# Defining and Testing Identifiability, Illustrated by a HIV model 

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Defining and Testing Identifiability, Illustrated by a HIV model
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## Basic mathematical definitions

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## Introduction

We question the notion of identifiability and its relation with the actual ability to achieve identification. We show that a boundedness hypothesis on the set of parameter is essential, that amounts in most pratical situation to bounding the magnitude of parameters.

In the formal definition of identifiability, one should be very carefull with the parentheses and the order of quantifiers, mostly concerning the external controls. However, under some good hypotheses, this is unimportant since a single trajectory may stand for all trajectory.

Without controls, a weaker property plays a similar le and is expressed by the non-vanishing of a Wrońskian determinant. Then, eliminating the state variables is enough to get an exhaustive summary and to test identifiability.

This method is close to criteria encountered in Wu st al $[7]$ for HIV models that we will consider as an application. Differential algebra also gives rigorous upper bounds on the minimal number of measurement times to test local and global identifiability

We conclude with references to some computer tools.

## Differential algebra

Founded by Joseph Ritt (1893-1951), differential algebra considers differential fields $\mathscr{F}$, i.e. fields equipped with a derivation, such as $\mathbb{Q}(t)$ with $d / d t$, systems of differenial polynomial equations, i.e. polynomials $\mathscr{F}\left\{x_{1}, \ldots, x_{n}\right\}$ in some variables $x_{i}$ and their derivatives $x_{i}^{\prime}, x_{i}^{\prime \prime \prime}, \ldots$. Charac teristic sets are particular sets of equations, that generalize the notion of normal form.

Any set of explicit equations is a characteristic set defining a prime differential ideal. A differential ideal is the set [ $\Sigma$ ] of all differential equations that are consequences of the initial system $\sigma$ and prime means that $P Q \in[\Sigma]$ implies $P \in \Sigma$ or $Q \in \Sigma$.

A differential Zariski open set is defined by an inequation $P(x) \neq 0$, where $P$ is a differential polynomial.

An abstract definition of identifiability

In practice, a structure or parametric model is most of the time given by explicit differential equations:

$$
x_{i}^{\prime}=f_{i}(x, u, \theta, t), \quad \mathrm{I} \leq i \leq n,
$$

where the $f_{i}$ are rational functions, $t$ is the time, satisfy ing $t^{\prime}=\mathrm{I}$, the $\theta_{j}, \mathrm{r} \leq j \leq s$ are constant parameters, so $\theta_{j}^{\prime}=0$, the $u_{k}, 1 \leq k \leq m$ are control functions, and the $x$ $\mathrm{I} \leq i \leq n$ are the state variables.

We need to complete this differential system with initial conditions that are rarely discussed:

$$
x_{i}(\mathrm{o})=c_{i} .
$$

We also need to define outputs:

$$
y_{\ell}=g_{\ell}(x, \theta, t), \quad 1 \leq \ell \leq r,
$$

where the $g_{\ell}$ are also rational functions. Let $X[c, \theta, u](t)$ be the unique solution defined by the control $u$, the parameters $\theta$ and the initial condition $c$. The differential Zariski open set $U$ is such that the functions $f$ and $g$ are defined on $U$.

Definition i. $-A$ function $H$ of the parameters $\theta$ is algebraically identifable if there exists a Zariski open set $V \subset U$ such that $\forall(c, \theta, u) \in V, \forall(\hat{c}, \hat{\theta}, u) \in U$ $g(X[c, \theta, u])=g(X[\hat{c}, \hat{\theta}, u])$ implies $H(\hat{\theta})=H(\theta)$, or equivalently if there exists a differential rational function $G$ such that $H(\theta)=G(u, y, t)$. It is locally identifiable if there exists a differential polynomial P of order O in $H(\theta)$ such that $P(H(\theta), u, y, t)=0$.

Algebraic identifiability means that $\theta_{i}$ is a differential rational function depending on the derivatives of the controls $u$ and the outputs $y$. Local identifiability mean that it is an algebraic function, so that the value is only locally unique.

One may notice that our definition explicitly in volves the initial condition $c$ which is too often omitted in the literature.

## Identifiability and identification

Assuming the mathematical model describes perfectly the actual behavior and the noise is zero, is identifiability th guarantee to achieve identification ? And if so, what ar the relations between the abstract mathematical identifiability and the potential succes of practical identification processes?

Consider the system corresponding to the following niversal equation. Assuming that the state $x$ is measured, it is identifiable (we will see why below)

$$
\begin{equation*}
x^{\prime}=\frac{b d}{1+d^{2}-\cos (b(t-a))} \cos \left(e^{t}\right) \tag{}
\end{equation*}
$$

$x(-a)=c$,
(2)
where $a, b, c$ and $d>0$ are real number
Boshernitzan [ $[\mathbb{1}]$ has shown the following result.
Theorem 2. - For any continuousfunction $F(t)$ defined on any compact interval $E$, for any $\epsilon>0$ there exist $a, b, c$ $d$ such that $\forall t \in E|X((a, b, c, d), t)-F|<\epsilon$.

This property implies that any (continuous) func tion may be approximated by four parameters in a ver specific model (II). In presence of noise, any perturbatio would give very different parameters, which is possible be cause the set of parameters is unbounded.

The technical details in the next theorem may be skept. Its meaning is that the value $\theta_{1}$ is unique in any bounded subset, when the precision $\epsilon$ is small enough. Other values may exist but for which the size of param eters go to infinity.

$\left\{\theta^{*} \in E \mid \nmid \forall_{0} \leq t \leq T\left\|g\left(X\left(\theta^{*}, t\right), \theta^{*}\right)-g(X(\theta, t), \theta)\right\|<\epsilon\right\}$

The proof of hhe theorem reliesie on the to topologicil propery that charac

## Irreducible systems

Hong et al. $[3$, prop. I] have noticed that we have an equiv alent definition of identifiability, where the unique value of $\theta$ is deduced from the knowledge of the input output behavior, i.e. the function that associate an output $y$ to all possible input of control functions $u$. This important property seems paradoxical as, in practice, one knows during the experiment only one vector of inputs and one vec tor of outputs.

In fact, in most cases, we can do as if we measured the output functions for all inputs $u$ and for all initial conditions $c$. We can give here a precise meaning to this genericity.

Definition 4. $-A$ system $\Sigma$ generating a prime differential ideal $\mathfrak{P}$ is irreducible if there exists a Zariski open set $U$ such that for any $(c, \theta, u) \in U P(X[c, \theta, u])=0$ implies $P \in \mathscr{P}$. Such solutions are called generic solutions of the ideal $\mathscr{P}$

One may prove that a system is irreducible if it admits no rational first integral.

One may eliminate state variables by computing characteristic set for the differential ideal associated to ou initial system, with an ordering such that the state variables $x_{i}$ and all their derivatives are greater than the $y_{y}$. It contains a set of differential equations $Q_{j}(y, u, \theta, t)$. We may chose to extend the order on derivatives to an order on monomials and assume that the coefficient of the main monomial in $Q_{j}$ is I . Then let $\rho_{j, i}(\theta), \mathrm{I} \leq \ell \leq p_{j}$ denote the coefficients of the $p_{j}$ monomials in $Q_{j}$.

Theorem 5. - If the system is irreducible, then the $\rho_{j, 1}$ are algebraically identifiable.

The main idea of the proof is easy: if some $\rho_{;, \ell}$ is not identifiable, then there exists $\vartheta \neq \theta$ such that $P:=$ $Q_{j}(\gamma, u, \theta, t)-Q_{j}(y, u, \vartheta, t) \notin \mathscr{P}$ but $P(X[c, \theta, u]) \in \mathscr{P}$. In other words, the independence of monomials implies the equality of their coefficients.

Controlability is a sufficient condition of irreducibility that is easy to check. But what can we do when there is no control?

## Using the Wrońskian

We borrow the following HIV model to Wu et al. wher it is assumed that only $V$ is measured.

$$
\begin{cases}T^{\prime}(t)=\lambda-\rho T(t)-\beta T(t) V(t), & , T(\mathrm{o})=T_{\mathrm{o}} \\ U^{\prime}(t)=+\beta T(t) V(t)-\partial U & , U(\mathrm{o})=U_{\mathrm{o}} \\ V^{\prime}(t)=N \delta U(t)-c V(t) & , V(\mathrm{o})=V_{\mathrm{o}} .\end{cases}
$$

Eliminating $T$ and $U$, one gets an equation of minimal or der in $V$ alone.
$P(, V):=V V^{(3)}(t)-\left(V^{\prime}-\rho V-\beta V^{2}\right)\left(V^{\prime \prime}+(\delta+c) V^{\prime}+\delta c V\right)$
$-N \lambda \delta \beta V+\partial c V^{\prime}+(\delta+c) V^{\prime \prime}=0 .(4)$
We see that only $N \lambda, \delta+c$ and $\delta c$ appear in this system. So only the vector of new parameters $\theta:=(\rho, \beta, \kappa:=$ $N \lambda \delta, \mu:=\delta+c, \nu:=\delta c)$ can be possibly identifiable. Thei local identifiability is deduced by Wuet al. from the nonvanishing of the Jacobian determinant:

$$
\left|\frac{\partial p^{k}}{\partial \theta_{j}}\right| 0 \leq k \leq 4,1 \leq j \leq s \mid,
$$

using the implicit function theorem.
Besides the main monomial $V V^{\prime \prime \prime}$, equation (4) contains the set of 12 monomials $M$ : $\left\{V^{\prime} V^{\prime \prime}, V^{\prime 2}, V V^{\prime}, V^{2} V^{\prime \prime}, V^{2} V^{\prime}, V^{3}, V V^{\prime \prime}, V V^{\prime}, V^{2}, V^{\prime \prime}, V^{\prime}, V\right\}$ By the proof of th 5 , if the coefficients $\rho_{i, j}(\theta)$ are not identifiable, there is a non-trivial relation with constant coefficients between these monomials : $P(\theta, V)-P(\vartheta, V)=0$. The existence of such a relation with $V$ analytic is equiva lent to the vanishing of the Wrońskian determinant:

$$
W(M):=\left|m^{k}\right| m \in M, o \leq k \leq \# M-\mathrm{x} \mid
$$

Theorem 6. - Let $\mathscr{B}$ is a characteristic set of the prime ideal $\mathcal{P}$ associated to the model $\Sigma$ for a ranking that elimi nates the state and let the $Q_{i}$ be the elements of $\mathscr{B}$ that do not depend on the state variables $x$. Let $M_{i}$ be the set of mono mials in $Q_{i}$, then for any $\mathrm{I} \leq i \leq r, i f W\left(M_{i}\right) \neq \mathrm{o}$, the coefficients $\rho_{i, j}(\theta)$ are identifable.

## Testing the non-vanishing of the Wrońskian

The expression of the derivatives in the Wrońskian may be huge. To test if the determinant is non-zero modulo the equations of the system, we need to reduce it, producing an increase of the size. An easy way to solve the problem is to replace $x_{i}(t)$ by power series solution, associated to random integer coefficients and initial conditions. One may also make computations modulo some great prime integer $p$. This gives a probabilistic answer. If one gets a non-zero evaluation the determinant is non-zero. If the evaluation is o , we can reduce the probability of failure by repeating the experiment and increasing the size of random integer coefficients.

The computation of power series can be done with a near linear asymptotic complexity algorithm, due to Sedoglavic [2]. Van der Hoeven's algorithm [6] has greater asymptotic complexity but is often better in practice.

The main idea of Sedoglavic's algorithm is to gen eralise Newton's method. A key ingredient is to be able to compute the derivative of $X[c, \theta]$ with respect to $\theta$ : $\partial X / \partial \theta_{i}$. It is solution to the linearized system

$$
\sum_{j=1}^{n} \frac{\partial P_{i}}{\partial x_{j}} d x_{j}, d x_{j}(\mathrm{o})=0 .
$$

With such a setting, one may also test local iden tifiability by computing the jacobian determinant $\left|\partial X[c, \theta]\left(t_{i}\right) / \partial \theta_{j}\right|$ using $s$ random observation points as proposed in Wu et al. Interval or ball arithmetic ca confirm it is non-zero, but the vanishing will remain du biou

## Minimal number of observation point

For an algebraically identifiable system, the $\theta_{i}$ can be expressed as differential rational functions of the outputs of order at most $n+s$, that is the number of initial con ditions $c$ and parameters $\theta$. This is indeed the order of th system, completed with the equations $\theta_{i}^{\prime}=\mathrm{o}$, and so the maximal order of any characteristic set. This suggests that $n+s+1$ generic observation points may be enough to compute the parameters. But one need to keep in mind that for some system, no identification is possible with noisy data without bounding the magnitude of parameters.

## Some implementation

Sedoglavic's method has been implemented in Maple and is available: $\qquad$
The algorithms in the Maple DifferentialAlgebr Package are also available in the C library BLAD: A Sage interface BMI is also available:
The free computer algebra system Mathemagix pro vides efficient implementations of up to date fast algo rithms for exact and approximate computations, including power series solutions of ODEs and ball arithmetics. An experimental package allows to compute a linearized system in a form that allows numerical integration in


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