

# Minimal graphs for hamiltonian extension

Christophe Picouleau

#### ▶ To cite this version:

Christophe Picouleau. Minimal graphs for hamiltonian extension. Open Journal of Discrete Applied Mathematics,  $2020,\ 10.30538/psrp-odam2020.0026$ . hal-02434695

# HAL Id: hal-02434695 https://hal.science/hal-02434695v1

Submitted on 17 Jan 2020

**HAL** is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

# Minimal graphs for hamiltonian extension

#### C. Picouleau \*

#### December 11, 2019

#### Abstract

For every  $n \geq 3$  we determine the minimum number of edges of graph with n vertices such that for any non edge xy there exits a hamiltonian cycle containing xy.

**Keywords**: 2-factor, hamiltonian cycle, hamiltonian path.

## 1 Introduction

For all graph theoretical terms and notations not defined here the reader is referred to [1]. We only consider simple finite loopless undirected graphs. For a graph G = (V, E) with |V| = n vertices, an edge is a pair of two connected vertices x, y, we denote it by  $xy, xy \in E$ ; when two vertices x, y are not connected this pair form the non-edge  $xy, xy \notin E$ . In G a 2-factor is a subset of edges  $F \subset E$  such that every vertex is incident to exactly two edges of F. Since G is finite a 2-factor consists of a collection of vertex disjoint cycles spanning the vertex set V. When the collection consists of an unique cycle the 2-factor is connected, so it is a hamiltonian cycle.

We intend to determine, for any integer  $n \geq 3$ , a graph G = (V, E), n = |V| with a minimum number of edges such that for every non-edge xy it is always possible to include the non-edge xy into a connected 2-factor, i.e., the graph  $G_{xy} = (V, E \cup \{xy\})$  has a hamiltonian cycle  $H, xy \in H$ . In other words for any non-edge xy of G there exits a hamiltonian path between x and y.

This problem is related to the minimal 2-factor extension studied in [3] in which the 2-factors are not necessary connected. It is also related to the problem of finding minimal graphs for non-edge extensions in the case of perfect matchings (1-factors) studied in [2].

**Definition 1.1** Let G = (V, E) be a graph and  $xy \notin E$  an non-edge. If  $G_{xy} = (V, E \cup \{xy\})$  has a hamitonian cycle that contains xy we shall say that xy has been extended (to a connected 2-factor, to an hamiltonian cycle).

**Definition 1.2** A graph G = (V, E) is connected 2-factor expandable or hamiltonian expandable (shortly expandable) if every non-edge  $xy \notin E$  can be extended.

<sup>\*</sup>Conservatoire National des Arts et Métiers, CEDRIC laboratory, Paris (France). Email: christophe.picouleau@cnam.fr

**Definition 1.3** An expandable graph G = (V, E) with |V| = n and a minimum number of edges is a minimum expandable graph. The size |E| of its edge set is denoted by  $Exp_h(n)$ .

The case where the 2-factor is not constrained to be hamiltonian is studied in [3]. In this context  $Exp_2(n)$  denotes the size of a minimum expandable graph with n vertices. It follows that  $Exp_h(n) \geq Exp_2(n)$ .

We use the following notations. For G = (V, E), N(v) is the set of neighbors of a vertex v,  $\delta(G)$  is the minimum degree of a vertex. A vertex with exactly k neighbors is a k-vertex. When  $P = v_i, \ldots, v_j$  is a sequence of vertices that corresponds to a path in G, we denote by  $\bar{P} = v_j, \ldots, v_i$  its mirror sequence (both sequences correspond to the same path).

We state our result.

**Theorem 1.1** The minimum size of a connected 2-factor expandable graph is:

$$Exp_h(3) = 2, Exp_h(4) = 4, Exp_h(5) = 6; Exp_h(n) = \lceil \frac{3}{2}n \rceil, n \ge 6$$

**Proof**: For  $n \geq 3$  we have  $Exp_h(n) \geq Exp_2(n)$ .

In [3] it is proved that the three graphs given by Fig. 1 are minimum for 2-factor extension. They are also minimum expandable for connected 2-factor extension.

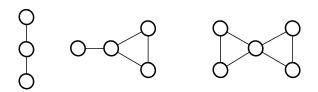


Figure 1:  $P_3$ , the paw, the butterfly.

Now let  $n \ge 6$ . From [3] we know the following when G a minimum expandable graph for the 2-factor extension:

- $\bullet$  G is connected;
- if  $\delta(G) = 1$  then  $Exp_2(n) \ge \frac{3}{2}n$ ;
- for  $n \geq 7$ , if u, v are two 2-vertices such that  $N(u) \cap N(v) \neq \emptyset$  then  $Exp_2(n) \geq \frac{3}{2}n$ ;

The graph given by Fig. 2 is minimum for 2-factor extension (see [3]). One can check that it is expandable for connected 2-factor extension. So we have  $Exp_h(6) = 9 = \frac{3}{2}n$ .



Figure 2: A minimum hamiltonian expandable graph with 6 vertices.

Suppose that G is a minimum expandable graph with  $n \geq 7$  and  $\delta(G) = 2$ . Let  $v \in V$  with d(v) = 2,  $N(v) = \{u_1, u_2\}$ . If  $u_1u_2 \notin E$  then  $u_1u_2$  cannot be expanded into a hamiltonian cycle. So  $u_1u_2 \in E$ . If  $d(u_1) = 2$  then  $u_2 \in N(u_1) \cap N(v)$  and  $Exp_h(n) \geq \frac{3}{2}n$ . So from now one we may assume  $d(u_1), d(u_2) \geq 3$ . Suppose that  $d(u_1) = d(u_2) = 3$ . Let  $N(u_1) = \{v, u_2, v_1\}, N(u_2) = \{v, u_1, v_2\}$ . If  $v_1 \neq v_2$  then  $u_1v_2$  is not expandable. If  $v_1 = v_2$  then  $v_1$  is not expandable. From now we can suppose that  $d(u_1) \geq 3, d(u_2) \geq 4$ . Moreover v is the unique 2-vertex in  $N(u_2)$ . It follows that every 2-vertex  $u \in V$  can be matched with a distinct vertex  $u_2$  with  $d(u_2) \geq 4$ . Then  $\sum_{v \in V} d(v) \geq 3n$  and thus  $m \geq \frac{3}{2}n$ .

When  $\delta(G) \geq 3$  we have  $\Sigma_{v \in V} d(v) \geq 3n$ . Thus for any expandable graph we have  $|E| = m \geq \frac{3}{2}n, n \geq 7$ .

For any even integer  $n \geq 8$  we define the graph  $G_n = (V, E)$  as follows. Let n = 2p,  $V = A \cup B$  where  $A = \{a_1, \ldots, a_p\}$  and  $B = \{b_1, \ldots, b_p\}$ . A (resp. B) induces the cycle  $C_A = (A, E_A)$  with  $E_A = \{a_1a_2, a_2a_3, \ldots, a_pa_1\}$  (resp.  $C_B = (B, E_B)$  with  $E_B = \{b_1b_2, b_2b_3, \ldots, b_pb_1\}$ . Now  $E = E_A \cup E_B \cup E_C$  with  $E_C = \{a_2b_2, a_3b_3, \ldots, a_{p-1}b_{p-1}, a_1b_p, a_pb_1\}$ . Note that  $G_n$  is cubic so  $m = \frac{3}{2}n$ . (see  $G_{10}$  in Fig. 3)

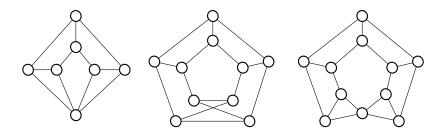


Figure 3: The graphs  $G_7, G_{10}, G_{11}$ , from the left to the right.

We show that  $G_n$  is expandable. First we consider a non-edge  $a_i a_j, p \geq j > i \geq 1$ . Note that the case of a non-edge  $b_i b_j$  is analogous. We have  $j \geq i + 2$  and since  $a_1 a_p \in E$  from symmetry we can suppose that j < p. Let  $P = a_j, a_{j-1}, \ldots, a_{i+1}, b_{i+1}, b_{i+2}, \ldots, b_{j+1}, a_{j+1}, a_{j+2}, b_{j+2}, \ldots, c_j$  where  $c_j$  is either  $a_p$  or  $b_p$  and let  $Q = a_i, b_i, b_{i-1}, a_{i-1}, \ldots, c_i$  where  $c_i$  is either  $a_1$  or  $b_1$ . From P and Q one can obtain an hamiltonian cycle containing  $a_i b_j$  whatever  $c_i$  and  $c_j$  are.

Now we consider a non-edge  $a_ib_j$ . Without loss of generality we assume  $j \geq i$ . Suppose first that j=i, so either i=1 or i=p. Without loss of generality we assume i=j=1:  $a_1,b_p,b_{p-1},\ldots,b_2,a_2,a_3,\ldots,a_p,b_1,a_1$  is an hamiltonian cycle. Now assume that j>i: Let  $P_j=b_j,b_{j-1},\ldots,b_{i+1},a_{i+1},a_{i+2},\ldots,a_{j+1},b_{j+1},b_{j+2},a_{j+2},\ldots,c_p$  where either  $c_p=a_p$  or  $c_p=b_p, P_i=a_i,b_i,b_{i-1},a_{i-1},a_{i-2},\ldots,c_1$  where either  $c_1=a_1$  or  $c_1=b_1$ . If  $c_p=a_p$  and  $c_1=a_1$  then  $P_j,b_1,b_p,P_i,a_j$  is an hamiltonian cycle. If  $c_p=a_p$  and  $c_1=b_1$  then  $P_j,a_1,b_p,P_i,a_j$  is an hamiltonian cycle. The two other cases are symmetric.

For any odd integer  $n = 2p + 1 \ge 7$  we define the graph  $G_n = (V, E)$  as follows. We set  $V = A \cup B \cup \{v_n\}$  where  $A = \{a_1, \ldots, a_p\}$  and  $B = \{b_1, \ldots, b_p\}$ .  $A \cup \{v_n\}$  (resp.  $B \cup \{v_n\}$ ) induces the cycle  $C_A = (A \cup \{v_n\}, E_A)$  with  $E_A = \{a_1a_2, a_2a_3, \ldots, a_pv_n, v_na_1\}$  (resp.  $C_B = (B \cup \{v_n\}, E_B)$  with  $E_B = \{b_1b_2, b_2b_3, \ldots, b_pv_n, v_nb_1\}$ . Now  $E = E_A \cup E_B \cup E_C$  with  $E_C = \{a_ib_i | 1 \le i \le p\} \cup \{a_1v_n, b_1v_n, a_pv_n, b_pv_n\}$ . Note that  $m = \lceil \frac{3}{2}n \rceil$ . (see  $G_7$  and  $G_{11}$  in Fig. 3)

We show that  $G_n$  is expandable. First, we consider a non-edge  $a_i a_j, p \geq j > i \geq 1$  (the case of a non-edge  $b_i b_j$  is analogous).  $a_i, a_{i+1}, \ldots, a_{j-1}, b_{j-1}, b_{j-2}, b_{j-3}, \ldots, b_i, b_{i-1}, a_{i-1}, a_{i-2}, b_{i-2}, \ldots, v_n, c_p, d_p, d_{p-1}, c_{p-1}, \ldots, c_j, d_j$ , where  $d_j = a_j$  and for any  $k, j \leq k \leq p$ , the ordered pairs  $c_k, d_k$  correspond to either  $a_k, b_k$  or  $b_k, a_k$ , is an hamiltonian cycle. Second, let a non-edge  $a_i b_j, p \geq j > i \geq 1$ . We use the same construction as above taking  $d_j = b_j$ .

#### Acknowledgements:

The author express its gratitude to Dominique de Werra for its constructive comments and remarks, which helped to improve the writing of this paper.

### References

- [1] J.A. Bondy, U. S. R. Murty, Graph Theory, Springer, (2008).
- [2] M.-C. Costa, D. de Werra, C. Picouleau, Minimal graphs for matching extensions, Discrete Applied Mathematics, 234 (2018), 47-55.
- [3] M.-C. Costa, D. de Werra, C. Picouleau, *Minimal graphs for 2-factor extension*, Discrete Applied Mathematics, accepted for publication.