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# Adaptive Domain Decomposition method for Saddle Point problem in Matrix Form 

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#### Abstract

We introduce an adaptive domain decomposition (DD) method for solving saddle point problems defined as a block two by two matrix. The algorithm does not require any knowledge of the constrained space. We assume that all sub matrices are sparse and that the diagonal blocks are the sum of positive semi definite matrices. The latter assumption enables the design of adaptive coarse space for DD methods, see [6.


## 1 Introduction

Solving saddle point problems with parallel algorithms is very important for many branches of scientific computing: fluid (see e.g. [14]) and solid mechanics, computational electromagnetism, inverse problem and optimization.

We are interested in domain decomposition (DD) methods since they are naturally well-fitted to modern parallel architectures. For specific systems of partial differential equations with a saddle point formulation, efficient DD methods have been designed, see e.g. [17, 13, 18] and [22] references therein.

Here as in [15, 2, 5], we consider the problem in the form of a two by two block matrix. Let $m$ and $n$ be two integers with $m<n$. Let $A n \times n$ SPD matrix and $B$ be a sparse $m \times n$ full rank matrix of constraints and $C$ a $m \times m$ non negative matrix (in particular, $C=0$ is allowed), we consider the following saddle point matrix:

$$
\mathcal{A}:=\left(\begin{array}{cc}
A & B^{T}  \tag{1}\\
B & -C
\end{array}\right) .
$$

When the kernel of matrix $B$ is known, very efficient multigrid methods have been designed in the context of finite element methods, see e.g. [4, 11, 10, 1, 19, 7.

Here we do not assume any knowledge on the kernel of matrix $B$. But in order to build a scalable method, we assume that all three matrices are sparse and that $A$ and $C$ are the sum of positive semi definite matrices. This is easily achieved in finite element or finite volume contexts for partial differential equations. The latter assumption enables the design of adaptive coarse space for DD methods, see 6].

## 2 Schur complement preconditioning

The sparse $n \times n$ SPD matrix $A$ is preconditioned by a two-level Schwarz type DD method :

$$
\begin{equation*}
M_{A}^{-1}:=R_{0}^{T}\left(R_{0} A R_{0}^{T}\right)^{-1} R_{0}+\sum_{i=1}^{N} R_{i}^{T}\left(R_{i} A R_{i}^{T}\right)^{-1} R_{i}, \tag{2}
\end{equation*}
$$

where $R_{0}$ is full rank $\operatorname{dim}\left(V_{0}\right) \times n$ where $V_{0}$ denotes the space spanned by the columns of $R_{0}^{T}$. The following assumptions are crucial to ensure the final method is scalable:

## Assumption 2.1 (dimension and structure of the coarse space)

- The coarse space dimension, $\operatorname{dim}\left(V_{0}\right)$, is $O(N)$ typically 10-20 times $N$.
- The coarse space is made of extensions by zero of local vectors.

Using the GenEO metod [20], it is possible to fix in advance two constants $0<\lambda_{m}<1<\lambda_{M}$ and then build a coarse space $V_{0}$ such that $M_{A}^{-1}$ is spectrally equivalent to $A^{-1}$ :

$$
\frac{1}{\lambda_{M}} M_{A}^{-1} \leqslant A^{-1} \leqslant \frac{1}{\lambda_{m}} M_{A}^{-1}
$$

The coarse space $V_{0}$ is made extensions by zero of local generalized eigenvalue problems and its dimension is typically proportional to the number of subdomains. This corresponds to Assumption 2.1.
Our aim is first to precondition the Schur complement $-S$ of matrix $\mathcal{A}$ (eq. (1)) where

$$
\begin{equation*}
S:=C+B A^{-1} B^{T} \tag{3}
\end{equation*}
$$

via the spectrally equivalent preconditioning of the spectrally equivalent Schur complement $M_{S}$,

$$
\begin{equation*}
M_{S}:=C+B M_{A}^{-1} B^{T} \tag{4}
\end{equation*}
$$

This is done in $\S 3$. Then preconditioner $M_{S}$ is itself by a spectrally equivalent preconditioner $N_{S}$, see eq. 12 in $\S 3.2$. Finally in $\S 4.2$ we will use $N_{S}$ to precondition the saddle point matrix $\mathcal{A}$.

Note that $M_{S}$ is by definition a sum of $N+2$ positive semi definite matrices

$$
\begin{equation*}
M_{S}:=B R_{0}^{T}\left(R_{0} A R_{0}^{T}\right)^{-1} R_{0} B^{T}+C+\sum_{i=1}^{N} B R_{i}^{T}\left(R_{i} A R_{i}^{T}\right)^{-1} R_{i} B^{T} \tag{5}
\end{equation*}
$$

Since $B$ is a sparse matrix, it is interesting to introduce, for all $0 \leqslant i \leqslant N$, $\tilde{R}_{i}$ the restriction operator on the support of $\Im\left(B R_{i}^{T}\right)$ so that $\tilde{R}_{i}^{T} \tilde{R}_{i} B R_{i}^{T}=$ $B R_{i}^{T}$. Then by defining for $0 \leqslant i \leqslant N$,

$$
\tilde{B}_{i}:=\tilde{R}_{i} B R_{i}^{T},
$$

the operator $M_{S}$ is rewritten as

$$
\begin{equation*}
M_{S}:=\tilde{R}_{0}^{T} \tilde{B}_{0}\left(R_{0} A R_{0}^{T}\right)^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0}+C+\sum_{i=1}^{N} \tilde{R}_{i}^{T} \tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T} \tilde{R}_{i} \tag{6}
\end{equation*}
$$

We consider a partition of unity on $H:=\mathbb{R}^{m}$ defined with local diagonal matrices $\left(\tilde{D}_{i}\right)_{1 \leqslant i \leqslant N} \in \mathbb{R}^{\operatorname{dim}\left(\operatorname{Im}\left(B R_{i}^{T}\right)\right) \times \operatorname{dim}\left(\operatorname{Im}\left(B R_{i}^{T}\right)\right)}$ :

$$
\sum_{i=1}^{N} \tilde{R}_{i}^{T} \tilde{D}_{i} \tilde{R}_{i}=I_{H}
$$

Remark 2.1 This partition of unity exists since

$$
B=\sum_{i=1}^{N} B R_{i}^{T} D_{i} R_{i}=\sum \tilde{R}_{i}^{T} \tilde{R}_{i}\left(B R_{i}^{T} D_{i} R_{i}\right)
$$

is full rank.

## 3 Preconditioning of $M_{S}$

We make the following assumption
Assumption 3.1 There exist symmetric positive semidefinite matrices $\left(\tilde{C}_{i}\right)_{1 \leqslant i \leqslant N}$ such that

$$
C=\sum_{i=1}^{N} \tilde{R}_{i}^{T} \tilde{C}_{i} \tilde{R}_{i}
$$

This assumption is not so restrictive. Indeed, if $C=0$, it is automatically satisfied, this corresponds to a minimization problem with constraints enforced exactly without penalization nor relaxation. Moreover, we have:

Lemma 3.1 If $C$ is a diagonal matrix, Assumption 3.1 is satisfied.
Proof If $C$ is a diagonal matrix, it suffices to take

$$
\tilde{C}_{i}:=\tilde{R}_{i} C \tilde{R}_{i}^{T} \tilde{D}_{i}
$$

which is a diagonal non negative matrix. Indeed, we have then:

$$
C \mathbf{P}=\sum_{i=1}^{N} C \tilde{R}_{i}^{T} \tilde{D}_{i} \tilde{R}_{i} \mathbf{P}=\sum_{i=1}^{N} \tilde{R}_{i}^{T}\left(\tilde{R}_{i} C \tilde{R}_{i}^{T} \tilde{D}_{i}\right) \tilde{R}_{i} \mathbf{P}
$$

Then, the operator $M_{S}$ is the sum of a non local but low rank matrix $S_{0}$ :

$$
S_{0}:=\tilde{R}_{0}^{T} \tilde{B}_{0}\left(R_{0} A R_{0}^{T}\right)^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0}
$$

and of $S_{1}$ which is a sum of $N$ local positive semi definite matrices:

$$
S_{1}:=\sum_{i=1}^{N} \tilde{R}_{i}^{T}\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right) \tilde{R}_{i}
$$

that is

$$
M_{S}=S_{0}+S_{1}
$$

Note that we may assume that $S_{1}$ is invertible whereas it does not make sense for $S_{0}$.

We consider next the construction of a preconditioner $M_{S_{1}}^{-1}$ to $S_{1}$.

### 3.1 Equivalent preconditioner of $S_{1}$

### 3.1.1 One-level DD

As in [6] chapter 7, we begin with a one-level Neumann-Neumann type DD method defined in terms of the Fictitious Space Lemma (FSL) [16, 8]. This study will be the basis for constructing the two-level preconditioner. Recall

$$
H:=\mathbb{R}^{m}
$$

and let $a$ be the following bilinear form:

$$
a: H \times H \rightarrow \mathbb{R} \quad a(\mathbf{P}, \mathbf{Q}):=\left(S_{1} \mathbf{P}, \mathbf{Q}\right) .
$$

Let

$$
H_{D}:=\Pi_{i=1}^{N} \mathbb{R}^{\operatorname{rank}\left(\tilde{B}_{i}\right)},
$$

and $b$ be the following bilinear form:

$$
b: H_{D} \times H_{D} \rightarrow \mathbb{R} \quad b(\mathcal{P}, \mathcal{Q}):=\sum_{i=1}^{N}\left(\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right) \mathbf{P}_{i}, \mathbf{Q}_{i}\right)
$$

We define $\mathcal{R}$ :

$$
\begin{array}{lcl}
\mathcal{R}: & H_{D} \rightarrow & H \\
& \left(\mathbf{P}_{i}\right)_{1 \leqslant i \leqslant N} \mapsto & \sum_{i=1}^{N} \tilde{R}_{i}^{T} \tilde{D}_{i} \mathbf{P}_{i},
\end{array}
$$

We now check the three assumptions of the FSL.
Surjectivity of $\mathcal{R}$ For any $\mathbf{P} \in H$, we have:

$$
\mathbf{P}=\sum_{i=1}^{N} \tilde{R}_{i}^{T} \tilde{D}_{i} \tilde{R}_{i} \mathbf{P},
$$

so that

$$
\begin{equation*}
\mathbf{P}=\mathcal{R}\left(\left(\tilde{R}_{i} \mathbf{P}\right)_{1 \leqslant i \leqslant N}\right) . \tag{7}
\end{equation*}
$$

Continuity of $\mathcal{R}$ On one hand, we have using $k_{0}$ the number of neighbours of a subdomain plus one, $k_{0}:=\max _{1 \leqslant i \leqslant N} \# \mathcal{O}(i)$ where $\mathcal{O}(i):=\{1 \leqslant j \leqslant$ $\left.N \mid \tilde{R}_{i} \tilde{D}_{i} S_{1} \tilde{D}_{j} \tilde{R}_{j}^{T} \neq 0\right\}:$

$$
\begin{aligned}
a(\mathcal{R}(\mathcal{P}), \mathcal{R}(\mathcal{P})) & =\left\|\left(\sum_{i=1}^{N} \tilde{R}_{i}^{T} \tilde{D}_{i} \mathbf{P}_{i}\right)\right\|_{a}^{2} \leqslant k_{0} \sum_{i=1}^{N}\left\|\tilde{R}_{i}^{T} \tilde{D}_{i} \mathbf{P}_{i}\right\|_{a}^{2} \\
& =k_{0}\left(\left(\sum_{j \in \mathcal{O}(i)} \tilde{R}_{j}^{T}\left(\tilde{C}_{j}+\tilde{B}_{j}\left(R_{j} A R_{j}^{T}\right)^{-1} \tilde{B}_{j}^{T}\right) \tilde{R}_{j}\right) \tilde{R}_{i}^{T} \tilde{D}_{i} \mathbf{P}_{i}, \tilde{R}_{i}^{T} \tilde{D}_{i} \mathbf{P}_{i}\right) .
\end{aligned}
$$

On the other hand, we have by definition:

$$
b(\mathcal{P}, \mathcal{P})=\sum_{i=1}^{N}\left(\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right) \mathbf{P}_{i}, \mathbf{P}_{i}\right)
$$

We can take:
$c_{R}:=\max _{1 \leqslant i \leqslant N} \max _{P_{i} \in \mathbb{R}^{r a n k\left(\tilde{B}_{i}\right)}} \frac{\left(\sum_{j \in \mathcal{O}(i)} \tilde{R}_{i} \tilde{R}_{j}^{T}\left(\tilde{C}_{j}+\tilde{B}_{j}\left(R_{j} A R_{j}^{T}\right)^{-1} \tilde{B}_{j}^{T}\right) \tilde{R}_{j} \tilde{R}_{i}^{T} \tilde{D}_{i} \mathbf{P}_{i}, \tilde{D}_{i} \mathbf{P}_{i}\right)}{\left(\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right) \mathbf{P}_{i}, \mathbf{P}_{i}\right)}$.
Stable decomposition Let $\mathbf{P} \in H$, we start from its decomposition (7) and estimate its $b$-norm

$$
b(\mathcal{P}, \mathcal{P})=\sum_{i=1}^{N}\left(\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right) \tilde{R}_{i} \mathbf{P}, \tilde{R}_{i} \mathbf{P}\right)=a(\mathbf{P}, \mathbf{P}),
$$

so that we can take $c_{T}=1$.

### 3.1.2 Two-level DD

In order to control the value of $c_{R}$ defined above, we introduce a two-level preconditioner similarly to what is done for Schur complement method in [6, $\S 7.8 .3$, page 197] or in [21]. In our case, the generalized eigenvalue value problem in each subdomain $1 \leqslant i \leqslant N$ to be solved to build the coarse space reads:

$$
\begin{array}{r}
\tilde{D}_{i}\left(\sum_{j \in \mathcal{O}(i)} \tilde{R}_{i} \tilde{R}_{j}^{T}\left(\tilde{C}_{j}+\tilde{B}_{j}\left(R_{j} A R_{j}^{T}\right)^{-1} \tilde{B}_{j}^{T}\right) \tilde{R}_{j} \tilde{R}_{i}^{T}\right) \tilde{D}_{i} \mathbf{P}_{i k}  \tag{8}\\
=\lambda_{i k}\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right) \mathbf{P}_{i k}
\end{array}
$$

It can be solved in $O(1)$ communications. The coarse space is defined as follows. Let $\tau_{S_{1}}$ be a user-defined threshold; for each subdomain $1 \leqslant i \leqslant N$, we introduce a subspace $W_{i} \subset \mathbb{R}^{\operatorname{rank}\left(\tilde{B}_{i}\right)}$ :

$$
\begin{equation*}
\tilde{W}_{i}:=\operatorname{Span}\left\{\mathbf{P}_{i k} \mid \lambda_{i k}>\tau_{S_{1}}\right\} \tag{9}
\end{equation*}
$$

Then the coarse space $\tilde{W}_{0}$ is defined by

$$
\tilde{W}_{0}:=\bigoplus_{1 \leqslant i \leqslant N} \tilde{R}_{i}^{T} \tilde{D}_{i} \tilde{W}_{i}
$$

Let $Z_{\tilde{P}_{1}}$ be a rectangular matrix whose columns span the coarse space $\tilde{W}_{0}$. Let $\tilde{P}_{0}$ be the $S_{1}$ orthogonal projection from $\mathbb{R}^{m}$ on $\tilde{W}_{0}$ whose formula is

$$
\begin{equation*}
\tilde{P}_{0}=Z_{S_{1}}\left(Z_{S_{1}}^{T} S_{1} Z_{S_{1}}\right)^{-1} Z_{S_{1}}^{T} S_{1} \tag{10}
\end{equation*}
$$

Remark 3.1 We forget for the moment the projection related to the possible kernel of $\left.\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right)$

Finally, the preconditioner for $S_{1}$ reads

$$
\begin{align*}
M_{S_{1}}^{-1} & :=Z_{S_{1}}\left(Z_{S_{1}}^{T} S_{1} Z_{S_{1}}\right)^{-1} Z_{S_{1}}^{T}+\left(I-\tilde{P}_{0}\right) \\
& \times\left(\sum_{i=1}^{N} \tilde{R}_{i}^{T} \tilde{D}_{i}\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right)^{\dagger} \tilde{D}_{i} \tilde{R}_{i}\right)\left(I-\tilde{P}_{0}^{T}\right) \tag{11}
\end{align*}
$$

Recall that we have for $\alpha:=\max \left(1, \frac{k_{0}}{\tau_{S_{1}}}\right)$ :

$$
\frac{1}{\alpha} M_{S_{1}}^{-1} \leqslant S^{-1} \leqslant M_{S_{1}}^{-1}
$$

Remark 3.2 Solving a linear system with a local Schur complement

$$
\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right) \mathbf{P}_{i}=\mathbf{G}_{i}
$$

amounts to solving an augmented sparse system of the form

$$
-\left(\begin{array}{lc}
R_{i} A R_{i}^{T} & \tilde{B}_{i}^{T} \\
\tilde{B}_{i} & -\tilde{C}_{i}
\end{array}\right)\binom{\mathbf{U}_{i}}{\mathbf{P}_{i}}=\binom{\mathbf{0}}{\mathbf{G}_{i}}
$$

### 3.2 Preconditioner for $M_{S}$

From the spectrally equivalent preconditioner $M_{S_{1}}$ to $S_{1}$, we define $N_{S}$ a spectrally equivalent preconditioner to $M_{S}$ and thus to $S$ as well:

$$
\begin{equation*}
N_{S}:=S_{0}+M_{S_{1}} . \tag{12}
\end{equation*}
$$

We now consider the application of the preconditioner, that is the solving of $N_{S} \mathbf{P}=\mathbf{G}$ in $\mathbf{P}$ for some right handsize $\mathbf{G} \in \mathbb{R}^{m}$. By factorizing $R_{0} A R_{0}^{T}=$ $L_{0} L_{0}^{T} \in \mathbb{R}^{\operatorname{dim}\left(V_{0}\right) \times \operatorname{dim}\left(V_{0}\right)}$, we have a Cholevsky factorization of $S_{0}$ :

$$
S_{0}=\tilde{R}_{0}^{T} \tilde{B}_{0} L_{0}^{T-1} L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0} .
$$

By Sherman-Morrison's technique, solving $N_{S} \mathbf{P}=\mathbf{G}$, amounts to solving

$$
\left(\begin{array}{cc}
M_{S_{1}} & \tilde{R}_{0}^{T} \tilde{B}_{0} L_{0}^{T-1} \\
L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0} & -I
\end{array}\right)\binom{\mathbf{P}}{\mathbf{y}}=\binom{\mathbf{G}}{\mathbf{0}}
$$

with $\mathbf{y}:=L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0} \mathbf{P} \in \mathbb{R}^{\operatorname{dim}\left(V_{0}\right)}$. This is equivalent to solving its schur complement:
$M_{A_{0}} \mathbf{y}:=\left(I+L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0} M_{S_{1}}^{-1} \tilde{R}_{0}^{T} \tilde{B}_{0} L_{0}^{T-1}\right) \mathbf{y}=L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0} M_{S_{1}}^{-1} \mathbf{G} \in \mathbb{R}^{\operatorname{dim}\left(V_{0}\right)}$.
Following 3.2.1, computing the entries of $M_{A_{0}}$ is scalable.

### 3.2.1 Scalability of the computation of $M_{A_{0}}$

The idea is that the computation of the entries of $M_{A_{0}}$ is scalable since $\tilde{R}_{0}^{T} \tilde{B}_{0}=B^{T} R_{0}$ has a sparse structure, see Assumption 2.1. We sketch here the idea by considering for instance the term corresponding to the coarse space in the formula for $M_{S_{1}}^{-1}$ (see eq. (11)):

$$
L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0} Z_{S_{1}}\left(Z_{S_{1}}^{T} S_{1} Z_{S_{1}}\right)^{-1} Z_{S_{1}}^{T} \tilde{R}_{0}^{T} \tilde{B}_{0} L_{0}^{T-1}
$$

Matrix $Z_{S_{1}}^{T} \tilde{R}_{0}^{T} \in \mathbb{R}^{\operatorname{dim}\left(\tilde{W}_{0}\right) \times \operatorname{dim}\left(V_{0}\right)}$ is the product of two DD coarse spaces so that its computation is $O(1)$.

For the other computations, we need extra assumptions 3.2 and 3.3 . Indeed, let's look at the sum over the subdomains in (11) and we isolate one subdomain $1 \leqslant i \leqslant N$ :
$L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0}\left(I-\tilde{P}_{0}\right) \tilde{R}_{i}^{T} \tilde{D}_{i}\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right)^{\dagger} \tilde{D}_{i} \tilde{R}_{i}\left(I-\tilde{P}_{0}^{T}\right) \tilde{R}_{0}^{T} \tilde{B}_{0} L_{0}^{T-1}$.
This can be decomposed into a four term sum, the first term being:

$$
L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0} \tilde{R}_{i}^{T} \tilde{D}_{i}\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right)^{\dagger} \tilde{D}_{i} \tilde{R}_{i} \tilde{R}_{0}^{T} \tilde{B}_{0} L_{0}^{T^{-1}},
$$

Assumption 3.2 We assume that, for all $1 \leqslant i \leqslant N$, the local matrix

$$
\tilde{D}_{i} \tilde{R}_{i} \tilde{R}_{0}^{T} \tilde{B}_{0}=\tilde{D}_{i} \tilde{R}_{i} B^{T} R_{0}
$$

has $O(1)$ non zero columns.
The second term being:

$$
-L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0} \tilde{R}_{i}^{T} \tilde{D}_{i}\left(\tilde{C}_{i}+\tilde{B}_{i}\left(R_{i} A R_{i}^{T}\right)^{-1} \tilde{B}_{i}^{T}\right)^{\dagger} \tilde{D}_{i} \tilde{R}_{i} \tilde{P}_{0}^{T} \tilde{R}_{0}^{T} \tilde{B}_{0} L_{0}^{T-1}
$$

We focus on the term

$$
\tilde{D}_{i} \tilde{R}_{i} \tilde{P}_{0}^{T} \tilde{R}_{0}^{T} \tilde{B}_{0}=\tilde{D}_{i} \tilde{R}_{i} S_{1} Z_{S_{1}}\left(Z_{S_{1}}^{T} S_{1} Z_{S_{1}}\right)^{-1} Z_{S_{1}}^{T} B^{T} R_{0}
$$

Assumption 3.3 We assume that, for all $1 \leqslant i \leqslant N$, the local matrix $\tilde{D}_{i} \tilde{R}_{i} S_{1} Z_{S_{1}}$ has $O(1)$ non zero columns

Using this assumption, the computation of second term is scalable. The two other terms can also be computed in a scalable way.

Note that, the resulting matrix $M_{A_{0}}$ being of small size, it can be factorized by a direct method.

## 4 Recap

### 4.1 Setup for the Schur complement preconditioner $M_{S}$

We have a setup phase which is composed of:

1. Build the two-level preconditioner $M_{A}^{-1}$ for $A$, see eq. (22),
2. Build the two-level preconditioner $M_{S_{1}}^{-1}$ for $S_{1}$, see eq. (11),
3. Compute the entries of matrix $M_{A_{0}}$, see eq. (13) and factorize it.

Once the setup is complete, applying preconditioner $N_{S}^{-1}$ can be performed following Algorithm 1 .

```
Algorithm \(1 N_{S}^{-1}\) matvec product
    INPUT: \(\mathbf{G} \in \mathbb{R}^{m}\) OUTPUT: \(\mathbf{P}=N_{S}^{-1} \mathbf{G}\)
    1. Compute \(=\mathbf{G}^{\prime}:=M_{S_{1}}^{-1} \mathbf{G}\).
    2. Compute the right handside \(L_{0}^{-1} \tilde{B}_{0}^{T} \tilde{R}_{0} \mathbf{G}^{\prime}\) of eq. (13).
    3. Solve eq. (13) in \(\mathbf{y}\).
    4. Compute \(\mathbf{P}:=M_{S_{1}}^{-1}\left(\mathbf{G}-\tilde{R}_{0}^{T} \tilde{B}_{0} L_{0}^{T-1} \mathbf{y}\right)\)
```


### 4.2 DD solver for the saddle point system

We now consider the solving of the saddle point problem:

$$
\left(\begin{array}{cc}
A & B^{T}  \tag{14}\\
B & -C
\end{array}\right)\binom{\mathbf{U}}{\mathbf{P}}=\binom{\mathbf{F}_{U}}{\mathbf{F}_{P}}
$$

The following three factor factorization, see e.g. [3]:

$$
\left(\begin{array}{cc}
A & B^{T} \\
B & -C
\end{array}\right)=\left(\begin{array}{cc}
I & 0 \\
B A^{-1} & I
\end{array}\right)\left(\begin{array}{cc}
A & 0 \\
0 & -\left(C+B A^{-1} B^{T}\right)
\end{array}\right)\left(\begin{array}{cc}
I & A^{-1} B^{T} \\
0 & I
\end{array}\right),
$$

shows that solving eq. (14) can be performed by solving sequentially linear systems with $A$ and one with $C+B A^{-1} B^{T}$. This leads to Algorithm 2 .

```
Algorithm 2 DD saddle point solver
    INPUT: \(\binom{\mathbf{F}_{U}}{\mathbf{F}_{P}} \in \mathbb{R}^{n+m} \quad\) OUTPUT: \(\binom{\mathbf{U}}{\mathbf{P}}\) the solution to 144 .
    1. Solve \(A \mathbf{G}_{U}=\mathbf{F}_{U}\) by a PCG with \(M_{A}^{-1}\) as a preconditioner
    2. Compute \(\mathbf{G}_{P}:=\mathbf{F}_{P}-B \mathbf{G}_{U}\)
    3. Solve \(\left(C+B A^{-1} B^{T}\right) \mathbf{P}=-\mathbf{G}_{P}\) by a PCG with \(N_{S}^{-1}\) as a preconditioner,
    see Algorithm 1 .
    4. Compute \(\mathbf{G}_{U}:=\mathbf{F}_{U}-B^{T} \mathbf{P}\)
    5. Solve \(A \mathbf{U}=\mathbf{G}_{U}\) by a PCG with \(M_{A}^{-1}\) as a preconditioner
```


## 5 Variants

In § (2), we start with a two level additive Schwarz method (ASM) eq. (2) as a preconditioner for matrix $A$ in (11). Another possibility is to start from a balancing Neumann-Neumann (BNN) or more generally a SORAS [9 or BDD-H [12] type method:

$$
\begin{equation*}
M_{A_{S O R A S}}^{-1}:=R_{0}^{T}\left(R_{0} A R_{0}^{T}\right)^{-1} R_{0}+\sum_{i=1}^{N} R_{i}^{T} D_{i} A_{i}^{R o b^{-1}} D_{i} R_{i} \tag{15}
\end{equation*}
$$

where for each subdomain $1 \leqslant i \leqslant N, A_{i}^{\text {Rob }}$ is a local Neumann matrix (BNN algorithm) or an arbitrary invertible matrix. In order to build the preconditioners for matrix $S$ and then $\mathcal{A}$, it is sufficient to replace matrices $\left(R_{i} A R_{i}^{T}\right)^{-1}$ with $A_{i}^{R o b^{-1}}$ in the above sections.

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