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# Second order balance property on Christoffel words

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**Abstract.** In this paper we study the *balance matrix* that gives the order of balance of any binary word. In addition, we define for Christoffel words a new matrix called *second order balance matrix*. This matrix gives more information on the balance property of a word that codes the number of occurrences of the letter 1 in successive blocks of the same length for the studied Christoffel word. By taking the maximum of the Second order balance matrix we define the second order of balance and we are able to order the Christoffel words according to these values. Our construction uses extensively the continued fraction associated with the slope of each Christoffel word, and we prove a recursive formula based on fine properties of the Stern-Brocot tree to construct second order matrices. Finally, we show that an infinite path on the Stern-Brocot tree, which minimizes the second order of balance is given by a path associated with the Fibonacci word.

**Index terms**— Balance property. Second order balance property. Christoffel words. Stern-Brocot tree. Continued fractions

## 1 Introduction

Balanced words appear in many developments of combinatorics on words and the balance property is considered as a fine tool to investigate the structure of words [17, 21]. As a typical example of infinite balanced words, Sturmian words could be constructed equivalently by discretizations of irrational slope lines in a square grid [11, 30], by billiard words in a square [2, 24] or by coding of irrational rotations on a unit circle with a partition in two intervals [11]. The finite balanced words are given by discretizations of rational slope lines in a square grid and have been studied in particular by Christoffel [9]. Interestingly, finite and infinite balanced words show up in specific optimization problems [1, 26, 30] and for example optimal schedules for job-shop problems with two tasks are exactly given by balanced words [14, 19, 30]. Furthermore, particular solutions of job-shop problems with  $k$  tasks sharing the same resource [1] are given by finite or infinite balanced words on a  $k$ -letters alphabet where the balanced property is checked on each letter of the alphabet [1, 23, 29]. More precisely, for  $k = 2$  the solutions of the job-shop problem is coded by an infinite word which is either a periodic balanced word or an aperiodic balanced word [14, 19]. The situation gets more complicated for  $k > 2$ , which leads to the famous Fraenkel's conjecture [12, 13]. It is restated in combinatorics on words terms: An infinite word on a  $k$ -letters alphabet balanced on each letter of the alphabet and with all letters frequencies pairwise distinct is given by an infinite periodic word constructed on a unique period word  $FR_k$  (up to a permutation of letter and

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circular permutation) by the recursive formula  $FR_k = FR_{k-1}kFR_{k-1}$  with  $FR_3 = 1213121$ . Many researchers have worked on the general problem of infinite balanced words on an alphabet with one letter [16] or equivalently to cover integers by Beatty sequences [13, 27]. The conjecture is proved for  $k = 3$  by Morikawa [20], for  $k = 4$  by Altman, Gaujal and Hordijk [27], for  $k = 5$  and  $k = 6$  by Tijdeman [29] and for  $k = 7$  by Barat et Varju [15] and the conjecture is still open for  $k > 7$ . Indeed, in order to investigate new discrete tools that allow us to deeply understand the structure of balanced words, we propose a second order balance property for Christoffel words that gives a refinement for the balance property. In fact, we define for Christoffel words a new matrix called *second order balance matrix* which gives information on the balance property of a word that codes the number of 1's of successive blocks of same length in the studied Christoffel word. Thus we investigate balance property on successive blocks instead of balance property on letters for second order balance property. The main idea, to go further in the resolution of the Fraenkel's conjecture, is to consider synchronization of the blocks instead of synchronization of letters and this is why we introduce the notion of second order of balance.

In Section 2, we recall some properties of the Christoffel words. In Section 3, we define the balance matrix which gives information on the number of occurrences of a given letter in all factors of a given binary circular word. This balance matrix gives us the order of balance, for binary words. Afterwards, in Section 4, we introduce the *second order balance matrix* of a given Christoffel word by computing the balance matrix for the rows of the associated balance matrix. We show in Section 5 that this matrix has many symmetries and is constructed by using properties of continued fractions and the Stern-Brocot tree. We present in Section 6 a recursive construction for the second order balance matrix by considering the properties of the continued fraction expansion for the slope of each Christoffel word. In Section 7, we show that an infinite path on the Stern-Brocot tree, which minimizes the second order of balance, is given by a path associated with the Fibonacci word. Section 8 is left for the perspectives of this work.

## 2 Notation and Christoffel words

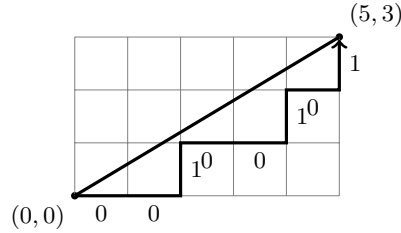
Let  $A$  be an alphabet of cardinality  $m$ , the word  $w$  is the concatenation of letters of this alphabet and we write  $w \in A^*$ , where  $A^*$  represents the set of all the words formed by the alphabet  $A$ . We denote by  $n = |w|$  the length of the word and by  $|w|_a$  the number of occurrences of the letter  $a$  in the word  $w$ . The notation  $w[i \dots j]$  refers to the factor of the word  $w$  from position  $i$  to position  $j$ . The notation  $w^\omega$  represents:  $w^\omega = ww \dots w \dots$  and named "circular word associated with  $w$ ". By convention,  $w^0 = \epsilon$  and a word  $w$  is said primitive if it is not the power of a nonempty word. Two words  $w$  and  $w'$  are *conjugate* of order  $k$  if and only if there exist  $u, v$  such that  $|u| = k$  with  $w = uv$  and  $w' = vu$  and we denote:  $w \equiv_k w'$ . When the exact value of  $k$  is not relevant, we simply write  $w \equiv w'$  and we say that the two words are conjugate. A positive integer  $p$  is a *period* of  $w$  if  $w[i] = w[i + p]$ ; for all  $1 \leq i \leq |w| - p$ . Given a word  $w = aw'$  where  $a$  is a letter, we note  $a^{-1}w = w'$  that is the removal of the letter  $a$  at the beginning of  $w$ . If  $w$  ends with letter  $a$ , then the notation  $wa^{-1}$  is defined accordingly. Let  $A = \{a_0, a_1, \dots, a_{m-1}\}$  be an alphabet, we let  $\bar{\cdot}$  be the anti-morphism such that:  $\bar{a}_0 = a_{m-1}, \bar{a}_1 = a_{m-2} \dots \bar{a}_i = a_{m-1-i}$ . A word  $w \in \{0, 1\}^*$  is  $k$ -balanced if and only if for all factors  $u, v$  of  $w$ , we get:  $|u| = |v| \implies ||u|_1 - |v|_1| \leq k$ . The word  $w$  is called balanced if  $k = 1$ .

Christoffel words [9] have many equivalent definitions and characterizations. The following geometrical definition is taken from [7] (see [5] for a self-contained survey). The lower Christoffel path of slope  $\frac{a}{b}$ , where  $a$  and  $b$  are relatively prime, is the path from  $(0, 0)$  to  $(b, a)$  in the integer lattice  $\mathbb{Z} \times \mathbb{Z}$  that satisfies the following conditions:

1. The path lies below the line segment that begins at the origin and ends at  $(b, a)$ .
2. The region enclosed by the path and the line segment contains no other points of  $\mathbb{Z} \times \mathbb{Z}$  besides those of the path.

We encode the lower Christoffel path (or simply Christoffel word) by means of a word in the alphabet  $A$  using 0 (resp. 1) for any unit horizontal (vertical) step. We get the *Christoffel word* of slope  $\frac{a}{b}$  denoted:  $C(\frac{a}{b})$ , see Figure 1. Equivalently, the Christoffel word  $w = C(\frac{a}{b})$  is obtained by calculating the elements of the sequence  $(r_i)_{0 \leq i \leq n}$ , where  $n = a + b$  as follows:  $r_i = ia \pmod n$ . Each letter,  $w[i]$ ,  $\forall 1 \leq i \leq n$ , of the word  $w$  and length  $n$  is obtained by computing:  $w[i] = \begin{cases} 0 & \text{if } r_{i-1} < r_i, \\ 1 & \text{otherwise.} \end{cases}$

*Example 1.* Let  $(a, b) = (3, 5)$ , the sequence  $(r_i)_{0 \leq i \leq 8} = (0, 3, 6, 1, 4, 7, 2, 5, 0)$  defines the Christoffel word  $C(\frac{3}{5}) = 00100101$ .



**Fig. 1:** Illustration of the geometrical definition of Christoffel words. The Christoffel path goes from  $(0,0)$  to  $(5,3)$  and  $C(\frac{3}{5}) = 00100101$ .

### 3 Balance matrix

In this section, we introduce a new matrix used to obtain the order of balance for any binary word in an explicit way. The  $i^{th}$  row of the matrix  $M$ ,  $M[i]$ , is seen as a word where each entry of the matrix is a letter. Given a word  $w \in A^*$  of length  $n$ , we let  $S_w$  be the  $n \times n$  matrix defined by  $S_w[i, j] = w[j] + \dots + w[i + j]$  over the circular word  $w$ . By definition, we have that  $w$  is  $\delta$ -balanced if

$$\delta = \max_i (\max(S_w[i]) - \min(S_w[i])).$$

The *balance matrix*  $B_w$ , is defined from  $S_w$  by subtracting the minimum value on each row,

$$B_w[i, j] = S_w[i, j] - \min(S_w[i]).$$

Obviously, we have that  $w$  is  $(\max B_w)$ -balanced.

It is clear that by construction, if  $w$  is not a sequence of 1's, then the first row of  $B_w$  is equal to  $w$ . If  $k < |w|$  is a period of  $w^\omega$ , again by construction we have that the row  $B_w[k]$  contains only zeros. The converse is also true, an integer  $k$  such that the row  $B_w[k]$  contains only zeros is a period of  $w^\omega$ . Consequently, a row of zeros is called a *period row*. Note that  $n = |w|$  is always a period of

$w^\omega$  and therefore we define the matrix  $B_w$  as a  $(n - 1) \times n$  matrix since the  $n$ -th row would not add any information.

For the sake of simplicity, when working with the Christoffel word  $C\left(\frac{a}{b}\right)$ , the balance matrix  $B_{C\left(\frac{a}{b}\right)}$  is simply denoted  $B_{\frac{a}{b}}$ .

*Example 2.* Let us consider the rational number  $\frac{5}{2}$  and  $n = 7$ . By writing the Christoffel word  $w = C\left(\frac{5}{2}\right) = 0110111$ , the elements of the balance matrix  $B_w$  and  $S_w$  are the following:

$$S_w = \begin{pmatrix} 0 & 1 & 1 & 0 & 1 & 1 & 1 \\ 1 & 2 & 1 & 1 & 2 & 2 & 1 \\ 2 & 2 & 2 & 2 & 3 & 2 & 2 \\ 2 & 3 & 3 & 3 & 3 & 3 & 3 \\ 3 & 4 & 4 & 3 & 4 & 4 & 3 \\ 4 & 5 & 4 & 4 & 5 & 4 & 4 \\ 5 & 5 & 5 & 5 & 5 & 5 & 5 \end{pmatrix}, \quad B_{\frac{5}{2}} = \begin{pmatrix} 0 & 1 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}.$$

### 3.1 Properties of the balance matrix

In this section, we present some properties of the matrix  $B_w$ , where  $w$  is a binary word of length  $|w| = n$ , allowing us to restrict the work to the upper half of the balance matrix since the lower part will be obtained by symmetry.

**Definition 1.** Two words  $w$  and  $w'$  are complementary if  $w \equiv_k \overline{w'}$  for some  $k$ . Moreover, a word is said to be autocomplementary if it is complementary to itself.

*Example 3.* The word  $w = 0110110010$  is an autocomplementary word since we have:  $w \equiv_5 \overline{w}$ . While  $w = 0000100$  and  $w' = 0111111$  are such that  $w \equiv_4 \overline{w'}$ .

We show some basic combinatorial properties of the matrix  $S_w$  that will be used in a further proof. All the proofs of this section can be found in the Appendix.

*Property 1.* For any binary word  $w$  of length  $n$ , and  $1 \leq i \leq n$ , the matrix  $S_w$  satisfies:

1.  $\max S_w[i] + \min S_w[n - i] = |w|_1$ ,
2.  $\min S_w[i] + \max S_w[n - i] = |w|_1$ .

In particular, if  $w$  is a 1-balanced Christoffel word, we have:

1.  $\min S_w[i] + \min S_w[n - i] = |w|_1 - 1$ ,
2.  $\max S_w[i] + \max S_w[n - i] = |w|_1 + 1$ ,

By applying the results of Property 1 and the definition of  $S_w$ , we can notice that the lower half of  $B_w$  is deduced from its upper half part, as we can see in Property 2.

*Property 2.* The balance matrix is such that for all  $1 \leq i < n$ ,  $B_w[i] \equiv_{n-i} \overline{B_w[n - i]}$ .

If  $B_{\frac{a}{b}}$  has a middle row, then this row is autocomplementary as we can see in Example 4.

**Corollary 1.** If  $n$  is an even number,  $B_{\frac{a}{b}}\left[\frac{n}{2}\right]$  is an autocomplementary row.

*Example 4.* Let us consider the rational number  $3/7$  with  $n = 10$ . The balance matrix  $B_{\frac{3}{7}}$  shows that  $B_{\frac{3}{7}}[5]$  is an autocomplementary row.

$$B_{\frac{3}{7}} = \begin{pmatrix} 0001001001 \\ 0011011011 \\ 0111111111 \\ 0001001000 \\ \mathbf{0011011001} \\ 0111111011 \\ 0001000000 \\ 0011001001 \\ 0111011011 \end{pmatrix}.$$

*Remark 1.* In the case where  $a = 1$ , the balance matrix is such that:  $B_{\frac{1}{b}}[i] = 0^{(n-i)}1^i$ .

### 3.2 Construction of the balance matrix for Christoffel words

In this section, we are interested in giving a direct construction for the balance matrix of a Christoffel word, by determining for each row of  $B_w$ , the positions of the letter 1. For that, we start by defining the set of positions for the letter 1 in  $C(\frac{a}{b})$ .

**Definition 2.** Let  $w$  be a Christoffel word of slope  $a/b$ . The set of decreasing positions of  $w$ , denoted  $D(a, b)$ , is the set of the positions of the occurrences of the letter 1 in  $w$ . More formally,  $D(a, b) = \{1 \leq i \leq n \mid w[i] = 1\}$ .

The following theorem is a reformulation of Paquin and Reutenauer's result ([22], Corollary 3.2), that describes the set of decreasing positions of a Christoffel word.

**Theorem 1.** Let  $\alpha$  be such that  $\alpha a \equiv -1 \pmod{n}$ , then the set  $D(a, b)$  is :

$$D(a, b) = \{(i\alpha \bmod n) + 1 \mid i = 1 \dots a\}.$$

*Example 5.* Let us consider the rational number  $a/b = 3/5$  with  $n = 8$ ,  $\alpha = 5$  and  $w = 00100101$ . The set of occurrences of the letter 1 is  $D(3, 5) = \{(5i \bmod 8) + 1 \mid i = 1 \dots 3\} = \{3, 6, 8\}$ .

Using Theorem 1, we can conclude that  $B_{\frac{a}{b}}$  and  $B_{\frac{b}{a}}$  are not equal but complementary as Lemma 1 shows. The proof can be found in the Appendix.

**Lemma 1.** The balance matrices  $B_{\frac{a}{b}}$  and  $B_{\frac{b}{a}}$  are conjugate in the sense that:

$$B_{\frac{a}{b}}[i] \equiv_{\alpha} \overline{B_{\frac{b}{a}}[i]} \quad \forall 1 \leq i < n; \text{ where } \alpha a \equiv -1 \pmod{n}.$$

For this part,  $w$  is a Christoffel word and we give a recursive construction of  $B_w$  by identifying for each row  $i$ , the set of positions of the 1's. This set is denoted by  $D_i$ , where for each row  $i$  in  $B_w$ ,  $D_i = \{j \mid B_w[i, j] = 1\}$ , with  $n = |w|$  and  $1 \leq j \leq n$ . For any set of integers  $S$  and any integer  $k$ , we denote  $S + k = \{a + k \mid a \in S\}$ .

**Theorem 2.** If  $a < b$  then the sets  $(D_i)_{1 \leq i \leq n-1}$  are recursively obtained as follows:  $D_1 = D(a, b)$  and for each  $i$  from 2 to  $n - 1$ :

$$D_i = \begin{cases} D_{i-1} \cup (D_1 - (i-1) \bmod n) & \text{if } i \notin D_1 \\ D_{i-1} \cap (D_1 - (i-1) \bmod n) & \text{if } i \in D_1 \end{cases} \quad (1)$$

The proof can be found in the Appendix.

*Example 6.* Let us consider the rational number  $2/7 \bmod 9$ . By calculating the set of decreasing values for each row of  $B_{\frac{2}{7}}$ , we get the following matrix:

$$\begin{aligned}
D_1 &= \{5, 9\} \\
D_2 &= D_1 \cup \{4, 8\} = \{4, 5, 8, 9\} \\
D_3 &= D_2 \cup \{3, 7\} = \{3, 4, 5, 7, 8, 9\} \\
D_4 &= D_3 \cup \{2, 6\} = \{2, 3, 4, 5, 6, 7, 8, 9\} \\
D_5 &= D_4 \cap \{1, 5\} = \{5\} \\
D_6 &= D_5 \cup \{4, 9\} = \{4, 5, 9\} \\
D_7 &= D_6 \cup \{3, 8\} = \{3, 4, 5, 8, 9\} \\
D_8 &= D_7 \cup \{2, 7\} = \{2, 3, 4, 5, 7, 8, 9\} \\
D_9 &= D_8 \cap \{1, 6\} = \{ \}.
\end{aligned}
\Rightarrow B_{\frac{2}{7}} = \begin{pmatrix} 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 \end{pmatrix}.$$

From this construction, we can get a relation between the number of occurrences of 1 in each row of the balance matrix of Christoffel words and the numerator of the slope related to this word.

**Lemma 2.** *For the balance matrix  $B_{\frac{a}{b}}$ , we have:  $|D_i| = i.a \bmod n$ , where  $|D_i|$  is the cardinal of the set  $D_i$ .*

**Note:** The proof can be found in the Appendix. This Lemma confirms that the period row is made only of zeros since  $|D_n| = n.a \bmod n = 0$ , hence we have no occurrences for the letter 1 in this row.

## 4 Second order balance matrix

Let  $w^\omega$  be a 1-balanced circular word associated with  $w$ ; by computing the balanced property on each row of  $B_w$ , we get a refinement of the balanced property for  $w$ . For any factor  $v$  of length  $j$ , we can find  $k_j$  or  $k_j + 1$  occurrences of the letter 1. The second order balance is the repartition of these blocks in a binary balanced word. This second order balance is computed via a matrix called the *second order balance matrix*. In other words, we are studying the balance of each row of  $B_w$ . For a pair of integers  $i, j$ , where  $1 \leq i, j \leq |w| - 1$ , we consider the word  $B_w[i]$  and we list all its factors of length  $j$ . Among these factors, we choose  $p$ , a factor that maximizes the number of occurrences of the letter 1 and  $q$ , a factor that minimizes it. The entry  $U_w[i, j]$  is given by  $|p|_1 - |q|_1$ . Equivalently, if  $L_\ell(w)$  is the restriction of the language of  $w^\omega$  to words of length  $\ell$ , then:

$$U_w[i, j] = \max_{v \in L_j(B_w[i])} |v|_1 - \min_{v \in L_j(B_w[i])} |v|_1.$$

In other words, we can define the second order balance matrix  $U_w$  by:

**Definition 3.** *Let  $w$  be a word such that  $w^\omega$  is 1-balanced or, equivalently, that  $B_w$  is a binary matrix. The second order balance matrix  $U_w = (u_{ij})_{1 \leq i, j \leq n-1}$  where  $U_w[i, j] = \max(B_{(B_w[i])}[j])$ .*

**Definition 4.** *The second order of balance of a circular 1-balanced word  $w$  is  $\delta^2(w) = \max(U_w)$ .*

Once again, in order to lighten the notation, when working with the Christoffel word  $C(\frac{a}{b})$ , the second order balance matrix  $U_{C(\frac{a}{b})}$  is simply denoted  $U_{\frac{a}{b}}$ .

For the rest of the paper, we let  $w$  be a Christoffel word of slope  $a/b$  and length  $n = a+b$ , such that:  $w = C(\frac{a}{b})$ . The second order balance matrix of a Christoffel word of slope  $\frac{a}{b}$ ,  $U_{\frac{a}{b}}$ , is of dimension  $(n-1) \times (n-1)$ .

*Example 7.* Let us consider the rational number  $\frac{a}{b} = \frac{3}{7}$  with  $n = 10$ . The balance matrix  $B_{\frac{3}{7}}$  was calculated previously and  $B_{\frac{3}{7}}[5] = [0, 0, 1, 1, 0, 1, 1, 0, 0, 1]$ . By computing the balance matrix for this word and taking the blocks of length 5 we get the 5<sup>th</sup> row of  $S_{B_{\frac{3}{7}}}[5]$  where the difference between the maximum and the minimum values of each row of  $S_{B_{\frac{3}{7}}}[5]$  determines the entries of  $U_{\frac{3}{7}}[5]$ . Hence, with these two blocks  $\overline{00|11011|001}$ , we obtain the element  $U_{\frac{3}{7}}[5, 5] = 4 - 1 = 3$  as we can see in the following second order balance matrix of  $\frac{3}{7}$ , where we also get  $\delta^2(C(\frac{3}{7})) = 3$ .

$$U_{\frac{3}{7}} = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 1 & 1 & 2 & 1 & 1 & 2 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 2 & 2 & 2 & 1 & 1 & 1 & 1 \\ 1 & 2 & 1 & 2 & \mathbf{3} & 2 & 1 & 2 & 1 & 1 \\ 1 & 1 & 1 & 2 & 2 & 2 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 1 & 1 & 2 & 1 & 1 & 2 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix}.$$

### Properties of the matrix $U_{\frac{a}{b}}$

Now we give some properties of the second order balance matrix in order to show that  $U_{\frac{a}{b}} = U_{\frac{b}{a}}$ . Hence, we can restrict our study to the irreducible fractions  $a/b$  with  $a < b$ . But before that, we prove the three symmetries that appear in this matrix. From Section 3.1, we have that the rows of the upper half of  $B_{\frac{a}{b}}$  are complementary to the rows of its lower half, which induces the symmetries in the matrix  $U_{\frac{a}{b}}$ . More precisely, the second order balance matrix  $U_{\frac{a}{b}}$  of dimension  $(n-1) \times (n-1)$ , has horizontal, vertical and diagonal symmetries. The axis of symmetry are at position  $\frac{n}{2}$  or between  $\frac{n-1}{2}$  and  $\frac{n+1}{2}$  depending on the parity of  $n$ .

*Property 3.* For any position  $(i, j)$ ,  $U_{\frac{a}{b}}[i, j] = U_{\frac{a}{b}}[n-i, j] = U_{\frac{a}{b}}[i, n-j] = U_{\frac{a}{b}}[n-i, n-j]$ .

*Proof.* By Lemma 2, we have that for any  $1 \leq i \leq n-1$ ,  $B_w[i] \equiv_{n-i} \overline{B_w[n-i]}$  where  $n = a + b$ . Consequently :

$$\begin{aligned} U_w[i, j] &= \max(B_{(B_{\frac{a}{b}}[i])}[j]) && \text{(def.)} \\ &= \max(B_{(B_{\frac{a}{b}}[n-i])}[j]) && \text{(autocomplementarity of } B_{\frac{a}{b}}) \\ &= \max(B_{(B_{\frac{a}{b}}[n-i])}[n-j]) && \text{(autocomplementarity of } B_{B_{\frac{a}{b}}[n-i]}) \\ &= \max(B_{(B_{\frac{a}{b}}[i])}[n-j]). \end{aligned}$$

Moreover,  $U_{\frac{a}{b}}$  has an extra diagonal symmetry;

*Property 4.* For the Christoffel word of slope  $a/b$ , we have  $(U_{\frac{a}{b}})^T = U_{\frac{a}{b}}$ .

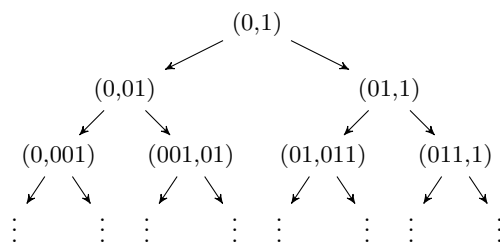
*Proof.* Let  $w = C(\frac{a}{b})$  and  $i < j$ , by definition:  $U_w[i, j] = \max(B_{(B_w[i])}[j])$ ; we recall that:  $B_w[i] = S_w[i] - \min(S_w[i])$  where  $S_w[i, j] = w[i] + \dots + w[i+j] = |w[i, j]|_1$ . Therefore  $B_{B_w[i]}[j] = S_{B_w[i]}[j] - \min(S_{B_w[i]}[j])$  with  $S_{B_w[i]}[j] = |w[j, \dots, j+i-1](j, i)|_1 = |w[j, \dots, j+i-1]|_1$ . While for  $U[j, i] = \max(B_{(B_w[j])}[i])$  we have:  $B_w[j] = S_w[j] - \min(S_w[j])$ . Therefore  $\forall i, j \in \{1, \dots, n-1\}$ ,  $B_{B_w[j]}[i] = S_{B_w[j]}[i] - \min(S_{B_w[j]}[i])$  with  $S_{B_w[j]}[i] = |w[i, \dots, j+i-1][i, j]|_1 = |w[j, \dots, j+i-1]|_1$  and  $U_w[i, j] = \max(B_{(B_w[i])}[j]) = \max(B_{(B_w[j])}[i]) = U[j, i]$ .



After those two properties, we are able to prove that  $U_{\frac{a}{b}} = U_{\frac{b}{a}}$ .

## 5 More about Christoffel words

Let  $w$  be a Christoffel word of length at least 2, the standard factorization is obtained by writing  $w = (w_1, w_2)$  in a unique way, where  $w_1, w_2$  are two Christoffel words by [7]. The *Christoffel tree* is an infinite tree whose vertices are all the standard factorizations of Christoffel words (see [5], Section 3.2). It uses the fact that given a standard factorization  $(w_1, w_2)$ , the pairs  $(w_1, w_1w_2)$  and  $(w_1w_2, w_2)$  are also standard factorizations. Let  $\phi_0, \phi_1$  be the two functions from  $A^* \times A^*$  into itself defined by:  $\phi_0(w_1, w_2) = (w_1, w_1w_2)$ ;  $\phi_1(w_1, w_2) = (w_1w_2, w_2)$ . We have that any Christoffel word can be obtained in a unique way by iteration of these two functions on  $(0, 1)$ . Consequently, the Christoffel tree is defined as follows: the root is  $(0, 1)$ , the Christoffel word of slope 1. Then each node  $(w_1, w_2)$  has two sons:  $\phi_0(w_1, w_2)$  on the left and  $\phi_1(w_1, w_2)$  on the right. See Figure 2 for an illustration.



**Fig. 2:** The first levels of the Christoffel tree.

We introduce and study the directive sequence of Christoffel words.

**Definition 5.** Let  $w = C(\frac{a}{b})$  be a non-trivial Christoffel word, the directive sequence of  $w$ , denoted  $\Delta(\frac{a}{b})$ , is the word  $\Delta(\frac{a}{b}) = i_1i_2 \cdots i_n \in A^*$  such that  $w = (\phi_{i_n} \circ \cdots \circ \phi_{i_2} \circ \phi_{i_1})(0, 1)$ .

Note that a directive sequence  $\Delta(\frac{a}{b}) = i_1 \cdots i_n$  describes the path from the root of the Christoffel tree to the Christoffel word  $C(\frac{a}{b})$  as follows : at step  $k$ , if  $i_k = 0$  then go left, otherwise, if  $i_k = 1$  then go right.

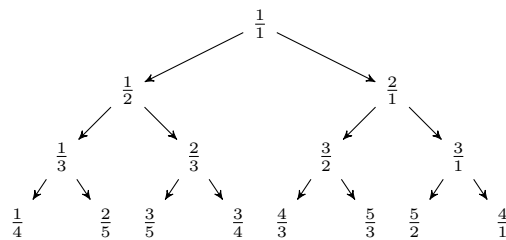
### 5.1 Stern-Brocot tree and continued fractions

In this section, we introduce the *Stern-Brocot tree* that contains all the reduced fractions  $\frac{a}{b}$ . It was first introduced by a German mathematician Moritz Abraham Stern and a French clockmaker Achille Brocot in the 19th century [8]. In order to construct recursively the Stern-Brocot tree, we need to introduce the *mediant* of two fractions  $\frac{a}{b}$  and  $\frac{c}{d}$ , that is  $\frac{a}{b} \oplus \frac{c}{d} = \frac{a+c}{b+d}$ . In addition to that, we have to define the recursive sequence  $s_i$ , that is obtained from  $s_{i-1}$ , by completing with the mediant of each two consecutive fractions in  $s_{i-1}$ , where  $s_0$  is given by:  $s_0 = (\frac{0}{1}, \frac{1}{0})$ .

Note that,  $\frac{1}{0}$ , is considered as a normal fraction:  $s_0 = (\frac{0}{1}, \frac{1}{0})$ ,  $s_1 = (\frac{0}{1}, \frac{1}{1}, \frac{1}{0})$ ,  $s_2 = (\frac{0}{1}, \frac{1}{2}, \frac{1}{1}, \frac{2}{1}, \frac{1}{0})$ ,  $s_3 = (\frac{0}{1}, \frac{1}{3}, \frac{1}{2}, \frac{2}{3}, \frac{1}{1}, \frac{3}{2}, \frac{2}{1}, \frac{3}{1}, \frac{1}{0})$ .

The mediant added in each new step to the sequence  $s_i$  are the fractions that appear on the  $i^{th}$  level of the Stern-Brocot tree. For example, on the third level, we have the fractions:  $\frac{1}{3}, \frac{2}{3}, \frac{3}{2}, \frac{3}{1}$  that are extracted from the sequence  $s_3$ . We call consecutive fractions, two fractions that belong to the same set  $s_i$  and are next to each other, like  $\frac{1}{1}$ , and  $\frac{2}{1}$  in  $s_2$ . Some properties about these consecutive fractions will be given in the next section.

In order to simplify the notation, we denote  $\frac{a}{b} \oplus \frac{c}{d}$  by  $\frac{a+c}{b+d}$  and  $\frac{a}{b} \oplus \dots \oplus \frac{a}{b}$  repeated  $p$  times by  $(\frac{a}{b})^p$ . The Christoffel tree is isomorphic to the Stern-Brocot tree where each vertex of the Christoffel tree of the form  $(u,v)$  is associated to the fraction  $\frac{|uv|_1}{|uv|_0}$ , see Figure 3.



**Fig. 3:** The first levels of the Stern-Brocot tree.

The continued fraction of a rational number  $\frac{a}{b} \geq 0$  is the sequence of integers  $\frac{a}{b} = [a_0, \dots, a_z]$ , with  $a_0 \geq 0$ ;  $a_i \geq 1$  for  $1 \leq i \leq z$  and if  $z \geq 2$  then  $a_z \geq 2$ .

$$\frac{a}{b} = a_0 + \frac{1}{a_1 + \frac{1}{\dots + a_z}}$$

In the following part, we will explain how to pass from the continued fraction of a rational number  $a/b$  to the Christoffel word of slope  $a/b$ . For that, we introduce the following theorem by Henry J.S in 1876. This theorem gives an additional characterization for the Christoffel words. In fact, he showed that the Christoffel word can be obtained by a recursive expression using the partial denominators of the rational number.

**Theorem 3.** [28] Let  $\frac{a}{b} = [a_0, a_1, \dots, a_z]$ , a word  $w = 0u1$  is a Christoffel word of slope  $\frac{a}{b}$  if and only if  $u01$  or  $u10$  is equal to  $s_{n+1}$ , where  $s_{n+1}$  is defined recursively by:  $s_{-1} = 0, s_0 = 1$  and  $s_{n+1} = s_n^{a_n} s_{n-1}$  for all  $n \geq 0$ .

Moreover, in 1987, Berstel introduced, in a report for LITP, the following relation between the partial denominators of a rational number and two particular matrices based on the Theorem 2.1 in George Raney's paper [25] in 1973, that was published in 1990 [3]. Using the two matrices  $A = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$  and  $B = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$ , and for the particular case of the rational number  $a/b$  where  $a < b$ , we have:

**Proposition 1.** [3] For  $a < b$ , we associate for the rational number  $\frac{a}{b} = [0, a_1, a_2, \dots, a_n]$ , the following matrix:  $M(\frac{a}{b}) = A^{a_1-1} B^{a_2} A^{a_3} \dots C^{a_{n-1}} D^{a_n-1}$ , where  $C$  and  $D$  are respectively  $A, B$  or  $B, A$  depending on the parity of  $n$ . We have:  $M(\frac{a}{b}) \cdot \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \begin{pmatrix} b-a \\ a \end{pmatrix}$ .

Furthermore, in 1993, Borel and Laubie, followed by de Luca in 1997 then Berthé, de Luca and Reutenauer in 2008, gave the following theorem, where they linked these two particular matrices to the Christoffel word of slope  $a/b$ . Let the function  $Pal$  be the iterative palindromic closure function defined recursively in [4] by  $Pal(ua) = (Pal(ua))^{(+)}$  and  $Pal(\epsilon) = \epsilon$ , where  $a$  is a letter and  $u$  a word and  $(w)^+$  is the palindromic closure of  $w$  i.e.  $(w)^+ = ww'$  with  $ww'$  is the shortest palindrome having  $w$  as a prefix.

**Theorem 4.** [4, 6, 7] Let  $w$  be a Christoffel word of slope  $a/b$ , there exists a unique word  $v$  such that  $w = 0Pal(v)1$ , if  $(w_1, w_2)$  is the standard factorization of  $w$  then we define the multiplicative monoid morphism  $\mu : \{0, 1\}^* \rightarrow SL_2(\mathbb{Z})$  such that:  $\mu(0) = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} = A$  and  $\mu(1) = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} = B$ ,

where:  $\mu(v) = \begin{pmatrix} |w_1|_0 & |w_2|_0 \\ |w_1|_1 & |w_2|_1 \end{pmatrix}$ .

In [3, 5, 18, 25] we can find results allowing us to write, in an explicit way, the relation between the continued fraction of a rational number and its directive sequence as we can see in the following theorem (see[4]), where its proof can be found in the Appendix.

**Theorem 5.** Let  $\frac{a}{b} = [a_0, \dots, a_z]$ , we have:  $\Delta(\frac{a}{b}) = 1^{a_0} 0^{a_1} 1^{a_2} \dots p^{a_z-1}$  where  $p \in \{0, 1\}$ . The Christoffel word of slope  $a/b$  is written:  $C(\frac{a}{b}) = 0w'1$ , where  $w'$  is a palindrome and  $w' = Pal(\Delta(\frac{a}{b}))$ .

## 6 Recursive construction of the second order balance matrix

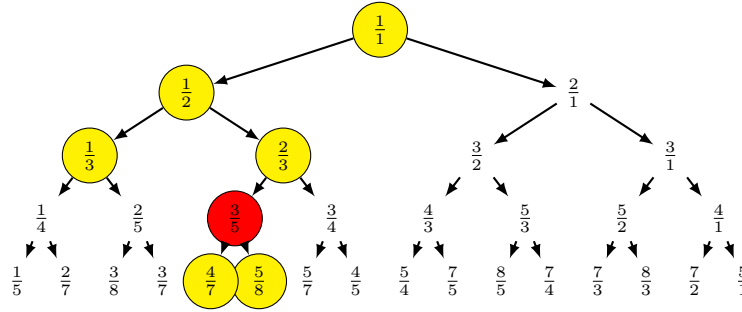
Due to the isomorphism and the recursive construction of the Stern-Brocot tree and the Christoffel tree, we can conclude that there must exist a recursive construction for the second order balance matrix. In order to show and to prove this recursivity, we let  $U_{\frac{a}{b}}$  where  $a/b = [a_0, \dots, a_z] = [a_0, a_1, \dots, a_z-1, 1]$  be the second order balance matrix of the Christoffel word  $C(\frac{a}{b})$ . We introduce the following terminology for some specific rational numbers on the Stern-Brocot tree and that will be used for the rest of the paper. See Figure 4 for an illustration.

**Definition 6.** Given  $\frac{a}{b} = [a_0, \dots, a_z]$ ,

1. the top branch fraction of  $\frac{a}{b}$ , denoted  $TBF(\frac{a}{b})$ , is the fraction  $[a_0, \dots, a_{z-1} + 1]$ ,
2. the first reduced fraction of  $\frac{a}{b}$ , denoted  $FRF(\frac{a}{b})$ , is the fraction  $[a_0, \dots, a_{z-1}]$ ,
3. the first extended fraction of  $\frac{a}{b}$ , denoted  $FEF(\frac{a}{b})$ , is the fraction  $[a_0, a_1, \dots, a_z + 1]$ ,
4. the first deviation fraction of  $\frac{a}{b}$ , denoted  $FDf(\frac{a}{b})$ , is the fraction  $[a_0, a_1, \dots, a_z - 1, 2]$ ,
5. the first parallel fraction of  $\frac{a}{b}$ , denoted  $FPF(\frac{a}{b})$ , is either  $[a_0, a_1, \dots, a_{z-1} - 1, 2]$  if  $a_{z-1} \neq 1$  or  $[a_0, a_1, \dots, a_{z-2} + 2]$  if  $a_{z-1} = 1$ .
6. the second unidirectional father of  $\frac{a}{b}$ , denoted  $SUF(\frac{a}{b})$ , is either  $[a_0, a_1, \dots, a_z - 2]$  if  $a_z > 2$  or  $[a_0, a_1, \dots, a_{z-2}]$  if  $a_z = 2$ .

Note that  $\text{SUF}(\frac{a}{b})$  is not defined for fractions  $\frac{1}{1}$ ,  $\frac{1}{2}$  and  $\frac{2}{1}$ . Using Theorem 5 we can get the directive sequence of each of these fractions.

*Example 8.* Let  $\frac{a}{b} = \frac{3}{5} = [0, 1, 1, 2]$ , from Definition 6, we get:  $\text{TBF}(3/5) = 2/3$ ,  $\text{FEF}(3/5) = 4/7$ ,  $\text{FDF}(3/5) = 5/8$ ,  $\text{FPF}(3/5) = 1/3$  and  $\text{SUF}(3/5) = 1/1$ . See Figure 4 for the positions of these fractions in the Stern-Brocot tree.



**Fig. 4:** Illustration of the specific rational numbers related to  $3/5 = [0, 1, 1, 2]$  on the Stern-Brocot tree. We have:  $\text{TBF}(3/5) = 2/3$ ,  $\text{FRF}(3/5) = 1/2$ ,  $\text{FEF}(3/5) = 4/7$ ,  $\text{FDF}(3/5) = 5/8$ ,  $\text{FPF}(3/5) = 1/3$  and  $\text{SUF}(3/5) = 1/1$ .

### General form of the second order balance matrix

To construct  $U_{\frac{a}{b}}$ , we start by placing 4 rows of separation that divide the matrix into 9 blocks. Due to the symmetries proved in Properties 3 and 4, it is sufficient to know three of these blocks to deduce the others. These blocks are denoted  $\alpha, \beta$  and  $\gamma$  and are represented in the matrix as follows :

$$U_{\frac{a}{b}} = \begin{pmatrix} \alpha & \cdot & \cdot \\ \gamma & \beta & \cdot \\ \cdot & \cdot & \cdot \end{pmatrix},$$

In the following section, we show that the blocks  $\alpha, \beta, \gamma$  are described by the second order balance matrices of simpler fractions. More precisely,  $\alpha$  is deduced from  $U_{\text{TBF}(\frac{a}{b})}$  while  $\beta$  is given by adding one to each entry of  $U_{\text{SUF}(\frac{a}{b})}$ . Finally, the construction of  $\gamma$  depends on the position of  $\frac{a}{b}$  in the Stern-Brocot tree, where the fraction is, relatively to its father, either a deviation (first deviation fraction) or an extension (first extended fraction).

#### 6.1 The construction of $U_{\frac{a}{b}}$

In this section, we start to explain how we can obtain the recursive construction of the second order balance matrix of the Christoffel word of slope  $a/b = [a_0, \dots, a_z]$ . Figure 6 displays the rational numbers needed for the construction. Since  $U_{\frac{a}{b}} = U_{\frac{b}{a}}$ , we reduce the work to the first half of the Stern-Brocot tree, that contains all the irreducible fractions with  $a \leq b$ . Given  $\frac{a}{b} = [a_0, \dots, a_z]$ , we consider separately the cases  $z < 2$  and  $z \geq 2$ .

**The trivial cases:  $z \in \{0, 1\}$**

For  $z = 0$ , we have  $U_{\frac{0}{1}} = U_{\frac{1}{0}} = [ ]$  and  $U_{\frac{1}{1}} = [1]$ . For  $z = 1$ , we have  $\frac{a}{b} = [0, a_1]$  which implies that  $a = 1$  and  $b = a_1$ . In this case, Remark 1, states that  $B_w[i] = 0^{b+1-i}1^i$  and from Definition 3, we have:  $U_w[i, j] = \max(B_{B_w[i]}[j])$ . Hence, we consider the first quarter of the matrix  $U_w$  which is sufficient from properties 3 and 4. Due to the diagonal symmetry, we will consider only  $j \leq i$  where  $i, j \leq \frac{n}{2}$  if  $n$  is even or  $i, j \leq \frac{n-1}{2}$  if  $n$  is odd.

**Proposition 2.** *If  $a = 1$ , then for all  $j \leq i$  and  $i, j \leq \frac{n}{2}$  we have:  $U_{\frac{a}{b}}[i, j] = j$ .*

*Proof.* Since  $B_w[i] = 0^{b+1-i}1^i$  and  $U_w[i, j] = \max(B_{B_w[i]}[j])$ , from Property 6 and with  $j < i$ , we have:  $B_{B_w[i]}[j] = 0^{b-i-j+2}1 \dots j^{i-j+2} \dots 1$ . Therefore,  $U_w[i, j] = \max(B_{B_w[i]}[j]) = j$ .

We can observe that this matrix can also be constructed using a recursive form where the first (resp. last) row and the first (resp. last) column are all 1's and in the middle we have the matrix of  $\text{SUF}(\frac{1}{b})$  where its elements are increased by 1 (see Figure 5).

$$U_{\frac{1}{b}} = \begin{pmatrix} \left( \begin{array}{c|cc} 1 & 1 & \dots & 1 & 1 \\ \hline 1 & & & & 1 \\ \cdot & & & & \cdot \\ \cdot & & & & \cdot \\ \cdot & & U_{[0, a_1-2]} + 1 & & \cdot \\ \cdot & & & & \cdot \\ \cdot & & & & \cdot \\ \hline 1 & 1 & \dots & 1 & 1 \\ \hline 1 & 1 & \dots & 1 & 1 \end{array} \right) \end{pmatrix}.$$

**Fig. 5:** The general form of the matrix  $U_{\frac{1}{b}}$ .

**The general case :  $z \geq 2$**

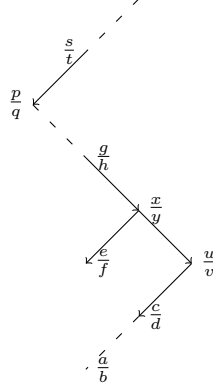
Now we assume that  $z \geq 2$ , in order to lighten the presentation we define the following fractions, let:  $\frac{u}{v} = \text{TBF}(\frac{a}{b})$ ,  $\frac{x}{y} = \text{FRF}(\frac{a}{b})$ ,  $\frac{c}{d} = \text{FDF}(\frac{u}{v})$ ,  $\frac{e}{f} = \text{FPF}(\frac{c}{d})$ ,  $\frac{g}{h} = \text{SUF}(\frac{u}{v})$ ,  $\frac{p}{q} = \text{TBF}(\frac{x}{y})$ ,  $\frac{s}{t} = \text{FRF}(\frac{u}{v})$ . See Figure 6 for an illustration of their relative positions on the Stern-Brocot tree.

**Lemma 3.** *The Christoffel words of slope  $\frac{a}{b}$ ,  $\frac{g}{h}$  and  $\frac{e}{f}$  can be written as follows:*

$$C(\frac{a}{b}) = \left(C(\frac{x}{y})\right)^{a_z-1} C(\frac{u}{v}) = \left(C(\frac{x}{y})\right)^{a_z} C(\frac{s}{t}), C(\frac{x}{y}) = C(\frac{g}{h})C(\frac{s}{t}) \text{ and } C(\frac{e}{f}) = C(\frac{g}{h})C(\frac{x}{y}).$$

*Proof.* This is a consequence of Theorem 5. The iteration of the application of functions  $\phi_0$  and  $\phi_1$  provides the equalities.

**Separation rows** For the rest of this section, the fractions mentioned in Figure 6 are used to prove the construction of  $U_{\frac{a}{b}}$ . In the following part we prove that separation rows allow the decomposition of  $U_{\frac{a}{b}}$  into 9 blocks. In Definition 7, we give a characteristic for the consecutive fractions of each sequence  $s_i$ , used to construct the Stern-Brocot tree.



**Fig. 6:** Position of the fractions  $\frac{u}{v}$ ,  $\frac{x}{y}$ ,  $\frac{e}{f}$ ,  $\frac{c}{d}$ ,  $\frac{g}{h}$ ,  $\frac{p}{q}$  and  $\frac{s}{t}$  relatively to  $\frac{a}{b}$ , for the case where  $a_{z-1} \geq 2$  and  $z$  odd. By definition, we have:  $\frac{a}{b} = [a_0 \dots, a_z]$ ,  $\frac{u}{v} = [a_0 \dots, a_{z-1} + 1]$ ,  $\frac{x}{y} = [a_0, \dots, a_{z-1}]$ ,  $\frac{p}{q} = [a_0, \dots, a_{z-2} + 1]$ ,  $\frac{s}{t} = [a_0, \dots, a_{z-2}]$ ,  $\frac{e}{f} = [a_0, \dots, a_{z-1} - 1, 2]$ ,  $\frac{c}{d} = [a_0, \dots, a_{z-1}, 2]$ ,  $\frac{g}{h} = [a_0, \dots, a_{z-1} - 1]$ . With respect to the Farey addition we have:  $\frac{u}{v} = \frac{x}{y} + \frac{s}{t}$  and  $\frac{c}{d} = \frac{x}{y} + \frac{u}{v}$ .

**Definition 7.** Let  $\frac{a}{b}$  and  $\frac{a'}{b'}$  be two consecutive fractions of a certain sequence  $s_i$ ,  $\forall i > 0$ . They respect the following relation:  $a'b - ab' = 1$ .

*Property 5.* Let  $\frac{a}{b}$  and  $\frac{a'}{b'}$  be two consecutive fractions in the Stern-Brocot tree, we have:

$$(a' + b').a = a'.(a + b) - 1.$$

The proof of this property is obtained by an arithmetic calculation based on Definition 7. In the following lemma, we prove that in each  $U_{\frac{a}{b}}$ , we have at least 2 rows and columns full of 1's.

**Lemma 4.** Let  $k = x + y$ , for the rational number  $\frac{a}{b}$  with  $n = a + b$ , we have:

$$U[i, k] = U[i, n - k] = U[k, j] = U[n - k, j] = 1, \forall i, j \in \{1, \dots, n - 1\}.$$

*Proof.* From Property 5, we have  $(x + y)a = x(a + b) - 1$ , using the consecutive fractions  $a/b$  and  $x/y$  from the Stern-Brocot tree, we get:  $ka = xn - 1$  hence  $ka \equiv -1 \pmod{n}$ . From Lemma 2, we have  $|D_i| = ia \pmod{n}$ , hence  $|D_k| = n - 1$  and  $|D_{n-k}| = 1$  since these two rows are complementary. Therefore, the words  $B_{\frac{a}{b}}[k]$  and  $B_{\frac{a}{b}}[n - k]$  are 1-balanced; hence  $U[k, j] = U[n - k, j] = 1, \forall j \in \{1, \dots, n - 1\}$ . Using the diagonal symmetry of  $U_{\frac{a}{b}}$ , we get:  $U[i, k] = U[i, n - k] = 1, \forall i \in \{1, \dots, n - 1\}$ .

*Remark 2.* From Lemma 4, we get that the row (respectively column)  $k$  and  $n - k$  contain only values of 1's. Therefore, we define the separation rows to be between the rows (resp. columns)  $(k; k + 1)$  and between the rows (resp. columns)  $(n - k - 1; n - k)$  which divide  $U_{\frac{a}{b}}$  into 9 blocks. See Figure 7.

Let  $\frac{a}{b} = \frac{5}{8}$ , the 9 blocks are represented as follows:

$$U_{\frac{5}{8}} = \begin{pmatrix} \begin{array}{cccc|cc|cc|cc} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 2 & 1 & 1 & 1 & 1 & 2 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 2 & 2 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 1 & 2 & 1 & 2 & 2 & 1 & 2 & 1 & 2 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 & 1 & 2 & 2 & 1 & 2 & 2 & 1 & 1 \\ 1 & 1 & 2 & 2 & 1 & 2 & 2 & 1 & 2 & 2 & 1 & 1 \\ \hline 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 1 & 2 & 1 & 2 & 2 & 1 & 2 & 1 & 2 & 1 \\ 1 & 1 & 1 & 1 & 1 & 2 & 2 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 2 & 1 & 1 & 1 & 1 & 2 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{array} \end{pmatrix}$$

Top left part of  $U_{\frac{3}{5}}$  ←
← Left central part of  $U_{\frac{3}{4}}$ 
↗  $U_{\frac{1}{2}} + 1$

**Fig. 7:** The decomposition of the matrix  $U_{\frac{5}{8}}$  into 9 blocks. We have  $\text{TBF}(\frac{5}{8}) = \frac{3}{5}$ ,  $\text{SUF}(\frac{5}{8}) = \frac{1}{2}$ ,  $\text{FPF}(\frac{5}{8}) = \frac{3}{4}$ ,  $\frac{x}{y} = \frac{2}{3}$  and  $k = 5$ .

Now we prove that each block  $\alpha, \beta$  and  $\gamma$  of  $U_{\frac{a}{b}}$  can be decomposed in some smaller second order balanced matrices. In fact, to construct the matrix  $U_{\frac{a}{b}}$  we first place the rows of separation. This shows that  $U_{\frac{a}{b}}$  is composed of 9 blocks where the three blocks  $\alpha_{k \times k}$ ,  $\beta_{(n-2k-1) \times (n-2k-1)}$  and  $\gamma_{(n-2k-1) \times k}$  are constructed while the others are obtained by symmetry. In the following part, we construct respectively each of the blocks  $\alpha, \gamma$  and  $\beta$  to get  $U_{\frac{a}{b}}$ .

**$\alpha$ -block:** Recall that  $z \geq 2$  and we let for the rest of this section  $z$  to be an odd number where the even case is obtained in a similar way. Let  $\frac{s}{t}$ ,  $\frac{g}{h}$  and  $\frac{x}{y}$  be the fractions defined at the beginning of Section 6.1. The  $\alpha$ -block is formed of  $k$  rows and columns where  $k = x + y = |\text{FRF}(\frac{a}{b})|$ .

**Lemma 5.** *The  $\alpha$ -block of  $U_{\frac{a}{b}}$  is exactly the first  $k - 1$  rows and columns of the matrix  $U_{\frac{u}{v}}$  where  $\frac{u}{v} = \text{TBF}(\frac{a}{b})$ .*

*Proof.* At the beginning, we show that the first  $k - 1$  rows of  $B_{\frac{a}{b}}$  are exactly the concatenation of  $a_z - 1$  times of the matrix  $B_{\frac{x}{y}}$  and the  $k - 1$  first rows of the matrix  $B_{\frac{u}{v}}$ . Then we show that the set of factors of these rows of length shorter than  $k$  is exactly the same for the rows of  $B_{\frac{u}{v}}$ .

By Lemma 3 and since  $z$  is odd, we have:  $C(\frac{a}{b}) = C(\frac{x}{y})^{a_z-1} C(\frac{u}{v})$  and  $C(\frac{u}{v}) = C(\frac{x}{y}) C(\frac{s}{t})$ . Therefore, the  $k$  first columns of  $B_{\frac{a}{b}}$  are given by  $B_{\frac{x}{y}}$ . Indeed, this part of  $B_{\frac{a}{b}}$  is constructed due to factors of lengths less or equal than  $k$  where only the first  $k$  positions are considered. This means that all these factors are included in the prefix of length  $2k$  of  $C(\frac{a}{b})$  which is exactly  $C(\frac{x}{y})^2$ . The same argument holds for each  $B_{\frac{a}{b}}[i, j]$ , for  $1 \leq i \leq k - 1$  and  $1 \leq j \leq n - (u + v)$ , which forms the  $a_z - 1$  blocks of  $B_{\frac{x}{y}}$ . Note that to obtain the last block of  $B_{\frac{x}{y}}$ , we use the fact that  $C(\frac{u}{v})$  has  $C(\frac{x}{y})$  as a prefix. For the last block, we observe that the entries of  $B_{\frac{a}{b}}[i, j]$  for  $1 \leq i \leq k - 1$  and  $n - (u + v) + 1 \leq j \leq n$ , are exactly the same as in  $B_{\frac{u}{v}}$  since, again, we only consider factors

up to length  $k$  and both words  $C\left(\frac{a}{b}\right)$  and  $C\left(\frac{u}{v}\right)$  start with  $C\left(\frac{x}{y}\right)$  of length  $k$ . Hence we have:  $B_{\frac{a}{b}}[i] = B_{\frac{x}{y}}[i]^{a_z-1} \cdot B_{\frac{u}{v}}[i]$ , for all  $1 \leq i \leq k-1$ . Finally, we can write that  $L_j(C\left(\frac{x}{y}\right)^\omega) \subseteq L_j(C\left(\frac{u}{v}\right)^\omega)$ , for  $j \leq k-1$  since  $C\left(\frac{x}{y}\right)$  ends with  $C\left(\frac{s}{t}\right)$  and  $C\left(\frac{u}{v}\right) = C\left(\frac{x}{y}\right)C\left(\frac{s}{t}\right)$ . Which implies

$$L_j\left(B_{\frac{a}{b}}^\omega[i]\right) = L_j\left(B_{\frac{u}{v}}^\omega[i]\right), \quad \forall 1 \leq i \leq k-1 \quad \text{and} \quad 1 \leq j \leq k-1.$$

**$\gamma$ -block:** The second block of  $U_{\frac{a}{b}}$  is the  $\gamma$ -block of dimension  $(n-2k-1) \times k$ . This block is located between the  $k+1$  and the  $(n-k-1)^{th}$  rows and bounded by the row of separation at the  $k^{th}$  column. This part of the second order balance matrix admits a recursive form. If we consider the concatenation of matrices as a vertical stacking, then  $\gamma\left(\frac{a}{b}\right)$  is given by

$$\gamma\left(\frac{a}{b}\right) = \begin{cases} \gamma\left(\frac{c}{d}\right) \left[\gamma\left(\frac{u}{v}\right) \gamma\left(\frac{c}{d}\right)\right]^{a_z-2} & \text{if } a_z \geq 3 \\ \gamma\left(\frac{e}{f}\right) & \text{if } a_z = 2 \end{cases} \quad (2)$$

We mentioned that we are considering the case where  $z \geq 2$  and with  $z$  an odd value. In fact, the parity of  $z$  determines if  $a/b$  is a right or left child. Furthermore, it helps us to know the value of  $\alpha$ , where  $\alpha$  is the value needed to obtain  $D(a, b)$ .

**Lemma 6.** Let  $\frac{a}{b} = [a_0, \dots, a_z]$ , and  $\text{FRF}\left(\frac{a}{b}\right) = \frac{x}{y}$  with  $x + y = k$ . The set  $D(a, b) = \{i\alpha + 1 \pmod n \mid 1 \leq i \leq a\}$  where  $n = a + b$  has:

$$\begin{cases} \alpha = k, & \text{if } z \text{ is even} \\ \alpha = n - k, & \text{if } z \text{ is odd.} \end{cases}$$

The proof can be found in the Appendix.

In order to prove the construction of the  $\gamma$ -block, we need to introduce some additional properties of the balance matrix  $B_{\frac{a}{b}}$  used in the technical proof of the  $\gamma$ -block.

*Property 6.* Let  $w = 0^x 1^y$  such that  $x, y > 0$ . For  $i \leq x$  and  $i \leq y$ , we have:  $B_w[i] = 0^{x-i+1} 12 \dots (i-2)(i-1) \cdot i^{y-i+1} \cdot (i-1)(i-2) \dots 21$ .

*Proof.* Let  $w = 0^x 1^y$ , for  $i = 2$ ,  $S_w[i, p] = |w^\omega[p, p+i-1]|_1$  for all  $1 \leq p \leq x+y$ , therefore  $S_w[2] = 0^{x-1} 12^{y-1} 1 = 0^{x-2+1} 12^{y-2+1} 1$  and since  $x-1 > 0$  then  $\min(S_w[2]) = 0$  and  $B_w[2] = S_w[2]$ . Suppose this is true for  $i \leq x$ , hence  $B_w[i] = 0^{x-i+1} 12 \dots i^{y-i+1} \dots 21$  prove it true for  $i+1$  such that  $i+1 \leq x$ . We have  $S_w[i+1, p] = 0$  for all  $1 \leq p \leq x-i$ , for  $p = x-i+1$ , we have  $S_w[i+1, p] = 1$  and for  $x-i+1 \leq p \leq y-x-i$ , we get  $S_w[i+1, p] = i+1$ . The values decrease until  $p = x+y$  and we have  $S_w[i+1, x+y] = 1$ . Since  $i+1 \leq x$ , then  $x-i \geq 1$  and  $\min S_w[i+1] = 0$ , therefore  $B_w[i] = 0^{x-i+1} 12 \dots i^{y-i+1} \dots 21$ .

**Lemma 7.** Let  $\frac{a}{b} = [a_0, \dots, a_z]$  and  $\frac{x}{y} = \text{FRF}\left(\frac{a}{b}\right)$  such that  $|C\left(\frac{x}{y}\right)| = x + y = k$ .

$$\begin{cases} B_{\frac{a}{b}}[k] = 0^\alpha 10^{n-\alpha-1}; & B_{\frac{a}{b}}[2k] = 0^{n-2k} 10^{k-1} 10^{k-1} \text{if } z \text{ is odd,} \\ B_{\frac{a}{b}}[n-k] = 0^\alpha 10^{n-\alpha-1}; & B_{\frac{a}{b}}[n-2k] = 0^{n-2k} 10^{k-1} 10^{k-1} \text{if } z \text{ is even.} \end{cases} \quad (3)$$



*Proof.* We prove the case where  $z$  is odd and the second case is obtained in a similar way. We let  $p_1, p_2$  be two palindromes such that:  $C(\frac{s}{t}) = 0p_11$ ,  $C(\frac{g}{h}) = 0p_21$  where  $C(\frac{x}{y}) = C(\frac{g}{h})C(\frac{s}{t})$ ;  $C(\frac{a}{b}) = (C(\frac{x}{y}))^{a_z-1}C(\frac{u}{v})$  and  $C(\frac{u}{v}) = C(\frac{x}{y})C(\frac{s}{t})$ . The central word of  $C(\frac{x}{y})$  is a palindrome then:  $p_210p_1 = p_101p_2$ . By the definition of the balance matrix, we have  $\forall 1 \leq i \leq n-1$  and  $\forall 1 \leq j \leq n$   $B_{\frac{a}{b}}[k, j] = S_{\frac{a}{b}}[k, j] - \min(S_{\frac{a}{b}}[k]) = |C(\frac{a}{b})[j \dots k + j - 1]|_1 - \min(S_{\frac{a}{b}}[k])$ . Summing up all this we get:

$$C(\frac{a}{b}) = C(\frac{x}{y})^{a_z-1} \cdot C(\frac{x}{y})C(\frac{s}{t}) = C(\frac{x}{y})^{a_z-1} \cdot 0p_210p_11 = C(\frac{x}{y})^{a_z-1} \cdot 0p_101p_210p_11.$$

We know that  $k = |p_1| + |p_2| + 4$  and  $|C(\frac{x}{y})|_1 = |p_1|_1 + |p_2|_1 + 2 = x$ . Therefore all the elements  $u \in L_k(C(\frac{a}{b})^\omega)$  have  $|u|_1 = |p_1|_1 + |p_2|_1 + 2$  except  $u' = 1p_210p_11$  which is of length  $k$  and at position  $n-k$  but has  $|u'|_1 = x+1$ . Hence by subtracting by  $\min(S_{\frac{a}{b}}[k])$ , we get  $B_{\frac{a}{b}}[k] = 0^{n-k}10^{k-1}$ . By lemma 6, we know that since  $z$  is odd,  $\alpha = n-k$ , which ends the proof.

Same reasoning is used to determine the row  $B_{\frac{a}{b}}[2k]$ , where  $2k = 2|C(\frac{x}{y})| = 2|p_1| + 2|p_2| + 8$ .

The following lemma states the recursive construction of the  $\gamma$ -block in a formal way.

**Lemma 8.** *The  $\gamma$ -block is obtained depending on the position of the fraction in the Stern-Brocot tree.*

- If  $a_z \geq 3$  then  $\gamma(\frac{a}{b})$  is obtained by stacking vertically  $\gamma(\frac{c}{d})$  over  $a_z - 2$  copies of  $\gamma(\frac{u}{v})$  and  $\gamma(\frac{c}{d})$ .
- Otherwise, if  $a_z = 2$ , then  $\gamma(\frac{a}{b})$  is given by the extension of  $\gamma(\frac{c}{f})$  to the  $k$ -th column of  $U_{\frac{c}{f}}$ .

where  $\frac{u}{v} = \text{TBF}(\frac{a}{b})$ ;  $\frac{c}{d} = \text{FDF}(\frac{u}{v})$ ;  $\frac{c}{f} = \text{FPF}(\frac{a}{b})$ .

The proof can be found in the Appendix.

**$\beta$ -block** The  $\beta$ -block is the center of  $U_{\frac{a}{b}}$ , the last block needed to complete the construction of the second order balance matrix. This block is of dimension  $(n-2k-1) \times (n-2k-1)$  and located between the rows (respectively columns)  $k$  and  $n-k$ .

**Lemma 9.** *The  $\beta$ -block of  $U_{\frac{a}{b}}$  is exactly the second order balance matrix of the fraction  $\frac{p}{q} = \text{SUF}(\frac{a}{b})$  in the Stern-Brocot tree where its elements are increased by 1.*

The proof can be found in the Appendix.

## 7 Perspectives

In the introduction, we stated that there were connections between the second order of balance and the Fraenkel conjecture, and we mentioned the importance of the second order of balance to construct balanced words on a  $k$ -letters alphabet. Hence, we have to show some hints on how to use the second order balance matrix  $U_w$  to have new information on the synchronisation of words. Of course the Fraenkel conjecture is far from being solved because of its own complexity nevertheless we have built a new tool to synchronize balanced words by considering synchronization of blocs instead of letters. The following example is based on the word  $FR_3 = 1213121$  on a three letter alphabet, which is obtained from the synchronization of the 3 Christoffel words of slope  $1/6$  associated with the word  $C_3 = **3**$ , of slope  $2/5$  associated with the word  $C_2 = *2***2*$

$$U_{\frac{1}{6}} = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 2 & 2 & 2 & 1 \\ 1 & 2 & 3 & 3 & 2 & 1 \\ 1 & 2 & 3 & 3 & 2 & 1 \\ 1 & 2 & 2 & 2 & 2 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix}, U_{\frac{2}{5}} = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 1 & 1 & 2 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 1 & 1 & 2 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix} \text{ and } U_{\frac{3}{4}} = U_{\frac{3}{4}} = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 2 & 2 & 1 & 1 \\ 1 & 1 & 2 & 2 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix}.$$

**Fig. 8:** The three matrices  $U_{\frac{1}{6}}$ ,  $U_{\frac{2}{5}}$  and  $U_{\frac{3}{4}}$  used for the synchronization of the triplet  $(1, 2, 4) \bmod 7$ .

and of slope  $3/4$  associated with the word  $C_1 = 1 * 1 * 1 * 1$ . Remark that in each position we have only one value and the other symbols are stars. We are able to read this synchronization on rows of matrices in Figure 8:

If we name  $M_{\frac{a}{b}}$  the set of indices of the rows with a maximal values in  $U_{\frac{a}{b}}$  and  $m_{\frac{a}{b}}$  the set of indices of the rows with minimal values in  $U_{\frac{a}{b}}$  (which are non previously chosen rows), we could extract the following set of indices:  $M_{\frac{1}{6}} = \{3, 4\}$  and  $m_{\frac{1}{6}} = \{1, 6\}$ ;  $M_{\frac{2}{5}} = \{2, 5\}$  and  $m_{\frac{2}{5}} = \{3, 4\}$ ;  $M_{\frac{3}{4}} = \{3, 4\}$  and  $m_{\frac{3}{4}} = \{2, 5\}$ . We can observe that starting with the elements of the set  $m_{\frac{3}{4}}$ , if we double the values and take modulo 7, we get the elements of  $m_{\frac{2}{5}}$ . The same note can be also given for the sets  $m_{\frac{2}{5}}$  and  $m_{\frac{1}{6}}$ .  $2 \cdot 2 = 4$ ;  $5 \cdot 2 = 10 \equiv_7 3$ , where  $\{3, 4\} = m_{\frac{2}{5}}$ .  $4 \cdot 2 = 8 \equiv_7 1$ ;  $3 \cdot 2 = 6$ , where  $\{1, 6\} = m_{\frac{1}{6}}$ . Except for the rows 1 and 6, we can see the synchronization between the sets  $m_{\frac{a}{b}}$  and  $M_{\frac{a}{b}}$ , more precisely, if a row appears as a minimum for the set  $m$  it appears as a maximum for the set  $M$ , which allows us to read all information for the synchronization in the second order balance matrix. This will be a keypoint for our next research, in order to study the synchronization of Christoffel words over a  $k$ -letters alphabet and to try to tackle the Fraenkel's conjecture.

In addition, many research problems are still open in this study. For example, how to extend the second order balance matrix construction for a binary word that are not Christoffel words?

We could also investigate the Stern-Brocot tree in order to find the structure of infinite paths that minimize the second order balance and find a combinatorial construction of these paths.

To end the perspectives, we could notice that considering the balance matrix is equivalent to compute balance property of a given word. This computation appears in another context in order to compute Parikh vectors that are  $k$ -dimensional vectors defined for a finite word  $w$  on a  $k$ -letters alphabet. The  $i$ th coordinate of this vector gives the number of occurrences of the letter  $i$  in the word  $w$ . The notion of relative Parikh vectors is introduced by Turek, in [?,?], in order to understand the structure of words. Nevertheless, the computation of the second order balance matrix using Parikh vectors is not straightforward and will be considered in a future article.

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## Appendix

### Proofs:

Property 1:

*Proof.* For (1), given  $i \in \{1, \dots, n-1\}$ , let  $j$  be such that  $S_w[i, j]$  is maximal on the  $i$ -th row. Since we have:  $S_w[i, j] + S_w[n-i, (i+j) \bmod n] = |w|_1$  is always true  $\forall j \leq n$ , then  $S_w[n-i, (i+j) \bmod n]$  is minimal on the  $(n-i)$ -th row, so that  $\max S_w[i] + \min S_w[n-i] = |w|_1$ . Equation (2) is shown in a similar way.

For (1) and (2), since  $w$  is 1-balanced, thus  $\max(S_w[i]) \leq \min(S_w[i]) + 1$ . On the other hand,  $\max(S_w[i]) \neq \min(S_w[i])$  because otherwise, it would imply that  $i$  is a period row in  $B_w$  which is impossible since  $w^\omega$  does not have a period smaller than  $n$  when  $w$  is a primitive Christoffel word.

Lemma 1:

*Proof.* Let  $a/b$  be a rational number with  $n = a + b$ , the Christoffel words  $C(\frac{a}{b})$  and  $C(\frac{b}{a})$  are conjugate and we have:  $D(a, b) = \{(i\alpha \bmod n) + 1; i = 1 \dots a\}$  and  $D(b, a) = \{(i\alpha' \bmod n) + 1; i = 1 \dots b\}$ . From Lemma 3.8 in [22]; we have  $C(\frac{a}{b}) \equiv_\alpha \overline{C(\frac{b}{a})}$ ; where  $a\alpha \equiv_n -1$ . Thus,  $B_{\frac{a}{b}}[1] \equiv_\alpha \overline{B_{\frac{b}{a}}[1]}$ . The other rows are deduced from the first one, hence the pattern is preserved and the same relation holds.

Theorem 2:

*Proof.* From the definition of  $D(a, b)$  and since  $B_w[1] = w$ , we have  $D_1 = D(a, b)$ . By construction,  $|D_1| = a$  and we write  $D_1 = \{d_1, d_2, \dots, d_a\}$ , where  $d_1 < d_2 < \dots < d_a$ .

Since  $a < b$ , the Christoffel word can be decomposed in blocks of the form:  $0^{d+1}1$  and  $0^d1$  where the first block is always of the form  $0^{d+1}1$ . We define an *intersection row*  $D_i$  to be a row such that  $i \in D_1$ . First, we show that the intersection rows are exactly the rows where the values of  $\min(S_w[i])$  and  $\max(S_w[i])$  increase.

Let  $D_i$  be an intersection row, since there is one more occurrence of the letter 1 in the prefix of length  $i$  than in the prefix of length  $i-1$ , we have  $S_w[i, 1] = S_w[i-1, 1] + 1$ . A particular case of Theorem 5 in [10] implies that  $\min(S_w[i]) = \min(S_w[i-1]) + 1$ . Since a Christoffel word has no period smaller than its length and is 1-balanced, we must have that  $\max(S_w[i]) = \min(S_w[i]) + 1$  for all  $i \in \{1, \dots, n-1\}$ . Thus  $\max(S_w[i]) = \max(S_w[i-1]) + 1$  if and only if  $D_i$  is an intersection row and  $\max(S_w[i]) = \max(S_w[i-1])$  otherwise. There are two cases to consider.

- If  $i \notin D_1$ , in this case  $\max S_w[i] = \max S_w[i-1]$  and  $\min S_w[i] = \min S_w[i-1] = S_w[i, 1] = S_w[i-1, 1]$ . Let  $j \in D_i$ , then by definition:

$$B_w[i, j] = S_w[i, j] - \min S_w[i] = S_w[i, j] - S_w[i, 1] = 1.$$

and

$$B_w[i-1, j] = S_w[i-1, j] - \min S_w[i-1] = S_w[i-1, j] - S_w[i-1, 1] \in \{0, 1\}.$$

Indeed,  $j \in D_i$  implies  $B_w[i, j] = 1$  but  $B_w[i-1, j]$  can be equal to 0 or 1.

- If  $B_w[i-1, j] = 1$ , then we have  $j \in D_{(i-1)}$ .

- If  $B_w[i-1, j] = 0$ ,

$$\begin{aligned}
S_w[i-1, j] - S_w[i-1, 1] &= 0, \\
S_w[i-1, j] - S_w[i, 1] &= 0, \\
S_w[i-1, j] - S_w[i, j] + 1 &= 0, \\
S_w[i, j] - S_w[i-1, j] &= 1, \\
|w[j, i+j-1]|_1 - |w[j, i+j-2]|_1 &= 1.
\end{aligned}$$

Therefore,  $w[i+j-1] = 1$  and  $B_w[1, i+j-1] = S_w[1, i+j-1] = w[i+j-1] = 1$ , hence  $i+j-1 \in D_1 \Leftrightarrow j \in D_1 - (i-1)$ .

From the above, we conclude that  $j \in D_i \Leftrightarrow j \in D_{(i-1)} \cup (D_1 - (i-1))$ .

- If  $D_i$  is an intersection row, then  $i \in D_1$  and  $S_w[i, 1] = S_w[i-1, 1] + 1$ ;

$$\begin{aligned}
j \in D_i &\Leftrightarrow S_w[i, j] = S_w[i, 1] + 1, \\
&= S_w[i-1, 1] + 2.
\end{aligned}$$

Let  $j \in D_i$ , by construction  $S_w[i-1, j]$  can be equal to  $S_w[i-1, 1]$  or to  $S_w[i-1, 1] + 1$ . By contradiction, suppose  $S_w[i-1, j] = S_w[i-1, 1]$ . In this case,  $S_w[i, j] = S_w[i-1, j] + 2$ , which is impossible since  $w$  is 1-balanced. Hence,  $S_w[i-1, j] = S_w[i-1, 1] + 1$  and in this case  $B_w[i-1, j] = 1$  and  $j \in D_{(i-1)}$ .

It remains to show that  $j \in D_1 - (i-1)$ . Since  $j \in D_i$  and  $B_w[i, j] = 1$ , we get  $\min(S_w[i]) = (S_w[i, j]) - 1$ . Knowing that  $\min(S_w[i]) = \min(S_w[i-1]) + 1$ , we get  $\min(S_w[i-1]) = S_w[i, j] - 2$  with  $B_w[i-1, j] = 1$  the following holds:

$$\begin{aligned}
S_w[i-1, j] - S_w[i, j] + 2 &= 1, \\
S_w[i, j] - S_w[i-1, j] &= 1, \\
|w[j, i+j-1]|_1 - |w[j, i+j-2]|_1 &= 1.
\end{aligned}$$

Therefore,  $w[i+j-1] = 1$  and  $j \in D_1 - (i-1)$ .

Lemma 2:

*Proof.* The set  $D_1$  contains exactly  $a$  elements, then  $|D_1| = 1 \cdot a = a$ . By induction, we let  $|D_{i-1}| = (i-1) \cdot a \bmod n$ , then we have to prove it true for  $D_i$ . From the previous lemma, we have if  $i \notin D_1$ ,  $D_i = D_{i-1} \cup (D_1 - (i-1))$  and since the two sets are distinct, we get:

$$|D_i| = |D_{i-1}| + |D_1 - (i-1)| = (i-1)a \bmod n + a = i \cdot a \bmod n.$$

When  $i \in D_1$ , we get:

$|D_i| = |D_{i-1}| + |D_1 - (i-1)| - |D_{i-1} \cap D_1 - (i-1)| = (i-1)a \bmod n + a - n = ia \bmod n$ . Since the union covers all the positions in the Christoffel word.

Theorem 5:

*Proof.* We divide the proof into two parts, the first one where  $a_0 = 0$  and the second one for  $a_0 > 0$ .

– For  $a_0 = 0$ :

From Theorem 4, the Christoffel word of slope  $a/b$  has a unique word  $v$  such that:  $w = 0Pal(v)1$  with  $\mu(v) = \begin{pmatrix} |w_1|_0 & |w_2|_0 \\ |w_1|_1 & |w_2|_1 \end{pmatrix}$ . In addition, from Proposition 1, for the rational number  $\frac{a}{b} = [0, a_1, \dots, a_z]$ , we have:  $M(\frac{a}{b}) = A^{a_1-1}B^{a_2} \dots C^{a_z-1}$  where  $C = A$  or  $C = B$  depending on the parity of  $z$  and with  $M(\frac{a}{b}) \cdot \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \begin{pmatrix} b-a \\ a \end{pmatrix}$ . If we multiply by  $A$  from the left side of the matrix  $M(\frac{a}{b})$ , we get:  $A.M(\frac{a}{b}) = A^{a_1}B^{a_2} \dots C^{a_z-1} = \mu(0^{a_1}1^{a_2} \dots p^{a_z-1})$ , where  $p = 0$  or  $p = 1$  depending on the parity of  $z$ .

$$A.M(\frac{a}{b}) \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \cdot \begin{pmatrix} b-a \\ a \end{pmatrix} = \begin{pmatrix} b \\ a \end{pmatrix} = \begin{pmatrix} |w|_0 \\ |w|_1 \end{pmatrix}.$$

On the other hand,  $\mu(v) \cdot \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \begin{pmatrix} |w_1|_0 + |w_2|_0 \\ |w_1|_1 + |w_2|_1 \end{pmatrix} = \begin{pmatrix} |w|_0 \\ |w|_1 \end{pmatrix}$ . This allows us to conclude that  $v = 1^{0a_1} \dots p^{a_z-1}$  where  $p \in \{0, 1\}$  depending on the parity of  $z$ . By Definition 5, we know that the directive sequence of a rational number  $a/b$  is the path followed in the Stern-Brocot tree to reach the rational number. Hence we have that  $w = 0Pal(\Delta(\frac{a}{b}))1$ .

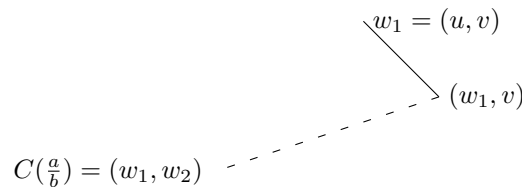
– For  $a_0 > 0$ ,

From Theorem 3, we are able to pass from the partial denominators of a rational number  $a/b = [a_0, \dots, a_z]$  to the Christoffel word of slope  $a/b$  by using Smith algorithm. In fact, the word obtained is of the form  $s_{n+1} = u01$  or  $s_{n+1} = u10$  where  $u$  is a palindrome and  $s_{n+1} = s_n^{a_n} s_{n-1}$ . By Theorem 4, we know that the Christoffel word can be written as:  $0Pal(v)1$ . Therefore, we can conclude that  $u = Pal(v)$ . If we calculate  $s_1 = s_0^{a_0} s_{-1} = 1^{a_0}0$ , then  $s_2 = s_1^{a_1} s_0 = (1^{a_0}0)^{a_1}1$  and we continue the iteration till  $s_{n+1}$ , we remark that  $u$  always starts with  $1^{a_0}$ . Hence, using the iterated palindromic closure function and since  $Pal(1^{a_0}) = 1^{a_0}$ , we can say that  $v$  always starts with  $1^{a_0}$ . Finally, from the previous case, we can get the general form of  $\Delta(\frac{a}{b})$  and say that  $v = \Delta(\frac{a}{b}) = 1^{a_0}0^{a_1} \dots p^{a_n-1}$ , where  $p \in \{0, 1\}$  depending on the parity of  $n$  and  $w = 0Pal(v)1 = 0Pal(\Delta(\frac{a}{b}))1$ .

Lemma 6:

*Proof.* Let  $\frac{a}{b} = [a_0, \dots, a_z]$  and  $D(a, b) = \{i\alpha + 1 \pmod n\}$  where  $\alpha.a \equiv_n -1$ . Using Lemma 3.5 in [5],  $C(\frac{a}{b}) = (w_1, w_2)$  is such that  $|w_1|.a \equiv_n 1 \Leftrightarrow (n - |w_1|).a \equiv_n -1$ . Therefore in  $D(a, b)$ , we have  $\alpha = n - |w_1|$ .

– If  $z$  is odd, then  $\Delta(\frac{a}{b}) = 1^{a_0}0^{a_1} \dots 1^{a_{z-1}}0^{a_z-1}$  and the last function used is  $\phi_0$  which means that the last path used to reach the fraction  $\frac{a}{b}$  on the Stern-Brocot tree is the left path. By the definitions of the two functions  $\phi_0$  and  $\phi_1$ , we get:



where  $w_1 = C(\text{FRF}(\frac{a}{b}))$  and  $|w_1| = k$ , therefore  $\alpha = n - k$ .

– Similarly, we prove the case  $z$  even.

Lemma 8:

*Proof.* By Theorem 1, we have that  $D(a, b) = \{i\alpha + 1 \pmod n; 1 \leq i \leq a\}$  with  $\frac{a}{b} = [a_0, \dots, a_z]$  and since  $z$  is odd then Lemma 6 states that  $\alpha = n - k$ . Let  $\frac{x}{y} = [a_0, \dots, a_{z-1}]$ , we recall that  $|C(\frac{x}{y})| = x + y = k$ , and  $|C(\frac{s}{t})| = s + t = k''$ . Let  $p_1, p_2$  be the two central words of  $C(\frac{s}{t})$  and  $C(\frac{g}{h})$  so that  $C(\frac{x}{y}) = 0p_210p_11 = 0p_101p_21$ , hence  $k'' = |p_1| + 2$ .

– **Case A:** If  $a_z \geq 3$ .

For  $a_z = 3$ , we have:

- $C(\frac{a}{b}) = \left(C(\frac{x}{y})\right)^2 C(\frac{u}{v}) = \left(C(\frac{x}{y})\right)^3 C(\frac{s}{t}); n = a + b$ .
- $C(\frac{c}{d}) = C(\frac{x}{y})C(\frac{u}{v}) = \left(C(\frac{x}{y})\right)^2 C(\frac{s}{t}); n' = c + d$ .
- $C(\frac{u}{v}) = C(\frac{x}{y})C(\frac{s}{t}); n'' = u + v$ , and  $C(\frac{g}{h}) = C(\frac{x}{y}).C(\frac{s}{t})$ , with  $k = x + y$ .

In order to simplify the proof, we give the following notation:

1.  $D_1 = D(a, b) = \{i\alpha + 1 \pmod n; 1 \leq i \leq a\}$ ; where  $\alpha = n - k$ ,
2.  $D'_1 = D(c, d) = \{i\alpha' + 1 \pmod n'; 1 \leq i \leq c\}$ ; where  $\alpha' = n' - k$ ,
3.  $D''_1 = D(u, v) = \{i\alpha'' + 1 \pmod n''; 1 \leq i \leq u\}$ ; where  $\alpha'' = k''$ ,
4.  $X = D(x, y)$  and  $S = D(s, t)$ .

We note that  $\alpha = 2k + k''$  and that by construction, we have:

- $X = D(g, h) \cup \{S + (k - k'')\}$  since  $C(\frac{x}{y})[k - k'' + 1 \dots k] = C(\frac{s}{t})$ .
- $C(\frac{u}{v}) = C(\frac{x}{y})C(\frac{s}{t})$ , hence  $D''_1 = X \cup \{S + k\}$ , knowing that  $\forall s \in S, s + k \leq n''$ .

Similarly, by noticing that  $k = \max(X)$  we get:

$$D'_1 = X \cup \{X + k\} \cup \{S + 2k\} \quad (4)$$

$$D_1 = X \cup \{X + k\} \cup \{X + 2k\} \cup \{S + 3k\}. \quad (5)$$

The following arithmetic relations exist between  $n, n'$  and  $n''$ :

- $n'' = k + k'' = \alpha'$ , and  $n = 3k + k'' = 2k + n'' = 2k + \alpha' = \alpha + k = k + n'$ ,
- $n' = 2k + k'' = k + n'' = k + \alpha' = \alpha$ , and  $n - n' = k = \alpha - \alpha' = n' - \alpha' = n' - n''$ .

Now, we prove that  $D_{k+1} = D_1 \cup \{\alpha\}$ . In fact, the  $C(\frac{a}{b})$  is a 1-balanced word with  $a < b$ , then:

- For  $i = a - 1$  and  $i = 1$ , we get respectively that  $k \in D_1$  and  $\alpha + 1 \in D_1$  which implies that  $k + 1 \notin D_1$  and  $\alpha = n' \notin D_1$ . Therefore,  $D'_1 \setminus \{n'\} \subseteq D_1$ .

Lemma 7 gives that  $B_{\frac{a}{b}}[k] = 0^\alpha 10^{n-\alpha-1}$ , hence  $D_k = \{\alpha + 1\}$ . Therefore, from Theorem 2 and since  $k + 1 \notin D_1$ ,  $D_{k+1} = D_k \cup \{D_1 - k\} = \{\alpha + 1\} \cup \{(i + 1)\alpha + 1 \pmod n, \forall 1 \leq i \leq a\}$ .

Equivalently,  $D_{k+1} = D_1 \cup \{\alpha\}$  since for  $i = a$ , we get  $(a + 1)\alpha + 1 \equiv_n \alpha$ . While for  $1 \leq i \leq a - 1$ , we get all the elements of  $D_1 \setminus \{\alpha + 1\}$ .

The same reasoning is repeated for  $C(\frac{c}{d})$ , and we get that:  $D'_{k+1} = D'_1 \cup \{\alpha'\} = D'_1 \cup \{n''\}$ .

By construction:  $D'_1 \setminus \{n'\} \subset D_1 \subset D_{k+1}$  and  $\{n'\} \subset D_{k+1}$ , therefore  $D'_1 \subset D_{k+1}$ . Since  $\alpha = 2k + k'' = 2|C(\frac{x}{y})| + |p_1| + 2$ , we can write:

$$B_{\frac{a}{b}}[k + 1] = C(\frac{x}{y})C(\frac{x}{y})0p_111p_21C(\frac{s}{t}) \quad \text{and} \quad B_{\frac{a}{b}}[k + 1] = C(\frac{x}{y})0p_111p_21C(\frac{s}{t}).$$

Therefore:  $B_{\frac{a}{b}}[k+1][k+1 \dots n] = B_{\frac{c}{d}}[k+1]$ . We have  $|C(\frac{x}{y})| = k$  and, by abuse of notation, we consider the rows of  $B_{\frac{x}{y}}$  in a circular way so that  $B_{\frac{x}{y}}[k+1] = B_{\frac{x}{y}}[1]$  which allows to write:  $B_{\frac{a}{b}}[i] = B_{\frac{x}{y}}[i]B_{\frac{c}{d}}[i]$  where this relation holds  $\forall k+1 \leq i \leq \alpha' = n - 2k$ .

Therefore, for all  $k+1 \leq i \leq \alpha'$  and  $1 \leq j \leq k$ , we have  $L_j(B_{\frac{a}{b}}^\omega[i]) = L_j(B_{\frac{c}{d}}^\omega[i])$  which gives  $U_{\frac{a}{b}}[i, j] = \max(B_{B_{\frac{a}{b}}[i]}[j]) = \max(B_{B_{\frac{c}{d}}[i]}[j]) = U_{\frac{c}{d}}[i, j]$ . Now we construct the middle part of the  $\gamma$ -block: Using Lemma 7 and Property 2, we have respectively  $B_{\frac{a}{b}}[2k] = B_{\frac{a}{b}}[n - \alpha'] = 0^{\alpha'} 10^{k-1} 10^{k-1}$  and  $B_{\frac{a}{b}}[\alpha'] = 01^{k-1} 01^{\alpha-1}$ . Hence  $D_{\alpha'} = \{2, \dots, k, k+2, \dots, n\}$ , and since  $n - n' = n' - \alpha'$ , we get  $n = 2n' - \alpha' = 2\alpha - \alpha'$ . We have  $\alpha' \notin D_1$  since for  $i = 2$ , we get  $\alpha' + 1 \in D_1$ . Therefore by Theorem 2 we get:  $D_{\alpha'+1} = D_{\alpha'} \cap \{D_1 - \alpha'\} = D_{\alpha'} \cap \{(i-2)\alpha + 1 \pmod n; \forall 1 \leq i \leq a\}$ . We have:

- The positions  $(k+1)$  and  $1$  are obtained for  $i = 1$  and  $2$  respectively and positions  $k$ , and  $n$  are obtained with  $i = a+1$  and  $i = a+2$  respectively.
- While for all  $3 \leq i \leq a$ , we get  $D_1 \setminus \{k, n\} \subseteq D_{\alpha'+1}$ .

Therefore,  $D_{\alpha'+1} = D_1 \setminus \{k, n\}$ .

Since  $\alpha'' = k'' \notin D_1'$ , then:  $B_{\frac{x}{y}}[k''] = 01^{n''-1} \Rightarrow D_{\alpha''} = \{2, \dots, n''\}$  and by Theorem 2 since  $\alpha'' + 1 \in D_1'$  then:

$$D_{\alpha''+1}'' = D_{\alpha''}'' \cap \{D_1'' - \alpha''\} = \{2, \dots, n''\} \cap \{(i-1)\alpha'' + 1 \pmod n''; \forall 1 \leq i \leq n''\}.$$

As a conclusion, we get  $D_{\alpha''+1}'' = D_1'' \setminus \{n''\} \subset D_1$  and then,  $D_{\alpha''+1}'' \setminus \{k\} \subset D_{\alpha'+1}$ , since the positions  $0$  and  $1$  are obtained for  $i = u+1$  and  $i = 1$  respectively. While for  $2 \leq i \leq u$ , we have the same positions as the ones in  $D_1'' \setminus \{0\} = D_1'' \setminus \{n''\} = D_1'' \setminus \{\alpha'\}$ .

We can write:  $B_{\frac{a}{b}}[\alpha' + 1] = 0p_2 10p_1 00p_2 10p_1 1C(\frac{x}{y})0p_1 0$  and  $B_{\frac{x}{y}}[\alpha'' + 1] = C(\frac{x}{y})0p_1 0$ .

Hence we get  $B_{\frac{a}{b}}[\alpha' + 1][2k+1, \dots, n] = B_{\frac{x}{y}}[\alpha'' + 1]$  and:

$$B_{\frac{a}{b}}[\alpha' + 1 + l] = 0p_2 10p_1 0C(\frac{x}{y})B_{\frac{x}{y}}[\alpha'' + 1 + l], \forall 0 \leq l \leq 2k - \alpha' - 1.$$

Therefore,  $L_j(B_{\frac{a}{b}}^\omega[i]) = L_j(B_{\frac{x}{y}}^\omega[i])$  for all  $\alpha' + 1 \leq i \leq 2k$  and  $1 \leq j \leq k$  since the element  $0p_1 00p_2 1$  of length  $k$  appears in both languages thus we have  $\forall \alpha' + 1 \leq i \leq 2k; 1 \leq j \leq k$ :

$$U_{\frac{a}{b}}[i, j] = \max(B_{B_{\frac{a}{b}}[i]}[j]) = \max(B_{B_{\frac{x}{y}}[i]}[j]) = U_{\frac{x}{y}}[i, j].$$

The rows from  $2k+1$  till  $n-k$  are obtained due to the symmetry of  $U_{\frac{a}{b}}$  and since  $2k > \frac{n}{2}$ .

Therefore, we get:  $\gamma(\frac{a}{b}) = \gamma(\frac{c}{d}) [\gamma(\frac{u}{v}) \gamma(\frac{c}{d})]^1$ , where  $1 = a_z - 2$ .

When  $a_z > 3$  and using the same reasoning we get  $\gamma(\frac{a}{b}) = \gamma(\frac{c}{d}) [\gamma(\frac{u}{v}) \gamma(\frac{c}{d})]^{a_z-2}$  since we have  $C(\frac{a}{b}) = (C(\frac{x}{y}))^{a_z-1} C(\frac{u}{v})$ , where  $D_1$  is given by:

$$D_1 = X \cup \{X+k\} \cup \{X+2k\} \dots \cup \{X+(a_z-1)k\} \cup \{S+a_z k\}.$$

– **Case B:** If  $a_z = 2$ .

For this part we have:  $C(\frac{a}{b}) = C(\frac{x}{y})C(\frac{u}{v})$  and we need the word  $C(\frac{e}{f}) = C(\frac{g}{h})C(\frac{x}{y})$ , where  $\frac{e}{f} = \text{FPF}(\frac{a}{b})$ . The same notation are used as in the case A, where  $n = a + b$ ,  $n' = e + f$ ,  $k = |C(\frac{x}{y})|$  and  $k' = |C(\frac{g}{h})| = |p_2| + 2$ . The sets  $D_1 = D(a, b)$  and  $D_1' = D(e, f)$  are calculated in the same way and we get:  $n = k + \alpha$ ,  $n' = k + k' = k' + \alpha'$  and  $k = \alpha'$ . Using the same previous technique, we get:



1.  $D_{k+1} = D_1 \cup \{\alpha\}$ ; and  $D'_{k'+1} = D'_1 \cup \{\alpha'\} = D'_1 \cup \{k\}$ ;
2.  $k+1 \notin D_1 \Rightarrow k+1 \notin D_1 \cup \{\alpha\}$ ; and  $\alpha' = k \in D_1$ .

The rest of the positions are obtained due to the concatenation of the word. We have  $D'_{k'+1} \setminus \{\alpha' + 1\} \subseteq D_{k+1}$ ; which means:  $\{D'_1 \cup \{\alpha'\}\} \setminus \{\alpha' + 1\} \subseteq D_1 \cup \{\alpha\}$ . We can write:  $C(\frac{a}{b}) = 0p_210p_11 \cdot 0p_210p_11 \cdot 0p_11 = 0p_210p_11 \cdot 0p_101p_21 \cdot 0p_11$ , and since  $\alpha = n - k = n - (|p_1| + |p_2| + 4)$ , we get:  $B_{\frac{a}{b}}[k+1] = 0p_210p_11 \cdot 0p_111p_21 \cdot 0p_11 = 0p_101p_21 \cdot 0p_111p_21 \cdot 0p_11$ . While for  $C(\frac{c}{f})$ , we have:  $C(\frac{c}{f}) = 0p_21 \cdot 0p_21 \cdot 0p_11 = 0p_21 \cdot 0p_101p_21$ , with  $\alpha' = n' - k' = n' - (|p_2| + 2)$ , hence  $B_{\frac{c}{f}}[k'+1] = 0p_21 \cdot 0p_111p_21 = 0p_101p_211p_21$ . Therefore we get:  $L_j(B_{\frac{a}{b}}^\omega[i]) = L_j(B_{\frac{c}{f}}^\omega[i])$ , for all  $k+1 \leq i \leq n-k-1$  and  $1 \leq j \leq k$ .

Thus  $\forall k+1 \leq i \leq n-k-1; 1 \leq j \leq k$ , we have:

$$U_{\frac{a}{b}}[i, j] = \max(B_{B_{\frac{a}{b}}[i]}[j]) = \max(B_{B_{\frac{c}{f}}[i]}[j]) = U_{\frac{c}{f}}[i, j].$$

Lemma 9:

*Proof.* By Definition 3,  $U_{\frac{a}{b}}$  studies the second order of balance for the Christoffel word of slope  $a/b$  where  $U_{\frac{a}{b}}[i, j] = \max_{u \in L_j(B_w[i])} |u|_1 - \min_{u \in L_j(B_w[i])} |u|_1$ . The elements of the  $\gamma$ -block are obtained for  $k+1 \leq i \leq n-k-1$  and  $1 \leq j \leq k$ , which refer to the words  $u \in L_j(B_w[i])$  where  $B_w[i, j] = S_w[i, j] - \min(S_w[i])$  and  $S_w[i, j] = |s_w[i, j]|_1 = |w^\omega[j, j+i-1]|_1$ . The elements of the  $\beta$ -block are obtained using the words  $u \in L_j(B_w[i])$  for  $k+1 \leq i \leq n-k-1$  and  $k+1 \leq j \leq n-k-1$ . We have 2 cases either  $a_z \geq 3$  or  $a_z = 2$ .

- If  $a_z \geq 3$ : In the previous section we proved that for  $a_z = 3$ , we have:

- $B_{\frac{a}{b}}[k+1] = C(\frac{x}{y})C(\frac{x}{y})0p_111p_21C(\frac{s}{t})$ .
- $\text{SUF}(\frac{a}{b}) = \frac{u}{v}$  with  $C(\frac{u}{v}) = 0p_210p_11 \cdot 0p_21$ .
- We have  $w' = 0p_111p_210 \in L_{k+1}(B_{\frac{a}{b}}^\omega[k+1])$ , and  $w'' = 0p_101p_210 \in L_{k+1}(B_{\frac{u}{v}}^\omega[1])$ .

We observe that  $|w'|_1 = |w''|_1 + 1$ . Since  $B_{\frac{a}{b}}[k+1] = C(\frac{x}{y})C(\frac{x}{y})w'p_21$ , then by calculating the  $L_j(B_{\frac{a}{b}}^\omega[k+1])$ ,  $\forall k+1 \leq j \leq n-k-1$ , the maximal value appears in the factor  $w'$  that causes that difference since  $|w'|_1$  is always increased by 1 comparing to  $|w''|_1$ . By studying the following rows of  $B_{\frac{a}{b}}$ , this same factor causes the increase of the maximal value.

Hence,  $\forall k+1 \leq i \leq n-k-1; k+1 \leq j \leq n-k-1$ :

$$U_{\frac{a}{b}}[i, j] = \max(B_{B_{\frac{a}{b}}[i]}[j]) = \max(B_{B_{\frac{u}{v}}[i-k]}[j-k]) + 1 = U_{\frac{u}{v}}[i-k, j-k] + 1.$$

While for  $a_z > 3$  and by concatenating with the  $C(\frac{x}{y})$  from the left side for each fraction, we always get that:

$$U_{\frac{a}{b}}[i, j] = U_{\frac{u}{v}}[i-k, j-k] + 1.$$

- If  $a_z = 2$ , we let  $\frac{a}{b} = [a_0, \dots, 2]$  and we have from the previous sections:

- $B_{\frac{a}{b}}[k+1] = C(\frac{x}{y})0p_111p_21C(\frac{s}{t}) = 0p_210p_110p_111p_210p_11$ .
- By Definition 6,  $\text{SUF}(\frac{a}{b}) = \frac{s}{t} = [0, \dots, a_{z-2}]$ , where  $C(\frac{s}{t}) = 0p_11$ .

Just like the previous case, the factor that causes the difference is the factor  $w' = 0p_111p_210$  that belongs to  $L_{k+1}(B_{\frac{a}{b}}^\omega[k+1])$ . Therefore, the maximal value appears in  $|w'|_1$  that gives the same value if we study the occurrence of 1 in all the factors of  $L_1(B_{\frac{s}{t}}^\omega[1])$ . By studying the following rows of  $B_{\frac{a}{b}}$ , the same factor causes the increase of the maximal value. Hence,  $\forall k+1 \leq i \leq n-k-1; k+1 \leq j \leq n-k-1$ :

$$U_{\frac{a}{b}}[i, j] = \max(B_{B_{\frac{a}{b}}[i]}[j]) = \max(B_{B_{\frac{s}{t}}[i-k]}[j-k]) + 1 = U_{\frac{s}{t}}[i-k, j-k] + 1.$$