

The physical impossibility of machine computations on sufficiently large integers inspires an open problem that concerns abstract computable sets $X \subseteq N$ and cannot be formalized in the set theory ZFC as it refers to our current knowledge on X

Slawomir Kurpaska, Apoloniusz Tyszka

▶ To cite this version:

Slawomir Kurpaska, Apoloniusz Tyszka. The physical impossibility of machine computations on sufficiently large integers inspires an open problem that concerns abstract computable sets $X \subseteq N$ and cannot be formalized in the set theory ZFC as it refers to our current knowledge on X. 2020. hal-01614087v8

HAL Id: hal-01614087 https://hal.science/hal-01614087v8

Preprint submitted on 19 Jun 2020 (v8), last revised 20 Sep 2023 (v34)

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

The physical impossibility of machine computations on sufficiently large integers inspires an open problem that concerns abstract computable sets $X \subseteq \mathbb{N}$ and cannot be formalized in the set theory *ZFC* as it refers to our current knowledge on X

Sławomir Kurpaska, Apoloniusz Tyszka

Abstract. Edmund Landau's conjecture states that the set \mathcal{P}_{n^2+1} of primes of the form $n^2 + 1$ is infinite. Let $\beta = (((24!)!)!)!$, and let Φ denote the implication: $\operatorname{card}(\mathcal{P}_{n^2+1}) < \omega \Rightarrow \mathcal{P}_{n^2+1} \subseteq (-\infty,\beta]$. We heuristically justify the statement Φ without invoking Landau's conjecture. The set $X = \{k \in \mathbb{N} : (\beta < k) \Rightarrow (\beta, k) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$ satisfies conditions (1)-(4). (1) There are a large number of elements of X and it is conjectured that X is infinite. (2) No known algorithm decides the finiteness/infiniteness of X. (3) There is a known algorithm that for every $n \in \mathbb{N}$ decides whether or not $n \in X$. (4) There is an explicitly known integer n such that $\operatorname{card}(X) < \omega \Rightarrow X \subseteq (-\infty, n]$. (5) There is an explicitly known integer n such that $card(X) < \omega \Rightarrow X \subseteq (-\infty, n]$ and some known definition of X is much simpler than every known definition of $X \setminus (-\infty, n]$. The following problem is open: Is there a set $X \subseteq \mathbb{N}$ that satisfies conditions (1)-(3) and (5)? The set $X = \mathcal{P}_{n^2+1}$ satisfies conditions (1)-(3). The set $X = \{k \in \mathbb{N} : \text{the number of } k \in \mathbb{N} \}$ digits of k belongs to \mathcal{P}_{n^2+1} contains $10^{10^{450}}$ consecutive integers and satisfies conditions (1)-(3). The statement Φ implies that both sets X satisfy condition (5).

Key words and phrases: complexity of a mathematical definition, computable set $X \subseteq \mathbb{N}$, current knowledge on X, explicitly known integer n bounds X from above when X is finite, infiniteness of X remains conjectured, known algorithm for every $n \in \mathbb{N}$ decides whether or not $n \in X$, large number of elements of X, mathematical statement that cannot be formalized in the set theory *ZFC*, no known algorithm decides the finiteness/infiniteness of X, physical impossibility of machine computations on sufficiently large integers.

1. Basic definitions and the goal of the article

Logicism is a programme in the philosophy of mathematics. It is mainly characterized by the contention that mathematics can be reduced to logic, provided that the latter includes set theory, see [3, p. 199].

Definition 1. Conditions (1)–(5) concern sets $X \subseteq \mathbb{N}$.

(1) There are a large number of elements of X and it is conjectured that X is infinite.

(2) No known algorithm decides the finiteness/infiniteness of X.

(3) There is a known algorithm that for every $n \in \mathbb{N}$ decides whether or not $n \in X$.

(4) *There is an explicitly known integer n such that* $card(X) < \omega \Rightarrow X \subseteq (-\infty, n]$.

(5) There is an explicitly known integer n such that $card(X) < \omega \Rightarrow X \subseteq (-\infty, n]$ and some known definition of X is much simpler than every known definition of $X \setminus (-\infty, n]$.

Definition 2. We say that an integer n is a threshold number of a set $X \subseteq \mathbb{N}$, if $\operatorname{card}(X) < \omega \Rightarrow X \subseteq (-\infty, n]$, cf. [8] and [9].

If a set $X \subseteq \mathbb{N}$ is empty or infinite, then any integer *n* is a threshold number of *X*. If a set $X \subseteq \mathbb{N}$ is non-empty and finite, then the all threshold numbers of *X* form the set $[\max(X), \infty) \cap \mathbb{N}$.

Edmund Landau's conjecture states that the set \mathcal{P}_{n^2+1} of primes of the form $n^2 + 1$ is infinite, see [5] and [6].

Definition 3. Let Φ denote the implication:

 $\operatorname{card}(\mathcal{P}_{n^2+1}) < \omega \Longrightarrow \mathcal{P}_{n^2+1} \subseteq (-\infty, (((24!)!)!)!)]$

Landau's conjecture implies the statement Φ . In Section 4, we heuristically justify the statement Φ without invoking Landau's conjecture.

Statement 1. There is no explicitly known threshold number of \mathcal{P}_{n^2+1} . It means that there is no explicitly known integer k such that $\operatorname{card}(\mathcal{P}_{n^2+1}) < \omega \Rightarrow \mathcal{P}_{n^2+1} \subseteq (-\infty, k]$.

Proving the statement Φ will falsify Statement 1. Statement 1 cannot be formalized in the set theory *ZFC* because it refers to the current mathematical knowledge. The same is true for Statements 2 and 3 and Open Problem 1 in the next sections. It argues against logicism as Open Problem 1 concerns abstract computable sets $X \subseteq \mathbb{N}$.

2. The physical impossibility of machine computations on sufficiently large integers inspires Open Problem 1

Definition 4. Let $\beta = (((24!)!)!)!$.

Lemma 1.
$$\beta \approx 10^{10} 10^{25.16114896940657}$$

Statement 2. The set $X = \{k \in \mathbb{N} : (\beta < k) \Rightarrow (\beta, k) \cap \mathcal{P}_{n^2+1} \neq \emptyset\}$ satisfies conditions (1)-(4).

Proof. Condition (1) holds as $X \supseteq \{0, ..., \beta\}$ and the set \mathcal{P}_{n^2+1} is conjecturally infinite. By Lemma 1, due to known physics we are not able to confirm by a direct computation that some element of \mathcal{P}_{n^2+1} is greater than β , see [2]. Thus condition (2) holds. Condition (3) holds trivially. Since the set

$$\{k \in \mathbb{N} : (\beta < k) \land (\beta, k) \cap \mathcal{P}_{n^2 + 1} \neq \emptyset\}$$

is empty or infinite, the integer β is a threshold number of X. Thus condition (4) holds.

In Statement 2,

$$\operatorname{card}(X) < \omega \Rightarrow X \subseteq (-\infty, \beta]$$

and the sets

$$\mathcal{X} = \{k \in \mathbb{N} : (\beta < k) \Rightarrow (\beta, k) \cap \mathcal{P}_{n^2 + 1} \neq \emptyset\}$$

and

$$\mathcal{X} \setminus (-\infty, \beta] = \{k \in \mathbb{N} : (\beta < k) \land (\beta, k) \cap \mathcal{P}_{n^2 + 1} \neq \emptyset\}$$

have definitions of similar complexity. The following problem arises:

Open Problem 1. *Is there a set* $X \subseteq \mathbb{N}$ *that satisfies conditions* (1)–(3) *and* (5)?

3. Number-theoretic statements Ψ_n

Let f(1) = 2, f(2) = 4, and let f(n + 1) = f(n)! for every integer $n \ge 2$. Let \mathcal{U}_1 denote the system of equations which consists of the equation $x_1! = x_1$. For an integer $n \ge 2$, let \mathcal{U}_n denote the following system of equations:

$$\begin{cases} x_1! = x_1 \\ x_1 \cdot x_1 = x_2 \\ \forall i \in \{2, \dots, n-1\} x_i! = x_{i+1} \end{cases}$$

The diagram in Figure 1 illustrates the construction of the system \mathcal{U}_n .

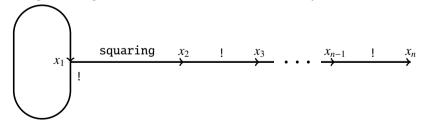


Fig. 1 Construction of the system \mathcal{U}_n

Lemma 2. For every positive integer n, the system \mathcal{U}_n has exactly two solutions in positive integers, namely $(1, \ldots, 1)$ and $(f(1), \ldots, f(n))$.

Let

$$B_n = \{x_i! = x_k : i, k \in \{1, \dots, n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\}\}$$

For a positive integer *n*, let Ψ_n denote the following statement: *if a system of equations* $S \subseteq B_n$ *has at most finitely many solutions in positive integers* x_1, \ldots, x_n , *then each such solution* (x_1, \ldots, x_n) *satisfies* $x_1, \ldots, x_n \leq f(n)$. The statement Ψ_n says that for subsystems of B_n with a finite number of solutions, the largest known solution is indeed the largest possible. The statements Ψ_1 and Ψ_2 hold trivially. There is no reason to assume the validity of the statement Ψ_9 , cf. Conjecture 1 in Section 4.

Theorem 1. For every statement Ψ_n , the bound f(n) cannot be decreased.

Proof. It follows from Lemma 2 because $\mathcal{U}_n \subseteq B_n$.

Theorem 2. For every integer $n \ge 2$, the statement Ψ_{n+1} implies the statement Ψ_n .

Proof. If a system $S \subseteq B_n$ has at most finitely many solutions in positive integers x_1, \ldots, x_n , then for every integer $i \in \{1, \ldots, n\}$ the system $S \cup \{x_i! = x_{n+1}\}$ has at most finitely many solutions in positive integers x_1, \ldots, x_{n+1} . The statement Ψ_{n+1} implies that $x_i! = x_{n+1} \leq f(n+1) = f(n)!$. Hence, $x_i \leq f(n)$.

Theorem 3. Every statement Ψ_n is true with an unknown integer bound that depends on *n*.

Proof. For every positive integer *n*, the system B_n has a finite number of subsystems.

4. A conjectural solution to Open Problem 1

Lemma 3. For every positive integers x and y, $x! \cdot y = y!$ if and only if

$$(x + 1 = y) \lor (x = y = 1)$$

Lemma 4. (Wilson's theorem, [1, p. 89]). For every integer $x \ge 2$, x is prime if and only if x divides (x - 1)! + 1.

Let \mathcal{A} denote the following system of equations:

$$\begin{cases} x_2! = x_3 \\ x_3! = x_4 \\ x_5! = x_6 \\ x_8! = x_9 \\ x_1 \cdot x_1 = x_2 \\ x_3 \cdot x_5 = x_6 \\ x_4 \cdot x_8 = x_9 \\ x_5 \cdot x_7 = x_8 \end{cases}$$

Lemma 3 and the diagram in Figure 2 explain the construction of the system \mathcal{A} .

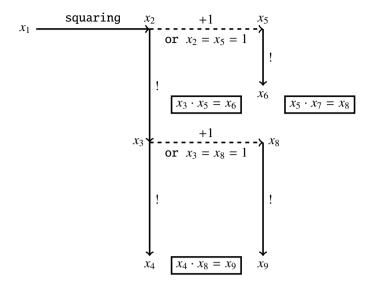


Fig. 2 Construction of the system \mathcal{A}

Lemma 5. For every integer $x_1 \ge 2$, the system \mathcal{A} is solvable in positive integers x_2, \ldots, x_9 if and only if $x_1^2 + 1$ is prime. In this case, the integers x_2, \ldots, x_9 are uniquely determined by the following equalities:

$$\begin{aligned} x_2 &= x_1^2 \\ x_3 &= (x_1^2)! \\ x_4 &= ((x_1^2)!)! \\ x_5 &= x_1^2 + 1 \\ x_6 &= (x_1^2 + 1)! \\ x_7 &= \frac{(x_1^2)! + 1}{x_1^2 + 1} \\ x_8 &= (x_1^2)! + 1 \\ x_9 &= ((x_1^2)! + 1)! \end{aligned}$$

Proof. By Lemma 3, for every integer $x_1 \ge 2$, the system \mathcal{A} is solvable in positive integers x_2, \ldots, x_9 if and only if $x_1^2 + 1$ divides $(x_1^2)! + 1$. Hence, the claim of Lemma 5 follows from Lemma 4.

Lemma 6. There are only finitely many tuples $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$, which solve the system \mathcal{A} and satisfy $x_1 = 1$. This is true as every such tuple (x_1, \ldots, x_9) satisfies $x_1, \ldots, x_9 \in \{1, 2\}$.

Proof. The equality $x_1 = 1$ implies that $x_2 = x_1^2 = 1$. Hence, for example, $x_3 = x_2! = 1$. Therefore, $x_8 = x_3 + 1 = 2$ or $x_8 = 1$. Consequently, $x_9 = x_8! \le 2$. \Box

Conjecture 1. The statement Ψ_9 is true when is restricted to the system \mathcal{A} .

Theorem 4. Conjecture 1 proves the following implication: if there exists an integer $x_1 \ge 2$ such that $x_1^2 + 1$ is prime and greater than f(7), then the set \mathcal{P}_{n^2+1} is infinite.

6

Proof. Suppose that the antecedent holds. By Lemma 5, there exists a unique tuple $(x_2, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^8$ such that the tuple (x_1, x_2, \ldots, x_9) solves the system \mathcal{A} . Since $x_1^2 + 1 > f(7)$, we obtain that $x_1^2 \ge f(7)$. Hence, $(x_1^2)! \ge f(7)! = f(8)$. Consequently,

$$x_9 = ((x_1^2)! + 1)! \ge (f(8) + 1)! > f(8)! = f(9)$$

Conjecture 1 and the inequality $x_9 > f(9)$ imply that the system \mathcal{A} has infinitely many solutions $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$. According to Lemmas 5 and 6, the set \mathcal{P}_{n^2+1} is infinite.

Theorem 5. Conjecture 1 implies the statement Φ .

Proof. It follows from Theorem 4 and the equality f(7) = (((24!)!)!)!.

Theorem 6. The statement Φ implies Conjecture 1.

Proof. By Lemmas 5 and 6, if positive integers x_1, \ldots, x_9 solve the system \mathcal{A} , then $(x_1 \ge 2) \land (x_5 = x_1^2 + 1) \land (x_5 \text{ is prime})$

or $x_1, \ldots, x_9 \in \{1, 2\}$. In the first case, Lemma 5 and the statement Φ imply that the inequality $x_5 \leq (((24!)!)!)! = f(7)$ holds when the system \mathcal{A} has at most finitely many solutions in positive integers x_1, \ldots, x_9 . Hence, $x_2 = x_5 - 1 < f(7)$ and $x_3 = x_2! < f(7)! = f(8)$. Continuing this reasoning in the same manner, we can show that every x_i does not exceed f(9).

Definition 5. Let $\mathcal{K} = \{k \in \mathbb{N} : the number of digits of k belongs to <math>\mathcal{P}_{n^2+1}\}$.

Lemma 7. $\operatorname{card}(\mathcal{K}) \ge 9 \cdot 10^9 \cdot 4^{747} \approx 10^{10} \times 10^{10$

Proof. The following PARI/GP ([4]) command

isprime(1+9*4^747,{flag=2})

returns %1 = 1. This command performs the APRCL primality test, the best deterministic primality test algorithm ([7, p. 226]). It rigorously shows that the number $(3 \cdot 2^{747})^2 + 1$ is prime. Since $9 \cdot 10^{9 \cdot 4^{747}}$ non-negative integers have $1 + 9 \cdot 4^{747}$ digits, the desired inequality holds. To establish the approximate equality, we ask Wolfram Alpha about $9 * (10^{\circ}(9 * 4^{\circ}747))$.

Statement 3. The sets $X = \mathcal{P}_{n^2+1}$ and $X = \mathcal{K}$ satisfy conditions (1)-(3). The statement Φ implies that both sets X satisfy condition (5).

Proof. Since the set \mathcal{P}_{n^2+1} is conjecturally infinite, Lemma 7 implies condition (1) for both sets \mathcal{X} . Condition (3) holds trivially for both sets \mathcal{X} . By Lemma 1, due to known physics we are not able to confirm by a direct computation that some element of \mathcal{P}_{n^2+1} is greater than $f(7) = (((24!)!)!)! = \beta$, see [2]. Thus condition (2) holds for both sets \mathcal{X} . Suppose that the statement Φ holds. This implies two facts:

$$\beta$$
 is a threshold number of $\mathcal{X} = \mathcal{P}_{n^2+1}$ (6)

and

$$\underbrace{9\dots9}_{\beta \text{ digits}} \text{ is a threshold number of } X = \mathcal{K}$$
(7)

Thus condition (4) holds for both sets X. The definition of \mathcal{P}_{n^2+1} is much simpler than the definition of $\mathcal{P}_{n^2+1} \setminus (-\infty, \beta]$. The definition of \mathcal{K} is much simpler than the definition of $\mathcal{K} \setminus (-\infty, \underbrace{9...9}_{\beta \text{ digits}}]$. The last three sentences imply that condition (5)

holds for both sets X.

Acknowledgment. Sławomir Kurpaska prepared two diagrams in *TikZ*. Apoloniusz Tyszka wrote the article.

References

- M. Erickson, A. Vazzana, D. Garth, *Introduction to number theory*, 2nd ed., CRC Press, Boca Raton, FL, 2016.
- [2] S. Lloyd, Ultimate physical limits to computation, Nature 406 (2000), 1047–1054, http: //doi.org/10.1038/35023282.
- [3] W. Marciszewski, *Logic, modern, history of,* in: *Dictionary of logic as applied in the study of language* (ed. W. Marciszewski), pp. 183–200, Springer, Dordrecht, 1981.
- [4] PARI/GP online documentation, http://pari.math.u-bordeaux.fr/dochtml/ html/Arithmetic_functions.html.
- [5] N. J. A. Sloane, The On-Line Encyclopedia of Integer Sequences, A002496, Primes of the form $n^2 + 1$, http://oeis.org/A002496.
- [6] Wolfram MathWorld, Landau's Problems, http://mathworld.wolfram.com/LandausProblems.html.
- [7] S. Y. Yan, Number theory for computing, 2nd ed., Springer, Berlin, 2002.
- [8] A. A. Zenkin, Super-induction method: logical acupuncture of mathematical infinity, Twentieth World Congress of Philosophy, Boston, MA, August 10–15, 1998, http: //www.bu.edu/wcp/Papers/Logi/LogiZenk.htm.
- [9] A. A. Zenkin, Superinduction: new logical method for mathematical proofs with a computer, in: J. Cachro and K. Kijania-Placek (eds.), Volume of Abstracts, 11th International Congress of Logic, Methodology and Philosophy of Science, August 20–26, 1999, Cracow, Poland, p. 94, The Faculty of Philosophy, Jagiellonian University, Cracow, 1999.

Sławomir Kurpaska Technical Faculty Hugo Kołłątaj University Balicka 116B, 30-149 Kraków, Poland E-mail: rtkurpas@cyf-kr.edu.pl

Apoloniusz Tyszka Technical Faculty Hugo Kołłątaj University Balicka 116B, 30-149 Kraków, Poland E-mail: rttyszka@cyf-kr.edu.pl