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# ON THE CONVERGENCE OF MIRROR DESCENT BEYOND STOCHASTIC CONVEX PROGRAMMING

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ABSTRACT. In this paper, we examine the convergence of mirror descent in a class of stochastic optimization problems that are not necessarily convex (or even quasi-convex), and which we call *variationally coherent*. Since the standard technique of “ergodic averaging” offers no tangible benefits beyond convex programming, we focus directly on the algorithm’s last generated sample (its “last iterate”), and we show that it converges with probability 1 if the underlying problem is coherent. We further consider a localized version of variational coherence which ensures local convergence of SMD with high probability. These results contribute to the landscape of nonconvex stochastic optimization by showing that (quasi-)convexity is not essential for convergence to a global minimum: rather, variational coherence, a much weaker requirement, suffices. Finally, building on the above, we reveal an interesting insight regarding the convergence speed of SMD: in problems with sharp minima (such as generic linear programs or concave minimization problems), SMD reaches a minimum point *in a finite number of steps* (a.s.), even in the presence of persistent gradient noise. This result is to be contrasted with existing black-box convergence rate estimates that are only asymptotic.

## 1. INTRODUCTION

Stochastic mirror descent (SMD) and its variants arguably comprise one of the most widely used families of first-order methods in stochastic optimization – convex and non-convex alike [3, 9, 10, 12, 24, 26–32, 34, 41, 46]. Heuristically, in the “dual averaging” (or “lazy”) incarnation of the method [34, 41, 46], SMD proceeds by aggregating a sequence of i.i.d. gradient samples and then mapping the result back to the problem’s feasible region via a specially constructed “mirror map” (the namesake of the method). In so doing, SMD generalizes and extends the classical stochastic gradient descent (SGD) algorithm (with Euclidean projections playing the role of the mirror map) [33, 35, 37], the exponentiated gradient method of [22], the matrix regularization schemes of [21, 26, 43], and many others.

Starting with the seminal work of Nemirovski and Yudin [32], the convergence of mirror descent has been studied extensively in the context of convex programming

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(including distributed and stochastic optimization problems) [3, 31, 34, 46], non-cooperative games/saddle-point problems [28, 31, 34], and monotone variational inequalities (VIs) [20, 30, 34]. In this monotone setting, it is customary to consider the so-called “ergodic average”  $\bar{X}_n = \sum_{k=1}^n \gamma_k X_k / \sum_{k=1}^n \gamma_k$  of the algorithm’s generated sample points  $X_n$ , with  $\gamma_n$  denoting the method’s step-size. The reason for this is that, by Jensen’s inequality, convexity guarantees that a regret-based analysis can lead to explicit convergence rates for  $\bar{X}_n$  [31, 34, 41, 46]. However, this type of averaging provides no tangible benefits in non-convex programs, so it seems more natural to focus directly on the algorithm’s last generated sample – its “last iterate”.

The long-term behavior of the last iterate of SMD was recently studied by Shamir and Zhang [42] and Nedic and Lee [29] in the context of strongly convex problems. In this case, the algorithm’s last iterate achieves the same value convergence rate as its ergodic average, so averaging is not more advantageous. Jiang and Xu [19] also examined the convergence of the last iterate of SGD in a class of (not necessarily monotone) variational inequalities that admit a unique solution, and they showed that it converges to said solution with probability 1. In a very recent paper [11], it was shown that in phase retrieval problems (a special class of non-convex problems that involve systems of quadratic equations), SGD with random initialization converges to global optimal solutions with probability 1. Finally, in general non-convex problems, Ghadimi and Lan [15, 16] showed that running SGD with a randomized stopping time guarantees convergence to a critical point in the mean, and they estimated the speed of this convergence. However, beyond these (mostly recent) results, not much is known about the convergence of the individual iterates of mirror descent in non-convex programs.

**Our contributions.** In this paper, we examine the asymptotic behavior of mirror descent in a class of stochastic optimization problems that are not necessarily convex (or even quasi-convex). This class of problems, which we call *variationally coherent* (VC), are related to a class of variational inequalities studied by Jiang and Xu [19] and, earlier, by Wang et al. [45] – though, importantly, we do not assume here the existence of a unique solution. Focusing for concreteness on the *dual averaging* (DA) variant of SMD (also known as “lazy” mirror descent) [34, 41, 46], we show that the algorithm’s last iterate converges to a global minimum with probability 1 under mild assumptions for the algorithm’s gradient oracle (unbiased i.i.d. gradient samples that are bounded in  $L^2$ ). This result can be seen as the “mirror image” of the analysis of [19] and reaffirms that (quasi-)convexity/monotonicity is not essential for convergence to a global optimum point: the weaker requirement of variational coherence suffices.

To extend the range of our analysis, we also consider a localized version of variational coherence which includes multi-modal functions that are not even *locally* (quasi-)convex near their minimum points (so, in particular, an eigenvalue-based analysis cannot be readily applied to such problems). Here, in contrast to the globally coherent case, a single, “unlucky” gradient sample could drive the algorithm away from the “basin of attraction” of a local minimum (even a locally coherent one), possibly never to return. Nevertheless, we show that, with overwhelming probability, the last iterate of SMD converges locally to minimum points that are locally coherent (for a precise statement, see [Section 5](#)).

Going beyond this “black-box” analysis, we also consider a class of optimization problems that admit *sharp* minima, a fundamental notion due to Polyak [35]. In stark contrast to existing ergodic convergence rates (which are asymptotic in nature), we show that the last iterate of SMD converges to sharp minima of variationally coherent problems in an almost surely *finite* number of iterations, provided that the method’s mirror map is surjective. As an important corollary, it follows that the last iterate of (lazy) stochastic gradient descent attains a solution of a stochastic linear program in a finite number of steps (a.s.). For completeness, we also derive a localized version of this result for problems with sharp local minima that are not globally coherent: in this case, convergence in a finite number of steps is retained but, instead of “almost surely”, convergence now occurs with overwhelming probability.

Important classes of problems that admit sharp minima are generic linear programs (for the global case) and concave minimization problems (for the local case). In both instances, the (fairly surprising) fact that SMD attains a minimizer in a finite number of iterations should be contrasted to existing work on stochastic linear programming which exhibits asymptotic convergence rates [2, 44]. We find this result particularly appealing as it highlights an important benefit of working with “lazy” descent schemes: “greedy” methods (such as vanilla gradient descent) always take a gradient step from the last generated sample, so convergence in a finite number of iterations is a priori impossible in the presence of persistent noise. By contrast, the aggregation of gradient steps in “lazy” schemes means that even a “bad” gradient sample might not change the algorithm’s sampling point (if the mirror map is surjective), so finite-time convergence *is* possible in this case.

Our analysis hinges on the construction of a primal-dual analogue of the Bregman divergence which we call the *Fenchel coupling*, and which tracks the evolution of the algorithm’s (dual) gradient aggregation variable relative to a target point in the problem’s (primal) feasible region. This energy function allows us to perform a quasi-Fejérian analysis of stochastic mirror descent and, combined with a series of (sub)martingale convergence arguments, ultimately yields the convergence of the algorithm’s last iterate – first as a subsequence, then with probability 1.

## 2. PROBLEM SETUP AND BASIC DEFINITIONS

**2.1. The main problem.** Let  $\mathcal{X}$  be a convex compact subset of a  $d$ -dimensional vector space  $\mathcal{V}$  with norm  $\|\cdot\|$ . Throughout this paper, we will focus on stochastic optimization problems of the general form

$$\begin{aligned} & \text{minimize} && f(x), \\ & \text{subject to} && x \in \mathcal{X}, \end{aligned} \tag{Opt}$$

where

$$f(x) = \mathbb{E}[F(x; \omega)] \tag{2.1}$$

for some stochastic objective function  $F: \mathcal{X} \times \Omega \rightarrow \mathbb{R}$  defined on an underlying (complete) probability space  $(\Omega, \mathcal{F}, \mathbb{P})$ . In terms of regularity, our blanket assumptions for (Opt) will be as follows:

**Assumption 1.**  $F(x, \omega)$  is continuously differentiable in  $x$  for almost all  $\omega \in \Omega$ .

**Assumption 2.** The gradient of  $F$  is uniformly bounded in  $L^2$ , i.e.,  $\mathbb{E}[\|\nabla F(x; \omega)\|_*^2] \leq V^2$  for some finite  $V \geq 0$  and all  $x \in \mathcal{X}$ .

*Remark.* In the above, gradients are treated as elements of the dual space  $\mathcal{Y} \equiv \mathcal{V}^*$  of  $\mathcal{V}$ , and  $\|y\|_* = \sup\{\langle y, x \rangle : \|x\| \leq 1\}$  denotes the dual norm of  $y \in \mathcal{Y}$ . We also note that  $\nabla F(x; \omega)$  refers to the gradient of  $F(x; \omega)$  with respect to  $x$ ; since  $\Omega$  is not assumed to carry a differential structure, there is no danger of confusion.

**Assumption 1** is a token regularity assumption which can be relaxed to account for nonsmooth objectives by using subgradient devices (as opposed to gradients). However, this would make the presentation significantly more cumbersome, so we stick with smooth objectives throughout. **Assumption 2** is also standard in the stochastic optimization literature: it holds trivially if  $F$  is uniformly Lipschitz (another commonly used condition) and, by the dominated convergence theorem, it further implies that  $f$  is smooth and  $\nabla f(x) = \nabla \mathbb{E}[F(x; \omega)] = \mathbb{E}[\nabla F(x; \omega)]$  is bounded. As a result, the solution set

$$\mathcal{X}^* = \arg \min f \tag{2.2}$$

of **(Opt)** is closed and nonempty (by the compactness of  $\mathcal{X}$  and the continuity of  $f$ ).

We briefly discuss below two important examples of **(Opt)**:

*Example 2.1* (Distributed optimization). An important special case of **(Opt)** with high relevance to statistical inference, signal processing and machine learning is when  $f$  is of the special form

$$f(x) = \frac{1}{N} \sum_{i=1}^N f_i(x), \tag{2.3}$$

for some family of functions (or training samples)  $f_i: \mathcal{X} \rightarrow \mathbb{R}$ ,  $i = 1, \dots, N$ . As an example, this setup corresponds to empirical risk minimization with uniform weights, the sample index  $i$  being drawn with uniform probability from  $\{1, \dots, N\}$ .

*Example 2.2* (Noisy gradient measurements). Another widely studied instance of **(Opt)** is when

$$F(x; U) = f(x) + \langle U, x \rangle \tag{2.4}$$

for some random vector  $U$  such that  $\mathbb{E}[U] = 0$  and  $\mathbb{E}[\|U\|_*^2] < \infty$ . This gives  $\nabla F(x; U) = \nabla f(x) + U$ , so **(Opt)** can be seen here as a model for deterministic optimization problems with noisy gradient measurements.

**2.2. Variational coherence.** We are now in a position to define the class of variationally coherent problems:

**Definition 2.1.** We say that **(Opt)** is *variationally coherent* if

$$\langle \nabla f(x), x - x^* \rangle \geq 0 \quad \text{for all } x \in \mathcal{X}, x^* \in \mathcal{X}^*, \tag{VC}$$

and there exists some  $x^* \in \mathcal{X}^*$  such that equality holds in **(VC)** only if  $x \in \mathcal{X}^*$ .

In words, **(VC)** states that solutions of **(Opt)** can be harvested by solving a (Minty) variational inequality – hence the term “variational coherence”. To the best of our knowledge, the closest analogue to this condition first appeared in the classical paper of Bottou [5] on online learning and stochastic approximation algorithms, but with the added assumptions that *a*) the problem **(Opt)** admits a unique solution  $x^*$ ; and *b*) an extra positivity requirement for  $\langle \nabla f(x), x - x^* \rangle$  in punctured neighborhoods of  $x^*$ . In the context of variational inequalities, a closely

related variant of (VC) has been used to establish the convergence of extra-gradient methods [14, 45] and stochastic gradient descent [19] in (Stampacchia) variational inequalities with a unique solution. By contrast, there is no uniqueness requirement in (VC), an aspect of the definition which we examine in more detail below.

We should also note that, as stated, (VC) is a non-random requirement for  $f$  so it applies equally well to *deterministic* optimization problems. Alternatively, by the dominated convergence theorem, (VC) can be written equivalently as

$$\mathbb{E}[\langle \nabla F(x; \omega), x - x^* \rangle] \geq 0, \quad (2.5)$$

so it can be interpreted as saying that  $F$  is variationally coherent “on average”, without any individual realization thereof satisfying (VC). Both interpretations will come in handy later on.

All in all, the notion of variational coherence will play a central role in our paper so a few examples are in order:

*Example 2.3* (Convex programming). If  $f$  is convex,  $\nabla f$  is *monotone* [39] in the sense that

$$\langle \nabla f(x) - \nabla f(x'), x - x' \rangle \geq 0 \quad \text{for all } x, x' \in \mathcal{X}. \quad (2.6)$$

By the first-order optimality conditions for  $f$ , it follows that  $\langle \nabla f(x^*), x - x^* \rangle \geq 0$  for all  $x \in \mathcal{X}$ . Hence, by monotonicity, we get

$$\langle \nabla f(x), x - x^* \rangle \geq \langle \nabla f(x^*), x - x^* \rangle \geq 0 \quad \text{for all } x \in \mathcal{X}, x^* \in \mathcal{X}^*. \quad (2.7)$$

By convexity, it further follows that  $\langle \nabla f(x), x - x^* \rangle < 0$  whenever  $x^* \in \mathcal{X}^*$  and  $x \in \mathcal{X} \setminus \mathcal{X}^*$ , so equality holds in (2.7) if and only if  $x \in \mathcal{X}^*$ . This shows that convex programs automatically satisfy (VC).

*Example 2.4* (Quasi-convex problems). More generally, the above analysis also extends to *quasi-convex* objectives, i.e., when

$$f(x') \leq f(x) \implies \langle \nabla f(x), x' - x \rangle \leq 0 \quad (\text{QC})$$

for all  $x, x' \in \mathcal{X}$  [6]. In this case, we have:

**Proposition 2.2.** *Suppose that  $f$  is quasi-convex and non-degenerate, i.e.,*

$$\langle \nabla f(x), z \rangle \neq 0 \quad \text{for all nonzero } z \in \text{TC}(x), x \in \mathcal{X} \setminus \mathcal{X}^*. \quad (2.8)$$

*Then,  $f$  is variationally coherent.*

*Remark.* The non-degeneracy condition (2.8) is *generic* in that it is satisfied by every quasi-convex function after an arbitrarily small perturbation leaving its minimum set unchanged. In particular, it is automatically satisfied if  $f$  is convex or pseudo-convex.

*Proof.* Take some  $x^* \in \mathcal{X}^*$  and  $x \in \mathcal{X}$ . Then, letting  $x' = x^*$  in (QC), we readily obtain  $\langle \nabla f(x), x - x^* \rangle \geq 0$  for all  $x \in \mathcal{X}, x^* \in \mathcal{X}^*$ . Furthermore, if  $x \notin \mathcal{X}^*$  but  $\langle \nabla f(x), x - x^* \rangle = 0$ , the gradient non-degeneracy condition (2.8) would be violated, implying in turn that, for any  $x^* \in \mathcal{X}^*$ , we have  $\langle \nabla f(x), x - x^* \rangle = 0$  only if  $x \in \mathcal{X}^*$ . This shows that  $f$  satisfies (VC).  $\square$

*Example 2.5* (Beyond quasi-convexity). A simple example of a function that is variationally coherent without even being quasi-convex is

$$f(x) = 2 \sum_{i=1}^d \sqrt{1 + x_i}, \quad x \in [0, 1]^d. \quad (2.9)$$

When  $d \geq 2$ , it is easy to see  $f$  is not quasi-convex: for instance, taking  $d = 2$ ,  $x = (0, 1)$  and  $x' = (1, 0)$  yields  $f(x/2 + x'/2) = 2\sqrt{6} > 2\sqrt{2} = \max\{f(x), f(x')\}$ , so  $f$  is not quasi-convex. On the other hand, to establish (VC), simply note that  $\mathcal{X}^* = \{0\}$  and  $\langle \nabla f(x), x - 0 \rangle = \sum_{i=1}^d x_i / \sqrt{1 + x_i} > 0$  for all  $x \in [0, 1]^d \setminus \{0\}$ .

*Example 2.6* (A weaker version of coherence). Consider the function

$$f(x) = \frac{1}{2} \prod_{i=1}^d x_i^2, \quad x \in [-1, 1]^d. \quad (2.10)$$

By inspection, it is easy to see that the minimum set of  $f$  is  $\mathcal{X}^* = \{x^* \in [-1, 1]^d : x_i^* = 0 \text{ for some } i = 1, \dots, d\}$ .<sup>1</sup> Since  $\mathcal{X}^*$  is not convex for  $d \geq 2$ ,  $f$  is not quasi-convex. On the other hand, we have  $\nabla f(x) = 2f(x) \cdot (1/x_1, \dots, 1/x_d)$ , so  $\langle \nabla f(x), x - 0 \rangle \geq 0$  for all  $x \in [-1, 1]^d$  with equality only if  $x \in \mathcal{X}^*$ . Moreover, for any  $x^* \in \mathcal{X}^*$  and all  $x \in \mathcal{X}$  sufficiently close to  $x^*$ , we have

$$\langle \nabla f(x), x - x^* \rangle = 2f(x) \sum_{i=1}^d \left[ 1 - \frac{x_i^*}{x_i} \right] = 2f(x) \left[ d - \sum_{i: x_i^* \neq 0} \frac{x_i^*}{x_i} \right] \geq 0. \quad (2.11)$$

We thus conclude that  $f$  satisfies the following weaker version of (VC):

**Definition 2.3.** We say that  $f: \mathcal{X} \rightarrow \mathbb{R}$  is *weakly coherent* if the following conditions are satisfied:

- a) There exists some  $p \in \mathcal{X}^*$  such that  $\langle \nabla f(x), x - p \rangle \geq 0$  with equality only if  $x \in \mathcal{X}^*$ .
- b) For all  $x^* \in \mathcal{X}^*$ ,  $\langle \nabla f(x), x - x^* \rangle \geq 0$  whenever  $x$  is close enough to  $x^*$ .

Our analysis also applies to problems satisfying these less stringent requirements, in which case the minimum set  $\mathcal{X}^* = \arg \min f$  of  $f$  need not even be convex.<sup>2</sup> For simplicity, we will first work with Definition 2.1 and relegate these considerations to Section 5.

**2.3. Stochastic mirror descent.** To solve (Opt), we will focus on the family of algorithms known as *stochastic mirror descent* (SMD), a class of first-order methods pioneered by Nemirovski and Yudin [32] and studied further by Beck and Teboulle [3], Nesterov [34], Lan et al. [24], and many others. Referring to [8, 41] for an overview, the specific variant of SMD that we consider here is usually referred to as *dual averaging* (DA) [28, 34, 46] or “*lazy*” mirror descent [41].

The main idea of the method is as follows: At each iteration, the algorithm takes as input an i.i.d. sample of the gradient of  $F$  at the algorithm’s current state. Subsequently, the method takes a step along this stochastic gradient in the dual space  $\mathcal{Y} \equiv \mathcal{V}^*$  of  $\mathcal{V}$  (where gradients live), the result is “mirrored” back to the problem’s feasible region  $\mathcal{X}$ , and the process repeats. Formally, this gives rise to the recursion

$$\begin{aligned} X_n &= Q(Y_n) \\ Y_{n+1} &= Y_n - \gamma_n \nabla F(X_n; \omega_n) \end{aligned} \quad (\text{SMD})$$

where:

<sup>1</sup>Linear combinations of functions of this type play an important role in training deep learning models – and, in particular generative adversarial networks (GANs) [17].

<sup>2</sup>Obviously, Definitions 2.1 and 2.3 coincide if  $\mathcal{X}^*$  is a singleton. This highlights the intricacies that arise in problems that do not admit a unique solution.

**Algorithm 1** Stochastic mirror descent (SMD)

---

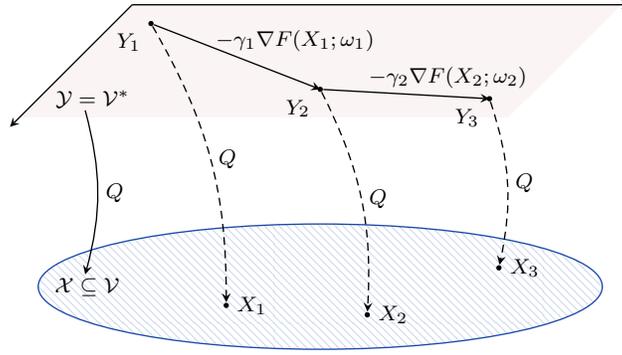
**Require:** mirror map  $Q: \mathcal{Y} \rightarrow \mathcal{X}$ ; step-size sequence  $\gamma_n > 0$

```

1: choose  $Y \in \mathcal{Y} \equiv \mathcal{V}^*$  # initialization
2: for  $n = 1, 2, \dots$  do
3:   set  $X \leftarrow Q(Y)$  # set state
4:   draw  $\omega \in \Omega$  # gradient sample
5:   get  $\hat{v} = -\nabla F(X; \omega)$  # get oracle feedback
6:   set  $Y \leftarrow Y + \gamma_n \hat{v}$  # update score variable
7: end for
8: return  $X$  # output

```

---

**Figure 1.** Schematic representation of SMD (Algorithm 1).

- (1)  $n = 1, 2, \dots$  denotes the algorithm's running counter.
- (2)  $Y_n \in \mathcal{Y}$  is a score variable that aggregates gradient steps up to stage  $n$ .
- (3)  $Q: \mathcal{Y} \rightarrow \mathcal{X}$  is the *mirror map* that outputs a solution candidate  $X_n \in \mathcal{X}$  as a function of the score variable  $Y_n \in \mathcal{V}^*$ .
- (4)  $\omega_n \in \Omega$  is a sequence of i.i.d. samples.<sup>3</sup>
- (5)  $\gamma_n > 0$  is the algorithm's step-size sequence, assumed in what follows to satisfy the Robbins–Monro summability condition

$$\sum_{n=1}^{\infty} \gamma_n^2 < \infty \quad \text{and} \quad \sum_{n=1}^{\infty} \gamma_n = \infty. \quad (2.12)$$

For a schematic illustration and a pseudocode implementation of (SMD), see Fig. 1 and Algorithm 1 respectively.

In more detail, the algorithm's mirror map  $Q: \mathcal{Y} \rightarrow \mathcal{X}$  is defined as

$$Q(y) = \arg \max_{x \in \mathcal{X}} \{ \langle y, x \rangle - h(x) \}, \quad (2.13)$$

where the *regularizer* (or *penalty function*)  $h: \mathcal{X} \rightarrow \mathbb{R}$  is assumed to be continuous and strongly convex on  $\mathcal{X}$ , i.e., there exists some  $K > 0$  such that

$$h(\tau x + (1 - \tau)x') \leq \tau h(x) + (1 - \tau)h(x') - \frac{1}{2}K\tau(1 - \tau)\|x' - x\|^2, \quad (2.14)$$

<sup>3</sup>The indexing convention for  $\omega_n$  means that  $Y_n$  and  $X_n$  are *predictable* relative to the natural filtration  $\mathcal{F}_n = \sigma(\omega_1, \dots, \omega_n)$  of  $\omega_n$ , i.e.,  $Y_{n+1}$  and  $X_{n+1}$  are both  $\mathcal{F}_n$ -measurable. To this history, we also attach the trivial  $\sigma$ -algebra as  $\mathcal{F}_0$  for completeness.

for all  $x, x' \in \mathcal{X}$  and all  $\tau \in [0, 1]$ . The mapping  $Q: \mathcal{V}^* \rightarrow \mathcal{X}$  defined by (2.13) is then called the *mirror map induced by  $h$* . For concreteness, we present below some well-known examples of regularizers and mirror maps:

*Example 2.7* (Euclidean regularization). Let  $h(x) = \frac{1}{2}\|x\|_2^2$ . Then,  $h$  is 1-strongly convex with respect to the Euclidean norm  $\|\cdot\|_2$ , and the induced mirror map is the closest point projection

$$\Pi(y) = \arg \max_{x \in \mathcal{X}} \{ \langle y, x \rangle - \frac{1}{2}\|x\|_2^2 \} = \arg \min_{x \in \mathcal{X}} \|y - x\|_2^2. \quad (2.15)$$

The resulting descent algorithm is known in the literature as (lazy) *stochastic gradient descent* (SGD) and we study it in detail in Section 6. For future reference, we also note that  $h$  is differentiable throughout  $\mathcal{X}$  and  $\Pi$  is *surjective* (i.e.,  $\text{im } \Pi = \mathcal{X}$ ).

*Example 2.8* (Entropic regularization). Let  $\Delta = \{x \in \mathbb{R}_+^d : \sum_{i=1}^d x_i = 1\}$  denote the unit simplex of  $\mathbb{R}^d$ . A widely used regularizer in this setting is the (negative) Gibbs entropy  $h(x) = \sum_{i=1}^d x_i \log x_i$ : this regularizer is 1-strongly convex with respect to the  $L^1$ -norm and a straightforward calculation shows that the induced mirror map is

$$\Lambda(y) = \frac{1}{\sum_{i=1}^d \exp(y_i)} (\exp(y_1), \dots, \exp(y_d)). \quad (2.16)$$

This example is known as *entropic regularization* and the resulting mirror descent algorithm has been studied extensively in the context of linear programming, online learning and game theory [1, 41]. For posterity, we also note that  $h$  is differentiable *only* on the relative interior  $\Delta^\circ$  of  $\Delta$  and  $\text{im } \Lambda = \Delta^\circ$  (i.e.,  $\Lambda$  is “essentially” surjective).

**2.4. Overview of main results.** To motivate the analysis to follow, we provide below a brief overview of our main results:

- *Global convergence:* If (Opt) is variationally coherent, the last iterate  $X_n$  of (SMD) converges to a global minimizer of  $f$  with probability 1.
- *Local convergence:* If  $x^*$  is a *locally coherent* minimum point of  $f$  (a notion introduced in Section 5), the last iterate  $X_n$  of SMD converges locally to  $x^*$  with high probability.
- *Sharp minima:* If  $Q$  is surjective and  $x^*$  is a *sharp* minimum of  $f$  (a fundamental notion due to Polyak which we discuss in Section 6),  $X_n$  reaches  $x^*$  in a *finite* number of iterations (a.s.).

### 3. MAIN TOOLS AND FIRST RESULTS

As a stepping stone to analyze the long-term behavior of (SMD), we derive in this section a recurrence result which is interesting in its own right. Specifically, we show that, if (Opt) is coherent, then, with probability 1,  $X_n$  visits any neighborhood of  $\mathcal{X}^*$  infinitely often; as a corollary, there exists a (random) subsequence  $X_{n_k}$  of  $X_n$  that converges to  $\arg \min f$  (a.s.). In what follows, our goal will be to state this result formally and to introduce the analytic machinery used for its proof (as well as the analysis of the subsequent sections).

**3.1. The Fenchel coupling.** The first key ingredient of our analysis will be the *Fenchel coupling*, a primal-dual variant of the Bregman divergence [7] that plays the role of an energy function for (SMD):

**Definition 3.1.** Let  $h: \mathcal{X} \rightarrow \mathbb{R}$  be a regularizer on  $\mathcal{X}$ . The induced *Fenchel coupling*  $F(p, y)$  between a base-point  $p \in \mathcal{X}$  and a dual vector  $y \in \mathcal{Y}$  is defined as

$$F(p, y) = h(p) + h^*(y) - \langle y, p \rangle \quad (3.1)$$

where  $h^*(y) = \max_{x \in \mathcal{X}} \{\langle y, x \rangle - h(x)\}$  denotes the convex conjugate of  $h$ .

By Fenchel’s inequality (the namesake of the Fenchel coupling), we have  $h(p) + h^*(y) - \langle y, p \rangle \geq 0$  with equality if and only if  $p = Q(y)$ . As such,  $F(p, y)$  can be seen as a (typically asymmetric) “distance measure” between  $p \in \mathcal{X}$  and  $y \in \mathcal{Y}$ . The following lemma quantifies some basic properties of this coupling:

**Lemma 3.2.** *Let  $h$  be a  $K$ -strongly convex regularizer on  $\mathcal{X}$ . Then, for all  $p \in \mathcal{X}$  and all  $y, y' \in \mathcal{Y}$ , we have:*

$$a) \quad F(p, y) \geq \frac{K}{2} \|Q(y) - p\|^2. \quad (3.2a)$$

$$b) \quad F(p, y') \leq F(p, y) + \langle y' - y, Q(y) - p \rangle + \frac{1}{2K} \|y' - y\|_*^2. \quad (3.2b)$$

**Lemma 3.2** (which we prove in [Appendix B](#)) shows that  $Q(y_n) \rightarrow p$  whenever  $F(p, y_n) \rightarrow 0$ , so the Fenchel coupling can be used to test the convergence of the primal sequence  $x_n = Q(y_n)$  to a given base point  $p \in \mathcal{X}$ . For technical reasons, it will be convenient to also make the converse assumption, namely:

**Assumption 3.**  $F(p, y_n) \rightarrow 0$  whenever  $Q(y_n) \rightarrow p$ .

**Assumption 3** can be seen as a “reciprocity condition”: essentially, it means that the sublevel sets of  $F(p, \cdot)$  are mapped under  $Q$  to neighborhoods of  $p$  in  $\mathcal{X}$  (cf. [Appendix B](#)). In this way, **Assumption 3** can be seen as a primal-dual analogue of the reciprocity conditions for the Bregman divergence that are widely used in the literature on proximal and forward-backward methods [9, 23]. Most common regularizers satisfy this technical requirement (including the Euclidean and entropic regularizers of [Examples 2.7](#) and [2.8](#) respectively).

**3.2. Main recurrence result.** To state our recurrence result, we require one last piece of notation pertaining to measuring distances in  $\mathcal{X}$ :

**Definition 3.3.** Let  $\mathcal{C}$  be a subset of  $\mathcal{X}$ .

- (1) The *distance* between  $\mathcal{C}$  and  $x \in \mathcal{X}$  is defined as  $\text{dist}(\mathcal{C}, x) = \inf_{x' \in \mathcal{C}} \|x' - x\|$ , and the corresponding  $\varepsilon$ -neighborhood of  $\mathcal{C}$  is

$$\mathbb{B}(\mathcal{C}, \varepsilon) = \{x \in \mathcal{X} : \text{dist}(\mathcal{C}, x) < \varepsilon\}. \quad (3.3a)$$

- (2) The (setwise) *Fenchel coupling* between  $\mathcal{C}$  and  $y \in \mathcal{Y}$  is defined as  $F(\mathcal{C}, y) = \inf_{x \in \mathcal{C}} F(x, y)$ , and the corresponding *Fenchel  $\delta$ -zone* of  $\mathcal{C}$  under  $h$  is

$$\mathbb{B}_F(\mathcal{C}, \delta) = \{x \in \mathcal{X} : x = Q(y) \text{ for some } y \in \mathcal{Y} \text{ with } F(\mathcal{C}, y) < \delta\}. \quad (3.3b)$$

We then have the following recurrence result for variationally coherent problems:

**Proposition 3.4.** *Fix some  $\varepsilon > 0$  and  $\delta > 0$ . If [\(Opt\)](#) is variationally coherent and [Assumptions 1–3](#) hold, the (random) iterates  $X_n$  of [Algorithm 1](#) enter  $\mathbb{B}(\mathcal{X}^*, \varepsilon)$  and  $\mathbb{B}_F(\mathcal{X}^*, \delta)$  infinitely many times (a.s.).*

**Corollary 3.5.** *With probability 1, there exists a subsequence  $X_{n_k}$  of  $X_n$  converging to a (random) minimum point  $x^*$  of [\(Opt\)](#).*

The proof of [Proposition 3.4](#) consists of three main steps which we outline below:

*Step 1: Martingale properties of  $Y_n$ .* First, let

$$v(x) = -\mathbb{E}[\nabla F(x; \omega)] = -\nabla f(x) \quad (3.4)$$

denote the negative gradient of  $f$  at  $x \in \mathcal{X}$ , and write

$$\hat{v}_n = -\nabla F(X_n; \omega_n) \quad (3.5)$$

for the corresponding oracle feedback at stage  $n$ . Then, [Algorithm 1](#) may be written in Robbins–Monro form as

$$Y_{n+1} = Y_n + \gamma_n \hat{v}_n = Y_n + \gamma_n [v(X_n) + U_n], \quad (3.6)$$

where

$$U_n = \nabla f(X_n) - \nabla F(X_n; \omega_n) \quad (3.7)$$

denotes the difference between the mean gradient of  $f$  at  $X_n$  and the  $n$ -th stage gradient sample.<sup>4</sup> By construction,  $U_n$  is a martingale difference sequence relative to the history (natural filtration)  $\mathcal{F}_n = \sigma(\omega_1, \dots, \omega_n)$  of  $\omega_n$ , i.e.,

$$\mathbb{E}[U_n | \mathcal{F}_{n-1}] = 0 \quad \text{for all } n. \quad (3.8a)$$

Furthermore, by [Assumption 2](#), it readily follows that  $U_n$  has uniformly bounded second moments, i.e., there exists some finite  $\sigma \geq 0$  such that

$$\mathbb{E}[\|U_n\|_*^2 | \mathcal{F}_{n-1}] \leq \sigma^2 \quad \text{for all } n, \quad (3.8b)$$

implying in turn that  $U_n$  is bounded in  $L^2$  (for a more detailed treatment, see [Appendix B](#)).

*Step 2: Recurrence of  $\varepsilon$ -neighborhoods.* Invoking the law of large numbers for  $L^2$ -bounded martingale difference sequences and using the Fenchel coupling as an energy function (cf. [Appendix B](#)), we show that if  $X_n$  remains outside  $\mathbb{B}(\mathcal{X}^*, \varepsilon)$  for sufficiently large  $n$ , we must also have  $F(\mathcal{X}^*, Y_n) \rightarrow -\infty$  (a.s.). This contradicts the positive-definiteness of  $F$ , so  $X_n$  must enter  $\mathbb{B}(\mathcal{X}^*, \varepsilon)$  infinitely often (a.s.).

*Step 3: Recurrence of Fenchel zones.* By reciprocity ([Assumption 3](#)),  $\mathbb{B}_F(\mathcal{X}^*, \delta)$  always contains an  $\varepsilon$ -neighborhood of  $\mathcal{X}^*$ . Since  $X_n$  enters  $\mathbb{B}(\mathcal{X}^*, \varepsilon)$  infinitely many times (a.s.), the same must hold for  $\mathbb{B}_F(\mathcal{X}^*, \delta)$ . Our claim and [Corollary 3.5](#) then follow immediately.

#### 4. GLOBAL CONVERGENCE UNDER COHERENCE

The convergence of a subsequence of  $X_n$  to the minimum set of [\(Opt\)](#) is one of the crucial steps in establishing our first main result:

**Theorem 4.1** (Almost sure global convergence). *Suppose that [\(Opt\)](#) is variationally coherent. Then, under [Assumptions 1–3](#),  $X_n$  converges with probability 1 to a (possibly random) minimum point of [\(Opt\)](#).*

**Corollary 4.2.** *If  $f$  is a non-degenerate quasi-convex (or pseudo-convex, or convex) function and [Assumptions 1–3](#) hold, the last iterate of [\(SMD\)](#) converges with probability 1 to a (possibly random) minimum point of [\(Opt\)](#).*

<sup>4</sup>Recall here that there is a one-step offset between  $X_n$  and  $\omega_{n+1}$  at the  $n$ -th iteration of SMD.

Before discussing the proof of [Theorem 4.1](#), it is important to note that most of the literature surrounding (SMD) and its variants (see e.g., [12, 31, 34, 46] and references therein) focuses on the so-called *ergodic average* of  $X_n$ , i.e.,

$$\bar{X}_n = \frac{\sum_{k=1}^n \gamma_k X_k}{\sum_{k=1}^n \gamma_k} \quad (4.1)$$

Despite the appealing “self-averaging” properties of  $\bar{X}_n$  in convex problems [31, 34], it is not clear how to extend the standard tools used to establish convergence of  $\bar{X}_n$  beyond convex/monotone problems (even to pseudo-convex programs). Since convergence of  $X_n$  automatically implies that of  $\bar{X}_n$ , [Theorem 4.1](#) simultaneously establishes the convergence of the last iterate of SMD and extends existing ergodic convergence results to a wider class of non-convex stochastic programs.

[Corollary 4.2](#) also extends the corresponding results of [29] for the convergence of the last iterate of (SMD) when  $f$  is (strongly) convex and  $h$  has Lipschitz-continuous gradients (so the induced Bregman divergence can be bounded from above by a quadratic surrogate of the primal norm). Our proof strategy is similar and relies on the following lemma, often attributed to Gladyshev [35, p. 49]:<sup>5</sup>

**Lemma 4.3** (Gladyshev). *Let  $a_n, n = 1, 2, \dots$ , be a sequence of nonnegative random variables such that*

$$\mathbb{E}[a_{n+1} \mid a_1, \dots, a_n] \leq (1 + \delta_n)a_n + \varepsilon_n, \quad (4.2)$$

where  $\delta_n$  and  $\varepsilon_n$  are nonnegative deterministic sequences with

$$\sum_{n=1}^{\infty} \delta_n < \infty \quad \text{and} \quad \sum_{n=1}^{\infty} \varepsilon_n < \infty. \quad (4.3)$$

Then,  $a_n$  converges (a.s.) to some random variable  $a_\infty \geq 0$ .

As shown below, this “quasi-Fejérian” monotonicity property plays a critical part in establishing the convergence of (SMD):

*Proof of [Theorem 4.1](#).* Let  $x^* \in \mathcal{X}^*$  be a minimum point of (Opt). Then, letting  $F_n = F(x^*, Y_n)$ , [Lemma 3.2](#) gives

$$\begin{aligned} F_{n+1} &= F(x^*, Y_{n+1}) = F(x^*, Y_n + \gamma_n \hat{v}_n) \\ &\leq F(x^*, Y_n) + \gamma_n \langle \hat{v}_n, X_n - x^* \rangle + \frac{\gamma_n^2}{2K} \|\hat{v}_n\|_*^2 \\ &= F_n + \gamma_n \langle v(X_n), X_n - x^* \rangle + \gamma_n \xi_n + \frac{\gamma_n^2}{2K} \|\hat{v}_n\|_*^2 \\ &\leq F_n + \gamma_n \xi_n + \frac{\gamma_n^2}{2K} \|\hat{v}_n\|_*^2, \end{aligned} \quad (4.4)$$

where we set  $\xi_n = \langle U_n, X_n - x^* \rangle$  in the third line and used the fact that  $f$  satisfies (VC) in the last one. Since  $Y_n$  is predictable relative to  $\mathcal{F}_n$  (i.e.,  $Y_n$  is  $\mathcal{F}_{n-1}$ -measurable), the process  $F_n = F(x^*, Y_n)$  is itself adapted to the shifted filtration

<sup>5</sup>We thank an anonymous reviewer for suggesting this approach. Our original proof strategy relied on the so-called “ODE method” of stochastic approximation [4] and was considerably more intricate.

$\mathcal{F}'_n = \sigma(\omega_1, Y_2, \dots, \omega_{n-1}, Y_n) = \mathcal{F}_{n-1}$ . Thus, taking conditional expectations and invoking [Assumption 2](#), the bound (4.4) becomes:

$$\begin{aligned} \mathbb{E}[F_{n+1} | \mathcal{F}'_n] &\leq F_n + \mathbb{E}[\xi_n | \mathcal{F}'_n] + \frac{\gamma_n^2}{2K} \mathbb{E}[\|\hat{v}_n\|_*^2 | \mathcal{F}'_n] \\ &= F_n + \mathbb{E}[\xi_n | \mathcal{F}_{n-1}] + \frac{\gamma_n^2}{2K} \mathbb{E}[\|\nabla F(X_n; \omega_n)\|_*^2 | \mathcal{F}_{n-1}] \\ &\leq F_n + \frac{\gamma_n^2 V^2}{2K}, \end{aligned} \tag{4.5}$$

where, in the last line, we used [Assumption 2](#) and the fact that  $U_n$  is a martingale difference sequence (so  $\mathbb{E}[\xi_n | \mathcal{F}_{n-1}] = 0$ ; for a detailed derivation, see the proof of [Proposition 3.4](#) in [Appendix B](#)). Hence, with  $\sum_{n=1}^{\infty} \gamma_n^2 < \infty$ , [Lemma 4.3](#) implies that  $F_n$  converges (a.s.) to some finite limit  $F_{\infty}$ .

Now, by [Proposition 3.4](#), there exists (a.s.) a subsequence  $Y_{n_k}$  of  $Y_n$  and some (possibly random)  $x^* \in \mathcal{X}^*$  such that  $\lim_{k \rightarrow \infty} F(x^*, Y_{n_k}) = 0$ . Since the limit  $\lim_{n \rightarrow \infty} F(x^*, Y_n)$  exists (a.s.), it follows that  $\lim_{n \rightarrow \infty} F(x^*, Y_n) = 0$ . This shows that, with probability 1,  $X_n = Q(Y_n)$  converges to some (random) minimum point  $x^*$  of  $(\text{Opt})$ , as claimed.  $\square$

In closing this section, we should note that the conclusion of [Theorem 4.1](#) also applies to problems that are “almost” coherent in the sense of [Example 2.6](#), i.e.,

- a) There exists a minimizer  $p \in \mathcal{X}^*$  such that  $\langle \nabla f(x), x - p \rangle \geq 0$  with equality only if  $x \in \mathcal{X}^*$ .
- b) For all  $x^* \in \mathcal{X}^*$ ,  $\langle \nabla f(x), x - x^* \rangle \geq 0$  whenever  $x$  is close enough to  $x^*$ .

Proving this more general result requires some of the machinery presented in the following section, so we relegate its discussion until all the requisite tools are in place.

## 5. CONVERGENCE UNDER LOCAL/WEAK COHERENCE

In this section, our goal is to extend the convergence analysis of the previous section to account for optimization problems that are only “locally” coherent. Building on [Definition 2.1](#), these are defined as follows:

**Definition 5.1.** Let  $\mathcal{C}$  be a closed set of local minimizers of  $f$ , viz.  $f(x) \geq f(x^*)$  for all  $x^* \in \mathcal{C}$  and all  $x$  sufficiently close to  $\mathcal{C}$ . We say that  $\mathcal{C}$  is *locally coherent* if there exists an open neighborhood  $U$  of  $\mathcal{C}$  such that

$$\langle \nabla f(x), x - x^* \rangle \geq 0 \quad \text{for all } x \in U, x^* \in \mathcal{C}, \tag{LVC}$$

and there exists some  $x^* \in \mathcal{C}$  such that equality holds in (LVC) only if  $x \in \mathcal{C}$ .

An immediate consequence of [Definition 5.1](#) is that locally coherent sets are isolated components of local minimizers of  $f$ . To see this, if  $\mathcal{C}$ ,  $U$  and  $x^*$  are as in [Definition 5.1](#) and  $x \in U$  is a local minimizer of  $f$ , we would have  $\langle \nabla f(x), z \rangle \geq 0$  for all tangent  $z \in \text{TC}(x)$ . Applying this to  $z = x^* - x$  gives  $\langle \nabla f(x), x - x^* \rangle \geq 0$ , so, by the definition of local coherence, we conclude that  $x \in \mathcal{C}$ .

We also note that although the minimum set of a globally coherent problem is *a fortiori* locally coherent, the converse need not hold. A concrete example of a function which is not globally coherent but which admits a locally coherent minimum

is the Rosenbrock test function

$$f(x) = \sum_{i=1}^d [100(x_{i+1} - x_i)^2 + (1 - x_i^2)], \quad x \in [-2, 2]^d, \quad (5.1)$$

which has seen extensive use in the literature as a non-convex convergence speed benchmark (cf. Section 7).<sup>6</sup> From this example, we see that the profile of  $f$  around a locally coherent set could be highly non-convex, possibly including a wide variety of valleys, talwegs and ridges; in fact, even *quasi*-convexity may fail to hold locally.

Now, in contrast to globally coherent optimization problems, an “unlucky” gradient sample could drive (SMD) out of the “basin of attraction” of a locally coherent set (the largest neighborhood  $U$  for which (LVC) holds), possibly never to return. For this reason, instead of focusing on global convergence results with probability one, we will focus on local convergence with high probability. Our main result along these lines is as follows:

**Theorem 5.2** (Local convergence with high probability). *Let  $\mathcal{C}$  be locally coherent for (Opt) and fix some confidence level  $\delta > 0$ . Then, under Assumptions 1–3, there exists an open neighborhood  $\mathcal{U}$  of  $\mathcal{C}$ , independent of  $\delta$ , such that*

$$\mathbb{P}(X_n \text{ converges to } \mathcal{C} \mid X_1 \in \mathcal{U}) \geq 1 - \delta, \quad (5.2)$$

*provided that the algorithm’s step-size sequence  $\gamma_n$  is small enough.*

*Remark.* As a concrete application of Theorem 5.2, fix any  $\beta \in (1/2, 1]$ . Then, for every confidence level  $\delta > 0$ , Theorem 5.2 implies that there exists some small enough  $\gamma > 0$  such that if Algorithm 1 is run with step-size  $\gamma_n = \gamma/n^\beta$ , Eq. (5.2) holds. We emphasize the interesting point here: the open neighborhood  $\mathcal{U}$  is fixed once and for all, and does not depend on the probability threshold  $\delta$ . That is, to get convergence with higher probability, it is *not* necessary to assume that  $X_1$  starts closer to  $\mathcal{C}$ : one need only use a smaller step-size sequence satisfying (2.12).

The key idea behind the proof of Theorem 5.2 is as follows: First, it suffices to consider the case where  $\mathcal{C}$  consists of a single local minimizer  $x^*$ ; the argument for the general follows the same techniques as in Section 4. Then, conditioning on the event that  $X_n$  remains sufficiently close to  $x^*$  for all  $n$ , convergence can be obtained by invoking Theorem 4.1 and treating (Opt) as a variationally coherent problem over a smaller subset of  $\mathcal{X}$  over which (LVC) holds. Therefore, to prove Theorem 5.2, it suffices to show that  $X_n$  remains close to  $x^*$  for all  $n$  with probability no less than  $1 - \delta$ . To achieve this, we rely again on the properties of the Fenchel coupling, and we decompose the stochastic errors affecting each iteration of the algorithm into a first-order  $\mathcal{O}(\gamma_n)$  martingale term and a second-order  $\mathcal{O}(\gamma_n^2)$  submartingale perturbation. Using Doob’s maximal inequality, we then show that the aggregation of both errors remains controllably small with probability at least  $1 - \delta$ .

We formalize all this below:

*Proof of Theorem 5.2.* We break the proof into three steps.

*Step 1: Controlling the martingale error.* Fix some  $\epsilon > 0$ . As in the proof of Theorem 4.1, let  $U_n = \nabla f(X_n) - \nabla F(X_n; \omega_n)$  and set  $\xi_n = \langle U_n, X_n - x^* \rangle$  where

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<sup>6</sup>Local coherence can be proved by a straightforward algebraic calculation (omitted for concision).

$x^* \in \mathcal{C}$  is such that (LVC) holds as an equality only if  $x \in \mathcal{C}$  (cf. Definition 5.1). We show below that there exists a step-size sequence  $(\gamma_n)_{n=1}^\infty$  such that

$$\mathbb{P}\left(\sup_n \sum_{k=1}^n \gamma_k \xi_k \leq \varepsilon\right) \geq 1 - \frac{\delta}{2}. \quad (5.3)$$

To show this, we start by noting that, as in the proof of Proposition 3.4, the aggregate process  $S_n = \sum_{k=1}^n \gamma_k \xi_k$  is a martingale relative to the natural filtration  $\mathcal{F}_n$  of  $\omega_n$ . Then, letting  $R = \sup_{x \in \mathcal{X}} \|x\|$ , we can bound the variance of each individual term of  $S_n$  as follows:

$$\begin{aligned} \mathbb{E}[\xi_k^2] &= \mathbb{E}[\mathbb{E}[\langle U_k, X_k - x^* \rangle^2 \mid \mathcal{F}_{k-1}]] \\ &\leq \mathbb{E}[\mathbb{E}[\|U_k\|_*^2 \|X_k - x^*\|^2 \mid \mathcal{F}_{k-1}]] \\ &= \mathbb{E}[\|X_k - x^*\|^2 \mathbb{E}[\|U_k\|_*^2 \mid \mathcal{F}_{k-1}]] \\ &\leq R^2 V^2, \end{aligned} \quad (5.4)$$

where the first inequality follows from the definition of the dual norm and the second one follows from (3.8b). Consequently, by Doob's maximal inequality (Theorem A.4 in Appendix A), we have:

$$\mathbb{P}\left(\sup_{0 \leq k \leq n} S_k \geq \varepsilon\right) \leq \mathbb{P}\left(\sup_{0 \leq k \leq n} |S_k| \geq \varepsilon\right) \leq \frac{\mathbb{E}[S_n^2]}{\varepsilon^2} \leq \frac{R^2 V^2 \sum_{k=1}^n \gamma_k^2}{\varepsilon^2}, \quad (5.5)$$

where the last inequality follows from expanding  $\mathbb{E}[|S_n|^2]$ , using Eq. (5.4), and noting that  $\mathbb{E}[\xi_k \xi_\ell] = \mathbb{E}[\mathbb{E}[\xi_k \xi_\ell \mid \mathcal{F}_{k \vee \ell - 1}]] = 0$  whenever  $k \neq \ell$ . Therefore, by picking  $\gamma_n$  so that  $\sum_{k=1}^\infty \gamma_k^2 \leq \varepsilon^2 \delta / (2R^2 V^2)$ , Eq. (5.5) gives

$$\mathbb{P}\left(\sup_{0 \leq k \leq t} S_k \geq \varepsilon\right) \leq \frac{R^2 V^2 \sum_{k=1}^n \gamma_k^2}{\varepsilon^2} \leq \frac{R^2 V^2 \sum_{k=1}^\infty \gamma_k^2}{\varepsilon^2} \leq \frac{\delta}{2} \quad \text{for all } n. \quad (5.6)$$

Since the above holds for all  $n$ , our assertion follows.

*Step 2: Controlling the submartingale error.* Again, fix some  $\varepsilon > 0$  and, with a fair amount of foresight, let  $R_n = (2K)^{-1} \sum_{k=1}^n \gamma_k^2 \|\hat{v}_k\|_*^2$ . By construction,  $R_n$  is a non-negative submartingale relative to  $\mathcal{F}_n$ . We again establish that there exists step-size sequence  $(\gamma_n)_{n=1}^\infty$  satisfying the summability condition (2.12) and such that

$$\mathbb{P}\left(\sup_n R_n \leq \varepsilon\right) \geq 1 - \frac{\delta}{2}. \quad (5.7)$$

To show this, Doob's maximal inequality for submartingales (Theorem A.3) gives

$$\mathbb{P}\left(\sup_{0 \leq k \leq n} R_k \geq \varepsilon\right) \leq \frac{\mathbb{E}[R_n]}{\varepsilon} \leq \frac{V^2 \sum_{k=1}^n \gamma_k^2}{2K\varepsilon}, \quad (5.8)$$

where we used the fact that  $\mathbb{E}[\|\nabla F(X_n; \omega_n)\|_*^2] \leq V^2$  for some finite  $V < \infty$ . Consequently, if we choose  $\gamma_n$  so that  $\sum_{k=1}^\infty \gamma_k^2 \leq K\delta\varepsilon/V^2$ , Eq. (5.8) readily gives

$$\mathbb{P}\left(\sup_{0 \leq k \leq n} R_k \geq \varepsilon\right) \leq \frac{V^2 \sum_{k=1}^\infty \gamma_k^2}{2K\varepsilon} \leq \frac{\delta}{2} \quad \text{for all } n \quad (5.9)$$

Since the above is true for all  $n$ , Eq. (5.7) follows.

*Step 3: Error aggregation.* To combine the above, fix some sufficiently small  $\bar{\varepsilon} > 0$  so that  $\mathbb{B}_F(x^*, 3\bar{\varepsilon}) \subset U$ , where  $U$  is the open neighborhood given in LVC.

Furthermore, let  $\mathcal{U} = \mathbb{B}_F(x^*, \bar{\varepsilon})$  and pick a step-size sequence  $\gamma_n$  satisfying (2.12) and such that

$$\sum_{n=1}^{\infty} \gamma_n^2 \leq \min \left\{ \frac{\delta \bar{\varepsilon}^2}{2R^2V^2}, \frac{K\delta\bar{\varepsilon}}{V^2} \right\}. \quad (5.10)$$

If  $X_1 \in \mathcal{U}$ , it follows that  $F(x^*, Y_1) < \bar{\varepsilon}$  by the definition of  $\mathbb{B}_F$  (cf. Definition 3.3). Then, by Eqs. (5.3) and (5.7), we get  $\mathbb{P}(\sup_n S_n \geq \bar{\varepsilon}) \leq \delta/2$  and  $\mathbb{P}(\sup_n R_n \geq \bar{\varepsilon}) \leq \delta/2$ . Consequently, with this choice of  $\gamma_n$ , it follows that

$$\mathbb{P}(\sup_n \max\{S_n, R_n\} \leq \bar{\varepsilon}) \geq 1 - \delta/2 - \delta/2 = 1 - \delta \quad (5.11)$$

Then, letting  $F_n = F(x^*, Y_n)$  and arguing as in the proof of Theorem 4.1, we may expand  $F_n = F(x^*, Y_n)$  to get

$$\begin{aligned} F_n &= F(x^*, Y_n + \gamma_n \hat{v}_n) \\ &\leq F_n + \gamma_n \langle v(X_n), X_n - x^* \rangle + \gamma_n \xi_n + \frac{\gamma_n^2}{2K} \|\nabla F(X_n; \omega_n)\|_*^2 \end{aligned} \quad (5.12)$$

with  $\xi_n = \langle U_n, X_n - x^* \rangle$  defined as above. Telescoping (5.12) then yields

$$F_n \leq F_1 + \sum_{k=1}^n \gamma_k \langle v(X_k), X_k - x^* \rangle + S_n + R_n \quad (5.13)$$

$$\leq \bar{\varepsilon} + \sum_{k=1}^n \gamma_k \langle v(X_k), X_k - x^* \rangle + \bar{\varepsilon} + \bar{\varepsilon}, \quad (5.14)$$

with probability at least  $1 - \delta$ . Therefore, with probability at least  $1 - \delta$ , we have

$$F(x^*, Y_n) \leq 3\bar{\varepsilon} + \sum_{k=1}^n \gamma_k \langle v(X_k), X_k - x^* \rangle. \quad (5.15)$$

Now, assume inductively that, for all  $k \leq n$ , we have  $F(x^*, Y_k) \leq 3\bar{\varepsilon}$  or, equivalently,  $X_k \in \mathbb{B}_F(x^*, 3\bar{\varepsilon})$ . In turn, this implies that  $\langle v(X_k), X_k - x^* \rangle \leq 0$  for all  $k \leq n$ , and hence, by (5.15), that  $F(x^*, Y_n) \leq 3\bar{\varepsilon}$  as well. Since the base case  $X_1 \in \mathcal{U} = \mathbb{B}_F(x^*, \bar{\varepsilon}) \subset \mathbb{B}_F(x^*, 3\bar{\varepsilon})$  is satisfied automatically, we conclude that  $X_n$  stays in  $\mathbb{B}_F(x^*, 3\bar{\varepsilon}) \subset \mathcal{U}$  for all  $n$  with probability at least  $1 - \delta$ . Our claim then follows by conditioning on this event and repeating the same steps as in the proof of Theorem 4.1.  $\square$

We close this section by revisiting the notion of weak coherence (Definition 2.3). In view of Definition 5.1, we see that weak coherence mixes elements of both global *and* local coherence: on the one hand, it posits the existence of a (global) minimizer  $p \in \mathcal{X}^*$  for which (VC) holds globally, thus satisfying the second part of Definition 2.1; on the other hand, minimizers other than  $p$  are only required to satisfy (VC) locally (though they need not be locally coherent themselves). From a stability viewpoint, this means that individual elements of a weakly coherent set may be stable but not necessarily attracting (even locally). However, taken as a whole, weakly coherent sets are *globally* attracting:

**Theorem 5.3.** *Suppose that (Opt) is weakly coherent. Then, under Assumptions 1–3,  $X_n$  converges with probability 1 to a (possibly random) minimum point of (Opt).*

*Proof.* The proof is essentially a combination of the proofs of Theorems 4.1 and 5.2, so we only provide the main arguments and omit the minor details.

The first observation is that the conclusion of [Proposition 3.4](#) only requires the first part of [Definition 2.3](#) (simply take  $x^* = p$  in the proof of [Proposition 3.4](#)). From this, we conclude that the generated sequence  $X_n$  admits (a.s.) a subsequence  $X_{n_k}$  converging to some  $x^* \in \mathcal{X}^*$ .

In view of this, arguing as in the proof of [Theorem 5.2](#) allows us to conclude that, with probability 1,  $X_n$  itself remains in some neighborhood  $\mathcal{U}$  of  $x^*$  such that  $\langle \nabla f(x), x - x^* \rangle \geq 0$  for all  $x \in \mathcal{U}$ . To see this, recall first that  $\sum_{n=1}^{\infty} \gamma_n^2 < \infty$ ; thus, by starting at some sufficiently large  $n_0$  (recall that  $X_n$  comes arbitrarily close to  $x^*$  infinitely many times), we can assume without loss of generality that (5.10) is satisfied for any  $\varepsilon, \delta > 0$ . Since  $\delta$  is arbitrary, this means that (5.15) actually holds with probability 1, and our claim follows.

The above shows that  $\langle \nabla f(X_n), X_n - x^* \rangle$  for all sufficiently large  $n$ . Thus, going back to the proof of [Theorem 4.1](#), it follows that the Fenchel coupling  $F_n = F(x^*, Y_n)$  converges. Since  $\liminf_{n \rightarrow \infty} F_n = 0$  (by [Assumption 3](#) and the fact that  $X_{n_k} = Q(Y_{n_k}) \rightarrow x^*$ ), we conclude that  $\lim_{n \rightarrow \infty} X_n = x^*$ , and our proof is complete.  $\square$

## 6. SHARP MINIMA AND APPLICATIONS

Given the randomness involved at each step, obtaining an almost sure (or high probability) bound for the convergence speed of the last iterate of SMD is fairly involved: indeed, in contrast to the ergodic rate analysis of SMD for convex programs, there is no intrinsic averaging in the algorithm's last iterate, so it does not seem possible to derive a precise black box convergence rate for  $X_n$ . Essentially, as in the analysis of [Section 5](#), a single “unlucky” gradient sample could violate any convergence speed estimate that is probabilistically independent of any finite subset of realizations.

Despite this difficulty, if SMD is run with a *surjective* mirror map, we show below that  $X_n$  reaches a minimum point of [\(Opt\)](#) in a *finite* number of iterations for a large class of optimization problems that admit *sharp* minima (see below). As we noted in [Section 2](#), an important example of a surjective mirror map is the standard Euclidean projection  $\Pi(y) = \arg \min_{x \in \mathcal{X}} \|y - x\|_2$ . The resulting descent method is the well-known *stochastic gradient descent* (SGD) algorithm (cf. [Algorithm 2](#) below), so our results in this section also provide new insights into the behavior of SGD.

**6.1. Definition and characterization.** The starting point of our analysis is Polyak's fundamental notion of a *sharp minimum* [[35](#), Chapter 5.2], which describes functions that grow at least linearly around their minimum points:

**Definition 6.1.** We say that  $x^* \in \mathcal{X}$  is a  $\rho$ -*sharp* (local) minimum of  $f$  if

$$f(x) \geq f(x^*) + \rho \|x - x^*\| \quad \text{for some } \rho > 0 \text{ and all } x \text{ sufficiently close to } x^*. \quad (6.1)$$

Polyak's original definition concerned global sharp minima of unconstrained (convex) optimization problems; by contrast, the above definition is tailored to local optima of constrained (and possibly non-convex) programs. In particular, [Definition 6.1](#) implies that sharp minima are isolated (local) minimizers of  $f$ , and they remain invariant under small perturbations of  $f$  (assuming of course that such a minimizer exists in the first place). In what follows, we shall omit the modifier “local” for concision and rely on the context to resolve any ambiguities.

Sharp minima admit a useful geometric interpretation in terms of the polar cone of  $\mathcal{X}$ . To state it, recall first the following basic facts from convex analysis:

**Definition 6.2.** Let  $\mathcal{X}$  be a closed convex subset of  $\mathbb{R}^d$ . Then:

- (1) The *tangent cone*  $\text{TC}(p)$  to  $\mathcal{X}$  at  $p$  is defined as the closure of the set of all rays emanating from  $p$  and intersecting  $\mathcal{X}$  in at least one other point.
- (2) The *dual cone*  $\text{TC}^*(p)$  to  $\mathcal{X}$  at  $p$  is the dual set of  $\text{TC}(p)$ , viz.  $\text{TC}^*(p) = \{y \in \mathbb{R}^d : \langle y, z \rangle \geq 0 \text{ for all } z \in \text{TC}(p)\}$ .
- (3) The *polar cone*  $\text{PC}(p)$  to  $\mathcal{X}$  at  $p$  is the polar set of  $\text{TC}(p)$ , viz.  $\text{PC}(p) = -\text{TC}^*(p) = \{y \in \mathbb{R}^d : \langle y, z \rangle \leq 0 \text{ for all } z \in \text{TC}(p)\}$ .

The above gives the following geometric characterization of sharp minima:

**Lemma 6.3.** *If  $x^* \in \mathcal{X}$  is a  $\rho$ -sharp minimum of  $f$ , we have*

$$\langle \nabla f(x^*), z \rangle \geq \rho \|z\| \quad \text{for all } z \in \text{TC}(x^*). \quad (6.2)$$

*In particular,  $\nabla f(x^*)$  belongs to the topological interior of  $\text{TC}^*(x)$ . Conversely, if (6.2) holds and  $f$  is convex,  $x^*$  is sharp.*

*Proof of Lemma 6.3.* For the direct implication, fix some  $x \in \mathcal{X}$  satisfying (6.1), and let  $z = x - x^* \in \text{TC}(x^*)$ . Then, by the definition of a sharp minimum, we get

$$f(x^* + \tau z) \geq f(x^*) + \rho \tau \|z\| \quad \text{for all } \tau \in [0, 1]. \quad (6.3)$$

In turn, this implies that

$$\frac{f(x^* + tz) - f(x^*)}{t} \geq \rho \|z\| \quad \text{for all sufficiently small } t > 0. \quad (6.4)$$

Hence, taking the limit  $\tau \rightarrow 0^+$ , we get  $\langle \nabla f(x^*), z \rangle \geq \rho \|z\|$ , and our claim follows from the definition of  $\text{TC}(x^*)$  as the closure of the set of all rays emanating from  $x^*$  and intersecting  $\mathcal{X}$  in at least one other point. As for the converse implication, simply note that  $f(x) - f(x^*) \geq \langle \nabla f(x^*), x - x^* \rangle \geq \rho \|x - x^*\|$  if  $f$  is convex.  $\square$

*Example 6.1 (Linear programs).* A first important class of examples of functions that admit sharp minima is that of generic linear programs.<sup>7</sup> Indeed, by definition, a linear function grows (exactly) linearly around its minimum points so, by genericity,

*Example 6.2 (Concave minimization).* For a non-convex class of examples, let  $f: \mathcal{X} \rightarrow \mathbb{R}$  be a strictly *concave* function defined over a convex polytope  $\mathcal{X}$  of  $\mathbb{R}^d$ . Concavity implies that  $f$  is superharmonic, i.e.,

$$\Delta f(x) = \sum_{i=1}^d \frac{\partial^2 f}{\partial x_i^2} \leq 0 \quad (6.5)$$

for all  $x \in \mathcal{X}$ .<sup>8</sup> By the minimum principle for superharmonic functions, the minimum of  $f$  over any connected region  $\mathcal{C}$  of  $\mathcal{X}$  is attained at the boundary of  $\mathcal{C}$ . Hence, by strict concavity, we conclude that the local minima of  $f$  are attained at 0-dimensional faces of  $\mathcal{X}$ , and they are *de facto* sharp (simply note that  $f$  is strictly concave along any ray of the form  $x^* + tz$ ,  $z \in \text{TC}(x^*)$ ).

<sup>7</sup>“Generic” means here that  $\mathcal{X}$  is a polytope,  $f: \mathcal{X} \rightarrow \mathbb{R}$  is affine, and  $f$  is constant only on the zero-dimensional faces of  $\mathcal{X}$ . Any linear program can be turned into a generic one after an arbitrarily small perturbation.

<sup>8</sup>We tacitly assume above that  $f$  is twice-differentiable but this conclusion still holds even if  $f$  is not differentiable.

Sharp minima have several other interesting and useful properties. First, by [Lemma 6.3](#), sharp minimum points are locally coherent. To see this, simply note that for all  $x \in \mathcal{X}$  sufficiently close to  $x^*$  (with  $x \neq x^*$ ), we have  $z = x - x^* \in \text{TC}(x^*)$  and  $\langle \nabla f(x^*), z \rangle \geq \rho \|z\| > 0$ . Consequently,  $\langle \nabla f(x^*), x - x^* \rangle > 0$ , implying by continuity that  $\langle \nabla f(x), x - x^* \rangle > 0$  for all  $x$  in some open neighborhood of  $x^*$  (excluding  $x^*$ ). In addition, if [\(Opt\)](#) is variationally coherent, then a sharp (local) minimum is globally sharp as well.

A second important property is that the dual cone  $\text{TC}^*(x^*)$  of a sharp minimum must necessarily have nonempty topological interior – since it contains  $\nabla f(x^*)$  by [Lemma 6.3](#). This implies that sharp minima can only occur at *corners* of  $\mathcal{X}$ : for instance, if a sharp minimum were an interior point of  $\mathcal{X}$ , the dual cone to  $\mathcal{X}$  at  $x^*$  would be a proper linear subspace of the ambient vector space, so it would have no topological content (see also [Example 6.2](#) above).

**6.2. Global convergence in a finite number of iterations.** We now turn to showing that, if a variationally coherent program admits a sharp minimum  $x^*$ , [Algorithm 1](#) reaches  $x^*$  in a finite number of iterations (a.s.). The interesting feature here is that convergence is guaranteed to occur in a *finite* number of iterations: specifically, there exists some (random)  $n_0$  such that  $X_n = x^*$  for all  $n \geq n_0$ . In general, this is a fairly surprising property for a first-order descent scheme, even if one considers the ergodic average  $n^{-1} \sum_{k=1}^n X_k$ : a priori, a single “bad” sample could kick  $X_n$  away from  $x^*$ , which is the reason why (ergodic) convergence rates are typically asymptotic.

The key intuition behind our analysis is that sharp minima must occur at corners of  $\mathcal{X}$  (as opposed to interior points). As a further key insight, when the solution of [\(Opt\)](#) occurs at a corner, noisy gradients may still play the role of a random disturbance; however, since they are applied to the dual process  $Y_n$ , a surjective mirror map would immediately project  $Y_n$  back to a corner of  $\mathcal{X}$  if  $Y_n$  has progressed far enough in the interior of the polar cone to  $\mathcal{X}$  at  $x$ . This ensures that the last iterate  $X_n$  of SMD will stay *exactly* at the optimal point, irrespective of the persistent noise entering [Algorithm 1](#).

Exploiting these insights and the structural properties of sharp minima, we have:

**Theorem 6.4.** *Suppose that [\(Opt\)](#) is variationally coherent. If  $f$  admits a (necessarily unique) sharp minimum  $x^*$ , and [Algorithm 1](#) is run with a surjective mirror map and [Assumptions 1–3](#) hold,  $X_n$  converges to  $x^*$  in a finite number of steps (a.s.). More precisely, we have*

$$\mathbb{P}(X_n = x^* \text{ for all sufficiently large } n) = 1 \quad (6.6)$$

**Corollary 6.5.** *Let  $f$  be a non-degenerate quasi-convex (or pseudo-convex, or convex) function and let  $x^*$  be a sharp minimum of  $f$ . Then, with assumptions as above,  $X_n$  reaches  $x^*$  in a finite number of steps (a.s.).*

The prototypical example of a surjective mirror map is the Euclidean projector  $\Pi(y) = \arg \min_{x \in \mathcal{X}} \|y - x\|_2$  induced by the quadratic regularization function  $h(x) = \|x\|_2^2/2$  (cf. [Example 2.7](#)). The resulting descent scheme is the well-known stochastic gradient descent (SGD) algorithm (see [Algorithm 2](#) for a pseudocode implementation), for which we obtain the following novel convergence result:

**Corollary 6.6.** *If [\(Opt\)](#) is a generic linear program, the last iterate  $X_n$  of SGD reaches its (necessarily unique) solution in a finite number of steps (a.s.).*

**Algorithm 2** Stochastic gradient descent (SGD)

---

**Require:** step-size sequence  $\gamma_n > 0$

```

1: choose  $Y \in \mathbb{R}^d$ ,  $X = \Pi(Y)$                                 # initialization
2: for  $n = 1, 2, \dots$  do
3:   get  $\hat{v} = -\nabla F(X; \omega)$                                     # oracle feedback
4:   set  $Y \leftarrow Y + \gamma_n \hat{v}$                                 # gradient step
5:   set  $X \leftarrow \Pi(Y)$                                         # set state
6: end for
7: return  $X$                                                     # output

```

---

With all this said and done, we proceed with the proof of [Theorem 6.4](#):

*Proof of [Theorem 6.4](#).* Since  $x^*$  is a  $\rho$ -sharp minimum, there exists a sufficiently small open neighborhood  $\mathcal{U}$  of  $x^*$  such that  $\langle \nabla f(x), z \rangle \geq \rho \|z\|/2$  for all  $z \in \text{TC}(x^*)$  and all  $x \in \mathcal{U}$ . By [Theorem 4.1](#),  $X_n$  converges to  $x^*$  (a.s.), so there exists some (random)  $n_0$  such that  $X_n \in \mathcal{U}$  for all  $n \geq n_0$ . In turn, this implies that  $\langle \nabla f(X_n), z \rangle \geq \rho \|z\|/2$  for all  $n \geq n_0$ . Thus, continuing to use the notation  $v(X_n) = -\nabla f(X_n)$  and  $U_n = \nabla f(X_n) - \nabla F(X_n; \omega_n)$ , we get for all  $z \in \text{TC}(x^*)$  with  $\|z\| \leq 1$ :

$$\begin{aligned}
\langle Y_n, z \rangle &= \left\langle Y_{n_0} + \sum_{k=n_0}^n \gamma_k \hat{v}_k, z \right\rangle \\
&= \langle Y_{n_0}, z \rangle + \sum_{k=n_0}^n \gamma_k \langle v(X_k), z \rangle + \sum_{k=n_0}^n \gamma_k \langle U_k, z \rangle \\
&\leq \|Y_{n_0}\|_* - \frac{\rho}{2} \sum_{k=n_0}^n \gamma_k + \sum_{k=n_0}^n \gamma_k \langle U_k, z \rangle, \tag{6.7}
\end{aligned}$$

where, in the last line, we used the fact that  $X_k \in \mathcal{U}$  for all  $k \geq n_0$ .

As we discussed in the proof of [Theorem 4.1](#),  $\gamma_n U_n$  is a martingale difference sequence relative to the natural filtration  $\mathcal{F}_n$  of  $\omega_n$ . Hence, by the law of large numbers for martingale differences (cf. [Theorem A.1](#) for  $p = 2$  and  $\tau_n = \sum_{k=0}^n \gamma_k$ ), we get

$$\lim_{n \rightarrow \infty} \frac{\sum_{k=n_0}^n \gamma_k U_k}{\sum_{k=n_0}^n \gamma_k} = 0 \quad (a.s.). \tag{6.8}$$

Thus, there exists some  $n^*$  such that  $\|\sum_{k=n_0}^n \gamma_k U_k\|_* \leq (\rho/4) \sum_{k=n_0}^n \gamma_k$  for all  $n \geq n^*$  (a.s.). [Eq. \(6.7\)](#) then implies

$$\begin{aligned}
\langle Y_n, z \rangle &\leq \|Y_{n_0}\|_* - \frac{\rho}{2} \sum_{k=n_0}^n \gamma_k + \sum_{k=n_0}^n \gamma_k \langle U_k, z \rangle \\
&\leq \|Y_{n_0}\|_* - \frac{\rho}{2} \sum_{k=n_0}^n \gamma_k + \frac{\rho}{4} \sum_{k=n_0}^n \gamma_k = \|Y_{n_0}\|_* - \frac{\rho}{4} \sum_{k=n_0}^n \gamma_k, \tag{6.9}
\end{aligned}$$

where we used the assumption that  $\|z\| \leq 1$ . Since  $\sum_{k=n_0}^n \gamma_k \rightarrow \infty$  as  $n \rightarrow \infty$ , we get  $\langle Y_n, z \rangle \rightarrow -\infty$  with probability 1.

To proceed, we claim that  $y^* + \text{PC}(x^*) \subseteq Q^{-1}(x^*)$  whenever  $Q(y^*) = x^*$ , i.e.,  $Q^{-1}(x^*)$  contains all cones of the form  $y^* + \text{PC}(x^*)$  for  $y^* \in Q^{-1}(x^*)$ . Indeed, note

first that  $x^* = Q(y^*)$  if and only if  $y^* \in \partial h(x^*)$ , where  $\partial h(x^*)$  is the set of all subgradients of  $h$  at  $x^*$  [38]. Therefore, it suffices to show that  $y^* + w \in \partial h(x^*)$  whenever  $w \in \text{PC}(x^*)$ . To that end, note that the definition of the polar cone gives

$$\langle w, x - x^* \rangle \leq 0 \quad \text{for all } x \in \mathcal{X}, w \in \text{PC}(x^*), \quad (6.10)$$

and hence

$$h(x) \geq h(x^*) + \langle y^*, x - x^* \rangle \geq h(x^*) + \langle y^* + w, x - x^* \rangle. \quad (6.11)$$

The above shows that  $y^* + w \in \partial h(x^*)$ , as claimed.

With  $Q$  surjective, the set  $Q^{-1}(x^*)$  is nonempty, so it suffices to show that  $Y_n$  lies in the cone  $y^* + \text{PC}(x^*)$  for some  $y^* \in Q^{-1}(x^*)$  and all sufficiently large  $n$ . To do so, simply note that  $Y_n \in y^* + \text{PC}(x^*)$  if and only if  $\langle Y_n - y^*, z \rangle \leq 0$  for all  $z \in \text{TC}(x^*)$  with  $\|z\| = 1$ . Since  $\langle Y_n, z \rangle$  converges to  $-\infty$  (a.s.), our assertion is immediate.  $\square$

**6.3. Local convergence in a finite number of iterations.** Our convergence analysis for locally coherent sets of minimizers (cf. Section 5), showed that SMD converges locally with high probability. Our last result in this section complements this analysis by showing that, with high probability, SMD converges locally to sharp local minima in a *finite* number of iterations:

**Theorem 6.7.** *Let  $x^*$  be a sharp (local) minimum of  $f$ , and fix some confidence level  $\delta > 0$ . If Algorithm 1 is run with a surjective mirror map and Assumptions 1–3 hold, there exists an open neighborhood  $\mathcal{U}$  of  $x^*$ , independent of  $\delta$ , such that*

$$\mathbb{P}(X_n = x^* \text{ for all sufficiently large } n \mid X_1 \in \mathcal{U}) \geq 1 - \delta, \quad (6.12)$$

provided that the algorithm’s step-size sequence  $\gamma_n$  is small enough.

Given that local minimizers of concave minimization programs are sharp, an application of Theorem 6.7 gives the following convergence result for SGD:

**Corollary 6.8.** *Suppose that  $f$  is strictly concave as in Example 6.2. Then, under assumptions Assumptions 1 and 2, the last iterate of SGD converges locally to a local minimum of (Opt) with arbitrarily high probability.*

*Proof of Theorem 6.7.* Under the stated assumptions, Theorem 5.2 implies that there exists an open neighborhood  $\mathcal{U}$  of  $x^*$  such that  $\mathbb{P}(\lim_{n \rightarrow \infty} X_n = x^* \mid X_0 \in \mathcal{U}) \geq 1 - \delta$ . In turn, this means that there exists some (random)  $n_0$  which is finite with probability at least  $1 - \delta$  and is such that  $\langle \nabla f(x_n), z \rangle \geq \rho \|z\|/2$  for all  $n \geq n_0$  (by the sharpness assumption). Our assertion then follows by conditioning on this event and proceeding as in the proof of Theorem 6.4.  $\square$

We close this section by noting that the convergence of  $X_n$  in a finite number of steps is a unique feature of *lazy* descent schemes with a surjective mirror map. For example, if we consider the *greedy* (or “*eager*”) projected descent scheme

$$X_{n+1} = \Pi(X_n - \gamma_n \nabla F(X_n; \omega_n)), \quad (6.13)$$

it is not possible to obtain a result similar to Theorems 6.4 and 6.7 without further assumptions on the stochasticity affecting (Opt). To see why, assume that  $x^*$  is a sharp minimum of (Opt) and  $X_n = x^*$  for some  $n$ . If the sampled gradient  $\nabla F(X_n; \omega_n)$  attains all directions with positive probability (more precisely, if the

unit vector  $\nabla F(X_n; \omega_n) / \|\nabla F(X_n; \omega_n)\|_*$  is supported on the entire unit sphere  $\mathbb{S}^d$  of  $\mathbb{R}^d$ , there exists some  $\delta > 0$  such that

$$\mathbb{P}(\nabla F(X_n; \omega_n) \notin \text{TC}^*(x^*)) \geq \delta \quad \text{for all } n. \quad (6.14)$$

We thus obtain

$$\mathbb{P}(X_{n+1} \neq x^* \mid X_n = x^*) \geq \delta \quad \text{for all } n, \quad (6.15)$$

implying in turn that  $X_n$  cannot converge to  $x^*$  in a finite number of iterations. We find this property of lazy descent schemes particularly appealing as it ensures very fast convergence in the presence of sharp minima.

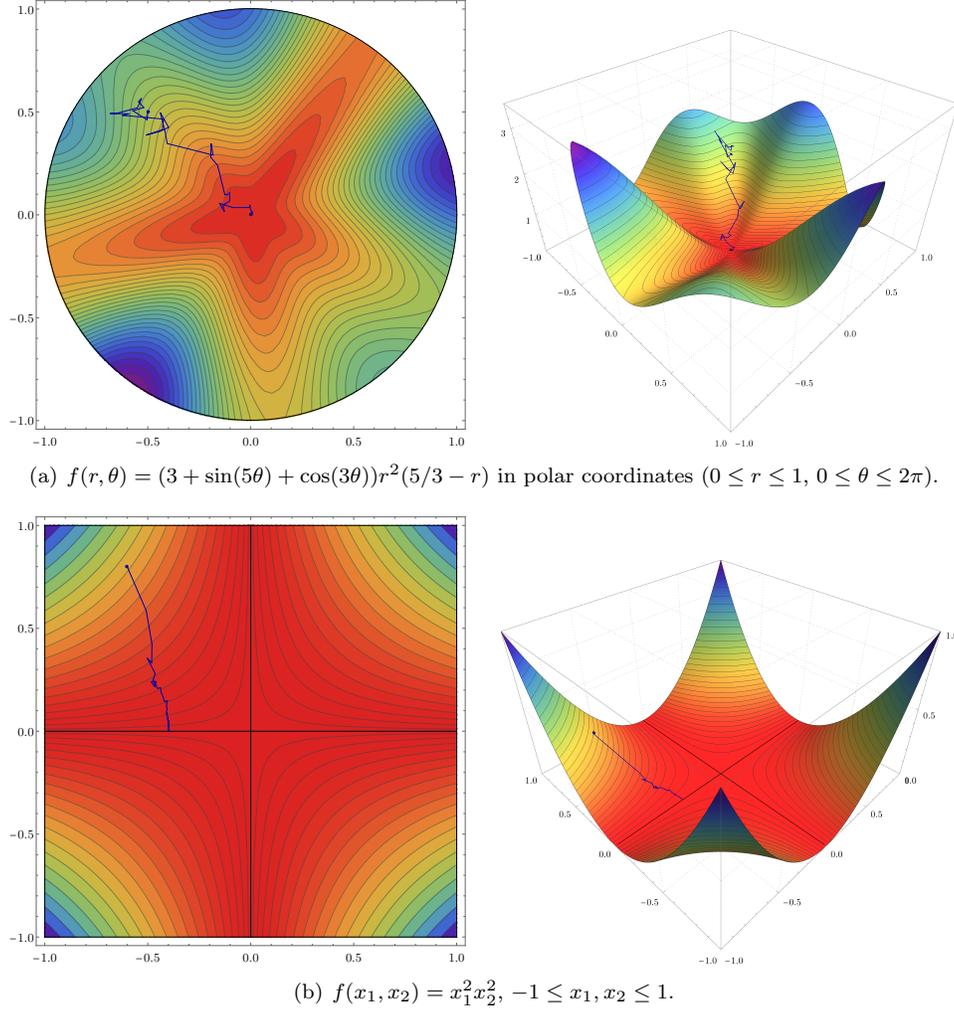
## 7. NUMERICAL EXPERIMENTS

In this section, we validate the theoretical analysis of the previous sections via a series of numerical experiments. As a first illustration of [Theorems 4.1](#) and [5.3](#), we begin by plotting the generated trajectories of (SMD) for two non-convex test functions satisfying the coherence requirements of [Definitions 2.1](#) and [2.3](#) respectively. Referring to [Fig. 2](#) for the detailed expressions, the specific setup is as in [Example 2.2](#) with  $U$  following a standard Gaussian distribution; (SMD) was then run with the Euclidean projector of [Example 2.7](#) and a step-size sequence  $\gamma_n \propto 1/n$ . In both cases, the (randomly) generated trajectories of (SMD) are seen to converge to a minimum point of (Opt), even when the problem's minimum set is non-convex (as in the second example plotted in [Fig. 2](#)).

To go beyond globally coherent problems, we also test the convergence of (SMD) against the widely used Rosenbrock benchmark of [Eq. \(5.1\)](#). This test function admits a unique global minimum point at  $x^* = (1, \dots, 1)$ ; however, this minimum is at the lowest point of a very flat and thin parabolic valley which is notoriously difficult for first-order methods to traverse [\[40\]](#). Because of the parabolic shape of this valley, the problem is not globally coherent (there are rays emanating from  $x^*$  along which  $f$  fails to be nondecreasing) but an easy algebraic calculation in the spirit of [Example 2.6](#) shows that  $x^*$  is locally coherent.

For illustration purposes, we first focus on a low-dimensional example with  $d = 2$  degrees of freedom and algorithmic parameters as in [Fig. 2](#). Despite the lack of global coherence, the simulated trajectories of (SMD) quickly reach the Rosenbrock valley and eventually converge to the minimum of  $f$ ; a typical such trajectory is shown in [Fig. 3](#). Subsequently, in [Fig. 4](#), we run a series of tests on the Rosenbrock function for  $d = 10^3$  and  $d = 10^4$  degrees of freedom. Because the calculation of the gradient becomes increasingly difficult as  $d$  grows large, we take the approach of [Example 2.1](#) and, at each iteration  $n = 1, 2, \dots$  of the algorithm, we randomly pick an integer between 1 and  $d$  and calculate the gradient of  $f_i(x) = 100(x_{i+1} - x_i)^2 + (1 - x_i)^2$ .

In so doing, we obtain the plots of [Fig. 4](#) where, for statistical significance, we report the findings of  $S = 100$  sample realizations. For comparison purposes, we also include in the figure the performance of the ergodic average  $\bar{X}_n$  of  $X_n$  as defined in [Eq. \(4.1\)](#). This sequence is the standard output of mirror descent/dual averaging schemes in convex problems; however, in our non-convex setting, this averaging offers no tangible benefits. This is seen clearly in [Fig. 4](#) where the convergence speed of  $\bar{X}_n$  is considerably slower than that of the algorithm's last iterate.

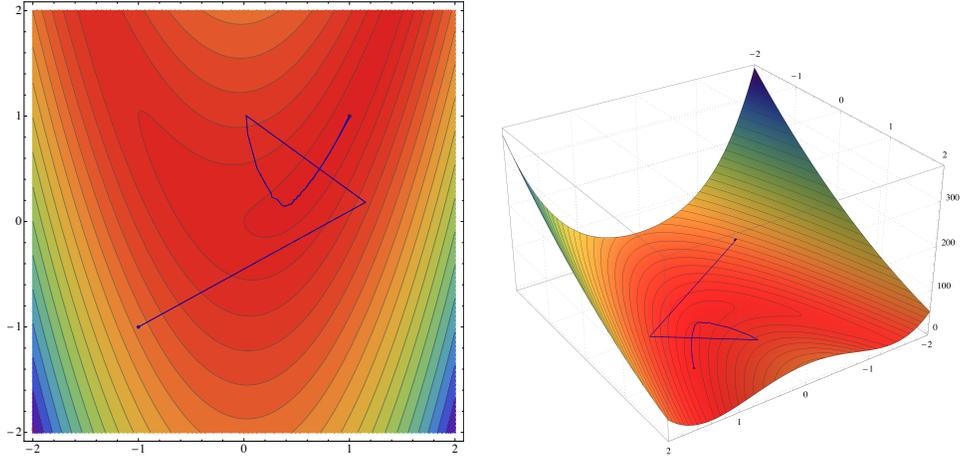


**Figure 2.** Convergence of SMD in a coherent problem with a unique minimizer (top) and a weakly coherent problem with a non-convex minimum set (bottom). In both cases, the minimum set of  $f$  is plotted in solid black.

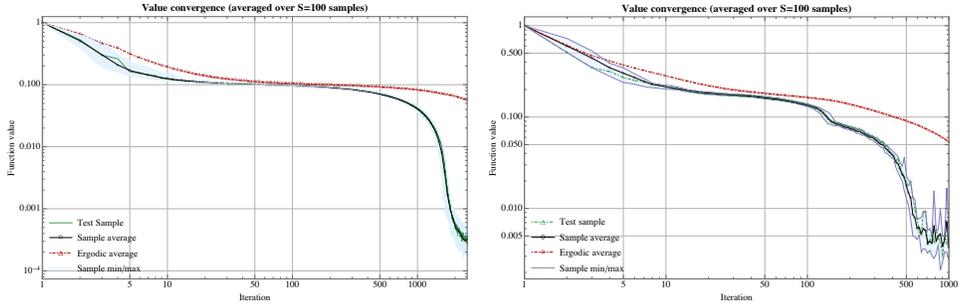
Finally, in Fig. 5, we examine the convergence rate of (SMD) for quadratic objective functions of the form

$$f(x) = \frac{1}{2} \sum_{i,j=1}^d Q_{ij} x_i x_j + \sum_{i=1}^d b_i x_i, \quad (7.1)$$

with  $Q = (Q_{ij})_{i,j=1}^d$  negative-definite (so  $f$  is concave). When  $x$  is constrained to lie on the unit simplex of  $\mathbb{R}^d$ , the minimization of  $f$  is related to the maximum weight clique problem [25]: this problem is well known to be NP-hard, so fast convergence to local minima of  $f$  is essential in order to get reasonable approximate solutions.



**Figure 3.** Convergence of the SMD algorithm in the Rosenbrock test with  $d = 2$  degrees of freedom.

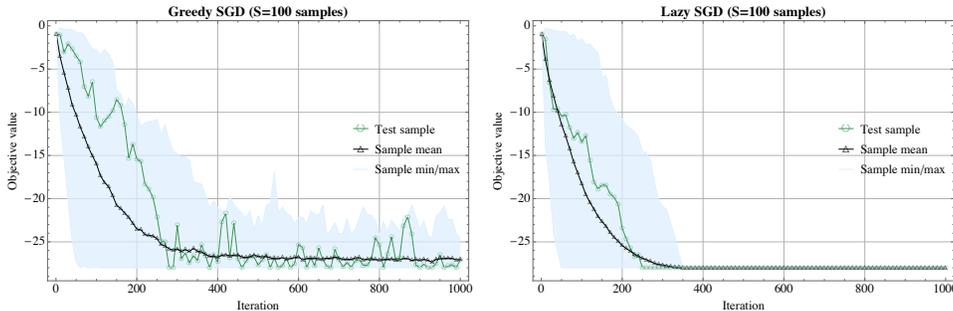


**Figure 4.** Convergence speed of SMD in the Rosenbrock benchmark for  $d = 10^3$  and  $d = 10^4$  degrees of freedom (left and right respectively). The lightly shaded envelope indicates the best and worst realizations of the algorithm over  $S = 100$  sample runs; the corresponding sample mean is represented by a solid black line. For comparison purposes, we also plot the performance of the ergodic average  $\bar{X}_n = \sum_{k=1}^n \gamma_k X_k / \sum_{k=1}^n \gamma_k$  of  $X_n$  (dashed red line). Due to lack of convexity, the ergodic average of  $X_n$  converges at a significantly slower rate.

Using again the stochastic setup of [Example 2.2](#), we ran both [Algorithm 2](#) and its greedy variant [\(6.13\)](#) for a concave quadratic objective of the form [\(7.1\)](#) with  $d = 100$  and randomly drawn  $Q$  and  $b$ . As can be seen in [Fig. 5](#), the lazy variant of SGD converges within a finite number of iterations whereas the greedy variant still oscillates within the allocated time window. This behavior is explained by [Theorem 6.7](#) and the discussion that follows: because the greedy variant essentially “remembers” only the last state, convergence within a finite number of iterations is not possible; by contrast, the dual averaging that takes place in the lazy variant allows  $X_n$  to converge in finite time, despite all the noise.

## 8. DISCUSSION

To conclude, we first note that our analysis can be extended to the study of stochastic variational inequalities with possibly non-monotone operators. The notion



**Figure 5.** Convergence of the lazy and greedy variants of SGD in a quadratic minimization problem of the form (7.1) with  $d = 100$ . The lightly shaded area traces the best and worst realizations of the algorithm over  $S = 100$  sample runs; the corresponding sample mean is drawn as a solid black line. In the dedicated runtime ( $n = 1000$  iterations), the greedy variant still hasn't converged; by contrast, even the worst realization of lazy SGD has converged within approximately 300 iterations.

of a variationally coherent problem still make sense for a variational inequality “as is”, and the Fenchel coupling can also be used to establish almost sure convergence to the solution set of a variational inequality. That said, there are several details that need to be adjusted along the way, so we leave this direction to future work.

Finally, we should also mention that another merit of SMD is that, at least for (strongly) convex optimization problems [13, 36], the algorithm is amenable to asynchronous parallelization. This is an increasingly desirable advantage, especially in the presence of large-scale datasets that are characteristic of “big data” applications requiring the computing power of a massive number of parallel processors. Although we do not tackle this question in this paper, the techniques developed here can potentially be leveraged to provide theoretical guarantees for certain non-convex stochastic programs when SMD is run asynchronously and in parallel.

#### APPENDIX A. ELEMENTS OF MARTINGALE LIMIT THEORY

In this appendix, we state for completeness some basic results from martingale limit theory which we use throughout our paper. The statements are adapted from [18] where we refer the reader for detailed proofs.

We begin with a strong law of large numbers for martingale difference sequences:

**Theorem A.1.** *Let  $M_n = \sum_{k=1}^n d_k$  be a martingale with respect to an underlying stochastic basis  $(\Omega, \mathcal{F}, (\mathcal{F}_n)_{n=1}^{\infty}, \mathbb{P})$  and let  $(\tau_n)_{n=1}^{\infty}$  be a nondecreasing sequence of positive numbers with  $\lim_{n \rightarrow \infty} \tau_n = \infty$ . If  $\sum_{n=1}^{\infty} \tau_n^{-p} \mathbb{E}[|d_n|^p | \mathcal{F}_{n-1}] < \infty$  for some  $p \in [1, 2]$  (a.s.), we have:*

$$\lim_{n \rightarrow \infty} \frac{M_n}{\tau_n} = 0 \quad (\text{a.s.}) \quad (\text{A.1})$$

The second result we use is Doob's martingale convergence theorem:

**Theorem A.2.** *If  $M_n$  is a submartingale that is bounded in  $L^1$  (i.e.  $\sup_n \mathbb{E}[|M_n|] < \infty$ ),  $M_n$  converges almost surely to a random variable  $M$  with  $\mathbb{E}[|M|] < \infty$ .*

The next result is also due to Doob, and is known as Doob's maximal inequality:

**Theorem A.3.** *Let  $M_n$  be a non-negative submartingale and fix some  $\varepsilon > 0$ . Then:*

$$\mathbb{P}(\sup_n M_n \geq \varepsilon) \leq \frac{\mathbb{E}[M_n]}{\varepsilon}. \quad (\text{A.2})$$

Finally, a widely used variant of Doob's maximal inequality is the following:

**Theorem A.4.** *With assumptions and notation as above, we have*

$$\mathbb{P}(\sup_n |M_n| \geq \varepsilon) \leq \frac{\mathbb{E}[M_n^2]}{\varepsilon^2}. \quad (\text{A.3})$$

## APPENDIX B. TECHNICAL PROOFS

In this appendix, we present the proofs that were omitted from the main text. We begin with the core properties of the Fenchel coupling:

*Proof of Lemma 3.2.* To prove the first claim, let  $x = Q(y) = \arg \max_{x' \in \mathcal{X}} \{\langle y, x' \rangle - h(x')\}$ , so  $y \in \partial h(x)$  from standard results in convex analysis [38]. We thus get:

$$F(p, y) = h(p) + h^*(y) - \langle y, p \rangle = h(p) - h(x) - \langle y, p - x \rangle. \quad (\text{B.1})$$

Since  $y \in \partial h(x)$  and  $h$  is  $K$ -strongly convex, we also have

$$h(x) + \tau \langle y, p - x \rangle \leq h(x + \tau(p - x)) \leq (1 - \tau)h(x) + \tau h(p) - \frac{1}{2}K\tau(1 - \tau)\|x - p\|^2 \quad (\text{B.2})$$

for all  $\tau \in [0, 1]$ , thereby leading to the bound

$$\frac{1}{2}K(1 - \tau)\|x - p\|^2 \leq h(p) - h(x) - \langle y, p - x \rangle = F(p, y). \quad (\text{B.3})$$

Our claim then follows by letting  $\tau \rightarrow 0^+$  in (B.3).

For our second claim, we start by citing a well-known duality principle for strongly convex functions [39, Theorem 12.60]: If  $f: \mathbb{R}^d \rightarrow \mathbb{R} \cup \{-\infty, +\infty\}$  is proper, lower semi-continuous and convex, its convex conjugate  $f^*$  is  $\sigma$ -strongly convex if and only if  $f$  is differentiable and satisfies

$$f(x') \leq f(x) + \langle \nabla f(x), x' - x \rangle + \frac{1}{2\sigma}\|x' - x\|^2 \quad (\text{B.4})$$

for all  $x, x' \in \mathbb{R}^d$ . In our case, if we extend  $h$  to all of  $\mathcal{V}$  by setting  $h \equiv +\infty$  outside  $\mathcal{X}$ , we have that  $h$  is  $K$ -strongly convex, lower semi-continuous and proper, so  $(h^*)^* = h$  [39, Theorem 11.1]. It is also easy to see that  $h^*$  is proper, lower semi-continuous and convex (since it is a point-wise maximum of affine functions by definition), so the  $K$ -strong convexity of  $h = (h^*)^*$  implies that  $h^*$  is differentiable and satisfies

$$h^*(y') \leq h^*(y) + \langle y' - y, \nabla h^*(y) \rangle + \frac{1}{2K}\|y' - y\|_*^2 \quad (\text{B.5})$$

$$= h^*(y) + \langle y' - y, Q(y) \rangle + \frac{1}{2K}\|y' - y\|_*^2 \quad (\text{B.6})$$

for all  $y, y' \in \mathcal{Y}$ , where the last equality follows from the fact that  $\nabla h^*(y) = Q(y)$ . Therefore, substituting the preceding inequality in the definition of the Fenchel coupling, we obtain

$$\begin{aligned} F(x, y') &= h(x) + h^*(y') - \langle y', x \rangle \\ &\leq h(x) + h^*(y) + \langle y' - y, \nabla h^*(y) \rangle + \frac{1}{2K}\|y' - y\|_*^2 - \langle y', x \rangle \\ &= F(x, y) + \langle y' - y, Q(y) - x \rangle + \frac{1}{2K}\|y' - y\|_*^2, \end{aligned} \quad (\text{B.7})$$

and our assertion follows.  $\square$

We now turn to the recurrence properties of SMD:

*Proof of Proposition 3.4.* Our proof proceeds step-by-step, as discussed in Section 3:

*Step 1: Martingale properties of  $Y_n$ .* By Assumption 2 and the fact that finiteness of second moments implies finiteness of first moments, we get  $\mathbb{E}[\|F(x; \omega_n)\|_*] < \infty$ . We then claim that  $U_n = \nabla f(X_n) - \nabla F(X_n; \omega_n)$  is an  $L^2$ -bounded martingale difference sequence with respect to the natural filtration of  $\omega_n$ . Indeed, we have:

- (1) Since  $X_n$  is  $\mathcal{F}_{n-1}$ -measurable and  $\omega_n$  is i.i.d., we readily get

$$\begin{aligned} \mathbb{E}[U_n | \mathcal{F}_{n-1}] &= \mathbb{E}[\nabla f(X_n) - \nabla F(X_n; \omega_n) | \mathcal{F}_{n-1}] \\ &= \mathbb{E}[\nabla f(X_n) - \nabla F(X_n; \omega_n) | \omega_1, \dots, \omega_{n-1}] \\ &= \nabla f(X_n) - \nabla f(X_n) \\ &= 0. \end{aligned} \tag{B.8}$$

- (2) Furthermore, by Assumption 2, the  $L^2$  norm of  $U$  satisfies

$$\begin{aligned} \mathbb{E}[\|U_n\|_*^2 | \mathcal{F}_{n-1}] &= \mathbb{E}[\|\nabla f(X_n) - \nabla F(X_n; \omega_n)\|_*^2 | \mathcal{F}_{n-1}] \\ &\leq 2\mathbb{E}[\|\nabla f(X_n)\|_*^2 | \mathcal{F}_{n-1}] + 2\mathbb{E}[\|\nabla F(X_n; \omega_n)\|_*^2 | \mathcal{F}_{n-1}] \\ &\leq 2\|\nabla f(X_n)\|_*^2 + 2V^2 \\ &= 2\|\mathbb{E}[\nabla F(X_n; \omega)]\|_*^2 + 2V^2 \\ &\leq 2\mathbb{E}[\|\nabla F(X_n; \omega)\|_*^2] + 2V^2 \\ &\leq \sigma^2, \end{aligned} \tag{B.9}$$

where we set  $\sigma^2 = 4V^2$ , and we used the dominated convergence theorem to interchange expectation and differentiation in the second line, and Jensen's inequality in the penultimate one.

*Step 2: Recurrence of  $\varepsilon$ -neighborhoods.* We proceed to show that every  $\varepsilon$ -neighborhood  $\mathbb{B}(\mathcal{X}^*, \varepsilon)$  of  $\mathcal{X}^*$  is recurrent under  $X_n$ . To do so, fix some  $\varepsilon > 0$  and assume ad absurdum that  $X_n$  enters  $\mathbb{B}(\mathcal{X}^*, \varepsilon)$  only a finite number of times, so there exists some finite  $n_0$  such that  $\text{dist}(\mathcal{X}^*, X_n) \geq \varepsilon$  for all  $n \geq n_0$ . Since  $\mathcal{X} \setminus \mathbb{B}(\mathcal{X}^*, \varepsilon)$  is compact,  $v(x) = -\nabla f(x)$  is continuous in  $x$ ; furthermore, letting  $x^*$  be such that  $\langle \nabla f(x), x - x^* \rangle = 0$  only if  $x \in \mathcal{X}^*$  (cf. Definition 2.1), it follows that there exists some  $c \equiv c(\varepsilon) > 0$  such that

$$\langle v(x), x - x^* \rangle \leq -c < 0 \quad \text{for all } x \in \mathcal{X} \setminus \mathbb{B}(\mathcal{X}^*, \varepsilon). \tag{B.10}$$

To proceed, let  $R = \max_{x \in \mathcal{X}} \|x\|$  and set  $\beta_n = \gamma_n^2 / (2K)$ . Then, letting  $F_n = F(x^*, Y_n)$  and  $\xi_n = \langle U_n, X_n - x^* \rangle$ , Lemma 3.2 yields

$$\begin{aligned} F_{n+1} &= F(x^*, Y_{n+1}) = F(x^*, Y_n + \gamma_n \hat{v}_n) \\ &\leq F(x^*, Y_n) + \gamma_n \langle v(X_n) + U_n, X_n - x^* \rangle + \beta_n \|\hat{v}_n\|_*^2 \\ &= F_n + \gamma_n \langle v(X_n), X_n - x^* \rangle + \gamma_n \xi_n + \beta_n \|\hat{v}_n\|_*^2 \\ &\leq F_n - \gamma_n c + \gamma_n \xi_n + \beta_n \|\hat{v}_n\|_*^2. \end{aligned} \tag{B.11}$$

Hence, letting  $\tau_n = \sum_{k=n_0}^n \gamma_k$  and telescoping from  $n_0$  to  $n$ , we get

$$\begin{aligned} F_{n+1} &\leq F_{n_0} - c \sum_{k=n_0}^n \gamma_k + \sum_{k=n_0}^n \gamma_k \xi_k + \sum_{k=n_0}^n \beta_k \|\hat{v}_k\|_*^2 \\ &= F_{n_0} - \tau_n \left[ c - \frac{\sum_{k=n_0}^n \gamma_k \xi_k}{\tau_n} \right] + \sum_{k=n_0}^n \beta_k \|\hat{v}_k\|_*^2 \end{aligned} \quad (\text{B.12})$$

We now proceed to bound each term of (B.12). First, by construction, we have

$$\mathbb{E}[\xi_n | \mathcal{F}_{n-1}] = \mathbb{E}[\langle U_n, X_n - x^* \rangle | \mathcal{F}_{n-1}] = \langle \mathbb{E}[U_n | \mathcal{F}_{n-1}], X_n - x^* \rangle = 0, \quad (\text{B.13})$$

where we used the fact that  $X_n$  is  $\mathcal{F}_{n-1}$ -measurable. Moreover, by Young's inequality, we also have

$$|\xi_n| = |\langle U_n, X_n - x^* \rangle| \leq \|U_n\|_* \|X_n - x^*\| \leq 2R \|U_n\|_*, \quad (\text{B.14})$$

where, as before,  $R = \max_{x \in \mathcal{X}} \|x\|$ . Eq. (B.9) then gives

$$\mathbb{E}[\xi_n^2 | \mathcal{F}_{n-1}] \leq (2R)^2 \mathbb{E}[\|U_n\|_*^2 | \mathcal{F}_{n-1}] \leq 4R^2 \sigma^2, \quad (\text{B.15})$$

implying in turn that  $\xi_n$  is an  $L^2$ -bounded martingale difference sequence. It then follows that  $\xi_n$  satisfies the summability condition

$$\sum_{n=n_0}^{\infty} \frac{\mathbb{E}[|\gamma_n \xi_n|^2 | \mathcal{F}_{n-1}]}{\tau_n^2} \leq 4R^2 \sigma^2 \sum_{n=n_0}^{\infty} \frac{\gamma_n^2}{\tau_n^2} < \infty, \quad (\text{B.16})$$

where the last inequality follows from the assumption that  $\gamma_n$  is square-summable. Thus, by the law of large numbers for martingale difference sequences (Theorem A.1), we get

$$\frac{\sum_{k=n_0}^n \gamma_k \xi_k}{\tau_n} \rightarrow 0 \quad \text{as } n \rightarrow \infty \text{ (a.s.)}, \quad (\text{B.17})$$

and, with  $\sum_{k=n_0}^{\infty} \gamma_k = \infty$ , we finally obtain

$$\lim_{n \rightarrow \infty} \tau_n \left[ c - \frac{\sum_{k=n_0}^n \gamma_k \xi_k}{\tau_n} \right] = \infty \quad (\text{a.s.}). \quad (\text{B.18})$$

For the last term of (B.12), let  $S_n = \sum_{k=n_0}^n \beta_k \|\hat{v}_k\|_*^2$ , so  $S_n$  is  $\mathcal{F}_n$ -measurable and nondecreasing. In addition, we have

$$\mathbb{E}[S_n] = \sum_{k=n_0}^n \beta_k \mathbb{E}[\|\hat{v}_k\|_*^2] \leq V^2 \sum_{k=n_0}^n \beta_k < \infty, \quad (\text{B.19})$$

with the last step following from (B.9). This implies that  $S_n$  is an  $L^1$ -bounded submartingale so, by Doob's submartingale convergence theorem (Theorem A.2),  $S_n$  converges almost surely to some random variable  $S_\infty$ , i.e., the last term of (B.12) is bounded. Hence, combining all of the above, we finally obtain

$$\limsup_{n \rightarrow \infty} F_n = -\infty \quad (\text{a.s.}), \quad (\text{B.20})$$

contradicting the positive-definiteness of the Fenchel coupling (cf. Lemma 3.2). We thus conclude that  $X_n$  enters  $\mathbb{B}(\mathcal{X}^*, \varepsilon)$  infinitely many times (a.s.), as claimed.

*Step 3: Recurrence of Fenchel zones.* Using the reciprocity of the Fenchel coupling (Assumption 3), we show below that every Fenchel zone  $\mathbb{B}_F(\mathcal{X}^*, \delta)$  of  $\mathcal{X}^*$  contains

an  $\varepsilon$ -neighborhood of  $\mathcal{X}^*$ . Then, given that  $X_n$  enters  $\mathbb{B}(\mathcal{X}^*, \varepsilon)$  infinitely often (per the previous step), it will also enter  $\mathbb{B}_F(\mathcal{X}^*, \delta)$  infinitely often.

To establish this claim, assume instead that there is no  $\varepsilon$ -ball  $\mathbb{B}(\mathcal{X}^*, \varepsilon)$  contained in  $\mathbb{B}_F(\mathcal{X}^*, \delta)$ . Then, for all  $k > 0$  there exists some  $y_k \in \mathcal{Y}$  such that  $\text{dist}(\mathcal{X}^*, Q(y_k)) = 1/k$  but  $F(\mathcal{X}^*, y_k) \geq \varepsilon$ . This produces a sequence  $(y_k)_{k=1}^\infty$  such that  $\text{dist}(\mathcal{X}^*, Q(y_k)) \rightarrow 0$  but  $F(\mathcal{X}^*, y_k) \geq \varepsilon$ . Since  $\mathcal{X}$  is compact and  $\mathcal{X}^*$  is closed, we can assume without loss of generality that  $Q(y_k) \rightarrow p$  for some  $p \in \mathcal{X}^*$  (at worst, we only need to descend to a subsequence of  $y_k$ ). We thus get

$$\varepsilon \leq F(\mathcal{X}^*, y_k) \leq F(p, y_k). \quad (\text{B.21})$$

However, since  $Q(y_k) \rightarrow p$ , [Assumption 3](#) gives  $F(p, y_k) \rightarrow 0$ , a contradiction. We conclude that  $\mathbb{B}_F(\mathcal{X}^*, \delta)$  contains an  $\varepsilon$ -neighborhood of  $\mathcal{X}^*$ , as required.  $\square$

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