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Melodic Pattern Segmentation of Polyphonic Music as a Set Partitioning Problem

Tsubasa Tanaka and Koichi Fujii

Abstract In polyphonic music, melodic patterns (motifs) are frequently imitated or repeated, and transformed versions of motifs such as inversion, retrograde, augmentations, diminutions often appear. Assuming that economical efficiency of reusing motifs is a fundamental principle of polyphonic music, we propose a new method of analyzing a polyphonic piece that economically divides it into a small number of types of motif. To realize this, we take an integer programming-based approach and formalize this problem as a set partitioning problem, a well-known optimization problem. This analysis is helpful for understanding the roles of motifs and the global structure of a polyphonic piece.

1 Motif Division

In polyphonic music like fugue-style pieces or J.S. Bach's *Inventions and Sinfonias*, melodic patterns (motifs) are frequently imitated or repeated. Although some motifs are easy to find, others are not. This is because they often appear implicitly and/or appear in the transformed versions such as inversion, retrograde, augmentations, diminutions. Therefore, motif analysis is useful to understand how polyphonic music is composed.

Simply speaking, we can consider the motifs that appear in a musical piece to be economical if the number of types of motif is small, the numbers of repetitions are large, and the lengths of the motifs are long. Assuming that this economical efficiency of motifs is a fundamental principle of polyphonic music, we propose a new method of analyzing a polyphonic piece that efficiently divides it into a small number of types of motif. Using this division, the whole piece is reconstructed with

Tsubasa Tanaka
IRCAM and Tokyo University of the Arts, e-mail: tsubasa.tanaka@ircam.fr

Koichi Fujii
NTT DATA Mathematical Systems Inc. e-mail: fujii@msi.co.jp

a small number of types of motif like the puzzle game *Tetris*[1] (In tetris, certain domains are divided with only seven types of piece). We call such a segmentation a *motif division*.

If a motif division is accomplished, it provide us a simple and higher-level representation whose atom is a motif, not a note, and it will be helpful to clarify the structures of polyphonic music. The representation may provide knowledge about how frequent and where each motif is used, the relationships between motifs such as causality and co-occurrence, which transformations are used, how the musical form is constructed by motifs, and how the long-term musical expectations are formed. This analysis may be useful for applications such as systems of music analysis, performance, and composition.

Studies about finding boundaries of melodic phrases are often based on human cognition. For example, [2] is based on grouping principles of gestalt psychology, and [3] is based on a short-term memory model. While these studies deal with relatively short range of perception and require small amounts of computational time, we focus on global configuration of motifs on the level of compositional planning. This requires us to solve an optimization problem that is hard to solve. To deal with this difficulty, we take an *integer programming*-based approach [4] and show that this problem can be formalize as a *set partitioning problem* [5]. This problem can be solved by integer programming solvers that use efficient algorithms such as the branch and bound method.

2 Transformation Group and Equivalence Classes of Motif

In this section, we introduce *equivalence classes* of motif derived from a group of motif transformations as the criterion of identicalness of motifs. These equivalence classes are used to formulate the motif division in Sect. 3.

Firstly, a motif is defined as an ordered correction of notes $[N_1, N_2, \dots, N_k]$ ($k > 0$), where N_i is the information for the i th note, comprising the combination of the pitch p_i , start position s_i , and end position e_i ($N_i = (p_i, s_i, e_i)$, $s_i < e_i \leq s_{i+1}$). Next, let \mathcal{M} be the set of every possible motif, and let T_p , S_t , R , I , A_r be one-to-one mappings (transformations) from \mathcal{M} to \mathcal{M} , where T_p is the transposition by pitch interval p , S_t is the shift by time interval t ($p, t \in \mathbb{R}$), R is the retrograde, I is the inversion, and A_r ($r > 0$) is the r -fold argumentation (diminution, in the case of $0 < r < 1$). These transformations generate a transformation group \mathcal{T} whose operation is the composition of two transformations and whose identity element is the transformation that does nothing. Each transformation in \mathcal{T} is a *strict imitation* that preserves the internal structures of the motifs.

Here, a binary relation between a motif m ($\in \mathcal{M}$) and $\tau(m)$ ($\tau \in \mathcal{T}$) can be defined. Due to the group structure of \mathcal{T} , this relation is an equivalence relation (i.e., it satisfies reflexivity, symmetry and transitivity[6]). Then, it derives equivalence

classes in \mathcal{M} . Because the motifs that belong to a same equivalence class share the same internal structure, they can be regarded as identical (or the same type)¹.

3 Formulation as a Set Partitioning Problem

A set partitioning problem, which is well known in the context of operations research, is an optimization problem defined as follows. Let N be a set that consists of n elements $\{N_1, N_2, \dots, N_n\}$, and let M be a family of sets $\{M_1, M_2, \dots, M_m\}$, where each M_j is a subset of N . If $\bigcup_{j \in X} M_j = N$ is satisfied, X , a subset of indexes of M , is called a cover, and the cover X is called a partition if $M_{j_1} \cap M_{j_2} = \emptyset$ is satisfied for different $j_1, j_2 \in X$. If a constant c_j called a cost is defined for each M_j , the problem of finding a partition X that minimizes the sum of the costs $\sum_{j \in X} c_j$ is called a set partitioning problem.

3.1 Condition of motif division

If N_i corresponds to each note of a musical piece to be analyzed and M_j corresponds to a motif, the problem of finding the most economically efficient motif division can be interpreted as a set partitioning problem. The index i starts from the first note of a voice to the last note of the voice and from the first voice to the last voice. $M_j (1 \leq j \leq m)$ corresponds to $[N_1]$, $[N_1, N_2]$, $[N_1, N_2, N_3]$, \dots , $[N_2]$, $[N_2, N_3]$, \dots in this order. The number of notes in a motif is less than a certain limit number (Fig. 1).

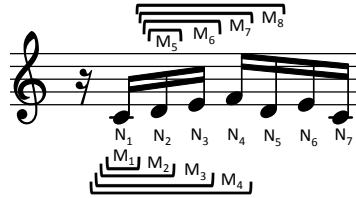


Fig. 1 Possible motifs of the first voice of J.S. Bach's *Invention No. 1*. (In the case where the maximum number of notes in a motif is 4.)

This information can be represented by the following matrix A :

¹ Although the criterion for identical motifs defined here only deals with strict imitations, we can define the criterion in different ways to allow more flexible imitations, such as by (1) defining an equivalence relation from the equality of a shape type [7, 8, 9, 10] and (2) defining a similarity measure and performing a clustering of motifs using methods such as k-medoid method [11] (the resulting clusters derive an equivalence relation). In any case, making equivalence classes from a certain equivalence relation is a versatile way to define the identicalness of the motifs.

$$A = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \dots \\ 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & \dots \\ 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & \dots \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 & \dots \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & \dots \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix} \quad (1)$$

where each row corresponds to each note N_i and each column corresponds to which notes are covered by each motif M_j . This matrix is the case where the maximum number of notes in a motif is 4. Representing the element of A as a_{ij} , the condition that the whole piece is exactly divided by a set of selected motifs can be described by the following constraints, which mean that each note N_i is covered by one of M_j once and only once:

$$\forall i \in \{1, 2, \dots, n\}, \sum_{j=1}^m a_{ij} x_j = 1, \quad (2)$$

where x_j is a 0-1 variable that represents whether or not M_j is used in the motif division. These conditions are equivalent to the condition of partitioning.

3.2 Objective Function

The purpose of motif division is to find the most efficient solution from the many solutions that satisfy the condition of partitioning. Then, we must define efficiency of motif division. We can consider that the average length (the number of notes) of motifs used in the motif division is one of the simplest barometers that represent the efficiency of motif division. Also, the number of motifs and that of the types of motif used in motif division will be efficient if they are small. In fact, the average length of motifs is inversely proportional to the number of motifs. Therefore, if the number of types of motif (denoted by P) is fixed, the number of motifs will be what we should minimize.

The number of motifs can be simply represented by $\sum_{j=1}^m x_j$. This is the cost function $\sum_{j=1}^m c_j x_j$ whose c_j is 1 for each j . We adopt this cost function. However, in the next subsection, we introduce additional variables and constraints to fix the number P .

3.3 Controlling the Number of Equivalence Classes

Let C be the set of equivalence classes of motif, which is derived from M , which is the set of all possible motif classes that can be found in a piece (only the motif classes whose number of notes is less than a certain number is included in M).

This means that M is derived from \mathcal{M} by a restriction. Let y_k be a 0-1 variable that represents whether or not one of the members of C_k appears in X (the set of selected motifs), where each element of C is denoted as $C_k (1 \leq k \leq l)$. This means that statement “ $y_k = 1 \Leftrightarrow \sum_{j \in C_k} x_j > 0$ ” must be satisfied. This statement can be represented by the following constraints that use $\sum_{j \in C_k} x_j$, the number of selected motifs that belong to C_k :

$$\forall k \in \{1, 2, \dots, l\}, y_k \leq \sum_{j \in C_k} x_j \leq Qy_k, \quad (3)$$

where Q is a constant that is sufficiently large. Then, the statement that the number of equivalence classes is P can be represented by the following constraint:

$$\sum_{k=1}^l y_k = P. \quad (4)$$

If P is small to a certain degree, the motif division will tend to be simple. However, if P is too small, covering whole piece with few motif classes will be difficult and one note motif will be used too many times. This will lead to a loss of the efficiency of motif division. Therefore, we should find good balance between the smallness of the objective function and the smallness of P . Because knowing which number is adequate for P in advance is difficult, we will solve the optimization problems for respective P in a certain range. Then, we will determine an adequate number for P , observing the solutions for respective P .

4 Result

We analyzed J.S. Bach’s *Invention No. 1* by solving the optimization problem described in the previous section. The maximum length of motif was set as 7. An IP solver *Numerical Optimizer 16.1.0.* and a branch and bound method was used for searching the solution. From the observation of solutions for various values for P , P was set as 13. It took less than one minute to obtain a solution for $P = 13$. Fig. 3 shows the result of motif division. The slurs represent the motifs and the one-note motifs don’t have a slur. Fig. 3 shows the representatives of 13 motif classes that are used in the motif division.

This result tells us many things. For example, 4th, 10th, and 11th motif classes in Fig. 3 are slightly different but can be regarded as the same motif, which corresponds to the subject of this piece. Searching for the domains where the subject doesn’t appear, we find that there are three domains whose durations are one and half bars (These are indicated by the big rectangles). The ends of these domains coincide with the places where the cadences exist. Therefore, we could detect three sections of this piece properly.

Fig. 2 The representatives of motif classes that appear in the motif division of J.S. Bach's *Invention No. 1*, in the case that $P = 13$. Some flats are replaced by sharps for the purpose of programming.

The last motif class in Fig. 3 is a leap of octave. This motif class appears in all of the cadence domains and is related to the ends of sections. It also co-occurs with 2nd motif class, which is a two-note motif, in the cadence domains.



Fig. 3 The motif classes that appear in the motif division of J.S. Bach's *Invention No. 1* (the number of motif classes was set at 13.).

The 12th motif class is a very characteristic one that includes a dotted note and a large leap. This motif class only appears before the cadence domains (the two motifs surrounded by the rounded rectangle). We can consider that this remarkable motif class plays an important role that tells listeners the end of the exposition of subject and the beginning of the cadence domain.

The 9th zigzag motif class and the motif classes that are one-way slow movements shown by the arrows in Fig. 3 only appear as the ascending form in the first 2 sections. In contrast, these motif classes appear only as the descending form in the final section. We interpret this contrast means that the ascending form creates a sense of continuation of the piece and the descending form creates a sense of conclusion. Thus, long-term musical expectations seems to be formed by the selections of transformation.

In such ways, motif division is useful to make us understand the roles of motifs and how global musical structures are formed.

5 Conclusion

In this paper, we formulated the problem of motif division, which decompose polyphonic music into a small number of motif classes, as a set partitioning problem, and we obtained the solution using an IP solver. It was shown that the motif division provides useful information to understand the roles of motifs and how global musical structures are constructed from the motifs.

Future tasks include construction of a program that automatically analyzes global structures utilizing the obtained motifs and automatic composition of new pieces that use the same motifs as the original piece using the result of the analysis program. To create a criterion for determining adequate value of P automatically is also a remaining problem.

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