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Number field lattices achieve Gaussian and Rayleigh channel capacity within a constant gap

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Abstract—This paper shows that a family of number field lattice codes simultaneously achieves a constant gap to capacity in Rayleigh fast fading and Gaussian channels. The key property in the proof is the existence of infinite towers of Hilbert class fields with bounded root discriminant. The gap to capacity of the proposed lattice codes is determined by the root discriminant. The comparison between the Gaussian and fading case reveals that in Rayleigh fading channels the *normalized minimum product distance* plays an analogous role to the Hermite invariant in Gaussian channels.

Index Terms—Lattice coding, Gaussian channel, fast fading channels, number theory

I. INTRODUCTION

The classical problem of achieving the capacity of the Gaussian channel using structured codes has seen significant recent advances. In particular, random lattice code ensembles have been shown to attain capacity [1, 2]. Good lattice code ensembles can be constructed by lifting linear codes over finite fields [4, 5] or using multilevel codes [6]; an explicit multilevel construction from polar codes was recently proposed in [7]. In this paper, we consider an alternative approach based on algebraic number theory. It is well-known that lattice constellations from number fields provide good performance on Gaussian and fading channels [8, 9]. As far as we know, the problem of achieving ergodic capacity with structured codes is still open in the case of fading channels.

In this work, we analyze the asymptotic behavior of algebraic lattices from number fields when the lattice dimension tends to infinity, and show that Hilbert class field towers with bounded root discriminants simultaneously reach a constant gap to capacity on both Gaussian and Rayleigh fading channels.

We note that the constant gap to capacity is achieved not only using ML decoding, but also with simple naive lattice decoding.

While we discuss specific number field lattices, our proofs do work for any ensemble of lattices with asymptotically good product distance. The larger the product distance, the smaller the gap to the capacity in the fast fading channel.

In the existing literature, the product distance is mostly seen as a rough tool to estimate the worst case pairwise error probability in the high SNR regime. Instead we will see that when we are allowed to decode and encode over a growing number

of time units the normalized product distance will play a role of an equal importance to the Hermite constant in Gaussian channels. We point out that the study of normalized product distance and Hermite invariant are both examples of the more general problem of finding the minima of homogeneous forms in the mathematical field of *geometry of numbers*. This seems to be a universal theme, where each fading channel model is linked to a natural problem in geometry of numbers. We will elaborate further on this topic in [3], where we also extend our capacity results to the MIMO context.

The families of number fields we consider were first brought to coding theory in [10], where the authors pointed out that the corresponding lattices have large Hermite constant. Our proof for the Gaussian channel is therefore an obvious corollary to this result. In [11] it was pointed out that these families of number fields provide the best known normalized product distance.

II. NOTATION AND PRELIMINARIES

We will use the notation \mathbb{F} for the field \mathbb{R} or \mathbb{C} . A lattice $L \subset \mathbb{F}^n$ has the form $L = \mathbb{Z}x_1 \oplus \mathbb{Z}x_2 \oplus \dots \oplus \mathbb{Z}x_k$, where the vectors x_1, \dots, x_k are linearly independent over \mathbb{R} , i.e., form a lattice basis.

Definition 1: Let $v = (v_1, \dots, v_n)$ be a vector in \mathbb{F}^n . The *Euclidean norm* of v is $\|v\|_E = \sqrt{\sum_{i=1}^n |v_i|^2}$. If L is a lattice in \mathbb{F}^n , the *minimum distance* $sv(L)$ of L is defined to be the infimum of the Euclidean norms of all non-zero vectors in the lattice.

Definition 2: Let $v = (v_1, \dots, v_n)$ be a vector in \mathbb{F}^n . We define the *product norm* of v as $n(v) = \prod_{i=1}^n |v_i|$.

Assuming that $n(v) \neq 0$ for all the non zero elements $v \in L$, we can define the *minimum product distance* $d_{p,\min}(L)$ of L to be the infimum of the product norms of all non-zero vectors in the lattice.

We will use the notation $\text{Vol}(L)$ for the volume of the fundamental parallelotope of the lattice L .

We denote by $\text{Nd}_{p,\min}(L)$ the *normalized minimum product distance* of the lattice L , i.e. here we first scale L to have a unit size fundamental parallelotope and then take $d_{p,\min}(L')$ of the resulting lattice L' . In the same way we can define the normalized shortest vector of L and denote it with $\text{Nsv}(L)$. The square of the normalized shortest vector is called the *Hermite invariant* of the lattice.

We then have the following scaling laws. If L is a full lattice in \mathbb{C}^n , then

$$\text{Nd}_{\text{p},\min}(L) = \frac{d_{\text{p},\min}(L)}{\text{Vol}(L)^{1/2}}, \quad \text{Nsv}(L) = \frac{\text{sv}(L)}{\text{Vol}(L)^{1/2n}}.$$

In the case of a real lattice $L \subset \mathbb{R}^n$ we have

$$\text{Nd}_{\text{p},\min}(L) = \frac{d_{\text{p},\min}(L)}{\text{Vol}(L)}, \quad \text{Nsv}(L) = \frac{\text{sv}(L)}{\text{Vol}(L)^{1/n}}.$$

These two concepts are related by the following simple and well known application of the arithmetic-geometric mean inequality.

Proposition 1: Let L be a lattice in \mathbb{F}^n . Then

$$\text{Nd}_{\text{p},\min}(L) \leq \frac{\text{Nsv}(\phi(L))^n}{n^{n/2}}.$$

The following Lemma [12] is useful in order to choose lattice constellations with prescribed minimum size.

Lemma 1: Let us suppose that L is a full lattice in \mathbb{F}^n and S a Jordan measurable bounded subset of \mathbb{F}^n . Then there exists $x \in \mathbb{F}^n$ such that

$$|(L+x) \cap S| \geq \frac{\text{Vol}(S)}{\text{Vol}(L)}.$$

III. LATTICE CODES FROM NUMBER FIELDS

In the following we will describe the standard method to build lattice codes from number fields [8]. We will denote the discriminant of a number field K with d_K . For every number field it is a non-zero integer.

A. Complex constellations

Let K/\mathbb{Q} be a totally complex extension of degree $2n$ and $\{\sigma_1, \dots, \sigma_n\}$ be a set of \mathbb{Q} -embeddings, such that we have chosen one from each complex conjugate pair. Then we can define a *relative canonical embedding* of K into \mathbb{C}^n by

$$\psi(x) = (\sigma_1(x), \dots, \sigma_n(x)).$$

The ring of algebraic integers \mathcal{O}_K has a \mathbb{Z} -basis $W = \{w_1, \dots, w_{2n}\}$ and $\psi(W)$ is a \mathbb{Z} -basis for the full lattice $\psi(\mathcal{O}_K)$ in \mathbb{C}^n .

Lemma 2: Let K/\mathbb{Q} be an extension of degree $2n$ and let ψ be the relative canonical embedding. Then

$$\text{Vol}(\psi(\mathcal{O}_K)) = 2^{-n} \sqrt{|d_K|}$$

$$\text{Nd}_{\text{p},\min}(\psi(\mathcal{O}_K)) = \frac{2^{\frac{n}{2}}}{|d_K|^{\frac{1}{4}}} \text{ and } \text{Nsv}(\psi(\mathcal{O}_K)) = \frac{\sqrt{2n}}{|d_K|^{1/4n}}.$$

We can now see that both the normalized product distance and Hermite invariant of the number field lattices depend only on the discriminant of the field. In order to find promising codes we need fields with as small discriminants as possible.

Martinet [13] proves the existence of an infinite tower of totally complex number fields $\{K_n\}$ of degree $2n$, where $2n = 5 \cdot 2^k$, such that

$$|d_{K_n}|^{\frac{1}{n}} = G^2, \quad (1)$$

for $G \approx 92.368$. For such fields K_n we have that

$$\text{Nd}_{\text{p},\min}(\psi(\mathcal{O}_{K_n})) = \left(\frac{2}{G}\right)^{\frac{n}{2}} \text{ and } \text{Nsv}(\psi(\mathcal{O}_{K_n})) = \frac{\sqrt{2n}}{\sqrt{G}}.$$

Given transmission power P , we require that every point \mathbf{s} in a finite code $\mathcal{C} \subset \mathbb{C}^n$ satisfies the average power constraint

$$\frac{1}{n} \sum_{i=1}^n |s_i|^2 = \frac{1}{n} \sum_{i=1}^n (\Re(s_i)^2 + \Im(s_i)^2) \leq P. \quad (2)$$

Let R denote the code rate in bits per complex channel use; equivalently, $|\mathcal{C}| = 2^{Rn}$. Let us now show how we can produce codes \mathcal{C} , having rate greater or equal to R , and satisfying the power constraint (2), from the number field lattices $\psi(\mathcal{O}_K)$, where K belongs to the Martinet family.

In the following we will use the notation $B(\sqrt{nP})$ for a $2n$ -dimensional ball of radius \sqrt{nP} in \mathbb{C}^n . Let us suppose that α is some energy normalization constant. According to Lemma 1, we can choose an element $x_R \in \mathbb{C}^n$ such that for $\mathcal{C} = B(\sqrt{nP}) \cap (x_R + \alpha\psi(\mathcal{O}_K))$ we have

$$|\mathcal{C}| \geq 2^{Rn} = \frac{\text{Vol}(B(\sqrt{nP}))}{\text{Vol}(\alpha\psi(\mathcal{O}_K))} = \frac{2^n C_n P^n}{\alpha^{2n} \sqrt{|d_K|}},$$

where $C_n = \frac{(\pi n)^n}{n!}$. We can now see that by using the energy normalization

$$\alpha^2 = \frac{2P(C_n)^{\frac{1}{n}}}{2^R |d_K|^{\frac{1}{2n}}} = \frac{2P(C_n)^{\frac{1}{n}}}{2^R G}$$

the code \mathcal{C} has rate R , or greater, and satisfies the average power constraint.

B. Real constellations

Let us now suppose that we have a degree n totally real extension K/\mathbb{Q} and that $\{\sigma_1, \dots, \sigma_n\}$ are the \mathbb{Q} embeddings of K . We define the canonical embedding of K into \mathbb{R}^n by

$$\psi(x) = (\sigma_1(x), \dots, \sigma_n(x)).$$

We then have that $\psi(\mathcal{O}_K)$ is an n -dimensional lattice in \mathbb{R}^n .

Lemma 3: Let K/\mathbb{Q} be a totally real extension of degree n and let ψ be the canonical embedding. Then

$$\text{Vol}(\psi(\mathcal{O}_K)) = \sqrt{|d_K|},$$

$$\text{Nd}_{\text{p},\min}(\psi(\mathcal{O}_K)) = \frac{1}{\sqrt{|d_K|}} \text{ and } \text{Nsv}(\psi(\mathcal{O}_K)) = \frac{\sqrt{n}}{|d_K|^{\frac{1}{2n}}}.$$

In the case of totally real fields [13] proves the existence of a family of fields of degree n , where $n = 2^k$, such that

$$|d_{K_n}|^{\frac{1}{n}} = G_1, \quad (3)$$

where $G_1 \approx 1058$. If K is a degree n field from this family,

$$\text{Nd}_{\text{p},\min}(\psi(\mathcal{O}_K)) = \frac{1}{G_1^{\frac{n}{2}}} \text{ and } \text{Nsv}(\psi(\mathcal{O}_K)) = \frac{\sqrt{n}}{\sqrt{G_1}}. \quad (4)$$

As in the case of complex constellations, we will consider finite codes $\mathcal{C} = B(\sqrt{nP}) \cap (x_R + \alpha\psi(\mathcal{O}_K))$, where x_R is chosen so that

$$|\mathcal{C}| \geq 2^{Rn} = \frac{\text{Vol}(B(\sqrt{nP}))}{\text{Vol}(\alpha\psi(\mathcal{O}_K))} = \frac{C_n^{\mathbb{R}} P^{n/2}}{\alpha^n \sqrt{|d_K|}},$$

and $C_n^{\mathbb{R}} = \frac{(\pi n)^{n/2}}{\Gamma(n/2+1)}$. We then have that the choice

$$\alpha^2 = \frac{P(C_n^{\mathbb{R}})^{\frac{2}{n}}}{2^{2R} |d_K|^{\frac{1}{n}}} = \frac{P(C_n^{\mathbb{R}})^{\frac{2}{n}}}{2^{2R} G_1},$$

yields a code of rate R satisfying the power constraint $(1/n) \sum_{i=1}^n s_i^2 \leq P$.

IV. NUMBER FIELD CODES IN THE GAUSSIAN CHANNEL

Let us now consider the question of the maximal rates we can achieve with the codes \mathcal{C} of the previous section, when we demand vanishing error probability when n grows to infinity.

A. Complex constellations

We consider a complex Gaussian channel model

$$\mathbf{y} = \mathbf{s} + \mathbf{w},$$

where $\mathbf{s} \in \mathcal{C}$, and $\forall i = 1, \dots, n$, the w_i are i.i.d. complex Gaussian random variables with variance $\sigma_h^2 = \sigma^2 = \frac{1}{2}$ per real dimension. (Thus, under the assumptions of the previous Section, the SNR is P). For this channel model we consider the codes \mathcal{C} of Section III-A. Let us denote with

$$d = \min_{\substack{\mathbf{s}, \bar{\mathbf{s}} \in \mathcal{C} \\ \mathbf{s} \neq \bar{\mathbf{s}}}} \|\mathbf{s} - \bar{\mathbf{s}}\|$$

the minimum Euclidean distance in the constellation. Then if ML decoding or naive lattice decoding (NLD)¹ is used, we have the *sphere bound*

$$P_e \leq \mathbb{P} \left\{ \|\mathbf{w}\|^2 \geq \left(\frac{d}{2} \right)^2 \right\}.$$

The minimum distance of the lattice is lower bounded by

$$d^2 \geq \alpha^2 \min_{x \in \mathcal{O}_K \setminus \{0\}} \|\psi(x)\|^2 = \alpha^2 \text{sv}(L)^2 = \alpha^2 n.$$

Thus, the error probability is bounded by

$$P_e \leq \mathbb{P} \left\{ \|\mathbf{w}\|^2 \geq \left(\frac{\alpha^2 n}{4} \right) \right\}.$$

Note that $2\|\mathbf{w}\|^2 \sim \chi^2(2n)$. For a random variable $Z \sim \chi^2(n)$, the following concentration result holds $\forall \epsilon > 0$ [14]:

$$\mathbb{P} \left\{ \frac{Z}{n} \geq 1 + \epsilon \right\} \leq 2e^{-\frac{n\epsilon^2}{16}}.$$

Consequently, the probability of the set of non-typical noise vectors vanishes exponentially fast:

$$\mathbb{P} \left\{ \frac{\|\mathbf{w}\|^2}{n} \geq 1 + \epsilon \right\} \leq 2e^{-\frac{n\epsilon^2}{8}}.$$

¹By naive lattice decoding, we mean the closest point search in the infinite shifted lattice $x_R + \alpha\psi(\mathcal{O}_K)$.

Therefore, $P_e \rightarrow 0$ when $n \rightarrow \infty$ provided that

$$2^R < \frac{PC_n^{\frac{1}{n}}}{(1+\epsilon)2G}.$$

As $C_n = \frac{(\pi n)^n}{n!}$, using Stirling's approximation we have $C_n \approx \frac{(\pi e)^n}{\sqrt{2\pi n}}$ for large n . We can conclude that $P_e \rightarrow 0$ for any rate

$$R < \log_2(P) - \log_2(2G(1+\epsilon)) + \log_2(\pi e).$$

Since the previous bounds hold $\forall \epsilon$, we get the following:

Proposition 2: Over the complex Gaussian channel, any rate

$$R < \log_2(P) - \log_2 \left(\frac{2G}{\pi e} \right)$$

is achievable with the code construction in Section III-A.

B. Real constellations

We consider a real Gaussian channel model

$$\mathbf{y} = \mathbf{s} + \mathbf{w},$$

where $\mathbf{s} \in \mathcal{C}$, and $\forall i = 1, \dots, n$, the w_i are i.i.d. real Gaussian random variables with variance $\sigma_h^2 = \sigma^2 = 1$. The finite codes we consider are those of section III-B.

Analogously to the complex case we have $d^2 \geq \alpha^2 \text{sv}(L)^2 = \alpha^2 n$ and

$$P_e \leq \mathbb{P} \left\{ \|\mathbf{w}\|^2 \geq \left(\frac{\alpha^2 n}{4} \right) \right\}.$$

For all $\epsilon > 0$, the error probability vanishes as long as

$$2^{2R} < \frac{P(C_n^{\mathbb{R}})^{\frac{2}{n}}}{4(1+\epsilon)G_1}.$$

Using Stirling's approximation $C_n^{\mathbb{R}} \approx \frac{(2\pi e)^{n/2}}{\sqrt{\pi n}}$, we get the following:

Proposition 3: Over the real Gaussian channel, any rate

$$R < \frac{1}{2} \log_2(P) - \frac{1}{2} \log_2 \left(\frac{2G_1}{\pi e} \right)$$

is achievable using the code construction in Section III-B.

V. NUMBER FIELD CODES IN THE FAST FADING CHANNEL

A. Complex fast Rayleigh fading channel

We consider a complex fast Rayleigh fading channel model

$$\mathbf{y} = \mathbf{h} \cdot \mathbf{s} + \mathbf{w},$$

where $\mathbf{s} \in \mathcal{C} \subset \mathbb{C}^n$, and $\forall i = 1, \dots, n$, the h_i, w_i are i.i.d. complex Gaussian random variables with variance $\sigma_h^2 = \sigma^2 = \frac{1}{2}$ per real dimension. Therefore, if \mathcal{C} is one of the lattice codes described in Section III-A, the SNR is equal to P .

The minimum distance in the received constellation is

$$d_{\mathbf{h}} = \min_{\substack{\mathbf{s}, \bar{\mathbf{s}} \in \mathcal{C} \\ \mathbf{s} \neq \bar{\mathbf{s}}}} \|\mathbf{h} \cdot (\mathbf{s} - \bar{\mathbf{s}})\|.$$

The ML and NLD error probabilities are both bounded by

$$P_e \leq \mathbb{P} \left\{ \|\mathbf{w}\|^2 \geq \left(\frac{d_{\mathbf{h}}}{2} \right)^2 \right\}.$$

From the arithmetic-geometric mean inequality, we get

$$\begin{aligned} d_{\mathbf{h}}^2 &\geq \alpha^2 \min_{x \in \mathcal{O}_K \setminus \{0\}} \|\mathbf{h} \cdot \psi(x)\|^2 = \\ &= \alpha^2 \min_{x \in \mathcal{O}_K \setminus \{0\}} \sum_{i=1}^n |h_i|^2 |\sigma_i(x)|^2 \geq \\ &\geq \alpha^2 \min_{x \in \mathcal{O}_K \setminus \{0\}} n \left(\prod_{i=1}^n |h_i|^2 |\sigma_i(x)|^2 \right)^{\frac{1}{n}}. \end{aligned}$$

Since $\prod_{i=1}^n |\sigma_i(x)| \geq 1$ for all $x \in \mathcal{O}_K \setminus \{0\}$, we have

$$d_{\mathbf{h}}^2 \geq \alpha^2 n \left(\prod_{i=1}^n |h_i|^2 \right)^{\frac{1}{n}}$$

Therefore we have the upper bound

$$P_e \leq \mathbb{P} \left\{ \frac{\|\mathbf{w}\|^2}{n} \geq \frac{\alpha^2}{4} \left(\prod_{i=1}^n |h_i|^2 \right)^{\frac{1}{n}} \right\}. \quad (5)$$

Fixing $\epsilon > 0$, the law of total probability implies that

$$P_e \leq \mathbb{P} \left\{ \frac{\|\mathbf{w}\|^2}{n} \geq 1 + \epsilon \right\} + \mathbb{P} \left\{ \frac{\alpha^2}{4} V_n < 1 + \epsilon \right\}, \quad (6)$$

where $V_n = \left(\prod_{i=1}^n |h_i|^2 \right)^{\frac{1}{n}}$.

Since the $|h_i|$ are Rayleigh distributed with parameter $\sigma_h^2 = \frac{1}{2}$, the random variables $X_i = |h_i|^2$ have exponential density $p_X(x) = e^{-x}$. To find a good upper bound for the error probability, we need to analyze the distribution of the random variable V_n which is a geometric average of exponential distributions.

Note that $\ln V_n = \frac{1}{n} \sum_{i=1}^n \ln X_i$. The random variables $Y_i = \ln X_i$ have density $p_Y(y) = e^{y-e^y}$ and mean

$$m_y = \mathbb{E}[\ln X] = \int_0^\infty (\ln x) e^{-x} dx = -\gamma,$$

where $\gamma \approx 0.577215$ is the Euler-Mascheroni constant. From the Chernoff bound [15, §2.1.6] for the zero-mean random variable $-\frac{1}{n} \sum_{i=1}^n \ln X_i - \gamma$, we get that $\forall \delta, v > 0$,

$$\mathbb{P} \left\{ \frac{1}{n} \sum_{i=1}^n \ln X_i \leq -(\delta + \gamma) \right\} \leq e^{-nv(\delta + \gamma)} (\mathbb{E}[e^{-vX}])^n \quad (7)$$

For a given $\delta > 0$, the optimal $v_\delta > 0$ that gives the tightest upper bound is the solution of the equation $\mathbb{E}[-\ln X e^{-v_\delta \ln X}] = (\delta + \gamma) \mathbb{E}[e^{-v_\delta \ln X}]$. We have

$$\begin{aligned} \mathbb{E}[e^{-v \ln X}] &= \int_0^\infty \frac{e^{-x}}{x^v} dx = \Gamma(1 - v), \\ \mathbb{E}[-\ln X e^{-v \ln X}] &= \int_0^\infty \frac{\ln x e^{-x}}{x^v} dx = -\Gamma(1 - v) \psi(1 - v), \end{aligned}$$

where $\psi(x) = \frac{d}{dx} \ln \Gamma(x)$ denotes the digamma function. Thus, $\psi(1 - v_\delta) = -(\delta + \gamma)$. Note that as $\delta \rightarrow 0$, also $v_\delta \rightarrow 0$ since $\psi(1) = -\gamma$. The Chernoff bound (7) thus gives

$$\begin{aligned} \mathbb{P} \{ \ln V_n \leq -(\delta + \gamma) \} &= \mathbb{P} \{ V_n \leq e^{-\delta} e^{-\gamma} \} \leq \\ &\leq e^{-nv_\delta(\gamma + \delta)} (\Gamma(1 - v_\delta))^n = e^{n(v_\delta \psi(1 - v_\delta) + \ln \Gamma(1 - v_\delta))} \end{aligned}$$

The mean value theorem for the function $\ln \Gamma(x)$ in the interval $[1 - v_\delta, 1]$ yields $|\ln \Gamma(1 - v_\delta)| \leq |\psi(\xi)| v_\delta$ for some $\xi \in (1 - v_\delta, 1)$. Since $\psi < 0$ in the interval $(0, 1)$, $|\psi(\xi)| \leq |\psi(1 - v_\delta)| = -\psi(1 - v_\delta)$, and so

$$v_\delta \psi(1 - v_\delta) + \ln \Gamma(1 - v_\delta) \leq 0.$$

Therefore $\forall \delta > 0$, $\mathbb{P} \{ \ln V_n \leq -(\delta + \gamma) \} \rightarrow 0$ as $n \rightarrow \infty$.

Let us now go back to Equation (6). As seen in the Gaussian case, the first term in the sum vanishes exponentially fast. The second term will tend to 0 when $n \rightarrow \infty$ provided that $\frac{4(1+\epsilon)}{\alpha^2} < e^{-(\delta+\gamma)}$. Therefore, $P_e \rightarrow 0$ provided that

$$2^R < \frac{PC_n^{\frac{1}{n}}}{2e^{\delta+\gamma}(1+\epsilon)d_K^{\frac{1}{2n}}} = \frac{PC_n^{\frac{1}{n}}}{2e^{\delta+\gamma}(1+\epsilon)G}.$$

Again using Stirling's approximation we have $C_n \approx \frac{(\pi e)^n}{\sqrt{2\pi n}}$ for large n , and the achievable rate is

$$R < \log_2(P) - \log_2 \left(\frac{2G(1+\epsilon)e^{\delta+\gamma}}{\pi e} \right)$$

Since the previous bounds hold for any choice of $\epsilon, \delta > 0$, we can state the following:

Proposition 4: Over the complex Rayleigh fading channel, any rate

$$R < \log_2(Pe^{-\gamma}) - \log_2 \left(\frac{2G}{\pi e} \right)$$

is achievable using the codes of Section III-A.

We can compare this result to the bound for Rayleigh channel capacity given in [16], equation (7):

$$C \geq \log_2(1 + Pe^{-\gamma}).$$

This is a lower bound, however it has been shown to be very tight for high SNR.

B. Real Rayleigh fast fading channel

We consider a real fast Rayleigh fading channel model [8]

$$\mathbf{y} = \mathbf{g} \cdot \mathbf{s} + \mathbf{w},$$

where $\mathbf{s} \in \mathcal{C}$, and $\forall i = 1, \dots, n$, the $g_i = |h_i|$ are Rayleigh distributed with parameter $\sigma_h^2 = \frac{1}{2}$, and w_i are i.i.d. real Gaussian random variables with variance $\sigma^2 = 1$. Note that the SNR is again P when using one of the real lattice constellations from Section III-B. The error probability estimate for this model proceeds exactly as in the case of the complex Rayleigh fading channel in Section V-A. A sufficient condition to have vanishing error probability when $n \rightarrow \infty$ is

$$2^{2R} < \frac{P(C_n^{\mathbb{R}})^{\frac{1}{n}}}{4e^{\delta+\gamma}(1+\epsilon)d_K^{\frac{1}{2n}}} \approx \frac{P(C_n^{\mathbb{R}})^{\frac{1}{n}}}{4e^{\delta+\gamma}(1+\epsilon)G_1}.$$

Since $C_n^{\mathbb{R}} \approx \frac{(2\pi e)^n}{\sqrt{\pi n}}$ for large n , and taking the supremum over all $\epsilon > 0$, we find the following:

Proposition 5: Over the real Rayleigh fading channel, any rate

$$R < \frac{1}{2} \log_2(Pe^{-\gamma}) - \frac{1}{2} \log_2 \left(\frac{2G_1}{\pi e} \right)$$

is achievable using the codes of Section III-B.

VI. DISCUSSION

Let us now draw some conclusions and highlight the similarities between Gaussian and fast-fading channels. We saw that there exists an ensemble of lattice codes from number fields that reach all rates satisfying

$$R < \frac{1}{2} \log_2(Pe^{-\gamma}) - \frac{1}{2} \log_2 \left(\frac{2G_1}{\pi e} \right)$$

in real fast fading channels and rates

$$R < \frac{1}{2} \log_2(P) - \frac{1}{2} \log_2 \left(\frac{2G_1}{\pi e} \right),$$

in Gaussian channel. According to (4) these results can be transformed into the following forms

$$R < \frac{1}{2} \log_2(Pe^{-\gamma}) - \frac{1}{2} \log_2 \left(\frac{2}{\pi e (\text{Nd}_{(p,\min)}(L))^{2/n}} \right) \quad (8)$$

$$R < \frac{1}{2} \log_2(P) - \frac{1}{2} \log_2 \left(\frac{2n}{\text{N}_{\text{sv}}(L)^2 \pi e} \right).$$

Here the normalized product distance and shortest vector play identical roles. The greater the distance, the smaller the gap to capacity. This is not only a property of these specific number field codes, but is true for any family of lattice codes. Indeed, while our proofs refer to specific number field codes, the performance only depends on the normalized product distances.

We can now see that in order to reach a constant gap to capacity in fast fading channel, at least with this method, we must have that $(\text{Nd}_{(p,\min)}(L_n))^{2/n}$ stays above some constant. According to Proposition 1 the product distance is upperbounded by the Hermite constant of the lattice. This result suggests that when n grows a lattice code must have a Hermite constant growing linearly with n in order to be good over the fast fading channel. However, we note that a good Hermite constant does not automatically guarantee a good performance in fast fading channels for general families of lattice codes.

Finally, let us consider how close to capacity this approach can bring us in an optimal scenario. If we consider totally real lattices from number fields, then the Odlyzko bound states that when $m \rightarrow \infty$ we have that $|d_K|^{1/m} \geq 60.8$. Assuming that we can reach this bound with an ensemble of lattice codes we have that any rate R satisfying

$$R < \frac{1}{2} \log_2(Pe^{-\gamma}) - \frac{1}{2} \log_2 \left(\frac{2 \cdot 60.8}{\pi e} \right)$$

is achievable. The Odlyzko bound does bound the achievable rate of number field codes, but if we consider all lattices we have a slightly weaker bound. For a full lattice in \mathbb{R}^n , a classical result of Minkowski gives us that $\text{Nd}_{p,\min}(L) \leq \frac{n!}{n^n}$. Assuming that we have an ensemble of lattice codes reaching this bound we have by Stirling's approximation and equation (8) that rates satisfying

$$R < \frac{1}{2} \log_2(Pe^{-\gamma}) - \frac{1}{2} \log_2 \left(\frac{2e}{\pi} \right),$$

are achievable. This result shows that with this method we will always have a gap to $\frac{1}{2} \log_2(Pe^{-\gamma})$ irrespective of the choice of lattice code. However, just like in the case of the Gaussian channel, these bounds do not represent the performance limits of lattice codes, because the method itself and the error probability bounds are suboptimal.

Remark 1: We note that the number field towers we used are not the best known possible. It was shown in [17] that one can construct a family of real fields such that $G_1 < 954.3$ and totally complex such that $G < 82.2$, but this choice would add some notational complications.

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