pseudolines* is a set Σ of pseudolines



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48 ■ **Figure 1** Obstructions to pseudolinearity.

20 such that, if ℓ, ℓ' are distinct elements of Σ , then $|\ell \cap \ell'| = 1$ and the intersection is a crossing
 21 point. More on pseudolines and their importance for studying geometric drawings of graphs
 22 can be found in [10, 11].

23 A drawing D of a graph G is *pseudolinear* if there is an arrangement of pseudolines
 24 consisting of a different pseudoline ℓ_e for each edge e of G and such that $D[e] \subseteq \ell_e$.

25 In the study of crossing numbers, restricting the drawing to either straight lines or
 26 pseudolines yields the rectilinear crossing number $\overline{\text{cr}}(K_n)$ or the pseudolinear crossing number
 27 $\tilde{\text{cr}}(K_n)$, respectively. Clearly $\overline{\text{cr}}(K_n) \geq \tilde{\text{cr}}(K_n)$ and the geometric methods prove that
 28 $\tilde{\text{cr}}(K_n) > H(n)$, for $n \geq 10$.

29 A *good drawing* is one where no edge self-intersects and any two edges share at most
 30 one point—either a crossing or a common end point—and no three edges share a common
 31 crossing. One somewhat surprising result is from Aichholzer et al.: *a good drawing of K_n
 32 in the plane is homeomorphic to a pseudolinear drawing if and only if it does not contain
 33 a non-planar drawing of K_4 whose crossing is incident with the unbounded face of the K_4
 34 [2].* There are equivalent characterizations in [5, 6]. These conditions can be shown to be
 35 equivalent to not containing the *B-configuration* depicted as the third drawing of the first
 36 row of Figure 1.

37 Twenty-five years earlier, Thomassen proved a similar theorem for drawings in which
 38 each edge is crossed only once [16]. The *B-* and *W-*configurations are shown as the third
 39 and fourth drawings in the first row of Figure 1. Thomassen’s theorem is: *if D is a planar
 40 drawing of a graph G in which each edge is crossed at most once, then D is homeomorphic
 41 to a rectilinear drawing of G if and only if D contains no *B-* or *W-*configuration.*

42 Thomassen presented in [16] the *clouds* (first column in Figure 1) as an infinite family of
 43 drawings that are minimally non-pseudolinear.

44 Shortly after Thomassen’s paper, Bienstock and Dean proved that if $\text{cr}(G) \leq 3$, then
 45 $\overline{\text{cr}}(G) = \text{cr}(G)$ [8]. They also exhibited examples based on overlapping *W-*configurations to
 46 show the result fails for $\text{cr}(G) = 4$; such graphs can have arbitrarily large rectilinear crossing
 47 number.

48 Despite the existence of infinitely many obstructions to pseudolinearity, we characterize
 49 them all.

51 ► **Theorem 1.** A good drawing of a graph G is pseudolinear if and only if it does not contain
 52 one of the infinitely many obstructions shown in Figure 1.

53 The drawings in Figure 1 are obtained from the *clouds* (first column) by replacing at most
 54 two crossings by vertices. The formal statement of Theorem 1 is Theorem 15 in Section 6;

55 also a more general version of this statement, Theorem 2, is discussed below. That there is a
 56 result such as ours is somewhat surprising, because stretching an arrangement of pseudolines
 57 to a rectilinear drawing has been shown by Mněv [14, 15] to be $\exists\mathbb{R}$ -hard. In particular,
 58 recognizing a drawing as being homeomorphic to a rectilinear drawing is NP-hard.

59 The natural setting for our characterization is strings embedded in the plane. An *arc* σ
 60 is the image $f([0, 1])$ of the compact interval $[0, 1]$ under a continuous map $f : [0, 1] \rightarrow \mathbb{R}^2$.
 61 Let $S(\sigma) = \{p \in \sigma : |f^{-1}(p)| \geq 2\}$ be the set of self-intersections of σ . A *string* is an arc σ
 62 for which $S(\sigma)$ is finite. If $S(\sigma) = \emptyset$, then σ is *simple*.

63 An intersection point between of two strings σ and σ' is *ordinary* if it is either an endpoint
 64 of σ or σ' , or is a *crossing* (a crossing is a non-tangential intersection point in $\sigma \cap \sigma'$ that
 65 is not an end of σ or σ'). A set Σ of strings is *ordinary* if Σ is finite and any two strings
 66 in Σ have only finitely many intersections, all of which are ordinary. All the sets of strings
 67 considered in this paper are ordinary.

68 If Σ is an ordinary set of strings, then its *planarization* $G(\Sigma)$ is the plane graph obtained
 69 from Σ by inserting vertices at each crossing between strings and also at the endpoints of
 70 every string in Σ . To keep track of the information given by the strings, we will always
 71 assume that each string Σ has a different color and that each edge in $G(\Sigma)$ inherits the color
 72 of the string including it.

73 If Σ is an ordinary set of strings, then, for a cycle C in $G(\Sigma)$ (which is a simple closed
 74 curve in \mathbb{R}^2) and a vertex $v \in V(C)$, v is a *rainbow* for C if all the edges incident with v and
 75 drawn in the closed disk bounded by C (including the two edges of C at v) have different
 76 colours. The reader can verify that, for each drawing in Figure 1, if we let Σ be the edges
 77 of the drawing, then the unique cycle in $G(\Sigma)$ has at most two rainbows. Our main result
 78 characterizes these cycles as the only possible obstructions:

79 ► **Theorem 2.** *An ordinary set of strings Σ can be extended to an arrangement of pseudolines
 80 if and only if every cycle C of $G(\Sigma)$ has at least three rainbows.*

81 Henceforth, we define any cycle C in $G(\Sigma)$ with at most two rainbows as an *obstruction*.
 82 A set of strings is *pseudolinear* if it has an extension to an arrangement of pseudolines.

83 Theorem 2 is our main contribution. In the next section, we show that the presence
 84 of an obstruction implies the set of ordinary strings is not pseudolinear. The converse is
 85 proved in Section 4 by extending, one small step at a time, the strings in Σ to get closer
 86 to an arrangement of pseudolines. After each extension, we must show that no obstruction
 87 has been introduced. This involves dealing with cycles in $G(\Sigma)$ that have precisely three
 88 rainbows (that we refer as *near-obstructions*). In Section 3 we show the key lemma that if G
 89 has two such near-obstructions that intersect nicely at a vertex v , then G has an obstruction.
 90 In Section 5 we present a polynomial-time algorithm for detecting obstructions and we argue
 91 why the proof of Theorem 2 implies a polynomial-time algorithm for extending a pseudolinear
 92 set of strings. Finally, in Section 6, we show how Theorem 1 follows from Theorem 2 and we
 93 present some concluding remarks.

94 **2 A set of strings with an obstruction is not extendible**

95 Let us start by showing the easy direction of Theorem 2:

96 ► **Lemma 3.** *If the underlying graph $G(\Sigma)$ of a set Σ of strings has an obstruction, then Σ
 97 is not pseudolinear.*

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98 Suppose that C is a cycle of $G(\Sigma)$ for some set of strings Σ . We define $\delta(C)$ as the set of
 99 vertices of C for which their two incident edges in C have different colours. In a set Σ of
 100 simple strings where no two intersect twice, $|\delta(C)| \geq 3$ for every cycle C of $G(\Sigma)$.

101 ► **Lemma 4.** *Let Σ be a set of simple strings where every pair intersect at most once. Suppose
 102 that C is an obstruction with $|\delta(C)|$ as small as possible. Let $S = x_0, x_1, \dots, x_\ell$ be a path
 103 of $G(\Sigma)$ representing a subsegment of some string $\sigma \in \Sigma$ such that $x_0x_1 \in E(C)$, $x_1 \in \delta(C)$
 104 and x_1 is not a rainbow of C . Then $V(C) \cap V(S) = \{x_0, x_1\}$.*

105 **Proof.** By way of contradiction, suppose that there is a vertex $x_r \in V(C) \cap V(S)$ with $r \geq 3$.
 106 Assume that $r \geq 3$ is as small as possible. Let P be the subpath of S connecting x_1 to x_r .
 107 Since $x_0x_1 \in E(C)$ and $x_1 \in \delta(C)$ and $P \subseteq \sigma$, $x_1x_2 \notin E(C)$. Because x_1 is not a rainbow
 108 for C and no two strings tangentially intersect at x_1 , the edge x_1x_2 is drawn in the closed
 109 disk bounded by C . By choice of r , P is an arc connecting x_1 to x_r in the interior of C .

110 Let C_1 and C_2 be the cycles obtained from the union of P and one of the two xy -subpaths
 111 in C . We may assume that $x_0x_1 \in E(C_1)$. Let $\rho(C)$ be either $\delta(C)$ or the set of rainbows
 112 in C . For $i = 1, 2$, let $Q_i = V(C_i) \setminus V(P)$. Then $\rho(C) \cap Q_i = \rho(C_i) \cap Q_i$. We see that
 113 $\rho(C_1) \setminus Q_1 \subseteq \{x_r\}$ and $\rho(C_2) \setminus Q_2 \subseteq \{x_1, x_r\}$.

114 For $\rho = \delta$, $|\delta(C_2)| \geq 3$, so $|\delta(C) \cap Q_2| \geq 1$. Since $x_1 \notin \delta(C_1)$, $|\delta(C_1)| \leq |\delta(C_1) \cap Q_2| +$
 115 $|\{x_r\}| \leq |\delta(C)| - 2 + |\{x_r\}| < |\delta(C)|$. Likewise, $|\delta(C) \cap Q_1| \geq 2$ and $x_1 \in \delta(C) \cap \delta(C_2)$.
 116 Therefore, $|\delta(C_2)| \leq |\delta(C)| - 2 + |\{x_r\}| < |\delta(C)|$. Thus, neither C_1 nor C_2 is an obstruction.

117 Now taking ρ to be the set of rainbows, the preceding paragraph shows $|\rho(C_1)| \geq 3$ and
 118 $|\rho(C_2)| \geq 3$. Therefore, $|\rho(C) \cap Q_1| = |\rho(C_1) \cap Q_1| \geq 2$ and $|\rho(C) \cap Q_2| = |\rho(C_2) \cap Q_2| \geq 1$.
 119 Thus, $|\rho(C)| \geq 3$, a contradiction. ◀

120 **Proof of Lemma 3.** By way of contradiction, suppose that Σ is pseudolinear and that $G(\Sigma)$
 121 has an obstruction C .

122 Consider an extension of Σ to an arrangement of pseudolines, and then cut off the two
 123 infinite ends of each pseudoline to obtain a set of strings Σ' extending Σ , and in which every
 124 pair of strings in Σ' cross once. In $G(\Sigma')$, there is a cycle C' that represents the same simple
 125 closed curve as C . Because C' is obtained from subdividing some edges of C and the colours
 126 of a subdivided edge are the same, C' has fewer than three rainbows. Therefore, we may
 127 assume that $\Sigma = \Sigma'$ and $C = C'$. Now, the ends of every string in Σ are degree-1 vertices in
 128 the outer face of $G(\Sigma)$.

129 As every string in Σ is simple and no two strings intersect more than once, $|\delta(C)| \geq 3$.
 130 We will assume that C is chosen to minimize $|\delta(C)|$.

131 Since C is an obstruction, there exists $x_1 \in \delta(C)$ such that x_1 is not a rainbow in
 132 C . Consider a neighbour x_0 of x_1 in C . Let $S = x_0, x_1, \dots, x_\ell$ be the path obtained by
 133 traversing the string σ extending x_0x_1 , such that x_ℓ is an end of σ . By Observation 4,
 134 $V(S) \cap V(C) = \{x_0, x_1\}$, and because x_ℓ is in the outer face of C , the segment of σ from x_1
 135 to x_ℓ has its relative interior in the outer face of C .

136 However, since x_1 is not a rainbow, there exists a string $\sigma' \in \Sigma$ including two edges
 137 at x_1 drawn in the disk bounded by C . Thus, σ and σ' tangentially intersect at x_1 , a
 138 contradiction. ◀

139 3 The key lemma

140 In this section we present the key lemma used in the proof of Theorem 2.

141 A plane graph G is *path-partitioned* if for $m \geq 1$, there exists a colouring $\chi : E(G) \rightarrow$
 142 $\{1, \dots, m\}$ such that for each $i \in \{1, \dots, m\}$, the edges in $\chi^{-1}(i)$ induce a path $P_i \subseteq G$ where

143 any two distinct paths P_i and P_j do not tangentially intersect. Indeed, every underlying
 144 planar graph $G(\Sigma)$ of a set of simple strings Σ is path-partitioned. Moreover, every path-
 145 partitioned plane graph can be obtained by subdividing a planarization of an ordinary set of
 146 simple strings. To extend the previously introduced notation we refer to each P_i as a string.
 147 The concepts of rainbow and obstruction naturally extend to the context of path-partitioned
 148 plane graphs.

149 Suppose that G is a path-partitioned plane graph. Given $v \in V(G)$, a *near-obstruction at*
 150 v is a cycle C with at most three rainbows and such that v is a rainbow of C . Understanding
 151 how near-obstructions behave is the key ingredient needed in the proof of Theorem 2:

152 ► **Lemma 5.** *Let G be a path-partitioned plane graph and let $v \in V(G)$. Suppose that C_1
 153 and C_2 are two near-obstructions at v such that the union of the closed disks bounded by C_1
 154 and C_2 contains a small open ball centered at v . Suppose that one of the following two holds:*

- 155 1. *no obstruction of G contains v ; or*
- 156 2. *the two edges of C_1 incident with v are the same as the two edges of C_2 incident with v .*

157 *Then G has an obstruction not including v .*

158 Given a plane graph G , a cycle $C \subseteq G$ and a vertex $v \in V(C)$, *the edges at v inside C* are
 159 *the edges of G incident with v drawn in the disk bounded by C .*

160 ► **Useful Fact.** *Let G be planar path-partitioned graph. Suppose that for two cycles C and
 161 C' , $v \in V(C) \cap V(C')$ is a vertex such that the edges at v inside C' are also edges at v inside
 162 C . If v is a rainbow for C , then v is a rainbow for C' .*

163 **Proof of Lemma 5.** By way of contradiction, suppose that G has no obstruction not includ-
 164 ing v . The “small ball” hypothesis implies that v is not in the outer face of the subgraph
 165 $C_1 \cup C_2$.

166 We claim that $|V(C_1) \cap V(C_2)| \geq 2$. Suppose not. Then C_1 and C_2 are edge-disjoint
 167 and $V(C_1) \cap V(C_2) = \{v\}$. For $i = 1, 2$, let e_i and f_i be the edges of C_i at v and let Δ_i
 168 be the closed disk bounded by C_i . From the “small ball” hypothesis it follows that (i) Δ_1
 169 contains the edges e_2 and f_2 ; and (ii) the points near v in the exterior of Δ_2 are contained
 170 in Δ_1 . These two properties imply that the path $C_2 - \{e_2, f_2\}$ intersects C_1 at least twice,
 171 and hence, $|V(C_1) \cap V(C_2)| \geq 2$.

172 From the last paragraph we know that $C_1 \cup C_2$ is 2-connected, and hence the outer face
 173 of $C_1 \cup C_2$ is bounded by a cycle C_{out} . We will assume that

- 174 (*) the cycles C_1 and C_2 satisfying the hypothesis of Lemma 5 are chosen so that the number
 175 of vertices of G in the disk bounded by C_{out} is minimal.

176 The Useful Fact applied to $C = C_{out}$ and to each $C' \in \{C_1, C_2\}$, shows that every vertex
 177 that is a rainbow in C_{out} is also a rainbow in each of the cycles in $\{C_1, C_2\}$ containing it.
 178 We can assume that C_{out} is not an obstruction or else we are done. We may relabel C_1 and
 179 C_2 so that two of the rainbows of C_{out} , say p and q , are also rainbows in C_1 . Neither p nor q
 180 is v because $v \notin V(C_{out})$. Because C_1 is a near-obstruction, p , q and v are the only rainbows
 181 of C_1 .

182 Since $v \notin V(C_{out})$, by following C_1 in the two directions starting at v , we find a path
 183 $P_v \subseteq C_1$ containing v in which only the ends u and w of P_v are in C_{out} (note that $u \neq v$
 184 because $\{p, q\} \subseteq V(C_1) \cap V(C_{out})$). As v is in the interior face of C_{out} , P_v is also in the
 185 interior of C_{out} . Let Q_{out}^1, Q_{out}^2 be the uw -paths of C_{out} . One of the two closed disks bounded

186 by $P_v \cup Q_{out}^1$ and $P_v \cup Q_{out}^2$ contains C_1 . By symmetry, we may assume that C_1 is contained
 187 in the first disk. Since $C_{out} \subseteq C_1 \cup C_2$, this implies that Q_{out}^2 is a subpath of C_2 .

188 Our desired contradiction will be to find three rainbows in C_2 distinct from v . We
 189 find the first: let $C_1 - (P_v)$ be the uw -path in C_1 distinct from P_v . The disk bounded
 190 by $(C_1 - (P_v)) \cup Q_{out}^2$ contains the one bounded by C_1 . The Useful Fact applied to $C =$
 191 $(C_1 - (P_v)) \cup Q_{out}^2$ and $C' = C_1$ implies that each vertex in $C_1 - (P_v)$ that is rainbow in
 192 $(C_1 - (P_v)) \cup Q_{out}^2$ is also rainbow in C_1 . Since C_1 has at most two rainbows in $C_1 - (P_v)$,
 193 namely p and q , $(C_1 - (P_v)) \cup Q_{out}^2$ has a third rainbow r_1 in the interior of Q_{out}^2 (else
 194 $(C_1 - (P_v)) \cup Q_{out}^2$ is an obstruction and we are done). Note that r_1 is also a rainbow for C_2 .

195 To find another rainbow in C_2 , consider the edge e_u of C_2 incident to u and not in Q_{out}^2 .
 196 We claim that either u is a rainbow in C_2 or that e_u is not included in the closed disk
 197 bounded by $P_v \cup Q_{out}^2$. Seeking a contradiction, suppose that u is not a rainbow of C_2 and
 198 that e_u is included in the disk. Then we can find two edges in the rotation at u , included in
 199 the disk bounded by $P_v \cup Q_{out}^2$, that belong to the same string σ . The vertex u is a rainbow
 200 in C_1 , as else, we would find a string σ' with two edges inside $Q_{out}^1 \cup P_v$, showing that σ and
 201 σ' tangentially intersect at u . As p and q are the only rainbows of C_1 in C_{out} , u is one of p
 202 and q . Therefore u is a rainbow in C_{out} , and hence, a rainbow in C_2 , a contradiction.

203 If u is a rainbow in C_2 , then this is the desired second one. Otherwise, e_u is not in the
 204 closed disk bounded by $P_v \cup Q_{out}^2$. Let $P_u \subseteq C_2$ be the path starting at u , continuing on e_u
 205 and ending on the first vertex u' in P_v that we encounter. Let C_u be the cycle consisting of
 206 P_u and the uu' -subpath uP_vu' of P_v .

207 \triangleright Claim 6. If P_u does not have a rainbow of C_u in its interior, then either C_u is an
 208 obstruction not containing v or:

- 209 (a) C_u and C_2 are near-obstructions at v satisfying the same conditions as C_1 and C_2 in
 210 Lemma 5; and
- 211 (b) the closed disk bounded by the outer cycle of $C_u \cup C_2$ contains fewer vertices than the
 212 disk bounded by C_{out} .

213 **Proof.** Suppose that all the rainbows of C_u are located in uP_vu' . If z is a rainbow of C_u ,
 214 then $z \in \{u, v, u'\}$, as otherwise z is a rainbow of C_1 distinct from p, q and v , a contradiction.
 215 Thus, if $v \notin V(C_u)$, then C_u is the desired obstruction. We may assume that $v \in V(C_u)$.

216 If $u' = w$, then $C_2 = P_u \cup Q_{out}^2$, violating the assumption that $v \in V(C_2)$. Thus $u' \neq w$.
 217 If $u' = v$, then the rainbows of C_u are included in $\{u, u'\}$, and hence C_u is an obstruction.
 218 However, the existence of C_u shows that both alternatives (1) and (2) in Lemma 5 fail:
 219 condition (1) fails because C_u contains v and (2) fails because the edge of P_u incident with v
 220 is in $E(C_2) \setminus E(C_1)$. Thus $u' \neq v$.

221 The previous two paragraphs show that C_u is a near-obstruction at v with rainbows u ,
 222 v and u' . Since the interior of C_u near v is the same as the interior of C_1 near v , the pair
 223 (C_u, C_2) satisfies the ‘‘small ball’’ hypothesis. Thus, (a) holds.

224 Let C'_{out} be the outer cycle of $C_u \cup C_2$. From the fact that $C_u \cup C_2 \subseteq C_1 \cup C_2$ it follows
 225 that the disk bounded by C_{out} includes the disk bounded by C'_{out} .

226 Since $p, q \in V(C_{out})$, p and q are in the disk bounded by C_{out} . If both p and q are in
 227 C_2 , then p, q and r_1 are rainbows in C_2 , and also distinct from v , contradicting that C_2 is a
 228 near-obstruction for v . If, say $p \notin V(C_2)$, then p is not in the disk bounded by C'_{out} , which
 229 implies (b). \blacktriangleleft

230 From Claim 6(b) and assumption (*) either C_u is the desired obstruction or P_u contains
 231 a rainbow r_2 of C_2 in its interior. We assume the latter as else we are done.

232 In the same way, the last rainbow r_3 comes by considering the edge of $C_2 - Q_{out}^2$ incident
 233 with w . It follows that v, r_1, r_2 and r_3 are four different rainbows in C_2 , contradicting the
 234 fact that C_2 is a near-obstruction. ◀

235 4 Proof of Theorem 2

236 In this section we prove that a set of strings with no obstructions can be extended to an
 237 arrangement of pseudolines.

238 **Proof of Theorem 2.** It was shown in Observation 3 that the existence of obstructions
 239 implies non-extendibility. For the converse, suppose that Σ is a set of strings for which $G(\Sigma)$
 240 has no obstructions.

241 We start by reducing to the case where the point set $\bigcup \Sigma$ is connected: iteratively add a
 242 new string in a face of $\bigcup \Sigma$ connecting two connected components of $\bigcup \Sigma$. No obstruction is
 243 introduced at each step (obstructions are cycles), and, eventually, the obtained set $\bigcup \Sigma$ is
 244 connected. An extension of the new set of strings contains an extension for the original set,
 245 thus we may assume that $\bigcup \Sigma$ is connected.

246 Our proof is algorithmic, and consists of repeatedly applying one of the three steps
 247 described below.

- 248 ■ **Disentangling Step.** If a string $\sigma \in \Sigma$ has an end a with degree at least 2 in $G(\Sigma)$,
 249 then we slightly extend the a -end of σ into one of the faces incident with a .
- 250 ■ **Face-Escaping Step.** If a string $\sigma \in \Sigma$ has an end a with degree 1 in $G(\Sigma)$, and is
 251 incident with an interior face, then we extend the a -end of σ until it intersects some point
 252 in the boundary of this face.
- 253 ■ **Exterior-Meeting Step.** Assuming that all the strings in Σ have their two ends in
 254 the outer face and these ends have degree 1 in $G(\Sigma)$, we extend the ends of two disjoint
 255 strings so that they meet in the outer face.

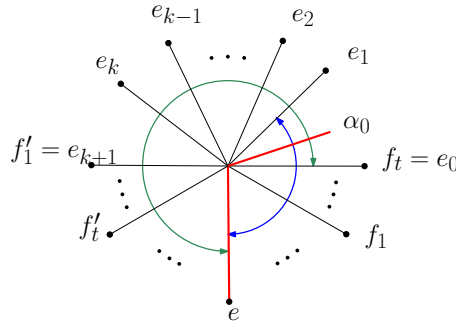
256 Each of these three steps either increases the number of pairs of strings that intersect, or
 257 increase the number crossings (recall that a crossing between σ and σ' is a non-tangential
 258 intersection point in $\sigma \cap \sigma'$ that is not an end of σ or σ'). Moreover, these steps can be
 259 performed as long as not all the strings have their ends in the outer face and they are pairwise
 260 crossing (in this case we extend their ends to infinity to obtain the desired arrangement
 261 of pseudolines). Henceforth, we will show that, if performed correctly, none of these steps
 262 introduces an obstruction. The proof for each step can be read independently.

263 ► **Lemma 7 (Disentangling Step).** *Suppose that $\sigma \in \Sigma$ has an end a with degree at least 2 in*
 264 *$G(\Sigma)$. Then we can extend the a -end of σ into one of the faces incident to a without creating*
 265 *an obstruction.*

266 **Proof.** A pair of different edges f and f' in $G(\Sigma)$ incident with a are *twins* if they belong to
 267 the same string in Σ . The edge $e \subseteq \sigma$ incident with a has no twin.

268 The fact that no pair of strings tangentially intersect at a tells us that if (f_1, f'_1) and
 269 (f_2, f'_2) are pairs of twins, then f_1, f_2, f'_1, f'_2 occur in this cyclic order for either the clockwise
 270 or counterclockwise rotation at a . Thus, we may assume that the counterclockwise rotation
 271 at a restricted to the twins and e is $e, f_1, \dots, f_t, f'_1, \dots, f'_t$, where (f_i, f'_i) is a twin pair for
 272 $i = 1, \dots, t$.

273 To avoid tangential intersections, the extension of σ at a must be in the angle between f_t
 274 and f'_1 not containing e . Let e_1, \dots, e_k be the counterclockwise ordered list of non-twin edges



287 ■ **Figure 2** Substrings included in the disk bounded by C_0 .

275 at a having an end in this angle (as depicted in Figure 2). We label $e_0 = f_t$ and $e_{k+1} = f'_1$.
 276 If there are no twins, then let $e_0 = e_{k+1} = e$.

277 Let us consider all the possible extensions: for $i \in \{0, \dots, k\}$, let Σ_i be the set of strings
 278 obtained from Σ by slightly extending the a -end of σ into the face containing the angle
 279 between e_i and e_{i+1} . Let α_i be the new edge at a extending σ in Σ_i (see α_0 in Figure 2).

280 Seeking a contradiction, suppose that, for each $i \in \{0, \dots, k\}$, $G(\Sigma_i)$ contains an obstruction
 281 C_i . Since α_i contains a degree-1 vertex, α_i is not in C_i . Hence C_i is a cycle of $G(\Sigma)$. Thus
 282 C_i is not an obstruction in $G(\Sigma)$ that becomes one in $G(\Sigma_i)$. This conversion has a simple
 283 explanation: in $G(\Sigma)$, C_i has exactly three rainbows, and one of them is a . After α_i is added,
 284 a is not a rainbow in C_i (witnessed by the edges e and α_i included in the new version of σ).

285 Recall from Section 3 that a *near-obstruction at a* is a cycle with exactly three rainbows,
 286 and one of them is a . Each of C_0, C_1, \dots, C_k is a near-obstruction at a in $G(\Sigma)$.

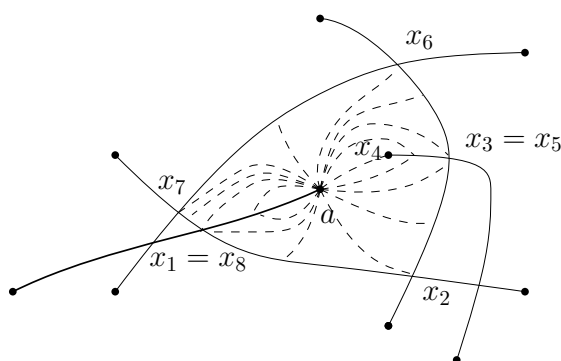
287 For a cycle $C \subseteq G$, let $\Delta(C)$ denote the closed disk bounded by C . Both e and α_0 are in
 288 $\Delta(C_0)$. Thus, either $\Delta(C_0) \supseteq \{e, f_1, f_2, \dots, f_t, e_1\}$ (blue bidirectional arrow in Figure 2) or
 289 $\Delta(C_0) \supseteq \{f_t, e_1, \dots, e_k, f'_1, f'_2, \dots, f'_t, e\}$ (green bidirectional arrow). We rule out the latter
 290 situation as the second list contains f_t and f'_t , and this would imply that a is not a rainbow
 291 for C_0 in $G(\Sigma)$.
 292

293 We just showed that $\{e, e_0, e_1\} \subseteq \Delta(C_0)$. By symmetry, $\{e_k, e_{k+1}, e\} \subseteq \Delta(C_k)$. Consider
 294 the largest index $i \in \{0, 1, \dots, k-1\}$ for which $\{e, e_0, \dots, e_{i+1}\} \subseteq \Delta(C_i)$. By the choice
 295 of i , and because $\{e, \alpha_{i+1}\} \subseteq \Delta(C_{i+1})$, $\{e, f'_t, \dots, f'_1, e_k, \dots, e_i\} \subseteq \Delta(C_{i+1})$. However, by
 296 applying Lemma 5 to the pair C_i and C_{i+1} , we obtain that $G(\Sigma)$ has an obstruction, a
 297 contradiction. ◀

298 ▶ **Lemma 8** (Face-Escaping Step). *Suppose that there is a string σ that has an end a with*
 299 *degree 1 in $G(\Sigma)$, and a is incident to an interior face F . Then there is an extension σ' of*
 300 *σ from its a -end to a point in the boundary of F such that the set $(\Sigma \setminus \{\sigma\}) \cup \{\sigma'\}$ has no*
 301 *obstruction.*

303 **Proof.** Let W be the closed boundary walk $(x_0, e_1, x_1, e_2, \dots, e_n, x_n)$ of F such that $x_0 =$
 304 $x_n = a$ and F is to the left as we traverse W (see Figure 3 for an illustration with $n = 9$).
 305 For $i = 1, \dots, n$ we let m_i be a point in the relative interior of e_i , and let P be the list of
 306 non-necessarily distinct points $m_1, x_1, m_2, x_2, \dots, m_n$, which are the potential ends for all
 307 the different extensions. For each $p \in P$, let Σ_p be the set of strings obtained from Σ by
 308 extending the a -end of σ by adding an arc α_p connecting a to p in F (see Figure 3). We
 309 assume that every two distinct arcs α_p and $\alpha_{p'}$ are internally disjoint.

310 Let f_p be the edge $e_1 \cup \alpha_p$ in $G(\Sigma_p)$; f_p has ends x_1 and p . Also, let $\sigma^p = \sigma \cup \alpha_p$. Seeking
 311 a contradiction, suppose that each $G(\Sigma_p)$ has an obstruction.



302 ■ **Figure 3** All possible extensions in the Face-Escaping Step.

312 ▷ **Claim 9.** Let $p \in P$. Then there exists an obstruction C_p in $G(\Sigma_p)$ including f_p . Moreover,

- 313 (1) if $p \in \sigma$, then C_p can be chosen so that all its edges are included in σ^p ; and
 314 (2) if $p \notin \sigma$, then every obstruction includes f_p .

315 **Proof.** First, if $p \in \sigma$, then the string σ^p self-intersects at p ; thus σ^p has a simple close curve
 316 including f_p . In this case let C_p be the cycle in $G(\Sigma_p)$ representing this simple closed curve
 317 without rainbows, and thus (1) holds.

318 Second, assume that $p \notin \sigma$ and let C_p be any obstruction of $G(\Sigma_p)$. For (2), we will show
 319 that $f_p \in E(C_p)$.

320 Seeking a contradiction, suppose that $f_p \notin E(C_p)$.

321 If $p = m_i$ for $i \in \{1, \dots, n\}$, since m_i is the only vertex whose rotation in $G(\Sigma)$ differs
 322 from its rotation in $G(\Sigma_{m_i})$, $m_i \in V(C_p)$. Consider the cycle C of $G(\Sigma)$ obtained from C_p
 323 by replacing the subpath (x_{i-1}, m_i, x_i) by the edge $x_{i-1}x_i$. For each vertex $v \in V(C)$ the
 324 colors of the edges of $G(\Sigma)$ at v included in the disk bounded by C are the same as in $G(\Sigma_p)$
 325 for the disk bounded by $V(C_p)$. Thus, C is an obstruction for $G(\Sigma)$, a contradiction.

326 Suppose now that p is one of x_1, \dots, x_{n-1} . The only vertex in $G(\Sigma)$ whose rotation is
 327 different in $G(\Sigma_p)$ is p . Therefore, p is a point that is a rainbow for C_p in $G(\Sigma)$, but not
 328 a rainbow in $G(\Sigma_p)$, witnessed by two edges included in σ^p . Since at least one of the two
 329 witnessing edges is in $G(\Sigma)$, $p \in \sigma$. This contradicts the assumption that $p \notin \sigma$. Hence
 330 $f_p \in E(C_p)$. ◀

331 Henceforth we assume that, for $p \in P$, C_p is an obstruction in $G(\Sigma_p)$ as in Claim 9.

332 More can be said about the obstructions in $G(\Sigma_p)$, but for this we need some terminology.
 333 If we orient an edge e in a plane graph, then the *sides* of e are either the points near e that
 334 are to the right of e , or the points near e to the left of e . For any cycle C of G through e ,
 335 exactly one side of e lies inside C . This is the side of e *covered* by C . For the next claim
 336 and in the rest of the proof we will assume that for $p \in P$, f_p is oriented from x_1 to p .

337 ▷ **Claim 10.** For $p \in P$ with $p \notin \sigma$, every obstruction in $G(\Sigma_p)$ covers the same side of f_p .

338 **Proof.** Suppose that for $p \in P$ there are obstructions C_p and C'_p covering both sides of f_p .
 339 Let G' be the plane graph obtained from $G(\Sigma_p)$ by subdividing f_p , and let v be the new
 340 degree-2 vertex inside f_p .

341 We consider the edge-colouring χ induced by the strings in Σ_p . Let χ' be a new colouring
 342 obtained from χ by replacing the colour of the edge vp by a new colour not used in χ . It is a

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343 routine exercise to verify that (i) χ' induces a path-partition in G' (defined in Section 3);
 344 and (ii) C_p and C'_p are near-obstructions for v with respect to χ' . By applying Lemma 5
 345 to $C_1 = C_p$ and $C_2 = C'_p$, we obtain an obstruction in G' not containing v . However, this
 346 implies the existence of an obstruction in $G(\Sigma)$, a contradiction. ◀

347 Recall that the boundary walk of F is $W = (x_0, e_1, \dots, e_n, x_n)$, with $x_0 = x_n = a$. Since
 348 x_1 and x_{n-1} are in σ , the extreme obstructions C_{x_1} and C_{x_2} cover the right of f_{x_1} and the
 349 left of $f_{x_{n-1}}$, respectively. Thus, there are two consecutive vertices x_{i-1}, x_i in $W - a$, such
 350 that the interior of $C_{x_{i-1}}$ covers the right of $f_{x_{i-1}}$ and the interior of C_{x_i} covers the left of
 351 f_{x_i} . Moreover, we may assume that the interior of C_{m_i} includes the left of f_{m_i} (otherwise
 352 we reflect our drawing).

353 The next claim (proved in the full version of this paper [4]) is the last ingredient to obtain
 354 a final contradiction.

355 ▷ **Claim 11.** Exactly one of the following holds:

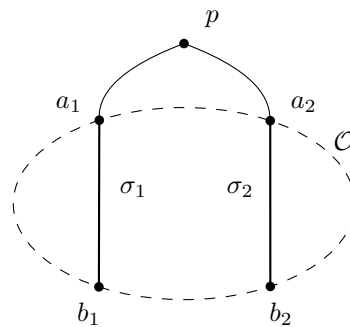
- 356 (a) $x_{i-1} \in \sigma$, $m_i \notin \sigma$ and $G(\Sigma_{m_i})$ has an obstruction covering the side of f_{m_i} not covered by
 357 C_{m_i} ; or
- 358 (b) $x_{i-1} \notin \sigma$ and $G(\Sigma_{x_{i-1}})$ has an obstruction covering the side of $f_{x_{i-1}}$ not covered by $C_{x_{i-1}}$.

359 Claims 10 and 11 contradict each other. Thus, for some $p \in P$, $G(\Sigma_p)$ has no obstructions.
 360 ◀

361 ► **Lemma 12 (Exterior-Meeting Step).** *If all the strings in Σ have their ends on the outer
 362 face of $G(\Sigma)$ and the ends have degree 1 in $G(\Sigma)$, then we can extend a pair disjoint strings
 363 so that they intersect without creating an obstruction.*

364 **Proof.** By following the outer boundary of $\bigcup \Sigma$, we obtain a simple closed curve \mathcal{O} containing
 365 all the ends of the strings in Σ , but otherwise disjoint from $\bigcup \Sigma$.

366 Suppose σ_1, σ_2 are two disjoint strings in Σ . For $i = 1, 2$, let a_i, b_i be the ends of σ_i ;
 367 since σ_1 and σ_2 do not cross, we may assume that these ends occur in the cyclic order $a_1, b_1,$
 368 b_2, a_2 . We extend the a_i -ends of σ_1 and σ_2 so that they meet in a point p in the outer face,
 369 and so that all the ends of σ_1 and σ_2 remain incident with the outer face (Figure 4). Let Σ'
 370 be the obtained set of strings.



371 ■ **Figure 4** Exterior-Meeting Step.

372 Seeking a contradiction, suppose that $G(\Sigma')$ has an obstruction C . Since $G(\Sigma)$ has no
 373 obstruction, $p \in V(C)$. Our contradiction will be to find three rainbows in C . Note that
 374 p is a rainbow. To obtain a second rainbow, traverse C starting from p towards a_1 . Let
 375 d_1 be the first vertex during our traversal that is not in the extended σ_1 , and let c_1 be its
 376 neighbour in σ_1 , one step before we reach d_1 . Since b_1 has degree one, $c_1 \neq b_1$.

377 ▷ **Claim 13.** The cycle C has a rainbow included in the closed disk Δ_1 bounded by σ_1 and
 378 the a_1b_1 -arc of \mathcal{O} disjoint from σ_2 .

379 **Proof.** First, suppose that $d_1 \notin \Delta_1$. In this case, c_1 is a rainbow because otherwise there
 380 would be a string σ that tangentially intersects σ_1 at c_1 . Thus, if $d_1 \notin \Delta_1$, then c_1 is the
 381 desired rainbow.

382 Second, suppose that $d_1 \in \Delta_1$. Let P_1 be the path of C starting at c_1 , continuing on the
 383 edge c_1d_1 , and ending at the first vertex we encounter in σ_1 . Since the cycle C' enclosed by
 384 $P_1 \cup \sigma_1$ is not an obstruction, there is one rainbow of C' that is an interior vertex of P_1 ; this
 385 is the desired rainbow of C . This concludes the proof of Claim 13. ◀

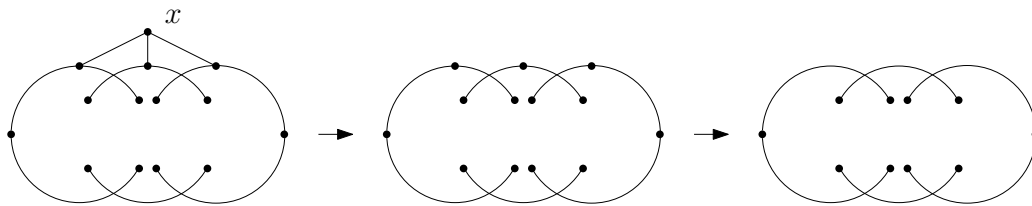
386 Considering σ_2 instead of σ_1 , Claim 13 yields a third rainbow in C inside an analogous
 387 disk Δ_2 disjoint from Δ_1 , contradicting that C is an obstruction. Hence Lemma 12 holds. ◀

388 We iteratively apply the Disentangling Step, Face-Escaping Step or Exterior-Meeting Step
 389 without creating obstructions. Each step increases the number of pairwise intersecting strings
 390 in Σ until we reach a stage where the strings are pairwise intersecting and all of them have
 391 their two ends in the unbounded face. From this we extend them into an arrangement of
 392 pseudolines. This concludes the proof of Theorem 2. ◀

393 **5 Finding obstructions and extending strings in polynomial time**

394 We start this section by describing an algorithm to detect obstructions. Henceforth, we
 395 assume that the input to the problem is the planarization $G(\Sigma)$ of an ordinary set of s strings
 396 Σ . For the running-time analysis, we assume that n and m are the number of vertices and
 397 edges in $G(\Sigma)$, respectively. Since $G(\Sigma)$ is planar, $m = O(n)$. Moreover, if Σ is pseudolinear,
 398 then $n \leq \binom{s}{2} + 2s = \binom{s+2}{2} - 1$. At the end of this section we explain how to extend Σ (if
 399 possible) in polynomial time.

401 Recall that each string in Σ receives a different colour; this induces an edge-colouring on
 402 $G(\Sigma)$ where each string is a monochromatic path. An *outer-rainbow* is a vertex $x \in V(G(\Sigma))$
 403 incident with the outer face and for which the edges incident with x have different colours.
 Next we describe the basic operation in our obstruction-detecting algorithm.



400 **Figure 5** From Σ to $\Sigma - x$.

404
 405 **Outer-rainbow deletion.** Given an outer-rainbow $x \in V(G(\Sigma))$, the instance $G(\Sigma - x)$ is
 406 defined by: first, removing x and the edges incident to x ; second, suppressing the degree-2
 407 vertices incident with edges of the same colour; and third, removing remaining degree-0
 408 vertices (Figure 5 illustrates this process). Edge colours are preserved.

409 It is easy to verify that $G(\Sigma - x)$ is the planarization of an arrangement of strings. The
 410 colours removed by this operation are those belonging to strings that are paths of length 1 in
 411 $G(\Sigma)$ incident with x . Our obstruction-detecting algorithm relies on the following property:

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412 (**) if x is an outer-rainbow of $G(\Sigma)$, then there is an obstruction in $G(\Sigma)$ not including x if
413 and only if there is an obstruction in $G(\Sigma - x)$.

414 This property holds because cycles in $G(\Sigma) - x$ and in $G(\Sigma - x)$ are in 1-1 correspon-
415 dence: two cycles correspond to each other if they are the same simple closed curve. This
416 correspondence is obstruction-preserving.

417 Let us now describe the two subroutines in our algorithm. For this, we remark that an
418 outer-rainbow of $G(\Sigma)$ is a rainbow for any cycle containing it.

419 **Subroutine 1.** *Detecting an obstruction through two outer-rainbows x and y .*

- 420 (1) Find a cycle C through x and y whose edges are incident with the outer face of $G(\Sigma)$. If
421 no such C exists, then output *No obstruction through x and y* . Else, go to Step 2.
- 422 (2) Find whether there is a third outer-rainbow $z \in V(C) \setminus \{x, y\}$. If such z exists, update
423 $G(\Sigma) \leftarrow G(\Sigma - z)$ and go to Step 1. If no such z exists, output C .

424 *Correctness and running-time of Subroutine 1:* If an obstruction through x and y exists, then
425 x and y are in the same block (some authors use the term ‘biconnected component’). Since
426 x and y are incident with the outer face, the outer boundary of the block containing x and y
427 is the cycle C from Step 1. This C can be found by considering outer boundary walk W of
428 $G(\Sigma)$ and then by finding whether x and y belong to the same non-edge block of W . Finding
429 W is $O(m)$ and computing the blocks of W via a DFS takes $O(m)$ time.

430 In Step 2, if there is a third outer rainbow z in C , then no obstruction through x and y
431 contains z . Property (**) justifies the update that takes $O(m)$ time.

432 A full run from Step 1 to Step 2 takes $O(m)$. Moving from Step 2 to Step 1 occurs $O(n)$
433 times. Thus, the total time for Subroutine 1 is $O(mn) = O(n^2)$.

434 **Subroutine 2.** *Detecting an obstruction through a single outer-rainbow x .*

- 435 (1) Find a cycle C through x whose edges are incident with the outer face of $G(\Sigma)$. If no
436 such C exists, output *No obstruction through x* . Else, go to Step 2.
- 437 (2) Find whether there is an outer-rainbow y in $V(C) \setminus \{x\}$. If no such y exists, output C .
438 Else, apply Subroutine 1 to x and y ; if there is an obstruction C' through x and y , then
439 output C' . Else, update $G(\Sigma) \leftarrow G(\Sigma - y)$ and go to Step 1.

440 *Correctness and running-time of Subroutine 2:* If $G(\Sigma)$ has an obstruction through x , then
441 there is a non-edge block in $G(\Sigma)$ containing x . The outer boundary of this block is a cycle
442 C through x having all edges incident with the outer face. As in Subroutine 1, Step 1 takes
443 $O(m)$ time.

444 Detecting the existence of y in Step 2 is $O(m)$ because to detect rainbows in C , each edge
445 incident with a vertex in $V(C)$ is verified at most twice. The update in Step 2 is justified
446 by Property (**). Since Step 2 may use Subroutine 1, Step 2 takes $O(n^2)$ time. As moving
447 from Step 2 to Step 1 occurs $O(n)$ times, the total running-time for Subroutine 2 is $O(n^3)$.

448 We are now ready for the algorithm to detect obstructions.

449 **Algorithm 1:** *Detecting obstructions in $G(\Sigma)$.*

- 450 (1) Find a cycle C having all edges incident with the outer face. If no such C exists, output
451 *No obstruction*. Else, go to step 2.
- 452 (2) Find whether there is an outer rainbow $x \in V(C)$. If not, output C . Else apply Subroutine
453 2 to x ; if there is an obstruction C' through x , output C' . Else, update $G(\Sigma) \leftarrow G(\Sigma - x)$
454 and go to Step 1.

455 *Correctness and running-time of Algorithm 1:* If $G(\Sigma)$ has an obstruction, then it has a
 456 non-trivial block whose outer boundary is a cycle C as in Step 1. As before, C and x as in
 457 Step 2 can be found in $O(m)$ steps. If C has not outer rainbow x , then C is an obstruction;
 458 Property (**) justifies the update in Step 2.

459 Since Step 2 may use Subroutine 2, a full run of Steps 1 and 2 takes $O(n^3)$ time. Since
 460 Step 2 goes to Step 1 $O(n)$ times, the running-time of Algorithm 1 is $O(n^4)$.

461 Algorithm 1 and the constructive proof of Theorem 2 imply the following result (proved
 462 in the full version of this paper [4]).

463 ► **Theorem 14.** *There is a polynomial-time algorithm to recognize and extend an ordinary*
 464 *set of strings that are extendible to an arrangement of pseudolines.*

465 6 Concluding remarks

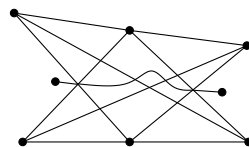
466 In this work we characterized in Theorem 2 sets of strings that can be extended into
 467 arrangements of pseudolines. Moreover, we showed that the obstructions to pseudolinearity
 468 can be detected in $O(n^4)$ time, where n is the number of vertices in the planarization of the
 469 set of strings.

470 An easy consequence of Theorem 2 is the following (presented before as Theorem 1). We
 471 prove this result in the full version of this paper [4].

472 ► **Theorem 15.** *Let D be a non-pseudolinear good drawing of a graph H . Then there is a*
 473 *subset S of edge-arcs in $\{D[e] : e \in E(H)\}$, such that each $\sigma \in S$ has a substring $\sigma' \subseteq \sigma$*
 474 *for which $\bigcup_{\sigma \in S} \sigma'$ is one of the drawings represented in Figure 1.*

475 Theorem 2 can also be applied to find a short proof that pseudolinear drawings of K_n
 476 are characterized by forbidding the B -configuration (see Theorem 2.5.1 in [3]). This implies
 477 the characterizations of pseudolinear drawings of K_n presented in [2, 5, 6].

478 A drawing is *stretchable* if it is homeomorphic to a rectilinear drawing. There are
 479 pseudolinear drawings that are not stretchable. For instance, consider the Non-Pappus
 480 configuration in Figure 6. Nevertheless, as an immediate consequence of Thomassen's main
 481 result in [16], pseudolinear and stretchable drawings are equivalent, under the assumption
 482 that every edge is crossed at most once.



483 ■ **Figure 6** Non-Pappus configuration.

484 ► **Corollary 16.** *A drawing D of a graph in which every edge is crossed at most once is*
 485 *stretchable if and only if it is pseudolinear.*

486 **Proof.** If a drawing D is stretchable then clearly it is pseudolinear. To show the converse,
 487 suppose that D is pseudolinear. Then D does not contain any obstruction, and in particular,
 488 neither of the B - and W -configurations in Figure 1 occurs in D . This condition was shown
 489 in [16] to be equivalent to being stretchable. ◀

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490 One can construct more general examples of pseudolinear drawings that are not stretchable
491 by considering non-stretchable arrangements of pseudolines. However, such examples seem to
492 inevitably have some edge with multiple crossings. This leads to a natural question.

493 ▷ **Question 17.** Is it true that if D is a pseudolinear drawing in which every edge is crossed
494 at most twice, then D is stretchable?

495 We believe that there are other instances where pseudolinearity characterizes stretchability
496 of drawings. A drawing is *near planar* if the removal of one edge produces a planar graph.
497 One instance, is the following result by Eades et al. that can be translated to the language
498 of pseudolines:

499 ► **Theorem 18.** [9] *A drawing of a near-planar graph is stretchable if and only if the drawing*
500 *induced by the crossed edges is pseudolinear.*

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