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# THE ROTATING NORMAL FORM IS REGULAR

JEAN FROMENTIN

ABSTRACT. Defined on Birman–Ko–Lee monoids, the rotating normal form has strong connections with the Dehornoy’s braid ordering. It can be seen as a process for selecting between all the representative words of a Birman–Ko–Lee braid a particular one, called *rotating* word. In this paper we construct, for all  $n \geq 2$ , a finite state automaton which recognizes the rotating words on  $n$  strands. As a consequence the language of rotating words on  $n$  strands is proved to be regular for any  $n \geq 2$ .

## 1. INTRODUCTION

Originally, the group  $B_n$  of  $n$ -strand braids was defined as the group of isotopy classes of  $n$ -strand geometric braids. An algebraic presentation of  $B_n$  was given by E. Artin in [1]

$$\left\langle \sigma_1, \dots, \sigma_{n-1} \mid \begin{array}{l} \sigma_i \sigma_j = \sigma_j \sigma_i \quad \text{for } |i - j| \geq 2 \\ \sigma_i \sigma_j \sigma_i = \sigma_j \sigma_i \sigma_j \quad \text{for } |i - j| = 1 \end{array} \right\rangle. \quad (1)$$

An  $n$ -strand braid is an equivalence class consisting of (infinitely many) words in the letters  $\sigma_i^{\pm 1}$ . The standard correspondence between elements of the presented group  $B_n$  and geometric braids consists in using  $\sigma_i$  as a code for the geometric braid where only the  $i$ th and the  $(i + 1)$ st strands cross, with the strands originally at position  $(i + 1)$  in front of the other.

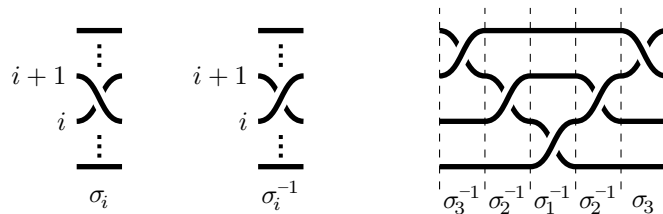


FIGURE 1. Interpretation of a word in the letters  $\sigma_i^{\pm 1}$  as a geometric braid diagram.

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In 1998, J.S. Birman, K.H. Ko, and S.J. Lee [3] introduced and investigated for each  $n$  a submonoid  $B_n^{+*}$  of  $B_n$ , which is known as the *Birman–Ko–Lee monoid*. The name “*dual braid monoid*” was subsequently proposed because several numerical parameters obtain symmetric values when they are evaluated on the positive braid monoid  $B_n^+$  and on  $B_n^{+*}$ , a correspondence that was extended to the more general context of Artin–Tits groups by D. Bessis [2] in 2003. The dual braid monoid  $B_n^{+*}$  is the submonoid of  $B_n$  generated by the braids  $a_{i,j}$  with  $1 \leq i < j \leq n$ , where  $a_{i,j}$  is defined by  $a_{i,j} = \sigma_i \dots \sigma_{j-1} \sigma_j \sigma_{j-1}^{-1} \dots \sigma_i^{-1}$ . In geometrical terms, the braid  $a_{i,j}$  corresponds to a crossing of the  $i$ th and  $j$ th strands, both passing behind the (possible) intermediate strands.

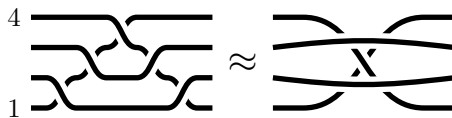


FIGURE 2. In the geometric braid  $a_{1,4}$ , the strands 1 and 4 cross under the strands 2 and 3.

By definition,  $\sigma_i$  equals  $a_{i,i+1}$  and, therefore, the positive braid monoid  $B_n^+$  is included in the monoid  $B_n^{+*}$ , a proper inclusion for  $n \geq 3$  since the braid  $a_{1,3}$  does not belong to the monoid  $B_3^+$ .

We denote by  $A_n$  the set  $\{a_{p,q} \mid 1 \leq p < q \leq n\}$ . The following presentation of the monoid  $B_n^{+*}$  is given in [3].

**Proposition 1.1.** *The monoid  $B_n^{+*}$  is presented by generators  $A_n$  and relations*

$$a_{p,q}a_{r,s} = a_{r,s}a_{p,q} \quad \text{for } [p,q] \text{ and } [r,s] \text{ disjoint or nested,} \quad (2)$$

$$a_{p,q}a_{q,r} = a_{q,r}a_{p,r} = a_{p,r}a_{p,q} \quad \text{for } 1 \leq p < q < r \leq n. \quad (3)$$

The interval  $[p,q]$  is said to be *nested* in  $[r,s]$  if the relation  $r < p < q < s$  holds.

Since [2] and [3] it is known that the dual braid monoid  $B_n^{+*}$  admits a Garside structure whose simple elements are in bijection with the non-crossing partitions of  $n$ . In particular, there exists a normal form associated with this Garside structure, the so-called greedy normal form.

The rotating normal form is another normal form on  $B_n^{+*}$ , it was introduced in [8, 9]. Roughly speaking, for every braid  $\beta \in B_n^{+*}$  the rotating normal form picks up a unique representative word on the letters  $A_n$  among all of these representing  $\beta$ . It can be seen as a map  $r_n$  from the dual braid monoid  $B_n^{+*}$  to the set of words  $A_n^*$ . The language of all  $n$ -rotating words, denoted by  $R_n$  is then the image of  $B_n^{+*}$  under the map  $r_n$ .

The aim of this paper is to construct for all  $n \geq 2$  an explicit finite state automaton which recognizes the language  $R_n$ . As a consequence we obtain that the language  $R_n$  of  $n$ -rotating words is regular.

The paper is divided as follow. In section 2 we recall briefly the construction of the rotating normal form and its useful already known properties. In third section we describe the left reversing process on dual braid monoids. In section 4 we give a syntactical characterization of  $n$ -rotating normal words. In fifth section we construct, for each  $n \geq 2$ , a finite state automaton which recognizes the language  $R_n$  of  $n$ -rotating normal words.

## 2. THE ROTATING NORMAL FORM

The main ingredient to define the rotating normal form is the Garside automorphism  $\phi_n$  of  $B_n^{+*}$  defined by  $\phi_n(\beta) = \delta_n \beta \delta_n^{-1}$  where  $\delta_n = a_{1,2} a_{2,3} \dots a_{n-1,n}$  is the Garside braid of  $B_n^{+*}$ . In terms of Birman–Ko–Lee generators, the map  $\phi_n$  can be defined by

$$\phi_n(a_{p,q}) = \begin{cases} a_{p+1,q+1} & \text{for } q \leq n-1, \\ a_{1,p+1} & \text{for } q = n. \end{cases} \quad (4)$$

Geometrically,  $\phi_n$  should be viewed as a rotation, which makes sense provided braid diagrams are drawn on a cylinder rather than on a plane rectangle.

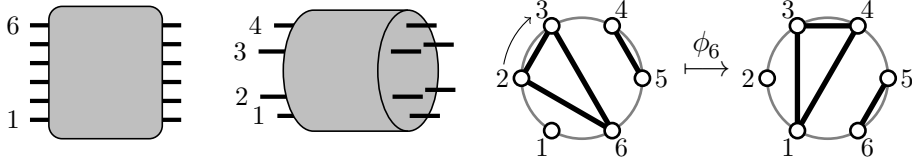


FIGURE 3. Rolling up the usual braid diagram helps us to visualize the symmetries of the braids  $a_{p,q}$ . On the resulting cylinder,  $a_{p,q}$  naturally corresponds to the chord connecting vertices  $p$  and  $q$ . With this representation,  $\phi_n$  acts as a clockwise rotation of the marked circles by  $2\pi/n$ .

For  $\beta$  and  $\gamma$  in  $B_n^{+*}$ , we say that  $\gamma$  is a *right-divisor* of  $\beta$ , if there exists a dual braid  $\beta'$  of  $B_n^{+*}$  satisfying  $\beta = \beta' \gamma$ .

**Definition 2.1.** For  $n \geq 3$  and  $\beta$  a braid of  $B_n^{+*}$ , the maximal braid  $\beta_1$  lying in  $B_{n-1}^{+*}$  that right divides the braid  $\beta$  is called the  $B_{n-1}^{+*}$ -tail of  $\beta$ .

Using basic Garside properties of the monoid  $B_n^{+*}$  we obtain the following result (Proposition 2.5 of [9]) which allow us to express each braid of  $B_n^{+*}$  as a unique finite sequence of braids lying in  $B_{n-1}^{+*}$ .

**Proposition 2.2.** Assume  $n \geq 3$ . For each nontrivial braid  $\beta$  of  $B_n^{+*}$  there exists a unique sequence  $(\beta_b, \dots, \beta_1)$  of braids of  $B_{n-1}^{+*}$  satisfying  $\beta_b \neq 1$  and

$$\beta = \phi_n^{b-1}(\beta_b) \cdot \dots \cdot \phi_n(\beta_2) \cdot \beta_1, \quad (5)$$

for each  $k \geq 1$ , the  $B_{n-1}^{+*}$ -tail of  $\phi_n^{b-k}(\beta_b) \cdot \dots \cdot \phi_n(\beta_{k+1})$  is trivial. (6)

Under the above hypotheses, the sequence  $(\beta_b, \dots, \beta_1)$  is called the  $\phi_n$ -splitting of the braid  $\beta$ . It is shown in [9] that Condition (6) can be replaced by

$$\text{for each } k \leq 1, \beta_k \text{ is the } B_{n-1}^{+*}\text{-tail of } \phi_n^{b-k}(\beta_b) \cdot \dots \cdot \phi_n(\beta_{k-1}) \cdot \beta_k. \quad (7)$$

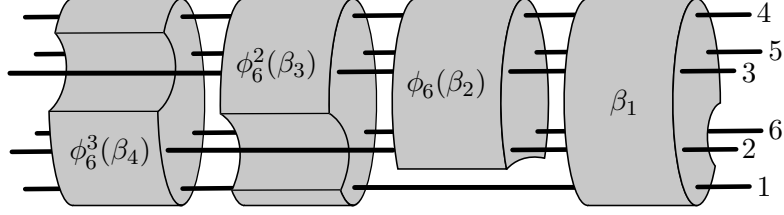


FIGURE 4. The  $\phi_6$ -splitting of a braid of  $B_6^{+*}$ . Starting from the right, we extract the maximal right-divisor that keeps the sixth strand unbraided, then extract the maximal right-divisor that keeps the first strand unbraided, etc.

**Example 2.3.** Consider the braid  $\beta = a_{1,2}a_{2,3}a_{1,2}a_{2,3}$  of  $B_3^{+*}$ . Using relations (3) on the underlined factors we obtain

$$\beta = a_{1,2}a_{2,3}\underline{a_{1,2}a_{2,3}} = a_{1,2}\underline{a_{2,3}a_{1,3}}a_{1,2} = a_{1,2}a_{1,3}a_{1,2}a_{1,2}$$

We decompose  $\beta$  as  $\phi_3(\gamma_1) \cdot \beta_1$  with  $\gamma_1 = \phi_3^{-1}(a_{1,2}a_{1,3}) = a_{1,3}a_{2,3}$  and  $\beta_1 = a_{1,2}a_{1,2}$ . The braid  $\phi_3(\gamma_1) = a_{1,2}a_{2,3}$  is exactly the one of (6) for  $n = 3$  and  $k = 1$ . As the word  $a_{1,3}a_{2,3}$  is alone in its equivalence class the braid  $\phi_3(\gamma_1)$  is not right divisible by  $a_{1,2}$  and so its  $B_2^{+*}$ -tail is trivial. Considering  $\gamma_1$  instead of  $\beta$  we obtain  $\gamma_1 = \phi_3(\gamma_2) \cdot \beta_2$  with  $\gamma_2 = \phi_3^{-1}(a_{1,3}a_{2,3}) = a_{2,3}a_{1,2}$  and  $\beta_2 = 1$ . The braid  $\phi_3(\gamma_2) = a_{1,3}a_{2,3}$  is the braid of (6) for  $n = 3$  and  $k = 2$  and it is always alone in its equivalence class, implying that its  $B_2^{+*}$ -tail is trivial. We express  $\gamma_2$  as  $\phi_3(\gamma_3) \cdot \beta_3$  with  $\gamma_3 = \phi_3^{-1}(a_{2,3}) = a_{1,2}$  and  $\beta_3 = a_{1,2}$ . Since  $\gamma_3$  equals  $a_{1,2}$  we obtain  $\gamma_4 = 1$  and  $\beta_4 = a_{1,2}$ . We conclude that the  $\phi_3$ -splitting of  $\beta$  is  $(a_{1,2}, a_{1,2}, 1, a_{1,2}^2)$ .

Before giving the definition of the rotating normal we fix some definitions about words.

**Definition 2.4.** A word on the alphabet  $A_n$  is an  $A_n$ -word. A word on the alphabet  $A_n^\pm = A_n \sqcup A_n^{-1}$  is an  $A_n^\pm$ -word. The braid represented by the  $A_n^\pm$ -word  $w$  is denoted by  $\overline{w}$ . For  $w, w'$  two  $A_n^\pm$ -word, we say that  $w$  is equivalent to  $w'$ , denoted by  $w \equiv w'$  if  $\overline{w} = \overline{w'}$  holds. The empty word is denoted by  $\varepsilon$ .

The  $n$ -rotating normal form is an injective map  $r_n$  from  $B_n^{+*}$  to the set of  $A_n$ -words defined inductively using the  $\phi_n$ -splitting.

**Definition 2.5.** For  $\beta \in B_2^{+*}$ , we define  $r_2(\beta)$  to be the unique word  $a_{1,2}^k$  representing  $\beta$ . The rotating normal form of a braid  $\beta \in B_n^{+*}$  with  $n \geq 3$  is

$$r_n(\beta) = \phi_n^{b-1}(r_{n-1}(\beta_b)) \cdot \dots \cdot \phi_n(r_{n-1}(\beta_2)) \cdot r_{n-1}(\beta_1),$$

where  $(\beta_b, \dots, \beta_1)$  is the  $\phi_n$ -splitting of  $\beta$ . A word  $w$  is said to be  $n$ -rotating if it is the  $n$ -rotating normal form of a braid of  $B_n^{+*}$ .

As the  $n$ -rotating normal form of a braid of  $B_{n-1}^{+*}$  is equal to its  $(n-1)$ -rotating normal form we can talk without ambiguities of the *rotating normal form* about a dual braid.

**Example 2.6.** We reconsider the braid  $\beta$  of Example 2.3. We know that the  $\phi_3$ -splitting of  $\beta$  is  $(a_{1,2}, a_{1,2}, 1, a_{1,2}^2)$ . Since  $r_2(1) = \varepsilon$ ,  $r_2(a_{1,2}) = a_{1,2}$  and  $r_2(a_{1,2}^2) = a_{1,2}^2$  we obtain

$$r_3(\beta) = \phi_3^3(a_{1,2}) \cdot \phi_3^2(a_{1,2}) \cdot \phi_3(\varepsilon) \cdot a_{1,2}^2 = a_{1,2}a_{1,3}a_{1,2}a_{1,2}.$$

Some properties of the rotating normal form have been established in [9]. Connections, established in [8] and [9], between the rotating normal form and the braid's ordering introduced by P. Dehornoy are based on these properties.

We finish this section with some already known or immediate properties about  $\phi_n$ -splittings and  $n$ -rotating words.

**Definition 2.7.** For every nonempty word  $w$ , the last letter of  $w$  is denoted by  $w^\#$ . For each nontrivial braid  $\beta$  in  $B_n^{+*}$ , we define the *last letter* of  $\beta$ , denoted  $\beta^\#$ , to be the last letter in the rotating normal form of  $\beta$ .

**Lemma 2.8** (Lemma 3.2 of [9]). *Assume  $n \geq 3$  and let  $(\beta_b, \dots, \beta_1)$  be a  $\phi_n$ -splitting*

- (i) *For  $k \geq 2$ , the letter  $\beta_k^\#$  is of type  $a_{\dots, n-1}$  unless  $\beta_k = 1$ .*
- (ii) *For  $k \geq 3$  and  $k = b$ , we have  $\beta_k \neq 1$ .*

The fact that  $\beta_b$  is not trivial is a direct consequence of the definition of  $\phi_n$ -splitting. As, for  $k \geq 2$ , the braid  $\beta' = \phi_n(\beta_{k+1}^\#)\beta_k$  is a right-divisor of  $\phi_n^{b-k}(\beta_b) \cdot \dots \cdot \beta_k$ , it must satisfy some properties. In particular, if  $\beta_{k+1}^\# = a_{p-1, n-1}$  holds then the  $B_{n-1}^{+*}$ -tail of  $\phi_n(a_{p,n}\beta_k)$  is trivial by (6).

**Definition 2.9.** We say that a letter  $a_{r,s}$  is an  $a_{p,n}$ -barrier if  $1 \leq r < p < s \leq n-1$  holds.

There exist no  $a_{p,n}$ -barrier with  $n \leq 3$  and the only  $a_{p,4}$ -barrier is  $a_{1,3}$ , which is an  $a_{2,4}$ -barrier. By definition, if the letter  $x$  is an  $a_{p,n}$ -barrier, then in the presentation of  $B_n^{+*}$  there exists no relation of the form  $a_{p,n} \cdot x = y \cdot a_{p,n}$  allowing one to push the letter  $a_{p,n}$  to the right through the letter  $x$ : so, in some sense,  $x$  acts as a barrier.

**Lemma 2.10** (Lemma 3.4 of [9]). *Assume that  $n \geq 3$ ,  $\beta$  is a braid of  $B_{n-1}^{+*}$  and the  $B_{n-1}^{+*}$ -tail of  $\phi_n(a_{p,n}\beta)$  is trivial for  $2 \leq p \leq n-2$ . Then the rotating normal form of  $\beta$  is not the empty word and it contains an  $a_{p,n}$ -barrier.*

**Lemma 2.11** (Lemma 3.5 of [9]). *Let  $(\beta_b, \dots, \beta_1)$  be a  $\phi_n$ -splitting of some braid of  $B_n^{+*}$  with  $n \geq 3$ . Then for each  $k \in [2, b-1]$  such that  $\beta_{k+1}^\#$  is not  $a_{n-2, n-1}$  (if any), the rotating normal form of  $\beta_k$  contains an  $\phi_n(\beta_{k+1}^\#)$ -barrier.*

### 3. LEFT REVERSING FOR DUAL BRAID MONOID

Left reversing process was introduced by P. Dehornoy in [5]. It is a powerful tool for the investigation of division properties in some monoids as stated by Proposition 3.6.

**Definition 3.1.** A monoid  $M$  defined by a presentation  $\langle S | R \rangle^+$  is *left complemented* if there exists a map  $f : S \times S \rightarrow S^*$  satisfying

$$R = \{f(x, y)x = f(y, x)y \mid (x, y) \in S^2\}$$

and if  $f(x, x) = \varepsilon$  holds for all  $x \in S$ .

As the relation  $x = x$  is always true for  $x \in S$  we say that  $M$  is left complemented even if  $x = x$  does not occur in  $R$  for  $x \in S$ .

The monoid  $B_3^{+*}$  with presentation of Proposition 1.1 is left complemented with respect to the map  $f$  given by

$$\begin{aligned} f(a_{1,2}, a_{2,3}) &= f(a_{1,2}, a_{1,3}) = a_{1,3} \\ f(a_{2,3}, a_{1,2}) &= f(a_{2,3}, a_{1,3}) = a_{1,2} \\ f(a_{1,3}, a_{1,2}) &= f(a_{1,3}, a_{2,3}) = a_{2,3} \end{aligned}$$

However the monoid  $B_4^{+*}$  with presentation of Proposition 1.1 is not left complemented. Indeed there is no relation of the form  $\dots a_{1,3} = \dots a_{2,4}$ . Hence the words  $f(a_{1,3}, a_{2,4})$  and  $f(a_{2,4}, a_{1,3})$  are not well defined.

In general for  $1 \leq p < r < q < s \leq n$ , the word  $f(a_{p,q}, a_{r,s})$  and  $f(a_{r,s}, a_{p,q})$  are not defined for the presentation of  $B_n^{+*}$  given in Proposition 1.1. In order to obtain a left complemented presentation of  $B_n^{+*}$  we must exhibit some extra relations from these given in Proposition 1.1.

By example, the relation  $a_{2,3}a_{1,4}a_{1,3} \equiv a_{3,4}a_{1,2}a_{2,4}$  holds and so we can consider  $f(a_{1,3}, a_{2,4})$  to be  $a_{2,3}a_{1,4}$ . However the relation  $a_{1,4}a_{2,3}a_{1,3} \equiv a_{3,4}a_{1,2}a_{2,4}$  is also satisfied and so  $f(a_{1,3}, a_{2,4}) = a_{2,3}a_{1,4}$  is an other valid choice.

**Lemma 3.2.** *For  $n \geq 2$ , the map  $f_n : A_n \times A_n \rightarrow A_n^*$  defined by*

$$f_n(a_{p,q}, a_{r,s}) = \begin{cases} \varepsilon & \text{for } a_{p,q} = a_{r,s}, \\ a_{p,s} & \text{for } q = r, \\ a_{s,q} & \text{for } p = r \text{ and } q > s, \\ a_{r,p} & \text{for } q = s \text{ and } p > r, \\ a_{r,q}a_{p,s} & \text{for } p < r < q < s, \\ a_{s,q}a_{r,p} & \text{for } r < p < s < q, \\ a_{r,s} & \text{otherwise.} \end{cases}$$

provides a structure of left complemented monoid to  $B_n^{+*}$ .

*Proof.* Direct computations using Proposition 1.1 establish  $f_n(x, y) \cdot x \equiv f_n(y, x) \cdot y$  for all  $(x, y) \in A_n^2$ .  $\square$

Our choice for  $f_n(a_{p,q}, a_{r,s})$  with  $p < r < q < s$  is well suited for the sequel and some proof would be invalid if we made an other one.

**Definition 3.3.** For  $w$  and  $w'$  two  $A_n^\pm$ -words, we say that  $w$  *left reverses in one step* to  $w'$ , denoted  $w \curvearrowright^1 w'$ , if we can obtain  $w'$  from  $w$  substituting a factor  $xy^{-1}$  (with  $x, y \in A_n$ ) by  $f_n(x, y)^{-1}f_n(y, x)$ . We say that  $w$  *left reverses* to  $w'$ , denoted by  $w \curvearrowright w'$ , if there exists a sequence  $w = w_1, \dots, w_\ell = w'$  of  $A_n^\pm$ -words such that  $w_k \curvearrowright^1 w_{k+1}$  for  $k \in [1, \ell - 1]$ .

**Example 3.4.** The word  $u = a_{1,2}a_{2,3}a_{1,2}a_{1,3}^{-1}$  left reverses to  $a_{2,3}a_{2,3}$  as the following left reversing sequence shows (left reversed factor are underlined)

$$a_{1,2}a_{2,3}\underline{a_{1,2}a_{1,3}^{-1}} \curvearrowright^1 a_{1,2}\underline{a_{2,3}a_{1,3}^{-1}}a_{2,3} \curvearrowright^1 \underline{a_{1,2}a_{1,2}^{-1}}a_{2,3}a_{2,3} \curvearrowright^1 a_{2,3}a_{2,3},$$

which is denoted by  $a_{1,2}a_{2,3}a_{1,2}a_{1,3}^{-1} \curvearrowright a_{2,3}a_{1,2}$ .

**Definition 3.5.** For  $w$  an  $A_n^\pm$ -word, we denote by  $D(w)$  and  $N(w)$  the unique  $A_n$ -word, if there exist, such that  $w \curvearrowright D(w)^{-1}N(w)$ . The word  $N(w)$  is the *left numerator* of  $w$  while the word  $D(w)$  is its *left denominator*.

Reconsidering Example 3.4, we obtain that the left denominator of  $u$  is  $D(u) = \varepsilon$  and that is left numerator its  $N(u) = a_{2,3}a_{2,3}$ .

A consequence of Example 8 and Proposition 3.5 of [6] based on [5] and [3] is that  $N(w)$  and  $D(w)$  exists for any  $A_n^\pm$ -word  $w$ . We obtain also the following result:

**Proposition 3.6.** *Let  $w$  be an  $A_n$ -word and  $a_{p,q}$  be in  $A_n$ . The braid  $\bar{w}$  is right divisible by  $a_{p,q}$  if and only if  $D(w \cdot a_{p,q}^{-1})$  is empty.*

Since the denominator of  $a_{1,2}a_{2,3}a_{1,2}a_{1,3}^{-1}$  is empty, the braid  $a_{1,3}$  right divides the braid  $a_{1,2}a_{2,3}a_{1,2}$ .

#### 4. CHARACTERIZATION OF ROTATING NORMAL WORDS

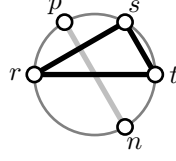
The aim of this section is to give a syntactical characterization of  $n$ -rotating words among  $A_n$ -words.

**Lemma 4.1.** *For  $\beta$  in  $B_n^{+*}$  there is equivalence between*

- (i) *an  $A_n$ -word representing  $\beta$  contains an  $a_{p,n}$ -barrier,*
- (ii) *every  $A_n$ -word representing  $\beta$  contains an  $a_{p,n}$ -barrier.*

*Proof.* Clearly (ii) implies (i). For (i)  $\Rightarrow$  (ii) it is sufficient to prove that relations of Proposition 1.1 preserve  $a_{p,n}$ -barriers. For relation (2) this is immediate since it conserves involved letters. If one of the three words  $a_{r,s}a_{s,t}$ ,  $a_{s,t}a_{r,s}$  and  $a_{r,t}a_{r,s}$  of relations (3) contains an  $a_{p,n}$ -barrier then the two others also, as illustrated by the following chord diagram (the chord of an  $a_{p,n}$ -barrier intersects properly the chord of the letter  $a_{p,n}$ .)





□

As a consequence, containing an  $a_{p,n}$ -barrier is not a word property but a braid one.

**Lemma 4.2.** *Assume that  $n \geq 3$ ,  $\beta$  belongs to  $B_{n-1}^{+*}$  and that the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta)$  is trivial. Then every  $A_{n-1}$ -word representing  $\beta$  ends with  $\beta^\#$ .*

*Proof.* Let  $u$  be an  $A_{n-1}$ -word representing  $\beta$ . As the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta)$  is trivial, the last letter  $u^\#$  of  $u$  not belongs to  $A_{n-2}$  and so  $u^\#$  is  $a_{p,n-1}$  for some integer  $p < n - 1$ . Assume now  $v$  is another  $A_{n-1}$ -word representing  $\beta$ . For the same reason as  $u$ , we have  $v^\# = a_{q,n-1}$  for some  $q < n - 1$ . Since the two braids  $a_{p,n-1}$  and  $a_{q,n-1}$  are right divisors of  $\beta$ , their left lcm is also a right divisor of  $\beta$ . Assume for a contradiction that  $p$  and  $q$  are different. The braid  $\beta$  is then right divisible by  $a_{p,q}a_{q,n-1}$ , which is the left lcm of  $a_{p,n-1}$  and  $a_{q,n-1}$ . Since  $a_{p,q}a_{q,n-1}$  is equivalent to  $a_{p,n-1}a_{p,q}$ , the braid  $a_{p,q}$  is also a right divisor of  $\beta$ . In particular  $a_{p+1,q+1}$ , with  $q + 1 < n$ , is a right divisor of  $\phi_n(\beta)$ , which is impossible since the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta)$  is supposed to be trivial. Therefore, every  $A_{n-1}$ -word representing  $\beta$  ends with the same letter, namely  $\beta^\#$ . □

We conclude that, under some hypotheses, the last letter of a word is a braid invariant.

**Definition 4.3.** For  $n \geq 3$  and  $2 \leq p \leq n - 1$ , we say that an  $n$ -rotating word  $w$  is an  $a_{p,n}$ -ladder if there exist a decomposition

$$w = v_0 x_1 v_1 \dots v_{h-1} x_h v_h,$$

a sequence  $p = j(0) < j(1) < \dots < j(h) = n - 1$  and a sequence  $i$  such that

- (i) for each  $k \leq h$ , the letter  $x_k$  is  $a_{i(k),j(k)}$  with  $i(k) < j(k-1) < j(k)$ ,
- (ii) for each  $k < h$ , the word  $v_k$  contains no  $a_{j(k),n}$ -barrier,

Condition (ii) is equivalent to: for each  $k \leq h$ , the letter  $x_k$  is an  $a_{j(k-1),n}$ -barrier of type  $a_{\dots,j(k)}$ .

An immediate adaptation of Proposition 3.9 of [9] is :

**Lemma 4.4.** *Assume that  $n \geq 3$ ,  $\beta$  belongs to  $B_{n-1}^{+*}$ , the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta)$  is trivial and  $\beta$  contains an  $a_{p,n}$ -barrier for some  $2 \leq p \leq n - 2$ . Then the normal form of  $\beta$  is an  $a_{p,n}$ -ladder.*

In order to obtain a syntactical characterization of  $n$ -rotating words we want a local version of condition (6) characterizing a  $\phi_n$ -splitting. The following result is the first one in this way.

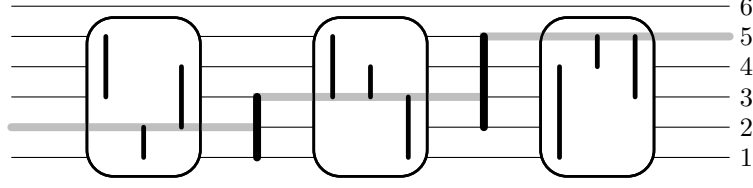


FIGURE 5. An  $a_{2,6}$ -ladder. The gray line starts at position 1 and goes up to position 5 using the bar of the ladder. The empty spaces between bars in the ladder are represented by a framed box. In such boxes the vertical line representing the letter  $a_{i,j}$  does not cross the gray line. The bar of the ladder are represented by black thick vertical lines.

**Proposition 4.5.** For  $\beta \in B_{n-1}^{+*}$  and  $p$  an integer satisfying  $2 \leq p \leq n-2$  there is equivalence between

- (i) the  $B_{n-1}^{+*}$ -tail of  $\phi_n(a_{p,n}\beta)$  is trivial,
- (ii) the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta)$  is trivial and  $\beta$  contains an  $a_{p,n}$ -barrier,
- (iii) the only  $A_n$ -letter that right divides  $a_{p,n}\beta$  is  $\beta^\#$ , which is of type  $a_{\dots,n-1}$ .

Our proof of Proposition 4.5 rests on the following Lemma.

**Lemma 4.6.** For  $n \geq 3$ ,  $u$  an  $A_{n-1}$ -word and  $p \in [1, n-1]$ , the left denominator  $D(ua_{p,n}^{-1})$  is not empty. More precisely,  $D(ua_{p,n}^{-1})^{-1}$  begins with  $a_{q,n}^{-1}$  satisfying  $q \leq p$ .

*Proof.* Assume that  $w_1, \dots, w_\ell$  is a reversing sequence from the word  $w_1 = ua_{p,n}^{-1}$  to the word  $D(w_1)^{-1}N(w_1)$ . For  $k \in [1, \ell]$  we denote by  $y_k$  the leftmost negative letter in  $w_k$ . Each reversing step consists in replacing a factor  $xy^{-1}$  of  $w_k$  by  $f_n(x, y)^{-1}f_n(y, x)$ . If for  $k \in [1, \ell]$  the reversed factor of  $w_k$  does not contains  $y_k^{-1}$  then  $y_{k+1}$  equals  $y_k$ . Assume now that the reversed factor is  $xy_k^{-1}$  with  $y_k = a_{r,n}$ . Lemma 3.2 implies

$$f_n(x, y_k) = f_n(a_{i,j}, a_{r,n}) = \begin{cases} a_{i,n} & \text{for } j = r, \\ a_{r,j}a_{i,n} & \text{for } i < r < j, \\ a_{r,n} & \text{otherwise,} \end{cases}$$

which gives in particular

$$xy_k^{-1} = a_{i,j}a_{r,n}^{-1} \curvearrowright \begin{cases} a_{i,n}^{-1}\dots & \text{for } i < r \leq j, \\ a_{r,n}^{-1}\dots & \text{otherwise.} \end{cases} \quad (8)$$

It follows that  $y_{k+1}$  is equal to  $a_{s,n}$  for some  $s \leq r$ . Eventually we obtain that  $ua_{p,n}^{-1}$  left reverses to  $a_{q,n}^{-1}\dots$  with the relation  $q \leq p$  and so the desired property on  $D(ua_{p,n})$  holds.  $\square$

*Proof of Proposition 4.5.* Assume (i). As the the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta)$  is also a right divisor of  $\phi_n(a_{p,n}\beta)$  the first statement of (ii) holds. The second statement is Lemma 2.10. Let us prove that (iii) implies (i). By hypothesis

the last letter of  $\beta$  is  $a_{q-1,n-1}$  for some  $q$ . As the only  $A_n$ -letter that right divides  $\phi_n(a_{p,n}\beta)$  is  $\phi_n(a_{q-1,n-1}) = a_{q,n}$ , the  $B_{n-1}^{+*}$ -tail of  $\phi_n(a_{p,n}\beta)$  must be trivial.

We now prove (ii)  $\Rightarrow$  (iii). Since the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta)$  is trivial, the letter  $\beta^\#$  must be of type  $a_{\dots,n-1}$ . We denote by  $w$  the rotating normal from of  $\beta$ . Let  $a_{r,s}$  be an  $A_n$ -letter different from  $\beta^\#$ . We will show that  $a_{r,s}$  cannot be a right divisor of  $a_{p,n}\beta$ . Assume first  $s \leq n-1$ . By Lemma 4.2,  $a_{r,s}$  is not a right divisor of  $\beta$ . Proposition 3.6 implies that the word  $D(w a_{r,s}^{-1})$  must be non empty. As the reversing of an  $A_{n-1}^\pm$ -word is also an  $A_{n-1}^\pm$ -word, there exists a letter  $a_{t,t'}$  with  $t' < n$  such that

$$a_{p,n} w a_{r,s}^{-1} \curvearrowright a_{p,n} a_{t,t'}^{-1} \dots,$$

holds. Clearly, the braid  $a_{t,t'}$  is not a right divisor of  $a_{p,n}$  (since we have  $t' < n$ ). Therefore, by Proposition 3.6, the left denominator of  $a_{p,n} w a_{r,s}^{-1}$  is not empty, and we conclude that  $a_{r,s}$  is not a right divisor of  $a_{p,n}\beta$ .

Assume now  $s = n$ . Hypotheses on  $\beta$  plus Lemma 4.4 imply that  $w$  is an  $a_{p,n}$ -ladder. Following Definition 4.3, we write

$$w = v_0 x_1 v_1 \dots v_{h-1} x_h v_h.$$

By Lemma 4.6, there exist two maps  $\eta$  and  $\mu$  from  $\mathbb{N}$  to itself such that

$$w a_{r,n}^{-1} = w_h a_{\eta(h),n}^{-1} \curvearrowright w'_h a_{\mu(h),n}^{-1} \dots \curvearrowright \dots \curvearrowright w_0 a_{\eta(0),n}^{-1} \dots \curvearrowright w'_0 a_{\mu(0),n}^{-1},$$

where for all  $k \in [0, h]$ ,

$$\begin{aligned} w_k &= v_0 x_1 v_1 \dots v_{k-1} x_k v_k, \\ w'_k &= v_0 x_1 v_1 \dots v_{k-1} x_k. \end{aligned}$$

By construction  $w_0$  is  $v_0$  while  $w'_0$  is the empty word. Lemma 4.6 implies

$$\mu(0) \leq \eta(0) \leq \mu(1) \leq \dots \leq \mu(h) \leq \eta(h) = r. \quad (9)$$

Following Definition 4.3 we write  $x_k = a_{i(k),j(k)}$ . We will now prove by induction

$$\text{for all } k \in [0, h-1], \mu(k+1) \leq j(k+1) \Rightarrow \eta(k) < j(k) \quad (10)$$

Let  $k \in [0, h-1]$  and assume  $\mu(k+1) \leq j(k+1)$ . Definition 4.3 (i) guarantees the relation  $i(k+1) < j(k) < j(k+1)$ . For  $\mu(k+1) \leq i(k+1)$  we have

$$\eta(k) \leq \mu(k+1) \leq i(k+1) < j(k),$$

and we are done in this case. The remaining case is  $\mu(k+1) > i(k+1)$ . By relation (8), with  $i = i(k+1)$ ,  $j = j(k+1)$  and  $r = \eta(k+1)$  we obtain

$$x_{k+1} a_{\mu(k+1),n}^{-1} = a_{i(k+1),j(k+1)} a_{\mu(k+1),n}^{-1} \curvearrowright a_{i(k+1),n}^{-1} v$$

with some  $A_n^\pm$ -word  $v$ . In particular, we have  $\eta(k) = i(k+1) < j(k)$  and (10) is established. For  $k = h-1$  the left hand member of (10) is satisfied since  $j(h)$  is equal to  $n-1$  and  $r \leq n-1$  holds by definition of  $r$ . Properties (9) and (10) imply  $\mu(k) < j(k)$  for all  $k \in [0, h-2]$ . In particular

we have  $\mu(0) < j(0) = p$  together with  $wa_{r,n}^{-1} \curvearrowright a_{\mu(0),n}^{-1} \cdots$ . As  $a_{\mu(0),n}$  can not be a right divisor of  $a_{p,n}$  it follows that the left denominator of  $a_{p,n} w a_{r,n}^{-1}$  is also non empty and so that  $a_{r,n}$  is not a right divisor of  $a_{p,n}\beta$ .  $\square$

As the reader can see, the case  $p = n - 1$  is excluded from Proposition 4.5. It is the aim of the following result.

**Proposition 4.7.** *For  $\beta$  a non-trivial braid of  $B_{n-1}^{+*}$  there is equivalence between*

- (i) *the  $B_{n-1}^{+*}$ -tail of  $\phi_n(a_{n-1,n}\beta)$  is trivial,*
- (ii) *the only  $A_n$ -letter right dividing  $a_{n-1,n}\beta$  is  $\beta^\#$  which is of type  $a_{\dots,n-1}$ .*

*Proof.* (ii)  $\Rightarrow$  (i) is similar as (iii)  $\Rightarrow$  (i) of Proposition 4.7. We now show that (i) implies (ii). Condition (i) implies in particular that the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta)$  is trivial. It follows that the last letter of  $\beta$  is of type  $a_{\dots,n-1}$ . Let  $w$  be the rotating normal form of  $\beta$  and  $a_{r,s}$  be an  $A_n$ -letter different from  $\beta^\#$ . For  $s \leq n - 1$  we follow proof of Proposition 4.7 to obtain that  $a_{r,s}$  is not a right divisor or the braid  $a_{n-1,n}$ . Assume now  $s = n$ . By Lemma 4.6 there exists  $q \leq r$  such that  $wa_{r,n}^{-1} \curvearrowright a_{q,n}^{-1} \cdots$  holds and so we obtain  $a_{n-1,n}wa_{r,n}^{-1} \curvearrowright a_{n-1,n}a_{q,n}^{-1} \cdots$ . As, for  $q \neq n - 1$  the braid  $a_{q,n-1}$  is not a right divisor of  $a_{n-1,n}$  it is sufficient to show  $q \neq n - 1$  for concluding that  $a_{r,n}$  not right divides  $a_{n-1,n}\beta$ . For  $r \leq n - 2$  it is obvious since  $q \leq r$  holds. Assume finally  $r = n - 1$ . We denote by  $a_{p,n-1}$  the last letter of  $\beta$ . By (8) we have  $a_{p,n-1}a_{n-1,n}^{-1} \curvearrowright a_{p,n}^{-1} \cdots$  and then Lemma 4.6 implies  $q \leq p < n - 1$ , as expected.  $\square$

**Theorem 4.8.** *A finite sequence  $(\beta_b, \dots, \beta_1)$  of braids in  $B_{n-1}^{+*}$  is the  $\phi_n$ -splitting of a braid of  $B_n^{+*}$  if and only if*

- (i) *for  $k \geq 3$  and  $k = b$ , the braid  $\beta_k$  is not trivial,*
- (ii) *for  $k \geq 2$ , the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta_k)$  is trivial,*
- (iii) *if, for  $k \geq 3$ , we have  $\beta_k^\# \neq a_{n-2,n-1}$  then  $\beta_{k-1}$  contains an  $\phi_n(\beta_k^\#)$ -barrier.*

*Proof.* Let  $(\beta_b, \dots, \beta_1)$  be the  $\phi_n$ -splitting of some braid of  $B_{n-1}^{+*}$ . Condition (i) is a consequence of Lemma 2.8.(ii). Condition (6) implies that the  $B_{n-1}^{+*}$ -tail of

$$\phi_n^{b-k}(\beta_b) \cdot \dots \cdot \phi_n(\beta_{k+1})$$

is trivial for  $k \geq 1$ . In particular the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta_{k+1})$  must be trivial for  $k \geq 1$ , which implies (ii). Condition (iii) is Lemma 2.11.

Conversely, let us prove that a sequence  $(\beta_b, \dots, \beta_1)$  of braids of  $B_{n-1}^{+*}$  satisfying (i), (ii) and (iii) is the  $\phi_n$ -splitting of some braid of  $B_n^{+*}$ . Condition (i) implies that  $\beta_b$  is not trivial. For  $k \geq 2$  we denote by  $\gamma_k$  the braid  $\phi_n^{b-k}(\beta_b) \cdot \dots \cdot \phi_n(\beta_{k+1}) \cdot \beta_k$ . For  $k \geq 3$  and  $k \geq 2$  whenever  $\beta_2 \neq 1$ , we first prove

$$\beta_k^\# \text{ is the only } A_n\text{-letter that right divides } \gamma_k. \quad (11)$$

We note that Condition (i) guarantees the existence of  $\beta_k^\#$  for  $k \geq 3$ . For  $k = b$ , Condition (ii) implies that the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta_b)$  is trivial. Hence,

by Lemma 4.2 the only  $A_{n-1}$ -letter that right divides  $\beta_b$  is  $\beta_b^\#$ . Since any right divisors of a braid of  $B_{n-1}^{+*}$  lie in  $B_{n-1}^{+*}$ , we have established (11) for  $k = b$ . Assume (11) holds for  $k \geq 4$  or  $k \geq 3$  whenever  $\beta_2 \neq 1$  and let us prove it for  $k - 1$ . By Condition (ii) there exists  $p$  such that  $\beta_k^\#$  is  $a_{p-1, n-1}$ . We denote by  $ua_{p-1, n-1}$  and  $v$  two  $A_n$ -words representing  $\gamma_k$  and  $\beta_{k-1}$  respectively. The braid  $\gamma_{k-1}$  is then represented by  $\phi_n(u)a_{p, n}v$ . Let  $y$  be an  $A_n$ -letter different from  $\beta_{k-1}^\#$ . Proposition 4.5 with Condition (iii) and Proposition 4.7 imply that  $y$  is not a right divisor of  $a_{p, n}\beta_{k-1}$ . Therefore, by Proposition 3.6 there exists an  $A_n$ -letter  $x$  different from  $a_{p, n}$  such that  $\phi_n(u)a_{p, n}vy^{-1} \curvearrowright \phi_n(u)a_{p, n}x^{-1} \dots$ . The word  $\phi_n(u)a_{p, n}$  represents  $\phi_n(\gamma_k)$ . By induction hypothesis  $x$  is not a right divisor of  $\phi_n(\gamma_k)$ . Then Proposition 3.6 implies that  $D(\phi_n(u)a_{p, n}x^{-1})$  is not empty. It follows  $D(\phi_n(u)a_{p, n}vy^{-1}) \neq \varepsilon$  and so always by Proposition 3.6, the letter  $y$  is not a right divisor of  $\gamma_{k-1}$ . Eventually we have established (11) for  $k \geq 3$ .

A direct consequence of (11) and Condition (ii) is that the only  $A_n$ -letter right dividing  $\phi_n(\gamma_k)$  is of type  $a_{\dots, n}$  and so the  $B_{n-1}^{+*}$ -tail of the braid  $\gamma_k$  is trivial for  $k \geq 3$  and for  $k = 2$  whenever  $\beta_k \neq 1$ . It remains to establish that the  $B_{n-1}^{+*}$  tail of  $\phi_n(\gamma_2)$  is also trivial whenever  $\beta_2$  is trivial. Assume  $\beta_2 = 1$ . Condition (iii) implies  $\beta_3^\# = a_{n-2, n-1}$ . By (11),  $a_{n-2, n-1}$  is the only  $A_n$ -letter that right divides the braid  $\gamma_3$ . Since  $\gamma_2 = \phi_n(\gamma_3)$ , the letter  $\phi_n^2(a_{n-2, n-1}) = a_{1, n}$  is the only letter right dividing  $\phi_n(\gamma_2)$ . In particular the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\gamma_2)$  is trivial.  $\square$

Conditions (i), (ii) and (iii) are easy to check if the braids  $\beta_1, \dots, \beta_b$  are given by their rotating normal forms.

**Corollary 4.9.** *Let  $(w_b, \dots, w_1)$  be a finite sequence of  $A_{n-1}$ -words, then the word*

$$\phi_n^{b-1}(w_b) \cdot \dots \cdot \phi_n(w_2) \cdot w_1, \quad (12)$$

*is  $n$ -rotating if the following conditions are satisfied*

- (i) for  $k \geq 1$ , the word  $w_k$  is  $(n - 1)$ -rotating,
- (ii) for  $k \geq 3$ , the word  $w_k$  ends by  $a_{p-1, n-1}$  for some  $p$ ,
- (iii) the word  $w_2$  is either empty (except for  $b = 2$ ) or ends by  $a_{p-1, n-1}$  for some  $p$ ,
- (iv) if, for  $k \geq 3$ , the word  $w_k$  ends by  $a_{p-1, n-1}$  with  $p \neq n - 1$  then the word  $w_{k-1}$  contains an  $a_{p, n}$ -barrier.

*Proof.* Assume that  $(w_b, \dots, w_1)$  satisfies Conditions (i)-(iv) and let us prove that the word  $w$  defined at (12) is rotating.

We denote by  $\beta_i$  (resp.  $\beta$ ) the braid represented by  $w_i$  (resp.  $w$ ). By Condition (i) and Definition 2.5, the word  $w$  is rotating if and only if  $(\beta_b, \dots, \beta_1)$  is a  $\phi_n$ -splitting. Conditions (ii) and (iii) imply Condition (i) of Theorem 4.8. Theorem 4.8.(iii) is a consequence of (ii) and (iv). We remark that the  $B_{n-1}^{+*}$ -tail of a braid  $\gamma$  is represented by a suffix of the rotating word of  $\gamma$ . If the  $B_{n-1}^{+*}$ -tail of  $\phi_n(\beta_k)$  is not trivial, then there exists  $a_{p, q}$  with  $1 \leq p < q < n$  that right divides  $\phi_n(\beta_k)$ . As  $\beta_k$  lies in  $B_{n-1}^{+*}$ , we have  $p \neq 1$

and therefore  $\beta_k$  is right divisible by  $a_{p-1,q-1}$  with  $q-1 \leq n-2$ . Assume the  $B_{n-1}^{+*}$ -tail of  $w_k$  is not trivial for  $k \geq 2$ . The previous remark implies that  $w_k$  must end with a letter  $a_{i,j}$  satisfying  $j \leq n-2$ , which is in contradiction with Conditions (ii) and (iii).  $\square$

It is not true that any decomposition of an  $n$ -rotating word as in (12) satisfies Conditions (i)–(iv) of Corollary 12. However we have the following result.

**Proposition 4.10.** *For every  $n$ -rotating word  $w$  with  $n \geq 3$  there exists a unique sequence  $(w_b, \dots, w_1)$  of  $(n-1)$ -rotating words such that  $w$  decompose as in (12) and Conditions (ii)–(iv) of 4.9 hold.*

*Proof.* By definition of a rotating normal word and by Lemma 2.8 such a sequence exists. Let us prove now the unicity. Assume  $w$  is a  $n$ -rotating normal word and that  $(w_b, \dots, w_1)$  and  $(w'_c, \dots, w'_1)$  are two different sequences of  $(n-1)$ -rotating normal words satisfying Conditions (ii) and (iii) of Corollary 4.9. Let  $k$  be the minimal integer satisfying  $w_k \neq w'_k$ . Since the sum of the word lengths of the two sequences are the same, we have  $k \leq \min\{b, c\}$ . Without loss of generality, we may assume that  $w'_k$  is a proper suffix of  $w_k$ , i.e.,  $w_k = u \cdot w'_k$ . By Conditions (ii) and (iii) of Corollary 4.9, the last letter  $x$  of  $u$  comes from the last letter of  $w'_{k+1}$  or  $w'_{k+2}$ . Hence the letter  $x$  is equal to  $a_{p-1,n-1}$  for some  $p$  and  $w_k$  admits either  $\phi_n(a_{p-1,n-1})w'_k = a_{p,n}w'_k$  or  $\phi_n^2(a_{p-1,n-1})w'_k = a_{1,p+1}w'_k$  as suffix. The first case is impossible since  $w_k$  is an  $A_{n-1}$ -word. The second case may occur only for  $k=1$  and  $w'_2 = \varepsilon$ . As  $w'_2$  is empty, the last letter of  $w'_3$ , which is  $x$ , is equal to  $a_{n-2,n-1}$ . This implies that  $w_k$  admits  $a_{1,n}u$  as suffix which is also impossible since it is an  $A_{n-1}$ -word.  $\square$

A direct consequence of Corollary 4.9 and Proposition 4.10 is

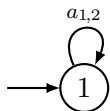
**Theorem 4.11.** *An  $A_n$ -word  $w$  is rotating if and only if it can be expressed as in (12) subject to Conditions (i)–(iv) of Corollary 4.9.*

## 5. REGULARITY

In this section we will show that the language of  $n$ -rotating words, denoted by  $R_n$  is regular, i.e., there exists a finite state automaton recognizing the  $n$ -rotating words. As the rotating normal form is defined using right division it is more natural for an automaton to read word from the right. For  $w = x_0 \dots x_k$  an  $A_n$ -word we will denote by  $\Pi(w)$  the word  $x_k \dots x_0$ . By Theorem 1.2.8 of [7] the language  $R_n$  is regular if and only if the language  $\Pi(R_n)$  is. In this section we will construct an automaton recognizing  $\Pi(R_n)$ .

For us a *finite state automaton* is a quintuplet  $(S \cup \{\otimes\}, A, \mu, Y, i)$  where  $S$  is the finite set of *states*,  $A$  is a finite *alphabet*,  $\mu : S \times A \rightarrow S$  is the *transition function*,  $Y \subseteq S$  is *accepting states* and  $i$  is the *initial state*. In this paper each automaton is equipped with an undraw dead state  $\otimes$  and all states except the dead one is accepting, i.e.,  $Y = S$  always holds.

Therefore an automaton will be briefly denoted by  $\mathcal{A} = (S, A, \mu, i)$ . To describe  $\mathcal{A}$  it is enough to describe  $\mu$  on  $(s, x) \in S \times A$  where  $\mu(s, x) \neq \otimes$  and  $s \neq \otimes$ . By example an automaton recognizing the language  $R_2$  is  $\mathcal{A}_2 = (\{1\}, \{a_{1,2}\}, \mu_2, 1)$  with  $\mu_2(1, a_{1,2}) = 1$ . The corresponding automaton diagram is :



The horizontal arrow points to the initial state.

**Proposition 5.1.** *An  $A_3$ -word  $x_b^{e_b} \cdot \dots \cdot a_{1,3}^{e_3} a_{2,3}^{e_2} a_{1,2}^{e_1}$  where*

$$x_b = \begin{cases} a_{1,2} & \text{if } b \equiv 1 \pmod{3}, \\ a_{2,3} & \text{if } b \equiv 2 \pmod{3}, \\ a_{1,3} & \text{if } b \equiv 3 \pmod{3}. \end{cases}$$

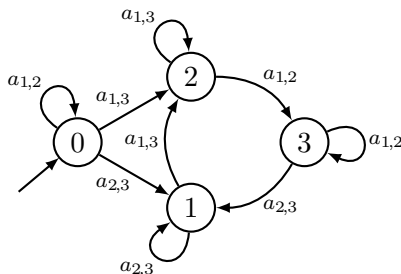
*is rotating if and only if  $e_k \neq 0$  for all  $k \geq 3$ .*

*Proof.* The 2-rotating words are powers of  $a_{1,2}$ . Let  $w$  be the word of the statement. Defining  $w_k$  to be  $a_{1,2}^{e_k}$ , we obtain

$$w = \phi_n^{b-1}(w_b) \cdot \dots \cdot \phi_n(w_2) \cdot w_1.$$

As there is no barrier in  $B_3^{+*}$ , the word  $w$  is rotating if and only if it satisfies Conditions (ii) and (iii) of Corollary 4.9, *i.e.*, the exponent  $e_k$  is not 0 for  $k \geq 3$ .  $\square$

As a consequence the following automaton recognizes the language  $\Pi(R_3)$ :



Unfortunately, for  $n \geq 4$  there is no so simple characterization of  $n$ -rotating words. We will describe an inductive construction for an automaton recognizing language  $\Pi(R_n)$ . The process will be illustrated on  $n = 4$ . The first step is to focus on  $n$ -rotating words ending with a letter of type  $a_{\dots n}$ .

**Definition 5.2.** We denote by  $R_n^*$  the language of  $n$ -rotating words which are empty or ends with a letter of type  $a_{p,n}$  for some  $p$ .

Before constructing an automaton  $\mathcal{A}_n$  recognizing the language  $\Pi(R_n)$ , we construct by induction on  $n \geq 3$  an automaton  $\mathcal{A}_n^*$  for the language  $\Pi(R_n^*)$ .

**Definition 5.3.** A *partial automaton* is a quadruplet  $P = (S, A, \mu, I)$  where  $S$ ,  $A$  and  $\mu$  are defined as for an automaton and  $I : A \rightarrow S$  is a map. The closure of a partial automaton  $P$  is the automaton  $\mathcal{A}(P) = (S \cup \{\circ\}, A, \mu^c, \circ)$  given by

$$\mu^c(s, x) = \begin{cases} I(x) & \text{if } s = \circ, \\ \mu(s, x) & \text{otherwise.} \end{cases}$$

A partial automaton is represented as an automaton excepted for the function  $I$ . For each  $x \in A$  we draw an arrow attached to state  $I(x)$  and labelled  $x$ . We say that a partial automaton recognizes a given language if its closure does.

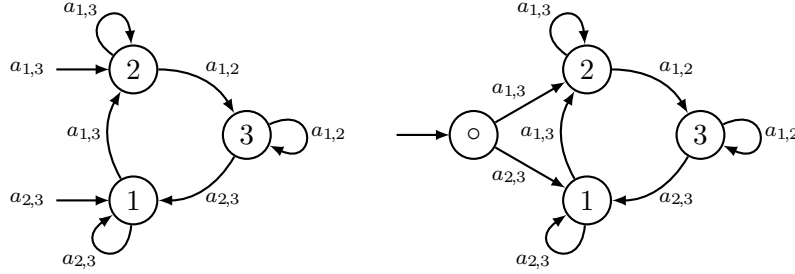


FIGURE 6. The partial automaton  $P_3$  and the corresponding closure which recognizes the language  $\Pi(R_3^*)$ .

We will now show how to construct by induction a partial automaton  $P_n$  recognizing  $\Pi(R_n^*)$  for  $n \geq 3$ . For  $n = 3$  this is already done by Figure 6. For the sequel we assume  $n \geq 4$  and that  $P_{n-1} = (S_{n-1}, A_{n-1}, \mu_{n-1}, I_{n-1})$  is a given partial automaton which recognizes the language  $\Pi(R_{n-1})$ .

We define  $S_n^0$  to be the set

$$S_n^0 = \{0\} \times (S_{n-1} \setminus \{\otimes\}) \times \mathcal{P}(\{a_{2,n}, \dots, a_{n-2,n}\}).$$

A state in  $S_n^0$  is then written  $(0, s, m)$ . For  $x = a_{i,j} \in A_{n-1}$  we denote by  $\text{bar}(x)$  the set  $\{a_{p,n} \mid i < p < j\}$ .

**Definition 5.4.** We define  $P_n^0 = (S_n^0 \cup \{\otimes\}, A_{n-1}, \mu_n^0, I_n^0)$  to be the partial automaton where for all  $x \in A_{n-1}$ ,

$$I_n^0(x) = \begin{cases} (0, I_{n-1}(x), \text{bar}(x)) & \text{if } I_{n-1}(x) \neq \otimes, \\ \otimes & \text{if } I_{n-1}(x) = \otimes. \end{cases}$$

and for all  $(0, s, m) \in S_n^0$  and for all  $x \in A_{n-1}$ ,

$$\mu_n^0((0, s, m), x) = \begin{cases} (0, \mu_{n-1}(s, x), m \cup \text{bar}(x)) & \text{if } \mu_{n-1}(s, x) \neq \otimes, \\ \otimes & \text{if } \mu_{n-1}(s, x) = \otimes. \end{cases}$$



**Proposition 5.5.** *The partial automaton  $P_n^0$  recognizes the language  $\Pi(R_{n-1}^*)$ . Moreover an accepted  $A_n$ -word  $\Pi(w)$  contains an  $a_{p,n}$ -barrier if and only if  $P_{n,0}$  has state  $(s, m)$  with  $a_{p,n} \in m$  after reading  $\Pi(w)$ .*

*Proof.* Let  $w$  be an  $A_n$ -word of length  $\ell$ ,  $\mathcal{A}$  and  $\mathcal{A}'$  be the closure of  $P_{n-1}$  and  $P_n^0$  respectively. We denote by  $s_k$  (resp.  $(s'_k, m_k)$ ) the state of automaton  $\mathcal{A}$  (resp.  $\mathcal{A}'$ ) after reading the  $k$ -th letter of  $\Pi(w)$ . If  $w$  does not contains an  $a_{p,n}$ -barrier then  $(s'_k, m_k)$  is equal to  $(s_k, \emptyset)$  for all  $k \in [1, \ell]$ . Hence  $\Pi(w)$  is accepted or not by the two automata and in particular  $m_\ell$  is the empty set. Assume now  $w$  contains an  $a_{p,n}$ -barrier. Let  $\ell'$  be the first occurrence on such a barrier in  $\Pi(w)$ . By construction of  $\mu_n$  we have  $(s'_k, m_k)$  with  $a_{p,n} \in m_k$  for  $k \geq \ell'$  except if  $s_k = \otimes$ . As  $\mathcal{A}$  (resp.  $\mathcal{A}'$ ) recognizes the word  $\Pi(w)$  if and only if  $s_\ell$  (resp.  $s'_\ell$ ) is different from  $\otimes$ , the word  $w$  is recognized or not by both automata. Moreover, in this case  $m_\ell$  contains  $a_{p,n}$ .  $\square$

As the only  $a_{p,4}$ -barrier in  $A_4$  is  $a_{1,3}$ , the partial automaton  $P_4^0$  is obtained from  $P_3$  by connecting edges labelled  $a_{1,3}$  to a copy of  $P_3$ , as illustrated on figure 5

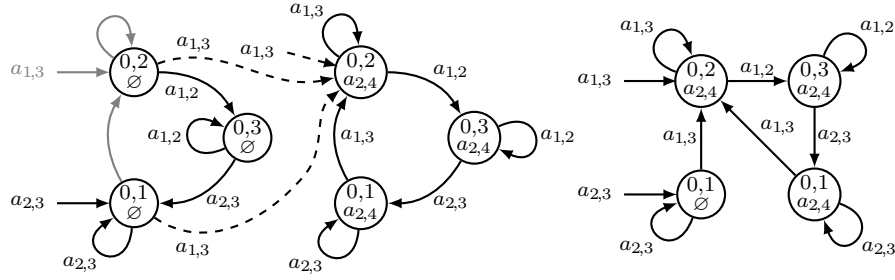


FIGURE 7. The partial automaton  $P_4^0$ . Obsolete transitions from  $P_3$  are in gray. New added transitions are dashed. The right partial automaton is  $P_4^0$  without inaccessible states.

For  $t = (0, s, m) \in S_n^0$  we define  $\phi_n^k(t)$  to be  $(k, s, m)$ . We also define  $S_n^k$  to be  $\phi_n^k(S_n^0)$  and

$$P_n^k = \left( S_n^k, \phi_n^k(A_{n-1}), \mu_n^k, I_n^k \right)$$

to be the partial automaton given by  $I_n^k(\phi_n^k(x)) = \phi_n^k(I_n^0(x))$  and

$$\mu_n^k((k, s, m), \phi_n^k(x)) = \phi_n^k(\mu_n^0((0, s, m), x))$$

with the convention  $\phi_n^k(\otimes) = \otimes$ . In other words,  $P_n^k$  is obtained from  $P_n^0$  by replacing the letter  $x$  by  $\phi_n^k(x)$  and state  $(0, s, m)$  by  $(k, s, m)$ . We obtain immediately that  $P_n^k$  recognizes the word  $\phi_n^k(\Pi(w))$  if and only if  $P_n^0$  recognizes  $\Pi(w)$ .

We can now construct the partial automaton  $P_n$  by plugging together  $n$  partial automaton  $P_n^k$  for  $k \in [0, n-1]$  together.

**Definition 5.6.** We define  $P_n = (S_n^* \cup \{\otimes\}, A_n, \mu_n^*, I_n)$ , with  $S_n^* = S_n^0 \sqcup \dots \sqcup S_n^k$  to be the partial automaton given by

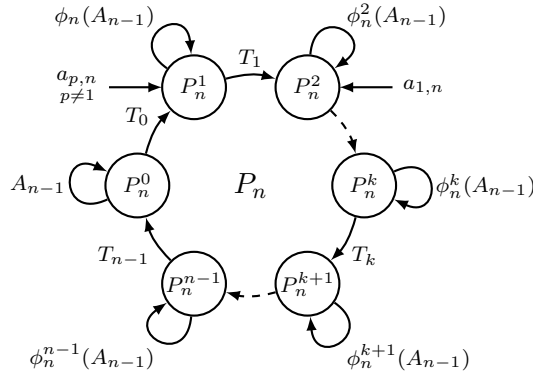
$$I_n(x) = \begin{cases} I_n^1(a_{p-1,n-1}) & \text{if } x = a_{p,n} \text{ with } p \neq 1, \\ I_n^2(a_{n-2,n-1}) & \text{if } x = a_{1,n}, \\ \otimes & \text{otherwise.} \end{cases}$$

and with transition function

$$\mu_n^*((k, s, m), \phi_n^k(x)) = \begin{cases} \mu_n^k((k, s, m), \phi_n^k(x)) & \text{if } x \in A_{n-1}, \\ I_n^{k+1}(\phi_n^k(x)) & \text{if } x = a_{n-1,n} \\ I_n^{k+1}(\phi_n^k(x)) & \text{if } x = a_{p,n} \text{ with } 2 \leq p \leq n-2 \\ & \text{and } a_{p,n} \in m, \\ \otimes & \text{otherwise} \end{cases}$$

with the convention  $I_n^n = I_n^0$ .

We summarize the construction of the partial automaton  $P_n$  on the following diagram.



An arrow labelled  $T_k$  represents the set of transitions  $\mu_n^*((k, s, m), \phi_n^k(a_{p,n}))$ .

**Lemma 5.7.** *The partial automaton  $P_n$  recognizes the language  $\Pi(R_n^*)$ .*

*Proof.* Let  $\mathcal{A}$  be the closure of  $P_n$  and  $w$  be a non empty  $A_n$ -word. There exists a unique sequence  $(w_b, \dots, w_1)$  of  $A_{n-1}$ -words such that  $w_b \neq \varepsilon$ ,  $w$  is equal to

$$\phi_n^{b-1}(w_b) \cdot \dots \cdot \phi_n(w_2) \cdot w_1$$

and for all  $i$ , the word  $\phi_n^i(w_i)$  is the maximal suffix of  $\phi_n^{i-1}(w_i) \cdot \dots \cdot \phi_n^1(w_i)$  belonging to  $\phi_n^i(A_{n-1})$ . By definition of  $I_n$ , the word  $\Pi(w)$  is accepted by  $P_n$  only if  $w$  ends by a letter  $a_{p,n}$  for some  $p$ . We assume now that  $w$  is such a word. Thus the first integer  $j$  such that  $w_j$  is non empty is 2 or 3. More precisely, we have  $j = 2$  if  $p > 1$  and  $j = 3$  if  $p = 1$  holds. In both cases, the reading of  $\Pi(w)$  starts by a state coming from  $P_n^j$ . The automaton reaches a state different from one of  $P_n^j$  if it goes to the state  $\otimes$  or if it reads a letter

outside of  $\phi_n^{j-1}(A_n)$ , *i.e.*, a letter of  $\phi_n^j(w_{j+1})$ . This is a general principle : after reading a letter of  $\phi_n^{i-1}(w_i)$  the automaton  $\mathcal{A}$  is in state  $(t, s, m)$  with  $t = i \bmod n$ . By construction of  $P_n^t$ , the word  $\phi_n^{i-1}(w_i)$  provides an accepted state if and only if  $w_i$  is a word of  $\Pi(R_{n-1})$ . At this point we have shown that  $\Pi(w)$  is accepted by  $\mathcal{A}$  only if  $w$  is empty or if  $w$  satisfies  $w^{\#} = a_{p,n}$  together with Conditions (i), (ii) and (iii) of Corollary 4.9. Let  $i$  be in  $[j, k-1]$ , and assume that  $\mathcal{A}$  is in an acceptable state  $(t, s, m)$  with  $t = i \bmod n$  after reading the word  $\Pi(\phi_n^{i-1}(w_i) \cdots \phi_n(w_2) \cdot w_1)$ . We denote by  $x$  the letter  $w_{i+1}^{\#}$ . By construction of  $w_{i+1}$  we have  $x \notin \phi_n^{i-1}(A_{n-1})$  and so  $x = \phi_n^i(a_{p,n})$  for some  $p$ . By definition of  $\mu_n^*$  we have  $\mu_n^*((t, s, m), \phi_n^i(a_{p,n})) \neq \otimes$  if and only if  $p = n-1$  or  $p \in [2, n-2]$  and  $a_{p,n} \in m$ . By construction of  $P_n^t$ , we have  $a_{p,n} \in m$  if and only if  $w_i$  contains an  $a_{p,n}$ , which corresponds to Condition (iv) of Corollary 4.9. Eventually, by Corollary 4.9, the word  $\Pi(w)$  is accepted by  $\mathcal{A}$  if and only if  $w \in R_n^*$ .  $\square$

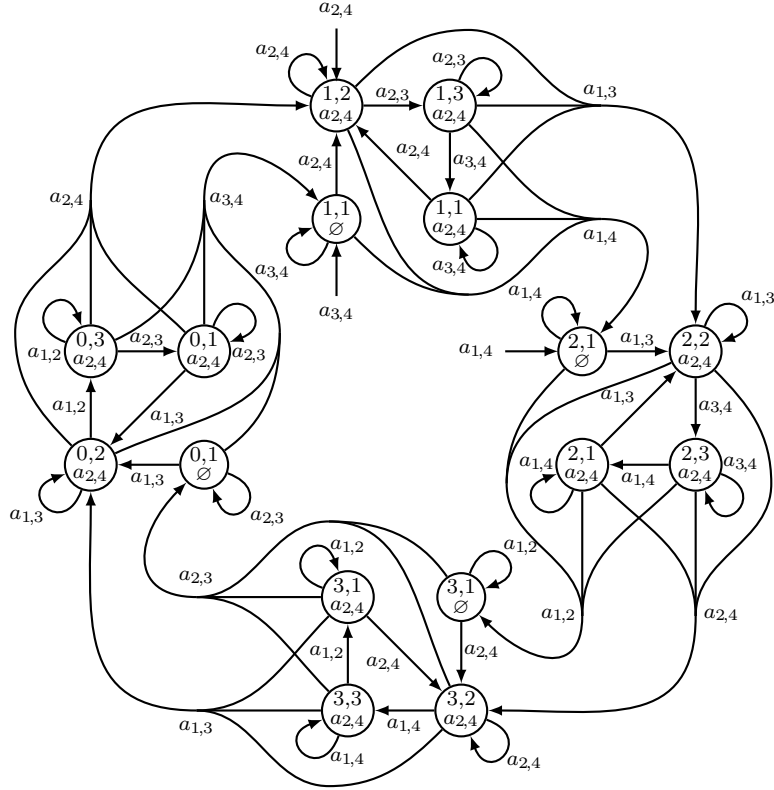


FIGURE 8. Partial automaton recognizing the language  $\Pi(R_4^*)$ .

Assume that an automaton  $\mathcal{A}_{n-1} = (S_{n-1} \cup \{\otimes\}, A_{n-1}, \mu_{n-1}, i)$  recognizing the language  $\Pi(R_{n-1})$  for  $n \geq 4$  is given. Using the partial automaton  $P_n = (S_n^* \cup \{\otimes\}, A_n, \mu_n^*, I_n)$  we construct the automaton  $\mathcal{A}_n = (S_n \cup \{\otimes\}, A_n, \mu_n, i)$

defined by  $S_n = S_{n-1} \sqcup S_n^*$  and

$$\mu_n(s, x) = \begin{cases} \mu_{n-1}(s, x) & \text{if } s \in S_{n-1} \text{ and } x \in A_{n-1}, \\ I_n(x) & \text{if } s \in S_{n-1} \text{ and } x \in A_n \setminus A_{n-1}, \\ \mu_n^*(s, x) & \text{if } s \in S_n^*. \end{cases}$$

**Proposition 5.8.** *If  $\mathcal{A}_{n-1}$  recognizes  $\Pi(R_{n-1})$ , the automaton  $\mathcal{A}_n$  recognizes the language  $\Pi(R_n)$ .*

*Proof.* Let  $w$  be an  $A_n$ -word,  $w_1$  be the maximal suffix of  $w$  which is an  $A_{n-1}$ -word and  $w'$  be the corresponding prefix. By Corollary 4.9, the word  $w$  is rotating if and only if  $w_1$  and  $w'$  are. By construction of  $\mathcal{A}_n$ , the automaton is in acceptable state after reading  $\Pi(w_1)$  if and only if  $w_1$  is an  $(n-1)$ -rotating word. Hence  $w$  is accepted only if  $w_1$  is rotating. Assume that it is the case. By Lemma 5.7 the automaton  $\mathcal{A}_n$  is always in an acceptable state after reading  $\Pi(w')$  if and only if the word  $w'$  is rotating. Eventually the word  $\Pi(w)$  is accepted by  $\mathcal{A}$  if and only if  $w_1$  and  $w'$  are both rotating, which is equivalent to  $w$  is rotating.  $\square$

By Proposition 5.8, the language  $\Pi(R_n)$  is regular and so we obtain:

**Theorem 5.9.** *The language of  $n$ -rotating words  $R_n$  is regular.*

#### FURTHER WORK

Using syntactical characterization of rotating words we have proved that the language of  $n$ -rotating words is regular. For  $W$  a finite state automaton, we denote by  $L(W)$  the language recognized by  $W$ . Following [4] and [7] we have the following definition:

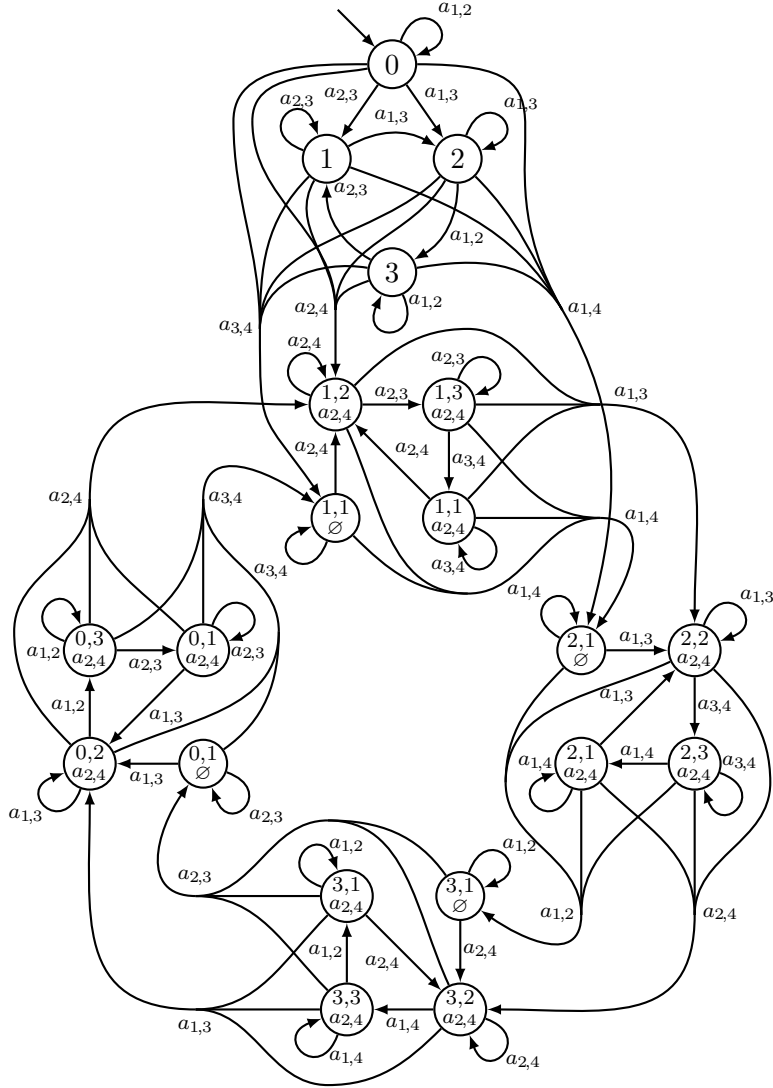
**Definition 5.10.** Let  $M$  be a monoid. A *right automatic structure*, resp. *left automatic structure*, on  $M$  consists of a set  $A$  of generators of  $M$ , a finite state automaton  $W$  over  $A$ , and finite state automata  $M_x$  over  $(A, A)$ , for  $x \in A \cup \{\varepsilon\}$ , satisfying the following conditions:

- (i) the map  $\pi : L(W) \rightarrow M$  is surjective.
- (ii) for  $x \in A \cup \{\varepsilon\}$ , we have  $(u, v) \in L(M_x)$  if and only if  $\overline{ux} = \overline{y}$ , resp.  $\overline{xu} = \overline{y}$ , and both  $u$  and  $v$  are elements of  $L(W)$ .

Naturally we can ask if the rotating normal form provides an left or right automatic structure for the dual braid monoid  $B_n^{+*}$ . Such a result may needs to obtain some syntactical properties on the word  $xw$  or  $wx$  where  $w$  is an  $n$ -rotating word and  $x$  is an  $A_n$ -generator. At this time no result have been obtained in this direction.

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FIGURE 9. Automaton  $\mathcal{A}_4$  for the language  $\Pi(R_4)$ .

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