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A Connected Version of the Graph Coloring Game

Clément Charpentier¹ Hervé Hocquard² Éric Sopena² Xuding Zhu³

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Abstract

The graph coloring game is a two-player game in which, given a graph G and a set of k colors, the two players, Alice and Bob, take turns coloring properly an uncolored vertex of G , Alice having the first move. Alice wins the game if and only if all the vertices of G are eventually colored. The game chromatic number of a graph G is then defined as the smallest integer k for which Alice has a winning strategy when playing the graph coloring game on G with k colors.

In this paper, we introduce and study a new version of the graph coloring game by requiring that, after each player's turn, the subgraph induced by the set of colored vertices is connected. The connected game chromatic number of a graph G is then the smallest integer k for which Alice has a winning strategy when playing the connected graph coloring game on G with k colors. We prove that the connected game chromatic number of every outerplanar graph is at most 5 and that there exist outerplanar graphs with connected game chromatic number 4.

Moreover, we prove that for every integer $k \geq 3$, there exist bipartite graphs on which Bob wins the connected coloring game with k colors, while Alice wins the connected coloring game with two colors on every bipartite graph.

Keywords: Coloring game; Marking game; Game coloring number; Game chromatic number.

MSC 2010: 05C15, 05C57, 91A43.

1 Introduction

All the graphs we consider in this paper are undirected, simple, and have no loops. For every such graph G , we denote by $V(G)$ and $E(G)$ its vertex set and edge set, respectively, by $\Delta(G)$ its maximum degree, and by $N_G(v)$ the set of neighbors of the vertex v in G .

The *graph coloring game* is a two-player game introduced by Steven J. Brams (reported by Martin Gardner in his column *Mathematical Games* in *Scientific American* in 1981 [9]) and rediscovered ten years later by Bodlaender [4]. Given a graph G and a set \mathcal{C} of k colors, the two players, Alice and Bob, take turns coloring properly an uncolored vertex of G , Alice having the first move. Alice wins the game if and only if all the vertices of G are eventually colored. In other words, Bob wins the game if and only if, at some step of the game, all the colors appear in the neighborhood of some uncolored vertex.

The *game chromatic number* $\chi_g(G)$ of G is then defined as the smallest integer k for which Alice has a winning strategy when playing the graph coloring game on G with k colors. The problem of determining the game chromatic number of several graph classes has attracted much interest in recent

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years (see [2] for a comprehensive survey of this problem), with a particular focus on planar graphs (see e.g. [6, 13, 14, 19–21]) for which the best known upper bound up to now is 17 [21].

Most of the known upper bounds on the game chromatic number of classes of graphs are derived from upper bounds on the game coloring number of these classes, a parameter defined through the so-called *graph marking game*, formally introduced by Zhu in [20]. This game is somehow similar to the graph coloring game, except that the players mark the vertices instead of coloring them, with no restriction. The *game coloring number* $\text{col}_g(G)$ of G is then defined as the smallest integer k for which Alice has a strategy such that, when playing the graph marking game on G , every unmarked vertex has at most $k - 1$ marked neighbors. It is worth noting here that the game coloring number is *monotonic*, which means that $\text{col}_g(H) \leq \text{col}_g(G)$ for every subgraph H of G , while this property does not hold for the game chromatic number [18].

Let G be a graph with $\text{col}_g(G) = k$ and consider the winning strategy of Alice for the marking game on G . Applying the same strategy for the coloring game on G , Alice ensures that each uncolored vertex has at most $k - 1$ colored neighbors, so that we get $\chi_g(G) \leq k$. Hence, the following inequalities clearly hold for every graph G .

Observation 1 *For every graph G , $\chi(G) \leq \chi_g(G) \leq \text{col}_g(G) \leq \Delta(G) + 1$.*

In this paper, we introduce and study a new version of the graph coloring game (resp. of the graph marking game), by requiring that, after each player’s turn, the subgraph induced by the set of colored (resp. marked) vertices is connected. In other words, on their turn, each player must color an uncolored vertex (resp. mark an unmarked vertex) having at least one colored (resp. marked) neighbor, except for Alice on her first move.

We call this new game the *connected graph coloring game* (resp. the *connected graph marking game*). We will denote by $\chi_{cg}(G)$ the *connected game chromatic number* of a graph G , that is, the smallest integer k for which Alice has a winning strategy when playing the connected graph coloring game on G with k colors, and by $\text{col}_{cg}(G)$ the *connected game coloring number* of G , that is, the smallest integer k for which Alice has a strategy such that, when playing the connected graph marking game on G , every unmarked vertex has at most $k - 1$ marked neighbors. It is not difficult to observe that, similarly to the ordinary case, the following inequalities hold for every graph G .

Observation 2 *For every graph G , $\chi(G) \leq \chi_{cg}(G) \leq \text{col}_{cg}(G) \leq \Delta(G) + 1$.*

It is proved in [18] that for every positive integer n , $\chi_g(K_{n,n} - M) = n$, where $K_{n,n} - M$ denotes the complete bipartite graph with n vertices in each part, minus a perfect matching. We prove in Section 2 that $\chi_{cg}(G) = 2$ for every nonempty bipartite graph G , which shows, since the graph $K_{n,n} - M$ is bipartite, that the difference $\chi_g(G) - \chi_{cg}(G)$ can be arbitrarily large.

One of the main open, and rather intriguing, question concerning the graph coloring game is the following: assuming that Alice has a winning strategy for the graph coloring game on a graph G with k colors, is it true that she has also a winning strategy with $k + 1$ colors? We will prove in Section 2 that the answer is “no” for the connected version of the coloring game. More precisely, we will prove that for every integer $k \geq 3$, there exist bipartite graphs on which Bob wins the connected coloring game with k colors, while Alice wins the connected coloring game with two colors on every bipartite graph.

The “connected variant” of other types of games on graphs have been considered in the literature. This is the case for instance for the domination game [5, 12], the surveillance game [7, 10], the graph searching game [1, 3, 8], or Hajnal’s triangle-free game [16, 17]. However, to our knowledge, the connected variant of the graph coloring game has not been considered yet.

Our paper is organized as follows. We consider bipartite graphs in Section 2, and outerplanar graphs in Section 3. We finally propose some directions for future research in Section 4.

2 Bipartite graphs

We consider the case of bipartite graphs in this section. We will prove that for every integer $k \geq 3$, there exist bipartite graphs on which Bob wins the connected coloring game with k colors, while Alice wins the connected coloring game with two colors on every bipartite graph.

It is easy to see that Alice always wins when playing the connected coloring game on a bipartite graph G with two colors: thanks to the connectivity constraint, the first move of Alice forces all the next moves to be consistent with a proper 2-coloring of G .

Theorem 3 *For every bipartite graph G , $\chi_{\text{cg}}(G) \leq 2$.*

Proof. Let G be any bipartite graph. The strategy of Alice is as follows. On her first move, she picks any vertex v of G and gives it color 1. From now on, each play will color some vertex having at least one of its neighbors already colored, so that, since G is bipartite, this eventually leads to a proper 2-coloring of G . \square

However, for every integer $k \geq 3$, there are bipartite graphs on which Bob wins the connected coloring game with k colors.

Theorem 4 *For every integer $k \geq 3$, there exists a bipartite graph G_k on which Bob wins the connected coloring game with k colors.*

Proof. Let H_k be any C_4 -free bipartite graph with minimum degree k^2 and let A and B denote the partite sets of H_k . Let now G_k be the bipartite graph obtained from H_k by adding, for each k -subset S of B , a new vertex v_S adjacent to all vertices of S .

We now define the strategy of Bob for playing the connected coloring game on G_k as follows. In his first moves (at most three, depending on the moves of Alice), Bob colors two vertices of A , say u and v , with two different colors. In his next two moves, Bob colors a neighbor u' of u in B with the same color as v , and a neighbor v' of v in B with the same color as u . Since the minimum degree of H_k is k^2 and H_k is C_4 -free, Alice cannot prevent Bob from doing so.

Now, Bob colors a k -subset $X \subseteq N_{H_k}(u) \cup N_{H_k}(v)$ containing u' and v' with k distinct colors. Again, Alice cannot prevent Bob from doing so since each move of Alice “forbids” at most k uncoloured vertices in $N_{H_k}(u) \cup N_{H_k}(v)$ (each vertex of G_k has at most k neighbors in this set).

After that, the vertex v_X cannot be colored and Bob wins the game. \square

3 Outerplanar graphs

We consider in this section the case of outerplanar graphs. An *outerplanar graph* is a graph that can be embedded on the plane in such a way that there are no edge crossings and all its vertices lie on the outer face.

Concerning the ordinary coloring game, Kierstead and Trotter proved in [14] that there exist outerplanar graphs with game chromatic number at least 6, and Guan and Zhu proved in [11] that the game chromatic number of every outerplanar graph is at most 7. This bound has then been proven to be tight by Kierstead and Yang in [15]. We will prove that the connected game chromatic number of every outerplanar graph is at most 5 and that there exist outerplanar graphs with connected game chromatic number 4.

Recall that an outerplanar graph is *maximal* if adding any edge makes it non outerplanar. It is not difficult to see that an outerplanar graph is maximal if and only if, in all its outerplanar embeddings, all faces are triangles, except possibly the outer face. Our goal in this section is to prove that the connected coloring number of every outerplanar graph is at most 5.

When playing the connected coloring game on a graph G , we will say that an uncolored vertex in G is *saturated* if each of the available colors appears in its neighborhood. Observe that Bob wins the

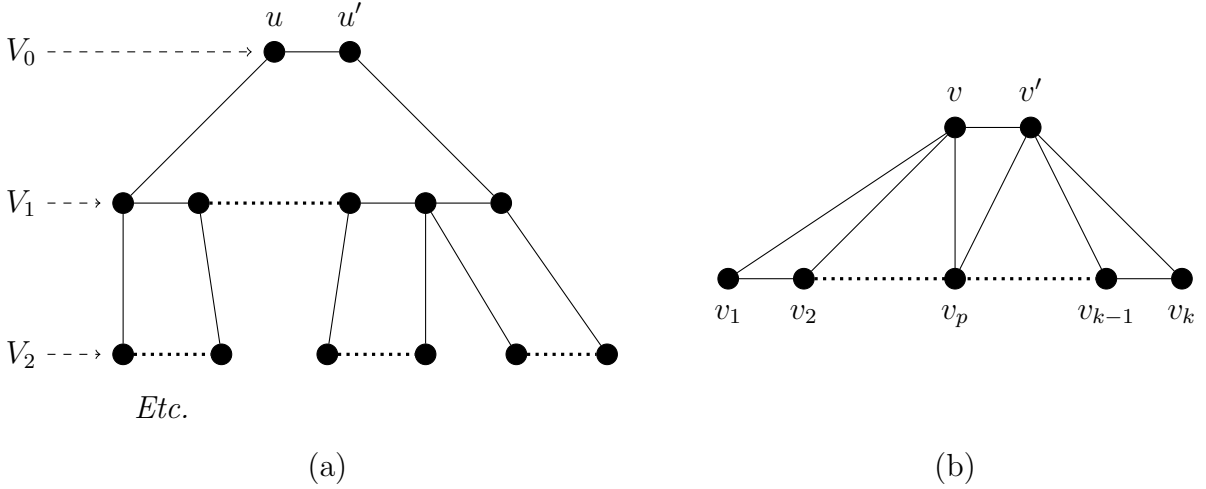


Figure 1: The structure of a maximal outerplanar graph.

connected coloring game on G if and only if he has a strategy such that, at some point in the game, an uncolored vertex in G becomes saturated. Similarly, when trying to prove that the connected game coloring number of some graph G is at most k , we will say that an unmarked vertex in G is saturated if it has at least k marked neighbors. Again, the connected game coloring number of G is at least $k + 1$ if and only if Bob has a strategy such that, at some point in the game, an unmarked vertex in G becomes saturated.

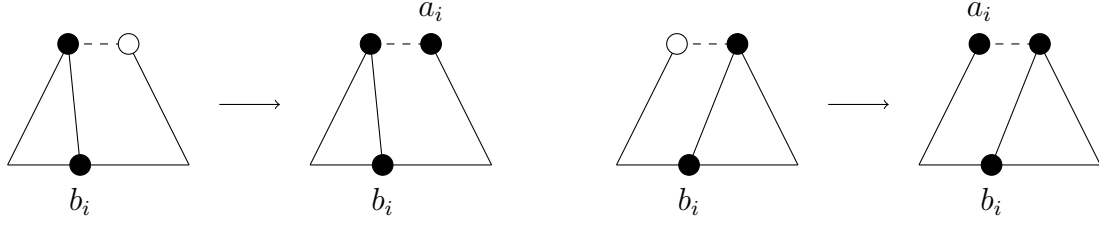
Finally, we will say that a vertex in G is *playable* if it is uncolored (resp. unmarked) and has at least one colored (resp. marked) neighbor. Moreover, when considering the connected marking game, we will say that a vertex is *threatened* if it is unmarked, has $k - 1$ differently marked neighbors and at least one playable neighbor. In that case, note that if Bob plays on a playable neighbor of any threatened vertex, then Alice loses the game. A winning strategy of Alice for the connected marking game thus consists in ensuring that, after each of her moves, the considered graph has no threatened vertex.

We first describe more precisely the structure of maximal outerplanar graphs, which will be used for defining the strategy of Alice. Let G be a maximal (embedded) outerplanar graph. An edge belonging to the outer face of G is an *outer edge* of G . Let us choose and fix any outer edge $e = uu'$ of G . The *distance* from any vertex v to the edge e is defined as $d_G(v, e) = \min \{d_G(v, u), d_G(v, u')\}$. For every integer $i \geq 0$, let V_i denote the set of vertices at distance i from e . It is not difficult to observe that the subgraph $G[V_i]$ of G induced by each set V_i is a linear forest, that is, a disjoint union of paths. In particular, $G[V_0]$ is the edge uu' and $G[V_1]$ is a single path. Therefore, G can be viewed as a “tree of trapezoids”, as illustrated in Figure 1(a).

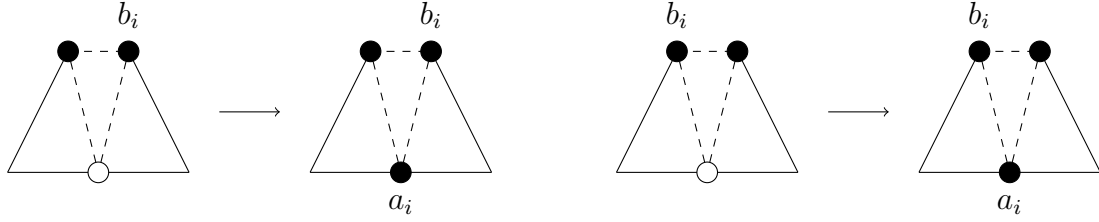
Each of these trapezoids has the structure depicted in Figure 1(b). Both v and v' belong to some V_i , $i \geq 0$, while v_1, \dots, v_k , $k \geq 2$, belong to V_{i+1} . The vertices v and v' are the *parents* of the *children vertices* v_1, \dots, v_k , the edge vv' is the *root edge* of the trapezoid, the unique vertex v_p , $1 \leq p \leq k$, which is joined by an edge to both v and v' is the *pivot* of the trapezoid. We will denote by $T_{vv'}$ the trapezoid whose root edge is vv' , and by $p(vv')$ the pivot of $T_{vv'}$. Note that each vertex v_i , $1 \leq i \leq k$, is a neighbor of at least one of its parents, and that only the pivot v_p is a neighbor of both its parents. Moreover, we will say that v_{i-1} , $2 \leq i \leq k$, is the *left neighbor* of v_i , while v_{i+1} , $1 \leq i \leq k - 1$, is the *right neighbor* of v_i .

Observe that if there is no trapezoid of the form $T_{vv'}$ or $T_{v'v}$ in G for some vertex v , then the degree of v is at most 4 (it is 4 only if v is the pivot of some trapezoid), so that v cannot be a threatened vertex. Note also that every vertex belongs to at most two root edges.

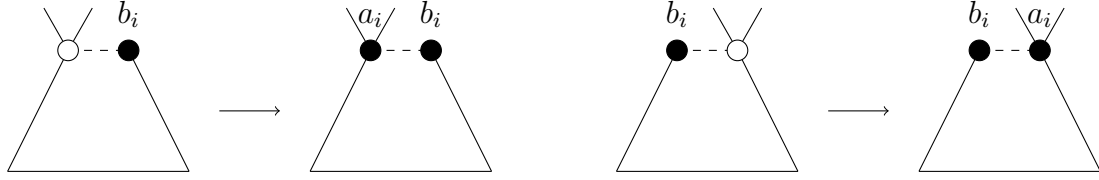
Based on the drawing of the outerplanar graph depicted in Figure 1(a), we can define a total



Rule R1 (the vertex a_i must be playable)



Rule R2 (a_i is a pivot and at least one of the dashed edges incident with a_i must exist)



Rule R3 (the vertex a_i must be playable)

Figure 2: The strategy of Alice on outerplanar graphs (Rules R1, R2 and R3).

ordering \leq_G of the vertices of G , obtained by listing the vertices of V_0 from left to right, then the vertices of V_1 from left to right, and so on. Finally, we will say that a vertex w_1 , belonging to a trapezoid $T_{v_1v'_1}$, lies above a vertex w_2 , belonging to a trapezoid $T_{v_2v'_2}$, if every shortest path from $\{v_2, v'_2\}$ to $\{u, u'\}$ goes through v_1 or v'_1 .

We now describe the strategy of Alice when playing the connected marking game on an outerplanar graph G . Let uu' be any outer edge of G , and G_m be any maximal outerplanar graph containing G as a subgraph, and such that uu' is also an outer edge of G_m . In the following, we assume that we are given a trapezoidal representation of G_m , starting from the edge uu' , as described above. Moreover, we can also assume that for every trapezoid $T_{vv'}$ of G_m , at least one of the edges $vp(vv')$, $v'p(vv')$ belongs to G . This will allow us to speak about children or parent vertices (with respect to G_m) even if the corresponding edges do not belong to G , and to use the ordering \leq_{G_m} of the vertices of G .

Let us denote by a_i , $i \geq 0$, the vertex marked by Alice on her $(i+1)$ -th move, and by b_i , $i \geq 1$, the vertex marked by Bob on his i -th move, so that the sequence of moves (that is, marked vertices) is $a_0, b_1, a_1, \dots, b_i, a_i, \dots$. Hence, a_0 is the vertex marked by Alice on her first move and, for every $i \geq 1$, a_i is the “response” of Alice to the move b_i of Bob.

The strategy of Alice will then consist in applying the first of the following rules that can be applied (see Figure 2 for an illustration of Rules R1, R2 and R3) for each of her moves.

R0: $a_0 := u$.

R1: If v is a playable unmarked parent of b_i , then $a_i := v$.

R2: If b_i belongs to a root edge vb_i or b_iv , v is marked and $p(vb_i)$ is playable, then $a_i := p(vb_i)$.

- R3: If b_i belongs to a root edge vb_i or b_iv , v is unmarked, v is a pivot and v is playable, then $a_i := v$.
- R4: If none of the above rules can be applied, and there are still unmarked vertices in G , then we let $a_i := w$, where w is the smallest (with respect to the ordering \leq_{G_m}) playable vertex.

Note that on his first move, Bob must mark either the vertex u' , in which case Alice will apply Rule R2 on her second move, or some neighbor $v \neq u'$ of u , in which case Alice will apply rule R1 and mark u' on her second move (recall that the edge uu' belongs to G). Moreover, if Bob marks a child vertex w of some trapezoid $T_{vv'}$, then at least one of v, v' must be marked (by the connectivity constraint), and Alice will immediately apply Rule R1 if one of them is unmarked and vv' is an edge in G . These remarks are summarized in the two following observations.

Observation 5 *After the second move of Alice, both vertices u and u' are marked.*

Observation 6 *After each move of Alice, if w is a marked child vertex of a trapezoid $T_{vv'}$ and vv' is an edge in G , then both v and v' are marked.*

We are now able to prove the main result of this section.

Theorem 7 *If G is an outerplanar graph, then $\text{col}_{\text{cg}}(G) \leq 5$.*

Proof. We assume that we are given an outerplanar embedding of G_m and its trapezoidal representation, as previously discussed. Clearly, it suffices to prove that if Alice applies the above described strategy, then, after each move of Alice, G contains no threatened vertex. This is clearly the case after the first and second move of Alice since, at that point, only one or three vertices have been marked, respectively.

Suppose to the contrary that, after Bob has marked the vertex b_i and Alice has marked the vertex a_i , $i \geq 2$, t is a threatened vertex in G , and that i is the smallest index with this property, which implies that a_i or b_i is a marked neighbor of t . Thanks to Observation 5, we know that both u and u' have been marked. Therefore, t is necessarily a child vertex of some trapezoid $T_{vv'}$ (we may have $vv' = uu'$). Let t^ℓ and t^r denote the left and right neighbors of t (in G_m), if they exist. Note that at least one of them must exist, since otherwise t would have at most two marked neighbors, and thus could not be a threatened vertex. Since t has four marked neighbors, at least one of t^ℓ, t^r must be marked, since otherwise no vertex lying below t could have been marked, due to the connectivity constraint, so that, again, t would have at most two marked neighbors. Thanks to Observation 6, we thus get that both v and v' are marked if vv' is an edge in G .

We now claim that neither $T_{t^\ell t}$ nor T_{tt^r} contains a marked child vertex which is a neighbor of t . Indeed, such a vertex, say w , cannot have been marked by Bob since, by Rule R1, Alice would have marked t just after Bob had marked the first such child vertex of the corresponding trapezoid. The vertex w has thus been marked by Alice which implies, since t is unmarked, that none of the edges tv, tv', tt^ℓ or tt^r belong to G (otherwise t would have been marked in priority by Alice), and that w is the only marked neighbor of t , so that t cannot be a threatened vertex.

Therefore, the four marked neighbors of w are necessarily v, v', t^ℓ and t^r . Hence, t is the pivot of $T_{vv'}$, which implies, since t is unmarked, that t^ℓ has been marked after v , and that t^r has been marked after v' , so that $b_i \in \{t^\ell, t^r\}$. (Note here that we cannot have $b_i \in \{v, v'\}$, since this would imply $a_i \in \{t^\ell, t^r\}$, contradicting the priority of rule R2.) But in each case, that is, $b_i = t^\ell$ or $b_i = t^r$, t would have been marked by Alice, thanks to Rule R3.

We thus get a contradiction in each case, which concludes the proof of Theorem 7. \square

Concerning the connected game chromatic number of outerplanar graphs, we can now prove the following.

Theorem 8 *If G is an outerplanar graph, then $\chi_{\text{cg}}(G) \leq 5$. Moreover, there exist outerplanar graphs with $\chi_{\text{cg}}(G) = 4$.*

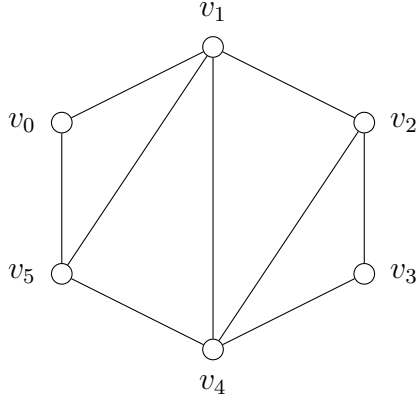


Figure 3: An outerplanar graph with connected game chromatic number 4.

Proof. From Observation 2 and Theorem 7, we get $\chi_{\text{cg}}(G) \leq \text{col}_{\text{cg}}(G) \leq 5$. For the second part of the statement, consider the outerplanar graph G depicted in Figure 3. We will prove that Bob has a winning strategy when playing the connected coloring game on G with three colors. Thanks to the symmetries in G , and up to permutation of colors, Alice has three possible first moves, that we consider separately.

1. If Alice colors v_0 with color 1, then Bob colors v_1 with color 2. Now, if Alice colors v_5 with color 3 then Bob colors v_2 with color 1 so that v_4 is saturated, while if Alice colors v_2 or v_4 with color 1 (resp. with color 3), then Bob colors v_3 with color 3 (resp. with color 1), so that v_4 or v_2 is saturated.
2. If Alice colors v_1 with color 1, then Bob colors v_0 with color 2, and the so-obtained configuration is similar to that of the previous case.
3. If Alice colors v_2 with color 1, then Bob colors v_3 with color 2. Now, if Alice colors v_4 with color 3 then Bob colors v_5 with color 2 so that v_1 is saturated, while if Alice colors v_1 with color 2 (note that using color 3 would saturate v_4), then Bob colors v_5 with color 3, so that v_4 is saturated.

This concludes the proof of Theorem 8. □

By Observation 2, the second part of the statement of Theorem 8 directly implies the following.

Corollary 9 *There exist outerplanar graphs with $\text{col}_{\text{cg}}(G) \geq 4$.*

4 Discussion

We have introduced in this paper a connected version of the graph coloring and graph marking games. We have proved in particular that the connected game coloring number of every outerplanar graph is at most 5, and that there exist infinitely many bipartite graphs on which Alice wins the connected coloring game with two colors but loses the game if the number of colors is at least three.

We conclude this paper by listing some open questions that should be considered for future work.

1. What is the optimal upper bound on the connected game coloring number and on the connected game chromatic number of outerplanar graphs? We know that both these values are either 4 or 5.
2. What is the optimal upper bound on the connected game coloring number and on the connected game chromatic number of planar graphs?

3. Does there exist, for every integers $k \geq 3$ and $p \geq 1$, a graph $G_{k,p}$ on which Alice wins the connected coloring game with k colors, while Bob wins the game with $k + p$ colors?
4. Is the connected game coloring number a monotonic parameter, that is, is it true that for every subgraph H of G , the inequality $\text{col}_{\text{cg}}(H) \leq \text{col}_{\text{cg}}(G)$ holds?
5. Does there exist a graph G for which $\chi_{\text{g}}(G) < \chi_{\text{cg}}(G)$? or $\text{col}_{\text{g}}(G) < \text{col}_{\text{cg}}(G)$? (That is, is it possible that the connectivity constraint is in favour of Bob?)

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