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# Approximate Capacity of the Gaussian Interference Channel with Noisy Channel-Output Feedback

Victor Quintero, Samir M. Perlaza, Iñaki Esnaola and Jean-Marie Gorce

Abstract—In this paper, an achievability region and a converse region for the two-user Gaussian interference channel with noisy channel-output feedback (G-IC-NOF) are presented. The achievability region is obtained using a random coding argument and three well-known techniques: rate splitting, superposition coding and backward decoding. The converse region is obtained using some of the existing perfect-output feedback outer-bounds as well as a set of new outer-bounds that are obtained by using genie-aided models of the original G-IC-NOF. Finally, it is shown that the achievability region and the converse region approximate the capacity region of the G-IC-NOF to within a constant gap in bits per channel use.

Index Terms—Capacity, Interference Channel, Noisy Channel-Output Feedback.

#### I. NOTATION

Throughout this paper,  $(\cdot)^+$  denotes the positive part operator, i.e.,  $(\cdot)^+ = \max(\cdot,0)$  and  $\mathbb{E}_X[\cdot]$  denotes the expectation with respect to the distribution of the random variable X. The logarithm function  $\log$  is assumed to be base 2.

#### II. SYSTEM MODEL

Consider the two-user G-IC-NOF in Fig. 1. Transmitter i, with  $i \in \{1,2\}$ , communicates with receiver i subject to the interference produced by transmitter j, with  $j \in \{1,2\} \setminus \{i\}$ . There are two independent and uniformly distributed messages,  $W_i \in \mathcal{W}_i$ , with  $\mathcal{W}_i = \{1,2,\ldots,2^{NR_i}\}$ , where N denotes the block-length in channel uses and  $R_i$  is the transmission rate in bits per channel use. At each block, transmitter i sends the codeword  $\mathbf{X}_i = (X_{i,1}, X_{i,2}, \ldots, X_{i,N})^{\mathsf{T}} \in \mathcal{X}_i^N$ , where  $\mathcal{X}_i$  and  $\mathcal{X}_i^N$  are respectively the channel-input alphabet and the codebook of transmitter i.

The channel coefficient from transmitter j to receiver i is denoted by  $h_{ij}$ ; the channel coefficient from transmitter i to receiver i is denoted by  $\overrightarrow{h}_{ii}$ ; and the channel coefficient from channel-output i to transmitter i is denoted by  $\overleftarrow{h}_{ii}$ . All channel coefficients are assumed to be non-negative real numbers. At a given channel use  $n \in \{\underbrace{1,2,\ldots,N}\}$ , the channel output at receiver i is denoted by  $\overleftarrow{Y}_{i,n}$ . During channel use n, the input-output relation of the channel model is given by:

$$\overrightarrow{Y}_{i,n} = \overrightarrow{h}_{ii} X_{i,n} + h_{ij} X_{j,n} + \overrightarrow{Z}_{i,n}, \tag{1}$$

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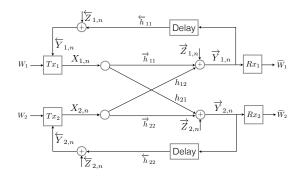


Fig. 1. Gaussian interference channel with noisy channel-output feedback at channel use n.

where  $\overrightarrow{Z}_{i,n}$  is a real Gaussian random variable with zero mean and unit variance that represents the noise at the input of receiver i. Let d>0 be the finite feedback delay in channel uses. At the end of channel use n, transmitter i observes  $\overleftarrow{Y}_{i,n}$ , which consists of a scaled and noisy version of  $\overrightarrow{Y}_{i,n-d}$ . More specifically,

$$\overleftarrow{Y}_{i,n} = \begin{cases} \overleftarrow{Z}_{i,n} & \text{for } n \in \{1,2,\dots,d\} \\ \overleftarrow{h}_{ii} \overrightarrow{Y}_{i,n-d} + \overleftarrow{Z}_{i,n}, & \text{for } n \in \{d+1,d+2,\dots,N\}, \end{cases} (2)$$

where  $\overleftarrow{Z}_{i,n}$  is a real Gaussian random variable with zero mean and unit variance that represents the noise in the feedback link of transmitter-receiver pair i. The random variables  $\overrightarrow{Z}_{i,n}$  and  $\overleftarrow{Z}_{i,n}$  are independent and identically distributed. In the following, without loss of generality, the feedback delay is assumed to be one channel use (d=1). The encoder of transmitter i is defined by a set of deterministic functions  $f_i^{(1)},\ldots,f_i^{(N)}$ , with  $f_i^{(1)}:\mathcal{W}_i\to\mathcal{X}_i$  and for all  $n\in\{2,\ldots,N\}$ ,  $f_i^{(n)}:\mathcal{W}_i\times\mathbb{R}^{n-1}\to\mathcal{X}_i$ , such that

$$X_{i,1} = f_i^{(1)}(W_i)$$
, and (3a)

$$X_{i,1} = f_i \quad (W_i), \text{ and}$$

$$X_{i,n} = f_i^{(n)} \left( W_i, \overleftarrow{Y}_{i,1}, \dots, \overleftarrow{Y}_{i,n-1} \right).$$
(3b)

The components of the input vector  $X_i$  are real numbers subject to an average power constraint:

$$\frac{1}{N} \sum_{n=1}^{N} \mathbb{E}\left(X_{i,n}^{2}\right) \le 1,\tag{4}$$

where the expectation is taken over the joint distribution of the message indexes  $W_1$ ,  $W_2$ , and the noise terms, i.e.,  $\overrightarrow{Z}_1$ ,  $\overrightarrow{Z}_2$ ,  $\overrightarrow{Z}_1$ , and  $\overleftarrow{Z}_2$ . The dependence of  $X_{i,n}$  on  $W_1$ ,  $W_2$ , and the previously observed noise realizations is due to the effect of feedback as shown in (2) and (3).

Assume that during a given communication, T blocks are transmitted. Hence, the decoder of receiver i is defined by a deterministic function  $\psi_i:\mathbb{R}_i^{NT}\to\mathcal{W}_i^T$ . At the end of the communication, receiver i uses the sequence  $(\overrightarrow{Y}_{i,1},\overrightarrow{Y}_{i,2},\ldots,\overrightarrow{Y}_{i,NT})$  to obtain an estimation of the message indices

$$(\widehat{W}_{i}^{(1)}, \widehat{W}_{i}^{(2)}, \dots, \widehat{W}_{i}^{(T)}) = \psi_{i}(\overrightarrow{Y}_{i,1}, \overrightarrow{Y}_{i,2}, \dots, \overrightarrow{Y}_{i,NT}), (5)$$

where  $\widehat{W}_i^{(t)}$  is an estimation of the message index sent during block  $t \in \{1, 2, \dots, T\}$ . The decoding error probability in the two-user G-IC-NOF during block t, denoted by  $P_e^{(t)}(N)$ , is given by

$$P_e^{(t)}\!(N) \!=\! \max\! \left( \Pr\!\left[ \widehat{W_1}^{(t)} \!\neq\! W_1^{(t)} \right], \Pr\!\left[ \left( \widehat{W_2}^{(t)} \!\neq\! W_2^{(t)} \right] \right) \!. \left( 6 \right)$$

The definition of an achievable rate pair  $(R_1, R_2) \in \mathbb{R}^2_+$  is given below.

Definition 1 (Achievable Rate Pairs): A rate pair  $(R_1, R_2) \in \mathbb{R}^2_+$  is achievable if there exists at least one pair of codebooks  $\mathcal{X}_1^N$  and  $\mathcal{X}_2^N$  with codewords of length N, and the corresponding encoding functions  $f_1^{(1)}, \ldots, f_1^{(N)}$  and  $f_2^{(1)}, \ldots, f_2^{(N)}$  such that the decoding error probability  $P_e^{(t)}(N)$  can be made arbitrarily small by letting the blocklength N grow to infinity, for all blocks  $t \in \{1, \ldots, T\}$ .

The two-user G-IC-NOF in Fig. 1 can be fully described by six parameters:  $\overrightarrow{SNR}_i$ ,  $\overrightarrow{SNR}_i$ , and  $INR_{ij}$ , with  $i \in \{1,2\}$  and  $j \in \{1,2\} \setminus \{i\}$ , which are defined as follows

$$\overrightarrow{SNR}_i = \overrightarrow{h}_{ii}^2, \tag{7}$$

$$INR_{ij} = h_{ij}^2$$
 and (8)

$$\overline{SNR}_i = \overline{h}_{ii}^2 \left( \overline{h}_{ii}^2 + 2 \overline{h}_{ii} h_{ij} + h_{ij}^2 + 1 \right).$$
(9)

#### III. MAIN RESULTS

This section introduces an achievable region (Theorem 1) and a converse region (Theorem 2), denoted by  $\underline{\mathcal{C}}_{G-IC-NOF}$  and  $\overline{\mathcal{C}}_{G-IC-NOF}$  respectively, for the two-user G-IC-NOF with fixed parameters  $\overline{SNR}_1$ ,  $\overline{SNR}_2$ ,  $\overline{INR}_{12}$ ,  $\overline{INR}_{21}$ ,  $\overline{SNR}_1$ , and  $\overline{SNR}_2$ . In general, the capacity region of a given multi-user channel is said to be approximated to within a constant gap according to the following definition.

Definition 2 (Approximation to within  $\xi$  units): A closed and convex set  $\mathcal{T} \subset \mathbb{R}^m_+$  is approximated to within  $\xi$  units by the sets  $\underline{\mathcal{T}}$  and  $\overline{\mathcal{T}}$  if  $\underline{\mathcal{T}} \subseteq \mathcal{T} \subseteq \overline{\mathcal{T}}$  and for all  $\mathbf{t} = (t_1, \dots, t_m) \in \overline{\mathcal{T}}$  then  $((t_1 - \xi)^+, \dots, (t_m - \xi)^+) \in \underline{\mathcal{T}}$ .

Denote by  $\mathcal{C}_{\mathrm{GIC-NOF}}$  the capacity region of the 2-user G-IC-NOF. The achievable region  $\underline{\mathcal{C}}_{\mathrm{G-IC-NOF}}$  and the converse region  $\overline{\mathcal{C}}_{\mathrm{G-IC-NOF}}$  approximate the capacity region  $\mathcal{C}_{\mathrm{GIC-NOF}}$  to within 4.4 bits per channel use (Theorem 3).

#### A. An Achievable Region for the Two-User G-IC-NOF

The description of the achievable region  $\underline{C}_{\mathrm{G-IC-NOF}}$  is presented using the constants  $a_{1,i}$ ; the functions  $a_{2,i}:[0,1]\to\mathbb{R}_+$ ,  $a_{l,i}:[0,1]^2\to\mathbb{R}_+$ , with  $l\in\{3,\ldots,6\}$ ; and  $a_{7,i}:[0,1]^3\to\mathbb{R}_+$ , which are defined as follows, for all  $i\in\{1,2\}$ , with  $j\in\{1,2\}\setminus\{i\}$ :

$$a_{1,i} = \frac{1}{2} \log \left( 2 + \frac{\overrightarrow{\text{SNR}_i}}{\text{INR}_{ji}} \right) - \frac{1}{2},$$
 (10a)

$$a_{2,i}(\rho) = \frac{1}{2}\log\left(b_{1,i}(\rho) + 1\right) - \frac{1}{2},$$
 (10b)

$$a_{3,i}(\rho,\mu) = \frac{1}{2} \log \left( \frac{\overleftarrow{\text{SNR}}_i \Big( \text{INR}_{ij} + 1 \Big) + b_{1,i}(\rho) + 1}{\overleftarrow{\text{SNR}}_i \Big( (1-\mu)b_{2,i}(\rho) + 2 \Big) + b_{1,i}(\rho) + 1} \right), \tag{10c}$$

$$a_{4,i}(\rho,\mu) = \frac{1}{2} \log \left( \left( 1 - \mu \right) b_{2,i}(\rho) + 2 \right) - \frac{1}{2},$$
 (10d)

$$a_{5,i}(\rho,\mu) = \frac{1}{2} \log \left( 2 + \frac{\overrightarrow{SNR}_i}{INR_{ji}} + \left( 1 - \mu \right) b_{2,i}(\rho) \right) - \frac{1}{2}, \tag{10e}$$

$$a_{6,i}(\rho,\mu) = \frac{1}{2} \log \left( \frac{\overrightarrow{\text{SNR}}_i}{\overrightarrow{\text{INR}}_{ji}} \left( \left( 1 - \mu \right) b_{2,j}(\rho) + 1 \right) + 2 \right) - \frac{1}{2}, \tag{10f}$$

$$a_{7,i}(\rho,\mu_1,\mu_2) = \frac{1}{2} \log \left( \frac{\overrightarrow{\text{SNR}}_i}{\overrightarrow{\text{INR}}_{ji}} \left( \left( 1 - \mu_i \right) b_{2,j}(\rho) + 1 \right) \right)$$

$$+\left(1-\mu_j\right)b_{2,i}(\rho)+2\left)-\frac{1}{2}.$$
 (10g)

where the functions  $b_{l,i}:[0,1]\to\mathbb{R}_+$ , with  $(l,i)\in\{1,2\}^2$  are defined as follows

$$b_{1,i}(\rho) = \overrightarrow{\mathrm{SNR}}_i + 2\rho\sqrt{\overrightarrow{\mathrm{SNR}}_i\mathrm{INR}_{ij}} + \mathrm{INR}_{ij}$$
 and (11a)

$$b_{2,i}(\rho) = \left(1 - \rho\right) INR_{ij} - 1, \tag{11b}$$

with  $j \in \{1, 2\} \setminus \{i\}$ .

Note that the functions in (10) and (11) depend on  $\overrightarrow{SNR}_1$ ,  $\overrightarrow{SNR}_2$ ,  $\overrightarrow{INR}_{12}$ ,  $\overrightarrow{INR}_{21}$ ,  $\overrightarrow{SNR}_1$ , and  $\overrightarrow{SNR}_2$ , however as these parameters are fixed in this analysis, this dependence is not emphasized in the definition of these functions. Finally, using this notation, Theorem 1 is presented on the next page.

*Proof:* The proof of Theorem 1 is presented in [1].

#### B. Comments on the Achievability

The achievable region is obtained using a random coding argument and combining three classical tools: rate splitting, superposition coding, and backward decoding. This coding scheme is described in [1] and it is specially designed for the two-user IC-NOF. Consequently, only the strictly needed number of superposition code-layers is used. Other achievable schemes, as reported in [2], can also be obtained as special cases of the more general scheme presented in [3]. However, in this more general case, the resulting code for the IC-NOF contains a handful of unnecessary superposing code-layers, which complicates the error probability analysis.

#### C. A Converse Region for the Two-User G-IC-NOF

The description of the converse region  $\overline{\mathcal{C}}_{\mathrm{G-IC-NOF}}$  is determined by the ratios  $\frac{\mathrm{INR}_{ij}}{\overline{\mathrm{SNR}}_j}$ , and  $\frac{\mathrm{INR}_{ji}}{\overline{\mathrm{SNR}}_j}$ , for all  $i \in \{1,2\}$ , with  $j \in \{1,2\} \setminus \{i\}$ . All relevant scenarios regarding these

Theorem 1: The capacity region  $\mathcal{C}_{\text{GIC-NOF}}$  contains the region  $\mathcal{C}_{\text{G-IC-NOF}}$  given by the closure of the set of all possible non-negative achievable rate pairs  $(R_1, R_2)$  that satisfy

$$R_1 \leq \min\left(a_{2,1}(\rho), a_{6,1}(\rho, \mu_1) + a_{3,2}(\rho, \mu_1), a_{1,1} + a_{3,2}(\rho, \mu_1) + a_{4,2}(\rho, \mu_1)\right),\tag{12a}$$

$$R_2 \leq \min\left(a_{2,2}(\rho), a_{3,1}(\rho, \mu_2) + a_{6,2}(\rho, \mu_2), a_{3,1}(\rho, \mu_2) + a_{4,1}(\rho, \mu_2) + a_{1,2}\right),\tag{12b}$$

$$R_1 + R_2 \leq \min \left( a_{2,1}(\rho) + a_{1,2}, a_{1,1} + a_{2,2}(\rho), a_{3,1}(\rho, \mu_2) + a_{1,1} + a_{3,2}(\rho, \mu_1) + a_{7,2}(\rho, \mu_1, \mu_2), a_{1,1} + a_{2,2}(\rho, \mu_1) + a_{1,2}(\rho, \mu_1)$$

$$a_{3,1}(\rho,\mu_2) + a_{5,1}(\rho,\mu_2) + a_{3,2}(\rho,\mu_1) + a_{5,2}(\rho,\mu_1), a_{3,1}(\rho,\mu_2) + a_{7,1}(\rho,\mu_1,\mu_2) + a_{3,2}(\rho,\mu_1) + a_{1,2}\Big), \quad (12c)$$

$$2R_1 + R_2 \leqslant \min\left(a_{2,1}(\rho) + a_{1,1} + a_{3,2}(\rho, \mu_1) + a_{7,2}(\rho, \mu_1, \mu_2),\right)$$
(12d)

$$a_{3,1}(\rho,\mu_2) + a_{1,1} + a_{7,1}(\rho,\mu_1,\mu_2) + 2a_{3,2}(\rho,\mu_1) + a_{5,2}(\rho,\mu_1), a_{2,1}(\rho) + a_{1,1} + a_{3,2}(\rho,\mu_1) + a_{5,2}(\rho,\mu_1)$$

$$R_1 + 2R_2 \leqslant \min \left( a_{3,1}(\rho, \mu_2) + a_{5,1}(\rho, \mu_2) + a_{2,2}(\rho) + a_{1,2}, a_{3,1}(\rho, \mu_2) + a_{7,1}(\rho, \mu_1, \mu_2) + a_{2,2}(\rho) + a_{1,2}, a_{3,1}(\rho, \mu_2) + a_{7,1}(\rho, \mu_1, \mu_2) + a_{7,1}(\rho, \mu_2) +$$

$$2a_{3,1}(\rho,\mu_2) + a_{5,1}(\rho,\mu_2) + a_{3,2}(\rho,\mu_1) + a_{1,2} + a_{7,2}(\rho,\mu_1,\mu_2),$$
(12e)

with 
$$(\rho, \mu_1, \mu_2) \in \left[0, \left(1 - \max\left(\frac{1}{INR_{12}}, \frac{1}{INR_{21}}\right)\right)^+\right] \times [0, 1] \times [0, 1]^2$$
.

ratios are described by two events denoted by  $S_{l_1,1}$  and  $S_{l_2,2}$ , where  $(l_1, l_2) \in \{1, \dots, 5\}^2$ :

$$S_{1,i}$$
:  $\overrightarrow{SNR}_{i} < \min(INR_{ij}, INR_{ji})$ , (13a)

$$S_{2,i}$$
:  $INR_{ji} \leqslant \overrightarrow{SNR}_j < INR_{ij}$ , (13b)

$$S_{3,i}$$
:  $INR_{ij} \leqslant \overrightarrow{SNR}_{j} < INR_{ji}$ , (13c)

$$S_{4,i}$$
:  $\max(INR_{ij}, INR_{ji}) \leqslant \overrightarrow{SNR}_{j} < INR_{ij}INR_{ji}$ , (13d)

$$S_{5,i}$$
:  $\overrightarrow{SNR}_i \ge \max(INR_{ij}, INR_{ij}, INR_{ij}INR_{ii})$ . (13e)

Note that for all  $i \in \{1, 2\}$ , the events  $S_{1,i}, S_{2,i}, S_{3,i}, S_{4,i}$ , and  $S_{5,i}$  are mutually exclusive. This observation shows that given any 4-tuple ( $SNR_1$ ,  $SNR_2$ ,  $INR_{12}$ ,  $INR_{21}$ ), there always exists one and only one pair of events  $(S_{l_1,1}, S_{l_2,2})$ , with  $(l_1, l_2) \in \{1, \dots, 5\}^2$ , that identifies a unique scenario. Note also that the pairs of events  $(S_{2,1}, S_{2,2})$  and  $(S_{3,1}, S_{3,2})$  are not feasible. In view of this, twenty-three different scenarios can be identified using the events in (13). Once the exact scenario is identified, the converse region is described using the functions  $\kappa_{l,i} : [0,1] \to \mathbb{R}_+$ , with  $(l,i) \in \{1,\ldots,3\} \times \{1,2\}$ ;  $\kappa_l: [0,1] \to \mathbb{R}_+, \text{ with } l \in \{4,5\}; \ \kappa_{6,l}: [0,1] \to \mathbb{R}_+, \text{ with }$  $l \in \{1, \ldots, 4\}$ ; and  $\kappa_{7,i,l} : [0,1] \to \mathbb{R}_+$ , with  $(i,l) \in \{1,2\}^2$ . These functions are defined as follows for all  $i \in \{1, 2\}$ , with  $j \in \{1, 2\} \setminus \{i\}$ 

$$\kappa_{1,i}(\rho) = \frac{1}{2} \log \left( b_{1,i}(\rho) + 1 \right), \tag{14a}$$

$$\kappa_{2,i}(\rho) = \frac{1}{2} \log \left( 1 + b_{5,j}(\rho) \right) + \frac{1}{2} \log \left( 1 + \frac{b_{4,i}(\rho)}{1 + b_{5,j}(\rho)} \right), \tag{14b}$$

$$\kappa_{3,i}(\rho) = \frac{1}{2} \log \left( \frac{\overleftarrow{\text{SNR}}_{j} \left( b_{4,i}(\rho) + b_{5,j}(\rho) + 1 \right)}{\left( b_{1,j}(\rho) + 1 \right) \left( b_{4,i}(\rho) + \left( 1 - \rho^{2} \right) \right)} + 1 \right)$$

$$+ \frac{1}{2} \log \left( b_{4,i}(\rho) + 1 \right), \tag{14c}$$

$$\kappa_{3,i}(\rho) = \frac{1}{2} \log \left( \frac{\overleftarrow{\text{SNR}}_{j} \left( b_{4,i}(\rho) + b_{5,j}(\rho) + 1 \right)}{\left( b_{1,j}(\rho) + 1 \right) \left( b_{4,i}(\rho) + \left( 1 - \rho^{2} \right) \right)} + 1 \right) + \frac{1}{2} \log \left( b_{4,i}(\rho) + 1 \right), \tag{14c}$$

$$\kappa_{4}(\rho) = \frac{1}{2} \log \left( 1 + \frac{b_{4,1}(\rho)}{1 + b_{5,2}(\rho)} \right) + \frac{1}{2} \log \left( b_{1,2}(\rho) + 1 \right), \tag{14d}$$

$$\kappa_{5}(\rho) = \frac{1}{2} \log \left( 1 + \frac{b_{4,2}(\rho)}{1 + b_{5,1}(\rho)} \right) + \frac{1}{2} \log \left( b_{1,1}(\rho) + 1 \right), \quad (14e)$$

$$\kappa_{6}(\rho) = \begin{cases}
\kappa_{6,1}(\rho) & \text{if } (S_{1,2} \vee S_{2,2} \vee S_{5,2}) \\
& \wedge (S_{1,1} \vee S_{2,1} \vee S_{5,1}) \\
\kappa_{6,2}(\rho) & \text{if } (S_{1,2} \vee S_{2,2} \vee S_{5,2}) \\
& \wedge (S_{3,1} \vee S_{4,1}) \\
\kappa_{6,3}(\rho) & \text{if } (S_{3,2} \vee S_{4,2}) \\
& \wedge (S_{1,1} \vee S_{2,1} \vee S_{5,1}) \\
\kappa_{6,4}(\rho) & \text{if } (S_{3,2} \vee S_{4,2}) \wedge (S_{3,1} \vee S_{4,1})
\end{cases}$$

$$(14f)$$

$$\kappa_{7,i}(\rho) = \begin{cases} \kappa_{7,i,1}(\rho) & \text{if } (S_{1,i} \vee S_{2,i} \vee S_{5,i}) \\ \kappa_{7,i,2}(\rho) & \text{if } (S_{3,i} \vee S_{4,i}) \end{cases}$$
(14g)

$$\kappa_{6,1}(\rho) = \frac{1}{2} \log \left( b_{1,1}(\rho) + b_{5,1}(\rho) INR_{21} \right) - \frac{1}{2} \log \left( 1 + INR_{12} \right)$$

$$+ \frac{1}{2} \log \left( 1 + \frac{b_{5,2}(\rho) \overleftarrow{SNR}_2}{b_{1,2}(\rho) + 1} \right)$$

$$+ \frac{1}{2} \log \left( b_{1,2}(\rho) + b_{5,1}(\rho) INR_{21} \right) - \frac{1}{2} \log \left( 1 + INR_{21} \right)$$

$$+ \frac{1}{2} \log \left( 1 + \frac{b_{5,1}(\rho) \overleftarrow{SNR}_1}{b_{1,1}(\rho) + 1} \right) + \log(2\pi e), \tag{15a}$$

$$\begin{split} \kappa_{6,2}(\rho) &= \frac{1}{2} \log \left( b_{1,1}(\rho) + b_{5,1}(\rho) \text{INR}_{21} \right) - \frac{1}{2} \log \left( 1 + \text{INR}_{12} \right) \\ &+ \frac{1}{2} \log \left( 1 + \frac{b_{5,2}(\rho)}{\overline{\text{SNR}}_2} \left( \text{INR}_{12} + \frac{\overleftarrow{\text{SNR}}_2 \left( b_{3,2} \right)}{b_{1,2}(\rho) + 1} \right) \right) \\ &- \frac{1}{2} \log \left( 1 + \frac{b_{5,1}(\rho) \text{INR}_{21}}{\overline{\text{SNR}}_2} \right) \\ &+ \frac{1}{2} \log \left( 1 + \frac{b_{5,1}(\rho) \overleftarrow{\text{SNR}}_1}{b_{1,1}(\rho) + 1} \right) \\ &+ \frac{1}{2} \log \left( b_{6,2}(\rho) + \frac{b_{5,1}(\rho) \text{INR}_{21}}{\overline{\text{SNR}}_2} \left( \overrightarrow{\text{SNR}}_2 + b_{3,2} \right) \right) \\ &- \frac{1}{2} \log \left( 1 + \text{INR}_{21} \right) + \log(2\pi e), \end{split} \tag{15b}$$

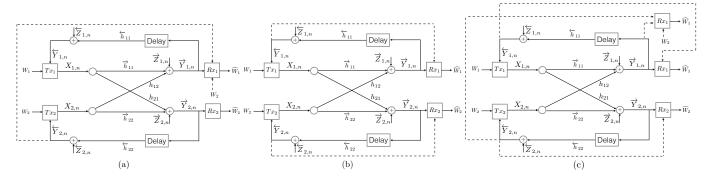


Fig. 2. Genie-Aided G-IC-NOF models for channel use n. (a) Model used to calculate the outer-bound on  $R_1$ ; (b) Model used to calculate the outer-bound on  $R_1 + R_2$ ; and (c) Model used to calculate the outer-bound on  $2R_1 + R_2$ 

$$\kappa_{6,3}(\rho) = \frac{1}{2} \log \left( b_{6,1}(\rho) + INR_{12} \right) \\
+ \frac{b_{5,1}(\rho)INR_{21}}{\overline{SNR}_{1}} \left( \overrightarrow{SNR}_{1} + b_{3,1} \right) \\
- \frac{1}{2} \log \left( 1 + INR_{12} \right) + \frac{1}{2} \log \left( 1 + \frac{b_{5,2}(\rho)\overrightarrow{SNR}_{2}}{b_{1,2}(\rho) + 1} \right) \\
+ \frac{1}{2} \log \left( b_{1,2}(\rho) + b_{5,1}(\rho)INR_{21} \right) - \frac{1}{2} \log \left( 1 + INR_{21} \right) \\
+ \frac{1}{2} \log \left( 1 + \frac{b_{5,1}(\rho)}{\overline{SNR}_{1}} \left( INR_{21} + \frac{\overrightarrow{SNR}_{1}(b_{3,1})}{b_{1,1}(\rho) + 1} \right) \right) \\
- \frac{1}{2} \log \left( 1 + \frac{b_{5,1}(\rho)INR_{21}}{\overline{SNR}_{1}} \right) + \log(2\pi e), \qquad (15c) \\
\kappa_{6,4}(\rho) = \frac{1}{2} \log \left( b_{6,1}(\rho) + \frac{b_{5,1}(\rho)INR_{21}}{\overline{SNR}_{1}} \left( \overrightarrow{SNR}_{1} + b_{3,1} \right) \right) \\
- \frac{1}{2} \log \left( 1 + INR_{12} \right) - \frac{1}{2} \log \left( 1 + INR_{21} \right) \\
+ \frac{1}{2} \log \left( 1 + \frac{b_{5,2}(\rho)}{\overline{SNR}_{2}} \left( INR_{12} + \frac{\overrightarrow{SNR}_{2}(b_{3,2})}{b_{1,2}(\rho) + 1} \right) \right) \\
- \frac{1}{2} \log \left( 1 + \frac{b_{5,1}(\rho)INR_{21}}{\overline{SNR}_{1}} \right) \\
- \frac{1}{2} \log \left( 1 + \frac{b_{5,1}(\rho)INR_{21}}{\overline{SNR}_{1}} \right) \\
+ \frac{1}{2} \log \left( b_{6,2}(\rho) + \frac{b_{5,1}(\rho)INR_{21}}{\overline{SNR}_{2}} \left( \overrightarrow{SNR}_{2} + b_{3,2} \right) \right) \\
+ \frac{1}{2} \log \left( 1 + \frac{b_{5,1}(\rho)}{\overline{SNR}_{1}} \left( INR_{21} + \frac{\overrightarrow{SNR}_{1}(b_{3,1})}{b_{1,1}(\rho) + 1} \right) \right) \\
+ \log(2\pi e), \qquad (15d)$$

and

$$\kappa_{7,i,1}(\rho) = \frac{1}{2} \log \left( b_{1,i}(\rho) + 1 \right) - \frac{1}{2} \log \left( 1 + INR_{ij} \right) 
+ \frac{1}{2} \log \left( 1 + \frac{b_{5,j}(\rho) \dot{S}NR_{j}}{b_{1,j}(\rho) + 1} \right) 
+ \frac{1}{2} \log \left( b_{1,j}(\rho) + b_{5,i}(\rho) INR_{ji} \right) 
+ \frac{1}{2} \log \left( 1 + b_{4,i}(\rho) + b_{5,j}(\rho) \right) - \frac{1}{2} \log \left( 1 + b_{5,j}(\rho) \right) 
+ 2 \log(2\pi e),$$
(16a)

$$\kappa_{7,i,2}(\rho) = \frac{1}{2} \log \left( b_{1,i}(\rho) + 1 \right) - \frac{1}{2} \log \left( 1 + INR_{ij} \right) 
- \frac{1}{2} \log \left( 1 + b_{5,j}(\rho) \right) + \frac{1}{2} \log \left( 1 + b_{4,i}(\rho) + b_{5,j}(\rho) \right) 
+ \frac{1}{2} \log \left( 1 + \left( 1 - \rho^2 \right) \frac{INR_{ji}}{\overrightarrow{SNR}_{j}} \left( INR_{ij} + \frac{\overleftarrow{SNR}_{j} \left( b_{3,j} \right)}{b_{1,j}(\rho) + 1} \right) \right) - \frac{1}{2} \log \left( 1 + \frac{b_{5,i}(\rho)INR_{ji}}{\overrightarrow{SNR}_{j}} \right) 
+ \frac{1}{2} \log \left( b_{6,j}(\rho) + \frac{b_{5,i}(\rho)INR_{ji}}{\overrightarrow{SNR}_{j}} \left( \overrightarrow{SNR}_{j} + b_{3,j} \right) \right) 
+ 2 \log(2\pi e),$$
(16b)

where the functions  $b_{l,i}$ , with  $(l,i) \in \{1,2\}^2$  are defined in (11);  $b_{3,i}$  are constants; and the functions  $b_{l,i} : [0,1] \to \mathbb{R}_+$ , with  $(l,i) \in \{4,5,6\} \times \{1,2\}$  are defined as follows, with  $j \in \{1,2\} \setminus \{i\}$ :

$$b_{3,i} = \overrightarrow{SNR}_i - 2\sqrt{\overrightarrow{SNR}_i INR}_{ii} + INR}_{ii},$$
 (17a)

$$b_{4,i}(\rho) = \left(1 - \rho^2\right) \overrightarrow{SNR}_i, \tag{17b}$$

$$b_{5,i}(\rho) = \left(1 - \rho^2\right) INR_{ij},\tag{17c}$$

$$b_{6,i}(\rho) = \overrightarrow{SNR}_i + INR_{ij} + 2\rho\sqrt{INR_{ij}} \left(\sqrt{\overrightarrow{SNR}_i} - \sqrt{INR_{ji}}\right)$$

$$+\frac{\overline{\text{INR}_{ij}}\sqrt{\overline{\text{INR}_{ji}}}}{\overrightarrow{\text{SNR}_{i}}}\left(\sqrt{\overline{\text{INR}_{ji}}}-2\sqrt{\overrightarrow{\text{SNR}_{i}}}\right). \quad (17d)$$

Note that the functions in (14), (15), (16) and (17) depend on  $\overrightarrow{SNR}_1$ ,  $\overrightarrow{SNR}_2$ ,  $\overrightarrow{INR}_{12}$ ,  $\overrightarrow{INR}_{21}$ ,  $\overrightarrow{SNR}_1$ , and  $\overrightarrow{SNR}_2$ . However, these parameters are fixed in this analysis, and therefore, this dependence is not emphasized in the definition of these functions. Finally, using this notation, Theorem 2 is presented below.

Theorem 2: The capacity region  $C_{GIC-NOF}$  is contained within the region  $\overline{C}_{G-IC-NOF}$  given by the closure of the set of non-negative rate pairs  $(R_1, R_2)$  that satisfy for all  $i \in \{1, 2\}$ , with  $j \in \{1, 2\} \setminus \{i\}$ :

$$R_i \leq \min \left( \kappa_{1,i}(\rho), \kappa_{2,i}(\rho) \right),$$
 (18a)

$$R_i \leqslant \kappa_{3,i}(\rho),$$
 (18b)

$$R_1 + R_2 \leqslant \min \left( \kappa_4(\rho), \kappa_5(\rho) \right),$$
 (18c)

$$R_1 + R_2 \leqslant \kappa_6(\rho), \tag{18d}$$

$$2R_i + R_i \leqslant \kappa_{7,i}(\rho), \tag{18e}$$

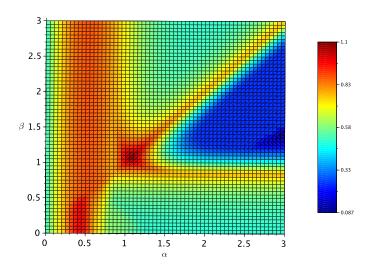


Fig. 3. Gap between the converse region  $\overline{\mathcal{C}}_{G-IC-NOF}$  and the achievable region  $\underline{\mathcal{C}}_{G-IC-NOF}$  of the two-user G-IC-NOF, under symmetric channel conditions, i.e.,  $\overrightarrow{SNR}_1 = \overrightarrow{SNR}_2 = \overrightarrow{SNR}$ ,  $INR_{12} = INR_{21} = INR_2$  and  $\overrightarrow{SNR}_1 = \overrightarrow{SNR}_2 = \overrightarrow{SNR}$ , as a function of  $\alpha = \frac{\log INR}{\log \overrightarrow{SNR}}$  and  $\beta = \frac{\log \overrightarrow{SNR}}{\log \overrightarrow{SNR}}$ .

with  $\rho \in [0,1]$ .

*Proof:* The proof of Theorem 2 is presented in [1].

#### D. Comments on the Converse Region

The outer bounds (18a) and (18c) correspond to the outer bounds for the case of perfect channel-output feedback [4]. The bounds (18b), (18d) and (18e) correspond to new outer bounds that generalize those presented in [2] for the two-user symmetric G-IC-NOF. These new outer-bounds were obtained using the genie-aided models shown in Fig. 2.

## E. A Gap Between the Achievable Region and the Converse Region

Theorem 3 describes the gap between the achievable region  $\underline{\mathcal{C}}_{\mathrm{G-IC-NOF}}$  and the converse region  $\overline{\mathcal{C}}_{\mathrm{G-IC-NOF}}$  using the approximation notion described in Definition 2.

Theorem 3: The capacity region of the two-user G-IC-NOF is approximated to within 4.4 bits per channel use by the achievable region  $\underline{\mathcal{C}}_{\mathrm{G-IC-NOF}}$  and the converse region  $\overline{\mathcal{C}}_{\mathrm{G-IC-NOF}}$ .

*Proof:* The proof of Theorem 3 is presented in [1].  $\blacksquare$  The gap, denoted by  $\delta$ , between the sets  $\overline{\mathcal{C}}_{G-IC-NOF}$  and  $\underline{\mathcal{C}}_{G-IC-NOF}$  can be approximated (Definition 2) as follows

$$\delta \leqslant \max\left(\delta_{R_1}, \delta_{R_2}, \frac{\delta_{2R}}{2}, \frac{\delta_{3R_1}}{3}, \frac{\delta_{3R_2}}{3}\right),$$
 (19)

where

$$\delta_{R_{1}} \triangleq \min\left(\kappa_{1,1}(\rho), \kappa_{2,1}(\rho), \kappa_{3,1}(\rho)\right) - \min\left(a_{2,1}(\rho), a_{6,1}(\rho, \mu_{1}) + a_{3,2}(\rho, \mu_{1}), a_{1,1} + a_{3,2}(\rho, \mu_{1}) + a_{4,2}(\rho, \mu_{1})\right),$$

$$(20a)$$

$$\delta_{R_{2}} \triangleq \min\left(\kappa_{1,2}(\rho), \kappa_{2,2}(\rho), \kappa_{3,2}(\rho)\right) - \min\left(a_{2,2}(\rho), a_{3,1}(\rho, \mu_{2}) + a_{6,2}(\rho, \mu_{2}), a_{3,1}(\rho, \mu_{2}) + a_{4,1}(\rho, \mu_{2}) + a_{1,2}\right),$$

$$(20b)$$

$$\begin{split} \delta_{2R} &\triangleq \min \left( \kappa_4(\rho), \kappa_5(\rho), \kappa_6(\rho) \right) - \min \left( a_{2,1}(\rho) + a_{1,2}, \\ a_{1,1} + a_{2,2}(\rho), \\ a_{3,1}(\rho, \mu_2) + a_{1,1} + a_{3,2}(\rho, \mu_1) + a_{7,2}(\rho, \mu_1, \mu_2), \\ a_{3,1}(\rho, \mu_2) + a_{5,1}(\rho, \mu_2) + a_{3,2}(\rho, \mu_1) + a_{5,2}(\rho, \mu_1), \\ a_{3,1}(\rho, \mu_2) + a_{7,1}(\rho, \mu_1, \mu_2) + a_{3,2}(\rho, \mu_1) + a_{1,2} \right), (20c) \end{split}$$

$$\delta_{3R_{1}} \triangleq \kappa_{7,1}(\rho) - \min\left(a_{2,1}(\rho) + a_{1,1} + a_{3,2}(\rho, \mu_{1}) + a_{7,2}(\rho, \mu_{1}, \mu_{2}), a_{3,1}(\rho, \mu_{2}) + a_{1,1} + a_{7,1}(\rho, \mu_{1}, \mu_{2}) + 2a_{3,2}(\rho, \mu_{1}) + a_{5,2}(\rho, \mu_{1}), a_{2,1}(\rho) + a_{1,1} + a_{3,2}(\rho, \mu_{1}) + a_{5,2}(\rho, \mu_{1})\right),$$
(20d)

$$\begin{split} \delta_{3R_2} &\triangleq \kappa_{7,2}(\rho) - \min \Big( a_{3,1}(\rho, \mu_2) + a_{5,1}(\rho, \mu_2) + a_{2,2}(\rho) \\ &+ a_{1,2}, a_{3,1}(\rho, \mu_2) + a_{7,1}(\rho, \mu_1, \mu_2) + a_{2,2}(\rho) + a_{1,2}, \\ &2 a_{3,1}(\rho, \mu_2) + a_{5,1}(\rho, \mu_2) + a_{3,2}(\rho, \mu_1) + a_{1,2} \\ &+ a_{7,2}(\rho, \mu_1, \mu_2) \Big). \end{split} \tag{20e}$$

Note that  $\delta_{R_1}$  and  $\delta_{R_2}$  represent the gap between the active achievable single-rate bound and the active converse single-rate bound;  $\delta_{2R}$  represents the gap between the active achievable sum-rate bound and the active converse sum-rate bound; and,  $\delta_{3R_1}$  and  $\delta_{3R_2}$  represent the gap between the active achievable weighted sum-rate bound and the active converse weighted sum-rate bound.

Finally, it is important to highlight that, as suggested in [2], [4], and [5], the gap between  $\mathcal{L}_{\text{G-IC-NOF}}$  and  $\overline{\mathcal{C}}_{\text{G-IC-NOF}}$  can be calculated more precisely. However, the choice in (19) eases the calculations at the expense of less precision.

Fig. 3 presents the exact gap existing between the achievable region  $\mathcal{C}_{\mathrm{G-IC-NOF}}$  and the converse region  $\overline{\mathcal{C}}_{\mathrm{G-IC-NOF}}$  for the case in which  $\overline{\mathrm{SNR}}_1 = \overline{\mathrm{SNR}}_2 = \overline{\mathrm{SNR}}$ ,  $\mathrm{INR}_{12} = \mathrm{INR}_{21} = \mathrm{INR}$ , and  $\overline{\mathrm{SNR}}_1 = \overline{\mathrm{SNR}}_2 = \overline{\mathrm{SNR}}$  as a function of  $\alpha = \frac{\log \mathrm{INR}}{\log \overline{\mathrm{SNR}}}$  and  $\beta = \frac{\log \overline{\mathrm{SNR}}}{\log \overline{\mathrm{SNR}}}$ . Note that in this case, the maximum gap is 1.1 bits per channel use and occurs when  $\alpha = \beta = 1.1$ .

#### IV. CONCLUSIONS

An achievable region and a converse region for the twouser G-IC-NOF have been introduced. It has been shown that these regions approximate the capacity region of the two-user G-IC-NOF to within 4.4 bits per channel use.

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