



Non-uniform cellular automata: Classes, dynamics, and decidability

Alberto Dennunzio, Enrico Formenti, Julien Provillard

► To cite this version:

Alberto Dennunzio, Enrico Formenti, Julien Provillard. Non-uniform cellular automata: Classes, dynamics, and decidability. Information and Computation, 2012, 10.1016/j.ic.2012.02.008 . hal-01297590

HAL Id: hal-01297590

<https://hal.science/hal-01297590>

Submitted on 4 Apr 2016

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Non-Uniform Cellular Automata: classes, dynamics, and decidability[☆]

Alberto Dennunzio^{*,a}, Enrico Formenti^{*,b}, Julien Provillard^b

^a *Università degli studi di Milano-Bicocca, Dipartimento di Informatica Sistemistica e Comunicazione, viale Sarca 336, 20126 Milano (Italy)*

^b *Université Nice-Sophia Antipolis, Laboratoire I3S, 2000 Route des Colles, 06903 Sophia Antipolis (France)*

Abstract

The dynamical behavior of non-uniform cellular automata is compared with the one of classical cellular automata. Several differences and similarities are pointed out by a series of examples. Decidability of basic properties like surjectivity and injectivity is also established. The final part studies a strong form of equicontinuity property specially suited for non-uniform cellular automata.

Key words: cellular automata, non-uniform cellular automata, decidability, symbolic dynamics

1. Introduction

A complex system is (roughly) defined by a multitude of simple individuals which cooperate to build a complex (unexpected) global behavior by local interactions. Cellular automata (CA) are often used to model complex systems when individuals are embedded in a uniform “universe” in which local interactions are the same for all. Indeed, a cellular automaton is made of identical finite automata arranged on a regular lattice. Each automaton updates its state by a local rule on the basis of its state and the one of a fixed set of neighbors. At each time-step, the same (here comes uniformity) local rule is applied to all finite automata in the lattice. For recent results on CA dynamics and an up-to-date bibliography see for instance [15, 21, 4, 18, 14, 7, 8, 1, 6, 13, 12, 11].

In a number of situations one needs a more general setting. One possibility consists in relaxing the uniformity constraint. This choice may result winning for example for

[☆]This is an extended and improved version of the paper [2] presented at LATA2009 conference.

*Corresponding author.

Email addresses: `dennunzio@disco.unimib.it` (Alberto Dennunzio), `enrico.formenti@unice.fr` (Enrico Formenti), `julien.provillard@i3s.unice.fr` (Julien Provillard)

Complexity design control. In many phenomena, each individual locally interacts with others but maybe these interactions depend on the individual itself or on its position in the space. For example, when studying the formation of hyper-structures in cells, proteins move in the cellular soup and do not behave just like billiard balls. They chemically interact each other only when they meet special situations (see for instance, [22]) or when they are at some special places (in rybosomes for instance). It is clear that one might simulate all those situations by a CA but the writing of a single local rule will be an excessive difficult task, difficult to control. A better option would be to write simpler (but different) local rules that are applied only at precise positions so that less constraints are to be taken into account at each time.

Structural stability. Assume that we are investigating the robustness of a system w.r.t. some specific property P . If some individuals change their “standard” behavior does the system still have property P ? What is the “largest” number of individuals that can change their default behavior so that the system does not change its overall evolution?

Reliability. CA are more and more used to perform fast parallel computations (beginning from [5], for example). Each cell of the CA can be implemented by a simple electronic device (FPGAs for example) [23]. Then, how reliable are computations w.r.t. failure of some of these devices? (Here failure is interpreted as a device which behaves differently from its “default” way).

The generalization of CA to non-uniform CA (ν -CA) has some interest in its own since the new model coincides with the set of continuous functions in Cantor topology. It is clear that the class of continuous functions is too large to be studied fruitfully. In the present paper, we present several sub-classes that are also interesting in applications. First of all, we show that several classical results about the dynamical behavior of CA are no longer valid in the new setting. Even when the analysis is restricted to smaller classes of non-uniform CA, the overall impression is that new stronger techniques will be necessary to study ν -CA. However, by generalizing the notion of De Bruijn graph, we could prove the decidability of basic set properties like surjectivity and injectivity. We recall that these property are often necessary conditions of many classical definitions of deterministic chaos.

Keeping on with surjectivity and injectivity, we give a partial answer about reliability and structural stability questions issued above. More precisely, we answer the following question: assuming to perturb some CA in some finite number of sites, if one knows that the corresponding ν -CA is surjective (resp., injective) does this imply that the original CA was surjective (resp., injective)?

The final part starts going more in deep with the study of the long-term dynamical behavior of ν -CA. Indeed, we show that under some conditions, if a ν -CA is a perturbed version of some equicontinuous or almost equicontinuous CA, then it shares the same dynamics.

Finally, we develop some complex examples showing that even small perturbations of an almost equicontinuous CA can give raise to sensitive to initial

conditions behavior or to equicontinuous dynamics.

We might conclude that breaking uniformity property in a CA may cause a dramatic change in the dynamical behavior.

2. Background

In this section, we briefly recall standard definitions about CA and dynamical systems. For introductory matter see [20], for instance. For all $i, j \in \mathbb{Z}$ with $i \leq j$ (resp., $i < j$), let $[i, j] = \{i, i+1, \dots, j\}$ (resp., $[i, j) = \{i, \dots, j-1\}$). Let $\mathbb{N}_+$ be the set of positive integers.

Configurations and CA. Let A be an alphabet. A *configuration* is a function from $\mathbb{Z}$ to A . The *configuration set* $A^{\mathbb{Z}}$ is usually equipped with the metric d defined as follows

$$\forall x, y \in A^{\mathbb{Z}}, d(x, y) = 2^{-n}, \text{ where } n = \min \{i \geq 0 : x_i \neq y_i \text{ or } x_{-i} \neq y_{-i}\} .$$

When A is finite, the set $A^{\mathbb{Z}}$ is a compact, totally disconnected and perfect topological space (i.e., $A^{\mathbb{Z}}$ is a Cantor space). For any pair $i, j \in \mathbb{Z}$, with $i \leq j$, and any configuration $x \in A^{\mathbb{Z}}$ we denote by $x_{[i, j]}$ the word $x_i \cdots x_j \in A^{j-i+1}$, i.e., the portion of x inside the interval $[i, j]$. Similarly, $u_{[i, j]} = u_i \cdots u_j$ is the portion of a word $u \in A^l$ inside $[i, j]$ (here, $i, j \in [0, l)$). In both the previous notations, $[i, j]$ can be replaced by $[i, j)$ with the obvious meaning. For any word $u \in A^*$, $|u|$ denotes its length. A *cylinder* of block $u \in A^k$ and position $i \in \mathbb{Z}$ is the set $[u]_i = \{x \in A^{\mathbb{Z}} : x_{[i, i+k)} = u\}$. Cylinders are clopen sets w.r.t. the metric d and they form a basis for the topology induced by d . A configuration x is said to be *a-finite* for some $a \in A$ if there exists $k \in \mathbb{N}$ such that $x_i = a$ for any $i \notin [-k, k]$. In the sequel, the collection of the *a-finite* configurations for a certain a will be simply called set of finite configurations.

A (one-dimensional) CA is a structure $\langle A, r, f \rangle$, where A is the *alphabet*, $r \in \mathbb{N}$ is the *radius* and $f : A^{2r+1} \rightarrow A$ is the *local rule* of the automaton. The local rule f induces a *global rule* $F : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ defined as follows,

$$\forall x \in A^{\mathbb{Z}}, \forall i \in \mathbb{Z}, \quad F(x)_i = f(x_{i-r}, \dots, x_{i+r}) . \quad (1)$$

Recall that F is a uniformly continuous map w.r.t. the metric d .

With an abuse of notation, a CA local rule f is extended to the function $f : A^* \rightarrow A^*$ which map any $u \in A^*$ of length l to the word $f(u)$ such that $f(u) = \epsilon$ (empty word), if $l \leq 2r$, and $f(u)_i = f(u_{[i, i+2r]})$ for each $i \in [0, l-2r)$, otherwise.

DTDS. A *discrete time dynamical system* (DTDS) is a pair (X, G) where X is a set equipped with a distance d and $G : X \mapsto X$ is a map which is continuous on X with respect to the metric d . When $A^{\mathbb{Z}}$ is the configuration space equipped with the above introduced metric and F is the global rule of a CA, the pair $(A^{\mathbb{Z}}, F)$ is a DTDS. From now on, for the sake of simplicity, we identify a CA with the dynamical system induced by itself or even with its global rule F .

Given a DTDS (X, G) , an element $x \in X$ is an *ultimately periodic point* if there exist $p, q \in \mathbb{N}$ such that $G^{p+q}(x) = G^q(x)$. If $q = 0$, then x is a *periodic point*, i.e., $G^p(x) = x$. The minimum p for which $G^p(x) = x$ holds is called *period* of x .

Recall that a DTDS (X, g) is *sensitive to the initial conditions* (or simply *sensitive*) if there exists a constant $\varepsilon > 0$ such that for any element $x \in X$ and any $\delta > 0$ there is a point $y \in X$ such that $d(y, x) < \delta$ and $d(G^n(y), G^n(x)) > \varepsilon$ for some $n \in \mathbb{N}$. A DTDS (X, G) is *positively expansive* if there exists a constant $\varepsilon > 0$ such that for any pair of distinct elements x, y we have $d(G^n(y), G^n(x)) \geq \varepsilon$ for some $n \in \mathbb{N}$. If X is a perfect set, positive expansivity implies sensitivity. Recall that a DTDS (X, g) is *(topologically) transitive* if for any pair of non-empty open sets $U, V \subseteq X$ there exists an integer $n \in \mathbb{N}$ such that $g^n(U) \cap V \neq \emptyset$. A DTDS (X, G) is *surjective* (resp., *injective*) iff G is surjective (resp., G is injective).

Let (X, G) be a DTDS. An element $x \in X$ is an *equicontinuity point* for G if $\forall \varepsilon > 0$ there exists $\delta > 0$ such that for all $y \in X$, $d(y, x) < \delta$ implies that $\forall n \in \mathbb{N}$, $d(G^n(y), G^n(x)) < \varepsilon$. For a CA F , the existence of an equicontinuity point is related to the existence of a special word, called *blocking word*. A word $u \in A^k$ is *s-blocking* ($s \leq k$) for a CA F if there exists an offset $j \in [0, k - s]$ such that for any $x, y \in [u]_0$ and any $n \in \mathbb{N}$, $F^n(x)_{[j, j+s-1]} = F^n(y)_{[j, j+s-1]}$. A word $u \in A^k$ is said to be *blocking* if it is *s-blocking* for some $s \leq k$. A DTDS is said to be *equicontinuous* if $\forall \varepsilon > 0$ there exists $\delta > 0$ such that for all $x, y \in X$, $d(x, y) < \delta$ implies that $\forall n \in \mathbb{N}$, $d(G^n(x), G^n(y)) < \varepsilon$. If X is a compact set, a DTDS (X, G) is equicontinuous iff the set E of all its equicontinuity points is the whole X . A DTDS is said to be *almost equicontinuous* if E is residual (i.e., E contains an infinite intersection of dense open subsets). In [19], Kůrka proved that a CA is almost equicontinuous iff it is non-sensitive iff it admits a *r-blocking* word.

Recall that two DTDS (X, G) and (X', G') are *topologically conjugated* if there exists a homeomorphism $\phi : X \rightarrow X'$ such that $G' \circ \phi = \phi \circ G$. In that case, (X, G) and (X', G') share some properties such as surjectivity, injectivity, transitivity.

3. Non-Uniform Cellular Automata

The meaning of (1) is that the same local rule f is applied to each site of the CA. Relaxing this constraint gives us the definition of a ν -CA. More formally one can give the following notion.

Definition 3.1 (Non-Uniform Cellular Automaton (ν -CA)).

A *Non-Uniform Cellular Automaton (ν -CA)* is a structure $(A, \{h_i, r_i\}_{i \in \mathbb{Z}})$ defined by a family of local rules $h_i : A^{2r_i+1} \rightarrow A$ of radius r_i all based on the same alphabet A .

Similarly to CA, one can define the global rule of a $\nu\text{-CA}$ as the map $H : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ given by the law

$$\forall x \in A^{\mathbb{Z}}, \forall i \in \mathbb{Z}, \quad H(x)_i = h_i(x_{i-r_i}, \dots, x_{i+r_i}) \quad .$$

From now on, we identify a $\nu\text{-CA}$ (resp., CA) with the discrete dynamical system induced by itself or even with its global rule H (resp., F).

It is well known that the Hedlund's Theorem [16] characterizes CA as the class of continuous functions commuting with the shift map $\sigma : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$, where $\forall x \in A^{\mathbb{Z}}, \forall i \in \mathbb{Z}, \sigma(x)_i = x_{i+1}$. It is straightforward to prove that a function $H : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ is the global map of a $\nu\text{-CA}$ iff it is continuous. i.e., in other words, iff the pair $(A^{\mathbb{Z}}, H)$ is a DTDS. Remark that the definition of $\nu\text{-CA}$ is by far too general to be useful. Therefore, we are going to focus our attention only over three special subclasses of $\nu\text{-CA}$.

Definition 3.2 (d $\nu\text{-CA}$). A $\nu\text{-CA}$ $(A, \{h_i, r_i\}_{i \in \mathbb{Z}})$ is a d $\nu\text{-CA}$ if there exist two naturals k, r and a rule $h : A^{2r+1} \rightarrow A$ such that $h_i = h$ for all integers i with $|i| > k$. In this case, we say that the given $\nu\text{-CA}$ has h as *default rule*.

Definition 3.3 (p $\nu\text{-CA}$). A $\nu\text{-CA}$ $(A, \{h_i, r_i\}_{i \in \mathbb{Z}})$ is a p $\nu\text{-CA}$ if there exist two naturals k, r , a *structural period* $p > 0$, and two sets $\{f_0, \dots, f_{p-1}\}$ and $\{g_0, \dots, g_{p-1}\}$ of rules of radius r such that for any integer i with $|i| > k$

$$h_i = \begin{cases} f_{i \bmod p} & \text{if } i > k \\ g_{i \bmod p} & \text{if } i < -k \end{cases}$$

If $p = 1$, we say that the given $\nu\text{-CA}$ has f_0 and g_0 as right and left default rules, respectively.

Definition 3.4 (r $\nu\text{-CA}$). A $\nu\text{-CA}$ $H = (A, \{h_i, r_i\}_{i \in \mathbb{Z}})$ is a r $\nu\text{-CA}$ if there exists an integer r such that $\forall i \in \mathbb{Z}, r_i = r$. In this case, we say that H has radius r .

The first two class restrict the number of positions at which non-default rules can appear, while the third class restricts the number of different rules but not the number of occurrences nor it imposes the presence of default rules. Some simple examples follow.

Example 1. Consider the $\nu\text{-CA}$ $H^{(1)} : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ defined as

$$\forall x \in A^{\mathbb{Z}}, \quad H^{(1)}(x)_i = \begin{cases} 1 & \text{if } i = 0 \\ 0 & \text{otherwise} \end{cases} \quad .$$

Remark that $H^{(1)}$ is a d $\nu\text{-CA}$ which cannot be a CA since it does not commute with σ . This trivially shows that the class of $\nu\text{-CA}$ is larger than the one of CA.

Example 2. Consider the $\nu\text{-CA}$ $H^{(2)} : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ defined as

$$\forall x \in A^{\mathbb{Z}}, \quad H^{(2)}(x)_i = \begin{cases} x_{i+1} & \text{if } i < 0 \\ x_0 & \text{if } i = 0 \\ x_{i-1} & \text{if } i > 0 \end{cases}$$

Remark that $H^{(2)}$ is a $p\nu\text{-CA}$ (with $p = 1$) but not a $d\nu\text{-CA}$.

Example 3. Consider the $\nu\text{-CA}$ $H^{(3)} : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ defined as

$$\forall x \in A^{\mathbb{Z}}, \quad H^{(3)}(x)_i = \begin{cases} 1 & \text{if } |i| \text{ is prime} \\ 0 & \text{otherwise.} \end{cases}$$

Remark that $H^{(3)}$ is a $r\nu\text{-CA}$ but not a $p\nu\text{-CA}$.

Example 4. Consider the $\nu\text{-CA}$ $H^{(4)} : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ defined as $\forall x \in A^{\mathbb{Z}}, H^{(4)}(x)_i = x_0$. Remark that $H^{(4)}$ is a $\nu\text{-CA}$ but not a $r\nu\text{-CA}$.

We give some relationships and properties involving the classes of $\nu\text{-CA}$ above introduced.

Proposition 3.1. $\mathcal{CA} \subsetneq d\nu\text{-CA} \subsetneq p\nu\text{-CA} \subsetneq r\nu\text{-CA} \subsetneq \nu\text{-CA}$, where $\mathcal{CA}$ is the set of all CA.

Proof. The inclusions $\subseteq$ easily follow from the definitions. For the strict inclusions refer to Examples 1 to 4. $\square$

Similarly to what happens in the context of CA one can prove the following.

Proposition 3.2. Any $r\nu\text{-CA}$ is topologically conjugated to a $r\nu\text{-CA}$ of radius 1.

Proof. Let H be a $r\nu\text{-CA}$ on the alphabet A . If H has radius $r = 1$ then the statement is trivially true. Otherwise, let $B = A^r$ and define $\phi : A^{\mathbb{Z}} \rightarrow B^{\mathbb{Z}}$ as $\forall i \in \mathbb{Z}, \phi(x)_i = x_{[ir, (i+1)r)}$. Then, it is not difficult to see that the $r\nu\text{-CA}$ $(B^{\mathbb{Z}}, H')$ of radius 1 defined as $\forall x \in A^{\mathbb{Z}}, \forall i \in \mathbb{Z}, H'(x)_i = h'_i(x_{i-1}, x_i, x_{i+1})$ is topologically conjugated to H via ϕ , where $\forall u, v, w \in B, \forall i \in \mathbb{Z}, \forall j \in \{0, \dots, r-1\}, (h'_i(u, v, w))_j = h_{ir+j}(u_{[j, r)} v w_{[0, j)})$. $\square$

Proposition 3.3. Any $p\nu\text{-CA}$ is topologically conjugated to a $p\nu\text{-CA}$ of radius and structural period 1.

Proof. The proof is similar to the one of Proposition 3.2. $\square$

Finally, the following result shows that every $r\nu\text{-CA}$ is a subsystem of a suitable CA.

Theorem 3.4. *Any $\nu\mathcal{CA}$ $H : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ of radius r is a subsystem of a CA, i.e., there exist a CA $F : B^{\mathbb{Z}} \rightarrow B^{\mathbb{Z}}$ on a suitable alphabet B and a continuous injection $\phi : A^{\mathbb{Z}} \rightarrow B^{\mathbb{Z}}$ such that $\phi \circ H = F \circ \phi$.*

Proof. Consider a $\nu\mathcal{CA}$ $H : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ of radius r . Remark that there are only $n = |A|^{|A|^{2r+1}}$ distinct functions $h_i : A^{2r+1} \rightarrow A$. Take a numbering $(f_j)_{1 \leq j \leq n}$ of these functions and let $B = A \times \{1, \dots, n\}$. Define the map $\phi : A^{\mathbb{Z}} \rightarrow B^{\mathbb{Z}}$ such that $\forall x \in A^{\mathbb{Z}}, \forall i \in \mathbb{Z}, \phi(x)_i = (x_i, k)$, where k is the integer for which $H(x)_i = f_k(x_{i-r}, \dots, x_{i+r})$. Clearly, ϕ is injective and continuous. Now, define a CA $F : B^{\mathbb{Z}} \rightarrow B^{\mathbb{Z}}$ using the local rule $f : B^{2r+1} \rightarrow B$ such that

$$f((x_{-r}, k_{-r}), \dots, (x_0, k_0), \dots, (x_r, k_r)) = (f_{k_0}(x_{-r}, \dots, x_r), k_0) .$$

It is not difficult to see that $\phi \circ H = F \circ \phi$. □

4. CA versus $\nu\mathcal{CA}$

In this section, we illustrate some differences in dynamical behavior between CA and $\nu\mathcal{CA}$. The following properties which are really specific for CA are lost in the larger class of $\nu\mathcal{CA}$.

- P1) *the set of ultimately periodic points is dense in $A^{\mathbb{Z}}$.*
- P2) *surjectivity $\Leftrightarrow$ injectivity on finite configurations.*
- P3) *surjectivity $\Leftrightarrow$ any configuration has a finite number of pre-images.*
- P4) *expansivity $\Rightarrow$ transitivity*
- P5) *expansivity $\Rightarrow$ surjectivity*
- P6) *injectivity $\Rightarrow$ surjectivity*

Some of the previous properties are not valid for the following $\nu\mathcal{CA}$.

Example 5. Let $A = \{0, 1\}$ and define the $\nu\mathcal{CA}$ $H^{(5)} : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ as

$$\forall x \in A^{\mathbb{Z}}, \forall i \in \mathbb{Z}, \quad H^{(5)}(x)_i = \begin{cases} x_i & \text{if } i = 0 \\ x_{i-1} & \text{otherwise} . \end{cases}$$

P1) is not valid for $H^{(5)}$.

Proof. Let $H = H^{(5)}$. For all naturals n , it holds that $H^n([01]_0) \subseteq [00^n 1]_0$. Since $[00^n 1]_0 \cap [00^m 1]_0 = \emptyset$ for any pair n, m of distinct naturals, no configuration belonging to $[01]_0$ is an ultimately periodic point. Therefore, the set of ultimately periodic points is not dense. □

P3) is not valid for $H^{(5)}$

Proof. We show that $H^{(5)}$ has no configuration with an infinite number of pre-images although it is not surjective. In particular, any configuration has either 0 or 2 pre-images.

First of all, $H^{(5)}$ is not surjective. Indeed, since $\forall x \in A^{\mathbb{Z}}, H^{(5)}(x)_0 = H^{(5)}(x)_1$, configurations in the set $B = \{x \in A^{\mathbb{Z}} : x_0 \neq x_1\}$ have no pre-image. Furthermore, any $x \in A^{\mathbb{Z}} \setminus B$ has y and z as unique pre-images, where y and z are configurations such that $\forall i \notin \{-1, 0\}, y_i = z_i = x_{i+1}, y_0 = z_0 = x_0, y_{-1} = 0; z_{-1} = 1$. $\square$

We stress that $H^{(5)}$ is not surjective, despite it is based on two local rules each of which generates a surjective CA (namely, the identity CA and the shift CA).

In order to explore other properties, we introduce an other ν -CA.

Example 6. Let $A = \{0, 1\}$ and define $H^{(6)} : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ by

$$\forall x \in A^{\mathbb{Z}}, \forall i \in \mathbb{Z}, \quad H^{(6)}(x)_i = \begin{cases} 0 & \text{if } i = 0 \\ x_{i-1} \oplus x_{i+1} & \text{otherwise} \end{cases},$$

where $\oplus$ is the xor operator.

P2) is not valid for $H^{(6)}$.

Proof. We prove that $H^{(6)}$ is injective on the 0-finite configurations but it is not surjective. It is evident that $H^{(6)}$ is not surjective. Let x, y be two finite configurations such that $H^{(6)}(x) = H^{(6)}(y)$. By contradiction, assume that $x_i \neq y_i$, for some $i \in \mathbb{Z}$. Without loss of generality, assume that $i \in \mathbb{N}$. Since $x_i \oplus x_{i+2} = H^{(6)}(x)_{i+1} = H^{(6)}(y)_{i+1} = y_i \oplus y_{i+2}$, it holds that $x_{i+2} \neq y_{i+2}$ and, by induction, $\forall j \in \mathbb{N}, x_{i+2j} \neq y_{i+2j}$. We conclude that $\forall j \in \mathbb{N}, x_{i+2j} = 1$ or $y_{i+2j} = 1$ contradicting the assumption that x and y are finite. $\square$

P4) and P5) are not valid for $H^{(6)}$.

Proof. Let $H = H^{(6)}$. H is not transitive since it is not surjective. We show that H is positively expansive. Let x and y be two distinct configurations and let $n = \min\{i \in \mathbb{N}, x_{[-i, i]} \neq y_{[-i, i]}\}$. If either $n = 0$ or $n = 1$, then $d(H^0(x), H^0(y)) \geq \frac{1}{2}$. Otherwise, it is clear that $H(x)_{n-1} = x_{n-2} \oplus x_n \neq y_{n-2} \oplus y_n = H(y)_{n-1}$ and $H(x)_{[0, n-2]} = H(y)_{[0, n-2]}$. Iterating the same reasoning one sees that $H^{n-1}(x)_1 \neq H^{n-1}(y)_1$. Hence $d(H^{n-1}(x), H^{n-1}(y)) \geq \frac{1}{2}$. Thus, H is positively expansive with expansivity constant $\frac{1}{2}$. $\square$

Consider now the ν -CA $H^{(2)}$ from Example 2.

P6) is not valid for $H^{(2)}$.

Proof. Concerning non-surjectivity, just remark that only configurations x such that $x_{-1} = x_0 = x_1$ have a pre-image. Let $x, y \in A^{\mathbb{Z}}$ with $H^{(2)}(x) = H^{(2)}(y)$. Then, we have $\forall i > 0, x_{i-1} = y_{i-1}$ and $\forall i < 0, x_{i+1} = y_{i+1}$. So $x = y$ and H is injective. $\square$

5. Basic Properties of $\nu\text{-CA}$ and Decidability

This section is centered on two fundamental properties, namely surjectivity and injectivity. Focusing on CA, both properties are strongly related to peculiar dynamical behaviors. Injectivity coincides with reversibility [16], while surjectivity is a necessary condition for almost all the widest known definitions of deterministic chaos (see [3], for instance).

In (1D) CA settings, the notion of De Bruijn graph is very handy to find fast decision algorithms for surjectivity, injectivity and openness [24]. Here, we extend this notion to work with $\nu\text{-CA}$ having period 1 and find decision algorithm for surjectivity. We stress that surjectivity is undecidable for two (or higher) dimensional $\nu\text{-CA}$, since surjectivity is undecidable for 2D CA [17].

Definition 5.1. Consider a $\nu\text{-CA}$ H of radius r and period $p = 1$ having f and g as right and left default rules. Let $k \in \mathbb{N}$ be the largest natural such that $h_k \neq f$ or $h_{-k} \neq g$. The *De Bruijn graph* of H is the triple $G = (V, E, \ell_G)$ where $V = A^{2r} \times \{-k, \dots, k+1\}$ and E is the set of pairs $((u, \alpha), (v, \beta)) \in V^2$ with label $\ell_G((u, \alpha), (v, \beta))$ in $A \times \{0, 1\}$ such that $u_{[1, 2r)} = v_{[0, 2r-1)}$ and one of the following conditions is verified

- a) $\alpha = \beta = -k$; in this case the label is $(g(u_0v), 0)$
- b) $\alpha + 1 = \beta$; in this case the label is $(h_\alpha(u_0v), 0)$
- c) $\alpha = \beta = k + 1$; in this case the label is $(f(u_0v), 1)$

By this graph, a configuration can be seen as a bi-infinite path on vertexes which passes once from a vertex whose second component is in $[-k+1, k]$ and infinite times through other vertices. The second component of vertices allows to single out the positions of local rules different from the default one. The image of a configuration is the sequence of first components of edge labels.

Lemma 5.1. *Surjectivity is decidable for $\nu\text{-CA}$ with structural period $p = 1$.*

Proof. Let H be a $\nu\text{-CA}$ with structural period $p = 1$ and let G be its De Bruijn graph. We prove that H is surjective iff G recognizes the language $(A \times \{0\})^*(A \times \{1\})^*$ when G is considered as the graph of a finite state automaton in which all states are both initial and final.

Let k be as in Definition 5.1 and denote by (w, s) any word of $(A \times \{0, 1\})^*$ with $w \in A^l$, $s \in \{0, 1\}^l$, for some $l \in \mathbb{N}$.

Assume that H is surjective and take $(w', s) \in (A \times \{0\})^n(A \times \{1\})^*$, for any $n \in \mathbb{N}$. Then, there exists $x \in A^{\mathbb{Z}}$ such that $H(x)_{[m, m+l)} = w'$, where $m = k + 1 - n$. Set $w = x_{[m-r, m+l+r)}$. Hence, the word (w', s) is the sequence of edge labels of the following vertex path on G :

$$(w_{[0, 2r)}, \alpha_0), \dots, (w_{[l, l+2r)}, \alpha_l)$$

where

$$\alpha_i = \begin{cases} -k & \text{if } m+i < -k \\ k+1 & \text{if } m+i > k \\ m+i & \text{otherwise} \end{cases}$$

For the opposite implication, assume that G recognizes $(A \times \{0\})^*(A \times \{1\})^*$. Take $y \in A^{\mathbb{Z}}$ and let $n > k$. Since G recognizes $(y_{[-n,n]}, 0^{n+k+1}1^{n-k})$, there exists $x \in A^{\mathbb{Z}}$ such that $H(x)_{[-n,n]} = y_{[-n,n]}$. Set $X_n = \{x \in A^{\mathbb{Z}}, x_{[n,n]} = y_{[-n,n]}\}$. For any $n \in \mathbb{N}$, X_n is non-empty and compact. Moreover, $X_{n+1} \subseteq X_n$. Therefore, $X = \bigcap_{n \in \mathbb{N}} X_n \neq \emptyset$ and $H(X) = \{y\}$. Hence, H is surjective. $\square$

In order to deal with injectivity, we introduce the following notion.

Definition 5.2. Consider a pv-CA of structural period 1 and let $G = (V, E, \ell_G)$ be its De Bruijn graph. The *product graph* P of H is a labeled graph $P = (V \times V, W, \ell_P)$ where $((u, \alpha), (v, \beta)), ((w, \gamma), (z, \delta)) \in W$ iff

$$\begin{cases} \alpha = \beta \text{ and } \gamma = \delta \\ ((u, \alpha), (w, \gamma)) \in E \text{ and } ((v, \beta), (z, \delta)) \in E \\ \ell_G((u, \alpha), (w, \gamma)) = \ell_G((v, \beta), (z, \delta)) \end{cases}$$

and $\ell_P : W \rightarrow A$ is defined as

$$\ell_P(((u, \alpha), (v, \beta)), ((w, \gamma), (z, \delta))) = \begin{cases} 0, & \text{if } u = v \text{ and } w = z \\ 1, & \text{otherwise.} \end{cases}$$

The *reduced product graph* D of P is the sub-graph of P made by all nodes and edges belonging to some bi-infinite path.

Let k be as in Definition 5.1. We can also consider D as the transition graph of a finite automaton with set of initial states $\{((u, -k), (v, -k)) : u, v \in A^{2^r}\}$ and set of final states $\{((u, k+1), (v, k+1)) : u, v \in A^{2^r}\}$. Denote by $L(D)$ the language recognized by this finite automaton.

Lemma 5.2. *Injectivity is decidable for pv-CA with structural period $p = 1$.*

Proof. Let H be a pv-CA with period $p = 1$ and let D its reduced product graph. We prove that H is injective if and only if $L(D) \subseteq 0^*$.

If $L(D) \not\subseteq 0^*$, there exists a word $w \in L(D)$ such that $w_i = 1$, for some i . By definition of D , this means that there are at least two distinct configurations which have the same image by H . Hence H is not injective.

If H is not injective, then there are two distinct configurations x and y such that $H(x) = H(y)$. Let $i \in \mathbb{Z}$ be such that $x_i \neq y_i$ and set $m = \max(|i|, k+1)$. For any $j \in \mathbb{Z}$, define $u^j = x_{[j-r, j+r]}$ et $v^j = y_{[j-r, j+r]}$. Then, the path on D

$$\begin{aligned} &(((u^{-m}, -k), (v^{-m}, -k)), \dots, ((u^{-k}, -k), (v^{-k}, -k)), \dots, ((u^0, 0), (v^0, 0)), \dots, \\ &((u^{k+1}, k+1), (v^{k+1}, k+1)), \dots, ((u^m, k+1), (v^m, k+1))) \end{aligned}$$

starts from an initial state, ends at a final state, and contains an edge labelled with 1. Hence, $L(D) \not\subseteq 0^*$. $\square$

Theorem 5.3. *Surjectivity and injectivity are decidable for $\nu\mathcal{CA}$.*

Proof. Let H be a $\nu\mathcal{CA}$ of radius r . If $p = 1$, by Lemma 5.1 (resp., Lemma 5.2) we can decide surjectivity (resp., injectivity) of H . Otherwise, without loss of generality, assume $p = r$. By Proposition 3.3, H is topologically conjugated to a $\nu\mathcal{CA}$ H' with structural period 1. By Lemma 5.1 (resp., Lemma 5.2) and since H is surjective (resp., injective) iff H' is as well, we can decide surjectivity (resp., injectivity) of H . $\square$

5.1. Injectivity and surjectivity: structural implications

We now study how informations (about surjectivity or injectivity) on the global rule H of a $\nu\mathcal{CA}$ with structural period 1 relate to properties of the composing local rules.

Proposition 5.4. *Let F and G be two $\mathcal{CA}$ of rules f and g , respectively. For any $\nu\mathcal{CA}$ H with period $p = 1$ having f and g as right and left default rules, it holds that*

1. H surjective $\Rightarrow F$ surjective and G surjective
2. H injective $\Rightarrow F$ surjective and G surjective
3. H injective on finite configurations $\Rightarrow F$ surjective and G surjective

Proof. Let k be the largest natural such that $h_k \neq f$ or $h_{-k} \neq g$. Without loss of generality, assume that F is not surjective.

1. There exists a block u which has no pre-image by f . Let y be any configuration belonging to $[u]_{k+1}$. By definition of H , there is no configuration $x \in A^{\mathbb{Z}}$ with $H(x) = y$.
2. By a theorem in [16], f admits a diamond, i.e., there exist $u, v, w \in A^+$ with $u \neq v$ of same length such that $f(uvw) = f(vvw)$. Build $x \in [uvw]_{k+1}$ and $y \in [vwv]_{k+1}$ such that $x_i = y_i$ for all i different from the cylinder positions. By definition of H , $H(x) = H(y)$.
3. the proof is similar to item 2.

$\square$

For $\text{dv}\mathcal{CA}$ a stronger result holds.

Proposition 5.5. *Let F be a $\mathcal{CA}$ of local rule f . For any $\text{dv}\mathcal{CA}$ H with default rule f , it holds that*

1. H injective $\Rightarrow F$ injective
2. H injective $\Rightarrow H$ surjective

Proof. Let k be the largest natural such that $h_k \neq f$ or $h_{-k} \neq f$. Fix a configuration $y \in A^{\mathbb{Z}}$ and for any $u \in A^{2k+1}$ let $y^u \in [u]_{-k}$ be the configuration such that $y_i^u = y_i$ for all $i \in \mathbb{Z}$, $|i| > k$. Define $Y = \{y^u : u \in A^{2k+1}\}$ and $X = F^{-1}(Y)$.

1. If H is injective then $|X| = |H(X)|$ and by Proposition 5.4 F is surjective. So $X \geq |A|^{2k+1}$. By definition of H , it holds that $H(X) \subseteq Y$. Hence,

$$|A|^{2k+1} \leq |X| = |H(X)| \leq |Y| = |A|^{2k+1}$$

which gives $|X| = |A|^{2k+1}$. Thus, F is injective.

2. if H is injective we also have $H(X) = Y$. Thus $y \in Y$ has a pre-image by H .

□

6. Dynamics

In order to study equicontinuity and almost equicontinuity, we introduce an intermediate class between $d\nu\text{-CA}$ and $r\nu\text{-CA}$.

Definition 6.1 (n -compatible $r\nu\text{-CA}$). A $r\nu\text{-CA}$ H is n -compatible with a local rule f if for any $k \in \mathbb{N}$, there exist two integers $k_1 > k$ and $k_2 < -k$ such that $\forall i \in [k_1, k_1 + n) \cup [k_2, k_2 + n)$, $h_i = f$.

In other words, a $\nu\text{-CA}$ is n -compatible with f if, arbitrarily far from the center of the lattice, there are intervals of length n in which the local rule f is applied.

The notion of blocking word and the related results cannot be directly restated in the context of $\nu\text{-CA}$ because some words are blocking just thanks to the uniformity of CA. To overcome this problem we introduce the following notion.

Definition 6.2 (Strongly blocking word). A word $u \in A^l$ is said to be *strongly s -blocking* ($0 < s \leq l$) for a CA F of local rule f if there exists an offset $d \in [0, l - s]$ such that for any $\nu\text{-CA}$ H with $\forall i \in [0, l)$, $h_i = f$ it holds that

$$\forall x, y \in [u]_0, \forall n \geq 0, \quad H^n(x)_{[d, d+s)} = H^n(y)_{[d, d+s)} \ .$$

Roughly speaking, a word is strongly blocking if it is blocking whatever be the perturbations involving the rules in its neighborhood. The following extends Proposition 5.12 in [20] to strongly r -blocking words.

Theorem 6.1. *Let F be a CA of local rule f and radius r . The following statements are equivalent:*

- (1) F is equicontinuous;

- (2) there exists $k > 0$ such that any word $u \in A^k$ is strongly r -blocking for F ;
(3) any $\nu\text{-CA}$ H of default rule f is ultimately periodic.

Proof. (1) $\Rightarrow$ (2). Suppose that F is equicontinuous. By [19, Th. 4], there exist $p > 0$ and $q \in \mathbb{N}$ such that $F^{q+p} = F^q$. As a consequence, we have that

$$\forall u \in A^*, |u| > 2(q+p)r \Rightarrow f^{p+q}(u) = f^q(u)_{[pr, |u| - (2q+p)r]} .$$

Let H be a $\nu\text{-CA}$ such that $h_j = f$ for each $j \in [0, (2p+2q+1)r)$. For any $x \in A^{\mathbb{Z}}$ and $i \in \mathbb{N}$, consider the following words:

$$\begin{aligned} s^{(i)} &= H^i(x)_{[0, qr)} \\ t^{(i)} &= H^i(x)_{[qr, (q+p)r)} \\ u^{(i)} &= H^i(x)_{[(q+p)r, (q+p+1)r)} \\ v^{(i)} &= H^i(x)_{[(q+p+1)r, (q+2p+1)r)} \\ w^{(i)} &= H^i(x)_{[(q+2p+1)r, (2q+2p+1)r)} . \end{aligned}$$

For all $i \in [0, q+p]$, $u^{(i)}$ is completely determined by $s^{(0)}t^{(0)}u^{(0)}v^{(0)}w^{(0)} = x_{[0, (2q+2p+1)r)}$ (see Figure 1). Moreover, for any natural i , we have

$$\begin{aligned} u^{(i+q+p)} &= f^{q+p}(s^{(i)}t^{(i)}u^{(i)}v^{(i)}w^{(i)}) \\ &= f^q(s^{(i)}t^{(i)}u^{(i)}v^{(i)}w^{(i)})_{[pr, (p+1)r)} \\ &= (t^{(i+q)}u^{(i+q)}v^{(i+q)})_{[pr, (p+1)r)} \\ &= u^{(i+q)} . \end{aligned}$$

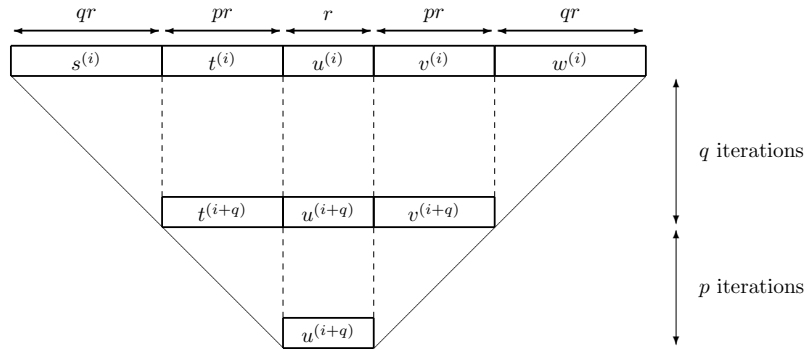


Figure 1: A strongly blocking word.

Summarizing, for all $i \in \mathbb{N}$, $u^{(i)}$ is determined by the word $x_{[0, (2q+2p+1)r)}$ which is then strongly r -blocking. Since x had been chosen arbitrarily, (2) is true.

(2) $\Rightarrow$ (3). Let H be a $\text{d}\nu\text{-CA}$ of default rule f and let n be such that $\forall i, |i| > n \Rightarrow h_i = f$. Consider an arbitrary configuration x and an integer i with $|i| > n + k$. By hypothesis, $x_{[i-k, i-1]}$ and $x_{[i+1, i+k]}$ are strongly blocking words and then the column $(H^j(x)_i)_{j \in \mathbb{N}}$ is blocked and completely determined by $x_{[i-k, i+k]}$.

For any $u \in A^{2k+1}$, consider the configuration $y = {}^\infty u {}^\infty \in [u]_{-k}$ (bi-infinite concatenation of u). There exist q_u and p_u such that $F^{p_u+q_u}(y) = F^{q_u}(y)$. Set

$$q = \max\{q_u : u \in A^{2k+1}\} \text{ and } p = \text{lcm}\{p_u : u \in A^{2k+1}\}$$

For any word $u \in A^{2k+1}$, the column blocked by u (when considered as a portion of a configuration in suitable positions) admits q and p as pre-period and period, respectively.

We have proved that for all integer i such that $|i| > n + k$, the sequence $(H^j(x)_i)_{j \in \mathbb{N}}$ is ultimately periodic with pre-period q and period p .

On the other hand, the sequence $(H^j(x)_{[-n-k, n+k]})_{j \in \mathbb{N}}$ is completely determined by $w = x_{[-m, m]}$ where $m = n + 2k + R$ and $R = \max\{r_i : i \in \mathbb{Z}\}$ (see Figure 2).

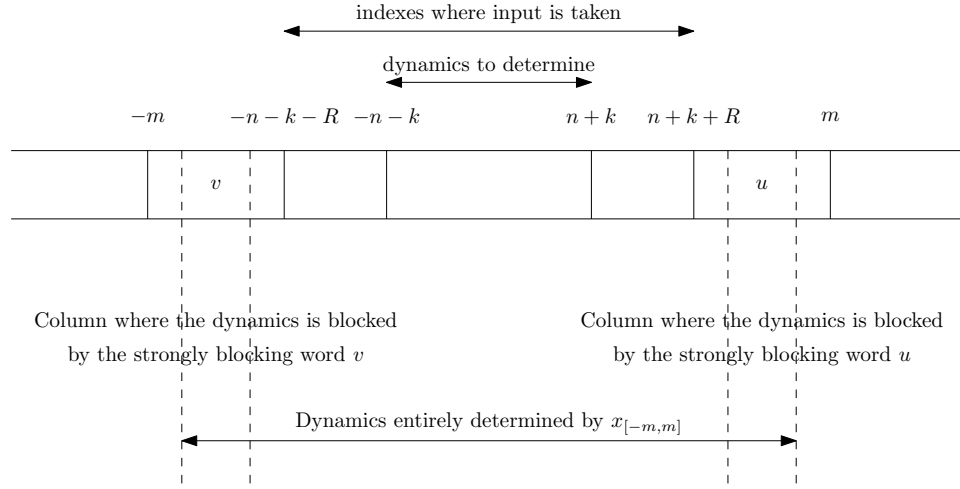


Figure 2: Dynamics of a $\text{d}\nu\text{-CA}$ in presence of strongly blocking words.

Moreover, there exist $\alpha_w > 0$ and $\beta_w \geq q$ such that

$$H^{\beta_w}(x)_{[-n-k, n+k]} = H^{\beta_w + p\alpha_w}(x)_{[-n-k, n+k]}$$

leading to $H^{\beta_w}(x) = H^{\beta_w + p\alpha_w}(x)$. Set now

$$q' = \max\{\beta_{w'} : w' \in A^{2m+1}\} \text{ and } p' = \text{lcm}\{p\alpha_{w'} : w' \in A^{2m+1}\}.$$

Hence, $\forall x \in A^{\mathbb{Z}}, H^{q'+p'}(x) = H^{q'}(x)$ and H is ultimately periodic.

(3) $\Rightarrow$ (1) Since F is a dv-CA of default rule f , F is ultimately periodic. By [19, Th. 4], F is equicontinuous. $\square$

Theorem 6.2. *Let F be a CA with local rule f admitting a strongly r -blocking word u . Let H be a rv-CA of radius r . If H is $|u|$ -compatible with f then H is almost equicontinuous.*

Proof. Let p and n be the offset and the length of u , respectively. For any $k \in \mathbb{N}$, consider the set $T_{u,k}$ of configurations $x \in A^{\mathbb{Z}}$ having the following property $\mathcal{P}$: there exist $l > k$ and $m < -k$ such that $x_{[l, l+n)} = x_{[m, m+n)} = u$ and $\forall i \in [l, l+n) \cup [m, m+n)$ $h_i = f$. Remark that $T_{u,k}$ is open, being a union of cylinders. Clearly, each $T_{u,k}$ is dense, thus the set $T_u = \bigcap_{k \geq n} T_{u,k}$ is residual. We claim that any configuration in T_u is an equicontinuity point. Indeed, consider an arbitrary configuration $x \in T_u$ and a real $\epsilon > 0$. Let $k \in \mathbb{N}$ be such that $2^{-k} \leq \epsilon$. Then, there exist $k_1 > k$ and $k_2 < -k - n$ satisfying $\mathcal{P}$ (see Figure 3).

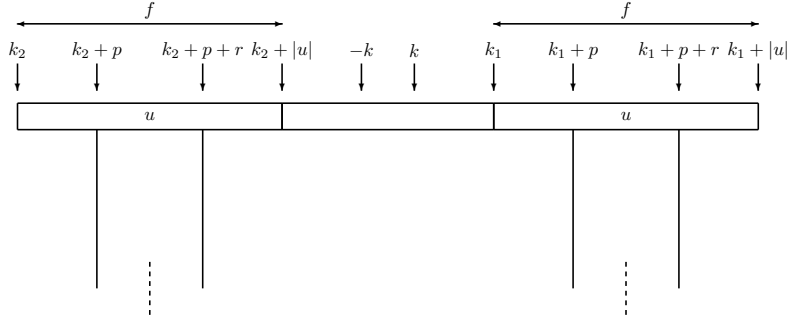


Figure 3: An equicontinuity point (see Theorem 6.2).

Fix $\delta = \min\{2^{-(k_1+n)}, 2^{-k_2}\}$ and let $y \in A^{\mathbb{Z}}$ be such that $d(x, y) < \delta$. Then $y_{[k_2, k_1+|u|)} = x_{[k_2, k_1+|u|)}$. Since u is r -blocking, $\forall t \in \mathbb{N}$, $H^t(x)$ and $H^t(y)$ are equal inside the intervals $[k_1 + p, k_1 + p + r]$ and $[k_2 + p, k_2 + p + r]$, then $d(H^t(x), H^t(y)) < \epsilon$. $\square$

In a similar manner one can prove the following.

Theorem 6.3. *Let F be an equicontinuous CA of local rule f . Let $k \in \mathbb{N}$ be as in item (2) of Theorem 6.1. Any rv-CA k -compatible with f is equicontinuous.*

6.1. Perturbing almost equicontinuous CA

In the sequel, we show how the loss of uniformity may lead to a dramatic change in the dynamical behavior of the automata network.

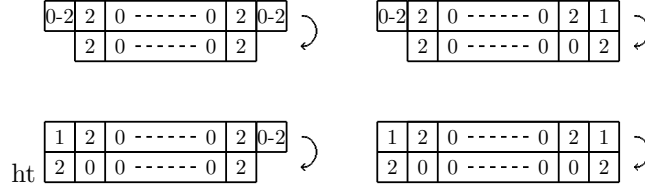


Figure 4: Evolution of words $20^i 2$ according to $F^{(9)}$.

Example 7 (An almost equicontinuous CA). Let $A = \{0, 1, 2\}$ and define a local rule $f : A^3 \rightarrow A$ as follows: $\forall x, y \in A$,

$$\begin{aligned} f(x, 0, y) &= \begin{cases} 1 & \text{if } x = 1 \text{ or } y = 1 \\ 0 & \text{otherwise} \end{cases} \\ f(x, 1, y) &= \begin{cases} 2 & \text{if } x = 2 \text{ or } y = 2 \\ 1 & \text{otherwise} \end{cases} \\ f(x, 2, y) &= \begin{cases} 0 & \text{if } x = 1 \text{ or } y = 1 \\ 2 & \text{otherwise} \end{cases} . \end{aligned}$$

We show that the CA defined in Example 7 is almost equicontinuous.

Proof. Just remark that the number of 0s inside the word $20^i 2$ is non-decreasing. Thus 202 is a 1-blocking word (see Figure 4). $\square$

The following example defines a ν -CA which is sensitive to the initial conditions although its default rule give rise to an almost equicontinuous CA.

Example 8 (A sensitive ν -CA with an almost equicontinuous default rule). Consider the $d\nu$ -CA $H^{(8)} : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$ defined as follows

$$\forall x \in A^{\mathbb{Z}}, \forall i \in \mathbb{Z}, \quad H^{(8)}(x)_i = \begin{cases} 1 & \text{if } i = 0 \\ f(x_{i-1}, x_i, x_{i+1}) & \text{otherwise} \end{cases} ,$$

where f and A are as in Example 7.

Remark that positive and negative cells do not interact each other under the action of $H^{(8)}$. Therefore, in order to study the behavior of $H^{(8)}$, it is sufficient to consider the action of $H^{(8)}$ on $A^{\mathbb{N}}$. In the sequel, we will simply note by H the map $H^{(8)}$.

In order to prove that $H^{(8)}$ is sensitive, we need some technical Lemmata.

Lemma 6.4. *For any $u \in A^*$, consider the sequence $(u^{(n)})_{n \in \mathbb{N}}$ defined as:*

$$\begin{cases} u^{(n+1)} = f(1u^{(n)}0) & \forall n \in \mathbb{N} \\ u^{(0)} = u2 \end{cases}$$

Then,

$$\exists m, \forall n \geq m, \quad u^{(n)} = 1^{|u|+1}$$

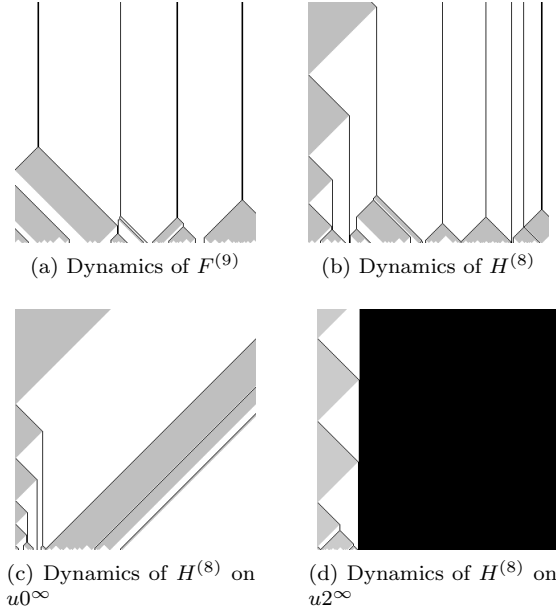


Figure 5: Space-time diagrams for $F^{(9)}$ and $H^{(8)}$.

Proof. We proceed in 4 steps.

1. First of all, we are going to show that there exists $n_0 \in \mathbb{N}$ such that $u^{(n_0)}$ does not contain any 1. In particular, we prove that there exists $n_0 \in \mathbb{N}$ such that $\forall n \leq n_0$ the integer

$$i^{(n)} = \min \left\{ i \leq |u| : u_i^{(n)} = 2 \text{ and } u_{[i+1, |u|]}^{(n)} \in \{0, 2\}^* \right\}$$

is well defined with the property $P(n) = (\forall k \in [0, n], i^{(k+1)} \leq i^{(k)})$ and $i^{(n_0)} = 0$. By definition, $i^{(0)}$ is well defined and the property $P(0)$ is true. Suppose now that, for some $n \in \mathbb{N}$, $i^{(n)}$ is well defined and the property $P(n)$ is true. We deal with the following cases.

- (a) If $i^{(n)} = 0$, then we set $n_0 = n$ and we are done.
- (b) If $i^{(n)} \neq 0$ and $u^{(n)}$ does not contain any 1, then we can write $u^{(n)} = 0^{i^{(n)}} 2w$ with $w \in \{0, 2\}^*$ and we have $u^{(n+1)} = 10^{i^{(n)}-1} 2w$. So, $i^{(n+1)} = i^{(n)}$ is well defined with $P(n+1)$ true and, having the element $n+1$ as a starting point, we fall in the next case.
- (c) If $i^{(n)} \neq 0$ and $u^{(n)}$ contains at least one 1, let $h \in [0, i^{(n)})$ be the greatest position in which it appears and set $s = i^{(n)} - h - 1$. We can write $u^{(n)} = v^{(n)} 10^s 2w$, for some $v^{(n)} \in A^h$ and $w \in \{0, 2\}^*$. Then, for each $j \in [1, s]$, we obtain $u^{(n+j)} = v^{(n+j)} 10^{s-j} 2w$, for

some $v^{(n+j)} \in A^{h+j}$. So, for each $j \in [1, s]$, $i^{(n+j)} = i^{(n)}$ is well defined with $P(n+j)$ true. Furthermore, it holds that $u^{(n+s+1)} = v^{(n+s+1)}20w$, for some $v^{(n+s+1)} \in A^{h+s+1}$, and $i^{(n+s+1)} = i^{(n+s)} - 1$ is well defined with $P(n+s+1)$ true, and we can reconsider the three cases having the element $n+s+1$ as starting point.

By considering iteratively the three cases, we are sure to reach a natural n_0 such that $i^{(n_0)} = 0$ since whenever we fall in the third case the value of $i^{(n)}$ decreases.

2. Proceeding by induction, we now show that

$$\forall n \geq n_0, \exists k \in \mathbb{N}, \exists v \in \{0, 2\}^*, \quad \text{s.t.} \quad u^{(n)} = 1^k v.$$

Clearly, this is true for $n = n_0$ with $k = 0$. Assume now that the statement is true for some $n \geq n_0$ and consider the following cases.

- If $v = \epsilon$, then $u^{(n+1)} = u^{(n)} = 1^k v$
- If $v_0 = 0$, then $u^{(n+1)} = 1^{k+1} v_{[1, |v|-1]}$
- If $v_0 = 2$ and $k \neq 0$, then $u^{(n+1)} = 1^{k-1} 20 v_{[1, |v|-1]}$
- If $v_0 = 2$ and $k = 0$, then $u^{(n+1)} = 0 v_{[1, |v|-1]}$

In all the cases, the statement is true for $n+1$.

As a consequence, we also have that the number $|u^{(n)}|_2$ of 2 inside $u^{(n)}$ is a (non strictly) decreasing sequence:

$$\forall n \geq n_0, \quad |u^{(n+1)}|_2 \leq |u^{(n)}|_2$$

Indeed, $u^{(n)}$ does not contain the block 121, which, transforming itself into 202, is the unique one able to increase the number of 2.

3. We now prove that there exists $n_1 \geq n_0$, such that $u^{(n_1)}$ no longer contains any 2, and then $u^{(n_1)} = 1 \cdots 10 \cdots 0$. This is assured by showing that

$$\forall n \geq n_0, |u^{(n)}|_2 > 0 \Rightarrow \exists s \in \mathbb{N}, |u^{(n+s)}|_2 < |u^{(n)}|_2$$

Let $n \geq n_0$ such that $|u^{(n)}|_2 > 0$. Since $u^{(n)} = 1^k v$ for some $k \in \mathbb{N}, v \in \{0, 2\}^*$, we can write $u^{(n)} = 1^k 0^h 2w$ for some $h \in \mathbb{N}, w \in \{0, 2\}^*$. Thus, we have $u^{(n+h)} = 1^{k+h} 2w$ and $u^{(n+h+i)} = 1^{k+h-i} 20^i w$ for each $i \in [1, h+k]$. So, $u^{(n+2h+k)} = 20^{h+k} w$ and, setting $s = 2h+k+1$, we obtain $u^{(n+s)} = 0^{s-h} w$, assuring that $|u^{(n+s)}|_2 < |u^{(n)}|_2$.

4. Since $u^{(n_1)} = 1^k 0^h$ for some $h, k \in \mathbb{N}$, it is easy to observe that $u^{(n_1+i)} = 1^{k+i} 0^{h-i}$, $i \in [1, h]$. In particular, setting $m = n_1 + h$, we obtain $u^{(m)} = 1^{|u|+1}$ and $\forall n \geq m, u^{(n)} = u^{(m)}$. This concludes the proof. $\square$

Lemma 6.5. *Consider the rewriting system on $A^* \times \{0, 1\}$ defined by the following rules:*

1. $(u0, 0) \xrightarrow{1} (u, 0)$
2. $(u1, 0) \xrightarrow{2} (u, 1)$
3. $(u2, 0) \xrightarrow{3} (f(1u20), 0)$
4. $(u0, 1) \xrightarrow{4} (f(1u0), 1)$
5. $(u1, 1) \xrightarrow{5} (u, 1)$
6. $(u2, 1) \xrightarrow{6} (f(1u2), 0)$
7. $(\epsilon, x) \xrightarrow{7} (\epsilon, 1)$

Starting from any $(u, x) \in A^* \times \{0, 1\}$, after a certain number m of rule applications, the system ultimately falls into $(\epsilon, 1)$.

Proof. This system is non ambiguous and then, for any $(u, x) \in A^* \times \{0, 1\}$, it (well) defines the sequence $(u^{(n)}, x^{(n)})_{n \in \mathbb{N}}$ such that

$$\begin{cases} (u^{(n)}, x^{(n)}) \rightarrow (u^{(n+1)}, x^{(n+1)}) & \forall n \in \mathbb{N} \\ (u^{(0)}, x^{(0)}) = (u, x) \end{cases}$$

where $\rightarrow$ is the unique possible application of a system rule. Consider the sequence $(l^{(n)})_{n \in \mathbb{N}} = (|u^{(n)}|)_{n \in \mathbb{N}}$. By definition, it is a (non strictly) decreasing sequence and then it converges to some $l \in \mathbb{N}$, or, equivalently, there exists $m \in \mathbb{N}$ such that $\forall n \geq m$, $l^{(n)} = l$. We show that $l = 0$ and this also prove the thesis. For a sake of argument, suppose that $l > 0$. Thus, there exists $k \in \mathbb{N}$ such that $\forall n \geq k$, $(u^{(n)}, x^{(n)}) \xrightarrow{3} (u^{(n+1)}, x^{(n+1)})$ since, except rule 7, rule 3 is the only one leaving $l^{(n)}$ unchanged. Furthermore, the sequence $(u^{(n+k)})_{n \in \mathbb{N}}$ verifies the hypothesis of Lemma 6.4 and so it is ultimately equal to 1^l , that is contrary to the fact that $\forall n \geq k$, $u_l^{(n)} = 2$, since rule 3 is always applied. $\square$

Lemma 6.6. *Let $\mathcal{F} = \{01, 12, 20, 22\}$. For any $x \in A^{\mathbb{N}}$ and any $i \in \mathbb{N}$, if no element of $\mathcal{F}$ appears inside $x_{[i, \infty)}$, then no element of $\mathcal{F}$ appears inside $H^{(8)}(x)_{[i+1, \infty)}$.*

Proof. The f -pre-images of words in $\mathcal{F}$ are :

- $f^{-1}(01) = \{0001, 1201, 2001\}$
- $f^{-1}(12) = \{0012, 0112, 1012, 1020, 1022, 1112, 2012\}$
- $f^{-1}(20) = \{0120, 0121, 0122, 0200, 0202, 0221, 1120, 1121, 1122, 2120, 2121, 2122, 2200, 2202, 2221\}$
- $f^{-1}(22) = \{0220, 0222, 2112, 2220, 2222\}$

So, if there exists $w \in \mathcal{F}$ appearing in $H^{(8)}(x)_{[i+1, \infty)}$, necessarily a word $u \in \mathcal{F}$ is inside $x_{[i, \infty)}$. $\square$

Lemma 6.7. *For any $u \in A^*$, there exists $n_0 \in \mathbb{N}$ s.t. $\forall n > n_0, H^n(u0^\infty)_1 = 1$.*

Proof. Consider the sequences $(u^{(n)}, x^{(n)})_{n \in \mathbb{N}}$ and $(l^{(n)})_{n \in \mathbb{N}}$ from Lemma 6.5 in which $(u^{(0)}, x^{(0)}) = (u_{[1, |u|-1]}, 0)$. Define the sequences $(k^{(n)})_{n \in \mathbb{N}}$ and $(y^{(n)})_{n \in \mathbb{N}}$ as follows:

$$\begin{cases} k^{(n+1)} = k^{(n)} + 1 & \forall n \in \mathbb{N} \text{ s.t. } (u^{(n)}, x^{(n)}) \xrightarrow{a} (u^{(n+1)}, x^{(n+1)}) \quad a = 3, 4, 6, 7 \\ k^{(n+1)} = k^{(n)} & \forall n \in \mathbb{N} \text{ s.t. } (u^{(n)}, x^{(n)}) \xrightarrow{a} (u^{(n+1)}, x^{(n+1)}) \quad a = 1, 2, 5 \\ k^{(0)} = 0 \end{cases},$$

and, $\forall n \in \mathbb{N}$, $y^n = H^{k^{(n)}}(u0^\infty)$, respectively. First of all, we are going to prove that the property

$$L(n) : \left(y_{[1, l^{(n)}+1]}^n = u^{(n)} x^{(n)} \right),$$

linking the dynamics of H with the one induced by the rewriting system, and the property

$$M(n) : \left(\forall w \in \mathcal{F}, \quad w \text{ does not appears inside } y_{[l^{(n)}+1, \infty)}^n \right),$$

are true for all $n \in \mathbb{N}$. We proceed by induction.

The properties $L(n)$ and $M(n)$ are clearly true for $n = 0$. Suppose now that they are valid for some $n \in \mathbb{N}$ and let the rewriting system evolve on $(u^{(n)}, x^{(n)})$ according to the rule a , for some $a = 1, \dots, 7$.

If $a \in \{1, 2, 5\}$, then $k^{(n+1)} = k^{(n)}$ and $l^{(n+1)} = l^{(n)} - 1$. Hence, $L(n+1)$ is true. Moreover, when the restriction passes from $[l^{(n)}+1, \infty)$ to $[l^{(n+1)}+1, \infty)$, the additional word inside the configuration $y^{n+1} = y^n$ is either 00, or 10, or 11, depending on the value of a . So, $M(n+1)$ is also true.

If $a = 7$, then $k^{(n+1)} = k^{(n)} + 1$ and $l^{(n+1)} = l^{(n)} = 0$. Since $L(n)$ and $M(n)$ are true, $y_{[0, 2]}^n \in \{100, 102, 110, 111\}$, and then $x^{(n+1)} = 1 = y_1^{n+1}$. So, $L(n+1)$ is valid. Moreover, no element $\mathcal{F}$ appears inside $y_{[1, \infty)}^n$ and neither inside $y_{[0, \infty)}^n$, and, hence, by Lemma 6.6, neither inside $y_{[1, \infty)}^{n+1}$. Thus, $M(n+1)$ is true.

If $a \in \{3, 4, 6\}$, then $k^{(n+1)} = k^{(n)} + 1$. By the fact that $L(n)$ is true, we have $u^{(n+1)} = f(y_{[0, l^{(n+1)}+1]}^n)$, or, equivalently, $u^{(n+1)} = y_{[1, l^{(n+1)}]}^{n+1}$. Since $M(n)$ is true, by Lemma 6.6, $y_{[l^{(n)}+2, \infty)}^{n+1}$ does not contain any element of $\mathcal{F}$. It remains to prove that $x^{(n+1)} = y_{l^{(n+1)}+1}^{n+1}$, and there is no word of $\mathcal{F}$ inside $y_{[l^{(n+1)}+1, l^{(n)}+2]}^{n+1}$.

If $a = 3$, $l^{(n+1)} = l^{(n)}$ and we have $y_{l^{(n+1)}+1}^{n+1} = f(20a)$ where necessarily $a \neq 1$ since $M(n)$ is true and $01 \in \mathcal{F}$. Hence, $y_{l^{(n+1)}+1}^{n+1} = 0 = x^{(n+1)}$ and $L(n+1)$ is true. For a sake of argument, assume that the word $w = y_{[l^{(n+1)}+1, l^{(n)}+2]}^{n+1} \in \mathcal{F}$. Necessarily, we obtain $w = 01$ and so, by definition of f , $y_{[l^{(n)}+2, l^{(n)}+3]}^n = 01$, that is a contradiction. Hence, $M(n+1)$ is true.

If $a = 4$ (resp., $a = 6$), $l^{(n+1)} = l^{(n)} - 1$ and we have $y_{l^{(n+1)}+1}^{n+1} = f(a01)$ (resp., $f(a21)$). So, $y_{l^{(n+1)}+1}^{n+1} = 0 = x^{(n+1)}$ and $L(n+1)$ is true. Since $M(n)$ is true, $y_{[l^{(n)}, l^{(n)}+3]}^n = 01bc$ (resp., $21bc$) with $bc \in \{00, 02, 10, 11\}$. By definition of f , it follows that $y_{[l^{(n+1)}, l^{(n+1)}+2]}^{n+1} = 111$ and then $M(n+1)$ is true.

Summarizing, we have proved that $L(n)$ and $M(n)$ are true for all $n \in \mathbb{N}$ and, in particular,

$$\forall n \in \mathbb{N}, \quad H^{k^{(n)}}(u0^\infty)_{[1, l^{(n)}+1]} = u^{(n)}x^{(n)}$$

Now, let m be the integer from Lemma 6.5. Recall that $\forall n \geq m$, $l^{(n)} = 0$ and $(u^{(n)}, x^{(n)}) = (\epsilon, 1)$. Thus, setting $n_0 = k^{(m)}$, we obtain

$$\forall n \geq n_0, \quad H^n(u0^\infty)_1 = 1.$$

□

Lemma 6.8. *For any $u \in A^*$ and any $n \in \mathbb{N}$, there exists $m > n$ such that $H^m(u2^\infty)_1 = 2$.*

Proof. For any $u \in A^*$ and any $n \in \mathbb{N}$, define the configuration $z^n = H^n(u2^\infty)$ and the integers $a^{(n)} = \max\{i \in \mathbb{N} : z_i^n = 1\}$ and $b^{(n)} = \min\{i \in \mathbb{N} : z_i^n = 2 \wedge \forall j > i, z_j^n \neq 1\}$. Remark that $a^{(n)}$ and $b^{(n)}$ are well defined and $\forall n \in \mathbb{N}$, $a^{(n)} < b^{(n)}$. We want to prove that $\forall n \in \mathbb{N}$, $(b^{(n)} > 1 \Rightarrow \exists k > n, b^{(k)} < b^{(n)})$. Let $n \in \mathbb{N}$ be such that $b^{(n)} > 1$. Since $z_{[a^{(n)}, b^{(n)}]}^n = 10^{b^{(n)}-a^{(n)}-1}2$ and $\forall i > b^{(n)}$, $z_i^n \neq 1$, by definition of f , it follows that $a^{(n+i)} = a^{(n)} + i$, $b^{(n+i)} = b^{(n)}$, for $i \in [0, b^{(n)} - a^{(n)} - 1]$. Hence, $b^{(n+b^{(n)}-a^{(n)})} = b^{(n)} - 1 < b^{(n)}$. □

We conclude stating that $H^{(8)}$ is sensitive.

Proof. It is a direct consequence of Lemmata 6.7 and 6.8. □

The following example shows that default rules individually defining almost equicontinuous CA can also constitute $\nu\text{-CA}$ that have a completely different behavior from the one in Example 8.

Example 9 (An equicontinuous $\nu\text{-CA}$ made by almost equicontinuous CA).

Let $A = \{0, 1, 2\}$ and define the local rule $f : A^3 \rightarrow A$ as: $\forall x, y, z \in A$, $f(x, y, z) = 2$ if $x = 2$ or $y = 2$ or $z = 2$, z otherwise. The CA F of local rule f is almost equicontinuous since 2 is a blocking word. The restriction of F to $\{0, 1\}^\mathbb{Z}$ gives the shift map which is sensitive. Thus F is not equicontinuous. Define now the following $d\nu\text{-CA}$ $H^{(9)}$:

$$\forall x \in A^\mathbb{Z}, \forall i \in \mathbb{Z}, \quad H^{(9)}(x)_i = \begin{cases} 2 & \text{if } i = 0 \\ f(x_{i-1}, x_i, x_{i+1}) & \text{otherwise} \end{cases}.$$

We now prove that $H^{(9)}$ is equicontinuous.

Proof. Let $n \in \mathbb{N}, x, y \in A^\mathbb{Z}$ be such that $x_{[-2n, 2n]} = y_{[-2n, 2n]}$. Since H is of radius 1, $\forall k \leq n$, $H^k(x)_{[-n, n]} = H^k(y)_{[-n, n]}$ and $\forall k > n$, $H^k(x)_{[-n, n]} = 2^{2n+1} = H^k(y)_{[-n, n]}$. So, H is equicontinuous. □

7. Conclusions

This paper have introduced $\nu\text{-CA}$, an extension of CA model obtained by relaxing the uniformity property (i.e., the fact that the same local rule is applied to all sites of the CA lattice). The study of how this change affects the dynamics of the systems has just started. We proved several results concerning basic set properties like injectivity and surjectivity. Moreover, we studied how informations about the $\nu\text{-CA}$ can determine properties on the underlying CA or vice-versa.

The study of $\nu\text{-CA}$ can be continued along several different directions. Of course, it would be interesting to progress in the analysis of the dynamical behavior. In particular we believe it would be worthwhile to study how information moves along the space-time diagrams and how the density of changes affects the entropy of the system.

It is well-known that CA cannot be used a random generator and, in general, they are poor (but fast) pseudo-random generators. Can $\nu\text{-CA}$ be a better tool in this context?

For some very recent results on the language complexity of local rule distributions see [9, 10].

8. Acknowledgements

This work has been supported by The French National Research Agency project EMC (ANR-09-BLAN-0164), by the “Coopération scientifique internationale Région Provence-Alpes-Côte d’Azur” Project “Automates Cellulaires, dynamique symbolique et décidabilité”, and by the PRIN/MIUR project “Automata and Formal Languages: Mathematical and Applicative Aspects”.

The authors warmly thank the anonymous referee for his valuable comments that helped to improve the paper.

References

- [1] L. Acerbi, A. Dennunzio, and E. Formenti. Conservation of some dynamical properties for operations on cellular automata. *Theoretical Computer Science*, 410:3685–3693, 2009.
- [2] G. Cattaneo, A. Dennunzio, E. Formenti, and J. Provillard. Non-uniform cellular automata. In Adrian Horia Dediu, Armand-Mihai Ionescu, and Carlos Martín-Vide, editors, *LATA*, volume 5457 of *Lecture Notes in Computer Science*, pages 302–313. Springer, 2009.
- [3] G. Cattaneo, E. Formenti, L. Margara, and G. Mauri. On the dynamical behavior of chaotic cellular automata. *Theoretical Computer Science*, 217:31–51, 1999.

- [4] J. Cervelle, A. Dennunzio, and E. Formenti. Chaotic behavior of cellular automata. In B. Meyers, editor, *Mathematical basis of cellular automata*, volume 1 of *Encyclopedia of Complexity and System Science*, pages 978–989. Springer Verlag, 2009.
- [5] P. Chaudhuri, D. Chowdhury, S. Nandi, and S. Chattopadhyay. *Additive Cellular Automata Theory and Applications*, volume 1. IEEE Press, 1997.
- [6] A. Dennunzio, P. Di Lena, E. Formenti, and L. Margara. On the directional dynamics of additive cellular automata. *Theoretical Computer Science*, 410:4823–4833, 2009.
- [7] A. Dennunzio and E. Formenti. Decidable properties of 2d cellular automata. In *Twelfth Conference on Developments in Language Theory (DLT 2008)*, volume 5257 of *Lecture Notes in Computer Science*, pages 264–275. Springer-Verlag, 2008.
- [8] A. Dennunzio, E. Formenti, and P. Kůrka. Cellular automata dynamical systems. In G. Rozenberg et al., editor, *Handbook of Natural Computing: Theory, Experiments, and Applications*, pages 25–75. Springer, 2010.
- [9] A. Dennunzio, E. Formenti, and J. Provillard. Computational complexity of rule distributions of non-uniform cellular automata. In *LATA*, volume 7183 of *Lecture Notes in Computer Science*, pages 204–215. Springer, 2012.
- [10] A. Dennunzio, E. Formenti, and J. Provillard. Local rule distributions, language complexity and non-uniform cellular automata. *Theoretical Computer Science*, 2012. Accepted (under revision).
- [11] A. Dennunzio, E. Formenti, and M. Weiss. Multidimensional cellular automata: closing property, quasi-expansivity, and (un)decidability issues. Preprint, CoRR, abs/0906.0857, 2009.
- [12] A. Dennunzio, E. Formenti, and M. Weiss. 2d cellular automata: dynamics and undecidability. In *Proceedings of 6th Conference on Computability in Europe (CiE 2010)*, 2010.
- [13] A. Dennunzio, B. Masson, and P. Guillon. Sand automata as cellular automata. *Theoretical Computer Science*, 410:3962–3974, 2009.
- [14] P. Di Lena and L. Margara. On the undecidability of the limit behavior of cellular automata. *Theoretical Computer Science*, 411:1075–1084, 2010.
- [15] E. Formenti and P. Kůrka. Dynamics of cellular automata in non-compact spaces. In B. Meyers, editor, *Mathematical basis of cellular automata*, Encyclopedia of Complexity and System Science. Springer Verlag, 2008.
- [16] G. A. Hedlund. Endomorphisms and automorphisms of the shift dynamical system. *Mathematical System Theory*, 3:320–375, 1969.

- [17] J. Kari. Reversibility and surjectivity problems of cellular automata. *Journal of Computer and System Sciences*, 48:149–182, 1994.
- [18] J. Kari. Tiling problem and undecidability in cellular automata. In B. Meyers, editor, *Mathematical basis of cellular automata*, Encyclopedia of Complexity and System Science. Springer Verlag, 2008.
- [19] P. Kůrka. Languages, equicontinuity and attractors in cellular automata. *Ergodic Theory & Dynamical Systems*, 17:417–433, 1997.
- [20] P. Kůrka. *Topological and Symbolic Dynamics*. Volume 11 of Cours Spécialisés. Société Mathématique de France, 2004.
- [21] P. Kůrka. Topological dynamics of one-dimensional cellular automata. In B. Meyers, editor, *Mathematical basis of cellular automata*, Encyclopedia of Complexity and System Science. Springer Verlag, 2008.
- [22] D. Bouchev L. B. Kier and G. A. Buck. Modeling biochemical networks: A cellular-automata approach. *Chemistry & Biodiversity*, 2:233–243, 2005.
- [23] T. Nakano, T. Morie, M. Nagata, and A. Iwata. A cellular automaton-type image extraction algorithm and its implementation using a fpga. In *IEEE Asia-Pacific Conference on Circuits and Systems*, volume 2, pages 197–200. IEEE Press, 2002.
- [24] K. Sutner. De Bruijn graphs and linear cellular automata. *Complex Systems*, 5:19–30, 1991.