



HAL
open science

Robust anisotropic power-functions-based filtrations for clustering

Claire Brécheteau

► **To cite this version:**

Claire Brécheteau. Robust anisotropic power-functions-based filtrations for clustering. 2019. hal-02397100v1

HAL Id: hal-02397100

<https://hal.science/hal-02397100v1>

Preprint submitted on 6 Dec 2019 (v1), last revised 25 Mar 2020 (v2)

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Robust anisotropic power-functions-based filtrations for clustering.

Claire BréchetEAU 

Laboratoire de Mathématiques Jean Leray, École Centrale Nantes, France
<https://www.math.sciences.univ-nantes.fr/~brecheteau/>

1 Abstract

2 We consider robust power-distance functions that approximate the distance function to a compact
3 set, from a noisy sample. We pay particular interest to robust power-distance functions that are
4 anisotropic, in the sense that their sublevel sets are unions of ellipsoids, and not necessarily unions
5 of balls. Using persistence homology on such power-distance functions provides robust clustering
6 schemes. We investigate such clustering schemes and compare the different procedures on synthetic
7 and real datasets. In particular, we enhance the good performance of the anisotropic method for
8 some cases for which classical methods fail.

2012 ACM Subject Classification Theory of computation → Unsupervised learning and clustering

Keywords and phrases Power functions, Filtrations, Hierarchical Clustering, Ellipsoids

Funding *Claire BréchetEAU*: [funding]

Acknowledgements I am extremely grateful to Samuel Tapie, for his suggestion to use tangency of ellipsoids at their intersection point, to derive the expression of their intersection radius.

Lines 500

9 1 Introduction

10 Often data can be represented as a point cloud \mathbb{X} in a Euclidean space \mathbb{R}^d . Grouping data
11 into clusters as homogeneous and well-separated as possible is the purpose of clustering.
12 When no label is known in advance, we talk about unsupervised clustering. Topological data
13 analysis (TDA) tools are designed to understand the shape of the data. Thereby, such tools
14 may help to understand the shape of clusters in which to group the data. In this paper, we
15 develop and study a TDA-based unsupervised clustering scheme. In addition, our method
16 detects and removes points that do not really belong to any cluster; the outliers.

17 Clustering datasets is of extreme importance in multiple domains including medicine and
18 social networks among others. The classical k -means method clusters data into isotropic
19 clusters. In particular, the trimmed version of k -means of [14] that removes outliers, supplies
20 balls-shaped clusters. These two algorithms have been extended by [2, 5] for Bregman-balls-
21 shaped clusters, see also `tlust` [17] for ellipsoidal clusters. Such methods are well-suited for
22 data generated according to mixtures of distributions whose sublevel-sets are Bregman balls
23 themselves. For more general datasets, for instance, a sample of points from a disconnected
24 manifold, these methods are no longer appropriate. Spectral clustering methods [26] perform
25 such tasks, but are not robust to outliers. DBSCAN [19] is an algorithm based on a fixed
26 upper-level set of an approximation of the density, and consequently, does not provide a
27 multiscale information. Via a dendrogram, the classical single-linkage hierarchical clustering
28 algorithm provides such a multiscale information. The dendrogram encodes information
29 about the connectivity of unions of balls centered at points in \mathbb{X} , or equivalently, of the
30 sublevel sets of the distance function to \mathbb{X} . For a fixed radius r , the Čech complex is a
31 simplicial complex defined as the collection of simplices (vertex, edge, triangle, tetrahedron)
32 for which the r -balls centered at the vertices have a non-empty common intersection. We



33 call 1-skeleton its subcomplex (a graph) that contains only vertices and edges. The non-
 34 decreasing family of such graphs indexed by $r \in \mathbb{R}$ is called a filtration. Single-linkage is
 35 a persistence-based method since is based on the persistence, prominence or equivalently
 36 lifetime of the connected components into this graph filtration, however, it is not robust
 37 to outliers. The algorithm ToMATo in [12] is robust and persistence-based. Indeed, it is
 38 based on a graph filtration built from a neighborhood graph and a (robust) distance-like
 39 function whose values guide the appearance of vertices and edges in the graph filtration.
 40 An example of robust distance function that Chazal et al. consider in [12] is given by the
 41 distance-to-measure (DTM) [11]. Note that the graph is a priori not intrinsic to the distance
 42 function, which may cause bad clustering. For instance, an edge that links two vertices
 43 with small distance-function value but intersects an area with large distance function value,
 44 may link two clusters that should not be. This problem was the cause of failure of the
 45 single-linkage method for data corrupted by outliers. Alternative filtrations that do not
 46 suffer from this problem are the DTM-filtration [1], or the power filtrations [7], based on the
 47 1-skeleton of the Čech filtration associated to the sublevel sets of a power distance function:
 48 a function of type $x \mapsto \min_{i \in I} \|x - m_i\|^2 + \omega_i$ for some $(m_i)_{i \in I}$ in \mathbb{R}^d and $(\omega_i)_{i \in I}$ in \mathbb{R} . Some
 49 approximations of the DTM that are power functions have been introduced and studied in
 50 the literature: the k -witnessed distance [18], the c -PDTM [6] whose sublevel sets are unions
 51 of c balls, and the c -PLM [4] whose sublevel sets are unions of c ellipsoids, with c possibly
 52 much smaller than the sample size. The last two functions are robust to outliers since their
 53 construction is based on the principle of trimmed least squares [25].

54 Contributions

55 By replacing balls with ellipsoids, we enlarge the notion of weighted Čech filtration into the
 56 anisotropic weighted Čech filtration. We derive an expression for the radius of intersection of
 57 two ellipsoids. We introduce a clustering algorithm based on persistence. Such a clustering
 58 algorithm can be run from any graph filtration, in particular, from the 1-skeleton of the
 59 anisotropic weighted Čech filtration, which corresponds to the filtration of sublevel sets of an
 60 anisotropic power function. We experiment this algorithm on the filtration of the c -PLM [4].

61 Practical interests

62 A clustering algorithm based on the persistence filtration of the sublevel sets of a power
 63 function is pertinent since unlike ToMATo, the graph is intrinsic to the distance function.
 64 So, no additional parameters are required for the algorithm. The main advantage of using
 65 an anisotropic power function is that its sublevel sets are ellipsoids. Much less ellipsoids are
 66 required than balls to Hausdorff-approximate a compact manifold with intrinsic dimension
 67 smaller than the ambient dimension. The clustering scheme can also be applied to decompose
 68 a set of points generated on a polygonal lines into segments. Once the ellipsoids computed, the
 69 persistence algorithm runs very fast, in c^2 , with c , the number of ellipsoids. Most importantly,
 70 the robustness of the persistence algorithm relies on the robustness of the distance function.
 71 The c -PLM [4] is robust to outliers. The guaranty for the clustering method follows from
 72 the $\|\cdot\|_\infty$ -distance closeness between the power distance function and the distance function
 73 to the underlying manifold \mathcal{X} , relatively to the minimal distance between the connected
 74 components of \mathcal{X} . Note that such a proximity condition is sufficient but not necessary, as
 75 illustrated by the different numerical examples, with the c -PLM.

76 Organisation of the paper

77 In Section 2, we recall the notions of power function and weighted Čech filtration, the
 78 filtration of the nerves of its sublevel sets, that we extend to anisotropic power functions.
 79 We prove some stability and approximation properties for such filtrations. Examples of
 80 robust power filtrations are also displayed. The main clustering algorithm, Algorithm 1 is
 81 given in Section 3. This algorithm applies to any filtration of graphs, including the graph
 82 filtrations obtained as the 1-skeleton of a weighted Čech filtration. We enumerate other types
 83 of filtrations that fit into this framework. Finally, we implement Algorithm 1 with the robust
 84 anisotropic aforementioned power function in Section 4. We compare this method to other
 85 clustering methods on synthetic and real datasets.

86 **2** Power-functions-based filtrations for robust clustering

87 In the sequel, we will recall the notion of filtration for subsets of \mathbb{R}^d and for simplicial
 88 complexes. We will consider a class of functions for which filtrations associated to sublevel
 89 sets are easily represented by filtrations of simplicial complexes, making the evolution of their
 90 connected components tractable: the power functions. In addition, we will give an example
 91 of robust power-functions [6] that can be built from a probability distribution or a pointset
 92 \mathbb{X} . Their sublevel sets are unions of c balls, with c possibly much smaller than the size of \mathbb{X} .
 93 Most importantly, we will also give an example of a robust anisotropic power-function, whose
 94 sublevel sets are unions of c ellipsoids [4]. Both of these power functions will be considered
 95 in the next sections for clustering purposes.

96 2.1 Generalities on filtrations

97 Consider T , a subset of \mathbb{R} that represents time. We equip \mathbb{R}^d with the Euclidean norm $\|\cdot\|$.
 98 A filtration indexed by T is a family $(V^t)_{t \in T}$ of subsets of \mathbb{R}^d , that is non-decreasing for the
 99 inclusion relation, in the sense that $\forall t \leq t', V^t \subset V^{t'}$. A typical example of filtration is given
 100 by the filtration of the sub-level sets of a function $f : \mathbb{R}^d \mapsto \mathbb{R} : (f^{-1}((-\infty, t]))_{t \in T}$. For any
 101 simplex S with vertex set given by \mathbb{X} , a finite set of points, the notion of filtration extends
 102 as follows. A filtration of simplicial complexes of S is a non-decreasing family $(S^t)_{t \in T}$ of
 103 subcomplexes of S , meaning that for every $t \leq t'$, any simplex of S^t is also a simplex of $S^{t'}$.

104 It is possible to compare two filtrations with the notion of interleaving. The interleaving
 105 pseudo-distance between two filtrations $(V^t)_{t \in T}$ and $(W^t)_{t \in T}$ is defined as the smallest $\epsilon > 0$
 106 such that $(V^t)_{t \in T}$ and $(W^t)_{t \in T}$ are ϵ -interleaved, meaning that: $\forall t \in T, V^t \subset W^{t+\epsilon}$ and
 107 $W^t \subset V^{t+\epsilon}$. The same definition extends to simplicial complexes. It should be noted that for
 108 two functions f and g such that $\|f - g\|_\infty \leq \epsilon$, their sub-level-sets filtrations are ϵ -interleaved.

109 We will see in Section 3 that the notion of interleaving is primordial, since it measures
 110 the difference of topology between two filtrations. In particular, for our purpose of clustering
 111 based on sub-level sets of two functions, clustering stability will be guaranteed from the
 112 closeness of the functions.

113 2.2 Power-functions-based filtrations

114 In this paper, we consider classes of functions whose sub-level sets filtration has a sparse
 115 representation, the power functions. The sublevel sets of these functions can be represented
 116 by simplicial complexes in so-called weighted Čech filtrations. We will consider two types of
 117 power functions, the isotropic and the anisotropic ones.

118 **2.2.1 The isotropic case**

119 An isotropic power function is a function $f_{\mathbf{m},\omega} : \mathbb{R}^d \rightarrow \mathbb{R}$ defined from an index set $I = \llbracket 1, c \rrbracket$,
 120 a family of centers $\mathbf{m} = (m_i)_{i \in I}$ in \mathbb{R}^d and a family of weights $\omega = (\omega_i)_{i \in I}$ in \mathbb{R} by
 121 $f_{\mathbf{m},\omega} : x \mapsto \min_{i \in I} \|x - m_i\|^2 + \omega_i$. A simple example of power function is the squared
 122 Euclidean distance function to a set of points \mathbb{X} , $d_{\mathbb{X}}^2 : x \in \mathbb{R}^d \mapsto \min_{m \in \mathbb{X}} \|x - m\|^2$. The
 123 sublevel sets of $f_{\mathbf{m},\omega}$, $V_{\mathbf{m},\omega}^t = f_{\mathbf{m},\omega}^{-1}((-\infty, t])$, are unions of at most c balls $\mathcal{B}_i^t = \overline{B}(m_i, \sqrt{t - \omega_i})$
 124 with $\overline{B}(m, r) = \{x \in \mathbb{R}^d \mid \|x - m\| \leq r\}$. Note that \mathcal{B}_i^t is empty for $t < \omega_i$. A simple
 125 computation enhances that two balls \mathcal{B}_i^t and \mathcal{B}_j^t intersect if and only if $t \geq t_{i,j}$ with
 126 $t_{i,j} = \frac{(\omega_j - \omega_i)^2 + 2(\omega_j + \omega_i)\|m_j - m_i\|^2 + \|m_j - m_i\|^4}{4\|m_j - m_i\|^2}$. The information of connectivity of $V_{\mathbf{m},\omega}^t$ can
 127 be encoded in a graph $\mathcal{G}_{\mathbf{m},\omega}^t$, composed of vertices indexed by I . The graph $\mathcal{G}_{\mathbf{m},\omega}^t$ contains
 128 vertices $i \in I$ that satisfy $\omega_i \leq t$ and for every $i, j \in I$, the edge between the vertices i and
 129 j is in $\mathcal{G}_{\mathbf{m},\omega}^t$ when $t_{i,j} \leq t$. Indeed, $\mathcal{G}_{\mathbf{m},\omega}^t$ and $V_{\mathbf{m},\omega}^t$ have the same number of connected
 130 components, and m_i and m_j are in the same connected component in $V_{\mathbf{m},\omega}^t$ if and only if i
 131 and j are in the same connected component in $\mathcal{G}_{\mathbf{m},\omega}^t$.

132 More generally, the topological information of $V_{\mathbf{m},\omega}^t$ (number of connected components,
 133 loops, voids etc.) can be encoded in the weighted Čech complex $\text{Cech}_{\mathbf{m},\omega}(t)$, defined as
 134 the nerve of the union of balls $(\mathcal{B}_i^t)_{i \in I}$: $\text{Cech}_{\mathbf{m},\omega}(t) = \{\sigma \subset I \mid \bigcap_{i \in \sigma} \mathcal{B}_i^t \neq \emptyset\}$, [1, 7, 3].
 135 According to the Nerve Lemma [20, Corollary 4G.3], any sublevel set $V_{\mathbf{m},\omega}^t$ is homotopic
 136 to $\text{Cech}_{\mathbf{m},\omega}(t)$ and thus contains the same topological information. For computational
 137 reasons, the weighted Vietoris-Rips filtration is frequently considered as a provably good
 138 surrogate for the weighted Čech filtration $(\text{Cech}_{\mathbf{m},\omega}(t))_{t \in T}$. The weighted Vietoris-Rips
 139 complex $\text{VR}_{\mathbf{m},\omega}(t)$ is the flag complex of $\mathcal{G}_{\mathbf{m},\omega}^t$ ($\mathcal{G}_{\mathbf{m},\omega}^t$ is the 1-skeleton of the weighted Čech
 140 complex): $\text{VR}_{\mathbf{m},\omega}(t) = \{\sigma \subset I \mid \forall i, j \in \sigma, \mathcal{B}_i^t \cap \mathcal{B}_j^t \neq \emptyset\}$. Indeed, as a direct consequence of
 141 [3, Theorem 3.2] which is a generalization of the non-weighted case in [15, Theorem 2.5.], if
 142 the weights in ω are non-negative, then these two filtrations are interleaved:

$$143 \quad \forall 0 < t' \leq \frac{d+1}{2d}t, \text{VR}_{\mathbf{m},\omega}(t') \subset \text{Cech}_{\mathbf{m},\omega}(t) \subset \text{VR}_{\mathbf{m},\omega}(t). \quad (1)$$

144 We will see that these notions can all be extended to anisotropic power functions.

145 **2.2.2 The anisotropic case**

146 Consider an index set $I = \llbracket 1, c \rrbracket$, families of centers $\mathbf{m} = (m_i)_{i \in I}$ in \mathbb{R}^d , of weights $\omega = (\omega_i)_{i \in I}$
 147 in \mathbb{R} and of matrices $\Sigma = (\Sigma_i)_{i \in I}$ in \mathcal{M}_d , the set of definite positive symmetric matrices.
 148 An anisotropic power function is a function $f_{\mathbf{m},\omega,\Sigma} : \mathbb{R}^d \rightarrow \mathbb{R}$ defined from I , \mathbf{m} , ω and Σ
 149 by $f_{\mathbf{m},\omega,\Sigma} : x \mapsto \min_{i \in I} \|x - m_i\|_{\Sigma_i^{-1}}^2 + \omega_i$. For any matrix $\Sigma \in \mathcal{M}_d$, $\|\cdot\|_{\Sigma^{-1}}$ denotes the
 150 Σ -Mahalanobis distance, defined for every $x \in \mathbb{R}^d$ by $\|x\|_{\Sigma^{-1}}^2 = x^T \Sigma^{-1} x$. The sublevel sets of
 151 $f_{\mathbf{m},\omega,\Sigma}$, $V_{\mathbf{m},\omega,\Sigma}^t = f_{\mathbf{m},\omega,\Sigma}^{-1}((-\infty, t])$, are unions of at most c ellipsoids $\mathcal{E}_i^t = \overline{B}_{\Sigma_i}(m_i, \sqrt{t - \omega_i})$
 152 with $\overline{B}_{\Sigma}(m, r) = \{x \in \mathbb{R}^d \mid \|x - m\|_{\Sigma^{-1}} \leq r\}$. Again, \mathcal{E}_i^t is empty for $t < \omega_i$. The value of
 153 $t_{i,j}$ above which the two ellipsoids \mathcal{E}_i^t and \mathcal{E}_j^t intersect is given below.

154 **► Proposition 1.** *Consider two ellipsoids $\mathcal{E}_i^t = \overline{B}_{\Sigma_i}(m_i, \sqrt{t - \omega_i})$ and $\mathcal{E}_j^t = \overline{B}_{\Sigma_j}(m_j, \sqrt{t - \omega_j})$
 155 with $\omega_i \leq \omega_j$ in \mathbb{R} , m_i and m_j in \mathbb{R}^d , $\Sigma_i = P_i D_i P_i^T$ and $\Sigma_j = P_j D_j P_j^T$ in \mathcal{M}_d , with two
 156 positive diagonal matrices D_i and D_j , and two orthogonal matrices P_i and P_j accordingly to
 157 the spectral theorem. Set $\tilde{\Sigma} = \sqrt{D_i} P_i^T \Sigma_j^{-1} P_i \sqrt{D_i}$, the orthogonal and diagonal matrices \tilde{P} and
 158 $\tilde{D} = \text{diag}(\lambda_1, \lambda_2, \dots, \lambda_d)$ such that $\tilde{\Sigma} = \tilde{P} \tilde{D} \tilde{P}^T$ and $\tilde{m} = \tilde{P}^T \sqrt{D_i^{-1}} P_i^T (m_j - m_i)$. Ellipsoids
 159 \mathcal{E}_i^t and \mathcal{E}_j^t intersect if and only if $t \geq t_{i,j}$, with $t_{i,j}$ defined as follows. If $\|\tilde{m}\| \leq \sqrt{\omega_j - \omega_i}$,*

160 then $t_{i,j} = \omega_j$. If $\|\tilde{m}\| > \sqrt{\omega_j - \omega_i}$, then $t_{i,j} = \omega_j + \sum_{k=1}^d \left(\frac{\lambda \tilde{m}_k}{\lambda + \lambda_k} \right)^2 \lambda_k$, where λ is the unique
 161 solution of the following equation:

$$162 \quad \sum_{k=1}^d \frac{\lambda_k - \lambda^2}{(\lambda + \lambda_k)^2} \lambda_k \tilde{m}_k^2 = \omega_j - \omega_i. \quad (2)$$

163 The proof of Proposition 1 is to be found in Section A. It is based on the fact that the
 164 ellipsoids \mathcal{E}_i^t and \mathcal{E}_j^t are tangent at their intersection point, and the corresponding gradients are
 165 orthogonal. In the context of isotropy (i.e. for $\Sigma_i = \Sigma_j = I_d$, the identity matrix of \mathbb{R}^d), $\tilde{m} =$
 166 $m_j - m_i$ and when $\|m_j - m_i\| > \sqrt{\omega_j - \omega_i}$, Equation (2) has a unique positive solution given by
 167 $\lambda = \frac{\omega_i - \omega_j + \|m_j - m_i\|^2}{\omega_j - \omega_i + \|m_j - m_i\|^2}$. The merging time is then $t_{i,j} = \frac{(\omega_j - \omega_i)^2 + 2(\omega_j + \omega_i)\|m_j - m_i\|^2 + \|m_j - m_i\|^4}{4\|m_j - m_i\|^2}$.
 168 This is consistent with the previous section. Furthermore, in the context of anisotropy, it is
 169 also possible to define $\mathcal{G}_{\mathbf{m},\omega,\Sigma}^t$, $\text{Cech}_{\mathbf{m},\omega,\Sigma}(t)$ and $\text{VR}_{\mathbf{m},\omega,\Sigma}(t)$, the anisotropic counterparts
 170 of $\mathcal{G}_{\mathbf{m},\omega}^t$, $\text{Cech}_{\mathbf{m},\omega}(t)$ and $\text{VR}_{\mathbf{m},\omega}(t)$. Note that the nerve lemma still applies, since unions of
 171 ellipsoids are contractible. Although this paper is mostly based on the study of connected
 172 components for the purpose of clustering, anisotropic weighted Čech and Vietoris-Rips
 173 filtrations are primordial to have a tractable good estimation of the topology of compact
 174 sets from good approximations of the compact sets with finite unions of ellipsoids. In fact,
 175 just as their isotropic counterparts (1), these filtrations are interleaved, provided that the
 176 eigenvalues of the matrices in Σ belong to some interval $[\lambda_{\min}, \lambda_{\max}]$ with λ_{\min} positive.

177 ► **Proposition 2.** *If ω is a set on non-negative weights in \mathbb{R} . Then, for every $t > 0$ and*
 178 *$0 < t' \leq \frac{\lambda_{\min}}{\lambda_{\max}} \frac{d+1}{2d} t$, we get that*

$$179 \quad \text{VR}_{\mathbf{m},\omega,\Sigma}(t') \subset \text{Cech}_{\mathbf{m},\omega,\Sigma}(t) \subset \text{VR}_{\mathbf{m},\omega,\Sigma}(t). \quad (3)$$

180 The condition of non-negative weights is not too restrictive since for general weights, the
 181 proposition can be applied with $\omega - \min_{i \in I} \omega_i$, $t - \min_{i \in I} \omega_i$ and $t' - \min_{i \in I} \omega_i$ instead of
 182 ω , t and t' . So, for general weights, the condition is actually: $\min_{i \in I} \omega_i < t' \leq \frac{\lambda_{\min}}{\lambda_{\max}} \frac{d+1}{2d} t +$
 183 $\left(1 - \frac{\lambda_{\min}}{\lambda_{\max}} \frac{d+1}{2d}\right) \min_{i \in I} \omega_i$. As noted in [15], when $\lambda_{\min} = \lambda_{\max}$ and the weights in ω are null,
 184 the term $\frac{\lambda_{\min}}{\lambda_{\max}} \frac{d+1}{2d}$ is optimal. Equality for the left inclusion is obtained for \mathbf{m} , the vertices
 185 of a regular d -simplex. The proof of Proposition 2 is available in Section B.

186 Often, less ellipsoids than balls are required to describe a compact set, for a fixed level
 187 of precision (e.g. for the Hausdorff distance). For instance, a segment in \mathbb{R}^2 , and more
 188 generally, any d' -dimensional submanifold in \mathbb{R}^d , with $d' < d$. For this reason, isotropic
 189 Čech and Vietoris-Rips filtrations are pertinent tools to compute and store the topological
 190 information about the compact set \mathcal{X} efficiently. The requisite condition is that we dispose
 191 of an anisotropic power function that is a good approximation of the squared Euclidean
 192 distance to the compact set \mathcal{X} , $d_{\mathcal{X}}^2$. Such examples of functions are discussed in the sequel.

193 2.3 Examples of filtrations based on robust power functions

194 2.3.1 Isotropic robust power functions

195 Set \mathbb{X} , a set of n points generated on the neighborhood of a compact subset \mathcal{X} of \mathbb{R}^d . In
 196 order to face the non robustness of the distance function to \mathbb{X} , Chazal et al. have introduced
 197 the notion of distance-to-measure (DTM), in [11]. The distance-to-measure function is a
 198 counterpart of the distance to the compact set \mathbb{X} that is robust to noise and in particular
 199 to outliers. Its robustness depends on some parameter $k \in \llbracket 1, n \rrbracket$, the number of nearest-
 200 neighbors X^1, X^2, \dots, X^k of x in \mathbb{X} that are used to estimate the distance function of

XX:6 Robust anisotropic power distance filtrations

201 x to \mathbb{X} . The distance-to-measure function $d_{\mathbb{X},k}$ with parameter $k \in \llbracket 1, n \rrbracket$ is defined by
 202 $d_{\mathbb{X},k}^2 : x \mapsto \frac{1}{k} \sum_{i=1}^k \|x - X^i\|^2 = \|x - m_{x,k}\|^2 + v_{x,k}$ with $m_{x,k} = \sum_{i=1}^k X^i$, the mean of the k
 203 nearest neighbours of x in \mathbb{X} and $v_{x,k} = \frac{1}{k} \sum_{i=1}^k \|X^i - m_{x,k}\|^2$ their variance. Note that $d_{\mathbb{X},1}$
 204 coincides with $d_{\mathbb{X}}$ and suffers from non robustness, whereas $d_{\mathbb{X},n}(x)$ is the distance of x to
 205 the barycenter of the point cloud \mathbb{X} , up to some factor, which is very robust, but very poor
 206 in terms of topological information. Guibas et al. noted in [18] that the distance-to-measure
 207 is actually a weighted power function, according to its alternative expression:

$$208 \quad d_{\mathbb{X},k}^2(x) = \inf_{y \in \mathbb{R}^d} \|x - m_{y,k}\|^2 + v_{y,k}. \quad (4)$$

209 Indeed, the mean distance between x and its k nearest neighbors is necessarily smaller than
 210 the mean distance between x and the k neighbors of any other point $y \in \mathbb{R}^d$. This infimum is
 211 actually a minimum over a set of c points $\mathbf{y} = (y_i)_{i \in \llbracket 1, c \rrbracket}$ in \mathbb{R}^d , with c of order $\binom{n}{k}$. In [18],
 212 Guibas et al. introduced and studied the k -witnessed distance, a power approximation of the
 213 DTM defined by replacing \mathbb{R}^d by the sample \mathbb{X} in (4). The sublevel sets of the k -witnessed
 214 distance are unions of n balls. An approximation of the DTM with c (possibly much smaller
 215 than n) balls has been introduced in [6]. This approximation, the c -PDTM, consists in
 216 replacing \mathbb{R}^d by a set $\mathbf{y}_{c,k}$ of c points in \mathbb{R}^d . This set $\mathbf{y}_{c,k}$ is obtained as a minimum for a
 217 “k-means”-type criterion [24] : $\mathbf{y}_{c,k} \in \arg \min_{\mathbf{y} \mid |\mathbf{y}|=c} \sum_{i=1}^n \min_{y \in \mathbf{y}} \|X_i - m_{y,k}\|^2 + v_{y,k}$, with
 218 $|\mathbf{y}|$, the cardinal of \mathbf{y} . Morally, the optimal set $\mathbf{y}_{c,k}$ is chosen such that on average, on \mathbb{X} , the
 219 power function $x \mapsto \min_{y \in \mathbf{y}} \|x - m_{y,k}\|^2 + v_{y,k}$ takes small values. Note that the graph of
 220 such a function is necessarily above the graph of the DTM. According to [6], for a sample on
 221 a regular d' -manifold, c can be chosen of order $n^{\frac{d'}{d'+4}}$, which is much smaller than n . For
 222 such a c , the m -PDTM is a good approximation of the distance to \mathcal{X} , in spite of noise.

223 2.3.2 An anisotropic robust power function

224 An anisotropic version of the k -PDTM has been introduced in [4], the k -power likelihood to
 225 measure (k -PLM). It consists in replacing Euclidean norms with Mahalanobis norms. For
 226 every $x \in \mathbb{R}^d$ and $\Sigma \in \mathcal{M}_d$, set X^1, X^2, \dots, X^k the k -nearest neighbors of x in \mathbb{X} , for the
 227 Σ^{-1} -Mahalanobis norm: $\|X^i - x\|_{\Sigma^{-1}} \leq \|X^j - x\|_{\Sigma^{-1}}$ for every $i \leq j$. Denote by $m_{x,\Sigma,k}$ their
 228 mean, and by $v_{x,\Sigma,k} = \frac{1}{k} \sum_{i=1}^k \|X^i - m_{x,\Sigma,k}\|_{\Sigma^{-1}}^2$ their variance, relative to the Σ -Mahalanobis
 229 norm. Set $\boldsymbol{\theta}_{c,k}$, a family of c pairs $(y, \Sigma) \in \mathbb{R}^d \times \mathcal{M}_d$ that minimizes (or which criterion is as
 230 close as possible to the optimal criterion, in case of non existence of a minimum) the following
 231 “k-means”-type criterion $R_{c,k}$ among all $\boldsymbol{\theta}$ s of cardinal c : $R_{c,k}(\boldsymbol{\theta}) = \sum_{i=1}^n \min_{(y,\Sigma) \in \boldsymbol{\theta}} \|X_i -$
 232 $m_{y,\Sigma,k}\|_{\Sigma^{-1}}^2 + v_{y,\Sigma,k} + \log(\det(\Sigma))$. The term $\log(\det(\Sigma))$ prevents optimal covariances matrices
 233 to be degenerated, with Σ^{-1} going to 0. In some sense, minimizing such a criterion boils
 234 down to fit Gaussian distributions to the data set \mathbb{X} , at best. The c -PLM is the power
 235 function defined from $\boldsymbol{\theta}_{c,k}$ by: $x \mapsto \min_{(y,\Sigma) \in \boldsymbol{\theta}_{c,k}} \|x - m_{y,\Sigma,k}\|_{\Sigma^{-1}}^2 + v_{y,\Sigma,k}$. A modification
 236 of the criterion $R_{c,k}$ has been introduced in [4], to remove some datapoints ($|\mathbb{X}| - sig$
 237 for some parameter sig), when \mathbb{X} is corrupted with outliers. The criterion is given by
 238 $R_{c,k,sig}(\boldsymbol{\theta}) = \min_{(i_1, i_2, \dots, i_{sig}) \in \llbracket 1, c \rrbracket} \sum_{j=1}^{sig} \min_{(y,\Sigma) \in \boldsymbol{\theta}} \|X_{i_j} - m_{y,\Sigma,k}\|_{\Sigma^{-1}}^2 + v_{y,\Sigma,k} + \log(\det(\Sigma))$.

239 In [4], it has been proved that Lloyd-type iterative algorithms [22] provide a local minimum
 240 $\tilde{\boldsymbol{\theta}}_{c,k}$ for the criterion $R_{c,k}$, but also a local minimum $\tilde{\boldsymbol{\theta}}_{c,k,sig}$ for the criterion $R_{c,k,sig}$. These
 241 algorithms run in $n \log(n)c \times it$, with it the number of iterations of the algorithm. Roughly,
 242 the algorithms consists, given $\boldsymbol{\theta} = (\mathbf{y}, \boldsymbol{\Sigma})$, to split the space \mathbb{R}^d into weighted $\boldsymbol{\Sigma}$ -curved
 243 Voronoi cells, to replace centers \mathbf{y} by the barycenters of the cells, and to update the set of
 244 matrices $\boldsymbol{\Sigma}$ by some close formula that depends on the points on the cell and the points in

245 the ellipsoids. To compute $\tilde{\theta}_{c,k,sig}$, a trimming step is added at each iteration. Actually,
 246 for the purpose of clustering, disposing of a local minimum is enough, as enhanced in the
 247 numerical illustration section. Indeed, the parameter *Threshold* of our clustering algorithm,
 248 Algorithm 1, will make possible the removing of bad centers in $\tilde{\theta}_{c,k}$ or in $\tilde{\theta}_{c,k,sig}$.

249 **3 Persistence-based clustering from power-functions-based filtrations**

250 **3.1 Persistence for power-functions-based filtrations**

251 Set $f_{\mathbf{m},\omega,\Sigma} : x \in \mathbb{R}^d \mapsto \min_{i \in I} \|x - m_i\|_{\Sigma_i^{-1}}^2 + \omega_i$, an anisotropic power-function indexed
 252 by a finite set $I = \llbracket 1, c \rrbracket$ and such that the ω_i s are sorted in non-decreasing order. As
 253 above-mentioned, the sublevel sets $V^t = f_{\mathbf{m},\omega,\Sigma}^{-1}((-\infty, t])$ are unions of at most c ellipsoids
 254 $\mathcal{E}_i^t = B_{\Sigma_i}(m_i, \sqrt{t - \omega_i})$, non empty as soon as $t \geq \omega_i$. In particular, each sublevel set of
 255 $f_{\mathbf{m},\omega,\Sigma}$ contains at most c connected components. Each connected component of V^t , V_i^t can
 256 be indexed by some element $i \in I$: the smallest i such that m_i belongs to the component.
 257 With a language abuse, we call connected component V_i , the family of connected components
 258 $(V_i^t)_{t \in T}$ that gets born at time $t = b_i = \omega_i$ and dies at a time $t = d_i$ when V_i^t merges
 259 with another connected component V_j^t for some $j \leq i$. Note that $d_1 = \infty$. The lifetime
 260 of the component V_i^t , $d_i - b_i$, is called persistence or prominence of the component i . We
 261 can represent this merging information with a barcode or a dendrogram. In these two
 262 representations, each line is associated to a component V_i , with length $d_i - b_i$, and begins at
 263 the height given by b_i . The dendrogram is obtained from the barcode by linking the bars
 264 associated to merging components, at a height given by the merging time.

265 When \mathbf{m} is a point set \mathbb{X} , $\Sigma_i = I_d$ and $\omega_i = 0$ for every i , clustering points accordingly
 266 to the connected components of V^t boils down to the classical single-linkage clustering
 267 procedure, with $t > 0$, calibrated in accordance with the dendrogram. Such a clustering
 268 procedure suffers from non-robustness to outliers. In this paper, we consider an adjacent
 269 procedure, based on the idea of ToMATo algorithm [12], that consists in taking into account
 270 the prominence of components. To be precise, in the clustering scheme, we decide that a
 271 component V_i cannot merge with another precise component V_j at a time t larger than $\omega_i + Stop$,
 272 for some fixed threshold *Stop*. In other words, components with large prominence will never
 273 die in this clustering procedure. This is the purpose of Algorithm 1 in the next section.

274 In order to better visualize the prominence of the components, we represent their lifetimes
 275 in a persistence diagram. A persistence diagram is a multiset of points $(b_i, d_i) \in \mathbb{R}^2$ that
 276 lie above the diagonal $b = d$. Each point (b_i, d_i) is associated to a connected component V_i .
 277 The notion of persistence diagram was actually introduced by Edelsbrunner et al. in [16], in
 278 the broader framework of homology, and allows to compute lifetimes of additional features
 279 such as loops, voids etc. It is defined for filtrations that are regular enough, on triangulable
 280 spaces such as \mathbb{R}^d . The proper notion of regularity is the notion of q -tameness [9]. In [7,
 281 Proposition 3.5], Buchet et al. proved that the distance-to-measure function is q -tame. The
 282 proof of [7] can be straightforwardly adjusted for the distance function to a compact set
 283 and most importantly, for anisotropic power functions, provided that the eigenvalues of the
 284 matrices Σ_i are all positive. A sketch of proof is given in the Appendix, in Section C.

285 Since distance to compact sets, distance-to-measure and anisotropic power functions
 286 are q -tame, the persistence diagrams associated to their filtrations are well defined. They
 287 can be compared with the bottleneck distance, a distance between two diagrams D and D'
 288 defined by the minimal value of $\max_{x \in D, y \in D'} |y - \phi(x)|_\infty$ among functions ϕ that pair points
 289 in D with points in D' , with some points potentially paired to diagonal points. Diagrams
 290 associated to interleaved filtrations are close, according to the following theorem.

291 ► **Theorem 3** (Stability of persistence diagrams [9, 10, 13]). *If two filtrations V and W are*
 292 *q -tame and ϵ -interleaved, then the persistence diagrams of these filtrations are ϵ -close in*
 293 *bottleneck distance.*

294 According to Proposition 3, the persistence diagram of any anisotropic power function
 295 $f_{\mathbf{m},\omega,\Sigma}$ that is $\epsilon - \|\cdot\|_\infty$ close to $d_{\mathcal{X}}$, will be ϵ -bottleneck close to the persistence diagram
 296 of the sublevel sets of $d_{\mathcal{X}}$. Consequently, prominence of the connected components of \mathcal{X}
 297 can be deduced from the diagram associated to $f_{\mathbf{m},\omega,\Sigma}$, for ϵ small enough. This bottleneck
 298 closeness occurs with large probability for a regular manifold \mathcal{X} for the k -PDTM built from
 299 a noisy sample from \mathcal{X} , according to [6]. No such result has been proved yet for the k -PLM
 300 function. Anyway, intuitively, its sublevel sets are good approximations of the manifold
 301 \mathcal{X} , with the advantage that they are made of less ellipsoids, and that these ellipsoids are
 302 oriented accordingly to the manifold, i.e. with large eigenvalues on the tangent space and
 303 small eigenvalues on its orthogonal. This will be confirmed in the numerical illustrations
 304 section.

305 By construction, the persistence diagram (for connected components) associated to the
 306 filtration of the sublevel sets of $f_{\mathbf{m},\omega,\Sigma}$ coincides with the persistence diagram associated to
 307 the anisotropic weighted Čech complex $\text{Cech}(f_{\mathbf{m},\omega,\Sigma})$. Consequently, we can forget about
 308 the ellipsoids and focus on the simplicial complex filtration, which can be computed and
 309 stored efficiently, in a $c \times c$ matrix $\text{Mat} = (t_{i,j})_{i,j \in I}$. Such a matrix would contain the times
 310 of appearance of vertices and of merging of connected components in $\text{Cech}(f_{\mathbf{m},\omega,\Sigma})$. The
 311 clustering scheme of this paper exposed just below is based on such a merging matrix Mat .

312 3.2 An algorithm for persistence-based clustering

313 Consider $(\mathcal{G}^t)_{t \in \mathbb{R}}$ a filtration of sub-graphs of \mathcal{G} , a graph with c nodes. Based on this filtration,
 314 we define an algorithm, strongly inspired from the ToMATo algorithm [12]. The clustering
 315 scheme is guided by the persistence of the connected components in $(\mathcal{G}^t)_{t \in \mathbb{R}}$, and preserves
 316 components with large prominence. We assume that the nodes of \mathcal{G} are labeled such that
 317 the node labeled i gets born before the node labeled j , when $i \leq j$. The procedure is as
 318 follows. A connected component gets born when an edge gets born, with the same label.
 319 A component will change of label at each time t for which it merges with a component
 320 with smaller label in \mathcal{G}^t , unless its prominence is larger than some parameter $Stop$. The
 321 prominence of an edge or a component is defined as the lifetime of the component in the
 322 filtration. In other words, the prominence of an edge is the elapsed time between the birth
 323 of the edge and the time t such that an edge with smaller index is present in its connected
 324 component in \mathcal{G}^t . The resulting clustering is given by the label of the edges at time $t = +\infty$.
 325 It will contain exactly labels of edges with a prominence larger than $Stop$. In this clustering
 326 scheme, it is also possible to decide that edges born after some time parameter $Threshold$
 327 are not relevant, and decide to remove them from the clustering procedure.

328 This procedure is implemented in Algorithm 1. This algorithm requires a merging matrix
 329 $\text{Mat} = (t_{i,j})_{i,j \in I}$, with $I = \llbracket 1, c \rrbracket$. We define its coefficients by $t_{i,i}$, the birth time of the
 330 node i in the filtration $(\mathcal{G}^t)_{t \in \mathbb{R}}$; for $i > j$, $t_{i,j}$ the birth time of the edge $[i, j]$ and for $i < j$,
 331 $t_{i,j} = \infty$. The vector $Color$ contains the resulting clustering, the vector $Birth$, the birth

332 time of the components and *Death* their death time. Note that *Death*[1] will always be $+\infty$.

333 **Algorithm 1** Persistence-based Clustering Algorithm

```

333 Data: Mat, Threshold, Stop
334 Result: Color, Birth, Death
335 Initialization ;
336  $c \leftarrow \max\{i \mid \text{Mat}[i,i] \leq \text{Threshold}\}$  ;
337  $\text{Mat} \leftarrow \text{Mat}[1:c,1:c]$  ;
338  $\text{Birth} \leftarrow [\text{Mat}[i,i] \text{ for } i \text{ in } 1:c]$  ;
339  $\text{Death} \leftarrow [\infty \text{ for } i \text{ in } 1:c]$  ;
340  $\text{indice} \leftarrow 1$  ;
341  $I \leftarrow 1$  ;
342  $\text{time} \leftarrow \text{Mat}[I,I]$  ;
343  $\text{Color} \leftarrow []$  ;
344 while  $\text{time} < \infty$  do
345   if  $\text{time} = \text{Mat}[I,I]$  then
346     Component  $I$  appears ;
347      $\text{indice} \leftarrow \text{indice} + 1$  ;
348      $\text{Mat}[I,I] \leftarrow \infty$  ;
349      $\text{Color}[I] \leftarrow I$  ;
350   else
351      $(\text{col\_max}, \text{col\_min}) \leftarrow (\max(\text{Color}[I], \text{Color}[J]), \min(\text{Color}[I], \text{Color}[J]))$ ;
352     if  $\text{time} - \text{Birth}[\text{col\_max}] \leq \text{Stop}$  then
353       Components  $\text{col\_max}$  and  $\text{col\_min}$  merge ;
354       Replace all entries  $\text{col\_max}$  by  $\text{col\_min}$  in  $\text{Color}$  ;
355        $\text{Death}[\text{col\_max}] \leftarrow \text{time}$  ;
356     else
357       Component  $\text{col\_max}$  will never die ;
358     end
359      $\text{Mat}[i,j] \leftarrow \infty$  for every  $i, j \leq \min(\text{indice}, c)$  such that
360      $(\text{Color}[i], \text{Color}[j]) \in \{(\text{col\_min}, \text{col\_max}), (\text{col\_max}, \text{col\_min})\}$ ;
361   end
362    $I, J \leftarrow \arg \min_{i,j \leq \min(\text{indice}, c)} \text{Mat}[i,j]$  ;
363    $\text{time} \leftarrow \text{Mat}[I,J]$ 
364 end

```

334 In practice, just as Chazal et al. [12], we recommend to run Algorithm 1 twice. A first
335 time to measure the prominence of the components, and a second time with the parameters
336 *Stop* and *Threshold* chosen in accordance with the prominence values. The selection of the
337 parameters *Stop* and *Threshold* for the second passage can be done from the persistence dia-
338 gram of the connected components given by the set of points $(\text{Birth}[i], \text{Death}[i])$, with *Birth*
339 and *Death* given by the first passage. If γ points are well-separated from the other points in
340 the diagram, then we will select *Stop* such that there are γ clusters after the second passage.
341 By well-separated, we mean that for some $b_{\min} < b_{\max}$, these points are above the line
342 $d = b + b_{\max}$ whereas the other points are below the line $d = b + b_{\min}$, with $b_{\max} - b_{\min}$ large
343 enough. In this context, we set $\text{Stop} = \frac{b_{\min} + b_{\max}}{2}$. In addition, if there is a group of points in
344 the diagram, that are below the line $d = b + b_{\min}$, on the very right of the diagram, then these
345 points correspond to components that can be considered as noise and can be removed. In
346 these points are separated to the others with a line $b = b_s$, then we can chose $\text{Threshold} = d_s$.

347

348 When the graph filtration $(\mathcal{G}_t)_{t \in T}$ corresponds to the filtration of the sublevel sets of
 349 some power function $f_{\mathbf{m}, \boldsymbol{\omega}, \boldsymbol{\Sigma}}$, the matrix Mat has coefficients given by $t_{i,i} = \omega_i$ and for $i > j$,
 350 $t_{i,j}$ given by Proposition 1: the intersecting time of the ellipsoids \mathcal{E}_i^t and \mathcal{E}_j^t . Recall that when
 351 the matrices are all equal to the identity, then $t_{i,j} = \frac{(\omega_j - \omega_i)^2 + 2(\omega_j + \omega_i)\|m_j - m_i\|^2 + \|m_j - m_i\|^4}{4\|m_j - m_i\|^2}$.

352 Giving a sense to a minimal prominence Stop^* for a good clustering is possible for distance
 353 functions. For instance, for the sublevel-sets filtration of $d_{\mathcal{X}}$, Stop^* is given by half of the
 354 minimal distance between two disjointed connected components of \mathcal{X} . Consequently, for any
 355 $\epsilon - \|\cdot\|_{\infty}$ -close approximation of $d_{\mathcal{X}}$, taking $\text{Stop} = \text{Stop}^* - \epsilon$ leads to a perfect clustering,
 356 provided that $2\epsilon < \text{Stop}^*$. A power function is not homogeneous to a distance function, but
 357 to the square of a distance function. Therefore, for positive weights $\boldsymbol{\omega}$, it could be more
 358 appropriate to consider the filtration of sublevel sets of $\sqrt{f_{\mathbf{m}, \boldsymbol{\omega}, \boldsymbol{\Sigma}}}$ instead of $f_{\mathbf{m}, \boldsymbol{\omega}, \boldsymbol{\Sigma}}$.

359 The power function used to approximate $d_{\mathcal{X}}^2$ from a set of points \mathbb{X} is based on a family \mathbf{m} of
 360 c centers, possibly different from \mathbb{X} . Algorithm 1 assigns labels to the centers in \mathbf{m} . Clustering
 361 points in \mathbb{X} is made accordingly to these labels and to the Voronoi decomposition of \mathbb{R}^d , based
 362 on \mathbf{m} , $\boldsymbol{\omega}$ and $\boldsymbol{\Sigma}$: $x \in \mathbb{X}$ has the same label as m_i if $\|x - m_i\|_{\Sigma_i^{-1}}^2 + \omega_i \leq \|x - m_j\|_{\Sigma_j^{-1}}^2 + \omega_j$
 363 for every j such that m_j and m_i are not born after parameter *Threshold*. This parameter
 364 removes centers in \mathbf{m} . To deal with outliers, it is also possible to remove points (i.e. assign a
 365 label 0) in the sample \mathbb{X} . For this, we remove points x for which $f_{\mathbf{m}^*, \boldsymbol{\omega}^*, \boldsymbol{\Sigma}^*}(x)$ is the largest,
 366 with $\mathbf{m}^*, \boldsymbol{\omega}^*, \boldsymbol{\Sigma}^*$ the parameters that were not removed by the algorithm. Indeed, $f_{\mathbf{m}^*, \boldsymbol{\omega}^*, \boldsymbol{\Sigma}^*}$
 367 is supposed to approximate $d_{\mathcal{X}}^2$. For the k -PLM function, the parameter *Threshold* is
 368 primordial. Indeed, algorithm that computes $\boldsymbol{\theta}_{c,k}$ actually converges to a local minimum of
 369 the criterion. Consequently, it happens that ellipsoids \mathcal{E}_i are actually far from the support.
 370 For such ellipsoids, the weight ω_i will necessarily be large, from its definition. Since other
 371 ellipsoids should have small weight, such bad ellipsoids will be removed for a suitable choice
 372 of parameter *Threshold* that should be calibrated from the persistence diagram.

373 3.3 Connection to other persistence-based clustering methods

374 In the sequel, we display different graph filtrations, to be used for persistence-based clustering,
 375 with Algorithm 1. For each of these filtrations, we give a summarize of the corresponding
 376 matrices Mat , in Table 1, with the convention that $t_{i,i} \leq t_{j,j}$ when $i \leq j$.

377 ToMATo Algorithm [12] is based on a graph filtration that is constructed from a graph \mathcal{G}
 378 and some function f defined on the nodes of the graph \mathcal{G} . Morally, \mathcal{G}^t is the sub-graph of \mathcal{G}
 379 that contains the node i is and only if $f(i) \leq t$, and the edge $[i, j]$ if and only if i and j are
 380 in \mathcal{G}^t . Given a set of points \mathbb{X} in \mathbb{R}^d , Chazal et al. have mostly studied this method for \mathcal{G} , a
 381 Rips graph of \mathbb{X} , and for $f(i)$, the DTM to \mathbb{X} at X_i .

382 The DTM-filtration [1] corresponds to the 1-skeleton of the nerve of the union of balls
 383 $(\bigcup_{x \in \mathbb{X}} \bar{B}(x, r_t(x)))_{t > 0}$ with $r_t(x) = -\infty$ for $t < d_{\mathbb{X},k}(x)$ and $r_t(x) = (t^p - d_{\mathbb{X},k}^p(x))^{\frac{1}{p}}$ for
 384 $t \geq d_{\mathbb{X},k}(x)$, for some $p \geq 1$ and with the convention that $\bar{B}(x, -\infty)$ is empty. In Table 1, we
 385 give the coefficients for $p = 1$. The DTM-filtration with $p = 2$ was actually introduced in [7],
 386 leading to what we call Power filtration, which is actually the sublevel-sets filtration of the
 387 square of a power distance. We also consider additional power-functions-based filtrations,
 388 from the k -witnessed distance [18], the c -PDTM [6] and the c -PLM.

389 **Table 1** Coefficients of Mat for the different methods, for $f = d_{\mathbb{X},k}$, the DTM to \mathbb{X} with number
390 of nearest neighbors parameter k .

Method	$t_{i,i}$	$t_{i,j}$ for $i < j$
ToMATo	$f(i)$	$\max(f(i), f(j))(\mathbb{1}_{[i,j] \in \mathcal{G}})^{-1}$
DTM-filtration	$f(i)$	$\left(\frac{\ X_i - X_j\ + f(i) + f(j)}{2} \right) \mathbb{1}_{\ X_i - X_j\ > f(i) - f(j) } + f(i) \mathbb{1}_{f(i) - f(j) \geq \ X_i - X_j\ }$
$f_{\mathbf{m},\omega}$	ω_i	$\frac{(\omega_j - \omega_i)^2 + 2(\omega_j + \omega_i)\ m_j - m_i\ ^2 + \ m_j - m_i\ ^4}{4\ m_j - m_i\ ^2}$
$\sqrt{f_{\mathbf{m},\omega}}$	$\sqrt{\omega_i}$	$\sqrt{\frac{(\omega_j - \omega_i)^2 + 2(\omega_j + \omega_i)\ m_j - m_i\ ^2 + \ m_j - m_i\ ^4}{4\ m_j - m_i\ ^2}}$
$f_{\mathbf{m},\omega,\Sigma}$	ω_i	Given by Proposition 1
Power filtration		$\sqrt{f_{\mathbf{m},\omega}}$ with $\mathbf{m} = \mathbb{X}$ and $\omega = (f^2(x))_{x \in \mathbb{X}}$
Witnessed		$\sqrt{f_{\mathbf{m},\omega}}$ with $(\mathbf{m}, \omega) = (m_{x,k}, v_{x,k})_{x \in \mathbb{X}}$
c -PDTM		$f_{\mathbf{m},\omega}$ with $(\mathbf{m}, \omega) = (m_{y,k}, v_{y,k})_{y \in \mathcal{Y}_{c,k}}$
c -PLM		$f_{\mathbf{m},\omega,\Sigma}$ with $(\mathbf{m}, \omega, \Sigma) = (m_{y,\Sigma,k}, v_{y,\Sigma,k} + \log(\det(\Sigma)), \Sigma)_{(y,\Sigma) \in \theta_{c,k}}$

4 Numerical Illustrations

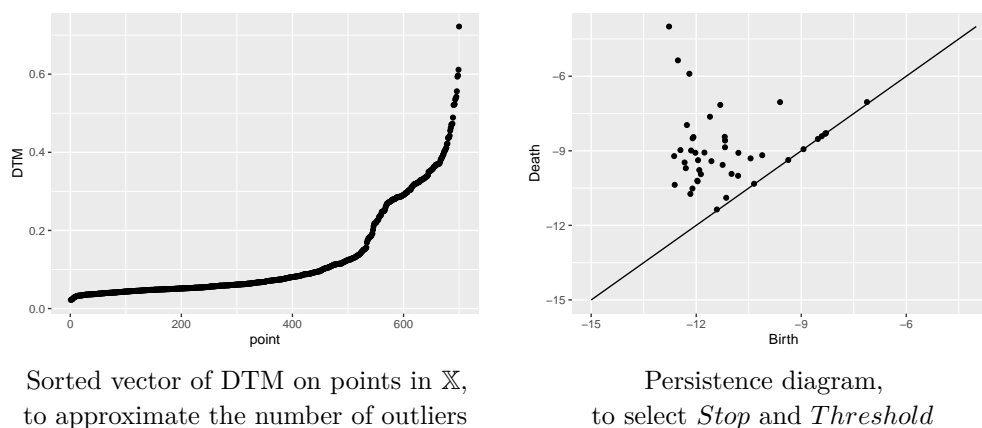
4.1 A complete illustration of the method

403 We consider a set of $N = 500$ points $\mathbb{Y} = (Y_i)_{i \in \llbracket 1, N \rrbracket}$ uniformly sampled on a set \mathcal{X} . We corrupt
404 this points with Gaussian noise : $\mathbb{Z} = (Z_i)_{i \in \llbracket 1, N \rrbracket}$, N variables of distribution $\mathcal{N}(0, \sigma^2)$, with
405 $\sigma = 0.02$. We observe a sample of N points $(X_i = Y_i + Z_i)_{i \in \llbracket 1, N \rrbracket}$, corrupted by $Nnoise = 200$
406 outliers $(X_i)_{i \in \llbracket N+1, N+Nnoise \rrbracket}$ uniformly sampled on the square $[-1.5, 2.5] \times [-1.5, 2.5]$. This
407 results in an $(N + Nnoise)$ -sample \mathbb{X} . In the sequel, we expose the clustering method based
408 on Algorithm 1 with the sublevel sets of an approximation of the c -PLM. The parameters are
409 set to $c = 50$ centers, $k = 10$ nearest neighbors, $sig = 520$ points to consider as signal, and
410 $it = 100$ iterations of the Lloyd-type algorithm to compute a local optimum of the criterion
411 $R_{c,k,sig}$, $\tilde{\theta}_{c,k,sig}$. We ran this Lloyd-type algorithm $ntimes = 10$ times and chose $\tilde{\theta}_{c,k,sig}$
412 that minimizes $R_{c,k,sig}$ among the 10 ones.

413 The criterion $sig = 520$ was chosen according to the following heuristic. Outliers have a
414 large DTM-value, in comparison to signal points. Therefore, we compute the vector of DTM
415 values, $[d_{\mathbb{X},k}(X_i), i \in \llbracket 1, N + Nnoise \rrbracket]$ and sort it by non-decreasing order. We represent
416 these sorted values and select sig as the point for which the curve's slope changes. The sorted
417 values are represented in Figure 1. Note that this step can also be performed by plotting the
418 values of $f_{\tilde{\theta}_{c,k}}$ for some $\tilde{\theta}_{c,k}$ obtained as a first approximation of the PLM, without trimming.

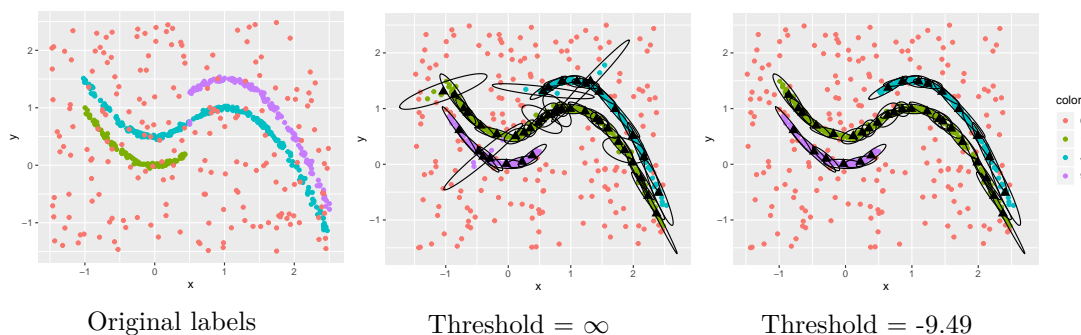
420 We ran Algorithm 1 a first time with the parameters $Threshold = \infty$ and $Stop = \infty$. We
421 obtained two vectors $Birth$ and $Death$ and represented the points $(Birth[i], Death[i])_{i \in \llbracket 1, c \rrbracket}$
422 in the persistence diagram, in Figure 1. From this diagram, it is possible to recover the
423 true number of clusters, 3. Indeed, 3 points are well-separated from the other ones with a
424 large band, parallel to the diagonal. Then, we selected $Stop = 5$ as the vertical distance
425 between a line parallel to the diagonal that separates the 3 points from the others, and
426 the diagonal. We ran Algorithm 1 a second time with the parameters $Threshold = \infty$ and
427 $Stop = 5$. The resulting clustering is represented in Figure 2. A sublevel set of the function
428 $f_{\tilde{\theta}_{c,k}}$ is represented by the union of ellipses in this figure. Note that some ellipses are badly
429 placed. Therefore, we use the parameter $Threshold$ to remove them. In the Figure 1, 7
430 points are on the right side, separated from the other points. We select $Threshold$ as the
431 abscissa of a vertical line separating these 7 points from the others: $Threshold = -9.49$.
432 We ran Algorithm 1 with the parameters $Threshold = -9.49$ and $Stop = 5$. The resulting
433 clustering is represented in Figure 2. The bad ellipses have been removed.

XX:12 Robust anisotropic power distance filtrations



419 ■ **Figure 1** Parameters selection heuristics

434 The color of a point x in Figure 2 is given by the label in $Color$ (returned by the algorithm)
 435 of its assigned center y_i in $\tilde{\theta}_{c,k}$ (the center y_i such that $f_{\tilde{\theta}_{c,k}}(x) = \|x - m_{y_i, \Sigma_i, k}\|_{\Sigma_i}^2 +$
 436 $v_{y_i, \Sigma_i, k} + \log(\det(\Sigma_i))$). Set $\tilde{\theta}'_{c,k, sig}$, the family $\tilde{\theta}_{c,k, sig}$ without the centered removed because
 437 of $Threshold$ in Algorithm 1. In the final clustering, we can assign the label 0 to the sig'
 438 points with smaller value for $f_{\tilde{\theta}'_{c,k, sig}}$, with sig' calibrated with the same heuristic as in Figure
 439 1, with $f_{\tilde{\theta}'_{c,k, sig}}$ in spite of the DTM. Note that for very large datasets, since computing $\tilde{\theta}_{c,k}$
 440 may take some time, it is possible to compute a $\tilde{\theta}_{c,k}$ from a random sub-sample of \mathbb{X} , to
 441 make the clustering procedure on the centers of $\tilde{\theta}_{c,k}$, and then to go back to \mathbb{X} by clustering
 442 points with the label of the center of its curved weighted Voronoi cell.



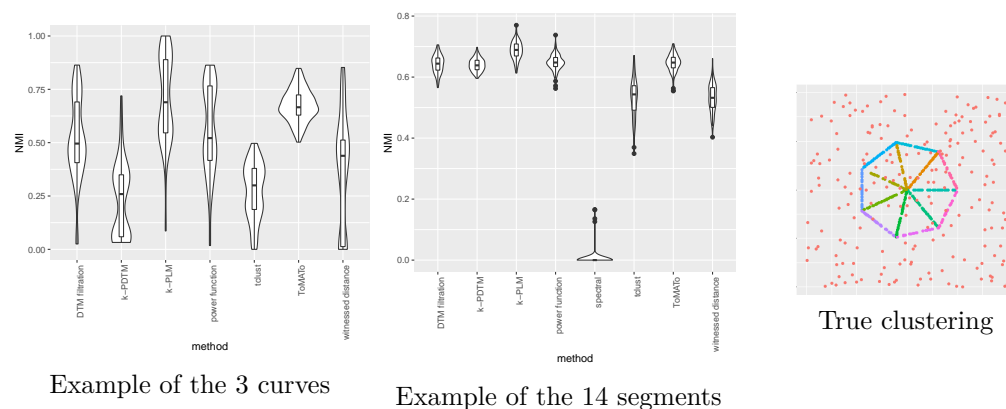
443 ■ **Figure 2** Two resulting clusterings, with ellipsoids

444 The false discovery rate FDR is defined as the proportion of signal points that are
 445 considered as outliers by the algorithm. In this example, for the two values of $Threshold$,
 446 we got exactly the same false discovery rate: $FDR = 0.022$. Nonetheless, the proportion of
 447 outliers that are considered as signal points by the algorithm equals 0.062 for $Threshold = \infty$
 448 and 0.03 for $Threshold = -9.49$. It means that the improved method that removes bad
 449 means is better to detect outliers. For both methods, we computed the normalized mutual
 450 information NMI, a classical tool that measures the clustering performance ($NMI = 1$ for
 451 perfect clustering, $NMI = 0$ for a terrible clustering). The NMI between the true labels (0,1,2
 452 or 3 ; 0 for outliers) and the label returned by the algorithm is 0.8052 when $Threshold = \infty$
 453 and 0.8587 when $Threshold = -9.49$. We also computed the NMI between the true labels

454 (1,2,3) and the labels returned by the algorithm for non-outliers points, considered as signal.
 455 For both value of *Threshold*, $NMI = 0.9766712$, this is an almost perfect clustering.

4.2 Comparison of the different methods on synthetic datasets

457 We compared different clustering methods on two generated datasets : the previous dataset
 458 with 3 curves, and datapoints from a polygonal curve of 14 segments, as in [8]. We set
 459 parameters to $N = 500$, $N_{noise} = 200$, $\sigma = 0.02$ is the standard deviation for the corrupting
 460 Gaussian distribution, $c = 50$, $k = 10$, $it = 100$ and at each step, *Threshold* is chosen such
 461 that 10 means are removed from centers $\tilde{\theta}_{c,k,sig}$ for the c -PLM. For the Tomato algorithm,
 462 the parameter $r = 0.12$ for the Rips graph is such that the graph is not connected (for the
 463 first example), thus there are more than 3 clusters. We used the function `dbscan` from the R
 464 packages `dbscan` [19], with parameters $eps = 0.15$ and $minPts = 10$; `tclust` and `specc` from
 465 the `tclust` [17] and `kernlab` [21] R packages. The shape of the violin for the c -PLM in the first
 466 example suggests that more than half of the clustering, among the 100, were almost perfect.



467 **Figure 3** Violin plots representing the NMI computed on signal points, detected as signal points.

4.3 Applications to real datasets

4.3.1 Recovering fleas species, based on 6 measurements

470 We picked the dataset flea from the R-package `tourr` [27], initially from [23]. This dataset
 471 contains records of 6 measurements for 74 males insects from the Palaeartic, from three
 472 different species : *Heptapotamica*, *Concinna*, *Heikertingeri*. The variables correspond to
 473 measurements on the tarsus, the aedeagus and the head. We normalized data so that the
 474 mean and variance of each of the 6 variables are respectively 0 and 1. In Table 2, we computed
 475 the NMI between the true species and the clustering returned by different methods. We ran
 476 each algorithm 10 times with at most 100 iterations. For every k -nearest-neighbours-based
 477 algorithm, we set $k = 10$. For ToMATO, we set $r = 1.9$ so that the graph is connected ;
 478 for the c -PLM and the c -DTM, $c = 50$ and for `dbscan`, $eps = 1.5$ and $minPts = 10$. The
 479 3-PDTM and 3-PLM methods consists in clustering data according to the weighted Voronoi
 480 cells given by the optimal centers and covariance matrices.

XX:14 Robust anisotropic power distance filtrations

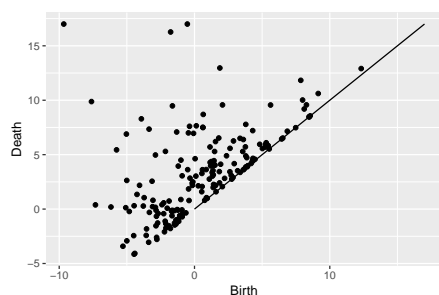
481 ■ **Table 2** NMI between clustering of fleas and their true specie

482 Without Algorithm 1	k -means 0.825	tclust 0.769	DBSCAN 0.647	Spectral clustering 1	3-PLM 1	3-PDTM 1
484 With Algorithm 1	ToMATo 0.628	Witnessed 0.906	power 1	DTM-filtration 1	c -PLM hier. 1	c -PDTM hier. 1

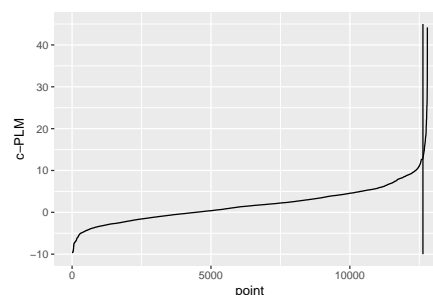
486 4.3.2 Clustering a earthquake dataset

487 We consider a set of 12790 points representing the longitude and latitude of measured earth-
 488 quakes of magnitude non smaller than 5.0, between the 01/01/1970 and the 01/01/2010. This
 489 dataset was picked from the website <http://earthquake.usgs.gov/earthquakes/eqarchives/epic/>.

490 We used Algorithm 1 with an approximation of the c -PLM based on a sub-sample of
 491 2000 points from the dataset, with parameters $c = 200$, $k = 10$ and for $it = 50$ iterations.
 492 We restricted matrices Σ to have eigenvalues smaller than 50 by thresholding them. The
 493 persistence diagram in Figure 4 suggests that the dataset has 4 or 10 clusters. Moreover, the
 494 curve of the sorted values of the c -PLM approximation on the pointset in Figure 4 suggests
 495 to keep $sig = 12250$ points as signal points. See Figure 5 for the corresponding clustering.

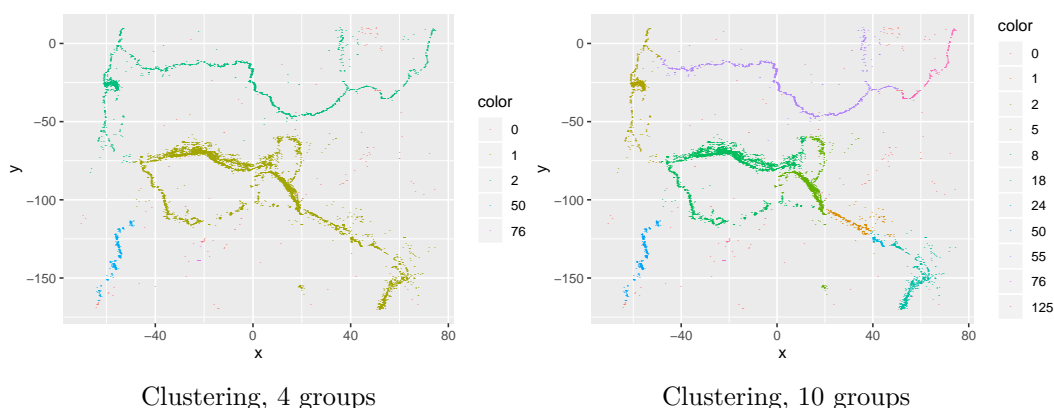


Persistence diagram



Number of signal points selection

496 ■ **Figure 4** Parameters selection heuristics.



Clustering, 4 groups

Clustering, 10 groups

497 ■ **Figure 5** Earthquake clustering with Algorithm 1, for the c -PLM function.

498 — References

- 499 1 Hirokazu Anai, Frédéric Chazal, Marc Glisse, Yuichi Ike, Hiroya Inakoshi, Raphaël Tinarrage,
500 and Yuhei Umeda. DTM-Based Filtrations. In Gill Barequet and Yusu Wang, editors, *35th*
501 *International Symposium on Computational Geometry (SoCG 2019)*, volume 129 of *Leibniz*
502 *International Proceedings in Informatics (LIPIcs)*, pages 58:1–58:15, Dagstuhl, Germany, 2019.
503 Schloss Dagstuhl–Leibniz-Zentrum fuer Informatik. URL: [http://drops.dagstuhl.de/opus/](http://drops.dagstuhl.de/opus/volltexte/2019/10462)
504 [volltexte/2019/10462](http://drops.dagstuhl.de/opus/volltexte/2019/10462), doi:10.4230/LIPIcs.SoCG.2019.58.
- 505 2 Arindam Banerjee, Srujana Merugu, Inderjit S. Dhillon, and Joydeep Ghosh. Clustering
506 with bregman divergences. *J. Mach. Learn. Res.*, 6:1705–1749, December 2005. URL: <http://dl.acm.org/citation.cfm?id=1046920.1194902>.
- 508 3 Gregory Bell, Austin Lawson, Joshua Martin, James Rudzinski, and Clifford Smyth. Weighted
509 persistent homology. *Involve*, 12(5):823–837, 2019. URL: [https://doi.org/10.2140/involve.](https://doi.org/10.2140/involve.2019.12.823)
510 [2019.12.823](https://doi.org/10.2140/involve.2019.12.823), doi:10.2140/involve.2019.12.823.
- 511 4 Claire Bréchet. Robust shape inference from a sparse approximation of the gaussian trimmed
512 loglikelihood. unpublished, 2018.
- 513 5 Claire Bréchet, Aurélie Fischer, and Clément Levrard. Robust bregman clustering. submit-
514 ted, 2018.
- 515 6 Claire Bréchet and Clément Levrard. A k-points-based distance for robust geometric
516 inference. submitted, 2017.
- 517 7 Mickaël Buchet, Frédéric Chazal, Steve Oudot, and Donald R Sheehy. Efficient and robust
518 persistent homology for measures. *Computational Geometry*, 58:70–96, 2016.
- 519 8 Mickael Buchet, Tamal K. Dey, Jiayuan Wang, and Yusu Wang. Declutter and Resample:
520 Towards Parameter Free Denoising. In Boris Aronov and Matthew J. Katz, editors, *33rd*
521 *International Symposium on Computational Geometry (SoCG 2017)*, volume 77 of *Leibniz*
522 *International Proceedings in Informatics (LIPIcs)*, pages 23:1–23:16, Dagstuhl, Germany, 2017.
523 Schloss Dagstuhl–Leibniz-Zentrum fuer Informatik.
- 524 9 F. Chazal, V. de Silva, M. Glisse, and S. Oudot. *The structure and stability of persistence*
525 *modules*. Springer International Publishing, 2016.
- 526 10 Frédéric Chazal, David Cohen-Steiner, Marc Glisse, Leonidas J. Guibas, and Steve Y. Oudot.
527 Proximity of persistence modules and their diagrams. In *Proceedings of the Twenty-fifth*
528 *Annual Symposium on Computational Geometry*, SCG '09, pages 237–246, New York, NY,
529 USA, 2009. ACM. URL: <http://doi.acm.org/10.1145/1542362.1542407>, doi:10.1145/
530 [1542362.1542407](http://doi.acm.org/10.1145/1542362.1542407).
- 531 11 Frédéric Chazal, David Cohen-Steiner, and Quentin Mérigot. Geometric Inference for Measures
532 based on Distance Functions. *Foundations of Computational Mathematics*, 11(6):733–751,
533 2011. URL: <https://hal.inria.fr/inria-00383685>, doi:10.1007/s10208-011-9098-0.
- 534 12 Frédéric Chazal, Leonidas J. Guibas, Steve Y. Oudot, and Primoz Skraba. Persistence-
535 based clustering in riemannian manifolds. *J. ACM*, 60(6):41:1–41:38, November 2013. URL:
536 <http://doi.acm.org/10.1145/2535927>, doi:10.1145/2535927.
- 537 13 David Cohen-Steiner, Herbert Edelsbrunner, and John Harer. Stability of persistence diagrams.
538 *Discrete & Computational Geometry*, 37(1):103–120, Jan 2007. URL: [https://doi.org/10.](https://doi.org/10.1007/s00454-006-1276-5)
539 [1007/s00454-006-1276-5](https://doi.org/10.1007/s00454-006-1276-5), doi:10.1007/s00454-006-1276-5.
- 540 14 Juan. A. Cuesta-Albertos, Alfonso Gordaliza, and Carlos Matrán. Trimmed k -means: an
541 attempt to robustify quantizers. *Ann. Statist.*, 25(2):553–576, 04 1997. URL: [https://doi.](https://doi.org/10.1214/aos/1031833664)
542 [org/10.1214/aos/1031833664](https://doi.org/10.1214/aos/1031833664), doi:10.1214/aos/1031833664.
- 543 15 Vin de Silva and Robert Ghrist. Coverage in sensor networks via persistent homology.
544 *Algebr. Geom. Topol.*, 7(1):339–358, 2007. URL: <https://doi.org/10.2140/agt.2007.7.339>,
545 doi:10.2140/agt.2007.7.339.
- 546 16 H. Edelsbrunner, D. Letscher, and A. Zomorodian. Topological persistence and simplification.
547 *Discrete Comput. Geom.*, 28:511–533, 2002.

- 548 17 Heinrich Fritz, Luis A. Garcia-Escudero, and Agustin Mayo-Iscar. tclust: An R package for
549 a trimming approach to cluster analysis. *Journal of Statistical Software*, 47(12):1–26, 2012.
550 URL: <http://www.jstatsoft.org/v47/i12/>.
- 551 18 Leonidas J. Guibas, Quentin Mérigot, and Dmitriy Morozov. Witnessed k-distance. In
552 *Proceedings of the Twenty-seventh Annual Symposium on Computational Geometry*, SoCG
553 '11, pages 57–64, New York, NY, USA, 2011. ACM. URL: [http://doi.acm.org/10.1145/](http://doi.acm.org/10.1145/1998196.1998205)
554 [1998196.1998205](http://doi.acm.org/10.1145/1998196.1998205), doi:10.1145/1998196.1998205.
- 555 19 Michael Hahsler, Matthew Piekenbrock, and Derek Doran. dbscan: Fast density-based
556 clustering with R. *Journal of Statistical Software*, 91(1):1–30, 2019. doi:10.18637/jss.v091.
557 i01.
- 558 20 Allen Hatcher. *Algebraic topology*. Cambridge Univ. Press, Cambridge, 2000. URL: [https://](https://cds.cern.ch/record/478079)
559 cds.cern.ch/record/478079.
- 560 21 Alexandros Karatzoglou, Alex Smola, Kurt Hornik, and Achim Zeileis. kernlab – an S4
561 package for kernel methods in R. *Journal of Statistical Software*, 11(9):1–20, 2004. URL:
562 <http://www.jstatsoft.org/v11/i09/>.
- 563 22 S. P. Lloyd. Least squares quantization in PCM. *IEEE Transactions on Information Theory*,
564 28:129–137, 1982.
- 565 23 Alexander A. Lubischew. On the use of discriminant functions in taxonomy. pages 455–477,
566 1962.
- 567 24 J. MacQueen. Some methods for classification and analysis of multivariate observations.
568 In *Proceedings of the Fifth Berkeley Symposium on Mathematical Statistics and Probability*,
569 *Volume 1: Statistics*, pages 281–297, Berkeley, Calif., 1967. University of California Press.
570 URL: <https://projecteuclid.org/euclid.bsmsp/1200512992>.
- 571 25 P. J. Rousseeuw and A. M. Leroy. *Robust Regression and Outlier Detection*. 1987.
- 572 26 Ulrike von Luxburg. A tutorial on spectral clustering. *Statistics and Computing*, 17(4):395–416,
573 Dec 2007.
- 574 27 Hadley Wickham, Dianne Cook, Heike Hofmann, and Andreas Buja. tourr: An R package for
575 exploring multivariate data with projections. *Journal of Statistical Software*, 40(2):1–18, 2011.
576 URL: <http://www.jstatsoft.org/v40/i02/>.

577 A Proof of Theorem 1

578 For ease of exposition, we set $i = 1$ and $j = 2$, with the assumption that $\omega_1 \leq \omega_2$. In order
579 to make the problem simpler, we first transform ellipsoid \mathcal{E}_1^t into a sphere by noting that

$$580 \|x - m_1\|_{\Sigma_1^{-1}}^2 = (x - m_1)^T P_1 D_1^{-1} P_1^T (x - m_1) = \left(\sqrt{D_1^{-1}} P_1^T (x - m_1) \right)^T \left(\sqrt{D_1^{-1}} P_1^T (x - m_1) \right).$$

581 Therefore, we set $m'_1 = \sqrt{D_1^{-1}} P_1^T m_1$ and $y = \sqrt{D_1^{-1}} P_1^T x - m'_1$. With this notation, it comes
582 that $y^T y \leq t - \omega_1$ is the new equation of \mathcal{E}_1^t , and $(y - m'_2)^T \tilde{\Sigma} (y - m'_2) \leq t - \omega_2$ the new
583 equation of \mathcal{E}_2^t , with $\tilde{\Sigma} = \sqrt{D_1} P_1^T \Sigma_2^{-1} P_1 \sqrt{D_1}$ and $m'_2 = \sqrt{D_1^{-1}} P_1^T m_2 - m'_1$. Set $\tilde{m} = \tilde{P}^T m'_2$
584 and $z = \tilde{P}^T y$, for \tilde{P} orthogonal and \tilde{D} diagonal matrices such that $\tilde{\Sigma} = \tilde{P} \tilde{D} \tilde{P}^T$. With these
585 notations, \mathcal{E}_1^t has equation $z^T z \leq t - \omega_1$ and \mathcal{E}_2^t has equation $(z - \tilde{m})^T \tilde{D} (z - \tilde{m}) \leq t - \omega_2$ in
586 some coordinate system.

587 The ellipsoid \mathcal{E}_1^t appears at time $t = \omega_1$, this is before \mathcal{E}_2^t that appears at time $t = \omega_2$.
588 Consequently, when $\|\tilde{m}\| \leq \sqrt{\omega_2 - \omega_1}$, the first time t for which \mathcal{E}_1^t and \mathcal{E}_2^t intersect is given
589 by $t_{1,2} = \omega_2$.

590 From now on, we may assume that $\|\tilde{m}\| > \sqrt{\omega_2 - \omega_1}$. Ellipsoids \mathcal{E}_1^t and \mathcal{E}_2^t merge at a
591 time $t_{1,2} > \omega_2$, and their intersection is given by some point $z \in \mathbb{R}^d$ that satisfies:

$$592 \blacksquare \sum_i (z_i - \tilde{m}_i)^2 \lambda_i = t_{1,2} - \omega_2$$

- 593 ■ $\sum_i z_i^2 = t_{1,2} - \omega_1$
 594 ■ $\exists \lambda > 0, \lambda z = -\tilde{D}(z - \tilde{m})$.

595 The last assumption comes from the fact that $\mathcal{E}_1^{t_{1,2}}$ and $\mathcal{E}_2^{t_{1,2}}$ are tangent at their intersection
 596 point z . Moreover, the tangent space of $\mathcal{E}_1^{t_{1,2}}$ at z is orthogonal to $2z$, the gradient of $\tilde{z} \mapsto \tilde{z}^T \tilde{z}$
 597 at z ; and the tangent space of $\mathcal{E}_2^{t_{1,2}}$ at z is orthogonal to $2\tilde{D}(z - \tilde{m})$. Then, necessarily, λ
 598 satisfies the following equation (c.f. Equation (2)):

$$599 \quad \sum_{i=1}^d \frac{\lambda_i - \lambda^2}{(\lambda + \lambda_i)^2} \lambda_i \tilde{m}_i^2 = \omega_2 - \omega_1.$$

600 The eigenvalues of $\tilde{\Sigma}$, $(\lambda_i)_{i \in [1,d]}$ are positive. For every $a > 0$, $f_a : \lambda \mapsto \frac{a - \lambda^2}{(\lambda + a)^2} a$ is decreasing,
 601 with $\sum_{i=1}^d f_{\lambda_i}(0) \tilde{m}_i^2 = \|\tilde{m}\|^2 > \omega_2 - \omega_1$ and $\lim_{\lambda \rightarrow +\infty} \sum_{i=1}^d f_{\lambda_i}(\lambda) \tilde{m}_i^2 = -\sum_{i=1}^d \lambda_i \tilde{m}_i^2 < 0$.
 602 Consequently, Equation (2) has a unique solution λ . It comes that the ellipsoids \mathcal{E}_1^t and \mathcal{E}_2^t
 603 intersect at time $t_{1,2} = \omega_2 + \sum_{i=1}^d \left(\frac{\lambda \tilde{m}_i}{\lambda + \lambda_i} \right)^2 \lambda_i$ for this unique λ .

604 **B** Proof of Theorem 2

605 Since the anisotropic weighted Vietoris-Rips complex $\text{VR}_{\mathbf{m}, \omega, \Sigma}(t)$ is the flag complex of
 606 the 1-skeleton of the weighted anisotropic Čech complex $\text{Cech}_{\mathbf{m}, \omega, \Sigma}(t)$, $\mathcal{G}_{\mathbf{m}, \omega, \Sigma}^t$, the second
 607 inclusion of (3) is trivially satisfied. We now focus on the first inclusion of (3). For every
 608 $i \in I$ and $t > 0$, we have that

$$609 \quad \bar{B}\left(m_i, \sqrt{\lambda_{\min}} \sqrt{t - \omega_i}\right) \subset \bar{B}_{\Sigma_i}\left(m_i, \sqrt{t - \omega_i}\right) \subset \bar{B}\left(m_i, \sqrt{\lambda_{\max}} \sqrt{t - \omega_i}\right). \quad (5)$$

610 Set $t > 0$ and $0 < t' \leq \frac{\lambda_{\min}}{\lambda_{\max}} \frac{d+1}{2d} t$. If $\sigma \in \text{VR}_{\mathbf{m}, \omega, \Sigma}(t')$, then, for every $i, j \in \sigma$, the intersection
 611 $B_{\Sigma_i}(m_i, \sqrt{t' - \omega_i}) \cap B_{\Sigma_j}(m_j, \sqrt{t' - \omega_j})$ is nonempty and according to (5), $B(m_i, \sqrt{\lambda_{\max}} \sqrt{t' - \omega_i}) \cap$
 612 $B(m_j, \sqrt{\lambda_{\max}} \sqrt{t' - \omega_j})$ is nonempty. According to the Vietoris-Rips theorem of [3, Theorem
 613 3.2], $\bigcap_{i \in \sigma} B(m_i, \sqrt{\lambda_{\max}} \sqrt{\frac{2d}{d+1}} \sqrt{t' - \omega_i})$ is nonempty and since the weights are non-negative
 614 and $\sqrt{\frac{\lambda_{\max}}{\lambda_{\min}} \frac{2d}{d+1}} \geq 1$, it comes that $\bigcap_{i \in \sigma} B(m_i, \sqrt{\lambda_{\min}} \sqrt{t - \omega_i})$ is nonempty, and finally,
 615 according to (5), $\sigma \in \text{Cech}_{\mathbf{m}, \omega, \Sigma}(t)$.

616 **C** Sketch of proof for the q -tameness of power functions

617 Any function f of this type is both continuous and proper (the pre-image of any compact set
 618 is included in a union of ellipsoids and thus compact).

619 Since \mathbb{R}^d is triangulable, there is some homeomorphism h between \mathbb{R}^d and some locally
 620 finite simplicial complex C , then for any fixed t , we can define a complex C_t that contains
 621 $f^{-1}((-\infty, t])$, and $f \circ h|_{C_t}$ is continuous. Just as [7, Proposition 3.5], the conclusion follows
 622 from [9, Theorem 2.22] that states that the filtration of the sublevel sets of a continuous
 623 function defined on a finite polyhedron is q -tame.