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Stochastic Online Shortest Path Routing: The Value of Feedback

M. Sadegh Talebi, Zhenhua Zou, Richard Combes, Alexandre Proutiere, and Mikael Johansson

Abstract—This paper studies online shortest path routing over multi-hop networks. Link costs or delays are time-varying and modeled by independent and identically distributed random processes, whose parameters are initially unknown. The parameters, and hence the optimal path, can only be estimated by routing packets through the network and observing the realized delays. Our aim is to find a routing policy that minimizes the regret (the cumulative difference of expected delay) between the path chosen by the policy and the unknown optimal path. We formulate the problem as a combinatorial bandit optimization problem and consider several scenarios that differ in where routing decisions are made and in the information available when making the decisions. For each scenario, we derive a tight asymptotic lower bound on the regret that has to be satisfied by any online routing policy. These bounds help us to understand the performance improvements we can expect when (i) taking routing decisions at each hop rather than at the source only, and (ii) observing per-link delays rather than end-to-end path delays. In particular, we show that (i) is of no use while (ii) can have a spectacular impact. Three algorithms, with a trade-off between computational complexity and performance, are proposed. The regret upper bounds of these algorithms improve over those of the existing algorithms, and they significantly outperform state-of-the-art algorithms in numerical experiments.

Index Terms—Shortest path routing, online combinatorial optimization, stochastic multi-armed bandits.

I. INTRODUCTION

In most real-world networks, link delays vary stochastically due to unreliable links and random access protocols (e.g. in wireless networks), mobility (e.g. in mobile ad-hoc networks), randomness of demand (e.g. in overlay networks for peer-to-peer applications), etc. In many cases, the associated parameters to links, e.g. the packet transmission success probabilities in wireless sensor networks, are initially unknown and must be estimated by transmitting packets and observing the outcomes. When designing routing policies, we therefore need to address a challenging trade-off between exploration and exploitation: on the one hand, it is important to route packets on new or poorly known links to explore the network and ensure that the optimal path is eventually found; on the other hand, it is critical that the accumulated knowledge on link parameters is exploited so that paths with low expected delays are preferred.

A preliminary version of this work was presented at the 2014 American Control Conference [1].

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When designing practical routing schemes, one is mostly concerned about the finite-time behaviour of the system and it is crucial to design algorithms that quickly learn link parameters so as to efficiently track the optimal path.

The design of such routing policies is often referred to as an online shortest path routing problem in the literature [2]–[6], and is a particular instance of a combinatorial Multi-Armed Bandit (combinatorial MAB) problem as introduced in [7]. In this paper, we study the *stochastic* version of this problem. More precisely, we consider a network, in which the transmission of a packet on a given link is successful with an unknown but fixed probability. A packet is sent on a given link repeatedly until the transmission is successful; the number of time slots to complete the transmission is referred to as the *delay* on this link. We wish to route N packets from a given source to a given destination in a minimum amount of time. A routing policy selects a path to the destination on a packet-by-packet basis. The path selection can be done at the source (source routing), or in the network as the packet progresses towards the destination (hop-by-hop routing). In the case of source routing, some feedback is available when the packet reaches the destination. This feedback can be either the end-to-end delay, or the delays on each link on the path from source to destination. In the MAB literature, the former type of feedback is referred to as *bandit* feedback, whereas the latter is called *semi-bandit* feedback. The routing policy then selects the path for the next packet based on the feedback gathered from previously transmitted packets. In the case of hop-by-hop routing, routing decisions are taken for each transmission, and the packet is sent over a link selected based on all transmission successes and failures observed so far (for the current packet, and all previously sent packets) on the various links.

The performance of a routing policy is assessed through its expected total delay, i.e., the expected time required to send all N packets to the destination. Equivalently, it can be measured through the notion of *regret*, defined as the difference between the expected total delay under the policy considered and the expected total delay of an oracle policy that would be aware of all link parameters, and would hence always send the packets on the best path. Regret conveniently quantifies the loss in performance due to the fact that link parameters are initially unknown and need to be learnt.

In this paper, we first address two fundamental questions: (i) what is the benefit of allowing routing decisions at every node, rather than only at the source? and (ii) what is the added value of feeding back the observed delay for every link that a packet has traversed compared to only observing

the end-to-end delay?¹ To answer these questions, we derive tight regret lower bounds satisfied by any routing policy in the different scenarios, depending on where routing decisions are made and what information is available to the decision-maker when making these decisions. By comparing the different lower bounds, we are able to quantify the value of having semi-bandit feedback rather than bandit feedback, and the improvements that can possibly be achieved by taking routing decisions hop by hop. We then propose routing policies in the semi-bandit feedback setting, and show that these policies outperform state-of-the-art online shortest path routing algorithms. More precisely, our contributions are the following:

1. Regret lower bounds. We derive tight asymptotic (when N grows large) regret lower bounds. The two first bounds concern source routing policies under bandit and semi-bandit feedback, respectively, whereas the third bound is satisfied by any hop-by-hop routing policy. As we shall see later, these bounds are tight in the sense that there exist policies that achieve them. As it turns out, the regret lower bounds for source routing policies with semi-bandit feedback and that for hop-by-hop routing policies are identical, indicating that taking routing decisions hop by hop does not bring any advantage. On the contrary, the regret lower bounds for source routing policies with bandit and semi-bandit feedback can be significantly different, illustrating the importance of having information about per-link delays.

2. Routing policies. In the case of semi-bandit feedback, we propose three online source routing policies, namely GEOCOMBUCB-1, GEOCOMBUCB-2, and KL-SR (KL-based Source-Routing). GEO refers to the fact that the delay on a given link is geometrically distributed, COMB stands for combinatorial, and UCB (Upper Confidence Bound) indicates that these policies are based on the same “optimism in face of uncertainty” principle as the celebrated UCB algorithm designed for classical MAB problems [8]. KL-SR already appears in the conference version of this paper [1]. Here we improve its regret analysis, and show that the latter scales at most as $\mathcal{O}(|E|H\Delta_{\min}^{-1}\theta_{\min}^{-2}\log(N))$,² where H denotes the length (number of links) of the longest path in the network from the source to the destination, θ_{\min} is the success transmission probability of the link with the worst quality, and Δ_{\min} is the minimal gap between the average end-to-end delays of the optimal and of a sub-optimal path (formal definitions of θ_{\min} and Δ_{\min} are provided in Section III-A). We further show that the regret under GEOCOMBUCB-1 and GEOCOMBUCB-2 scales at most as $\mathcal{O}(|E|\sqrt{H}\Delta_{\min}^{-1}\theta_{\min}^{-2}\log(N))$. The tradeoff between computational complexity and performance (regret) of online routing policies is certainly hard to characterize, but our policies provide a first insight into such a trade-off. Furthermore, they exhibit better regret upper bounds than that of the CUCB (Combinatorial UCB) algorithm [9], which is, to our knowledge, the state-of-the-art algorithm for stochastic

¹The effect of different feedback in the adversarial setting was studied in, e.g., [3], [4].

²This improves over the regret upper bound scaling as $\mathcal{O}(\Delta_{\max}|E|H^3\Delta_{\min}^{-1}\theta_{\min}^{-3}\log(N))$ derived in [1], where Δ_{\max} denotes the maximal gap between the average end-to-end delays of the optimal and of a sub-optimal path.

Algorithm	Regret	Complexity
CUCB [9]	$\mathcal{O}\left(\frac{ E H}{\Delta_{\min}\theta_{\min}^3}\log(N)\right)$	$\mathcal{O}(V E)$
GEOCOMBUCB-1	$\mathcal{O}\left(\frac{ E \sqrt{H}}{\Delta_{\min}\theta_{\min}^2}\log(N)\right)$	$\mathcal{O}(\mathcal{P})$
GEOCOMBUCB-2	$\mathcal{O}\left(\frac{ E \sqrt{H}}{\Delta_{\min}\theta_{\min}^2}\log(N)\right)$	$\mathcal{O}(\mathcal{P})$
KL-SR	$\mathcal{O}\left(\frac{ E H}{\Delta_{\min}\theta_{\min}^2}\log(N)\right)$	$\mathcal{O}(V E)$

TABLE I
COMPARISON OF VARIOUS ALGORITHMS FOR SHORTEST PATH ROUTING UNDER SEMI-BANDIT FEEDBACK.

online shortest path routing. Furthermore, we conduct numerical experiments, showing that our routing policies perform significantly better than CUCB. The Thompson Sampling (TS) algorithm of [10] is applicable to the shortest path problem, but its analysis for general topologies is an open problem. While TS performs slightly better than our algorithms on average, its regret sometimes has a large variance according to our experiments. The regret guarantees of various algorithms, and their computational complexity are summarized in Table I.

The remaining of the paper is organized as follows. In Section II we review the literature related to MAB problems and to online shortest path problems. In Section III, we introduce the network model and formulate our online routing problem. Fundamental performance limits (regret lower bounds) are derived in Section IV. We propose online routing algorithms and evaluate their performance in Section V. Finally, Section VI concludes the paper and provides future research directions. All the proofs are presented in the Appendix.

II. RELATED WORK

Stochastic MAB problems have been introduced by Robbins [11]. In the classical setting, in each round, a decision maker pulls an arm from a set of available arms and observes a realization of the corresponding reward, whose distribution is unknown. The performance of a policy is measured through its regret, defined as the difference between its expected total reward and the optimal reward the decision maker could collect if she knew the reward distributions of all arms. The goal is to find an optimal policy with the smallest regret. This classical stochastic MAB problem was solved by Lai and Robbins in their seminal paper [12], where they derived the asymptotic (when the time horizon is large) lower bound of regret satisfied by any algorithm, and proposed an optimal algorithm that matches the lower bound.

Online shortest path routing problems fall into the class of combinatorial MAB problems. In these MAB problems, arms are subsets of a set of basic actions (in routing problems, a basic action corresponds to a link), and most existing studies concern the adversarial setting where the successive rewards of each arm are arbitrary, see e.g. [7], [13]–[15] for algorithms for generic combinatorial problems, and [2], [4] for efficient algorithms for routing problems. Stochastic combinatorial MAB problems have received little attention so far. Usually they are investigated in the semi-bandit feedback

setting [9], [16]–[18]. Some papers deal with problems where the set of arms exhibits very specific structures, such as fixed-size sets [19], matroid [20], and permutations [21].

In the case of online shortest path routing problems, as a particular instance of a combinatorial MAB, one could think of modeling each path as an arm, and applying sequential arm selection policies as if arms would yield independent rewards. Such policies would have a regret scaling as $|\mathcal{P}| \log(N)$ where $|\mathcal{P}|$ denotes the number of possible paths from the source to the destination. However, since $|\mathcal{P}|$ grows exponentially with the length H of the paths, treating paths as independent arms would lead to a prohibitive regret. In contrast to classical MAB in [12] where the random rewards from various arms are *independent*, in online routing problems, the end-to-end delays (i.e., the rewards) of the various paths are inherently correlated, since paths may share the same links. It may then be crucial to exploit these correlations, i.e., the structure of the problem, to design efficient routing algorithms which in turn may have a regret scaling as $C \log(N)$ where C is much smaller than $|\mathcal{P}|$.

Next we summarize existing results for generic stochastic combinatorial bandits that could be applied to online shortest path routing. In [9], the authors present CUCB, an algorithm for generic stochastic combinatorial MAB problems under semi-bandit feedback. When applied to the online routing problem, the best regret upper bound for CUCB presented in [9] scales as $\mathcal{O}(\frac{|E|H}{\Delta_{\min} \theta_{\min}^3} \log(N))$ (see Appendix J for details).

This upper bound constitutes the best existing result for our problem, where the delay on each link is geometrically distributed. It is important to note that most proposed algorithms for combinatorial bandits [16]–[18] deal with bounded rewards, i.e., here bounded delays, and are not applicable to geometrically distributed delays. In [17], the authors consider the case where the rewards of basic actions (here links) can be arbitrarily correlated and bounded, and show that the regret under CUCB is $\mathcal{O}(\frac{|E|H}{\Delta_{\min}} \log(N))$. They also prove that this regret scaling has order-optimal regret in terms of $|E|$ and H^3 . In other words, the dependence of their regret upper bound on $|E|$ and H cannot be improved in general. This order-optimality does not contradict our regret upper bound (scaling as $\mathcal{O}(\frac{|E|\sqrt{H}}{\Delta_{\min}} \log(N))$), because [17] considers possibly dependent delays across links. Interestingly, to prove that a regret of $\mathcal{O}(\frac{|E|H}{\Delta_{\min}} \log(N))$ cannot be beaten, they artificially create an instance of the problem where the rewards of the basic actions of the same arm are identical. In other words, they consider a classical bandit problem where the rewards of the various arms are either 0 or equal to H . This clearly highlights the fact that the approach of [17] cannot be directly applied to our routing problem where delays are unbounded. For bounded rewards, the results of [17] have been recently improved in [18] when the rewards are *independent* across basic actions (links). There, the authors propose an algorithm whose regret scales at most

³ A policy π is order-optimal in terms of $|E|$ and H , if it satisfies the following: for all problem instances, $R^\pi(N) = \mathcal{O}(C_1 g(|E|, H) \log(N))$ with C_1 independent of $|E|$, H , and N , and there exists a problem instance and a constant $C_2 > 0$, independent of $|E|$, H , and N , such that $\liminf_{N \rightarrow \infty} R^\pi(N) / \log(N) \geq C_2 g(|E|, H)$ for all uniformly good algorithm π' .

as $\mathcal{O}(\frac{|E|\sqrt{H}}{\Delta_{\min}} \log(N))$. Wen et al. [22] study combinatorial problems under semi-bandit feedback and provide algorithms with $\mathcal{O}(\sqrt{N})$ regret. Gopalan et al. [10] study TS [23] for learning problems with complex arms and provide implicit regret upper bounds with $\mathcal{O}(\log(N))$ regret.

Stochastic online shortest path routing problems have been addressed in [5], [24], [25]. Liu and Zhao [24] consider routing with bandit (end-to-end) feedback and propose a forced-exploration algorithm with $\mathcal{O}(|E|^3 H \log(N))$ regret in which a random barycentric spanner⁴ path is chosen for exploration. He et al. [5] consider routing under semi-bandit feedback, where the source chooses a path for routing and a possibly different path for probing. Our model coincides with the coupled probing/routing case in their paper, for which they derive an asymptotic lower bound on the regret growing logarithmically with time. As we shall see later, their lower bound is not tight.

Finally, it is worth noting that the papers cited above considered source-routing only. To the best of our knowledge, this paper is the first to consider online routing problems with hop-by-hop decisions. Such a problem can be formulated as a classical Markov Decision Process (MDP), in which the states are the packet locations and the actions are the outgoing links of each node. However, most studies consider MDP problems under stricter assumptions than ours and/or targeted different performance measures. Burnetas and Katehakis [26] derive the asymptotic lower bound on the regret and propose an optimal index policy. Their result can be applied only to the so-called ergodic MDP [27], where the induced Markov chain by any policy is irreducible and consists of a single recurrent class. In hop-by-hop routing, however, the policy that routes packets on a fixed path results in a Markov chain with reducible states that are not in the chosen path. [28], [29] study general MDPs and present algorithms with finite-time regret upper bounds scaling as $\mathcal{O}(\log(T))$. Nevertheless, these algorithms perform badly when applied to hop-by-hop routing due to loose confidence intervals. [28] also presents non-asymptotic, but problem independent (minimax) regret lower bounds scaling as $\Omega(\sqrt{T})$. This latter bound does not contradict our problem-dependent lower bounds that grow logarithmically.

III. ONLINE SHORTEST PATH ROUTING PROBLEMS

A. Network Model

The network is modeled as a directed graph $G = (V, E)$ where V is the set of nodes and E is the set of links. Each link $i \in E$ may, for example, represent an unreliable wireless link. Without loss of generality, we assume that time is slotted and that one slot corresponds to the time to send a packet over a single link. At time t , $X_i(t)$ is a binary random variable indicating whether a transmission on link i at time t is successful. $(X_i(t))_{t \geq 1}$ is a sequence of i.i.d. Bernoulli variables with initially unknown mean θ_i . Hence if a packet is sent on link i repeatedly until the transmission

⁴A barycentric spanner is a set of paths from which the delay of all other paths can be computed as its linear combination with coefficients in $[-1, 1]$ [2].

Policy Set	Routing Type	Feedback
Π_1	Source-routing	Bandit
Π_2	Source-routing	Semi-bandit
Π_3	Hop-by-hop	Semi-bandit

TABLE II
VARIOUS POLICY SETS FOR ONLINE SHORTEST PATH ROUTING.

is successful, the time to complete the transmission (referred to as the delay on link i) is geometrically distributed with mean $1/\theta_i$. Let $\theta = (\theta_i, i \in E)$ be the vector representing the packet successful transmission probabilities on the various links. We consider a single source-destination pair $(s, d) \in V^2$, and denote by $\mathcal{P} \subseteq \{0, 1\}^{|E|}$ the set of loop-free paths from s to d in G , where each path $p \in \mathcal{P}$ is a $|E|$ -dimensional binary vector; for any $i \in E$, $p_i = 1$ if and only if i belongs to p . Let H denote the maximum length of the paths in \mathcal{P} , i.e., $H = \max_{p \in \mathcal{P}} \sum_{i \in E} p_i$. For brevity, in what follows, for any binary vector z , we write $i \in z$ to denote $z_i = 1$. Moreover, we use the convention that $z^{-1} = (z_i^{-1})_i$.

For any path p , $D_\theta(p) = \sum_{i \in p} \frac{1}{\theta_i}$ is the average packet delay through path p given link success rates θ . The path with minimal delay is: $p^* \in \arg \min_{p \in \mathcal{P}} D_\theta(p)$. Moreover, for any path $p \in \mathcal{P}$, we define $\Delta_p = D_\theta(p) - D_\theta(p^*) = (p - p^*)^\top \theta^{-1}$. Let $\Delta_{\min} = \min_{\Delta_p \neq 0} \Delta_p$. We let $\theta_{\min} = \min_{i \in E} \theta_i$ and assume that $\theta_{\min} > 0$. Finally define $D^* = D_\theta(p^*)$ and $D^+ = \max_{p \in \mathcal{P}} D_\theta(p)$ the delays of the shortest and longest paths, respectively.

The analysis presented in this paper can be easily extended to more general link models, provided that the (single-link) delay distributions are taken within one-parameter exponential families of distributions.

B. Online Routing Policies and Feedback

We assume that the source is fully backlogged (i.e., it always has packets to send), and that the parameter θ is initially unknown. Packets are sent successively from s to d over various paths, and the outcome of each packet transmission is used to estimate θ , and in turn to learn the path p^* with the minimum average delay. After a packet is sent, we assume that the source gathers feedback from the network (essentially per-link or end-to-end delays) before sending the next packet.

Our objective is to design and analyze online routing policies, i.e., policies that take routing decisions based on the feedback received for the packets previously sent.

We consider and compare three different types of online routing policies, depending (i) on where routing decisions are taken (at the source or at each node), and (ii) on the received feedback (per-link or end-to-end path delay). Table II lists different policy sets for the three types of online routing policies considered.

- **Policy Set Π_1** : The path used by a packet is determined at the source based on the observed end-to-end delays for previous packets. More precisely, for the n -th packet, let $p^\pi(n)$ be the path selected under policy π , and let $D^\pi(n)$ denote the corresponding end-to-end delay. Then $p^\pi(n)$ depends on $p^\pi(1), \dots, p^\pi(n-1), D^\pi(1), \dots, D^\pi(n-1)$.

- **Policy Set Π_2** : The path used by a packet is determined at the source based on the observed per-link delays for previous packets. In other words, under policy π , $p^\pi(n)$ depends on $p^\pi(1), \dots, p^\pi(n-1), (d_i^\pi(1), i \in p^\pi(1)), \dots, (d_i^\pi(n-1), i \in p^\pi(n-1))$, where $d_i^\pi(k)$ is the delay experienced on link i for the k -th packet (if this packet uses link i at all).
- **Policy Set Π_3** : Routing decisions are taken at each node in an adaptive manner. At a given time t , the packet is sent over a link selected based on all successes and failures observed on the various links before time t .

In the case of source-routing policies (in $\Pi_1 \cup \Pi_2$), if a transmission on a given link fails, the packet is retransmitted on the same link until it is successfully received (per-link delays are geometric random variables). On the contrary, in the case of hop-by-hop routing policies (in Π_3), the routing decisions at a given node can be adapted to the observed failures on a given link. For example, if transmission attempts on a given link failed, one may well decide to switch link and select a different next-hop node.

C. Performance Metrics and Objectives

1) *Regret*: Under any reasonably smart routing policy, the parameter θ will eventually be estimated accurately and the minimum delay path will be discovered with high probability after sending a large number of packets. Hence, to quantify the performance of a routing policy, we examine its transient behavior. More precisely, we use the notion of *regret*, a performance metric often used in MAB literature [12]. The regret $R^\pi(N)$ of policy π up to the N -th packet is the expected difference of delays for the first N packets under π and under the policy that always selects the best path p^* for transmission:

$$R^\pi(N) := \mathbb{E} \left[\sum_{n=1}^N D^\pi(n) \right] - N D_\theta(p^*),$$

where $D^\pi(n)$ denotes the end-to-end delay of the n -th packet under policy π and the expectation $\mathbb{E}[\cdot]$ is taken with respect to the random transmission outcomes and possible randomization in the policy π . The regret quantifies the performance loss due to the need to explore sub-optimal paths to learn the path with minimum delay.

2) *Objectives*: The goal is to design online routing policies in Π_1, Π_2 , and Π_3 that minimize regret over the first N packets. As it turns out, there are policies in any Π_j , $j = 1, 2, 3$, whose regrets scale as $\mathcal{O}(\log(N))$ when N grows large, and no policy can have a regret scaling as $o(\log(N))$.

Our objective is to derive, for each $j = 1, 2, 3$, an asymptotic regret lower bound $c_j(\theta) \log(N)$ for policies in Π_j , and then propose simple policies whose regret upper bounds asymptotically approach that of the *optimal* algorithm, i.e., an algorithm whose regret matches the lower bound in Π_j . As we shall discuss later, there exists an algorithm whose regret asymptotically matches these lower bound. Therefore, by comparing $c_1(\theta)$, $c_2(\theta)$, and $c_3(\theta)$, we can quantify the potential performance improvements taking routing decisions at each hop rather than at the source only, and observing per-link delays rather than end-to-end delays.

IV. FUNDAMENTAL PERFORMANCE LIMITS

In this section, we provide fundamental performance limits satisfied by *any* online routing policy in Π_1 , Π_2 , or Π_3 . Specifically, we derive asymptotic (when N grows large) regret lower bounds for our three types of policies. These bounds are obtained exploiting some results and techniques used in the control of Markov chains [30], and they are *tight* in the sense that there exist algorithms achieving these performance limits.

A. Regret Lower Bounds

We restrict our attention to the so-called *uniformly good* policies, under which the number of times sub-optimal paths are selected until the transmission of the n -th packet is $o(n^\alpha)$ when $n \rightarrow \infty$ for any $\alpha > 0$ and for all θ . We know from [30, Theorem 2] that such policies exist.

1) *Source-Routing with Bandit Feedback*: Denote by $\psi_\theta^p(k)$ the probability that the delay of a packet sent on path p is k slots, and by $h(p)$ the length (or number of links) of path p . The end-to-end delay is the sum of several independent random geometric variables. If we assume that $\theta_i \neq \theta_j$ for $i \neq j$, we have [31], for all $k \geq h(p)$,

$$\psi_\theta^p(k) = \sum_{i \in \mathcal{P}} \left(\prod_{j \in \mathcal{P}, j \neq i} \frac{\theta_j}{\theta_j - \theta_i} \right) \theta_i (1 - \theta_i)^{k-1},$$

i.e., the path delay distribution is a weighted average of the individual link delay distributions where the weights can be negative but always sum to one.

The next theorem provides the fundamental performance limit of online routing policies in Π_1 .

Theorem 4.1: For all θ and for any uniformly good policy $\pi \in \Pi_1$, $\liminf_{N \rightarrow \infty} \frac{R^\pi(N)}{\log(N)} \geq c_1(\theta)$, where $c_1(\theta)$ is the infimum of the following optimization problem:

$$\inf_{x \geq 0} \sum_{p \in \mathcal{P}} x_p \Delta_p \quad (1)$$

$$\text{subject to: } \inf_{\lambda \in B_1(\theta)} \sum_{p \neq p^*} x_p \sum_{k=h(p)}^{\infty} \psi_\theta^p(k) \log \frac{\psi_\theta^p(k)}{\psi_\lambda^p(k)} \geq 1,$$

with

$$B_1(\theta) = \left\{ \lambda : \{\lambda_i, i \in p^*\} = \{\theta_i, i \in p^*\}, \min_{p \in \mathcal{P}} D_\lambda(p) < D_\lambda(p^*) \right\}.$$

The variables $x_p, p \in \mathcal{P}$ solving (1) have the following interpretation: for $p \neq p^*$, $x_p \log(N)$ is the asymptotic number of packets that needs to be sent (up to the N -th packet) on sub-optimal path p under optimal routing strategies in Π_1 . Hence, x_p determines the optimal rate of *exploration* of sub-optimal path p . $B_1(\theta)$ is the set of *bad* network parameters: if $\lambda \in B_1(\theta)$, then the end-to-end delay distribution along the optimal path p^* is the same under θ or λ (hence by observing the end-to-end delay on path p^* , we cannot distinguish λ or θ), and p^* is not optimal under λ .

It is important to observe that in the definition of $B_1(\theta)$, the equality $\{\lambda_i, i \in p^*\} = \{\theta_i, i \in p^*\}$ is a set equality, i.e., order does not matter (e.g., if $p^* = \{1, 2\}$, the equality means that either $\lambda_1 = \theta_1, \lambda_2 = \theta_2$ or $\lambda_1 = \theta_2, \lambda_2 = \theta_1$).

2) *Source-Routing with Semi-Bandit (Per-Link) Feedback*: We now consider routing policies in Π_2 that make decisions at the source, but have information on the individual link delays. Let $\text{KLG}(u, v)$ denote the KL divergence number between two geometric random variables with parameters u and v :

$$\text{KLG}(u, v) := \sum_{k \geq 1} u(1-u)^{k-1} \log \frac{u(1-u)^{k-1}}{v(1-v)^{k-1}}.$$

Theorem 4.2: For all θ and for any uniformly good policy $\pi \in \Pi_2$, $\liminf_{N \rightarrow \infty} \frac{R^\pi(N)}{\log(N)} \geq c_2(\theta)$, where $c_2(\theta)$ is the infimum of the following optimization problem:

$$\inf_{x \geq 0} \sum_{p \in \mathcal{P}} x_p \Delta_p \quad (2)$$

$$\text{subject to: } \inf_{\lambda \in B_2(\theta)} \sum_{p \neq p^*} x_p \sum_{i \in \mathcal{P}} \text{KLG}(\theta_i, \lambda_i) \geq 1,$$

with

$$B_2(\theta) = \left\{ \lambda : \lambda_i = \theta_i, \forall i \in p^*, \min_{p \in \mathcal{P}} D_\lambda(p) < D_\lambda(p^*) \right\}.$$

The variables $x_p, p \in \mathcal{P}$ solving (2) have the same interpretation as that given previously in the case of bandit feedback. Again $B_2(\theta)$ is the set of parameters λ such that the distributions of link delays along the optimal path are the same under θ and λ , and p^* is not the optimal path under λ . The slight difference between the definitions of $B_1(\theta)$ and $B_2(\theta)$ comes from the difference of feedback (bandit vs. semi-bandit). It is also noted that $B_2(\theta) \subset B_1(\theta)$. We stress that by [30, Theorem 2], the asymptotic regret lower bounds of Theorems 4.1-4.2 are tight, namely there exists policies that achieve these regret bounds.

Remark 4.1: Of course, we know that $c_1(\theta) \geq c_2(\theta)$, since the lower bounds we derive are tight and getting per-link delay feedback can be exploited to design smarter routing policies than those we can devise using end-to-end delay feedback (i.e., $\Pi_1 \subset \Pi_2$).

Remark 4.2: The asymptotic lower bound proposed in [5] has a similar expression to ours, but the set $B_2(\theta)$ is replaced by $B'_2(\theta) = \bigcup_{i \in E} \{\lambda : \lambda_j = \theta_j, \forall j \neq i, \min_{p \in \mathcal{P}} D_\lambda(p) < D_\lambda(p^*)\}$. Note that $B'_2(\theta) \subset B_2(\theta)$, which implies that the lower bound derived in [5] is smaller than ours. In other words, we propose a regret lower bound that improves that in [5]. Furthermore, our bound is tight (it cannot be improved further).

The proof of Theorems 4.1 and 4.2 leverage techniques from [30] developed for the control of Markov chains, and are presented in Appendix A. Theorem 4.2 can be seen as a direct consequence of [30, Theorem 1] (the problem can be easily mapped to a controlled Markov chain). In contrast, the proof of Theorem 4.1 requires a more clever mapping due to the different nature of feedback. To prove Theorem 4.1, we establish Lemma 2, a property for geometric random variables.

3) *Hop-by-hop Routing*: Finally, we consider routing policies in Π_3 . These policies are more involved to analyze as the routing choices may change at any intermediate node in the network, and they are also more complex to implement. Surprisingly, the next theorem states that the regret lower

bound for hop-by-hop routing policies is the same as that derived for strategies in Π_2 (source-routing with semi-bandit feedback). In other words, we cannot improve the performance by taking routing decisions at each hop.

Theorem 4.3: For all θ and for any uniformly good rule $\pi \in \Pi_3$, $\liminf_{N \rightarrow \infty} \frac{R^\pi(N)}{\log(N)} \geq c_3(\theta) = c_2(\theta)$.

The proof of Theorem 4.3 is more involved than those of previous theorems, since in the hop-by-hop case, the chosen path could change at intermediate nodes. To overcome this difficulty, we introduce another notion of regret corresponding to the achieved throughput (i.e., the number of packets successfully received by the destination per unit time), which we refer to as the *throughput regret*. The proof uses the results of [30] for throughput regret, but also relies on Lemma 4, which provides an asymptotic relationship between $R^\pi(N)$ and the throughput regret.

As shown in [30, Theorem 2], the asymptotic regret lower bounds derived in Theorems 4.1-4.2-4.3 are *tight* in the sense that one can design actual routing policies achieving these regret bounds (although these policies might well be extremely complex to compute and impractical to implement). Hence from the fact that $c_1(\theta) \geq c_2(\theta) = c_3(\theta)$, we conclude that:

- The optimal source-routing policy with semi-bandit feedback asymptotically achieves a lower regret than the optimal source-routing policy with bandit feedback;
- The optimal hop-by-hop routing policy asymptotically obtains the same regret as the optimal source-routing policy with semi-bandit feedback.

B. Numerical Example

There are examples of network topologies where the above asymptotic regret lower bounds can be explicitly computed. One such example is the line network; see e.g. Figure 1(a). Notice that in line networks, the optimal routing policy consists in selecting the best link in each hop. The following lemma is immediate:

Lemma 1: For any line network with H hops, we have:

$$c_1(\theta) \geq \sum_{i \notin p^*} \frac{\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}}}{\max_{p: i \in p} \sum_{k=H}^{\infty} \psi_{\theta}^p(k) \log \frac{\psi_{\theta}^p(k)}{\psi_{\theta}^i(k)}},$$

$$c_2(\theta) = c_3(\theta) = \sum_{i \notin p^*} \frac{\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}}}{\text{KLG}(\theta_i, \theta_{\zeta(i)})},$$

where $\zeta(i)$ is the best link on the same hop as link i and ϑ^i is a vector of link parameters defined as $\vartheta_j^i = \theta_j$ if $j \neq i$, and $\vartheta_i^i = \theta_{\zeta(i)}$.

Proposition 4.4: There exist problem instances in line networks, for which the regret of any uniformly good policy in $\Pi_2 \cup \Pi_3$ is $\Omega\left(\frac{|E|-H}{\Delta_{\min} \theta_{\min}^2} \log(N)\right)$.

For line networks, both $c_1(\theta)$ and $c_2(\theta)$ scale linearly with the number of links in the network. In Figure 1(b), we plot the lower bound of the ratio $\frac{c_1(\theta)}{c_2(\theta)}$ (based on the previous lemma) averaged over various values of θ (we randomly generated 10^6 link parameters θ) as a function of the network size H in a

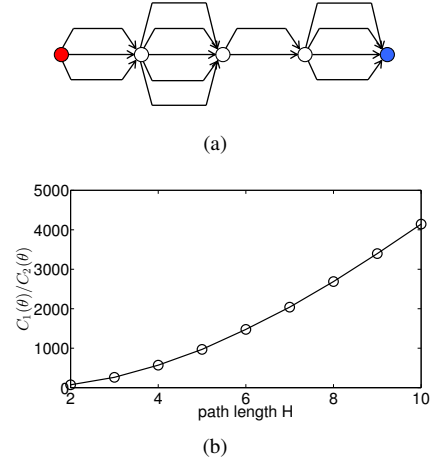


Fig. 1. The line network: (a) Topology, (b) Semi-bandit vs. bandit feedback: lower bound on the average ratio between the two corresponding asymptotic regret lower bounds ($c_1(\theta)/c_2(\theta)$).

simple line network, which has two links in the first hop and one link in the rest of hops and hence $|E| = H + 1$. These results suggest that collecting semi-bandit feedback (per-link delays) can significantly improve the performance of routing policies. The gain is significant even for fairly small networks – the regret is reduced by at least a factor 1500 on average in 6-hop networks when collecting per-link delays.

V. ROUTING POLICIES FOR SEMI-BANDIT FEEDBACK

Theorems 4.1-4.2-4.3 indicate that within the first N packets, the total amount of packets routed on a sub-optimal path p should be of the order of $x_p^* \log(N)$ where x_p^* is the optimal solution of the optimization problems in (1) and (2). In [30], the authors present policies that achieve the regret bounds of Theorems 4.1-4.2-4.3 (see [30, Theorem 2]). These policies suffer from two problems: firstly, they are computationally infeasible for large problems since their implementation involves solving in each round a semi-infinite linear program [32] similar to those providing the regret lower bounds (defined in (1) and (2)). Secondly, these policies have no finite-time performance guarantees, and numerical experiments suggest that their finite-time performance on typical problems is rather poor.

In this section, we present online routing policies for semi-bandit feedback, which are simple to implement, yet approach the performance limits identified in the previous section. We further analyze their regret, and show that they outperform existing algorithms. To present our policies, we introduce additional notations. Under a given policy, we let $t_i(n)$ be the total number of transmission attempts (including retransmissions) on link i before the n -th packet is sent. We define $\hat{\theta}_i(n)$ the empirical success rate of link i estimated over the transmissions of the first $(n-1)$ packets. We define the corresponding vectors $t(n) = (t_i(n))_{i \in E}$ and $\hat{\theta}(n) = (\hat{\theta}_i(n))_{i \in E}$.

Note that the proposed policies and regret analysis presented in this section directly apply for generic combinatorial optimization problems with linear objective function and geometrically distributed rewards.

Index	Type	Computation	Algorithm
b_p	Path	Line search	GEOCOMBUCB-1
c_p	Path	Explicit	GEOCOMBUCB-2
ω_i	Edge	Line search	KL-SR

TABLE III
SUMMARY OF INDEXES.

A. Path and Link Indexes

The proposed policies rely on indexes attached either to individual links or paths. Next we introduce three indexes used in our policies. They depend on the round, i.e., on the number n of packets already sent, and on the estimated link parameters $\hat{\theta}(n)$. The three indexes and their properties (i.e., in which policy they are used, and how one can compute them) are summarized in Table III. Let $n \geq 1$ and assume that the n -th packet is to be sent. The indexes are defined as follows.

1) *Path Indexes:* Let $\lambda \in (0, 1]^{|E|}$, $t \in \mathbb{N}^{|E|}$, and $n \in \mathbb{N}$. The first path index, denoted by $b_p(n, \lambda, t)$ for path $p \in \mathcal{P}$, is motivated by the index defined in [18]. $b_p(n, \lambda, t)$ is defined as the infimum of the following optimization problem:

$$\begin{aligned} & \inf_{u \in (0,1]^{|E|}} p^\top u^{-1} \\ \text{subject to: } & \sum_{i \in p} t_i \text{KL}(\lambda_i, u_i) \leq f_1(n), \\ & u_i \geq \lambda_i, \quad \forall i \in E, \end{aligned}$$

where $f_1(n) = \log(n) + 4H \log(\log(n))$, and for all $a, b \in [0, 1]$, $\text{KL}(a, b)$ is the KL-divergence number between two Bernoulli distributions with respective means a and b , i.e., $\text{KL}(a, b) = a \log(a/b) + (1-a) \log((1-a)/(1-b))$.

The second index is denoted by $c_p(n, \lambda, t)$ and defined for path $p \in \mathcal{P}$ as:

$$c_p(n, \lambda, t) = p^\top \lambda^{-1} - \sqrt{\sum_{i \in p} \frac{2f_1(n)}{t_i \lambda_i^3}}.$$

The next theorem provides generic properties of the two indexes b_p and c_p .

Theorem 5.1: (i) For all $n \geq 1$, $p \in \mathcal{P}$, $\lambda \in (0, 1]^{|E|}$, and $t \in \mathbb{N}^{|E|}$, we have $b_p(n, \lambda, t) \geq c_p(n, \lambda, t)$.

(ii) There exists a constant $K_H > 0$ depending on H only such that, for all $p \in \mathcal{P}$ and $n \geq 2$:

$$\mathbb{P}[b_p(n, \hat{\theta}(n), t(n)) \geq p^\top \theta] \leq K_H n^{-1} (\log(n))^{-2}.$$

Corollary 5.2: We have:

$$\begin{aligned} & \sum_{n \geq 1} \mathbb{P}[b_{p^*}(n, \hat{\theta}(n), t(n)) \geq p^{*\top} \theta^{-1}] \\ & \leq 1 + K_H \sum_{n \geq 2} n^{-1} (\log(n))^{-2} < \infty. \end{aligned}$$

2) *Link Index:* Our third index is a link index. For $n, t \in \mathbb{N}$ and $\lambda \in (0, 1]$, the index $\omega_i(n, \lambda, t)$ of link $i \in E$ is defined as:

$$\omega_i(n, \lambda, t) = \min \left\{ \frac{1}{u} : u \in [\lambda, 1], \quad t \text{KL}(\lambda, u) \leq f_2(n) \right\},$$

where $f_2(n) = \log(n) + 4 \log(\log(n))$.

Algorithm 1 GEOCOMBUCB

for $n \geq 1$ **do**

Select path $p(n) \in \arg \min_{p \in \mathcal{P}} \xi_p(n)$ (ties are broken arbitrarily), where $\xi_p(n) = b_p(n)$ for GEOCOMBUCB-1, and $\xi_p(n) = c_p(n)$ for GEOCOMBUCB-2.

Collect feedback on links $i \in p(n)$, and update $\hat{\theta}_i(n)$ for $i \in p(n)$.

Algorithm 2 KL-SR

for $n \geq 1$ **do**

Select path $p(n) \in \arg \min_{p \in \mathcal{P}} p^\top \omega(n)$ (ties are broken arbitrarily).

Collect feedback on links $i \in p(n)$, and update $\hat{\theta}_i(n)$ for $i \in p(n)$.

B. Routing policies

We present three routing policies, referred to as GEOCOMBUCB-1, GEOCOMBUCB-2 and KL-SR, respectively. For the transmission on the n -th packet, GEOCOMBUCB-1 (resp. GEOCOMBUCB-2) selects the path p with the lowest index $b_p(n) := b_p(n, \hat{\theta}(n), t(n))$ (resp. $c_p(n) := c_p(n, \hat{\theta}(n), t(n))$). KL-SR was initially proposed in [1] and for the transmission of the n -th packet, it selects the path $p(n) \in \arg \min_{p \in \mathcal{P}} p^\top \omega(n)$, where $\omega(n) = (\omega_i(n), i \in E)$ and $\omega_i(n) := \omega_i(n, \hat{\theta}_i(n), t_i(n))$. The pseudo-code of GEOCOMBUCB and KL-SR are presented in Algorithm 1 and Algorithm 2, respectively.

In the following theorems, we provide a finite time analysis of the GEOCOMBUCB and KL-SR policies and show the optimality of KL-SR in line networks. Define $\varepsilon = (1 - 2^{-\frac{1}{4}}) \frac{\Delta_{\min}}{D^+}$.

Theorem 5.3: For all $N \geq 1$, under policies $\pi \in \{\text{GEOCOMBUCB-1}, \text{GEOCOMBUCB-2}\}$ we have:

$$R^\pi(N) \leq \frac{16|E|\sqrt{H}f_1(N)}{\Delta_{\min}\theta_{\min}^2} + 2D^+ \left(2K_H + \sum_{i \in E} \frac{1}{\varepsilon^2 \theta_i^2} \right).$$

Hence $R^\pi(N) = \mathcal{O} \left(\frac{|E|\sqrt{H}}{\Delta_{\min}\theta_{\min}^2} \log(N) \right)$ when $N \rightarrow \infty$.

Theorem 5.4: For all $N \geq 1$, under policy $\pi = \text{KL-SR}$ we have:

$$R^\pi(N) \leq \frac{360|E|Hf_2(N)}{\Delta_{\min}\theta_{\min}^2} + 2D^+ \left(4H + \sum_{i \in E} \frac{1}{\varepsilon^2 \theta_i^2} \right).$$

Hence $R^\pi(N) = \mathcal{O} \left(\frac{|E|H}{\Delta_{\min}\theta_{\min}^2} \log(N) \right)$ when $N \rightarrow \infty$.

The index b_p is an extension of the KL-based index of [18] to the case of geometrically distributed rewards. However the proof of Theorem 5.3 is novel and uses the link between b_p and c_p established in Theorem 5.1. The proof of Theorem 5.3 uses some of ideas from [18]. The proof of Theorem 5.4 is completely different from the regret analysis of KL-SR in [1]; it relies on Lemma 8, which provides a tight lower bound for the index ω_i , and borrows some ideas from [17, Theorem 5].

Remark 5.1: Theorem 5.4 holds even when the delays on the various links are not independent as in [17].

The proposed policies have better performance guarantees than existing routing algorithms. Indeed, as shown in Appendix J, the best regret upper bound for the CUCB algorithm [9] is $R^{\text{CUCB}}(N) = \mathcal{O}\left(\frac{|E|H}{\Delta_{\min}\theta_{\min}^3} \log(N)\right)$, which constitutes a weaker performance guarantee than those of our routing policies. The numerical experiments presented in the next section will confirm the superiority of GEOCOMBUCB and KL-SR over CUCB. The next proposition states that KL-SR is asymptotically optimal in line networks.

Proposition 5.5: In line networks, the regret under $\pi = \text{KL-SR}$ satisfies $\limsup_{N \rightarrow \infty} \frac{R^\pi(N)}{\log(N)} \leq c_2(\theta)$. Hence, $R^\pi(N) = \mathcal{O}\left(\frac{|E|H}{\Delta_{\min}\theta_{\min}^3} \log(N)\right)$ when $N \rightarrow \infty$.

Remark 5.2: When the link parameters smoothly evolve over time, we can modify the proposed routing policies so that routing decisions are based on past choices and observations over a sliding window consisting of a fixed number of packets, as considered in [33] and [34].

C. Implementation

Next we discuss the implementation of our routing policies, and give simple methods to compute $b_p(n, \lambda, t)$, $c_p(n, \lambda, t)$, $\omega_i(n, \lambda, t)$ given p, i, n, λ and t . The path index c_p is explicit and easy to compute. The link index ω_i is also straightforward as it amounts to finding the roots of a strictly convex and increasing function in one variable (note that $v \mapsto \text{KL}(u, v)$ is strictly convex and increasing for $v \geq u$). Hence, the index ω_i can be computed by a simple line search. The path index $b_p(n, \lambda, t)$ can also be computed using simple line search, as shown below.

Define $I_p(\lambda) = \{i \in p : \lambda_i \neq 1\}$, and for $\gamma > 0$, define:

$$F(\gamma, \lambda, n, t) = \sum_{i \in I_p(\lambda)} t_i \text{KL}(\lambda_i(n), g(\gamma, \lambda_i, t_i)), \text{ with}$$

$$g(\gamma, \lambda_i, t_i) = \frac{1}{2\gamma t_i} \left(\gamma \lambda_i t_i - 1 + \sqrt{(1 - \gamma \lambda_i t_i)^2 + 4\gamma t_i} \right).$$

Proposition 5.6: (i) $\gamma \mapsto F(\gamma, \lambda, n, t)$ is strictly increasing, and $F(\mathbb{R}^+, \lambda, n, t) = \mathbb{R}^+$. (ii) If $I_p(\lambda) = \emptyset$, $b_p(n, \lambda, t) = \sum_{i \in E} p_i$. Otherwise, let γ^* is the unique solution to $F(\gamma, \lambda, n, t) = f_1(n)$. Then,

$$b_p(n, \lambda, t) = \sum_{i \in E} p_i - |I_p(\lambda)| + \sum_{i \in I_p(\lambda)} g(\gamma^*, \lambda_i, t_i).$$

As stated in Proposition 5.6, proven in Appendix I, γ^* can be computed efficiently by a simple line search, and b_p is easily deduced. We thus have efficient methods to compute the three indexes. To implement our policies, we then need to find in each round, the path maximizing the index (or the sum of link indexes along the path for KL-SR). KL-SR can be implemented (in a distributed fashion) using the Bellman-Ford algorithm, and its complexity is $\mathcal{O}(|V||E|)$ in each round. GEOCOMBUCB-1 and GEOCOMBUCB-2 are more computationally involved than KL-SR and have complexity $\mathcal{O}(|\mathcal{P}|)$ in each round.

D. Numerical Experiments

In this section, we conduct numerical experiments to compare the performance of the proposed source-routing policies to that of the CUCB algorithm [9] and TS applied to our online routing problem. The CUCB algorithm is an index policy in Π_2 (the set of source-routing policies with semi-bandit feedback), and selects path $p(n)$ for the transmission of the n -th packet:

$$p(n) \in \arg \min_{p \in \mathcal{P}} \sum_{i \in p} \frac{1}{\hat{\theta}_i(n) + \sqrt{1.5 \log(n)/t_i(n)}}.$$

We consider a grid network whose topology is depicted in Figure 2(a), where the node in red (resp. blue) is the source (resp. the destination). In this network, there are $\binom{6}{3} = 20$ possible paths from the source to the destination. Let us compare these algorithms in terms of their per-packet complexity. The complexity of GEOCOMBUCB-1 and GEOCOMBUCB-2 is $\mathcal{O}(|\mathcal{P}|)$, whereas that of KL-SR, CUCB, and TS is $\mathcal{O}(|V||E|)$.

In Figures 2(b)-(c), we plot the regret against the number of the packets N under the various routing policies, and for two sets of link parameters θ . For each set, we choose a value of θ_{\min} and generate the values of θ_i independently, uniformly at random in $[\theta_{\min}, 1]$. The results are averaged over 100 independent runs, and the 95% confidence intervals are shown using the grey area around curves. The three proposed policies outperform CUCB, and GEOCOMBUCB-1 attains the smallest regret amongst the proposed policies. The comparison between GEOCOMBUCB-2 and KL-SR is more subtle and depends on the link parameters: while in Figure 2(b) KL-SR significantly outperforms GEOCOMBUCB-2, they attain regrets growing similarly for the link parameter of Figure 2(c). Yet there are some parameters for which KL-SR is significantly outperformed by GEOCOMBUCB-2. KL-SR seems to perform better than GEOCOMBUCB-2 in scenarios where Δ_{\min} is large. TS performs slightly better than GEOCOMBUCB-1 on average. Its regret, however may not be well concentrated around the mean for some link parameters, as in Figure 2(c). Furthermore, the regret analysis of TS for shortest-path routing with general topologies is an open problem.

E. A distributed hop-by-hop routing policy

Motivated by the Bellman-Ford implementation of KL-SR algorithm, we propose KL-HHR, a distributed routing policy which is a hop-by-hop version of KL-SR algorithm and hence belongs to the set of policies Π_3 . We first introduce the necessary notations. For any node $v \in V$, we let \mathcal{P}_v denote the set of loop-free paths from node v to the destination. For any time slot τ , we denote by $n(\tau)$ the packet number that is about to be sent or already in the network. For any edge i , let $\hat{\theta}_i(\tau)$ be the empirical success rate of edge i up to time slot τ , that is $\hat{\theta}_i(\tau) = s_i(n(\tau))/t'_i(\tau)$, where $t'_i(\tau)$ denotes the total number of transmission attempts on link i up to time slot τ . Moreover, with slight abuse of notation, we denote the index of link i at time τ by $\omega_i(\tau, \hat{\theta}_i(\tau))$. Note that by definition $t'_i(\tau) \geq t_i(n)$ and $\hat{\theta}_i(\tau)$ is a more accurate estimate of θ_i than $\hat{\theta}_i(n(\tau))$. We define $J_v(\tau)$ as the minimum cumulative index

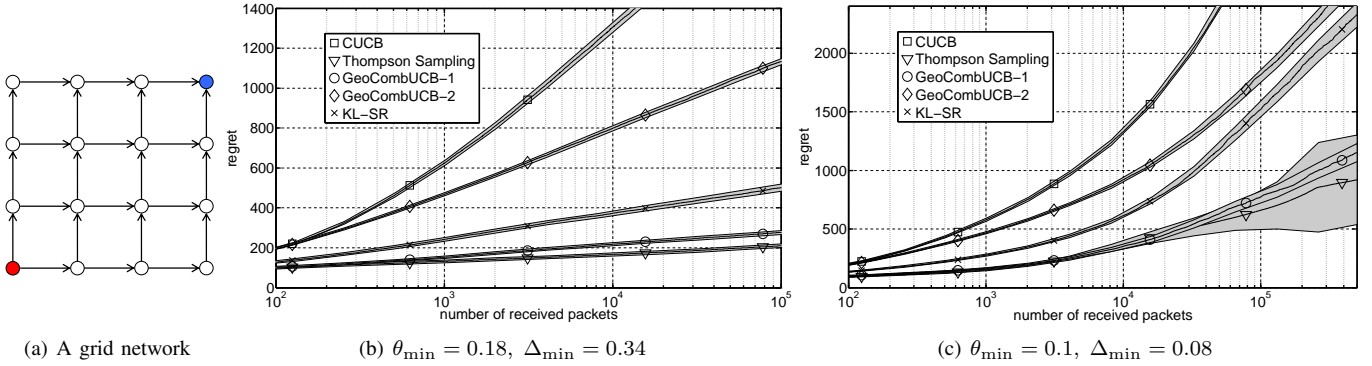


Fig. 2. Network topology, and regret versus number of received packets.

from node v to the destination:

$$J_v(\tau) = \min_{p \in \mathcal{P}_v} \sum_{i \in p} \omega_i(\tau, \tilde{\theta}_i(\tau)).$$

We note that $J_v(\tau)$ can be computed using Bellman-Ford algorithm. KL-HHR works based on the following idea: at time τ if the current packet is at node v , it will be sent to node v' with $(v, v') \in E$ such that $\omega_{(v, v')}(\tau, \tilde{\theta}_v(\tau)) + J_{v'}(\tau)$ is minimal over all outgoing edges of node v . The pseudo-code of KL-HHR is given in Algorithm 3.

Algorithm 3 KL-HHR for node v

for $\tau \geq 1$ **do**

 Select link $(v, v') \in E$, where

$$v' \in \arg \min_{w \in V: (v, w) \in E} \left(\omega_{(v, w)}(\tau, \tilde{\theta}_v(\tau)) + J_w(\tau) \right).$$

 Update index of the link (v, v') .

We compare the performance of KL-HHR and KL-SR through numerical experiments. We consider a grid network whose topology is depicted in Figure 3(a), in which there are 40 links and 413 possible paths from the source (in red) to the destination (in blue). Figures 3(b)-(c) display the regret against the number of the packets N under KL-SR and KL-HHR for two sets of link parameters θ . The values of θ_i are generated similarly to the previous experiments and the results are averaged over 100 independent runs. As expected, KL-HHR outperforms KL-SR in both scenarios, and the difference is significant when θ_{\min} is small. The reason is that KL-HHR can change routing decisions dynamically at intermediate nodes, and does not waste transmissions on bad links when they are discovered. It is noted, however that, irrespective of the value of θ_{\min} , the regret of both KL-HHR and KL-SR grow similarly when the number of received packets grows large.

The regret analysis of KL-HHR is beyond the scope of this paper, and is left for future work.

VI. CONCLUSIONS AND FUTURE WORK

We have studied online shortest path routing problems in networks with stochastic link delays. The distributions of these

delays are initially unknown, and have to be estimated by actual packet transmissions. Three types of routing policies are analyzed: source-routing with semi-bandit feedback, source-routing with bandit feedback, and hop-by-hop routing. Tight asymptotic lower bounds on the regret for the three types of policies are derived. By comparing these bounds, we observed that semi-bandit feedback significantly improves performance while hop-by-hop decisions do not. Finally, we proposed several simple routing policies for semi-bandit feedback that outperform alternatives from the literature both in theory and in numerical experiments. As future work, we plan to propose practical algorithms with provable performance bounds for hop-by-hop routing and source-routing with bandit feedback. Furthermore, we would like to study the effect of delayed feedback on the performance as studied in, e.g., [35].

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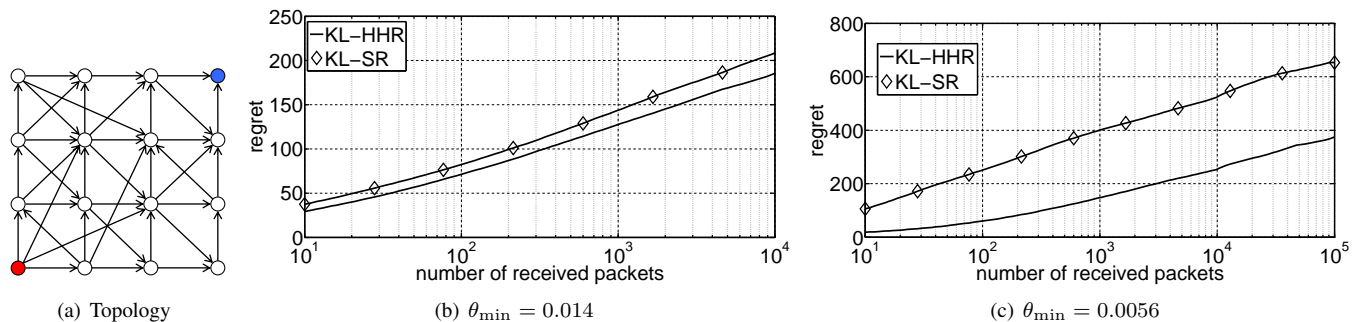


Fig. 3. Network topology, and regret versus number of received packets.

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APPENDIX A PROOFS OF THEOREMS 4.1, 4.2 AND 4.3

To derive the asymptotic regret lower bounds, we apply the techniques used by Graves and Lai [30] to investigate efficient adaptive decision rules in controlled Markov chains. We recall here their general framework. Consider a controlled Markov chain $(X_t)_{t \geq 0}$ on a countable state space \mathcal{S} with a control set \mathcal{U} . The transition probabilities given control $u \in \mathcal{U}$ are parameterized by θ taking values in a compact metric space Θ : the probability to move from state x to state y given the control u and the parameter θ is $P(x, y; u, \theta)$. The parameter θ is not known. The decision maker is provided with a finite set of stationary control laws $G = \{g_1, \dots, g_K\}$ where each control law g_j is a mapping from \mathcal{S} to \mathcal{U} : when control law g_j is applied in state x , the applied control is $u = g_j(x)$. It is assumed that if the decision maker always selects the same control law g , the Markov chain is irreducible with respect to some maximum irreducibility measure and has stationary distribution π_θ^g . The reward obtained when applying control u in state x is denoted by $r(x, u)$, so that the expected reward achieved under control law g is $\mu_\theta(g) = \sum_x r(x, g(x))\pi_\theta^g(x)$. There is an optimal control law given θ whose expected reward is denoted by $\mu_\theta^* = \max_{g \in G} \mu_\theta(g)$. Now the objective

of the decision maker is to sequentially apply control laws so as to maximize the expected reward up to a given time horizon N . The performance of the decision making scheme can be quantified through the notion of regret which compares the expected reward to that obtained by always applying the optimal control law.

A. Source Routing with Bandit Feedback – Theorem 4.1

To prove Theorem 4.1, we construct a controlled Markov chain as follows. The state space is \mathbb{N} , the control set is the set of paths \mathcal{P} , and the parameter $\theta = (\theta_i, i \in E)$ defines the success rates on the various links. The parameter θ takes value in the compact space $\Theta = [\varepsilon, 1]^{|E|}$ for ε arbitrarily close to zero. The set of control laws are stationary and each of them corresponds to a given path, i.e., $G = \mathcal{P}$. A transition in the Markov chain occurs at time epochs where a new packet is sent. The state after a transition records the end-to-end delay of the packet. Hence the transition probabilities are $P(k, l; p, \theta) = \psi_\theta^p(l)$, and do not depend on the starting state. The cost (the opposite of reward) at state l is simply equal to the delay l . Let us fix θ , and denote by p^* the corresponding optimal path. For any two sets of parameters θ and λ , we define the KL information number under path (or control law) p as:

$$I^p(\theta, \lambda) = \sum_{l=h(p)}^{\infty} \psi_\theta^p(l) \log \frac{\psi_\theta^p(l)}{\psi_\lambda^p(l)}. \quad (3)$$

We have that $I^p(\theta, \lambda) = 0$ if and only if the delays over path p under parameters θ and λ have the same distribution. By Lemma 2, proven at the end of this subsection, this occurs if and only if the two following sets are identical: $\{\theta_i, i \in p\}$, $\{\lambda_i, i \in p\}$. We further define $B_1(\theta)$ as the set of bad parameters λ such that under λ , p^* is not the optimal path, and such that θ and λ are statistically not distinguishable (they lead to the same delay distribution along path p^*). Then:

$$B_1(\theta) = \left\{ \lambda : \{\lambda_i, i \in p^*\} = \{\theta_i, i \in p^*\}, \min_{p \in \mathcal{P}} D_\lambda(p) < D_\lambda(p^*) \right\}.$$

By [30, Theorem 1], we conclude that the delay regret scales at least as $c_1(\theta) \log(N)$ where

$$c_1(\theta) = \inf \left\{ \sum_{p \in \mathcal{P}} x_p \Delta_p : x \geq 0, \inf_{\lambda \in B_1(\theta)} \sum_{p \neq p^*} x_p I^p(\theta, \lambda) \geq 1 \right\},$$

where $I^p(\theta, \lambda)$ is given in (3). \square

Lemma 2: Consider $(X_i)_i$ independent with $X_i \sim \text{Geo}(\theta_i)$ and $0 < \theta_i \leq 1$. Consider $(Y_i)_i$ independent with $Y_i \sim \text{Geo}(\lambda_i)$ and $0 < \lambda_i \leq 1$. Define $\bar{X} = \sum_i X_i$ and $\bar{Y} = \sum_i Y_i$. Then $\bar{X} \stackrel{d}{=} \bar{Y}$ if and only if $(\theta_i)_i = (\lambda_i)_i$ up to a permutation⁵.

Proof. If $(\theta_i)_i = (\lambda_i)_i$, up to a permutation then $\bar{X} \stackrel{d}{=} \bar{Y}$ by inspection. Assume that $\bar{X} \stackrel{d}{=} \bar{Y}$. Define $z_m = \min_i (\min(1/(1-\theta_i), 1/(1-\lambda_i)))$. For all z such that $|z| < z_m$ we have $\mathbb{E}[z^{\bar{X}}] = \mathbb{E}[z^{\bar{Y}}]$ so that

$$\prod_i \frac{\theta_i}{1 - (1 - \theta_i)z} = \prod_i \frac{\lambda_i}{1 - (1 - \lambda_i)z}.$$

⁵The symbol $\stackrel{d}{=}$ denotes equality in distribution.

Hence:

$$P_X(z) := \prod_i \theta_i (1 - (1 - \lambda_i)z) = \prod_i \lambda_i (1 - (1 - \theta_i)z) := P_Y(z).$$

Both $P_X(z)$ and $P_Y(z)$ are polynomials and are equal on an open set. So they are equal everywhere, and the sets of their roots are equal $\{1/(1 - \theta_i), i\} = \{1/(1 - \lambda_i), i\}$. So $(\theta_i)_i = (\lambda_i)_i$ up to a permutation as announced. \square

B. Source Routing with Semi-bandit Feedback – Theorem 4.2

The proof of Theorem 4.2 is similar to that of Theorem 4.1, except that here we have to account for the fact that the source gets feedback on per-link basis. To this end, we construct a Markov chain that records the delay on each link of a path. The state space is $\mathbb{N}^{|E|}$. Transitions occur when a new packet is sent from the source, and the corresponding state records the observed delays on each link of the chosen path, and the components of the state corresponding to links not involved in the path are set equal to 0. For example, the state $(0, 1, 4, 0, 7)$ indicates that the path consisting of links 2, 3, and 5 has been used, and that the per-links delays are 1, 4, and 7, respectively. The cost of a given state is equal to the sum of its components (total delay). Now assume that path $p = (i_1, \dots, i_{h(p)})$ is used to send a packet, then the transition probability to a state whose i_k -th component is equal to d_k , $k = 1, \dots, h(p)$ (the other components are 0) is $\prod_{k=1}^{h(p)} q_\theta(i_k, d_k)$, where $q_\theta(i, m) = \theta_i (1 - \theta_i)^{m-1}$ for any link i and any delay m . Now the KL information number of (θ, λ) under path p is given by

$$I^p(\theta, \lambda) = \sum_{i \in p} \text{KLG}(\theta_i, \lambda_i), \quad (4)$$

since KL divergence is additive for independent random variables. Hence, under semi-bandit feedback, we have $I^p(\theta, \lambda) = 0$ if and only if $\theta_i = \lambda_i$ for all $i \in p$. The set $B_2(\theta)$ of bad parameters is defined as:

$$B_2(\theta) = \left\{ \lambda : \lambda_i = \theta_i \forall i \in p^*, \min_{p \in \mathcal{P}} D_\lambda(p) < D_\lambda(p^*) \right\}.$$

Applying [30, Theorem 1] gives:

$$c_2(\theta) = \inf \left\{ \sum_{p \in \mathcal{P}} x_p \Delta_p : x \geq 0, \inf_{\lambda \in B_2(\theta)} \sum_{p \neq p^*} x_p I^p(\theta, \lambda) \geq 1 \right\},$$

where $I^p(\theta, \lambda)$ is given in (4). \square

C. Hop-by-hop Routing – Theorem 4.3

This case is more involved. We first define another notion of regret corresponding to the achieved throughput (i.e., the number of packets successfully received by the destination per unit time). The throughput regret is introduced to ease the analysis, since computing the throughput regret is easier in the hop-by-hop case. Define $\mu_\theta(p)$ as the average throughput on path p given link success rates θ : $\mu_\theta(p) = 1/D_\theta(p)$. The *throughput regret* $S^\pi(T)$ of π over time horizon T is: $S^\pi(T) := T\mu_\theta(p^*) - \mathbb{E}[N^\pi(T)]$, where $N^\pi(T)$ is the number of packets received up to time T under policy π . Lemma 4, stated at the end of the proof, provides the relation between asymptotic bound on $R^\pi(N)$ and $S^\pi(N)$.

Now we are ready to prove Theorem 4.3. We let the state of the Markov chain be the packet location. The action is the selected outgoing link. The transitions between two states take one time slot – the time to make a transmission attempt. Hence, the transition probability between state x and y with the action of using link i is denoted by (where $y \neq x$) $P_\theta^i(x, y) = \theta_i$ if link i connects node x and y and is zero otherwise. On the other hand, the probability of staying at the same state is the transmission failure probability on link i if link i is an outgoing link, that is $P_\theta^i(x, x) = 1 - \theta_i$ if link i is an outgoing link, and is zero otherwise.

We assume that the packet is injected at the source immediately after the previous packet is successfully delivered, and we are interested in counting the number of successfully delivered packets. In order not to count the extra time slot we will spend at the destination, we use a single Markov chain state to represent both the source and the destination.

We give a reward of 1 whenever the packet is successfully delivered to the destination. Let $r(x, y, i)$ be the immediate reward after the transition from node x to node y under the action i , i.e., $r(x, y, i) = 1$ if y is the destination node and is zero otherwise (see Figure 4 for an example). Hence $r(x, i)$ (i.e., the reward at state x with action i) is

$$r(x, i) = \begin{cases} \theta_i & \text{if link } i \text{ connects node } x \text{ and the destination;} \\ 0 & \text{otherwise.} \end{cases}$$

The stationary control law prescribes the action at each state, i.e., the outgoing link at each node. A stationary control law of this Markov chain is then a path p in the network, and we assign arbitrary actions to the nodes that are not on the path p . The maximal irreducibility measure is then to assign measure zero to the nodes that are not on the path p , and a counting measure to the nodes on the path p . The Markov chain is irreducible with respect to this maximal irreducibility measure, and the stationary distribution of the Markov chain under path p is,

$$\pi_\theta^p(x) = \frac{1}{\theta_{p(x)} \sum_{i \in p} \frac{1}{\theta_i}} \mathbb{1}\{\text{if node } x \text{ is on the path } p\},$$

where $p(x)$ denotes the link we choose at node x . The long-run average reward of the Markov chain under control law p is $\sum_x \pi_\theta^p(x) r(x, p(x)) = \frac{1}{\sum_{i \in p} \frac{1}{\theta_i}} = \mu_\theta(p)$. The optimal control law is then p^* with long run average reward $\mu_\theta(p^*)$.

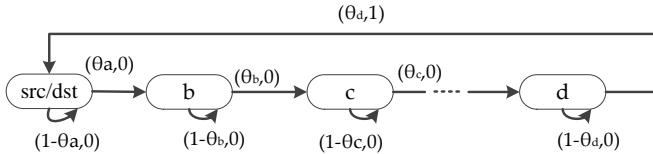


Fig. 4. A Markov chain example under a control law p where the values in the parenthesis respectively denote the transition probability and the reward.

The throughput regret of a policy $\pi \in \Pi_3$ for this controlled Markov chain at time T is

$$S^\pi(T) = T\mu_\theta(p^*) - \mathbb{E}_\theta \left[\sum_{t=1}^T r(x_t, \pi(t, x_t)) \right], \quad (5)$$

where x_t is the state at time t and $\pi(t, x_t)$ is the corresponding action for state x_t at time t . To this end, we construct a controlled Markov chain that corresponds to the hop-by-hop routing in the network. Now define $I^p(\theta, \lambda)$ as the KL information number for a control law p :

$$\begin{aligned} I^p(\theta, \lambda) &= \sum_x \pi_\theta^p(x) \sum_y P_\theta^{p(x)}(x, y) \log \frac{P_\theta^{p(x)}(x, y)}{P_\lambda^{p(x)}(x, y)} \\ &= \sum_x \pi_\theta^p(x) \left(\theta_{p(x)} \log \frac{\theta_{p(x)}}{\lambda_{p(x)}} + (1 - \theta_{p(x)}) \log \frac{1 - \theta_{p(x)}}{1 - \lambda_{p(x)}} \right) \\ &= \mu_\theta(p) \sum_{i \in p} \frac{\text{KL}(\theta_i, \lambda_i)}{\theta_i} = \mu_\theta(p) \sum_{i \in p} \text{KLG}(\theta_i, \lambda_i), \end{aligned} \quad (6)$$

where we used Lemma 3 in the last equality. Since $I^p(\theta, \lambda) = 0$ if and only if $\theta_i = \lambda_i$ for all $i \in p$, the set $B_2(\theta)$ of bad parameters is:

$$\begin{aligned} B_2(\theta) &= \left\{ \lambda : \lambda_i = \theta_i \forall i \in p^*, \max_{p \in \mathcal{P}} \mu_\lambda(p) > \mu_\lambda(p^*) \right\} \\ &= \left\{ \lambda : \lambda_i = \theta_i, \forall i \in p^*, \min_{p \in \mathcal{P}} D_\lambda(p) < D_\lambda(p^*) \right\}. \end{aligned}$$

Applying [30, Theorem 1], we get: $\liminf_{T \rightarrow \infty} S^\pi(T) / \log(T) \geq c'_3(\theta)$, where

$$c'_3(\theta) = \inf \left\{ \sum_{p \in \mathcal{P}} x_p \Delta_p : x \geq 0; \inf_{\lambda \in B_2(\theta)} \sum_{p \neq p^*} x_p \mu_\theta(p) I^p(\theta, \lambda) \geq 1 \right\},$$

where $I^p(\theta, \lambda)$ is given in (6). By Lemma 4, $c_3(\theta) \geq c'_3(\theta) / \mu_\theta(p^*)$. Lastly, observe that $\mu_\theta(p^*) - \mu_\theta(p) = \mu_\theta(p^*) \mu_\theta(p) (D_\theta(p) - D_\theta(p^*))$. It then follows that $c'_3(\theta) / \mu_\theta(p^*) = c_2(\theta)$ and therefore $c_3(\theta) \geq c_2(\theta)$. On the other hand, $c_3(\theta) \leq c_2(\theta)$ since $\Pi_2 \subset \Pi_3$. As a result, $c_3(\theta) = c_2(\theta)$ and the proof is completed. \square

The following two lemmas prove useful in the proof of Theorem 4.3. Lemma 3 follows from a straightforward calculation, and relates the KL-divergence between two geometric distributions to that of corresponding Bernoulli distributions. Lemma 4 provides the connection between the throughput regret $S^\pi(T)$ and delay regret $R^\pi(N)$ and its proof is provided in the next section.

Lemma 3: For any $u, v \in (0, 1]$, we have:

$$\text{KLG}(u, v) = \frac{\text{KL}(u, v)}{u}. \quad (7)$$

Proof. We have:

$$\begin{aligned} \text{KLG}(u, v) &= \sum_{i=1}^{\infty} \left[\log \frac{u(1-u)^{i-1}}{v(1-v)^{i-1}} \right] u(1-u)^{i-1} \\ &= \sum_{i=1}^{\infty} (\log \frac{u}{v}) u(1-u)^{i-1} \\ &\quad + \sum_{i=1}^{\infty} (i-1) (\log \frac{1-u}{1-v}) u(1-u)^{i-1} \\ &= \log \frac{u}{v} + (\log \frac{1-u}{1-v}) \frac{1-u}{u} = \frac{\text{KL}(u, v)}{u}. \end{aligned}$$

\square

Lemma 4: For any $\pi \in \Pi_i$, $i = 1, 2, 3$, and any $\beta > 0$ we have:

$$\liminf_{T \rightarrow \infty} \frac{S^\pi(T)}{\log(T)} \geq \beta \implies \mu_\theta(p^*) \liminf_{N \rightarrow \infty} \frac{R^\pi(N)}{\log(N)} \geq \beta.$$

APPENDIX B
PROOF OF LEMMA 4

Proof. Define $\mu^* = \mu_\theta(p^*)$ and $r_t = \sum_{n=1}^t (D^\pi(n) - D^*)$. Since $T \leq \sum_{n=1}^{N^\pi(T)+1} D^\pi(n)$ and $\mu^* = \frac{1}{D^*}$:

$$T\mu^* - N^\pi(T) \leq 1 + \sum_{n=1}^{N^\pi(T)+1} (\mu^* D^\pi(n) - 1) = 1 + \mu^* r_{N^\pi(T)+1}.$$

Since r_t is a submartingale, $N^\pi(T)$ is a stopping time and $N^\pi(T) \leq T$ a.s., Doob's optional stopping theorem gives:

$$\mathbb{E}(r_{N^\pi(T)+1}) \leq \mathbb{E}(r_{T+1}) = R^\pi(T+1).$$

Taking expectations above yields:

$$\frac{S^\pi(T)}{\log(T)} \leq \frac{1}{\log(T)} + \mu^* \frac{R^\pi(T+1)}{\log(T)},$$

and letting $T \rightarrow \infty$ proves the result since $\frac{\log(T)}{\log(T+1)} \rightarrow 1$. \square

APPENDIX C
PROOF OF LEMMA 1

A. Lower bound for $c_1(\theta)$

Let us first decompose the set $B_1(\theta)$. Observe that $\min_{p \in \mathcal{P}} D_\lambda(p) < D_\lambda(p^*)$, implies that *at least one* sub-optimal link i should have a higher success probability than the link $\zeta(i)$ under the parameter λ . Hence, we decompose $B_1(\theta)$ into sets where the link i is better than the link $\zeta(i)$ under parameter λ . For any $i \notin p^*$, define

$$A_i(\theta) = \left\{ \lambda : \{\lambda_j, j \in p^*\} = \{\theta_j, j \in p^*\}, \lambda_i > \theta_{\zeta(i)} \right\}.$$

Then, $B_1(\theta) = \bigcup_{i \neq \zeta(i)} A_i(\theta)$ and Eq. (1) reads

$$c_1(\theta) = \inf_{x \geq 0} \sum_{p \in \mathcal{P}} x_p \Delta_p$$

subject to: $\inf_{\lambda \in A_i(\theta)} \sum_{p \neq p^*} x_p I^p(\theta, \lambda) \geq 1, \forall i \notin p^*$.

Let $i \notin p^*$. Consider ϑ^i with $\vartheta_i^i = \theta_{\zeta(i)}$ and $\vartheta_j^i = \theta_j$ for $j \neq i$. Since $\vartheta^i \in A_i(\theta)$, we have

$$\begin{aligned} \inf_{\lambda \in A_i(\theta)} \sum_{p \neq p^*} x_p I^p(\theta, \lambda) &\leq \sum_{p \neq p^*} x_p I^p(\theta, \vartheta^i) \\ &\leq \max_{p \neq p^*} I^p(\theta, \vartheta^i) \sum_{p \neq p^*} x_p. \end{aligned}$$

Moreover, we have that

$$\begin{aligned} \sum_{p \in \mathcal{P}} x_p \Delta_p &= \sum_{p \in \mathcal{P}} x_p \sum_{i \in p} \left(\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}} \right) \\ &= \sum_{i \notin p^*} \left(\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}} \right) \sum_{p: i \in p} x_p. \end{aligned}$$

Putting these together yields

$$c_1(\theta) \geq \inf_{x \geq 0} \sum_{i \notin p^*} \left(\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}} \right) \sum_{p: i \in p} x_p$$

subject to: $(\max_{p: i \in p} I^p(\theta, \vartheta^i)) \sum_{p: i \in p} x_p \geq 1, \forall i \notin p^*$.

Introducing $z_i = \sum_{p: i \in p} x_p$ for any i , we rewrite the above problem as:

$$c_1(\theta) \geq \inf_{z \geq 0} \sum_{i \notin p^*} \left(\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}} \right) z_i$$

subject to: $z_i \geq (\max_{p: i \in p} I^p(\theta, \vartheta^i))^{-1}, \forall i \notin p^*$,

thus giving:

$$c_1(\theta) \geq \sum_{i \notin p^*} \frac{\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}}}{\max_{p: i \in p} I^p(\theta, \vartheta^i)},$$

where $I^p(\cdot, \cdot)$ is given by (3).

B. Derivation of $c_2(\theta)$

Let us first decompose the set $B_2(\theta)$. We argue that $\min_{p \in \mathcal{P}} D_\lambda(p) < D_\lambda(p^*)$ implies that *at least one* sub-optimal link i should have a higher success probability than the link $\zeta(i)$ under parameter λ .

We let $A_i(\theta)$ be the set where link i is better than the link $\zeta(i)$ under parameter λ :

$$A_i(\theta) = \left\{ \lambda : (\lambda_j = \theta_j, \forall j \in p^*), \lambda_i > \theta_{\zeta(i)} \right\}.$$

Hence, $B_2(\theta) = \bigcup_{i \notin p^*} A_i(\theta)$. Note $\text{KLG}(u, v) = 0$ if and only if $u = v$ and it is monotone increasing in v in the range $v > u$. Thus, for any $\lambda \in A_i(\theta)$, the infimum is obtained when $\lambda_i = \theta_{\zeta(i)}$ and $\lambda_j = \theta_j \forall j \neq i$, so that

$$\begin{aligned} \inf_{\lambda \in A_i(\theta)} \sum_{p \neq p^*} x_p \sum_{i \in p} \text{KLG}(\theta_i, \lambda_i) &\geq 1 \\ \iff \text{KLG}(\theta_i, \theta_{\zeta(i)}) \sum_{p: i \in p} x_p &\geq 1. \end{aligned}$$

Defining $z_i = \sum_{p: i \in p} x_p$ for any i and recalling that $\sum_{p \in \mathcal{P}} x_p \Delta_p = \sum_{i \notin p^*} \left(\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}} \right) \sum_{p: i \in p} x_p$, we rewrite problem (2) as

$$\begin{aligned} \inf_{z \geq 0} \sum_{i \notin p^*} \left(\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}} \right) z_i \\ \text{subject to: } \text{KLG}(\theta_i, \theta_{\zeta(i)}) z_i \geq 1, \forall i \notin p^*, \end{aligned}$$

which gives

$$c_2(\theta) = \sum_{i \notin p^*} \frac{\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}}}{\text{KLG}(\theta_i, \theta_{\zeta(i)})}$$

and concludes the proof. \square

APPENDIX D
PROOF OF PROPOSITION 4.4

Proof. Consider a problem instance with line topology in which $\theta_i = \alpha$ for all $i \notin p^*$, and $\theta_i = \alpha + \alpha^2$ for all $i \in p^*$ for some $\alpha \in (0, 0.36]$. Hence, $\theta_i < 0.5$ for all $i \in p^*$. For any uniformly good policy $\pi \in \Pi_2 \cup \Pi_3$, by Lemma 1 we have that:

$$\begin{aligned} \liminf_{N \rightarrow \infty} \frac{R^\pi(N)}{\log(N)} &\geq \sum_{i \notin p^*} \frac{1}{\text{KLG}(\theta_i, \theta_{\zeta(i)})} \left(\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}} \right) \\ &\geq \sum_{i \notin p^*} \frac{1}{2(\theta_{\zeta(i)} - \theta_i)} = \sum_{i \notin p^*} \frac{1}{2\theta_i \theta_{\zeta(i)} (\theta_i^{-1} - \theta_{\zeta(i)}^{-1})} \\ &= \frac{|E| - H}{2\alpha(\alpha + \alpha^2)(\alpha^{-1} - (\alpha + \alpha^2)^{-1})} \\ &= \frac{|E| - H}{2\alpha(\alpha + \alpha^2)\Delta_{\min}} \geq \frac{|E| - H}{4\alpha^2\Delta_{\min}} = \frac{|E| - H}{4\theta_{\min}^2\Delta_{\min}}, \end{aligned}$$

where in the second inequality we used Lemma 3 and $\text{KL}(u, v) \leq \frac{(u-v)^2}{v(1-v)} \leq \frac{2(u-v)^2}{v}$ for $v \leq 0.5$. This implies that the regret of any uniformly good policy $\pi \in \Pi_2 \cup \Pi_3$ for this problem instance is at least $\Omega\left(\frac{|E| - H}{\Delta_{\min}\theta_{\min}^2} \log(N)\right)$. \square

APPENDIX E
PROOF OF THEOREM 5.1

We first recall two results. Lemma 5 is a concentration inequality derived in [36, Theorem 2]. Lemma 6, proven in [37, Lemma 6], is a local version of Pinsker's inequality for the KL-divergence between two Bernoulli distributions.

Lemma 5: There exists a number $K_H > 0$ that only depends on H such that for all p and $n \geq 2$:

$$\mathbb{P}\left[\sum_{i \in p} t_i(n) \text{KL}(\hat{\theta}_i(n), \theta_i) \geq f_1(n)\right] \leq K_H n^{-1} (\log(n))^{-2}.$$

Lemma 6 ([37, Lemma 2]): For $0 \leq u < v \leq 1$ we have:

$$\text{KL}(u, v) \geq \frac{1}{2v}(u - v)^2.$$

Next we prove the theorem.

Statement (i): Let $p \in \mathcal{P}$, $n \in \mathbb{N}$, $t \in \mathbb{N}^{|E|}$, and $u, \lambda \in (0, 1]^{|E|}$ with $u_i \geq \lambda_i$ for all i . By Cauchy-Schwarz inequality we have:

$$\begin{aligned} p^\top \lambda^{-1} - p^\top u^{-1} &= \sum_{i \in p} \frac{u_i - \lambda_i}{u_i \lambda_i} = \sum_{i \in p} \frac{\sqrt{t_i}(u_i - \lambda_i)}{\sqrt{u_i}} \frac{1}{\lambda_i \sqrt{t_i} u_i} \\ &\leq \sqrt{\sum_{i \in p} \frac{t_i(u_i - \lambda_i)^2}{u_i}} \sqrt{\sum_{i \in p} \frac{1}{t_i u_i \lambda_i^2}} \\ &\leq \sqrt{\sum_{i \in p} \frac{t_i(u_i - \lambda_i)^2}{u_i}} \sqrt{\sum_{i \in p} \frac{1}{t_i \lambda_i^3}}, \end{aligned}$$

where we used $u_i \geq \lambda_i$ for all i in the last step. Using Lemma 6, it then follows that

$$p^\top \lambda^{-1} - p^\top u^{-1} \leq \sqrt{\sum_{i \in p} 2t_i \text{KL}(\lambda_i, u_i)} \sqrt{\sum_{i \in p} \frac{1}{t_i \lambda_i^3}}.$$

Thus, $\sum_{i \in p} t_i \text{KL}(\lambda_i, u_i) \leq f_1(n)$ implies:

$$p^\top \lambda^{-1} - p^\top u^{-1} \leq \sqrt{\sum_{i \in p} \frac{2f_1(n)}{t_i \lambda_i^3}},$$

or equivalently, $p^\top u^{-1} \geq c_p(n, \lambda, t)$. Hence, by definition of $b_p(n, \lambda, t)$, we have $b_p(n, \lambda, t) \geq c_p(n, \lambda, t)$.

Statement (ii): If $\sum_{i \in p} t_i(n) \text{KL}(\hat{\theta}_i(n), \theta_i) \leq f_1(n)$, then we have $b_p(n, \hat{\theta}(n), t(n)) \leq p^\top \theta^{-1}$ by definition of b_p . Therefore, using Lemma 5, there exists K_H such that for all $n \geq 2$ we have:

$$\begin{aligned} \mathbb{P}[b_p(n, \hat{\theta}(n), t(n)) > p^\top \theta^{-1}] &\leq \mathbb{P}\left[\sum_{i \in p} t_i(n) \text{KL}(\hat{\theta}_i(n), \theta_i) \geq f_1(n)\right] \\ &\leq K_H n^{-1} (\log(n))^{-2}, \end{aligned}$$

which concludes the proof. \square

APPENDIX F
PROOF OF THEOREM 5.3

A. Preliminary

Define $a = (1 - 2^{-\frac{1}{4}})$ and $\varepsilon = a \frac{\Delta_{\min}}{D^+} < a$. For $s \in \mathbb{N}^{|E|}$ and $p \in \mathcal{P}$ define $h(s) = \sum_{i \in p} \frac{1}{s_i}$. Define $s_i(n) = t_i(n) \hat{\theta}_i(n)$ the number of packets routed through link i before the n -th packet is sent and $s(n) = (s_i(n))_{i \in E}$. To ease notation define $h(n) = h(s(n))$. We will use the following technical lemma.

Lemma 7: Consider $S \subset \mathbb{N}$, $(s(n))_n$ an integer sequence such that $s(n) \neq s(n')$ for all $(n, n') \in S$, $n \neq n'$. Consider a constant $C > 0$, and a positive function δ , such that $\min_{n \in S} \delta(s(n)) \geq \delta_{\min}$. Then:

$$Z := \sum_{n \in S} \delta(s(n)) \mathbb{1}\{s(n) \leq C\delta(s(n))^{-2}\} \leq \frac{2C}{\delta_{\min}}.$$

Proof. If $s(n) \leq C\delta(s(n))^{-2}$, we have $\delta(s(n)) \leq \sqrt{C/s(n)}$, and $s(n) \leq C\delta_{\min}^{-2}$. So:

$$Z \leq \sum_{n \in S} \sum_{t=1}^{C\delta_{\min}^{-2}} \mathbb{1}\{s(n) = t\} \sqrt{\frac{C}{t}} \leq \sum_{t=1}^{C\delta_{\min}^{-2}} \sqrt{\frac{C}{t}},$$

using the fact that $\sum_{n \in S} \mathbb{1}\{s(n) = t\} \leq 1$. Using the inequality $\sum_{t=1}^T t^{-\frac{1}{2}} \leq \int_1^T t^{-\frac{1}{2}} dt \leq 2\sqrt{T}$ yields the result. \square

B. Proof of the Theorem

For any n , introduce the following events:

$$A_n = \left\{ \sum_{i \in p^*} t_i(n) \text{KL}(\hat{\theta}_i(n), \theta_i) > f_1(n) \right\},$$

$$B_{n,i} = \{p_i(n) = 1, |\hat{\theta}_i(n) - \theta_i| \geq \varepsilon \theta_i\}, \quad B_n = \bigcup_{i \in E} B_{n,i},$$

$$F_n = \{\Delta_{p(n)} \leq (1-a)^{-2} \theta_{\min}^{-1} \sqrt{2f_1(N)h(n)}\}.$$

We first prove that $p(n) \neq p^*$ implies: $n \in A_n \cup B_n \cup F_n$. Consider n such that $p(n) \neq p^*$ and $A_n \cap B_n$ does not occur. By design of the algorithm, $\xi_{p(n)}(n) \leq \xi_{p^*}(n)$, and $\xi_{p^*}(n) \leq D^*$ since A_n does not occur. By Theorem 5.1

we have $c_{p(n)}(n) \leq \xi_{p(n)}(n)$. Hence $c_{p(n)}(n) \leq D^*$. This implies:

$$p(n)^\top \hat{\theta}(n)^{-1} - \sqrt{\sum_{i \in p} \frac{2f_1(n)}{s_i(n)\hat{\theta}_i(n)^2}} \leq D^*,$$

so that:

$$\Delta_{p(n)} \leq p(n)^\top \theta^{-1} - p(n)^\top \hat{\theta}(n)^{-1} + \sqrt{\sum_{i \in p(n)} \frac{2f_1(n)}{s_i(n)\hat{\theta}_i(n)^2}}.$$

Since B_n does not occur $\hat{\theta}(n)^{-1} \geq \theta^{-1}/(1+\varepsilon)$ and:

$$\begin{aligned} p(n)^\top \theta^{-1} - p(n)^\top \hat{\theta}(n)^{-1} &\leq \frac{p(n)^\top \theta^{-1} \varepsilon}{(1+\varepsilon)} \leq D^+ \varepsilon \\ &= a \Delta_{\min} \leq a \Delta_{p(n)}. \end{aligned}$$

Also $\hat{\theta}_i(n) \geq \theta_{\min}(1-a)$, and $f_1(n) \leq f_1(N)$ so:

$$\sum_{i \in p(n)} \frac{2f_1(n)}{s_i(n)\hat{\theta}_i(n)^2} \leq \frac{2f_1(N)h(n)}{(1-a)^2 \theta_{\min}^2}.$$

Hence:

$$\Delta_{p(n)} \leq a \Delta_{p(n)} + \frac{\sqrt{2f_1(N)h(n)}}{(1-a)\theta_{\min}},$$

and $\Delta_{p(n)} \leq (1-a)^{-2} \theta_{\min}^{-1} \sqrt{2f_1(N)h(n)}$ and $n \in F_n$.

The regret $R^\pi(N)$ is upper bounded by:

$$\mathbb{E} \left(\sum_{n=1}^N \Delta_{p(n)} \right) \leq \mathbb{E} \left(\sum_{n=1}^N \Delta_{p(n)} (\mathbb{1}\{A_n\} + \mathbb{1}\{B_n\} + \mathbb{1}\{F_n\}) \right). \quad R^\pi(N) \leq \frac{2|E|\sqrt{HU}}{\Delta_{\min}} + D^+ \left(4K_H + 2\varepsilon^{-2} \sum_{i \in E} \theta_i^{-2} \right).$$

Set A: Using corollary 5.2, and $K_H \geq 1$ we have:

$$\sum_{n \geq 1} \mathbb{P}(A_n) \leq 1 + K_H \sum_{n \geq 2} n^{-1} (\log(n))^{-2} \leq 4K_H. \quad (8)$$

Set B: Define $\tau_i(n) = \sum_{n'=1}^n \mathbb{1}\{B_{n',i}\}$. Since $B_{n',i}$ implies $p_i(n') = 1$, we have $s_i(n) \geq \tau_i(n)$. Applying [34, Lemma B.1], we have $\sum_{n=1}^N \mathbb{P}(B_{n,i}) \leq 2(\varepsilon \theta_i)^{-2}$. A union bound yields:

$$\sum_{n=1}^N \mathbb{P}(B_n) \leq 2\varepsilon^{-2} \sum_{i \in E} \theta_i^{-2}. \quad (9)$$

Set F: Define $U = \frac{4f_1(N)}{(1-a)^4 \theta_{\min}^2}$. Define the set

$$S_n = \{i \in p(n) : s_i(n) \leq HU \Delta_{p(n)}^{-2}\}$$

and events:

$$\begin{aligned} G_n &= \{|S_n| \geq \sqrt{H}\}, \\ L_n &= \{|S_n| < \sqrt{H}, \min_{i \in p(n)} s_i(n) \leq \sqrt{HU} \Delta_{p(n)}^{-2}\}. \end{aligned}$$

Assume that neither G_n nor L_n occurs, then:

$$\begin{aligned} h(n) &= \sum_{i \in p(n), i \in S_n} \frac{1}{s_i(n)} + \sum_{i \in p(n), i \notin S_n} \frac{1}{s_i(n)} \\ &\leq \frac{|S_n| \Delta_{p(n)}^2}{\sqrt{HU}} + \frac{(H - |S_n|) \Delta_{p(n)}^2}{HU} < \frac{2\Delta_{p(n)}^2}{U}, \end{aligned}$$

since $|S_n| < \sqrt{H}$. Hence $\Delta_{p(n)}^2 > Uh(n)/2$ and F_n does not occur. So $F_n \subset G_n \cup L_n$. Further decompose G_n and L_n as:

$$\begin{aligned} G_{i,n} &= G_n \cap \{i \in p(n), s_i(n) \leq HU \Delta_{p(n)}^{-2}\}, \\ L_{i,n} &= L_n \cap \{i \in p(n), s_i(n) \leq \sqrt{HU} \Delta_{p(n)}^{-2}\}. \end{aligned}$$

Applying Lemma 7 twice, we get:

$$\sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\{G_{i,n}\} \leq \frac{HU}{\Delta_{\min}}, \quad \sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\{L_{i,n}\} \leq \frac{\sqrt{HU}}{\Delta_{\min}}.$$

We have

$$\sum_{i \in E} \mathbb{1}\{G_{i,n}\} = |S_n| \mathbb{1}\{G_n\} \geq \sqrt{H} \mathbb{1}\{G_n\}.$$

So:

$$\sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\{G_n\} \leq \frac{1}{\sqrt{H}} \sum_{n=1}^N \sum_{i \in E} \Delta_{p(n)} \mathbb{1}\{G_{i,n}\} \leq \frac{|E|\sqrt{HU}}{\Delta_{\min}}.$$

Further:

$$\sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\{L_n\} \leq \sum_{n=1}^N \sum_{i \in E} \Delta_{p(n)} \mathbb{1}\{L_{i,n}\} \leq \frac{|E|\sqrt{HU}}{\Delta_{\min}}.$$

Since $\mathbb{1}\{F_n\} \leq \mathbb{1}\{G_n\} + \mathbb{1}\{L_n\}$ we get:

$$\mathbb{E} \left(\sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\{F_n\} \right) \leq \frac{2|E|\sqrt{HU}}{\Delta_{\min}}. \quad (10)$$

Combining (8), (9) and (10) with $\Delta_{p(n)} \leq D^+$, yields the announced result:

APPENDIX G

PROOF OF THEOREM 5.4

The proof technique is similar to the analysis of [17, Theorem 5].

A. Preliminary

For $s \in \mathbb{N}^{|E|}$ and $p \in \mathcal{P}$ define $h'(s) = (\sum_{i \in p} \frac{1}{\sqrt{s_i}})^2$, and as before $s_i(n) = t_i(n)\hat{\theta}_i(n)$ and $s(n) = (s_i(n))_{i \in E}$, and $h'(n) = h'(s(n))$. We will use the following technical lemma.

Lemma 8: For all $n, t \in \mathbb{N}$, $\lambda \in (0, 1]$, and $i \in E$:

$$\omega_i(n, \lambda, t) \geq \frac{1}{\lambda} - \sqrt{\frac{2f_2(n)}{t\lambda^3}}.$$

Proof. Let $i \in E$, $n, t \in \mathbb{N}$ and $u, \lambda \in (0, 1]$ with $u \geq \lambda$. We have:

$$\frac{1}{\lambda} - \frac{1}{u} = \sqrt{\frac{t(u-\lambda)^2}{u}} \cdot \frac{1}{\sqrt{tu\lambda^2}} \leq \sqrt{2t\text{KL}(\lambda, u)} \cdot \frac{1}{\sqrt{t\lambda^3}},$$

where the second inequality follows from Lemma 6 and $u \geq \lambda$.

Hence, $t\text{KL}(\lambda, u) \leq f_2(n)$ implies: $\frac{1}{u} \geq \frac{1}{\lambda} - \sqrt{\frac{2f_2(n)}{t\lambda^3}}$. The above holds for all $u \in [\lambda, 1]$, and by definition of $\omega_i(n, \lambda, t)$:

$$\omega_i(n, \lambda, t) \geq \frac{1}{\lambda} - \sqrt{\frac{2f_2(n)}{t\lambda^3}}.$$

□

B. Proof of the theorem

For any n , we define the following events:

$$A_{n,i} = \left\{ t_i(n) \text{KL}(\hat{\theta}_i(n), \theta_i) > f_2(n) \right\}, \quad A_n = \bigcup_{i \in p^*} A_{n,i},$$

$$B_{n,i} = \{p_i(n) = 1, |\hat{\theta}_i(n) - \theta_i| \geq \varepsilon \theta_i\}, \quad B_n = \bigcup_{i \in E} B_{n,i},$$

$$F_n = \{\Delta_{p(n)} \leq (1-a)^{-2} \theta_{\min}^{-1} \sqrt{2f_2(N)h'(n)}\}.$$

We show that $p(n) \neq p^*$ implies: $n \in A_n \cup B_n \cup F_n$. Consider n such that $p(n) \neq p^*$ and $A_n \cup B_n$ does not occur. By design of the algorithm, $p(n)^\top \omega(n) \leq (p^*)^\top \omega(n)$, and $(p^*)^\top \omega(n) \leq D^*$ since A_n does not occur. Hence $p(n)^\top \omega(n) \leq D^*$. By Lemma 8, for all i :

$$\omega_i(n) \geq \frac{1}{\hat{\theta}_i(n)} - \sqrt{\frac{2f_2(n)}{s_i(n)\hat{\theta}_i(n)^2}}.$$

Summing over $i \in p(n)$ we get:

$$\Delta_{p(n)} \leq p(n)^\top \theta^{-1} - p(n)^\top \hat{\theta}(n)^{-1} + \sum_{i \in p(n)} \sqrt{\frac{2f_2(n)}{s_i(n)\hat{\theta}_i(n)^2}}.$$

As before, when B_n does not occur we have

$$p(n)^\top \theta^{-1} - p(n)^\top \hat{\theta}(n)^{-1} \leq a \Delta_{p(n)}.$$

Furthermore $\hat{\theta}_i(n) \geq \theta_{\min}(1-a)$ and $f_2(n) \leq f_2(N)$ so that:

$$\sum_{i \in p(n)} \sqrt{\frac{2f_2(n)}{s_i(n)\hat{\theta}_i(n)^2}} \leq \sum_{i \in p(n)} \sqrt{\frac{f_2(N)}{s_i(n)\theta_{\min}^2(1-a)^2}},$$

Hence:

$$\Delta_{p(n)} \leq a \Delta_{p(n)} + \frac{\sqrt{2f_2(N)h'(n)}}{(1-a)\theta_{\min}}$$

and $\Delta_{p(n)} \leq (1-a)^{-2} \theta_{\min}^{-1} \sqrt{2f_2(N)h'(n)}$ so that $n \in F_n$.

The regret $R^\pi(N)$ is upper bounded by:

$$\mathbb{E}\left(\sum_{n=1}^N \Delta_{p(n)}\right) \leq \mathbb{E}\left(\sum_{n=1}^N \Delta_{p(n)} (\mathbb{1}\{A_n\} + \mathbb{1}\{B_n\} + \mathbb{1}\{F_n\})\right).$$

Set A: By [38, Theorem 10] and a union bound:

$$\mathbb{P}(A_n) \leq \sum_{i \in p^*} \mathbb{P}(A_{n,i}) \leq H[f_2(n) \log(n)] e^{1-f_2(n)}.$$

Hence:

$$\sum_{n=1}^N \mathbb{P}(A_n) \leq H \left(1 + e \sum_{n \geq 2} [f_2(n) \log(n)] e^{-f_2(n)}\right) \leq 8|H|. \quad (11)$$

Set B: As in the proof of Theorem 5.3:

$$\sum_{n=1}^N \mathbb{P}(B_n) \leq 2\varepsilon^{-2} \sum_{i \in E} \theta_i^{-2}. \quad (12)$$

Set F: Define $U' = 2H^2 f_2(N)(1-a)^{-4} \theta_{\min}^{-2}$. Similarly to the proof of [17, Theorem 5], consider $\alpha, \beta > 0$, for $\ell \in \mathbb{N}$ define $\alpha_\ell = \left(\frac{1-\beta}{\sqrt{\alpha-\beta}}\right)^2 \alpha^\ell$ and $\beta_\ell = \beta^\ell$. Introduce set $S_{\ell,n}$ and events $G_{\ell,n}$:

$$S_{\ell,n} = \{i \in p(n), s_i(n) \leq U' \alpha_\ell \Delta_{p(n)}^{-2}\},$$

$$G_{\ell,n} = \{|S_{\ell,n}| \geq \beta_\ell H\} \cap \{|S_{j,n}| < \beta_j H, j = 1, \dots, \ell-1\}.$$

If $\overline{\cup_{\ell \geq 1} G_{\ell,n}} = \{|S_{\ell,n}| < H\beta_\ell, \ell \geq 1\}$ then:

$$\begin{aligned} \sum_{\ell \geq 1} \frac{|S_{\ell-1,n}| - |S_{\ell,n}|}{\sqrt{\alpha_\ell}} &= \frac{|S_{0,n}|}{\sqrt{\alpha_1}} + \sum_{\ell \geq 1} |S_{\ell,n}| \left(\frac{1}{\sqrt{\alpha_{\ell+1}}} - \frac{1}{\sqrt{\alpha_\ell}} \right) \\ &< \frac{H\beta_0}{\sqrt{\alpha_1}} + \sum_{\ell \geq 1} H\beta_\ell \left(\frac{1}{\sqrt{\alpha_{\ell+1}}} - \frac{1}{\sqrt{\alpha_\ell}} \right) \\ &= H \sum_{\ell \geq 1} \frac{\beta_\ell - \beta_{\ell-1}}{\sqrt{\alpha_\ell}} \leq H, \end{aligned}$$

since $\frac{1}{\sqrt{\alpha_{\ell+1}}} - \frac{1}{\sqrt{\alpha_\ell}} \geq 0$. Now:

$$|\{i : s_i(n) \in U' \Delta_{p(n)}^{-2} [\alpha_\ell, \alpha_{\ell-1}]\}| = |S_{\ell-1,n}| - |S_{\ell,n}|$$

so that:

$$\sqrt{h'(n)} \leq \sum_{\ell \geq 1} \frac{(|S_{\ell-1,n}| - |S_{\ell,n}|) \Delta_{p(n)}}{\sqrt{\alpha_\ell} \sqrt{U'}} < H \frac{\Delta_{p(n)}}{\sqrt{U'}}.$$

Hence $\Delta_{p(n)}^2 > h'(n)U'H^{-2}$, and F_n does not occur. Therefore $F_n \subset \cup_{\ell \geq 1} G_{\ell,n}$ and:

$$\sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\{F_n\} \leq \sum_{n=1}^N \sum_{\ell \geq 1} \Delta_{p(n)} \mathbb{1}\{G_{\ell,n}\}.$$

Further decompose $G_{i,\ell}$ as:

$$G_{i,\ell,n} = G_{\ell,n} \cap \{i \in p(n), s_i(n) \leq U' \alpha_\ell \Delta_{p(n)}^{-2}\}.$$

Observe that:

$$\mathbb{1}\{G_{\ell,n}\} \leq \frac{|S_{\ell,n}|}{H\beta_\ell} \mathbb{1}\{G_{\ell,n}\} = \frac{1}{H\beta_\ell} \sum_{i \in E} \mathbb{1}\{G_{i,\ell,n}\}.$$

Applying Lemma 7, we get:

$$\begin{aligned} \sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\{G_{i,\ell,n}\} &\leq \sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\left\{s_i(n) \leq \frac{U' \alpha_\ell}{\Delta_{p(n)}^2}\right\} \\ &\leq \frac{2U' \alpha_\ell}{\Delta_{\min}}. \end{aligned}$$

Putting it together:

$$\sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\{F_n\} \leq \frac{2|E|U'}{H\Delta_{\min}} \sum_{\ell \geq 1} \frac{\alpha_\ell}{\beta_\ell} \leq \frac{90|E|U'}{H\Delta_{\min}}, \quad (13)$$

by choosing $\alpha = 0.15$ and $\beta = 0.24$ so that $\sum_{\ell \geq 1} \frac{\alpha_\ell}{\beta_\ell} \leq 45$. Combining (11), (12) and (13) with $\Delta_{p(n)} \leq D^+$, yields the result:

$$R^\pi(N) \leq \frac{90|E|U'}{H\Delta_{\min}} + D^+ \left(8H + 2\varepsilon^{-2} \sum_{i \in E} \theta_i^{-2}\right).$$

□

APPENDIX H
PROOF OF PROPOSITION 5.5

In the line network, KL-SR simply chooses the link with the smallest index on each hop. Hence, on each hop, KL-SR is equivalent to the KL-UCB algorithm for a classical MAB with geometrically distributed rewards. By [38, Theorem 1 and Lemma 6], the regret of KL-SR on the m -th hop asymptotically grows as:

$$\sum_{i \in E_m \setminus p^*} \frac{\log(N)}{\text{KLG}(\theta_i, \theta_{\zeta(i)})} \left(\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}} \right),$$

where E_m denotes the set of links in the m -th hop. Since decisions at various hops are decoupled, the regret due to all hops satisfies

$$\begin{aligned} \limsup_{N \rightarrow \infty} \frac{R^{\text{KL-SR}}(N)}{\log(N)} &\leq \sum_{m=1}^H \sum_{i \in E_m \setminus p^*} \frac{\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}}}{\text{KLG}(\theta_i, \theta_{\zeta(i)})} \\ &= \sum_{i \notin p^*} \frac{\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}}}{\text{KLG}(\theta_i, \theta_{\zeta(i)})} = c_2(\theta). \end{aligned}$$

Furthermore, using Lemma 3 and Lemma 6 we have for any $i \notin p^*$:

$$\frac{\frac{1}{\theta_i} - \frac{1}{\theta_{\zeta(i)}}}{\text{KLG}(\theta_i, \theta_{\zeta(i)})} = \frac{\theta_{\zeta(i)} - \theta_i}{\theta_{\zeta(i)} \text{KL}(\theta_i, \theta_{\zeta(i)})} \leq \frac{2}{\theta_{\zeta(i)} - \theta_i}.$$

Moreover, in line networks $\Delta_{\min} = \min_{i \notin p^*} (\theta_i^{-1} - \theta_{\zeta(i)}^{-1})$. Thus,

$$\begin{aligned} c_2(\theta) &\leq \sum_{i \notin p^*} \frac{2}{\theta_{\zeta(i)} - \theta_i} = \sum_{i \notin p^*} \frac{2}{\theta_i \theta_{\zeta(i)} (\theta_i^{-1} - \theta_{\zeta(i)}^{-1})} \\ &\leq \frac{|E| - H}{\Delta_{\min}} \cdot \frac{2}{\min_{i \notin p^*} \theta_i \theta_{\zeta(i)}} \leq \frac{2(|E| - H)}{\Delta_{\min} \theta_{\min}^2}, \end{aligned}$$

which completes the proof. \square

APPENDIX I
PROOF OF PROPOSITION 5.6

The proof is similar to that of [18, Theorem 4]. Note that if $i \notin I_p(\lambda)$, then the optimal solution satisfies $u_i = 1$ since $\text{KL}(1, v) = \infty$ unless $v = 1$. Thus, if $I_p(\lambda) = \emptyset$, then $u_i = 1, \forall i \in E$, and $b_p(n, \lambda, t) = \sum_{i \in p} p_i$.

If $I_p(\lambda) \neq \emptyset$, let $i \in I_p(\lambda)$. Computing b_p involves solving a convex optimization problem with one inequality constraint which must hold with equality since $u_i \mapsto \text{KL}(\lambda_i, u_i)$ is monotone increasing for $u_i \geq \lambda_i$. Since $\frac{d}{du_i} \text{KL}(\lambda_i, u_i) = \frac{u - \lambda}{u(1-u)}$, the Karush-Kuhn-Tucker conditions are:

$$\begin{aligned} \frac{1}{u_i^2} - \gamma t_i \frac{u_i - \lambda_i}{u_i(1-u_i)} &= 0, \\ \sum_{i \in I_p(\lambda)} t_i \text{KL}(\lambda_i, u_i) - f_1(n) &= 0. \end{aligned}$$

with $\gamma > 0$ the Lagrange multiplier. The first equation is the quadratic equation:

$$u_i^2 + u_i \left(\frac{1}{\gamma t_i} - \lambda_i \right) - \frac{1}{\gamma t_i} = 0.$$

Solving for u_i , we obtain $u_i(\gamma) = g(\gamma, \lambda_i, t_i)$ and replacing in the second equation, we obtain $F(\gamma, n, \lambda, t) = f_1(n)$. The results then follow directly. \square

APPENDIX J
REGRET UPPER BOUND FOR CUCB

CUCB (see [9]) uses the following link index:

$$\gamma_i(n) = \frac{1}{\hat{\theta}_i(n) + \sqrt{1.5 \log(n)/t_i(n)}}, \quad \forall i \in E$$

Define $a = (1 - 2^{-\frac{1}{2}})$ and $\varepsilon = a \frac{\Delta_{\min}}{D^*} < a$. For any $s \in \mathbb{N}^{|E|}$ and $p \in \mathcal{P}$ define $h'(s) = (\sum_{i \in p} \frac{1}{\sqrt{s_i}})^2$, and as in the proof of Theorem 5.4, $s_i(n) = t_i(n) \hat{\theta}_i(n)$ and $s(n) = (s_i(n))_{i \in E}$, and $h'(n) = h'(s(n))$. We have that:

$$\begin{aligned} p(n)^\top \gamma(n) &= \sum_{i \in p(n)} \frac{1}{\hat{\theta}_i(n) + \sqrt{1.5 \hat{\theta}_i(n) \log(n)/s_i(n)}} \\ &= \sum_{i \in p(n)} \frac{1}{\hat{\theta}_i(n)} - \sum_{i \in p(n)} \frac{\sqrt{1.5 \log(n)/(s_i(n) \hat{\theta}_i(n)^3)}}{1 + \hat{\theta}_i(n)^{-\frac{1}{2}} \sqrt{1.5 \log(n)/s_i(n)}} \\ &\geq p(n)^\top \hat{\theta}(n)^{-1} - \sum_{i \in p(n)} \sqrt{\frac{1.5 \log(n)}{s_i(n) \hat{\theta}_i(n)^3}}. \end{aligned} \quad (14)$$

For any n , introduce the following events:

$$A_{n,i} = \left\{ |\hat{\theta}_i(n) - \theta_i| > \sqrt{1.5 \log(n)/t_i(n)} \right\}, \quad A_n = \bigcup_{i \in p^*} A_{n,i},$$

$$B_{n,i} = \{p_i(n) = 1, |\hat{\theta}_i(n) - \theta_i| \geq \varepsilon \theta_i\}, \quad B_n = \bigcup_{i \in E} B_{n,i},$$

$$F_n = \{\Delta_{p(n)} \leq (1-a)^{-\frac{5}{2}} \theta_{\min}^{-\frac{3}{2}} \sqrt{2 \log(N) h'(n)}\}.$$

We show that if $p(n) \neq p^*$ then $A_n \cup B_n \cup F_n$ occurs. Consider n such that $p(n) \neq p^*$ and $A_n \cup B_n$ does not occur. By design of the algorithm, $p(n)^\top \gamma(n) \leq (p^*)^\top \gamma(n)$, and $(p^*)^\top \gamma(n) \leq D^*$ since A_n does not occur. Hence $p(n)^\top \gamma(n) \leq D^*$. \square

When B_n does not occur, $(1-a)\theta_{\min} \leq \hat{\theta}_i(n) \leq (1+\varepsilon)\theta_i$ and $p(n)^\top \theta^{-1} - p(n)^\top \hat{\theta}(n)^{-1} \leq a \Delta_{p(n)}$. Hence, using (14), we get

$$\begin{aligned} \Delta_{p(n)} &= p(n)^\top \theta^{-1} - D^* \leq p(n)^\top \theta^{-1} - p(n)^\top \gamma(n) \\ &\leq a \Delta_{p(n)} + (1-a)^{-\frac{3}{2}} \theta_{\min}^{-\frac{3}{2}} \sqrt{1.5 \log(N) h'(n)} \end{aligned}$$

so that

$$\Delta_{p(n)} \leq (1-a)^{-\frac{5}{2}} \theta_{\min}^{-\frac{3}{2}} \sqrt{1.5 \log(N) h'(n)}$$

and thus $n \in F_n$.

The regret $R^\pi(N)$ is upper bounded by:

$$\mathbb{E} \left(\sum_{n=1}^N \Delta_{p(n)} \right) \leq \mathbb{E} \left(\sum_{n=1}^N \Delta_{p(n)} (\mathbb{1}\{A_n\} + \mathbb{1}\{B_n\} + \mathbb{1}\{F_n\}) \right).$$

Set A: Using a Chernoff bound and a union bound, we have that $\mathbb{P}(A_n) \leq 2Hn^{-2}$ (see, e.g., [9, Lemma 3]). Hence

$$\sum_{n=1}^N \mathbb{P}(A_n) \leq \sum_{n=1}^N \frac{2H}{n^2} \leq \frac{2\pi^2 H}{3}. \quad (15)$$

Set B: As in the proof of Theorem 5.3:

$$\sum_{n=1}^N \mathbb{P}(B_n) \leq 2\epsilon^{-2} \sum_{i \in E} \theta_i^{-2}. \quad (16)$$

Set F: Define $U' = 2H^2 f_2(N)(1-a)^{-\frac{5}{2}} \theta_{\min}^{-3}$. By the same technique as the proof of Theorem 5.4 we get

$$\sum_{n=1}^N \Delta_{p(n)} \mathbb{1}\{F_n\} \leq \frac{278H|E|\log(N)}{\Delta_{\min} \theta_{\min}^3}. \quad (17)$$

Putting (15), (16), and (17) together, we obtain

$$R^\pi(N) \leq \frac{278H|E|\log(N)}{\Delta_{\min} \theta_{\min}^3} + 2D^+ \left(\frac{\pi^2 H}{3} + \sum_{i \in E} \frac{1}{(\epsilon \theta_i)^2} \right).$$

□



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