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Using inconsistency measures for estimating reliability

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A B S T R A C T

Any decision taken by an agent requires some knowledge of its environment. Communication with other agents is a key issue for assessing the overall quality of its own knowledge. This assessment is a challenge itself as the agent may receive information from unknown agents. The aim of this paper is to propose a framework for assessing the reliability of unknown agents based on communication. We assume that information is represented through logical statements and logical inconsistency is the underlying notion of reliability assessment. In our context, assessing consists of ranking the agents and representing reliability through a total preorder.

The overall communication set is first evaluated with the help of inconsistency measures. Next, the measures are used for assessing the contribution of each agent to the overall inconsistency of the communication set. After stating the postulates specifying the expected properties of the reliability preorder, we show through a representation theorem how these postulates and the contribution of the agent are interwoven. We also detail how the properties of the inconsistency measures influence the properties of the contribution assessment. Finally we describe how to aggregate different reliability preorders, each of them may be based on different inconsistency measures.

Keywords:

Logic

Inconsistency

Reliability assessment

1. Introduction

To be able to act or deliberate, any rational agent must acquire knowledge of its environment. It gets it by merging information provided by its own sensors and/or by merging information communicated by other agents. Merging basic information is a key issue for any agent as it is the underlying rationale for decision making and it contributes to justify the agent's epistemic state. Techniques for merging raw information have been studied in an extensive way. These techniques usually assume that all information provided by the sources (i.e. agents) should be considered as a whole. Two different approaches have been studied: the first one considers sources in an equal way and has led to merging techniques such as majority merging, negotiation, arbitration merging or distance-based merging for solving conflict between contradicting information [21,8,27,9]. The second one distinguishes sources through a reliability criterion. Taking sources reliability into account provides rationales for discounting or ignoring pieces of information whose source is not considered as sufficiently reliable. Some promote a quantitative model of reliability: information sources are associated with a reliability level represented by a number used by the merging operator. According to the belief function theory, the reliability level of a source is a number between 0 and 1. This number is then used by the discounting rule in order to weaken the importance of infor-

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mation provided by this source [31]. Some others promote a qualitative approach to reliability and consider that information sources are ranked according to their reliability. This order or pre-order is then used by the merging operator. In [5], the author defines a merging operator which assumes that the sources are totally ordered: if s is said to be more reliable than s' and together provide contradicting information, then information provided by s is privileged; while information provided by s' which does not contradict information of s is also considered as acceptable. The same idea is followed by [22] for reasoning about more complex beliefs and in [24] for revising a belief base. All these works assume that the reliability of the sources is given as a parameter (quantitative or qualitative), they do not address the question of how to build up this reliability.

In this present paper our aim is to address the key question of how to build a reliability preorder of information sources, in a context where sources are unknown: no extra information about sources is available and information provided by the sources is only qualitative (i.e., statements). We adopt a qualitative point of view to represent reliability: the relative reliability of information sources is represented by a total preorder. We propose to consider a phase, before the information merging phase, during which information sources are observed in order to obtain a reliability preorder. The purpose of this phase is to analyze the inconsistency of information reported by the different sources w.r.t. some trusted knowledge.

Our main goal is thus to show that the relative reliability of information sources can be estimated from the inconsistency of reported information. Two different approaches can be followed. The first approach consists in using an ad-hoc model for reported information and in developing new inconsistency measures. The second approach consists in modeling reported information in a conventional way in order to use well known inconsistency measures.

In a recent paper [6], we followed the first approach. Reported information was modeled by pairs: $\langle agent, formula \rangle$, $formula$ representing a piece of information communicated by $agent$. For instance, the set $\{\langle a, p \rangle, \langle b, \neg p \rangle, \langle b, q \rangle\}$ represented the fact that agent a had reported p , agent b had reported $\neg p$ and had also reported q . The main notions (inconsistency, minimal inconsistent subsets, inconsistency measures...) available in the literature have been adapted to this model.

In this paper, our very motivation is to show an original application of inconsistency measures, i.e. *reliability estimation*. Our starting point is the existing inconsistency measures. Hereafter, we simplify the representation of reported information so that we can re-use these existing inconsistency measures for elaborating agent's reliability.

Our original contributions consist in (i) characterizing the individual contribution of each agent to the overall inconsistency of a set of reported information and (ii) introducing postulates which characterize the expected properties of the reliability preorder; Based on these axiomatic perspective on reliability assessment, we show (i) how the properties of the inconsistency measure influence the properties of the contributions measures and (ii) how postulates about reliability and properties of agent contribution are related through a representation theorem. Finally, we show how to aggregate several preorders possibly obtained through different inconsistency measures; namely we show how the overall aggregated preorder may satisfy the reliability postulates if the initial preorders also satisfy these postulates.

This paper is organized as follows. Section 2 and Section 3 introduce the main notions needed to assess reliability of agents. They introduce inconsistent communication sets and focus on measuring the inconsistency in communication sets. Based on the inconsistency measures, Section 4 shows how to assess the individual contribution of an agent to the overall inconsistency of a communication set. Contribution is first characterized in an axiomatic way and next two possible contribution functions instantiating the expected properties are detailed. Some implementation and complexity considerations are also addressed. Section 5 proposes a set of postulates which axiomatically characterize reliability preorders and show through two representation theorems how these postulates and the agent contributions are related. Still in Section 5, we present two possible solutions for building a reliability preorder compliant with these postulates. Section 6 considers the aggregation of several reliability preorders and shows how Arrow's condition for aggregation and our postulates interplay. Finally, Section 7 concludes the paper and discusses future work.

2. Inconsistent communication sets

This section introduces communication sets and focuses on their inconsistency.

2.1. Preliminaries

Let \mathcal{L} be a propositional language of formulas defined over a finite set of propositional symbols \mathcal{P} , propositional constants \top, \perp and the logical connectives \wedge, \vee, \neg . We use p, q, r, \dots to denote the propositional symbols and Greek letters ϕ, ψ, \dots to denote formulas of the classical propositional logic defined over \mathcal{L} . An interpretation i is a total function from \mathcal{P} to $\{0, 1\}$ from which an assignment to $\{0, 1\}$ is generated for all the formulas of \mathcal{L} defined in the usual way of classical logic. As usual, $i(\top) = 1$ and $i(\perp) = 0$. Interpretation i is a *model* of formula ϕ iff $i(\phi) = 1$. *Tautologies* are formulas which are interpreted by 1 in any interpretation. We write $\models \phi$ when ϕ is a tautology. A formula is *consistent* iff it has at least one model. Otherwise it is *inconsistent*.

A *communication base*¹ K is a finite (possibly empty) set of formulas of \mathcal{L} . $At(K)$ denotes the set of propositional symbols appearing in formulas which belong to K . A communication base is *consistent* iff the conjunction of its formulas is

¹ This term corresponds to the term "knowledge base", widely used in the community of belief merging.

consistent. Otherwise, it is *inconsistent*. For a communication base K , $MI(K)$ is the set of minimal inconsistent subsets of K , i.e., $MI(K) = \{K' \subseteq K \mid K' \text{ is inconsistent and } \forall K'' \subset K' \text{ } K'' \text{ is consistent}\}$. $MC(K)$ is the set of maximal consistent subsets of K , i.e., $MC(K) = \{K' \subseteq K \mid K' \text{ is consistent and } \forall K'' \text{ s.t. } K' \subset K'' \text{ } K'' \text{ is inconsistent}\}$. If $MI(K) = \{M_1, \dots, M_n\}$ then $Problematic(K) = M_1 \cup \dots \cup M_n$, and $Free(K) = K \setminus Problematic(K)$. The set of formulas in K that are inconsistent is given by the function $Selfcontradiction(K) = \{\phi \in K \mid \phi \text{ is inconsistent}\}$. Notice that these definitions are provided by [10].

Finally, as shown in [10], a three-valued logic can be used to give a semantics to inconsistent formulae. The three values are T, F, B where T and F correspond to the classical values 1, 0 respectively and the additional truth value B stands for *both* and represents inconsistency. Assuming that the three values are ordered by: $F <_t B <_t T$, the valuation of formulae in an interpretation i is given by: $i(\top) = T$, $i(\perp) = F$, $i(\neg\phi) = B \iff i(\phi) = B$, $i(\neg\phi) = T \iff i(\phi) = F$, $i(\phi \wedge \psi) = \min_{\leq_t}(i(\phi), i(\psi))$, $i(\phi \vee \psi) = \max_{\leq_t}(i(\phi), i(\psi))$. Interpretation i is a model of K if no formula in K is assigned the truth value F . We write $i \models_3 K$. $Binarybase(i) = \{p \in \mathcal{P} \mid i(p) = T \text{ or } i(p) = F\}$ and $Conflictbase(i) = \{p \in \mathcal{P} \mid i(p) = B\}$. Notice that for $i \models_3 K$, $Conflictbase(i)$ returns the subset of propositional symbols directly involved in inconsistency according to i .

2.2. Communication sets

From now, we assume a finite set of agents $A = \{a_1, \dots, a_n\}$ ($n \geq 1$). For each agent a_i , $K(a_i)$ is an \mathcal{L} -formula which is the conjunction of all the pieces of information reported by a_i . $K(a_i)$ is called *report of agent a_i* . The report of an agent which provides no information is any tautology or, for short, \top . Given $K(a_1), \dots, K(a_n)$, the multi-set of formulas $\Psi = \{K(a_1), \dots, K(a_n)\}$ is called *communication set*.

Given $\Psi = \{K(a_1), \dots, K(a_n)\}$ and $C = \{a_{i_1}, \dots, a_{i_m}\} \subseteq A$, we define $\Psi(C)$ by $\Psi(C) = \{K(a_{i_1}), \dots, K(a_{i_m})\}$.

Consider now two situations, one in which agent reports are $K(a_1) \dots K(a_n)$ and a second one in which they report $K'(a_1) \dots K'(a_n)$. Let $\Psi = \{K(a_1), \dots, K(a_n)\}$ and $\Psi' = \{K'(a_1), \dots, K'(a_n)\}$ be the corresponding communication sets. We write $\Psi \equiv \Psi'$ (Ψ and Ψ' are *equivalent*) iff $\forall a \in A, \models K(a) \leftrightarrow K'(a)$. That is, a 's report in Ψ is equivalent to a 's report in Ψ' . We write $\Psi \rightrightarrows \Psi'$ (Ψ and Ψ' are *weakly equivalent*) iff $\forall a \in A, \exists b, \exists c \in A$ such that $\models K(a) \leftrightarrow K'(b)$ and $\models K'(a) \leftrightarrow K(c)$. That is, we relax here the constraint that report of agent a should be equivalent both in Ψ and Ψ' ; instead we only require some other agent, possibly different from a , reports equivalent information.

2.3. IC-inconsistent communication sets

In the context of a communication set Ψ , consistency will be evaluated with respect to some integrity constraints IC which is a consistent formula of \mathcal{L} . IC has to be viewed as information taken for granted or certain. Thus we say that $\Psi = \{K(a_1), \dots, K(a_n)\}$ is *IC-inconsistent* iff $\Psi \cup \{IC\}$ is inconsistent; otherwise Ψ is *IC-consistent*. In the following, we adapt the definitions given in the preliminaries to the case of communication sets.

Definition 1.

- $\Psi \perp IC$ is the multi-set of minimal IC -inconsistent subsets of Ψ i.e the multi-set of $X \subseteq \Psi$ such that X is IC -inconsistent and $\forall X' \subset X, X'$ is IC -consistent.
- $\Psi \top IC = \{X \subseteq \Psi : X \text{ is } IC\text{-consistent and } \forall X' \subset X, X' \text{ is } IC\text{-inconsistent}\}$ is the set of maximal IC -consistent sub-multisets of Ψ .
- $Problematic(\Psi) = \bigcup_{M \in \Psi \perp IC} M$.
- $Free(\Psi) = \Psi \setminus Problematic(\Psi)$.
- $Selfcontradiction(\Psi) = \{K(a) \in \Psi : K(a) \wedge IC \text{ is inconsistent}\}$.

2.4. IC-inconsistent sets of agents

We finally introduce the notion of minimal IC -inconsistent subsets of agents and the notion of problematic agents which are agents whose reports are problematic:

Definition 2.

- $A \perp IC = \{X \subseteq A : \Psi(X) \in \Psi \perp IC\}$.
- $Problematic(A) = \{a \in A : K(a) \in Problematic(\Psi)\}$.

Example 1. Consider $A = \{a, b, c\}$, $IC = \neg q$, $K(a) = p$, $K(b) = \neg p \wedge q$, $K(c) = r \wedge s$. i.e., a has reported p , b has reported $\neg p$ and q , and c has reported r and s . Here, $\Psi = \{p, \neg p \wedge q, r \wedge s\}$. Thus, $\Psi \perp IC = \{\{\neg p \wedge q\}\}$, $\Psi \top IC = \{\{p, r \wedge s\}\}$, $A \perp IC = \{\{b\}\}$ and $Problematic(A) = \{b\}$.

Example 2. Assume $A = \{a, b, c\}$, $IC = \neg(p \wedge q)$, $K(a) = K(b) = p$, $K(c) = q$. Here, $\Psi = \{p, p, q\}$. Thus, $\Psi \perp IC = \{\{p, q\}, \{p, p, q\}\}$. $Problematic(\Psi) = \{p, q\}$ and $Free(\Psi) = \emptyset$. $A \perp IC = \{\{a, c\}, \{b, c\}\}$ and $Problematic(A) = \{a, b, c\}$.

3. Measuring the IC-inconsistency of communication sets

Understanding the nature of inconsistency of a set of formulas is an important topic which aroused a great amount of research during the past decade. The purpose is to analyze to which degree a set of formulas is inconsistent. A lot of inconsistency measures have been proposed according to different points of view: some of them are based on minimal inconsistent subsets [15,16,25,20] or on maximal consistent subsets [10,1], some consider paraconsistent models such as three valued logic [17,10], some consider probabilistic functions over the underlying propositional language [29], some consider model distance [11] and finally some are proof-based [19]. For a detailed review of these inconsistency measures we refer the reader to [30]. In this section, we first review some well-known inconsistency measures then we show our requirements to adapt them to our context.

3.1. Inconsistency measures for sets of formulas

According to [10], an inconsistency measure on sets of formulas is a function I which assigns any set of formulas K to an element of \mathbb{R}_+ satisfying at least the following three properties:

- **Consistency:** $I(K) = 0$ iff K is consistent.
- **Monotony:** if $K \subseteq K'$, then $I(K) \leq I(K')$.
- **Free formula independence:** If $\phi \in \text{Free}(K)$ then $I(K) = I(K \setminus \{\phi\})$.

That is, the measure of inconsistency of a set of formulas is null iff this set is not inconsistent. The measure of inconsistency of a set of formulas does not decrease if we add more formulas. Finally, removing a formula that does not cause any contradiction does not change the inconsistency measure.

Let us now briefly review some inconsistency measures for sets of formulas, introduced in the literature [15–17,10]

- $I_D(K) = 0$ if K is consistent; 1 otherwise.
- $I_C(K) = |MI(K)|$.
- $I_M(K) = (|MC(K)| + |\text{Selfcontradictions}(K)|) - 1$.
- $I_P(K) = |\text{Problematic}(K)|$.
- $I_Q(K) = 0$ if K is consistent; $\sum_{K' \in MI(K)} \frac{1}{|K'|}$ otherwise.
- $I_B(K) = \min\{|\text{Conflictbase}(i)| \mid i \models_3 K\}$.

$I_D(K)$ is a trivial measure which assigns 0 to any consistent set of formulas and 1 to any inconsistent set. It does not quantify how much K is inconsistent. $I_C(K)$ counts the number of minimal inconsistent subsets of K . $I_M(K)$ counts the number of maximal consistent subsets together with the number of contradictory formulas but 1 to make $I_M(K) = 0$ when K is consistent. $I_P(K)$ counts the number of formulas in minimal inconsistent subsets of K . $I_Q(K)$ computes the weighted sum of the minimal inconsistent subsets of K , where the weight is the inverse of the size of the minimal inconsistent subset, so that smaller inconsistent subsets are regarded as more inconsistent than larger ones. $I_B(K)$ returns the minimum number of propositional symbols that have to be set to value B in order to get a three valued model of K . Measures $I_C(K)$, $I_P(K)$ and $I_Q(K)$ assume that inconsistency is rooted in minimal inconsistent subsets: removing any formula from a minimal inconsistent subset is sufficient to produce a maximal consistent subset. To some extent $I_M(K)$ can be viewed as a variant of these inconsistency measures. It is interesting to note that while $I_C(K)$ give the same score to sets of formulas containing the same number of minimal inconsistent subsets, $I_Q(K)$ is able to differentiate their level of inconsistency considering that the smaller is the size of a minimal inconsistent subset, the bigger is the amount of inconsistency. Measures I_B and I_{LP_m} follow a different approach assessing the severity of inconsistency by considering the number of propositional symbols involved in inconsistency.

J. Grant and A. Hunter showed in [10] that inconsistency measures can be used to order several sets of formulas from the least inconsistent one to the most inconsistent one. They defined two inconsistency measures I_x and I_y as being *order-compatible* if for all sets of formulas K_1 and K_2 , $I_x(K_1) < I_x(K_2)$ iff $I_y(K_1) < I_y(K_2)$. They proved that I_C , I_M , I_P , I_Q and I_B are pairwise order-incompatible and we add that I_D is also order-incompatible with each of them. This result is interesting as it shows that each inconsistency measure gives a particular insight on the inconsistency of a set of formulas.

3.2. IC-inconsistency measures for communication sets

In order to assign a degree of IC-inconsistency to a communication set, we adapt the inconsistency measures. I_D , I_C , I_M , I_P and I_Q are typically syntax-based inconsistency measures in the sense they mainly take into consideration the formulas composing K . This aspect is quite important in our context. However we need to introduce an extra property that IC-inconsistency measures should satisfy. This property is named *syntax weak-independence*. It states that IC-inconsistency measures should not be dependent of the syntax of IC and that two weakly equivalent communication sets should have the same IC-inconsistency measure. In other words, the measure does not depend on the sources of the communication.

Definition 3. Let IC and IC' be two integrity constraint and Ψ and Ψ' be two communication sets on A . Function $I_{IC} : 2^{A \times L} \rightarrow \mathbb{R}_+$ is a *syntax weak-independent IC-inconsistency measure* iff it satisfies the following properties:

- **Consistency:** $I_{IC}(\Psi) = 0$ iff Ψ is IC -consistent.
- **Monotony:** If $\Psi \subseteq \Psi'$ then $I_{IC}(\Psi) \leq I_{IC}(\Psi')$.
- **Free formula independence:** If $\phi \in \text{Free}(\Psi)$ then $I_{IC}(\Psi) = I_{IC}(\Psi \setminus \{\phi\})$.
- **Syntax weak-independence:**
 1. for all IC' if $\models IC \leftrightarrow IC'$ then $I_{IC}(\Psi) = I_{IC'}(\Psi)$.
 2. for all Ψ' if $\Psi \equiv \Psi'$ then $I_{IC}(\Psi) = I_{IC}(\Psi')$.

According to this definition, the measure of IC -inconsistency of a communication set is null iff this communication set is not IC -inconsistent. The measure of IC -inconsistency of a communication set does not decrease if we add more communications. Removing a report which does not cause any contradiction does not change the IC -consistency measure of the communication set. Finally, the measure of IC -inconsistency of a communication set does not depend on the syntax on the integrity constraints and two weakly equivalent communication sets get the same measure of IC -inconsistency.

The different measures introduced for sets of formulas can now be redefined for communication sets as follows:

Definition 4. Let Ψ be a communication set and IC an integrity constraint. The IC -inconsistency measures $I_{IC}^D, I_{IC}^C, I_{IC}^M, I_{IC}^P, I_{IC}^Q$ and I_{IC}^B are defined as follows:

- $I_{IC}^D(\Psi) = 0$ if Ψ is IC -consistent; 1 otherwise.
- $I_{IC}^C(\Psi) = |\Psi \perp IC|$.
- $I_{IC}^M(\Psi) = (|\Psi \top IC| + |\text{Selfcontradictions}(\Psi)|) - 1$.
- $I_{IC}^P(\Psi) = |\text{Problematic}(\Psi)|$.
- $I_{IC}^Q(\Psi) = 0$ if Ψ is IC -consistent; $\sum_{K \in \Psi \perp IC} \frac{1}{|K|}$ otherwise.
- $I_{IC}^B(\Psi) = \min\{|\text{Conflictbase}(i)| \mid i \models \bigwedge_{K \in \Psi} K \wedge IC\}$.

Proposition 1. $I_{IC}^D, I_{IC}^C, I_{IC}^M, I_{IC}^P, I_{IC}^Q$ and I_{IC}^B are *syntax weak-independent inconsistency measures*.

(All proofs are detailed in the appendix section, before references)

$I_{IC}^D(\Psi)$ is still a trivial measure, while each other measure gives insight on the contribution of reports to inconsistency. $I_{IC}^C(\Psi)$ counts the number of minimal IC -inconsistent subsets of reports. $I_{IC}^M(\Psi)$ counts the number of the maximal IC -consistent subsets of reports together with the number of self contradictory reports but 1 to make $I_{IC}^M(\Psi) = 0$ when Ψ is IC -consistent. $I_{IC}^P(\Psi)$ counts the number of problematic reports. $I_{IC}^Q(\Psi)$ adds the inverse of the sizes of the minimal IC -inconsistent subsets of reports, so that smaller IC -inconsistent subsets of reports are regarded as more inconsistent than larger ones. $I_{IC}^B(\Psi)$ returns the minimum number of propositional symbols that have to be set to value B in order to get a three valued model of $\Psi \cup IC$.

4. Contribution of an agent to inconsistency

In this section, we aim at characterizing the contribution of an agent to the overall inconsistency of a communication set.

4.1. Contribution function

Evaluating the contribution of an agent to the overall inconsistency of a communication set should be relative to an inconsistency measure. As previously stressed, inconsistency measures provides different perspectives on inconsistency and the contribution of an agent may then differ w.r.t. some inconsistency measure. Several possible solutions are possible for assessing this contribution but, as for the inconsistency measure, some constraints should also be satisfied. At first, the contribution of an agent should be null if it does not contribute to any inconsistency. Second, an agent not involved in any inconsistency should not influence the assessment of the contribution of the other agents; finally contribution should be syntax independent. Syntax-weak independence is not relevant here as illustrated by the following example: suppose two sets Ψ and Ψ' such that they are weakly equivalent; in that context, it may be the case that some agent a is problematic w.r.t. Ψ while it is not w.r.t. Ψ' . Hence their contribution should not be the same.

Definition 5. Consider a set of agents A , a communication set Ψ on A , an integrity constraint IC and an IC -inconsistency measure I_{IC} . Function $\text{Cont}^{\Psi, I_{IC}}$ is a syntax independent contribution function if it associates to any agent $a \in A$ a positive real number $\text{Cont}^{\Psi, I_{IC}}(a)$ so that:

- **Consistency:** $Cont^{\Psi, I_C}(a) = 0$ iff $a \notin Problematic(A)$.
- **Free agent independence:** if $b \notin Problematic(A)$ then $Cont^{\Psi, I_C}(a) = Cont^{\Psi \setminus \{K(b)\}, I_C}(a)$.
- **Syntax independence:**
 1. for all IC' if $\models IC \leftrightarrow IC'$ then $Cont^{\Psi, I_C}(a) = Cont^{\Psi, I_{C'}}(a)$.
 2. for all Ψ' if $\Psi \equiv \Psi'$ then $Cont^{\Psi, I_C}(a) = Cont^{\Psi', I_C}(a)$.

4.2. Assessing the contribution

Hereafter, we adapt two well known metrics in order to define the contribution of an agent to the communication set inconsistency, namely the Shapley value and the Banzhaf index. They are well known measures for assessing the power of an agent in a voting procedure. Our context is similar: Shapley value enables to assess the personal contribution of an agent to overall inconsistency of a communication base; in other words it measures the importance of this agent in a coalitional game defined by function I_C [17]. The Banzhaf index assesses the pivotal role of an agent in the definition of an inconsistent set. It corresponds to the Banzhaf score as shown in [18]. These functions are respectively denoted $Cont_s^{\Psi, I_C}$ and $Cont_b^{\Psi, I_C}$ and whenever it's clear, we denote them $Cont_s$ and $Cont_b$ for short. The Shapley based contribution function is defined as follows.

Definition 6. Consider a set of agents A , a communication set Ψ on A , an integrity constraint IC and an IC -inconsistency measure I_C . Function $Cont_s$ associates any agent a with a positive real number $Cont_s(a)$ so that:

$$Cont_s(a) = \sum_{\substack{C \subseteq A \\ C \neq \emptyset}} \frac{(|C| - 1)! (|A| - |C|)!}{|A|!} (I_C(\Psi(C)) - I_C(\Psi(C \setminus \{a\}))).$$

The following proposition connects the inconsistency measure and the contribution function. It states that if the inconsistency measure is syntax weak independent then the contribution function is syntax independent.

Proposition 2. Function $Cont_s$ is a syntax independent contribution function if I_C is a syntax-weak independent IC -inconsistency measure.

Let us now consider the Banzhaf based contribution:

Definition 7. Consider a set of agents A , a communication set Ψ on A , an integrity constraint IC and an IC -inconsistency measure I_C . Function $Cont_b$ associates any agent a with a positive real number $Cont_b(a)$ so that:

$$Cont_b(a) = |\{C : I_C(\Psi(C \cup \{a\})) \neq 0 \text{ and } I_C(\Psi(C)) = 0\}|.$$

As for the Shapley value, function $Cont_b$ is syntax independent as long as the inconsistency measure is syntax weak independent.

Proposition 3. Function $Cont_b$ is a syntax independent contribution function if I_C is a syntax-weak independent IC -inconsistency measure.

4.3. Implementation and complexity considerations

4.3.1. Implementation of $Cont_s$

Algorithm 1 implements function $Cont_s$. Its parameters are a communication set Ψ , an integrity constraint IC and a list of inconsistency measures $I_C list$. It shows that if we only consider inconsistency measures based on minimal inconsistent subsets, then the Shapley values for all the inconsistency measures can be computed without adding extra calls to a SAT solver. It is composed of two steps. First step (Line 3) is a call to Function *InconsistencyMeasures* (detailed in **Algorithm 2**) which computes, for each coalition of agents, the inconsistency measures requested in $I_C list$ and stores the results in a tuple of arrays indexed by the coalition number. Second step (Lines 4 to 16) computes, for each agent, the Shapley values for each inconsistency measure by applying simple operations on the arrays computed at the previous step. Clearly, the second step is not computationally hard except that it contains an internal loop on each coalition which is exponential with respect to the number of agents. This second step is precisely in $O(n * 2^{n-1}/2)$, n being the number of agents and considering that each agent is involved in half of the possible coalitions. It doesn't require any call to a SAT solver.

Algorithm 2 is composed of a preparation step and an execution step. The preparation step (Lines 3 to 5) is executed only if inconsistency measures based on minimal inconsistent subsets are requested and computes the set of minimal inconsistent subsets for the coalition composed of the whole set of agents ($\Psi \perp IC$). Computing the set of minimal inconsistent subsets of a belief base is a hard problem which requires to use a SAT solver. It is shown in [26,12] that this problem is DP-complete.

Algorithm 1 Contribution function $Cont_s$.

```
1: function  $Cont_s(\Psi, IC, I_{IC}list)$ 
2:    $n = |\Psi|$ 
3:    $(I_{IC}^C, I_{IC}^P, I_{IC}^B, Ratio) = InconsistencyMeasures(\Psi, IC, I_{IC}list)$ 
4:   for  $j = 0$  to  $n - 1$  do
5:     for  $i = 1$  to  $2^n - 1$  do ▷ for each agent
6:       if  $Bit(j, BinString(i)) = 1$  then ▷ for each coalition
7:         switch  $I_{IC}list$  do
8:           case  $I_{IC}^C$ 
9:              $Cont_s^{I_{IC}^C}[j] = Cont_s^{I_{IC}^C}[j] + Ratio[i] * (I_{IC}^C[i] - I_{IC}^C[i - 2^j])$ 
10:          case  $I_{IC}^P$ 
11:             $Cont_s^{I_{IC}^P}[j] = Cont_s^{I_{IC}^P}[j] + Ratio[i] * (I_{IC}^P[i] - I_{IC}^P[i - 2^j])$ 
12:          case  $I_{IC}^B$ 
13:             $Cont_s^{I_{IC}^B}[j] = Cont_s^{I_{IC}^B}[j] + Ratio[i] * (I_{IC}^B[i] - I_{IC}^B[i - 2^j])$ 
14:          end if
15:        end for
16:      end for
17:    return  $(Cont_s^{I_{IC}^C}, Cont_s^{I_{IC}^P}, Cont_s^{I_{IC}^B})$ 
18:  end function
```

Algorithm 2 Computing inconsistency measures for each coalition.

```
1: function  $InconsistencyMeasures(\Psi, IC, I_{IC}list)$ 
2:    $n = |\Psi|$ 
3:   switch  $I_{IC}list$  do
4:     case  $I_{IC}^C$  or  $I_{IC}^P$  or  $I_{IC}^Q$ 
5:        $\Psi \perp IC = Compute_{\Psi \perp IC}(\Psi, IC)$ 
6:     for  $i = 0$  to  $2^n - 1$  do ▷ for each coalition
7:       switch  $I_{IC}list$  do
8:         case  $I_{IC}^C$  or  $I_{IC}^P$  or  $I_{IC}^Q$ 
9:            $MI[i] = MIC(i, \Psi \perp IC)$ 
10:        switch  $I_{IC}list$  do
11:          case  $I_{IC}^C$ 
12:             $I_{IC}^C[i] = |MI[i]|$ 
13:          case  $I_{IC}^P$ 
14:             $I_{IC}^P[i] = Compute_{I_P}(MI[i])$ 
15:          case  $I_{IC}^B$ 
16:             $I_{IC}^B[i] = Compute_{I_B}(C_i)$ 
17:           $Ratio[i] = \frac{(|C_i| - 1)!(n - |C_i|)!}{n!}$ 
18:        end for
19:      return  $(I_{IC}^C, I_{IC}^P, I_{IC}^B, Ratio)$ 
20:    end function
```

Fortunately, algorithms such as those published in [12] make the computation of minimal inconsistent subsets feasible in real-life applications. The execution step (Lines 6 to 18) computes, for each coalition of agents, the requested inconsistency measures plus the ratio based on the coalition size. We use function MIC (detailed later) which allows to extract for any coalition of agents the corresponding set of minimal inconsistent subsets out of $\Psi \perp IC$ without any more call to a SAT solver. Thanks to this result computing I_{IC}^C , I_{IC}^P or I_{IC}^Q is very easy at this level. On the other hand computing I_{IC}^B still requires to compute the three valued models of the current coalition, which is computationally hard. The execution sub-step is in $O(2^{n-1})$ and is not computationally hard if we only compute inconsistency measures based on minimal inconsistent subsets. Note that in [12], the set of maximal consistent subsets is computed before producing the set of minimal inconsistent subsets. Consequently, inconsistency measure I_{IC}^M can also be computed without extra complexity. It is not the case for inconsistency measure I_{IC}^B .

In summary, computing $Cont_s$ is $C_{preparation\ step} + C_{execution\ step} + C_{second\ step}$ where $C_{preparation\ step}$ is DP-complete, $C_{execution\ step}$ is in $O(2^{n-1})$ and $C_{second\ step}$ is in $O(n * 2^{n-1}/2)$. In [16] the authors show that in the specific case using inconsistency measure $I_C(K)$, it is possible to compute the Shapley value in polynomial time. This result can directly be applied for computing $Cont_s$ using inconsistency measure I_{IC}^C .

Let us detail some hints used by Algorithm 1. Suppose there are n agents named $0, \dots, n - 1$. Then there are 2^n possible coalitions of agents we name $0, 1, \dots, 2^n - 1$ with 0 being the empty coalition and $2^n - 1$ the whole set of agents. We introduce two functions $BinString$ and Bit . $BinString(i)$ returns a binary string corresponding to the binary representation of coalition i . $Bit(j, BinString(i))$ returns the bit of weight 2^j in $BinString(i)$. Thus, agent j is involved in coalition i iff $Bit(j, BinString(i)) = 1$.

Let us now detail some hints used by Algorithm 2. Line 5 calls function $Compute_{\Psi \perp IC}(\Psi, IC)$ which returns $\Psi \perp IC$. As previously mentioned, this is DP-complete. One can notice that once $\Psi \perp IC$ is computed, we immediately get $\Psi' \perp IC$ for

any Ψ' included in Ψ i.e removing the report of some agent from Ψ does not create new inconsistent subset. This remark allows us to drastically reduce the complexity for producing the minimal inconsistent subsets of each coalition. Line 9 is a call to function $MIC(i, \Psi \perp IC)$ which returns the set of minimal inconsistent subsets of Ψ which are made of reports emitted by agents in coalition i , that is $\{X \mid X \subseteq \Psi(Agents(i)) \text{ and } X \in \Psi \perp IC\}$ ($Agents(i)$ returns the agents involved in coalition i , i.e. agents whose bit is set to 1 in $BinString(i)$).

4.3.2. Implementation of $Cont_b$

Algorithm 3 provides an implementation of function $Cont_b$. It takes as parameter a communication set Ψ and an integrity constraint IC . This algorithm is similar to **Algorithm 1** and has the same computational complexity. Again, most of the complexity is in the computation of the inconsistency measures for each possible coalition. However it is interesting to notice that the Banzhaf-based contribution uses only two inconsistency values, namely > 0 and $= 0$. Consequently, the drastic inconsistency measure, whose computation is the cheapest one, should be preferred to compute the Banzhaf-based contribution of the agents.

Algorithm 3 Contribution function $Cont_b$.

```

1: function  $Cont_b(\Psi, IC, I_{IC})$ 
2:    $n = |\Psi|$ 
3:    $(I_{IC}) = InconsistencyMeasures(\Psi, IC, I_{IC})$ 
4:   for  $j = 0$  to  $n - 1$  do
5:     for  $i = 1$  to  $2^n - 1$  do
6:       if  $Bit(j, BinString(i)) = 1$  then
7:         if then  $(I_{IC}[i] > 0 \text{ and } I_{IC}[i - 2^j] = 0)$ 
8:            $Cont_b^{I_{IC}}[j] = Cont_b^{I_{IC}}[j] + 1$ 
9:         end if
10:      end if
11:    end for
12:  end for
13:  return  $(Cont_b^{I_{IC}})$ 
14: end function

```

▷ for each agent
▷ for each coalition

5. Assessing reliability

We now consider the question of assessing reliability. In the following, we represent reliability as a preorder over the set of agents and statement $a \leq b$ stands for b is at least as reliable as a . The following postulates axiomatically characterize any reliability preorder based on what the agents have reported in the context of a given integrity constraint.

Given a set of agents A , an integrity constraint IC and a communication set Ψ , the total preorder representing the relative reliability of agents in A is denoted $\leq_{IC}^{A, \Psi}$. This preorder is characterized by the following postulates which show that reliability should be rooted in inconsistency:

- P1** $\leq_{IC}^{A, \Psi}$ is a total preorder on A .
- P2** If $\Psi \equiv \Psi'$ then $\leq_{IC}^{A, \Psi} = \leq_{IC}^{A, \Psi'}$.
- P3** If $\models IC \leftrightarrow IC'$ then $\leq_{IC}^{A, \Psi} = \leq_{IC'}^{A, \Psi}$.
- P4** If $a \notin Problematic(A)$ then $\forall b, c \in A$, if $b \leq_{IC}^{A \setminus \{a\}, \Psi \setminus \{K(a)\}} c$ then $b \leq_{IC}^{A, \Psi} c$.
- P5** If Ψ is IC -consistent then $\leq_{IC}^{A, \Psi}$ is the equality preorder.
- P6** If Ψ is IC -inconsistent then $\forall a \in Problematic(A), \forall b \notin Problematic(A), a <_{IC}^{A, \Psi} b$.
- P7** If $\{a_1, \dots, a_k\} \in A \perp IC$ for $k \geq 2$, then $\exists i, j$ such that $j \neq i, a_i <_{IC}^{A, \Psi} a_j$.

Postulate **P1** specifies that the reliability preorder is a total preorder. **P2** and **P3** deal with syntax independence. More precisely, if we consider two equivalent communication sets or if we consider two equivalent integrity constraints, then we get the same total preorder on agents. **P4** states that reliability is assessed with respect to inconsistency: in other words, agents which do not cause inconsistency issues have no influence. A typical example is an agent which reports a tautology or which reports no information, then it should not have influence on the relative reliability of other agents. **P5**, **P6** and **P7** focus on consistency of information provided by agents in A . Postulate **P5** considers the case when Ψ set is not IC -inconsistent. In such a case, the agents are considered as equally reliable. **P6** and **P7** consider the case where Ψ is IC -inconsistent. According to **P6**, any agent which is responsible of the IC inconsistency of Ψ is considered as strictly less reliable than any other agent which is not. According to **P7**, the agents of a minimal IC -inconsistent subset cannot be equally reliable: at least one of these agents is strictly less reliable than the others. This is coherent with the way we understand reliability: if some agents are equally reliable, then after merging we will believe, with the same strength, information they will provide. However, it is generally assumed [7,23] that graded belief satisfies the modal logic axiom which states that belief should be consistent: that is, two contradictory pieces of information cannot be believed with the same strength. Consequently, agents who are involved in a minimal IC -inconsistent set cannot be equally reliable.

5.1. Building reliability from contribution

It is clear that the more an agent contributes to the overall inconsistency of a communication set, the less is should be reliable. That is, a source is considered strictly more (resp, equally) reliable than a second one iff its contribution to the global inconsistency is strictly smaller than (resp equal to) the contribution of the other. Notice that this principle only takes care of the first six postulates as shown by the following postulates.

Theorem 1. Given a set of agents A , an integrity constraint IC and a communication set Ψ , the reliability preorder $\leq_{IC}^{A,\Psi}$ satisfies **P1–P6** iff there exists a syntax independent Contribution Function $Cont$ such that for any two agents a and b :

$$a \leq_{IC}^{A,\Psi} b \quad \text{iff} \quad Cont(a) \geq Cont(b).$$

The following example illustrates that **P7** does not hold.

Example 3. Consider $\Psi = \{ \langle a, p \rangle, \langle b, \neg p \rangle, \langle c, q \rangle \}$ and $IC = \top$. Then $\{a, b\}$ is a minimal IC -conflicting set of agents but Shapley-based Contribution Function gives: $Cont_s^{\Psi, IC}(a) = Cont_s^{\Psi, IC}(b)$ and consequently $a = b$ which violates **P7**.

This example shows that a tie-breaking rule is missing: if all agents involved in the definition of a minimal inconsistent subset have an equal contribution then one of them should still be considered as less reliable. For voting rules, a classical way to handle tie-breaking is to consider an additional lexicographic order over the set of agents in order to always get a winner [2]. In our context, the contribution function should prevent the cases where all agent involved in an inconsistent coalition have similar contribution. The following constraint translates this principle in formal terms:

Definition 8. Let $Cont$ be a syntax independent Contribution Function. $Cont$ is said to be *tie-free* if it satisfies the following constraint:

$$\forall C \in A \perp IC, \exists a, b \in C, \text{ s.t. } Cont(a) \neq Cont(b).$$

The immediate question is “does such function exist?”. To get the answer, let us revisit our Shapley and Banzhaf based contribution functions and built on top of them a *tie-free* function. The idea is similar tie-breaking: for any inconsistent set of agents, one agent contribution is slightly increased so that there is no more inconsistent set of agents where all agents have the same contribution. As we have no information about the agents involved in the communication set, the modified contribution may be chosen in an arbitrary way. Hereafter we just consider a lexicographic order for this choice (as in the voting rules).

Definition 9. Let $Cont$ be a Contribution function defined either by [Definitions 6 and 7](#). Function $Cont_t$ is then defined as follows:

- If $Problematic(A) = \emptyset$ then $Cont_t = Cont$.
- Otherwise, let ϵ be a positive real number which is strictly smaller than the smallest difference between two contributions: $0 < \epsilon < \min(\{Cont(a) - Cont(b) \mid Cont(a) \geq Cont(b)\})$.
 1. Let S_{tie} be the set of agents with equal contribution t :

$$S_{tie} = \{a \mid \exists C \in A \perp IC \text{ such that } |C| > 1 \text{ and } a \in C \text{ and } \forall x, y \in C, Cont(x) = Cont(y)\}.$$

Let S be the minimal subset of S_{tie} w.r.t. the lexicographic order such that (i) $\forall C \in A \perp IC, \exists a \in S \cap C$ and (ii) $\forall C \in A \perp IC, C \setminus S \neq \emptyset$. Set $Cont_t(a) = Cont(a) + \epsilon$ for any $a \in S$.

2. $Cont_t(b) = Cont(b)$ for any agent $b \notin S$.

The following proposition shows that this function is a *tie-free* one.

Proposition 4. $Cont_t$ is tie-free.

We have proved that there exist Contribution functions which are tie-free, we can rephrase the previous theorem so that **P7** holds.

Theorem 2. Given a set of agents A , an integrity constraint IC and a communication set Ψ , the reliability preorder $\leq_{IC}^{A,\Psi}$ satisfies **P1–P7** iff there exists a syntax independent Contribution Function $Cont$ which is tie-free such that for any two agents a and b :

$$a \leq_{IC}^{A,\Psi} b \quad \text{iff} \quad Cont(a) \geq Cont(b).$$

Example 4. Take again $A = \{a, b\}$ and $K(a) = p$, $K(b) = \neg p$ and $IC = \top$. $\{a, b\} \in A \perp IC$ and, according to the Shapley based Contribution function, it holds that $Cont_s(a) = Cont_s(b)$. Let us consider the tie-free Contribution function based on $Cont_s$. According to the lexicographic agent a then, $Cont_t(a) > Cont_t(b)$ and we get: $a < b$.

This second theorem shows at first that there exists a reliability preorder which satisfy all the postulates. Second, it shows that one of the key issue in the inconsistency definition of reliability is Postulate **P7**. It forces to rank the agents involved in a minimal IC -inconsistent subset. We propose to handle this issue by constraining the contribution function. Remind that all agents are unknown and we do not have extra information for setting the choice. Hence this function may lead to arbitrary choices. However, even if we do not know any extra information, we have seen that we can build up a reliability preorder by first assessing the overall inconsistency of a communication set; secondly by computing the “responsibility” degree or contribution of each agent in the overall inconsistency and thirdly by rooting the reliability of the agents into their contributions.

6. Reliability aggregation

In the previous section, we have shown that assessing reliability of agents may be achieved by choosing a particular inconsistency measure and a particular contribution function, each specific pair $\langle measure, contribution \rangle$ assessing the sources in a different way. The choice of a specific pair is a challenge since we have no information about the sources. A solution may be to assess reliability with respect to several pairs and to merge all the resulting preorders to obtain a “consolidated” one. The immediate question is then: if each of the reliability preorders satisfies postulates **P1–P7**, can we then obtain an aggregated preorder which also satisfies these postulates? The aim of this section is to tackle this issue.

This problem is clearly connected to the field of preferences aggregation. Arrow impossibility theorem [3] states that it is not possible to aggregate preferences while guaranteeing Universality, Non-Dictatorship, Unanimity and Independence to Irrelevant Alternatives. The question is then to evaluate how Arrow’s conditions interplay with the reliability postulates.

6.1. Arrow’s conditions

In the following, we first define our aggregation operator and next revisit the Arrow’s conditions. Let \oplus be our aggregation operator defined as follows:

Definition 10. A n -ary reliability aggregation operator \oplus is a function which associates n total preorders on A , respectively \leq_1, \dots, \leq_n , with a total preorder on A denoted \leq_{\oplus} .

Next, we rephrase the classical Arrow’s conditions:

Definition 11. Let \oplus be a n -ary reliability aggregation operator. We consider the following properties:

Universality the domain of \oplus is the set of all possible n -tuple of reliability total preorders.

Non-dictatorship $\nexists i \in \{1 \dots n\}$ such that $\forall \leq_1 \dots \leq_n$ total preorders on A , $\leq_{\oplus} = \leq_i$.

Unanimity Let $a \in A$ and $b \in A$. If $\forall i \in \{1 \dots n\}$ $a \leq_i b$ then $a \leq_{\oplus} b$.

Independence of irrelevant alternatives (IIA) Let a and b be two agents. Let (\leq_1, \dots, \leq_n) and $(\leq'_1, \dots, \leq'_n)$ be two sets of n total preorders. If $\forall i \in \{1 \dots n\}$, $(a \leq_i b \Leftrightarrow a \leq'_i b)$ then $(a \leq_{\oplus} b \Leftrightarrow a \leq'_{\oplus} b)$.

6.2. Influence of arrow’s conditions

The following propositions exhibit the interplay between Arrow’s conditions and the seven postulates characterizing the consistency-based reliability assessment. Namely, assuming that each \leq_i satisfies our postulates, which Arrow’s condition(s) are required so that the postulates are also satisfied by the resulting preorder.

Propositions 5 to 11 all use the following elements: let Ψ be a communication set on A and IC an integrity constraint; let I_{IC}^i ($i = 1 \dots n$) be n syntax weak-independent IC -inconsistency measures. For the sake of conciseness, \leq_i ($i = 1 \dots n$) stands for $\leq_{IC}^{A, \Psi}$ i th reliability preorder.

Proposition 5 (P1 $^{\oplus}$). If \oplus satisfies the condition of **Universality** then

If all \leq_i satisfies **P1**, then \leq_{\oplus} satisfies **P1** (i.e., \leq_{\oplus} is a total preorder).

This proposition states that if the input is itself a tuple of total preorders **Universality** will guarantee that the aggregation operator \oplus will return a total preorder.

Proposition 6 (P2 $^{\oplus}$). If all \leq_i satisfy **P2** then \leq_{\oplus} satisfies **P2** (i.e., if $\Psi \equiv \Psi'$ then $\leq_{\oplus} = \leq'_{\oplus}$).

Table 1
Interplay between postulates and Arrow's conditions.

Sufficient conditions	Satisfied postulates
Universality	{ P1 }
Dictatorship	{ P7 }
Unanimity	{ P5, P6 }
IIA	{ P4 }

This proposition is an immediate consequence of Postulate **P2**: if each reliability assessment operator is syntax independent (w.r.t. communication set) then the resulting preorder is also syntax independent.

Proposition 7 (P3[⊕]). If all \leq_i satisfy **P3** then \leq_{\oplus} satisfies **P3** (i.e., for any IC' s.t. $\models IC \leftrightarrow IC'$, $\leq_{\oplus} = \leq'_{\oplus}$).

This proposition is a consequence of Postulate **P3**: if each \leq_i is syntax independent w.r.t. some integrity constraint then the resulting preorder is also syntax independent. Notice that the last two propositions do not involve any Arrow conditions as opposed to the other ones.

Proposition 8 (P4[⊕]). If \oplus satisfies the condition of **Independence of Irrelevant Alternatives** then

If all \leq_i satisfy **P4** then \leq_{\oplus} satisfies **P4** (i.e., if \leq^{-a} denotes the preorder $\leq_{IC}^{A \setminus \{a\}, \Psi \setminus K(a)}$, if $a \notin \text{Problematic}(A)$, then $\forall b, c \in A$, if $b \leq_{\oplus}^{-a} c$ then $b \leq_{\oplus} c$).

This proposition states that **P4** (no influence of agent a who only reports tautologies) can only be preserved if **IIA** holds. That is, if all preorders are unchanged after excluding agent a , then aggregation also produces a similar result if **IIA** also holds.

Proposition 9 (P5[⊕]). If \oplus satisfies the condition of **Unanimity** then

If all \leq_i satisfy **P5** then \leq_{\oplus} satisfies **P5** (i.e., if Ψ is IC-consistent then \leq_{\oplus} is the equality preorder).

Postulate **P5** states that consistency leads to a flat ordering. Consequently, if there is no conflict among agents, every \leq_i has to be flat and the aggregation produces a flat order as long as **Unanimity** holds for the operator \oplus .

Proposition 10 (P6[⊕]). If \oplus satisfies the condition of **Unanimity** then

If all \leq_i satisfy **P6** then \leq_{\oplus} satisfies **P6** (i.e., if Ψ is IC-inconsistent then $\forall a \in \text{Problematic}(A)$, $\forall b \notin \text{Problematic}(A)$, $a <_{\oplus} b$).

Proposition 10 stresses up a second time the key role of **Unanimity** condition. If Postulate **P6** holds for every \leq_i then preorders \leq_i unanimously states that for any $a \in \text{Problematic}(A)$ and $b \in A \setminus \text{Problematic}(A)$, b is more reliable than a . $a <_{\oplus} b$ also holds only if **Unanimity** holds.

Proposition 11 (P7[⊕]). If \oplus does **not** satisfy the condition of **Non-dictatorship** then

If all \leq_i satisfy **P7** then \leq_{\oplus} satisfies **P7** (i.e., if $\{a_1, \dots, a_k\} \in A \perp IC$ for $k \geq 2$, then $\exists i, j$ such that $j \neq i$ and $a_i <_{\oplus} a_j$).

This proposition shows the role of the **Non-dictatorship** condition. Postulate **P7** states that at least one agent involved in a conflicting set of agents should be ranked with a lower reliability. If each \leq_i have decreased the reliability of a different agent then, if the **Non-dictatorship** property holds, no agent will be decreased by the aggregation procedure \oplus . Hence, the constraint enforces by **P7** will hold only if the **Non-dictatorship** property is not satisfied, i.e. there is a dictator.

To sum up, as shown by **Table 1** four Arrow's conditions have an influence on the fulfillment of the postulates. Notice that these conditions are sufficient conditions for satisfying the postulates.

With the help of **Table 1**, we are fully informed on what postulates or Arrow's conditions one should give up during the definition of an aggregation operator. The table shows that reliability aggregation is different from preference aggregation: the postulate are not compatible with Arrow's conditions. It also means that some underlying priority should be considered in the definition of the operator: if we give priority to the fulfillment of the postulates then fulfillment of Arrow's conditions is not that important. In the next section, we illustrate this issue by considering a simple aggregation procedure satisfying all the postulates but abandoning one Arrow's condition.

6.3. Lexicographic-based aggregation

The proposed aggregation procedure is a lexicographic procedure [14], widely used in preference aggregation and belief merging. In [2], the authors show how preference relations can be aggregated. First consider a hierarchy among n preorders (with no loss of generality, \leq_1 is the most important while \leq_n is the least important). According to this hierarchy, if all k first preorders state $a =_k b$ and at $k + 1$, $a <_{k+1} b$ then a will be considered as less reliable than b in the resulting preorder. No divergence entails that a and b are considered as equal in the resulting preorder. Let \oplus^L denote the reliability aggregation operator defined as follows:

Definition 12. If \leq_1, \dots, \leq_n are n reliability preorders on a set of agents A and if \leq^L denotes the preorder $\oplus^L(\leq_1 \dots \leq_n)$, then \leq^L is defined by:

- $a <^L b$ iff $\exists k = 1 \dots n \forall j \in \{1 \dots k - 1\} a =_j b$ and $a <_k b$.
- $a =^L b$ else.

Example 5. Assume the following total preorders defined over the set of agents $A = \{a, b, c, d, e\}$.

$$a =_1 b <_1 c <_1 d =_1 e.$$

$$a <_2 b =_2 c <_2 d =_2 e.$$

$$e =_3 b <_3 c =_3 d =_3 a.$$

We obtain the following aggregated preorder $a <^L b^L c <^L e <^L d$. For instance, $a <^L b$ holds because $a =_1 b$ and $a <_2 b$.

It is well known that among the four Arrow's conditions, **Non-dictatorship** does not hold for Lexicographic aggregation (see [2]). I.e.,

Theorem 3. [14] \oplus^L satisfies **Universality, Unanimity, Independence of Irrelevant Alternatives** and does not satisfy **Non-dictatorship**.

Corollary 1. If all preorders $<_i$ satisfy Postulates **P1–P7**, then Preorder $<^L$ also satisfies Postulates **P1–P7**.

It means that this reliability aggregation operator is a good candidate to aggregate reliability preorders as long as these individual preorders also satisfy **P1–P7**.

Example 6. Consider agent a whose aim is to assess the relative reliability of three communicating agents b, c and d . With the specific pair $\langle \text{measure}, \text{contribution} \rangle$ it chose, a gets the preorder: $b <_s c =_s d$, i.e., b has to be considered as strictly less reliable than c and d who have to be considered as equally reliable. Assume that before applying this method, a had an a priori about the three agents and thought that d was more reliable than the two others which were equally reliable, i.e., $b =_a c <_a d$ (where $<_a$ denotes the a priori preorder). This a priori information can be used by a to decide whether c or d is the least reliable. More precisely, by aggregating the two preorders and by giving priority to $<_s$, a gets: $b <^L c <^L d$.

6.4. Avoiding dictatorship condition

If we want to consider each inconsistency measure in an equal way, we should avoid the dictatorship condition and instead go towards voting rules for the definition of a non-dictator reliability aggregation operator. An immediate consequence is that postulate **P7** will not be satisfied in the aggregated preorder. Another consequence is that Arrow's theorem forces us to give up an other condition since we want the non-dictatorship condition to be satisfied. It means that some other postulates will also have to be abandoned depending on the condition that will be given up (**IIA** or **Unanimity**).

Let us consider the classical preferential voting rules such as Condorcet, Borda and Copeland [13], which enforce the non-dictatorship condition. All these rules usually consider strict and total preorders as input and provide a preorder as output. In our context, we have to consider a variant where input is a set of preorders that may not be strict. We also consider that the aggregation method should take into account the relative position of each agent (w.r.t. other agents) in all preorders: agent a may be poorly ranked according to some inconsistency measure while it may obtain a good rank according to some other inconsistency measure. Hence the rank of an agent should be balanced by its position in the overall preorder. It means that a counting based procedure is more adapted than a pairwise comparison method, which argues for a Borda based aggregation procedure instead of an aggregation procedure based on Condorcet or Copeland.

The rule proposed hereafter rephrases the Borda-based rule initially proposed in [13, Chapter 13] which handles total preorders (that may be strict). As mentioned, we want to take into account the relative position in the preorder ("good" or "poor" position). Consequently, for any agent a and preorder \leq_i , the rule computes a score which takes into account its relative position w.r.t. the median agent associated to preorder \leq_i .

Let us detail the computation of the score of agent a , w.r.t. \leq_i . Set of agent A is first partitioned w.r.t. \leq_i such that in each partition all agents are equally reliable; let t be the number of partitions. The partitions are ranked w.r.t. \leq_i such that P_1 is the partition containing the most reliable agents and P_t contains the less reliable agent. Let m be the median agent; the score of each agent is set according to its relative position with m w.r.t. the partitions. w.r.t. some preorder \leq :

$$\text{score}(a, \leq) = (k - j) \text{ iff } a \in P_j \text{ and } m \in P_k.$$

The median agent m receives a score equal to zero, while each agent strictly more reliable than m gets a positive score and agents strictly less reliable than m receive a negative score.

Notice, that the score also gives privileges to preorders which are informative (i.e. discriminating agents), that is if the number of partitions is large then the range of the score is also large.

Definition 13. If \leq_1, \dots, \leq_n are n reliability preorders on a set of agents A and if \leq^B denotes the preorder $\oplus^B(\leq_1 \dots \leq_n)$, then \leq^B is defined by:

$$a \leq^B b \iff \sum_{i=1}^n \text{score}(a, \leq_i) \leq \sum_{i=1}^n \text{score}(b, \leq_i).$$

Example 7. Let us revisit the previous example Assume the following total preorders defined over the set of agents $A = \{a, b, c, d, e\}$. We present the partition and the associated score

$$\begin{array}{ll} \{a, b\} <_1 \{c\} <_1 \{d, e\} & [1 \ 0 \ -1] \\ \{a\} <_2 \{b, c\} <_2 \{d, e\} & [1 \ 0 \ -1] \\ \{e, b\} <_3 \{c, d, a\} & [1 \ 0] \end{array}$$

We get the following scores:

$$\text{score}(a) = 2, \quad \text{score}(b) = 2, \quad \text{score}(c) = 0, \quad \text{score}(d) = -2, \quad \text{score}(e) = -1.$$

We then obtain the following aggregating preorder $a =^B b <^B c <^B e <^B d$. The key difference with the previous ranking is the equal rank for a and b : as they are both rank first twice in this example, the absence of dictator entails to not distinguish them.

We remind that among the four Arrow's conditions, **IIA** does not hold for Borda aggregation.

Theorem 4. [14] \oplus^B satisfies **Universality, Unanimity, Non-dictatorship** and does not satisfy **Independence of Irrelevant Alternatives**.

Corollary 2. If all preorders $<_i$ satisfy Postulates **P1–P7**, then preorder $<^B$ only satisfies Postulates **P1–P3** and **P5–P6**.

As shown by the corollary, the main consequence is that even if all individual preorders satisfy Postulates **P1–P7**, the properties corresponding to **P4** and **P7** do not hold in the resulting preorder. For postulate **P4**, it means that an agent only reporting tautologies may influence the output, i.e. the overall reliability. Second, for postulate **P7**, it means that agents involved in some minimal *IC*-inconsistent subset may not be discriminated with the aggregated reliability ranking while they are supposed.

7. Conclusion

This work proposes to assess the relative reliability of some agents by analyzing the inconsistency of information they report w.r.t some trusted knowledge. We have defined postulates stating what the relative reliability preorder should be. Then we have shown through two representation theorems how such preorders are related to the contribution of each agent to the overall inconsistency of a communication set. Implementation and complexity questions have also been considered. Finally, we have shown how aggregating reliability preorders is constrained by Arrow's impossibility theorem.

This framework may be extended in several ways. First, inconsistency measures should deserve more attention. Recent work on this topic [11,4] shows promising results such as giving a weight to the inconsistency itself. A second issue concerns our key principle for reliability assessment which only considers inconsistency: the more an agent is connected to inconsistency, the less it is reliable. We could extend this principle by considering not only conflicts between agents but also agreement [28]. We could also extend it by considering that an agent producing more information than others should be rewarded. The problem here will be how to quantify the amount of information produced by an agent which is not contradicted. To this end some information measures have already been proposed in [10]. The main question will be to revise the postulates to take into account these extensions. The third issue concerns the one shot dimension of the assessment process: iteration should be possible and reliability assessment should then be viewed as a refinement process.

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Appendix A. Proofs

Proof of Proposition 1. First, if $\models IC \leftrightarrow IC'$ then $\Psi \perp IC = \Psi \perp IC'$, $\Psi \top IC = \Psi \top IC'$ and $i \models_3 \bigwedge_{K \in \Psi} K \wedge IC$ iff $i \models_3 \bigwedge_{K \in \Psi} K \wedge IC'$. Secondly, if $\Psi \equiv \Psi'$ then $|\Psi \perp IC| = |\Psi' \perp IC|$, $|\Psi \top IC| = |\Psi' \top IC|$ and $i \models_3 \bigwedge_{K \in \Psi} K \wedge IC$ iff $i \models_3 \bigwedge_{K \in \Psi'} K \wedge IC$.

Proof of Proposition 2. *Consistency.* Assume $a \notin \text{Problematic}(A)$, then $K(a) \in \text{Free}(\Psi)$. Then $\forall C \ I_{IC}(\Psi(C)) = I_{IC}(\Psi(C \setminus \{a\}))$ since I_{IC} satisfies Free formula Independence. This proves that if $a \notin \text{Problematic}(A)$ then $\text{Cont}_s(a) = 0$.

Assume $a \in \text{Problematic}(A)$: if $K(a) \in \text{Problematic}(\Psi)$ then $\exists M \in \Psi \perp IC \ K(a) \in M$. Let's write $M = \{K(a), K(b_1), \dots, K(b_k)\}$. Then $\Psi(\{a, b_1, \dots, b_k\})$ is IC -inconsistent while $\Psi(\{b_1, \dots, b_k\})$ is IC -consistent. Consequently, $I_{IC}(\Psi(\{a, b_1, \dots, b_k\})) > I_{IC}(\Psi(\{b_1, \dots, b_k\})) = 0$ as I_{IC} satisfies the *Consistency* property and is monotonic. Thus $\text{Cont}_s(a) \neq 0$. This proves that if $\text{Cont}_s(a) = 0$, then $a \notin \text{Problematic}(A)$.

Free agent independence. Assume $b \notin \text{Problematic}(A)$. Then, $\forall C \subseteq A$, $I_{IC}(\Psi(C)) = I_{IC}(\Psi(C \setminus \{b\}))$ since I_{IC} satisfies Free formula Independence. Hence for any agent a , $I_{IC}(\Psi(C)) - I_{IC}(\Psi(C \setminus \{a\})) = I_{IC}(\Psi(C \setminus \{b\})) - I_{IC}(\Psi(C \setminus \{b\} \setminus \{a\}))$. Consequently $\text{Cont}_s^{\Psi, I_{IC}}(a) = \text{Cont}_s^{\Psi - \{K(b)\}, I_{IC}}(a)$.

Syntax independence. Assume IC' such that $\models IC \leftrightarrow IC'$. Then $\forall C$, $I_{IC}(\Psi(C)) = I_{IC'}(\Psi(C))$ since I_{IC} satisfies the syntax weak independence property. Hence for any agent a , $I_{IC}(\Psi(C)) - I_{IC}(\Psi(C \setminus \{a\})) = I_{IC'}(\Psi(C)) - I_{IC'}(\Psi(C \setminus \{a\}))$. Consequently $\text{Cont}_s^{\Psi, I_{IC}}(a) = \text{Cont}_s^{\Psi, I_{IC'}}(a)$.

Assume $\Psi \equiv \Psi'$. Then $\forall C \ \Psi(C) \equiv \Psi'(C)$. Thus $\forall C \ \Psi(C) \equiv \Psi'(C)$. Since I_{IC} satisfies the syntax weak independence property, this implies that $\forall C \ I_{IC}(\Psi(C)) = I_{IC}(\Psi'(C))$. Hence for any agent a , $I_{IC}(\Psi(C)) - I_{IC}(\Psi(C \setminus \{a\})) = I_{IC}(\Psi'(C)) - I_{IC}(\Psi'(C \setminus \{a\}))$. Then $\text{Cont}_s^{\Psi, I_{IC}}(a) = \text{Cont}_s^{\Psi', I_{IC}}(a)$.

Hence Cont_s is a syntax independent contribution function.

Proof of Proposition 3. The proof is similar to the proof of Proposition 2.

Consistency. Assume $a \notin \text{Problematic}(A)$, then $K(a) \in \text{Free}(\Psi)$. Then $\forall C \ I_{IC}(\Psi(C)) = I_{IC}(\Psi(C \setminus \{a\}))$ since I_{IC} satisfies Free formula Independence. Hence there is no C where a is pivotal. This proves that if $a \notin \text{Problematic}(A)$ then $\text{Cont}_b(a) = 0$.

Assume $a \in \text{Problematic}(A)$: if $K(a) \in \text{Problematic}(\Psi)$ then $\exists M \in \Psi \perp IC \ K(a) \in M$. Let's write $M = \{K(a), K(b_1), \dots, K(b_k)\}$. Then $\Psi(\{a, b_1, \dots, b_k\})$ is IC -inconsistent while $\Psi(\{b_1, \dots, b_k\})$ is IC -consistent. Consequently, $I_{IC}(\Psi(\{a, b_1, \dots, b_k\})) > I_{IC}(\Psi(\{b_1, \dots, b_k\})) = 0$ as I_{IC} satisfies the *Consistency* and *Monotonicity* property. Thus $\text{Cont}_b(a) \neq 0$ as a is pivotal for coalition $C = \{b_1, \dots, b_k\}$. This proves that if $\text{Cont}_b(a) = 0$, then $a \notin \text{Problematic}(A)$.

Free agent independence. Let $b \notin \text{Problematic}(A)$. Then $\forall C \ I_{IC}(\Psi(C)) = I_{IC}(\Psi(C \setminus \{b\}))$ since I_{IC} satisfies Free formula Independence. Hence for any agent a , $I_{IC}(\Psi(C)) - I_{IC}(\Psi(C \setminus \{a\})) = I_{IC}(\Psi(C \setminus \{b\})) - I_{IC}(\Psi(C \setminus \{b\} \setminus \{a\}))$. Consequently $\text{Cont}_b^{\Psi, I_{IC}}(a) = \text{Cont}_b^{\Psi - \{K(b)\}, I_{IC}}(a)$.

Syntax-weak independence.

Syntax independence. Let IC' be such that $\models IC \leftrightarrow IC'$. Then $\forall C \ I_{IC}(\Psi(C)) = I_{IC'}(\Psi(C))$ since I_{IC} satisfies the syntax weak independence property. Hence for any agent a , $I_{IC}(\Psi(C)) - I_{IC}(\Psi(C \setminus \{a\})) = I_{IC'}(\Psi(C)) - I_{IC'}(\Psi(C \setminus \{a\}))$. Consequently $\text{Cont}_b^{\Psi, I_{IC}}(a) = \text{Cont}_b^{\Psi, I_{IC'}}(a)$.

Assume $\Psi \equiv \Psi'$. Then $\forall C$, $\Psi(C) \equiv \Psi'(C)$, thus $\forall C$, $\Psi(C) \equiv \Psi'(C)$, thus $\forall C$, $I_{IC}(\Psi(C)) = I_{IC}(\Psi'(C))$ since I_{IC} satisfies the syntax weak independence property. Hence for any agent a , $I_{IC}(\Psi(C)) - I_{IC}(\Psi(C \setminus \{a\})) = I_{IC}(\Psi'(C)) - I_{IC}(\Psi'(C \setminus \{a\}))$. Then $\text{Cont}_b^{\Psi, I_{IC}}(a) = \text{Cont}_b^{\Psi', I_{IC}}(a)$.

Consequently, Cont_b is syntax independent.

A.1. Theorems about reliability preorders

Proof of Theorem 1. Let us first prove that Postulates **P1–P6** hold. In the following, whenever it's clear \leq denotes the reliability preorder instead of full notation $\leq_{IC}^{A, \Psi}$.

- Function $Cont$ outputs numerical values and thus preorder \leq is total **P1** is satisfied.
- If $\Psi \equiv \Psi'$, then $\forall C \subseteq A \ \Psi(C) \equiv \Psi'(C)$. As $Cont$ satisfies syntax independence, then, $\forall x \in A \ Cont^{\Psi, IC}(x) = Cont^{\Psi', IC}(x)$. Thus, **P2** holds.
- If $\models IC \leftrightarrow IC'$ then, because $Cont$ is syntax independent, $\forall a \subseteq A, \ Cont^{\Psi, IC}(a) = Cont^{\Psi, IC'}(a)$. Consequently, **P3** holds.
- Assume a is not problematic: $K(a) \in Free(\Psi)$. Next function $Cont$ is free agent independent, $\forall x, y \in A \ Cont^{\Psi, IC}(x) = Cont^{\Psi \setminus \{K(a)\}, IC}(x)$ and $Cont^{\Psi, IC}(y) = Cont^{\Psi \setminus \{K(a)\}, IC}(y)$. Then, if $x \leq_{IC}^{A \setminus \{a\}, \Psi \setminus \{K(a)\}} y$ then it must be the case that $x \leq y$. This shows that **P4** is satisfied by \leq .
- If Ψ is IC -consistent, then by consistency property, $\forall x \in A, \ Cont(x) = 0$ and then $\forall x, y \in A, \ x = y$. This shows that **P5** is satisfied.
- Suppose Ψ is IC -inconsistent. Let a be in A .
 - Assume $a \notin Problematic(A)$. Then by consistency property for Function $Cont$, $Cont(a) = 0$.
 - Assume $a \in Problematic(A)$. Then by consistency property for Function $Cont$, $Cont(a) > 0$.
This shows that **P6** holds.

Let us focus on the other direction: assume a preorder $\leq_{IC}^{A, \Psi}$, or \leq for short such that **P1–P6** hold. Next, we build a Function $Cont$ as following: we partition the set of agents w.r.t. preorder \leq as follows:

- $A_0 = \min(A, \leq)$
- $A_i = \min(A - \cup_{0 \leq j < i} A_j, \leq)$ for any $i > 0$.

Notice, that as A is finite, then the number of partitions is also finite. Assume there are N partitions and then A_0 the less reliable agents while A_{N-1} contains the most reliable ones. For any agent $a \in A_k$, $Cont(a) = N - k - 1$.

Due to postulate **P1**, the immediate consequence of this definition is that for any agents a and b :

$$\text{If } a \leq_{IC}^{A, \Psi} b \text{ then } Cont(a) \geq Cont(b)$$

Hence, the first requirement stating that for any agent a , $Cont$ provides a positive real number is fulfilled.

Let us now check that Function $Cont$ is actually syntax independent.

First we check that $Cont$ is consistent. Postulate **P5** entails that there is a unique partition $A_0 = A$ and consequently $Cont(a) = 0$ for all agent a . Assume now, that some agents are problematic. Then there is at least 2 partitions: suppose there are $N > 1$ partitions. Due to postulate **P6**, only agents a such that $a \notin Problematic(A)$ belong to A_{N-1} . Hence, for all agents $a \notin Problematic(A)$, it holds that $Cont(a) = 0$ and **P6** also enforces that for all agents $a \in Problematic(A)$, $Cont(a) > 0$.

Then we check that $Cont$ is free agent independent. Assume a such $a \notin Problematic(A)$. Suppose that A is partitioned in N partitions. **P6** entails that a belongs to A_{N-1} . Due to postulate **P4**, we get that the partition of $A - \{a\}$ w.r.t. $\leq_{IC}^{A - \{a\}, \Psi - K(a)}$ is equal to the partition of A w.r.t. $\leq_{IC}^{A, \Psi}$ except for A_{N-1} which does not contain a . Hence for every other agent $b \neq a$, we get that $Cont^{\Psi, IC}(b) = Cont^{\Psi - \{K(b)\}, IC}(b)$ and thus $Cont$ is free agent independent.

Finally, we check that $Cont$ is syntax independent. Assume IC' such that $\models IC \leftrightarrow IC'$. Then **P3** entails that preorders $\leq_{IC}^{A, \Psi}$ and $\leq_{IC'}^{A, \Psi}$ are identical. Then partitions will be the same and consequently, for any agent a , $Cont^{\Psi, IC}(a) = Cont^{\Psi, IC'}(a)$. In a similar way, assume Ψ' such that $\Psi \equiv \Psi'$. **P2** entails that preorders $\leq_{IC}^{A, \Psi}$ and $\leq_{IC}^{A, \Psi'}$ are identical. Then partitions will be the same and consequently, for any agent a , $Cont^{\Psi, IC}(a) = Cont^{\Psi', IC}(a)$.

Proof of Proposition 4. If the communication set is consistent then $A \perp IC = \emptyset$ thus the constraint of tie-free contribution function holds.

If the communication set is inconsistent then we have two cases: either S_{tie} is empty or not. If S_{tie} is empty then *tie-free* holds.

Now suppose that S_{tie} is not empty. Let S be the minimal subset of S_{tie} w.r.t. the lexicographic order such that (i) $\forall C \in A \perp IC, \exists a \in S \cap C$ and (ii) $\forall C \in A \perp IC, C \setminus S \neq \emptyset$. Let us first show that such S exists. Assume it is not the case, then it entails that either condition (i) or (ii) does not hold. Suppose (i) holds but not (ii). Hence, it means there exists a coalition C such that there is no agent a which belongs to $C \setminus S$. Hence $C = S$, and S is then not minimal. Contradiction. Other cases can be considered in a similar way.

Consequently, the definition of $Cont_t$ entails that for every coalition C where agents have an equal contribution: (i) one agent has its contribution increased as it belongs to S and C and (ii) one agent has its contribution unchanged as it belongs to $C \setminus S$. Next the second step guarantees that the contribution is not increased for all other members of the set. Hence, it means that in each minimal IC -inconsistent subset of agents, there is always two agents which have unequal contribution, i.e., the constraint of tie-free contribution function holds.

Proof of Theorem 2. We only focus on Postulate **P7** as the proof for the other postulates is similar to the proof of [Theorem 1](#).

Let us first prove that the postulate is satisfied: suppose that $\{a_1, \dots, a_k\} \in A \perp IC$ for $k \geq 2$, then as $Cont$ is tie-free, it entails $\exists i, j$, such $Cont(a_i) \neq Cont(a_j)$. Then by definition of $<_{IC}^{A, \Psi}$, $a_i <_{IC}^{A, \Psi} a_j$ or $a_j <_{IC}^{A, \Psi} a_i$ holds and **P7** holds.

Now consider the other direction. Assume that we built up a function $Cont$ similar to the one considered in the proof of [Theorem 1](#). Due to **P7**, for any inconsistent coalition C such that $|C| > 1$, there always exist $a, b \in C$ such that $a < b$. Hence,

it means that there is no partition A_i such that all agents of an inconsistent coalition belong to A_i . Consequently, for any inconsistent coalition $C \in A \perp IC$, its members belong at least to two partitions. Then it means that there is no $C \in A \perp IC$ such that $|C| > 1$ where $\forall a, b \in C$, $Cont(a) = Cont(b)$. Consequently $Cont$ satisfies the tie-free constraint.

A.2. Propositions on arrow's conditions and postulates

We do not detail the proofs of [Propositions 6 and 7](#), they are obvious.

Proof of Proposition 5. Operator \oplus is defined as an operator with n total preorders as input. Postulate **P1** guarantees that each \leq_i is total. Second, **Universality** entails that all possible input $\leq_1 \cdots \leq_n$ will give an output. Definition of \oplus states that this output is a total preorder.

Proof of Proposition 8. Let a be such that $\models K(a)$. If every \leq_i satisfies **P4**, then for any x and y in A if $(x \leq_i^{-a} y)$ then $(x \leq_i y)$. Consequently, for any x and y in A , $(x \leq_i y) \iff (x \leq_i^{-a} y)$. Thus, if \leq_{\oplus} satisfies **Independence of Irrelevant Alternatives**, then it holds that $(x \leq_{\oplus} y) \iff (x \leq_{\oplus}^{-a} y)$. Thus, for any x and y in A if $(x \leq_i^{-a} y)$ then $(x \leq_i y)$, i.e., \leq_{\oplus} is compatible with \leq_{\oplus}^{-a} .

Proof of Proposition 9. Suppose that A is not IC -conflicting. Now assume that **P5** holds for all \leq_i . Then $\forall i = 1..n$, \leq_i is the equality preorder. **Unanimity** states that for all $a, b \in A$, if $\forall i \in \{1..n\} a \leq_i b$ then $a \leq_{\oplus} b$. Consequently, if **Unanimity** holds then \leq_{\oplus} is the equality preorder and **P5** holds for \leq_{\oplus} .

Proof of Proposition 10. Let $a \in Problematic(A)$. Assume **P6** holds for all \leq_i . That is $\forall i \in \{1..n\}$ and $\forall a \in Problematic(A)$, $\forall b \notin Problematic(A)$, $a <_{IC}^{A, \Psi} b$. If **Unanimity** holds then, it entails that $a <_{\oplus} b$ as $a \leq_i b \forall i \in \{1..n\}$. Consequently **P6** holds for \leq_{\oplus} .

Proof of Proposition 11. If \oplus does not satisfy the property of **Non-dictatorship** then $\exists i_0 \in \{1..n\}$ such that $\forall \leq_1 \dots \leq_n$ total preorders on A , $\oplus(\leq_1, \dots, \leq_n) = \leq_{i_0}$. Assume preorder \leq_{i_0} satisfies **P7** and suppose $\{a_1 \dots a_k\} \in A \perp IC$. Consequently, $\exists j, j' \in \{1 \dots k\}$ such that $j \neq j'$ and $a_j <_{i_0} a_{j'}$. Hence, as $\oplus(\leq_1, \dots, \leq_n) = \leq_{i_0}$, it also holds that $a_j <_{\oplus} a_{j'}$. Consequently **P7** holds for \leq_{\oplus} .

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