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# An exact method for solving the two-machine flow-shop problem with time delays

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## 1 Introduction

We address the flow-shop scheduling problem with two machines and time delays, which is denoted by  $F2|l_j|C_{max}$ .  $F2|l_j|C_{max}$  can be described as follows. Given a set  $J = \{1, 2, \dots, n\}$  of  $n$  jobs and two machines  $(M_1, M_2)$ , each job  $j \in J$  has two operations  $O_{1,j}$  and  $O_{2,j}$ . The operation  $O_{1,j}$  (resp.  $O_{2,j}$ ) must be executed without preemption during  $p_{1,j}$  (resp.  $p_{2,j}$ ) time units on  $M_1$  (resp.  $M_2$ ). In addition to the resource constraints of the machines and the non-preemption of jobs, a schedule  $S$  is considered feasible if, for each job  $j \in J$ , a time delay of at least  $l_j$  time units must elapse between the completion of  $O_{1,j}$  and the start of  $O_{2,j}$ . The goal is to find a feasible schedule  $S$  that minimizes the makespan, i.e, the maximum job completion time.

The  $F2|l_j|C_{max}$  problem is NP-hard in the strong sense even with unit-time operations [5]. Therefore, it has been the scope of a variety of investigations. As far we know, [2] and [4] proposed lower bound methods. Moreover, [2] investigated heuristic approaches where he introduced four constructive heuristics and a Tabu Search algorithm. It should be noted that [2] implemented an exact method based on the branch-and-bound method of [1], which is originally made for the job-shop problem. Another branch-and-bound method was proposed by [3] for the unit-time operations case.

The objective of this paper consists in introducing a branch-and-bound algorithm for  $F2|l_j|C_{max}$ . To our aim, we propose a set of lower bounds and an upper bound. In addition, a set of dominance rules is proposed in order to reduce the search space.

## 2 The Branch-and-Bound Algorithm

In this section, we present an exact method for  $F2|l_j|C_{max}$  based on a branch-and-bound enumeration scheme. At first, let us introduce the following observation.

**Observation 1** *Let  $\sigma$  be a fixed job sequence on  $M_1$  of all jobs, then schedules in which the job sequence  $\sigma$  is fixed first on  $M_1$  and the jobs are executed on  $M_2$  according to the nondecreasing order of their arrival times are dominant.*

According to this observation, our branch-and-bound enumerates job sequences on  $M_1$  as follows. At a given node  $N_{\sigma_1}$  of the search tree of the branch-and-bound, a partial

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job sequence  $\sigma_1$  of  $|\sigma_1|$  jobs is fixed on  $M_1$ . In order to reduce the number of explored nodes, we invoke at each node  $N_{\sigma_1}$  the following features:

- A preprocessing procedure that aims to fix some precedence relationships between jobs on each machine.
- Several lower bounds of [2] and [4] that have been adapted to be used inside the branch-and-bound search tree.
- An upper bound based on a local search technique applied on the current sub-sequence.
- Three dominance rules that allow to discard nodes from additional expansions in order to reduce the computational time of the proposed branch-and-bound algorithm.

If these features fail to prune the current node  $N_{\sigma_1}$ , then for each unscheduled job  $j$  a *son node* is created from  $N_{\sigma_1}$ , in which  $j$  is fixed after  $\sigma_1$ . Note that we adopted the depth-first node selection strategy.

### 3 Computational results

A computer simulation of the branch-and-bound algorithm, which was carried out on a set of 360 instances of [2], shows that our branch-and-bound method outperforms the state of the art exact method. In particular, we manage to solve 358 instances among 360 possible ones.

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