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Covering with Clubs: Complexity and Approximability

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Abstract. Finding cohesive subgraphs in a network is a well-known problem in graph theory. Several alternative formulations of cohesive subgraph have been proposed, a notable example being s -club, which is a subgraph where each vertex is at distance at most s to the others. Here we consider the problem of covering a given graph with the minimum number of s -clubs. We study the computational and approximation complexity of this problem, when s is equal to 2 or 3. First, we show that deciding if there exists a cover of a graph with three 2-clubs is NP-complete, and that deciding if there exists a cover of a graph with two 3-clubs is NP-complete. Then, we consider the approximation complexity of covering a graph with the minimum number of 2-clubs and 3-clubs. We show that, given a graph $G = (V, E)$ to be covered, covering G with the minimum number of 2-clubs is not approximable within factor $O(|V|^{1/2-\varepsilon})$, for any $\varepsilon > 0$, and covering G with the minimum number of 3-clubs is not approximable within factor $O(|V|^{1-\varepsilon})$, for any $\varepsilon > 0$. On the positive side, we give an approximation algorithm of factor $2|V|^{1/2} \log^{3/2} |V|$ for covering a graph with the minimum number of 2-clubs.

1 Introduction

The quest for modules inside a network is a well-known and deeply studied problem in network analysis, with several application in different fields, like computational biology or social network analysis. A highly investigated problem is that of finding cohesive subgroups inside a network which in graph theory translates in highly connected subgraphs. A common approach is to look for cliques (i.e. complete graphs), and several combinatorial problems have been considered, notable examples being the Maximum Clique problem ([11, GT19]), the Minimum Clique Cover problem ([11, GT17]), and the Minimum Clique Partition problem ([11, GT15]). This last is a classical problem in theoretical computer science, whose goal is to partition the vertices of a graph into the minimum number of cliques. The Minimum Clique Partition problem has been deeply studied since the seminal paper of Karp [15], studying its complexity in several graph classes [5,6,21,9].

In some cases, asking for a complete subgraph is too restrictive, as interesting highly connected graphs may have some missing edges due to noise in the data considered or because some pair may not be directly connected by an edge in the subgraph of interest. To overcome this limitation of the clique approach, alternative definitions of highly connected graphs have been proposed, leading to the concept of *relaxed clique* [16]. A relaxed clique is a graph $G = (V, E)$ whose vertices satisfy a property which is a relaxation of the clique property. Indeed, a clique is a subgraph whose vertices are all at distance one from each other and have the same degree (the size of the clique minus one). Different definitions of relaxed clique are obtained by modifying one of the properties of clique, thus leading to distance-based relaxed cliques, degree-based relaxed cliques, and so on (see for example [16]).

In this paper, we focus on a distance-based relaxation. In a clique all the vertices are required to be at distance at most one from each other. Here this constraint is relaxed, so that the vertices have to be at distance at most s , for an integer $s \geq 1$. A subgraph whose vertices are all distance at most s is called an s -club (notice that, when $s = 1$, an s -club is exactly a clique). The identification of s -clubs inside a network has been applied to social networks [19,1,18,20,23], and biological networks [3]. Interesting recent studies have shown the relevance of finding s -clubs

in a network [18,20], in particular focusing on finding 2-clubs in real networks like DBLP or a European corporate network.

Contributions to the study of s -clubs mainly focus on the Maximum s -Club problem, that is the problem of finding an s -club of maximum size. Maximum s -Club is known to be NP-hard, for each $s \geq 1$ [4]. Even deciding whether there exists an s -club larger than a given size in a graph of diameter $s + 1$ is NP-complete, for each $s \geq 1$ [3]. The Maximum s -Club problem has been studied also in the approximability and parameterized complexity framework. A polynomial-time approximation algorithm with factor $|V|^{1/2}$ for every $s \geq 2$ on an input graph $G = (V, E)$ has been designed [2]. This is optimal, since the problem is not approximable within factor $|V|^{1/2-\varepsilon}$, on an input graph $G = (V, E)$, for each $\varepsilon > 0$ and $s \geq 2$ [2]. As for the parameterized complexity framework, the problem is known to be fixed-parameter tractable, when parameterized by the size of an s -club [22,17,7]. The Maximum s -Club problem has been investigated also for structural parameters and specific graph classes [13,12].

In this paper, we consider a different combinatorial problem, where we aim at covering the vertices of a network with a set of subgraphs. Similar to Minimum Clique Partition, we consider the problem of covering a graph with the minimum number of s -clubs such that each vertex belongs to an s -club. We denote this problem by Min s -Club Cover, and we focus in particular on the cases $s = 2$ and $s = 3$. We show some analogies and differences between Min s -Club Cover and Minimum Clique Partition. We start in Section 3 by considering the computational complexity of the problem of covering a graph with two or three s -clubs. This is motivated by the fact that Clique Partition is known to be in P when we ask whether there exists a partition of the graph consisting of two cliques, while it is NP-hard to decide whether there exists a partition of the graph consisting of three cliques [10]. As for Clique Partition, we show that it is NP-complete to decide whether there exist three 2-clubs that cover a graph. On the other hand, we show that, unlike Clique Partition, it is NP-complete to decide whether there exist two 3-clubs that cover a graph. These two results imply also that Min 2-Club Cover and Min 3-Club Cover do not belong to the class XP for the parameter "number of clubs" in a cover.

Then, we consider the approximation complexity of Min 2-Club Cover and Min 3-Club Cover. We recall that, given an input graph $G = (V, E)$, Minimum Clique Partition is not approximable within factor $O(|V|^{1-\varepsilon})$, for any $\varepsilon > 0$, unless $P = NP$ [24]. Here we show that Min 2-Club Cover has a slightly different behavior, while Min 3-Club Cover is similar to Clique Partition. Indeed, in Section 4 we prove that Min 2-Club Cover is not approximable within factor $O(|V|^{1/2-\varepsilon})$, for any $\varepsilon > 0$, unless $P = NP$, while Min 3-Club Cover is not approximable within factor $O(|V|^{1-\varepsilon})$, for any $\varepsilon > 0$, unless $P = NP$. In Section 5, we present a greedy approximation algorithm that has factor $2|V|^{1/2} \log^{3/2} |V|$ for Min 2-Club Cover, which almost match the inapproximability result for the problem. We start the paper by giving in Section 2 some definitions and by formally defining the problem we are interested in.

2 Preliminaries

Given a graph $G = (V, E)$ and a subset $V' \subseteq V$, we denote by $G[V']$ the subgraph of G induced by V' . Given two vertices $u, v \in V$, the distance between u and v in G , denoted by $d_G(u, v)$, is the length of a shortest path from u to v . The diameter of a graph $G = (V, E)$ is the maximum distance between two vertices of V . Given a graph $G = (V, E)$ and a vertex $v \in V$, we denote by $N_G(v)$ the set of neighbors of v , that is $N_G(v) = \{u : \{v, u\} \in E\}$. We denote by $N_G[v]$ the close neighborhood of v , that is $N_G[v] = N_G(v) \cup \{v\}$. Define $N_G^l(v) = \{u : u \text{ has distance at most } l \text{ from } v\}$, with $1 \leq l \leq 2$. Given a set of vertices $X \subseteq V$ and l , with $1 \leq l \leq 2$, define $N_G^l(X) = \bigcup_{u \in X} N_G^l(u)$. We may omit the subscript G when it is clear from the context. Now, we give the definition of s -club, which is fundamental for the paper.

Definition 1. *Given a graph $G = (V, E)$, and a subset $V' \subseteq V$, $G[V']$ is an s -club if it has diameter at most s .*

Notice that an s -club must be a connected graph. We present now the formal definition of the Minimum s -Club Cover problem we are interested in.

Minimum s-Club Cover (Min s-Club Cover)

Input: a graph $G = (V, E)$ and an integer $s \geq 2$.

Output: a minimum cardinality collection $\mathcal{S} = \{V_1, \dots, V_h\}$ such that, for each i with $1 \leq i \leq h$, $V_i \subseteq V$, $G[V_i]$ is an s -club, and, for each vertex $v \in V$, there exists a set V_j , with $1 \leq j \leq h$, such that $v \in V_j$.

We denote by **s-Club Cover(h)**, with $1 \leq h \leq |V|$, the decision version of **Min s-Club Cover** that asks whether there exists a cover of G consisting of at most h s -clubs.

Notice that while in **Minimum Clique Partition** we can assume that the cliques that cover a graph $G = (V, E)$ partition V , hence the cliques are vertex disjoint, we cannot make this assumption for **Min s-Club Cover**. Indeed, in a solution of **Min s-Club Cover**, a vertex may be covered by more than one s -club, in order to have a cover consisting of the minimum number of s -clubs. Consider the example of Fig. 1. The two 2-clubs induced by $\{v_1, v_2, v_3, v_4, v_5\}$ and $\{v_1, v_6, v_7, v_8, v_9\}$ cover G , and both these 2-clubs contain vertex v_1 . However, if we ask for a partition of G , we need at least three 2-clubs. This difference between **Minimum Clique Partition** and **Min s-Club Cover** is due to the fact that, while being a clique is a hereditary property, this is not the case for being an s -club. If a graph G is an s -club, then a subgraph of G may not be an s -club (for example a star is a 2-club, but the subgraph obtained by removing its center is not anymore a 2-club).

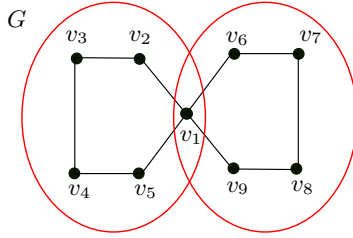


Fig. 1. A graph G and a cover consisting of two 2-clubs (induced by the vertices in the ovals). Notice that the 2-clubs of this cover must both contain vertex v_1 .

3 Computational Complexity

In this section we investigate the computational complexity of **2-Club Cover** and **3-Club Cover** and we show that **2-Club Cover(3)**, that is deciding whether there exists a cover of a graph G with three 2-clubs, and **3-Club Cover(2)**, that is deciding whether there exists a cover of a graph G with two 3-clubs, are NP-complete.

3.1 2-Club Cover(3) is NP-complete

In this section we show that **2-Club Cover(3)** is NP-complete by giving a reduction from the **3-Clique Partition** problem, that is the problem of computing whether there exists a partition of a graph $G^p = (V^p, E^p)$ in three cliques. Consider an instance $G^p = (V^p, E^p)$ of **3-Clique Partition**, we construct an instance $G = (V, E)$ of **2-Club Cover(3)** (see Fig. 2). The vertex set V is defined as follows:

$$V = \{w_i : v_i \in V^p\} \cup \{w_{i,j} : \{v_i, v_j\} \in E^p \wedge i < j\}$$

The set E of edges is defined as follows:

$$E = \{\{w_i, w_{i,j}\}, \{w_i, w_{h,i}\} : v_i \in V^p, w_i, w_{i,j}, w_{h,i} \in V\} \cup \{\{w_{i,j}, w_{i,l}\}, \{w_{i,j}, w_{h,i}\}, \{w_{h,i}, w_{z,i}\} : w_{i,j}, w_{i,l}, w_{h,i}, w_{z,i} \in V\}$$

Before giving the main results of this section, we prove a property of G .

Lemma 2. Let $G^p = (V^p, E^p)$ be an instance of 3-Clique Partition and let $G = (V, E)$ be the corresponding instance of 2-Club Cover(3). Then, given two vertices $v_i, v_j \in V^p$ and the corresponding vertices $w_i, w_j \in V$:

- if $\{v_i, v_j\} \in E^p$, then $d_G(w_i, w_j) = 2$
- if $\{v_i, v_j\} \notin E^p$, then $d_G(w_i, w_j) \geq 3$

Proof. Notice that $N_G(w_i) = \{w_{i,z} : \{v_i, v_z\} \in E^p \wedge i < z\} \cup \{w_{h,i} : \{v_i, v_h\} \in E^p \wedge h < i\}$. It follows that $w_j \in N_G^2(w_i)$ if and only if there exists a vertex $w_{i,j}$ (or $w_{j,i}$), which is adjacent to both w_i and w_j . But then, by construction, $w_j \in N_G^2(w_i)$ if and only if $\{v_i, v_j\} \in E^p$. \square

We are now able to prove the main properties of the reduction.

Lemma 3. Let $G^p = (V^p, E^p)$ be a graph input of 3-Clique Partition and let $G = (V, E)$ be the corresponding instance of 2-Club Cover(3). Then, given a solution of 3-Clique Partition on $G^p = (V^p, E^p)$, we can compute in polynomial time a solution of 2-Club Cover(3) on $G = (V, E)$.

Proof. Consider a solution of 3-Clique Partition on $G^p = (V^p, E^p)$, and let $V_1^p, V_2^p, V_3^p \subseteq V^p$ be the sets of vertices of G^p that partition V^p . We define a solution of 2-Club Cover(3) on $G = (V, E)$ as follows. For each d , with $1 \leq d \leq 3$, define

$$V_d = \{w_j \in V : v_j \in V_d^p\} \cup \{w_{i,j} : v_i \in V_d^p\}$$

We show that each $G[V_d]$, with $1 \leq d \leq 3$, is a 2-club. Consider two vertices $w_i, w_j \in V_d$, with $1 \leq i < j \leq |V|$. Since they correspond to two vertices $v_i, v_j \in V^p$ that belong to a clique of G^p , it follows that $\{v_i, v_j\} \in E^p$ and $w_{i,j} \in V_d$. Thus $d_{G[V_d]}(w_i, w_j) = 2$. Now, consider the vertices $w_i \in V_d$, with $1 \leq i \leq |V|$, and $w_{h,z} \in V_d$, with $1 \leq h < z \leq |V|$. If $i = h$ or $i = z$, assume w.l.o.g. $i = h$, then by construction $d_{G[V_d]}(w_i, w_{i,z}) = 1$. Assume that $i \neq h$ and $i \neq z$ (assume w.l.o.g. that $i < h < z$), since $w_{h,z} \in V_d$, it follows that $w_h \in V_d$. Since $w_i, w_h \in V_d$, it follows that $w_{i,h} \in V_d$. By construction, there exist edges $\{w_{i,h}, w_{h,z}\}, \{w_i, w_{i,h}\}$ in E^p , thus implying that $d_{G[V_d]}(w_i, w_{h,z}) = 2$. Finally, consider two vertices $w_{i,j}, w_{h,z} \in V_d$, with $1 \leq i < j \leq |V|$ and $1 \leq h < z \leq |V|$. Then, by construction, $w_i \in V_d$ and $w_h \in V_d$. But then, $w_{i,h}$ belongs to V_d , and, by construction, $\{w_{i,j}, w_{i,h}\} \in E$ and $\{w_{h,z}, w_{i,h}\} \in E$. It follows that $d_{G[V_d]}(w_{i,j}, w_{h,z}) = 2$.

We conclude the proof observing that, by construction, since V_1^p, V_2^p, V_3^p partition V^p , it holds that $V = V_1 \cup V_2 \cup V_3$, thus $G[V_1], G[V_2], G[V_3]$ covers G . \square

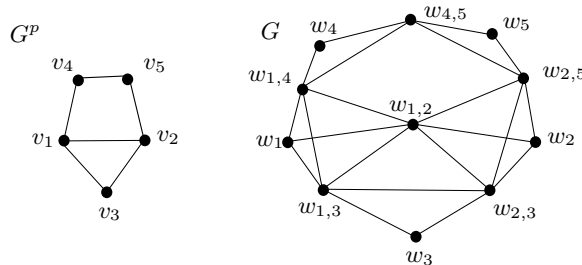


Fig. 2. An example of a graph G^p input of 3-Clique Partition and the corresponding graph G input of 2-Club Cover(3).

Based on Lemma 2, we can prove the following result.

Lemma 4. Let $G^p = (V^p, E^p)$ be a graph input of 3-Clique Partition and let $G = (V, E)$ be the corresponding instance of 2-Club Cover(3). Then, given a solution of 2-Club Cover(3) on $G = (V, E)$, we can compute in polynomial time a solution of 3-Clique Partition on $G^p = (V^p, E^p)$.

Proof. Consider a solution of 2-Club Cover(3) on $G = (V, E)$ consisting of three 2-clubs $G[V_1]$, $G[V_2]$, $G[V_3]$. Consider a 2-club $G[V_d]$, with $1 \leq d \leq 3$. By Lemma 2, it follows that, for each $w_i, w_j \in V_d$, $\{v_i, v_j\} \in E$. As a consequence, we can define three cliques $G^p[V_1^p]$, $G^p[V_2^p]$, $G^p[V_3^p]$ in G^p as follows. For each d , with $1 \leq d \leq 3$, V_d^p is defined as:

$$V_d^p = \{v_i : w_i \in V_d\}$$

Next, we show that $G[V_d^p]$, with $1 \leq d \leq 3$, is indeed a clique. By Lemma 2 if $w_i, w_j \in V_d$ then it holds $\{v_i, v_j\} \in E$, thus by construction $\{v_i, v_j\} \in E^p$ and $G[V_d^p]$ is a clique in G^p . Moreover, since $V_1 \cup V_2 \cup V_3 = V$, then $V_1^p \cup V_2^p \cup V_3^p = V^p$. Notice that V_1^p, V_2^p, V_3^p may not be disjoint, but, starting from V_1^p, V_2^p, V_3^p , it is easy to compute in polynomial time a partition of G^p in three cliques. \square

Now, we can prove the main result of this section.

Theorem 5. 2-Club Cover(3) is NP-complete.

Proof. By Lemma 3 and Lemma 4 and from the NP-hardness of 3-Clique Partition [15], it follows that 2-Club Cover(3) is NP-hard. The membership to NP follows easily from the fact that, given three 2-clubs of G , it can be checked in polynomial time whether they are 2-clubs and cover all vertices of G . \square

3.2 3-Club Cover(2) is NP-complete

In this section we show that 3-Club Cover(2) is NP-complete by giving a reduction from a variant of Sat called 5-Double-Sat. Recall that a literal is positive if it is a non-negated variable, while it is negative if it is a negated variable.

Given a collection of clauses $\mathcal{C} = \{C_1, \dots, C_p\}$ over the set of variables $X = \{x_1, \dots, x_q\}$, where each $C_i \in \mathcal{C}$, with $1 \leq i \leq p$, contains exactly five literals and does not contain both a variable and its negation, 5-Double-Sat asks for a truth assignment to the variables in X such that each clause C_i , with $1 \leq i \leq p$, is *double-satisfied*. A clause C_i is double-satisfied by a truth assignment f to the variables X if there exist a positive literal and a negative literal in C_i that are both satisfied by f . Notice that we assume that there exist at least one positive literal and at least one negative literal in each clause C_i , with $1 \leq i \leq p$, otherwise C_i cannot be double-satisfied. Moreover, we assume that each variable in an instance of 5-Double-Sat appears both as a positive literal and a negative literal in the instance. Notice that if this is not the case, for example a variable appears only as a positive literal, we can assign a true value to the variable, as defining an assignment to false does not contribute to double-satisfy any clause. First, we show that 5-Double-Sat is NP-complete, which may be of independent interest.

Theorem 6. 5-Double-Sat is NP-complete.

Proof. We reduce from 3-Sat, where given a set X_3 of variables and a set \mathcal{C}_3 of clauses, which are a disjunction of 3 literals (a variable or the negation of a variable), we want to find an assignment to the variables such that all clauses are satisfied. Moreover, we assume that each clause in \mathcal{C}_3 does not contain a positive variable x and its negation \bar{x} , since such a clause is obviously satisfied by any assignment. The same property holds also for the instance of 5-Double-Sat we construct.

Consider an instance (X_3, \mathcal{C}_3) of 3-Sat, we construct an instance (X, \mathcal{C}) of 5-Double-Sat as follows. Define $X = X_3 \cup X_N$, where $X_3 \cap X_N = \emptyset$ and X_N is defined as follows:

$$X_N = \{x_{C,i,1}, x_{C,i,2} : C_i \in \mathcal{C}_3\}$$

The set \mathcal{C} of clauses is defined as follows:

$$\mathcal{C} = \{C_{i,1}, C_{i,2} : C_i \in \mathcal{C}_3\}$$

where $C_{i,1}, C_{i,2}$ are defined as follows. Consider $C_i \in \mathcal{C}_3 = (l_{i,1} \vee l_{i,2} \vee l_{i,3})$, where $l_{i,p}$, with $1 \leq p \leq 3$ is a literal, that is a variable (a positive literal) or a negated variable (a negative literal), the two clauses $C_{i,1}$ and $C_{i,2}$ are defined as follows:

- $C_{i,1} = l_{i,1} \vee l_{i,2} \vee l_{i,3} \vee x_{C,i,1} \vee \overline{x_{C,i,2}}$
- $C_{i,2} = l_{i,1} \vee l_{i,2} \vee l_{i,3} \vee \overline{x_{C,i,1}} \vee x_{C,i,2}$

We claim that (X_3, \mathcal{C}_3) is satisfiable if and only if (X, \mathcal{C}) is double-satisfiable.

Assume that (X_3, \mathcal{C}_3) is satisfiable and let f be an assignment to the variables on X that satisfies \mathcal{C}_3 . Consider a clause C_i in \mathcal{C}_3 , with $1 \leq i \leq |\mathcal{C}_3|$. Since it is satisfied by f , it follows that there exists a literal $l_{i,p}$ of C_i , with $1 \leq p \leq 3$, that is satisfied by f . Define an assignment f' on X that is identical to f on X_3 and, if $l_{i,p}$ is positive, then assigns value false to both $x_{C,i,1}$ and $x_{C,i,2}$, if $l_{i,p}$ is negative, then assigns value true to both $x_{C,i,1}$ and $x_{C,i,2}$. It follows that both $C_{i,1}$ and $C_{i,2}$ are double-satisfied by f' .

Assume that (X, \mathcal{C}) is double-satisfied by an assignment f' . Consider two clauses $C_{i,1}$ and $C_{i,2}$, with $1 \leq i \leq |\mathcal{C}|$, that are double-satisfied by f' , we claim that there exists at least one literal of $C_{i,1}$ and $C_{i,2}$ not in X_N which is satisfied. Assume this is not the case, then, if $C_{i,1}$ is double-satisfied, it follows that $x_{C,i,1}$ is true and $x_{C,i,2}$ is false, thus implying that $C_{i,2}$ is not double-satisfied. Then, an assignment f that is identical to f' restricted to X_3 satisfies each clause in \mathcal{C} .

Now, since 3-Sat is NP-complete [15], it follows that 5-Double-Sat is NP-hard. The membership to NP follows from the observation that, given an assignment to the variables on X , we can check in polynomial-time whether each clause in \mathcal{C} is double-satisfied or not. \square

Let us now give the construction of the reduction from 5-Double-Sat to 3-Club Cover(2). Consider an instance of 5-Double-Sat consisting of a set \mathcal{C} of clauses C_1, \dots, C_p over set $X = \{x_1, \dots, x_q\}$ of variables. We assume that it is not possible to double-satisfy all the clauses by setting at most two variables to true or to false (this can be easily checked in polynomial-time).

Before giving the details, we present an overview of the reduction. Given an instance (X, \mathcal{C}) of 5-Double-Sat, for each positive literal x_i , with $1 \leq i \leq q$, we define vertices $x_{i,1}^T, x_{i,2}^T$ and for each negative literal $\overline{x_i}$, with $1 \leq i \leq q$, we define a vertex x_i^F . Moreover, for each clause $C_j \in \mathcal{C}$, with $1 \leq j \leq p$, we define a vertex $v_{C,j}$. We define other vertices to ensure that some vertices have distance not greater than three and to force the membership to one of the two 3-clubs of the solution (see Lemma 7). The construction implies that for each i with $1 \leq i \leq q$, $x_{i,1}^T$ and x_i^F belong to different 3-clubs (see Lemma 8); this corresponds to a truth assignment to the variables in X . Then, we are able to show that each vertex $v_{C,j}$ belongs to the same 3-club of a vertex $x_{i,1}^T$, with $1 \leq i \leq q$, and of a vertex x_h^F , with $1 \leq h \leq q$, adjacent to $v_{C,j}$ (see Lemma 10); these vertices correspond to a positive literal x_i and a negative literal $\overline{x_h}$, respectively, that are satisfied by a truth assignment, hence C_j is double-satisfied.

Now, we give the details of the reduction. Let (X, \mathcal{C}) be an instance of 5-Double-Sat, we construct an instance $G = (V, E)$ of 3-Club Cover(2) as follows (see Fig. 3). The vertex set V is defined as follows:

$$V = \{r, r', r_T, r'_T, r_T^*, r_F, r'_F\} \cup \{x_{i,1}^T, x_{i,2}^T, x_i^F : x_i \in X\} \cup \{v_{C,j} : C_j \in \mathcal{C}\} \cup \{y_1, y_2, y\}$$

The edge set E is defined as follows:

$$\begin{aligned} E = & \{\{r, r'\}, \{\{r', r_T\}, \{r', r_T^*\}\{r', r_F\}\} \cup \{\{r_T, x_{i,1}^T\} : x_i \in X\} \\ & \cup \{\{r_F, x_i^F\} : x_i \in X\} \cup \{\{r'_T, x_{i,1}^T\} : x_i \in X\} \cup \{\{r'_F, x_i^F\} : x_i \in X\} \cup \\ & \{\{x_{i,1}^T, x_{i,2}^T\} : x_i \in X\} \cup \{\{r_T^*, x_{i,2}^T\}, \{y_1, x_{i,2}^T\} : x_i \in X\} \cup \\ & \{\{x_{i,2}^T, x_j^F\} : x_i, x_j \in X, i \neq j\} \cup \{\{x_{i,1}^T, v_{C,j}\} : x_i \in C_j\} \cup \{\{x_i^F, v_{C,j}\} : \overline{x_i} \in C_j\} \cup \\ & \{\{v_{C,j}, y\} : C_j \in \mathcal{C}\} \cup \{\{y, y_2\}, \{y_1, y_2\}, \{y_1, r'_T\}, \{y_1, r'_F\}\} \end{aligned}$$

We start by proving some properties of the graph G .

Lemma 7. *Consider an instance (\mathcal{C}, X) of 5-Double-Sat and let $G = (V, E)$ be the corresponding instance of 3-Club Cover(2). Then, (1) $d_G(r', y) > 3$, (2) $d_G(r, y) > 3$, (3) $d_G(r, v_{C,j}) > 3$, for each j with $1 \leq j \leq p$, and (4) $d_G(r, r'_F) > 3$, $d_G(r, r'_T) > 3$.*

Proof. We start by proving (1). Notice that any path from r' to y must pass through r_T , r_T^* or r_F . Each of r_T , r_T^* or r_F is adjacent to vertices $x_{i,1}^T$, $x_{i,2}^T$ and x_i^F , with $1 \leq i \leq q$ (in addition to r'), and none of these vertices is adjacent to y , thus concluding that $d_G(r', y) > 3$. Moreover, observe that for each vertex $v_{C,j}$, with $1 \leq j \leq p$, there exists a vertex $x_{i,1}^T$, with $1 \leq i \leq q$, or x_h^F , with $1 \leq h \leq q$, that is adjacent to $v_{C,j}$, with $1 \leq j \leq p$, thus $d_G(r', v_{C,j}) = 3$, for each j with $1 \leq j \leq p$. As a consequence of (1), it follows that (2) holds, that is $d_G(r, y) > 3$. Since $d_G(r', v_{C,j}) = 3$, for each j with $1 \leq j \leq p$, it holds (3) $d_G(r, v_{C,j}) > 3$.

Finally, we prove (4). Notice that $N_G^2(r) = \{r', r_T^*, r_T, r_F\}$ and that none of the vertices in $N_G^2(r)$ is adjacent to r'_F and r'_T , thus $d_G(r, r'_F) > 3$. \square

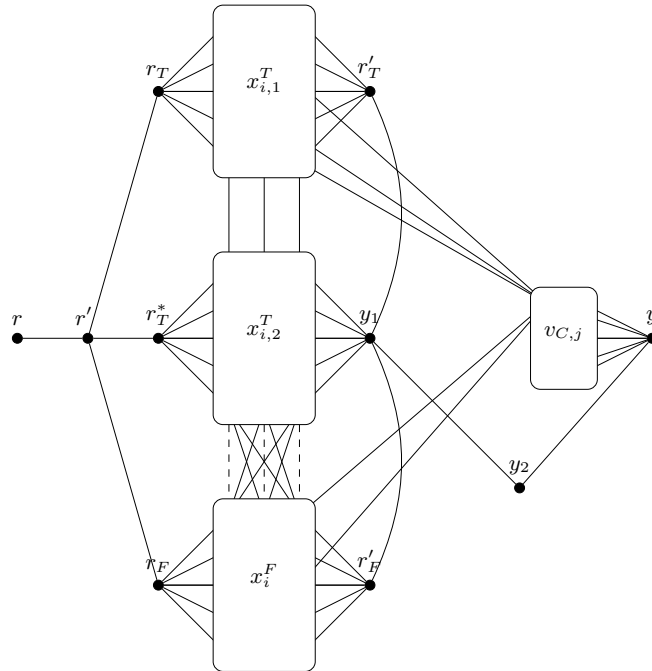


Fig. 3. Schematic construction for the reduction from 5-Double-Sat to 3-Club Cover(2).

Consider two sets $V_1 \subseteq V$ and $V_2 \subseteq V$, such that $G[V_1]$ and $G[V_2]$ are two 3-clubs of G that cover G . As a consequence of Lemma 7, it follows that r and r' are in exactly one of $G[V_1]$, $G[V_2]$, w.l.o.g. $G[V_1]$, while r'_T , r'_F , y and $v_{C,j}$, for each j with $1 \leq j \leq p$, belong to $G[V_2]$ and not to $G[V_1]$.

Next, we show a crucial property of the graph G built by the reduction.

Lemma 8. *Given an instance (\mathcal{C}, X) of 5-Double-Sat, let $G = (V, E)$ be the corresponding instance of 3-Club Cover(2). Then, for each i with $1 \leq i \leq q$, $d_G(x_{i,1}^T, x_i^F) > 3$.*

Proof. Consider a path π of minimum length that connects $x_{i,1}^T$ and x_i^F , with $1 \leq i \leq q$. First, notice that, by construction, the path π after $x_{i,1}^T$ must pass through one of these vertices: r_T , r'_T , $x_{i,2}^T$ or $v_{C,j}$, with $1 \leq j \leq p$.

We consider the first case, that is the path π after $x_{i,1}^T$ passes through r_T . Now, the next vertex in π is either r' or $x_{h,1}^T$, with $1 \leq h \leq q$. Since both r' and $x_{h,1}^T$ are not adjacent to x_i^F , it follows that in this case the path π has length greater than three.

We consider the second case, that is the path π after $x_{i,1}^T$ passes through r'_T . Now, after r'_T , π passes through either y_1 or $x_{h,1}^T$, with $1 \leq h \leq q$. Since both y_1 and $x_{h,1}^T$ are not adjacent to x_i^F , it follows that in this case the path π has length greater than three.

We consider the third case, that is the path after $x_{i,1}^T$ passes through $x_{i,2}^T$. Now, the next vertex of π is either r_T^* or y_1 or x_h^F , with $1 \leq h \leq q$ and $h \neq i$. Since r_T^* , y_1 and x_h^F are not adjacent to x_i^F , it follows that in this case the path π has length greater than three.

We consider the last case, that is the path after $x_{i,1}^T$ passes through $v_{C,j}$, with $1 \leq j \leq p$. We have assumed that x_i and \bar{x}_i do not belong to the same clause, thus by construction x_i^F is not incident in $v_{C,j}$. It follows that after $v_{C,j}$, the path π must pass through either y or $x_{h,1}^T$, with $1 \leq h \leq q$, or x_z^F , $1 \leq z \leq q$ and $z \neq i$. Once again, since y , $x_{h,1}^T$ and x_z^F are not adjacent to x_i^F , it follows that also in this case the path π has length greater than three, thus concluding the proof. \square

Now, we are able to prove the main results of this section.

Lemma 9. *Given an instance (\mathcal{C}, X) of 5-Double-Sat, let $G = (V, E)$ be the corresponding instance of 3-Club Cover(2). Then, given a truth assignment that double-satisfies \mathcal{C} , we can compute in polynomial-time two 3-clubs that cover G .*

Proof. Consider a truth assignment f on the set X of variables that double-satisfies \mathcal{C} . In the following we construct two 3-clubs $G[V_1]$ and $G[V_2]$ that cover G . The two sets V_1, V_2 are defined as follows:

$$V_1 = \{r, r', r_T, r_T^*, r_F\} \cup \{x_{i,1}^T, x_{i,2}^T : f(x_i) = \text{false}\} \cup \{x_i^F, : f(x_i) = \text{true}\}$$

$$V_2 = \{r'_T, r'_F, y, y_1, y_2\} \cup \{x_{i,1}^T, x_{i,2}^T : f(x_i) = \text{true}\} \cup \{x_i^F : f(x_i) = \text{false}\} \cup \{v_{C,j} : 1 \leq j \leq p\}$$

Next, we show that $G[V_1]$ and $G[V_2]$ are indeed two 3-clubs that cover G . First, notice that $V_1 \cup V_2 = V$, hence $G[V_1]$ and $G[V_2]$ cover G . Next, we show that both $G[V_1]$ and $G[V_2]$ are indeed 3-clubs.

Let us first consider $G[V_1]$. By construction, $d_{G[V_1]}(r, x_{i,1}^T) = 3$ and $d_{G[V_1]}(r, x_{i,2}^T) = 3$, for each i with $1 \leq i \leq q$, and $d_{G[V_1]}(r, x_i^F) = 3$, for each i with $1 \leq i \leq q$. Moreover, $d_{G[V_1]}(r', x_{i,1}^T) = 2$ and $d_{G[V_1]}(r', x_{i,2}^T) = 2$, for each i with $1 \leq i \leq q$, and $d_{G[V_1]}(r', x_i^F) = 2$, for each i with $1 \leq i \leq q$. As a consequence, it holds that r_T, r'_T and r_F have distance at most three in $G[V_1]$ from each vertex $x_{i,1}^T$, from each vertex $x_{i,2}^T$, and from each vertex x_i^F . Since r, r_T, r_T^* and r_F are in $N(r')$, it follows that r, r', r_T, r_T^* and r_F are at distance at most 2 in $G[V_1]$. Hence, we focus on vertices $x_{i,1}^T$, with $1 \leq i \leq q$, $x_{h,2}^T$, with $1 \leq h \leq q$ and x_j^F , with $1 \leq j \leq q$. Since there exists a path that passes through $x_{i,1}^T, r_T, x_{h,1}^T$ and $x_{h,2}^T$, vertices $x_{i,1}^T, x_{h,1}^T$ are at distance at most two in $G[V_1]$, while $x_{i,1}^T, x_{h,2}^T$ are at distance at most three in $G[V_1]$ (if $i = h$ they are at distance one). Vertices $x_{h,2}^T$ and x_j^F are at distance one in $G[V_1]$, since $h \neq j$ and $\{x_{h,2}^T, x_j^F\} \in E$ by construction. Finally, $x_{i,1}^T$ and x_j^F are at distance two in $G[V_1]$, since there exists a path that passes through $x_{i,1}^T, x_{i,2}^T$ and x_j^F in $G[V_1]$, as $i \neq j$. It follows that $G[V_1]$ is a 3-club.

We now consider $G[V_2]$. We recall that, for each i with $1 \leq i \leq q$, if $x_{i,1}^T, x_{i,2}^T \in V_2$, then $x_i^F \in V_1$. Furthermore, we recall that we assume that each x_i appears as a positive and a negative literal in the instance of 5-Double-Sat, thus each vertex $x_{i,1}^T$, with $1 \leq i \leq q$, and each vertex x_h^F , with $1 \leq h \leq q$, are connected to some $v_{C,j}$, with $1 \leq j \leq p$.

First, notice that vertex y is at distance at most three in $G[V_2]$ from each vertex of V_2 , since it has distance one in $G[V_2]$ from each vertex $v_{C,j}$, with $1 \leq j \leq p$, thus distance two from $x_{i,1}^T$, with $1 \leq i \leq q$, and x_h^F , with $1 \leq h \leq q$, and three from $x_{i,2}^T$, with $1 \leq i \leq q$, r'_T and r'_F . Since y is adjacent to y_2 , it has distance one from y_2 and two from y_1 .

Now, consider a vertex $v_{C,j}$, with $1 \leq j \leq p$. Since f double-satisfies \mathcal{C} , it follows that there exist two vertices in V_2 , $x_{i,1}^T$, with $1 \leq i \leq q$, and x_z^F , with $1 \leq z \leq q$, which are connected to $v_{C,j}$. It follows that $v_{C,j}$ has distance 2 in $G[V_2]$ from r'_T and from r'_F , and at most 3 from each $x_{h,1}^T \in V_2$, with $1 \leq h \leq q$, and from each $x_z^F \in V_2$, with $1 \leq z \leq q$. Furthermore, notice that, since $v_{C,j}$ is adjacent to x_z^F and x_z^F is adjacent to each $x_{h,2}^T \in V_2$, with $1 \leq h \leq q$ and $h \neq z$, then

$v_{C,j}$ has distance at most two in $G[V_2]$ from each $x_{h,2}^T \in V_2$. Finally, since $v_{C,j}$ is adjacent to y , it has distance two and three respectively, from y_2 and y_1 , in $G[V_2]$.

Consider a vertex $x_{i,1}^T \in V_2$, with $1 \leq i \leq q$. We have already shown that it has distance at most three in $G[V_2]$ from any $v_{C,j}$, with $1 \leq j \leq p$, and two from y . Since $x_{i,1}^T$ is adjacent to r'_T , it has distance at most two from each other vertex $x_{h,1}^T$, with $1 \leq h \leq q$, and three from each other vertex $x_{h,2}^T$ of $G[V_2]$. Moreover, it has distance two from y_1 and three from y_2 and r'_F . Since $x_{i,2}^T$ is adjacent to every vertex $x_z^F \in V_2$, with $1 \leq z \leq q$, as $z \neq i$, it follows that $x_{i,1}^T$ has distance at most two from every vertex $x_z^F \in V_2$.

Consider a vertex $x_{i,2}^T \in V_2$, with $1 \leq i \leq q$. We have already shown that it has distance at most two from each $v_{C,j}$ in $G[V_2]$. Since it is connected to $x_{i,1}^T$, it has distance three from y and two from r'_T in $G[V_2]$. By construction $x_{i,2}^T$ is adjacent to every vertex $x_z^F \in V_2$, with $1 \leq z \leq q$, $x_{i,2}^T$ has distance at most two from r'_F in $G[V_2]$. Moreover, $x_{i,2}^T$ has distance two from each vertex $x_{h,2}^T$ in $G[V_2]$, with $1 \leq i \leq q$, since by construction they are both adjacent to y_1 . Since $x_{i,2}^T$ is adjacent to y_1 , thus it has distance at most two from y_2 in $G[V_2]$.

Consider a vertex x_h^F , with $1 \leq h \leq q$. It has distance one from r'_F in $G[V_2]$, and thus distance two from y_1 and three from y_2 in $G[V_2]$. Moreover, x_h^F is adjacent to each $x_{i,2}^T \in V_2$, with $1 \leq i \leq q$, thus it has distance two from each $x_{i,1}^T$ and distance three from r'_T in $G[V_2]$. Since by construction there exists at least one $v_{C,j}$, with $1 \leq j \leq p$, adjacent to x_h^F , thus x_h^F has distance two from y and three from each $v_{C,z}$ in $G[V_2]$.

Finally, we consider vertices r'_T, r'_F, y_1 and y_2 . Notice that it suffices to show that these vertices have pairwise distance at most three in $G[V_2]$, since we have previously shown that any other vertex of V_2 has distance at most three from these vertices in $G[V_2]$. Since $r'_T, r'_F, y_2 \in N(y_1)$, they are all at distance at most two. It follows that $G[V_2]$ is a 3-club, thus concluding the proof. \square

Lemma 10. *Given an instance (\mathcal{C}, X) of 5-Double-Sat, let $G = (V, E)$ be the corresponding instance of 3-Club Cover(2). Then, given two 3-clubs that cover G , we can compute in polynomial time a truth assignment that double-satisfies \mathcal{C} .*

Proof. Consider two 3-clubs $G[V_1], G[V_2]$, with $V_1, V_2 \subseteq V$, that cover G . First, notice that by Lemma 7 we assume that $r, r' \in V_1 \setminus V_2$, while $y, r'_T, r'_F \in V_2 \setminus V_1$ and $v_{C,j} \in V_2 \setminus V_1$, for each j with $1 \leq j \leq p$. Moreover, by Lemma 8 it follows that for each i with $1 \leq i \leq q$, $x_{i,1}^T$ and x_i^F do not belong to the same 3-club, that is exactly one belongs to V_1 and exactly one belongs to V_2 .

By construction, each path of length at most three from a vertex $v_{C,j}$, with $1 \leq j \leq p$, to r'_F must pass through some x_h^F , with $1 \leq h \leq q$. Similarly, each path of length at most three from a vertex $v_{C,j}$, with $1 \leq j \leq p$, to r'_T must pass through some $x_{i,1}^T$. Assume that $v_{C,j}$, with $1 \leq j \leq p$, is not adjacent to a vertex $x_{i,1}^T \in V_2$, with $1 \leq i \leq q$ ($x_h^F \in V_2$, with $1 \leq h \leq p$ respectively). It follows that $v_{C,j}$ is only adjacent to y and to vertices x_w^F , with $1 \leq w \leq q$ ($x_{u,1}^T$, with $1 \leq u \leq q$, respectively) in $G[V_2]$. In the first case, notice that y is adjacent only to $v_{C,z}$, with $1 \leq z \leq p$, and y_2 , none of which is adjacent to r'_T (r'_F , respectively), thus implying that this path from $v_{C,j}$ to r'_T (to r'_F , respectively) has length at least 4. In the second case, x_w^F ($x_{u,1}^T$, respectively) is adjacent to r'_F , r_F , $v_{C,j}$ and $x_{i,2}^T$ (r'_T , r_T , $v_{C,j}$, $x_{u,2}^T$, respectively), none of which is adjacent to r'_T (r'_F , respectively), implying that also in this case the path from $v_{C,j}$ to r'_T (to r'_F , respectively) has length at least 4. Since $r'_T, r'_F, v_{C,j} \in V_2$, it follows that, for each $v_{C,j}$, the set V_2 contains a vertex $x_{i,1}^T$, with $1 \leq i \leq q$, and a vertex x_h^F , with $1 \leq h \leq q$, connected to $v_{C,j}$.

By Lemma 8 exactly one of $x_{i,1}^T, x_i^F$ belongs to V_2 , thus we can construct a truth assignment f as follows: $f(x_i) := \text{true}$, if $x_{i,1}^T \in V_2$, $f(x_i) := \text{false}$, if $x_i^F \in V_2$. The assignment f double-satisfies each clause of \mathcal{C} , since each $v_{C,j}$ is connected to a vertex $x_{i,1}^T$, for some i with $1 \leq i \leq q$, and a vertex x_h^F , for some h with $1 \leq h \leq q$. \square

Based on Lemma 9 and Lemma 10, and on the NP-completeness of 5-Double-Sat (see Theorem 6), we can conclude that 3-Club Cover(2) is NP-complete.

Theorem 11. *3-Club Cover(2) is NP-complete.*

Proof. By Lemma 9 and Lemma 10, and from the NP-hardness of 5-Double-Sat (see Theorem 6), it follows that 3-Club Cover(2) is NP-hard. The membership in NP follows easily from the fact that, given two 3-clubs, it can be checked in polynomial time whether are 3-clubs and cover all vertices of G . \square

4 Hardness of Approximation

In this section we consider the approximation complexity of Min 2-Club Cover and Min 3-Club Cover and we prove that Min 2-Club Cover is not approximable within factor $O(|V|^{1/2-\varepsilon})$, for each $\varepsilon > 0$, and that Min 3-Club Cover is not approximable within factor $O(|V|^{1-\varepsilon})$, for each $\varepsilon > 0$. The proof for Min 2-Club Cover is obtained with a reduction very similar to that of Section 3.1, except from the fact that we reduce Minimum Clique Partition to Min 2-Club Cover.

Corollary 12. *Unless $P = NP$, Min 2-Club Cover is not approximable within factor $O(|V|^{1/2-\varepsilon})$, for each $\varepsilon > 0$.*

Proof. We present a preserving-factor reduction from Minimum Clique Partition to Min 2-Club Cover. Let $G^p = (V^p, E^p)$ be a graph input of Minimum Clique Partition, we compute in polynomial time a corresponding instance $G = (V, E)$ of Min 2-Club Cover as in Section 3.1. In what follows we prove the following results that are useful for the reduction.

Lemma 13. *Let $G^p = (V^p, E^p)$ be a graph input of Minimum Clique Partition and let $G = (V, E)$ be the corresponding instance of Min 2-Club Cover. Then, given a solution of Minimum Clique Partition on $G^p = (V^p, E^p)$ consisting of k cliques, we can compute in polynomial time a solution of Min 2-Club Cover on $G = (V, E)$ consisting of k 2-clubs.*

Proof. Consider a solution of Minimum Clique Partition on $G^p = (V^p, E^p)$ where $\{V_1^p, V_2^p, \dots, V_k^p\}$ is the set of k cliques that partition V^p . We define a solution of Min 2-Club Cover on $G = (V, E)$ consisting of k 2-clubs as follows. For each $d, 1 \leq d \leq k$, let

$$V_d = \{w_j \in V : v_j \in V_d^p\} \cup \{w_{i,j} : v_i \in V_d^p \wedge i < j\}$$

As for the proof of Lemma 9, it follows that for each d , $G[V_d]$ is a 2-club. Furthermore, $G[V_1], \dots, G[V_k]$ cover each vertex of V , as each $v_i \in V^p$ is covered by one of the cliques $V_1^p, V_2^p, \dots, V_k^p$. \square

Lemma 14. *Let $G^p = (V^p, E^p)$ be a graph input of Minimum Clique Partition and let $G = (V, E)$ be the corresponding instance of Min 2-Club Cover. Then, given a solution of Min 2-Club Cover on $G = (V, E)$ consisting of k 2-clubs, we can compute in polynomial time a solution of Minimum Clique Partition on $G^p = (V^p, E^p)$ with k cliques.*

Proof. Consider the 2-clubs $G[V_1], \dots, G[V_k]$ that cover G . As for the proof of Lemma 10, the result follows from the fact that by Lemma 2, given $w_i, w_j \in V_d$, for each d with $1 \leq d \leq k$, it holds that $\{v_i, v_j\} \in E$. As a consequence, we can define a solution of Minimum Clique Partition on $G^p = (V^p, E^p)$ consisting of k cliques as follows, for each $d, 1 \leq d \leq k$:

$$V_d^p = \{v_i : w_i \in V_d\}$$

\square

The inapproximability of Min 2-Club Cover follows from Lemma 13 and Lemma 14, and from the inapproximability of Minimum Clique Partition, which is known to be inapproximable within factor $O(|V^p|^{1-\varepsilon'})$ [24] (where $G^p = (V^p, E^p)$ is an instance of Hence Min 2-Club Cover is not approximable within factor $O(|V^p|^{1-\varepsilon'})$, for each $\varepsilon' > 0$, unless $P = NP$, hence Min 2-Club Cover is not approximable within factor $O(|V^p|^{(1-\varepsilon')})$. By the definition of $G = (V, E)$, it holds $|V| = |V^p| + |E^p| \leq |V^p|^2$ hence, for each $\varepsilon > 0$, Min 2-Club Cover is not approximable within factor $O(|V|^{1/2-\varepsilon})$, unless $P = NP$. \square

Next, we show that Min 3-Club Cover is not approximable within factor $O(|V|^{1-\varepsilon})$, for each $\varepsilon > 0$, unless $P = NP$, by giving a preserving-factor reduction from Minimum Clique Partition.

Consider an instance $G^p = (V^p, E^p)$ of Minimum Clique Partition, we construct an instance $G = (V, E)$ of Min 3-Club Cover by adding a pendant vertex connected to each vertex of V^p . Formally, $V = \{u_i, w_i : v_i \in V^p\}$, $E = \{\{u_i, w_i\} : 1 \leq i \leq |V^p|\} \cup \{\{u_i, u_j\} : \{v_i, v_j\} \in E^p\}$.

We prove now the main properties of the reduction.

Lemma 15. *Let $G^p = (V^p, E^p)$ be an instance of Minimum Clique Partition and let $G = (V, E)$ be the corresponding instance of Min 3-Club Cover. Then, given a solution of Minimum Clique Partition on $G^p = (V^p, E^p)$ consisting of k cliques, we can compute in polynomial time a solution of Min 3-Club Cover on $G = (V, E)$ consisting of k 3-clubs.*

Proof. Consider a solution of Minimum Clique Partition on $G^p = (V^p, E^p)$, consisting of the cliques $\{G^p[V_{c,1}], G^p[V_{c,2}], \dots, G^p[V_{c,k}]\}$. Then, for each i , with $1 \leq h \leq k$, define the following subset $V_h \subseteq V$:

$$V_h = \{u_j, w_j \in V : v_j \in V_h^p\}$$

Since $V_1^p, V_2^p \dots V_k^p$ partition V^p , it follows that $V_1, V_2 \dots V_k$ partition (hence cover) G . Now, we show that each $G[V_h]$, with $1 \leq h \leq k$, is a 3-club. First, notice that since $G[V_h^p]$, is a clique, then the set $\{u_j : u_j \in V_h\}$ induces a clique in G . Then, it follows that, for each $u_i, w_j, w_z \in V_h$, $d_{G[V_h]}(u_i, w_j) \leq 2$ and $d_{G[V_h]}(w_j, w_z) \leq 3$, thus concluding the proof. \square

Lemma 16. *Let $G^p = (V^p, E^p)$ be a graph input of Minimum Clique Partition and let $G = (V, E)$ be the corresponding instance of Min 3-Club Cover. Then, given a solution of Min 3-Club Cover on $G = (V, E)$ consisting of k 3-clubs, we can compute in polynomial time a solution of Minimum Clique Partition on $G^p = (V^p, E^p)$ consisting of k cliques.*

Proof. Consider the k 3-clubs $G[V_1], \dots, G[V_k]$ that cover G . First, we show that for each $V_h, 1 \leq h, \leq k, \forall w_i, w_j \in V_h$, with $1 \leq i, j \leq |V^p|$, it holds that $u_i, u_j \in V_h$. Indeed, notice that $N(w_i) = \{u_i\}$ and $N(w_j) = \{u_j\}$, and by the definition of a 3-club we must have $d_{G[V_h]}(w_i, w_j) \leq 3$, it follows that $u_i, u_j \in V_h$. Hence, we can define a set of cliques of G^p . For each V_h , with $1 \leq h \leq k$, define a set V_h^p :

$$V_h^p = \{v_i : w_i \in V_h\}$$

Notice that each $V_h^p, 1 \leq h \leq k$, induces a clique in G^p , as by construction if $v_i, v_j \in V_h^p$, then $w_i, w_j \in V_h$, and this implies $\{v_i, v_j\} \in E^p$. Notice that the cliques V_1^p, \dots, V_k^p may overlap, but starting from V_1^p, \dots, V_k^p , we can easily compute in polynomial time a clique partition of G^p consisting of at most k cliques. \square

Lemma 15 and Lemma 16 imply the following result.

Theorem 17. *Min 3-Club Cover is not approximable within factor $O(|V|^{1-\varepsilon})$, for each $\varepsilon > 0$, unless $P = NP$.*

Proof. The result follows from Lemma 15 and Lemma 16, as these results imply that we have defined a factor-preserving reduction, and from the inapproximability of Minimum Clique Partition, which is known to be inapproximable within factor $O(|V^p|^{1-\varepsilon})$, for each $\varepsilon > 0$, unless $P = NP$ [24] (where $G^p = (V^p, E^p)$ is an instance of Minimum Clique Partition). Thus, Min 3-Club Cover is not approximable within factor $O(|V^p|^{1-\varepsilon})$, for each $\varepsilon > 0$, unless $P = NP$, and since it holds $|V| = 2|V^p|$, Min 3-Club Cover is not approximable within factor $O(|V|^{1-\varepsilon})$, unless $P = NP$. \square

5 An Approximation Algorithm for Min 2-Club Cover

In this section, we present an approximation algorithm for Min 2-Club Cover that achieves an approximation factor of $2|V|^{1/2} \log^{3/2} |V|$. Notice that, due to the result in Section 4, the approxi-

mation factor is almost tight. We start by describing the approximation algorithm, then we present the analysis of the approximation factor.

Algorithm 1: Club-Cover-Approx

Data: a graph G
Result: a cover \mathcal{S} of G

- 1 $V' := V$; /* V' is the set of uncovered vertices of G , initialized to V */
- 2 $\mathcal{S} := \emptyset$;
- 3 **while** $V' \neq \emptyset$ **do**
- 4 Let v be a vertex of V such that $|N[v] \cap V'|$ is maximum;
- 5 Add $N[v]$ to \mathcal{S} ;
- 6 $V' := V' \setminus N[v]$;

Club-Cover-Approx is similar to the greedy approximation algorithm for Minimum Dominating Set and Minimum Set Cover. While there exists an uncovered vertex of G , the Club-Cover-Approx algorithm greedily defines a 2-club induced by the set $N[v]$ of vertices, with $v \in V$, such that $N[v]$ covers the maximum number of uncovered vertices (notice that some of the vertices of $N[v]$ may already be covered). While for Minimum Dominating Set the choice of each iteration is optimal, here the choice is suboptimal. Notice that indeed computing a maximum 2-club is NP-hard.

Clearly the algorithm returns a feasible solution for Min 2-Club Cover, as each set $N[v]$ picked by the algorithm is a 2-club and, by construction, each vertex of V is covered. Next, we show the approximation factor yielded by the Club-Cover-Approx algorithm for Min 2-Club Cover.

First, consider the set V_D of vertices $v \in V$ picked by the Club-Cover-Approx algorithm, so that $N[v]$ is added to \mathcal{S} . Notice that $|V_D| = |\mathcal{S}|$ and that V_D is a dominating set of G , since, at each step, the vertex v picked by the algorithm dominates each vertex in $N[v]$, and each vertex in V is covered by the algorithm, so it belongs to some $N[v]$, with $v \in V_D$.

Let D be a minimum dominating set of the input graph G . By the property of the greedy approximation algorithm for Minimum Dominating Set, the set V_D has the following property [14]:

$$|V_D| \leq |D| \log |V| \tag{1}$$

The size of a minimum dominating set in graphs of diameter bounded by 2 (hence 2-clubs) has been considered in [8], where the following result is proven.

Lemma 18 ([8]). *Let $H = (V_H, E_H)$ be a 2-club, then H has a dominating set of size at most $1 + \sqrt{|V_H| + \ln(|V_H|)}$.*

The approximation factor $2|V|^{1/2} \log^{3/2} |V|$ for Club-Cover-Approx is obtained by combining Lemma 18 and Equation 1.

Theorem 19. *Let OPT be an optimal solution of Min 2-Club Cover, then Club-Cover-Approx returns a solution having at most $2|V|^{1/2} \log^{3/2} |V| |OPT|$ 2-clubs.*

Proof. Let D be a minimum dominating set of G and let OPT be an optimal solution of Min 2-Club Cover. We start by proving that $|D| \leq 2|OPT| |V|^{1/2} \log^{1/2} |V|$. For each 2-club $G[C]$, with $C \subseteq V$, that belongs to OPT , by Lemma 18 there exists a dominating set D_C of size at most $1 + \sqrt{|C| + \ln(|C|)} \leq 2\sqrt{|C| + \ln(|C|)}$. Since $|C| \leq |V|$, it follows that each 2-club $G[C]$ that belongs to OPT has a dominating set of size at most $2\sqrt{|V| + \ln(|V|)}$. Consider $D' = \bigcup_{C \in OPT} D_C$. It follows that D' is a dominating set of G , since the 2-clubs in OPT covers G . Since D' contains $|OPT|$ sets D_C and $|D_C| \leq 2\sqrt{|V| + \ln(|V|)}$, for each $G[C] \in OPT$, it follows that $|D'| \leq 2|OPT| \sqrt{|V| + \ln(|V|)}$. Since D is a minimum dominating set, it follows that $|D| \leq |D'| \leq 2|OPT| (\sqrt{|V| + \ln(|V|)})$. By Equation 1, it holds $|V_D| \leq 2|D| \log |V|$ thus $|V_D| \leq 2|V|^{1/2} \ln^{1/2} |V| \log |V| |OPT| \leq 2|V|^{1/2} \log^{3/2} |V| |OPT|$. \square

6 Conclusion

There are some interesting direction for the problem of covering a graph with s -clubs. From the computational complexity point of view, the main open problem is whether **2-Club Cover(2)** is NP-complete or is in P. Moreover, it would be interesting to study the computational/parameterized complexity of the problem in specific graph classes, as done for **Minimum Clique Partition** [5,6,21,9].

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