



HAL
open science

Semantics and Validation of Shapes Schemas for RDF

Iovka Boneva, Jose G Labra Gayo, Eric G Prud 'Hommeaux

► **To cite this version:**

Iovka Boneva, Jose G Labra Gayo, Eric G Prud 'Hommeaux. Semantics and Validation of Shapes Schemas for RDF. ISWC2017 - 16th International semantic web conference, Oct 2017, Vienna, Austria. hal-01590350

HAL Id: hal-01590350

<https://hal.science/hal-01590350v1>

Submitted on 19 Sep 2017

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Semantics and Validation of Shapes Schemas for RDF

Iovka Boneva¹, Jose E. Labra Gayo², and Eric G. Prud'hommeaux³

¹ Univ. Lille - CRISStAL - F-59000 Lille, France. iovka.boneva@univ-lille.fr

² University of Oviedo, Spain. labra@uniovi.es

³ W3C and Stata Center, MIT. eric@w3.org

Abstract. We present a formal semantics and proof of soundness for shapes schemas, an expressive schema language for RDF graphs that is the foundation of Shape Expressions Language 2.0. It can be used to describe the vocabulary and the structure of an RDF graph, and to constrain the admissible properties and values for nodes in that graph. The language defines a typing mechanism called shapes against which nodes of the graph can be checked. It includes an algebraic grouping operator, a choice operator and cardinality constraints for the number of allowed occurrences of a property. Shapes can be combined using Boolean operators, and can use possibly recursive references to other shapes.

We describe the syntax of the language and define its semantics. The semantics is proven to be well-defined for schemas that satisfy a reasonable syntactic restriction, namely stratified use of negation and recursion. We present two algorithms for the validation of an RDF graph against a shapes schema. The first algorithm is a direct implementation of the semantics, whereas the second is a non-trivial improvement. We also briefly give implementation guidelines.

1 Introduction

RDF's distributed graph model encouraged adoption for publication and manipulation of e.g. social and biological data. Coding errors in data stores like DBpedia have largely been handled in a piecemeal fashion with no formal mechanism for detecting or describing schema violations. Extending uptake into environments like medicine, business and banking requires structural validation analogous to what is available in relational or XML schemas.

While OWL ontologies can be used for limited structural validation, they are generally used for formal models of reusable classes and predicates describing objects in some domain. Applications typically consume and produce graphs composed of precise compositions of such ontologies. A company's human resources records may leverage terms from FOAF and Dublin Core, but only certain terms, composed into specific structures, and subject to additional use-specific constraints. We would no more want to impose the constraints of a single human resources application suite on FOAF and Dublin Core than we would want to assert that such applications need to consume all ontologically valid permutations

of FOAF and Dublin Core entities. Further, open-world constraints on OWL ontologies make it impossible to use conventional OWL tools to e.g. detect missing properties. Shape expression schemas (ShEx 1.0) [6, 8] were introduced as a high level language in which it is easy to mix terms from arbitrary ontologies. They provide a schema language in which one can define structural constraints (arc labels, cardinalities, datatypes, etc.) and since version 2.0 (ShEx 2.0)⁴, mix them using Boolean connectives (disjunction, conjunction and negation).

A schema language for any data format has several uses: communicating to humans and machines the form of input/output data; enabling machine-verification of data for production, publication, or consumption; driving query and input interfaces; static analysis of queries. In this, ShEx provides a similar role as relational and XML schemas. A ShEx schema validates nodes in a graph against a schema construct called a *shape*. In XML, validating an element against an XML Schema⁵ type or element or Relax NG⁶ production recursively tests nested elements against constituent rules. In ShEx, validating a node in a graph against a shape recursively tests the nodes which are the object of triples constrained in that shape. An essential difference however is that unlike trees, graphs can have cycles and recursive definitions can yield infinite computation. Moreover, ShEx 2.0 includes a negation operator, and it is well known that mixing recursion with negation can lead to incoherent semantics.

Contributions. In this paper we present *shapes schemas*, a schema language that is the foundation of ShEx 2.0 (Sect. 2). The precise relationship between shapes schemas and ShEx 2.0 is given at the end of Sect. 2. We formally define the semantics of shapes schemas and show that it is sound for schemas that mix recursion and negation in a stratified manner (Sect. 3). We then propose two algorithms for validating an RDF graph node against a shapes schema. Both algorithms are shown to be correct w.r.t. the semantics (Sect. 4). We finally discuss future research directions and conclude (Sect. 5).

Related Work. In [8] we gave semantics for ShEx 1.0. The latter does not use Boolean operators end because of negation, the extension to ShEx 2.0 (and thus to shapes schemas) is non trivial.

Closest to shapes schemas is the SHACL⁷ language both in terms of purpose and expressiveness. SHACL also defines named constraints called shapes to be checked on RDF graph nodes. Unlike ShEx, SHACL is not completely independent from the RDF Schema vocabulary: `rdfs:Classes` play a particular role there as a shape can be required to hold for all the nodes that are instances of some `rdfs:Class`. Therefore validation in SHACL requires partial RDF Schema entailment in order to discover all `rdfs:Classes` of a node. Regarding expressiveness, the main differences between SHACL and shapes schemas are that SHACL

⁴ Shape Expressions Language 2.0. <http://shex.io/shex-semantics/index.html>

⁵ W3C XML Schema. <http://www.w3.org/XML/Schema>

⁶ RELAX NG home page. <http://relaxng.org>

⁷ Shapes Constraint Language (SHACL). <https://www.w3.org/TR/shacl/>

allows to define constraints based on property paths and for comparison of values; SHACL does not have the algebraic operators some-of and each-of and uses Boolean connectives for defining complex shapes; finally SHACL does not define the semantics of recursive shapes.

Ontology languages such as OWL, description logics or RDF Schema are not meant to define (complex) constraints on the data and we do not compare shapes schemas with them. Proposals were made for using OWL with a closed world assumption in order to express integrity constraints [9, 5]. They associate alternative semantics with the existing OWL syntax and can be misleading for users.

Some approaches use SPARQL to express constraints on graphs (SPIN⁸, RDFUnit [3]), or compile a domain specific language into SPARQL queries [2]. SPARQL allows to express complex constraints but does not support recursion. While SPARQL constraints can be validated by standard SPARQL engines, they are harder to write and maintain compared to high-level schemas like ShEx and SHACL.

Description Set Profiles⁹ is a constraint language that uses an RDF vocabulary to define templates and constrain the value and cardinality of properties. It does not have any equivalent of the each-of algebraic operator, and was not designed to be human-readable.

Introductory Example. Let `is:` be a namespace prefix from some ontology, `ex:` be the prefix used in the example schema and instance, and `foaf:` and `xsd:` be the standard FOAF and XSD prefixes, respectively. The schema S_0 is as follows

```

<UserShape>      → foaf:name @<StringValue> ; foaf:mbox @<IRIValue> [0;1]
<ProgShape>      → ex:expertise @<IRIValue> [0;*] ; ex:experience @<ExpValueSet>
<ClientShape>    → ex:clientNbr @<IntValue> | ex:clientAffil @<AnythgShape>
<IssueShape>     → is:reportedBy @<ClientAndUser> ;
                  is:reproducedBy @<ProgShape> [1;5] ;
                  is:relatedTo @<IssueShape> [0;*]
<AnythgShape>    → IRI @<AnythgShape> [0;*]
<ClientAndUser> → DefClient ; IRI – {ex:clientNbr, ex:clientAffil} @<AnythgShape> [0;*]
                  AND DefUser ; IRI – {foaf:name, foaf:mbox} @<AnythgShape> [0;*]
<StringValue>   → xsd:string
<IRIValue>       → IRI
<ExpValueSet>   → {ex:senior, ex:junior}
<IntValue>       → xsd:integer

```

where Def_{Client} is the definition of `<ClientShape>`, and similarly for Def_{User} , and `–` in the definition of `<ClientAndUser>` is the set difference operator. The schema S_0 defines four shapes intended to describe users, programmers, clients and issues, respectively. `<UserShape>` requires that a node has one `foaf:name` property with string value, and an optional `foaf:mbox` that is an IRI. The optional mailbox is specified by the cardinality constraint `[0;1]`. Other cardinality constraints used in S_0 are `[0;*]` for zero or more, and `[1;5]` for one up to five. When no cardinality is given, the default is “exactly one”. A `<ProgShape>` node has zero or more

⁸ SPIN - Modeling Vocabulary. <http://www.w3.org/Submission/spin-modeling/>

⁹ Description Set Profiles: A constraint language for Dublin Core Application Profiles. <http://dublincore.org/documents/dc-dsp/>

ex:expertise properties with values that are IRIs, and one ex:experience property whose value is one among ex:senior and ex:junior. A $\langle \text{ClientShape} \rangle$ has either a ex:clientNbr that is an integer, or a ex:clientAffil(iation) with unconstrained value (i.e. $\langle \text{AnythgShape} \rangle$), but not both. Finally, an issue ($\langle \text{IssueShape} \rangle$) is reported by somebody who is client and user, is reproduced by one to five programmers, and can be related to zero or more issues.

The shapes in \mathbf{S}_0 whose name contains Value specify the set of allowed values for a node. This can be the set of all values of some literal datatype (e.g. string, integer), the set of all nodes of some kind (e.g. IRI), or an explicitly given set (e.g. $\langle \text{ExpValueSet} \rangle$). $\langle \text{AnythgShape} \rangle$ is satisfied by every node. It states that the node can have zero or more outgoing triples whose predicates can be any IRI, and whose objects match $\langle \text{AnythgShape} \rangle$. Finally, $\langle \text{ClientAndUser} \rangle$ uses a conjunction to require that a node has both the client and the user properties. Its definition is a bit technical. The right hand side of the conjunction states that the node must have a foaf:name and an optional foaf:mbox (Def_{User}). Moreover (the ; operator), the node can have any number ($[0;^*]$) of properties that can be any IRI except for foaf:name and foaf:mbox and whose value is unconstrained. The latter is necessary in order to allow the “client” properties required by the left hand side of the conjunction.

Graph \mathbf{G}_0 here after is described by schema \mathbf{S}_0 . Nodes ex:issue1 and ex:issue2 have shape $\langle \text{IssueShape} \rangle$; ex:fatima and ex:emin are $\langle \text{ClientAndUser} \rangle$; ex:ren and ex:noa have shape $\langle \text{ProgShape} \rangle$.

```

ex:issue1
  is:reportedBy ex:fatima ;
  is:reproducedBy ex:ren , ex:noa ;
  is:relatedTo ex:issue2 .
ex:issue2
  is:reportedBy ex:emin ;
  is:reproducedBy ex:ren ;
  is:relatedTo ex:issue1 .
ex:fatima ex:clientNbr 1 ;
  foaf:name "Fatima Smith".
ex:ren ex:expertise ex:semweb ;
  ex:experience ex:senior .
ex:noa ex:experience ex:junior .
ex:emin ex:clientAffil "ABC" ;
  foaf:name "Emin V. Petrov" ;
  foaf:mbox <mailto:evp@example.org> .

```

The RDF Graph Model. As usual, we assume three disjoint sets: IRI a set of IRIs, Lit a set of literals, and Blank a set of blank nodes. An *RDF graph* is a set of triples over $\text{IRI} \cup \text{Blank} \times \text{IRI} \times \text{IRI} \cup \text{Lit} \cup \text{Blank}$. For a triple (s, p, o) in some graph, s is called its subject, p is called its predicate, and o is called its object. We denote $\text{Nodes}(\mathbf{G})$ the set of *nodes* of the graph \mathbf{G} , that is, the elements that appear in a subject or object position in some triple of \mathbf{G} . The *neighbourhood* of node n in graph \mathbf{G} is the set of triples in \mathbf{G} that have n as subject, and is denoted $\text{neigh}_{\mathbf{G}}(n)$ or simply $\text{neigh}(n)$ when \mathbf{G} is clear from the context. We use disjoint union on sets of triples, denoted \uplus : if N, N_1, N_2 are sets of triples, $N = N_1 \uplus N_2$ means that $N_1 \cup N_2 = N$ and $N_1 \cap N_2 = \emptyset$.

2 Shapes Schemas

A shapes schema \mathbf{S} defines a set of named shapes. A shape is a description of the graph structure that can be visited starting from a particular node. It can

<pre> ShExpr ::= ValueDescr NeigDescr ShapeAnd ShapeOr ShapeNot ValueDescr ::= a subset of IRI ∪ Lit ∪ {_b} NeigDescr ::= TExpr ShapeAnd ::= ShExpr 'AND' ShExpr ShapeOr ::= ShExpr 'OR' ShExpr ShapeNot ::= 'NOT' ShExpr </pre>	<pre> TExpr ::= TriplePattern 'EMPTY' SomeOfExpr EachOfExpr RepetExpr TriplePattern ::= PropSet '@'ShapeRef SomeOfExpr ::= TExpr ' ' TExpr EachOfExpr ::= TExpr ';' TExpr RepetExpr ::= TExpr '[' <i>min</i> ';' <i>max</i> ']' PropSet ::= a subset of IRI ShapeRef ::= a shape label in L </pre>
(a) Shape expressions.	(b) Triple expressions.

Fig. 1: The grammar for shape expressions and triple expressions.

talk about the value of the node itself and about its neighbourhood. Shapes can use (Boolean combinations of) other shapes and can be recursive.

Formally, a shapes schema \mathbf{S} is a pair $(\mathbf{L}, \mathbf{def})$, where \mathbf{L} is a set of *shape labels* used as names of shapes and \mathbf{def} is a function that with every shape label associates a shape expression. In examples, we write $L \rightarrow S$ as short for $\mathbf{def}(L) = S$ (for a shape label L and a shape expression S).

Shape Expressions. The grammar for shape expressions is given on Fig. 1a. A *shape expression* (\mathbf{ShExpr}) is a Boolean combination of two atomic components: value description and neighbourhood description. A neighbourhood description ($\mathbf{NeigDescr}$) defines the expected neighbourhood of a node and is given by a triple expression (\mathbf{TExpr} , see below). A value description ($\mathbf{ValueDescr}$) is a set that declares the admissible values for a node. The set can contain IRIs, literals, and the special constant $_b$ to indicate that the node can be a blank node. ShEx 2.0 proposes concrete syntax for different kinds of value description sets (literal datatypes, regex patterns to be matched by IRIs, intervals, etc.). Here we focus on defining the semantics so the concrete syntax for such sets is irrelevant. A $\mathbf{ValueDescr}$ can be an arbitrary set with the unique assumption that it has a finite representation for which membership can be effectively computed.

Triple expressions. Triple expressions describe the expected neighbourhood of a node. They are inspired by regular expressions likewise DTDs and XML Schema for XML. A triple expression will be matched by the neighbourhood of a node in a graph, similarly to type definitions in XML Schema that are matched by the children of some node. The main difference is that the neighbourhood of a node in an RDF graph is a (unordered) set, whereas the children of a node in an XML document form a sequence.

The grammar for triple expressions (\mathbf{TExpr}) is given on Fig. 1b, in which *min* is a natural, and *max* is a natural or the special value $*$. The basic triple expression is a triple pattern and it constrains triples. A triple expression composed of each-of (separated by a $;$), some-of (separated by a $|$) and repetition operators is satisfied if some distribution of the triples in the neighborhood of a node exactly satisfies the expression. Sect. 3.1 defines this and draws the analogy with

regular expressions. In examples, we omit the braces for singleton PropSets, e.g. we write `foaf:name @<StringValue>` instead of `{foaf:name} @<StringValue>`.

Example 1 (Shape expressions, triple expressions). In schema \mathbf{S}_0 from the introductory example, the definitions of the five shapes with name `...Shape...` are triple expressions and collectively make use of all the operators: each-of (`:`), some-of (`()`), repetition. All shapes with name `...Value...` are defined by atomic ValueDescrs. The definition of `<ClientAndUser>` is a ShapeAnd expression. \square

Relationship between Shapes Schemas and ShEx 2.0. Shapes schemas slightly generalizes ShEx 2.0 and thus allows for a more concise definition of syntax and semantics. For readers familiar with ShEx 2.0 we now explain how shapes schemas differ from ShEx 2.0. First, TriplePattern uses a set of properties, whereas the analogous triple constraint in ShEx 2.0 uses a single property. This slight generalization allows to encode the CLOSED and EXTRA constructs of ShEx 2.0. In shapes schemas, triple expressions are always closed (whereas in ShEx 2.0 they are non closed by default) but an expression E can be made non-closed by transforming it into $E;P @\langle\text{AnythgShape}\rangle[0;*]$, where P is the set of all IRIs not mentioned as properties in E , and `<AnythgShape>` is as defined in the introductory example. The EXTRA modifier is encoded in a similar way, using sets of properties in triple patterns and negation. Second, a ValueDescr is an arbitrary set of values that can be IRIs, literals or blank nodes, whereas the analogous node constraint in ShEx 2.0 defines a set of allowed values using a combination of elementary constraints such as XSD datatypes, facets, numerical intervals, node kinds. Using an arbitrary set of values allows to get rid of unnecessary (w.r.t. defining the semantics) details. Third, ShEx 2.0 allows to use shape labels in shape definitions; this is syntactic sugar and is equivalent to replacing the label by its definition. Finally, in shapes schemas we omit inverse properties which would make the proofs longer without representing any additional challenge w.r.t. the semantics.

3 Semantics of Shapes Schemas

A shape defines the structure of a graph when visited starting from a node that has that shape. In this section we give a precise meaning of the following statement

SHAPES_SEM: node n in graph \mathbf{G} has shape (or type) L from schema \mathbf{S}^{10}

To give a sound definition for SHAPES_SEM is not trivial because of the presence of recursion. It also requires to make a design choice that we explain now.

Example 2 (Simple recursive schema). Let schema \mathbf{S}_1 and graph \mathbf{G}_1 be:
¹⁰ “type” is used as synonym of “shape”, esp. in the notion of *typing* to be introduced shortly. The use of “type” must not be confused with `rdf:type` from RDF Schema. Shapes schemas are totally independent from the RDF Schema vocabulary.

$\langle \text{IssueSh} \rangle \rightarrow \text{is:reportedBy } @\langle \text{Str} \rangle ; \text{is:relatedTo } @\langle \text{IssueSh} \rangle [0;^*]$
 $\langle \text{Str} \rangle \rightarrow \text{xsd:string}$

$\langle i1 \rangle \text{ is:reportedBy "Ren" ; is:relatedTo } \langle i2 \rangle .$
 $\langle i2 \rangle \text{ is:reportedBy "Bob" ; is:relatedTo } \langle i1 \rangle .$

Ex. 2 captures the essence of recursion. If $\langle i1 \rangle$ has shape $\langle \text{IssueSh} \rangle$ then $\langle i2 \rangle$ also has shape $\langle \text{IssueSh} \rangle$. If on the other hand $\langle i1 \rangle$ does not have shape $\langle \text{IssueSh} \rangle$, then neither does $\langle i2 \rangle$. This illustrates two important aspects of the semantics of shapes schemas. First, whether a node has some shape cannot be defined independently of the shapes of the other nodes in the graph. The consequence of this apparently simple fact is that we need a global statement about which nodes satisfy which shapes; we call this a typing. A typing must be correct, i.e. coherent with itself. Second, in the above example there is a (design) choice to make. Clearly, there are two acceptable alternatives: either (1) both $\langle i1 \rangle$ and $\langle i2 \rangle$ have shape $\langle \text{IssueSh} \rangle$, or (2) none of them does. Such choice is well known for recursive languages: (1) corresponds to a maximal solution, and (2) to a minimal solution. Both choices can lead to sound semantics. In shapes schemas we choose the maximal solution. This is justified by applications: in the above example we *do want* to consider $\langle i1 \rangle$ as a valid $\langle \text{IssueSh} \rangle$. It would not be the case with semantics based on a minimal solution.

3.1 Typing and Correct Typing

The semantics is based on the notion of *typing*: this is a set of couples that associate a node of an RDF graph with a shape label (a type). In the sequel we consider a graph \mathbf{G} and a schema $\mathbf{S} = (\mathbf{L}, \mathbf{def})$.

Definition 1 (node-type association, typing). A node-type association is a couple (n, L) in $\text{Nodes}(\mathbf{G}) \times \mathbf{L}$. A typing of \mathbf{G} by \mathbf{S} is a set of node-type associations.

Example 3. With \mathbf{S}_1 and \mathbf{G}_1 from Ex. 2, the following are typings

$\text{typing}_1 = \{(\langle i1 \rangle, \langle \text{IssueSh} \rangle), (\langle i2 \rangle, \langle \text{IssueSh} \rangle), ("Ren", \langle \text{Str} \rangle), ("Bob", \langle \text{Str} \rangle)\}$

$\text{typing}_2 = \{("Ren", \langle \text{Str} \rangle), ("Bob", \langle \text{Str} \rangle)\}$

$\text{typing}_3 = \{(\langle i1 \rangle, \langle \text{IssueSh} \rangle), (\langle i2 \rangle, \langle \text{IssueSh} \rangle)\}$

$\text{typing}_4 = \emptyset$.

A typing is correct if, intuitively, it contains an evidence for every node-type association in it. In the above example typing_1 and typing_2 are correct, whereas typing_3 is not correct as it contains e.g. the association $(\langle i1 \rangle, \langle \text{IssueSh} \rangle)$ but does not contain the association $(\text{"Ren"}, \langle \text{Str} \rangle)$ that is required for $\langle i1 \rangle$ to have type $\langle \text{IssueSh} \rangle$. The empty typing (typing_4) is always correct.

Definition 2 (correct typing). Let $\text{typing} \subseteq \text{Nodes}(\mathbf{G}) \times \mathbf{S}$. We say that typing is a correct typing if for any $(n, L) \in \text{typing}$, it holds $\text{typing}, n \vdash \mathbf{def}(L)$, where \vdash is the relation defined on Fig. 2a.

$\text{se-value-descr1} \frac{n \in V}{\text{typing}, n \vdash V}$	$\text{te-tpattern} \frac{\begin{array}{l} N = \{(subj, pred, obj)\} \\ pred \in P \\ (obj, L) \in \text{typing} \end{array}}{\text{typing}, N \models P @L}$
$\text{se-value-descr2} \frac{\begin{array}{l} _b \in V \\ n \in \text{Blank} \end{array}}{\text{typing}, n \vdash V}$	$\text{te-empty} \frac{N = \emptyset}{\text{typing}, N \models \text{EMPTY}}$
$\text{se-neig-descr} \frac{\text{typing}, \text{neigh}(n) \models E}{\text{typing}, n \vdash E}$	$\text{te-some-of1} \frac{\text{typing}, N \models E_1}{\text{typing}, N \models E_1 E_2}$
$\text{se-shape-and} \frac{\begin{array}{l} \text{typing}, n \vdash S_1 \\ \text{typing}, n \vdash S_2 \end{array}}{\text{typing}, n \vdash S_1 \text{ AND } S_2}$	$\text{te-some-of2} \frac{\text{typing}, N \models E_2}{\text{typing}, N \models E_1 E_2}$
$\text{se-shape-or1} \frac{\text{typing}, n \vdash S_1}{\text{typing}, n \vdash S_1 \text{ OR } S_2}$	$\text{te-each-of} \frac{\begin{array}{l} N = N_1 \uplus N_2 \\ \text{typing}, N_1 \models E_1 \\ \text{typing}, N_2 \models E_2 \end{array}}{\text{typing}, N \models E_1; E_2}$
$\text{se-shape-or2} \frac{\text{typing}, n \vdash S_2}{\text{typing}, n \vdash S_1 \text{ OR } S_2}$	$\text{te-repet} \frac{\begin{array}{l} N = N_1 \uplus \dots \uplus N_k \\ \min \leq k \leq \max \\ \text{typing}, N_i \models E \text{ for all } 0 \leq i \leq k \end{array}}{\text{typing}, N \models E[\min; \max]}$
$\text{se-shape-not} \frac{\text{typing}, n \not\vdash S}{\text{typing}, n \vdash \text{NOT } S}$	
(a) Node satisfies a shape expression.	(b) Set of triples matches a triple expression.

Fig. 2: Definitions of the \vdash and \models relations.

Discussion on \vdash . For a shape expression S , the definition of $\text{typing}, n \vdash S$ on Fig. 2a is by recursion on the structure of S . In Rules `se-value-descr`, V is a subset of $\text{IRI} \cup \text{Lit} \cup \{_b\}$ defining a `ValueDescr`. A node n satisfies the value description V if n belongs to the set V , or if n is a blank node and $_b$ is in V . The other base case is Rule `se-neig-descr`, in which E is a `TExpr` representing a neighbourhood description. A node n satisfies the `NeigDescr` E if the neighbourhood of n matches the triple expression E . The matching relation \models is defined on Fig. 2b and discussed below. The remaining four rules are for the Boolean operators. The rules for `AND` and `OR` are as one would expect. Regarding negation, a node satisfies a `ShapeNot` expression if it does not satisfy its sub-expression, as stated by Rule `se-shape-not`. The premise of that rule is $\text{typing}, n \not\vdash S$ and means that (using the inference rules on Fig. 2a) it is impossible to construct a proof for $\text{typing}, n \vdash S$.

Discussion on \models . For a set of triples N , a typing typing and a `TExpr` E , we say that N matches E with typing , and we write $\text{typing}, N \models E$, as defined recursively on the structure of E on Fig. 2b. Note that the \models relation is defined for an arbitrary set of triples N . In practice, N will be (a subset of) the neighbourhood of some node. In the basic Rule `te-tpattern`, $P@L$ is a `TriplePattern` with P a set of IRIs and L a shape label. A singleton set of triples $\{(subj, pred, obj)\}$ matches

the triple pattern if the predicate $pred$ belongs to P and the object has type L in $typing$. The other basic rule is Rule te-empty: an empty set of triples satisfies the EMPTY triple expression.

The remaining rules are about the composed triple expressions. A set of triples matches a **SomeOfExpr** if it matches one of its sub-expressions (Rules te-some-of). The semantics of a **EachOfExpr** is a bit more complex. A set N matches an each-of triple expression $E_1; E_2$ if N is the disjoint union of two sets N_1 and N_2 , and N_1 matches the sub-expressions E_1 , and N_2 matches the sub-expression E_2 . Let us make a parallel between regular expressions and triple expressions. The each-of operator is analogous to concatenation. Recall that a string w matches a regular expression $R_1 \cdot R_2$ (where \cdot is concatenation) whenever w can be “split” into two strings w_1 and w_2 such that their concatenation gives w ($w = w_1 \cdot w_2$), and w_1 matches R_1 , and w_2 matches R_2 . In the case of triple expressions, the set of triples N is “split” into two disjoint sets N_1 and N_2 : disjoint union on sets is analogous to concatenation on words. Following the same analogy, repetition in triple expressions corresponds to Kleene star (the star operator) in regular expressions, with the difference that it allows to express arbitrary intervals for the number of allowed repetitions, whereas Kleene star is always $[0, *]$. So, in Rule te-repet, a set of triples N matches a repetition triple expression $E[min; max]$ if N can be split as the disjoint union of k sets N_1, \dots, N_k such that k is within the interval bound $[min; max]$ and each of these sets matches the sub-expression E . Note that $k = 0$ is possible only when $N = \emptyset$.

The laws of the Boole algebra can be used to put a shape expression in disjunctive normal form in which only atomic sub-expressions **ValueDescr** and **NeigDescr** are negated. From now on we consider only shape expressions in disjunctive normal form. Note also that the each-of and some-of operators are associative and commutative and we use them as operators of arbitrary arity, as e.g. in schema **S**₀ from the introductory example.

3.2 Stratified Negation

Because of the presence of recursion and negation, the notion of correct typing is not sufficient for defining sound semantics of shapes schemas.

Example 4 (Negation and recursion). Let schema **S**₂ and graph **G**₂ below:

$$\begin{array}{ll} \langle L1 \rangle \rightarrow \text{NOT}(\text{ex:p } \langle L2 \rangle) & \langle n1 \rangle \text{ ex:p } \langle n2 \rangle . \\ \langle L2 \rangle \rightarrow \text{NOT}(\text{ex:p } \langle L1 \rangle) & \langle n2 \rangle \text{ ex:p } \langle n1 \rangle . \end{array}$$

These two typings of **G**₂ by **S**₂ are both correct: $typing_5 = \{(\langle n1 \rangle, \langle L1 \rangle)\}$ and $typing_6 = \{(\langle n2 \rangle, \langle L2 \rangle)\}$ □

The two typings in Ex. 4 strongly contradict each other. In order to prove that node $\langle n1 \rangle$ has shape $\langle L1 \rangle$ (in $typing_5$), we need to prove that $\langle n1 \rangle$ does **not** have shape $\langle L2 \rangle$. The latter however does hold in $typing_6$. Such strong contradictions are possible only in presence of negation. In comparison, in Ex. 2

we also have two contradicting typings, but none of them uses in its proof a negative statement that is positive in the other typing.

This problem is well known in logic programming e.g. in Datalog, see Chapter 15 in [1] for an overview. The literature considers several solutions for defining coherent semantics in this case, among which the most popular are negation-as-failure, stratified negation and well-founded semantics. For instance, well-founded semantics would answer *undefined* to the question “does n have shape L ” whenever there exist two proofs that contradict each-other on that fact. We exclude this solution for two reasons: it is not helpful for users, and it might require to compute all possible typings which is costly. We opt for stratification semantics instead. It imposes a syntactic restriction on the use of recursion together with negation, so that schemas as the one on Ex. 4 are not allowed. This is a reasonable restriction because negation in ShEx is expected to be used mainly locally, e.g. to forbid some property in the neighbourhood of a node.

We now define of stratified negation. The *dependency graph* of \mathbf{S} is a graph whose set of nodes is \mathbf{L} , and that has two kinds of edges labelled dep^- and dep^+ defined by (recall that shape expressions in disjunctive normal form):

- There is a *negative dependency* edge $dep^-(L_1, L_2)$ from L_1 to L_2 iff the shape label L_2 appears in $\mathbf{def}(L_1)$ under an occurrence of the NOT operator;
- There is a *positive dependency* edge $dep^+(L_1, L_2)$ from L_1 to L_2 iff the shape label L_2 appears in $\mathbf{def}(L_1)$ but never under an occurrence of NOT.

Definition 3 (schema with stratified negation). *A schema $\mathbf{S} = (\mathbf{L}, \mathbf{def})$ is with stratified negation if there exists a natural number k and a mapping $strat$ from \mathbf{L} to the interval $[1; k]$ such that for all shape labels L_1, L_2 :*

- if $dep^-(L_1, L_2)$, then $strat(L_1) > strat(L_2)$;
- if $dep^+(L_1, L_2)$, then $strat(L_1) \geq strat(L_2)$.

The mapping $strat$ is called a *stratification* of \mathbf{S} . A well known property of stratified negation is that the dependency graph does not have a cycle that goes through a negative dependency edge. This intuitively means that if shape L_1 depends negatively on shape L_2 , then L_2 does not (transitively) depend on L_1 . Positive interdependence is allowed in an unrestricted manner, as in \mathbf{S}_1 from Ex. 2. \mathbf{S}_2 from Ex. 4 is not with stratified negation because $dep^-(\langle L1 \rangle, \langle L2 \rangle)$ and $dep^-(\langle L2 \rangle, \langle L1 \rangle)$

Example 5 (Stratification). Let schema \mathbf{S}_3 below.

$\langle L1 \rangle \rightarrow \text{NOT}(\text{ex:a } @\langle L2 \rangle ; \text{ex:b } @\langle \text{Str} \rangle)$	$\langle L3 \rangle \rightarrow \text{ex:c } @\langle L2 \rangle$
$\langle L2 \rangle \rightarrow \text{ex:c } @\langle L3 \rangle$	$\langle \text{Str} \rangle \rightarrow \text{xsd:string}$

The dependency graph contains the edges $dep^-(\langle L1 \rangle, \langle L2 \rangle)$, $dep^-(\langle L1 \rangle, \langle \text{Str} \rangle)$, $dep^+(\langle L2 \rangle, \langle L3 \rangle)$, $dep^+(\langle L3 \rangle, \langle L2 \rangle)$. The unique loop is around $\langle L2 \rangle$ and $\langle L3 \rangle$ and it goes through positive dependencies only, so the schema is stratified. A stratification should be such that $\langle \text{Str} \rangle$ and $\langle L2 \rangle$ are on stratum strictly

lower than $\langle L1 \rangle$, and $\langle L2 \rangle$ and $\langle L3 \rangle$ are on the same stratum. One possible stratification is $\langle L1 \rangle$ on stratum 2 and the other three shape labels on stratum 1. Another one is $\langle L2 \rangle$ and $\langle L3 \rangle$ on stratum 1, $\langle Str \rangle$ on stratum 2, and $\langle L1 \rangle$ on stratum 3. The latter is called a most refined stratification as none of the stratums can be split.

3.3 Maximal Correct Typing

Recall from Ex. 3 that both $typing_1$ and $typing_4$ are correct. Note that $\langle i1 \rangle$ has shape $\langle IssueSh \rangle$ according to $typing_1$ but not according to $typing_4$. Then what is the correct answer of SHAPES_SEM for $\langle i1 \rangle$ and $\langle IssueSh \rangle$? Does $\langle i1 \rangle$ have shape $\langle IssueSh \rangle$ at the end? This section provides an answer to that question. In one sentence: we trust $typing_1$ because it is greater; actually it is the greatest (maximal) typing. The comparison is based on set inclusion.

The following Lemma 1 establishes that a maximal typing always exists in absence of negation. The proof is based on Lemma 2 in [8] that can be easily extended for the richer schemas we have here.

Lemma 1. *Let \mathbf{S} be a schema that does not use the negation operator NOT. Then for all graphs \mathbf{G} , there exists a correct typing $typing_g$ of \mathbf{G} by \mathbf{S} such that for every $typing'$, if $typing'$ is a correct typing of \mathbf{G} by \mathbf{S} , then $typing' \subseteq typing_g$.*

The typing $typing_g$ can be computed as the union of all correct typings $typing'$.

Let us now define a maximal typing in presence of negation. Let $strat$ be a stratification of \mathbf{S} that has k strata, with $k \geq 1$. For any $1 \leq i \leq k$, the schema \mathbf{S}_i is the restriction of \mathbf{S} that uses only the shape labels whose stratum is less than i . Formally, $\mathbf{S}_i = (\mathbf{L}_i, \mathbf{def}_i)$ with $\mathbf{L}_i = \{L \in \mathbf{L} \mid strat(L) \leq i\}$, and their respective definitions $\mathbf{def}_i(L) = \mathbf{def}(L)$. Remark that if \mathbf{S} is stratified, then \mathbf{S}_1 is negation-free.

For a set of labels $\mathbf{L}_i \subseteq \mathbf{L}$, $typing|_{\mathbf{L}_i} = \{(n, L) \in typing \mid L \in \mathbf{L}_i\}$ is the restriction of $typing$ on the labels from \mathbf{L}_i .

Definition 4 (stratification-maximal correct typing). *Let $\mathbf{S} = (\mathbf{L}, \mathbf{def})$ be a schema, \mathbf{G} be a graph, and $strat$ be a stratification of \mathbf{S} with k stratums (for $k \geq 1$). For any $1 \leq i \leq k$, let $typing_i$ be the typing of \mathbf{G} by \mathbf{S}_i , defined by:*

- $typing_1$ is the maximal correct typing of \mathbf{G} by \mathbf{S}_1 , as defined in Lemma 1;
- for any $1 \leq i < k$, $typing_{i+1}$ is the union of all correct typings $typing'$ of \mathbf{G} by \mathbf{S}_{i+1} s.t. $typing'|_{\mathbf{L}_i} = typing_i$.

The stratification-maximal correct typing of \mathbf{G} by \mathbf{S} with stratification $strat$ is $Typing(\mathbf{G}, \mathbf{S}, strat) = typing_k$.

$Typing(\mathbf{G}, \mathbf{S}, strat)$ from the above definition is indeed a correct typing for \mathbf{G} by \mathbf{S} , as shown in the following proposition that is the core of the proof of soundness for the semantics of shapes schemas.

Proposition 1. *For any schema \mathbf{S} , any stratification $strat$ of \mathbf{S} and any graph \mathbf{G} , $Typing(\mathbf{G}, \mathbf{S}, strat)$ is a correct typing of \mathbf{G} by \mathbf{S} .*

Proof. Goes by induction on the number of stratums. The base case (1 stratum) is Lemma 1. For the induction case and stratum $i + 1$, by induction hypothesis $typing_i$ is correct for \mathbf{G} and \mathbf{S}_i . It is enough to show that if $typing'$ and $typing''$ are two correct typings for \mathbf{G} by \mathbf{S}_{i+1} and $typing'_{\mathbf{L}_i} = typing''_{\mathbf{L}_i} = typing_i$, then their union $typing = typing' \cup typing''$ is correct for \mathbf{G} by \mathbf{S}_{i+1} . Let $(n, L) \in typing$ and suppose that $(n, L) \in typing'$. Because $typing'$ is correct, we have $typing', n \vdash \mathbf{def}(L)$. We will show that (*) the proof for $typing', n \vdash \mathbf{def}(L)$ can be used as a proof for $typing, n \vdash \mathbf{def}(L)$. If $\mathbf{def}(L)$ does not contain a negation of a triple expression, then (*) easily follows from the definition of \vdash .

So suppose $\mathbf{def}(L)$ contains a negation operator on top of the triple expressions E_1, \dots, E_l . That is, (recall that shape expressions are in disjunctive normal form), $\text{NOT } E_j$ is a sub-expression of $\mathbf{def}(L)$ for every $1 \leq j \leq l$. Then the proof for $typing', n \vdash \mathbf{def}(L)$ contains applications of Rule se-shape-not for $\text{NOT } E_j$ that witness that there does not exist a proof for $typing', n \vdash E_j$, for every $1 \leq j \leq l$. We need to show that a proof $typing, n \vdash E_j$ cannot exist. Suppose by contradiction that P is a proof for $typing, n \vdash E_j$, for some $1 \leq j \leq l$. Let \mathbf{L}' be the set of all shape labels that appear in E_j , then P uses only node-type associations with labels from \mathbf{L}' . That is, $typing_{\mathbf{L}'}, n \vdash E_j$ holds. As E_j is negated in $\mathbf{def}(L)$, we have $\mathbf{L}' \subseteq \mathbf{L}_i$, so $typing_{\mathbf{L}_i}, n \vdash E_j$ also holds. But $typing_{\mathbf{L}_i} = typing_i \subseteq typing'$. Contradiction. \square

Lemma 2 below establishes that $Typing(\mathbf{G}, \mathbf{S}, strat)$ does not depend on the stratification being chosen. This allows to define the maximal correct typing (Def. 5) and to give a precise meaning of SHAPES_SEM (Def. 6) which was the objective of this section.

Lemma 2. *Let $\mathbf{S} = (\mathbf{L}, \mathbf{def})$ be a schema and \mathbf{G} be a graph. Let $strat_1$ and $strat_2$ be two stratifications of \mathbf{S} . Then $Typing(\mathbf{G}, \mathbf{S}, strat_1) = Typing(\mathbf{G}, \mathbf{S}, strat_2)$.*

Proof. (Idea) The proof uses a classical technique as e.g. for stratified Datalog. There exists a unique (up to permutation on the numbering of stratums) most refined stratification $strat_{\text{ref}}$ such that for any other stratification $strat'$, each stratum of $strat'$ can be obtained as a union of stratums of $strat_{\text{ref}}$. Then we show that for any stratification $strat'$, $Typing(\mathbf{G}, \mathbf{S}, strat') = Typing(\mathbf{G}, \mathbf{S}, strat_{\text{ref}})$.

Definition 5 (maximal correct typing). *Let $\mathbf{S} = (\mathbf{L}, \mathbf{def})$ be a schema and \mathbf{G} be a graph. The maximal correct typing of \mathbf{G} by \mathbf{S} is denoted $Typing(\mathbf{G}, \mathbf{S})$ and is defined as $Typing(\mathbf{G}, \mathbf{S}, strat)$ for some stratification $strat$ of \mathbf{S} .*

Definition 6 (shapes_sem). *Let $\mathbf{S} = (\mathbf{L}, \mathbf{def})$ be a schema and \mathbf{G} be a graph. We say that node n (of \mathbf{G}) has shape L (from \mathbf{S}) if $(n, L) \in Typing(\mathbf{G}, \mathbf{S})$.*

4 Validation

In Sect. 3 we have given a declarative semantics of the shapes language. We now consider the related computational problem. We are again interested by the SHAPES_SEM statement (as defined in Sect. 3), i.e. checking whether a given node has a given shape.

Input: \mathbf{G} : a graph, $\mathbf{S} = (\mathbf{L}, \mathbf{def})$: a schema, $strat$ a stratification for \mathbf{S} with k strata

Output: $Typing(\mathbf{G}, \mathbf{S})$

```

1   $typing \leftarrow \emptyset$ ;
2  for  $i$  from 1 to  $k$  do
    // Add all node-type associations for stratum  $i$ 
3  foreach  $n$  in  $\mathbf{Nodes}(\mathbf{G})$  do
4      foreach  $L$  in  $\mathbf{L}_i$  do
5          | add  $(n, L)$  to  $typing$ ;
    // Refine w.r.t the types on stratum  $i$ 
6   $changing \leftarrow true$ ;
7  while  $changing$  do
8       $changing \leftarrow false$ ;
9      foreach  $(n, L)$  in  $typing$  s.t.  $L \in \mathbf{L}_i$  do
10         if not  $typing, n \vdash \mathbf{def}(L)$  then
11             | remove  $(n, L)$  from  $typing$ ;
12             |  $changing \leftarrow true$ ;
13 return  $typing$ 

```

Algorithm 1: The algorithm $refine(\mathbf{G}, \mathbf{S}, strat)$.

4.1 Refinement Algorithm

Algorithm 1 computes $Typing(\mathbf{G}, \mathbf{S}, strat)$. The i -th iteration of the loop on line 2 computes $typing_i$ from Def. 4. The algorithm is correct thanks to Lemma 2 from [8] applied to every stratum i . According to that lemma, the maximal typing defined as the union of all correct typings (i.e. $typing_i$) can be computed by iteratively removing unsatisfied node-type associations (done on line 11) until a fixed point is reached (detected when $changing$ remains *false*). The advantage of the *refine* algorithm is that once $Typing(\mathbf{G}, \mathbf{S})$ is computed, testing whether node n has shape L is done with no additional cost by testing whether (n, L) belongs to $Typing(\mathbf{G}, \mathbf{S})$. The drawback is that it considers all node-type associations which is not always necessary, as shown here after.

4.2 Recursive Algorithm

Algorithm 2 allows to check whether node n has shape L without constructing $Typing(\mathbf{G}, \mathbf{S})$. The idea is to visit only a sufficiently large portion of $Typing(\mathbf{G}, \mathbf{S})$.

Example 6 (Motivation of the prove algorithm). Considering schema \mathbf{S}_3 from Ex. 5 and graph \mathbf{G}_3 below:

```

ex:n1 ex:a ex:n2          ex:n2 ex:c ex:n3 .
ex:n1 ex:b 4 .          ex:n3 ex:c ex:n2 .

```

We want to check whether $ex:n1$ has shape $\langle L1 \rangle$. Remark that the neighbor nodes of $ex:n1$ are $ex:n2$ and 4, whereas the shape labels on which the definition of $\langle L1 \rangle$ depends are $\langle L2 \rangle$ and $\langle Str \rangle$. Any correct proof for $typing, ex:n1 \vdash \langle L1 \rangle$ (or for $typing, ex:n1 \not\vdash \langle L1 \rangle$) would have as leaves either applications of Rule *se-value-descr* that do not depend on *typing*, or applications of Rule *te-tpattern* that

Input: n : node in \mathbf{G} , L : label in \mathbf{L} , Hyp : a stack over $\mathbf{Nodes}(G) \times \mathbf{L}$

Output: *true* if n has label L , *false* otherwise

```

1  Hyp = Hyp.push((n, L));
2  Dep =  $\emptyset$ ;
3  foreach ( $n', L'$ ) in  $dep(n, L) \setminus Hyp$  do
4  |   if prove( $n', L', Hyp$ ) then
5  |   |   Dep = Dep  $\cup$   $\{(n', L')\}$ ;
6  result = Dep  $\cup$  Hyp,  $n \vdash \mathbf{def}(L)$  ;
7  Hyp = Hyp.pop();
8  return result ;

```

Algorithm 2: $prove(n, L, Hyp)$. Graph \mathbf{G} and schema \mathbf{S} are global variables.

uses node-type associations (n, L) where n is a neighbor of $ex:n1$ and L' is a label such that $dep^+(\langle L1 \rangle, L')$ or $dep^-(\langle L1 \rangle, L')$.

Assume schema $\mathbf{S} = (\mathbf{L}, \mathbf{def})$ and graph \mathbf{G} . For a shape label L in \mathbf{S} and a node n in \mathbf{G} , we denote $dep(n, L)$ the set of node-type associations (n', L') s.t. n' is a neighbor of n (that is, $(n, p, n') \in neigh(n)$ for some $|R| p$) and L' appears as a shape reference in $\mathbf{def}(L)$. Algorithm 2 uses this easy to show property: $typing, n \models \mathbf{def}(L)$ iff $typing \cap dep(n, L), n \models \mathbf{def}(L)$. In order to check whether n has shape L , Algorithm 2 will (recursively) check whether n' has shape L' for all (n', L') in $dep(n, L)$. The parameter Hyp is a stack of node-type associations that is also seen (on line 3) as the set of node-type associations it contains. Dep is a set of node-type associations.

Example 7 (Execution trace of the prove algorithm). Here is the tree of recursive calls generated during the evaluation of $prove(ex:n1, \langle L1 \rangle, [])$ for graph \mathbf{G}_3 and schema \mathbf{S}_3 , where $[]$ is the empty stack. The returned value is given on the right. $prove(ex:n1, \langle L1 \rangle, [])$ generates four recursive calls that correspond to $dep(ex:n1, \langle L1 \rangle)$. The call for $ex:n3$ and $\langle L3 \rangle$ does not generate any recursive call: $dep(ex:n3, \langle L3 \rangle)$ contains only $(ex:n2, \langle L2 \rangle)$ which is on the stack.

$prove(ex:n1, \langle L1 \rangle, [])$	<i>true</i>
$prove(ex:n2, \langle L2 \rangle, [(ex:n1, \langle L1 \rangle)])$	<i>true</i>
$prove(ex:n3, \langle L3 \rangle, [(ex:n1, \langle L1 \rangle), (ex:n2, \langle L2 \rangle)])$	<i>true</i>
$prove(ex:n2, \langle Str \rangle, [(ex:n1, \langle L1 \rangle)])$	<i>false</i>
$prove(4, \langle L2 \rangle, [(ex:n1, \langle L1 \rangle)])$	<i>false</i>
$prove(4, \langle Str \rangle, [(ex:n1, \langle L1 \rangle)])$	<i>false</i>

The correctness of the *prove* algorithm is stated by the following:

Proposition 2 (Correctness of the prove algorithm). *For any node n and any shape label L , the evaluation of $prove(n, L, [])$ terminates and returns *true* if $(n, L) \in Typing(\mathbf{G}, \mathbf{S})$ and *false* otherwise.*

Proof (Sketch). For termination: the recursion cannot be infinite-breadth as *prove* generates a finite number of recursive calls on line 4. Infinite-depth recursion is also impossible because *Hyp* is a call stack and the condition on line 3 prevents from (recursively) calling *prove* with the same node and label.

The proof of correctness goes by induction on the stratum of L using the most refined stratification $strat$. For every stratum i we show that whenever Hyp contains only node-type associations (n', L') with $strat(L') > i$, and for any L s.t. $strat(L) = i$, $Typing, n \vdash \mathbf{def}(L)$ iff $prove(n, L, Hyp)$ returns true. For the \Rightarrow direction, the main argument is that if $Typing, n \vdash \mathbf{def}(L)$ then also $Typing \cup Hyp, n \vdash \mathbf{def}(L)$. This is not true in general because of negation, but is true if Hyp is on stratum $\geq strat(L)$ as in this case no type in Hyp is negated in $\mathbf{def}(L)$. For the \Leftarrow direction, we need to show that if $prove(n, L, Hyp)$ returns true then $(n, L) \in Typing(\mathbf{G}, \mathbf{S})$. The problematic case is when $prove(n, L, Hyp)$ returns true whereas n does not have label L . Such error necessarily comes from the fact that on line 6 the algorithm used some $(n', L') \in Hyp \setminus Typing(\mathbf{G}, \mathbf{S})$ in the proof for $Dep \cup Hyp, n \vdash \mathbf{def}(L)$. Consequently, $strat(L) = strat(L')$, and because we consider the most refined stratification, it follows that L and L' mutually depend on each other in the dependency graph of \mathbf{G} . Then we need to distinguish two cases. Either all shape labels on stratum i only depend on each others, as for instance $\langle L2 \rangle$ and $\langle L3 \rangle$ from Ex. 5. In that case $prove(n, L, Hyp)$ returns true based only on hypotheses in Hyp , which is correct w.r.t. the semantics based on maximal solution: if nothing outside stratum i allows to disprove that n has label L , then it is indeed the case. The other possibility is that a shape label L' on stratum i depends also on shapes from lower stratum, as $\langle lsueSh \rangle$ from Ex. 2 that depends on $\langle Str \rangle$. Then the test on line 6 of the call of $prove$ with L' will take this dependency into account and return true only if all conditions, including those that depend on the lower stratum, are satisfied. \square

4.3 On Implementation of the Validation Algorithms

Both algorithms use a test for $typing, n \vdash \mathbf{def}(L)$, which non trivial part is the test of the \models relation required in Rule se-neig-descr. The latter is equivalent to checking whether a word (a string) matches a regular expression disregarding the ordering of the letters of the word. Here the word is over the alphabet of triple patterns that occur in the triple expression. In [4] we presented an algorithm for this problem based on regular expression derivatives. In [8] we gave another algorithm for so called deterministic single-occurrence triple expressions. That algorithm can be extended to general expressions, and was used in several of the implementations of ShEx available as open source¹¹.

The $prove$ algorithm was presented in a form that is easier to understand but not optimized. An implementation could reduce considerably the search space of the algorithm by exploring only relevant node-shape associations from $dep(n, L)$. For instance, in Ex. 7 checking 4 against L2 is useless first because 4 is accessible from $ex:n1$ by $ex:b$ whereas $\langle L2 \rangle$ in the schema is accessible from $\langle L1 \rangle$ by $ex:a$.

A more involved version of the $prove$ algorithm could memorize portion of $Typing(\mathbf{G}, \mathbf{S})$ to be reused. This however should be done carefully: one should not memorize all node-shape associations (n, L) for which the algorithm returned true, as some of these can be false positives as discussed in the proof of Prop. 2.

¹¹ A list of the available ShEx implementations can be found on <http://shex.io/>

5 Conclusion

In this paper we introduced shapes schemas that formalize the semantics of ShEx 2.0 and we showed that the semantics of ShEx 2.0 is sound. We also presented two algorithms for validating an RDF graph against a shapes schema.

ShEx and the underlying formalism presented here are still evolving, and there are several promising directions some of which are already being explored: introduce operators for value comparison, use property paths in triple patterns, define an RDF transformation language based on ShEx. We also plan to consider several heuristics and optimizations as the ones discussed in Sect. 4.3 in order to accelerate the validation of shapes schemas. These will be validated on examples. Another open problem is error reporting in ShEx: how to give useful feedback for correcting validation errors. We also plan to explore the exact relationship between shapes schemas and SHACL and establish whether shapes schemas can be encoded in SPARQL extended with recursion as the one defined in [7].

Acknowledgments. This work was partially supported by CPER Nord-Pas de Calais/FEDER DATA Advanced data science and technologies 2015-2020, ANR project DataCert ANR-15-CE39-0009.

References

1. S. Abiteboul, R. Hull, and V. Vianu. *Foundations of Databases*. Addison-Wesley, 1995.
2. P. M. Fischer, G. Lausen, A. Schätzle, and M. Schmidt. RDF constraint checking. In *Proceedings of the Workshops of the EDBT/ICDT 2015 Joint Conference*. CEUR-WS.org, 2015.
3. D. Kontokostas, P. Westphal, S. Auer, S. Hellmann, J. Lehmann, R. Cornelissen, and A. Zaveri. Test-driven evaluation of linked data quality. In *Proceedings of the 23rd International Conference on World Wide Web (WWW '14)*, 2014.
4. J. E. Labra Gayo, E. Prud'hommeaux, I. Boneva, S. Staworko, H. R. Solbrig, and S. Hym. Towards an RDF Validation Language Based on Regular Expression Derivatives. In *Proceedings of the Workshops of the EDBT/ICDT 2015 Joint Conference*. CEUR-WS.org, 2015.
5. B. Motik, I. Horrocks, and U. Sattler. Adding Integrity Constraints to OWL. In *OWL: Experiences and Directions 2007 (OWLED 2007)*, 2007.
6. E. Prud'hommeaux, J. E. Labra Gayo, and H. R. Solbrig. Shape expressions: an RDF validation and transformation language. In *Proceedings of the 10th International Conference on Semantic Systems, SEMANTICS 2014*. ACM, 2014.
7. J. L. Reutter, A. Soto, and D. Vrgoč. Recursion in SPARQL. In *14th International Semantic Web Conference*, 2015.
8. S. Staworko, I. Boneva, J. E. Labra Gayo, S. Hym, E. G. Prud'hommeaux, and H. R. Solbrig. Complexity and Expressiveness of ShEx for RDF. In *18th International Conference on Database Theory (ICDT)*. Schloss Dagstuhl - Leibniz-Zentrum fuer Informatik, 2015.
9. J. Tao, E. Sirin, J. Bao, and D. L. McGuinness. Integrity constraints in OWL. In *Proceedings of the 24th AAAI Conference on Artificial Intelligence*. AAAI, 2010.