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Slawomir Kurpaska, Apoloniusz Tyszka. The physical impossibility of machine computations on sufficiently large integers inspires an open problem that concerns abstract computable sets  $X\subseteq N$  and cannot be formalized in the set theory ZFC as it refers to our current knowledge on X. 2020. hal-01614087v9

## HAL Id: hal-01614087 https://hal.science/hal-01614087v9

Preprint submitted on 19 Jul 2020 (v9), last revised 20 Sep 2023 (v34)

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The physical impossibility of machine computations on sufficiently large integers inspires an open problem that concerns abstract computable sets  $X \subseteq \mathbb{N}$  and cannot be formalized in the set theory ZFC as it refers to our current knowledge on X

Sławomir Kurpaska, Apoloniusz Tyszka

**Abstract.** Edmund Landau's conjecture states that the set  $\mathcal{P}_{n^2+1}$  of primes of the form  $n^2 + 1$  is infinite. Let  $\beta = (((24!)!)!)!$ , and let  $\Phi$ denote the implication:  $\operatorname{card}(\mathcal{P}_{n^2+1}) < \omega \Rightarrow \mathcal{P}_{n^2+1} \subseteq (-\infty, \beta]$ . We heuristically justify the statement Φ without invoking Landau's conjecture. The set  $X = \{k \in \mathbb{N} : (\beta < k) \Rightarrow (\beta, k) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$  satisfies conditions (1)-(4). (1) There are a large number of elements of X and it is conjectured that X is infinite. (2) No known algorithm decides the finiteness/infiniteness of X. (3) There is a known algorithm that for every  $n \in \mathbb{N}$  decides whether or not  $n \in X$ . (4) There is an explicitly known integer n such that  $\operatorname{card}(X) < \omega \Rightarrow X \subseteq (-\infty, n]$ . (5) There is an explicitly known integer n such that  $card(X) < \omega \Rightarrow X \subseteq (-\infty, n]$  and some known definition of X is much simpler than every known definition of  $X \setminus (-\infty, n]$ . The following problem is open: Is there a set  $X \subseteq \mathbb{N}$  that satisfies conditions (1)-(3) and (5)? The set  $X = \mathcal{P}_{n^2+1}$  satisfies conditions (1)-(3). The set  $X = \{k \in \mathbb{N} : \text{the number of } \}$ digits of k belongs to  $\mathcal{P}_{n^2+1}$  contains  $10^{10^{450}}$  consecutive integers and satisfies conditions (1)-(3). The statement  $\Phi$  implies that both sets X satisfy condition (5).

**Key words and phrases:** complexity of a mathematical definition, computable set  $X \subseteq \mathbb{N}$ , current knowledge on X, explicitly known integer n bounds X from above when X is finite, infiniteness of X remains conjectured, known algorithm for every  $n \in \mathbb{N}$  decides whether or not  $n \in X$ , large number of elements of X, mathematical statement that cannot be formalized in the set theory ZFC, no known algorithm decides the finiteness/infiniteness of X, physical impossibility of machine computations on sufficiently large integers.

### 1. Basic definitions and the goal of the article

Logicism is a programme in the philosophy of mathematics. It is mainly characterized by the contention that mathematics can be reduced to logic, provided that the latter includes set theory, see [3, p. 199].

**Definition 1.** Conditions (1)-(5) concern sets  $X \subseteq \mathbb{N}$ .

- (1) There are a large number of elements of X and it is conjectured that X is infinite.
- (2) No known algorithm decides the finiteness/infiniteness of X.
- (3) There is a known algorithm that for every  $n \in \mathbb{N}$  decides whether or not  $n \in X$ .
- (4) There is an explicitly known integer n such that  $card(X) < \omega \Rightarrow X \subseteq (-\infty, n]$ .
- (5) There is an explicitly known integer n such that  $card(X) < \omega \Rightarrow X \subseteq (-\infty, n]$  and some known definition of X is much simpler than every known definition of  $X \setminus (-\infty, n]$ .

**Definition 2.** We say that an integer n is a threshold number of a set  $X \subseteq \mathbb{N}$ , if  $\operatorname{card}(X) < \omega \Rightarrow X \subseteq (-\infty, n]$ , cf. [8] and [9].

If a set  $X \subseteq \mathbb{N}$  is empty or infinite, then any integer n is a threshold number of X. If a set  $X \subseteq \mathbb{N}$  is non-empty and finite, then the all threshold numbers of X form the set  $[\max(X), \infty) \cap \mathbb{N}$ .

Edmund Landau's conjecture states that the set  $\mathcal{P}_{n^2+1}$  of primes of the form  $n^2 + 1$  is infinite, see [5] and [6].

**Definition 3.** *Let*  $\Phi$  *denote the implication:* 

$$\operatorname{card}(\mathcal{P}_{n^2+1}) < \omega \Rightarrow \mathcal{P}_{n^2+1} \subseteq (-\infty, (((24!)!)!)!]$$

Landau's conjecture implies the statement  $\Phi$ . In Section 4, we heuristically justify the statement  $\Phi$  without invoking Landau's conjecture.

**Statement 1.** There is no explicitly known threshold number of  $\mathcal{P}_{n^2+1}$ . It means that there is no explicitly known integer k such that  $\operatorname{card}(\mathcal{P}_{n^2+1}) < \omega \Rightarrow \mathcal{P}_{n^2+1} \subseteq (-\infty, k]$ .

Proving the statement  $\Phi$  will falsify Statement 1. Statement 1 cannot be formalized in the set theory ZFC because it refers to the current mathematical knowledge. The same is true for Statements 2 and 3 and Open Problem 1 in the next sections. It argues against logicism as Open Problem 1 concerns abstract computable sets  $X \subseteq \mathbb{N}$ .

# 2. The physical impossibility of machine computations on sufficiently large integers inspires Open Problem 1

**Definition 4.** *Let*  $\beta = (((24!)!)!)!$ .

Lemma 1. 
$$\beta \approx 10^{10} 10^{10} 10^{25.16114896940657}$$

*Proof.* We ask Wolfram Alpha at http://wolframalpha.com.

**Statement 2.** The set  $X = \{k \in \mathbb{N} : (\beta < k) \Rightarrow (\beta, k) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$  satisfies conditions (1)-(4).

*Proof.* Condition (1) holds as  $X \supseteq \{0, \dots, \beta\}$  and the set  $\mathcal{P}_{n^2+1}$  is conjecturally infinite. By Lemma 1, due to known physics we are not able to confirm by a direct computation that some element of  $\mathcal{P}_{n^2+1}$  is greater than  $\beta$ , see [2]. Thus condition (2) holds. Condition (3) holds trivially. Since the set

$$\{k \in \mathbb{N} : (\beta < k) \land (\beta, k) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$$

is empty or infinite, the integer  $\beta$  is a threshold number of X. Thus condition (4) holds.

In Statement 2,

$$card(X) < \omega \Rightarrow X \subseteq (-\infty, \beta]$$

and the sets

$$\mathcal{X} = \{k \in \mathbb{N} : (\beta < k) \Rightarrow (\beta, k) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$$

and

$$X \setminus (-\infty, \beta] = \{k \in \mathbb{N} : (\beta < k) \land (\beta, k) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$$

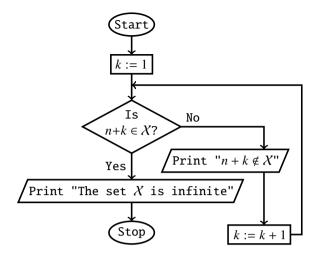
have definitions of similar complexity. The following problem arises:

**Open Problem 1.** *Is there a set*  $X \subseteq \mathbb{N}$  *that satisfies conditions* (1)-(3) *and* (5)?

**Theorem 1.** Assume that for every positive integers b and s, at some future day, machine computations will be possible on every integers from the interval [-b,b] and this will be possible with the speed of s FLOPS. These assumptions contradict the current paradigm of physics, although they alone have no consequences in mathematics formalized in ZFC. We claim that our assumptions alone imply that no set  $X \subseteq \mathbb{N}$  will satisfy conditions (1)–(4) forever.

*Proof.* The proof goes by contradiction. Since conditons (2) – (4) will hold forever, the algorithm in Figure 1 never terminates and sequentially prints the following sentences:

$$n+1 \notin X$$
,  $n+2 \notin X$ ,  $n+3 \notin X$ , ... (T)



**Fig. 1** Algorithm whose execution never terminates if the set X is finite

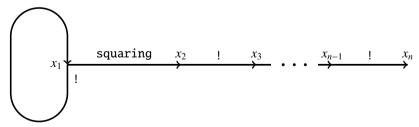
The sentences from the sequence (T) and our assumptions alone imply that for every explicitly given integer m > n, at some future day, a computer will be able to confirm in 1 second or less that  $(n, m] \cap X = \emptyset$ . Thus, at some future day, the empirical evidence will suggest that the set X is finite, contrary to the conjecture in condition (1).

### 3. Number-theoretic statements $\Psi_n$

Let f(1) = 2, f(2) = 4, and let f(n + 1) = f(n)! for every integer  $n \ge 2$ . Let  $\mathcal{U}_1$  denote the system of equations which consists of the equation  $x_1! = x_1$ . For an integer  $n \ge 2$ , let  $\mathcal{U}_n$  denote the following system of equations:

$$\begin{cases} x_1! &= x_1 \\ x_1 \cdot x_1 &= x_2 \\ \forall i \in \{2, \dots, n-1\} x_i! &= x_{i+1} \end{cases}$$

The diagram in Figure 2 illustrates the construction of the system  $\mathcal{U}_n$ .



**Fig. 2** Construction of the system  $\mathcal{U}_n$ 

**Lemma 2.** For every positive integer n, the system  $\mathcal{U}_n$  has exactly two solutions in positive integers, namely (1, ..., 1) and (f(1), ..., f(n)).

Let

$$B_n = \{x_i! = x_k : i, k \in \{1, \dots, n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\}\}$$

For a positive integer n, let  $\Psi_n$  denote the following statement: if a system of equations  $S \subseteq B_n$  has at most finitely many solutions in positive integers  $x_1, \ldots, x_n$ , then each such solution  $(x_1, \ldots, x_n)$  satisfies  $x_1, \ldots, x_n \le f(n)$ . The statement  $\Psi_n$  says that for subsystems of  $B_n$  with a finite number of solutions, the largest known solution is indeed the largest possible. The statements  $\Psi_1$  and  $\Psi_2$  hold trivially. There is no reason to assume the validity of the statement  $\Psi_9$ , cf. Conjecture 1 in Section 4.

**Theorem 2.** For every statement  $\Psi_n$ , the bound f(n) cannot be decreased.

*Proof.* It follows from Lemma 2 because  $\mathcal{U}_n \subseteq B_n$ .

**Theorem 3.** For every integer  $n \ge 2$ , the statement  $\Psi_{n+1}$  implies the statement  $\Psi_n$ .

*Proof.* If a system  $S \subseteq B_n$  has at most finitely many solutions in positive integers  $x_1, \ldots, x_n$ , then for every integer  $i \in \{1, \ldots, n\}$  the system  $S \cup \{x_i! = x_{n+1}\}$  has at most finitely many solutions in positive integers  $x_1, \ldots, x_{n+1}$ . The statement  $\Psi_{n+1}$  implies that  $x_i! = x_{n+1} \le f(n+1) = f(n)!$ . Hence,  $x_i \le f(n)$ .

**Theorem 4.** Every statement  $\Psi_n$  is true with an unknown integer bound that depends on n.

*Proof.* For every positive integer n, the system  $B_n$  has a finite number of subsystems.

### 4. A conjectural solution to Open Problem 1

**Lemma 3.** For every positive integers x and y,  $x! \cdot y = y!$  if and only if

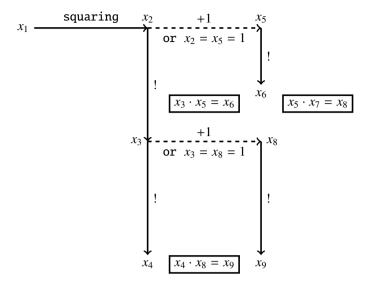
$$(x + 1 = y) \lor (x = y = 1)$$

**Lemma 4.** (Wilson's theorem, [1, p. 89]). For every integer  $x \ge 2$ , x is prime if and only if x divides (x - 1)! + 1.

Let  $\mathcal{A}$  denote the following system of equations:

$$\begin{cases} x_2! &= x_3 \\ x_3! &= x_4 \\ x_5! &= x_6 \\ x_8! &= x_9 \\ x_1 \cdot x_1 &= x_2 \\ x_3 \cdot x_5 &= x_6 \\ x_4 \cdot x_8 &= x_9 \\ x_5 \cdot x_7 &= x_8 \end{cases}$$

Lemma 3 and the diagram in Figure 3 explain the construction of the system  $\mathcal{A}$ .



**Fig. 3** Construction of the system  $\mathcal{A}$ 

**Lemma 5.** For every integer  $x_1 \ge 2$ , the system  $\mathcal{A}$  is solvable in positive integers  $x_2, \ldots, x_9$  if and only if  $x_1^2 + 1$  is prime. In this case, the integers  $x_2, \ldots, x_9$  are uniquely determined by the following equalities:

$$x_{2} = x_{1}^{2}$$

$$x_{3} = (x_{1}^{2})!$$

$$x_{4} = ((x_{1}^{2})!)!$$

$$x_{5} = x_{1}^{2} + 1$$

$$x_{6} = (x_{1}^{2} + 1)!$$

$$x_{7} = \frac{(x_{1}^{2})! + 1}{x_{1}^{2} + 1}$$

$$x_{8} = (x_{1}^{2})! + 1$$

$$x_{9} = ((x_{1}^{2})! + 1)!$$

*Proof.* By Lemma 3, for every integer  $x_1 \ge 2$ , the system  $\mathcal{H}$  is solvable in positive integers  $x_2, \ldots, x_9$  if and only if  $x_1^2 + 1$  divides  $(x_1^2)! + 1$ . Hence, the claim of Lemma 5 follows from Lemma 4.

**Lemma 6.** There are only finitely many tuples  $(x_1, ..., x_9) \in (\mathbb{N} \setminus \{0\})^9$ , which solve the system  $\mathcal{A}$  and satisfy  $x_1 = 1$ . This is true as every such tuple  $(x_1, ..., x_9)$  satisfies  $x_1, ..., x_9 \in \{1, 2\}$ .

*Proof.* The equality  $x_1 = 1$  implies that  $x_2 = x_1^2 = 1$ . Hence, for example,  $x_3 = x_2! = 1$ . Therefore,  $x_8 = x_3 + 1 = 2$  or  $x_8 = 1$ . Consequently,  $x_9 = x_8! \le 2$ .  $\square$ 

**Conjecture 1.** The statement  $\Psi_9$  is true when is restricted to the system  $\mathcal{A}$ .

**Theorem 5.** Conjecture 1 proves the following implication: if there exists an integer  $x_1 \ge 2$  such that  $x_1^2 + 1$  is prime and greater than f(7), then the set  $\mathcal{P}_{n^2+1}$  is infinite.

*Proof.* Suppose that the antecedent holds. By Lemma 5, there exists a unique tuple  $(x_2, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^8$  such that the tuple  $(x_1, x_2, \ldots, x_9)$  solves the system  $\mathcal{A}$ . Since  $x_1^2 + 1 > f(7)$ , we obtain that  $x_1^2 \ge f(7)$ . Hence,  $(x_1^2)! \ge f(7)! = f(8)$ . Consequently,

$$x_9 = ((x_1^2)! + 1)! \ge (f(8) + 1)! > f(8)! = f(9)$$

Conjecture 1 and the inequality  $x_9 > f(9)$  imply that the system  $\mathcal{A}$  has infinitely many solutions  $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$ . According to Lemmas 5 and 6, the set  $\mathcal{P}_{n^2+1}$  is infinite.

**Theorem 6.** Conjecture 1 implies the statement  $\Phi$ .

*Proof.* It follows from Theorem 5 and the equality f(7) = (((24!)!)!)!.

**Theorem 7.** The statement  $\Phi$  implies Conjecture 1.

*Proof.* By Lemmas 5 and 6, if positive integers  $x_1, \ldots, x_9$  solve the system  $\mathcal{A}$ , then

$$(x_1 \ge 2) \land (x_5 = x_1^2 + 1) \land (x_5 \text{ is prime})$$

or  $x_1, \ldots, x_9 \in \{1, 2\}$ . In the first case, Lemma 5 and the statement  $\Phi$  imply that the inequality  $x_5 \leq (((24!)!)!)! = f(7)$  holds when the system  $\mathcal{F}$  has at most finitely many solutions in positive integers  $x_1, \ldots, x_9$ . Hence,  $x_2 = x_5 - 1 < f(7)$  and  $x_3 = x_2! < f(7)! = f(8)$ . Continuing this reasoning in the same manner, we can show that every  $x_i$  does not exceed f(9).

**Definition 5.** Let  $\mathcal{K} = \{k \in \mathbb{N} : \text{the number of digits of } k \text{ belongs to } \mathcal{P}_{n^2+1} \}.$ 

**Lemma 7.**  $card(\mathcal{K}) \ge 9 \cdot 10^9 \cdot 4^{747} \approx 10^{10}450.6930560314272$ .

*Proof.* The following PARI/GP ([4]) command

 $isprime(1+9*4^747, \{flag=2\})$ 

returns %1 = 1. This command performs the APRCL primality test, the best deterministic primality test algorithm ([7, p. 226]). It rigorously shows that the number  $(3 \cdot 2^{747})^2 + 1$  is prime. Since  $9 \cdot 10^9 \cdot 4^{747}$  non-negative integers have  $1 + 9 \cdot 4^{747}$  digits, the desired inequality holds. To establish the approximate equality, we ask Wolfram Alpha about  $9 * (10^{\circ}(9 * 4^{\circ}747))$ .

**Statement 3.** The sets  $X = \mathcal{P}_{n^2+1}$  and  $X = \mathcal{K}$  satisfy conditions (1)–(3). The statement  $\Phi$  implies that both sets X satisfy condition (5).

*Proof.* Since the set  $\mathcal{P}_{n^2+1}$  is conjecturally infinite, Lemma 7 implies condition (1) for both sets  $\mathcal{X}$ . Condition (3) holds trivially for both sets  $\mathcal{X}$ . By Lemma 1, due to known physics we are not able to confirm by a direct computation that some element of  $\mathcal{P}_{n^2+1}$  is greater than  $f(7) = (((24!)!)!)! = \beta$ , see [2]. Thus condition (2) holds for both sets  $\mathcal{X}$ . Suppose that the statement  $\Phi$  holds. This implies two facts:

$$\beta$$
 is a threshold number of  $X = \mathcal{P}_{n^2+1}$  (6)

and

$$\underbrace{9...9}_{\beta \text{ digits}} \text{ is a threshold number of } \mathcal{X} = \mathcal{K}$$
 (7)

Thus condition (4) holds for both sets X. The definition of  $\mathcal{P}_{n^2+1}$  is much simpler than the definition of  $\mathcal{P}_{n^2+1} \setminus (-\infty, \beta]$ . The definition of  $\mathcal{K}$  is much simpler than the definition of  $\mathcal{K} \setminus (-\infty, \underbrace{9...9}_{\beta \text{ digits}}]$ . The last three sentences imply that condition (5)

holds for both sets X.

**Acknowledgment.** Sławomir Kurpaska prepared three diagrams in *TikZ*. Apoloniusz Tyszka wrote the article.

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