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Towards deadlock free Sokoban

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Abstract. We present algorithms to compute and match maze specific deadlock tables at Sokoban. They enable to use a greedy search that can solve problems that are not solvable using IDA*.

1 Introduction

Sokoban is a PSPACE-complete single agent search problem [3]. The presence of deadlocks makes Sokoban hard to solve even non optimally. Identifying deadlock states is therefore an important step towards solving Sokoban. In this paper we describe algorithms that find maze specific deadlocks as well as efficient algorithms to match the deadlocks during search.

The second section details deadlocks, the third section explains how they can be efficiently matched, the fourth section deals with the retrograde analysis of deadlocks, the fifth section deals with search algorithm using deadlocks, the sixth section gives experimental results.

2 Maze specific deadlocks

In Rolling Stone [8, 6, 7] small 5x4 deadlock patterns are used to find simple deadlocks. Moreover a specific search algorithm is used to identify conflict among limited subsets of stones and dynamically create penalty patterns.

We use a different approach since we compute maze specific deadlocks containing a predefined number of stones.

Basic one stone and two stones deadlocks can be defined as:

- stones in corners,
- stones on a border between two corners,
- two adjoining stones along a wall.

More complex deadlocks are computed using retrograde analysis. Retrograde analysis has already been successfully used in two player [13, 11, 14, 1, 2, 12] and single agent games [9, 4, 10, 5].
3 Forest of deadlocks

A deadlock is represented as an ordered set of locations of stones followed by a location of the man. Two deadlocks can have the same locations of stones and a different location for the man.

We represent sets of deadlocks as a different forest for each maze. Each tree of the forest is rooted at a location of the maze, the deadlocks in the tree all contain a stone at that location. The children of the root contain the different second locations of deadlock in the set of deadlocks starting at the root, and so on.

A deadlock can be matched on a screen if a subset of all the stones in the screen corresponds to the set of stones of the deadlock and if the man on the screen can see the man of the deadlock inside the deadlock screen.

In order to verify all deadlocks for a screen, the tree associated to each stone in the screen is verified. The verification consists in checking recursively for each child that the corresponding stone is present in the screen.

The next figures explain how a tree is built for three deadlock with the same first location.

The first screen is a two stones deadlock containing locations 16 and 21 and the man at location 0. These two locations and the location of the man are used to build the leftmost tree. Then the second deadlock containing locations 16, 29 and 32 and the man at location 0 is added to the tree, resulting in the tree in the middle. The third deadlock (16, 29, 33 and man at 44) is then added and gives the rightmost tree.
Nodes contain a stone and possibly several man locations. The pseudo-code for verifying a tree is:

**deadlock tree verification process**

```plaintext
bool known_state (deadlock, man)
    if stone in state
        for deadlockMan in men
            if deadlockMan can see man in deadlock
                return true
        push stone in deadlock
        for c in children
            if c->known_state (deadlock, man)
                return true
        pop stone from deadlock
    return false
```

A state is a maze containing stones at different locations and a man at a given location, the goal of the tree verification is to find if a subset of the stones of a state, is a deadlock. At each node of the tree, the algorithm verifies that the state contains the stone of the node. If a node contains deadlocks, it contains men locations. If one of the men locations is visible from the man in the state, a deadlock is found. Otherwise the verification continues for all children of the node.

All the trees rooted to the locations of the stones of a state are verified.

### 4 Retrograde analysis of deadlocks

Retrograde analysis starts with finding all the possible combinations of three stones. In a first pass, for each combination, it verifies that a subset is not already a deadlock. Then it tries all possible moves and verifies that they all lead to a known deadlock. When it is the case the combination is added to the set of deadlocks. In the subsequent passes, it goes through all the found deadlocks, finds all possible previous states. For each previous state, it checks if all moves lead to a known deadlock, and when it is the case the previous state is added to the set of deadlocks.

The pseudo-code for the first pass is:

**Combinations deadlock search process**

```plaintext
list combinationsDeadlocksSearch (nbStones)
    list = empty list
    state = first_combination (nbStones)
    while true
        if state is not a known deadlock
            isDeadlock = 1
        for child in nextStates (state)
```
if child is not a known deadlock  
    isDeadlock = 0  
break  
if isDeadlock equals 1  
    add state to deadlocksTree  
    push state in list  
if state equals last_combination (nbStones)  
    break  
state = next_combination (state)  
return list  

The pseudo code for subsequent passes is:

Subsequent passes deadlock search process

list subsequentDeadlocksSearch (nbStones, states)  
list = empty_list  
for state in states  
    for move in possiblePreviousMoves (state)  
        parent = apply (state, move)  
        if parent is not a known deadlock  
            isDeadlock = 1  
        for child in nextStates (parent)  
            if child is not a known deadlock  
                isDeadlock = 0  
                break  
        if isDeadlock equals 1  
            add parent to deadlocksTree  
            add parent to list  
return list  

Finding all possible combinations in the first pass can be time consuming when the number of stones grows. Many possible combinations do not lead to a deadlock, moreover a first pass deadlock is often a state where a move leads to a smaller deadlock. In order to rapidly compute the first pass, we only consider states where a move leads to a smaller deadlock. These states can be efficiently found by taking a smaller deadlock, finding all previous states, and adding the missing stones in every possible location.  

The pseudo code for the first pass using the smaller deadlocks is:

Smaller deadlock search process

list smallerDeadlocksSearch (nbStones)  
list = empty list  
for P = 1; P < nbStones && P <= Pmax; P++  
    previous_states = empty list  
    for deadlock in deadlocksTree with nbStones-P stones  
        push deadlock in previous_states
for state in previous_states
    for move in possiblePreviousMoves (state)
        parent = apply (state, move)
        freeStates = free_states (parent)
        subState = first_combination (P, freeStates)
        while true
            newState = parent + subState
            if newState is not a known deadlock
                isDeadlock = 1
                for child in nextStates (newState)
                    if child is not a known deadlock
                        isDeadlock = 0
                        break
                if isDeadlock equals 1
                    add newState to deadlocksTree
                    push newState in list
                if subState equals last_combination (P, freeStates)
                    break
                subState = next_combination (subState, freeStates)
        return list

However the first pass using the smaller deadlock does not find all possible deadlocks. For example a few deadlocks do not have any possible moves and are not found by the smaller deadlock search.

Here are some examples of deadlocks that are not found:

```
################
#..@ # ### #@ # #@ #
#.. # # # # ## # # ##
#.. # #### # # # # #
#.. ## # ## # # # #
#.. # # $$ ## ##########$$ ###
######$##$ $ # #.... ## $ $ # #.... ## $ $ #
# $ $ # ##... # ##... $ #
# # # #.... ########## #.... ##########
################
```

Only very few deadlocks have no possible moves. However the algorithm also misses other deadlocks that are based on even smaller deadlocks that the one used.

The retrograde analysis algorithm we have used starts its first passes with all possible combinations of stones until a predefined number of stones (maxStonesCombination), and then continues with first passes based on smaller deadlocks:
Global deadlock search process

```
nbStones = 1
while not end of deadlock search
    if nbStones < maxStonesCombination
        newStates = combinationsDeadlocksSearch (nbStones)
    else
        newStates = smallerDeadlocksSearch (nbStones)
    while size of newStates is not 0
        otherStates = subsequentDeadlocksSearch (nbStones, newStates)
        newStates = otherStates
    nbStones++
```

Another way of looking at Sokoban problems is to start from the goal position and going backward to the start position pulling stones. In this case, the usual deadlocks never appear since the goal position is always reachable by pushing back stones. However other kinds of deadlocks appears which stop the backward search: the pull deadlocks.

Here are example of one stone, two stones and 6 stones pull deadlocks for screens 1 and 13:

```
#####
# @
# #
# #
# #
# #
# #
```
```
#####
# # # #
# # # #
# $ # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
```
```
#####
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
# # # # # # # # # # # # # # # # # # # # #
```

In the first one stone pull deadlock, the stone cannot be pulled. In the second two stones pull deadlock pulling the stones reaches a position where no more pulling moves are possible. In the last six stone pull deadlock a more complex retrograde analysis finds that all moves lead to position where no more moves are available, and the man is trapped in a local area.

5 Searching with deadlocks

In this section, we present the algorithms we use to solve problems. The first algorithm is a greedy search. It consists in developing the state that has the minimal associated heuristic. A greedy search is pointless in Sokoban without
deadlock verification since it is trapped in promising deadlock states for a very long time. When deadlocks can be detected a greedy search avoids such states and can find solutions.

The pseudo code for the greedy algorithm with deadlock verification is:

**Greedy algorithm**

```plaintext
void greedy (initialState)
    currentGreedy = manhattan (initialState)
    push initialState in sortedStates [currentGreedy]
    while not end of search
        while size of sortedStates [currentGreedy] is 0
            currentGreedy ++
        state = front of sortedStates [currentGreedy]
        pop front of sortedStates [currentGreedy]
        for child in nextStates (state)
            if child is without remaining stones
                return
            if child is known as deadlock
                continue
            h = manhattan (child)
            push child in front of sortedStates [h]
            if h < currentGreedy
                currentGreedy = h
```

The other algorithm we have tested is IDA*. We used a very simple admissible heuristic which consists in summing for all stones the distance to the closest goal location. Moreover we do not place stones inside the goal area and assume they are at their goal location as soon as they enter the area.

The greedy algorithm is very similar for pull deadlocks, except that the heuristic used is a greedy one, the closest pair of stone and start location is chosen, their distance is added to the heuristic, then the start location is removed from the possible locations and the process is repeated until all pairs have been chosen.

6 Results

In order to compute the deadlock tables, all combinations are tried until 4 stones, then for more stones, a smaller search first pass with \( P_{max} = 2 \) is performed, followed by usual passes. Table 1 gives the time used to build deadlock tables and the number of deadlocks found for screen 13. A distinction is made between the first pass and subsequent passes. For 5 stones, the combination search takes 40 minutes and the smaller search takes 10 seconds, the number of deadlock found in the first pass of the smaller search is higher than in the combination search since it the order in which it considers deadlocks is not the same and as it checks a deadlock as soon as it is found. For 10 stones, the smaller search uses 2h14m for the first pass and 1 minute for subsequent passes.
<table>
<thead>
<tr>
<th>screen 13</th>
<th>with combination</th>
<th>with smaller $P_{max} = 2$</th>
</tr>
</thead>
<tbody>
<tr>
<td>time nb passes</td>
<td>1st next</td>
<td>time nb passes</td>
</tr>
<tr>
<td>1</td>
<td>&lt;1s</td>
<td>2 40 0</td>
</tr>
<tr>
<td>2</td>
<td>&lt;1s</td>
<td>4 49 6</td>
</tr>
<tr>
<td>3</td>
<td>1s</td>
<td>5 59 20</td>
</tr>
<tr>
<td>4</td>
<td>40s</td>
<td>6 160 117</td>
</tr>
<tr>
<td>5</td>
<td>40m</td>
<td>7 417 345</td>
</tr>
<tr>
<td>6</td>
<td></td>
<td></td>
</tr>
<tr>
<td>7</td>
<td></td>
<td></td>
</tr>
<tr>
<td>8</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>screen 14</th>
<th>with combination</th>
<th>with smaller $P_{max} = 2$</th>
</tr>
</thead>
<tbody>
<tr>
<td>time nb passes</td>
<td>1st next</td>
<td>time nb passes</td>
</tr>
<tr>
<td>1</td>
<td>&lt;1s</td>
<td>2 36 0</td>
</tr>
<tr>
<td>2</td>
<td>&lt;1s</td>
<td>8 64 38</td>
</tr>
<tr>
<td>3</td>
<td>1s</td>
<td>8 68 38</td>
</tr>
<tr>
<td>4</td>
<td>41s</td>
<td>8 167 49</td>
</tr>
<tr>
<td>5</td>
<td>43m</td>
<td>10 436 248</td>
</tr>
<tr>
<td>6</td>
<td></td>
<td></td>
</tr>
<tr>
<td>7</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Table 1. Time needed and number of deadlocks found with deadlock search.

Table 2. Time needed and number of deadlocks found with deadlock search.

Table 2 gives similar information for the deadlocks of screen 14.

Table 3 gives the number of nodes required to solve screen 13 and screen 16 with a greedy search and different deadlock tables. Using deadlock tables with only three stones it does not solve the problem 16 in one million nodes. Using four stones deadlock tables, it solves screen 16 in 56941 nodes. The number of nodes decreases with the number of deadlock stones until 8 where it reaches its minimum at 623 nodes. Concerning screen 13, the problem can be solved with the 8 stones deadlock tables, computing more deadlock tables is of little practical interest given the computing times of the deadlock tables greater than 8 stones given in table 1.
Table 3. Greedy search with different numbers of deadlock tables.

<table>
<thead>
<tr>
<th>Screen 13</th>
<th>Screen 16</th>
</tr>
</thead>
<tbody>
<tr>
<td>Stones</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>3</td>
</tr>
<tr>
<td>6</td>
<td>4</td>
</tr>
<tr>
<td>7</td>
<td>5</td>
</tr>
<tr>
<td>8</td>
<td>6</td>
</tr>
<tr>
<td>9</td>
<td>7</td>
</tr>
<tr>
<td>10</td>
<td>8</td>
</tr>
<tr>
<td>11</td>
<td>9</td>
</tr>
<tr>
<td>12</td>
<td>10</td>
</tr>
<tr>
<td>Nodes</td>
<td></td>
</tr>
<tr>
<td>&gt;1M</td>
<td>&gt;1M</td>
</tr>
<tr>
<td>&gt;1M</td>
<td>56941</td>
</tr>
<tr>
<td>1414</td>
<td>7278</td>
</tr>
<tr>
<td>1407</td>
<td>1986</td>
</tr>
<tr>
<td>1372</td>
<td>1496</td>
</tr>
<tr>
<td>1371</td>
<td>623</td>
</tr>
<tr>
<td>1371</td>
<td>623</td>
</tr>
<tr>
<td>1371</td>
<td>623</td>
</tr>
</tbody>
</table>

Table 4 compares the different search algorithms on different screens. The greedy algorithm with deadlock tables solves most of the problems we have tested. It solves problems 13, 15, 16, 18, 20, 22, 24, 28, 29, 31, 32 and 44 which are not solved by Rolling Stone. The first column of the table is the number of the screen, the second column is the number of stones of the screen, the third column gives the number of deadlock stones in the tables (– if equals to the number of stones), the fourth column gives the number of nodes used by a greedy search without deadlock tables (search is stopped after 1M nodes, so 1M means the search has failed), the fifth column gives the number of remaining stones for the position in the greedy search that has the minimum such number, the sixth column gives the number of nodes used by a greedy search with deadlock tables, the seventh column the related number of remaining stones, the eighth column the number of IDA* without deadlock tables, the ninth column the related number of remaining stones, the tenth column the number of IDA* with deadlock tables and the last column the related number of remaining stones.

Our current greedy search does not manage the placement of the stones in the goal area, as soon as a stone is in the goal area it is removed and considered as placed in the right location of the goal area. For many problems this simplification does not change the solution of the problem, however in some screens there can be interactions between the goal area and the solution or there can be multiple goal areas (for example in screen 18 there are three goal areas so our solution is not valid, in screen 15 there are two stones in the goal areas in the start position and they interact with the solution so our greedy solution is also not valid). Our simplification poses no problem in screens with a simple goal area such as screens 13, 20, 24, 28 and 31 which are solved by greedy search with deadlock tables and not by Rolling Stone.

Deadlock tables were added to Rolling Stone. The IDA* search was stopped as soon as a deadlock was encountered. The results are given in table 5. The first column gives the screen number, the second column gives the number of nodes used by the original Rolling Stone, the third column gives the number of nodes used by Rolling Stone equipped with deadlock tables. Deadlock tables reduce the number of nodes for almost all problems, problem 8 being an exception. A
A possible explanation is that in this problem some penalty patterns could not be found by Rolling Stone due to deadlock tables cuts and could not be reused in other parts of the search.

Table 6 gives the number of pull deadlocks and their computation times for screen 13.

A greedy search with pull deadlocks can be trapped in position that are close to the goal position but that are not solvable. Here is an example for screen 1 where the search is one move away from the start but is in a hopeless position, moreover the man cannot move to its start position:

```
#####  #####
#  #  #  #
#$  #  #  #
### $## ### $##
# $ $ # # $ $ #
### # ## # ###### ### # ## # ######
# # ## ##### ..# # # ## ##### ..#
# $ $ ..# # $ $ ..#
##### ### #@## ..# ##### ### # ## ..#
# ######### # #########
####### #######
```

Table 6 gives the number of pull deadlocks and their computation times for screen 13.

A greedy search with pull deadlocks can be trapped in position that are close to the goal position but that are not solvable. Here is an example for screen 1 where the search is one move away from the start but is in a hopeless position, moreover the man cannot move to its start position:

```
#####
#  #
#$  #
### $##
# $ #
### # ## #
# # # # # #
#$ # # #
### $##
```

7 Conclusion

We have presented how to use retrograde analysis to compute maze specific deadlock tables at Sokoban, we have also presented how to use deadlock trees to match the computed deadlocks efficiently as well as a greedy search algorithm. Moreover we have presented a new way to solve Sokoban problems pulling stones out from the goal area towards their start positions as well as the related pull deadlocks.

These algorithms have been tested on many Sokoban screens including difficult ones. Computing maze specific deadlock tables has enabled to successfully use a greedy search at Sokoban and to solve problems that are not solved by IDA*.

In future works, we will improve greedy search so as to cope with the placement of stones in the goal area. We will also try to improve the retrograde analysis efficiency and completeness.

References

Table 4. Number of nodes evaluated and resulting top level node needed to solve problems.

<table>
<thead>
<tr>
<th>Screen</th>
<th>MZ</th>
<th>DL</th>
<th>greedy basicDL</th>
<th>DL</th>
<th>basicDL</th>
<th>DL</th>
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<tbody>
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<td>1</td>
<td>6</td>
<td>–</td>
<td>643 (0) 400 (0)</td>
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<td>1M (1)</td>
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<tr>
<td>2</td>
<td>10</td>
<td>–</td>
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<td>1M (6)</td>
<td>1M (4)</td>
<td></td>
</tr>
<tr>
<td>3</td>
<td>11</td>
<td>–</td>
<td>1M (10) 490 (0)</td>
<td>1M (3)</td>
<td>723410 (0)</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td>20</td>
<td>–</td>
<td>1M (20) 1981 (0)</td>
<td>1M (20)</td>
<td>1M (5)</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>12</td>
<td>–</td>
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<td>1M (4)</td>
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<td></td>
</tr>
<tr>
<td>6</td>
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<td>–</td>
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<td>1M (1)</td>
<td>150 (0)</td>
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<td>7</td>
<td>11</td>
<td>–</td>
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<td>1M (2)</td>
<td>40888 (0)</td>
<td></td>
</tr>
<tr>
<td>8</td>
<td>18</td>
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<td>1M (12)</td>
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<td></td>
</tr>
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<td>–</td>
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<td>1M (2)</td>
<td></td>
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<td>14</td>
<td>11</td>
<td>1M (14) 62345 (0)</td>
<td>1M (13)</td>
<td>1M (13)</td>
<td></td>
</tr>
<tr>
<td>12</td>
<td>15</td>
<td>13</td>
<td>1M (14) 1440 (0)</td>
<td>1M (10)</td>
<td>1M (10)</td>
<td></td>
</tr>
<tr>
<td>13</td>
<td>16</td>
<td>12</td>
<td>1M (16) 1371 (0)</td>
<td>1M (7)</td>
<td>1M (6)</td>
<td></td>
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<tr>
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<td>18</td>
<td>11</td>
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<td>1M (14)</td>
<td>1M (9)</td>
<td></td>
</tr>
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<td>15</td>
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<td>1M (9) 769 (0)</td>
<td>1M (7)</td>
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<td>15</td>
<td>14</td>
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<td>1M (10)</td>
<td>1M (7)</td>
<td></td>
</tr>
<tr>
<td>18</td>
<td>11</td>
<td>10</td>
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<td>1M (4)</td>
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<td>22</td>
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<td>1M (17)</td>
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<td>10</td>
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<td>1M (18)</td>
<td>1M (15)</td>
<td></td>
</tr>
<tr>
<td>28</td>
<td>20</td>
<td>12</td>
<td>1M (20) 9691 (0)</td>
<td>1M (15)</td>
<td>1M (8)</td>
<td></td>
</tr>
<tr>
<td>29</td>
<td>16</td>
<td>11</td>
<td>1M (9) 396 (0)</td>
<td>1M (4)</td>
<td>1M (2)</td>
<td></td>
</tr>
<tr>
<td>31</td>
<td>20</td>
<td>10</td>
<td>1M (19) 1800 (0)</td>
<td>1M (20)</td>
<td>1M (11)</td>
<td></td>
</tr>
<tr>
<td>32</td>
<td>15</td>
<td>13</td>
<td>1M (11) 646 (0)</td>
<td>1M (5)</td>
<td>1M (5)</td>
<td></td>
</tr>
<tr>
<td>44</td>
<td>9</td>
<td>–</td>
<td>1M (7) 908532 (0)</td>
<td>1M (6)</td>
<td>1M (5)</td>
<td></td>
</tr>
</tbody>
</table>

Table 5. Total nodes searched with Rolling Stone.

<table>
<thead>
<tr>
<th>screen</th>
<th>without deadlock tables</th>
<th>with deadlock tables</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1042</td>
<td>384</td>
</tr>
<tr>
<td>2</td>
<td>7530</td>
<td>6366</td>
</tr>
<tr>
<td>3</td>
<td>12902</td>
<td>7181</td>
</tr>
<tr>
<td>4</td>
<td>50369</td>
<td>50369</td>
</tr>
<tr>
<td>5</td>
<td>43294</td>
<td>20483</td>
</tr>
<tr>
<td>6</td>
<td>5118</td>
<td>4028</td>
</tr>
<tr>
<td>7</td>
<td>28460</td>
<td>25954</td>
</tr>
<tr>
<td>8</td>
<td>336495</td>
<td>408123</td>
</tr>
<tr>
<td>9</td>
<td>436801</td>
<td>343591</td>
</tr>
<tr>
<td>11</td>
<td>3091658</td>
<td>2323731</td>
</tr>
<tr>
<td>12</td>
<td>559216</td>
<td>556353</td>
</tr>
<tr>
<td>13</td>
<td>&gt;20M</td>
<td>&gt;20M</td>
</tr>
<tr>
<td>44</td>
<td>&gt;20M</td>
<td>&gt;20M</td>
</tr>
</tbody>
</table>

Table 6. Time needed and number of pull deadlocks found with pull deadlock search.

<table>
<thead>
<tr>
<th>screen</th>
<th>with combination</th>
<th>with smaller $P_{max} = 2$</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>time nb passes 1st next</td>
<td>time nb passes 1st next</td>
</tr>
<tr>
<td>1</td>
<td>&lt;1s</td>
<td>1 0 –</td>
</tr>
<tr>
<td>2</td>
<td>&lt;1s</td>
<td>3 35 2</td>
</tr>
<tr>
<td>3</td>
<td>8s</td>
<td>2 78 0</td>
</tr>
<tr>
<td>4</td>
<td>10m51s</td>
<td>2 234 0</td>
</tr>
<tr>
<td>5</td>
<td>16s</td>
<td>10 195 400</td>
</tr>
<tr>
<td>6</td>
<td>76s</td>
<td>10 714 1023</td>
</tr>
</tbody>
</table>


