



Sparse multivariate factorization by mean of a few bivariate factorizations

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Sparse multivariate factorization by mean of a few bivariate factorizations.

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Abstract

We describe an algorithm to factor sparse multivariate polynomials using $O(d)$ bivariate factorizations where d is the number of variables. This algorithm is implemented in the Giac/Xcas computer algebra system.

1 Introduction

To my knowledge, there are three classes of algorithms to factor multivariate polynomials over \mathbb{Z} :

- reduction to bivariate factorization and Hensel lifting (Von zur Gathen and Kaltofen [3], Bernardin and Monagan [1])
- evaluation of one variable at a sufficiently large z , factorization and reconstruction by writing coefficients in basis z with symmetric remainders (heuristic factorization)
- Kronecker-like substitution (replace one of the variable by another one to a sufficiently large power).

Bivariate factorization may be obtained by partial differential equation (Gao [2]) or interpolation or one of the previous method.

We present here a method that is adapted to sparse factors, where the previous method would require too many ressources. For example Hensel lifting does not work if a leading coefficient of the factors vanishes once evaluated to 0 at other variables. The usual trick to avoid this is to translate the origin, but this will densify the polynomial to be factored.

2 The algorithm

2.1 Main idea

Let $n \geq 2$ and $P(x, x_1, \dots, x_n)$ be a sparse square-free polynomial that we want to factor, assume that the factorization is :

$$P = P_1 \dots P_k$$

The basic idea is replace all variables x_1, \dots, x_n with t, t, \dots, t and factor the substituted bivariate polynomial $P_{t, \dots, t}$, then compare with the factorization of $P_{t^2, t, \dots, t}$ (where x_1, \dots, x_n are substituted by t^2, t, t, \dots, t in P) or with $P_{t^3, t, \dots, t}$ or etc. If the factorization is sparse enough, there is a good chance that the factors will be similar (same number of monomials, same pattern in x , same value for the coefficients), and the monomial power differences in t will give us the x_1 contribution to the monomials. Doing the same for x_2, \dots, x_n will give us the reconstruction.

The details are a little more complicated, because we must take care of the content of the substituted polynomials $P_{t, t, \dots, t}$ and of the order of the monomials having the same x powers in a given factor. The next example that motivated the implementation in Giac/Xcas will demonstrate the main idea, problems and solutions.

2.2 Example

The following example was discussed on the `sage-devel` list, it was obtained with a random generation command returning 2 polynomials in 5 variables. We make the product and try to factor it back. It was not factored by Sage 7.4 (with Singular 4 inside), but was reported to be factored by magma in 3s.

```
A:=3732480000000000*a^25*b^9*c^25*d^21*e^21 +
1866240000000000*a^20*b^9*c^25*d^24*e^21 +
3732480000000000*a^20*b^4*c^25*d^22*e^21 +
1244160000000000*a^20*b^4*c^28*d^21*e^18 +
3732480000000*a^16*b^4*c^25*d^21*e^20 +
1866240000000*a^16*b^4*c^26*d^21*e^18 +
1866240000000*a^13*b^6*c^25*d^21*e^17 +
12441600000*a^13*b^5*c^25*d^21*e^12 +
3110400000*a^13*b^7*c^23*d^16*e^12 +
12441600000*a^13*b^4*c^25*d^16*e^13 +
3110400000*a^16*b^4*c^20*d^16*e^12 +
622080000*a^13*b*c^21*d^16*e^12 +
233280000*a^13*b*c^20*d^17*e^8 + 77760000*a^13*b*c^15*d^18*e^8 +
25920000*a^13*b*c^15*d^14*e^10 + 25920000*a^13*b^4*c^15*d^10*e^8 +
17280000*a^8*b*c^15*d^14*e^8 + 3240000*a^8*b^4*c^15*d^6*e^8 +
216000*a^4*b^3*c^15*d^6*e^8 + 216000*a^4*b*c^10*d^9*e^8 +
86400*a^4*b*c^10*d^8*e^7 + 32400*a^7*b*c^10*d^3*e^7 +
2700*a^4*b^4*c^10*d^3*e^3 + 675*a^6*b*c^7*d^3*e^3 +
1125*a^5*b*c^2*d^3*e^3 + 135*b^5*c^2*d^3*e^3 +
```

$$27*c^2*d^6*E^3 + 12*c^7 + 9*c^3*d^3 + a^2;$$

[illegible]

The smallest partial degree of the product is $9+19$ in b , therefore b will be our x variable, while a, c, d, E are our x_1, \dots, x_4 variables. A and B are irreducible, we set $P = AB$.

Factoring $P(x, t, t, t, t)$ gives

$$\begin{aligned} & t^5 * \\ & (37324800000000 * b^9 * t^{90} + 186624000000000 * b^9 * t^{88} \\ & + 3110400000 * b^7 * t^{62} + 186624000000 * b^6 * t^{74} + \\ & 12441600000 * b^5 * t^{69} + 135 * b^5 * t^6 + \\ & 37324800000000 * b^4 * t^{86} + 12441600000000 * b^4 * t^{85} + \\ & 373248000000 * b^4 * t^{80} + 1866240000000 * b^4 * t^{79} \\ & + 12441600000 * b^4 * t^{65} + 3110400000 * b^4 * t^{62} + 25920000 * b^4 * t^{44} \\ & + 3240000 * b^4 * t^{35} + 2700 * b^4 * t^{18} + \\ & 216000 * b^3 * t^{31} + \\ & 622080000 * b * t^{60} + 233280000 * b * t^{56} + 77760000 * b * t^{52} \\ & + 25920000 * b * t^{50} + 17280000 * b * t^{43} + 216000 * b * t^{29} \\ & + 86400 * b * t^{27} + 32400 * b * t^{25} + 675 * b * t^{17} + 1125 * b * t^{11} \\ & + 27 * t^9 + 12 * t^5 + 9 * t^4 + 1) * \dots \end{aligned}$$

Factoring $P(x, t^2, t, t, t)$ gives

$$t^7_*$$

It is clearly a similar factorization, the number of monomials differ only by 1 ($12441600000 * b^4 * t^{65} + 3110400000 * b^4 * t^{62}$ is grouped in one monomial in the second factorization), and the order is not the same for the coefficient of b . Note that there is also a content term in t . In fact, we just got an unlucky evaluation at $x_1 = a = t^2$, $x_1 = a = t^3$ is also unlucky, while $x_1 = a = t^4$ returns 30 monomials like t .

In order to compare the two factorizations, we must solve these 3 problems: content, number of monomials, and monomials ordering.

We assume that the factors of the bivariate factorization are x -distincts, that is the distribution of non-zero coefficients in powers of x are not the same. This way, we can isolate the same factor in two bivariate factorizations, and we will now reconstruct the true factor.

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$5t^3)/(t^{90} + 5t^{88}) = t^{63}$. For $P_{t^4, \dots, t}$, we multiply the factor by $t^{4 \times 22 + 31 + 54 + 41}(t^{20} + 5t^3)/(t^{161} + 5t^{144}) = t^{73}$.

Solving the number of monomials is done like for any modular reconstruction: if an evaluation with x_k replaced by t^j contains less term than the previous one, we ignore it (unlucky evaluation), if an evaluation contains more terms than a previous one, we throw what we had before and restart from this one. If the number of monomials is the same, we also check that the non-zero partials degrees in x are the same.

Keeping the right order of monomials is more original: it can be done by comparing first the x power, then comparing the coefficient of the monomial (it is impossible to insure the same ordering by sorting with powers of t). In order to do that we must insure that in the factor to be reconstructed the coefficients of the monomials of the same power of x are all distincts. If this is not the case, we can dilate some variables by a constant factor and retry (in our implementation we dilate all variables except x randomly by ± 1 or ± 2).

If we have two matching factors for evaluations at $t..t, t^j, t..t$ and $t..t, t^{j'}, t..t$ then the power of x_k in a monomial is the difference of powers in t of the same monomial in the two factorizations divided by $j - j'$. For example the first monomial in $P_{t, \dots, t}$ multiplied by t^{63} is $37324800000000 * b^9 * t^{90+63}$, the corresponding monomial in $P_{t^4, \dots, t}$ multiplied by t^{73} is $37324800000000 * b^9 * t^{161+73}$, that's a power contribution for $x_1 = a$ of $(234 - 153)/3 = 27$. Indeed the leading coefficient of A is a^{25} , multiplied by a^2 inside the leading coefficient of Q in b is a^{27} .

2.4 Implementation

This algorithm is implemented in C++ in the file `ezgcd.cc` of the source code of Giac/Xcas, in the function `try_sparse_factor_bi`. It factors the polynomial in the example in less than 2s (without this function, the factorization was impracticable).

We hope it will help other open-source softwares implement more efficient sparse multivariate factorization algorithms!

References

- [1] L. Bernardin and M. B. Monagan. Efficient multivariate factorization over finite fields. In *International Symposium on Applied Algebra, Algebraic Algorithms, and Error-Correcting Codes*, pages 15–28. Springer, 1997.
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- [3] J. von zur Gathen and E. Kaltofen. Factoring sparse multivariate polynomials. *Journal of Computer and System Sciences*, 31(2):265–287, 1985.