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Distributed Hypothesis Testing Over a Noisy Channel

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Abstract—A coding and testing scheme is presented for the distributed hypothesis testing problem over a noisy channel. The coding scheme combines the Shimokawa-Han-Amari hypothesis testing scheme with Borade’s unequal error protection (UEP) channel coding. The type-II error exponent of our scheme consists of three competing error exponents: two of them coincide with the exponents found by Shimokawa-Han-Amari for distributed hypothesis testing over a noiseless link (with the rate be replaced by the mutual information between channel input and output), and the third includes Borade’s miss-detection exponent for UEP over a noisy channel. Depending on the problem setup, any of the three exponents can be active. When testing against conditional independence, only the two Shimokawa-Han-Amari exponents are active, and the scheme achieves the optimal type-II error exponent found by Sreekuma and Gündüz.

I. INTRODUCTION

Consider a distributed hypothesis testing problem where a sensor describes its collected information to a remote decision center over a noisy channel. The decision center decides on a binary hypothesis ($\mathcal{H} = 0$ or $\mathcal{H} = 1$) that determines the joint probability distribution underlying its own observation and the information observed at the sensor. The goal of the communication is to maximize the type-II error (deciding $\hat{\mathcal{H}} = 0$ when $\mathcal{H} = 1$) exponent under a constrained type-I error (deciding $\hat{\mathcal{H}} = 1$ when $\mathcal{H} = 0$).

The special case of this problem where communication takes place over a noiseless link was studied in [1]–[4]. These works present achievable type-II error exponents for general joint probability distributions underlying the two hypotheses and the optimal type-II error exponent for the special case called “testing against conditional independence” [4]. Distributed hypothesis testing problems over noiseless networks with multiple sensors or decision centers or with relays have been considered in [4]–[8]. The work most closely related to this paper is by Sreekumar and Gündüz [9]. It proves that the optimal type-II error exponent for “testing against conditional independence” over a noisy channel, coincides with the optimal type-II error exponent of the same test over a noiseless link of rate equal to the capacity of the noisy channel. Their result is based on a joint hypothesis-testing and channel-coding scheme, see also [9, Remark 6] for a discussion on this.

In this work, we propose a coding scheme for distributed hypothesis testing over a noisy channel with general probability distributions. The coding and testing scheme applies separate hypothesis testing and channel coding by combining the Shimokawa-Han-Amari (SHA) hypothesis-testing scheme [3]

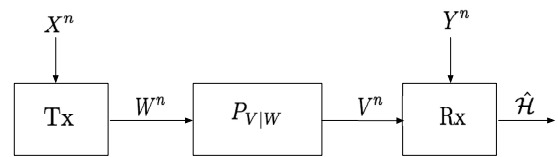


Fig. 1. Hypothesis testing over a noisy channel

with Borade’s unequal error protection (UEP) channel coding [12]. The idea is to reinforce the protection of the message that the SHA scheme produces to indicate that the transmitter decides on the alternative hypothesis $\mathcal{H} = 1$. Our analysis in general shows three competing error exponents, two of them coincide with the two competing error exponents obtained for testing over a noiseless link [3] when the communication rate is replaced by the mutual information between input and output of the channel. The third error exponent depends again on this mutual information, and on Borade’s *miss-detection exponent* [12] for channel coding with UEP. In the special case of “testing against conditional independence”, recover the optimal exponent by Sreekuma and Gündüz [9]. In this case, our third error exponent is never active and the overall type-II error exponent depends on the noisy channel only through its capacity.

Notation: We mostly follow the notation in [10]. Moreover, we use $\text{tp}(\cdot)$ to denote the *joint type* of a tuple. For a joint type π_{AB} over alphabets $\mathcal{A} \times \mathcal{B}$, we denote by $I_{\pi_{AB}}(A; B)$ the mutual information of a pair of random variables (A, B) with probability mass function (pmf) π_{AB} . Similarly for entropy, conditional entropy, and conditional mutual information. When it is unambiguous, we may abbreviate π_{AB} by π . We also abbreviate *independent and identically distributed* by i.i.d.

II. SYSTEM MODEL

Consider the distributed hypothesis testing problem in Fig. 1, where a transmitter observes source sequence X^n and a receiver source sequence Y^n . Under the null hypothesis:

$$\mathcal{H} = 0: (X^n, Y^n) \text{ i.i.d. } \sim P_{XY}, \quad (1)$$

and under the alternative hypothesis:

$$\mathcal{H} = 1: (X^n, Y^n) \text{ i.i.d. } \sim Q_{XY}. \quad (2)$$

for two given pmfs P_{XY} and Q_{XY} . The transmitter can communicate with the receiver over n uses of a discrete

memory channel $(\mathcal{W}, \mathcal{V}, P_{V|W})$ where \mathcal{W} denotes the finite channel input alphabet and \mathcal{V} the finite channel output alphabet. Specifically, the transmitter feeds inputs

$$W^n = f^{(n)}(X^n) \quad (3)$$

to the channel, where $f^{(n)}$ denotes the chosen (possibly stochastic) encoding function

$$f^{(n)} : \mathcal{X}^n \rightarrow \mathcal{W}^n. \quad (4)$$

Based on the sequence of channel outputs V^n and the source sequence Y^n , the receiver decides on the hypothesis \mathcal{H} . That means, it produces the guess

$$\hat{\mathcal{H}} = g^{(n)}(V^n, Y^n), \quad (5)$$

by means of a decoding function

$$g^{(n)} : \mathcal{V}^n \times \mathcal{Y}^n \rightarrow \{0, 1\}. \quad (6)$$

Definition 1: For each $\epsilon \in (0, 1)$, an exponent θ is said ϵ -achievable, if for each sufficiently large blocklength n , there exist encoding and decoding functions $(f^{(n)}, g^{(n)})$ such that the corresponding type-I and type-II error probabilities at the receiver

$$\alpha_n \triangleq \Pr[\hat{\mathcal{H}} = 1 | \mathcal{H} = 0], \quad (7)$$

$$\beta_n \triangleq \Pr[\hat{\mathcal{H}} = 0 | \mathcal{H} = 1], \quad (8)$$

satisfy

$$\alpha_n \leq \epsilon, \quad (9)$$

and

$$-\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \log \beta_n \geq \theta. \quad (10)$$

The goal is to maximize the type-II error exponent θ .

III. CODING AND TESTING SCHEME

We describe a coding and testing scheme for the general distributed hypothesis testing problem over a noisy channel. The analysis of the scheme is postponed to Section V.

Preparations: Choose a large positive integer n , an auxiliary distribution P_T over \mathcal{W} , a conditional channel input distribution $P_{W|T}$, and a conditional source distribution $P_{S|X}$ over a finite auxiliary alphabet \mathcal{S} so that

$$I(S; X) < I(S; Y) + I(V; W|T), \quad (11)$$

where the mutual informations in (11) are calculated according to the following joint distribution

$$P_{SXYWVT} = P_{S|X} \cdot P_{XY} \cdot P_T \cdot P_{W|T} \cdot P_{V|W}. \quad (12)$$

Then choose a sufficiently small $\mu > 0$ and nonnegative rates (R, R') so that

$$R + R' = I(X; S) + \mu \quad (13)$$

$$R < I(V; W|T) \quad (14)$$

$$R' < I(S; Y). \quad (15)$$

Code Construction: Construct a random codebook

$$\mathcal{C}_S = \{S^n(m, \ell) : m \in \{1, \dots, \lfloor 2^{nR} \rfloor\}, \ell \in \{1, \dots, \lfloor 2^{nR'} \rfloor\}\},$$

by independently drawing all codewords i.i.d. according to $P_S(s) = \sum_{x \in \mathcal{X}} P_X(x) P_{S|X}(s|x)$.

Generate a sequence T^n i.i.d. according to P_T . Construct a random codebook

$$\mathcal{C}_W = \{W^n(m) : m \in \{1, \dots, \lfloor 2^{nR} \rfloor\}\}$$

superpositioned on T^n where each codeword is drawn independently according to $P_{W|T}$ conditioned on T^n . Reveal the realizations of the codebooks and the sequence T^n to all terminals.

Transmitter: Given that it observes the source sequence $X^n = x^n$, the transmitter looks for a pair (m, ℓ) that satisfies

$$(s^n(m, \ell), x^n) \in \mathcal{T}_{\mu/2}^n(P_{SX}). \quad (16)$$

If successful, it picks one of these pairs uniformly at random and sends the codeword $w^n(m)$ over the channel. Otherwise it sends the sequence of inputs t^n over the channel.

Receiver: Assume that $V^n = v^n$ and $Y^n = y^n$ and that the “time-sharing sequence” $T^n = t^n$. The receiver first looks for an index $m' \in \{1, \dots, \lfloor 2^{nR} \rfloor\}$ so that

$$(w^n(m'), v^n, t^n) \in \mathcal{T}_{\mu}^n(P_{WVT}). \quad (17)$$

If it is not successful, it declares $\hat{\mathcal{H}} = 1$. Otherwise, it randomly picks one of the indices ℓ' that satisfy

$$H_{\text{tp}}(s^n(m', \ell'), y^n)(S|Y) = \min_{\tilde{\ell} \in \{1, \dots, \lfloor 2^{nR'} \rfloor\}} H_{\text{tp}}(s^n(m', \tilde{\ell}), y^n)(S|Y), \quad (18)$$

and checks whether

$$(s^n(m', \ell'), y^n) \in \mathcal{T}_{\mu}^n(P_{SY}). \quad (19)$$

If successful, it declares $\hat{\mathcal{H}} = 0$. Otherwise, it declares $\hat{\mathcal{H}} = 1$.

IV. AN ACHIEVABLE ERROR EXPONENT

The coding and testing scheme described in the previous section allows to establish the following theorem.

Theorem 1: Every error exponent $\theta \geq 0$ that satisfies the following condition (33) is achievable:

$$\theta \leq \max_{P_{S|X}, P_{TW}: I(S; X|Y) \leq I(W; V|T)} \min \{\theta_1, \theta_2, \theta_3\}, \quad (20)$$

where

$$\theta_1 = \min_{\substack{\tilde{P}_{SXY}: \\ \tilde{P}_{SX} = P_{SX} \\ \tilde{P}_{SY} = P_{SY}}} D(\tilde{P}_{SXY} || Q_{XY} P_{S|X}), \quad (21)$$

$$\theta_2 = \min_{\substack{\tilde{P}_{SXY}: \\ \tilde{P}_{SX} = P_{SX} \\ \tilde{P}_Y = P_Y \\ H(S|Y) \leq H_{\tilde{P}}(S|Y)}} \left[D(\tilde{P}_{SXY} || P_{S|X} Q_{XY}) + I(V; W|T) - I(S; X|Y) \right], \quad (22)$$

$$\theta_3 = D(P_Y || Q_Y) + I(V; W|T) - I(S; X|Y)$$

$$+ \sum_{t \in \mathcal{W}} P_T(t) \cdot D(P_{V|T=t} \| P_{V|W=t}), \quad (23)$$

and all expressions are calculated with respect to the joint distribution in (12).

Proof: Based on the scheme in Section V. ■

Lemma 1: It suffices to consider the auxiliary random variable S over an alphabet \mathcal{S} that is of size $|\mathcal{S}| = |\mathcal{X}| + 2$. For the special case of $P_Y = Q_Y$, it suffices to consider $|\mathcal{S}| = |\mathcal{X}| + 1$.

Proof: Based on Carathéodory's theorem. Omitted. ■

Our coding and testing scheme combines the SHA hypothesis testing scheme for a noiseless link [3] with Borade's UEP channel coding that protects the 0-message (which indicates that the transmitter decides on $\mathcal{H} = 1$) better than the other messages [11], [12]. In fact, since here we are only interested in the type-II error exponent, the receiver should decide on $\mathcal{H} = 0$ only if the transmitter also shares this opinion.

The expressions in Theorem 1 show three competing error exponents. In (21) and (22), we recognize the two competing error exponents of the SHA scheme for the noiseless setup: θ_1 is the exponent associated to the event that the receiver reconstructs the correct binned codeword and θ_2 is associated to the event that either the binning or the noisy channel introduces a decoding error. The exponent θ_3 in (23) is new and can be associated to the event that the specially protected 0-message is wrongly decoded. We remark in particular that θ_3 contains the term

$$E_{\text{miss}} := \sum_{t \in \mathcal{W}} P_T(t) \cdot D(P_{V|T=t} \| P_{V|W=t}), \quad (24)$$

which represents the largest possible *miss-detection exponent* for a single specially protected message at a given rate $I(W; V|T)$ [12, Th. 34].

Which of the three exponents $\theta_1, \theta_2, \theta_3$ is smallest depends on the source and channel parameters and the choice of $P_{S|X}$ and P_W . Notice that the third error exponent θ_3 is inactive for channels with large miss-detection exponent (24), such as binary symmetric channels with small cross-over probability, or for sources where

$$\min_{\substack{\tilde{P}_{SXY}: \\ \tilde{P}_{SX}=P_{SX} \\ \tilde{P}_Y=P_Y}} D(\tilde{P}_{SXY} \| P_{S|X} Q_{XY}) = D(P_Y \| Q_Y), \quad (25)$$

This is the case for example when “testing against conditional independence” [4] where both terms are 0.

Corollary 1 (Lemma 5 in [9]): Consider the “testing against independence” setup where

$$Y = (\bar{Y}, Z), \quad (26)$$

and $Q_{X\bar{Y}Z}$ decomposes as

$$Q_{X\bar{Y}Z} = P_{XZ} \cdot P_{\bar{Y}|Z}. \quad (27)$$

Error exponent $\theta \geq 0$ is achievable if,

$$\theta \leq \max_{\substack{P_{S|X}, \tilde{P}_W: \\ I(S; X|Z) \leq I(W; V)}} I(S; \bar{Y}|Z), \quad (28)$$

where mutual informations are calculated with respect to the joint law $P_{X\bar{Y}Z} P_{S|X} P_W P_{V|W}$.

Proof: Fix independent random variables T and W and a random variable S so that

$$I(S; X|Z) \leq I(W; V|T) = I(W; V). \quad (29)$$

Then, Theorem 1 specializes to:

$$\begin{aligned} \theta_1 &= \min_{\substack{\tilde{P}_{SX\bar{Y}Z}: \\ \tilde{P}_{SX}=P_{SX} \\ \tilde{P}_{S\bar{Y}Z}=P_{S\bar{Y}Z}}} D(\tilde{P}_{SX\bar{Y}Z} \| Q_{X\bar{Y}Z} P_{S|X}) \\ &= \min_{\substack{\tilde{P}_{SX\bar{Y}Z}: \\ \tilde{P}_{SX}=P_{SX} \\ \tilde{P}_{S\bar{Y}Z}=P_{S\bar{Y}Z}}} D(\tilde{P}_{SX\bar{Y}Z} \| P_{XZ} P_{\bar{Y}|Z} P_{S|X}) \\ &= D(P_{S\bar{Y}Z} \| P_Z P_{\bar{Y}|Z} P_{S|Z}) \\ &= I(S; \bar{Y}|Z). \end{aligned}$$

Moreover, exponents θ_2 and θ_3 cannot be smaller than $I(S; \bar{Y}|Z)$ because of the nonnegativity of the KL-divergence and the mutual information and because

$$\begin{aligned} I(V; W) - I(S; X) + I(S; \bar{Y}, Z) \\ &= I(V; W) - I(S; X|Z) + I(S; \bar{Y}|Z) \\ &\geq I(S; \bar{Y}|Z), \end{aligned} \quad (30)$$

where the inequality holds by (29). ■

Notice that the error exponent in Corollary 1 is optimal [9].

We now present an example and evaluate the largest type-II error exponents attained by our scheme. We also show that depending on the choice of the model parameters, a different error exponent θ_1, θ_2 , or θ_3 is active.

Example 1: Let under the null hypothesis

$$\begin{aligned} \mathcal{H} = 0: \quad X &\sim \text{Bern}(p_0), \quad Y = X \oplus N_0, \\ N_0 &\sim \text{Bern}(q_0), \end{aligned} \quad (31)$$

for N_0 independent of X . Under the alternative hypothesis:

$$\mathcal{H} = 1: \quad X \sim \text{Bern}(p_1), \quad Y \sim \text{Bern}(p_0 \star q_0), \quad (32)$$

with X and Y independent. Assume that $P_{V|W}$ is a binary symmetric channel (BSC) with cross-over probability $r \in [0, 1/2]$.

For this example, $P_Y = Q_Y$ and Theorem 1 simplifies to:

$$\theta \leq \max_{P_{S|X}, P_{TW}: I(S; X|Y) \leq I(W; V|T)} \min \{\theta_1, \theta_2, \theta_3\}, \quad (33)$$

where

$$\theta_1 \leq D(P_X \| Q_X) + I(S; Y), \quad (34)$$

$$\theta_2 \leq D(P_X \| Q_X) + I(V; W|T) + I(S; Y) - I(S; X), \quad (35)$$

$$\theta_3 \leq \sum_{t \in \mathcal{W}} P_T(t) D(P_{V|T=t} \| P_{V|W=t}) + I(V; W|T) + I(S; Y) - I(S; X). \quad (36)$$

Depending on the parameters of the setup and the choice of the auxiliary distributions, either of the exponents θ_1, θ_2 , or θ_3 is active. For example, when the cross-over probability of the BSC is large, $r \geq 0.4325$,

$$D(P_X \| Q_X) \geq \sum_{t \in \mathcal{W}} P_T(t) D(P_{V|T=t} \| P_{V|W=t}) + I(V; W|T), \quad (37)$$

and irrespective of the choice of the random variables S, T, W the exponent θ_3 is smaller than θ_1 and θ_2 . It is then optimal to choose S constant and (T, W) so as to maximize the sum $\sum_{t \in \mathcal{W}} P_T(t) D(P_{V|T=t} \| P_{V|W=t}) + I(V; W|T)$. In particular, for a scenario with parameters $p_0 = 0.1, q_0 = 0.25, p_1 = 0.2$ and $r = \frac{4}{9}$ one obtains numerically that the optimal error exponent achieved by our scheme is $\theta = 0.0358$.

In contrast, when the cross-over probability of the BSC is small, the miss-detection exponent (24) is large and the exponent θ_3 is never active irrespective of the choice of the auxiliary random variable S . The overall exponent is then determined by the smaller of θ_1 and θ_2 , and in particular by a choice S, X, W that makes the two equal. In this case, for a scenario with parameters $p_0 = 0.2, q_0 = 0.3, p_1 = 0.4$, and $r = 0.1$, the largest exponent achieved by our scheme is $\theta = 0.19$.

V. PROOF OF THEOREM 1

The proof of the theorem is based on the scheme in Section III. Before analyzing this scheme, notice that by the functional representation lemma, there exists a function γ over appropriate domains and for each time $t \in \{1, \dots, n\}$ a random variable ϕ_t over a finite alphabet Φ so that the time- t channel input and output satisfy:

$$V_t = \xi(W_t, \phi_t). \quad (38)$$

Let \mathcal{P}^n be the set of all types over the product alphabets $\mathcal{S}^n \times \mathcal{S}^n \times \mathcal{W}^n \times \mathcal{W}^n \times \mathcal{W}^n \times \mathcal{V}^n \times \Phi^n \times \mathcal{X}^n \times \mathcal{Y}^n$, and let \mathcal{P}_μ^n be the subset of types $\pi_{SS'TWW'V\phi XY} \in \mathcal{P}^n$ that simultaneously satisfy the following conditions:

$$|\pi_{SX} - P_{SX}| \leq \mu/2, \quad (39a)$$

$$|\pi_{S'Y} - P_{S'Y}| \leq \mu, \quad (39b)$$

$$|\pi_{TW'V} - P_{TW'V}| \leq \mu, \quad (39c)$$

$$\pi_{V|\phi TW} = \mathbb{1}\{V = \xi(\phi, T, W)\}, \quad (39d)$$

$$H_{\pi_{S'Y}}(S|Y) \leq H_{\pi_{SY}}(S|Y). \quad (39e)$$

We first analyze the type-I error probability averaged over the random code construction. Let (M, L) be the indices of the codeword chosen at the transmitter, if they exist, and define the following events:

$$\mathcal{E}_{\text{Tx}}: \{\exists(m, \ell): (S^n(m, \ell), X^n) \in \mathcal{T}_{\mu/2}^n(P_{SX})\} \quad (40)$$

$$\mathcal{E}_{\text{Rx}}^{(1)}: \{(S^n(M, L), Y^n) \notin \mathcal{T}_\mu^n(P_{SY})\} \quad (41)$$

$$\mathcal{E}_{\text{Rx}}^{(2)}: \{\exists m' \neq M: (T^n, W^n(m'), V^n) \in \mathcal{T}_\mu^n(P_{TW}P_{V|W})\} \quad (42)$$

$$\mathcal{E}_{\text{Rx}}^{(3)}: \{\exists \ell' \neq L:$$

$$H_{\text{tp}(s^n(M, \ell'), y^n)}(S|Y) = \min_{\tilde{\ell}} H_{\text{tp}(s^n(M, \tilde{\ell}), y^n)}(S|Y)\}. \quad (43)$$

With these definitions, we obtain for all sufficiently small values of μ and sufficiently large blocklengths n :

$$\alpha_n \leq \Pr[\mathcal{E}_{\text{Tx}}] + \Pr[\mathcal{E}_{\text{Rx}}^{(1)} | \mathcal{E}_{\text{Tx}}^c] + \Pr[\mathcal{E}_{\text{Rx}}^{(2)} | \mathcal{E}_{\text{Tx}}^c, \mathcal{E}_{\text{Rx}}^{(1)c}] + \Pr[\mathcal{E}_{\text{Rx}}^{(3)} | \mathcal{E}_{\text{Rx}}^{(1)c}, \mathcal{E}_{\text{Tx}}^c] \quad (44)$$

$$\leq \epsilon/4 + \epsilon/4 + \epsilon/4 + \epsilon/4 = \epsilon, \quad (45)$$

where the first summand of (44) can be upper bounded by means of the covering lemma [10] and the rate constraint (15); the second by means of the Markov lemma [10]; the third by means of the packing lemma [10] and the rate constraint (14); and the fourth by following similar steps as in analysis of the type-I error probability in [5, Appendix H].

Now, consider the type-II error probability. Let $\mathcal{P}_{\mu,0}^n$ be the subset of types $\pi_{S'TW'V\phi XY}$ over the alphabets $\mathcal{S}^n \times \mathcal{W}^n \times \mathcal{W}^n \times \Phi^n \times \mathcal{V}^n \times \mathcal{X}^n \times \mathcal{Y}^n$ that satisfy (39b), (39c), and

$$\pi_{V|\phi T} = \mathbb{1}\{V = \xi(T, \phi)\}. \quad (46)$$

Define for each pair $(m, m') \in \{1, \dots, [2^{nR}]\}^2$ and $(\ell, \ell') \in \{1, \dots, [2^{nR'}]\}^2$ the set:

$$\mathcal{A}(m, m', \ell, \ell') := \{(\varphi^n, x^n, y^n) : \text{tp}(S^n(m, \ell), S^n(m', \ell')), W^n(m), W^n(m'), \varphi^n, \xi^n(W^n(m), \varphi^n), x^n, y^n) \in \mathcal{P}_\mu^n\};$$

and for each $m' \in \{1, \dots, [2^{nR}]\}$ and $\ell' \in \{1, \dots, [2^{nR'}]\}$ the set:

$$\mathcal{A}(0, m', \ell') := \{(\varphi^n, x^n, y^n) : \text{tp}(S^n(m', \ell'), T^n, W^n(m'), \varphi^n, \xi^n(T^n, \varphi^n), x^n, y^n) \in \mathcal{P}_{\mu,0}^n\}. \quad (47)$$

By $\xi^n(W^n(m), \varphi^n)$, here we mean the component-wise application of the function $\xi(\cdot, \cdot)$ defined in (38) to the n -length sequences $W^n(m)$ and φ^n .

Define the region $\mathcal{A}_{\text{Rx},n} \subseteq \Phi^n \times \mathcal{X}^n \times \mathcal{Y}^n$

$$\mathcal{A}_{\text{Rx},n} \triangleq \bigcup_{m, m'} \bigcup_{\ell, \ell'} \mathcal{A}(m, m', \ell, \ell') \cup \bigcup_{m', \ell'} \mathcal{A}(0, m', \ell'), \quad (48)$$

where m and m' take value in $\{1, \dots, [2^{nR}]\}$ and ℓ and ℓ' in $\{1, \dots, [2^{nR'}]\}$. Notice that $\mathcal{A}_{\text{Rx},n}$ is deterministic for a given codebook, but random in the analysis here.

Since $\mathcal{A}_{\text{Rx},n}$ includes the acceptance region at the receiver, the average (over the random codebooks) type-II error probability is upper bounded as:

$$\mathbb{E}_{\mathcal{C}}[\beta_n] \leq \Pr[(\phi^n, X^n, Y^n) \in \mathcal{A}_{\text{Rx},n} | \mathcal{H} = 1]. \quad (49)$$

We can then write:

$$\mathbb{E}_{\mathcal{C}}[\beta_n] \leq \Pr[(\phi^n, X^n, Y^n) \in$$

$$\begin{aligned}
 & \bigcup_{m,m'} \bigcup_{\ell,\ell'} \mathcal{A}(m, m', \ell, \ell') \cup \bigcup_{m',\ell'} \mathcal{A}(0, m', \ell') | \mathcal{H} = 1 \Big] \\
 \leq & \Pr \left[(\phi^n, X^n, Y^n) \in \bigcup_{(m,\ell)} \mathcal{A}(m, m, \ell, \ell) | \mathcal{H} = 1 \right] \\
 & + \Pr \left[(\phi^n, X^n, Y^n) \in \bigcup_{(m,\ell) \neq (m',\ell')} \mathcal{A}(m, m', \ell, \ell') | \mathcal{H} = 1 \right] \\
 & + \Pr \left[(\phi^n, X^n, Y^n) \in \bigcup_{m',\ell'} \mathcal{A}(0, m', \ell') | \mathcal{H} = 1 \right]. \quad (50)
 \end{aligned}$$

In a similar way as in [5], it can be shown that for sufficiently large blocklengths n , the first probability in (50) is upper bounded as:

$$\Pr \left[(\phi^n, X^n, Y^n) \in \bigcup_m \bigcup_\ell \mathcal{A}(m, m, \ell, \ell) | \mathcal{H} = 1 \right] \leq 2^{-n\theta_{1,\mu}}, \quad (51)$$

where

$$\theta_{1,\mu} := \min_{\pi_{SXY} \in \mathcal{P}_\mu^n} D(\pi_{SXY} || P_{S|X} Q_{XY}) - \delta(\mu) \quad (52)$$

for a function $\delta(\mu)$ that goes to zero as $\mu \rightarrow 0$. Moreover, for sufficiently large n , the second probability in (50) is upper bounded as:

$$\Pr \left[(\phi^n, X^n, Y^n) \in \bigcup_{(m,\ell) \neq (m',\ell')} \mathcal{A}(m, m', \ell, \ell') \right] \leq 2^{-n\theta_{2,\mu}}, \quad (53)$$

where

$$\begin{aligned}
 \theta_{2,\mu} := & \min_{\pi_{S'WVXY} \in \mathcal{P}_\mu^n} D(\pi_{SXY} || P_{S|X} Q_{XY}) \\
 & + I(S; Y) + I(V; W|T) - I(S; X) - \delta'(\mu), \quad (54)
 \end{aligned}$$

for a function $\delta'(\mu)$ that goes to zero as $\mu \rightarrow 0$.

The last term in (50) is upper bounded for sufficiently large blocklength n :

$$\begin{aligned}
 & \Pr \left[(\phi^n, X^n, Y^n) \in \bigcup_{m',\ell'} \mathcal{A}(0, m', \ell') | \mathcal{H} = 1 \right] \\
 \leq & \sum_{m',\ell'} \Pr \left[(\phi^n, X^n, Y^n) \in \mathcal{A}(0, m', \ell') | \mathcal{H} = 1 \right] \\
 \leq & \sum_{m',\ell'} \sum_{\pi_{S'TW'\phi VXY} \in \mathcal{P}_{\mu,0}^n} \\
 & \Pr \left[\text{tp} \left(S^n(m', \ell'), T^n, W^n(m'), \phi^n, \xi^n(T^n, \phi^n), X^n, Y^n \right) \right. \\
 & \quad \left. = \pi_{S'TW'\phi VXY} \Big| \mathcal{H} = 1 \right] \\
 \leq & \sum_{m',\ell'} \sum_{\pi_{S'TW'\phi VXY} \in \mathcal{P}_{\mu,0}^n} 2^{-nD(\pi_{S'TW'\phi VXY} || P_S P_{TW} P_\phi \pi_{V|\phi T} Q_{XY})}
 \end{aligned}$$

where the last inequality holds by the way the random codebooks are generated and because given $\mathcal{H} = 1$, the sources X^n, Y^n are i.i.d. $\sim Q_{XY}$. Define now

$$\tilde{\theta}_{3,\mu} := \min_{\pi_{S'TW'\phi VXY} \in \mathcal{P}_{\mu,0}^n} D(\pi_{S'TW'\phi VXY} || P_S P_{TW} P_\phi \pi_{V|\phi T} Q_{XY})$$

$$- R - R' - \mu \quad (55)$$

and notice that there exist functions $\delta''(\mu)$ that $\rightarrow 0$ as $\mu \rightarrow 0$ and so that the following inequalities hold:

$$\begin{aligned}
 \tilde{\theta}_{3,\mu} & \stackrel{(a)}{\geq} \min_{\pi_{S'TW'\phi VXY} \in \mathcal{P}_{\mu,0}^n} \left[D(\pi_{TW'\phi V} || P_{TW} P_\phi \pi_{V|\phi T}) \right. \\
 & \quad \left. + D(\pi_{XY} || Q_{XY}) + \mathbb{E}_{\pi_{XY}} [D(\pi_{S'|XY} || P_S)] \right] \\
 & \quad - I(S; X) - 2\mu \\
 & \stackrel{(b)}{\geq} \min_{\pi_{S'TW'\phi VXY} \in \mathcal{P}_{\mu,0}^n} \left[D(\pi_{TW'V} || P_{TW} P_{V|W=T}) \right] \\
 & \quad + D(P_Y || Q_Y) + I(S; Y) - I(S; X) - \delta''(\mu) \\
 & \stackrel{(c)}{\geq} \mathbb{E}_{P_{TW}} [D(P_{V|W} || P_{V|W=T})] \\
 & \quad + D(P_Y || Q_Y) + I(S; Y) - I(S; X) - \delta''(\mu) \\
 & = D(P_Y || Q_Y) + I(V; W|T) + I(S; Y) - I(S; X) \\
 & \quad + \sum_t P_T(t) \cdot D(P_{V|T=t} || P_{V|W=t}) - \delta''(\mu) \\
 & := \theta_{3,\mu}. \quad (56)
 \end{aligned}$$

All three inequalities are based on the data processing inequality for KL-divergences; (a) also uses (13); and (b) and (c) also use the continuity of KL-divergences and that all types in $\mathcal{P}_{0,\mu}^n$ satisfy (39b), (39c), and (46). Thus, for sufficiently large n :

$$\Pr \left[(\phi^n, X^n, Y^n) \in \bigcup_{m',\ell'} \mathcal{A}(0, m', \ell') | \mathcal{H} = 1 \right] \leq 2^{-n\theta_{3,\mu}}. \quad (57)$$

Combining (50), (51), (53), and (57), taking $\mu \rightarrow 0$ and $n \rightarrow \infty$, the proof can be established by standard arguments.

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